Introduction to Network Security

For this section, the slides are from selected sections of the following textbook: Data communications and networking, 4th edition, Forouzan, McGrawHill

- Chapters 30 & 31

Cryptography Basics

Let us introduce the issues involved in cryptography. First, we need to define some terms; then we give some taxonomies.

Topics discussed in this section:

Definitions

Two Categories:

Symmetric-Key vs. Asymmetric Key Systems

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Figure 30.1 Cryptography components

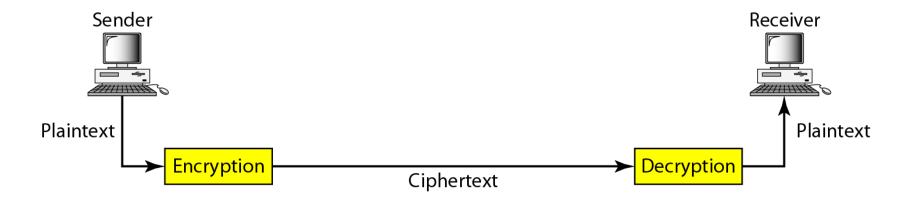


Figure 30.2 Categories of cryptography

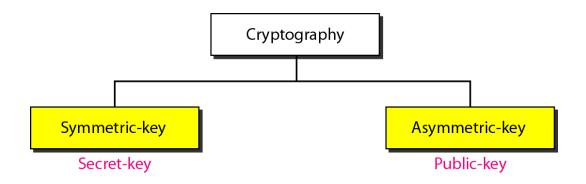
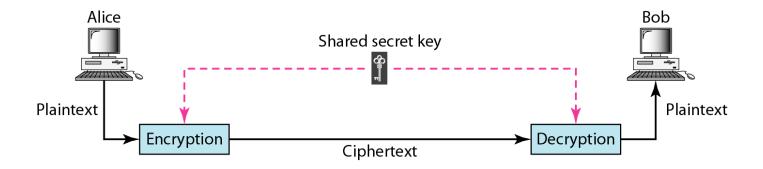


Figure 30.3 Symmetric-key cryptography





In symmetric-key cryptography, the same key is used by the sender (for encryption) and the receiver (for decryption).

The key is shared.

Figure 30.4 Asymmetric-key cryptography

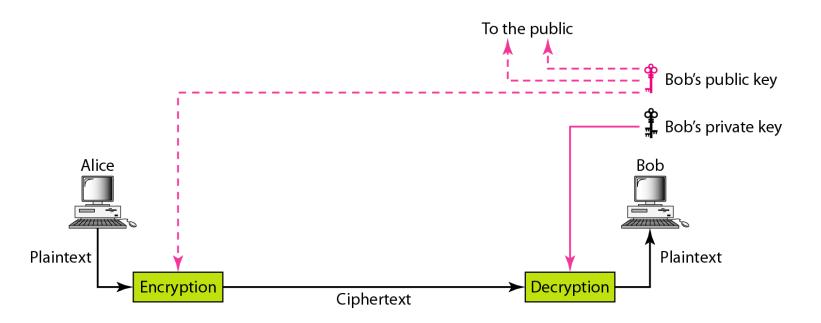
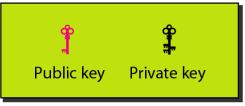


