Mereology

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1 Introduction

 $\begin{array}{l} \textbf{theory} \ \textit{Mereology} \\ \textbf{imports} \ \textit{Main} \\ \textbf{begin} \end{array}$

This is a presentation in Isabelle/HOL of *Classical Extensional Mereology*. The presentation is based on those in "Parts" by Peter Simons [3] and "Parts and Places" by Roberto Casati and Achille Varzi [1]. Some corrections and important proofs are from [2]

Please note that this is an extremely ROUGH DRAFT.

2 Ground Mereology

Ground Mereology (M) introduces parthood as a primitive relation amongst individuals. It's assumed that parthood is a partial ordering relation - that is reflexive, symmetric and transitive [1], p. 36:

```
typedecl i — the type of individuals
locale M =
  fixes P:: i \Rightarrow i \Rightarrow bool (infix \leq 50)
  assumes R: x \leq x — reflexivity of parthood
    and AS: x \leq y \longrightarrow y \leq x \longrightarrow x = y — antisymmetry of parthood
    \textbf{and} \quad T \hbox{:} \ x \leq y \longrightarrow y \leq z \longrightarrow x \leq z - \text{transitivity of parthood}
begin
The following relations are defined in terms of parthood [1], p. 36-7:
definition PP:: i \Rightarrow bool (infix < 50)
  where x < y \equiv x \le y \land \neg y \le x — proper parthood
definition O:: i \Rightarrow i \Rightarrow bool(O)
  where O \ x \ y \equiv \exists \ z. \ z \le x \land z \le y — overlap
definition D:: i \Rightarrow i \Rightarrow bool(D)
  where D x y \equiv \neg O x y — disjointness
definition U:: i \Rightarrow i \Rightarrow bool(U)
  where U x y \equiv \exists z. x \leq z \land y \leq z—underlap
As are the following operations on individuals [1], p. 43-5:
definition S:: i \Rightarrow i \Rightarrow i \text{ (infix } + 52)
  where x + y \equiv THE \ z. \forall \ w. O \ w \ z \longleftrightarrow O \ w \ x \lor O \ w \ y — sum or fusion
definition T:: i \Rightarrow i \Rightarrow i \text{ (infix} \times 53)
 where x \times y \equiv THE z. \forall w. O w z \longleftrightarrow O w x \land O w y — product or intersection
definition u:: i (u)
  where u \equiv THE z. \forall w. P w z — universe
definition M:: i \Rightarrow i \Rightarrow i \text{ (infix } -51)
  where x - y \equiv THE z. \forall w. O w z \longleftrightarrow O w x \land \neg O w y — difference
definition C:: i \Rightarrow i (\neg)
  where \neg x \equiv (u - x) — complement
And the operations of general sum and product [1], p. 46:
definition \sigma:: (i \Rightarrow bool) \Rightarrow i (\sigma)
  where \sigma F \equiv THE z. (\forall y. O y z \longleftrightarrow (\exists x. F x \land O x y))
abbreviation \sigma x:: (i \Rightarrow bool) \Rightarrow i (binder \sigma [8] 9)
  where \sigma x. F x \equiv \sigma F — general sum or fusion of the Fs
definition \pi:: (i \Rightarrow bool) \Rightarrow i (\pi)
  where \pi F \equiv THE z. (\forall x. F x \longrightarrow z \le x) — general products [1], p. 46
abbreviation \pi x:: (i \Rightarrow bool) \Rightarrow i (binder \pi [8] 9)
  where \pi x. F x \equiv \pi F — general sum or product of the Fs
```

Note that the symbols for part, proper part, sum, product, difference and complement are distinguished by bold font.

end

3 Minimal Mereology

Minimal mereology (MM) adds to ground mereology the axiom of weak supplementation [1], p. 39:

```
locale MM = M + assumes WS: x < y \longrightarrow (\exists z. z < y \land D z x) — weak supplementation
```

