

# Tutorial-2: Floating point and Computer Logic

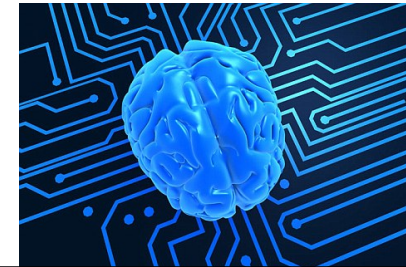
AbdulWahab Kabani  
PhD Student – Computer Science



[ 1 ]

## Why is this topic important?

- Understand Computer Circuits will help us understand how computers perform operations
- Build foundation for the next chapters



[ 2 ]

## Outline

- 32 IEEE Floating Point
- Computer Logic
  - Combinational Circuits vs. Sequential Circuits
  - Gates
  - Combinational Circuits
  - Sequential Circuits
- Conclusion



[ 3 ]

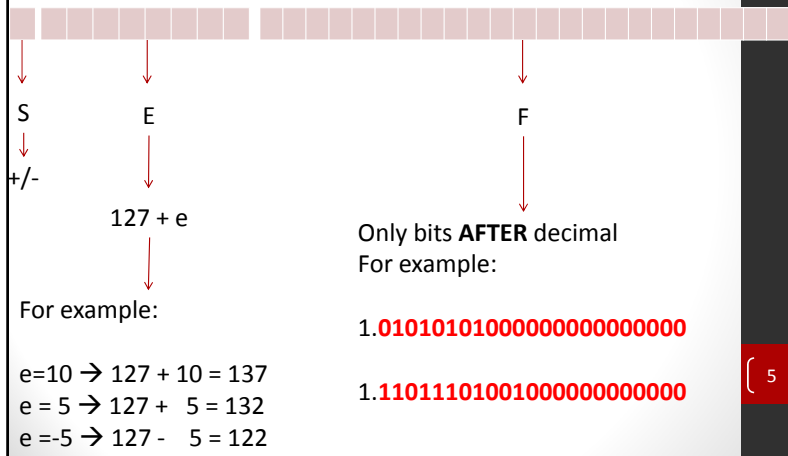
## Outline

- **32 IEEE Floating Point**
- Computer Logic
  - Combinational Circuits vs. Sequential Circuits
  - Gates
  - Combinational Circuits
  - Sequential Circuits
- Conclusion

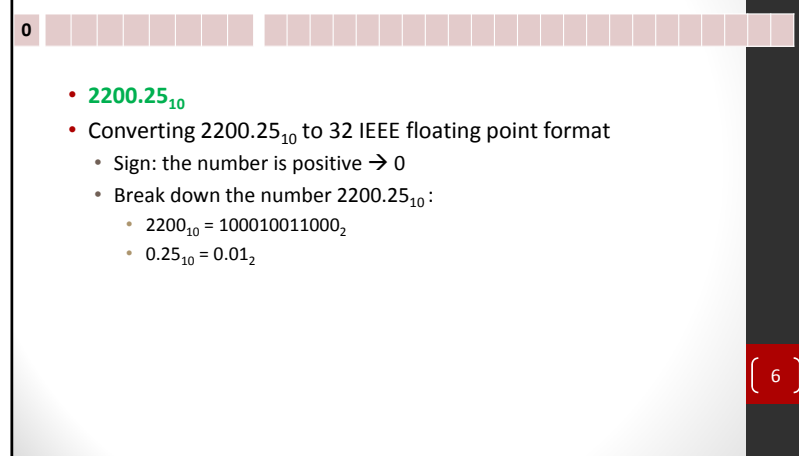


[ 4 ]

## IEEE Floating-Point (32 bit)



## IEEE Floating-Point (32 bit)



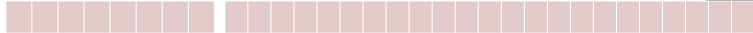
## IEEE Floating-Point (32 bit)