Figure 30.5 Keys used in cryptography

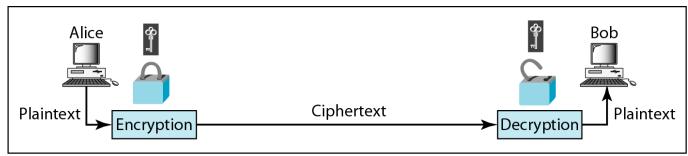


Symmetric-key cryptography

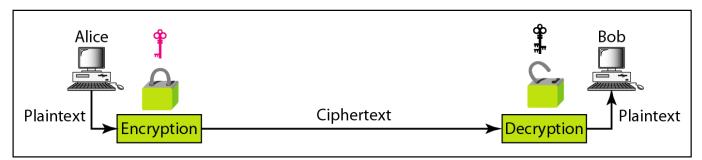


Asymmetric-key cryptography

Figure 30.6 Comparison between two categories of cryptography



a. Symmetric-key cryptography



b. Asymmetric-key cryptography

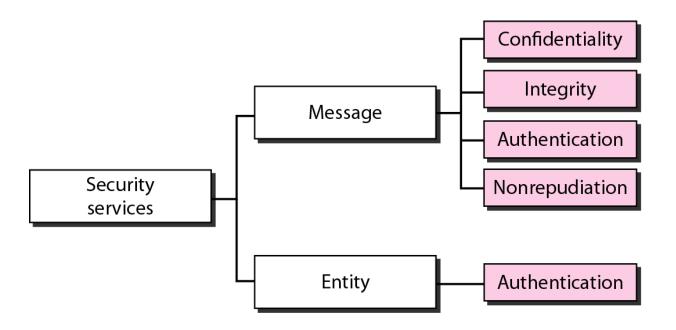
SECURITY SERVICES

Network security can provide five services. Four of these services are related to the message exchanged using the network. The fifth service provides entity authentication or identification.

Topics discussed in this section:

Message Confidentiality
Message Integrity
Message Authentication
Message Nonrepudiation
Entity Authentication

Figure 31.1 Security services related to the message or entity



MESSAGE CONFIDENTIALITY

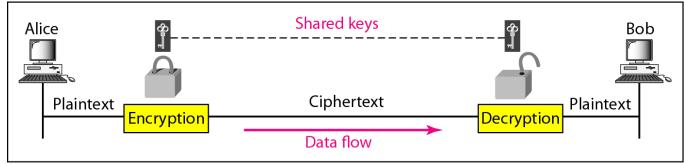
The concept of how to achieve message confidentiality or privacy has not changed for thousands of years. The message must be encrypted at the sender site and decrypted at the receiver site. This can be done using either symmetric-key cryptography or asymmetric-key cryptography.

<u>Topics discussed in this section:</u>

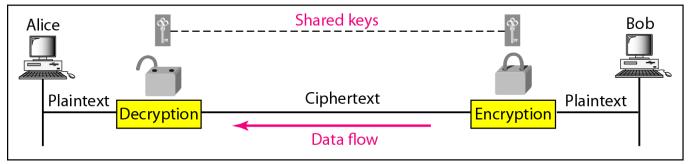
Confidentiality with Symmetric-Key Cryptography Confidentiality with Asymmetric-Key Cryptography

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Figure 31.2 Message confidentiality using symmetric keys in two directions

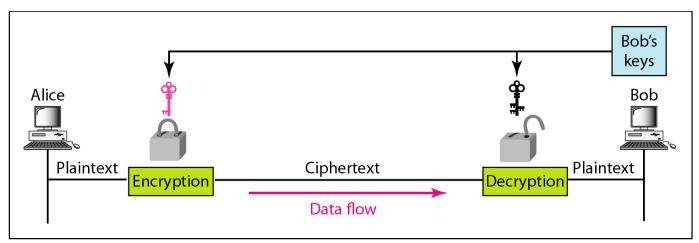


a. A shared secret key can be used in Alice-Bob communication

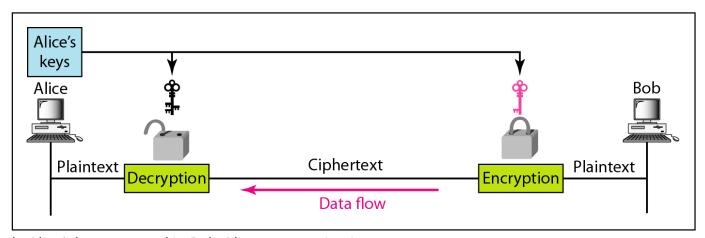


b. A different shared secret key is recommended in Bob-Alice communication

Figure 31.3 Message confidentiality using asymmetric keys



a. Bob's keys are used in Alice-Bob communication



b. Alice's keys are used in Bob-Alice communication

MESSAGE INTEGRITY

Encryption and decryption provide secrecy, or confidentiality, but not integrity. However, on occasion we may not even need secrecy, but instead must have integrity.