Weak supplementation is sometimes stated with parthood rather than proper parthood in the consequent. The following lemma in ground mereology shows that the two definitions are equivalent, given anti-symmetry:

```
lemma (in M) (x < y \longrightarrow (\exists z. z < y \land D z x)) \longleftrightarrow (x < y \longrightarrow (\exists z. z \le y \land D z x))
by (metis\ AS\ D\text{-}def\ O\text{-}def\ PP\text{-}def\ R)
```

The following two lemmas are weaker supplementation principles taken from Simons [3], p. 27. The names *company* and *strong company* are from Varzi's *Stanford Encyclopedia of Philosophy* entry on mereology [4].

```
lemma (in MM) C: x < y \longrightarrow (\exists z. z \neq x \land z < y) by (metis WS D-def O-def R) — company lemma (in MM) SC: x < y \longrightarrow (\exists z. z < y \land \neg z \leq x) by (metis WS D-def O-def R) — strong company
```

Minimal Mereology is not strong enough to proved the *Proper Parts Principle*, according to which if x has a proper part, and every proper part of x is a part of y, then x is a part of y [3] p. 28:

```
lemma (in MM) PPP: \exists z. z < x \Longrightarrow \forall z. z < x \longrightarrow z \le y \Longrightarrow x \le y—proper parts principle nitpick [user-axioms] oops
```

The proper parts principle is Simons way of expressing *extensionality*, which is not a theorem of Minimal Mereology either:

```
lemma (in M) E: (\exists z. z < x \lor z < y) \longrightarrow (\forall z. z < x \longleftrightarrow z < y) \longrightarrow x = y — extensionality nitpick oops
```

The failure of weak supplementation to entail the proper parts principle or extensionality motivates a stronger axiom, to which we turn in the next section.

4 Extensional Mereology

Extensional Mereology (EM) adds the axiom of strong supplementation [1], p. 39:

```
locale EM = M + assumes SS: \neg x \le y \longrightarrow (\exists z. z \le x \land D z y) — strong supplementation
```

Extensional Mereology (EM is so called because it entails the proper parts principle [3] p. 29:

```
lemma (in EM) PPP: \exists z. z < x \Longrightarrow \forall z. z < x \longrightarrow z \le y \Longrightarrow x \le y by (metis SS D-def R O-def PP-def T)
```

And thus extensionality proper [1] p. 40:

```
lemma (in EM)
E: (\exists \ z. \ z < x \lor z < y) \longrightarrow (\forall \ z. \ z < x \longleftrightarrow z < y) \longrightarrow x = y - \text{extensionality}
by (metis R O-def D-def AS PP-def SS)
```

In the context of the other axioms, strong supplementation entails weak supplementation [3], p. 29:

```
lemma (in M) SStoWS:

assumes SS: \bigwedge x. \bigwedge y. \neg x \leq y \longrightarrow (\exists z. z \leq x \land D z y)

— assumes strong supplementation

shows WS: \bigwedge x. \bigwedge y. x < y \longrightarrow (\exists z. z < y \land D z x)

— shows weak supplementation

by (metis AS D-def O-def PP-def R assms)
```

But not vice versa:

```
lemma (in M) WStoSS:
assumes WS: \bigwedge x. \bigwedge y. \ x < y \longrightarrow (\exists \ z. \ z < y \land D \ z \ x)
— assumes weak supplementation
shows SS: \bigwedge x. \bigwedge y. \neg x \le y \longrightarrow (\exists \ z. \ z \le x \land D \ z \ y)
— shows strong supplementation
nitpick oops
```

So Extensional Mereology is stronger than Minimal Mereology [1] p. 43:

sublocale $EM \subseteq MM$ **using** T SS SSto WS **by** (metis MM.intro MM-axioms.intro M-axioms)

sublocale $MM \subseteq EM$ nitpick oops

```
lemma (in MM) assumes PPP: \exists z. z < x \Longrightarrow \forall z. z < x \longrightarrow z \le y \Longrightarrow x \le y shows SS: \neg x \le y \longrightarrow (\exists z. z \le x \land D z y) nitpick oops
```

