

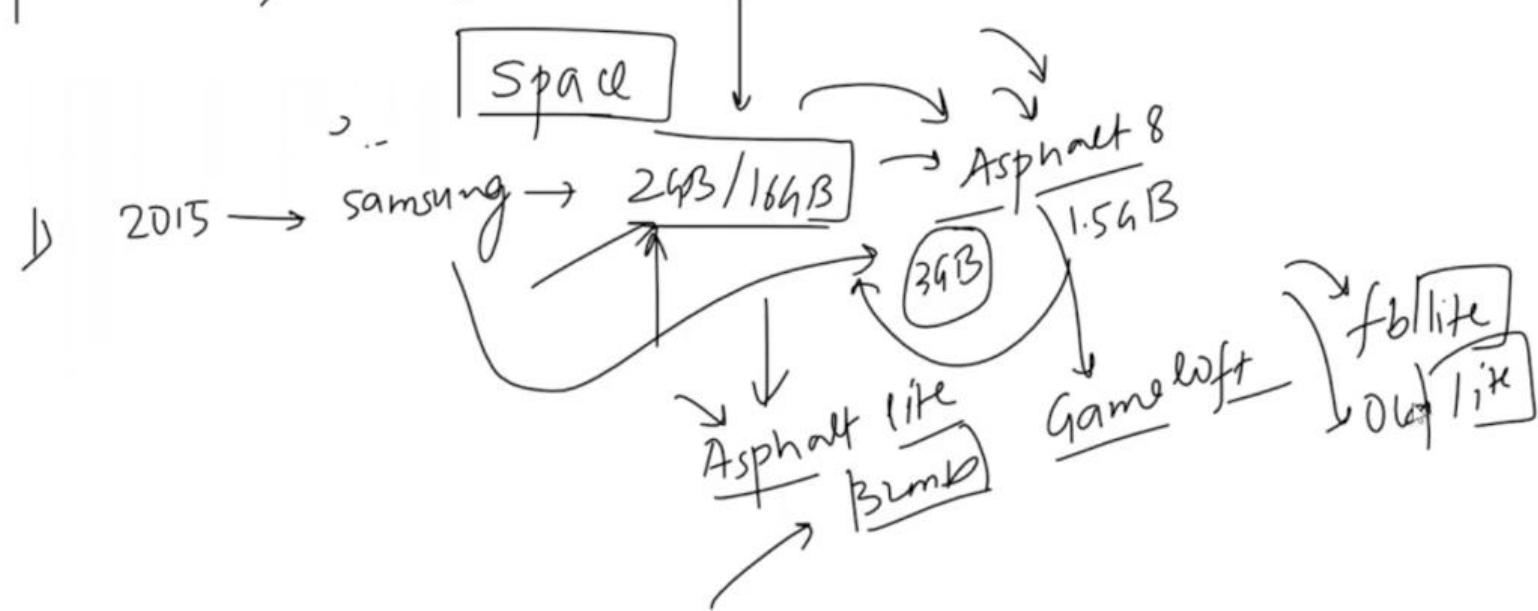


What is efficiency in programming?



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+ Prachile



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<https://en.wikipedia.org/wiki/Kolkata>

Kolkata - Wikipedia

According to the 2011 Indian census, **Kolkata** is the seventh-most populous city in India, with a population of 45 lakh (4.5 million) residents within the city ...

State: West Bengal Country: India
GDP/PPP: \$150.1 billion(GDP PPP 2022) Telephone code: +91 33
Kolkata district · Kolkata metropolitan area · Culture of Kolkata · Port of Kolkata



Kolkata

City in West Bengal

Kolkata (formerly Calcutta) is the capital of India's West Bengal state. Founded as an East India Company trading post, it was India's capital under the British Raj from 1773–1911. Today it's known for its grand colonial architecture, art galleries and cultural festivals. It's also home to Mother House, headquarters of the Missionaries of Charity, founded by Mother Teresa, whose tomb is on site. — Google

Area: 206.1 km²

Elevation: 9.14 m

Weather: 29 °C, Wind S at 13 km/h, 81% Humidity weather.com

Population: 1.49 crores (2020)

Local time: Saturday, 7:13 pm

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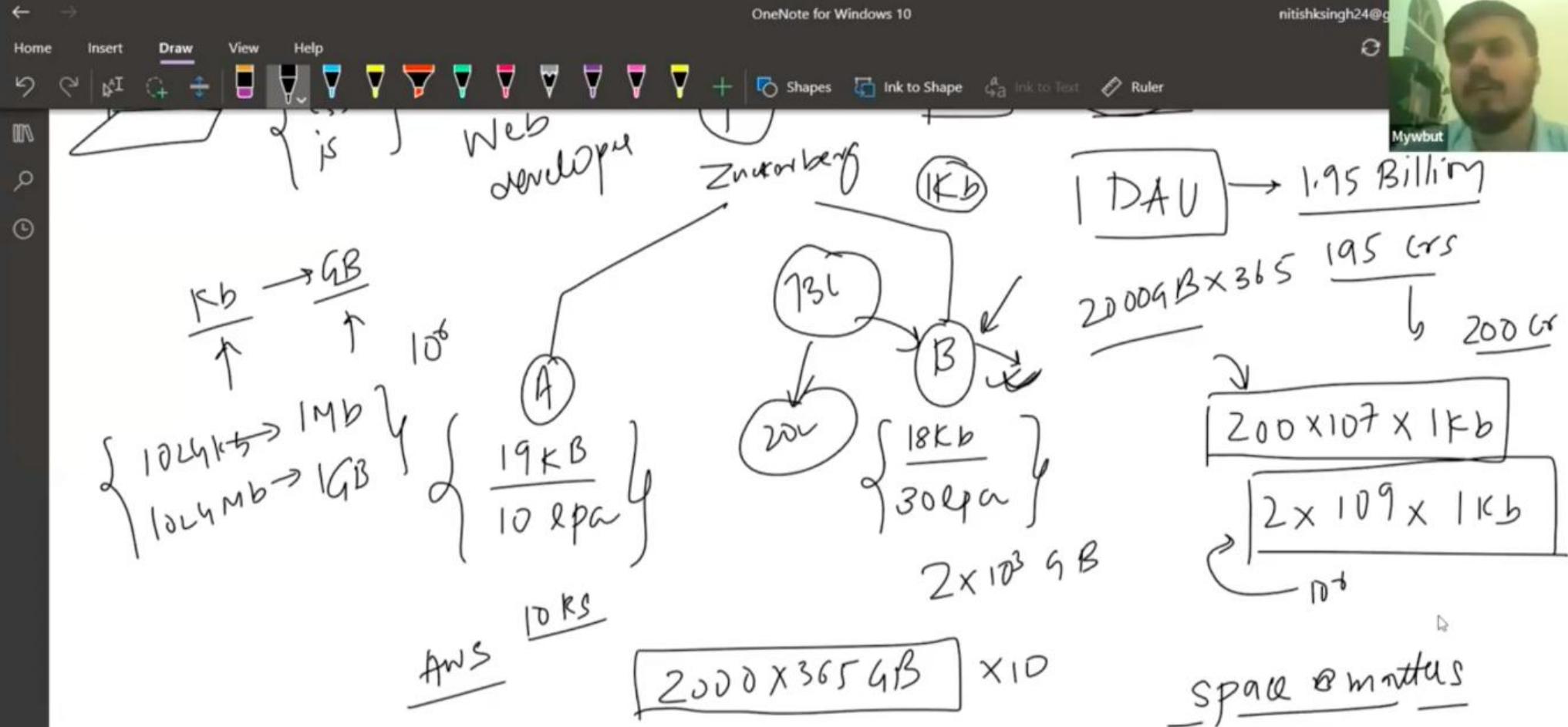


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Our focus - Time





Techniques

1. Measuring **time** to execute
2. **Counting** operations involved
3. Abstract notion of **order of growth**





1. Measuring Time



etheth.py - C:\Users\91842\Desktop\etheth.py (3.6.0)

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```
import time
```

```
start = time.time()
```

```
for i in range(1,101):
```

```
    print(i)
```

```
print(time.time() - start)]
```



I

Python 3.6.0 Shell

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```
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98  
99  
100
```

```
0.5520379543304443
```

```
>>>
```



Data Structures and Algorithms using Python | Mega Video | DSA in Python in 1 video



etheth.py - C:\Users\91842\Desktop\etheth.py (3.6.0)
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```
7 import time  
8  
7 start = time.time()  
8 i = 1  
8 while i<10:  
8     print(i)  
8     i+=1  
8  
8 print(time.time() - start)
```

0
>



26:17 / 11:41:29 • How to calculate order of growth >



Python 3.6.0 Shell

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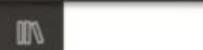
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0.4904956817626953

>>>



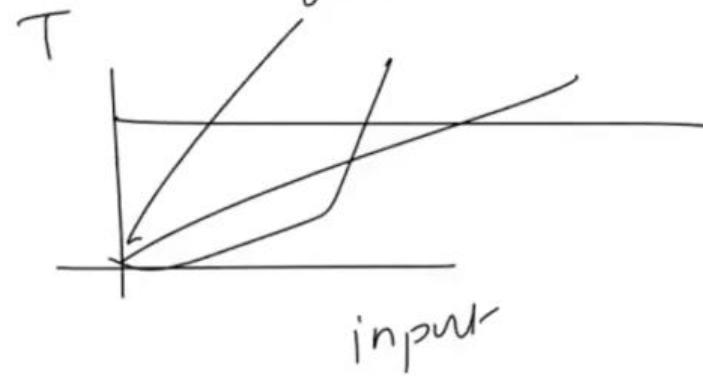
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Ans

2000 x 365 GB

 x 10space @ minutes

T → input
↓ relationship





Problems with this approach

Different time for different algorithm	✓
Time varies if implementation changes	✗
Different machines different time	✗
Does not work for extremely small input	✗
Time varies for different inputs, but can't establish a relationship	✗





2. Counting Operations



COUNTING OPERATIONS



- assume these steps take **constant time**:
 - mathematical operations
 - comparisons
 - assignments
 - accessing objects in memory
- then count the number of operations executed as function of size of input

