

COMPILER DESIGN

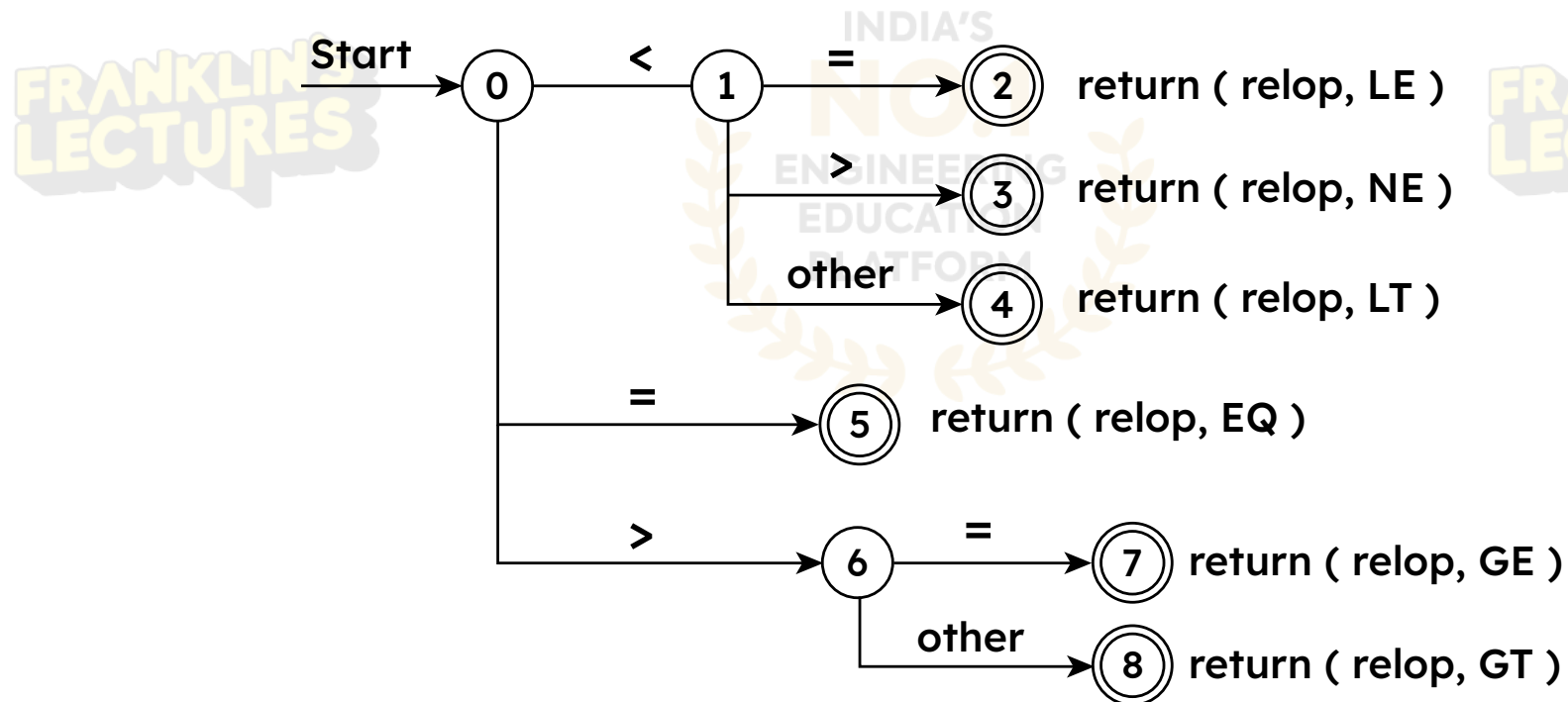
Module 1

Part 2

CST302

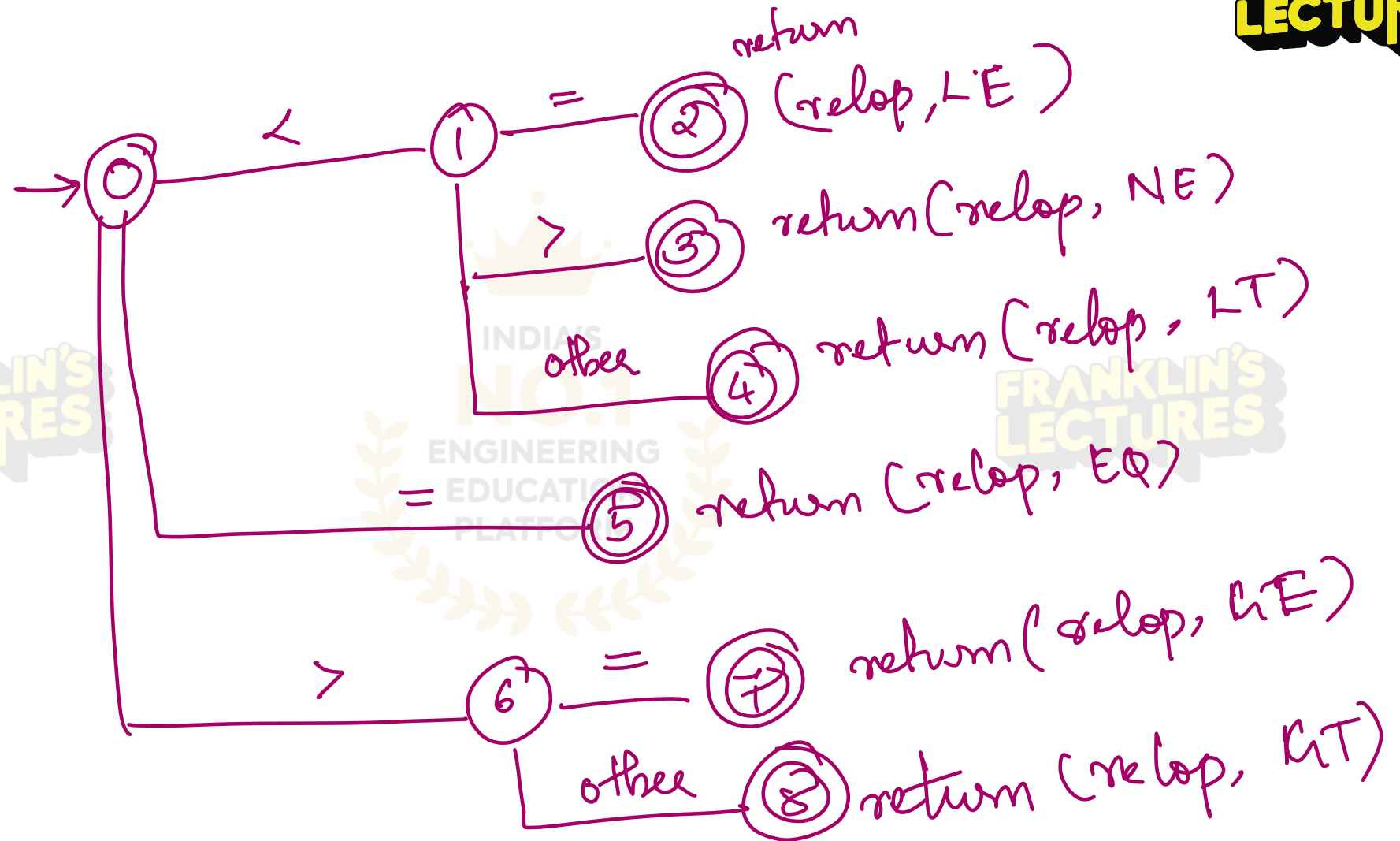
Q1. Draw the transition diagram for the regular definition

relop \longrightarrow <|<=|=|<>|>=|>



$< | < = | < >$ $> | > =$

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Q2. With an example source language explain tokens, lexemes & pattern,

Token: Are the terminal symbol in the grammar for the source language. which are generally represented using bold phase names.

Eg:- *id*, *const*

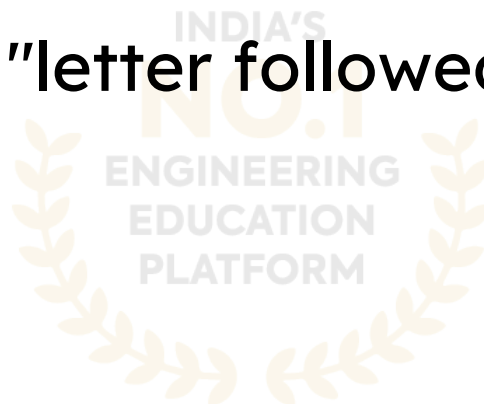
lexemes: A lexeme is a sequence of characters in the source pgm that is matched by a pattern for a token

eg: consider the statement, *const* Pi = 3.14

Here P_i is a lexeme for the token identifier.

Pattern: A pattern is a rule describing a set of lexemes that can represent a particular token in the source program

eg:- pattern for identifier is "letter followed by letters and digits"



Q3. Describe Input buffering scheme in lexical analyzer.

The process of Storing a block of input data in buffer to avoid costly access to secondary storage each time.

