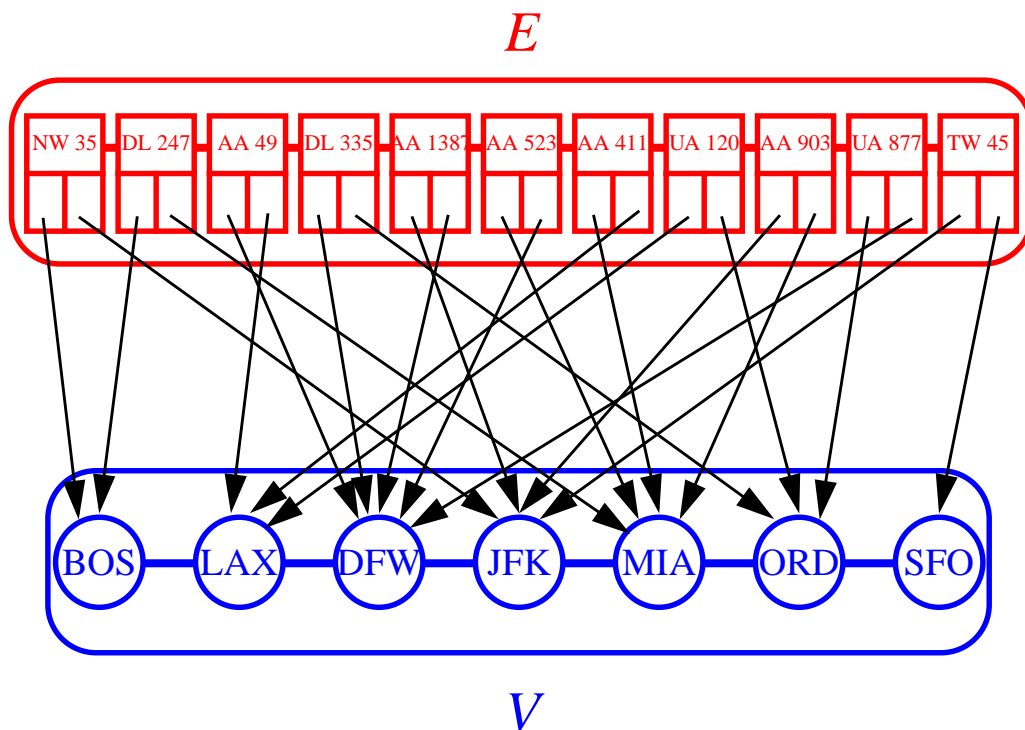


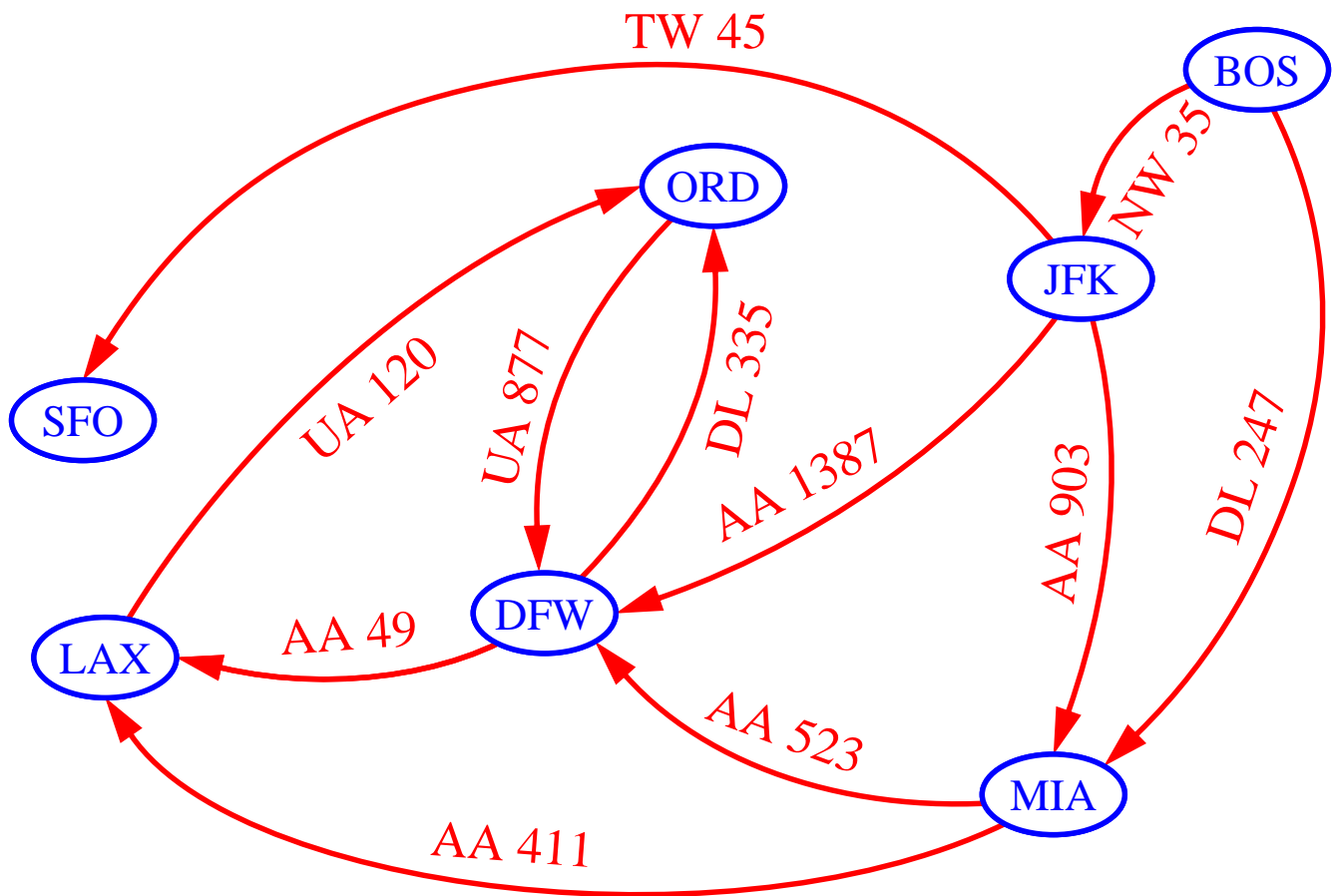
DATA STRUCTURES FOR GRAPHS

- Edge list
- Adjacency lists
- Adjacency matrix



Data Structures for Graphs

