# **Assignment3 Report**

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### **File Tree**

```
bonus
├─ data.bin
 ├─ main.cu
├— Makefile
 - snapshot0.bin
├─ snapshot1.bin
├— snapshot2.bin
 ├─ snapshot3.bin
- user_program.cu
 ├─ virtual_memory.cu
 └─ virtual_memory.h
- report.pdf
- source
 ├─ data.bin
 ├─ main.cu
 ├─ Makefile
 ├─ snapshot.bin
 - user_program.cu
  - virtual_memory.cu
  └─ virtual_memory.h
```

### Under **source** directory:

data.bin: input data to the virtual memory
main.cu: main file for the program
Makefile: help compile the program
snapshot.bin: data output from the virtual memory
user\_program.cu: you can change it for another program reading or writing to the virtual memory
virtual\_memory.h: header file for the virtual\_memory.cu
virtual\_memory.cu: our implementation of read / write operations on the memory

### Under **bonus** directory:

data.bin: input data to the virtual memory

main.cu: main file for the program

Makefile: help compile the program

snapshot0.bin: data output from the virtual memory from after the thread0 user program terminates

snapshot1.bin: data output from the virtual memory from after the thread0 user program terminates

snapshot2.bin: data output from the virtual memory from after the thread0 user program terminates

snapshot3.bin: data output from the virtual memory from after the thread0 user program terminates

user\_program.cu: you can change it for another program reading or writing to the virtual memory

virtual\_memory.h: header file for the virtual\_memory.cu

virtual\_memory.cu: our implementation of read / write operations on the memory

# 1. Running Environment

For the **homework & bonus** program:

### **Operating System**

cat /etc/redhat-release

(base) [cuhksz@TC-301-12 .csc3150]\$ cat /etc/redhat-release
CentOS Linux release 7.6.1810 (Core)

#### **CUDA Version**

nvcc --version

```
(base) [cuhksz@TC-301-12 .csc3150]$ nvcc --version nvcc: NVIDIA (R) Cuda compiler driver Copyright (c) 2005-2019 NVIDIA Corporation Built on Fri_Feb__8_19:08:17_PST_2019 Cuda compilation tools, release 10.1, V10.1.105
```

```
cat /usr/local/cuda/version.txt
```

```
(base) [cuhksz@TC-301-12 .csc3150]$ cat /usr/local/cuda/version.txt CUDA Version 10.1.105
```

#### **GPU Information**

```
lspci | grep -i vga
```

```
(base) [cuhksz@TC-301-12 .csc3150]$ lspci | grep -i vga
00:02.0 <mark>VGA</mark> compatible controller: Intel Corporation UHD Graphics 630 (Desktop)
01:00.0 <mark>VGA</mark> compatible controller: N<u>V</u>IDIA Corporation GP106 [GeForce GTX 1060 6GB] (rev a1)
```

**NOTICE**: the machine is located at TC301 room, the 12th machine.

# 2. Execution Steps

#### Homework

```
cd <DIRCT_TO_PROJ>
                                            # come to the project
directory
cd source
                                            # go to the source file
make all
                                            # compile all the files
./test
                                            # execute the program
make clean
                                            # clean all the compiled files
diff data.bin snapshot.bin
                                            # compare the input and output
data bin
diff <(xxd data.bin) <(xxd snapshot.bin)</pre>
                                            # compare to view dirty data
```

### **Bonus Task**

```
# come to the project
cd <DIRCT_TO_PROJ>
directory
cd bonus
                                             # go to the bonus file
make all
                                             # compile all the files
./test
                                             # execute the program
make clean
                                             # clean all the compiled files
diff data.bin snapshot0.bin
                                             # compare the input and output
data bin
diff data.bin snapshot1.bin
diff data.bin snapshot2.bin
diff data.bin snapshot3.bin
diff <(xxd data.bin) <(xxd snapshot0.bin) # compare to view dirty data</pre>
diff <(xxd data.bin) <(xxd snapshot1.bin)</pre>
diff <(xxd data.bin) <(xxd snapshot2.bin)</pre>
diff <(xxd data.bin) <(xxd snapshot3.bin)</pre>
```

# 3. Progame Design

### **Homework Program**

#### **Overall Ideas**

In this program we are going to implement the **write**, **read**, **and snapshot** functions to the virtual memory.

