

Deriving a Symbolic Executor for Definitional Interpreters Suitable for the Study of Heuristics

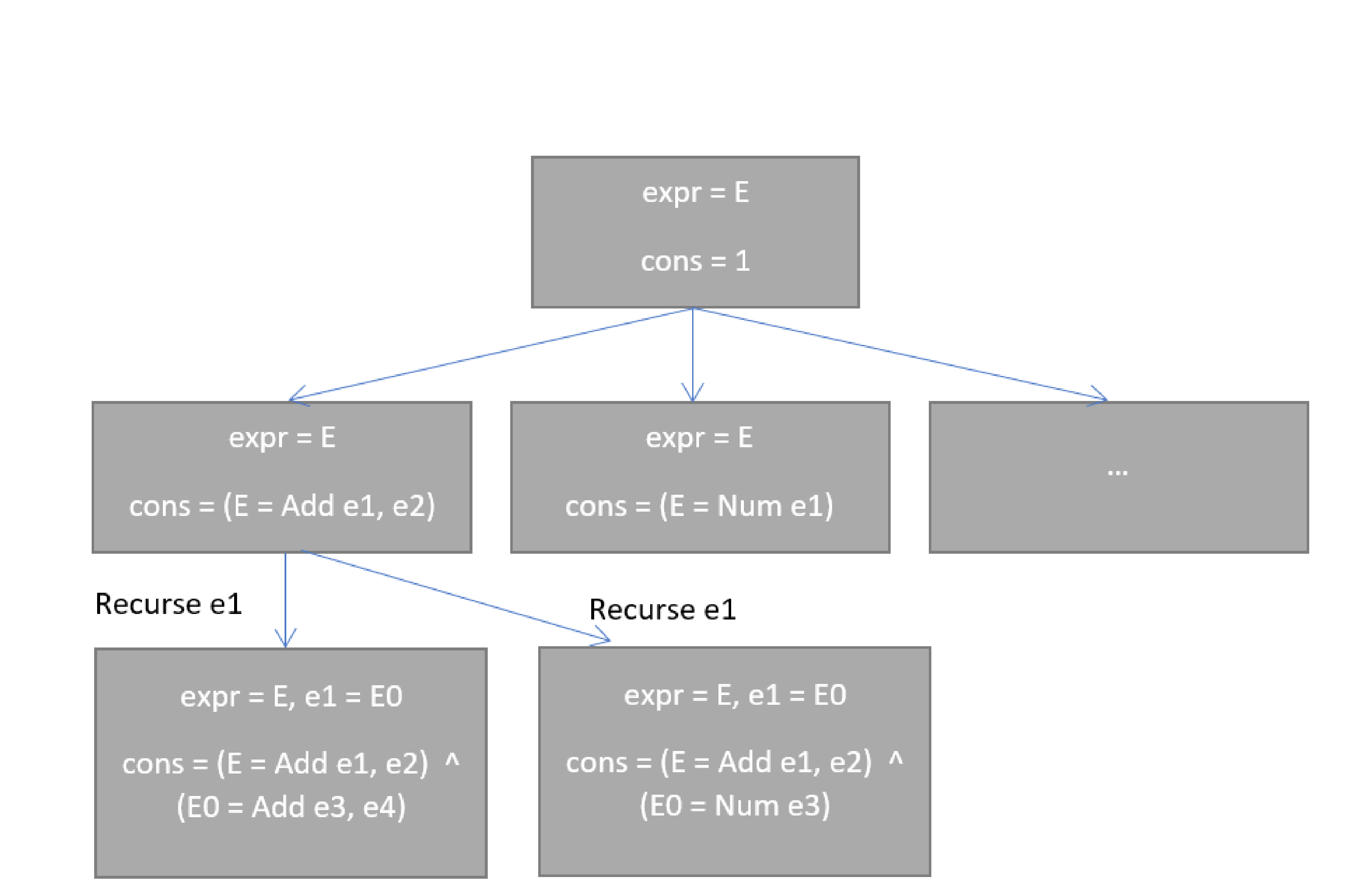
Laura Pircalaboiu, Casper Bath Poulsen & Cas van der Rest
Research Project CSE3000

Introduction

- Student submissions in courses are hard to manually evaluate
 - Unit testing insufficient
 - Mensing approach to symbolic execution- effective, but we want to extend it
 - Use as starting point: interpreters like defined in PLAI [2].
- The goal of this research project is to determine whether any two given definitional interpreters are equal or not using symbolic execution.

We want to create a simple and extensible approach to symbolic execution.

Intuition: Building Execution Trees



Method: Intermediate Representation

How do we encode interpreters to ensure consistency in our results and extensibility of the approach?



```
data Expr = Num Int |
  Add Expr Expr

eval :: Expr -> Int
eval (Num i) = i
eval (Add e1 e2) =
  eval e1 + eval e2

eval =
  ([e = Num(i)].
   return t)
  + ([e != Num(i),
      e = Add(e1, e2)]
     .recurse e1
    as i1
     .recurse e2 as i2
     .return +(i1, i2))
```

Results

		True	False
Actual	True	4	0
	False	10	20

We compared 8 interpreters that belong to 3 equivalence classes, for a total of 34 test cases. The results are reported in the confusion matrix above.