```
def c_to_f(c):  
    return c*9.0/5 + 32  
  
def mysum(x):  
    total = 0  
    for i in range(x+1):  
        total += i  
    return total
```

Annotations on the code:

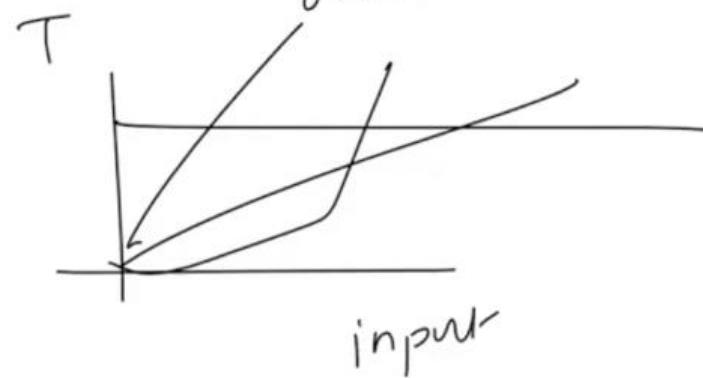
- Red box around `c*9.0/5 + 32` labeled **3 ops**
- Red box around `total = 0` labeled **1 op**
- Red box around `for i in range(x+1):` labeled **loop x times**
- Red box around `total += i` labeled **1 op**
- Red box around `return total` labeled **1 op**
- Red box around `mysum → 1+3x ops`



Home Insert Draw View Help

Ans
$$\boxed{2000 \times 365 \text{ GB}}$$
 $\times 10$ space @ minutes

T \rightarrow input
 \downarrow relationship



$$\boxed{1 + 3x} \rightarrow \begin{array}{ll} \text{input} & \\ x = 10 & \\ 31 & x = 20 \\ 61 & x = 30 \\ 91 & \end{array}$$



Problems with this approach

Different time for different algorithm	✓
Time varies if implementation changes	✗
Different machines different time	✓
No clear definition of which operation to count	✗
Time varies for different inputs, but can't establish a relationship	✓





What do we want

1. We want to evaluate the algorithm
2. We want to evaluate scalability
3. We want to evaluate in terms of input size



DIFFERENT INPUTS CHANGE HOW THE PROGRAM RUNS



- a function that searches for an element in a list

```
def search_for_elmt(L, e):  
    for i in L:  
        if i == e:  
            return True  
    return False
```

- when e is **first element** in the list \rightarrow BEST CASE
- when e is **not in list** \rightarrow WORST CASE
- when **look through about half** of the elements in list \rightarrow AVERAGE CASE





3. Orders of Growth



ORDERS OF GROWTH



Goals:

- want to evaluate program's efficiency when **input is very big**
- want to express the **growth of program's run time** as input size grows
- want to put an **upper bound** on growth – as tight as possible
- do not need to be precise: "**order of**" not "**exact**" growth
- we will look at **largest factors** in run time (which section of the program will take the longest to run?)
- **thus, generally we want tight upper bound on growth, as function of size of input, in worst case**

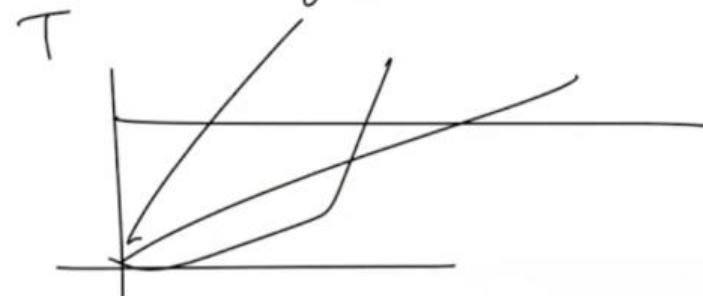


Ans

$$[2000 \times 365 \text{ GB}]$$

 $\times 10$ space @ minutes

T

 $T \rightarrow \text{input}$
 \downarrow
relationship $T \rightarrow \text{input}$
 \rightarrow
mathematical

input

$$[1 + 3x] \rightarrow \text{input}$$

$x = 10$

$$\downarrow$$

 $31 \quad x = 20$

$$61 \quad x = 30$$

 \downarrow

$$T = 1 + 3x \rightarrow \text{input}$$



MEASURING ORDER OF GROWTH: BIG OH NOTATION



- Big Oh notation measures an **upper bound on the asymptotic growth**, often called order of growth
- **Big Oh or $O()$** is used to describe worst case
 - worst case occurs often and is the bottleneck when a program runs
 - express rate of growth of program relative to the input size
 - evaluate algorithm **NOT** machine or implementation



EXACT STEPS vs O()



