# TP2-KYBER

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# 1 Trabalho prático 2 - Estruturas Criptográficas

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### Grupo 12

## 1.1 Kyber

A última técnica a implementar tem como objetivo a implementação de um KEM IND-CPA e um PKE IND-CCA, do protótipo **Crystalis Kyber**. De seguida, são apresentadas em duas seções (PKE-IND-CPA e PKE-IND-CCA) os resultados da implementação de cada ténica. Esta resolução foi construída com base no documento *kyber.pdf* mais recente do site disponibilizado pela equipa docente.

#### 1.1.1 KEM-IND-CPA

Esta versão permite obter uma segurança do tipo IND-CPA (Chosen Plaintext Attacks). Em primeiro lugar foram desenvolvidas funções auxiliares de implementação de aritmética, encode, decode entre outras:

- *parse*: recebe como *input* um conjunto de *bytes* e retorna o polinómio correspondente a esse conjunto;
- XOF: corresponde a uma função do tipo extendable output function, utilizando SHAKE-128;
- PRF: corresponde a uma função do tipo pseudorandom function, utilizando o SHAKE-256;
- G: função de hash construída com base no SHA3-512;
- *encode*: recebe um polinómio como argumento e retorna um *byte array* correspondente ao polinómio;
- decode: função inversa da encode recebe um byte array e retorna o polinómio correspondente;
- compress: comprime um polinómio e retorna os bytes correspondentes;
- *decompress*: função inversa da *compress*. Descomprime *bytes* e retorna o polinómio correspondente;
- transposta: calcula a transposta de uma matriz.

As funções principais centram-se na geração de chaves, cifragem da mensagem passada como parâmetro e posterior decifragem do texto cifrado, obtendo deste modo, a mensagem original:

