



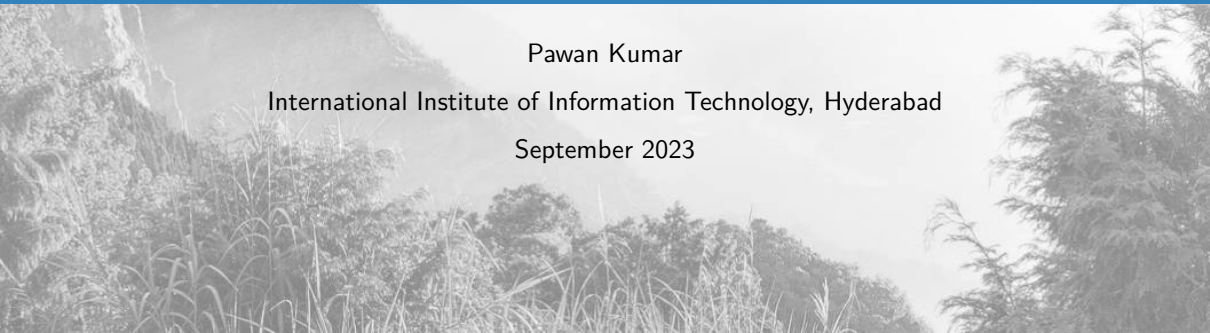
# Graph Theory

Monsoon 2023

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International Institute of Information Technology, Hyderabad

September 2023



## 0 Outline

| 1

### ① Graph Models

# 1 Outline

| 2

## ① Graph Models

Graphs and Digraphs

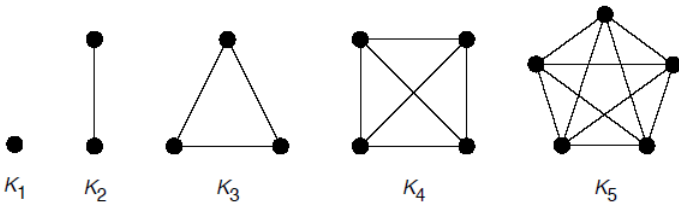
Applications of Graph Models

### Definition (Complete Graphs)

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### Example

What is a smallest possible simple graph that is not bipartite?

### Theorem

*A bipartite graph cannot have any self-loops.*

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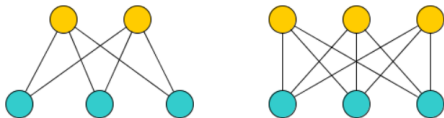
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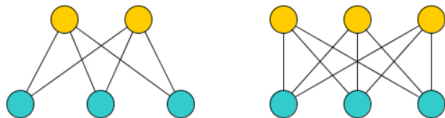


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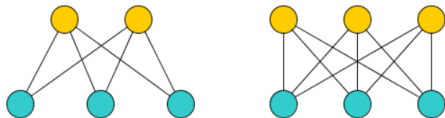


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Answer the following:

- 1 What are  $m, n$  in above graphs?
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- 3 What could be the applications of bipartite graphs?

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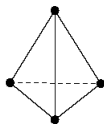
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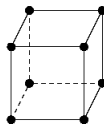
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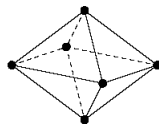
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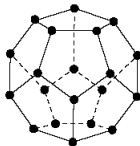
Tetrahedron



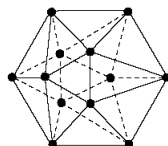
Cube



Octahedron



Dodecahedron

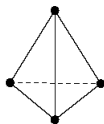


Icosahedron

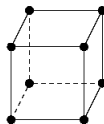
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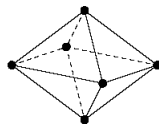
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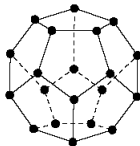
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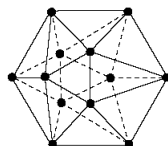
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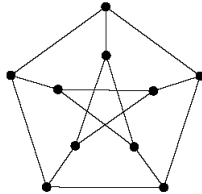
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### Example

The Petersen graph is the 3-regular graph represented by the line drawing in Figure below. Because it possesses a number of interesting graph-theoretic properties, the Petersen graph is frequently used both to illustrate established theorems and to test conjectures.

