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Exercise sheet 7 – Processor architecture

Goals:

- Pipelining
- Instruction Scheduling

Exercise 7.1: Pipelining

(a) Given is a sequence of instructions and a 5-stage-pipeline. State the procedure and be careful, the instructions are not ordered perfectly to fully utilise the pipeline.

Instruction sequence:

1115010	action sequence:	
Nr.	Instruction	Comment
(1)	ADD R2, R1, R1	; R1 = R1 + R2
(2)	ADD R1, R3, R3	; R3 = R1 + R3
(3)	ADD R4, R5, R5	; R5 = R4 + R5
(4)	CMP R5, #0, R6	; R6 = cmp(R5, 0)
(5)	BNE R6, M	; Jump to M if result != 0 BNE = branch not equal
(6)	INST1	; some random instruction 1
(7)	INST2	; some random instruction 2
(8)	M: INST100	; some random instruction 100

Five-stage-pipeline:

Stage	Operation
1.	Fetch instruction
2.	Load operands
3.	Execute instruction
4.	Memory access
5.	Save result into register

Proposal for solution:

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Fetch instruction	Load operands	Execute instruction	Memory access	Save result into register
(1) ADD R2,R1,R1				
(2) ADD R1,R3,R3	(1) ADD R2,R1,R1			
(2) ADD R1,R3,R3	(*1)	(1) ADD R2,R1,R1		
(2) ADD R1,R3,R3			(1) ADD R2,R1,R1	
(2) ADD R1,R3,R3				(1) ADD R2,R1,R1 //write R1
(3) ADD R4,R5,R5	(2) ADD R1,R3,R3			
(4) CMP R5, #0,R6	(3) ADD R4,R5,R5	(2) ADD R1,R3,R3		
(4) CMP R5, #0,R6	(*1)	(3) ADD R4,R5,R5	(2) ADD R1,R3,R3	
(4) CMP R5, #0,R6			(3) ADD R4,R5,R5	(2) ADD R1,R3,R3 //write R3
(4) CMP R5, #0,R6				(3) ADD R4,R5,R5 //write R5
(5) BNE R6, M	(4) CMP R5, #0,R6			
(5) BNE R6, M	(*1)	(4) CMP R5, #0,R6		
(5) BNE R6, M			(4) CMP R5, #0,R6	
(5) BNE R6, M				(4) CMP R5, #0,R6 //write R6
(6) INST 1 (*2)	(5) BNE R6, M			
(7) INST 2 (*2)		(5) BNE R6, M		
			(5) BNE R6, M	
			(*3)	(5) BNE R6, M //jump to INST100
8) INST 100	(0) 11.07 1.00			
INST 101	(8) INST 100			
		(8) INST 100	(0) 11.07 100	
			(8) INST 100	(0) INIOT 100
				(8) INST 100

(b) Suggest ideas for a better pipeline utilisation of exercise 7.1a.

Assume there are additional instructions before the given program excerpt.

Proposal for solution: It is not possible to avoid data-dependency itself, but you can rearrange the execution order of the instructions that do not depend on each other. For example while the program waits for $ADD\ R2,R1,R1$ to finish and to load $ADD\ R1,R3,R3$ it could already load $ADD\ R4,R5,R5$. This is called *instruction scheduling*.

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Fetch instruction	Load operands	Execute instruction	Memory access	Save result into register
1) ADD R2,R1,R1				
3) ADD R4,R5,R5	(1) ADD R2,R1,R1			
2) ADD R1,R3,R3	(3) ADD R4,R5,R5	(1) ADD R2,R1,R1		
2) ADD R1,R3,R3	(*1)	(3) ADD R4,R5,R5	(1) ADD R2,R1,R1	
2) ADD R1,R3,R3			(3) ADD R4,R5,R5	(1) ADD R2,R1,R1 //write R1
4) CMP R5, #0,R6	(2) ADD R1,R3,R3			(3) ADD R4,R5,R5 //write R5
5) BNE R6, M	(4) CMP R5, #0,R6	(2) ADD R1,R3,R3		
5) BNE R6, M	(*1)	(4) CMP R5, #0,R6	(2) ADD R1,R3,R3	
5) BNE R6, M			(4) CMP R5, #0,R6	(2) ADD R1,R3,R3 //write R1
5) BNE R6, M				(4) CMP R5, #0,R6 //write R6
(6) INST 1 (*2)	(5) BNE R6, M			
7) INST 2 (*2)		(5) BNE R6, M		
			(5) BNE R6, M	
			(*3)	(5) BNE R6, M
8) INST 100				
NST 101	(8) INST 100			
		(8) INST 100		
			(8) INST 100	
				(8) INST 100

(*1) Data dependency

- (*2) Assumption: The next instructions (INST 1, INST 2) are already loaded into the pipeline
- speculative loading of instructions if the jump is not taken, execution can be proceeded with less empty cycles
- (*3) Assumption: The jump is taken
- The plant is the

Another order which additionally can reduce the overall time by one time step:

Fetch instruction	Load operands	Execute instruction	Memory access	Save result into register
(3) ADD R4,R5,R5				
(1) ADD R2,R1,R1	(3) ADD R4,R5,R5			
(4) CMP R5, #0,R6	(1) ADD R2,R1,R1	(3) ADD R4,R5,R5		
(4) CMP R5, #0,R6	(*1)	(1) ADD R2,R1,R1	(3) ADD R4,R5,R5	
(4) CMP R5, #0,R6			(1) ADD R2,R1,R1	(3) ADD R4,R5,R5 //write R5
(2) ADD R1,R3,R3	(4) CMP R5, #0,R6			(1) ADD R2,R1,R1 //write R1
(5) BNE R6, M	(2) ADD R1,R3,R3	(4) CMP R5, #0,R6		
(5) BNE R6, M	(*1)	(2) ADD R1,R3,R3	(4) CMP R5, #0,R6	
(5) BNE R6, M			(2) ADD R1,R3,R3	(4) CMP R5, #0,R6 //write R6
(6) INST 1 (*2)	(5) BNE R6, M			(2) ADD R1,R3,R3 //write R1
(7) INST 2 (*2)		(5) BNE R6, M		
			(5) BNE R6, M	
			(*3)	(5) BNE R6, M //jump to INST100
(8) INST 100				
INST 101	(8) INST 100			
		(8) INST 100		
			(8) INST 100	
				(8) INST 100