5 Closure Mereology

Closure Mereology adds to Ground Mereology the axioms of sum closure and product closure [1] p. 43:

```
locale CM = M + assumes SC: U \times y \longrightarrow (\exists z. \forall w. O \times z \longleftrightarrow (O \times x \vee O \times y))— sum closure assumes PC: O \times y \longrightarrow (\exists z. \forall w. w \leq z \longleftrightarrow (w \leq x \land w \leq y))— product closure
```

Combining Closure Mereology with Minimal Mereology yields the theory known as Closure Minimal Mereology *CMM* whereas combining Closure Mereology with Extensional Mereology obtains *Closed Extensional Mereology CEM* [1] p. 43:

```
locale CMM = CM + MM
locale CEM = EM + CM
```

In Closed Minimal Mereology, the product closure axiom and weak supplementation jointly entail strong supplementation. The proof verified here is from Pontow [2] p. 200:

```
lemma (in CMM) SS: \neg x \leq y \longrightarrow (\exists z. z \leq x \land D z y)
proof fix x y
  assume \neg x \leq y
 \mathbf{show}\ (\exists\ z.\ z\leq x\,\wedge\,D\,\,z\,\,y)
 proof cases
    assume D x y
    thus (\exists z. z \leq x \land D z y) using R by auto
    assume \neg D x y
    hence O x y using D-def by simp
   hence \exists z. \forall w. w \leq z \longleftrightarrow (w \leq x \land w \leq y) using PC by simp
    then obtain z where z: \forall w. w \leq z \longleftrightarrow (w \leq x \land w \leq y)..
    hence z < x using PP-def R \leftarrow x \le y by auto
    hence (\exists w. w < x \land D w z) using WS by simp
    then obtain w where w < x \wedge D w z..
    hence w \le x \land D \ w \ y \ \mathbf{by} \ (meson \ D\text{-}def \ O\text{-}def \ PP\text{-}def \ T \ z)
    thus (\exists z. z \leq x \land D z y)..
  qed
qed
```

Because Strong Supplementation is provable in Closed Minimal Mereology, it follows that Closed Extensional Mereology and Closed Minimal Mereology are the same theory [1] p. 44:

```
sublocale CEM \subseteq CMM by (simp\ add:\ CMM.intro\ CM-axioms\ MM-axioms) sublocale CMM \subseteq CEM by (simp\ add:\ CEM.intro\ CM-axioms\ EM.intro\ EM-axioms.intro\ M-axioms\ SS)
```

Closure Mereology with Universe (CMU) is obtained by adding an axiom ensuring existence of a universe [1] p. 44:

```
locale CMU = CM +
assumes U: \exists z. \forall x. x \leq z— universe
```

And adding Extensional Mereology (or Minimal Mereology) to this theory results in Closed Extensional Mereology with Universe *CEMU*:

```
locale CEMU = EM + CMU
```

In Closure Extensional Mereology with Universe, it's possible to derive a strengthening of the sum axiom, since everything underlaps everything else:

```
lemma (in CEMU) EU: U \times y using U-def U by auto — everything underlaps lemma (in CEMU) SSC: (\exists z. \forall w. O w z \longleftrightarrow (O w x \lor O w y)) using EU SC by simp
```

— strengthened sum closure

6 General Mereology

General Mereology is obtained from Ground Mereology by adding the axiom of fusion or unrestricted composition [1] p. 46:

```
locale GM = M + assumes F: (\exists x. Fx) \longrightarrow (\exists z. \forall y. Oyz \longleftrightarrow (\exists x. Fx \land Oxy)) — fusion or unrestricted composition
```

