

[15] 1. State whether each of the following is **true** or **false**. **Justify your answer!**

- a. Any finite language over the alphabet $\{a, b, c, \dots, z\}$ is context free.

True.

Since it is the finite language, it must be regular.

So it is context-free.

- b. Every context free language can be recognized by a deterministic push-down automaton (PDA).

False.

$$\{a^i b^j c^k \mid i, j, k \geq 0 \text{ and } i=j \text{ or } j=k\} \text{ is}$$

non-deterministic

- c. Union of any context free language and a regular language is context free.

True.

Since a regular language must be context-free.

So the union of two context-free language is still context-free.

- d. Pumping lemma is used to prove that a language is regular or context free.

False.

It is used to prove the language is not

- c. Union of any context free language and a regular language is context free.

True. Since a regular language must be context-free.
So the union of two context-free language is still context-free.

- d. Pumping lemma is used to prove that a language is regular or context free.

False. It is used to prove the language is not regular or context-free.

- e. The set of context free languages is closed under operation of intersection.

False.

$$L_1 = \{a^i b^j c^k \mid i < j\}$$

$$L_2 = \{a^i b^j c^k \mid i < k\}$$

but $L_1 \cap L_2$ is not context free.

[16] 2. Consider the following context-free grammar G

$$S \rightarrow 0S1 \mid 1S0 \mid SS \mid M$$

$$M \rightarrow R1 \mid 1R \mid R$$

$$R \rightarrow 1R \mid \epsilon$$

Provide brief answers to the following questions:

[2] a. List all variables and all terminals of G.

Variables: $\{S, M, R\}$ 2

Terminals: $\{1, 0\}$

[2] b. List two strings in $L(G)$.

01, 10 2

[2] c. List two strings not in $L(G)$.

00, 0 2

[2] d. List two sentential forms of G.

0S1, 1S0 2

[2] d. List two sentential forms of G.

OSI, ISO

2

Open with

[4] d. Give two **different** leftmost derivations of the string 1110001 in G.

$S \rightarrow SS \rightarrow SOS \rightarrow 1S00S \rightarrow 11S000S \rightarrow 111M000S \rightarrow 111R000S \rightarrow 111000S$
 $\rightarrow 111000M \rightarrow 111000R \rightarrow 111000E1 \rightarrow 1110001$

$S \rightarrow SS \rightarrow SOS \rightarrow 1S00S \rightarrow 11M00S \rightarrow 111R00S \rightarrow 111E00S \rightarrow 11100S$
 $\rightarrow 111000S1 \rightarrow 111000M1 \rightarrow 111000R1 \rightarrow 111000E1 \rightarrow 1110001$

2

[4] e. Give a description of $L(G)$ in English.

$L(G)$ consists of 0 and 1, and the number of 1's is not less than the number of 0's

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[10] 3. Convert grammar G from the previous question into Chomsky normal form (the grammar is repeated here for convenience):

$$S \rightarrow 0S1 \mid 1S0 \mid SS \mid M$$

$$M \rightarrow R1 \mid 1R \mid R$$

$$R \rightarrow 1R \mid \epsilon$$

Eliminate $S \rightarrow S$

$$S_0 \rightarrow S \mid \epsilon$$

$$S \rightarrow 0S1 \mid 1S0 \mid SS \mid 01 \mid 10 \mid 1R \mid R1 \mid 1$$

$$M \rightarrow 1R \mid R1 \mid 1$$

$$R \rightarrow 1R \mid 1$$

Eliminate $S_0 \rightarrow S$

$$S_0 \rightarrow 0S1 \mid 1S0 \mid SS \mid 01 \mid 10 \mid 1R \mid R1 \mid 1 \mid \epsilon$$