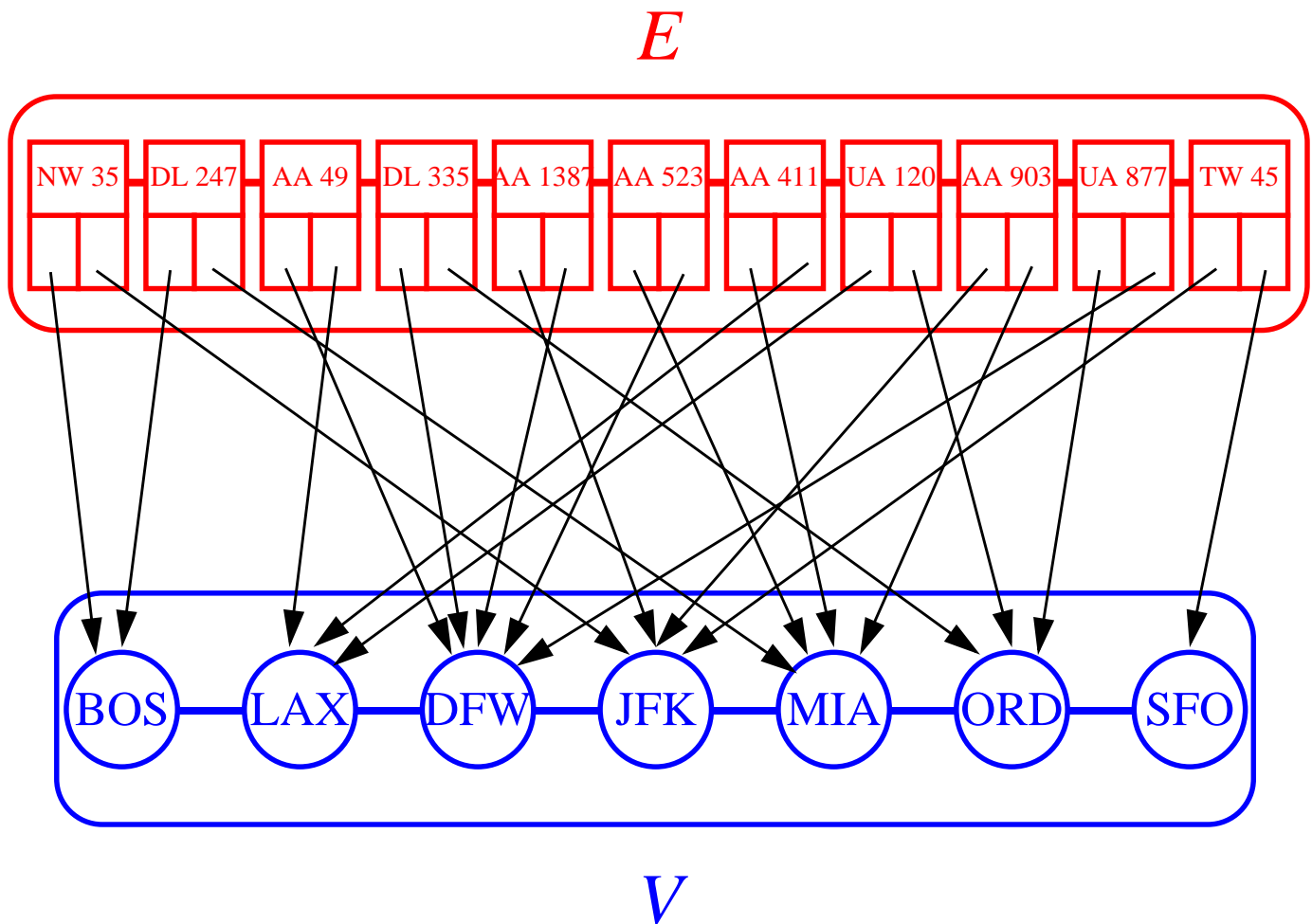
- A Graph! How can we represent it?
- To start with, we store the **vertices** and the **edges** into two containers, and each edge object has references to the vertices it connects.