- *gerar\_chaves*: com recurso às funções anteriormente apresentadas, ocorre a geração das chaves pública e secreta. A primeira é essencial para a cifragem da mensagem e a segunda para a decifragem do criptograma;
- *cifragem*: tem como objetivo principal a cifragem de uma mensagem. Desta forma, recebe como parâmetros a chave pública, a mensagem e *coins* (*bytes* aleatórios) e dá como *output* o texto cifrado;
- *decifragem*: tem como objetivo decifrar um criptograma, obtendo como resultado o texto limpo correspondente. Recebe como argumentos a chave secreta e o criptograma.

```
[17]: import os
  import random as rn
  from sympy import ntt
  from cryptography.hazmat.primitives import hashes
  import numpy
  from sympy import intt
  import gzip
  import struct
```

```
[18]: # constantes Kyber
      n = 256
      q = 343576577
     k = 2
      n1 = 3
      n2 = 2
      du, dv = 10, 4
      # criação dos anéis
      Z.<w> = ZZ[]
      R.<w> = QuotientRing(_Z ,_Z.ideal(w^n - 1))
      _Q.<_W> = GF(q)[]
      Rq.<w> = QuotientRing(_Q , _Q.ideal(w^n + 1))
      # tamanho necessário para o decompress
      def tamanho(stringB, numberS):
          count = 10
          auxCount = 1
          while i < len(stringB):</pre>
              if numberS == auxCount:
                  i = i + 10
                  while (i < len(stringB)) and (stringB[i] != 31 or stringB[i + 1] !=_{\sqcup}
       4139 or stringB[i + 2] != 8 or stringB[i + 3] != 0 ):
                      count = count + 1
                      i = i + 1
```

```
auxCount = auxCount + 1
        i = i + 1
        if (i + 10) < len(stringB) and (stringB[i] == 31 and stringB[i + 1] == 1
 4139 and stringB[i + 2] == 8 and stringB[i + 3] == 0):
            auxCount = auxCount + 1
        if auxCount > numberS:
            break
    return count
# input: conjunto de bytes; output: polinómio do conjunto de bytes inserido;
def parse(str_bytes):
   result = []
    for i in str_bytes:
        result.append(i)
    return Rq(result)
# XOF, com o SHAKE-128
def XOF(p,i,j):
    digest = hashes.Hash(hashes.SHAKE128(int(32)))
    digest.update(p)
    digest.update(bytes(i))
    digest.update(bytes(j))
    r = digest.finalize()
    return r
# pseudorandom function com SHAKE-256
def PRF(s,b):
    digest = hashes.Hash(hashes.SHAKE256(int(32)))
    digest.update(s)
    digest.update(bytes(b))
    r = digest.finalize()
    return r
# função de hash com SHA3-512;
def G(d):
    digest = hashes.Hash(hashes.SHA512())
    digest.update(bytes(d))
    r = digest.finalize()
    return r
# input: polinómio; output: byte array;
def encode(poly):
    byt=b''
    aux=1
    countX=0
    for j in poly:
```

```
if(j>255):
            aux=2
        if (j > 65025):
            aux = 3
        if (j > 16581375):
            aux = 4
        if (j > 4228250625):
            aux = 5
        byt = byt+ int((_Z(j))).to_bytes( aux, 'big')
        byt = byt +"/-n-/".encode()
        countX =countX +1
    return byt
# input: byte array; output: polinómio correspondente ao byte array do input;
def decode(byt):
    listaCoef = []
    byteAux = b''
    listAux = []
    desc=0
    while desc <byt.__len__():</pre>
        if byt[desc] == 47 and byt[desc+1] == 45 and byt[desc+2] == 110 and _{\sqcup}
 \hookrightarrowbyt[desc+3]==45 and byt[desc+4] == 47 :
            desc = desc+4
            listaCoef.append(int.from_bytes(byteAux, 'big'))
            byteAux = b''
        else:
            byteAux = byteAux + bytearray([int(_Z(byt[desc]))])
        desc = desc+1
    return listaCoef
# comprime um polinómio em bytes
def compress(polinomio):
    polinomioB= encode(polinomio)
    compress = gzip.compress(polinomioB)
    return compress
# descomprime os bytes e transfora-os num polinómio
def decompress(compress):
    unpack = gzip.decompress(compress)
    return Rq(decode(unpack))
# calcula a transposta de uma matriz
def transposta(matrix):
    zipped_rows = zip(*matrix)
    transpose_matrix = [list(row) for row in zipped_rows]
    return transpose_matrix
```

```
# geração do par de chaves (publickey, secretkey)
def gerar_chaves():
    d = bytearray(os.urandom(32))
    p = G(d)[:32]
    teta = G(d)[-32:]
    N = 0
    A = [[0 \text{ for } x \text{ in } range(k-1)] \text{ for } y \text{ in } range(k-1)]
    for i in range(0,k-1):
        for j in range(0,k-1):
            A[i][j] =parse(XOF(p,i,j))
    s = \prod
    for i in range(0,k-1):
        s.append(parse(PRF(teta,N)))
        N = N + 1
    e = []
    for i in range(0,k-1):
        e.append(parse(PRF(teta,N)))
        N = N + 1
    s1 = Rq(T.ntt(s[0]))
    e1 = Rq(T.ntt(e[0]))
    t = A[0][0].lift() * s1.lift() + e1.lift()
    pk = encode(t) + p
    sk = encode(s1)
    return pk, sk
# cifra uma mensagem m
def cifragem(pk, m, coins):
    N = 0
    t2 = pk[:len(pk)-32]
    t= decode(t2)
    p = pk[-32:]
    A = [[0 \text{ for } x \text{ in } range(k-1)] \text{ for } y \text{ in } range(k-1)]
    for i in range(0,k-1):
        for j in range(0,k-1):
             A[i][j] =parse(XOF(p,i,j))
    AT = transposta(A)
    r = []
    for i in range(0,k-1):
        r.append(parse((PRF(bytearray(r),bytearray([N])))))
        N = N + 1
    e1 = []
    for i in range(0,k-1):
        e1.append(parse(PRF(bytearray(r[i]),bytearray([N]))))
        N = N + 1
    e2 = parse(PRF(bytearray(r[0].list()),N))
    r1= Rq(T.ntt(r[0]))
    u= Rq(T.ntt_inv(A[0][0].lift()*r1.lift()))+e1[0]
```

```
v = Rq(T.ntt(Rq(t).lift() * r1.lift())) + e2 + decompress(m)
c1 = compress(u)
c2 = compress(v)
c = c1 + c2
return c

# decifragem do criptograma
def decifragem(sk, ciphertext):
    u = decompress(ciphertext[:tamanho(ciphertext,1)])
    v = decompress(ciphertext[-tamanho(ciphertext,2):])
s1 = Rq(decode(sk))
m = compress(v.lift() - (T.ntt_inv(s1.lift()*Rq(T.ntt(u)).lift())))
return m
```

De seguida são apresentados os resultados obtidos com recurso às funções anteriores. Neste caso, estamos a considerar uma mensagem fixa. Note-se que houve uma dificuldade acrescida nesta implementação pelo que, deparamos-nos com alguns erros não identificados que impossibilitam o correto funcionamento do programa.

