



**NITTE**  
EDUCATION TRUST

**NMAM INSTITUTE  
OF TECHNOLOGY**

REPORT  
ON  
CD Mini Project  
Carried out on

## **Compiler for Finding the Largest Number using IF statement**

*Submitted to*

**NMAM INSTITUTE OF TECHNOLOGY, NITTE**  
(An Autonomous Institution under VTU, Belagavi)

*In partial fulfillment of the requirements for the award of the*

Degree of Bachelor of  
Engineering in

**Computer Science and Engineering**

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## CERTIFICATE

*This is to certify that Mr. Abdeali bearing USN 4NM19CS002, Mr. Akash L M bearing USN 4NM19CS012, Mr. Aniket Ajit Pai bearing USN 4NM19CS020 of seventh semester B.E., bonafide students of NMAM Institute of Technology, Nitte, have completed CD mini project on “ Designing the compiler for the hypothetical program to find the largest number using if statement ” during October 2022 -December 2022 fulfilling the partial requirements for the award of degree of Bachelor of Engineering in **Computer Science and Engineering** at NMAM Institute of Technology, Nitte.*

*Name and Signature of Mentor  
(Dr Arista P Rodrigues)*

*Signature of HOD*

## **ACKNOWLEDGEMENT**

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# **1. ABSTRACT**

The purpose of this project is to design lexical analyzer and syntax analyzer for a Context Free Grammar. The two stages are the integral part of the Analysis phase of a compilation process which involves identifying the tokens of the given program and using these tokens to identify if each of them is syntactically proper based on given production rules.

The main program takes in two parts namely the source program which we need to process and the grammar rules to parse the program. The objective of the project is to generate the parsed sequence which can be further given for the later stages of the compiler.

## 2. INTRODUCTION

### COMPILER

A compiler is a software that takes a program written in a high-level language and translates it into an equivalent program in a target language. Most specifically a compiler takes a computer program and translates it into an object program. Some other tools associated with the compiler are responsible for making an object program into executable form

**Source program** – It is normally written in a high-level programming language. It contains a set of rules, symbols and special words used to construct a computer program.

**Target program** – It is normally the equivalent program in machine code. It contains the binary representation of the instructions that the hardware of the computer can perform.

**Error Message** – A message issued by the compiler due to detection of syntax errors in the source program.

Compilation is a large process. It is often broken into stages. An efficient compiler must preserve semantics of the source program and it should create an efficient version of the target language.

### PHASES OF COMPILERS

Typically, a compiler includes several functional parts. For example, a conventional compiler may include a lexical analyzer that looks at the source program and identifies successive “tokens” in the source program. A conventional compiler also includes a parser or syntactical analyzer, which takes as an input a grammar defining the language.

## **CLASSIFICATION OF COMPILER PHASES**

There are two major parts of a compiler phases: Analysis and Synthesis. In analysis phase, an intermediate representation is created from the given source program that contains:

- Lexical Analyzer
- Syntax Analyzer
- Semantic Analyzer

In the synthesis phase, the equivalent target program is created from this intermediate representation. This contains:

- Intermediate code Generator
- Code Optimization
- Code Generation

## **LEXICAL ANALYZER**

Lexical analyzer takes the source program as an input and produces a string of tokens or lexemes. Lexical Analyzer reads the source program character by character and returns the tokens of the source program. The process of generation and returning the tokens is called lexical analysis. Representation of lexemes in the form of tokens as:  
<token-name, attribute-value>

## **SYNTAX ANALYSER**

A Syntax Analyzer creates the syntactic structure (generally a parse tree) of the given program. In other words, a Syntax Analyzer takes the output of a lexical analyzer (list of tokens) and produces a parse tree. A syntax analyzer is also called a parser. The parser checks if the expression made by the tokens is syntactically correct.

## **SEMANTIC ANALYSER**

Semantic analyzer takes the output of syntax analyzer. Semantic analyzer checks a source program for semantic consistency with the language definition. It also gathers type information for use in intermediate-code generation.

## **INTERMEDIATE CODE GENERATION**

After semantic analysis, the compiler generates an intermediate code of the source code for the target machine. It represents a program for some abstract machine. It is in between the high-level language and the machine language.

## **CODE OPTIMISER**

The code optimizer takes the code produced by the intermediate code generator. The code optimizer reduces the code (if the code is not already optimized) without changing the meaning of the code. The optimization of code is in terms of time and space.