- $2200_{10} = 100010011000_2$ 
    - $2200/2 = 1100$  R = 0
    - $1100/2 = 550$  R = 0
    - $550/2 = 275$  R = 0
    - $275/2 = 137$  R = 1
    - $137/2 = 68$  R = 1
    - $68/2 = 34$  R = 0
    - $34/2 = 17$  R = 0
    - $17/2 = 8$  R = 1
    - $8/2 = 4$  R = 0
    - $4/2 = 2$  R = 0
    - $2/2 = 1$  R = 0
    - $1/2 = 0$  R = 1
- [ 7 ]

## IEEE Floating-Point (32 bit)

- $0.25_{10} = 0.01_2$ 
    - $0.25 * 2 = 0.5$  WP = 0
    - $0.5 * 2 = 1.0$  WP = 1
- [ 8 ]

## IEEE Floating-Point (32 bit)

0 

- **2200.25<sub>10</sub>**
- Converting 2200.25<sub>10</sub> to 32 IEEE floating point format
  - Sign: the number is positive → 0
  - Break down the number 2200.25<sub>10</sub>:
    - 2200<sub>10</sub> = 100010011000<sub>2</sub>
    - 0.25<sub>10</sub> = 0.01<sub>2</sub>
  - Combine Both numbers:
    - 100010011000.01<sub>2</sub>

[ 9 ]

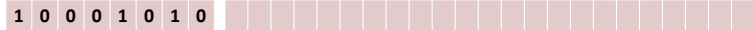
## IEEE Floating-Point (32 bit)

• **100010011000.01**  
 11 10 9 8 7 6 5 4 3 2 1

- Unbiased exponent = 11
- Remember our bias: 127
- 127+11 = 138

[ 10 ]

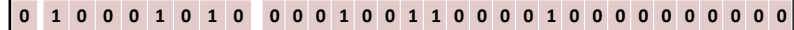
## IEEE Floating-Point (32 bit)

0 1 0 0 0 1 0 1 0 

- **2200.25<sub>10</sub>**
- Converting 2200.25<sub>10</sub> to 32 IEEE floating point format
  - Sign: the number is positive → 0
  - Break down the number 2200.25<sub>10</sub>:
    - 2200<sub>10</sub> = 100010011000<sub>2</sub>
    - 0.25<sub>10</sub> = 0.01<sub>2</sub>
  - Combine Both numbers:
    - 100010011000.01
  - 138 in binary is 10001010

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## IEEE Floating-Point (32 bit)

0 1 0 0 0 1 0 1 0 0 0 0 1 0 0 1 1 0 0 0 0 1 0 0 0 0 0 0 0 0 0 0 

- **2200.25<sub>10</sub>**
- Converting 2200.25<sub>10</sub> to 32 IEEE floating point format
  - Sign: the number is positive → 0
  - Break down the number 2200.25<sub>10</sub>:
    - 2200<sub>10</sub> = 100010011000<sub>2</sub>
    - 0.25<sub>10</sub> = 0.01<sub>2</sub>
  - Combine Both numbers:
    - 100010011000.01
  - 138 in binary is 10001010
  - Don't store the 1 at the beginning of 1.0001001100001

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## IEEE Floating-Point (32 bit)

- 0x7F400000<sub>16</sub>
- 7 F 4 0 0 0 0 0<sub>16</sub>
- 0111 1111 0100 0000 0000 0000 0000

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## IEEE Floating-Point (32 bit)

0

- 0x7F400000
- 7 F 4 0 0 0 0 0<sub>16</sub>
- 0111 1111 0100 0000 0000 0000 0000

[ 14 ]

## IEEE Floating-Point (32 bit)

0 1 1 1 1 1 1 0

- 0x7F400000
- 7 F 4 0 0 0 0 0<sub>16</sub>
- 0111 1111 0100 0000 0000 0000 0000

[ 15 ]