$$S \rightarrow 0S1 \mid 1S0 \mid SS \mid 01 \mid 10 \mid 1R \mid R1 \mid 1$$

$$M \rightarrow 1R \mid R1 \mid 1$$

$$R \rightarrow 1R \mid 1$$

④ Add Rules

$$U \rightarrow 1$$

$$X \rightarrow 0S$$

$$Y \rightarrow S0$$

$$Z \rightarrow 0$$

$$S_0 \rightarrow XU \mid UY \mid SS \mid 01 \mid 10 \mid UR \mid RU \mid 1 \mid \epsilon$$

$$S \rightarrow XU \mid UY \mid SS \mid 01 \mid 10 \mid UR \mid RU \mid 1$$

$$M \rightarrow UR \mid RU \mid 1$$

$$R \rightarrow UR \mid 1$$

$$U \rightarrow 1$$

$$X \rightarrow 0S$$

$$Y \rightarrow S0$$

$$Z \rightarrow 0$$

①

$$S_0 \rightarrow S$$

$$S \rightarrow 0S1 \mid 1S0 \mid SS \mid M$$

$$M \rightarrow R1 \mid 1R \mid R$$

$$R \rightarrow 1R \mid \epsilon$$

② Eliminate $R \rightarrow \epsilon$

$$S_0 \rightarrow S$$

$$S \rightarrow 0S1 \mid 1S0 \mid SS \mid M$$

$$M \rightarrow R1 \mid 1R \mid R \mid 1 \mid \epsilon$$

$$R \rightarrow 1R \mid 1$$

Eliminate $M \rightarrow \epsilon$

$$S_0 \rightarrow S$$

$$S \rightarrow 0S1 \mid 1S0 \mid SS \mid M \mid \epsilon$$

$$M \rightarrow R1 \mid 1R \mid R \mid 1$$

$$R \rightarrow 1R \mid 1$$

Eliminate $S \rightarrow \epsilon$

$$S_0 \rightarrow S \mid \epsilon$$

$$S \rightarrow 0S1 \mid 1S0 \mid SS \mid 01 \mid 10 \mid S \mid M$$

$$M \rightarrow 1R \mid R1 \mid R \mid 1$$

$$R \rightarrow 1R \mid 1$$

③ Eliminate $M \rightarrow R$

$$S_0 \rightarrow S \mid \epsilon$$

$$S \rightarrow 0S1 \mid 1S0 \mid SS \mid 01 \mid 10 \mid S \mid M$$

$$M \rightarrow 1R \mid R1 \mid 1$$

$$R \rightarrow 1R \mid 1$$

Eliminate $S \rightarrow M$

$$S_0 \rightarrow S \mid \epsilon$$

$$S \rightarrow 0S1 \mid 1S0 \mid SS \mid 01 \mid 10 \mid S \mid 1R \mid R1 \mid 1$$