Commonly used buffering methods are,

1. One buffer scheme :- Here only one buffer is used to store lexemes. If the the lexeme crosses buffer boundary, the buffer has to be re filled, there by over-writing the buffer.
2. Two buffer scheme: Here two buffers are used to store the input. Here buffers 1 and 2 are scanned alternatively. When the end of the current buffer is reached, the other buffer is filled. Hence the problems that occur in one buffer Scheme is solved.

overwriting.

Begin $i = i+1 ; j := j + 1$

Fig : one buffer scheme

Var $i, j : \text{integer} ; \dots j := j+1$ e

eof

Buffer 1

nd

Fig : Two buffer scheme

Buffer 2

}

Q4. Construct a regular expression to denote a language L over $\Sigma = \{0,1\}$ accepting all strings of 0's and 1's that do not contain Substring 011.

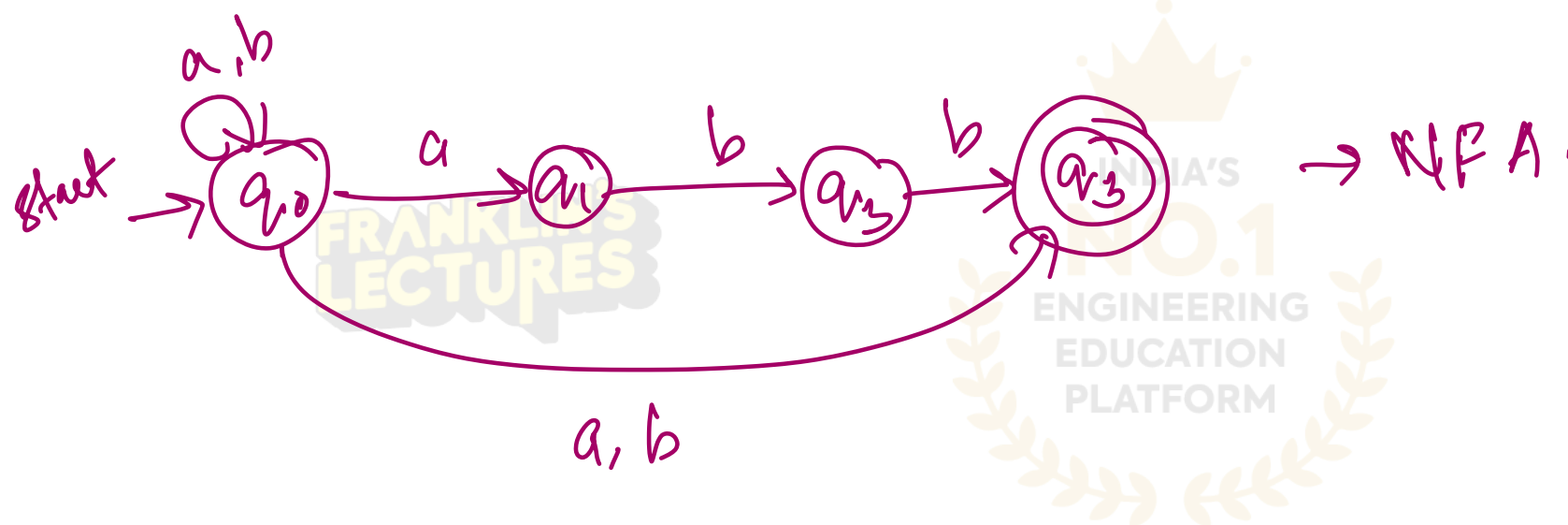
$$\Sigma = \{0,1\}$$

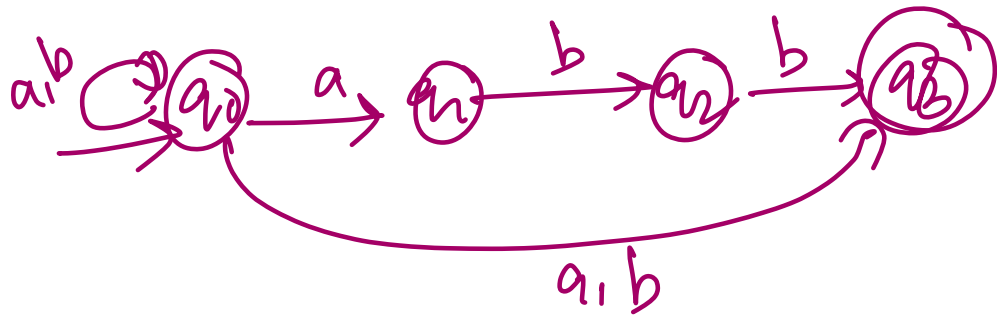
Ans: $1^* + 1^*0(10+0)^*(\epsilon+1)$

Solution:- construct DFA, then apply Aden's theorem to get regular expression.

