# A reference GLL implementation



### Acknowledgements



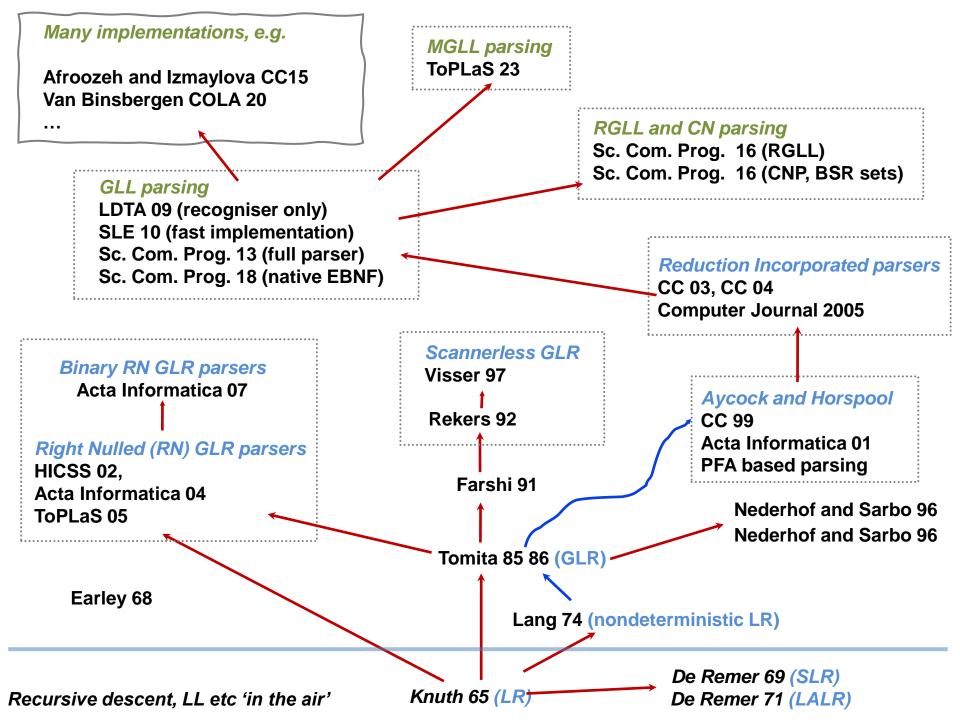
### The Leverhulme Trust

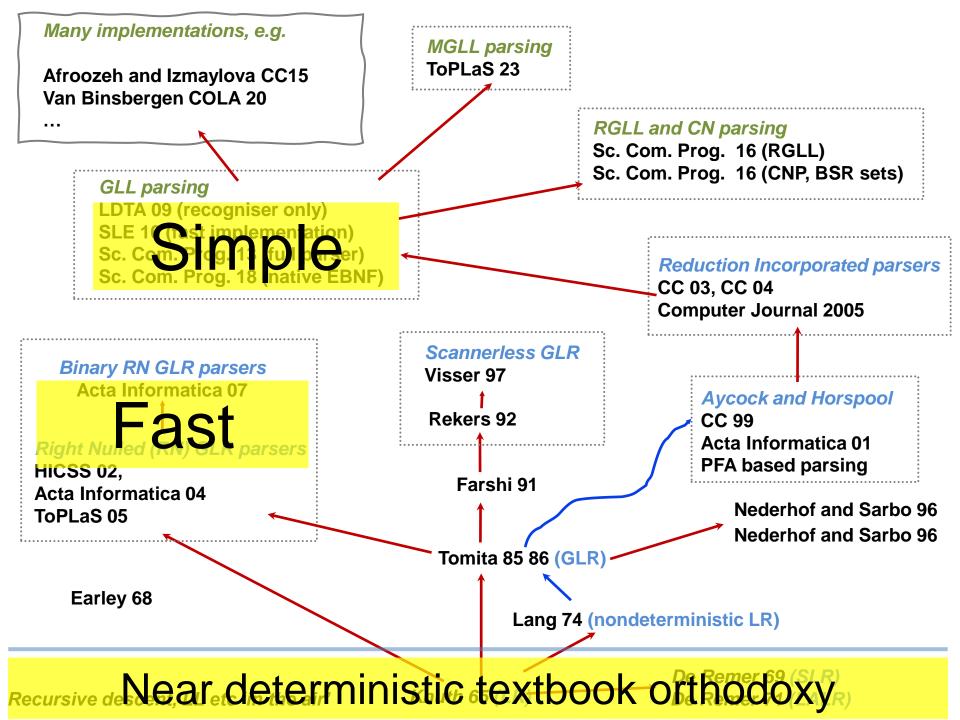


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LALR/LL parsing: fine for linear things Flips over if you drive it too hard



Singleton derivation parsers: OK on special surfaces Always follows one track, ignoring other routes

