RISC-V Instruction Set Architecture I

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Lecture 3

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Competing ISAs

RISC-V ISA

- x-86 Architecture CISC, register-memory, used in desktop computers, servers, proprietary
- ARM Architecture RISC, register-register, majority of embedded and cellular market, proprietary
- RISC-V Architecture RISC, register-register, upcoming, open-source

RISC-V ISA

RISC-V ISA

 RISC-V is a modular architecture with one base architecture (mandatory) and multiple extensions (optional)

Table: RISC-V Base architectures.

ISA	Description				
RV32I	Base Integer Instruction Set, 32-bit				
RV64I	Base Integer Instruction Set, 64-bit				
RV128I	Base Integer Instruction Set, 128-bit				
RV32E	Base Integer Instruction Set for Embedded				
	Systems, 32-bit (Subset of RV32I)				

RISC-V ISA Cont'd

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Table: List of RISC-V ISA extensions.

Extension	Description
M	Standard Extension for Integer Multiplication, Division
Α	Standard Extension for Atomic Instructions
F	Standard Extension for Single-Precision Floating-Point
D	Standard Extension for Double-Precision Floating-Point
Q	Standard Extension for Quad-Precision Floating-Point
L	Standard Extension for Decimal Floating-Point
C	Standard Extension for Compressed Instructions
В	Standard Extension for Bit Manipulation
J	Standard Extension for Dynamically Translated Languages
Т	Standard Extension for Transactional Memory
Р	Standard Extension for Packed-SIMD Instructions
V	Standard Extension for Vector Operations
N	Standard Extension for User-Level Interrupts

RISC-V ISA Cont'd

RISC-V ISA

- Instruction Encoding \sim variable length, e.g. RV32IC will have 32-bit base and 16-bit compressed extension
- Addressing Modes ~ four addressing modes, register, immediate, displacement (including PC-relative)
- Instruction Operands \sim number of operands can be from zero to three

- General Purpose Registers (GPRs) (32 registers)
 - x0. x1. x31
 - 32-bit wide integer registers
 - x0 is hard-wired to zero
- Control and Status Registers (CSRs)
 - User Level CSRs.
 - Supervisor Level CSRs
 - Hypervisor Level CSRs
 - Machine Level CSRs
- Program counter (PC)

RV32I Instruction Set

- Three groups of instructions and corresponding tree
 - Data Processing Instructions
 - Register-Register instructions (R-type)
 - Register-immediate instructions (I-type)
 - Register-immediate instructions (U-type)

RV32I Instruction Set

- Three groups of instructions and corresponding tree
 - Data Processing Instructions
 - Register-Register instructions (R-type)
 - Register-immediate instructions (I-type)
 - Register-immediate instructions (U-type)
 - Memory Access Instructions
 - Load (I-type)
 - Store (S-type)

RV32I Instruction Set

- Three groups of instructions and corresponding tree
 - Data Processing Instructions

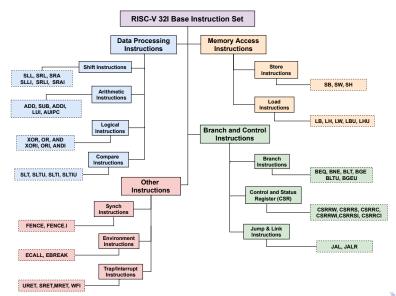
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- Register-Register instructions (R-type)
- Register-immediate instructions (I-type)
- Register-immediate instructions (U-type)
- Memory Access Instructions
 - Load (I-type)
 - Store (S-type)
- Flow Control Instructions
 - Unconditional jump and link (J-type)
 - Unconditional jump via register and link (I-type)
 - Conditional branches (B-type)



RV32I Instruction Set Cont'd

RISC-V ISA



RISC-V Base Instruction Formats: R, I, S, U

R Format



I Format



S Format



U Format





RISC-V Base Instruction Formats: B Format

- Similar to S Format
 - S Format

imm[11:5]	rs2	rs1	func3	imm[4:0]	opcode
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B Format

imm[12:10:5]	rs2	rs1	func3	imm[4:1:11]	opcode

RISC-V Base Instruction Formats: J Format

- Similar to U Format
 - U Format

imm[31:12]	rd	opcode
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J Format

imm[20]:imm[10:1]:imm[11]:imm[19:12]	rd	opcode
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Data Processing Instructions: R Type

• Format: R-type

7	5	5	3	5	7
func7	rs2	rs1	func3	rd	opcode

Table: R type data processing instructions.