## **CODE GENERATION**

This produces the target language in a specific architecture. The target program is normally an object file containing the machine codes. Memory locations are selected for each of the variables used by the program.

## **SYMBOL TABLE**

It is a data-structure maintained throughout all the phases of a compiler. All the identifiers' names along with their types are stored here. The symbol table makes it easier for the compiler to quickly search the identifier record and retrieve it. The symbol table is also used for scope management.



## LEXICAL ANALYSIS

Lexical analysis is the first phase of a compiler. It takes the modified source code from language preprocessors that are written in the form of sentences. The lexical analyzer breaks these syntaxes into a series of tokens, by removing any whitespace or comments in the source code. If the lexical analyzer finds a token invalid, it generates an error. The lexical analyzer works closely with the syntax analyzer. It reads character streams from the source code, checks for legal tokens, and passes the data to the syntax analyzer when it demands. The main purpose of lexical analysis is to make life easier for the subsequent syntax analysis phase.

### **Token:**

Token is a sequence of characters that can be treated as a single logical entity. Typical tokens are,

- 1) Identifiers
- 2) keywords
- 3) operators
- 4) special symbols
- 5) constants

### **Pattern:**

A set of strings in the input for which the same token is produced as output. This set of strings is described by a rule called a pattern associated with the token.

### **Lexeme:**

A lexeme is a sequence of characters in the source program that is matched by the pattern for a token. We can implement lexical analyses using lex tools

### 3. PROBLEM STATEMENT & METHODOLOGY

**Q. Find the largest number using if statement**

```
int main()
begin
    int n1, n2, n3;
    if( expr relop expr )begin
        printf( n1);
    end
    if ( expr relop expr )begin
        printf( n2);
    end
    if( expr relop expr )begin
        printf( n3);
    end
end
```

#### **Algorithm to construct LL(1) Parsing Table:**

**Step 1:** First check for left recursion in the grammar, if there is left recursion in the grammar remove that and go to step 2.

**Step 2:** Calculate First() and Follow() for all non-terminals.

1. **First()**: If there is a variable, and from that variable, if we try to derive all the strings then the beginning Terminal Symbol is called the First.

2. **Follow()**: What is the Terminal Symbol which follows a variable in the process of derivation.

**Step 3:** For each production  $A \rightarrow \alpha$ . (A tends to alpha)

1. Find First( $\alpha$ ) and for each terminal in First( $\alpha$ ), make entry  $A \rightarrow \alpha$  in the table.

2. If  $\text{First}(\alpha)$  contains  $\epsilon$  (epsilon) as a terminal then, find the  $\text{Follow}(A)$  and for each terminal in  $\text{Follow}(A)$ , make entry  $A \rightarrow \alpha$  in the table.

3. If the  $\text{First}(\alpha)$  contains  $\epsilon$  and  $\text{Follow}(A)$  contains  $\$$  as terminal, then make entry  $A \rightarrow \alpha$  in the table for the  $\$$ .

### **Construction of SLR parsing table :**

1. Construct  $C = \{ I_0, I_1, \dots, I_n \}$ , the collection of sets of LR(0) items for  $G'$ .

2. State  $i$  is constructed from  $I_i$ . The parsing actions for state  $i$  are determined as follow :

- If  $[A \rightarrow ?.a]$  is in  $I_i$  and  $\text{GOTO}(I_i, a) = I_j$ , then set  $\text{ACTION}[i, a]$  to “shift  $j$ ”.

Here  $a$  must be terminal.

- If  $[A \rightarrow ?.]$  is in  $I_i$ , then set  $\text{ACTION}[i, a]$  to “reduce  $A \rightarrow ?$ ” for all  $a$  in  $\text{FOLLOW}(A)$ ; here  $A$  may not be  $S'$ .

- If  $[S \rightarrow S.]$  is in  $I_i$ , then set action  $[i, \$]$  to “accept”. If any conflicting actions are generated by the above rules, we say that the grammar is not SLR.

3. The goto transitions for state  $i$  are constructed for all nonterminals  $A$  using the rule:

if  $\text{GOTO}(I_i, A) = I_j$  then  $\text{GOTO}[i, A] = j$ .