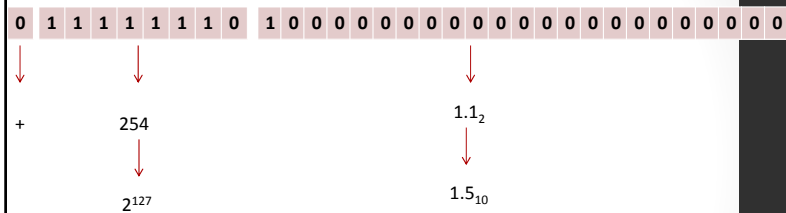
## IEEE Floating-Point (32 bit)

0 1 1 1 1 1 1 0 1 0

- 0x7F400000
- 7 F 4 0 0 0 0 0<sub>16</sub>
- 0111 1111 0100 0000 0000 0000 0000

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## IEEE Floating-Point (32 bit)



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## IEEE Floating-Point (32 bit)



- The number is:  $+1.5 \times 2^{127}$
- $2^{127} = 10^x$
- Solve for x by taking the log  $\rightarrow X = 38.23081$
- $2^{127} = 10^x = 10^{38.23081}$
- $= 10^{38} \times 10^{0.23081} = 10^{38} \times 1.7014$
- $+1.5 \times 2^{127} = 1.5 \times 10^{38} \times 1.7014 = +2.5521 \times 10^{38}$

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## Outline

- **32 IEEE Floating Point**
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  - Gates
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## Combinational vs. Sequential

- **Combinational :**
  - Gates are the building blocks
  - Output depends **only** on its current inputs



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## Combinational vs. Sequential

- **Sequential:**

- Flip-Flops and latches are the building blocks
- Output depends on current inputs + **Current State**



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## Combinational vs. Sequential

	Combinational	Sequential
<b>Building Block</b>	Gates	Flip-Flops and latches (set of gates)
<b>Depends on</b>	Input	Input + Current State



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## Outline

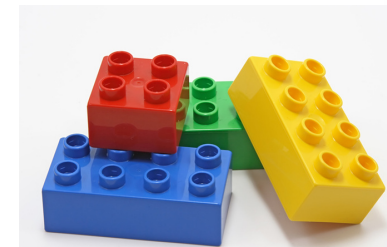
- **32 IEEE Floating Point**
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## Gates

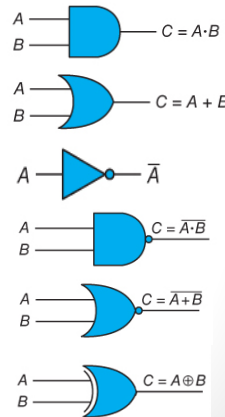
- Our Building Blocks
- Also, used to build flip-flops
- Come in many flavors:
  - AND
  - OR
  - NOT
  - NAND (Not AND)
  - NOR (Not OR)
  - XOR



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## Gates

- **AND:**
  - True only if both input are True
- **OR:**
  - True if at least one input is True
- **NOT:**
  - inverts the sign
- **NAND:**
  - The Opposite of an AND gate
- **NOR:**
  - The Opposite of an OR gate
- **XOR:**
  - True only if the inputs are different



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## Gates

- Time to use our blocks to construct digital circuits!



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## Outline

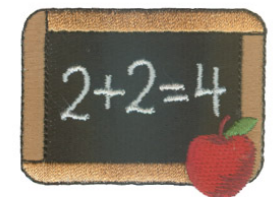
- **32 IEEE Floating Point**
- Computer Logic
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## Addition Circuit

- Let's construct a circuit that will perform addition

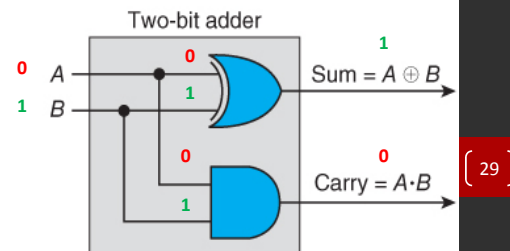


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## Half Adder

- A circuit that adds 1 bit to another
- All possible input/output:
  - $0+0=0$
  - $0+1=1$
  - $1+0=1$
  - $1+1=0$  (with a carry of 1)

