NPTEL MOOC

PROGRAMMING, DATA STRUCTURES AND ALGORITHMS IN PYTHON

Week 1, Lecture 3

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Algorithm for gcd(m,n)

- * To find the largest common factor, start at the end and work backwards
- * Let i run from min(m,n) to 1
- * First common factor that we find will be gcd!

- * Suppose d divides both m and n, and m > n
- * Then m = ad, n = bd
- * So m-n = ad bd = (a-b)d
- * d divides m-n as well!
- * So gcd(m,n) = gcd(n,m-n)

- * Consider gcd(m,n) with m > n
- * If n divides m, return n
- * Otherwise, compute gcd(n,m-n) and return that value

```
def gcd(m,n):
   \# Assume m >= n
   if m < n:
     (m,n) = (n,m)
   if (m\%n) == 0:
     return(n)
   else:
     diff = m-n
     # diff > n? Possible!
     return(gcd(max(n,diff),min(n,diff))
```

Euclid's algorithm, again

```
def gcd(m,n):
 if m < n: # Assume m >=
    (m,n) = (n,m)
  while (m%n) != 0:
    diff = m-n
    # diff > n? Possible!
    (m,n) = (max(n,diff),min(n,diff))
  return(n)
```

Even better

- * Suppose n does not divide m
- * Then m = qn + r, where q is the quotient, r is the remainder when we divide m by n
- * Assume d divides both m and n
- * Then m = ad, n = bd
- * So ad = q(bd) + r
- * It follows that r = cd, so d divides r as well

- * Consider gcd(m,n) with m > n
- * If n divides m, return n
- * Otherwise, let r = m%n
- * Return gcd(n,r)

```
def gcd(m,n):
  if m < n: # Assume m
    (m,n) = (n,m)
  if (m\%n) == 0:
    return(n)
  else:
    return(gcd(n, m%n)) # m%n < n, always!
```

Euclid's algorithm, revisited

```
def gcd(m,n):
  if m < n: # Assume m
    (m,n) = (n,m)
  while (m%n) != 0:
    (m,n) = (n,m\%n) # m\%n < n, always!
  return(n)
```

Efficiency

- * Can show that the second version of Euclid's algorithm takes time proportional to the number of digits in m
- * If m is 1 billion (10⁹), the naive algorithm takes billions of steps, but this algorithm takes tens of steps