

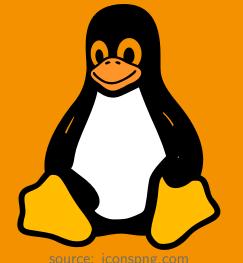
Stort: 8:01

Prof. Dr. Florian Künzner

Technical University of Applied Sciences Rosenheim, Computer Science

Bitte: Vornanc + Nadyhane

=> sellständig unserennen



OS 9 – Communication 2

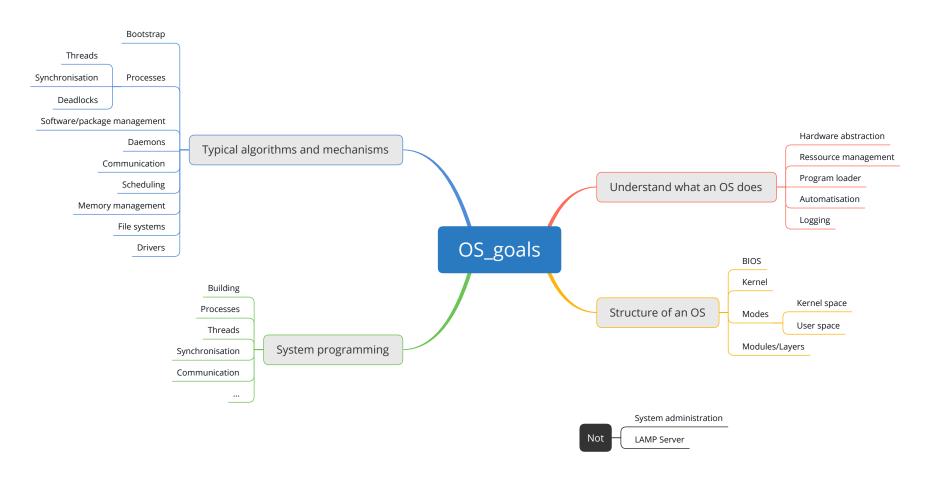
The lecture is based on the work and the documents of Prof. Dr. Ludwig Frank

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Goal



Goal



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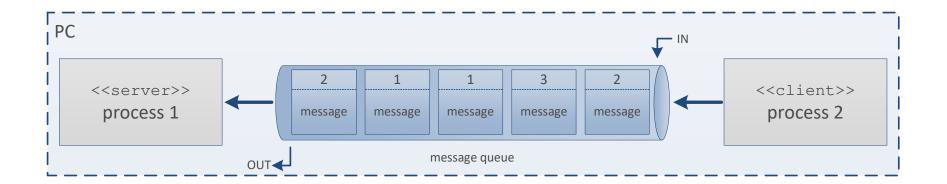
OS::Communication

- Message queue
- Shared memory
- Process communication summary

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Message queue



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Message queue

Message queue concept

- Queue to store messages
- Inter-process communication (IPC) between processes on one PC.
- Messages have priority/type.
- Internal stored as a linked list.
- Send into queue does not require an active receiver
- Read from queue does not require an active sender
- Max queue size (default: 16 KiB on Linux).
- Max message size (default: 8 KiB on Linux).

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Message queue

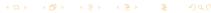
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Summary

Message queue

Message structure

```
1 struct message {
2    long priority;    //priority or type
3    char message[64]; //buffer for message bytes
4 };
```

Message priority/type

- The lower the number, the higher the priority
- The priority can be interpreted as a type (each type has its own number)

Message queue usage scenarios (PRIO_FETCH_FLAG)

- Read message after message (FIFO principle) (priority == 0)
- Read only message with specific priority/type (priority == N)
- Read the messages with the highest priority/type first, up to a certain number (priority == -N)

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 A highest prio

 31 lover prio

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Summary

Message queue: Pseudo C code

```
struct message { //structure for messages
  long priority;
  char message[64];
};
```

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Summary

Message queue: Pseudo C code

```
struct message { //structure for messages
     long priority;
     char message[64];
   };
   void receiver() {
     //create message queue
     msgget(...);
9
10
     //receive message
     //blocks if message queue is empty
12
13
     msgrcv(...);
14
15
     //... work with message
16
17
     //remove message queue
     msgctl(...);
18
19
```