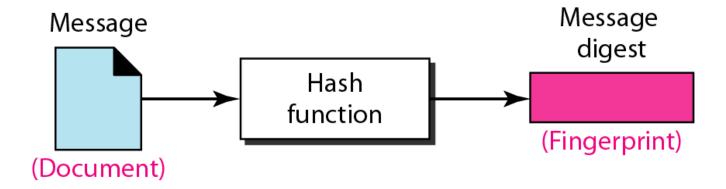
Topics discussed in this section:

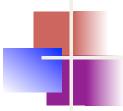
Document and Fingerprint
Message and Message Digest
Creating and Checking the Digest
Hash Function Criteria
Hash Algorithms: SHA-1



To preserve the integrity of a document, both the document and the fingerprint are needed.

Figure 31.4 Message and message digest





The message digest needs to be kept secret.

Figure 31.5 Checking integrity

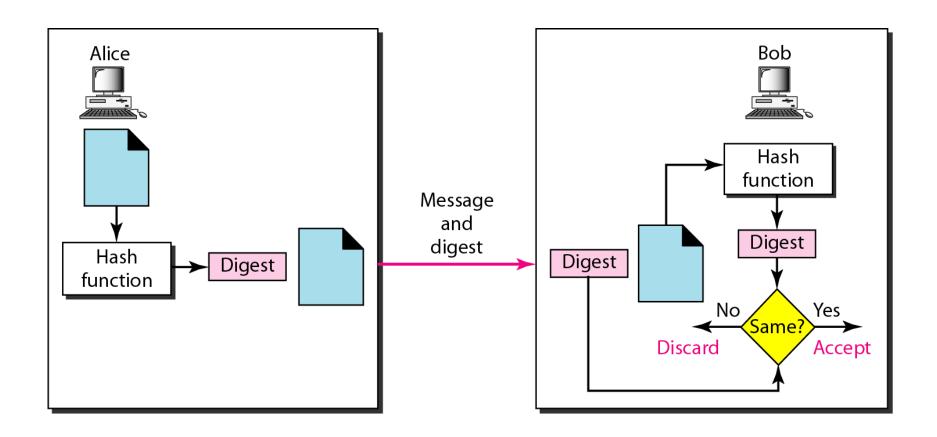
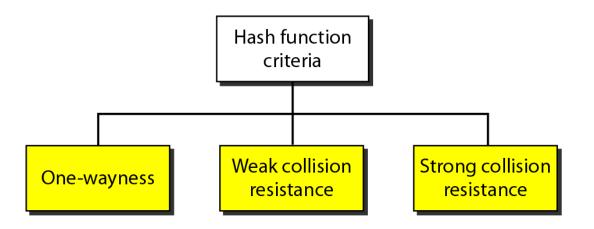


Figure 31.6 Criteria of a hash function



Example 31.1

Can we use a conventional lossless compression method as a hashing function?

Solution

We cannot. A lossless compression method creates a compressed message that is reversible. You can uncompress the compressed message to get the original one.

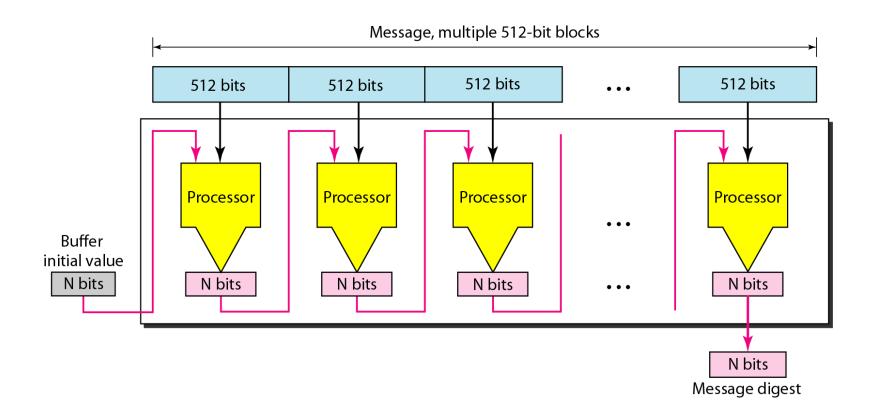
Example 31.2

Can we use a checksum method as a hashing function?