Texto cifrado =  $b'\x1f\x8b\x08\x00\xf3\xc8mb\x02\xff5UyT\xd4U\x14\x1ef\x06fa\x16\x06\xf5\x90,\xa1\x89\xec)\xc3"t\x12\xd9\x86MBb$ 

 $x8a\x01X\x95\xcf\xf6\x93\xcf\xef\xe3\xccq\x98\xd2\x97\x94\x0cr\x1b\x90\x93E\xd6n\xf0$  $\x0.07g$  $x99\\x92\\xef\\xf1\\xeej\\x8f\\x8c|%jW\\xb07\\xeb\\x871\\x88k\\xe4\\xd3F&\\xeb\\x8a\\t\\x83\\xb4$  $\xcd\xf3\x11\xbch\xd0\xf2?\xc45\xd7\x90\x8d\x9b\xaa\xb87\t\xc9C\x17\xa3\x18\xa4\$  $x95\xe3TEd0\xfc\x7f$\x9f4\xe0@\x16\x02kZ\xa7\x86\x14\xe7\x8e\xf70\xe8\xfc\xc29\xe3$  $04\xd6\xd8\I\x0c\xead\xe4e\xf8\x03\ndg\xda\x0c\xb2\xc2\x9btA\xbb\xc9u\x84L<\x$  $xab\x8e\n\xefL\xf4) $\xe7\x19\x11\x83\x84\xfd\xfb>\x04\xaa\x8asB\xc2\xb6\xdc\x8b$  $\x88\x01\xc1\xc5\xd2\xc5\xd2\xc7:L\xc2P0\xcb[\xbb\xc1\x1c\xe9s\x04\xcc\xc]$  $xc2\xf6\xfb\x88\Lp\x08,\n\x974H\xa3j|\r\xb6\xd2G\x9f\x80\xaf\xcd\xd8\x0bo\xe4\x8$  $d/xa1/xbf*x8cZ/xc0wz1'/xce/x15;/xbf/x81/xbbI5/xbd/x88/xe0b/xf2(^{8HQ/xb1/xbf})$  $xd6\xa3\xb6c\xc7\x86\xec\x01Gu\x9c\n\xe7G:H\x1f\x1dy5F@j\x9d\xe7\x08\xd4.Cq$\x1e$ \xfc@<\xb5\xd0\x95\xbd&\x94\xe2[ \xb7T\x07\xdc\xa8\r\x8a\$\$\xb0P"m\xfe\xe8\x83L\x  $948\xd4\xf7g\xaa\xcf|?5\x04\xa7\xdbe\x1f\x12<\xd9\xff\x1e8\xde=\x99\x0f\x11o\x18$  $\xdf\xc3\xfb\xfa\x7fP\x82\x9c\xc6L\xb4\x90\xa8\''\x9aLE9\x8b\xe0\xc8\xb9r\x07/\xeq$  $8\xdf\x12\xea\xcc\xba-\xebQ\#\xdf\x0eTK\xae\xb7\xa9\n\xaf\xab\xa5>\xd4\|\w\x921\xb7$  $R\x1bk\x1aj\x1d\xef\&w\xd4\xef\x94\xf8- \xbbh<\xa5+S\x99\x82W9\x8cJ\xeaz\xb9\xff\$  $x0e[V\x0b\r\x12\xe7H\xb77\xf4\x8b\xb3A\x93\xca\xe4\xbe\xb4R8\xe6\xcb\xa4\x1f\x9b$  $8f} xa40 xf9. % xa2 xf0 xa3 xd2 y x8e x86 xe3 x88 x9b { x0e xddz xae xf4*r x8c xc8}$  $\x82\xfc\\x07G\x16\x836\x1aiH\\xad\xa4\xd7q\\xaf9\xd7\x13\xc7\x07C\x92\xd0qe\x967$  $\x18\x14\xd7\xf7\x1dq.\xf3y\x04\xccb\xc2P\xb4\xfd\x85\x02\x94\xe3\x85bL\x00\xff/$  $N k_1d: xa6\\xc0\\xdc?P\\xb8\\x10\\xc3\\xed\\x02H$\\xa8\\x7f\\x84\\xae\\x93\\x05\\x94\\xa2c\\xe$  $40\times0e+\times08\timesc3\timese9\timesd65t$ y.\xce\xf1v\r\x87\xe2F\xc3\xd0+\xf4\xd4\x8bm\xf8n\xc4\  $x9a\xebe\xc8<\x9e\xe3G\xce^\x15\x96\xbe\xbe\xdd,mWn\xeaF1\xc8T\xc5\x98^\xc1\xee@$  $xd4\xba|\x13\\x83\x1b\xb0\xcf\%\xf7e.\x98\xfe \x96A]\xd60\x11\xd7\xc1,\x8c\xb4N$  $X\x82\x06\x81\xfdT\xb8\xe6u\xf9\xf3!)\xb4ys\x06\xb2\xac\\x98\x02Y\xc1m\xfa\xec\$ xb8\xb9\x05\$\xb6\xd6j?\r\x81\xd6\xf6B,\x05Y\x82Z\x87dOH\xc0/&JwN\x85\xb3V\x8e\x7  $f\xd3\xbc\x85>Wn\xb4\x13\xba\<\xa5m\x89V\xc9\xc3/\xc4\xdb3M\r\xcc\xe5E\xd5$  $xe2<xaf\\x02#\\x93\\xe6J\\xcb\\x99S7\\x89\\x95#\\xba\\xaa\\xf9\\x1c\\xb9%\\xb5\\x1e\\xc6\\xe7mF$  $0 = 5 \times 3 \times 4 \times 1 \times 6 \times 1 \times$  $xc3\\x9e\\x0cx\\x07\\x17\\x1e\\xcf\\xe2C\\xd6s\\JB\\xb4\\xb2\\xa0K\\xaf"\\x8e\\xb7\\xb3cQ\\xa4\\xde$  $8\xd1\xe5\x81\xd9\x86\xfeT\xd4u\xbb\xacA\x95\x94\x87\x8d,\x0f\xf3\xa4Z\x8e\x07$  $v\xc0b\xdaL\xf1k\xc6\x18\x16A\x8ay\xc6\xda\xc2\xd7P\xd4\xff\x03?ak\x9a\x9d\t\x0$  $0\x00\x1f\x8b\x00\xf3\xc8mb\x02\xff5VyP\xd4e\x18\xde\x85ewY\xf6\xde\x05\x83\$  $x19\x8e\x81(\xcc\xb8D\x11+\x8e\x8c\x151\x8e\x00\x1bE\&\Y\n,.Ga\x91!\xd1H\x08\x14$  $E!\x12\x87\#\x96\xc3Q@ABG\x11\xec0\x84d\x85i(0!a\'\x8e\x80\x81P\x04D\xd0f\xf6}\xf$  $6\xbfg\xbe\xef\xf3\xbe\xcf{|\xcb\xe0Y\x9e\xb9\xee\xea\x9c\xe4\xec\xca\x1}$  $\x0^x80\x63\x84\x6\xa6\xa13^x848\xa2\x9b\xb7\x08\t\xa3\xfdv\xe$  $3\x17aG\xa1$lz;\x8f\x90Y\x9a\xf5s=2\xd2Io\xe2\xf2\xf1\x94\x1b!\xf9\xfeBs\xba\x8c$ 

