

AIWC for the Masses: Supporting Architecture-Independent Workload Characterization on OpenMP, OpenACC, CUDA and OpenCL

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Abstract—The next-generation of supercomputers will feature a diverse mix of accelerator devices. These accelerators span an equally wide range of hardware properties. Unfortunately, achieving good performance on these devices has historically required multiple programming languages with a separate implementation for each device. In the present day this results in the fragmentation of implementation – where an increasing amount of a programmer’s effort is expended to migrate codes between languages in order to use a new device. We have previously shown that presenting the characteristics of a code in a architecture-independent fashion is useful to predict execution times. From examining these highly accurate predictions we propose that Architecture-Independent Workload Characterization (AIWC) metrics are also useful in determining the suitability of a code for potential accelerators. To this end, we extend the usability of AIWC by supporting additional programming languages common to accelerator-based High-Performance Computing (HPC). We use two compilers to perform source-to-source level translation from CUDA-to-OpenCL, OpenMP-to-OpenCL and OpenACC-to-OpenCL and extend the usefulness of the AIWC tool by evaluating the base execution behaviour on these outputs. Essentially, we examine how AIWC metrics change between OpenMP, OpenACC, CUDA and OpenCL implementations of identical applications commonly used in scientific benchmarking.

Index Terms—architecture independent workload characterization, portability, OpenCL, CUDA

I. INTRODUCTION

High-Performance Computing (HPC) today is dominated by closed, proprietary software models. Despite OpenCL having existed for over a decade, it has generated little traction/adoption by the scientific programming community at large. Fragmentation between implementations of scientific codes already exists. Many lines of legacy code have already been reimplemented in accelerator directive-based languages such as OpenACC and OpenMP, while some computationally intensive kernels have been

implemented in CUDA. Unfortunately, many of these codes must frequently target new accelerators, as supercomputers are rapidly updated, and the accelerators used on a node are just as quickly replaced – often by different vendors and likely another class of accelerator (MIC, GPU or FPGA). We can easily imagine a world where developers at research laboratories will have full-time jobs rewriting the same codes on loop, taking years optimising and rewriting kernels for the upcoming system. This demand will only grow as the HPC center’s dependence on heterogeneous accelerator architectures increases, pushed by energy-efficiency requirements in the era of exascale computing, and is untenable.

Workload characterization is an important tool to examine the essential behaviour of a code and its underlying structure, and identifying the performance-limiting factors – an understanding of the latter is critical when considering the portability between accelerators. We perform these characterizations by examining the SPIR-V execution traces on an abstract OpenCL device. For instance, a code with regular memory accesses and predictable branching that is highly parallel (utilizing a large number of threads) is a suitable candidate for selection for a GPU type of accelerator. Conversely, inherently serial tasks are more suited for CPU devices which commonly offer a higher clock-speed. In the past we have used AIWC and Workload Characterization to perform accurate run-time predictions of OpenCL codes over multiple accelerators – motivated by the goal of automatically scheduling kernels to the most appropriate accelerator on an HPC system based on these essential characteristics. We have also shown that these characteristics are useful in guiding a developer’s efforts in optimizing performance on accelerators by outlining the potential bottlenecks of the implementation of an algorithm (in the amount of parallelism available, memory access patterns and scaling over problem sizes, etc). Unfortunately both motivations will be undone without industry adoption and community support for SPIR-V (and the OpenCL

runtime).

However Recent advances with SYCL, HIP/ROCM, and oneAPI may increase adoption of open-source models and present an alternative to the OpenCL-only characterization approach. These frameworks are also based on the SPIR-V representation, and support the OpenCL runtime and memory model, and can be leveraged by the same AIWC and Workload Characterization tools. Meanwhile, compilers are maturing and are increasingly able to provide source-to-source translations and code transformations, to generate low-level device-optimized code, and to allow implementations in one language to be mapped to another. The goal of this paper is to extend AIWC support for all languages: (1) firstly by assessing the AIWC feature-space outputs of contemporary source-to-source compiler/translator tools, and (2) secondly by highlighting the differences against a native OpenCL implementation to identify any potential inefficiencies of the translation by these tools.

In summary, in this paper we extend the use of AIWC to evaluate the current state-of-the-art in source-to-source translation to primarily examine the feasibility of leveraging existing language implementations of codes (as an example of those currently popular) to automatically map back to OpenCL. We summarize the common languages used on accelerators and survey the community’s interest in each. This is done to motivate our interest in tooling and offering support for evaluating the back-end generation of SPIR-V codes. Related work is discussed, followed by our methodology highlighting our selection of source-to-source translation tools used in our experiments. Finally we present results, close with a summary of our findings, and the explore directions for future-work.

II. ACCELERATOR PROGRAMMING FRAMEWORKS AND THEIR ADOPTION

CUDA – NVIDIA’s proprietary software model – has been around since 2007. It has been widely adopted to repurpose GPUs from rendering in gaming workloads to highly-parallel compute intensive tasks common to scientific codes. Unfortunately, it is a single-vendor language, thus CUDA codes are executed solely on NVIDIA devices. There has been some recent activity by AMD with ROCm, which aims to automatically translate CUDA codes to their HIP (Heterogeneous-Compute Interface for Portability) framework – an AMD defined standard and modified clang compiler.

OpenCL (Open Compute Language) was introduced shortly after CUDA (in 2009) as a standard agreed upon by accelerator vendors. This ideally allowing a code to be written once and executed on any (OpenCL compliant) device. Vendors typically each develop their own-backend OpenCL runtime. This can result in greater variation in performance since vendors can choose there extent of

support and optimization. To remedy this POCL (Portable Computing Language) is an Open Source initiative providing an implementation of the OpenCL standard. It is analogous to HIP but for OpenCL instead of CUDA – both leverage Clang and LLVM, Clang for the front-end compilation of codes and LLVM for the kernel compiler implementation. POCL offers backends to NVIDIA devices (via a mapping of LLVM to CUDA/PTX), HSA (Heterogeneous System Architecture) GPUs such as those offered by ARM and AMD, along with CPUs and ASIPs (Application-Specific Instruction-set Processor).

OpenMP [1] exists as one of the most widely-used tools in high-performance computing, in-part because of it’s high-level directive-based approach and broad availability. While OpenMP has traditionally been used to generate multi-threaded code on CPU devices, the recent addition of offloading directives in the OpenMP 4.X+ standards has extended OpenMP to also support accelerator devices, primarily GPUs. Similarly to OpenACC, the OpenMP offloading directives provide a high-level alternative to low-level offloading models like OpenCL and CUDA, and provide existing OpenMP programmers a familiar entrance to accelerator computing.