$$\underline{110011} = \frac{1^* + 1^*0(10+0)^*(\epsilon+1)}{110X}$$

Q5. Draw the DFA for the regular expression $(a/b)^* \underline{abb} / a + b$

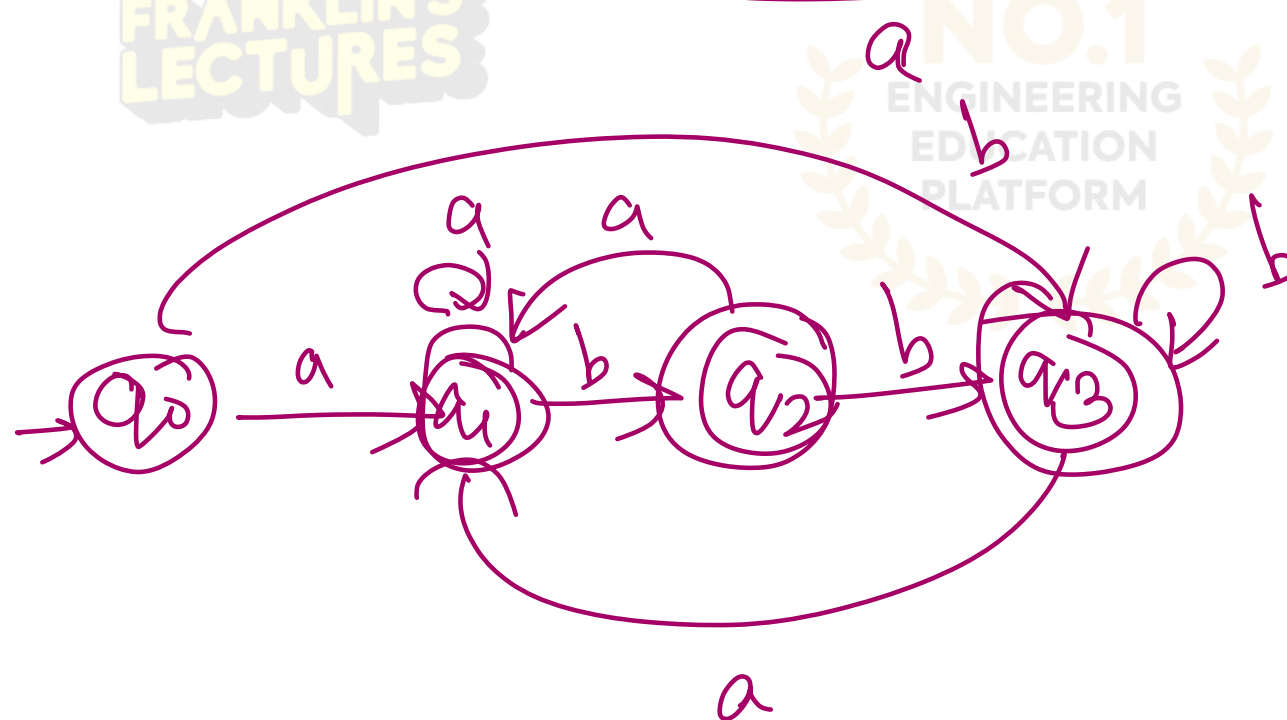
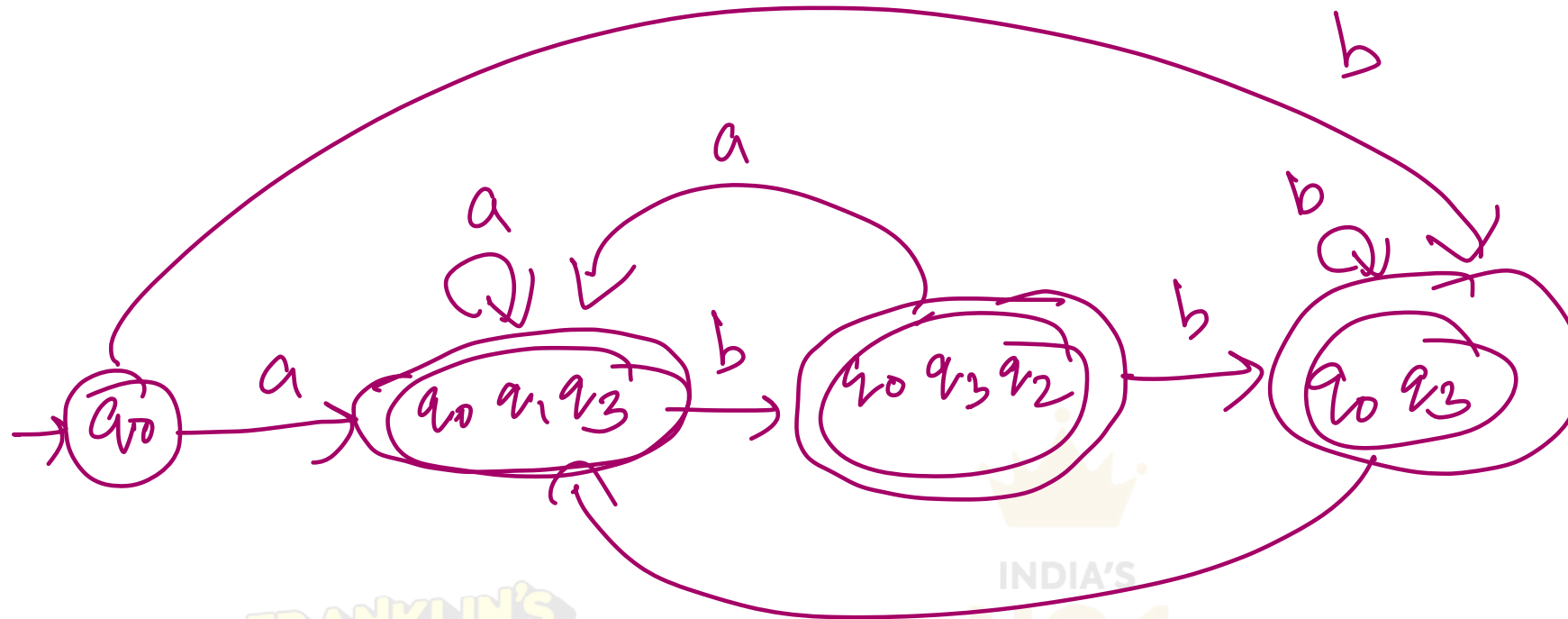




