

# CS & IT ENGINEERING

Data Structure



Tree  
Chapter- 5  
Lec- 04



By- Pankaj Sharma sir

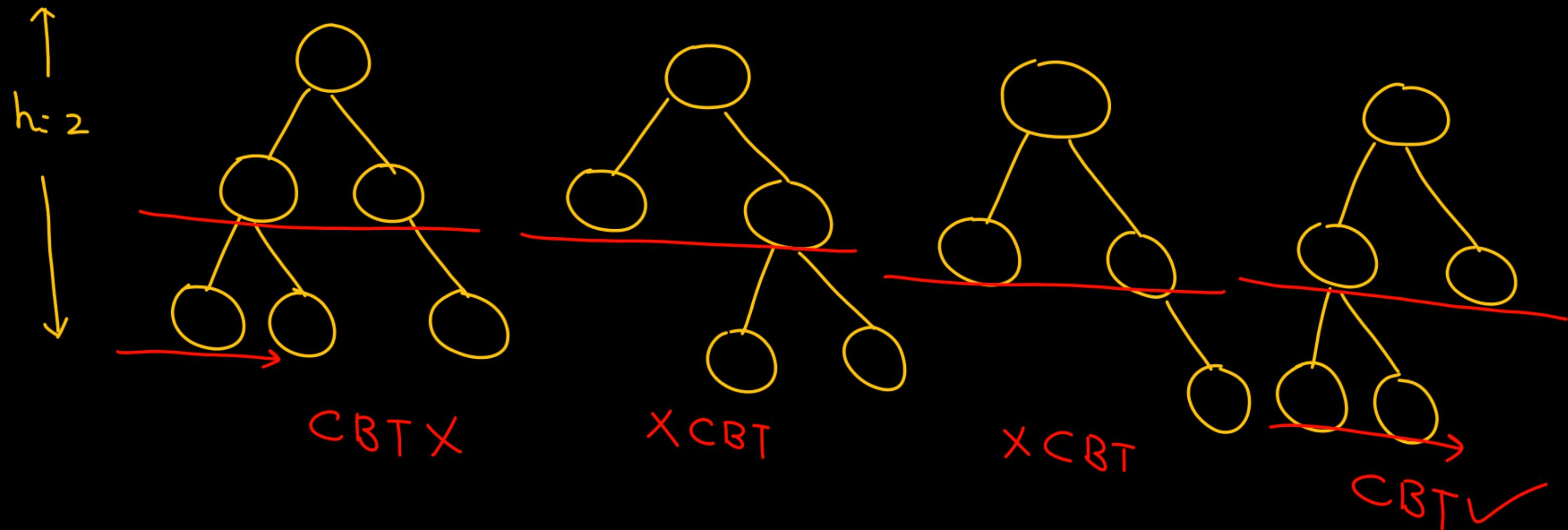
## TOPICS TO BE COVERED

Tree-IV

Complete binary tree

CBT is a binary tree which is Full till second last level.

& nodes at last level are filled from left to right.

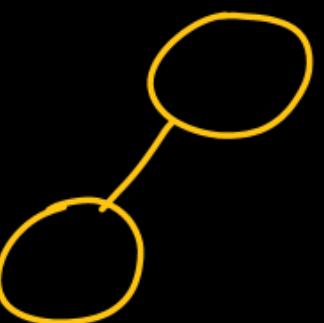


1. Structure of a CBT with 1 node

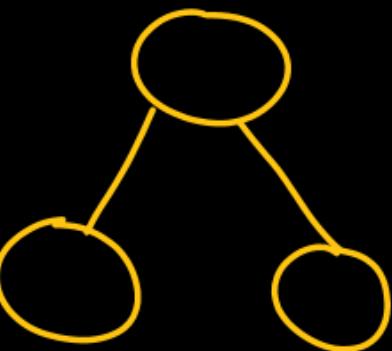


The structure of a

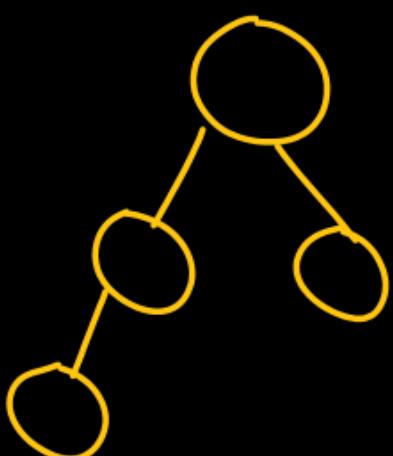
2. Structure of a CBT with 2 node



3. Structure of a CBT with 3 node



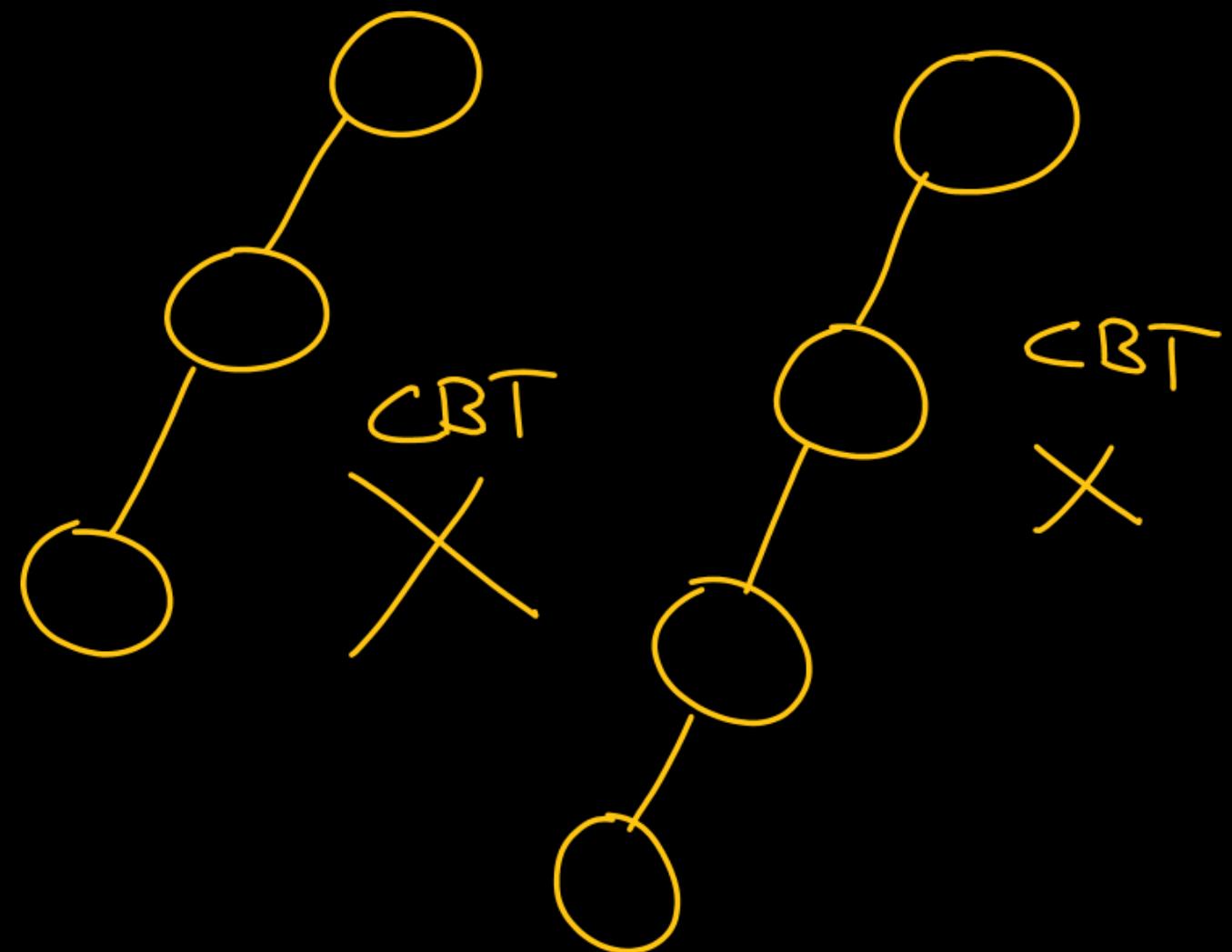
4. Structure of a CBT with 4 node

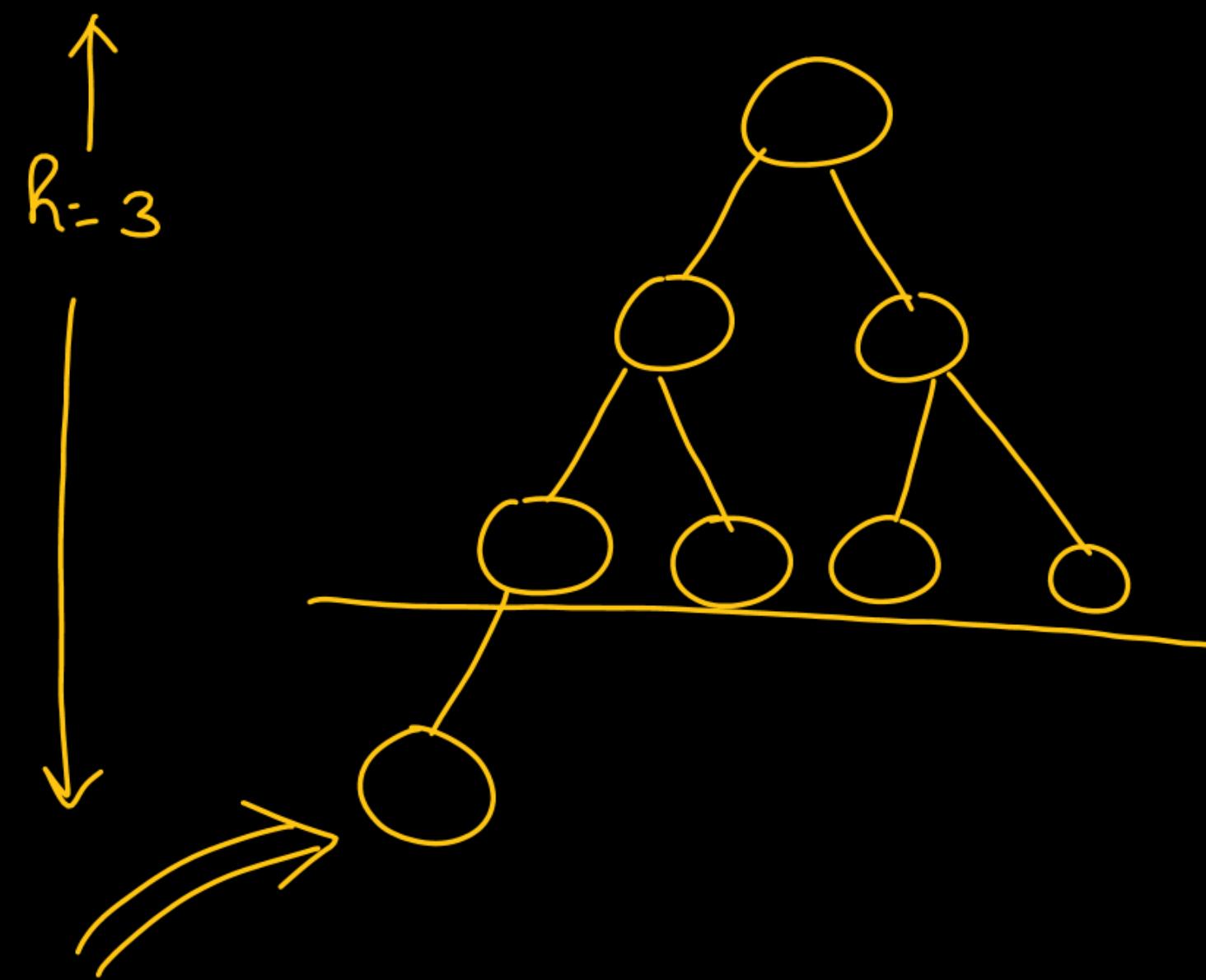


CBT with  $K$  nodes  
is always fix.