OpenACC [2] has been developed as a high-level directive-based alternative to low-level accelerator programming models. Lower-level approaches like CUDA and OpenCL typically require the programmer to explicitly manage mappings of parallelism and memory to device threads and memory, and require detailed knowledge about the target device. In contrast, OpenACC allows programmers to augment existing programs with directives to generically expose available parallelism, shifting the burden of thread and memory mapping to the underlying compiler. Because of its straightforward API and implications for performance portability, OpenACC has become an attractive option for domain scientists interested in exploring accelerator computing.

The confidentiality of most super-computer scale scientific codes means it is unknown what percentage of kernels are developed in CUDA, OpenCL, OpenMP or OpenACC. If we assume the open-source community adoption of accelerator programming frameworks reflects their popularity used by the scientific HPC community, a quick survey of GitHub can help indicate their adoption. As of January 2020, the number of repositories including CUDA in the title was 18k, OpenCL was in 8K repositories, OpenMP 5K repositories and OpenACC was featured < 400 repositories. If we compare the lines of code as a measure of language popularity, CUDA was listed more than 8M times, OpenCL occurred 2M times, OpenMP 2M times while OpenACC was featured just ~124K times. From the newer frameworks, SYCL is either in the title or description of 144 repositories, 402k lines of code in the search contained SYCL. At the time of writing, OneAPI is less than 6 months old which

biases this comparison, nonetheless there are 56 repositories and 4k lines. This survey was performed by searching for “CUDA”, “OPENCL”, “OPENMP”, “OPENACC”, “SYCL” and “ONEAPI” in GitHub¹. We are unable to be more precise around actual repositories using accelerator programming frameworks because OpenCL, OpenMP and OpenACC are not tracked as languages by GitHub – CUDA is listed as a language and has ~5.9k repositories listed as such. Based on these results, we suspect that OpenCL isn’t in the pole position to winning the race to adoption, however, translator tools may allow us to ensure OpenCL is used regardless – as a backend runtime to SPIR-V intermediate representation.

Mapping back to a common OpenCL runtime is an obvious compromise, since it supports the greatest range of accelerator devices, and by supporting multiple front-end languages and allowing developers the freedom to use desired language, still increases the impact of OpenCL. This common backend holds the potential to avoid the fragmentation and repeated reimplementations as systems are updated and accelerators replaced. Having efficient tools to enable this mapping is paramount. We use AIWC to evaluate the similarities and differences between outputs of translation tools on functionally equivalent kernel codes and show that it can be used to guide understanding of the mapping between versions and potential improvements to these tools.

III. RELATED WORK

A. SPIR-V

- LLVM common IR of all kernel codes

B. AIWC

- Brief about how AIWC currently operates on the Oclgrind LLVM OpenCL device simulator.
- The collected metrics have shown to accurately represent the codes characteristics and the suitability of accelerator devices with precise execution time predictions
- Offers insights around optimization of a code for accelerators by presenting interesting summaries for the developer to consider.
- Unfortunately, Oclgrind – and thus also AIWC – only support the OpenCL programming language by simulating an ideal (and abstract) OpenCL device.
- Abstract OpenCL device simulation enables that AIWC already architecture independent

C. OpenARC

The Open Accelerator Research Compiler (OpenARC) [3] has been developed as a research-oriented OpenACC and OpenMP compiler. OpenARC performs source-to-source translations and code transformations to generate low-level device-optimized code, like CUDA or OpenCL, specific to a targeted device. OpenARC’s primary strength is its ability to enable rapid prototyping of novel ideas, features, and API extensions for emerging technologies.

In this work, we leverage OpenARC’s OpenACC to OpenCL and OpenMP to OpenCL translations and use the output with AIWC to characterize the workloads resultant translation. This integration allows us to extend AIWC to characterize high-level codes written with OpenACC and OpenMP.

D. Coriander

At first glance we may be hopeful that AMD’s HIP could be used for all CUDA to OpenCL backend translations, however, sadly, it does not use OpenCL as a back-end, instead choosing to use LLVM to generate kernels for the HSA Runtime and Direct-To-ISA skipping intermediate layers such as PTX, HSAIL, or SPIR-V. It is unknown why this decision was made and by skipping SPIR we are unable to perform the analysis with AIWC over the output since it skips any potential abstract “ideal”/“uniform” device, useful to checking errors, validating memory accesses and performing workload characterization that AIWC and Oclgrind provides. Instead, we chose Coriander for the functionality of CUDA to OpenCL translation. Unlike OpenARC it skips source-to-source level translation and instead produces LLVM-IR/SPIR-V directly.

****TODO, summarize Coriander****

IV. METHODOLOGY

The Rodinia Benchmark Suite was selected since they offer a number of scientific benchmarks each implemented in all the targeted programming frameworks we compare. We consider the gaussian elimination benchmark, which is composed of two kernels. Coriander was used to convert a subset of the Rodinia Benchmark suite from CUDA to OpenCL translation while OpenARC was used for both OpenMP to OpenCL and OpenACC to OpenCL translation.

Thus for each benchmarks kernel we present a comparison of the AIWC feature-spaces of the baseline OpenCL against CUDA, OpenACC and OpenMP. We also discuss the required changes made to each implementation to get the closest approximation of work between versions.

Figure 1 presents a summary of the workflow and the various representations generated to interoperate between

¹www.github.com

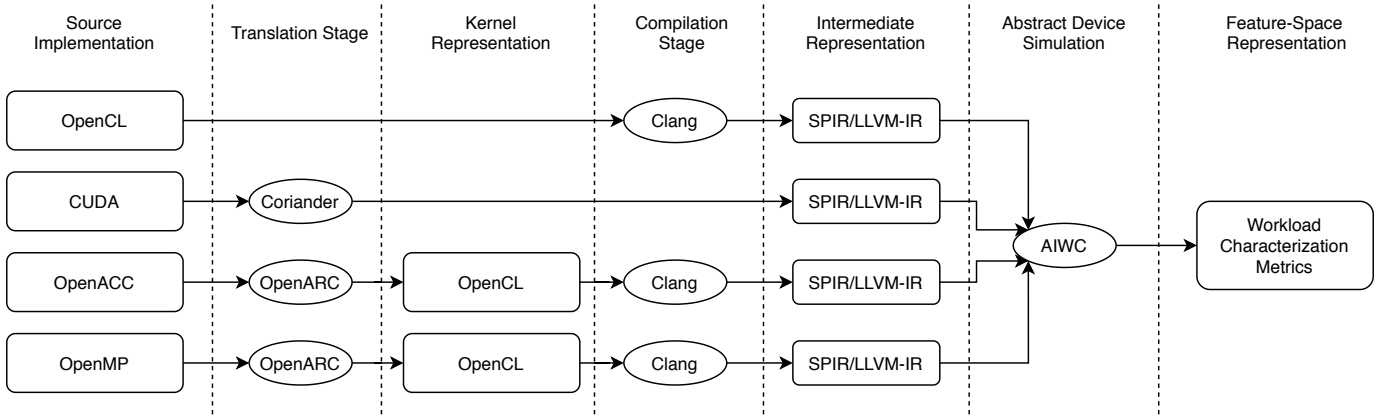


Figure 1: Workflow of using translators and compilers to interoperate with AIWC.

