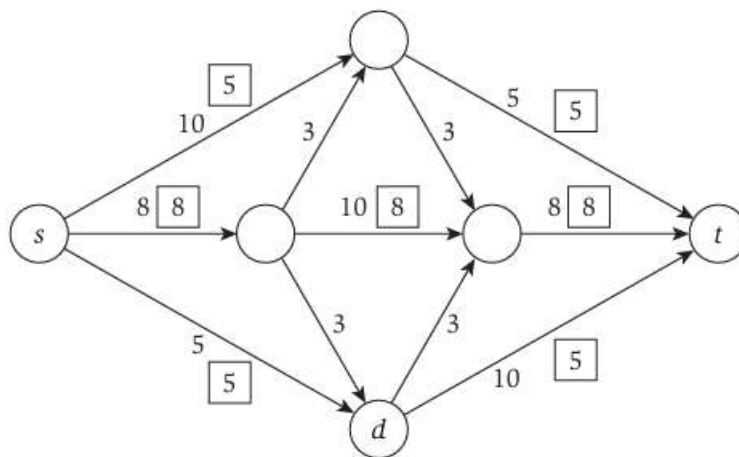


## NETWORK FLOW EXAMPLES SOLVED

Book: Algorithm-Design-by-Jon-Kleinberg-Eva-Tardos.pdf

Chapter 7: Network Flow Exercise Question 2 [PDF pg. 441]

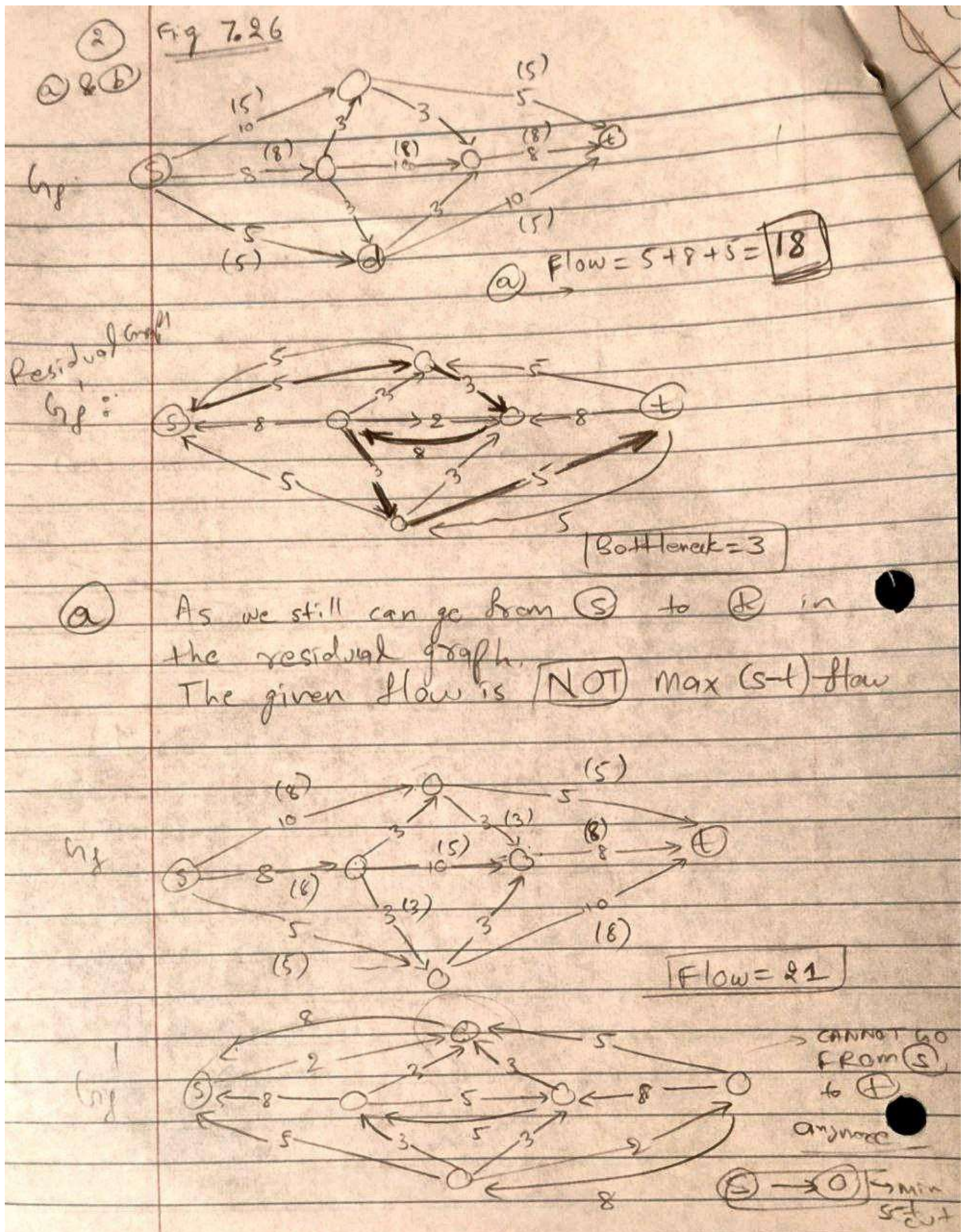
2. Figure 7.26 shows a flow network on which an  $s$ - $t$  flow has been computed. The capacity of each edge appears as a label next to the edge, and the numbers in boxes give the amount of flow sent on each edge. (Edges without boxed numbers—specifically, the four edges of capacity 3—have no flow being sent on them.)
- (a) What is the value of this flow? Is this a maximum  $(s,t)$  flow in this graph?
- (b) Find a minimum  $s$ - $t$  cut in the flow network pictured in Figure 7.26, and also say what its capacity is.