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Summary

Message queue: Pseudo C code

```
struct message { //structure for messages
      long priority;
      char message [64];
   };
   void receiver() {
                                                        void sender() {
                                                           //open existing message queue
     //create message queue
                                                     21
     msgget(...);
                                                     22
                                                           msgget(...);
                                                     23
                                                     24
                                                           //prepare message
10
                                                     25
11
     //receive message
                                                     26
                                                           //send message
     //blocks if message queue is empty
msgrcv(...);
                                                     27
                                                           //blocks if message queue is full
12
13
                                                     28
                                                           msgsnd(...);
                                                     29
14
15
     //... work with message
                                                     30
                                                           //close not needed
16
                                                     31
17
     //remove message queue
                                                     32
     msgctl(...);
                                                     33
18
                                                     34
19
```

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Message queue: Linux commands

Command Description ipcs Show information on IPC facilities

ipcs -q Shows active message queues in the system

ipcmk Make various IPC resources

ipcmk -Q Create a message queue

ipcrm Remove certain IPC resources

ipcrm -q 1 Remove message queue with id 1

ipcrm -Q 2 Remove message queue with key 2

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Summary

Message queue

C example

Slide 9 of 20

^{*}Please find the source file(s) in the repository.

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Questions?

All right? \Rightarrow

(4)

Question? \Rightarrow

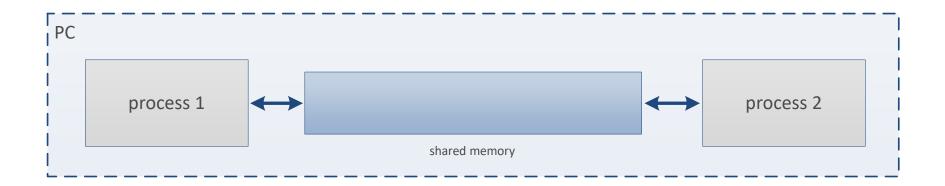
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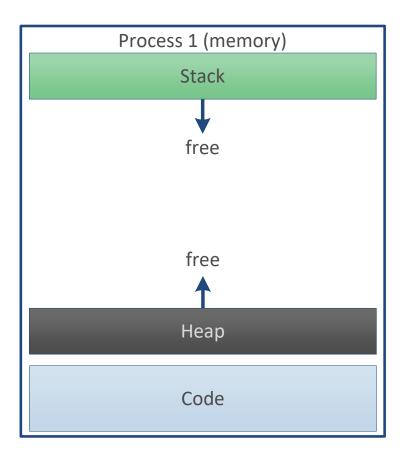


Summary



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Shared memory



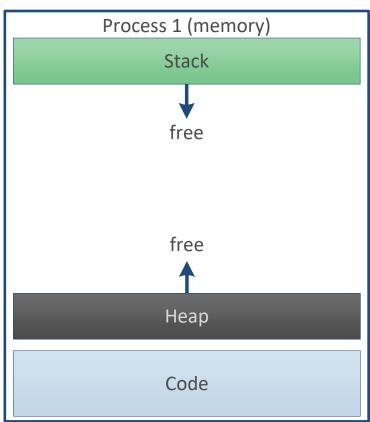


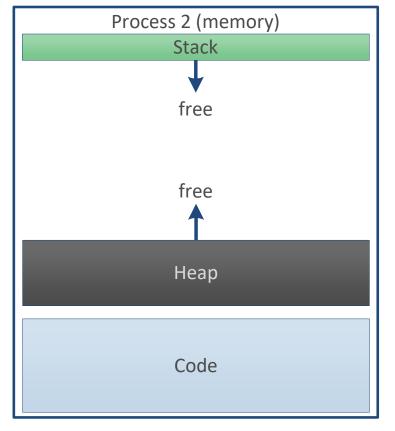
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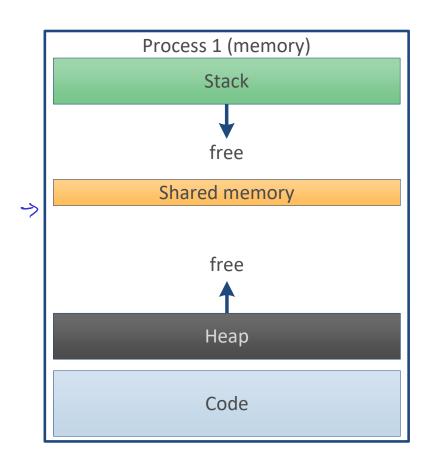