Solution

We can. A checksum function is not reversible; it meets the first criterion. However, it does not meet the other criteria.

Figure 31.7 Message digest creation

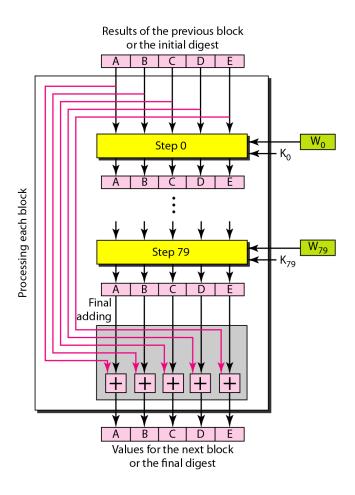




SHA-1 hash algorithms create an N-bit message digest out of a message of 512-bit blocks.

SHA-1 has a message digest of 160 bits (5 words of 32 bits).

Figure 31.8 Processing of one block in SHA-1



MESSAGE AUTHENTICATION

A hash function per se cannot provide authentication. The digest created by a hash function can detect any modification in the message, but not authentication.

<u>Topics discussed in this section:</u>

MAC

Figure 31.9 MAC, created by Alice and checked by Bob

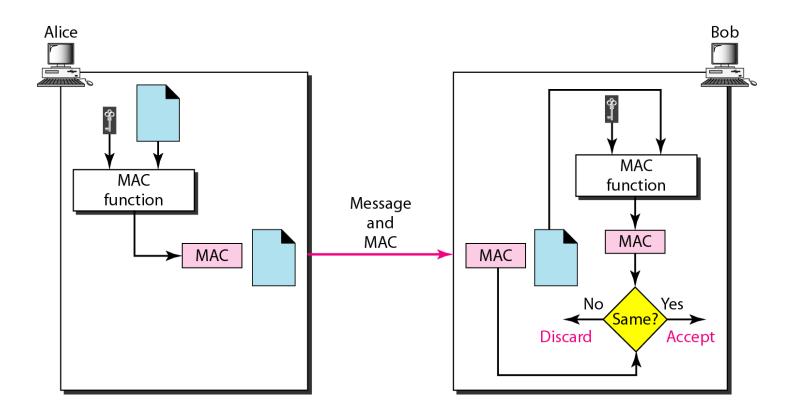
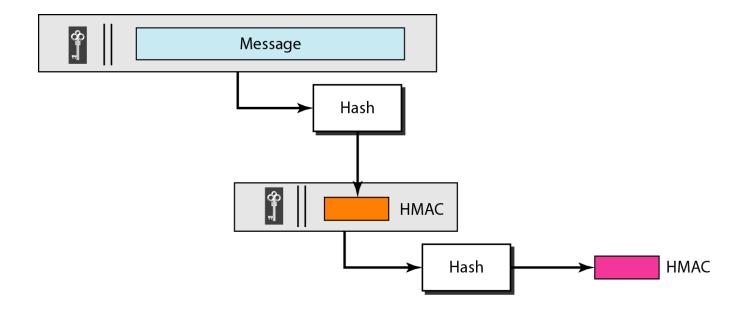


Figure 31.10 HMAC



DIGITAL SIGNATURE

When Alice sends a message to Bob, Bob needs to check the authenticity of the sender; he needs to be sure that the message comes from Alice and not Eve. Bob can ask Alice to sign the message electronically. In other words, an electronic signature can prove the authenticity of Alice as the sender of the message. We refer to this type of signature as a digital signature.