The memory system in this program is consist of 3 parts, which are disk, physical memory for data, physical memory for page table. We can use virtual page number as the index of disk and physical page number as the index of physical memory. The page table is inverted, in which the index is physical page number and the value is {valid, dirty, pid, count, virtual page number}.

In the three kinds of operation, we first get to the page table, check the valid entry that whether the target virtual page number appears in the value. If yes, we get a **page hit**. Otherwise, we get into a **page fault**. To solve page fault, we check whether the page table is full. If there are positions in the page table, we insert the target page from disk into the physical memory and update the page table entry. If the table is full, we swap the least recently used page, write back to disk based on write back policy, and insert the new page into physical memory, update the page table.

### **How to Make a Good Write Back Policy**

Instead of write back the page in physical memory to disk in every page swap, we should figure out a better solution. We maintain a dirty bit for the page table entry. Once there is a **write** to the page, we update this bit as 1. When there is a page swap, we check the dirty bit of the page table entry. If the dirty bit is 1, we write it back into the disk. Otherwise, we replace the page directly.

### How to count the page fault

If the memory operation (read / write) cannot find a corresponding entry in the page table, we can conclude that the target page is not in the physical memory. Before doing the next step, we add 1 to the **global** number count of page fault.

### **Bonus Program**

#### **Overall Ideas**

The overall ideas are the same as the one of required task. The major difference is that we simulate the memory access of 4 processes by CUDA threads. We let this 4 threads to invoke the  ${\bf user\_program()}$  function one by one. That is, only when the ith thread's  ${\bf user\_program()}$  returns, the i+1th thread invoke its  ${\bf user\_program()}$ .

### **How to Isolate Memory Address Space of Different Processes**

Since the memory address space of different processes should be different, we maintain the pid value in page table using **2-bit** to verify the page owners. Only when the process id (thread id indeed) and the virtual page number are checked consistent with the target page in a valid page table entry, we conclude it as a page hit.

### **Thread Scheduling**

To make the 4 threads execute in sequence, we make use of the CUDA barrier function **syncthreads()** for thread synchronization.

### 4. Test Results

# **Required Homework - Provided User Program**

### Compilation

```
cd <DIREC_TO_PROJ>
cd source
make all
```

```
(base) [cuhksz@TC-301-12 hw3]$ ls
       report.pdf
(base) [cuhksz@TC-301-12 hw3]$ cd s*
(base) [cuhksz@TC-301-12 source]$ ls
data.bin main.cu Makefile snapshot.bin user_program.cu virtual_memory.cu virtual_memory.h
(base) [cuhksz@TC-301-12 source]$ make all
nvcc -c -rdc=true user_program.cu
nvcc -c -rdc=true virtual_memory.cu
nvcc -c -rdc=true main.cu
main.cu(89): warning: conversion from a string literal to "char *" is deprecated
main.cu(109): warning: conversion from a string literal to "char *" is deprecated
main.cu(89): warning: conversion from a string literal to "char *" is deprecated
main.cu(109): warning: conversion from a string literal to "char *" is deprecated
nvcc -arch=sm_37 -std=c++11 -cudart=shared -rdc=true user_program.o virtual_memory.o main.o -o test
(base) [cuhksz@TC-301-12 source]$ ls
data.bin main.o snapshot.bin user_program.cu virtual_memory.cu virtual_memory.o
main.cu Makefile test
                                user_program.o virtual_memory.h
```

### **Execution**

```
./test
```

```
(base) [cuhksz@TC-301-12 source]$ ./test input size: 131072 pagefault number is 8193
```

### **Clean the Compiled Files**

```
make clean
```

```
(base) [cuhksz@TC-301-12 source]$ make clean
rm -rf test *.o
(base) [cuhksz@TC-301-12 source]$ ls
data.bin main.cu Makefile snapshot.bin user_program.cu virtual_memory.cu virtual_memory.h
```

#### **Result Validation**

### The Number of Page Faults

In the test user program, we firstly invoke  $\mathbf{vm\_write()}$  to write 4K pages into the virtual memory, which will introduce 4K page faults. The writing order is from page 0 to page 4K.