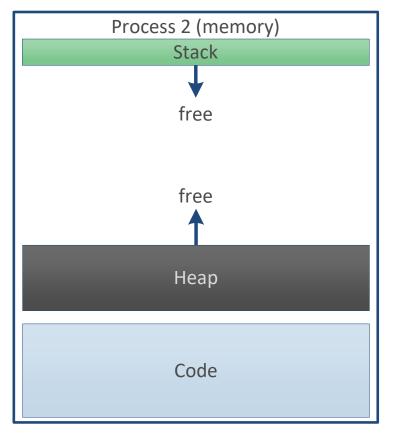


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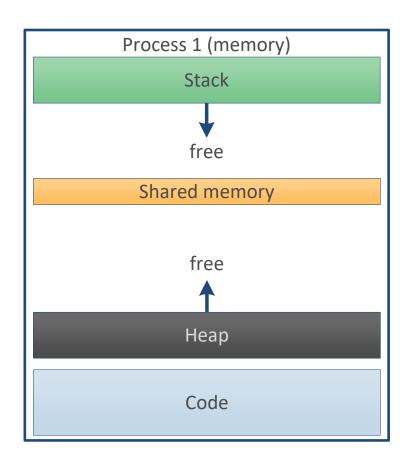


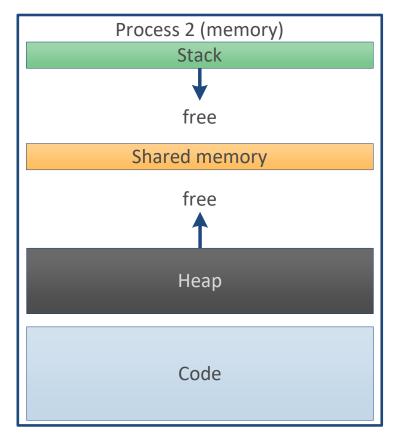


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Summary

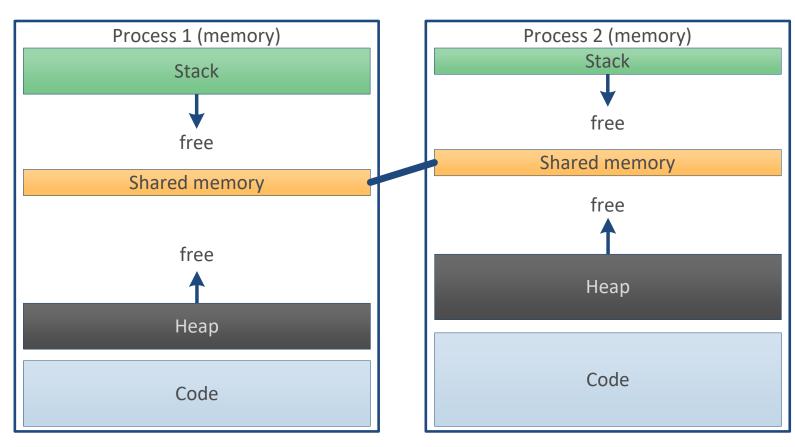




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Summary



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Summary

Shared memory

Shared memory concept

- Shared memory area between processes
- Inter-process communication (IPC) between processes on one PC.
- It is a plain memory area with a certain size
- Access needs to be synchronised (e.g. semaphore)
- Access is very fast (comparable with own memory access)

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Summary

Shared memory: Pseudo C code

```
seminit(READY TO WRITE, 1); //declare and initialise semaphore
seminit(READY TO READ, 0); //declare and initialise semaphore
```

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Summary

Shared memory: Pseudo C code

```
seminit(READY TO WRITE, 1); //declare and initialise semaphore
   seminit(READY TO_READ, 0); //declare and initialise semaphore
   void receiver() {
     //create shared memory
     shmget(...);
     //attach the shared memory
     shared mem address = shmat(...);
9
10
11
     //copy data from shared memory
     P(READY TO READ);
13
     copy(data, shared mem address); //data = sm
14
     V(READY TO WRITE);
15
16
17
     //... work with data
18
     work with(data);
19
     //detach shared memory
20
     shmdt(...);
     //remove shared memory
```

shmctl(...);