```
def fact_iter(n):
    """assumes n an int >= 0"""
    answer = 1
    while n > 1:
        answer *= n
        n -= 1
    return answer
```

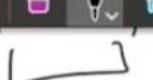
answer = answer * n
temp = n-1
n = temp
 $1 + 5n + 1$

- computes factorial
- number of steps:
- worst case asymptotic complexity: $O(n)$
 - ignore additive constants
 - ignore multiplicative constants



[←](#) [→](#)

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mathematics

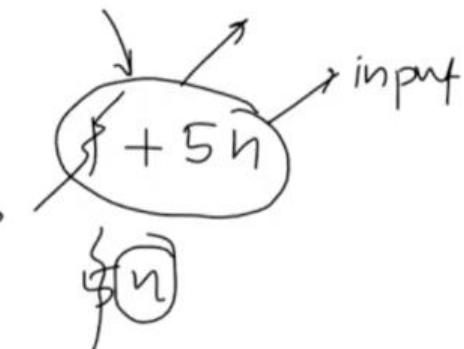
$$T = i^2 + 2i + 2$$



✓ { for X
 =

 for
 for
 for

O(1) big Oh
 α)
f(n)



So the idea is simple

$$n^2 + 2n + 2$$

$$n^2 + 100000n + 3^{1000}$$

$$\log(n) + n + 4$$

$$0.0001 * n * \log(n) + 300n$$

$$2n^{30} + 3^n$$





\cancel{K} - n

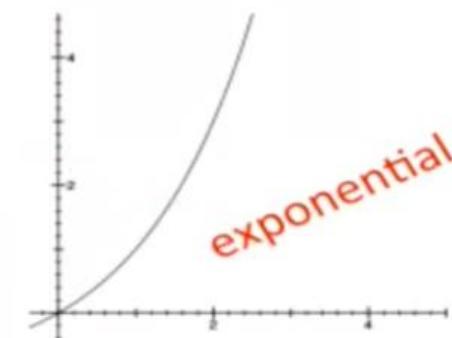
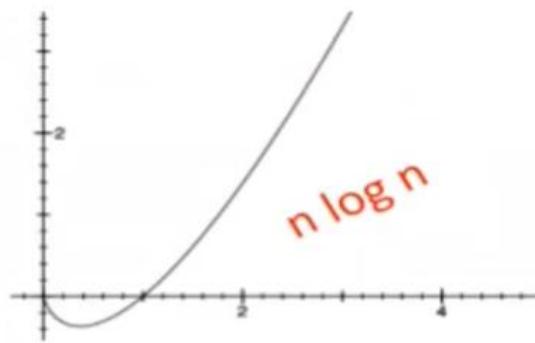
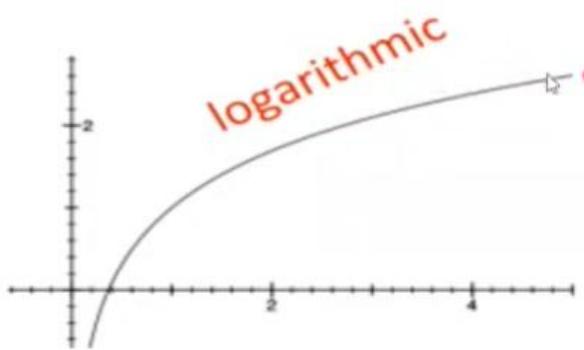
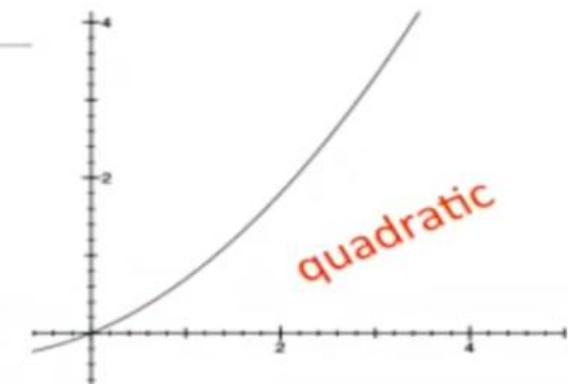
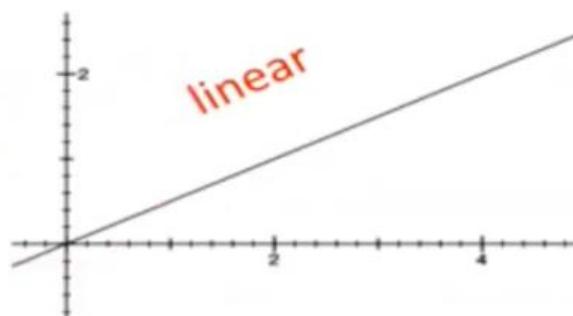
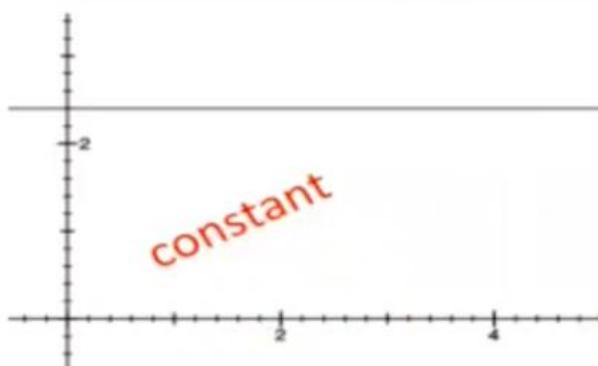
$$n^2 + 2n + \cancel{\{ } \quad n^2 + \cancel{\} } n$$

$$\overbrace{n^2 + n}^{2} \quad n^2$$

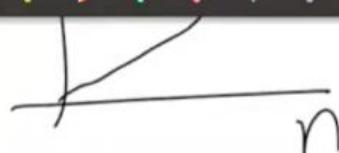
$O(n^2)$



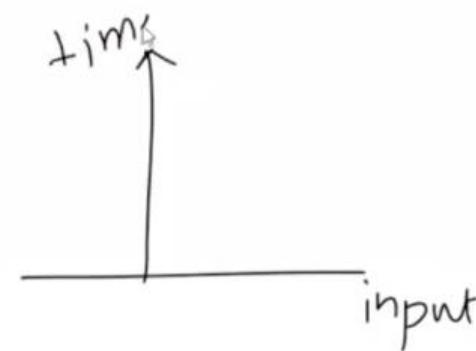
TYPES OF ORDERS OF GROWTH

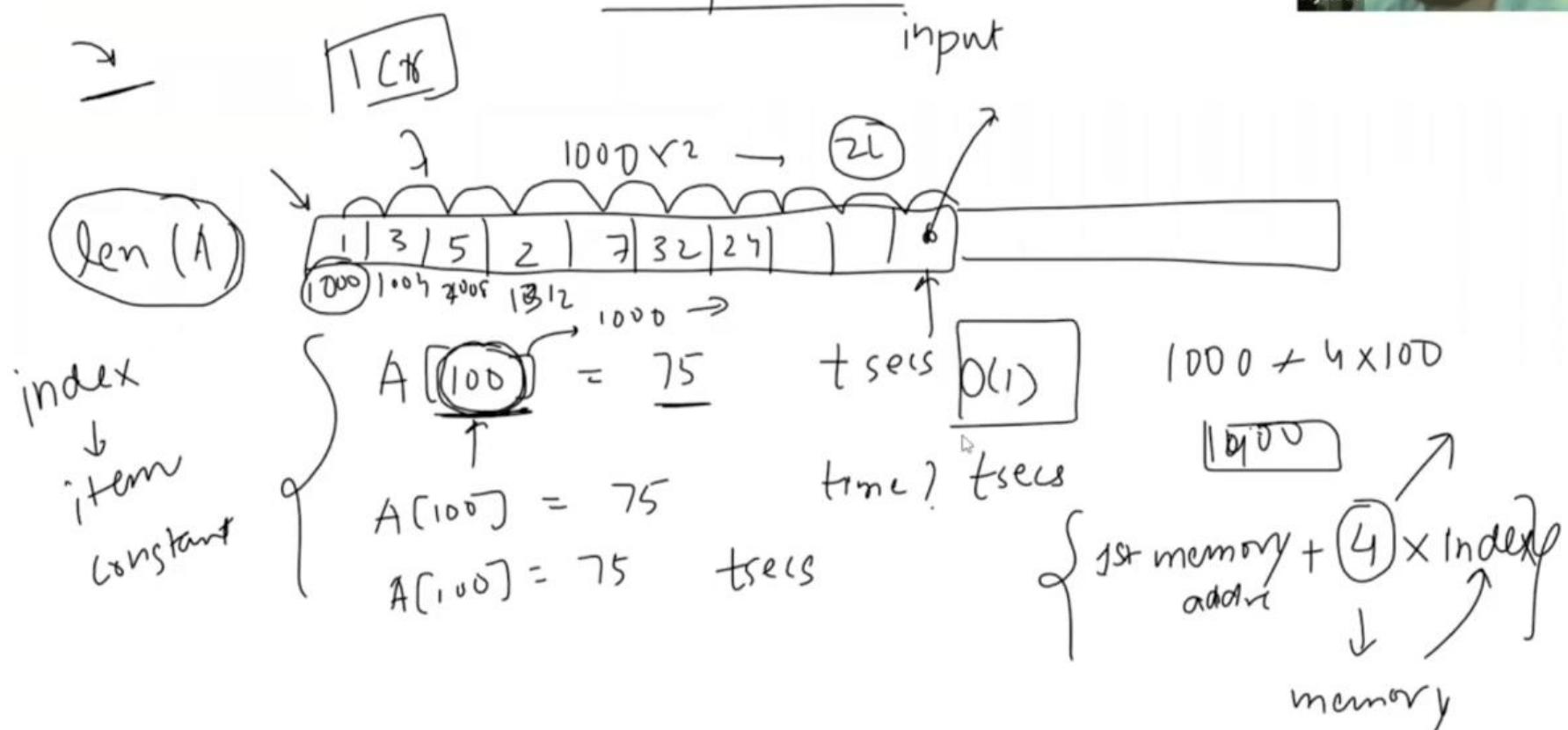


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n





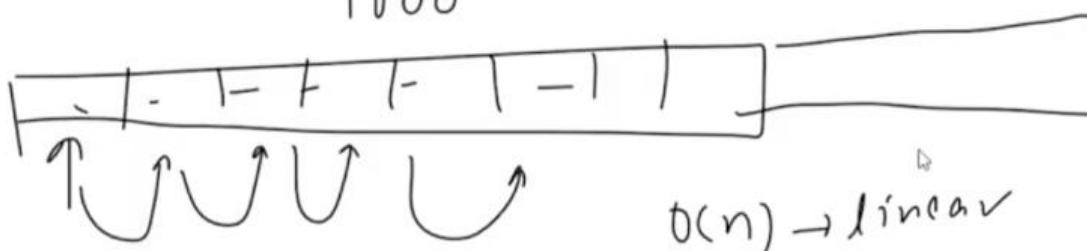


Linear search

1000m

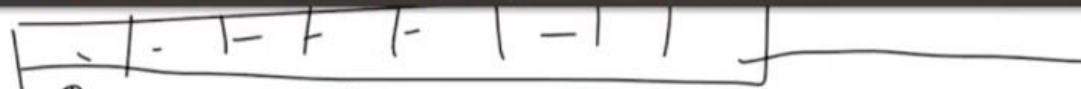
T →

36





36

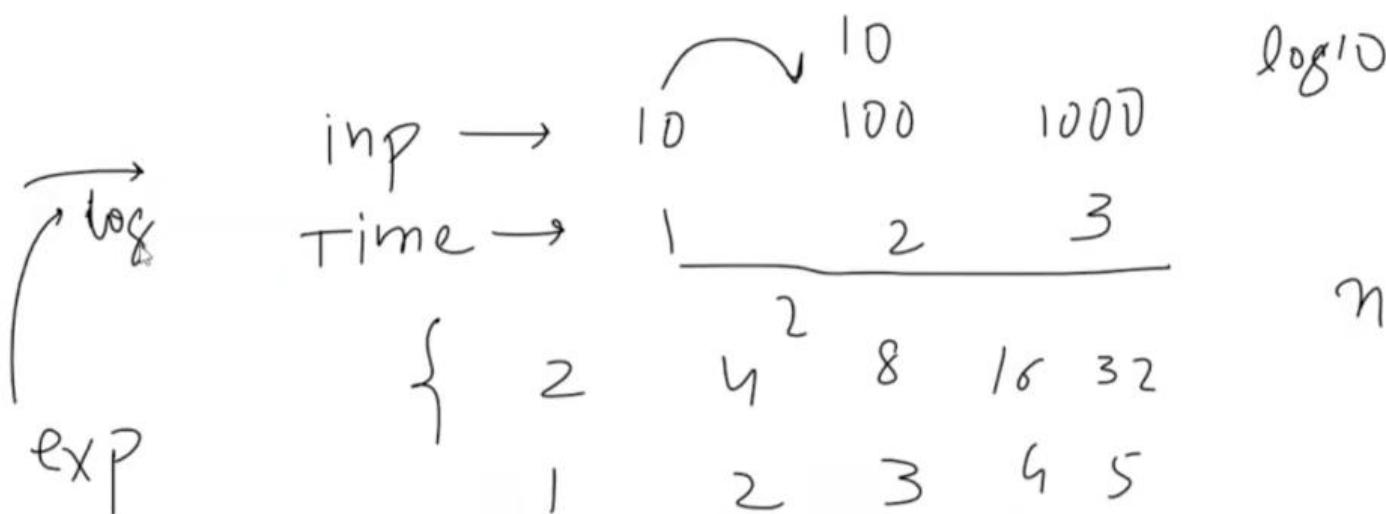
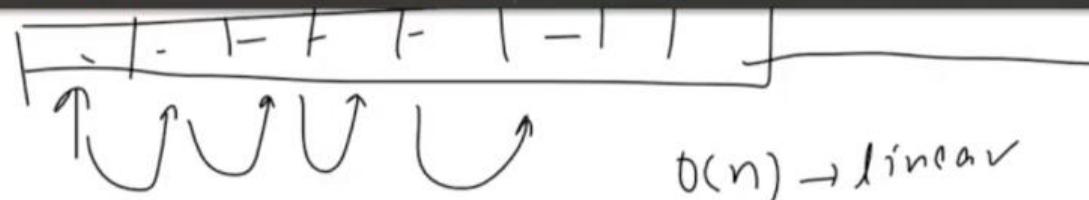


$\uparrow \downarrow \uparrow \downarrow \uparrow \downarrow$ $O(n) \rightarrow \text{linear}$

inp →	10	100	1000	$\log 10$
Time →	1	2	3	
	2	4	8	16 32
	1	2	3	4 5



36



$$n > n \log n > n^2$$

inp ²	1	2	3	4
T	1	10	100	1000

Time complexity

n, log n, n log n, n^2, 2^n

O(1) > log n > n > n log n > n^2 > 2^n

Complexity Growth



CLASS	n=10	= 100	= 1000	= 10000	= 100000
O(1)	1	1		1	1
O(log n)	1	2		3	6
O(n)	10	100		1000	1000000
O(n log n)	10	200		3000	6000000
O(n^2)	100	10000		1000000	1000000000000
O(2^n)	1024	12676506 00228229 40149670 3205376	1071508607186267320948425049060 0018105614048117055336074437503 8837035105112493612249319837881 5695858127594672917553146825187 1452856923140435984577574698574 8039345677748242309854210746050 6237114187795418215304647498358 1941267398767559165543946077062 9145711964776865421676604298316 52624386837205668069376		Good luck!!



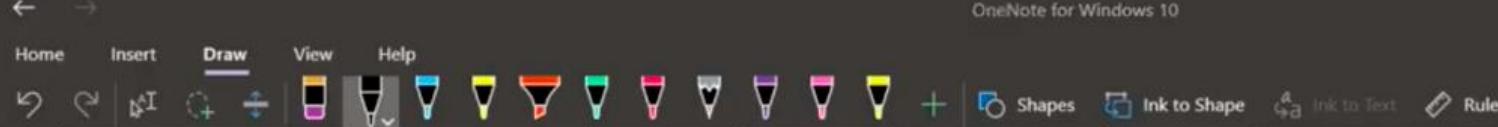
```
*python-codes.py - C:\Users\91842\Desktop\python-codes.py (3.6.0)*
File Edit Format Run Options Window Help
#####
##### Problem 1 #####
CL L = [1,2,3,4,5]
O sum = 0
for i in L:
    sum = sum + i
print(sum)
O product = 1
for i in L:
    product = product*i
print(product)
#####
##### Problem 2 #####
O L = [1,2,3,4,5]
for i in L:
    for j in L:
        print('({},{})'.format(i,j))
#####
##### Problem 4 #####
Linear Search
```



1
6
1000000
6000000
0000000

od luck!!





T I b 100 1000

$$\left\{ \begin{array}{l} O(n) \\ + \\ O(n) \end{array} \right\}_n O(n) + O(n) = O(n+n) \\ = O(2n) \\ \boxed{= O(n)}$$

← → *python-codes.py - C:\Users\91842\Desktop\python-codes.py (3.6.0)*
Home Insert File Edit Format Run Options Window Help
L = [1,2,3,4,5]



```
sum = 0
for i in L:
    sum = sum +i
print(sum)
```

```
product = 1
for i in L:
    product = product*i
print(product)
```

```
##### Problem 2 #####

```

```
L = [1,2,3,4,5]
```

```
for i in L:
    for j in L:
        print('({},{})'.format(i,j))
```

```
##### Problem 4 #####
Linear Search
```

```
##### Problem 5 #####

```



T I b 100 1000

$$\left\{ \begin{array}{l} O(n) \\ + \\ O(n) \end{array} \right. n \quad \begin{array}{l} O(n) + O(n) = O(n+n) \\ = O(2n) \\ = O(n) \end{array}$$

$O(n)$ \times $O(n)$ $O(n) \times O(n) = O(n \times n) = O(n^2)$

↑
inf

A hand-drawn diagram illustrating the addition of two linear functions. Two straight lines are drawn, both labeled $O(n)$. They intersect at a point labeled n . The equation $O(n) + O(n) = O(n+n) = O(2n) = O(n)$ is written next to the lines. Below this, another hand-drawn diagram shows two curves representing quadratic functions, labeled $O(n^2)$, with an arrow pointing upwards towards infinity, labeled 'inf'.

python-codes.py - C:\Users\91842\Desktop\python-codes.py (3.6.0)
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for j in L:

 print('{},{})'.format(i,j))



Problem 4

Linear Search

Problem 5

