

CS & IT ENGINEERING

Theory of Computation

Push Down Automata



Lecture No. 4



By- DEVA Sir



01 PDA & DPDA

02

→ Identify DPDA/PDA?

03

→ Construction

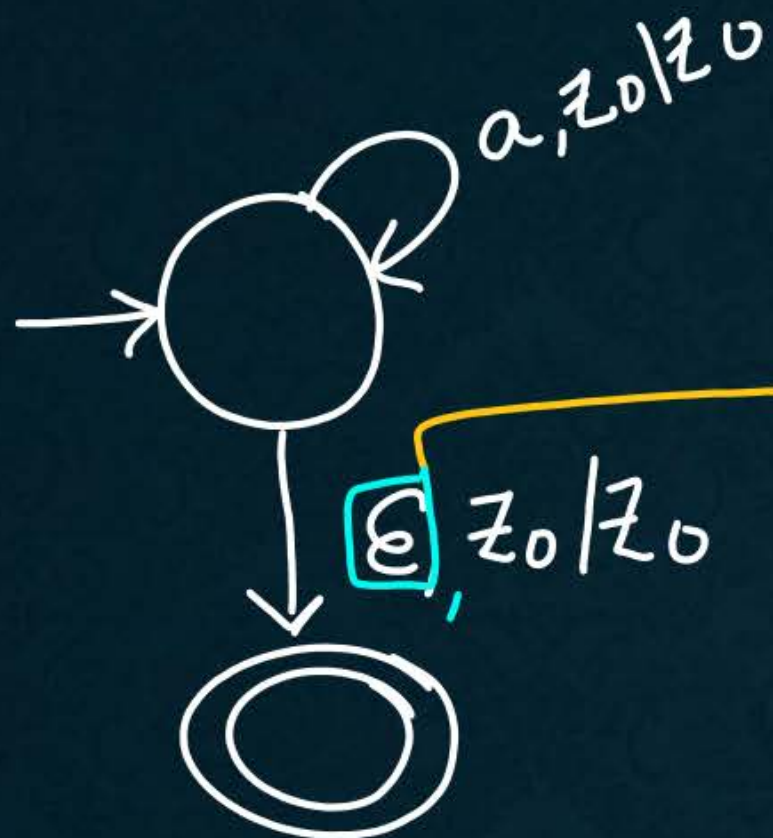
04

→ Identify language accepted by m/c?

05

Identify DPDA or not.
PDA or not.

①

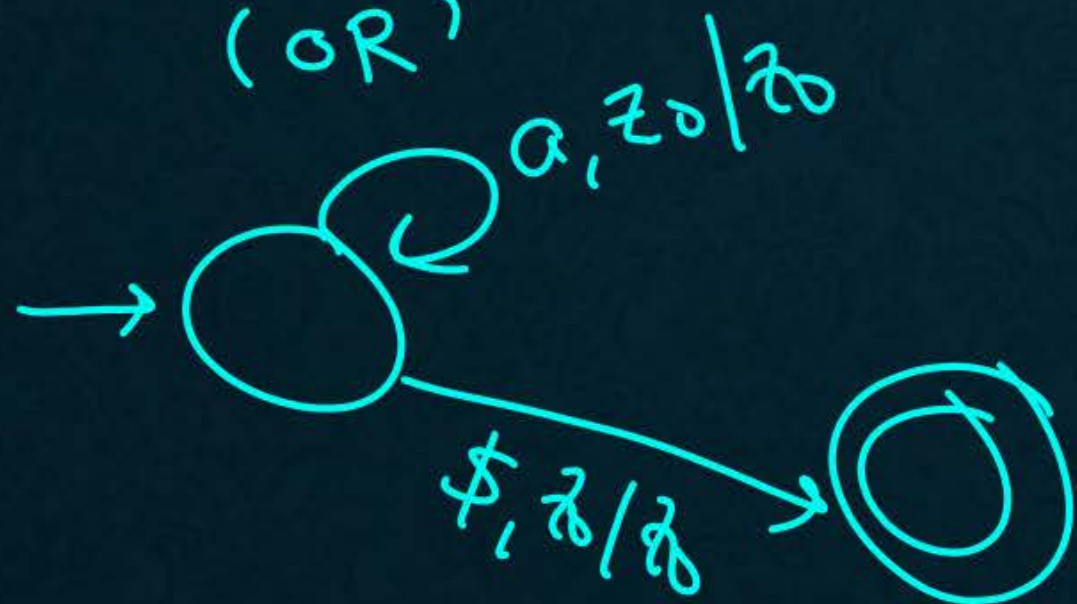


special, it represents end of i/p

DPDA ✓

PDA ✓

(OR)