23

24 }

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Summary

Shared memory: Pseudo C code

```
seminit(READY TO WRITE, 1); //declare and initialise semaphore
   seminit(READY TO READ, 0); //declare and initialise semaphore
   void receiver() {
                                                  25 void sender() {
                                                        //get existing shared memory
     //create shared memory
                                                  26
                                                        shmget(...);
     shmget(...);
                                                  27
     //attach the shared memory
                                                  28
                                                        //attach the shared memory
     shared mem address = shmat(...);
                                                        shared mem address = shmat(...);
                                                  29
                                                  30
9
                                                  31
                                                        //... prepare data
10
                                                  32
                                                        data = prepare data();
                                                  33
11
12
     //copy data from shared memory
                                                  34
                                                        //copy data into shared memory
     P(READY TO READ);
                                                  35
                                                      P(READY TO WRITE);
13
     copy(data, shared_mem_address); //data =
                                                        copy(shared mem address, data); //sm = data
14
                                                sm 36
     V(READY TO WRITE);
                                                       - V(READY TO READ);
15
                                                  38
16
17
     //... work with data
                                                  39
18
     work with(data);
                                                  40
19
                                                  41
20
     //detach shared memory
                                                  42
                                                        //detach shared memory
     shmdt(...);
                                                        shmdt(...);
                                                  43
```

shmctl(...);

23

24 }

//remove shared memory

44

45

46

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Summary

Shared memory: Linux commands

Command Description

ipcs Show information on IPC facilities

ipcs -m Shows active shared memory in the system

ipcmk Make various IPC resources

ipcmk -M 8 Create a shared memory with 8 bytes

ipcrm Remove certain IPC resources

ipcrm -m 1 Remove shared memory with id 1

ipcrm -M 2 Remove shared memory with key 2



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Summary

C example

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Summary

Questions?

All right? \Rightarrow

V

Question? \Rightarrow

*

and use chat

speak after I ask you to



Summary

Process communication summary

Mechanism	data	store	access contr.	remote	bidirect.	fast	prio.	sync req.
signal								
unix socket								
network socket								
message queue								
shared memory								



Summary

Process communication summary

Mechanism	data	store	access contr.	remote	bidirect.	fast	prio.	sync req.
signal						X		
unix socket								
network socket								
message queue								
shared memory								

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Summary

Process communication summary

Mechanism	data	store	access contr.	remote	bidirect.	fast	prio.	sync req.
signal						X		
unix socket	X		X		Χ	Χ		
network socket								
message queue								
shared memory								



Summary

Process communication summary

Mechanism	data	store	access contr.	remote	bidirect.	fast	prio.	sync req.
signal						Χ		
unix socket	Χ		X		Χ	Χ		
network socket	Χ			X	X			
message queue								
shared memory								



Summary

Process communication summary

Mechanism	data	store	access contr.	remote	bidirect.	fast	prio.	sync req.
signal						Χ		
unix socket	Χ		X		Χ	Χ		
network socket	Χ			X	X			
message queue								
shared memory								



Summary

Process communication summary

Mechanism	data	store	access contr.	remote	bidirect.	fast	prio.	sync req.
signal						Χ		
unix socket	X		X		Χ	X		
network socket	Χ			X	Χ			
message queue	Χ	X	X			Χ	Χ	
shared memory								



Summary

Process communication summary

Mechanism	data	store	access contr.	remote	bidirect.	fast	prio.	sync req.
signal						Χ		
unix socket	X		X		Χ	X		
network socket	Χ			X	X			
message queue	Χ	X	X			Χ	X	
shared memory	XX	XX	X		X	XX		Χ

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Question? \Rightarrow

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Summary

Summary and outlook

Summary

- Message queue
- Shared memory
- Process communication summary



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Summary and outlook

Summary

- Message queue
- Shared memory
- Process communication summary

Outlook

- Deadlocks
- Deadlock analysis
- Deadlock prevention