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	a	b
* $[q_0]$	$q_0 q_1 q_3$	$q_0 q_3$
$[q_0 q_1 q_3]$	$q_0 q_1 q_3$	$q_0 q_3 q_2$
$[q_0 q_3]$	$q_0 q_1 q_3$	$q_0 q_3$
$[q_0 q_3 q_2]$	$q_0 q_1 q_3$	$q_0 q_3$

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Q6. State the role of lexical analyzer. Identify the lexemes and their corresponding tokens in the following statement

`Printf("Simple Interest = %f\n", Si);`

Role of lexical analyzer:-

- * Scans the input from left to right and group the characters into tokens.

Eliminates white spaces and comments

correlates error messages with the source program

lexeme: Printf, token: ID,

lexeme: C, token: left parenthesis

lexeme: "Simple Interest = %f\n

token: LITERAL

lexeme: , token: comma.

lexeme: , token: right parenthesis

lexeme: Si, token: ID

lexeme:), token: right parenthesis

lexeme: ;, token: semicolon.

Q7. Explain any three tools that help a Programmer in building a compiler efficiently.

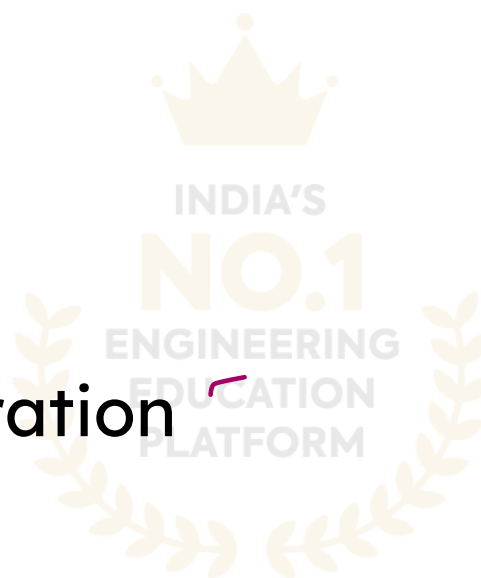
→ compiler construction tools.

Explain about any three compiler construction tools. (Refer QnNo.9)

Q8. Explain the different phases in design of compiler.

- 1) Lexical analysis
- 2) Syntactic analysis
- 3) Semantic analysis.
- 4) Intermediate code generation
- 5) code optimization
- 6) code generation.

Explain each phase with an example.



Q9. Explain compiler writing tools

construction-



Compiler writing tools are the tools that have been created for automatic design of specific compiler components. These tools use specialized languages and algorithms. The following is a list of compiler construction tools.