4. All entries not defined by rules 2 and 3 are made error.

## 4. RESULT

### Code to generate tokens:

```
import re
output = [ ]
def isValidDelimiter(ch):
    if( ch == ' ' or ch == '+' or ch == '-' or ch == '*' or ch == '/' or ch == ';' or ch == ':'
    or ch == '>' or ch == '<' or ch == '=' or ch == '(' or ch == ')' or ch == '[' or ch == ']'):
        return True
    else:
        return False

def isValidOperator(ch):
    if(ch == '+' or ch == '-' or ch == '*' or ch == '/' or ch == '>' or ch == '<' or ch ==
    '='):
        return True
    else:
        return False
    " " "regex = '^[-\^*>=<=]$(\)\{\}\[ ]'
    print(str)
    if(re.search(regex, str)):
        return True
    else:
        return False" " "

def isValidIdentifier(str):
    " " "if(str[0] == '0' or str[0] == '1' or str[0] == '2' or str[0] == '3' or str[0] == '4' or
    str[0] == '5' or str[0] == '6' or str[0] == '7' or str[0] == '8' or str[0] == '9' or
    isValidDelimiter(str[0]) == True):
        return False
    else:
        return True" " "
    regex = '^([A-Za-z_][A-Za-z0-9_]*)'
    if(re.search(regex, str)):
        return True
    else:
        return False

def isValidKeyword(str):
    lst = ["if", "else", "while", "break", "begin", "end", "printf", "int", "main"]
    if str in lst:
        return True
    else:
        return False

def isValidInteger(str):
```

```

    return bool(re.search(r"^\d+$", str))

def isRealNumber(str):
    return bool(re.search(r"^\d+\.\d+$", str))

def isRealNumber(str):
    return bool(re.search(r"^\d+\.\d+$", str))

def takeSub(str, left, right):
    return str[left:right]
    """print(str)
    lst = list(str)
    newlst = []
    for i in range(left, right+1):
        newlst.append(lst[i])
    #sub[right-left+1] = '\0'
    return str(lst)"""

def detectTokens(str):
    left = 0
    right = 0
    length = len(str)

    # print(length)
    try:
        while (right<= length and left <= right):
            #print(left)
            if isValidDelimiter(str[right])==False :
                right=right+1

            if (isValidDelimiter(str[right])==True and left == right):
                if (isValidOperator(str[right]) ==True):
                    print(("|   Operator : "+ '\t'+ sub + '\t' + "|").expandtabs(10))
                    output.append(str[right])
                    if str[right]==")" or str[right]=="(" or str[right] == '+' or str[right] == '-' or
str[right] == '*' or str[right] == '/' or str[right] == ',' or str[right] == ';':
                        # print("Here ", str[right])
                        output.append(str[right])
                        right=right+1
                        left = right
                    elif (isValidDelimiter(str[right]) == True and left != right or (right == length
and left != right)):
                        sub = takeSub(str, left, right)
                        # print("sub",sub)
                        # right=right+1
                        if isValidKeyword(sub)==True:
                            print(("|   Keyword :"+ '\t'+ sub + '\t' + "|").expandtabs(10))
                            output.append(sub)
                        elif isValidInteger(sub) == True:
                            print(("|   Integer :"+ '\t'+ sub + '\t' + "|").expandtabs(10))
                            output.append(sub)
                        elif (isRealNumber(sub) == True):

```

```

        print("|    Real Number : "+ '\t'+ sub + '\t' + "|").expandtabs(10))
        output.append(sub)
    elif (isvalidIdentifier(sub) == True and isValidDelimiter(str[right - 1]) ==
False):
        print("|    Identifier : "+ '\t'+ sub + '\t' + "|").expandtabs(10))
        for i in range(len(sub)):
            output.append(sub[i])
    elif (isvalidIdentifier(sub) == False and isValidDelimiter(str[right - 1]) ==
False):
        print("Invalid Identifier : "+ '\t'+ sub + '\t' + "|").expandtabs(10))
        output.append(sub)
    left = right

except:
    print("|    keyword : "+ '\t'+ "end" + '\t' + "|").expandtabs(10))
    output.append("end")

f = open("InputProg.txt", "r")
str = f.readline()

print("\n
\n.....
.....")

print("CODE : ", str)
print(".....")

print("All Tokens are :")
try:
    detectTokens(str)
except :
    pass

def listToString(instr):
    emptystr=""
    for ele in instr:
        emptystr += ele
        emptystr += " "
    return emptystr

temp =listToString(output)

print(".....
.....")
print(listToString(output))
print(".....
.....")

f2 = open("intermediate.txt", "w")
f2.write(temp)

