

6.857  
**NETWORK AND COMPUTER SECURITY**  
**RECITATION 2**

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1. FAST EXPONENTIATION

1.1. **Euler's Theorem.** *Theorem:* Take any element  $a \in \mathbb{Z}_n^*$ , then for all  $n \in \mathbb{N}$ , we have  $a^{\phi(n)} \equiv 1 \pmod{n}$ , where  $\phi(n)$  is the order of the group.

*Corollary:*  $\forall n \in \mathbb{N}$  and  $\forall a \in \mathbb{Z}_n^*$ , we have  $a^d \equiv a^{d \bmod \phi(n)} \pmod{n}$ .

For any prime  $p$ , we have  $\phi(p) = p - 1$ . As a consequence, we obtain *Fermat's Little Theorem:*  $a^{\phi(p)} \equiv a^{p-1} \equiv 1 \pmod{p}$ .

1.2. **Exercises with Euler's Theorem.**

$$(1) \quad 3^{25} \pmod{13} = 3^{25 \bmod 12} = 3 \pmod{13}$$

$$(2) \quad 7^{10} \pmod{10} \equiv 7^{10 \bmod 4} \pmod{10} \equiv 9 \pmod{10}$$

Because we know that  $10 = (2)(5)$  so that  $\phi(10) = \phi(2)\phi(5) = 4$ .

2. CHINESE REMAINDER THEOREM

*Theorem:* For all  $m_1, \dots, m_r < N$  where  $\gcd(m_i, m_j) = 1$  for all  $i, j$ ,  $m = m_1 \dots m_r < N$ , and for all  $a_1, \dots, a_r < N$ , there exists an integer  $y$  such that

$$(3) \quad y \equiv a_1 \pmod{m_1}$$

$$(4) \quad y \equiv a_2 \pmod{m_2}$$

$$(5) \quad \vdots$$

$$(6) \quad y \equiv a_r \pmod{m_r}$$

Moreover,  $y$  is easy to find (can be found in  $\theta(\log N)$ ).

2.1. **Proof of CRT.**

(1)  $n_i = \frac{m}{m_i}$  for all  $i$ . So for instance  $n_i = m_1 m_2 \dots m_{i-1} m_{i+1} \dots m_r$ .

(2) Compute  $b_i$  which is the inverse of  $n_i \pmod{m_i}$ . So compute  $b_i n_i \equiv 1 \pmod{m_i}$ . We know that if  $\gcd(c, d) = 1$ , then  $cx + dy \equiv 1 \pmod{d}$ , which implies  $cx \equiv 1 \pmod{d}$ . But we know that  $\gcd(n_i, m_i) = 1$  because all of the factors of  $n_i$  are  $m_j$  where  $j \neq i$  and we have pairwise  $\gcd(m_i, m_j) = 1$ .

(3)  $y = \sum_i n_i b_i a_i \pmod{m}$

Proof: Straightforward, see wikipedia.

2.2. **Example.** Find  $y$  such that  $y \equiv 6 \pmod{7}$  and  $y \equiv 8 \pmod{11}$ .

We can do this by using CRT. We know  $a_1 = 6, a_2 = 8, m_1 = 7, m_2 = 11$ . So  $n_1 = 11, n_2 = 7$ . Now we can compute the inverses,  $b_1 = 2$  because  $2(11) \equiv 1 \pmod{7}$  and  $b_2 = 8$  because  $8(7) \equiv 1 \pmod{11}$ . So we know that  $y = (11)(2)(6) + (7)(8)(8) \pmod{77} = 41 \pmod{77}$ .