$$M \rightarrow 1R \mid R1 \mid 1$$

$$R \rightarrow 1R \mid 1$$

[6] 4. Convert grammar G from the previous question to an equivalent PDA

$$S \rightarrow 0S1 \mid 1S0 \mid SS \mid M$$

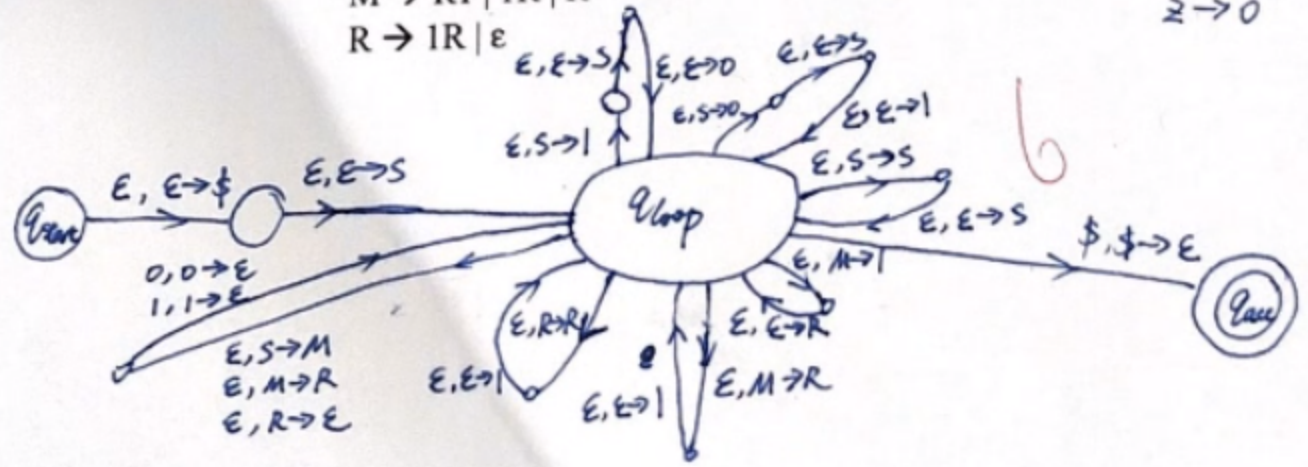
$$M \rightarrow R1 \mid 1R \mid R$$

$S_0 \rightarrow S$
 $S \rightarrow OS1 \mid IS0 \mid SS \mid M/E$
 $M \rightarrow R1 \mid IR \mid R \mid$
 $R \rightarrow IR \mid$

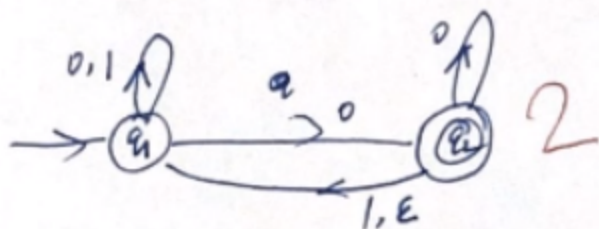
Eliminate $S \rightarrow M$
 $S_0 \rightarrow S \mid \epsilon$
 $S \rightarrow OS1 \mid IS0 \mid SS \mid 01 \mid 10 \mid S \mid IR \mid N \mid$
 $M \rightarrow IR \mid R \mid$
 $R \rightarrow IR \mid$

$X \rightarrow OS$
 $Y \rightarrow SO$
 $Z \rightarrow O$
 $S_0 \rightarrow XV \mid UY \mid SS \mid 01 \mid 10 \mid UR \mid RV \mid \mid \epsilon$
 $S \rightarrow XV \mid UY \mid SS \mid 01 \mid 10 \mid UR \mid RV \mid$
 $M \rightarrow UR \mid RV \mid$
 $R \rightarrow UR \mid$
 $U \rightarrow$
 $X \rightarrow \bar{S}$
 $Y \rightarrow S$
 $Z \rightarrow O$

[6] 4. Convert grammar G from the previous question to an equivalent PDA
 $S \rightarrow OS1 \mid IS0 \mid SS \mid M$
 $M \rightarrow R1 \mid IR \mid R$
 $R \rightarrow IR \mid \epsilon$



[8] 5. The formal description of a NFA is $((\{q_1, q_2\}, \{0,1\}, \delta, q_1, \{q_2\}))$, where δ is given by the following table. Give the context free grammar that defines the same language as the one accepted by this NFA.



| | 0 | 1 | ϵ |
|-------|----------------|-----------|-------------|
| q_1 | $\{q_1, q_2\}$ | $\{q_1\}$ | \emptyset |
| q_2 | $\{q_2\}$ | $\{q_1\}$ | $\{q_1\}$ |

$$\begin{aligned}
 S &\rightarrow q_1 \\
 q_1 &\rightarrow 0q_1 \mid 0q_2 \mid 1q_1 \\
 q_2 &\rightarrow 1q_1 \mid 0q_2 \mid \epsilon q_1
 \end{aligned}$$