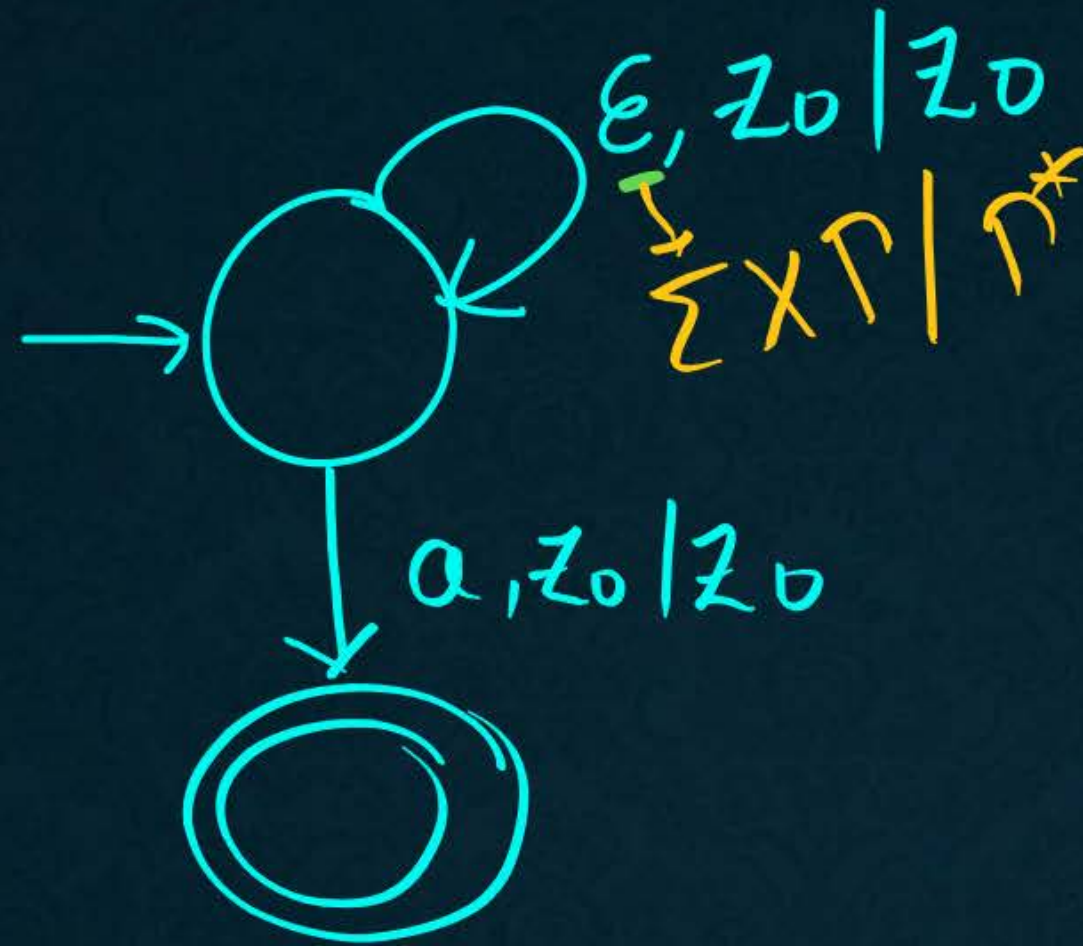


abab\$
ε

end of i/p

In DPDA
We assume
special transition