```
def intToStr(i):
    digits = '0123456789'
    if i == 0:
        return '0'
    result = ""
    while i > 0:
        result = digits[i%10] + result
        i = i//10
    return result
```

Problem 6

n = 1000

Home Insert Draw View Help



Shapes Ink to Shape Ink to Text Ruler



inf

logn

	10 times				
<u>inp</u>	<u>123</u>	1230	12300	123000	
T	3 ↓ 1+	4 ↓ 1+	5 ↓ 1+	6 ↓ 1+	

← → *python-codes.py - C:\Users\91842\Desktop\python-codes.py (3.6.0)*
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Home Insert
undo redo |
i = i//10
return result



Problem 6

```
n = 1000
int i,j,k = 0
2,4,8(10)
2,4,8,16,32,64(100) I
2,4,8,16,32,64,128,256,512(1000)
```

```
for(i=n/2;i<=n;i++){
    for(j=2;j<=n;j=j*2){
        k=k+n/2
    }
}
```

Problem 7

Binary Search



inf

log

10 times

$$\begin{array}{c} \text{inp} \\ \hline 123 \\ T \\ \hline 3 \\ \uparrow 1+ \end{array} \quad \begin{array}{c} 1230 \\ 4 \\ \downarrow \\ 1+ \end{array} \quad \begin{array}{c} 12300 \\ 5 \\ \downarrow \\ 1+ \end{array} \quad \begin{array}{c} 123000 \\ 6 \\ \downarrow \\ 1+ \end{array}$$

{ log } → mult

$$\begin{array}{r} \text{inp} \quad 10 \\ \text{out} \quad 3 + ? \\ \hline \text{add} \end{array} \quad \begin{array}{r} 100 \\ 6 + 3 \\ \hline 9 \end{array}$$



out
3 + 5
6 + 5
9
→ add



$$O\left(\frac{n}{2}\right)$$

x

$$O(10^m)$$

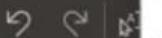
$$O\left(\frac{n}{2} \times 10^m\right)$$

$$= O\left(\frac{1}{2} n 10^m\right)$$

$$O(n 10^m)$$

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```
2,4,8(10)
2,4,8,16,32,64(100)
2,4,8,16,32,64,128,256,512(1000)
```

```
for(i=n/2;i<=n;i++){
    for(j=2;j<=n;j=j*2){
        k=k+n/2
    }
}
```

```
##### Problem 7 #####
Binary Search
```

```
##### Problem 8 #####
```

```
L = [1,2,3,4,5]
```

```
for i in range(0,len(L)):
    for j in range(i+1,len(L)):
        print('{},{}'.format(L[i],L[j]))
```



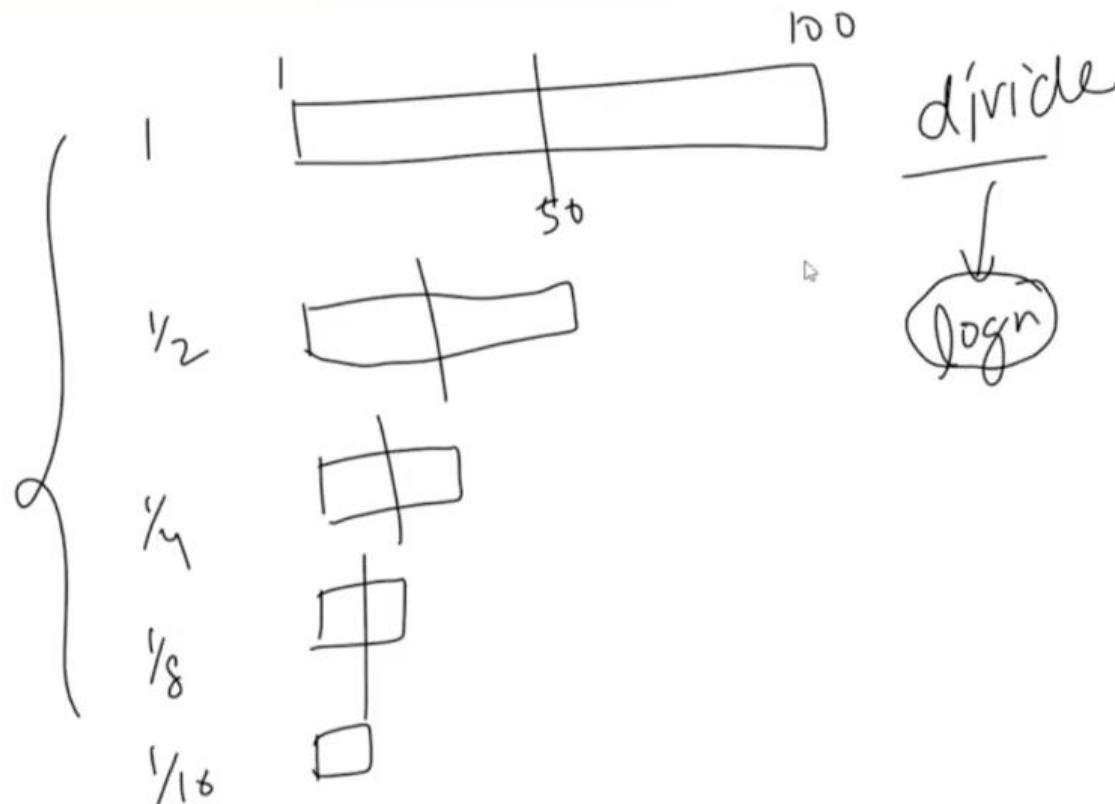


← →

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undo redo | Shapes Ink to Shape Ink to Text Ruler

|

 $O(n \log n)$ 

```
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File Edit Format Run Options Window Help
Home Insert
}
}

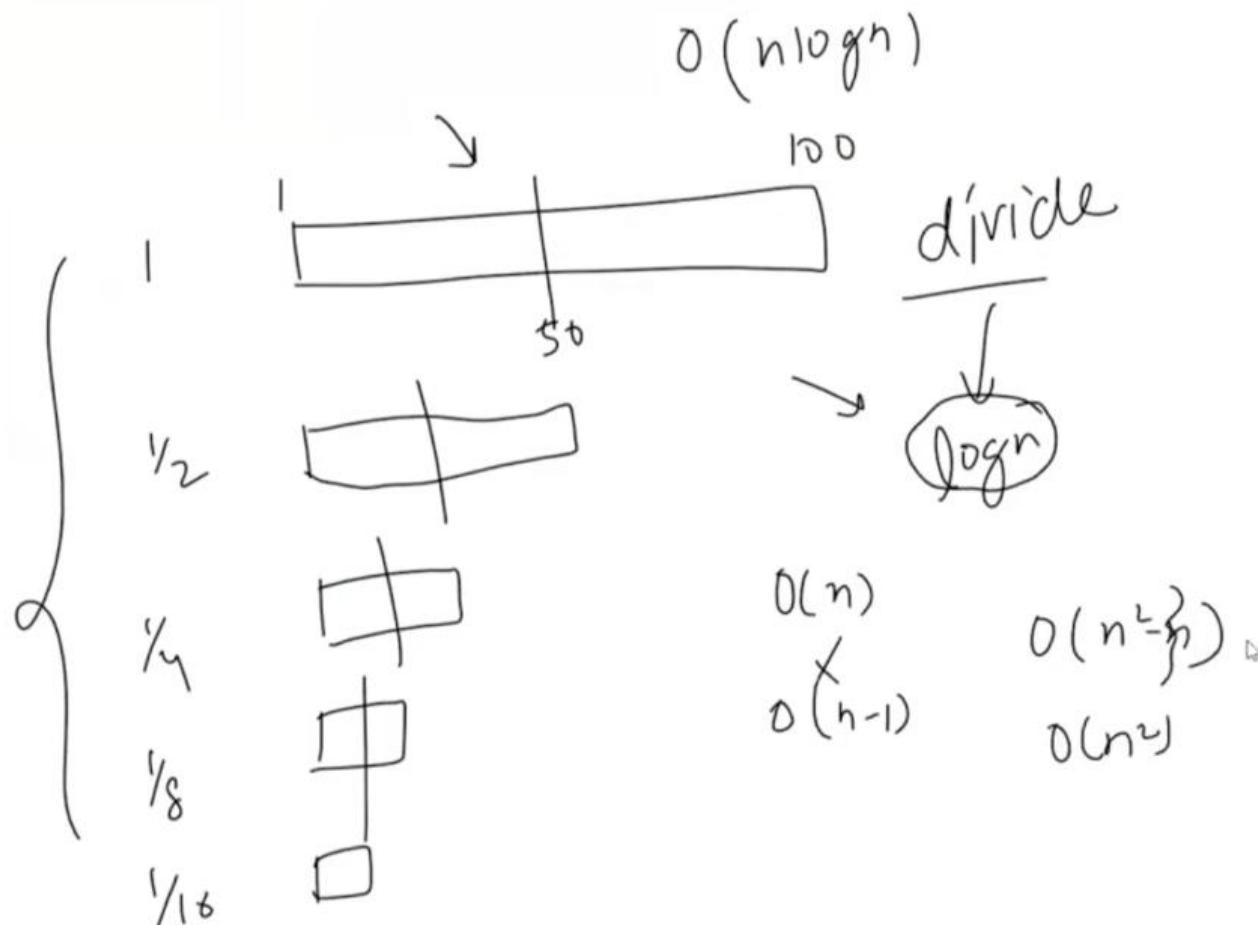
#####
##### Problem 7 #####
Binary Search

#####
##### Problem 8 #####
L = [1,2,3,4,5]
for i in range(0,len(L)):
    for j in range(i+1,len(L)):
        print('({},{})'.format(L[i],L[j]))

#####
##### Problem 9 #####
A = [1,2,3,4]
B = [2,3,4,5,6]

for i in A:
    for j in B:
        if i < j:
            print('({},{})'.format(i,j))
```





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```
##### Problem 8 #####
#
```

```
L = [1,2,3,4,5]
```

```
for i in range(0,len(L)):  
    for j in range(i+1,len(L)):  
        print('({},{})'.format(L[i],L[j]))
```

```
##### Problem 9 #####
#
```

```
O(ab)  
A = [1,2,3,4]  
B = [2,3,4,5,6]
```

```
for i in A:  
    for j in B:  
        if i < j:  
            print('({},{})'.format(i,j))
```

```
##### Problem 10 #####
#
```

```
A = [1,2,3,4]  
B = [2,3,4,5]
```



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O(ab)

A = [1,2,3,4]
B = [2,3,4,5,6]

```
for i in A:  
    for j in B:  
        if i < j:  
            print('({},{})'.format(i,j))
```

Problem 10

A = [1,2,3,4]
B = [2,3,4,5]

```
for i in A:  
    for j in B:  
        for k in range(100000):  
            print('({},{})'.format(i,j))
```

1

Problem 11

L = [1,2,3,4,5]

```
for i in range(0,len(L)//2):  
    other = len(L) - i -1
```



python-codes.py - C:\Users\91842\Desktop\python-codes.py (3.6.0)

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```
        print('{},{})'.format(i,j))
```

```
##### Problem 10 #####
```

```
A = [1,2,3,4]  
B = [2,3,4,5]
```

```
for i in A:  
    for j in B:  
        for k in range(100000):  
            print('{},{})'.format(i,j))
```

```
##### Problem 11 #####
```

```
L = [1,2,3,4,5]
```

```
for i in range(0,len(L)//2):  
    other = len(L) - i - 1  
    temp = L[i]  
    L[i] = L[other]  
    L[other] = temp
```

```
print(L)
```

```
##### Problem 12 #####
```