A	B	Sum	Carry
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

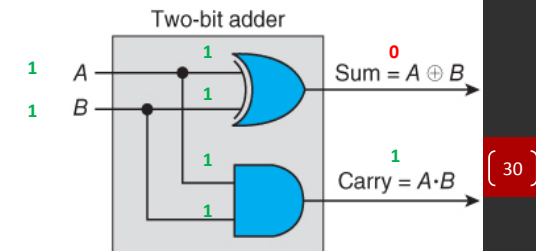


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## Half Adder

- A circuit that adds 1 bit to another
- All possible input/output:
  - $0+0=0$
  - $0+1=1$
  - $1+0=1$
  - $1+1=0$  (with a carry of 1)

A	B	Sum	Carry
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1



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## Half Adder

- What if we want to add three bits?
  - Use a Full Adder

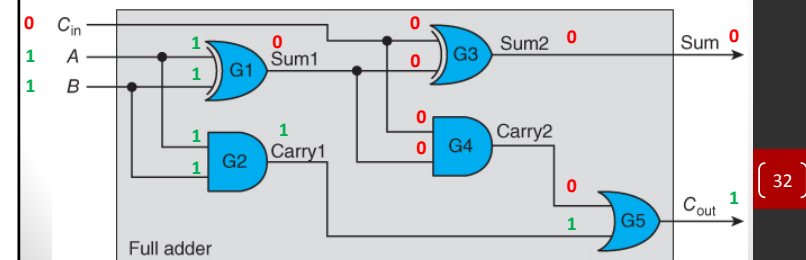


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## Full Adder

- 2 Half-Adders combined with an OR gate

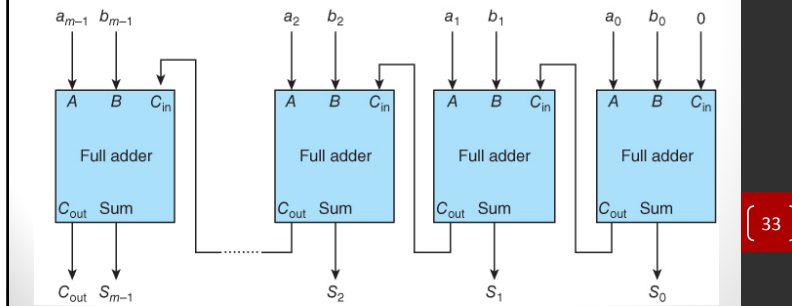
A	B	C	Sum	Carry
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1



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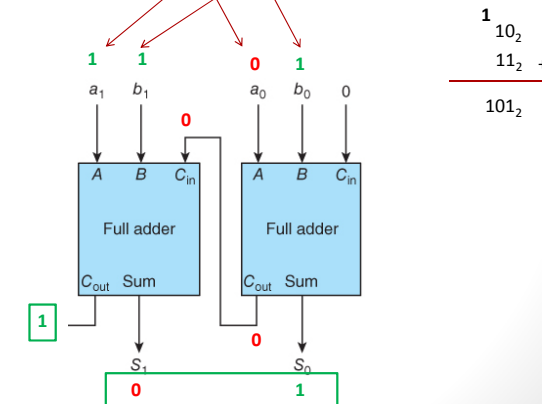