requires to represent
end of i/p

2



PDA ✓

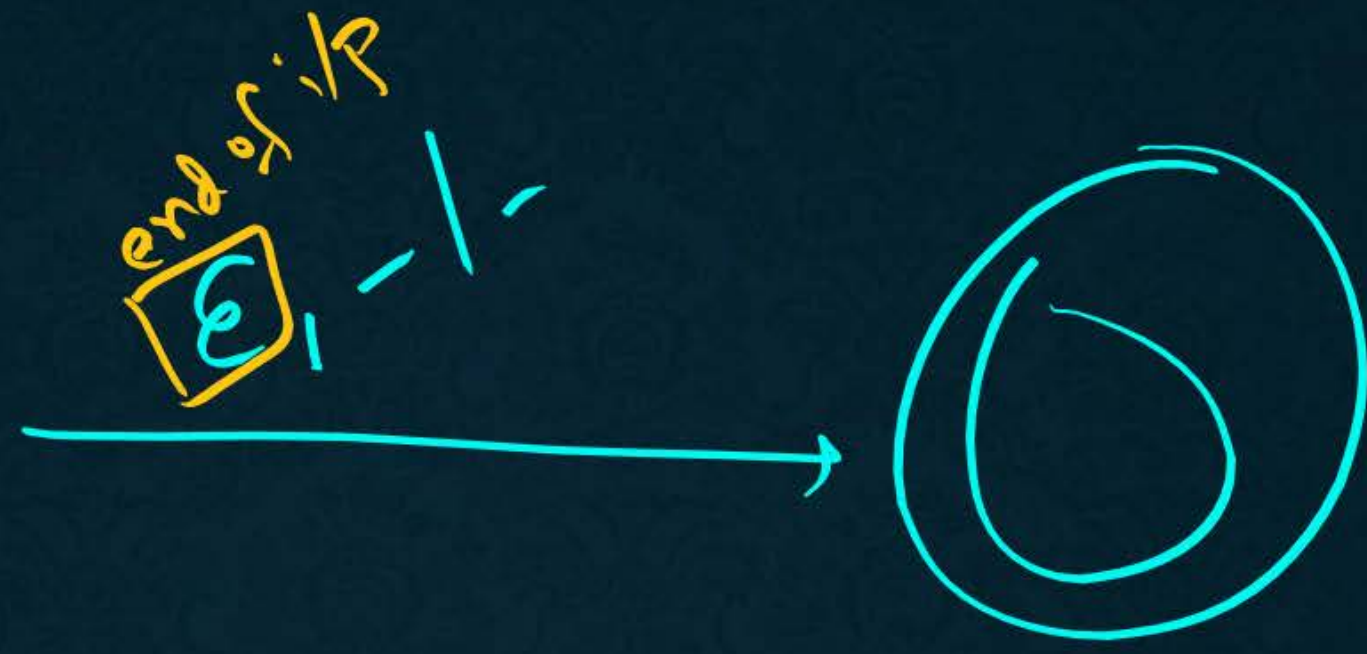