 $\xba\xa0\xa5\#q\x8dM\#\x14\xf2\xfa\xe2\x08\t\x86\xf2\x95\x84\xb8\xaa53\xc4\xf4\xf0$  $84\$\x9f7x\xc3\x8b=/\x907g\xb8t\x0e|\x17\x8d\x9e\x02)\x1fh\x08\xf1n4\xfe\x03/2\xe$  $3N\times9\times11[\times2]\times06/\times5\times311[\times2]\times06/\times5\times11[\times2]\times111[\times2]\times06/\times111[\times2]\times11[\times2]\times11$  $x8d\xcc\x9a\x8a\xa3\x1d\xbeJ_\x8d\xfc\x98\x8a\)\x9b\xaa\r\x9a\x2\x96\Y\xc5\v\xd$  $5\x03\xe3\x97\x0b\xdb\xc1\xa0-\xae\xd1#VV\xe0g\x08V\xec\x80\xb0\x05\x81\xbe\x942$  $3\xe2t\x08\xf2\xoc\xef\xc8\x85\x1f\xe6\x91\xe9\xe0\x0fm\xad\xd7\xa3c\xe5\xc9\xd0$  $\x1e\xb8=\x04\x0f\xde\x19\x9f\xc9\x18.\x1d\xb1B\x130@(\xfd]8\xcfN\xfc$\x14^d\xf$  $48 \times 0 \times 0 = 1 \times 48 \times 0 \times 0 = 1 \times 48 \times 0 \times 0 = 1 \times 10 \times 0 = 1 \times 10 \times 0 \times 0 = 1 \times 10 \times 0 = 1$  $xcb\xd5\x16\x18\xa4\xa9\xf4\x84\xd9\x11\n\x89\xf5t\xf0"xK\x17\xd0\x83\xdc\x01[T]$  $Eh\times x = x^{x} x b f x b x 1 f x a 1 x c a x 9 c x c d x x 1 x c a (x 0 4 x d 7 x a b 0 x 9 b x 8 9 x 1 d x 1 x c a x 6 x b 1 x c a x 6 x b 1 x a 1 x c a x 6 x b 1 x a 1 x c a x 6 x b 1 x a 1 x c a x 6 x b 1 x a 1 x c a x 6 x b 1 x a 1 x c a x 6 x b 1 x a 1 x c a x 6 x b 1 x a 1 x$  $9\xedPH\xbe\xb4\x02\xfd\x84\xa2\x17\xb0\xe6\xfe\xf8\xd8\xfb\xef\xfb\x90\x95\x07$  $xf4x81x85\{v\\xbe\\x00g\\x1evT5\\x9f\\x01\\x1a\\x1a\\xe6\\x13V\\x1f\\x12\\x1c\\x1c\\xa5:\\x9a$  $xf3\x8b[\xab1$f\xdc\xf6@\xe28re/\x81\xc9<\x06\x84v\xaeR\xf9\xb5q\xf3\xf0)\xcb\xc$  $3\x19v\xca9D\xf4\x8b\x96\x93\xe8\xec\xb2\xfdHq\xd9\x19\x03\xca\xefdV''\x9c-\x8d\xeq$  $xeb\\xe0\\xab\\xd2\\x07\\x03/*\\xcfF\\x9b\\xf0:J$\\xc4V\\x9fd(u\\xf2qw\\xd4\\x90;u\\x0c$\\x9b\\x$  $9c^x05Iz_+$x82] xb1x17x04q xedb x04w xa2xb9x1f,?xf8xa3x9cxf2-xbegxc9$ \xd2\x9eH{J\xd4Q\x83\x89\xe20>,\x81\x96\xb6\xff{\xb0\xed\x99\x10\x80\xe3\xb5\x93  $\x85\x0f\x02\x94\xc2 \xaa\xf6r\x0bU\x83\xd9\xca.\xc5W\xea\xaa}p\xa1\x8bg\x07\x8e$  $\x98) \times 14YV \times b \times 86 \times c2K2 \times 9e \times d0d \times 9ad8 + x0c \times f1 \times 82 \times 8e \times 7x9f \times bas] P \times ef \times de$  $\xef\x90\x15\xa7\x82S\x8b\xca\x1ce~I)\xa7\xf2\r\x191\xec\xc1\x96\xe4\xb9h\xbf\xa$  $2\x90\xe2\xab\xd3\xd1\xca\x98aM\x0b5R\xc8_\xd5\x19\x16\xca\x1b\x8bPx\xef\xd3sH$  $b\xaa\xe5*\xceJF\xff\x04q0\xdd0\xb8\xcd\x9d\x1bEb)N\xd4\xdaF\xd7\x870z\x92\xf0\x$  $df\x13\x0ctMi\xd8\x97\xd7Pp\xa9M\x00\xd6*7\xdf\x9d,1.\x8bF\xff\x88\xda\xfck\x15$  $f\x13\xdb\x9d\x7f\xb0\xef\x00\x0c\x8f*E?\x89\x9e\xb5\xc4\xc4uQXS\\\x91\x12/\x90$  $\xf88\x0b0\xa5Y}\xf0\x88\xe1E\xf3@\xb7K\xef5\xa2\x1dd\x91I\xab\xf8E\xb7\xd5\x07p$  $f\&1\\x822\\xd5.\\x95\\xf7i\\x95\\xc1\\xb5\\xf3\\x06\\x8e\\x18\\'\\xcc\\xa9L\\x1d\\xbf\\x1e\\n" SD$  $7\x99\x1db(s\x90\x04\xb1\'0\xa90,\xe5<\xfd\xd4\xd8we\x17\xee<x\xb1t\xa7\xae+\xa6$  $\x05\xe3\x16\no(\xcb\xa2_g\xf8\xdb\xe1\xdb\x81\xd9c\x87\xc9\#\x0b\xef^\xxda\xd0$  $\label{lem:decomposition} $D \times d2 \times 14 \times bU = 1 \times bU$  $a6\x95\x03\xb5\xa4q\xedCl2\xb33\x91\xd8nB\x81k\x18P\xd6aX%\xf3L\x0f\x00\xdb\x8a}$  $3\xb4\xa4i\xf8\xcf\xc5\xd3\x8da\xe5\x88}\x1b1k\x82\xe1\x1c\xbc\x03\xf2u\xf6\rz\x$  $f4?\xd9\x1d\xbb^\xb9\t\x00\x00'$ 