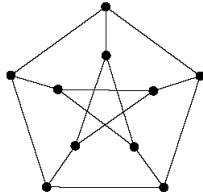
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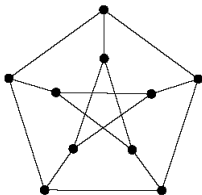


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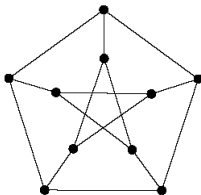


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- 3 Is peterson graph bipartite? If yes, then show the partition of vertices.

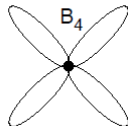
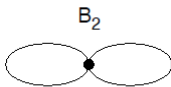


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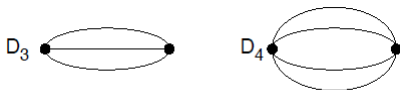
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A **path graph**  $P$  is a simple graph with  $|V_P| = |E_P| + 1$  that can be drawn so that all of its vertices and edges lie on a single straight line. A path graph with  $n$  vertices and  $n - 1$  edges is denoted  $P_n$ .

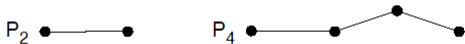
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A **cycle graph** is a single vertex with a self-loop or a simple graph  $C$  with  $V_C = E_C$  that can be drawn so that all of its vertices and edges lie on a single circle. An  $n$ -vertex cycle graph is denoted  $C_n$ .

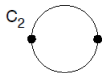
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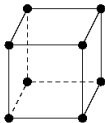
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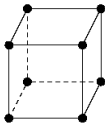
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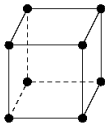
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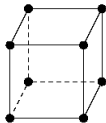


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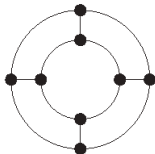
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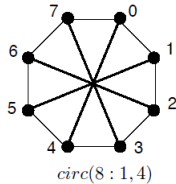
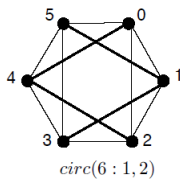
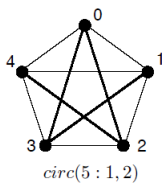
### Definition (circulant)

To the group of integers  $Z_n = \{0, 1, \dots, n-1\}$  under addition modulo  $n$  and a set  $S \subset \{1, \dots, n-1\}$ , we associate the **circulant graph**  $\text{circ}(n : S)$  whose vertex set is  $Z_n$ , such that two vertices  $i$  and  $j$  are adjacent if and only if there is a number  $s \in S$  such that  $i + s = j \bmod n$  or  $j + s = i \bmod n$ . In this regard, the elements of the set  $S$  are called **connections**.



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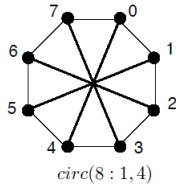
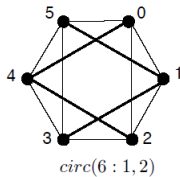
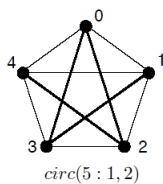


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Notice that circulant graphs are simple graphs.

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Notice that circulant graphs are simple graphs. Circulant graphs are a special case of **Cayley graphs**, which are themselves derived from a special case of **voltage graphs**.

## Definition (Intersection and Interval Graphs)

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A simple graph  $G$  with vertex-set  $V_G = \{v_1, v_2, \dots, v_n\}$  is an **intersection graph** if there exists a family of sets  $F = \{S_1, S_2, \dots, S_n\}$  such that vertex  $v_i$  is adjacent to  $v_j$  if and only if  $i \neq j$  and  $S_i \cap S_j \neq \emptyset$ .