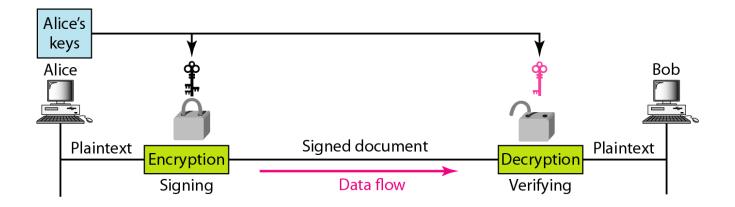
Topics discussed in this section:

Comparison
Need for Keys
Process



A digital signature needs a public-key system.

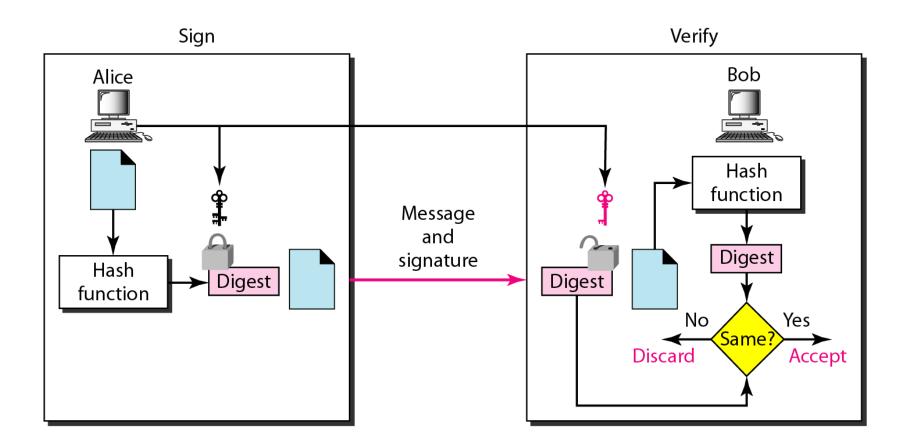
Figure 31.11 Signing the message itself in digital signature





In a cryptosystem, we use the private and public keys of the receiver; in digital signature, we use the private and public keys of the sender.

Figure 31.12 Signing the digest in a digital signature



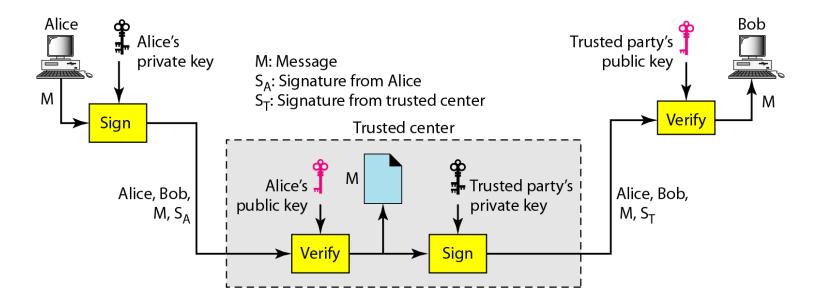


A digital signature today provides message integrity.



Digital signature provides message authentication.

Figure 31.13 Using a trusted center for nonrepudiation





Nonrepudiation can be provided using a trusted party.

ENTITY AUTHENTICATION

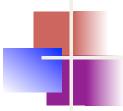
Entity authentication is a technique designed to let one party prove the identity of another party. An entity can be a person, a process, a client, or a server. The entity whose identity needs to be proved is called the claimant; the party that tries to prove the identity of the claimant is called the verifier.

Topics discussed in this section:

Passwords Challenge-Response



In challenge-response authentication, the claimant proves that she knows a secret without revealing it.



The challenge is a time-varying value sent by the verifier; the response is the result of a function applied on the challenge.