translators, compilers and AIWC. This shows how different source implementations of the same algorithm can generate equivalent workload characterization metrics. The entire workflow is composed of several stages and representations – broadly progressing from source code written in various languages, computational intensive regions are translated into OpenCL kernels, compiler with Clang into SPIR, executed on the Oclgrind OpenCL device simulator with the AIWC plugin, AIWC then presents and stores these metrics as a set of features describing the characteristics of the code. OpenCL can skip the translation stage and kernel representation since the kernel is already presented in OpenCL-C. OpenARC is used to perform source-to-source translation from OpenACC and OpenMP to OpenCL-C kernels. Note, coriander does not perform source-to-source translation effectively skipping the kernel representation and compilation stages, it still uses the same version of Clang (from LLVM 3.9.0) to produce the SPIR however this is used from within the tool effectively using clang behind-the-scenes for the compilation stage however the details are hidden from the user and are thus omitted in the diagram. Our hypothesis is that these characteristics should be largely independent of language used to implement it, although it is expected that different compiler optimizations and translation strategies will subtly change the metrics. Part of the goal of this paper is to examine the magnitudes of change imposed by language, compiler and translator. Ultimately, the similarities between metrics despite the original implementation can be used to evaluate the similarities in compiler and translator toolchains.

V. MANUAL MODIFICATIONS

Source code changes include C++ programs are unsupported as input programs, as such we need to change the `bfs` application extension. Most of the translation was autonomous, only one explicit typecast for a user defined struct needed to be clarified use ‘struct’ keyword to refer to the type. Use of the preprocessor also doesn’t work

within OpenARC thus `TRANSFER_GRAPH_NODE` which was `typedefed` to 1 was explicitly replaced.

- Steps taken to perform code translation from openacc to opencl kernel and wrapper

Use of Coriander was largely autonomous except it lacks support for `cudaMemcpyToSymbol` function calls.

Deprecated calls in the code-base to `cudaThreadSynchronize` also required replacing with `cudaDeviceSynchronize`. As such, the translation effort was almost entirely focused on replacing these operations with `.` Manually replacing certain includes – for instance `<math.h>` to `<cmath>` – were needed for the translation effort since Coriander is tied to LLVM 3.9 and is hard-coded to only support C++11. **TODO:** examine bfs (maybe fix invalid memory accesses) **TODO:** list which demos weren’t able to build using Coriander

- Coriander and AIWC were required to both be statically linked against the same version of LLVM due to an error which occurs when two separate instances of LLVM are registered into the same vtable. **do I mean vtable?** error “‘phi-node-folding-threshold’ registered more than once” These changes have been provided in as an artefact with Docker incase you wish to try for yourself ² Also many of the CUDA codes would crash during the translation to OpenCL – thus, only a small subset of the Rodinia Benchmark suite are presented, nonetheless this work forms a proof-of-concept and these issues within Coriander can be fixed.

VI. RESULTS

We now present the results of this comparison over Gaussian Elimination (GE). The GE benchmark was provided by The Rodinia Benchmark Suite [4] with an OpenACC, CUDA and OpenCL implementation. We found

²<https://github.com/ANU-HPC/coriander-and-oclgrind>

the existing implementations to lack a perfect mapping between versions, in particular our work modifies the partitioning of work to ensure an equivalent division of work is allocated between versions. We also built the OpenMP version – with 4.0 accelerator directives – based on the OpenACC version. All implementations have been divided into two discrete computationally intensive regions/kernels – known as **Fan1** and **Fan2**. **Fan1** calculates the multiplier matrix while **Fan2** modifies the matrix A into the Lower-Upper Decomposition. For our experiments the application was evaluated over a fixed dataset of a 4x4 matrix, where the same data was applied to all the implementations.

A. Fan1

4 work-items were used for all implementations, however the way parallelism is expressed differ between languages and for an apples to apples comparison in generated AIWC features required code changes to the kernel. These changes are listed in full in the associated Jupyter artefact

Add url once this is pushed to github

, for convenience a summary

In OpenCL

The initial CUDA implementation had a variation in parallelism due to the way it is expressed in CUDA compared to OpenCL. To this end, it was adjusted to more closely mirror the behaviour of the OpenCL version offered in the benchmark. The Block size was explicitly set to the MAXBLOCKSIZE (512 threads), our change: `block_size = (Size % MAXBLOCKSIZE == 0) ? MAXBLOCKSIZE : Size;` states that if we have smaller work to do than the max block size, just run 1 block of that size, which mirrors the way OpenCL expresses parallelism of this benchmark – i.e. the `global workgroup size` is the total number of threads to execute run in teams of `local workgroup size`. Thus, the CUDA implementation went from 512 workitems being invoked (where only 4 of them did any meaningful work) to 4 workitems being run.

In OpenMP 4 threads are explicitly requested by setting the `OMP_NUM_THREADS=4` environment variable at runtime while work-items in OpenACC was manually modified to support the same number of parallelism as the other versions. **mention other omp changes**

OpenARC uses `workers` and `gangs` variables to express parallelism in the OpenACC to OpenCL setting. To this end, we added these variables and the MAXBLOCKSIZE to be 512 to be equivalent to the CUDA version of the Gaussian Elimination benchmark. `gangs = (Size % MAXBLOCKSIZE == 0) ? MAXBLOCKSIZE : Size;` is set to be analogous to `block_size` (`block_size = (Size % MAXBLOCKSIZE == 0) ? MAXBLOCKSIZE : Size;`) which we added to the CUDA version, similarly, `workers = (Size/gangs)`

`+ (!(Size%gangs)? 0:1);` is identical to the CUDA version of `grid_size` (`grid_size = (Size/block_size) + (!(Size%block_size)? 0:1);`). Finally, the OpenACC pragmas were modified to explicitly use the `workers` and `gangs` variables: from `#pragma acc parallel loop present(m,a)` to `#pragma acc kernels loop independent gang(gangs) worker(workers)`.