```

## Productions:

Grammer Productions:

```
-----  
S->['int', "S"]  
S'->['main', '(', ')', 'Y']  
Y->['begin', 'X', 'end']  
X->['D', 'B']  
D->['int', 'A', ',', 'A', ',', 'A', ';']  
B->['IF']  
A->['a', 'G']  
G->['a', 'G']  
IF->['if', '(', 'A', 'OP', 'A', ')', 'C']  
C->['Y']  
OP->['<']  
-----
```

## Output of tokenizer:

```
CODE :  int main ( ) begin int a1 , b , c ; if ( a > b ) begin printf ( a ) ; end end
```

All Tokens are :

Keyword :	int
Keyword :	main
Keyword :	begin
Keyword :	int
Identifier :	a1
Identifier :	b
Identifier :	c
Keyword :	if
Identifier :	a
Operator :	a
Identifier :	b
Keyword :	begin
Keyword :	printf
Identifier :	a
Keyword :	end
keyword :	end

```
int main ( ) begin int a 1 , b , c ; if ( a > b ) begin printf ( a ) ; end end
```

### Output for invalid string:

```
Code -> int main ( ) begin int abcd1213 , b , c ; if ( a > b ) begin printf ( a )
end end
```

	Buffer	Stack	Action
\$ end end	a (printf begin) b > a (if ; c, b, 321 d c b a int begin) (	main int	S \$ T[S][int] = S->int S'
\$ end end	a (printf begin) b > a (if ; c, b, 321 d c b a int begin) (	main int	int S' \$ Matched:int
\$ end end	a (printf begin) b > a (if ; c, b, 321 d c b a int begin) (	main	S' \$ T[S'] [main] = S' ->main ( ) Y
\$ end end	a (printf begin) b > a (if ; c, b, 321 d c b a int begin) (	main	main ( ) Y \$ Matched:main
\$ end end	a (printf begin) b > a (if ; c, b, 321 d c b a int begin) (	( ) Y \$	Matched:(
\$ end end	a (printf begin) b > a (if ; c, b, 321 d c b a int begin) (	Y \$	Matched:)
\$ end end	a (printf begin) b > a (if ; c, b, 321 d c b a int begin) (	begin X end \$	T[V][begin] = Y->begin X end Matched:begin
\$ end end	a (printf begin) b > a (if ; c, b, 321 d c b a int	X end \$	T[X][int] = X->B
\$ end end	a (printf begin) b > a (if ; c, b, 321 d c b a int	D B end \$	T[D][int] = D->int A, A, A ;
\$ end end	a (printf begin) b > a (if ; c, b, 321 d c b a int	int A, A, A ; B end \$	Matched:int
\$ end end	a (printf begin) b > a (if ; c, b, 321 d c b a	A, A, A ; B end \$	T[A][a] = A->A G
\$ end end	a (printf begin) b > a (if ; c, b, 321 d c b a	a G, A, A ; B end \$	Matched:a
\$ end end	a (printf begin) b > a (if ; c, b, 321 d c b a	a G, A, A ; B end \$	T[G][b] = G->b G
\$ end end	a (printf begin) b > a (if ; c, b, 321 d c b a	b G, A, A ; B end \$	Matched:b
\$ end end	a (printf begin) b > a (if ; c, b, 321 d c	G, A, A ; B end \$	T[G][c] = G->c G
\$ end end	a (printf begin) b > a (if ; c, b, 321 d c	c G, A, A ; B end \$	Matched:c
\$ end end	a (printf begin) b > a (if ; c, b, 321 d	G, A, A ; B end \$	T[G][d] = G->d G
\$ end end	a (printf begin) b > a (if ; c, b, 321 d	d G, A, A ; B end \$	Matched:d
\$ end end	a (printf begin) b > a (if ; c, b, 321	G, A, A ; B end \$	T[G][i] = G->i G
\$ end end	a (printf begin) b > a (if ; c, b, 321	1 G, A, A ; B end \$	Matched:1
\$ end end	a (printf begin) b > a (if ; c, b, 32	G, A, A ; B end \$	T[G][2] = G->2 G
\$ end end	a (printf begin) b > a (if ; c, b, 32	2 G, A, A ; B end \$	Matched:2
\$ end end	a (printf begin) b > a (if ; c, b, 3	G, A, A ; B end \$	T[G][3] = G->3 G
\$ end end	a (printf begin) b > a (if ; c, b, 3	3 G, A, A ; B end \$	Matched:3
\$ end end	a (printf begin) b > a (if ; c, b,	G, A, A ; B end \$	T[G][.] = G->#
\$ end end	a (printf begin) b > a (if ; c, b,	A, A, A ; B end \$	Matched:.
\$ end end	a (printf begin) b > a (if ; c, b,	A, A ; B end \$	T[A][b] = A->b G
\$ end end	a (printf begin) b > a (if ; c, b	b G, A ; B end \$	Matched:b
