


Data structures and analysis of algorithms (CS340a)  
Final Examination (Three Hour)

Problem 1. (10 marks) In big  $O$  notation for worst case, what is the time complexity of:  
a) mergesort.

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- c) Maximum-Cost Spanning Tree.
  - d) Single-Source Longest Paths in a weighted directed graph with non-positive weights.
  - e) All-pair shortest paths.

Problem 2. (5 marks) a). (4 marks) Prove, by using the definitions of  $O$  and  $\Omega$ , the following:

$$n = O(n - \log_2(n))$$
$$2^n - n = \Omega(2^n)$$

Problem 3. (10 marks) (Short answer. Your answer should be yes or no. If the answer is no, explain why)

- a) Dijkstra algorithm for Single-Source Shortest Paths works even if there are negative weights.
- b) Dijkstra method for finding Single-Source Shortest Paths can be used to find the Single-Source *Longest* Paths assuming that all the weights are non-negative.
- c) Minimum-Cost Spanning Tree for a graph will not change if *all* the weights are multiplied by a constant  $c$ .
- d) Graphs  $G_1$  and  $G_2$  have the same transitive closure iff  $G_1 = G_2$ .
- e) If 3SAT is not in  $P$  then Graph 3-coloring is in  $P$ .

Problem 4. (15 marks) A directed graph  $G = (V, E)$  is given in the figure.

- a) Show an adjacency list presentation of the graph.
- b) Find a DFS tree starting from vertex  $v$  using the adjacency list representation in a).
- c) Show the shortest-path tree from  $v$  to all the other vertices.
- d) Now assume that the same graph is undirected, show the minimum cost spanning tree of  $G$ .

Problem 5. (15 marks) Given a connected, undirected, weighted graph  $G = (V, E)$ , define the cost of a spanning tree to be the maximum weight among the weights associated with the edges of the spanning tree.

Design an efficient algorithm to find the spanning tree of  $G$  which maximize above defined cost. What is the complexity of your algorithm.

Problem 6. (15 marks) Let  $G = (V, E)$  be a connected, undirected, weighted graph with non-negative weights. Let  $v$  be a vertex of  $G$ . Given a spanning tree  $T$  of  $G$ , denote by  $c(v, w)$  the summation of the weights along the simple path from  $v$  to  $w$  on  $T$  ( $c(v, v) = 0$ ). We then define the cost of spanning tree  $T$  by

$$\text{cost}(T) = \max_{w \in V} c(v, w).$$

Design an efficient algorithm to find the spanning tree of  $G$  which minimize this cost.

[REDACTED]

Problem 8. (15 marks)

a) Let  $G = (V, E)$  be an undirected graph such that, for any vertex  $v$ , degree  $d(v)$  (number of edges incident to  $v$ ) is bounded by 3. Given such a graph and an integer  $k$  we want to determine if  $G$  contains a clique of size  $\geq k$ . Show that this problem is in  $P$ . Does this contradict to the fact that Clique Problem is  $NP$ -complete.

b) The composite numbers problem is defined as follows:

Given a positive integer  $K$ , determine if there are integers  $m, n > 1$  such that  $K = m \cdot n$ .

There is an  $O(\sqrt{K})$  time algorithm solving this problem. Is this problem in  $P$ ? Why?