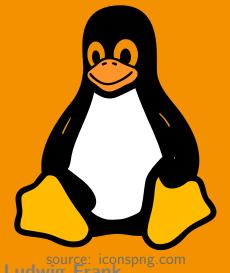


Prof. Florian Künzner

Technical University of Applied Sciences Rosenheim, Computer Science

OS 4 – System boot; OS architecture



The lecture is based on the work and the documents of Prof. Dr. Ludwig Frank

BIOS/UEFI

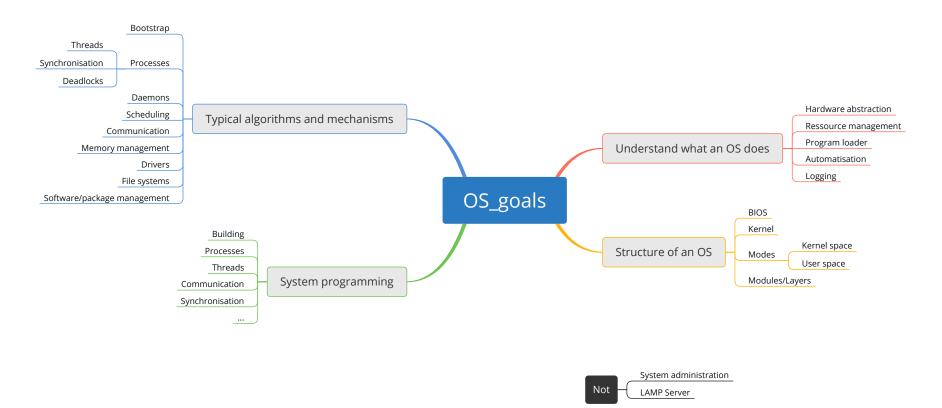
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Goal

Goal



Technische Hochschule Rosenheim Technical University of Applied Sciences

Goal

Goal

OS::System boot; OS architecture

- Boot procedure
- BIOS/UEFI
- MBR/GPT
- Boot loader: Grub2
- init/systemd
- OS architecture



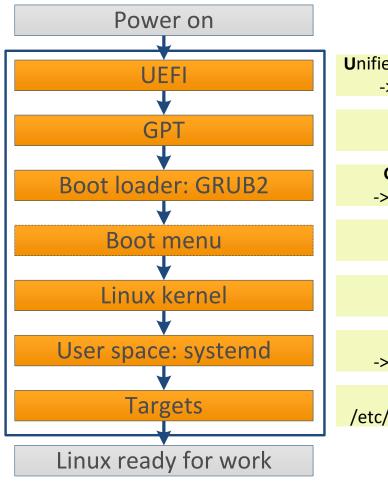
Boot

How does a PC boot into the operating system?

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Boot procedure: UEFI (typically)



Unified Extensible Firmware Interface
-> executes EFI program files

GUID **P**artition **T**able

-> describes disk partitions

Grand Unified Bootloader 2

-> executes boot menu/kernel

Boot menu

-> choose kernel for boot

Kernel

-> executes /bin/systemd

Systemd

-> executes runlevel programs

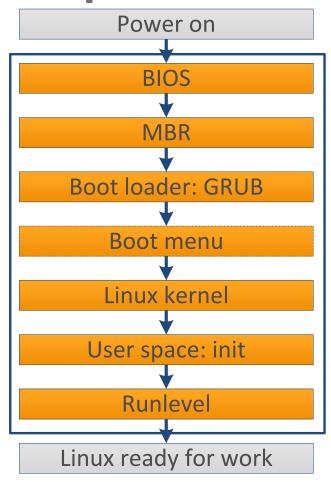
Starts **Targets**

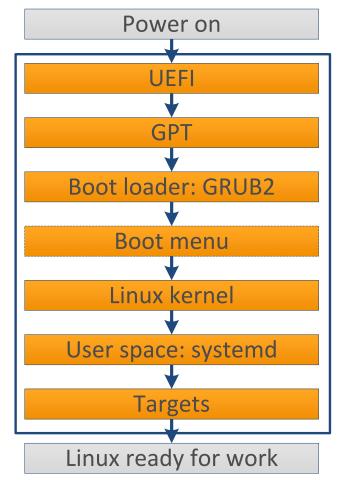
/etc/systemd/system/default.target

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Boot procedure: BIOS vs. UEFI





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UEFI - Unified ext. firmware interface

- Since 2010 in consumer products
- Graphical user interface (with mouse)
- Fast boot: cache + hibernation (Windows only)
- Secure boot:
 - Protection against malware
 - Prevents against execution of unsigned code
- Network boot
- Modular interface for applications and devices (EFI drivers)
- Supported modes:
 - UEFI mode: Requires an EFI partition on boot device
 - BIOS mode: Old way of booting