1) Max. no. of nodes in a CBT of height  $h$  =  $2^{h+1} - 1$

Min. no. of nodes in a CBT of height  $h$  = ?





$h = 3$

#Nodes

 $2^0$  $2^1$ 

.

 $2^2$ 

|level|

0

1

2

.

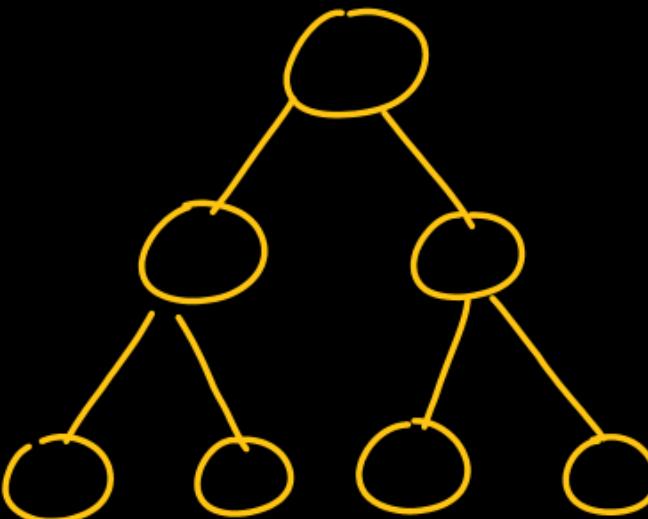
.

 $2^{h-1}$ 

1

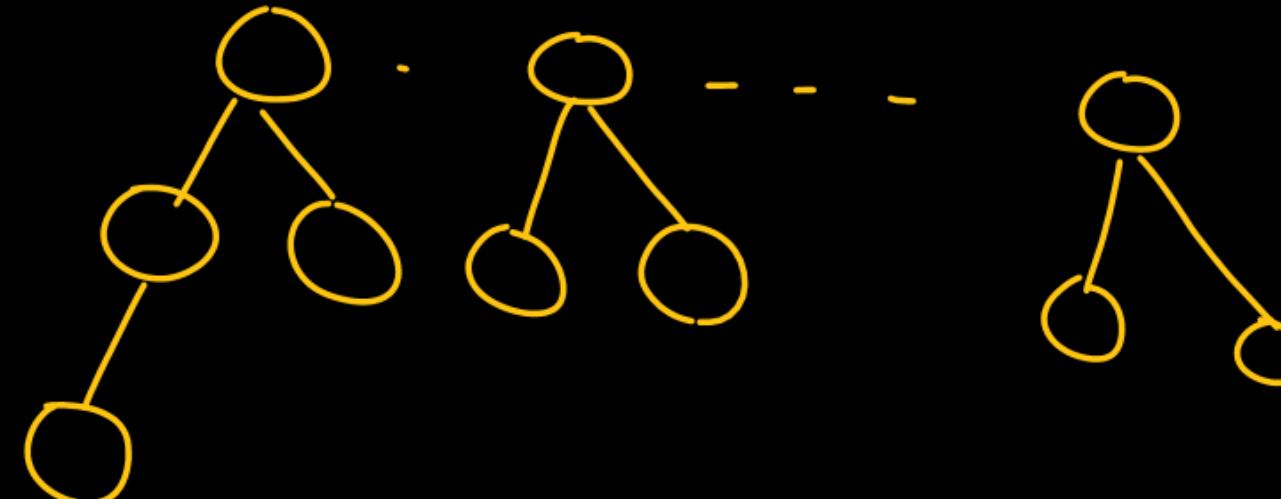
 $h-1$ 

h



$$S = \frac{a(r^n - 1)}{r - 1}$$

$a = 1$   
 $r = 2$   
 $n = h$



$$n_{\min} = (2^0 + 2^1 + 2^2 + \dots + 2^{h-1}) + 1 = \left( \frac{2^h - 1}{2 - 1} \right) + 1 = 2^h - 1 + 1 = 2^h$$

$$x \leq 4 \quad \text{max}$$

$$x \geq 2 \quad \text{min}$$

$$2 \leq x \leq 4$$

$$2^h \leq n \leq 2^{h+1} - 1$$

$$2^h \leq n$$

$$\log_2 2^h \leq \log_2 n$$

$$h \log_2 2 \leq \log_2 n$$

$$h \leq \frac{\log_2 n}{\log_2 2}$$

$$h \leq \log_2 n$$

$$n \leq 2^{h+1} - 1$$

$$(n+1) \leq 2^{h+1}$$

$$\log(n+1) \leq (h+1) \log_2 2$$

$$\frac{\log(n+1)}{\log 2} \leq h+1$$

$$(h+1) \geq \log_2(n+1)$$

$$h \geq \log_2(n+1) - 1$$

$$\log_2(n+1) - 1 \leq h \leq \log_2(n)$$

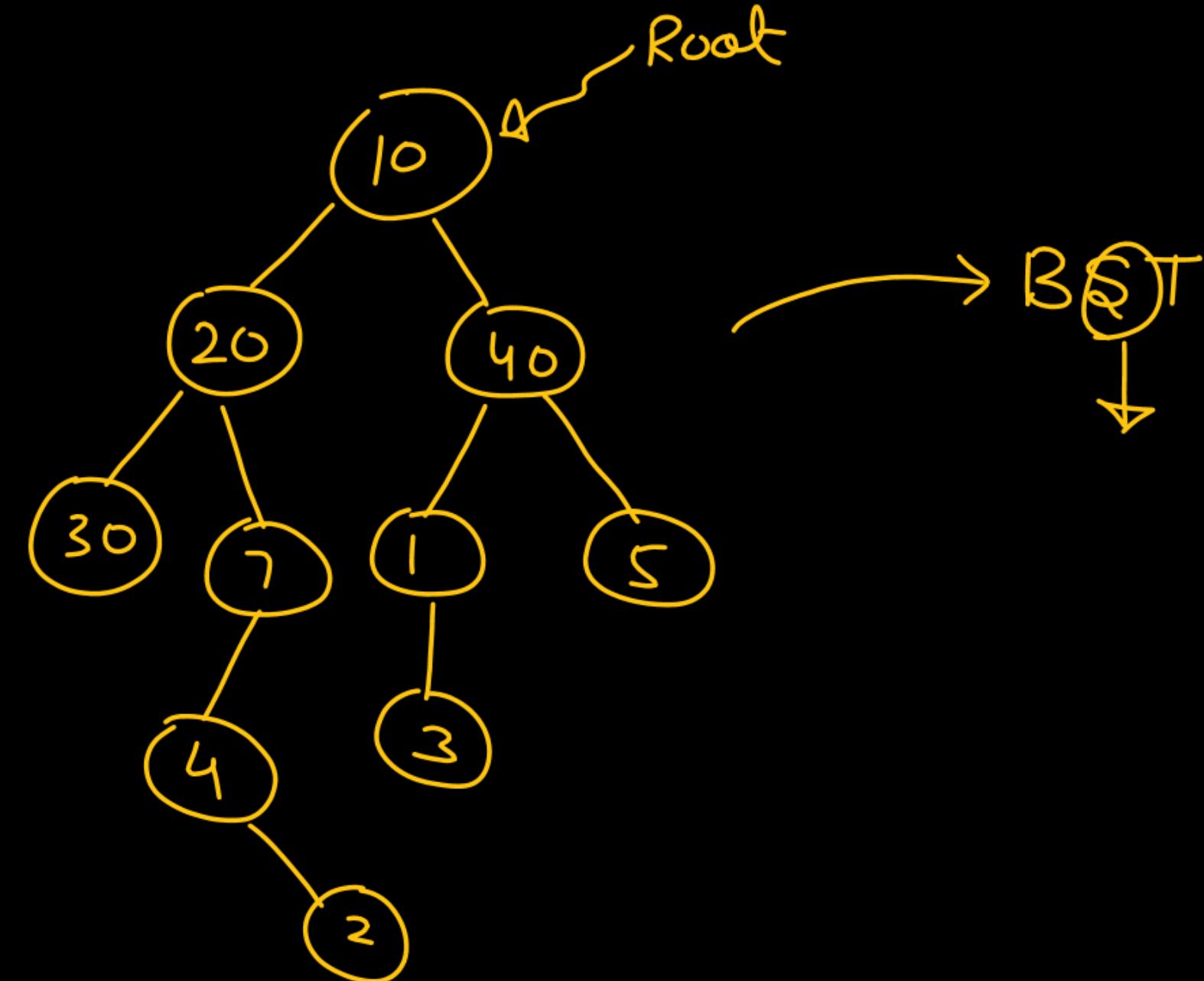
## Binary Search Tree

Why?

Given, Binary tree and a key and we need to find whether the key is present in the tree or not ?

