

6.857
NETWORK AND COMPUTER SECURITY
LECTURE 5

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1. HASH FUNCTION APPLICATIONS

1.1. **Password Storage.** System stores $h(pw)$ rather than pw itself. System might also store username, salt, etc.

1.2. **File Modification Detector.** You want to monitor to detect when files have been changed. For each file, store $h(F)$ securely. You can check to see if the files have been modified by recomputing the hash. Provides detection (not prevention).

1.3. **Digital Signatures (hash and sign).**

- PK_A is Alice's public key (for signature verification).
- SK_A is Alice's secret key (for signing).
- Signing: $\sigma = \text{sign}(SK_A, m)$ and σ is Alice's signature on message.
- Verification: $\text{verify}(M, \sigma, PK_A) \in \{\text{true}, \text{false}\}$.

Idea: computing $h(m)$ is fast, so sign $h(m)$ instead of signing m . We do $\text{sign}(m, SK_A) = \text{sign}(h(m), SK_A)$ if $h(m) = h(m')$.

Problem is that if $h(m) = h(m')$, then asking Alice to sign m , her signature σ is also a signature for m' .

1.4. **Commitments.** Alice has value x which is her bid. She computes $C(x)$ and gives auctioneer $C(x)$, which is her sealed bid. When bidding is over, Alice should be able to open $C(x)$ to reveal x .

Want these properties:

- Binding: Alice should not be able to open $C(x)$ in more than one way.
- Secrecy: Anyone seeing $C(x)$ should have no information about x .
- Non-malleability: Anyone seeing $C(x)$ shouldn't be able to come up with a related bid, e.g. $C(x+1)$.

Let's try $C(x) = h(\text{username}||x)$.