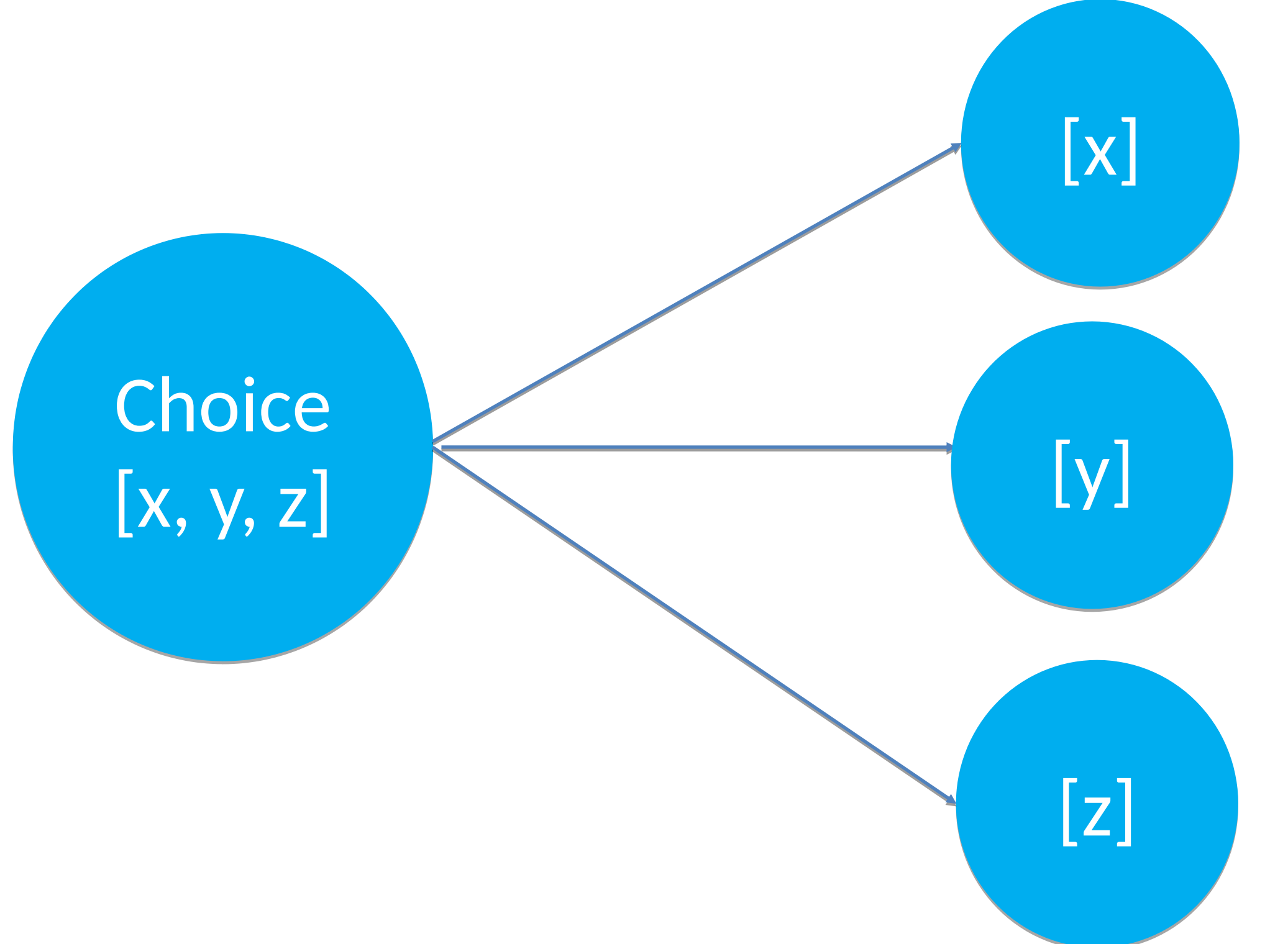
Discussion & Future Work

The approach works for finding trivial bugs, such as wrong order of variables or typos, but gives false negatives in the case of equivalent interpreters that have a different branch order.

- Possible future improvements are:
- Extensions to the programming languages
 - Usage of Heuristics and/or branch pruning
 - Ability to run two interpreters in (real) lock-step

Method: Small Step Transition Function

Idea adopted from Mensing et al. [1]:



References

1. A. D. Mensing, "From Definitional Interpreter to Symbolic Executor," en, p. 10, 2019.
- 2.S. Krishnamurthi, "Programming Languages: Application and Interpretation," en, p. 207.