← → *python-codes.py - C:\Users\91842\Desktop\python-codes.py (3.6.0)*
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L = [1,2,3,4,5]



```
for i in range(0,len(L)//2):
    other = len(L) - i -1
    temp = L[i]
    L[i] = L[other]
    L[other] = temp

print(L)
```

Problem 12

```
def factorial(n):

    if n == 1:
        return 1
    else:
        return n*factorial(n-1)
    I
print(factorial(5))
```

Problem 13

```
def fib(n):

    if n ==1 or n == 0:
```





Home

Insert

Draw

View

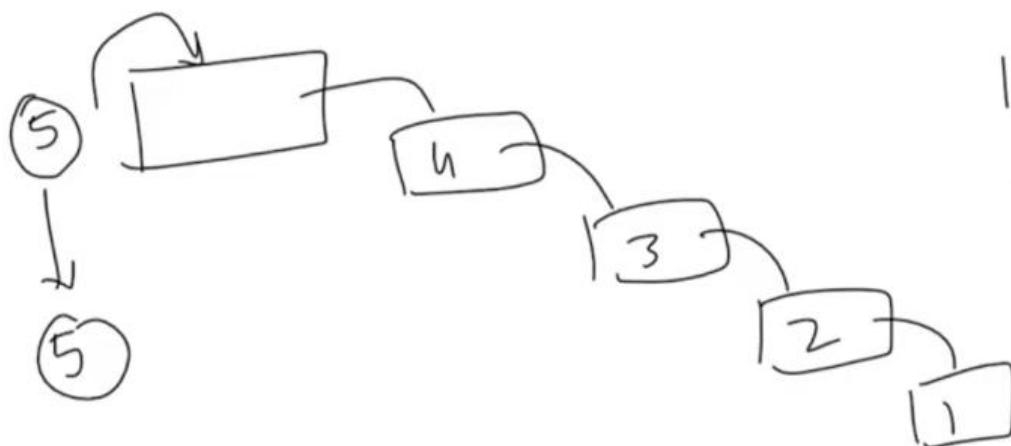
Help



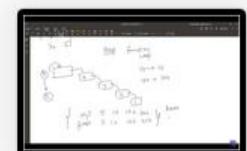
1/16



loop functioning
copy

 $10 \rightarrow 10$ $100 \rightarrow 100$

{ inp 5 10 100 200 }
 focus 5 10 100 200 } $O(n)$ linear



python-codes.py - C:\Users\91842\Desktop\python-codes.py (3.6.0)

File Edit Format Run Options Window Help



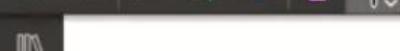
```
##### Problem 12 #####
def factorial(n):
    if n == 1:
        return 1
    else:
        return n*factorial(n-1)

print(factorial(5))

##### Problem 13 #####
def fib(n):
    if n ==1 or n == 0:
        return 1
    else:
        return fib(n-1) + fib(n-2)

##### Problem 14 #####
def power(num):
    if num < 1:
        return 0
```

Home Insert Draw View Help



fails

5

10

100

200

VVV

fib(1)

fib(2)

add

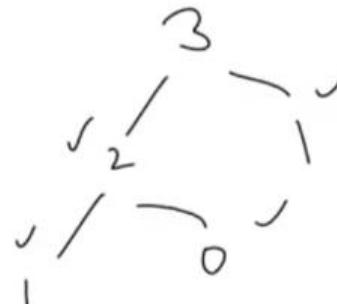
→

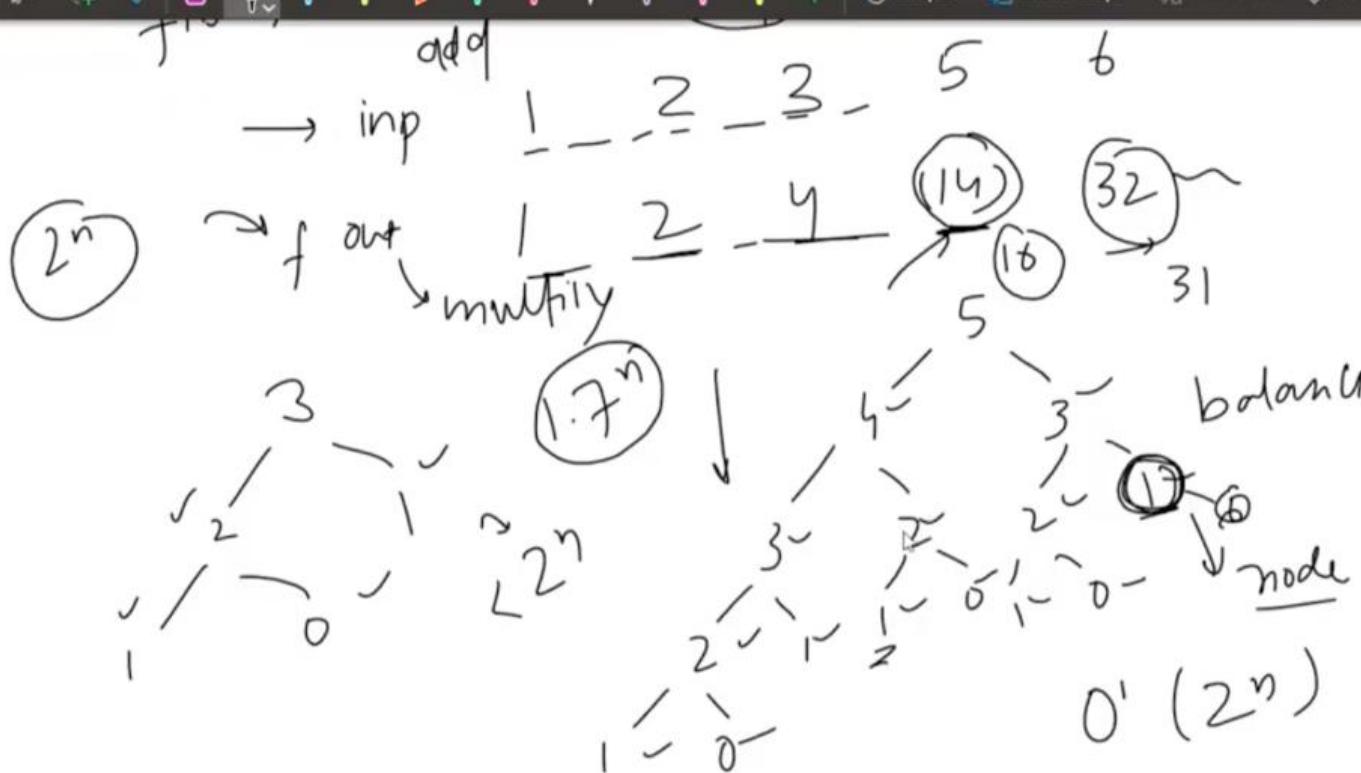
inp 1 2 3

5 6

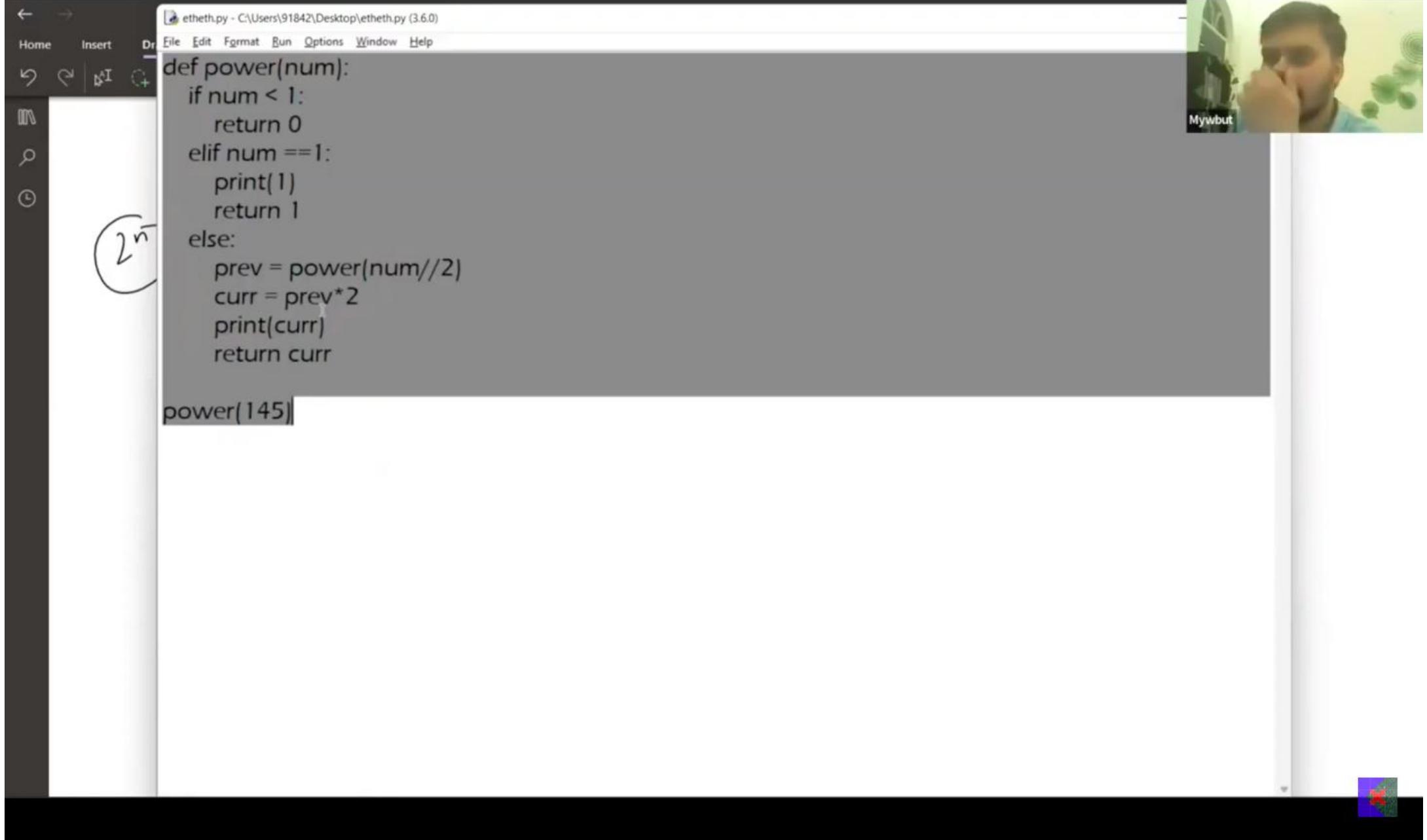
f out 1 2 4 (14) 32 ~

multiply





$1) \underline{2^n} \quad 2) \underline{(2^n)} \quad 3) \underline{> 2^n}$



```
def power(num):
    if num < 1:
        return 0
    elif num == 1:
        print(1)
        return 1
    else:
        prev = power(num//2)
        curr = prev*2
        print(curr)
        return curr

power(145)
```