key=100

$O(n)$

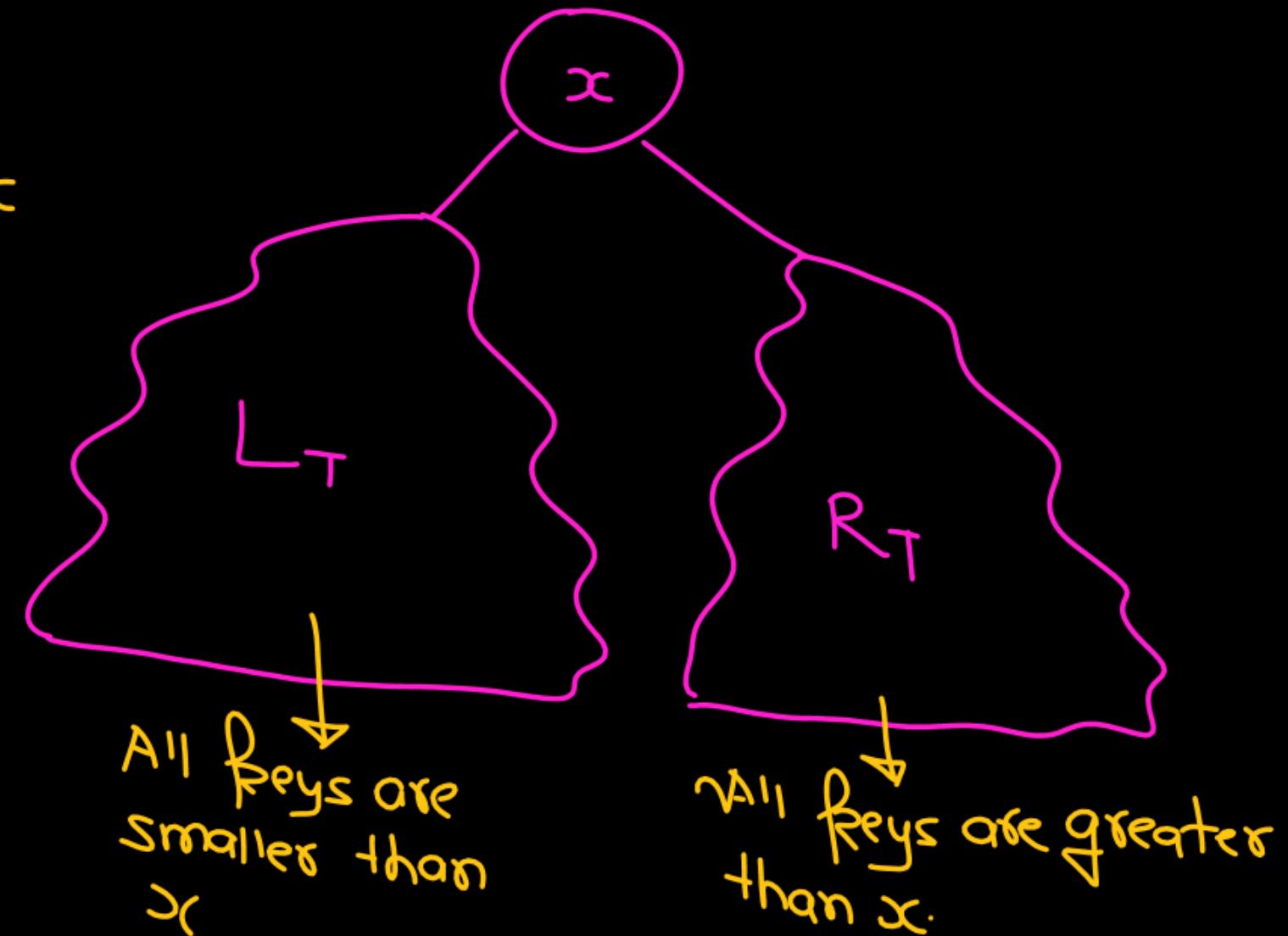


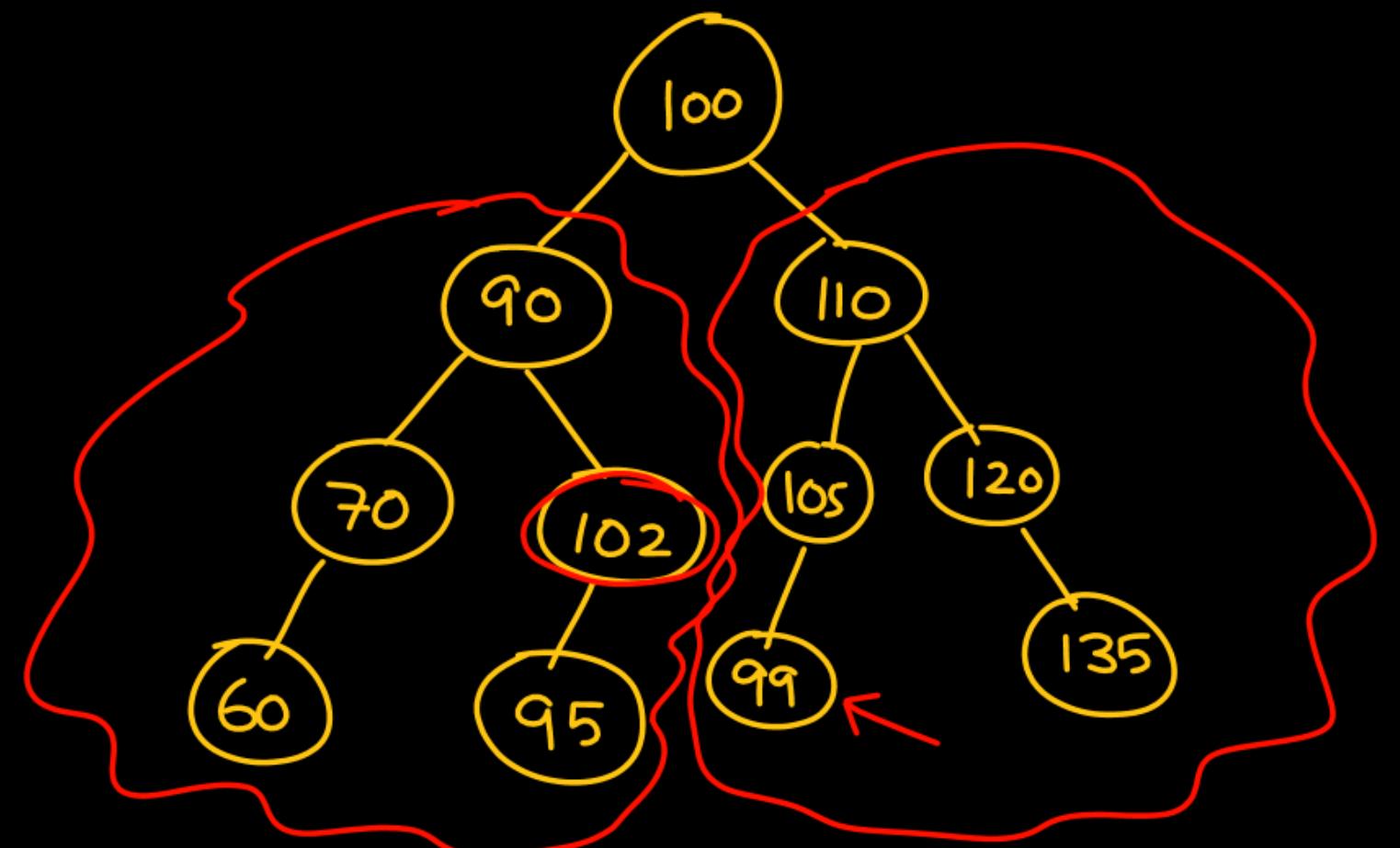
# Binary Search Tree

A binary tree in which **every node** satisfies following property:

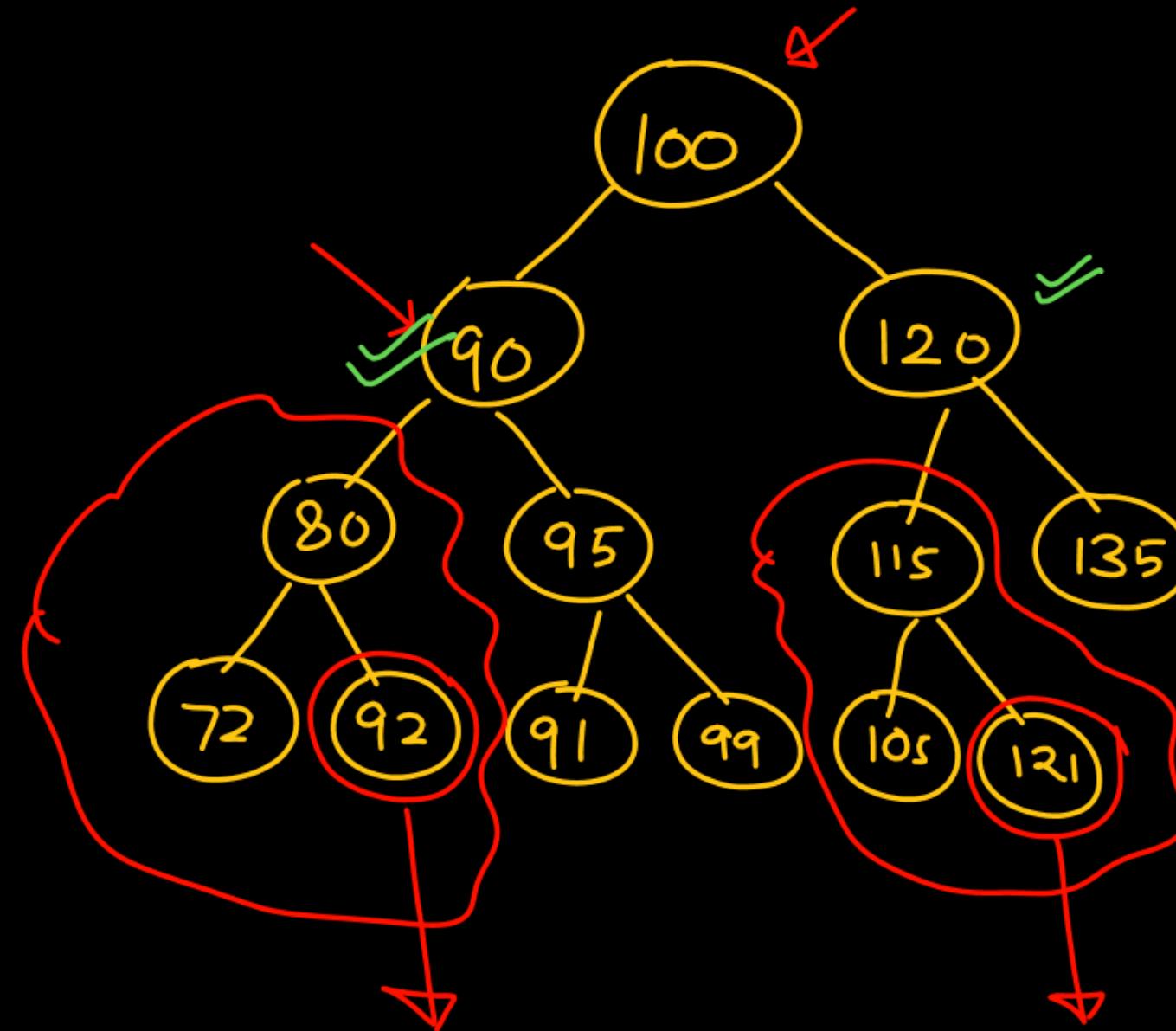
All the keys in the left subtree of a node are smaller than the node value.