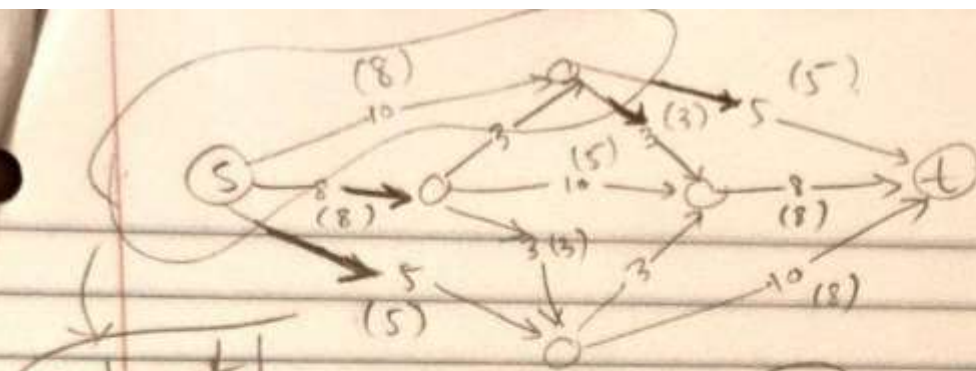
**Figure 7.26** What is the value of the depicted flow? Is it a maximum flow? What is the minimum cut?

There's a misprint, these nodes are  $\{a, b, c, d\}$

Solution:



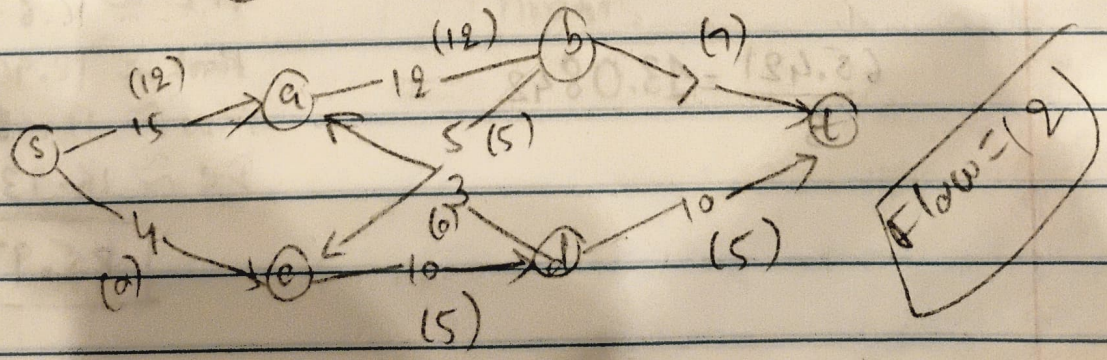
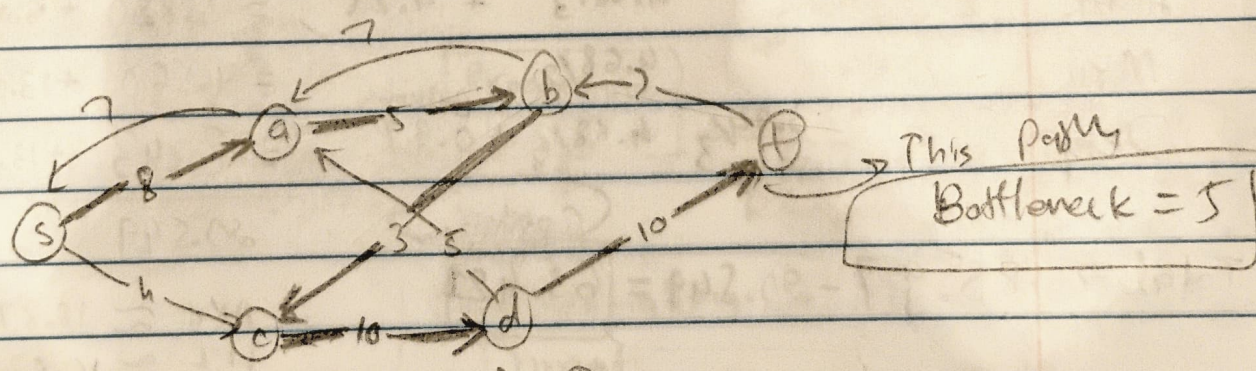
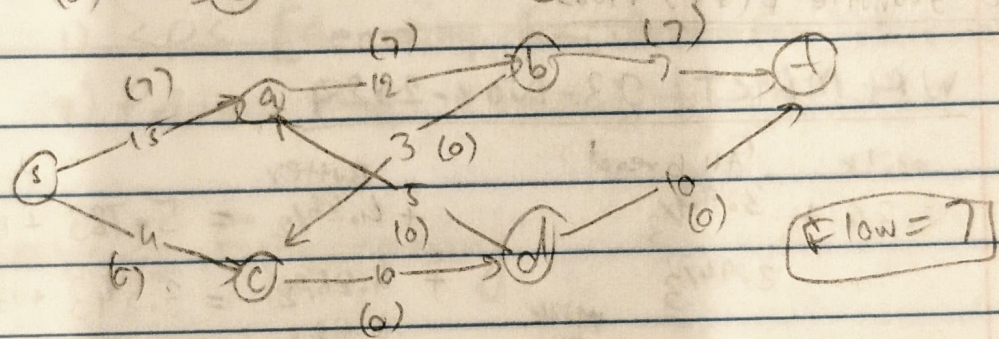
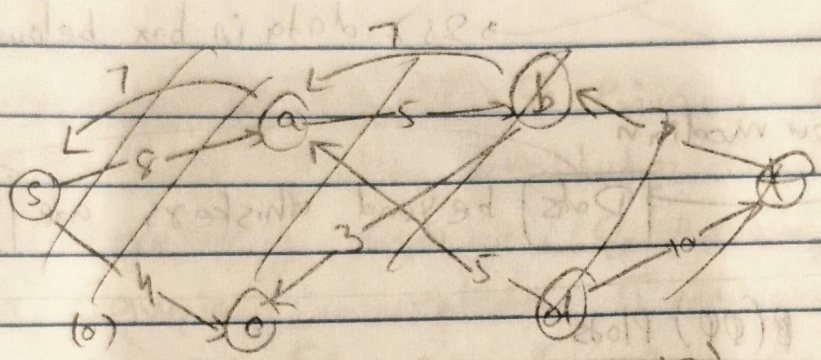
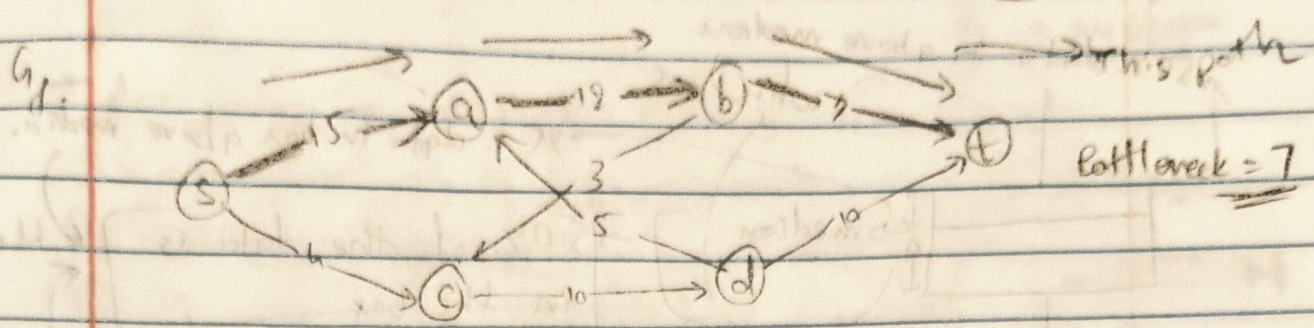
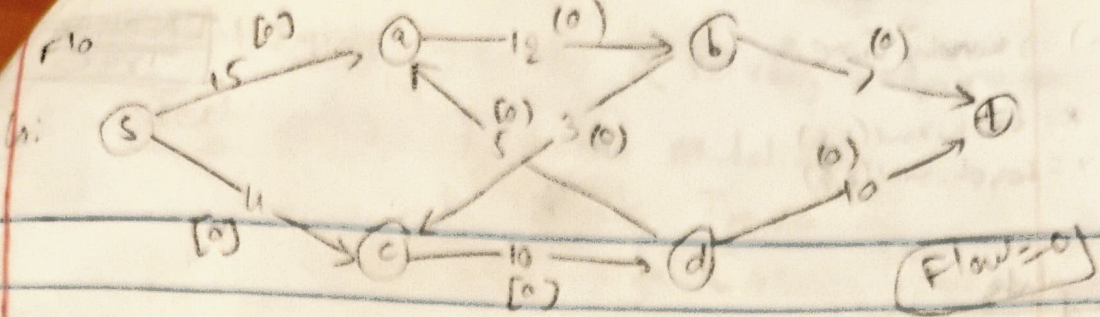