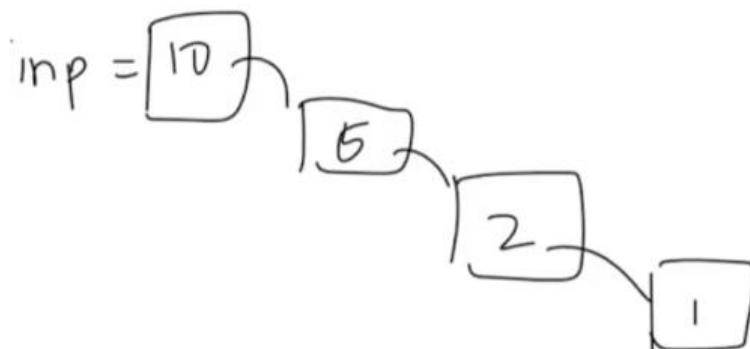


[←](#) [→](#)

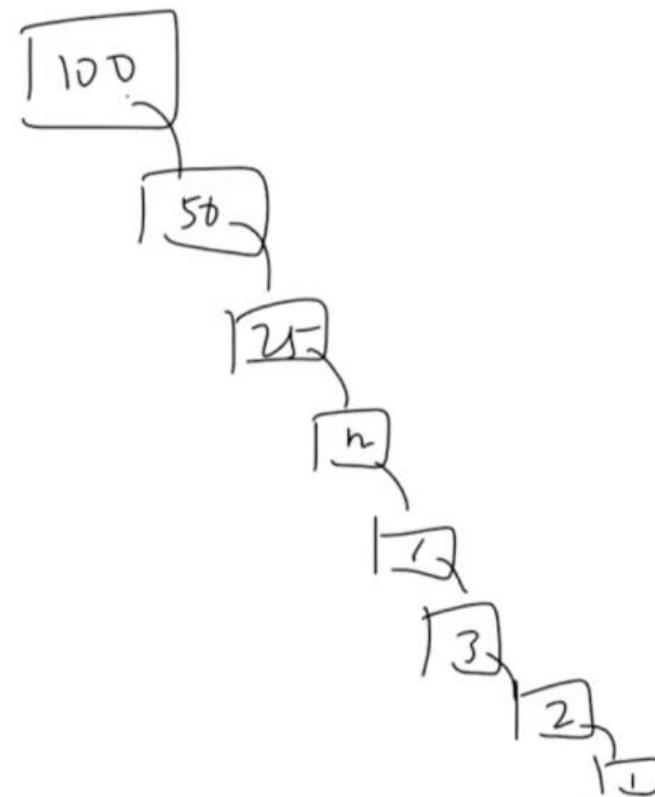
Home Insert Draw View Help

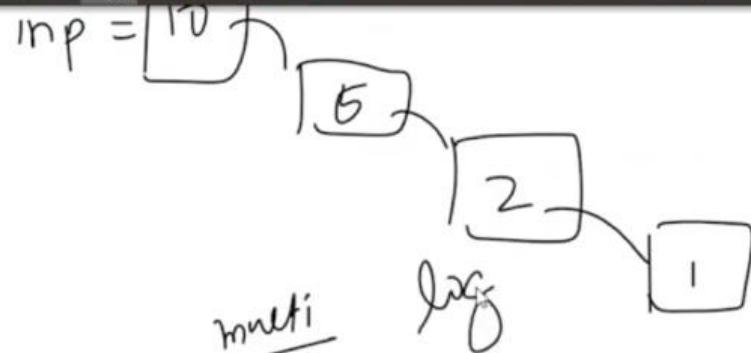


1) 2^n 2) (2^n) 3) $> 2^t$



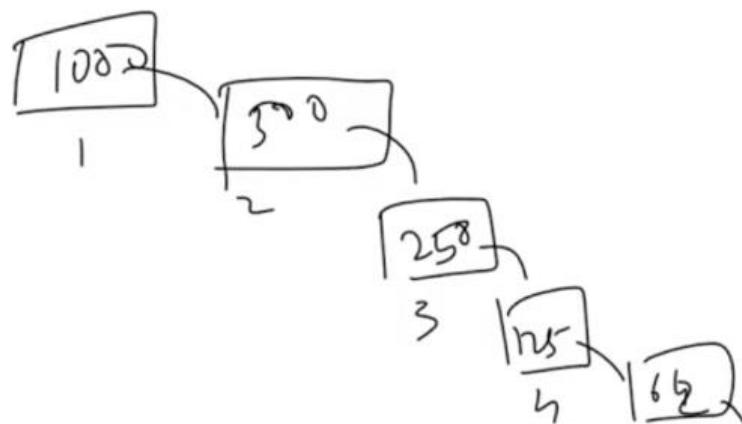
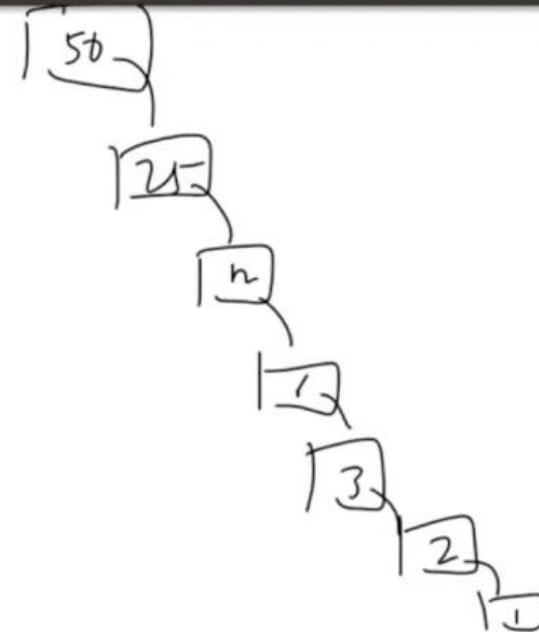
inp → 10 100
focuses n 8





{ Inp → 10 100 1000
 folds → n 8 12

Annotation: "add" under the "folds" column.



← → *python-codes.py - C:\Users\91842\Desktop\python-codes.py (3.6.0)*
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```
def power(num):
    if num < 1:
        return 0
    elif num == 1:
        print(1)
        return 1
    else:
        prev = power(num//2)
        curr = prev*2
        print(curr)
        return curr
```



```
##### Problem 15 #####
```

```
def mod(a,b):
    if b <= 0:
        return -1
    div = a//b
    return a - div *b
```

```
mod(5,3)
```

```
##### Problem 16 #####
```



python-codes.py - C:\Users\91842\Desktop\python-codes.py (3.6.0)

File Edit Format Run Options Window Help

Home Insert

Problem 15

```
def mod(a,b):
    if b <= 0:
        return -1
    div = a//b
    return a - div *b

mod(5,3)
```

Problem 16

```
def sum_digits(num):
    sum = 0
    while (num > 0):
        sum+=num%10
        num/=10

    return sum

sum_digits(123)I
```

python-codes.py - C:\Users\91842\Desktop\python-codes.py (3.6.0)

File Edit Format Run Options Window Help

Home Insert

return a - div *b

mod(5,3)

Mywbut