In	struction	Operation	Type	Illustration
SI	_L	Shift Left Logical	R	SLL rd,rs1,rs2
SI	RL	Shift Right Logical	R	SRL rd,rs1,rs2
SI	RA	Shift Right Arithmetic	R	SRA rd,rs1,rs2
Α	DD	Addition	R	ADD rd,rs1,rs2
SI	JB	Subtraction	R	SUB rd,rs1,rs2
X	OR	XOR	R	XOR rd,rs1,rs2
0	R	OR	R	OR rd,rs1,rs2
Α	ND	AND	R	AND rd,rs1,rs2

Data Processing Instructions: R Type Cont'd

Table: R type data processing instructions.

Instruction	Operation	Туре	Illustration
MUL	Multiply	R	MUL rd,rs1,rs2
MULH	Multiply upper Half	R	MULH rd,rs1,rs2
MULHSU	Multiply Half Sign/Uns	R	MULHSU rd,rs1,rs2
MULHU	Multiply upper Half Uns	R	MULHU rd,rs1,rs2
DIV	Divide	R	DIV rd,rs1,rs2
DIVU	Divide Unsigned	R	DIVU rd,rs1,rs2

Format: R-type

7	5	5	3	5	7
func7	rs2	rs1	func3	rd	opcode

The destination register: rd ← rs1(func3, func7)rs2

Data Processing Instructions: I Type

• Format: I-type

12	5	3	5	7
imm[11:0]	rs1	func3	rd	opcode

Table: I type data processing instructions.

Instruction	Operation	Туре	Illustration
ADDI	Add immediate	I	ADDI rd, rs1, imm
ANDI	AND immediate	ı	ANDI rd, rs1, imm
ORI	OR immediate	- 1	ORI rd, rs1, imm
XORI	XOR immediate	1	XORI rd, rs1, imm
SLTI	Set less than immediate	1	SLTI rd, rs1, imm
SLTIU	Set less than immediate unsigned	I	SLTIU rd, rs1, imm

Data Processing Instructions: I Type Cont'd

Table: I type data processing instructions.

Instruction	Operation	Туре	Illustration
SLLI	Shift left logical immediate	ı	SLLI rd, rs1, shamt
SRLI	Shift right logical immediate	ı	SRLI rd, rs1, shamt
SRAI	Shift right arithmetic immediate	ı	SRAI rd, rs1, shamt

• Format: I-type

7 : 5	5	3	5	7
imm[11:5] : imm[4:0]	rs1	func3	rd	opcode

• The shift amount: imm[4:0] ← shamt

Data Processing Instructions: U Type

• Format: U-type

20	5	7
imm[31:12]	rd	opcode

Table: U type data processing instructions.

Instruction	truction Operation		Illustration
LUI	Load Upper Immediate	U	LUI rd, imm
AUIPC	Add Upper Immediate to PC	U	AUIPC rd, imm

- U-imm = imm[31:12] : 12'b0
- opcode = LUI : $rd \leftarrow U$ -imm
- opcode = AUIPC : $rd \leftarrow pc + U$ -imm

Memory Access Instructions: I Type

• Format: I-type

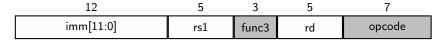


Table: I type memory access instructions.