Substituting $x = a \lor x = b$ for F x in the fusion axiom allows the derivation of an unrestricted version of sum closure GS, and so of course sum closure itself, as follows:

```
lemma (in GM) FS: (\exists \ x. \ (x = a \lor x = b)) \longrightarrow (\exists \ z. \ \forall \ y. \ O \ y \ z \longleftrightarrow (\exists \ x. \ (x = a \lor x = b) \land O \ x \ y)) using F solve-direct. lemma (in GM) GFS: (\exists \ z. \ \forall \ y. \ O \ y \ z \longleftrightarrow (\exists \ x. \ (x = a \lor x = b) \land O \ x \ y)) using FS by blast lemma (in M) GFStoGS: assumes GFS: (\exists \ z. \ \forall \ y. \ O \ y \ z \longleftrightarrow (\exists \ x. \ (x = a \lor x = b) \land O \ x \ y)) shows (\exists \ z. \ \forall \ w. \ O \ w \ z \longleftrightarrow (O \ w \ a \lor O \ w \ b)) by (metis \ O - def \ GFS) lemma (in GM) GS: (\exists \ z. \ \forall \ w. \ O \ w \ z \longleftrightarrow (O \ w \ x \lor O \ w \ y)) using GFS \ GFStoGS by simp lemma (in GM) S: U \ x \ y \longrightarrow (\exists \ z. \ \forall \ w. \ O \ w \ z \longleftrightarrow (O \ w \ x \lor O \ w \ y)) using GS \ by simp
```

But product closure cannot be derived:

```
lemma (in GM) T: O x y \longrightarrow (\exists z. \forall w. w \le z \longleftrightarrow (w \le x \land w \le y)) nitpick [show-all] oops
```

It follows that General Mereology does not encompass Closure Mereology, contrary to Simons [3] p. 36 and Casati and Varzi [1] p. 46 (this point is discussed in detail by Pontow [2]:

```
sublocale GM \subseteq CM nitpick [show-all] oops
```

It's possible to prove from fusion in General Mereology that there is something that overlaps everything:

```
lemma (in GM) \exists z. \forall x. O x z — something overlaps everything proof — have (\exists x. x = x) \longrightarrow (\exists z. \forall y. O y z \longleftrightarrow (\exists x. x = x \land O x y)) using F by fast hence \exists z. \forall y. O y z \longleftrightarrow (\exists x. x = x \land O x y) by simp hence \exists z. \forall y. O y z \longleftrightarrow (\exists x. O x y) by simp thus ?thesis by (metis O-def R) qed
```

But it doesn't follow that there is a universe, let alone a unique universe. If for example, there is just an infinite ascending chain, then everything overlaps everything else, but there isn't a particular thing which everything is a part of, since for anything in particular, the things above it are not part of the chain:

```
lemma (in \mathit{GM}) \mathit{U} \colon \exists \ \mathit{z} . \ \forall \ \mathit{x} . \ \mathit{x} \leq \mathit{z} \ \text{nitpick oops}
```

The existence of differences is not guaranteed either:

```
lemma (in GM) D: (\exists w. w \le x \land \neg O w y) \longrightarrow (\exists z. \forall w. w \le z \longleftrightarrow (w \le x \land \neg O w y)) nitpick oops
```

Call the combination of General Mereology with weak supplementation General Minimal Mereology, or *GMM*:

```
locale GMM = MM + GM — General Minimal Mereology
```

Although Strong Supplementation can be derived from Weak Supplementation in Closed Minimal Mereology via the product axioms, it cannot be derived in General Minimal Mereology, since the product axiom itself still cannot be derived in General Minimal Mereology:

```
lemma (in GMM) SS: \neg x \leq y \longrightarrow (\exists z. z \leq x \land D z y) nitpick oops lemma (in GMM) T: O x y \longrightarrow (\exists z. \forall w. w \leq z \longleftrightarrow (w \leq x \land w \leq y)) nitpick [show-all] oops
```

Nor can the existence of a universe or differences be proved in General Minimal Mereology:

```
lemma (in GMM) U: \exists z. \forall x. x \leq z nitpick oops lemma (in GMM) D: (\exists w. w \leq x \land \neg O w y) \longrightarrow (\exists z. \forall w. w \leq z \longleftrightarrow (w \leq x \land \neg O w y)) nitpick oops
```

7 Classical Extensional Mereology

Classical Extensional Mereology GEM is simply Extensional Mereology combined with General Mereology [1] p. 46:

```
locale GEM = EM + GM
```

The presence of strong supplementation in Classical Extensional Mereology enables the derivation of product closure from fusion. The following proof is from Pontow [2] pp. 202-3.

