Mygo	Best	Avascoge	Moses	Space complenity.
Bulle	O(n)	0 (na)	0 (n2)	00)
Selection	O(n3)	0 (ona)	0(n3)	3(1)
guseeksn	0(2)	o (na)	o (na)	000
menge	o (in log n)	O(n log n)	o (n log n)	o(n)
auek	o(n log n)	on log n	0 (nz)	O (logn) (erecursion)
Heap	O (n log n)	O (n log h)	O(n logn)	000
Counting	(n+n)	(MAN)O	D(N+W)	(M4W)
Radin	O(NW)	·O(nw)	O (NK)	(m+m)
Buckel	(M+W) 0	(M+N)0	6(n2)	(n)
Shell	o(n dog n	("(m (pos)") o (	o'(n (dog n)	000

Soul A Bulle sout 3.4 @dep:8-> vol. 91.5 the simplest

sositive algo. 91's works on repetably swapping of adjacont elements until they are not in indended onder. called as ludde sout.

Dis-advarlage: ->

97's not subalele ton læge data sets. The averge and wount case complenity of bullelle sout is - 6 (n2). membress of element.

@ word case is -o(n) Bullele sout use where-Despace complexity -> 6(1)

1) complenets does not matter.

And simple and shoulcode es prigaled.

BALGO 8-20 Coulde soll (arrey) begin

> 204 all array element MET Jose & [:] rought;

([1:] rose, [:] rose of power

end posy.

Reples ruly

end levelile

Le nordeni

"n c 3-7 implemelion

```
//bubble sort implemetion
 2
 3
    #include<stdio.h>
 5
    void print(int a[], int n)
        int i;
 8
        for(i = 0; i < n; i++) {
 9
            printf("%d ", a[i]);
10
11
        printf("\n");
12
13
14
    void bubble(int a[], int n)
15
16
17
        int i, j, temp;
        for(i = 0; i < n-1; i++) { // Outer loop iterates over the entire array</pre>
18
            for(j = 0; j < n-i-1; j++) { // Inner loop compares adjacent elements</pre>
19
                 if(a[j] > a[j+1]) { // If current element is greater than the next
20
                     // Swap the elements
21
22
                     temp = a[j];
                     a[j] = a[j+1];
23
24
                     a[j+1] = temp;
25
26
27
28
29
    int main()
30
31
32
        int a[5] = \{10, 25, 32, 13, 20\};
33
        int n = sizeof(a) / sizeof(a[0]);
34
        printf("Before sorting:\n");
35
        print(a, n);
36
37
        bubble(a, n);
38
39
        printf("After sorting:\n");
40
        print(a, n);
41
42
        return 0;
43
44 }
45
```

Defination 8% mage sould is a souling oppered. This will be very nelptul and engage interesting.

alposithen as it uses the divide and conquere opposition to sout clements.

Time :-> + Bost case | Bug case

Best case | Aug case | Wood case 1

O(n log n) | O(n log n) | O(n log n) |

all are same.

space 3-20 o(n),

and end is last element of array.

MERGIE\_SORT (aver, beg, end)
if beg < end
set mid = (beg + end)/2.

The second second

MERGIE\_SORT (aver, beg, mid)
MERGIE\_SORT (aver, mid+1, end)

MERGIE (aver, beg, nied, end)