- Additional structures can be used to perform efficiently the methods of the Graph ADT

Edge List

- The **edge list** structure simply stores the vertices and the edges into unsorted sequences.
- Easy to implement.
- Finding the edges incident on a given vertex is inefficient since it requires examining the entire edge sequence

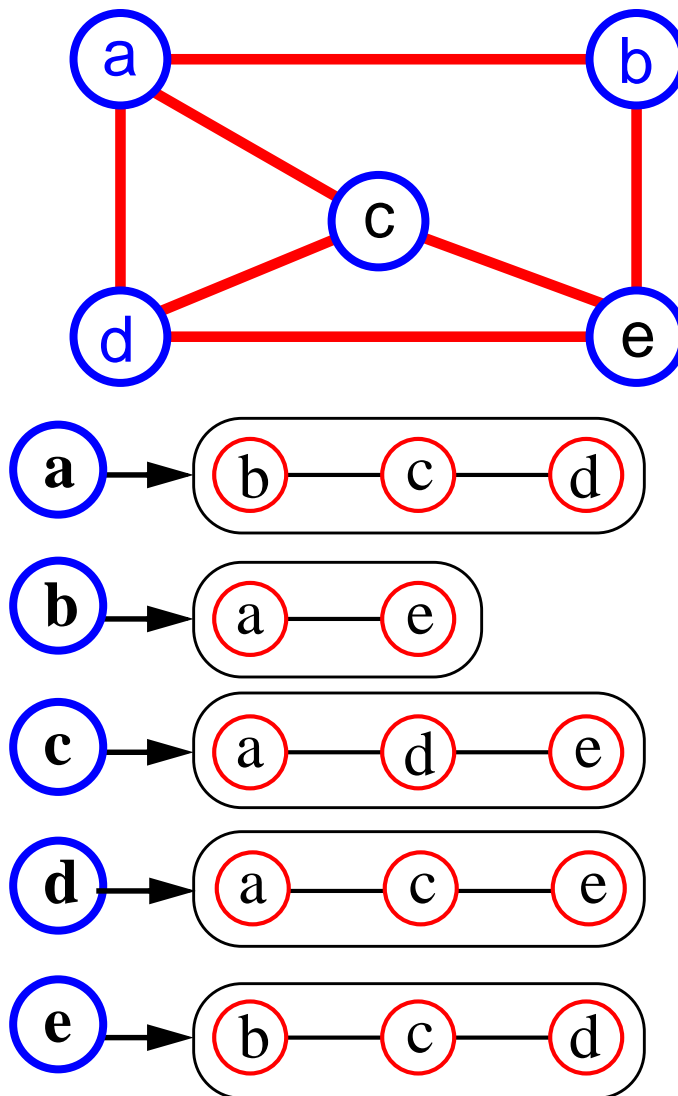


Performance of the Edge List Structure

Operation	Time
size, isEmpty, replaceElement, swap	$O(1)$
numVertices, numEdges	$O(1)$
vertices	$O(n)$
edges, directedEdges, undirectedEdges	$O(m)$
elements, positions	$O(n+m)$
endVertices, opposite, origin, destination, isDirected	$O(1)$
incidentEdges, inIncidentEdges, outIncidentEdges, adjacentVertices, inAdjacentVertices, outAdjacentVertices, areAdjacent, degree, inDegree, outDegree	$O(m)$
insertVertex, insertEdge, insertDirectedEdge, removeEdge, makeUndirected, reverseDirection, setDirectionFrom, setDirectionTo	$O(1)$
removeVertex	$O(m)$

Adjacency List (traditional)

- **adjacency list of a vertex v :**
sequence of vertices adjacent to v
- represent the graph by the adjacency lists of all the vertices