DPDA ✗

3



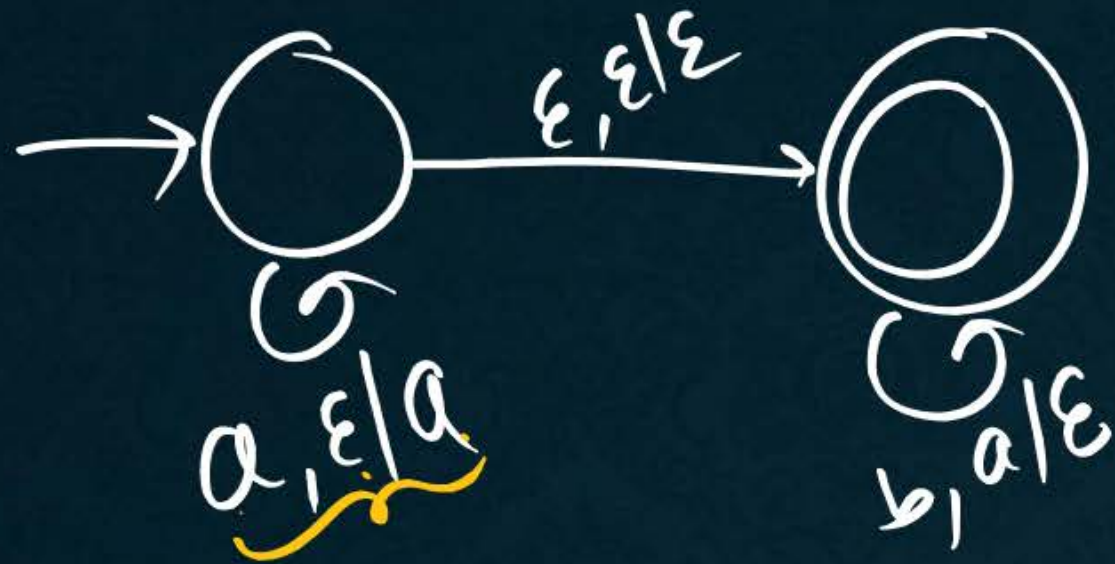
DPDA ✓

PDA ✓



It is allowed in DPAA