## Full Adder



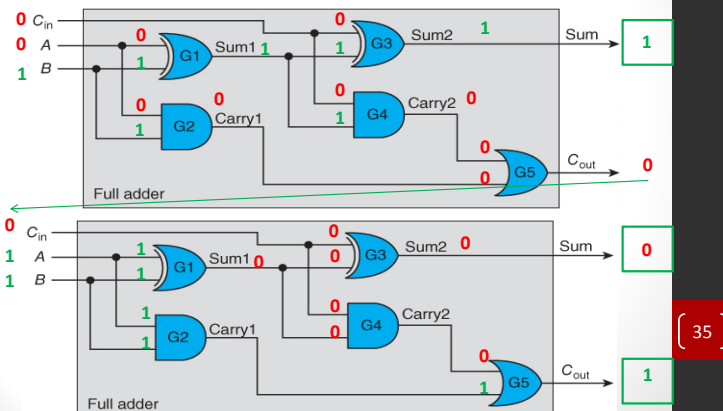
## Full Adder

- Say we have want to add  $10_2 + 11_2 = 2 + 3 = 5 = 101_2$



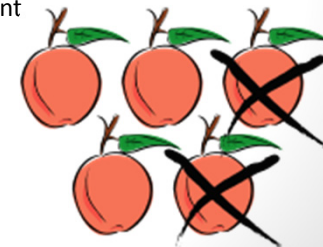
## Full Adder

- $10_2 + 11_2 = 101_2$



## Full Adder/Subtractor

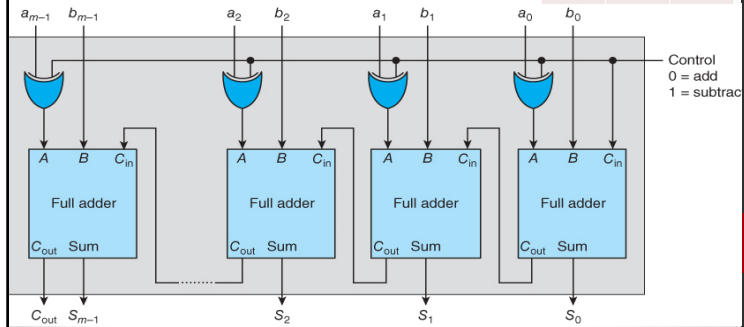
- How can we subtract?
  - Recall Two's complement
  - Subtraction is addition with the two's complement
  - Example:  $011_2 - 010_2 = 3 - 2 = 1$ 
    - N = 3 bits
    - Get the Two's complement
    - $100_2 - 010_2 = 110_2 = -2_{10}$
    - $011_2 + 110_2 = 001_2 = 1$



## Full Adder/Subtractor

- Why Does this work?
- Let's review the truth table of the XOR

A	B	XOR
0	0	0
0	1	1
1	0	1
1	1	0

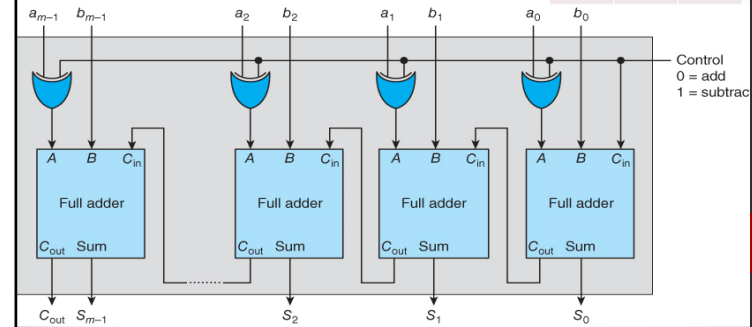


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## Full Adder/Subtractor

- Bit  $a_1$  will be flipped if the control is 1

A	B	XOR
0	0	0
0	1	1
1	0	1
1	1	0



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## Decoder

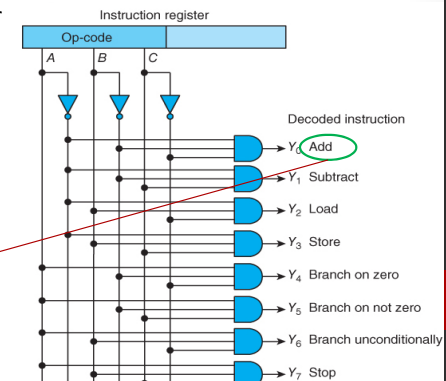
- Say we want a device that does a specific task if the user enters a certain code
- Something like a remote control. You enter a code and go to channel. Few inputs, lots of output (channels)
- Or maybe, Bar code: few input, lots of products
- Computer Instructions