The proof begins by substitutions $z \le a \land z \le b$ for F in the fusion axiom, to give the existence of a sum of all the parts of a and b:

```
lemma (in GM) FP: (\exists z. (z \le a \land z \le b)) \longrightarrow (\exists z. \forall w. O w z \longleftrightarrow (\exists x. (x \le a \land x \le b) \land O x w)) using F solve-direct.
```

Three lemmas are helpful before proceeding with the proof proper. First, strong supplementation is needed to proceed from the fact that z is a sum of the Fs to the fact that z is the sum of the Fs:

```
lemma (in EM) atothesum:

assumes asum: \forall y. O y z \longleftrightarrow (\exists x. F x \land O x y)

shows thesum: (THE \ v. \ \forall y. O y \ v \longleftrightarrow (\exists x. F x \land O x y)) = z

proof (rule \ the\text{-}equality)

show \forall y. O y z \longleftrightarrow (\exists x. F x \land O x y) using asum.

show \land v. \ \forall y. O y \ v = (\exists x. F x \land O x y) \Longrightarrow v = z by (metis \ SS \ AS \ D\text{-}def \ O\text{-}def \ R \ asum)

qed
```

Using this lemma, we can show that if something overlaps z just in case it overlaps an F, then it is the sum of the Fs:

```
lemma (in EM) UGS: (\forall y. \ O\ y\ z \longleftrightarrow (\exists\ x.\ F\ x \land O\ x\ y)) \longrightarrow (\sigma\ x.\ F\ x) = z proof assume (\forall\ y.\ O\ y\ z \longleftrightarrow (\exists\ x.\ F\ x \land O\ x\ y)) hence (THE\ v.\ \forall\ y.\ O\ y\ v \longleftrightarrow (\exists\ x.\ F\ x \land O\ x\ y)) = z using atothesum by simp thus (\sigma\ v.\ F\ v) = z using \sigma-def by blast qed
```