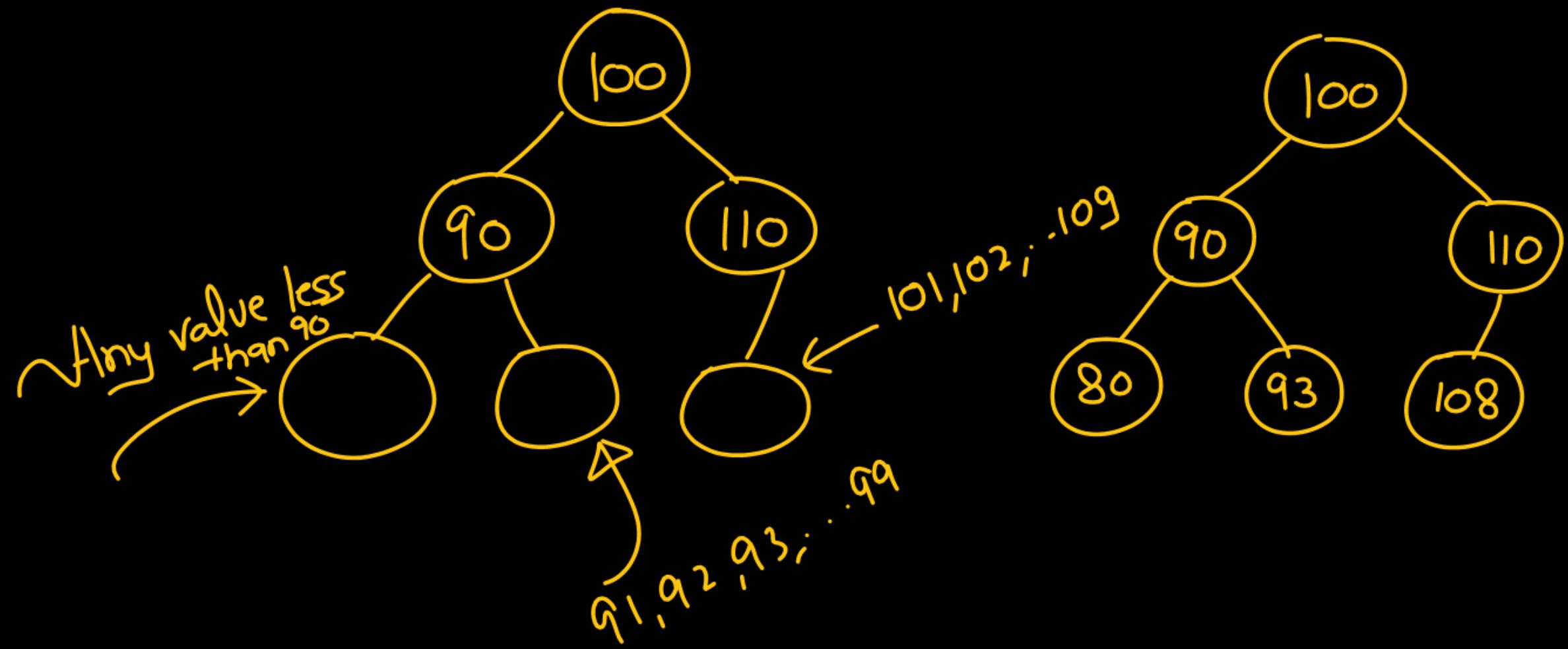
All the keys in the right subtree of a node are greater than the node value.





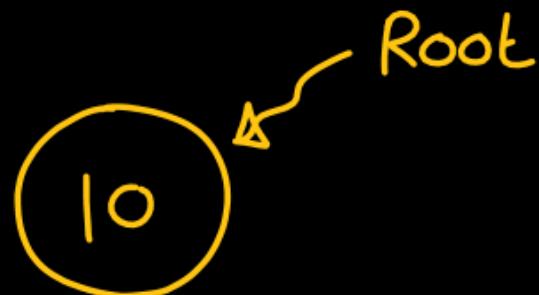
Is it a BST?



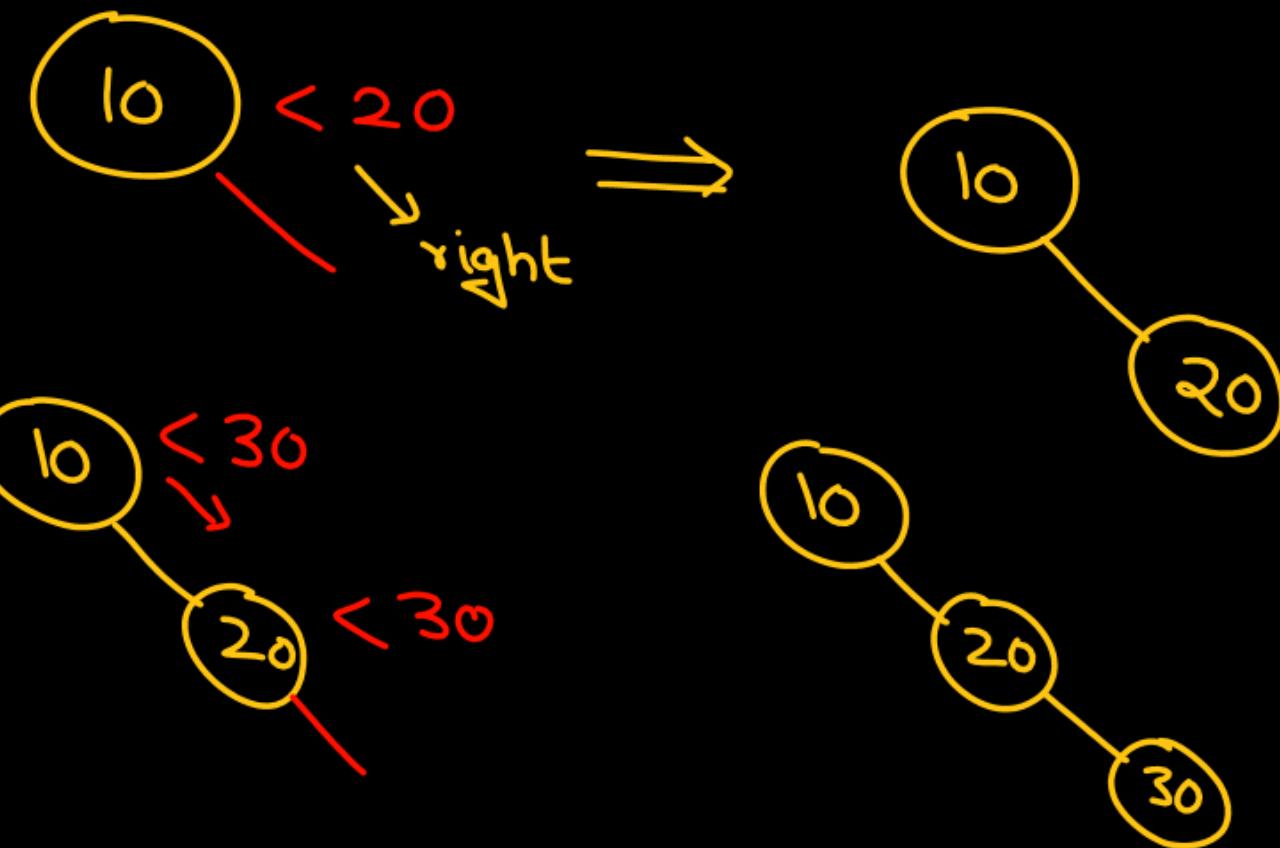


↳ Construct a binary search tree by inserting Keys 10, 20, 30 in order

a]

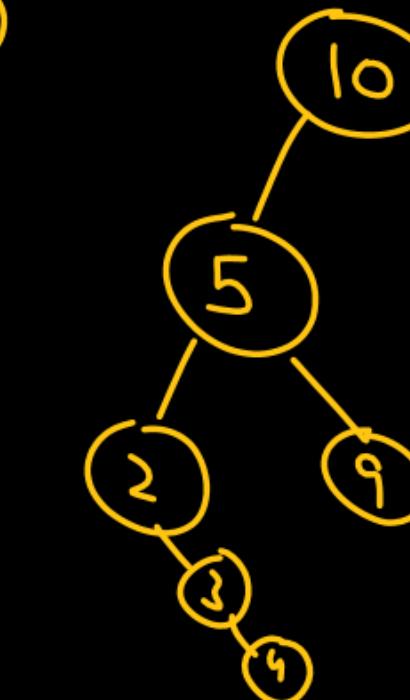
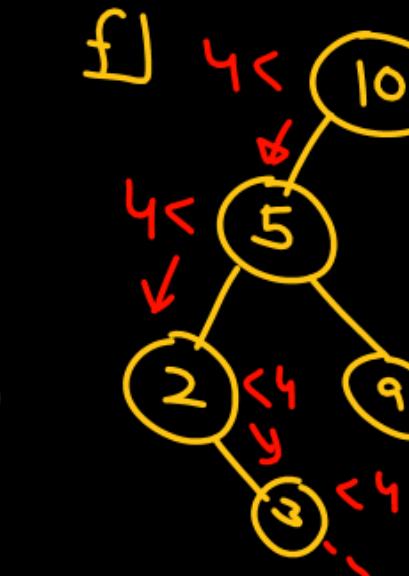
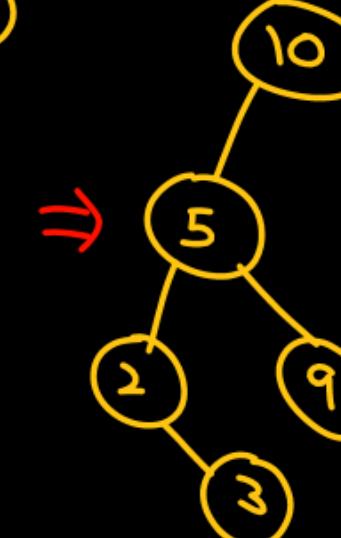
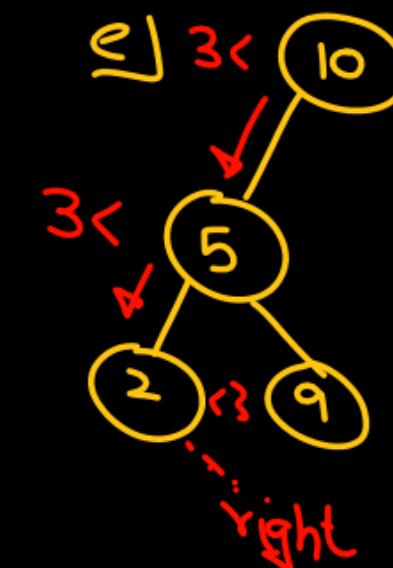
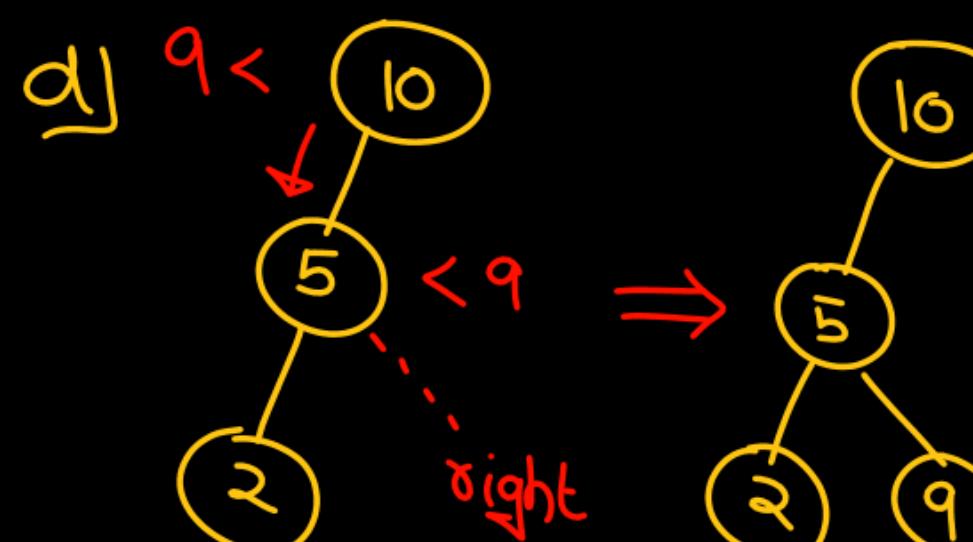
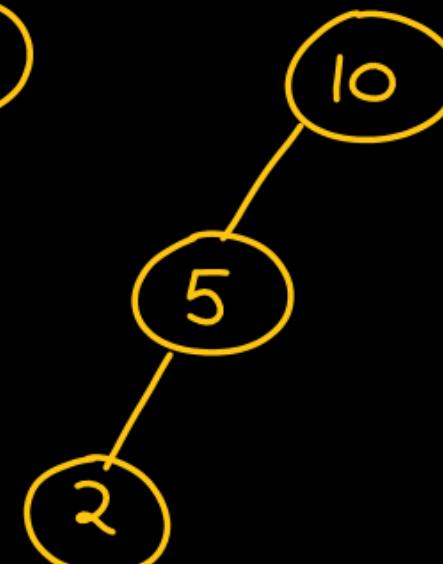
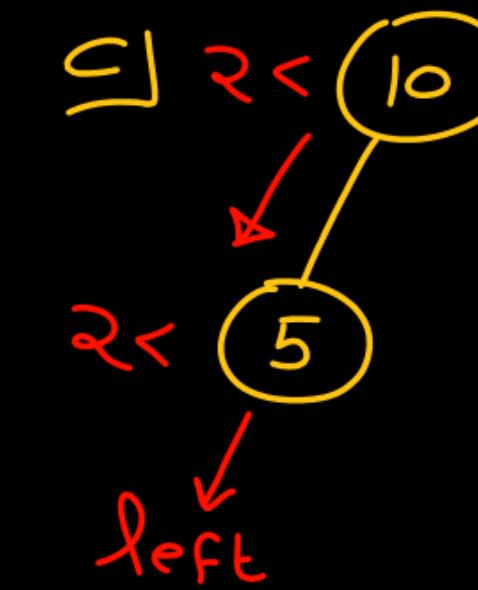
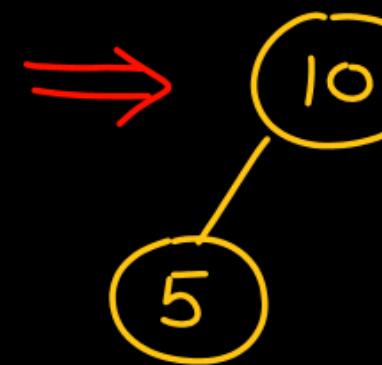