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## Decoder

- In all cases, you go from a small set of possible inputs to a large set of possible outputs.
- Example: 3-8 Decoder
  - 3 input channels
  - 8 possible outputs
  - 3 input  $\rightarrow 2^3$  outputs



You can put a set of full adders here to create an addition circuit

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## Sequential Circuits

- Flip-Flops and latches are the building blocks
- Output depends on current inputs + **Current State**



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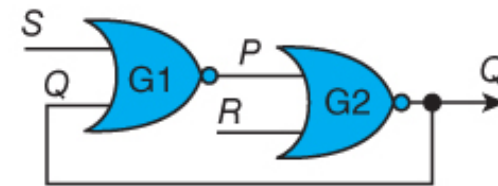
## Sequential Circuits

- Three basic types:
  - RS latches
  - D flip-flops
  - JK flip-flops
- These types are the building blocks of sequential Circuits
- What are they composed of?

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## RS latches

- Convention:
  - Q is the current state
  - $Q^+$  is the next (new) state
  - R is reset
  - S is set
- This circuit employs *feedback*
  - The value of p is determined by Q (current state)



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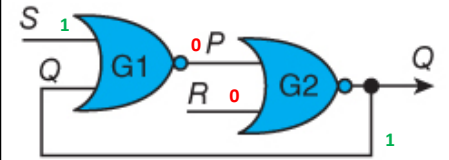
## RS latches

- In Sequential Circuits, it is useful to assume an initial state
- Don't try to think how the previous state was generated



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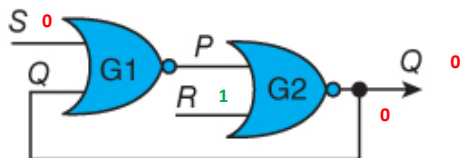
## RS latches



Inputs			Output
R	S	Q	Q <sup>+</sup>
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	?
1	1	1	?

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## RS latches



Inputs			Output
R	S	Q	Q <sup>+</sup>
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	?
1	1	1	?

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## RS latches

Inputs		Output	Description
R	S	Q <sup>+</sup>	
0	0	Q	No change
0	1	1	Set output to 1
1	0	0	Reset output to 0
1	1	X	Forbidden

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## RS latches

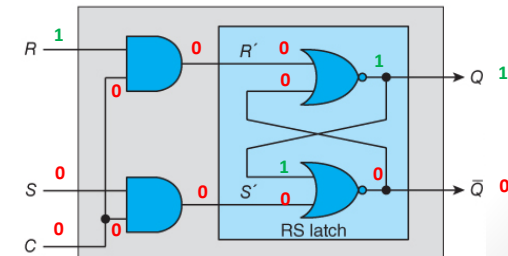
- How can we control **when** to set/reset an RS latch?
  - Use a clock
  - If the clock is **0**, nothing happens to the latch
  - If the clock is **1**, the latch can be set/reset/maintain its state



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## Clocked RS Flip-Flop

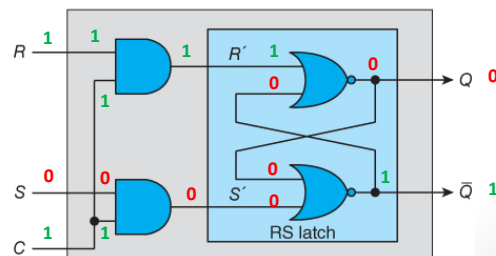
- When clock = 0, nothing happens to the latch
- Assume Q = 1



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## Clocked RS Flip-Flop

- When clock = 1, nothing happens to the latch
- Assume Q = 1



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## What is the difference between a latch and flip-flop?