(*1) Data dependency

- (*2) Assumption: The next instructions (INST 1, INST 2) are already loaded into the pipeline
- speculative loading of instructions
 if the jump is not taken, execution can be proceeded with less empty cycles
- (*3) Assumption: The jump is taken
- → INST 1 and INST 2, which are already loaded into the pipeline have to be dropped → the pipeline has to rebuild completely, beginning at INST 100 (jump marker M)

Exercise 7.2: Pipeline-Simulator

The pipeline simulator is taken from http://euler.vcsu.edu/curt.hill/mips.html.

- (a) Update the CA exercises repository with git pull.
- (b) Change into the CA_exercises/sheet_07_pipelining/PipelineSimulator directory.

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- (c) To compile the pipeline-simulator run javac GUI.java from a terminal, or use the Makefile. Hint: Make sure that you have properly installed the JDK and JRE. Within the VM it is already installed.
- (d) Start the simulator with: java GUI
- (e) In *mips.html* you can find some information on how to work with the simulator.
- (f) Create a little assembly program, based on exercise 7.1a and run the simulator. Hint: There is no one-to-one mapping possible, because not all required instructions are available.

Proposal for solution: Currently there is no solution to load a program into the simulator than loading the instructions one by one.

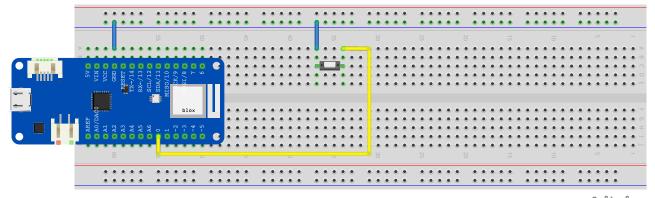
```
O: ADDI $1 $0 1
                     //can be used as an initialisation step for the registers
   0: ADDI $5 $0 3
                     //can be used as an initialisation step for the registers
   0: ADD $2 $1 $1
   1: ADD $3 $2 $2
     ADD $4 $3 $3
     ADD
          $5 $4 $4
     ADD
          $6 $5 $5
  5: ADD $7 $6 $6
9
   6: ADDI $1 $1 0
10
   7: BNE $10 $1 0
```

This programs run infinite, as the statement in line 7 never becomes true. You can check out the current values in the register-view in the right upper corner of the simulator window.

Exercise 7.3: Instruction scheduling (coding)

We want to write an Arduino sketch which counts the number of button presses. If the button is pressed, an interrupt occurs which calls an ISR.

(a) Prepare the wiring with the Arduino MKR WiFi 1010 as follows:



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- (b) Please double check your wiring with the lecturer, before you connect the Arduino MKR WiFi 1010.
- (c) Make sure that you have installed the Arduino IDE (https://www.arduino.cc/en/software).
- (d) Make sure you have installed the board SDK:
 - Tools -> Board: -> Boards Manager...
 - Install (latest version): Arduino SAMD Boards (32-bits ARM Cortex-MO+)

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- (e) Open the skeleton file from CA_exercises/sheet_07_pipelining/io_interrupt/io_interrupt.ino with the Arduino IDE.
- (f) Follow the TODOs in the code. Some configuration depends on your wiring of the I/O pins. Hint: The Arduino reference contains descriptions of the used functions: https://www.arduino.cc/reference/en.

```
Proposal for solution:
   //PIN configuration
   const int BUTTON PIN = 0;
2
   //Global variables
4
   volatile uint8_t counter = 0;
5
8
    * setup() is called once on startup/reset of the Arduino
9
10
   void setup(){
11
     //Configure serial connection
12
     Serial.begin(9600);
13
14
     //configure PINs
15
     pinMode(BUTTON PIN, INPUT PULLUP);
16
17
     //Configure interrupt
18
     attachInterrupt(digitalPinToInterrupt(BUTTON_PIN), isr_button_pressed, RISING); //TO
19
20
21
^{22}
    * loop() is called as fast as possible.
23
24
    * As you can see, there is no call to a function
25
    * changing the LED state in the main-loop
26
27
   void loop(){
28
     //TODO: transmit the counter value with the serial connection
29
     Serial.print("The value is: ");
30
     Serial.println(counter);
31
32
     delay(1000); //wait for 1000 ms -> 1 sec
33
   }
34
35
36
    * isr_button_pressed() = Interrupt service routine
37
    * Change here the state of the LED
38
    */
39
   void isr_button_pressed(){
40
     static unsigned long time_prev_rising_edge = 0;
41
42
     if (millis() - time_prev_rising_edge > 50) { //only react on rising edges every 50 m
43
       counter++;
44
45
       time_prev_rising_edge = millis();
46
47
   }
48
```

(g) Configure the board within the Arduino IDE: Tools -> Board: -> Arduino MKR WiFi 1010.

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- (h) Compile (verify) your sketch within the Arduino IDE. If it compiles then upload it your sketch.
- (i) Open the 'Serial Monitor' to see the printed strings and do some debugging with the text based logging.
- (j) Press the button to test your sketch. Does it work as expected?