```
##### Problem 16 #####
def sum_digits(num):
    sum = 0
    while (num > 0):
        sum+=num%10
        num/=10
    return sum
sum_digits(123)
##### Problem 17 #####
T(n) = { 3T(n-1) if n>0
         { 1, otherwise
##### Problem 18 #####
T(n) = { 2T(n-1) -1 if n>0
         { 1, otherwise
##### Problem 19 #####
```



$$T_3 = \boxed{15} \rightarrow \boxed{2} | \boxed{3} | \boxed{2} - \boxed{1}$$

6 7 8 9

$$T_n = \begin{cases} 3T(n-1) & \text{if } n > 0 \\ 1 & \text{otherwise} \end{cases}$$

$T \rightarrow \text{time}$
 $n \rightarrow \text{input}$

$$\begin{array}{l} n=5 \quad 9t \\ \hline n=4 \quad 3t \\ \hline n=3 \rightarrow t \text{ sec} \end{array}$$

$$f(x) = \begin{cases} 2f(x-1) & \text{if } x > 0 \\ 1 & \text{otherwise} \end{cases}$$

$f(x)$ x $f(x-1)$ $2f(x-1)$ $\{$ $n=6 \quad 27t$
exp inp otherwise add $n=7 \quad 81t$

$\begin{matrix} 3 & 4 & 5 & 6 & 7 \\ + & 3t & 9t & 27t & 81t \end{matrix}$

Time
multi

Home Insert Draw View Help



Shapes

Ink to Shape

Ink to Text

Ruler



Mywbut

$$T(n) = \begin{cases} 3T(n-1) & \text{if } n > 0 \\ 1 & \text{otherwise} \end{cases}$$

$$\begin{aligned} T(n) &= 3\underbrace{T(n-1)}_{\substack{\downarrow \\ (3^n)}} \\ &= 3[3T(n-2)] = 3^2\underbrace{T(n-2)}_{\substack{\downarrow \\ (3^{n-1})}} \\ &= 3^2[3T(n-3)] = 3^3\underbrace{T(n-3)}_{\substack{\downarrow \\ (3^{n-2})}} \\ &\quad \vdots \\ &= 3^n\underbrace{T(n-n)}_{\substack{\downarrow \\ (3^0)}} \\ &= 3^n \times 1 \quad \leftarrow = 3^n T(0) \end{aligned}$$



Batch 6 Session 10

Wednesday, March 30, 2022 6:33 PM

DS Algo → Algorithmic complexity → Big Oh → Time complexity

{ Data structures

↳ What are DS? Need?

→ Linear Vs Non-linear DS

→ Arrays

Referential

Dynamic
Arrays

| DS
↓
{ way to store and
organize data
efficiently } Y



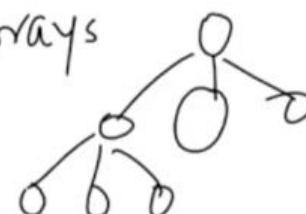
→ What are DS? Need?

→ Linear Vs Non-linear DS

→ Arrays

Referential

Dynamic
arrays

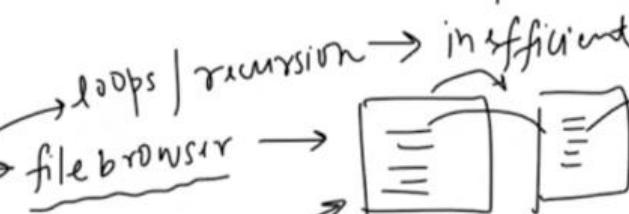


{ way to store and organize data efficiently

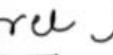
Need

[DS]

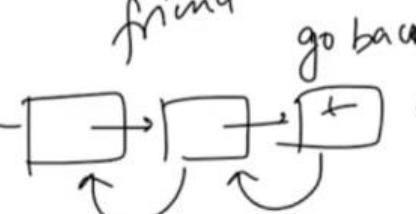
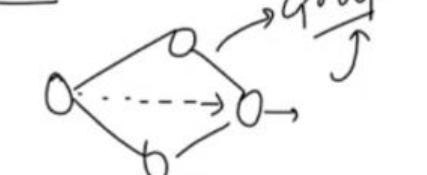
OS → file browser



DS → tree



graph TD



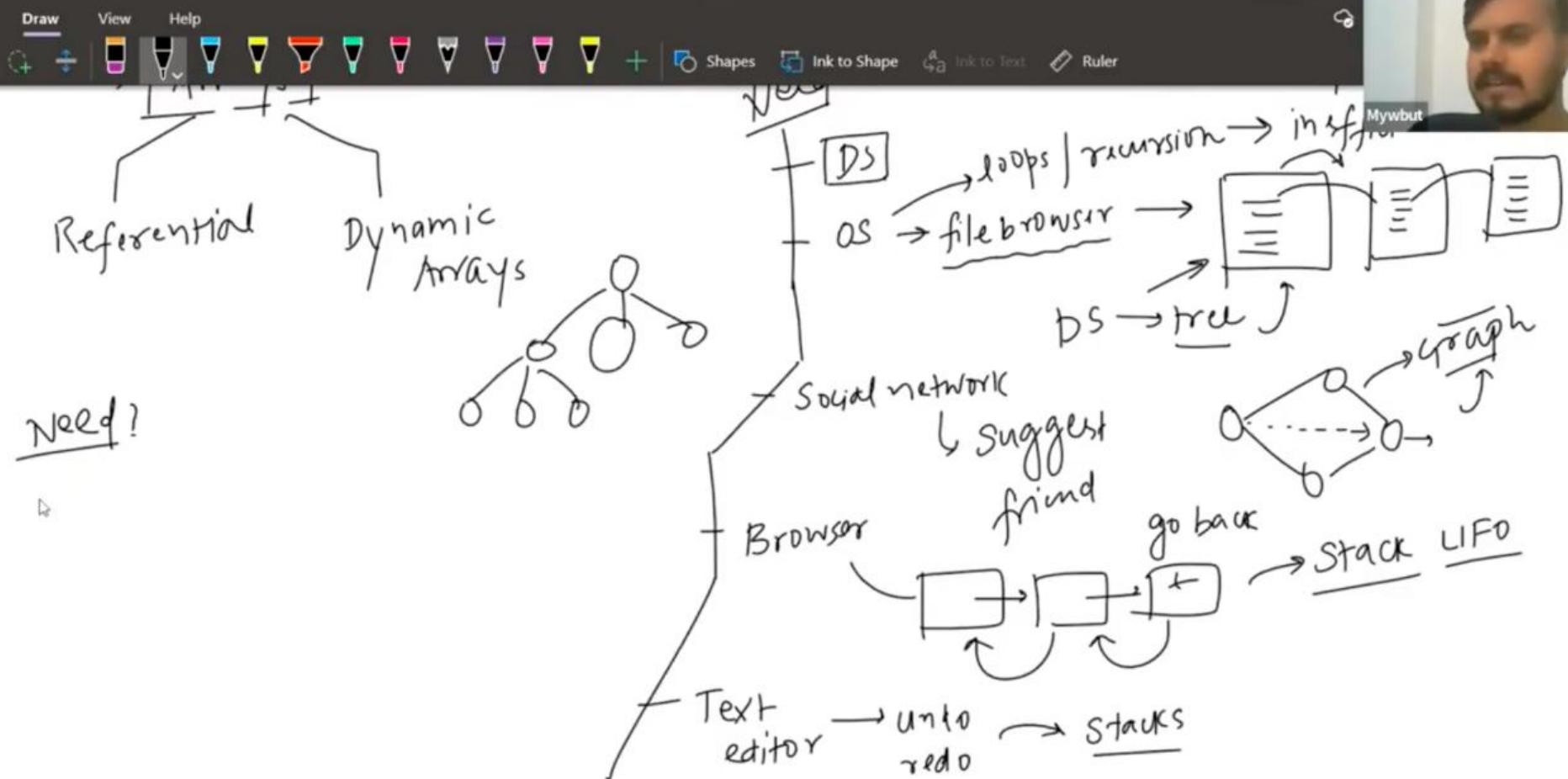
stack LIFO

Social network

suggest friend

Browser

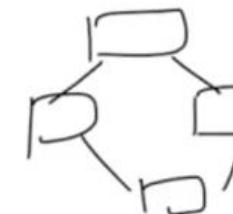
go back



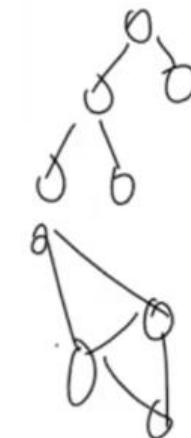


Linear

Vs Non Linear



Tree



Graph

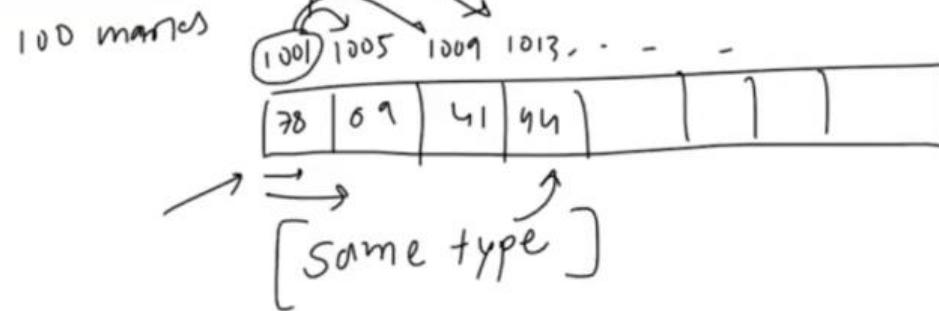
{
 Array -
 Linked List -
 Stacks ✓
 Queue ✓
 Hashing ✓



{ Queue
+ Hashing }

Arrays

Linear DS
store used to store
multiple item of
same type
in continuous
memory location





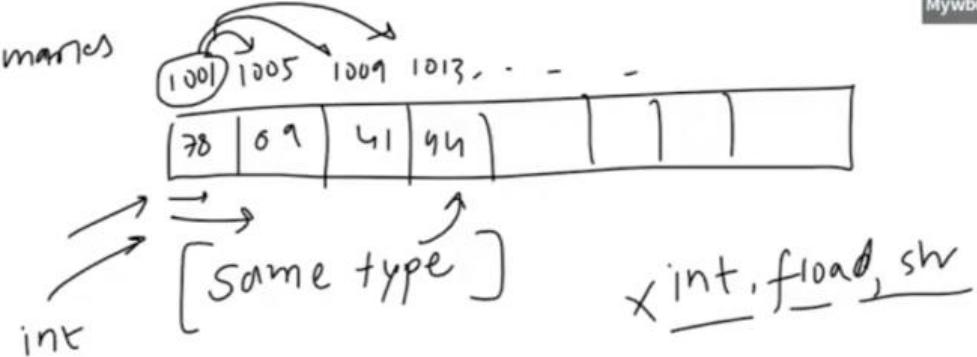
Linear DS
store used to store
multiple item of
same type
in contiguous
memory location

int A[5000]
(2)

C/C++ | Java

int A[50]
to do
task

100 marks



Disadvantage

→ fixed size (memory waste) →

→ homogeneous [lack of flexibility] →

multiple int
in same type
continuous
memory location

int A[5000]
(2)

C/C++ | Java

int A[50]
to-do
task

int [Same type]

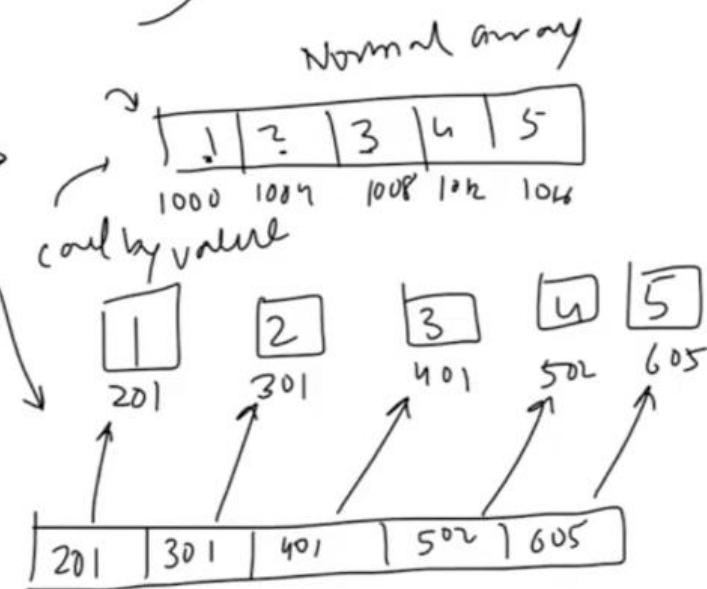
int, float, shr

Disadvantage

→ fixed size (memory waste) →

→ homogeneous [lack of flexibility] →

Referential arrays



multiple int
in same type
continuous
memory location

int A[5000]
(2)

C/C++ | Java

int A[50]
to-do
task

int [Same type]

int, float, shr

Disadvantage

→ fixed size (memory waste) →

→ homogeneous [lack of flexibility] →

Referential arrays

hetro

4-int

1-str

homoge

Normal array

1	2	3	4	5
1000	1004	1008	1012	1016

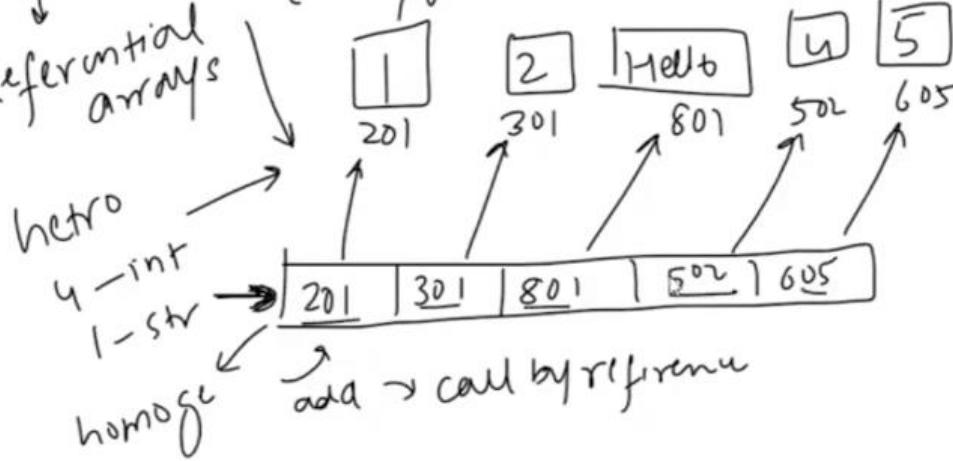
call by value

[1]

[2]

[Hello]

[4]
[5]



multiple int in continuous memory location

int A[5000]
(2)

C/C++ | Java

int A[50]
to-do
task

Python (Refer)
list

int [Same type]

int, float shr

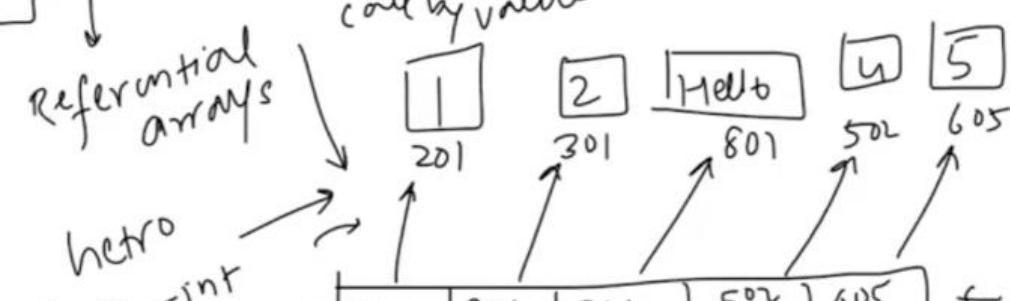
Normal array

1	2	3	4	5
1000	1004	1008	1012	1016

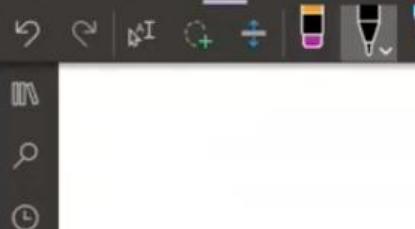
call by value

referential
arrays

hetro
4-int
1-str
homoge
add → call by reference



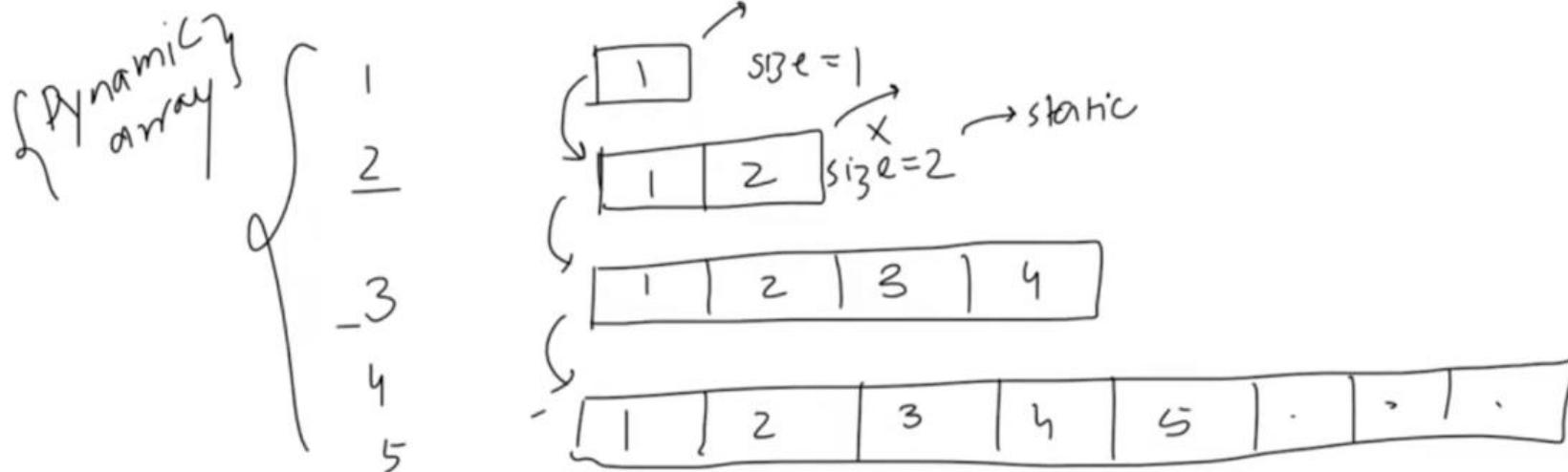
Home Insert Draw View Help

Interview
simp

↑ speed ↗ (n) ↗ NASA
{ website ↗ software
app's ↗ game ↗ control }

(Dynamic Array) ↗ static → fixed
↳ size → adjust

Python ↗ dynamic
list ↗ array



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Linear Data Struct → Linked Lists

Linear Data Struct → Array

How → collection of nodes

object

64 int first

Head

d 700 + 900
↑ ↑
500 700

Null ↑ tail

Quick Tip You've done pretty good work here - don't let it go to waste! Save your board here.

1 / 1 +