b]



c]

Q. Construct binary search tree by inserting keys  $\overrightarrow{10, 5, 2, 9, 3, 4}$  in order.



No. of BST when insertion order of keys are fixed

Const. BST by inserting keys 10, 20, 30  $\Rightarrow$  1

Const. BST by inserting keys 10, 5, 2, 9, 3, 4  $\Rightarrow$  1

Const. BST by inserting keys 1, 2, 3  
(in any order)

(i)  $\overrightarrow{1, 2, 3}$

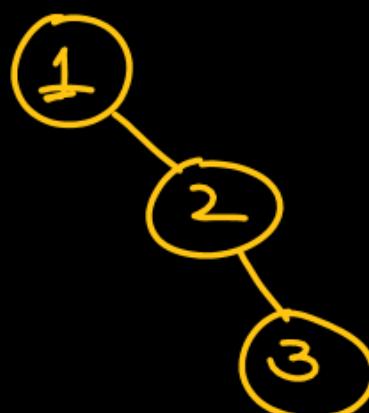
6 orders possible  
(ii)  $1, 3, 2$   
(iii)  $2, 1, 3$

(iv)  $2, 3, 1$

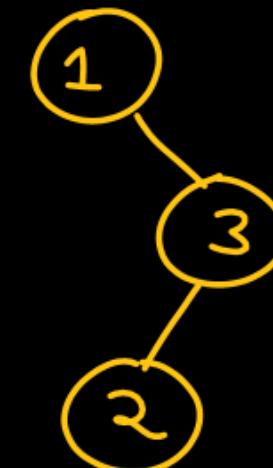
(v)  $3, 1, 2$

(vi)  $3, 2, 1$

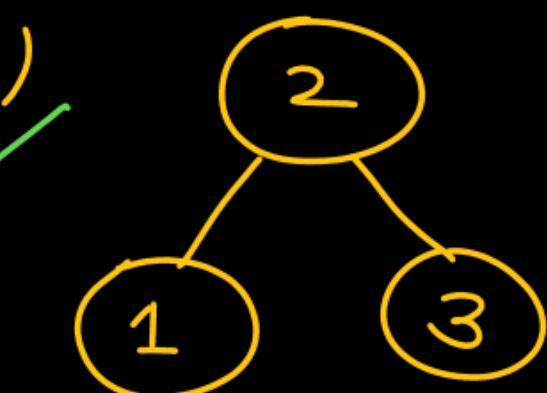
(i)



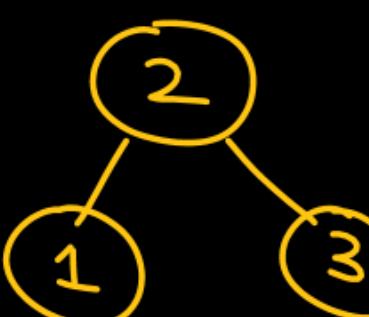
(ii)



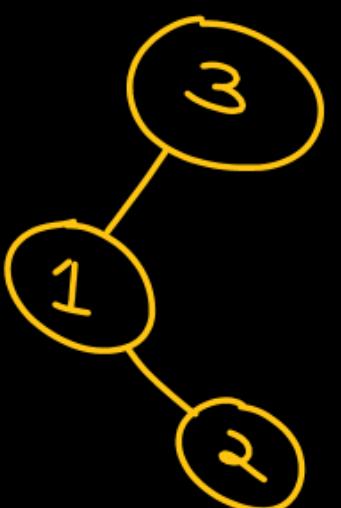
(iii)



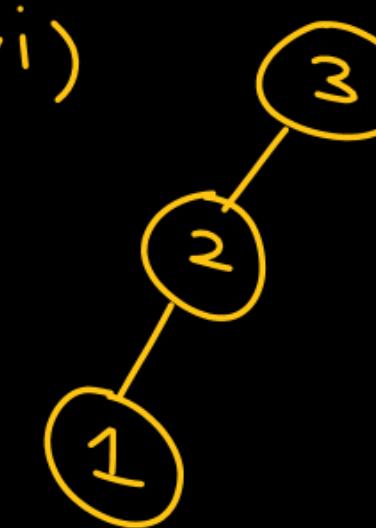
(iv)



(v)



(vi)

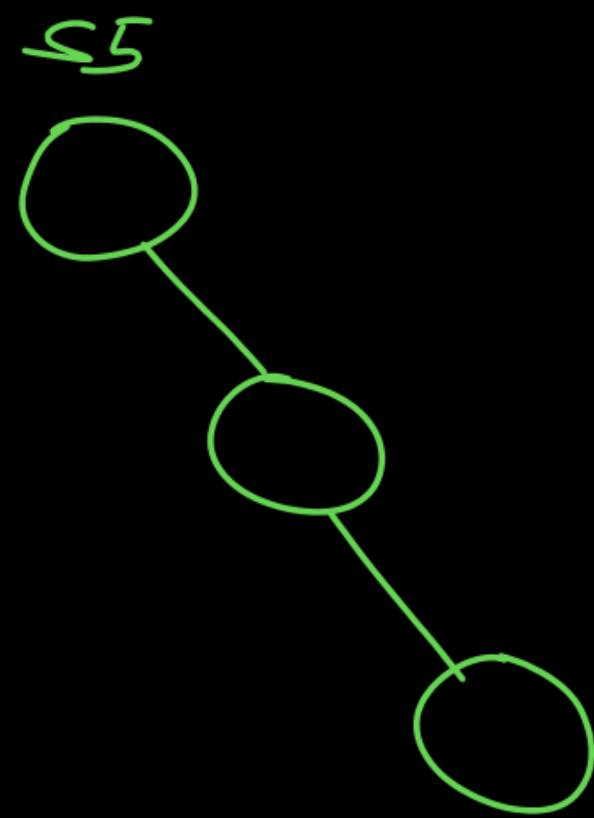
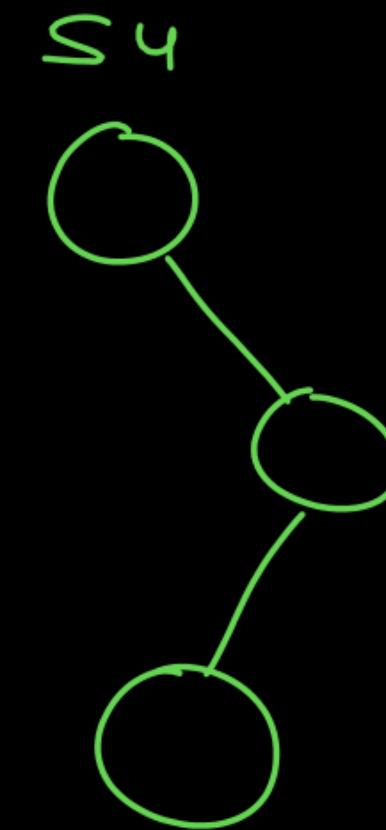
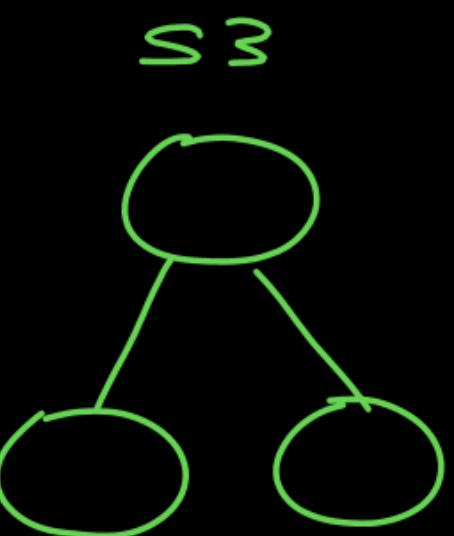
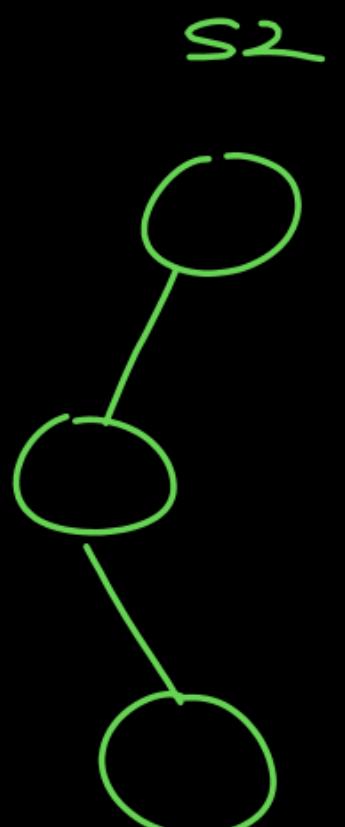
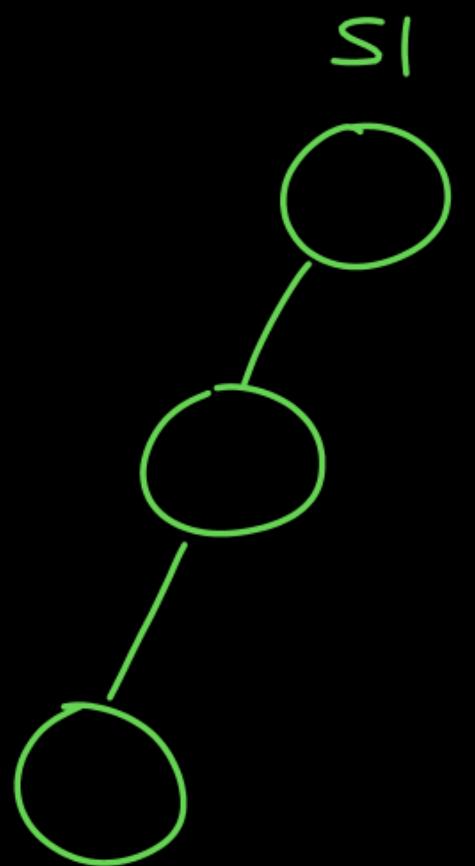