\$ end end	a (printf begin) b > a (if ; c,	G, A ; B end \$	T[G][_] = G->#
\$ end end	a (printf begin) b > a (if ; c,	A ; B end \$	Matched:;
\$ end end	a (printf begin) b > a (if ; c	A ; B end \$	T[A][c] = A->c G
\$ end end	a (printf begin) b > a (if ; c	c G ; B end \$	Matched:c
\$ end end	a (printf begin) b > a (if ;	G ; B end \$	T[G][:] = G->#
\$ end end	a (printf begin) b > a (if ;	B end \$	Matched:;
\$ end end	a (printf begin) b > a (if	B end \$	T[B][f] = B->f
\$ end end	a (printf begin) b > a (if	I f ( A OP A ) C end \$	T[I f][i f] = I f->i f ( A OP A ) C
\$ end end	a (printf begin) b > a (if	I f ( A OP A ) C end \$	Matched:if
\$ end end	a (printf begin) b > a (	( A OP A ) C end \$	Matched:(
\$ end end	a (printf begin) b > a (	A OP A ) C end \$	T[A][a] = A->a G

```

$ end end) a (printf begin b > a (if ; c , b , 3 2 1 d c G , A , A ; B end $ T[G][c] = G->c G Matched:c
$ end end) a (printf begin b > a (if ; c , b , 3 2 1 d c C , G , A , A ; B end $ Matched:d
$ end end) a (printf begin b > a (if ; c , b , 3 2 1 d G , A , A ; B end $ T[G][d] = G->d G Matched:d
$ end end) a (printf begin b > a (if ; c , b , 3 2 1 d d , G , A , A ; B end $ Matched:d
$ end end) a (printf begin b > a (if ; c , b , 3 2 1 G , A , A ; B end $ T[G][1] = G->1 G Matched:1
$ end end) a (printf begin b > a (if ; c , b , 3 2 1 1 G , A , A ; B end $ T[G][2] = G->2 G Matched:2
$ end end) a (printf begin b > a (if ; c , b , 3 2 2 G , A , A ; B end $ T[G][3] = G->3 G Matched:3
$ end end) a (printf begin b > a (if ; c , b , 3 3 G , A , A ; B end $ Matched:3
$ end end) a (printf begin b > a (if ; c , b , 3 G , A , A ; B end $ T[G][_] = G-># Matched:_
$ end end) a (printf begin b > a (if ; c , b , A , A ; B end $ Matched:_
$ end end) a (printf begin b > a (if ; c , b , A , A ; B end $ T[A][b] = A->b G Matched:b
$ end end) a (printf begin b > a (if ; c , b , b , G , A ; B end $ Matched:b
$ end end) a (printf begin b > a (if ; c , G , A ; B end $ T[G][_ ] = G-># Matched:_
$ end end) a (printf begin b > a (if ; c , A , A ; B end $ Matched:_
$ end end) a (printf begin b > a (if ; c , A ; B end $ T[A][c] = A->c G Matched:c
$ end end) a (printf begin b > a (if ; c C , G ; B end $ Matched:c
$ end end) a (printf begin b > a (if ; c G , B end $ T[G][_ ] = G-># Matched:_
$ end end) a (printf begin b > a (if ; B end $ Matched:_
$ end end) a (printf begin b > a (if IF end $ T[B][if] = B->if Matched:if
$ end end) a (printf begin b > a (if IF end $ T[IF][if] = IF->if (A OP A) C Matched:if
$ end end) a (printf begin b > a (if (A OP A) C end $ Matched:(
$ end end) a (printf begin b > a (A OP A) C end $ T[A][a] = A->a G Matched:a
$ end end) a (printf begin b > a G OP A) C end $ Matched:a
$ end end) a (printf begin b > G OP A) C end $ T[G][>] = G->> Matched:>
$ end end) a (printf begin b > OP A) C end $ T[OP][>] = OP->> Matched:>
$ end end) a (printf begin b > A ) C end $ Matched:A
$ end end) a (printf begin b > A ) C end $ T[A][b] = A->b G Matched:b
$ end end) a (printf begin G ) C end $ T[G][()] = G-># Matched:_
$ end end) a (printf begin ) C end $ Matched:_
$ end end) a (printf begin C end $ T[C][begin] = C->Y Matched:Y
$ end end) a (printf begin Y end $ T[Y](y) = Y->ybegin X end Matched:begin
$ end end) a (printf begin X end $ T[X](printf) = x->printf ( A ) Matched:printf
$ end end) a (printf printf ( A ) ; end end $ Matched:printf
$ end end) a ( A ) ; end end $ Matched:(
$ end end) a A ; end end $ T[A][a] = A->a G Matched:a
$ end end) a G ; end end $ T[G][_] = G-># Matched:_
$ end end) ; end end $ Matched:_