Attr: CC BY-SA 3.0 Deed by Wikimedia commons user Morio 2011 Japanese GP: Jenson Button (McLaren) during race



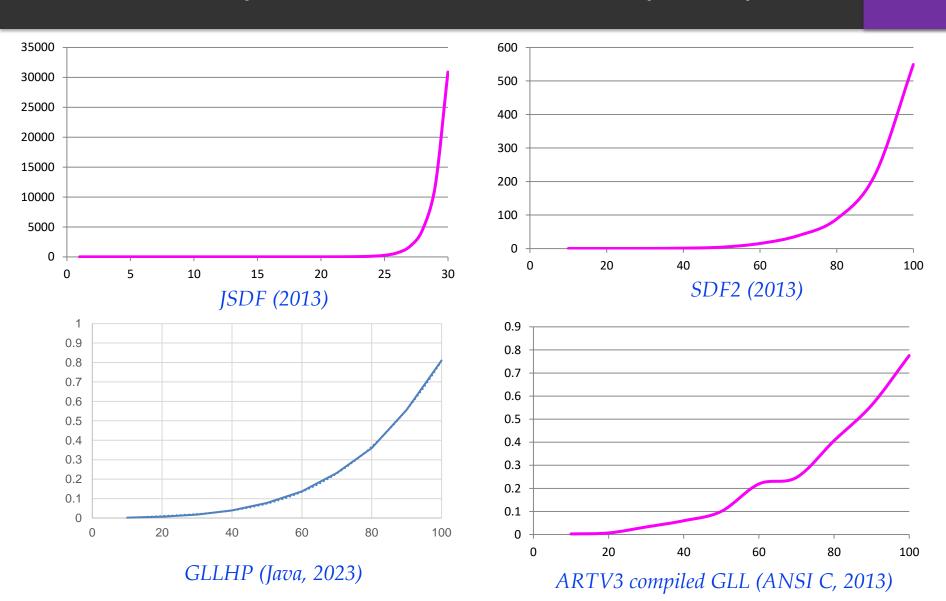
General parsers: go anywhere Slow and bulky

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https://www.youtube.com/watch?app=desktop&v=Cz1BpbsbFkA

### Worst case: parse $b^n$ with $S := b \mid S S \mid S S$



### Paper overview

Notation and compiled vs interpreted parsers Grammar representation Parser context

Backtracking recursive descent (neophyte friendly) From compiled to state based interpretation

Baseline – gllBL (Java API user friendly)

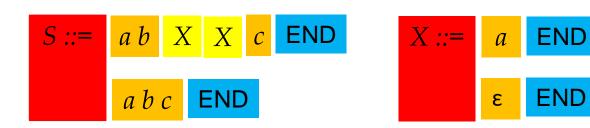
Memory management and hash pools Memory efficient baseline – glIHP (hash pool)

Variants to be characterised

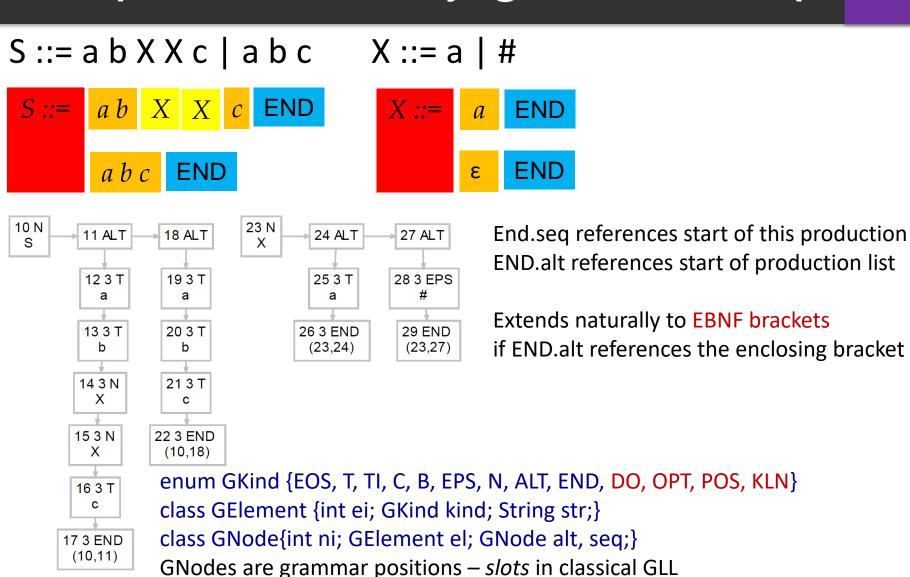
### GLL control flow fragments

A CFG is a nondeterministic specification of a language GLL provides one style of sequentialisation for the general parsing problem

The basic idea is to split the grammar specification up into a set of GLL fragments roughly corresponding to the basic blocks in a recursive descent parser