The resulting source code of each of the four implementations is listed in Table 1.

The baseline OpenCL kernel code for Fan1 is presented in Listing 1-i, CUDA in , OpenACC in , and OpenMP in.

write summary

In summary, the OpenACC implementation went from 64 workitems being invoked (where only 4 of them did any meaningful work) to 4 workitems being run.

The AIWC metrics of this kernel are presented in figure~2.

Metrics (along the x-axis) have been grouped by category and is indicated by colour. These categories outline the overall type of characteristic being measured by each metric. The blue metrics (Opcode and Total Instruction Count) show the “Compute” category (which denote the amount of work to be done per thread and the diversity of the instruction sets required), metrics in green present “Parallelism” type metrics (these metrics are broadly around number of threads available in the workload, the amount of independence between threads and whether vectorization/SIMD is appropriate), “Memory” are presented in beige (and are included to collect the spread, proximity and raw number of memory addresses accessed), while purple metrics indicate “Control” (the predictability of branching during control flow of the workload). A full list and description of these metrics is available [5] but for brevity is not further discussed in this paper.

The y-axis presents the absolute count of each AIWC metric. The bars have been coloured according to Implementation (as shown in the legend) with CUDA in green, OpenACC in blue, OpenCL in tan and OpenMP in grey. Each metric has the four implementations grouped together, thus Figure~2 gives a visual inspection of the feature-space comparison of each metric between all implementations. It’s expected that OpenCL should be the lowest count – or the lowest overhead – of all the implementations regardless of metric, since it serves as the baseline; a compiler doing source-to-source translations would have to be doing additional optimizations to result in lower counts than the OpenCL baseline.

Figure~3 shows the same comparison of Fan1 implementations but with normalization against the baseline OpenCL counts, and is done to show the relative difference between each implementation – allowing a closer inspection of the differences. A flat-line at 0% is ideal since it shows no difference between metrics captured by AIWC, a

i: OpenCL

```
__kernel void Fan1(__global float *m, __global float *a,
    ↪ const int size, const int t){
    int gid = get_local_id(0) + get_group_id(0) *
        ↪ get_local_size(0);
    if (gid < size-1-t) {
        m[size*(gid+t+1)+t] = a[size*(gid+t+1)+t] / a[size
            ↪ *t+t];
    }
}
```

ii: CUDA

```
__global__ void Fani(float *m, float *a, int size, int t)
{
    int gid = threadIdx.x + blockIdx.x * blockDim.x;

    if(gid < size-1-t){
        m[size*(gid+t+1)+t] = a[size*(gid+t+1)+t] / a[size
            ↪ *t+t];
    }
}
```

iii: OpenACC

```
void Fan1(float *m, float *a, int size, int t)
{
    int i;
    #pragma acc kernels loop independent gang(fan1_gangs)
    ↪ worker(fan1_workers)

    for (i=0; i < size-1-t; i++)
        m[size*(i+t+1)+t] = a[size*(i+t+1)+t] / a[size*t+t]
        ↪ ];
}
```

iv: OpenMP

```
void Fan1(float *m, float *a, int size, int t)
{
    int i;
    #pragma omp target teams distribute parallel for
        ↪ num_teams(fan1_teams) num_threads(fan1_threads)
        ↪ )
    for (i=0; i < size-1-t; i++)
        m[size*(i+t+1)+t] = a[size*(i+t+1)+t] / a[size*t+t]
        ↪ ];
}
```

Listing 1: All source code implementations of the Fan1 kernel separated by language.

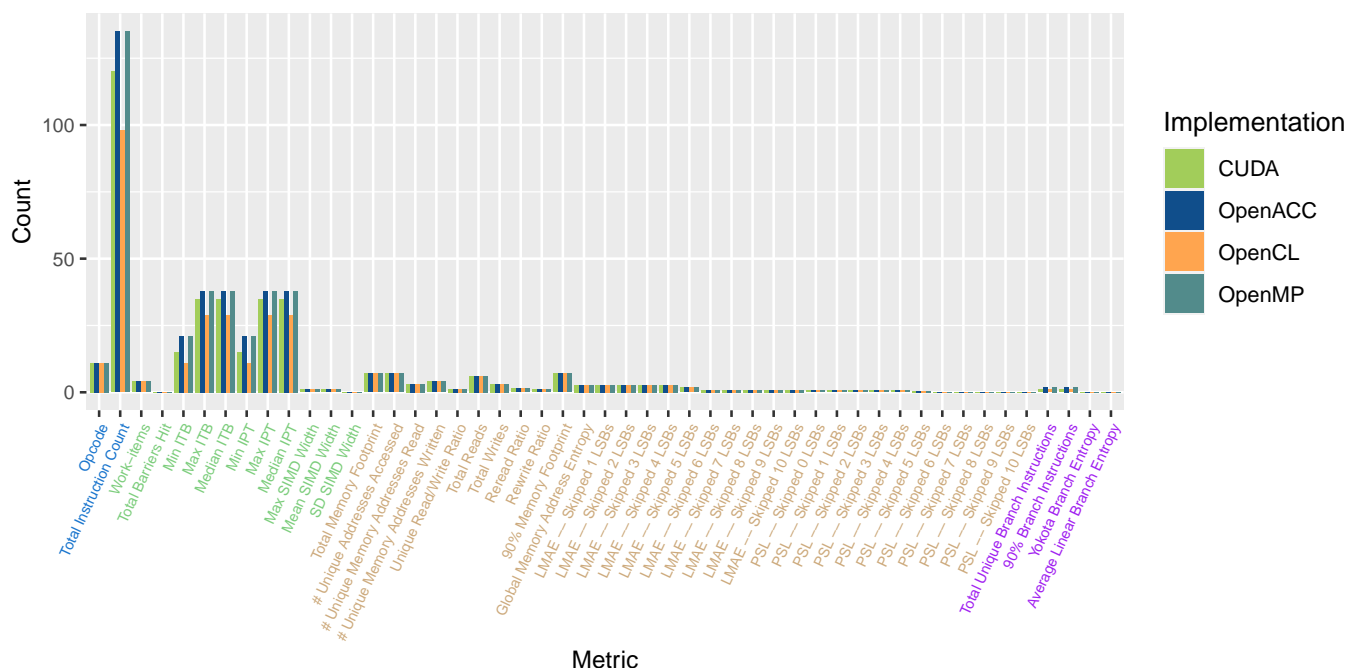


Figure 2: Absolute difference in AIWC metrics between each translated Fan1 kernel implementation.

perfect translation between implementations results in the same instructions being executed, operating on the same sequence of locations in memory, under the same degree of parallelism and identical AIWC metrics will ensue. In other words, if the applications workload characteristics are identical between languages the translator is doing an excellent job in preserving the structure (in terms of memory accesses, parallelism and compute work) of the code regardless of language. The implementations have been separated by colour and grouped into metrics for contrast. Firstly, the Opcode diversity metric is the same between all implementations however the number of instructions executed differ – the CUDA translation has 24% more instructions than OpenCL, while OpenACC and OpenMP increase this count by 37%. To understand the reason, we must examine the generated SPIR and associated traces, presented in Listing 3 and 5 respectively, and is discussed in Sections ?? and ??.

fix these references and add summary of the kernel representation subsection

We see all “Memory” (beige) metrics (on the x-axis) do not indicate any difference of any implementations against the OpenCL case – this is good as it ensures that all the same frequency of memory accesses, the type (whether a read or write), the locations and order of memory accesses are preserved and are equivalent in all implementations, and shows an indistinguishable amount of work has occurred.