```



### Output for invalid string with error recovery:

```
Code -> int main ( ) begin int abcd1213 , b , c ; if ( a > b ) begin printf ( a )
end end
```

Buffer	Stack	Action
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c b a int begin	main int	\$ S \$ T[S][int] = S->int S'
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c b a int begin	main int	int S' \$ Matched:int
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c b a int begin	main	S' \$ T[S'] [main] = S' ->main ( ) Y
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c b a int begin	main	main ( ) Y \$ Matched:main
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c b a int begin	( ) Y \$	Matched:(
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c b a int begin	Y \$	Matched:)
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c b a int begin	Y \$	T[V][begin] = Y->begin X end
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c b a int begin	begin X end \$	Matched:begin
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c b a int	X end \$	T[X][int] = X->D B
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c b a int	D B end \$	T[D][int] = D->int A , A , A ;
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c b a int	int A , A , A ; B end \$	Matched:int
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c b a	A , A , A ; B end \$	T[A][a] = A->a G
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c b a	a G , A , A ; B end \$	Matched:a
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c b	G , A , A ; B end \$	T[G][b] = G->b G
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c b	b G , A , A ; B end \$	Matched:b
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c	G , A , A ; B end \$	T[G][c] = G->c G
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d c	c G , A , A ; B end \$	Matched:c
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d	G , A , A ; B end \$	T[G][d] = G->d G
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1 d	d G , A , A ; B end \$	Matched:d
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1	G , A , A ; B end \$	T[G][1] = G->1 G
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2 1	1 G , A , A ; B end \$	Matched:1
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2	G , A , A ; B end \$	T[G][2] = G->2 G
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3 2	2 G , A , A ; B end \$	Matched:2
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3	G , A , A ; B end \$	T[G][3] = G->3 G
\$ end end ) a ( printf begin ) b > a ( if ; c , b , 3	3 G , A , A ; B end \$	Matched:3
\$ end end ) a ( printf begin ) b > a ( if ; c , b ,	G , A , A ; B end \$	T[G][,] = G->#
\$ end end ) a ( printf begin ) b > a ( if ; c , b ,	, A , A ; B end \$	Matched:,
\$ end end ) a ( printf begin ) b > a ( if ; c , b	A , A ; B end \$	T[A][b] = A->b G
\$ end end ) a ( printf begin ) b > a ( if ; c , b	b G , A , B end \$	Matched:b
\$ end end ) a ( printf begin ) b > a ( if ; c ,	G , A ; B end \$	T[G][,] = G->#
\$ end end ) a ( printf begin ) b > a ( if ; c ,	A ; B end \$	Matched:,
\$ end end ) a ( printf begin ) b > a ( if ; c	A ; B end \$	T[A][c] = A->c G
\$ end end ) a ( printf begin ) b > a ( if ; c	c G ; B end \$	Matched:c
\$ end end ) a ( printf begin ) b > a ( if ;	G ; B end \$	T[G][:] = G->#
\$ end end ) a ( printf begin ) b > a ( if ;	; B end \$	Matched:;
\$ end end ) a ( printf begin ) b > a ( if	B end \$	T[B][if] = B->I F
\$ end end ) a ( printf begin ) b > a ( if	I F end \$	T[I F][if] = I F->if ( A O P A ) C