More details: https://www.marksei.com/bios-uefi-explained

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MBR - Master Boot Record

MBR (512 bytes)

Master boot code (446 bytes)

Partition table (64 bytes)

1st 2nd 3rd 4th entry entry

Magic Number: 0x55, 0xAA

- Supports up to 2 TiB disks
- Supports up to 2 TiB partitions
- No safety (no checksum)
- Supports 4 primary partitions
- Supports one extended partition (not bootable)

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GPT - GUID Partition Table

GPT

Protective MBR (512 bytes)

Primary GUID partition table header (~512 bytes)

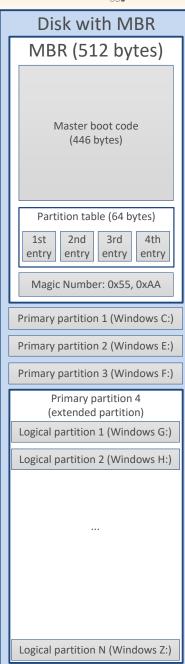
Primary GUID partition entry (each 128 bytes) array

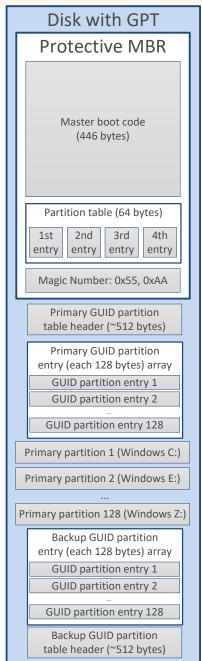
GUID partition entry 1

GUID partition entry 2

GUID partition entry 128

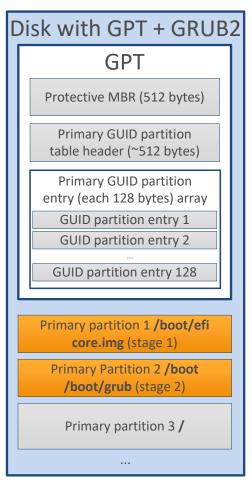
- Supports up to 18 EiB disks
- Supports up to 18 EiB partitions
- Safety (checksum, backup)
- Version number
- Supports 128 primary partitions
- Does not have a boot code
- UEFI boots from an EFI partition







GRUB2 - GRand Unified Bootloader 2



- Boots in stages:
 - stage1 (core.img): Loads directly stage 2 (/boot/grub), usually from /boot partition. Contains file system drivers.
 - **stage2** (/boot/grub): Loads the default configuration file and any other modules needed.
- Themes, **graphical menus**, scripting support
- Uses UUIDs to identify disks
- Supports additionally: PowerPC
- Supports LLVM and RAID
- Boots live CD images from hard drive
- Automated search for other OS (like Windows)

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User space: systemd

- systemd is the first process started by the kernel
- The mother of all subprocesses
- Has always PID 1
- Looks in /etc/systemd/system/default.target for default target
- **Executes** the default **target** (unit) scripts
- **Starts** the user space **processes** on boot:
 - Daemons (crond, httpd, sshd, inetd, syslogd, ...)
 - Terminals
 - Graphical desktop
- Speed-up the boot: starts processes in parallel (where possible)



User space: systemd (commands)

- Get default target: systemctl get-default
- Change runlevel: systemctl isolate poweroff.target
- Control a daemon:
 /etc/systemd/system/<daemon>.service
 - service daemon start
 - service daemon stop
 - service daemon reload
 - service daemon restart
 - service daemon status
 - ...
- Enable a daemon: systemctl enable daemon.service
- Disable a daemon: systemctl disable daemon.service



User space: systemd (example)

1 [Unit]
2 Description=Demo daemon that does something useful
3
4 [Service]
5 Type=forking
6 ExecStart=/usr/sbin/daemon
7 ExecStopPost=/bin/rm /var/run/daemon.pid
8
9 [Install]
10 WantedBy=multi-user.target

More details: https://www.digitalocean.com/community/tutorials/understanding-systemd-units-and-unit-files

Goal Boot procedure BIOS/UEFI MBR/GPT Boot loader User space OS architecture Summary

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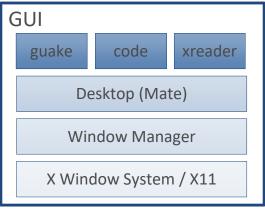
Init runlevel vs. system targets