The “Total Unique Branch Instructions” and “90% Branch Instructions” are doubled in both the OpenACC and OpenMP versions compared to OpenCL and CUDA. This is due to the absolute number of branch instructions are doubled –

reference SPIR code block

We see no causes where the compiler improves beyond the initial OpenCL baseline.

B. Kernel Representation

Listing 2 presents the OpenCL kernels generated in the Kernel Representation stage. The workflow from Figure 1 shows how the mix of translators interoperate with AIWC, and justifies why CUDA and OpenCL implementations are excluded from this comparison. Namely, no translation is needed for the OpenCL implementation, while Coriander operating on the CUDA implementation does not generate any kernel representation form and only offers an intermediate representation – which are discussed in Section VI-C. Of the two translated kernels presented in Listing 2-i and 2-ii can be compared directly to the hand-coded OpenCL kernel presented in Listing 1-i – they are the translated OpenARC output of OpenACC and OpenMP (from 1-iii and 1-iv) respectively. We see the pragmas are preserved in the translated output regardless of whether

OpenACC or OpenMP are used, however the OpenMP pragma is expressed in terms of OpenACC – the number of threads and number of teams are rewritten as workers and gangs – this is due to an intermediate step in OpenARC which converts OpenMP into OpenACC so it can directly use the OpenACC to OpenCL functionality. Regarding generated OpenCL kernels, both versions are equivalent, sharing the same logic (identical instructions at the same line numbers). Both have the same number of lines in the generated kernels – although the lines in the OpenMP based version are longer because of longer variable names.

When compared to the OpenCL hand-coded version shown in Listing 1-i, both generated kernels have a fundamental difference in structure. There is the same check (in the form of an `if`-statement) to ensure work isn’t occurring beyond the defined global boundary – expressed as global work size in OpenCL. However, there is an added `for`-loop than exists in the OpenCL baseline (Listing 1-i). OpenARC expresses all pragma based acceleration as using both local and global workgroups – this makes sense as many kernels use local workgroups to utilise shared memory and ensure good memory access patterns (in the form of cache reuse on many hardware architectures) – but the Fan1 base-line kernel doesn’t. This artifact of translation explains many of the differences in the AIWC metrics when comparing the OpenACC and OpenMP to OpenCL and is discussed further in both the Intermediate-Representation analysis, in Section VI-C, and trace analysis, in Section VI-D.

C. Intermediate-Representation

A comparison between generated LLVM/SPIR is presented in Listings 3 and 4. Both identify differences in SPIR between Coriander (for CUDA implementations) and OpenARC (OpenACC and OpenMP) compiler outputs against the OpenCL based native version. The similarities between OpenMP and OpenACC implementations of the Fan1 kernel – along with using the same compiler/translator toolchain – means that the generated SPIR are identical and thus consolidated into a single Listing (4-ii).

D. Trace Analysis

To examine these differences in actual execution based on the LLVM-IR codes we added the printing of the name of each executed instruction thereby giving a trace of each implementation. This was achieved by adding:

```
if(workItem->getGlobalID()[0]==0){
    printf("%s\n",opcode_name.c_str());
}
```

to the function `instructionExecuted` to AIWC (in `src/plugins/WorkloadCharacterisation.cpp`) which is triggered as a callback when the `Oclgrind` simulator executes each instruction. Since `oclgrind` is a multithreaded program – to the extent that each OpenCL workitem is run

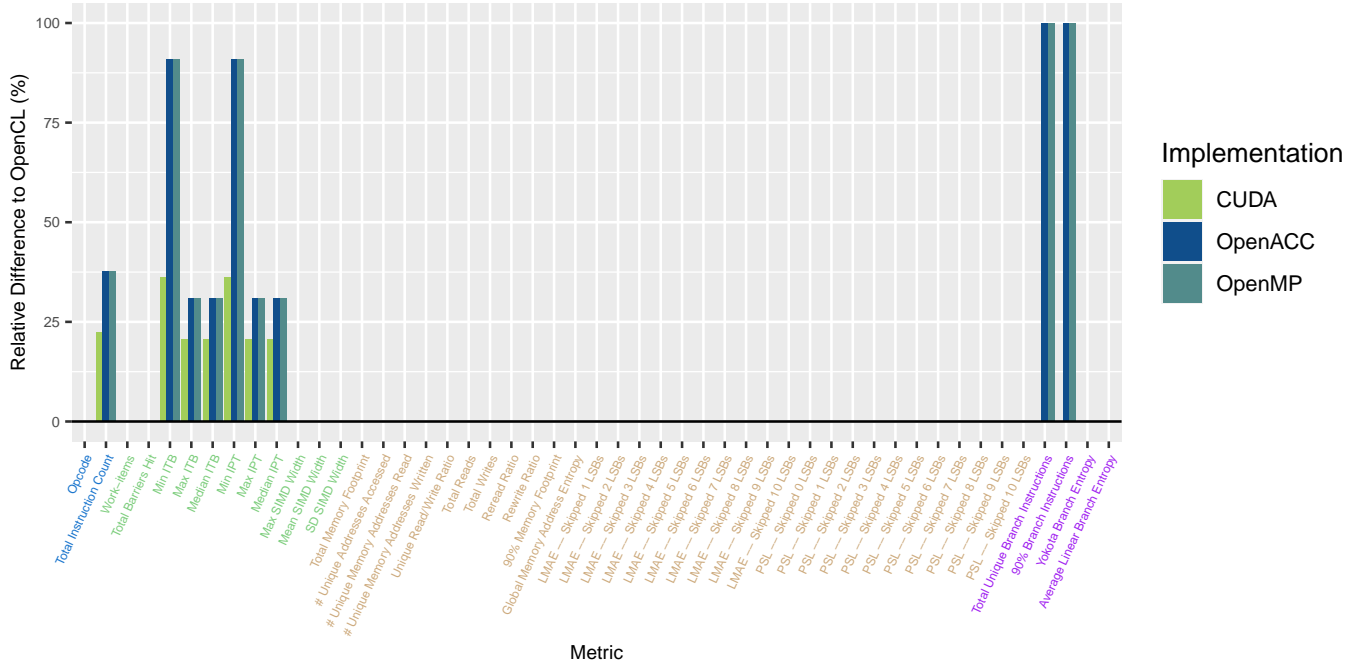


Figure 3: Relative difference in AIWC metrics between each translated Fan1 kernel implementation against the baseline OpenCL.

