

Date Structures and Analysis of Algorithms

Midterm Examination (Two Hours)

1. (10 marks)

a). (5 marks) Fill in the relationship between each pair of the functions below.
(Use Ω , O , Θ , or none of the above)

$$n = ___ n + n/\log_2(n) \quad (1)$$

$$n + \sqrt{n} = ___ n \quad (2)$$

$$n * \log_2(n^2) = ___ n^{\log_2(n)} \quad (3)$$

$$(2/3)^n = ___ n \quad (4)$$

$$2^n = ___ (3/2)^n \quad (5)$$

b). (5 marks) Answer following questions.

- What is the worst case time complexity of constructing Huffman coding tree?
- What is the average case time complexity of quicksort?
- What is the worst case time complexity of heap construction (heapification) ?
- What is the worst case time complexity of sequence comparison algorithm?
- What is the lower bound for sorting algorithm based on comparison?

2. (10 marks)

a). (4 marks) Prove, by using the definitions of O and Ω , the following:

$$\sqrt{n} = O(\sqrt{n} - 5)$$

$$\log_2(\max\{f(n), g(n)\}) = \Omega(\log_2(f(n) * g(n)))$$

b). (6 marks) Find a closed form for the following recursive equation.

$$T(n) = T(n/2) + 1; \text{ if } n > 1 \text{ and } n \text{ is of power of } 2.$$

$$T(1) = 1;$$

3. (10 marks)

a) (5 marks) For a given string $S = p_1p_2...p_n$, the definition for $next(i)$, $2 \leq i \leq n$, is as follows:

$next(i)$ = the maximum j ($0 < j < i - 1$) such that $p_1p_2...p_j = p_{i-j}p_{i-j+1}...p_{i-1}$,

0 if no such j exists.

($next(1)$ is defined as -1.)

Compute the $next$ function for string "abababc".

b) (5 marks) Given two strings, $A = a_1, a_2, ...a_m$ and $B = b_1, b_2, ...b_n$, we define a matrix $C[0..m, 0..n]$ as follows: ($1 \leq i \leq m, 1 \leq j \leq n$)

$$C[0, 0] = 0, C[0, j] = j, C[i, 0] = i.$$

$$C[i, j] = \min\{C[i - 1, j] + 1, C[i, j - 1] + 1, C[i - 1, j - 1] + r_{ij}\} \text{ where } r_{ij} = 0, \text{ if } a_i = b_j, \text{ or } = 1, \text{ if } a_i \neq b_j.$$

Compute the C matrix for strings $A = ababab$, and $B = babaaa$. What is the meaning of the value in $C[m, n]$?

4. (10 marks)

Given a set of characters $A = \{a_1, a_2, \dots, a_n\}$, where $n > 2$. Let f_i and f_j be the frequencies of a_i and a_j , and d_i and d_j be the depths of a_i and a_j in a Huffman encoding tree for A . Prove or disprove (by counter example) the following statement.

- a. (5 marks) if $d_i = d_j$ then $f_i = f_j$.
- b. (5 marks) if $f_i > f_j$ then $d_i \leq d_j$.

5. (10 marks) Suppose that you are given a sorted list of n elements followed by $f(n)$ randomly ordered elements. How would you sort the entire list if

- a. (3 marks) $f(n) = 2$?
- b. (4 marks) $f(n) = \sqrt{n}$?
- c. (3 marks) How large can $f(n)$ be for the entire list to be sorted in $O(n)$ time?