Flow out of min (s-t) cut =  $5 + 3 + 8 + 5$   
 $\boxed{= 21}$

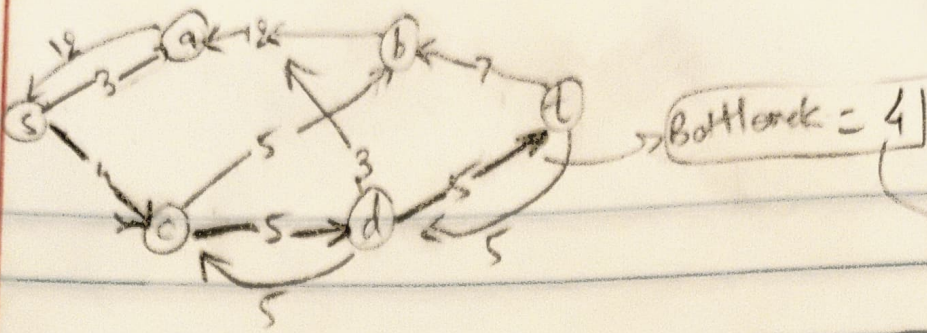
- 2 a 18, No it is not max s-t flow
- b 21  $\leftarrow$  max s-t flow capacity



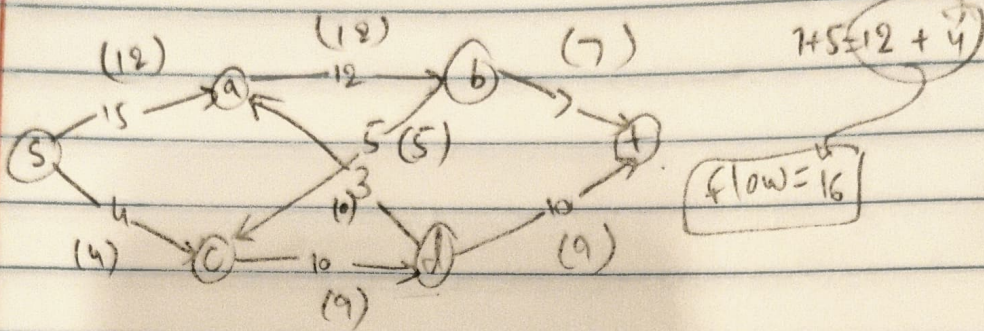




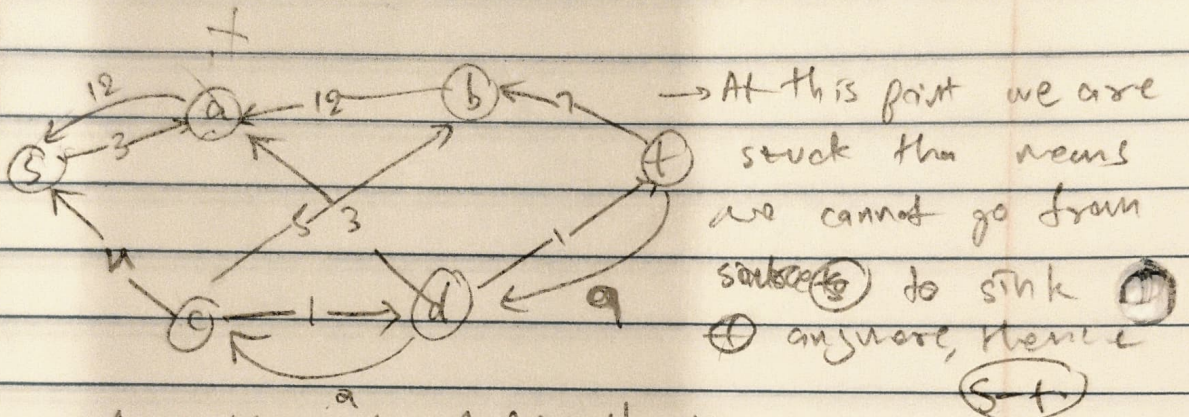
g:



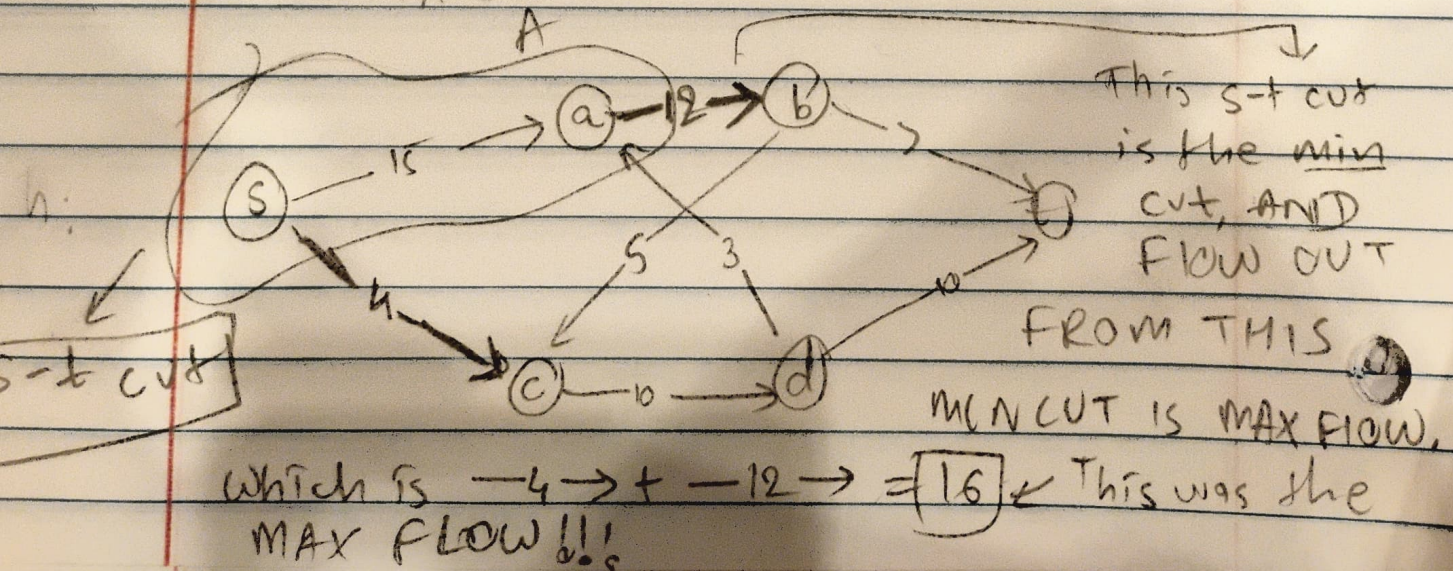
h:



h<sub>g</sub>:



We found the optimal PATH. Also, the minimum cut will include <sup>all</sup> nodes ~~from~~ reachable from source s ~~to all nodes~~. In h<sub>g</sub> i.e. residual graph h<sub>g</sub>. It is clear that we can only reach a from s, hence the s-t cut will look like





amount equal to capacity  
and remove flow part.

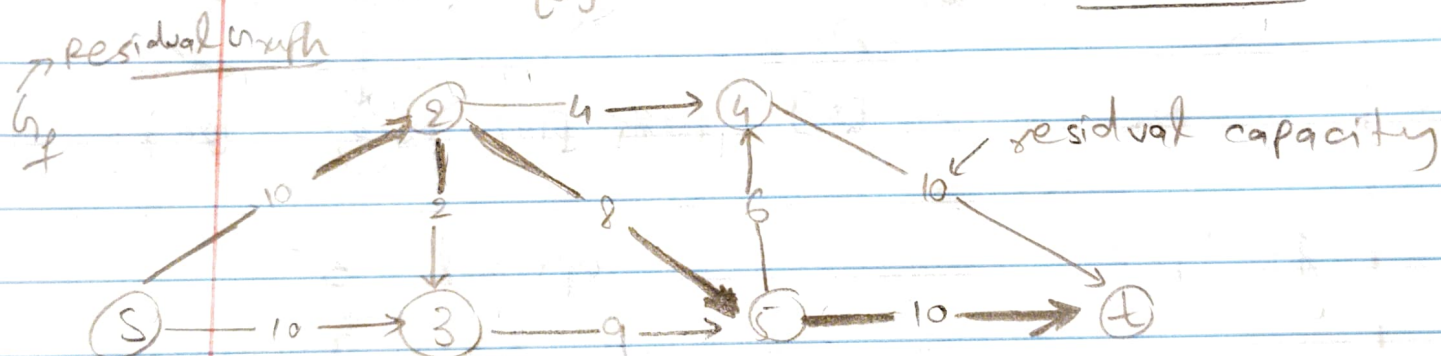
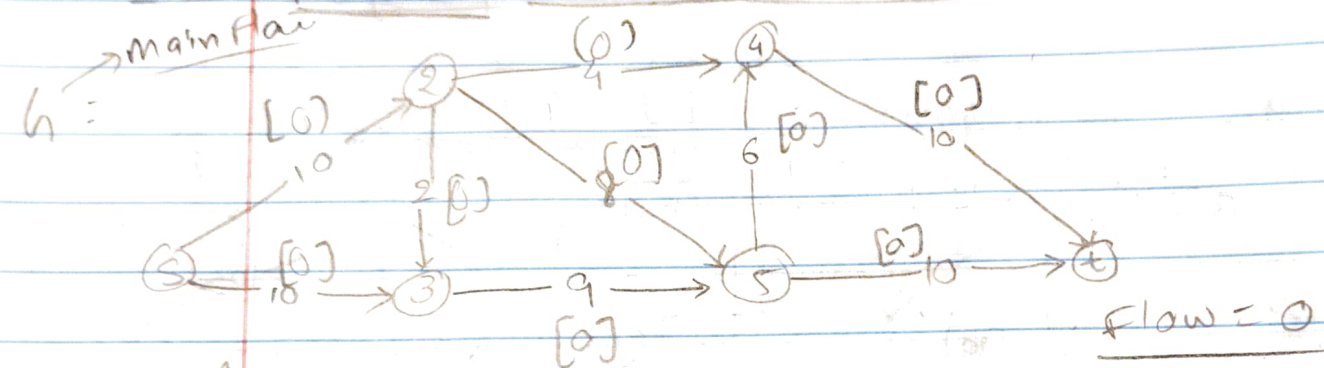
## Residual graph