- 1) **Scanner generator** :- They generate lexical analyzers automatically from the language specification written using regular expressions. It generate finite automata to recognize the regular expression. An example of this tool is **LEX**.

- 2) Parser generator: - They produce syntax analyzers from context free grammars. Many parser generator utilize powerful parsing algorithm that are too complex to be carried out by hand. An example for this tool is YACC (yet another compiler compiler)
- 3) Syntax directed translation Engine:-
These engine have routines to traverse the parse tree & to produce intermediate code. one or more translations are associated with each node of the parse tree.

- 4) Automatic code generators: - These tools convert the intermediate language into machine language using a collection of rules. Here template matching process is used. An intermediate language statement is replaced by its equivalent machine language statement.
- 5) Data flow Engines:- It is used in code Optimization. These tools performs good code optimization using "data-flow analysis" means the gathering of information about how values are transmitted from one part of a program to each other part.

Q10. Explain how the regular expressions and finite automata are used for the Specification and recognition of a tokens.

Scanners are special pattern matching processors. For representing Pattern of strings of characters, regular expressions are used. For example the pattern for the token identifier is, letter-followed by letters and digits. This can be represented by using the regular expression given below.

letter (letter / digit)*

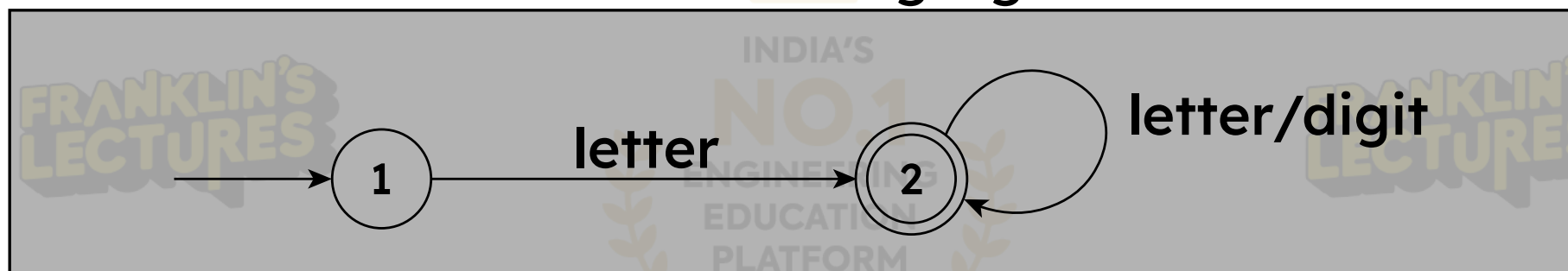
- The tokens obtained during lexical analysis are recognized using a finite automata. Finite automata can be used to describe the process of recognizing patterns in input strings and so they can be used to construct Scanners. A regular definition for an identifier is given below,

letter \rightarrow a|b|c|.....z|A|B|....|Z

digit \rightarrow 0|1|2|.....9

identifier \rightarrow letter(letter/digit)*

- There an identifier consists of a letter followed by any number of letters or digits. The finite automata for the identities can thus be represented as shown in the following figure



finite automata for recognizing an identifier

Q11. Develop a lexical analyzer for the token identifier.

Lexical analyzer convert a sequence of characters into a sequence of tokens. A program that performs lexical analysis may be termed as lexical analyzer/Scanner. The regular definition for an identifier is given below,

letter \rightarrow a|b|c|.....z|A|B|....|Z

digit \rightarrow 0|1|2|.....9

identifier \rightarrow letter(letter/digit)*

Regular expression for identifier token is Letter (letter/digit)*
lexical analyzer uses finite automata to recognize the regular expression. The finite automata used by lexical analyzer to identify the identifier token is given below.