i: OpenACC

```
__kernel void Fan1_kernel0(__global float * a, __global
    ↪ float * m, int Size, int fan1_gangs, int
    ↪ fan1_workers, int t)
{
    int _ti_100_501;
    int lwpriv__i;
    _ti_100_501=get_global_id(0);
    #pragma acc kernels loop gang(fan1_gangs) worker(
        ↪ fan1_workers) independent copyin(Size, t) present
        ↪ (a[0:(Size*Size)], m[0:(Size*Size)]) private(i)
    if (_ti_100_501<(get_num_groups(0)*fan1_workers))
    {
        for (lwpriv__i=(_ti_100_501+0); lwpriv__i<((Size-1)-t); (
            ↪ lwpriv__i+=(get_num_groups(0)*fan1_workers)))
        {
            m[((Size*((lwpriv__i+t)+1))+t)]=a[((Size*((lwpriv__i+t)
                ↪ +1))+t)]/a[((Size*(t))+t)]];
        }
    }
}
```

ii: OpenMP

```
__kernel void Fan1_kernel0(int lfpriv__Size, int
    ↪ lfpriv__t, __global float * a, __global float * m
    ↪ , int fan1_teams, int fan1_threads)
{
    int _ti_100_501;
    int lwpriv__i;
    _ti_100_501=get_global_id(0);
    #pragma acc parallel loop num_workers(fan1_threads) gang
        ↪ worker present(a[0:(Size*Size)], m[0:(Size*Size)
        ↪ ]) private(i) firstprivate(Size, t) num_gangs(
        ↪ fan1_teams)
    if (_ti_100_501<(get_num_groups(0)*fan1_threads))
    {
        for (lwpriv__i=(_ti_100_501+0); lwpriv__i<((lfpriv__Size
            ↪ -1)-lfpriv__t); (lwpriv__i+=(get_num_groups(0)*
            ↪ fan1_threads)))
        {
            m[((lfpriv__Size*((lwpriv__i+lfpriv__t)+1))+lfpriv__t)]=
                ↪ a[((lfpriv__Size*((lwpriv__i+lfpriv__t)+1))+
                ↪ lfpriv__t)]/a[((lfpriv__Size*lfpriv__t)+lfpriv__t
                ↪ )]);
        }
    }
}
```

Listing 2: OpenCL kernel representation comparison to translator generated OpenACC and OpenMP kernels of Fan1.

i: OpenCL

ii: CUDA

```

1 ; Function Attrs: nounwind
2 define void @Fan1(float* nocapture %m_dev, float* ←
    ↪ nocapture readonly %a_dev, float* nocapture ←
    ↪ readnone %b_dev, i32 %size, i32 %t) ←
    ↪ local_unnamed_addr #0 !kernel_arg_addr_space !1←
    ↪ !kernel_arg_access_qual !2 !kernel_arg_type !3←
    ↪ !kernel_arg_base_type !3 !kernel_arg_type_qual←
    ↪ !4 {
3   %call = tail call i32 @i32 (i32, ...) bitcast (i32 (...)*) ←
    ↪ @get_local_id to i32 (i32, ...)*(i32 0) #3
4   %call1 = tail call i32 (i32, ...) bitcast (i32 (...)*)←
    ↪ @get_group_id to i32 (i32, ...)*(i32 0) #3
5   %call2 = tail call i32 (i32, ...) bitcast (i32 (...)*)←
    ↪ @get_local_size to i32 (i32, ...)*(i32 0) ←
    ↪ #3
6   %mul = mul nsw i32 %call2, %call1
7   %add = add nsw i32 %mul, %call
8   %sub = add nsw i32 %size, -1
9   %sub3 = sub i32 %sub, %t
10  %cmp = icmp slt i32 %add, %sub3
11  br i1 %cmp, label %1, label %4
12
13 ; <label>:1: ; preds = %0
14  %add4 = add i32 %t, 1
15  %add5 = add i32 %add4, %add
16  %mul6 = mul nsw i32 %add5, %size
17  %idx.ext = sext i32 %mul6 to i64
18  %add.ptr = getelementptr inbounds float, float* ←
    ↪ %a_dev, i64 %idx.ext
19  %idx.ext7 = sext i32 %t to i64
20  %add.ptr8 = getelementptr inbounds float, float* %add←
    ↪ .ptr, i64 %idx.ext7
21  %2 = load float, float* %add.ptr8, align 4, !tbaa !5
22  %mul9 = mul nsw i32 %t, %size
23  %idx.ext10 = sext i32 %mul9 to i64
24  %add.ptr11 = getelementptr inbounds float, float* ←
    ↪ %a_dev, i64 %idx.ext10
25  %add.ptr13 = getelementptr inbounds float, float* ←
    ↪ %add.ptr11, i64 %idx.ext7
26  %3 = load float, float* %add.ptr13, align 4, !tbaa !5
27  %div = fdiv float %2, %3, !fpmath !9
28  %add.ptr18 = getelementptr inbounds float, float* ←
    ↪ %m_dev, i64 %idx.ext
29  %add.ptr20 = getelementptr inbounds float, float* ←
    ↪ %add.ptr18, i64 %idx.ext7
30  store float %div, float* %add.ptr20, align 4, !tbaa ←
    ↪ !5
31  br label %4
32
33 ; <label>:4: ; preds = %1, %0
34  ret void
35 }