 $\label{thm:limpo} Texto limpo = b'\x1f\x8b\x08\x00\xf3\xc8mb\x02\xff5V{P\xd4e\x14e\x7f\x0b\xbb\xee\x8be\x17\x88\x08\x08\x8b\x14y#\x08\x08\x94\x81\x85\xa2c;\xe8\x08\xab\x18\x0e\xa1\xc0\x08!\x0fMM\xd8\x14\xe3\xfd\x88\xe1\xd5\x103P\x10J<\x84\x8cYSy\x19P)e\x19 o)^\x8a\x1bDE1" \xcd\xfc\xee\xd9\xbf8s\xbf\xefw\xee\xb9\xe7\xde\xef.\x9cV{C''x8} 78\x07\'\x1da\xe6\xac\x07!Qnq\x1d\x8b\xf4\xb89\x1a\n\t\x9c\x8b\'Xd\xa1\x8c@\xc4\xcbz\x95.5\xa6\xe8\xb1 c\xa5\x84\x8e\x0cC[[\t1F\xeb\xad\x08\xc90\x84\xac\xb3\x88]$ 

 $\x0^{xc5}\xc40\xc8\xa0v\xd9\x98B\x1a\xdb!\xa2\x12\xfd\x1e\x12\x0el\n $\xb3U\x8f\$ b5\xe4\xb6\x13\xc7\xc6\xe6x\x1c\x06\xa4\xfcJH?\xf6q6\xd0\xe1\xe2\xdd\x84x\*\x8f;\  $xa8E=\\xae@\\xd6\\x1b\\xa1\\xb9,\\x1a\\xd9r\\x00\\xb4a~T\\x00\\xb7mG\#\\x81F\\xd7\\x0f\\Q\\xc0\\xcd$  $\x8c\xdfQ\xc0\xa9\xab\xa1\xf8\xb2\xf9)\xbc\xce\xb9P\x8d\x98M\xe9I\x16\xe9\xca\x$  $d\x85D\xcd\x82\rU\xe7\xafxF\\x9a\xe0\xdf \xf1\xd0\xa0?!\#\xff\x81\xd7Y\xf4j\xf7\$  $x10\xca\xfd31\x91E\x85\xe1\xa7\xc0\xb91\x0f\xd7\rT\xa7}\xc0\w|\x8a,\xefJ>\x81\xeq$  $f5 \times 1[(cDU \times 87 \times fb# \times c1S \times dx1c \times d8BH \times dc \times 2x9cFE \times d6 \times f4i \times 89 \times e6*B3)$  $6\xcFRm\x9c\x9d+d\x05\x7f\xe5"Eb\x8b1\xc2\xbcz\xe5qx\x12\xdb\xac\x0f=\xed\xdeNh$  $y\xe2W\xdf\x00\xcd\xb9\xa2bY\x1f\xe7\x1f\xd4\x9et\x9b\x9a\xcf\xe99\xb1\x8ck\;\x1$  $0\x89^xd5\xd6,\xe8\x91u\xeb\xe0^xfb\xca+\xc8\xb0Pb\xo1\xef\xa5\xcf3\xd1^xfb\x$ b2?PKPx\x07\x9c\xbc\$\x0c\x01\xea\xb7\x13\xc3\x86/d4&\\M\xaf%\xa5\_,\x98\x84\xa7V\ t\xdf=\x15\x88\x19\xd6xv\x81\xae\x04s#RL\xccRQ\xf6\x83\x94\x88\xe9\x9c\xd9H\  $x1a^x04\x00\x0b\x3\x1f\x98\x90\xb5\xde^x90\xc0\xcb\x9a4b\xd1\xb00\x1f\x8a\xa8$  $\xcoB9\xb2\xdaj\xoc\x0b\xbf\&\x0b]\x15M;\x07\xe2\xf4\xdd\xbc\x04i^\x04\xc5\x98\x$  $e4\xbe^xVx18F\xac\xba\xdf\xdd;\x82\xebn\xae\xb42\xb8\x8b\x8d\xf4\xde\xf4\x06\x$  $02\x1e\x00gy\xbe\x85\x64\xc9E\x8f<\xca\xb5\#\x18\xf8-\x99\xca\xa4\x8b\xd2X0>$  $\x8f\#\xfeC0\xe8\x18\xae\x97y\xc2\xa5.