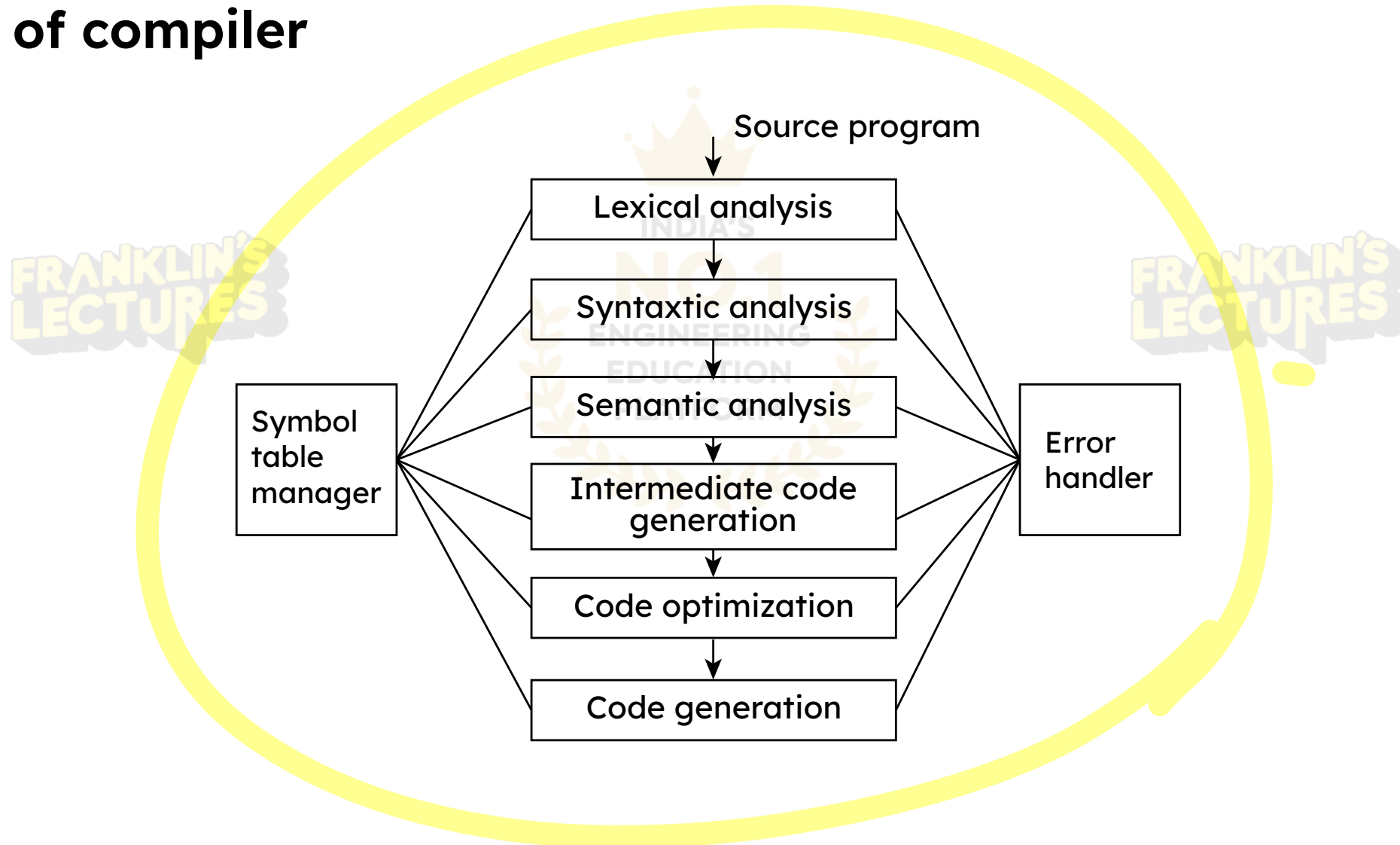


Q12. Trace the output after each phase of the compiler for the assignment statement $a = b + c * 10$, if variables given are of float type.

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a, b, c — variable
float type.

Phases of compiler



$$a = b + c * 10$$

$$a = b + c * 10$$

↓

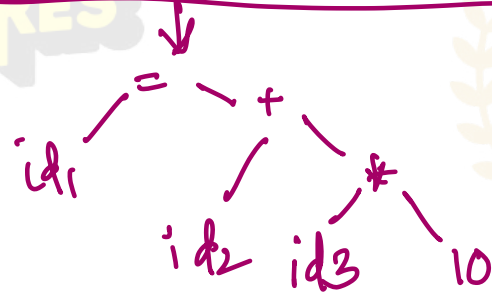
Lexical Analysis

↓

$$id_1 = id_2 + id_3 * 10$$

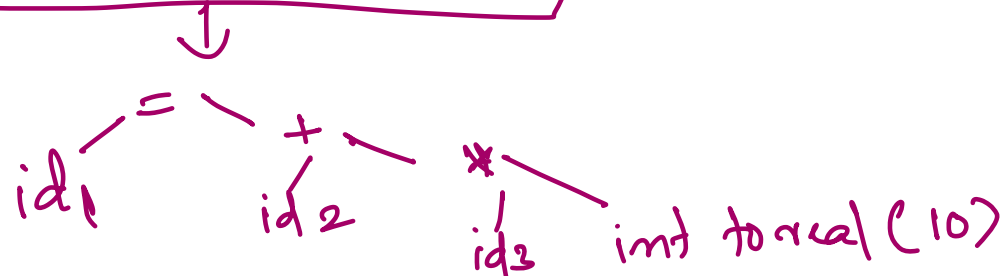
↓

Syntactic Analysis



↓

Semantic Analysis



↓
Intermediate code generation
↓

temp1 = int to real (10)