3 keys  $\Rightarrow$  5 BST

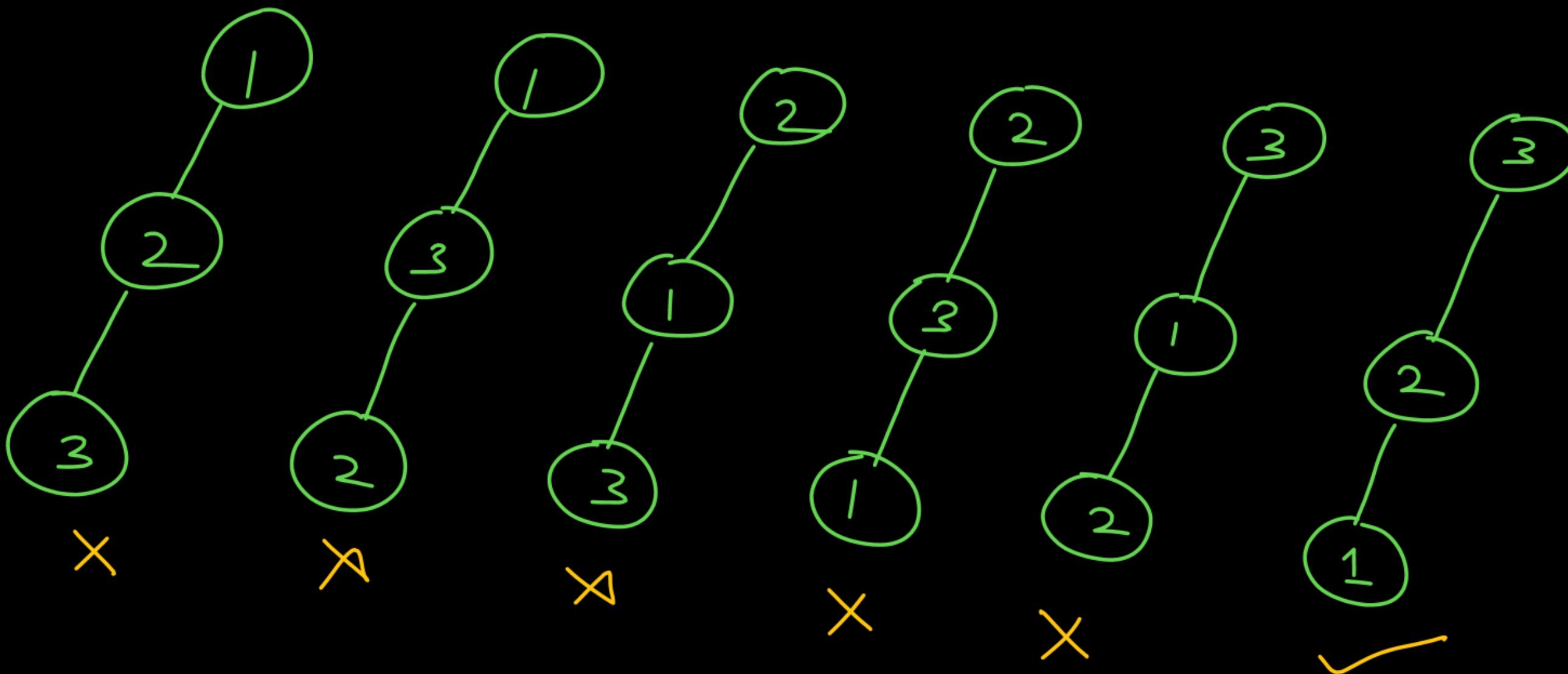
# unlabelled binary tree with 3 nodes = 5

# structure with n nodes =  $\frac{2^n C_n}{n+1}$

3 node



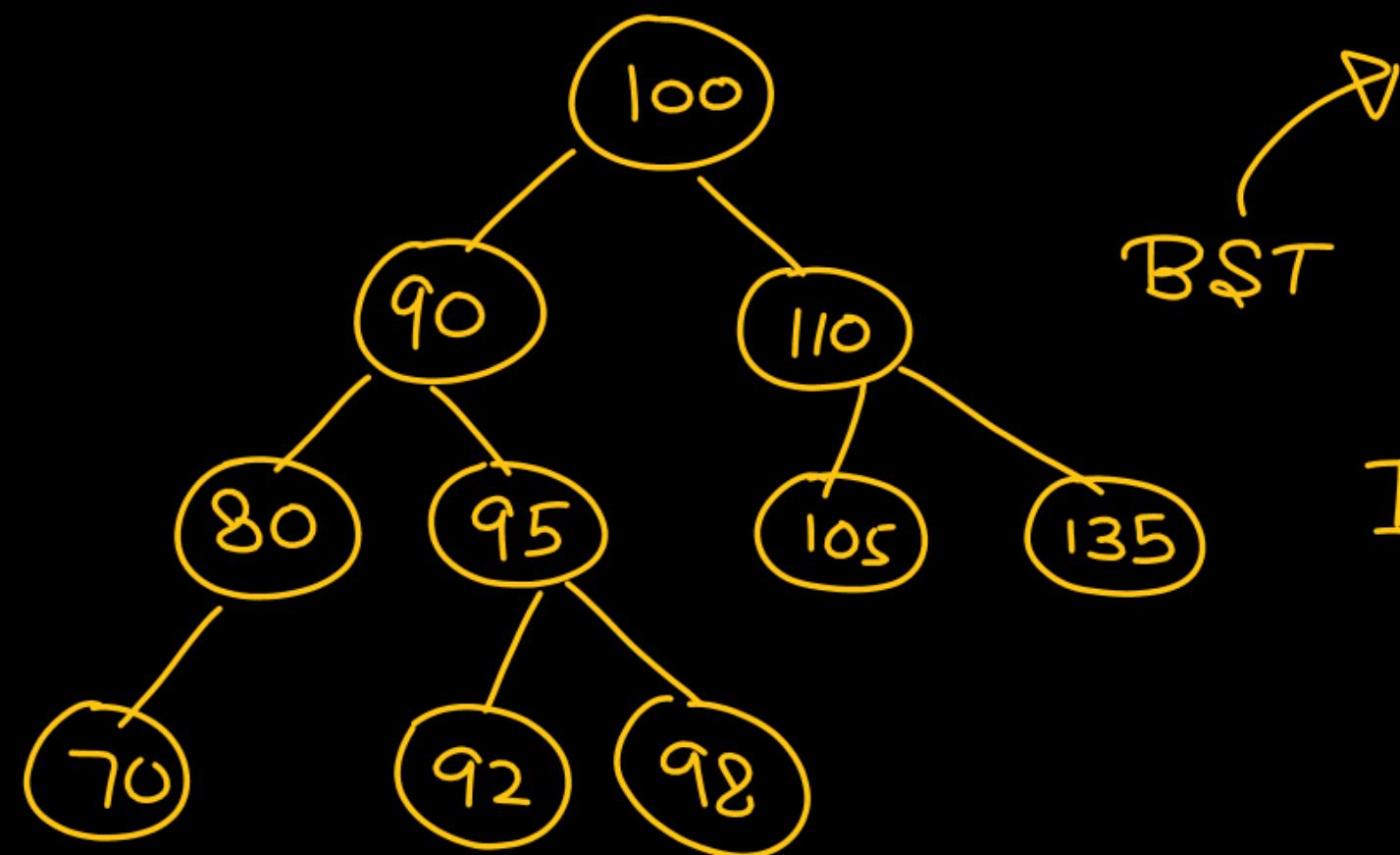
*S1*



$$\# \text{ BST with } n \text{ keys} = \frac{2^n C_n}{n+1}$$

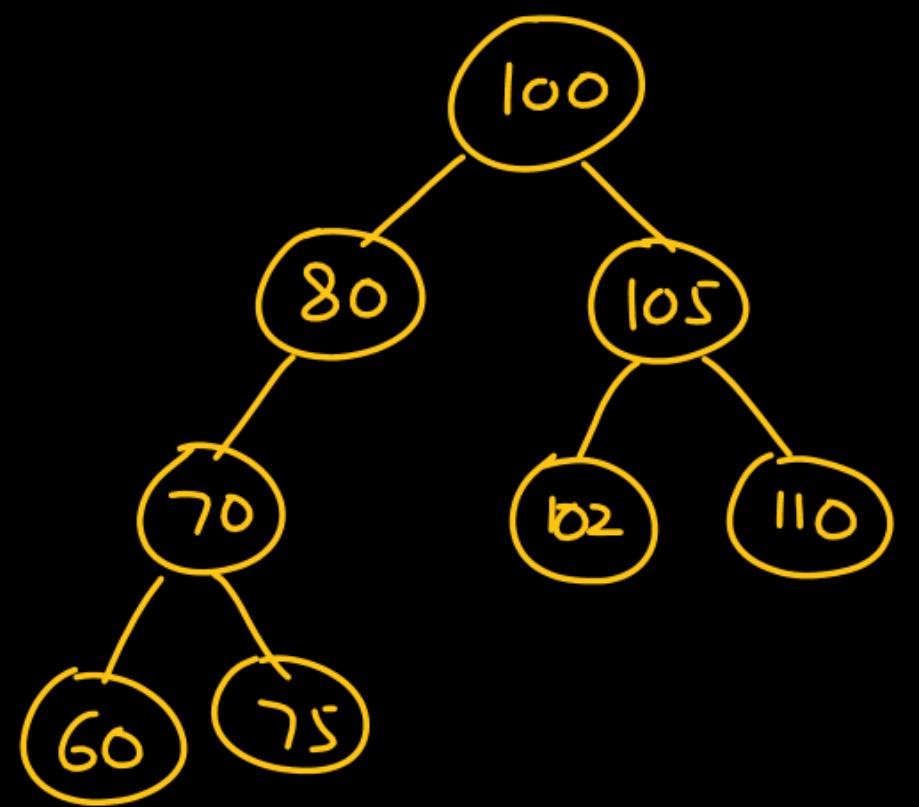
2 \* \* \* \*

Inorder traversal of a BST is always increasing  
order of keys.