→ All such edges with positive residual capacity will make the Residual graph

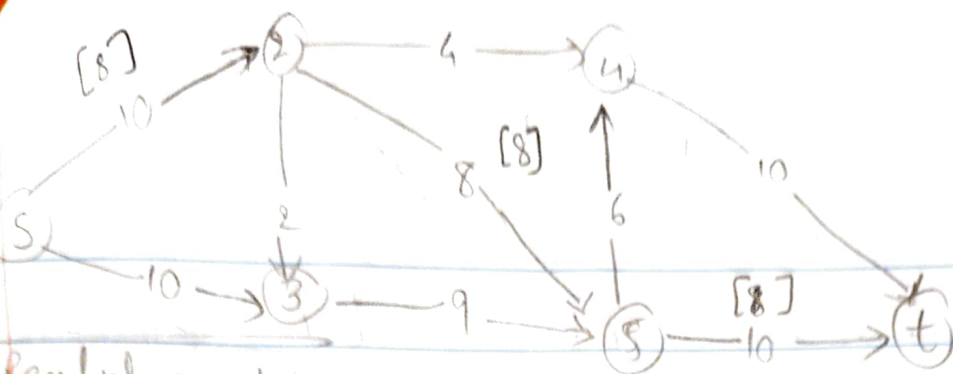
→ Now we can run an algorithm on this residual graph, ~~into~~ to find MAX FLOW.



## FORD-FULKERSON ALGORITHM

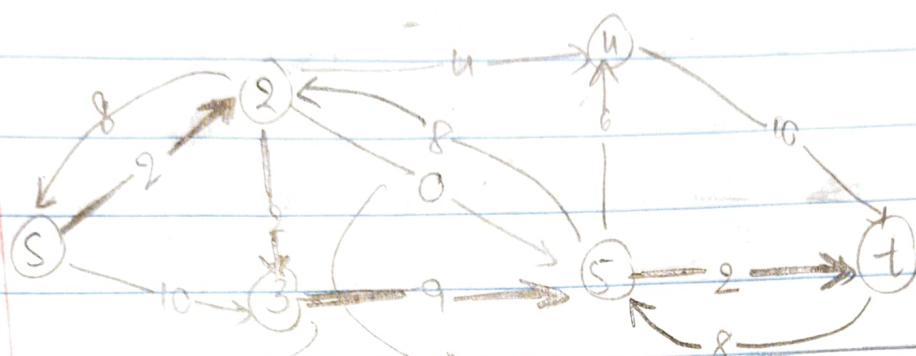


choose this path, (you can choose any!)  
The bottleneck = edge with lowest capacity as that is the MAX Flow we can send in this path = 8 in this path, So take this and pass flow of 8 through this path in the main flow.