Figure 31.14 Challenge/response authentication using a nonce

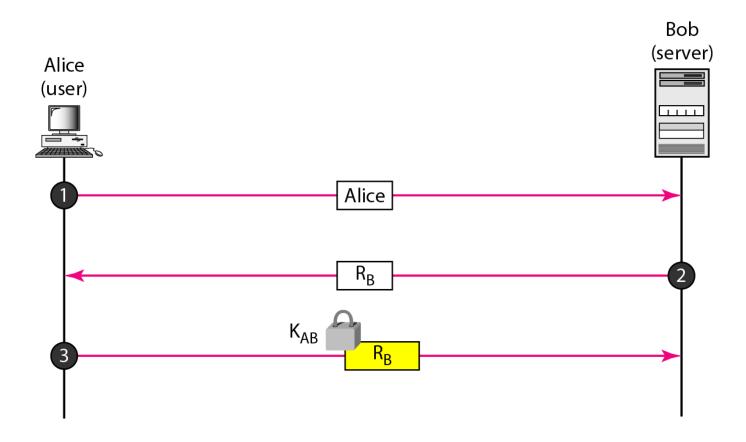


Figure 31.15 Challenge-response authentication using a timestamp

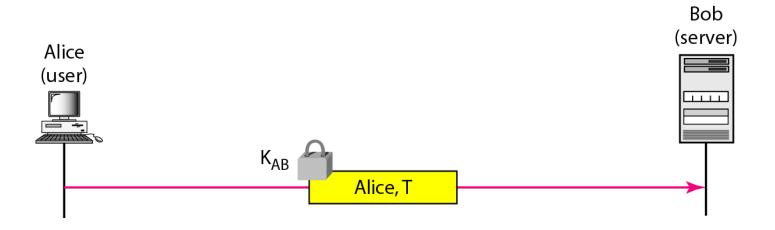


Figure 31.16 Challenge-response authentication using a keyed-hash function

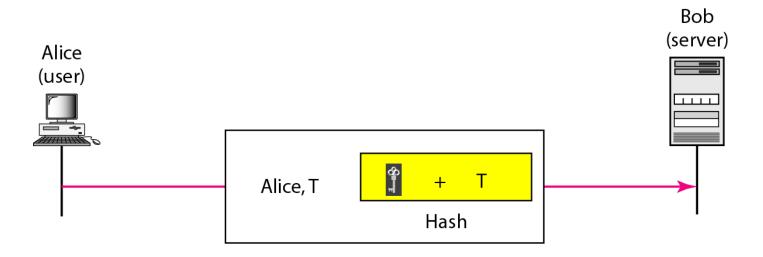


Figure 31.17 Authentication, asymmetric-key

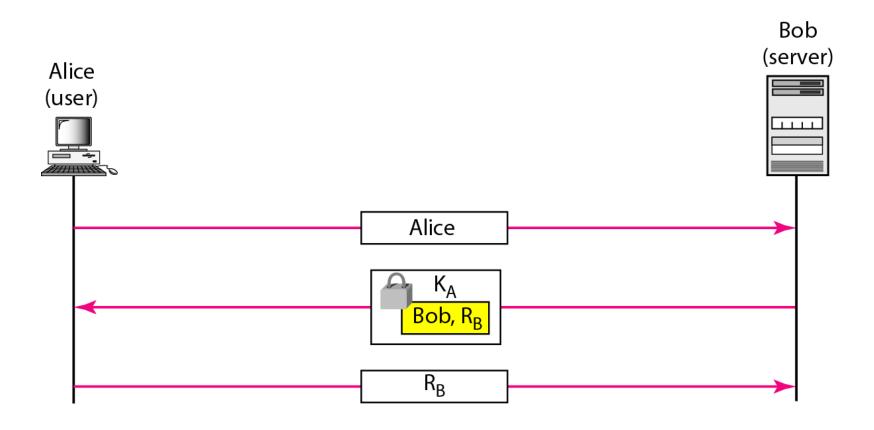
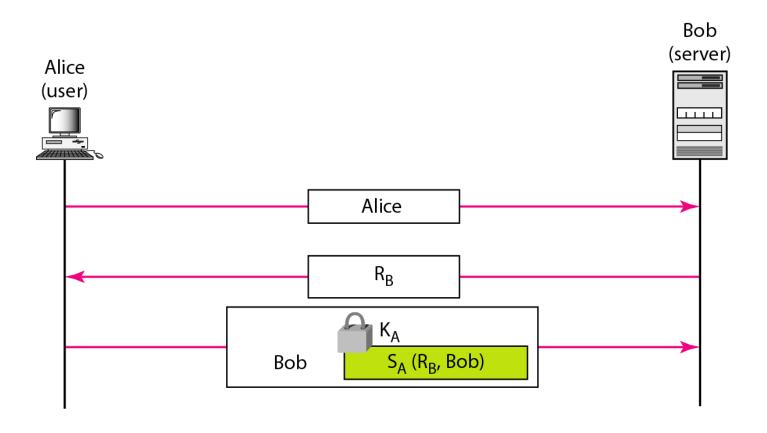