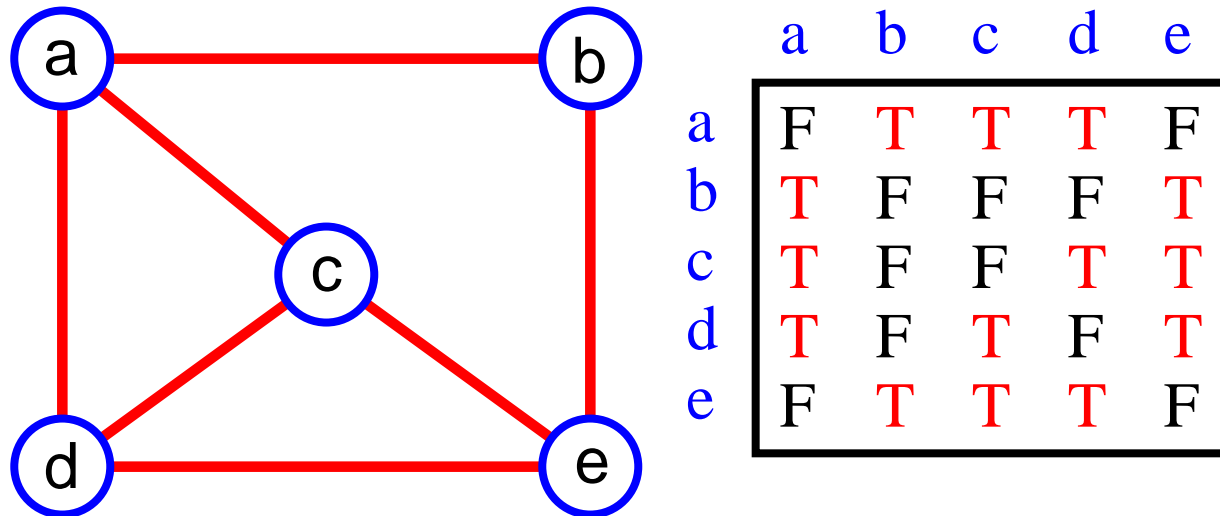


- Space = $\Theta(\mathbf{N} + \sum \text{deg}(\mathbf{v})) = \Theta(\mathbf{N} + \mathbf{M})$

Performance of the Adjacency List Structure

Operation	Time
size, isEmpty, replaceElement, swap	$O(1)$
numVertices, numEdges	$O(1)$
vertices	$O(n)$
edges, directedEdges, undirectedEdges	$O(m)$
elements, positions	$O(n+m)$
endVertices, opposite, origin, destination, isDirected, degree, inDegree, outDegree	$O(1)$
incidentEdges(v), inIncidentEdges(v), outIncidentEdges(v), adjacentVertices(v), inAdjacentVertices(v), outAdjacentVertices(v)	$O(\deg(v))$
areAdjacent(u, v)	$O(\min(\deg(u), \deg(v)))$
insertVertex, insertEdge, insertDirectedEdge, removeEdge, makeUndirected, reverseDirection,	$O(1)$
removeVertex(v)	$O(\deg(v))$

Adjacency Matrix (traditional)



- matrix M with entries for all pairs of vertices
- $M[i,j] = \text{true}$ means that there is an edge (i,j) in the graph.
- $M[i,j] = \text{false}$ means that there is no edge (i,j) in the graph.
- There is an entry for every possible edge, therefore:
 $\text{Space} = \Theta(N^2)$

Adjacency Matrix (

- The adjacency matrix structures augments the edge list structure with a matrix where each row and column corresponds to a vertex.

	0	1	2	3	4	5	6
0	∅	∅	NW 35	∅	DL 247	∅	∅
1	∅	∅	∅	AA 49	∅	DL 335	∅
2	∅	AA 1387	∅	∅	AA 903	∅	TW 45
3	∅	∅	∅	∅	∅	UA 120	∅
4	∅	AA 523	∅	AA 411	∅	∅	∅
5	∅	UA 877	∅	∅	∅	∅	∅
6	∅	∅	∅	∅	∅	∅	∅

BOS DFW JFK LAX MIA ORD SFO
 0 1 2 3 4 5 6

- The space requirement is $O(n^2 + m)$

Performance of the Adjacency Matrix Structure

Operation	Time
size, isEmpty, replaceElement, swap	$O(1)$
numVertices, numEdges	$O(1)$
vertices	$O(n)$
edges, directedEdges, undirectedEdges	$O(m)$
elements, positions	$O(n+m)$
endVertices, opposite, origin, destination, isDirected, degree, inDegree, outDegree	$O(1)$
incidentEdges, inIncidentEdges, outIncidentEdges, adjacentVertices, inAdjacentVertices, outAdjacentVertices,	$O(n)$
areAdjacent	$O(1)$
insertEdge, insertDirectedEdge, removeEdge, makeUndirected, reverseDirection, setDirectionFrom, setDirectionTo	$O(1)$
insertVertex, removeVertex	$O(n^2)$