Then we invoke **vm\_read()** to read 1k+1 pages in the virtual memory, from page 4K to 3k-1, which introduces 1 page faults.

After all, we invoke  $\mathbf{vm\_snapshot()}$  to get a snapshot of the virtual memory, which invokes  $\mathbf{vm\_read()}$  from 0 to 4K, introducing 4K page faults.

In sum, there are 4K+4K+1=8193 pages faults, **consistent** with our output results.

#### The snapshot of virtual memory

We compare the input data bin and the snapshot data bin.

```
diff data.bin snapshot.bin

(base) [cuhksz@TC-301-12 source]$ diff snapshot.bin data.bin
(base) [cuhksz@TC-301-12 source]$
```

They are **identical**! Our results for the provided user program is **correct**!

### **Required Homework - My User Program**

Here are some of my user programs for test.

In the first two examples we want to figure out **whether the page fault counter works well**. Therefore, we implement two user programs, in which we change the number of page fault introduced.

The result is correct:

```
(base) [cuhksz@TC-301-12 source]$ ./test
input size: 131072
pagefault number is 8192
```

```
// program2: expected 4096+4096+1+1024 = 9217 page faults

__device___ void user_program(VirtualMemory *vm, uchar *input, uchar *results,
    int input_size) {
    for (int i = 0; i < input_size; i++)
        vm_write(vm, i, input[i]);

    for (int i = input_size - 1-32768; i >= input_size - 32769-32768; i--)
    int value = vm_read(vm, i);

    vm_snapshot(vm, results, 0, input_size);
}
```

The result is correct:

```
(base) [cuhksz@TC-301-12 source]$ ./test
input size: 131072
pagefault number is 9217
```

In the last user program we want to test whether the write back policy works well.

In this version of user program, we try to write 0-15 to virtual address 0-15 and let these 16 pages written back into the memory, introducing 8195 page faults. The following is the test result:

```
diff <(xxd data.bin) <(xxd snapshot.bin) # compare to view dirty data</pre>
```

```
[cuhksz@TC-301-18 source]$ ./test
input size: 131072
pagefault number is 8195
[cuhksz@TC-301-18 source]$ diff <(xxd data.bin) <(xxd snapshot.bin)
1c1
< 00000000: 370c 6051 385a 250b 132a 2a01 0c07 052f 7.`Q8Z%..**.../
---
> 00000000: 0001 0203 0405 0607 0809 0a0b 0c0d 0e0f ......
```

In the result picture we can see that the first 16 address positions the correct value 0-15. And the page fault number is correct. The write back policy functions well.

### **Bonus Task - Provided User Program**

### Compilation

```
cd <DIREC_TO_PROJ>
cd bonus
make all
```

```
[cuhksz@TC-301-18 bonus]$ make all
nvcc -c -rdc=true user_program.cu
nvcc -c -rdc=true virtual_memory.cu
nvcc -c -rdc=true main.cu
main.cu(135): warning: conversion from a string literal to "char *" is deprecated
main.cu(155): warning: conversion from a string literal to "char *" is deprecated
main.cu(156): warning: conversion from a string literal to "char *" is deprecated
main.cu(157): warning: conversion from a string literal to "char *" is deprecated
main.cu(158): warning: conversion from a string literal to "char *" is deprecated
main.cu(158): warning: conversion from a string literal to "char *" is deprecated
main.cu(155): warning: conversion from a string literal to "char *" is deprecated
main.cu(156): warning: conversion from a string literal to "char *" is deprecated
main.cu(157): warning: conversion from a string literal to "char *" is deprecated
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main.cu(158): warning: conversion from a string literal to "char *" is deprecated
main.cu(158): warning: conversion from a string literal to "char *" is deprecated
main.cu(158): warning: conversion from a string literal to "char *" is deprecated
```

### **Execution**

```
./test
```

```
[cuhksz@TC-301-18 bonus]$ ./test
input size: 131072
pagefault number is 32772 _
```

### **Clean the Compiled Files**

```
make clean
```

#### **Result Validation**

### The Number of Page Faults

We let 4 threads to execute the **user\_program()** function one by one. Therefore, we firstly get the page fault number after one thread finishes its function.