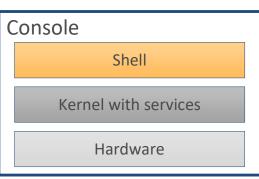
init runlevel	systemd target	systemd target aliases	Description
	halt.target		Halts the system without powering it down.
0	poweroff.target	runlevel0.target	Halts the system and turns the power off.
S	emergency.target		Single user mode. No services are running; filesystems are not mounted.
1	rescue.target	runlevel1.target	A base system including mounting the filesystems with only the most basic services running and a rescue shell on the main console.
2		runlevel2.target	Multiuser, without NFS but all other non-GUI services running.
3	multi-user.target	runlevel3.target	All services running but command line interface (CLI) only.
4		runlevel4.target	Unused.
5	graphical.target	runlevel5.target	Multiuser with a GUI.
6	reboot.target	runlevel6.target	Reboot.
	default.target		This target is always aliased with a symbolic link to either multi- user.target or graphical.target. systemd always uses the de- fault.target to start the system.

More details: https://opensource.com/article/17/2/linux-boot-and-startup



Linux high level overview (1)





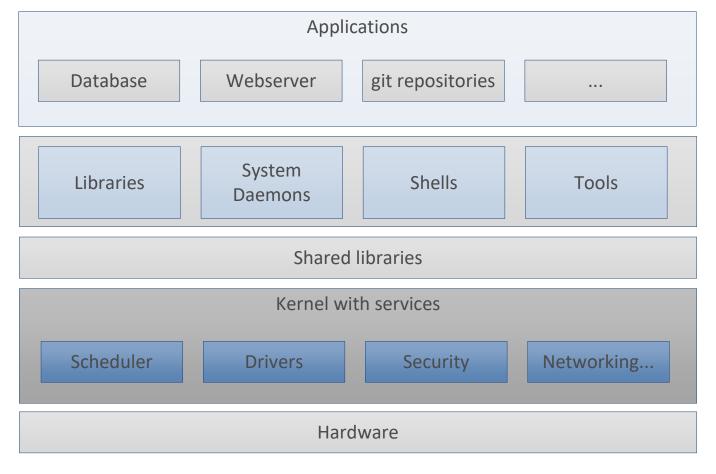
OS tasks

- Execute (graphical) user applications
- Provide desktop environment
- Draw windows (graphical elements)
- Provide shells
- Manage resources
- Support, abstract and virtualise hardware

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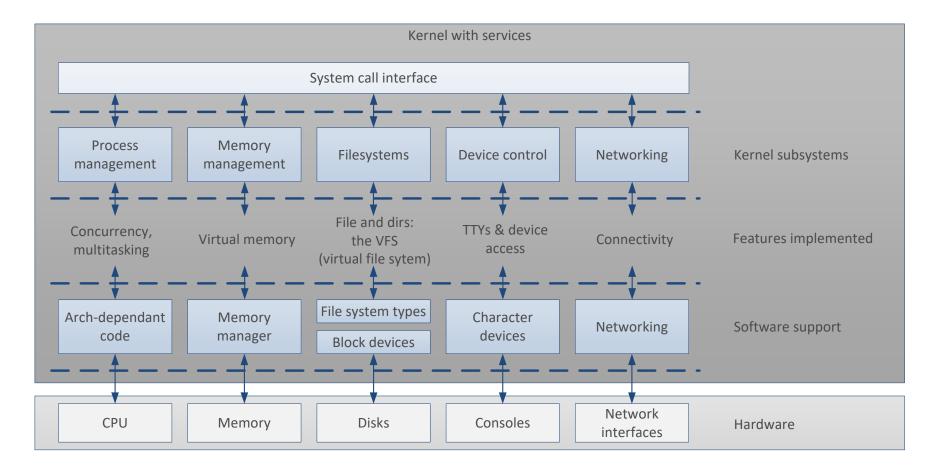
Linux high level overview (2)