- Both are bi-stable
- Flip-Flops are **clocked**



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## D flip-flops

- Similar to the clocked RS flip-flops
- However, no Reset/set Inputs
- Inputs:
  - D: Data
  - Clock
- Similar to a dam



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## D flip-flops

- If clock (C) is 0, no data (D) flows
- If clock (C) is 1, data (D) flows through the flipflop
- Unlike dams, the flow is either 0, or 1 (no quantity concept)



[ 54 ]

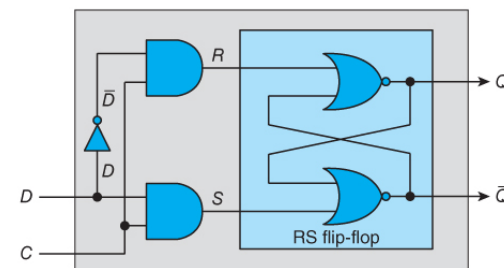
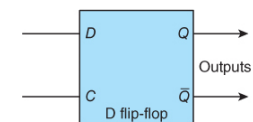
## D flip-flops

- How to prepare a D-Flip-Flops?
- You need:
  - RS flip-flop
  - 2 AND gates
  - D (Data) and its opposite (inverter)



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## D flip-flops



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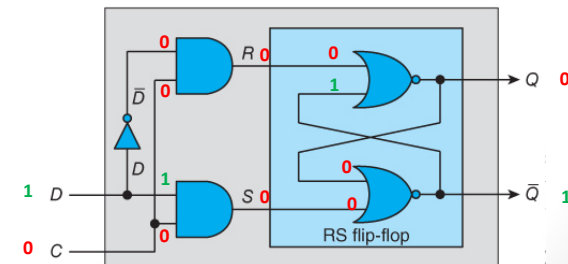
## D flip-flops

Inputs		Output	Description
C	D	$Q^+$	
0	0	Q	No change
0	1	Q	No change
1	0	0	D is copied to output
1	1	1	D is copied to output

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## D flip-flops

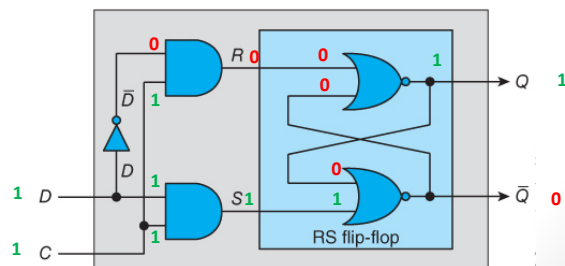
- When clock is 0, the state remain the state
- Assume  $Q_{\text{initial}} = 0$



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## D flip-flops

- When clock is 1, the state changes and have the input of D
- Assume  $Q_{\text{initial}} = 0$



[ 59 ]

## JK Flip-Flop

Inputs		Output	Description
R	S	$Q^+$	
0	0	Q	No change
0	1	1	Set output to 1
1	0	0	Reset output to 0
1	1	X	Forbidden

Inputs		Output	Description
K	J	$Q^+$	
0	0	Q	No change
0	1	1	Set output to 1
1	0	0	Reset output to 0
1	1	$Q'$	Opposite State

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  - **Gates**
  - **Combinational Circuits**
  - **Sequential Circuits**
- **Conclusion**



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## Conclusion

- Introduced an Example of the 32 bit IEEE Floating point format
- Talked about digital circuits
- Combinational and Sequential
- Should be able to understand a circuit



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## Help

- **Graduate Teaching Assistants:**
  - Abdulwahab Kabani, [akabani5@uwo.ca](mailto:akabani5@uwo.ca)
  - Mike Molnar, [mmolnar2@uwo.ca](mailto:mmolnar2@uwo.ca)
  - Sakif Pritom, [spritom@uwo.ca](mailto:spritom@uwo.ca)
- Send us an email or see us after class
- Consultation Room : Middlesex College - 342



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