+\tN\xeb\xf7\x1a\x83\xb6t\xdbs\xa0\xd9\xb0f$  $\xf8ZU\xb8\t\xcan\xa8,\xc1\xa2R\xd2\x0c2r\x81)\x89\x8d*\xbd\x8e\xfb\xee\xeaVp\xf$  $8\xeaF\xc3\xaa\xfcG\xc5tmN\x8d\x1d\&\x9cJ9\x88A\xba\xb2\xb40;H}x0\x92\xfd\xf2\xfb$  $\xd4hsK1K\x99Mz\x03\xe5<g5H\x80\x11\xd3\xb6T4\xfakM\x7f\x84\xdc\xb2\xecgd\x95n^{}$  $x0fC@\\xe2\\x85W+M\\x8f\\xc1\\x0b\\x158\\x1e\\xfc\\x19_v$\\xf3!\\xe2N\\x05$\\xf2\\x9a\\xda\\xad1$  $\xc30\xaa+\xe0Z\x83\xf0,\xd1\x15\end{2}e?\r\\xb7I\x14"\xc9\xb4z\x81B\xca!\x03\x84\xa4$  $k/xa3; xe1x0f\xb1oEi\xd3<x9c::xc3xxc9x8c\xc5x00xa6xdf\x0f\xfd\x91g''\$  $3\x1d\x0e\xcd\xe2\x9b\xe8\xf1h\x94/\xdeh\x08\x8fZ6\x93\xe7@\x1f\xfeg\xac\xdd\xd9$  $KJz\\x1aLd\\xcf\\x1ar\\x87\\xee\\xa3\\x15\\xcc\\x0bl\\xf1\\x03\\x94\\x1d\\xe5\\xe6\\xdcso\\x81\\x$  $fdE\xee1\xa1m\xb4\xfe\rb\xb5;\xac\xe54\xd4\xa7\x06U\x91\x9f\x05\xa9\xef@\xfc$  $\xe6\x89y\xa8\x10\xcf\x98 \xd1B;\x96\xb6$L\x9d\x8f\x98C2\x8d\xaf\xaey:\x89\xe5V\$  $x8a\xb0\xd5\xc4\x81c\xf0_:Tk\x066\xa7V\x11\xdd\x0f\xfdA\x859\x08\xea\x85w\xf2\xe$  $d\xf1xv\xbc\x9b\xd8\xcbR\xbb\xf8HH2U\x0c\xe3U\x98\xa9\x8b!\xa6-\x06,0\xc4\x$  $f3Z\x6f\x62\xf1\xef\x85\xe0\xf6\x91=0\x91\xe6\xb49\xe7S\x0ei7\xdd\xc9N\x98\x15\x$  $ec2\x07\xc1a\x1d\x18"y\xc7<\xb2\xcaJ\x8e6\xb0|\xcf"\x17\xffX\x12\xde\x81\xec\xe$  $e06\xb2\xa1\xe7\xfb\xb0\\\xe\xa8\xe8?\x98\xdaA\xda\x8e\xdc\x927h\xa9\xf3L\xd6\\\\xe\xparxe0$  $x88L] \times 9f \times 8c \times 1.05 \times 1.0$  $bb\xdbz_{\x0}\x1a\x12X0\xe6\x02\xdb\xf8\x8f\xad\xf2\xc2\xce\x91\x86\x94\xf9d\x$ x9b\t\x00\x00'