With this lemma in hand, we can proceed with a final lemma the proof from Pontow [2] pp. 202-3, according to which if there is an F, then everything is part of the sum of the Fs just in case every part of it overlaps with an F.

```
lemma (in GEM) PS: (\exists x. Fx) \longrightarrow (\forall y. y \le (\sigma v. Fv) \longleftrightarrow (\forall w. w \le y) \longrightarrow (\exists v. Fv \land Ovw)))
proof
assume (\exists x. Fx)
hence \exists z. \forall y. Oyz \longleftrightarrow (\exists x. Fx \land Oxy) using F by simp
then obtain z where z: \forall y. Oyz \longleftrightarrow (\exists x. Fx \land Oxy)..
```

```
hence \sigma: (\sigma \ v. \ F \ v) = z \ using \ UGS \ by \ simp
  show \forall y. y \leq (\sigma v. F v) \longleftrightarrow (\forall w. w \leq y \longrightarrow (\exists v. F v \land O v w))
  proof
    \mathbf{fix} \ y
    \mathbf{show}\ y \leq (\sigma\ v.\ F\ v) \longleftrightarrow (\forall\ w.\ w \leq y \longrightarrow (\exists\ v.\ F\ v \land\ O\ v\ w))
       proof
         assume y \leq (\sigma \ v. \ F \ v)
         hence y \leq z using \sigma by simp
         hence O y z using O-def R by auto
         hence (\exists x. Fx \land Oxy) using z by simp
        thus (\forall w. w \leq y \longrightarrow (\exists v. F v \land O v w)) by (metis O - def R T \lor y \leq z) z)
         \mathbf{assume}\ (\forall\ w.\ w\leq y\longrightarrow (\exists\ v.\ F\ v\ \wedge\ O\ v\ w))
         hence y \le z using z by (meson D-def SS)
         thus y \leq (\sigma \ v. \ F \ v) using \sigma by simp
       qed
    qed
  qed
Continuing to follow the proof from [2] pp. 204, we can prove the Product
Axiom proper:
lemma (in GEM) T: O \times y \longrightarrow (\exists z. \forall w. w \leq z \longleftrightarrow (w \leq x \land w \leq y))
proof
  assume O x y
  hence ez: (\exists z. (z \leq x \land z \leq y)) using O-def by simp
  hence (\exists z. \forall w. O \ w \ z \longleftrightarrow (\exists v. (v \le x \land v \le y) \land O \ v \ w)) using FP by
  then obtain z where (\forall w. O w z \longleftrightarrow (\exists v. (v \le x \land v \le y) \land O v w))..
  hence \sigma zxy: (\sigma \ v. \ v \leq x \land v \leq y) = z \text{ using } UGS \text{ by } simp
have gragra: (\forall \ s. \ s \leq (\sigma \ v. \ v \leq x \land v \leq y) \longleftrightarrow (\forall \ w. \ w \leq s \longrightarrow (\exists \ v. \ v \leq x \land v \leq y \land O \ v \ w))) using PS ez by simp
  have \forall w. w \leq z \longleftrightarrow (w \leq x \land w \leq y)
  proof
    \mathbf{fix}\ w
    show w \le z \longleftrightarrow (w \le x \land w \le y)
       proof
         assume w \leq z
         hence w \leq (\sigma \ v. \ v \leq x \land v \leq y) using \sigma zxy by simp
          hence dadada: (\forall t. t \leq w \longrightarrow (\exists v. v \leq x \land v \leq y \land O v t)) using
gragra by simp
         have \forall t. t \leq w \longrightarrow (O t x \land O t y)
         proof
            \mathbf{show}\ t \leq w \longrightarrow (O\ t\ x \wedge O\ t\ y)
              proof
                assume t \leq w
                hence (\exists \ \overline{v}.\ v \leq x \land v \leq y \land O\ v\ t) using dadada by simp
                thus O t x \wedge O t y using O-def T by blast
              qed
```

```
\begin{array}{l} \operatorname{qed} \\ \operatorname{thus} \ w \leq x \wedge w \leq y \ \operatorname{using} \ SS \ T \ D\text{-}def \ \operatorname{by} \ meson \\ \operatorname{next} \\ \operatorname{assume} \ w \leq x \wedge w \leq y \\ \operatorname{thus} \ w \leq z \ \operatorname{using} \ O\text{-}def \ R \ \sigma zxy \ gragra \ \operatorname{by} \ fastforce \\ \operatorname{qed} \\ \operatorname{qed} \\ \operatorname{thus} \ (\exists \ z. \ \forall \ w. \ w \leq z \longleftrightarrow (w \leq x \wedge w \leq y)).. \\ \operatorname{qed} \\ \operatorname{qed} \end{array}
```

It follows that General Extensional Mereology is stronger than Closed Extensional Mereology

sublocale $GEM \subseteq CEM$ using $CEM.intro\ CM.intro\ CM-axioms-def\ EM-axioms$ $M-axioms\ S\ T$ by blast

Likewise, substituting x = x for F x in fusion allows the derivation of the existence of a universe:

```
lemma (in GM) selfidentity:

(\exists \ x. \ x = x) \longrightarrow (\exists \ z. \ \forall \ y. \ O \ y \ z \longleftrightarrow (\exists \ x. \ x = x \land O \ x \ y))

using F by fast

lemma (in GM) (\exists \ z. \ \forall \ y. \ O \ y \ z \longleftrightarrow (\exists \ x. \ O \ x \ y)) using selfidentity by simp

lemma (in GEM) U: \exists \ z. \ \forall \ x. \ x \le z

using selfidentity by (metis \ D-def \ O-def \ SS)
```