# Interpreter friendly grammar rep



### The six sets

#### **Threads**

Thread descriptors

#### GLL BL declarations

class Descriptor {GNode gn; int i; SNode sn; DNode dn;}

#### Stack

GSS nodes GSS edges Pops

Pops are distributed over the GSS nodes (Set of pairs in classical GLL)

#### **Derivations**

SPPF nodes
SPPF packed nodes

class SPPFN {GNode gn; int li; int ri; Set<SPPFPN> packNS;}

class SPPFPN {GNode gn; int pivot; SPPFN IC; SPPFN rC;}

### Baseline gll

```
void gllBL() {
 initialise();
 nextDescriptor: while (dequeueDesc())
 while (true) {
  switch (gn.elm.kind) {
   case B,T,TI,C: if (input[i] == qn.elm.ei)
           {du(1); i++; gn = gn.seq; break;}
           else continue nextDescriptor;
   case N: call(gn); continue nextDescriptor;
   case EPS: du(0); gn = gn.seq; break;
   case END: ret(); continue nextDescriptor;
   case DO: gn = gn.alt; break;
   case ALT:
     for (GNode tmp = gn; tmp != null; tmp = tmp.alt)
       queueDesc(tmp.seq, i, sn, dn);
     continue nextDescriptor;
}}}
```

## Control flow sequentialisation

```
void call(GNode gn) {
GSSN gssN = gssFind(gn.seq, i);
GSSE gssE = new GSSE(sn, dn);
if (!gssN.edges.contains(gssE)) {
 gssN.edges.add(gssE);
 for (SPPFN rc : gssN.pops)
  queueDesc(gn.seq, rc.ri, sn,
       sppfUpdate(gn.seq, dn, rc));
for (GNode p = rules(gn).alt;
     p != null; p = p.alt)
 queueDesc(p.seq, i, gssN, null);
```

```
void ret() {
 if (sn.equals(gssRoot)) { // Stack base
  if (accepting(gn) &&
     (i == input.length - 1)) {
    sppfRootNode =
    sppf.get(new SPPFN(
    startNonterminal, 0, input.length - 1));
    accepted = true;
  } else
    rightmostParseIndex = sppfWidestIndex();
  return; // End of parse
 sn.pops.add(dn);
 for (GSSE e : sn.edges)
  queueDesc(sn.gn, i, e.dst,
         sppfUpdate(sn.gn, e.sppfnode, dn));
```

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     p != null; p = p.alt)
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```

### Memory management

Allocate pool blocks using Java runtime

One hash table per set; hand-code offsets to set element fields

No safety net, and anathema to OO programmers

```
/*Constant field offsets. Offset zero links to next in hash chain */
protected final int gssNode gn = 1; // Key
protected final int gssNode i = 2; // Key
protected final int gssNode edgeList = 3;
protected final int gssNode popList = 4;
protected final int gssNode SIZE = 5; // Key size 2
protected final int gssEdge src = 1; // Key
protected final int gssEdge dst = 2; // Key
/* Lookup key <a,b>
   If not found, allocate allocationSize and load <a,b> to offsets (1,2) */
 protected void find(int[] hashBuckets, int hashBucketCount, int allocationSize,
                     int a, int b) {
   hash(hashBucketCount, a, b);
   findIndex = hashBuckets[hashResult];
   do {
     findBlockIndex = findIndex >> poolAddressOffset;
     findOffset = findIndex & poolAddressMask;
```

### Characterising the variants

Target: fast BRNGLR parsing (ToPLaS05, Acta Inf.07)

Parenthesised BNF and FBNF (SLE 10)

Lookahead (SCP 13)

EBNF (SCP 18)

Generated parsers: ANSI C; Java fragmentation issue

RGLL (SCP 16)

**CNP (SCP 16)** 

BSR sets (SCP 16)

MGLL and multiparsing (ToPLaS23)

# Good enough for GNU?

GEG - no more than 10% slow down against g++ -Ofast gcc parser went hand-crafted around 2004

```
One (large - 117 Kbyte) example:
g++ 18.8s
gllHP on ANSI-C 1.35s
gllHP on ANSI-C++ 7.02s
```

### Now add these optimisations:

Lookahead and iteration for EBNF BSR sets instead of full SPPF representation Generated C parser

### Repos

Code and small examples from the paper <a href="https://github.com/AJohnstone2007/referenceImplementation">https://github.com/AJohnstone2007/referenceImplementation</a>

Java 18 and SML grammars with large test sets (from ToPLaS23) <a href="https://github.com/AJohnstone2007/referenceLanguageCorpora">https://github.com/AJohnstone2007/referenceLanguageCorpora</a>

Research papers mentioned on slide 3 <a href="https://pure.royalholloway.ac.uk/en/persons/adrian-johnstone/publications">https://pure.royalholloway.ac.uk/en/persons/adrian-johnstone/publications</a>

Future investigations

Linear handling of deterministic sub languages (cf Scott McPeak CC04 – GLR/LR parser)

Threaded multicore

**Threaded**