Inorder : 70 80 90 92 95 98  
100 105 110 135

Pre : 100 80 70 60 75 105 102 110



Pre : 100 80 70 60 75 105 102 110

Given the preorder of a BST as : 100, 80, 70, 60, 75, 105, 102, 110  
the postorder of BST is \_\_\_\_\_

a)

b)

c)

d)

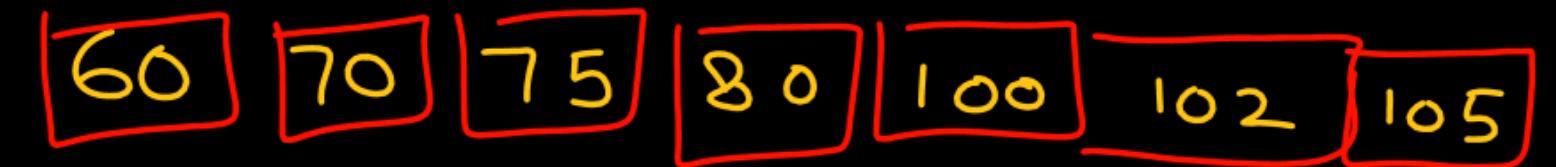
Pre : 100 80 70 60 75 105 102 110

Given the preorder of a BST as : 100, 80, 70, 60, 75, 105, 102, 110  
the postorder of BST is —

In: 60 70 75 80 100 102 105 110

Pre : 100 80 70 60 75 105 102 110

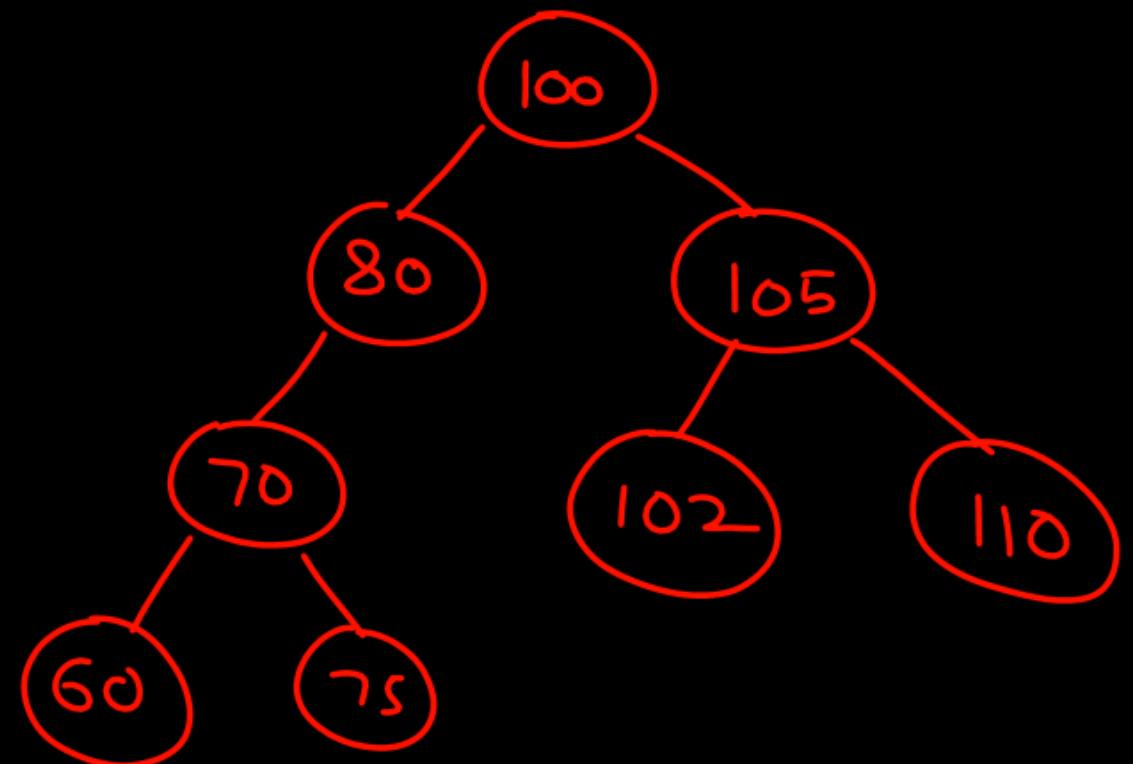
In:



110

Pre:

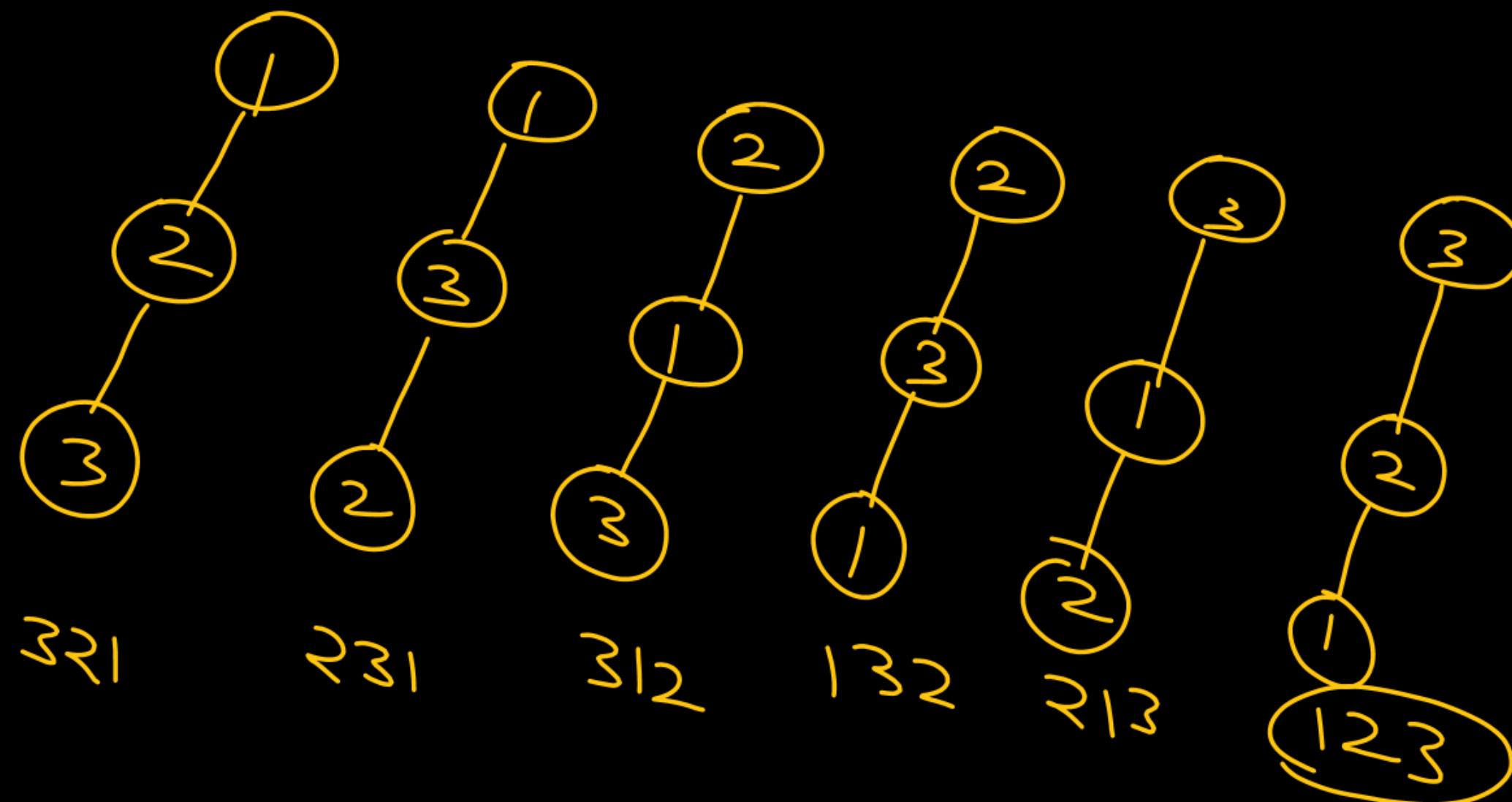
100 80 70 60 75 105 102 110  
→



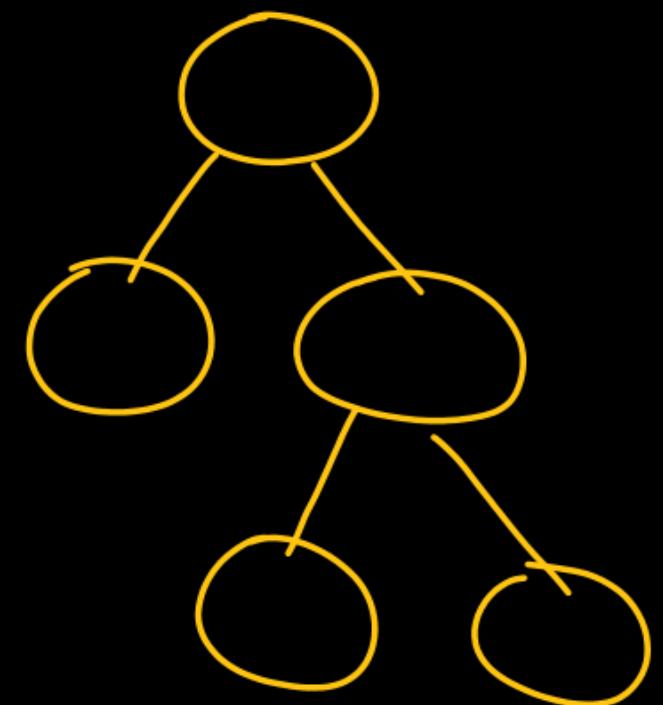
Q BST

Inorder :