It follows that Classical Extensional Mereology is also stronger than Closed Extensional Mereology with Universe:

```
sublocale GEM \subseteq CEMU
proof
show \exists z. \forall x. x \leq z using U by simp
qed
```

The existence of differences is also derivable in General Extensional Mereology. Like, the proof of the product axiom, the proof of the existence of differences is quite involved. It can be found in Pontow [2] p. 209.

```
lemma (in GM) FD: (\exists x. x \le a \land \neg O x b) \longrightarrow (\exists z. \forall y. O y z \longleftrightarrow (\exists x. (x \le a \land \neg O x b) \land O x y)) using F solve-direct.
```

```
lemma (in GEM) D: (\exists \ w. \ w \le x \land \neg \ O \ w \ y) \longrightarrow (\exists \ z. \ \forall \ w. \ w \le z \longleftrightarrow (w \le x \land \neg \ O \ w \ y)) proof assume (\exists \ w. \ w \le x \land \neg \ O \ w \ y) hence (\exists \ z. \ \forall \ w. \ O \ w \ z \longleftrightarrow (\exists \ v. \ (v \le x \land \neg \ O \ v \ y) \land \ O \ v \ w)) using FD by simp then obtain \Sigma where \Sigma: \forall \ w. \ O \ w \ \Sigma \longleftrightarrow (\exists \ v. \ (v \le x \land \neg \ O \ v \ y) \land \ O \ v \ w). have \forall \ w. \ w \le \Sigma \longleftrightarrow (w \le x \land \neg \ O \ w \ y) proof fix w show w \le \Sigma \longleftrightarrow (w \le x \land \neg \ O \ w \ y)
```

```
proof
        assume w \leq \Sigma
        have \forall z. z \leq w \longrightarrow Ozx
           \mathbf{show}\ z \leq w \,\longrightarrow\, O\ z\ x
            proof
               assume z \leq w
                hence \exists s0. (s0 \le x \land \neg O s0 y) \land O s0 z  using M.T M-axioms
O-def R \Sigma \langle w \leq \Sigma \rangle by blast
               thus O z x using M.T M-axioms O-def by blast
           qed
           hence w \leq x using SS D-def by blast
           have \forall v. v \leq w \longrightarrow \neg v \leq y by (metis O-def M.T M-axioms \Sigma \land w \leq
\Sigma)
               hence \neg O w y using O-def by simp
               thus w \leq x \land \neg O w y  using \langle w \leq x \rangle  by blast
               assume w \leq x \land \neg O w y
               show w \leq \overline{\Sigma} by (meson D-def O-def R SS \Sigma \langle w \leq x \land \neg O w y \rangle)
             qed
            thus (\exists z. \forall w. w \leq z \longleftrightarrow (w \leq x \land \neg O w y))..
         qed
```

8 Atomism

An atom is an individual with no proper parts:

```
definition (in M) A:: i \Rightarrow bool(A)
where A x \equiv \neg (\exists y. y < x)
```

Each theory discussed above can be augmented with an axiom stating that everything has an atom as a part, viz. [1], p. 48:

```
locale AM = M + assumes A: \forall x. \exists y. A y \land y \leq x — atomicity locale AMM = AM + MM locale AEM = AM + EM locale ACM = AM + CM locale ACEM = AM + CEM locale AGM = AM + GM locale AGEM = AM + GEM locale AGEM = AM + GEM
```

It follows in AEM that if something is not part of another, there is an atom which is part of the former but not part of the later:

```
lemma (in AEM)

ASS: \neg x \leq y \longrightarrow (\exists z. A z \land (z \leq x \land \neg O z y))
```

```
proof assume \neg x \leq y hence (\exists w. w \leq x \land D w y) using SS by simp then obtain w where w: w \leq x \land D w y. hence \exists z. A z \land z \leq w using A by simp then obtain z where z: A z \land z \leq w.. hence A z \land (z \leq x \land \neg O z y) by (meson \ D\text{-}def \ O\text{-}def \ T \ w) thus \exists z. A z \land (z \leq x \land \neg O z y).. qed
```