```
graph LR; Head[Head] --> d[d]; d --> n1[n1 700]; n1 --> n2[n2 900]; n2 --> n3[n3 1]; n3 -- Null --> tail[tail]
```

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Linked Lists

Linear Data Struct → **Array**

Node → collection of nodes

Array ? α

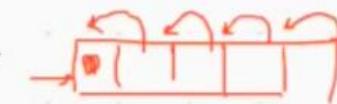


(LL)

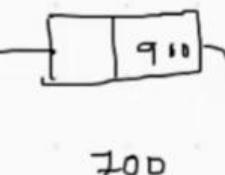
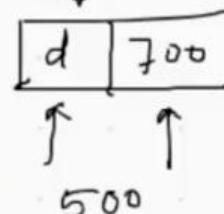


object

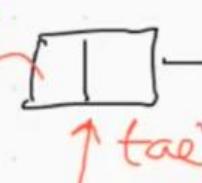
for
int first



{ Head }



Null



Quick Tip



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X



Array

In an array write(insert)/delete operation is very difficult because during this operation its time complexity is $O(n)$.

Means if we increase the size of an array and we try to perform the write(insert)/delete operation it will increase the time of execution. I.e there will be linear relationship.

But reading operation is easy, its time complexity is $O(1)$.

In array sometimes most of the memory is not getting used it gets wasted

So we will use this when we are making read heavy applications.

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Linked Lists

Linear Data Struct → Array

[Hn] → collection of nodes

Array ? α LL

+ LL write + $O(1)$

object {str int float}

{Head} d 700 → 900

500 Null tail

{insert delete}

700 1 900

{write $\rightarrow O(n)$ }

Quick Tip You've done pretty good work here - don't let it go to waste! Save your board here.

1 / 1 +

linked list

In a linked list writing(Inserting) & (Deleting) operation is easy its time complexity is constant. $O(1)$. If we use linked list it not acquiring much memory but in array sometimes most of the memory is not getting used it gets wasted. But reading operation is hard, its time complexity is $O(n)$.

With help off this linked list we can make other Data Structure like stack, queue etc.

So we will use this when we are making write heavy applications

```
*linked-list.py - C:\Users\91842\Desktop\linked-list.py (3.6... File Edit Format Run Options Window Help  
# sum of all the elements at the odd pos  
  
# 3. What is the output of following function?  
# 1->2->3->4->5  
  
def fun(head):  
    if(head==None):  
        return  
    if head.next.next!= None:  
        print(head.data, " ", end="")  
        fun(head.next)  
    print(head.data, " ",end="")  
  
# 1234321  
  
# 5. Consider the input_linked_list below  
# input_linked_list: 1->4->9->16  
# What will be the value of the elements  
# by passing the head of the input_link  
  
# iter 1  
-> h-> 1 next_node-> 4  
-> h.d->5 h-> 4 next_node-> 9  
  
# 5 22 25
```

Feb Schedule | InfyTQ Syllabus | AWW App | On | linked-list-demo | +

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1 2 3 -

1 2 3 4 3 2 1

$h=1$

$h=2$

$h=3$

$h=4$

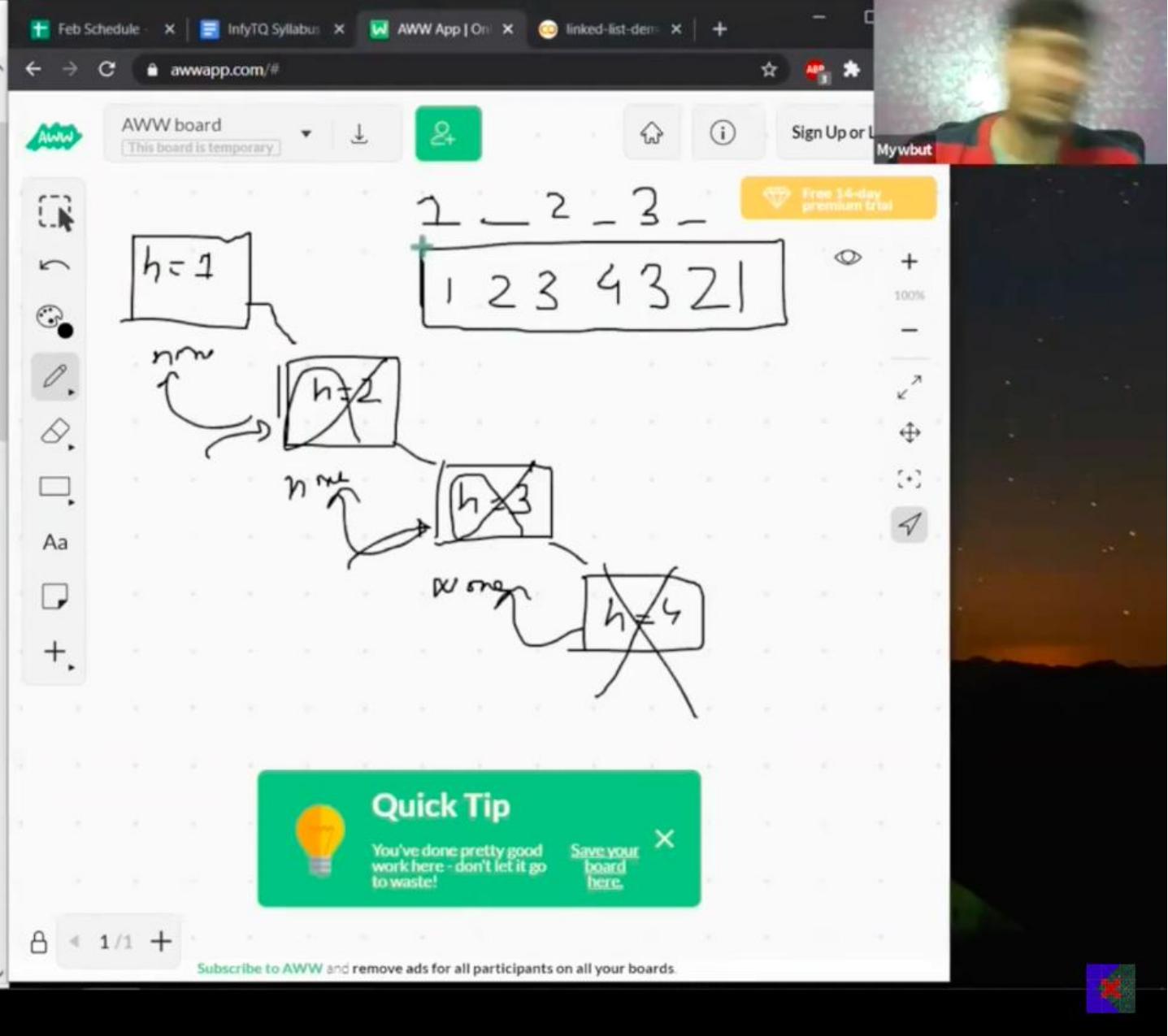
$h=5$

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Time complexity

Space complexity

Stacks

Quick Tip

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Save your board here

1. for i in HWo
H s.push(i)

2. while(!s.isemp)
S+=s.pop

10
Hello

11
olleH

11
Hello

11
olleH

11
Hello

11
olleH