Instruction	Operation	Туре	Illustration
LB	Load Byte	I	LB rd, rs1, imm
LH	Load Halfword	- 1	LH rd, rs1, imm
LW	Load Word	- 1	LW rd, rs1, imm
LBU	Load Byte Unsigned	I	LBU rd, rs1, imm
LHU	Load Half Unsigned	I	LHU rd, rs1, imm

Memory Access Instructions: I Type Cont'd

• Format: I-type



- I-imm = signExtend(imm[11:0])
- opcode = LOAD: rd ← mem[rs1 + l-imm]
- funct3 = LW/LB/LBU/LH/LHU

Memory Access Instructions: S Type

Format: S-type

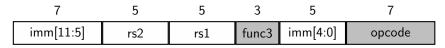


Table: S type memory access instructions.

Instruction	Operation	Туре	Illustration
SB	Store Byte	I	SB rs1, rs2, imm
SH	Store Halfword	I	SH rs1, rs2, imm
SW	Store Word	I	SW rs1, rs2, imm

Memory Access Instructions: S Type Cont'd

• Format: S-type

7	5	5	3	5	7
imm[11:5]	rs2	rs1	func3	imm[4:0]	opcode

- S-imm = signExtend(imm[11:5], imm[4:0])
- opcode = STORE: $mem[rs1 + S-imm] \leftarrow rs2$
- funct3 = SW/SB/SH

Memory Access Instructions: Addressing Modes

- PC-relative Addressing: Memory address is the sum of PC and an offset (e.g. auipc, jal and branch instructions). The offset is specified in the instruction.
- Register-offset Addressing: Memory address is the sum of a Register and an offset (e.g. jalr, addi and all memory instructions). The offset is specified in the instruction.
- Absolute Addressing: Memory address is a constant (32-bit) (e.g the lui instruction). Though this addressing can be interpreted as register-offset addressing with x0 being the register (x0-offset).

Flow Control Instructions: J Type

Format: J-type

Table: J type flow control instructions.

Instruction	Operation	Type Illustration	
JAL	Jump and Link	J	JAL rd, imm

- J-imm = signExtend(inst[31], inst[19:12], inst[20], inst[30:21], 1'b0)
- opcode = JAL: $rd \leftarrow pc + 4$; $pc \leftarrow pc + J$ -imm
- Jump = ± 1 MB range

Flow Control Instructions: I Type

Format: I-type

12	5	3	5	7
imm[11:0]	rs1	func3	rd	opcode

Table: I type flow control instructions.

Instruction	Operation	Туре	Illustration
JALR	Jump and Link	I	JALR rd, rs1, imm

- I-imm = signExtend(imm[11:0])
- opcode = JALR: rd \leftarrow pc + 4; pc \leftarrow ((rs1 + I-imm) & \sim (0x01))

Flow Control Instructions: B Type

Conditional branches: B-type

1 : 6	5	5	3	4 : 1	7
imm[12:10:5]	rs2	rs1	func3	imm[4:1:11]	opcode

Table: B type flow control instructions.

Instruction	Operation	Type	Illustration
BEQ	Branch Equal	В	BEQ rs1, rs2, imm
BNE	Branch Unequal	В	BNE rs1, rs2, imm
BLT	Branch Lessthan	В	BLT rs1, rs2, imm
BGE	Branch Greater equal	В	BGE rs1, rs2, imm
BLTU	Branch Lessthan Unsigned	В	BLTU rs1, rs2, imm
BGEU	Branch Greater equal Unsigned	В	BGEU rs1, rs2, imm

Flow Control Instructions: B Type Cont'd

• Format: B-type

- B-imm = signExtend(inst[31], inst[7], inst[30:25], inst[11:8], 1'b0)
- opcode = BRANCH: $pc \leftarrow compare(funct3, rs1, rs2)$? pc + B-imm : pc + 4
- funct3 = BEQ/BNE/BLT/BLTU/BGE/BGEU

Suggested Reading

- Read User Manual for the instruction set and its architecture [Waterman et al., 2016b].
- For Control and Status register description consult Privileged architecture manual [Waterman et al., 2016a].

Acknowledgment

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References



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