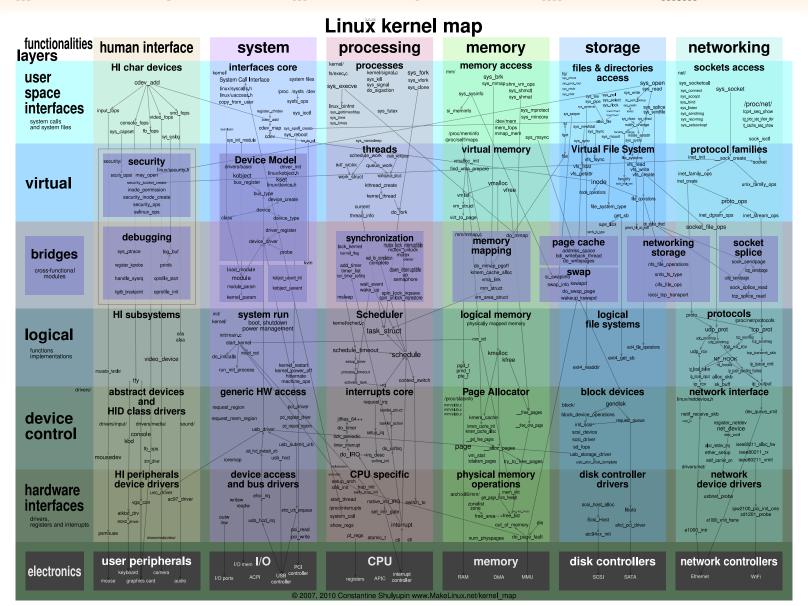


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OS overview: Linux kernel



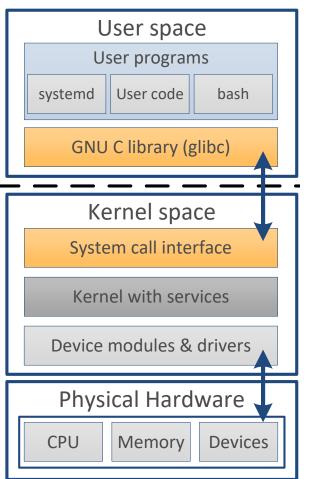


Source: http://www.makelinux.net/kernel_map

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Protection: user vs. kernel space



User space

- All code outside the kernel
- Also called "userland"
- **Restricted** (encapsulated) access to the hardware
- Can **only** use a **subset of CPU instructions**
- Can **only** access the **assigned memory** addresses
- **Crashes** in a user process: **only stops these process**.

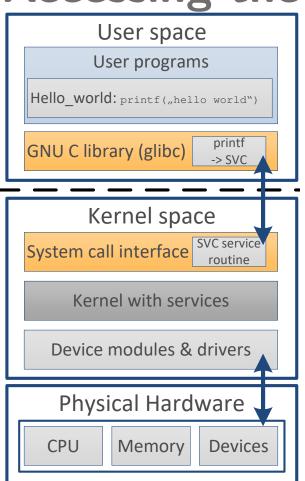
Kernel space

- Complete and unrestricted access to the hardware
- Can execute any CPU instruction
- Can access any memory address
- **Crashes** in kernel are "catastrophic": **system stop!**

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Accessing the kernel: supervisor call



Supervisor call (SVC)

- **CPU** instruction to give control to the OS/kernel
- **Requests** for an **OS service**:
 - Start process
 - Allocate memory
 - File open/close/read/write/...
 - Print something on terminal (TTY) or screen
 - Send data over network
- SVCs are numbered
- The calling **process is interrupted** while the kernel executes the SVC.

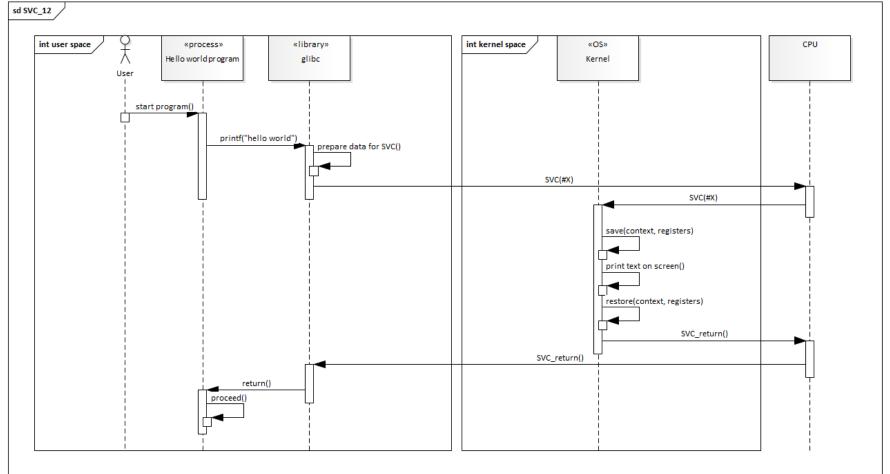
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Supervisor call (SVC) example sequence



Goal Boot procedure BIOS/UEFI MBR/GPT Boot loader User space OS architecture Summary

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Summary and outlook

Summary

- Boot procedure
- BIOS/UEFI
- MBR/GPT
- Boot loader: Grub2
- init/systemd
- OS architecture

Outlook

- Processes vs. Threads
- Parallelisation