```

```

1 ; Function Attrs: nounwind
2 define void @_Z4Fan1PfS_ii(float* nocapture, float* ←
    ↪ nocapture readonly, i32, i32) ←
    ↪ local_unnamed_addr #1 {
3   %5 = tail call i32 @llvm.nvvm.read.ptx.sreg.tid.x() ←
    ↪ #3, !range !5 @thread_id
4   %6 = tail call i32 @llvm.nvvm.read.ptx.sreg.ctaid.x()←
    ↪ #3, !range !6 @block_id
5   %7 = tail call i32 @llvm.nvvm.read.ptx.sreg.ntid.x() ←
    ↪ #3, !range !7 @block_dim
6   %8 = mul i32 %7, %6
7   %9 = add i32 %8, %5
8   %10 = add nsw i32 %2, -1
9   %11 = sub i32 %10, %3
10  %12 = icmp slt i32 %9, %11
11  br i1 %12, label %13, label %30
12
13 ; <label>:13: ; preds = %4
14  %14 = add i32 %3, 1
15  %15 = add i32 %14, %9
16  %16 = mul nsw i32 %15, %2
17  %17 = sext i32 %16 to i64
18  %18 = getelementptr inbounds float, float* %1, i64 ←
    ↪ %17
19  %19 = sext i32 %3 to i64
20  %20 = getelementptr inbounds float, float* %18, i64 ←
    ↪ %19
21  %21 = load float, float* %20, align 4, !tbaa !8
22  %22 = mul nsw i32 %3, %2
23  %23 = sext i32 %22 to i64
24  %24 = getelementptr inbounds float, float* %1, i64 ←
    ↪ %23
25  %25 = getelementptr inbounds float, float* %24, i64 ←
    ↪ %19
26  %26 = load float, float* %25, align 4, !tbaa !8
27  %27 = fdiv float %21, %26
28  %28 = getelementptr inbounds float, float* %0, i64 ←
    ↪ %17
29  %29 = getelementptr inbounds float, float* %28, i64 ←
    ↪ %19
30  store float %27, float* %29, align 4, !tbaa !8
31  br label %30
32
33 ; <label>:30: ; preds = %13, %4
34  ret void
35 }

```

Listing 3: OpenCL compared to the CUDA implementations generated LLVM-IR/SPIR of the Fan1 kernel.

on a separate pthread – we only print the log if it occurs on the first thread. The default Gaussian Elimination test data is run on 4 threads and calls the `Fan1` and `Fan2` kernels three (3) times. For this analysis we only store the traces of first execution of the `Fan1` kernel. These traces were then piped from each of the implementations.

The differences between traces are shown in Listing 5. The OpenCL trace is shown in Listing 5-i and presents the baseline progression of instructions expected, ii is the CUDA trace, OpenACC in iii and OpenMP in iv. Each

trace should be read as the LLVM instruction executed over time as we proceed down the Listing. Blank lines have been inserted to align common instructions in the trace between implementations, this is to present the clearest difference between traces. Instructions of interest have also been coloured – red indicates added instructions no apparent in the baseline OpenCL trace, blue instructions show a reordering of instructions between traces and olive shows substitution (or deviation) of instructions. The CUDA trace shows that each

i: OpenCL

```

1 ; Function Attrs: nounwind
2 define void @Fan1(float* nocapture %m_dev, float* ←
    → nocapture readonly %a_dev, float* nocapture ←
    → readnone %b_dev, i32 %size, i32 %t) ←
    → local_unnamed_addr #0 !kernel_arg_addr_space !1←
    → !kernel_arg_access_qual !2 !kernel_arg_type !3←
    → !kernel_arg_base_type !3 !kernel_arg_type_qual←
    → !4 {
3   %call = tail call i32 (i32, ...) bitcast (i32 (...)*) ←
    → @get_local_id to i32 (i32, ...)*(i32 0) #3
4   %call1 = tail call i32 (i32, ...) bitcast (i32 (...)*) ←
    → @get_group_id to i32 (i32, ...)*(i32 0) #3
5   %call2 = tail call i32 (i32, ...) bitcast (i32 (...)*) ←
    → @get_local_size to i32 (i32, ...)*(i32 0) ←
    → #3
6   %mul = mul nsw i32 %call2, %call1
7   %add = add nsw i32 %mul, %call
8   %sub = add nsw i32 %size, -1
9   %sub3 = sub i32 %sub, %t
10  %cmp = icmp slt i32 %add, %sub3
11  br i1 %cmp, label %1, label %4
12
13 ; <label>:1: ; preds = %0
14  %add4 = add i32 %t, 1
15  %add5 = add i32 %add4, %add
16  %mul6 = mul nsw i32 %add5, %size
17  %idx.ext = sext i32 %mul6 to i64
18  %add.ptr = getelementptr inbounds float, float* ←
    → %a_dev, i64 %idx.ext
19  %idx.ext7 = sext i32 %t to i64
20  %add.ptr8 = getelementptr inbounds float, float* %add←
    → .ptr, i64 %idx.ext7
21  %2 = load float, float* %add.ptr8, align 4, !tbaa !5
22  %mul9 = mul nsw i32 %t, %size
23  %idx.ext10 = sext i32 %mul9 to i64
24  %add.ptr11 = getelementptr inbounds float, float* ←
    → %a_dev, i64 %idx.ext10
25  %add.ptr13 = getelementptr inbounds float, float* ←
    → %add.ptr11, i64 %idx.ext7
26  %3 = load float, float* %add.ptr13, align 4, !tbaa !5
27  %div = fdiv float %2, %3, !fpmath !9
28  %add.ptr18 = getelementptr inbounds float, float* ←
    → %m_dev, i64 %idx.ext
29  %add.ptr20 = getelementptr inbounds float, float* ←
    → %add.ptr18, i64 %idx.ext7
30  store float %div, float* %add.ptr20, align 4, !tbaa ←
    → !5
31  br label %4
32
33 ; <label>:4: ; preds = %1, %0
34  ret void
35 }