In the test user program, we firstly invoke  $\mathbf{vm\_write()}$  to write 4K pages into the virtual memory, which will introduce 4K page faults. The writing order is from page 0 to page 4K.

Then we invoke  $\operatorname{vm\_read}()$  to read 1k+1 pages in the virtual memory, from page 4K to 3k-1, which introduces 1 page faults.

After all, we invoke **vm\_snapshot()** to get a snapshot of the virtual memory, which invokes **vm\_read()** from 0 to 4K, introducing 4K page faults.

In sum, there are 4K + 4K + 1 = 8193 pages faults, **consistent** with our output results. For 4 threads, there are 8193 \* 4 = 32772.

#### The snapshot of virtual memory

We compare the input data bin and the snapshot data bin.

```
diff data.bin snapshot0.bin
diff data.bin snapshot1.bin
diff data.bin snapshot2.bin
diff data.bin snapshot3.bin
```

```
[cuhksz@TC-301-18 bonus]$ diff data.bin snapshot0.bin
[cuhksz@TC-301-18 bonus]$ diff data.bin snapshot1.bin
[cuhksz@TC-301-18 bonus]$ diff data.bin snapshot2.bin
[cuhksz@TC-301-18 bonus]$ diff data.bin snapshot3.bin
[cuhksz@TC-301-18 bonus]$
```

They are **identical**! Our results for the provided user program is correct!

### **Bonus Task - Provided User Program**

Here are some of my user programs for test.

In the first two examples we want to figure out **whether the page fault counter works well**. Therefore, we implement two user programs, in which we change the number of page fault introduced.

The result is correct:

```
[cuhksz@TC-301-18 bonus]$ sh test.sh

nvcc -c -rdc=true user_program.cu

nvcc -arch=sm_37 -std=c++11 -cudart=shared -rdc=true user_program.o virtual_memory.o main.o -o test

input size: 131072

pagefault number is 32768
```

```
// program2: expected 4096+4096+1+1024 = 9217 page faults for each
thread: 36,868

__device___ void user_program(VirtualMemory *vm, uchar *input, uchar
*results,
    int input_size) {
    for (int i = 0; i < input_size; i++)
        vm_write(vm, i, input[i]);

    for (int i = input_size - 1-32768; i >= input_size - 32769-32768; i--)
    int value = vm_read(vm, i);

    vm_snapshot(vm, results, 0, input_size);
}
```

The result is correct:

```
[cuhksz@TC-301-18 bonus]$ ./test.sh
nvcc -arch=sm_37 -std=c++11 -cudart=shared -rdc=true user_program.o virtual_memory.o main.o -o test
input size: 131072
pagefault number is 36868
```

In the last user program we want to test whether the write back policy works well.

In this version of user program, we try to write 0-15 to virtual address 0-15 and let these 16 pages written back into the memory, introducing 8195 page faults. The following is the test result:

```
[cuhksz@TC-301-18 bonus]$ ./test.sh
nvcc -c -rdc=true user_program.cu
nvcc -arch=sm_37 -std=c++11 -cudart=shared -rdc=true user_program.o virtual_memory.o main.o -o test
input size: 131072
pagefault number is 32780
Binary files data.bin and snapshot0.bin differ
Binary files data.bin and snapshot1.bin differ
Binary files data.bin and snapshot2.bin differ
Binary files data.bin and snapshot2.bin differ
Binary files data.bin and snapshot3.bin differ
```

```
diff <(xxd data.bin) <(xxd snapshot0.bin) # compare to view dirty data
diff <(xxd data.bin) <(xxd snapshot1.bin)
diff <(xxd data.bin) <(xxd snapshot2.bin)
diff <(xxd data.bin) <(xxd snapshot3.bin)</pre>
```

In the result picture we can see that the first 16 address positions the correct value 0-15. And the page fault number is correct. The write back policy functions well.

# 5. Conclusion

According to the tests results, we can conclude that the design and code implementation of these two programs are successful.

In these tasks, I learnt:

- We should make good use of the space of page table to include more useful information in small space if possible.
- Page swap is really complicated and requires design to reduce it or make it more efficient. LRU can reduce the number of page swap based on the locality. A good write back policy can reduce the overhead of page swap.
- It is important to make the memory address space isolated among processes to avoid memory access error.