\$ end end ) a ( printf begin ) b > a ( if	if ( A O P A ) C end \$	Matched:if
\$ end end ) a ( printf begin	Y end \$	T[Y][begin] = Y->begin X end
\$ end end ) a ( printf begin	begin X end \$	Matched:begin
\$ end end ) a ( printf	X end \$	T[X][printf] = X->printf ( A ) ;
\$ end end ) a ( printf	printf ( A ) ; end \$	Matched:printf
\$ end end ) a (	( A ) ; end \$	Matched:(
\$ end end ) a (	A ) ; end \$	T[A][a] = A->a G

```
$ end end ) a ( printf begin b > a ( if ; c , b , 3 2 1 d c b      G , A , A ; B end $      T[G][b] = G->b G      Matched:b
$ end end ) a ( printf begin b > a ( if ; c , b , 3 2 1 d c b      b G , A , A ; B end $
$ end end ) a ( printf begin b > a ( if ; c , b , 3 2 1 d c b      G , A , A ; B end $      T[G][c] = G->c G      Matched:c
$ end end ) a ( printf begin b > a ( if ; c , b , 3 2 1 d c      c G , A , A ; B end $
$ end end ) a ( printf begin b > a ( if ; c , b , 3 2 1 d      G , A , A ; B end $      T[G][d] = G->d G      Matched:d
$ end end ) a ( printf begin b > a ( if ; c , b , 3 2 1 d      d G , A , A ; B end $
$ end end ) a ( printf begin b > a ( if ; c , b , 3 2 1      G , A , A ; B end $      T[G][1] = G->1 G      Matched:1
$ end end ) a ( printf begin b > a ( if ; c , b , 3 2 1      1 G , A , A ; B end $
$ end end ) a ( printf begin b > a ( if ; c , b , 3 2      G , A , A ; B end $      T[G][2] = G->2 G      Matched:2
$ end end ) a ( printf begin b > a ( if ; c , b , 3 2      2 G , A , A ; B end $
$ end end ) a ( printf begin b > a ( if ; c , b , 3      G , A , A ; B end $      T[G][3] = G->3 G      Matched:3
$ end end ) a ( printf begin b > a ( if ; c , b , 3      3 G , A , A ; B end $
$ end end ) a ( printf begin b > a ( if ; c , b , 3      G , A , A ; B end $      T[G][,] = G->#      Matched:,
$ end end ) a ( printf begin b > a ( if ; c , b ,      A , A ; B end $
$ end end ) a ( printf begin b > a ( if ; c , b      A , A ; B end $      T[A][b] = A->b G      Matched:b
$ end end ) a ( printf begin b > a ( if ; c , b      b G , A ; B end $
$ end end ) a ( printf begin b > a ( if ; c ,      G , A ; B end $      T[G][_] = G->#      Matched:,
$ end end ) a ( printf begin b > a ( if ; c ,      A , A ; B end $
$ end end ) a ( printf begin b > a ( if ; c      A ; B end $      T[A][c] = A->c G      Matched:c
$ end end ) a ( printf begin b > a ( if ; c      c G ; B end $
$ end end ) a ( printf begin b > a ( if ;      G ; B end $      T[G][] = G->#      Matched;;
$ end end ) a ( printf begin b > a ( if ;      ; B end $
$ end end ) a ( printf begin b > a ( if ;      B end $      T[b][if] = B->if
$ end end ) a ( printf begin b > a ( if      IF end $      T[IF][if] = IF->if ( A OP A ) C      Matched:if
$ end end ) a ( printf begin      if ( A OP A ) C end $
$ end end ) a ( printf begin      Y end $      T[Y][begin] = Y->begin X end      Matched:begin
$ end end ) a ( printf      begin X end end $
$ end end ) a ( printf      X end end $      T[X][printf] = X->printf ( A ) ;      Matched:printf
$ end end ) a ( printf      printf ( A ) ; end end $      Matched:(
$ end end ) a (      ( A ) ; end end $
$ end end ) a (      A ) ; end end $      T[A][a] = A->a G      Matched:a
$ end end ) a      a G ) ; end end $      T[G]() = G->#      Matched:)
$ end end )      G ) ; end end $
$ end end )      ) ; end end $

Invalid String! Unmatched terminal symbols
Missing:      ;

$ end end $      ; end end $      Matched;;
$ end end $      end end $      Matched:end
$ end $      end $      Matched:end
$      $      Valid

valid String!
```