Residual graph

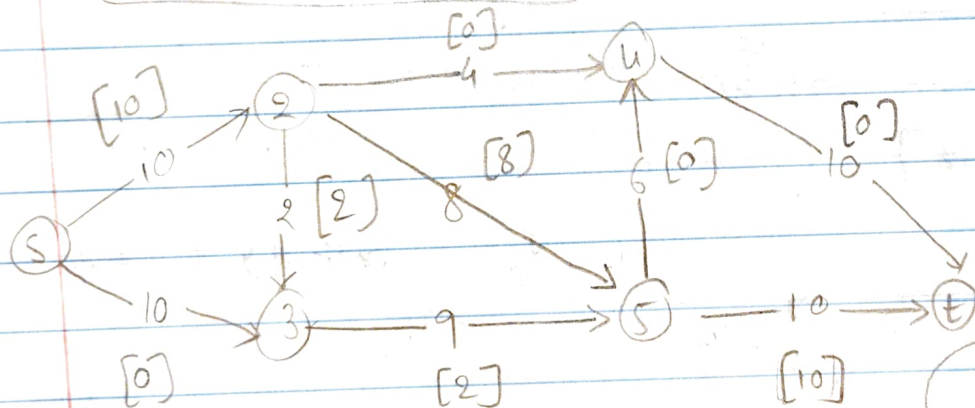
Flow = 8



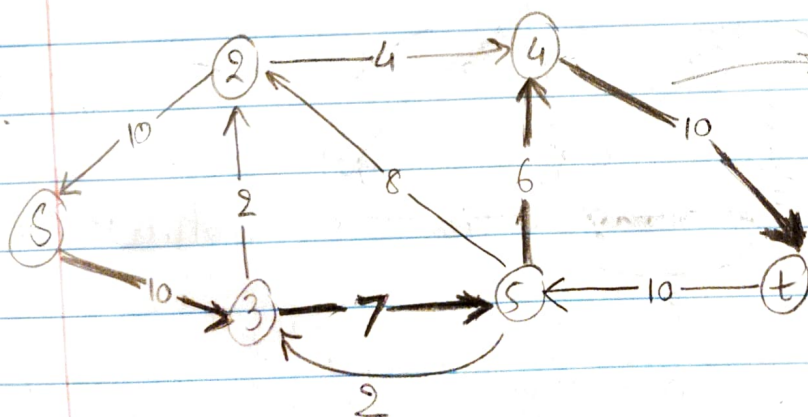
Afterwards, let's choose this path

No need to draw edge if cap = 0, that means it DOES NOT EXIST

Bottleneck = 2 → so take 2 flow on this path



Flow = 10

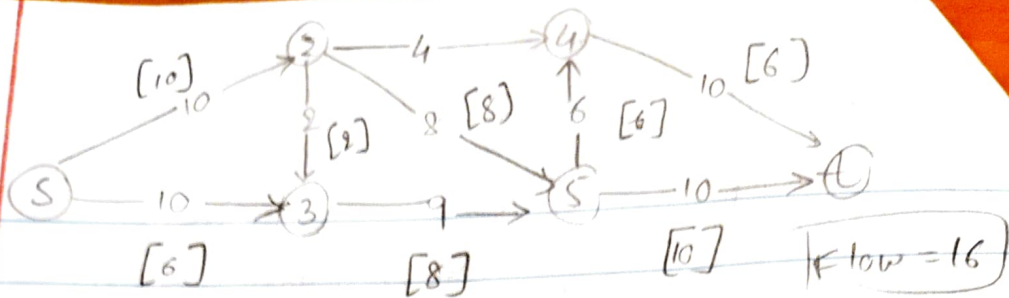


Next path

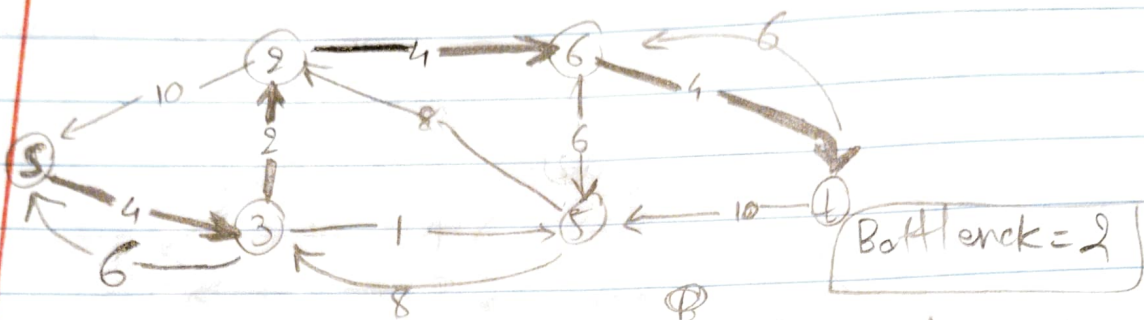
Bottleneck = 6

Add flow of 6 on this path in original 5.



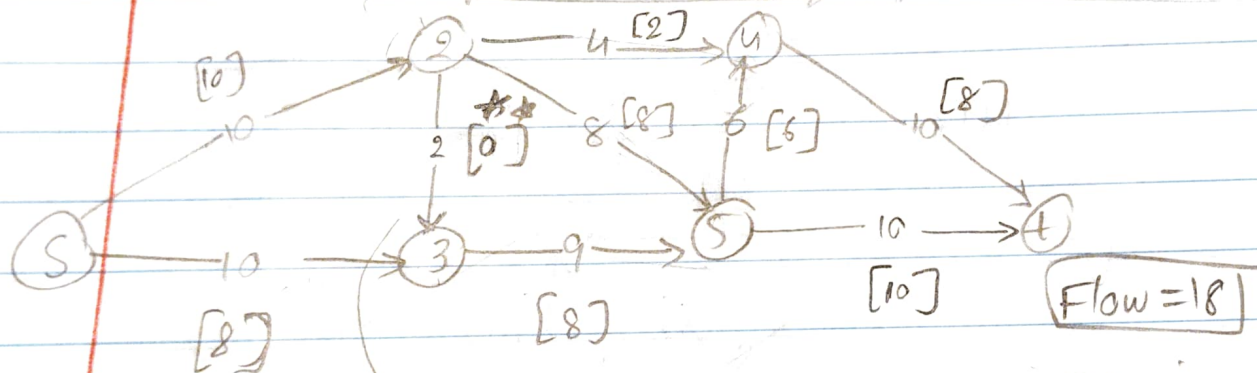


Update  
Gf:

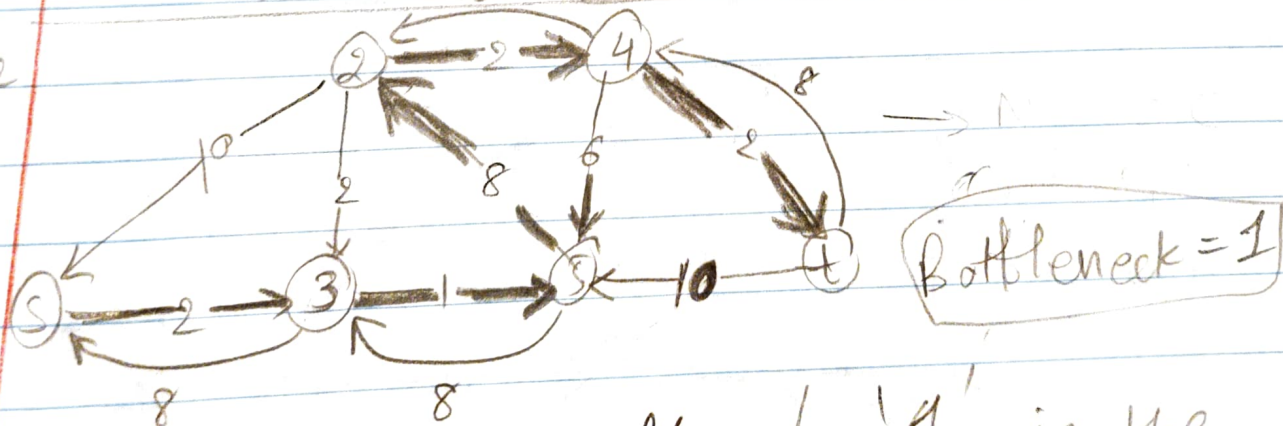


⊕ Add flow of '2' on this

path in the original 'G'

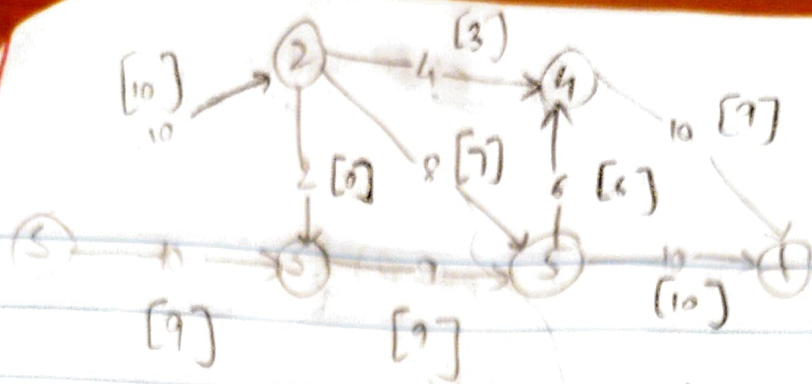


→ This edge is actually reverse in the 'Gf', hence it is telling us to decrease the flow by 'bottleneck amount'



→ Increase flow by '1' in the main flow graph 'G'



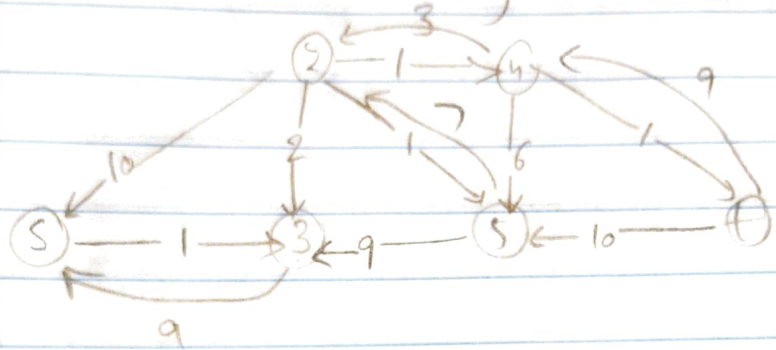


Flow = 19

→ This is reverse edge in  $G_f$ , hence decrement

flow by bottleneck = 1, Flow = 19 - 1 = 18

Update  $G_f$



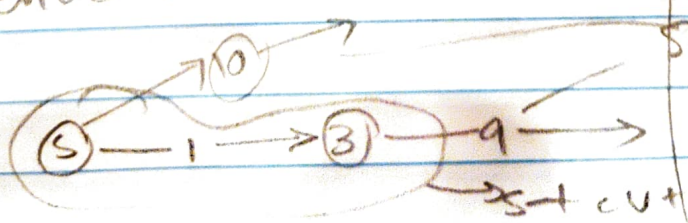
→ NOW WE CANNOT GO FROM SOURCE (s) to SINK (t)

→ When this occurs, it means we found our optimal flow OR MAX FLOW OF THE SYSTEM Hence max flow of this example = 19

s-t cut

→ Cut: All the nodes at this point of time that can be reached by sink (s), are the "part" of (s-t) cut, which is NOTHING BUT THE MINIMUM CUT OR MAXIMUM FLOW → (minimum cut required to restrain maximum flow of the system)

In our case, only (3) is reachable from (s) Hence the s-t cut is



See the flow from s-t cut = 10 + 9 = 19 i.e. the MAX FLOW