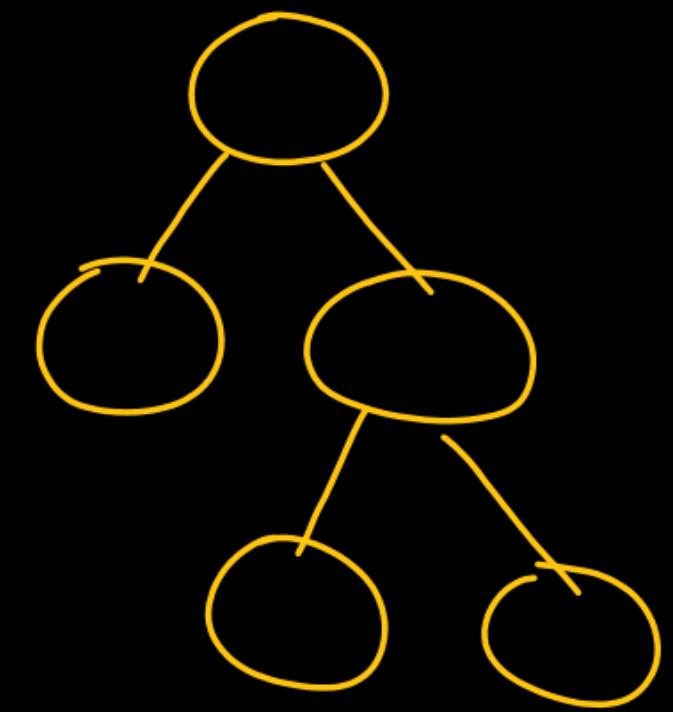
1, 2, 3 Keys



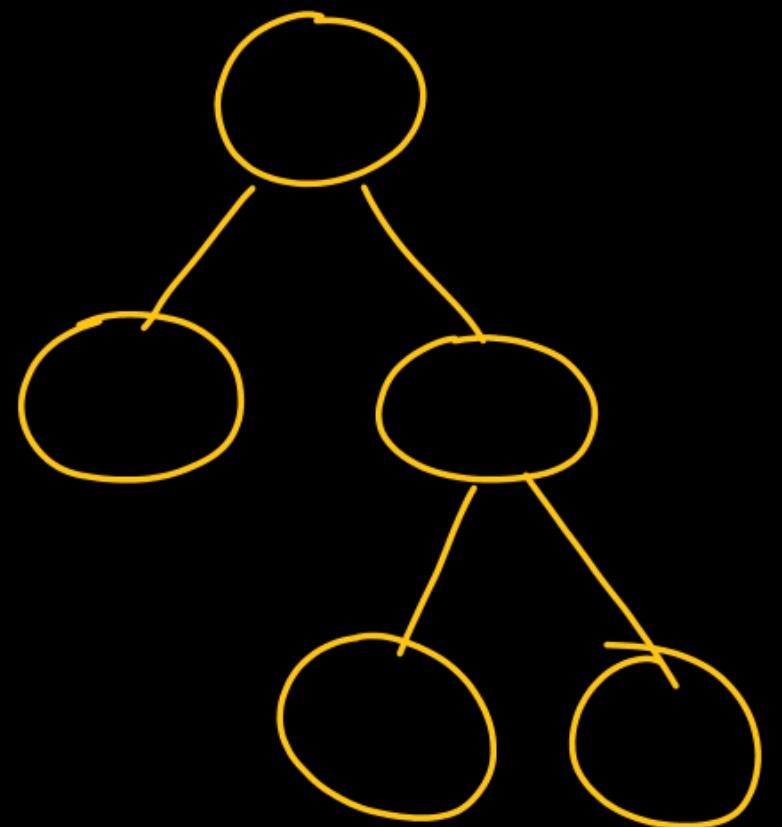
Given a binary tree structure with  $n$  nodes (unlabelled)



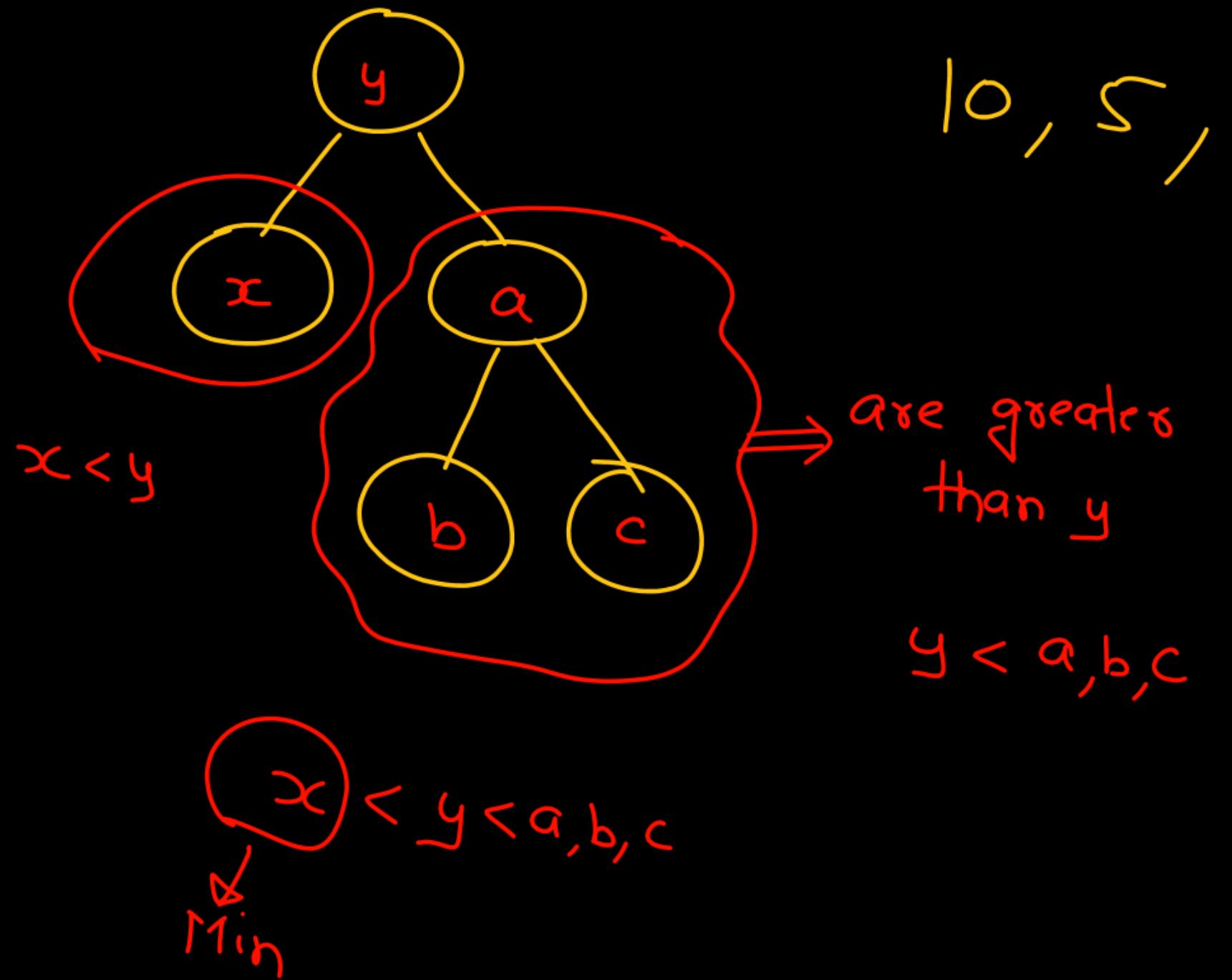
Given a binary tree structure with  $n$  nodes (unlabelled)



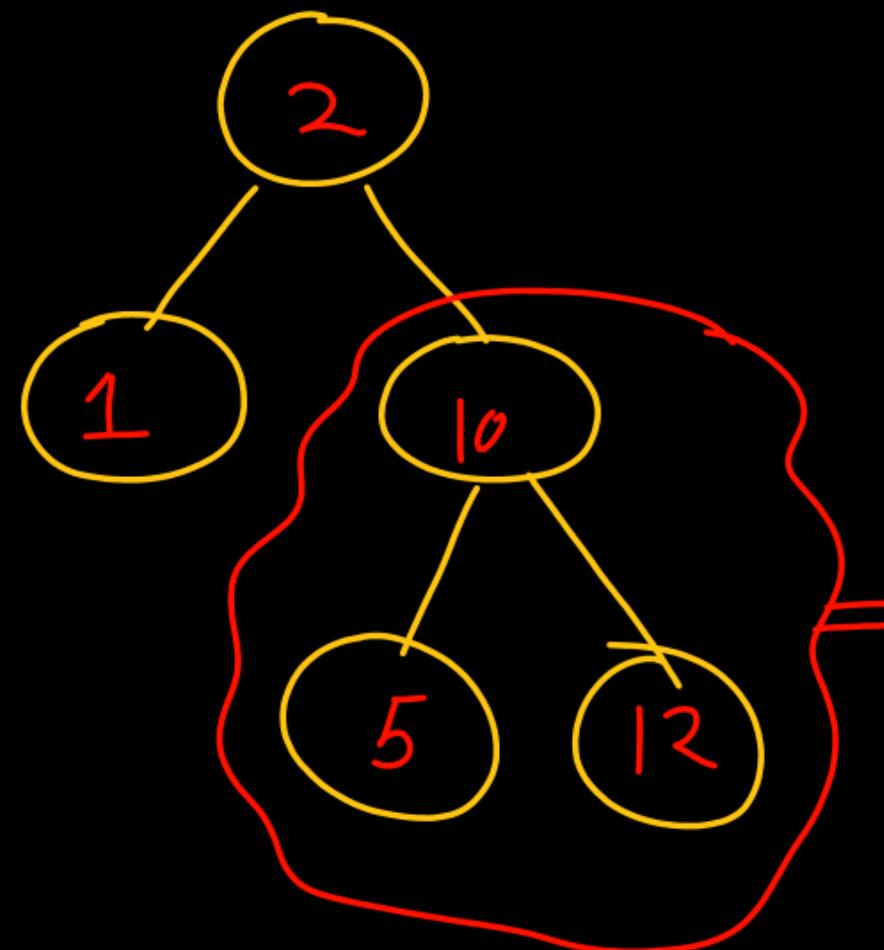
and  $n$  keys are also given  
how many BSTs are  
possible



10, 5, 12, 2, 1



10, 5, 12, 2, 1

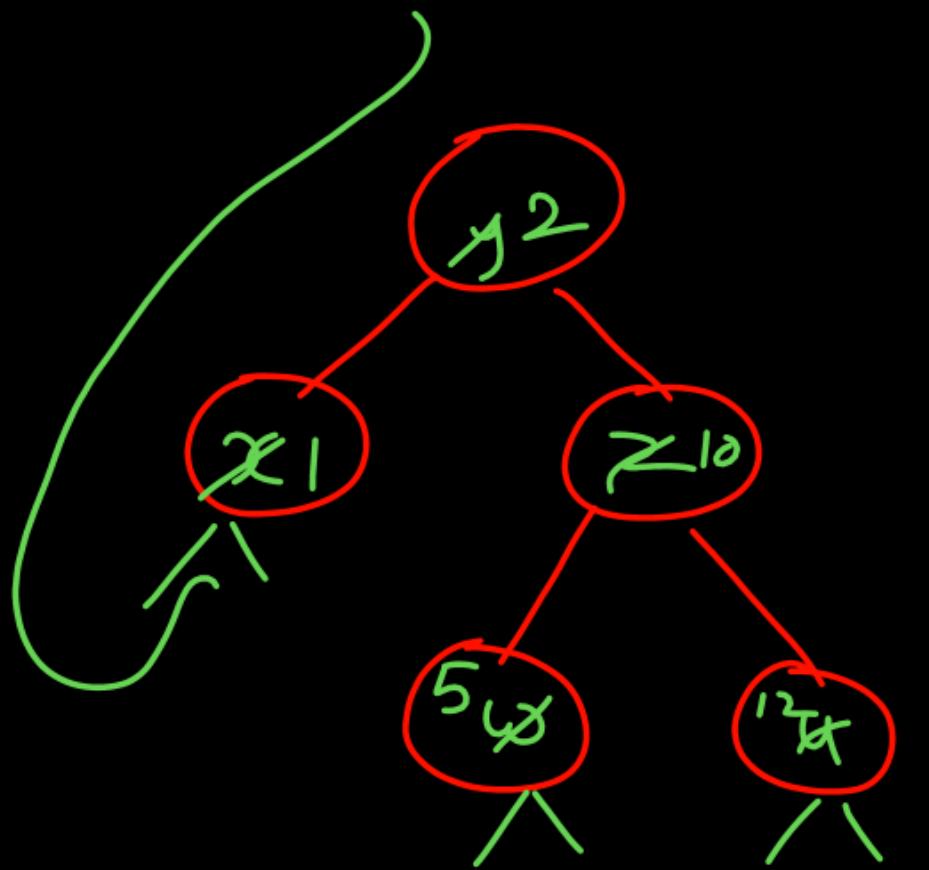


10, 5, 12, 2, 1

are greater  
than y

$y < a, b, c$

$x < y < a, b, c$   
Min



10, 5, 12, 2, 1  
→  
1 2 5 10 12

$x, y, w, z, q$