## Output for valid string without if statement:

Code -> int main ( ) begin int a , b , c ; end

```
Validate String => int main ( ) begin int a , b , c ; end

Buffer                                     Stack                                     Action
$ end ; c , b , a int begin ) ( main int   S $                               T[S][int] = S->int S'
$ end ; c , b , a int begin ) ( main int   int S' $                             Matched:int
$ end ; c , b , a int begin ) ( main       S' $                               T[S'] [main] = S' ->main ( ) Y
$ end ; c , b , a int begin ) ( main       main ( ) Y $                         Matched:main
$ end ; c , b , a int begin ) (           ( ) Y $                             Matched:(
$ end ; c , b , a int begin ) (           ) Y $                             Matched:)
$ end ; c , b , a int begin ) (           Y $                               T[Y][begin] = Y->begin X end
$ end ; c , b , a int begin ) (           begin X end $                         Matched:begin
$ end ; c , b , a int begin ) (           X end $                             T[X][int] = X->D B
$ end ; c , b , a int begin ) (           D B end $                         T[D][int] = D->int A , A , A ;
$ end ; c , b , a int begin ) (           int A , A , A ; B end $             Matched:int
$ end ; c , b , a int begin ) (           A , A , A ; B end $               T[A][a] = A->a G
$ end ; c , b , a int begin ) (           a G , A , A ; B end $             Matched:a
$ end ; c , b , a int begin ) (           G , A , A ; B end $               T[G][,] = G->#
$ end ; c , b , a int begin ) (           , A , A ; B end $                 Matched:,
$ end ; c , b , a int begin ) (           A , A ; B end $                   T[A][b] = A->b G
$ end ; c , b , a int begin ) (           b G , A ; B end $                 Matched:b
$ end ; c , b , a int begin ) (           G , A ; B end $                   T[G][,] = G->#
$ end ; c , b , a int begin ) (           , A ; B end $                     Matched:,
$ end ; c , b , a int begin ) (           A ; B end $                       T[A][c] = A->c G
$ end ; c , b , a int begin ) (           c G ; B end $                     Matched:c
$ end ; c , b , a int begin ) (           G ; B end $                       T[G][;] = G->#
$ end ; c , b , a int begin ) (           ; B end $                         Matched;;
$ end ; c , b , a int begin ) (           B end $                           T[B][end] = B->#
$ end ; c , b , a int begin ) (           end $                             Matched:end
$ end ; c , b , a int begin ) (           $                               Valid
Valid String!
```

## Output for valid string without statements:

Code -> int main ( ) begin end

```
Validate String => int main ( ) begin end

Buffer                                     Stack                                     Action
$ end begin ) ( main int                   S $                               T[S][int] = S->int S'
$ end begin ) ( main int                   int S' $                             Matched:int
$ end begin ) ( main                       S' $                               T[S'] [main] = S' ->main ( ) Y
$ end begin ) ( main                       main ( ) Y $                         Matched:main
$ end begin ) (                           ( ) Y $                             Matched:(
$ end begin ) (                           ) Y $                             Matched:)
$ end begin ) (                           Y $                               T[Y][begin] = Y->begin X end
$ end begin ) (                           begin X end $                         Matched:begin
$ end begin ) (                           X end $                             T[X][end] = X->D B
$ end begin ) (                           D B end $                         T[D][end] = D->#
$ end begin ) (                           B end $                           T[B][end] = B->#
$ end begin ) (                           end $                             Matched:end
$ end begin ) (                           $                               Valid
Valid String!
```

## 5. CONCLUSION

The project is mainly aimed at implementing Lexical and Syntax Analyzer which are the first two stages in Compiler Analysis Phase. We were successful in doing the same. Similarly, we can also develop remaining stages of the compiler such that our final outcome would be an assembly level language, which can be easily understood by the computer. The two phases of the compiler are separated to increase simplicity, improve the efficiency and increase the portability.

In lexical analyses, when we give a program statement as input, the input is parsed. If it is parsed successfully, the success message is displayed in the output. Otherwise, the error message is displayed.

The compiler technology is applied in various computer fields such as HLL implementation, program translation, and computer architecture.

In the future, we may experience complex compiler technologies that will be integrated with various computer applications.

## 6. REFERENCES

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