Figure 31.18 Authentication, using digital signature



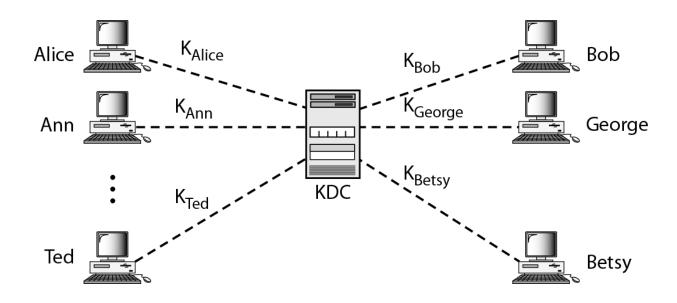
KEY MANAGEMENT

So far, we have not yet discussed how secret keys in symmetric-key cryptography and how public keys in asymmetric-key cryptography are distributed and maintained. In this section, we touch on these two issues. We first discuss the distribution of symmetric keys; we then discuss the distribution of asymmetric keys.

Topics discussed in this section:

Symmetric-Key Distribution Public-Key Distribution

Figure 31.19 KDC





A session symmetric key between two parties is used only once.

Figure 31.30 Creating a session key between Alice and Bob using KDC

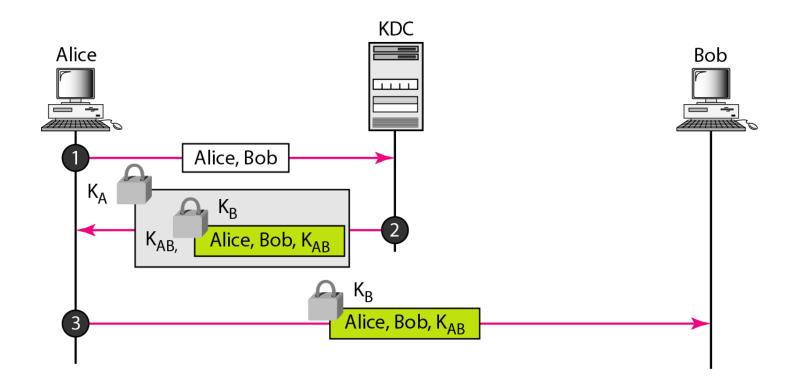


Figure 31.21 Kerberos servers

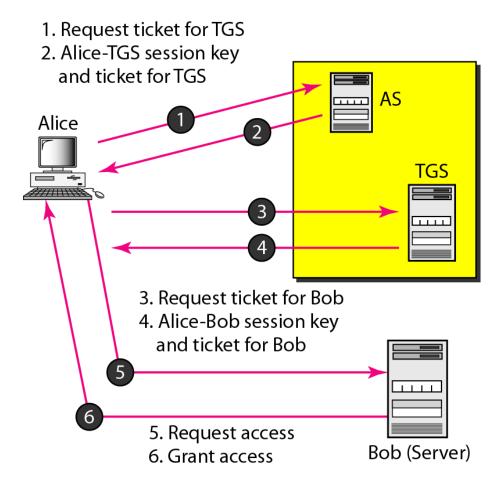
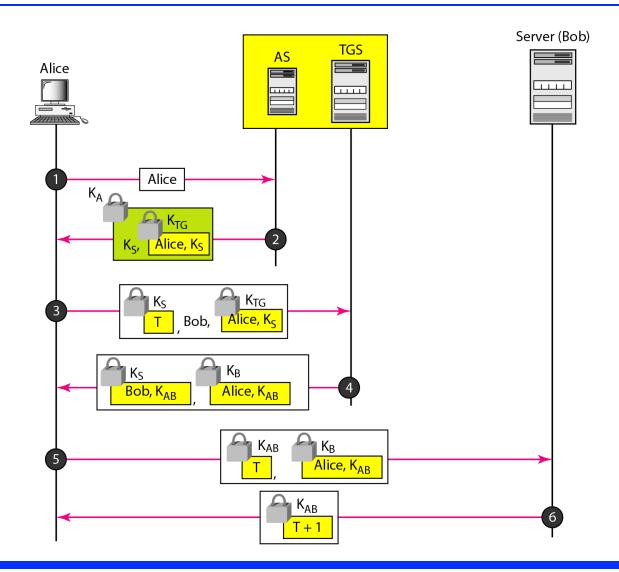