end of of

End MERGIE\_SORT

```
//merge sort implemetion
 2
    #include <stdio.h>
 4
 5
    void merge(int arr[], int left, int mid, int right) {
        int n1 = mid - left + 1; // Size of the left subarray
 6
        int n2 = right - mid;  // Size of the right subarray
 8
 9
        // Create temporary arrays
10
        int L[n1], R[n2];
11
        // Copy data to temporary arrays L[] and R[]
12
        for (int i = 0; i < n1; i++)
13
            L[i] = arr[left + i];
14
        for (int j = 0; j < n2; j++)
15
16
            R[j] = arr[mid + 1 + j];
17
        // Merge the temporary arrays back into arr[left..right]
18
        int i = 0; // Initial index of the first subarray
19
        int j = 0; // Initial index of the second subarray
20
        int k = left; // Initial index of the merged subarray
21
22
        while (i < n1 && j < n2) {
23
            if (L[i] <= R[j]) {
24
                arr[k] = L[i];
25
26
                i++;
27
            } else {
                arr[k] = R[j];
28
29
                j++;
30
31
            k++;
32
33
        // Copy the remaining elements of L[], if any
34
        while (i < n1) {
35
            arr[k] = L[i];
36
37
            i++;
38
            k++;
39
40
        // Copy the remaining elements of R[], if any
41
        while (j < n2) {
42
            arr[k] = R[j];
44
            j++;
            k++;
45
46
47
48
49
    void mergeSort(int arr[], int left, int right) {
        if (left < right) {</pre>
50
            int mid = left + (right - left) / 2; // Avoid overflow for large indices
51
52
            // Recursively sort the first and second halves
53
            mergeSort(arr, left, mid);
54
            mergeSort(arr, mid + 1, right);
55
56
            // Merge the sorted halves
57
58
            merge(arr, left, mid, right);
59
60
61
    void printArray(int arr[], int size) {
62
63
        for (int i = 0; i < size; i++)
            printf("%d ", arr[i]);
64
        printf("\n");
65
66
67
    int main() {
68
69
        int arr[] = {12, 11, 13, 5, 6, 7};
70
        int n = sizeof(arr) / sizeof(arr[0]);
71
        printf("Given array is:\n");
72
        printArray(arr, n);
73
74
        mergeSort(arr, 0, n - 1);
75
76
77
        printf("\nSorted array is:\n");
        printArray(arr, n);
78
79
        return 0;
80
81
```

## 3) Quiek Soul 65

Dependent of the divide and conquer approach.

It is highly estimated and widely used in paietice due to 115 avanage-ease performance

@ wage 5-6

anier soul wonks by selecting a pivot element, pandid oning the army into two parts, and recursively southly the subarrays.

## Englands &

O (n log n) | O (n log n) | O (ng)

The state of the s

- Best and avarage case occur when the pivot divides the away evenly.
- e would coose occur when the pivot sesults highly unbalanced partitions.

## Space o (logn) (for recursive calls)

Co Super Co

the starting inden, and end is the last inden of array.

Canich-sout (auro beg, end)

set Pivot = Povot :4:0n(

aver, eg, end)

Canick-sout (asa, cep, Avet-1)
Canick-sout (asa, pivot+1, end)

End of 97 End auren\_sout. Partition Function 3->

Partiolism (and, Seq, end)

Set : Anot = and Eend?

Set := Seq - 1

For i = Seq +0 end -1

'f area [i] = Frot

Set := i+1

Set := i+1

Set := i+1

Send ap it

End of ton.

Since for I was about a server to server to server to the s

End Parbition.

```
//quick sort implemetion
    #include <stdio.h>
 4
    // Function to partition the array
    int partition(int arr[], int low, int high) {
        int pivot = arr[high]; // Pivot element (last element)
        int i = low - 1; // Index of smaller element
 8
10
        for (int j = low; j < high; j++) {
            // If the current element is smaller than or equal to the pivot
11
            if (arr[j] <= pivot) {</pre>
12
13
                i++;
14
                // Swap arr[i] and arr[j]
15
                int temp = arr[i];
16
                arr[i] = arr[j];
17
                arr[j] = temp;
18
19
20
21
        // Swap arr[i + 1] and arr[high] (pivot)
        int temp = arr[i + 1];
22
        arr[i + 1] = arr[high];
23
24
        arr[high] = temp;
25
        return (i + 1); // Return the partition index
26
27
28
    // Function to implement Quick Sort
29
    void quickSort(int arr[], int low, int high) {
30
        if (low < high) {</pre>
31
            // Partition the array and get the pivot index
32
            int pi = partition(arr, low, high);
33
34
            // Recursively sort elements before and after the partition
35
            quickSort(arr, low, pi - 1); // Left subarray
36
            quickSort(arr, pi + 1, high); // Right subarray
37
38
39
40
    // Function to print the array
41
    void printArray(int arr[], int size) {
       for (int i = 0; i < size; i++)
            printf("%d ", arr[i]);
44
        printf("\n");
45
46
47
48
    int main() {
49
        int arr[] = {10, 80, 30, 90, 40, 50, 70};
        int n = sizeof(arr) / sizeof(arr[0]);
50
51
        printf("Given array is:\n");
52
        printArray(arr, n);
53
54
55
        quickSort(arr, 0, n - 1);
56
        printf("\nSorted array is:\n");
57
        printArray(arr, n);
58
59
60
        return 0;
61 }
```