temp2 = id3 * temp1

temp3 = id2 + temp2

id1 = temp3

↓
Code optimization
↓

temp1 = id3 * 10.0

id1 = id2 + temp1

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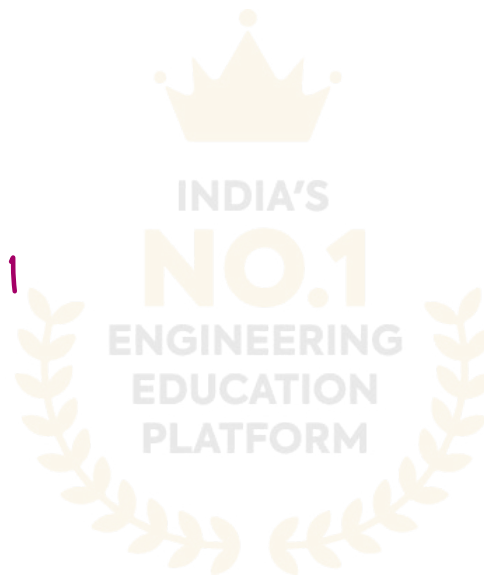
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↓
Code generation

↓
MOVF id3, R2
MULF #10.0, R2
MOVF id2, R1
ADDF R2, R1
MOVF R1, id1

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**Q13. For a source language statement,
 $a = b * c - 2$ where a, b, c are float
Variables, $*$ and $-$ represents
multiplications and subtraction on
same data type, show the input &
output for each of the compiler
Phases.**

$$a = b * c - 2 \quad \longrightarrow \quad \text{source program}$$


Lexical Analysis

२

$$id_1 = id_2 * id_3 - 2$$

2

1. Syntactic Analysis



Handwritten diagram showing a tree structure with nodes d_2 , id_3 , and 2 . The root node is d_2 , which has two children: id_3 and 2 . The node id_3 has a child 2 . The root node is labeled d_2 .

Semantic Analysis

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1. of L. A

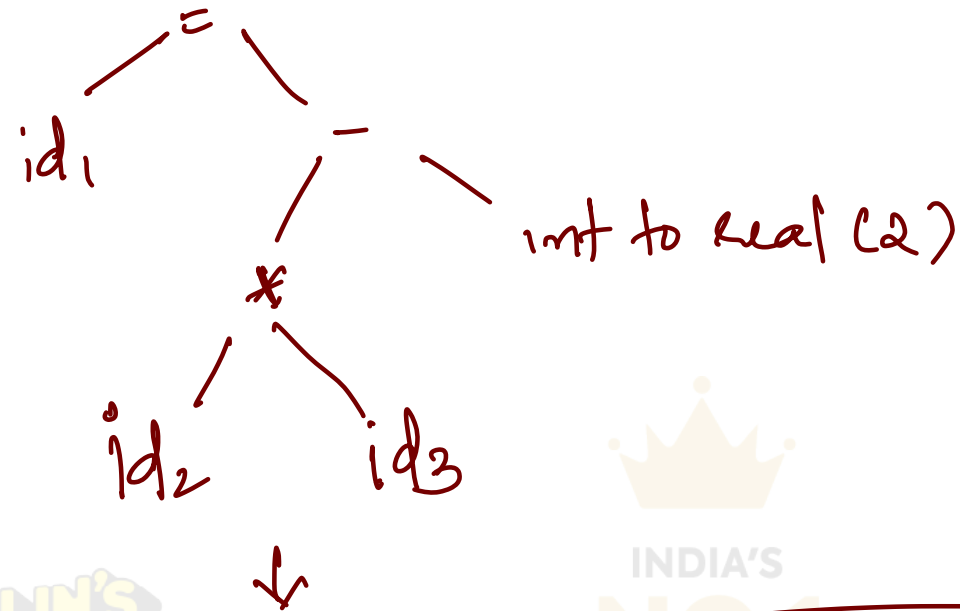
o/p syn. A.



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intermediate code generation

temp1 = int to real (2)
temp2 = id2 * id3
temp3 = temp2 - temp1
id1 = temp3

sem. A o/p

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o/p of ICC.

↓
code optimization

temp1 = id2 * id3
id1 = temp1 - 2.0

↓

code generation

MOVF id2, R2
MOVF id3, R1
MULF R1, R2
SUBF #2.0, R2
MOVF R2, id1

o/p of code opt.

target program / code

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