

Experiment No. 1

Selection sort.

Aim -

Implementation of selection sort.

Introduction -

The selection sort algo is a simple, yet effective sorting algo. A selection-based sorting algo is described as in place comparison-based algo that divides the list into 2 parts, the sorted part on the left & the unsorted part on the right. Initially, the sorted selection is empty, & the unsorted selection contains the entire list.

Example -

$\text{arr}[] = 64, 25, 12, 22, 2 \quad //$

// find the minimum element in $\text{arr}[0 \dots 4]$

// And place it at beginning

12 25 12 22 64

// Find the minimum element in $\text{arr}[1 \dots 4]$

// And place it at beginning of $\text{arr}[1 \dots 4]$

12 25 22 64

// Find the minimum element in $\text{arr}[2 \dots 4]$

// And place it at beginning of $\text{arr}[2 \dots 4]$

12 22 25 64

// Find the minimum element in $\text{arr}[3 \dots 4]$

// And place it at beginning of $\text{arr}[3 \dots 4]$

12 22 25 24

Time complexity -

Let's assume $T[n]$ defines the running time to solve the problem of size n . In selection sort, after each iteration, one element gets sorted & problem size reduces by one element for each problem of size, inner loop iterates n times. So recurrence eqn for selection sort can be written as

$$T(n) = T(n-1) + n$$

put $n = n-1$ in above eqn

$$T(n-1) = T(n-2) + (n-1)$$

put value of $T(n-1)$ in previous eqn of $T(n)$

$$T(n) = T(n-2) + (n-1) + n$$

put $n = n-2 \neq n$ in above eqn

$$T(n-2) = T(n-3) + n-2$$

put value of $T(n-2)$ in previous eqn of $T(n)$

$$\therefore T(n) = T(n-3) + (n-3) + (n-1) + n$$

After k iterations

$$T(n-k) = T(n-k-1) + (n-k)$$

$$T(n) = T(n-k) + (n-k+1) + (n-k+2) + \dots + (n-1) + n$$

when k approaches to n

$$T(n) = T(0) + 1 + 2 + 3 + \dots + (n-1) + n$$

$T(0) = 0$, because it is running time of problem of size zero, & no effort it is needed to solve problem.

$$T(n) = 1 + 2 + 3 + \dots + n$$

$$= \frac{n(n+1)}{2} = \frac{n^2 + n}{2} = \frac{n^2}{2} + \frac{n}{2}$$

$$= O \left(\max \left(\frac{n^2}{n}, \frac{n}{n} \right) \right) = O \left(\frac{n^2}{2} \right)$$

$$T(n) = O(n^2)$$

Best case

$$O(n^2)$$

Avg case

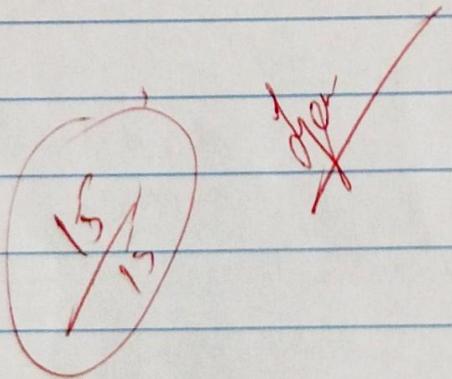
$$O(n^2)$$

worst case

$$O(n^2)$$

Conclusion -

Hence we have successfully implemented algo & analyzed the complexity.



```
import sys  
A = [64, 25, 12, 22, 11]  
  
for i in range (len(A)):  
  
    min_idx = i  
  
    for j in range(i+1, len(A)):  
  
        if A[min_idx] > A[j]:  
  
            min_idx = j  
  
    A[i], A[min_idx] = A[min_idx], A[i]  
  
print ("Sorted array")  
  
for i in range(len(A)):  
  
    print ("%d" %A[i],end=" ")
```

PROBLEMS OUTPUT DEBUG CONSOLE TERMINAL

```
SyntaxError: invalid syntax
* (base) computer@computer:~$ /usr/bin/python3.9 "/home/computer/selec
  File "/home/computer/selection sort.py", line 11
    A[i], A[min_idx] = A[min_idx]: A[i]
                                ^
SyntaxError: invalid syntax
* (base) computer@computer:~$ /usr/bin/python3.9 "/home/computer/selec
  Sorted array
* 1112222564(base) computer@computer:~$ /usr/bin/python3.9 "/home/compu
  Sorted array
* 1112222564(base) computer@computer:~$ /usr/bin/python3.9 "/home/compu
  Sorted array
* 11 12 22 25 64 (base) computer@computer:~$ []
```

Experiment No.2

Merge Sort

Aim - Implementation of merge sort.

Theory -

Algo:-

- 1) If it is only one element in the list it is already sorted, return
- 2) Divide the list recursively into two halves until it can no more divided.
- 3) Merge the smaller list into new list in sorted order.

Example -

Mergesort (arr[], l, r)

if $r > l$

1. Find the middle point to divide the array into two halves
 $\text{middle } m = l + (r-1)/2$
2. Call the mergesort for 1st half:
 call mergesort $m = (\text{arr}, l, m)$
3. Call mergesort for 2nd half:
 call mergesort $(\text{arr}, m+1, r)$
4. Merge the 2 halves sorted in step 2 & 3
 call merge (arr, l, m, r)

Time complexity -

Best	Average	Worst
$\Omega(n \log(n))$	$\Theta(n \log(n))$	$\Omega(n \log(n))$

Time complexity -

Merge sort is recursive algo & time complexity can be expressed as following recurrence relation.



$$T(n) = 2T(n/2) + O(n)$$

Conclusion -

Hence, we have successfully implemented algorithm & analyzed the complexity.

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```
#Python program to perform Merge Sort
#Function to merge and sort elements

def merge_sort(arr):
    if len(arr)>1:
        m=len(arr)//2
        Left=arr[:m]
        Right=arr[m:]
        merge_sort(Left)
        merge_sort(Right)

    i=j=k=0

    while i<len(Left) and j<len(Right):
        if Left[i]<Right[j]:
            arr[k]=Left[i]
            i+=1
        else :
            arr[k]=Right[j]
            j+=1
        k+=1

    while i<len(Left):
        arr[k]=Left[i]
        i+=1
    k+=1

#Inserting an no of elements
n =int(input("Enter no of elements in an array"))

array=[]

#Taking input
print("Enter elements in array")
for i in range(0,n):
    b=int(input())
    array.append(b)

print('Given Array is',array)
merge_sort(array)
print("Sorted array is ",array)
```

~~Q~~

```
* (base) computer@computer-ThinkCentre:~/Downloads/SIMPLE_INVENTORY_SYSTEM_IN_PYTHON_WITH_SOURCE_CODE/Simple_Inventory PYTHON/Simple Inventory Systems /bin/python3 "/home/computer/Downloads/SIMPLE_INVENTORY_SYSTEM_IN_PYTHON_WITH_SOURCE_CODE/Simple_Inventory PYTHON/Simple Inventory System/Merge Sort.py"
Enter no of elements in an array5
Enter elements in array
23
576
12
68
99
Given Array is [23, 576, 12, 68, 99]
Sorted array is [12, 23, 68, 99, 576]
```

o (base) computer@computer-ThinkCentre:~/Downloads/SIMPLE_INVENTORY_SYSTEM_IN_PYTHON_WITH_SOURCE_CODE/Simple_Inventory PYTHON/Simple Inventory Systems █