Figure 31.22 Kerberos example





In public-key cryptography, everyone has access to everyone's public key; public keys are available to the public.

Figure 31.23 Announcing a public key

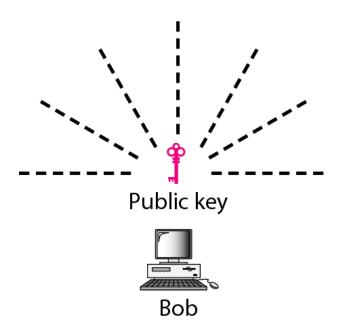


Figure 31.24 Trusted center

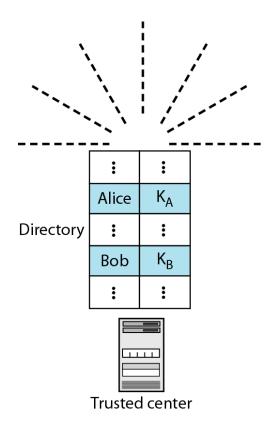


Figure 31.25 Controlled trusted center

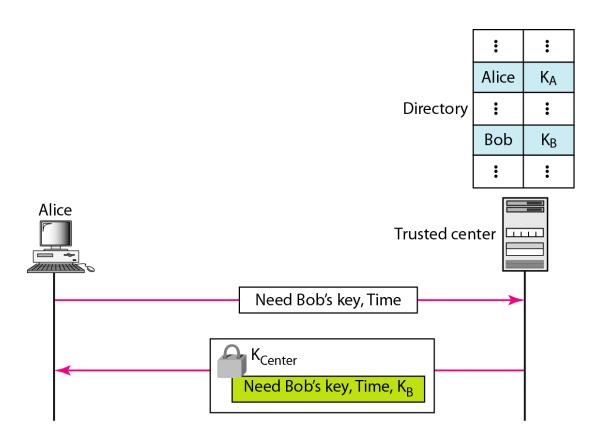


Figure 31.26 Certification authority

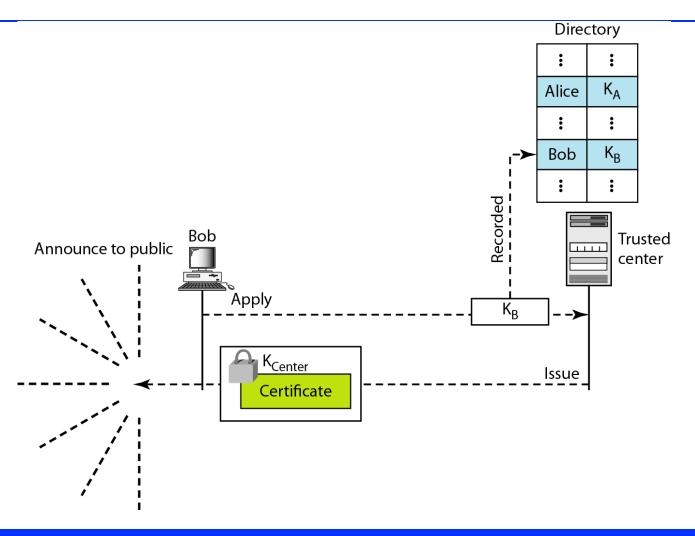


Figure 31.27 PKI hierarchy