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A simple graph is an **interval graph** if it is an intersection graph corresponding to a family of intervals on the real line.

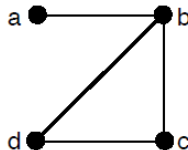
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$$a \leftrightarrow (1, 3) \quad b \leftrightarrow (2, 6) \quad c \leftrightarrow (5, 8) \quad d \leftrightarrow (4, 7)$$



## Example (Interval Graph)

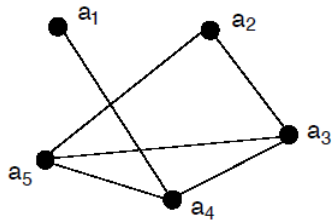


### Example (Interval Graph)

(Archeology) Suppose that a collection of artifacts was found at the site of a town known to have existed from 1000 BC to AD 1000. In the graph shown, the vertices correspond to artifact types, and two vertices are adjacent if some grave contains artifacts of both types. It is reasonable to assume that artifacts found in the same grave have overlapping time intervals during which they were in use. If the graph is an interval graph, then there is an assignment of subintervals of the interval  $(-1000; 1000)$  (by suitable scaling, if necessary) that is consistent with the archeological find.

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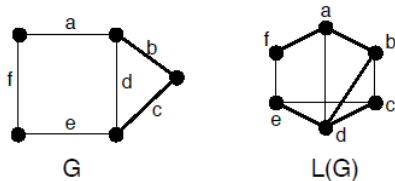
The **line graph**  $L(G)$  of a graph  $G$  has a vertex for each edge of  $G$ , and two vertices in  $L(G)$  are adjacent if and only if the corresponding edges in  $G$  have a vertex in common.

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Answer the following:

- 1 Verify that  $L(G)$  is the intersection graph.

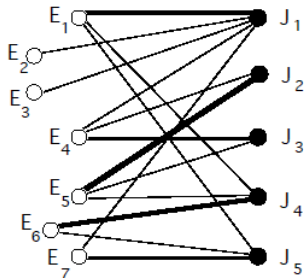
## Example (Personnel-Assignment Problem)

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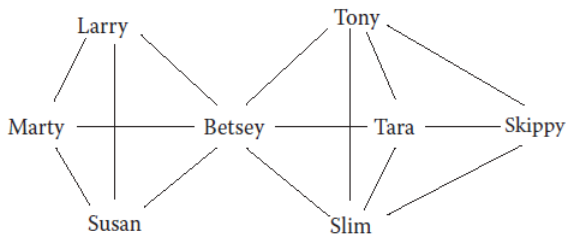
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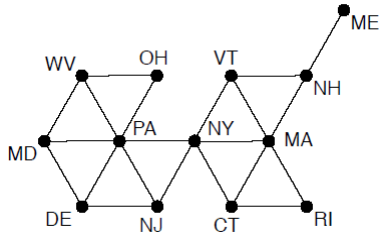
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## Example (Geometric Polyhedra)

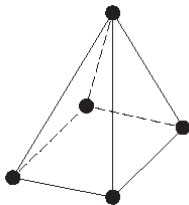
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The vertex and edge configuration of any polyhedron in 3-space forms a simple graph, which topologists call its 1-skeleton. The 1-skeletons of the platonic solids, appearing in the previous section, are regular graphs. Figure below shows a polyhedron whose 1-skeleton is not regular.