### 1.2 Kyber CPAPKE

Começou-se com a implementação do algoritmo que permite uma segurança do tipo IND-CPA, isto é, contra Chosen Plaintext Attacks.

Sendo assim, começou-se com a implementação de funções auxiliares:

- parse: Através de um conjunto de bytes, devolve um polinómio.
- XOF: Corresponde a uma função do tipo extendable output function.
- PRF: Corresponde a uma função do tipo pseudorandom function.
- G: Realiza o hash através do SHA3-512;
- encode: Devolve um array de bytes através de um polinómio.
- decode: Possui um comportamento contrário ao encode.
- compress: Realiza a compressão de um polinómio.
- decompress: Função oposta ao compress.
- *transposta*: Calcula a transposta de uma matriz.

De seguida, realizou-se a construção das funções principais:

- gerar\_chaves: Realiza a geração de chaves privadas e públicas.
- cifragem: Realiza a cifragem da mensagem através da chave pública.
- decifragem: Decifra um criptograma, devolvendo um texto limpo.

No entanto, houveram dificuldades na construção da função ntt e  $ntt\_inversa$  o que impossibilitou o bom funcionamento destas funções.

```
import os
import random as rn
from sympy import ntt
from cryptography.hazmat.primitives import hashes
import numpy
from sympy import intt
import gzip
import struct
```

```
[2]: n = 256
    q = 343576577
    T = NTT(n,q)
    k = 2
    n1 = 3
    n2 = 2
    du, dv = 10, 4

    _Z.<w> = ZZ[]
    R.<w> = QuotientRing(_Z ,_Z.ideal(w^n - 1))

    _Q.<w> = GF(q)[]
    Rq.<w> = QuotientRing(_Q , _Q.ideal(w^n + 1))
```

```
def ntt():
    return
def ntt_inv():
   return
def tamanho(stringB, numberS):
    contador = 10
    auxContador = 1
    i = 0
    while i < len(stringB):</pre>
        if numberS == auxContador:
            i = i + 10
            while (i < len(stringB)) and (stringB[i] != 31 or stringB[i + 1] !=u
 \hookrightarrow139 or stringB[i + 2] != 8 or stringB[i + 3] != 0 ):
                countador = countador + 1
                i = i + 1
            auxContador = auxContador + 1
        i = i + 1
        if (i + 10) < len(stringB) and (stringB[i] == 31 and stringB[i + 1] ==__
 \hookrightarrow139 and stringB[i + 2] == 8 and stringB[i + 3] == 0):
            auxContador = auxContador + 1
        if auxContador > numberS:
            break
    return countador
def parse(str_bytes):
    result = []
    for i in str_bytes:
        result.append(i)
    return Rq(result)
def XOF(p,i,j):
    digest = hashes.Hash(hashes.SHAKE128(int(32)))
    digest.update(p)
    digest.update(bytes(i))
    digest.update(bytes(j))
    r = digest.finalize()
    return r
def PRF(s,b):
    digest = hashes.Hash(hashes.SHAKE256(int(32)))
    digest.update(s)
    digest.update(bytes(b))
    r = digest.finalize()
    return r
```

```
def G(d):
    digest = hashes.Hash(hashes.SHA512())
    digest.update(bytes(d))
    r = digest.finalize()
    return r
def encode(poly):
    byte=b''
    aux=1
    countadorX=0
    for j in poly:
        if(j>255):
            aux=2
        if (j > 65025):
            aux = 3
        if (j > 16581375):
            aux = 4
        if (j > 4228250625):
            aux = 5
        byte = byte+ int((_Z(j))).to_bytes( aux, 'big')
        byte = byte +"/-n-/".encode()
        countadorX =countadorX +1
    return byt
def decode(byt):
