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/*
 * The background pageout daemon, started as a kernel thread
 * from the init process.
 * This basically executes once a second, trickling out pages
 * so that we have _some_ free memory available even if there
 * is no other activity that frees anything up. This is needed
 * for things like routing etc, where we otherwise might have
 * all activity going on in asynchronous contexts that cannot
 * page things out.
 * If there are applications that are active memory-allocators
 * (most normal use), this basically shouldn't matter.
int kswapd(void *unused)
    struct task struct *tsk = current;
    kswapd_process = tsk;
    tsk->session = 1;
    tsk->pgrp = 1;
    strcpy(tsk->comm, "kswapd");
    sigfillset(&tsk->blocked);
     * Tell the memory management that we're a "memory allocator",
     * and that if we need more memory we should get access to it
     * regardless (see "__get_free_pages()"). "kswapd" should
     * never get caught in the normal page freeing logic.
     * (Kswapd normally doesn't need memory anyway, but sometimes
     * you need a small amount of memory in order to be able to
     * page out something else, and this flag essentially protects
     * us from recursively trying to free more memory as we're
     * trying to free the first piece of memory in the first place).
    tsk->flags |= PF_MEMALLOC;
    while (1) {
        /*
          * Wake up once a second to see if we need to make
          * more memory available.
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* If we actually get into a low-memory situation,
         * the processes needing more memory will wake us
         * up on a more timely basis.
         */
        do {
            if (nr_free_pages >= freepages.high)
                break:
            if (!do_try_to_free_pages(GFP_KSWAPD))
                break:
        } while (!tsk->need_resched);
        run_task_queue(&tq_disk);
        tsk->state = TASK_INTERRUPTIBLE;
        schedule_timeout(HZ);
    }
}
 * We need to make the locks finer granularity, but right
 * now we need this so that we can do page allocations
 * without holding the kernel lock etc.
 * We want to try to free "count" pages, and we need to
 * cluster them so that we get good swap-out behaviour. See
 * the "free_memory()" macro for details.
 */
static int do_try_to_free_pages (unsigned int gfp_mask)
{
    int priority;
    int count = SWAP_CLUSTER_MAX;
    lock_kernel();
    /* Always trim SLAB caches when memory gets low. */
    kmem_cache_reap(gfp_mask);
    priority = 6;
    do {
        while (shrink_mmap(priority, gfp_mask)) {
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if (!--count)
                 goto done;
        }
        /* Try to get rid of some shared memory pages.. */
        if (gfp_mask & __GFP_IO) {
             while (shm_swap(priority, gfp_mask)) {
                 if (!--count)
                     goto done;
             }
        }
        /* Then, try to page stuff out.. */
        while (swap_out(priority, gfp_mask)) {
             if (!--count)
                 goto done;
        }
        shrink_dcache_memory(priority, gfp_mask);
    } while (--priority \geq = 0);
done:
    unlock_kernel();
    return priority >= 0;
}
int shrink_mmap(int priority, int gfp_mask)
{
    static unsigned long clock = 0;
    unsigned long limit = num_physpages;
    struct page * page;
    int count;
    count = limit >> priority;
    page = mem_map + clock;
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do {
    int referenced:
    /* This works even in the presence of PageSkip because
     * the first two entries at the beginning of a hole will
     * be marked, not just the first.
     */
    page++;
    clock++;
    if (clock >= max_mapnr) {
        clock = 0;
        page = mem_map;
    if (PageSkip(page)) {
        /* next_hash is overloaded for PageSkip */
        page = page->next_hash;
        clock = page - mem_map;
    }
    referenced = test_and_clear_bit(PG_referenced, &page->flags);
    if (PageLocked(page))
        continue:
    if ((gfp_mask & __GFP_DMA) && !PageDMA(page))
        continue;
    /* We can't free pages unless there's just one user */
    if (atomic_read(&page->count) != 1)
        continue;
    count--;
     * Is it a page swap page? If so, we want to
     * drop it if it is no longer used, even if it
     * were to be marked referenced..
     */
    if (PageSwapCache(page)) {
        if (referenced && swap_count(page->offset) != 1)
             continue;
        delete_from_swap_cache(page);
        return 1;
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}
        if (referenced)
            continue;
        /* Is it a buffer page? */
        if (page->buffers) {
            if (buffer_under_min())
                continue;
            if (!try_to_free_buffers(page))
                continue;
            return 1;
        }
        /* is it a page-cache page? */
        if (page->inode) {
            if (pgcache_under_min())
                continue;
            remove_inode_page(page);
            return 1;
        }
    \} while (count > 0);
    return 0;
}
 * This must be called only on pages that have
 * been verified to be in the swap cache.
void delete_from_swap_cache(struct page *page)
    long entry = page->offset;
   remove_from_swap_cache (page);
    swap_free (entry);
}
static inline void remove_from_swap_cache(struct page *page)
    if (!page->inode) {
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printk ("VM: Removing swap cache page with zero inode hash "
            "on page %08lx\n", page_address(page));
        return;
    }
    if (page->inode != &swapper_inode) {
        printk ("VM: Removing swap cache page with wrong inode hash"
            "on page %08lx\n", page_address(page));
    }
    PageClearSwapCache (page);
   remove_inode_page(page);
}
void remove_inode_page(struct page *page)
         remove_page_from_hash_queue(page);
         remove_page_from_inode_queue(page);
         page_cache_release(page);
}
int shm_swap (int prio, int gfp_mask)
{
    pte_t page;
    struct shmid_kernel *shp;
    unsigned long swap_nr;
    unsigned long id, idx;
    int loop = 0;
    int counter;
    counter = shm_rss >> prio;
    if (!counter || !(swap_nr = get_swap_page()))
        return 0;
 check_id:
    shp = shm_segs[swap_id];
    if (shp == IPC\_UNUSED \parallel shp == IPC\_NOID \parallel shp->u.shm\_perm.mode &
SHM_LOCKED) {
        next_id:
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swap\_idx = 0;
       if (++swap_id > max_shmid) {
           swap\_id = 0;
           if (loop)
               goto failed;
           loop = 1;
       goto check_id;
   }
  id = swap_id;
check_table:
  idx = swap\_idx++;
  if (idx >= shp->shm_npages)
       goto next_id;
  page = __pte(shp->shm_pages[idx]);
  if (!pte_present(page))
       goto check_table;
  if((gfp_mask&__GFP_DMA)&&
      !PageDMA(&mem_map[MAP_NR(pte_page(page))]))
       goto check_table;
  swap_attempts++;
  if (--counter < 0) { /* failed */
       failed:
       swap_free (swap_nr);
       return 0;
   }
  if (atomic_read(&mem_map[MAP_NR(pte_page(page))].count) != 1)
       goto check_table;
  shp->shm_pages[idx] = swap_nr;
  rw_swap_page_nocache (WRITE, swap_nr, (char *) pte_page(page));
  free_page(pte_page(page));
  swap_successes++;
  shm_swp++;
  shm_rss--;
  return 1;
```

}