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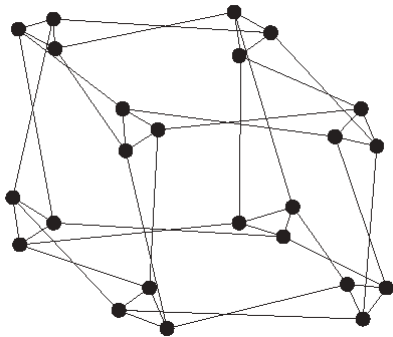
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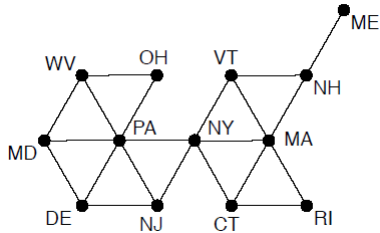
## Example (Roadways between States)

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## Example (Chemical Molecules)

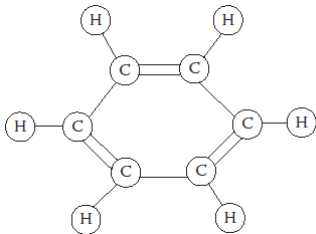


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The benzene molecule shown in Figure below has double bonds for some pairs of its atoms, so it is modeled by a non-simple graph. Since each carbon atom has valence 4, corresponding to four electrons in its outer shell, it is represented by a vertex of degree 4; and since each hydrogen atom has one electron in its only shell, it is represented by a vertex of degree 1.

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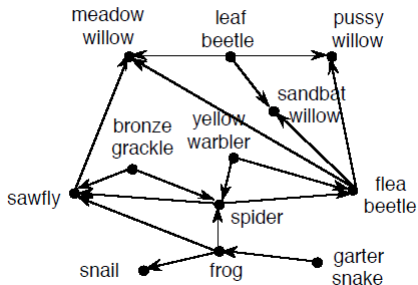
## Example (Ecosystems)

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The feeding relationships among the plant and animal species of an ecosystem are called a food web and may be modeled by a simple digraph. The food web for a Canadian willow forest is illustrated in Figure below. Each species in the system is represented by a vertex, and a directed edge from vertex  $u$  to vertex  $v$  means that the species corresponding to  $u$  feeds on the species corresponding to  $v$ .

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## Example (Activity-Scheduling Networks)

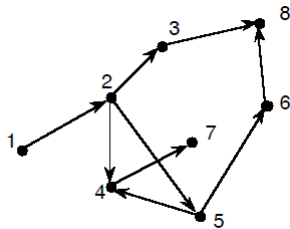
### Example (Activity-Scheduling Networks)

In large projects, often there are some tasks that cannot start until certain others are completed. Figure below shows a digraph model of the precedence relationships among some tasks for building a house. Vertices correspond to tasks. An arc from vertex  $u$  to vertex  $v$  means that task  $v$  cannot start until task  $u$  is completed. To simplify the drawing, arcs that are implied by transitivity are not drawn. This digraph is the cover diagram of a partial ordering of the tasks.

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	Activities
1	Foundation
2	Walls and Ceilings
3	Roof
4	Electrical wiring
5	Windows
6	Siding
7	Paint interior
8	Paint exterior





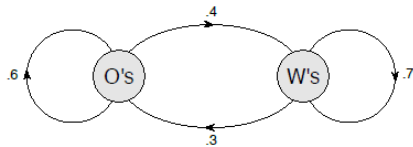
## Example (Markov Diagrams)

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Suppose that the inhabitants of some remote area purchase only two brands of breakfast cereal, O's and W's. The consumption patterns of the two brands are encapsulated by the transition matrix shown in Figure below. For instance, if someone just bought O's, there is a 0.4 chance that the person's next purchase will be W's and a 0.6 chance it will be O's. In a Markov process, the transition probability of going from one state to another depends only on the current state. Here, states "O" and "W" correspond to whether the most recent purchase was O's or W's, respectively. The digraph model for this Markov process, called a Markov diagram, is shown in Figure below. Each arc is labeled with the transition probability of moving from the state at the tail vertex to the state at the head. Thus, the probabilities on the outgoing edges from each vertex must sum to 1. This Markov diagram is an example of a weighted graph.

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	O's	W's
O's	.6	.4
W's	.3	.7