```

ii: OpenACC and OpenMP

```

1 ; Function Attrs: nounwind
2 define void @Fan1_kernel0(float* nocapture readonly %a,←
    → float* nocapture %m, i32 %Size, i32 ←
    → %fan1_gangs, i32 %fan1_workers, i32 %t) ←
    → local_unnamed_addr #0 !kernel_arg_addr_space !1←
    → !kernel_arg_access_qual !2 !kernel_arg_type !3←
    → !kernel_arg_base_type !3 !kernel_arg_type_qual←
    → !4 {
3   %call = tail call i32 (i32, ...) bitcast (i32 (...)*) ←
    → * @get_global_id to i32 (i32, ...)*(i32 0) ←
    → #3
4   %call1 = tail call i32 (i32, ...) bitcast (i32 ←
    → (...)* @get_num_groups to i32 (i32, ...)*(←
    → i32 0) #3
5   %mul = mul nsw i32 %call1, %fan1_workers
6   %cmp = icmp slt i32 %call, %mul
7   br i1 %cmp, label %.preheader, label %.loopexit
8
9   .preheader: ; preds = %0
10  %sub = add nsw i32 %Size, -1
11  %sub2 = sub i32 %sub, %t
12  %cmp338 = icmp slt i32 %call, %sub2
13  br i1 %cmp338, label %.lr.ph, label %.loopexit
14
15  .lr.ph: ; preds = %.preheader
16  %add4 = add i32 %t, 1
17  %mul8 = mul nsw i32 %t, %Size
18  %add9 = add nsw i32 %mul8, %t
19  %idxprom10 = sext i32 %add9 to i64
20  %arrayidx11 = getelementptr inbounds float, float* ←
    → %a, i64 %idxprom10
21  br label %1
22
23  ; <label>:1: ; preds = %.lr.ph, %1
24  %lwpriv__i.039 = phi i32 [ %call, %.lr.ph ], [ ←
    → %add20, %1 ]
25  %add5 = add i32 %add4, %lwpriv__i.039
26  %mul6 = mul nsw i32 %add5, %Size
27  %add7 = add nsw i32 %mul6, %t
28  %idxprom = sext i32 %add7 to i64
29  %arrayidx = getelementptr inbounds float, float* %a←
    → , i64 %idxprom
30  %2 = load float, float* %arrayidx, align 4, !tbaa ←
    → !5
31  %3 = load float, float* %arrayidx11, align 4, !tbaa←
    → !5
32  %div = fdiv float %2, %3, !fpmath !9
33  %arrayidx17 = getelementptr inbounds float, float* ←
    → %m, i64 %idxprom
34  store float %div, float* %arrayidx17, align 4, !←
    → tbaa !5
35  %call18 = tail call i32 (i32, ...) bitcast (i32 ←
    → (...)* @get_num_groups to i32 (i32, ...)*(←
    → i32 0) #3
36  %mul19 = mul nsw i32 %call18, %fan1_workers
37  %add20 = add nsw i32 %mul19, %lwpriv__i.039
38  %cmp3 = icmp slt i32 %add20, %sub2
39  br i1 %cmp3, label %1, label %.loopexit.loopexit
40
41  .loopexit.loopexit: ; preds = %1
42  br label %.loopexit
43
44  .loopexit: ; preds = %.loopexit.loopexit, %←
    → preheader, %0
45  ret void
46 }

```

Listing 4: OpenCL compared to the OpenMP and OpenACC implementations generated LLVM-IR/SPIR of the Fan1 kernel.

i: OpenCL	ii: CUDA	iii: OpenACC	iv: OpenMP
call	getelementptr bitcast call trunc	call trunc sext call	call trunc sext call
call	call trunc call trunc	sext call	sext call
mul	mul	mul	mul
add	add		
trunc			
add	add		
sub	sub		
icmp	icmp	icmp	icmp
br	br	br	br
	getelementptr bitcast		
add	add	add sub	add sub
add	add	add	add
mul	mul	mul	mul
sext	sext	add sext	add sext
getelementptr	getelementptr	getelementptr	getelementptr
sext	sext	br	br
load	load	phi icmp	phi icmp
		br	br
		add	add
mul	mul	mul	mul
sext	sext	add sext	add sext
getelementptr	getelementptr	getelementptr	getelementptr
getelementptr	getelementptr	load	load
load	load	load	load
fdiv	fdiv	fdiv	fdiv
getelementptr	getelementptr	getelementptr	getelementptr
getelementptr	getelementptr		
store	store	store	store
		zext	zext
		add	add
		trunc	trunc
		br	br
		phi	phi
		icmp	icmp
		br	br
br	br	br	br
ret	ret	ret	ret

Listing 5: Trace of instructions executed by thread 0 of Fan1 kernel for each of the language implementations.

tie in instructions added with each memory lookup from the SPIR

. The OpenACC trace has no instruction reordering but has instructions added to compensate for the different control flow (looping) to support the workitems in a workgroup logic – as was described in Section ?? . There is no difference in traces between OpenMP and OpenACC traces because it uses the same OpenARC compiler toolchain.

E. Fan2

Thus we can use example to illustrate how AIWC metrics highlight discrepancies between languages and implementations and how it can be used to guide sourcecode changes. Two separate loops in the **Fan2** function were consolidated into one, to mirror how the task is performed in the OpenCL and CUDA implementations of the algorithm.

VII. CONCLUSIONS AND FUTURE WORK

This work extends the applicability of AIWC by supporting multiple languages and have demonstrated it's usefulness

to evaluate the overhead and complexities of the OpenCL output from two source code translation tools. We believe this methodology can be boardly more useful in the future development of translators.

Since AIWC metrics are based on an abstract/ideal OpenCL device simulator there are architecture-specific optimizations that may occur when you target the LLVM-IR/SPIR to a specific accelerator however this is expected to be consistant regardless of the compiler front-end, and the former issue is not a subject for this work.

We will apply the methodology of generating AIWC feature-spaces to other languages, specifically OpenMP and OpenACC – since this is the typical means of accelerating conventional HPC workloads using many-core CPUs. The ability to examine the characteristics of these kernels in large code-bases allow optimization work to occur on these kernels, effectively guiding a developers hand to optimize a code by providing strategies to minimise certain AIWC metrics shown to be advantageous to specific accelerators. Additionally, examining the AIWC features of existing code-bases may facilitate identifying a better accelerator match as they become available.

Thus, this work identifies the potential overheads when translating between functionally identical implementations of kernels written in different languages by examining differences in their respective AIWC metrics. We also offer a methodology which uses AIWC over a number of tests kernels, and against an OpenCL implementation as a base-line, is an option to assess the suitability of any changes made to the translator. This work will assist the improvement of the translation tools on offer, increasing the adoption of SPIR-V and the use of the OpenCL runtime behind-the-scenes, resulting in less fragmentation between software models and languages on contemporary HPC systems.

This paper shows that the visual inspection of AIWC metrics facilitates a high-level (and quick) overview of computational characteristics of kernels, and we’ve found that how they change has enabled us to compare the execution behaviour of codes in response to two different compilers performing translation. We’ve seen how source-code modifications in our selected benchmark kernels change these features – in our instance to more closely resemble an OpenCL baseline. We believe the same methodology will be useful for compiler engineers to evaluate their own translators – especially with the increasing use of LLVM as a backend, on which AIWC and this approach is based. This is useful since it is abstracted to an ideal OpenCL device and, as such, is free from micro-architecture and architectural details. We propose this methodology will also encourage application developers of scientific codes to take a deeper-dive into their codes, and more generally, our future work will examine how AIWC metrics provide a developer insights around the suitability of the kernel code

when initially selecting for, then optimizing on, a specific accelerators.

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