    listaCoeficiente = []
    byteAux = b''
    desc=0
    while desc <byt.__len__():</pre>
        if byt[desc] == 47 and byt[desc+1] == 45 and byt[desc+2] == 110 and
 \hookrightarrowbyt[desc+3]==45 and byt[desc+4] == 47 :
            desc = desc+4
            listaCoeficiente.append(int.from_bytes(byteAux, 'big'))
            byteAux = b''
        else:
            byteAux = byteAux + bytearray([int(_Z(byt[desc]))])
        desc = desc+1
    return listaCoeficiente
def compress(polinomio):
    polinomioB= encode(polinomio)
    compress = gzip.compress(polinomioB)
    return compress
def decompress(compress):
    unpack = gzip.decompress(compress)
```

```
return Rq(decode(unpack))
def transposta(matrix):
    zipped_rows = zip(*matrix)
    transpose_matrix = [list(row) for row in zipped_rows]
    return transpose_matrix
def gerar_chaves():
    d = bytearray(os.urandom(32))
    p = G(d)[:32]
    teta = G(d)[-32:]
    A = [[0 \text{ for } x \text{ in } range(k-1)] \text{ for } y \text{ in } range(k-1)]
    for i in range(0,k-1):
        for j in range(0,k-1):
            A[i][j] =parse(XOF(p,i,j))
    s = []
    for i in range(0,k-1):
        s.append(parse(PRF(teta,N)))
        N = N + 1
    e = []
    for i in range(0,k-1):
        e.append(parse(PRF(teta,N)))
        N = N + 1
    s1 = Rq(T.ntt(s[0]))
    e1 = Rq(T.ntt(e[0]))
    t = A[0][0].lift() * s1.lift() + e1.lift()
    pk = encode(t) + p
    sk = encode(s1)
    return pk, sk
def cifragem(pk, m, coins):
    N = 0
    t2 = pk[:len(pk)-32]
    t= decode(t2)
    p = pk[-32:]
    A = [[0 \text{ for } x \text{ in } range(k-1)] \text{ for } y \text{ in } range(k-1)]
    for i in range(0,k-1):
        for j in range(0,k-1):
             A[i][j] =parse(XOF(p,i,j))
    AT = transposta(A)
    r = []
    for i in range(0,k-1):
        r.append(parse((PRF(bytearray(r),bytearray([N])))))
        N = N + 1
    e1 = []
    for i in range(0,k-1):
```

```
e1.append(parse(PRF(bytearray(r[i]),bytearray([N]))))
        N = N + 1
    e2 = parse(PRF(bytearray(r[0].list()),N))
    r1= Rq(T.ntt(r[0]))
    u= Rq(T.ntt_inv(A[0][0].lift()*r1.lift()))+e1[0]
    v = Rq(T.ntt(Rq(t).lift() * r1.lift())) + e2 + decompress(m)
    c1 = compress(u)
    c2 = compress(v)
    c = c1 + c2
    return c
def decifragem(sk, ciphertext):
    u = decompress(ciphertext[:tamanho(ciphertext,1)])
    v = decompress(ciphertext[-tamanho(ciphertext,2):])
    s1 = Rq(decode(sk))
    m = compress(v.lift() - (T.ntt_inv(s1.lift()*Rq(T.ntt(u)).lift())))
    return m
```

```
NameError Traceback (most recent call last)

/tmp/ipykernel_4688/1247484197.py in <cell line: 3>()

1 n = Integer(256)
2 q = Integer(343576577)

----> 3 T = NTT(n,q)
4 k = Integer(2)
5 n1 = Integer(3)

NameError: name 'NTT' is not defined
```

#### 1.3 Resultados obtidos

```
texto_cifrado = cifragem(pk, compress(m),coins)
print("Texto cifrado = ", texto_cifrado, "\n")
texto = decifragem(sk, texto_cifrado)
print("Texto limpo = ", texto)
```

```
NameError Traceback (most recent call last)

/tmp/ipykernel_4688/3069925417.py in <cell line: 1>()

----> 1 pk, sk = gerar_chaves()

2

3 m = Rq([Integer(1), Integer(0), Integer(1), Integer(0), Integer(1), Integer(1), Integer(1), Integer(0), Integer(0), Integer(1), Integer(1), Integer(1), Integer(0), Integer(1), Inte
```