Moreover, in Minimal Mereology this lemma entails both strong supplementation and atomism, so it serves as an alternative characterisation of Atomistic Extensional Mereology:

```
lemma (in MM)
  assumes ASS: \neg x \leq y \longrightarrow (\exists z. A z \land (z \leq x \land \neg O z y))
  shows SS: \neg x \leq y \longrightarrow (\exists z. z \leq x \land D z y) using D-def assms by blast
lemma (in M)
  assumes ASS: \forall x. \forall y. \neg x \leq y \longrightarrow (\exists z. A z \land (z \leq x \land \neg O z y))
  shows A: \forall x. \exists y. A y \land y \leq x
proof
  \mathbf{fix} \ x
  show \exists y. A y \land y \leq x
  proof cases
    assume A x
    hence A \ x \land x \le x  using R  by simp
    thus \exists y. A y \land y \leq x..
    assume \neg A x
    hence \exists y. y < x \text{ using } A\text{-}def \text{ by } simp
    then obtain y where y: y < x..
    hence \neg x \leq y \text{ using } PP\text{-}def \text{ by } simp
    hence \exists z. A z \land (z \leq x \land \neg O z y) using ASS by blast
    thus \exists y. A y \land y \leq x \text{ by } blast
  qed
qed
```

So the axiom of Atomistic Strong Supplementation could be used in place of the two axioms of Atomicity and Strong Supplementation [1]

For the same reason that the Product axiom and Strong Supplementation do not follow from the Fusion Axiom in General Mereology, and so General Mereology is strictly weaker than Classical Extensional Mereology, the Product Axiom and Strong Supplementation still do not follow from the Fusion Axiom in Atomistic General Mereology, and so Atomistic General Mereology is also strictly weaker than Atomistic Classical Extensional Mereology [2], p. 206:

```
lemma (in AGM) T: O x y \longrightarrow (\exists z. \forall w. w \leq z \longleftrightarrow (w \leq x \land w \leq y)) nitpick
```

```
[user-axioms] oops lemma (in AGM) SS: \neg x \leq y \longrightarrow (\exists z. z \leq x \land D z y) nitpick [user-axioms] oops
```

Alternatively, each theory discussed above can be augmented with an axiom stating that there are no atoms, viz.:

```
locale XAM = M +
assumes XA: \neg (\exists x. A x) — atomlessness
locale XAMM = XAM + MM
locale XAEM = XAM + EM
locale XACM = XAM + CM
locale XACEM = XAM + CEM
locale XACEM = XAM + GEM
locale XAGEM = XAM + GM
locale XAGEM = XAM + GEM
```

Pontow notes that the question of whether the Fusion Axiom entails the Product and Strong Supplementation axioms in Atomless General Mereology is open [2]. Nitpick does not find a countermodel (since an infinite countermodel is needed?) and sledgehammer fails to find a proof, so this problem remains open for now:

```
lemma (in XAGM) T \colon O \ x \ y \longrightarrow (\exists \ z. \ \forall \ w. \ w \le z \longleftrightarrow (w \le x \land w \le y)) oops lemma (in XAGM) SS \colon \neg \ x \le y \longrightarrow (\exists \ z. \ z \le x \land D \ z \ y) oops
```

9 Consistency

I conclude by proving the consistency of all the theories mentioned.

```
lemma (in M) False nitpick [show-all] oops lemma (in MM) False nitpick [show-all] oops lemma (in EM) False nitpick [show-all] oops
```

end

References

- [1] R. Casati and A. C. Varzi. *Parts and Places: The Structures of Spatial Representation*. MIT Press, Cambridge, Mass., 1999.
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- [3] P. Simons. Parts: A Study in Ontology. Clarendon Press, Oxford, 1987.
- [4] A. Varzi. Mereology. In E. N. Zalta, editor, *The Stanford Encyclopedia of Philosophy*. Metaphysics Research Lab, Stanford University, winter 2016 edition, 2016.