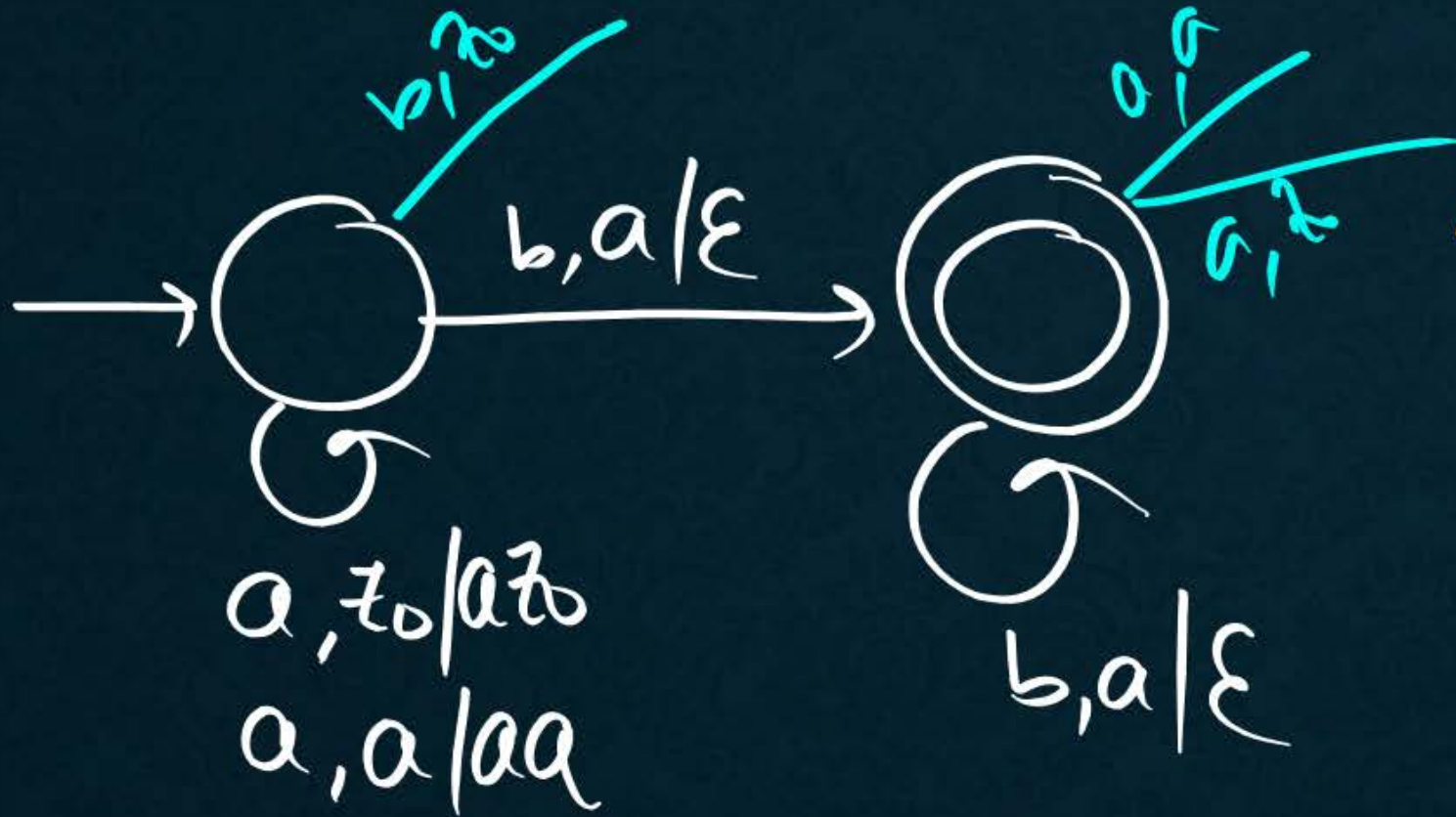
④



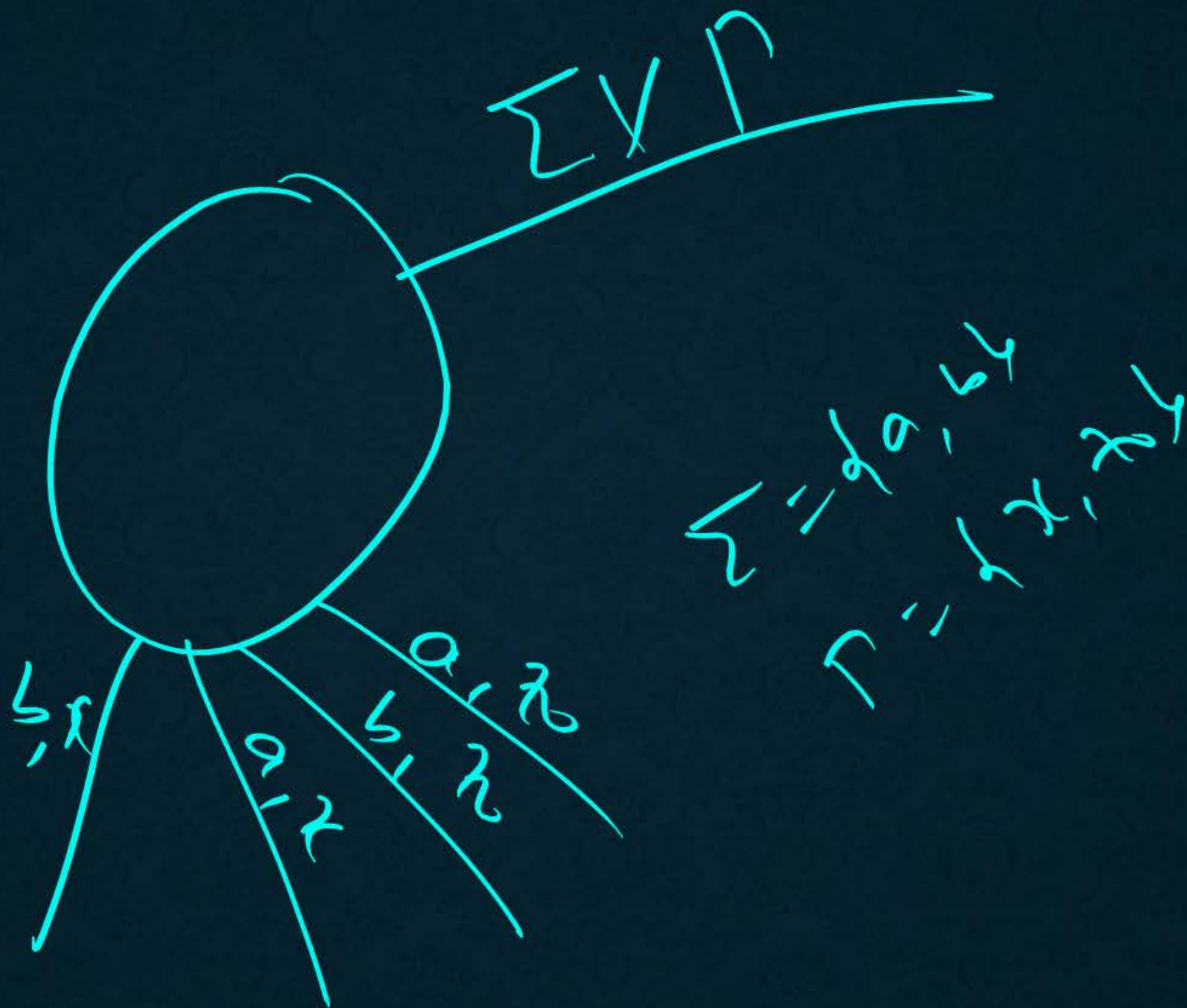
PDA ✓

DPDA ✗

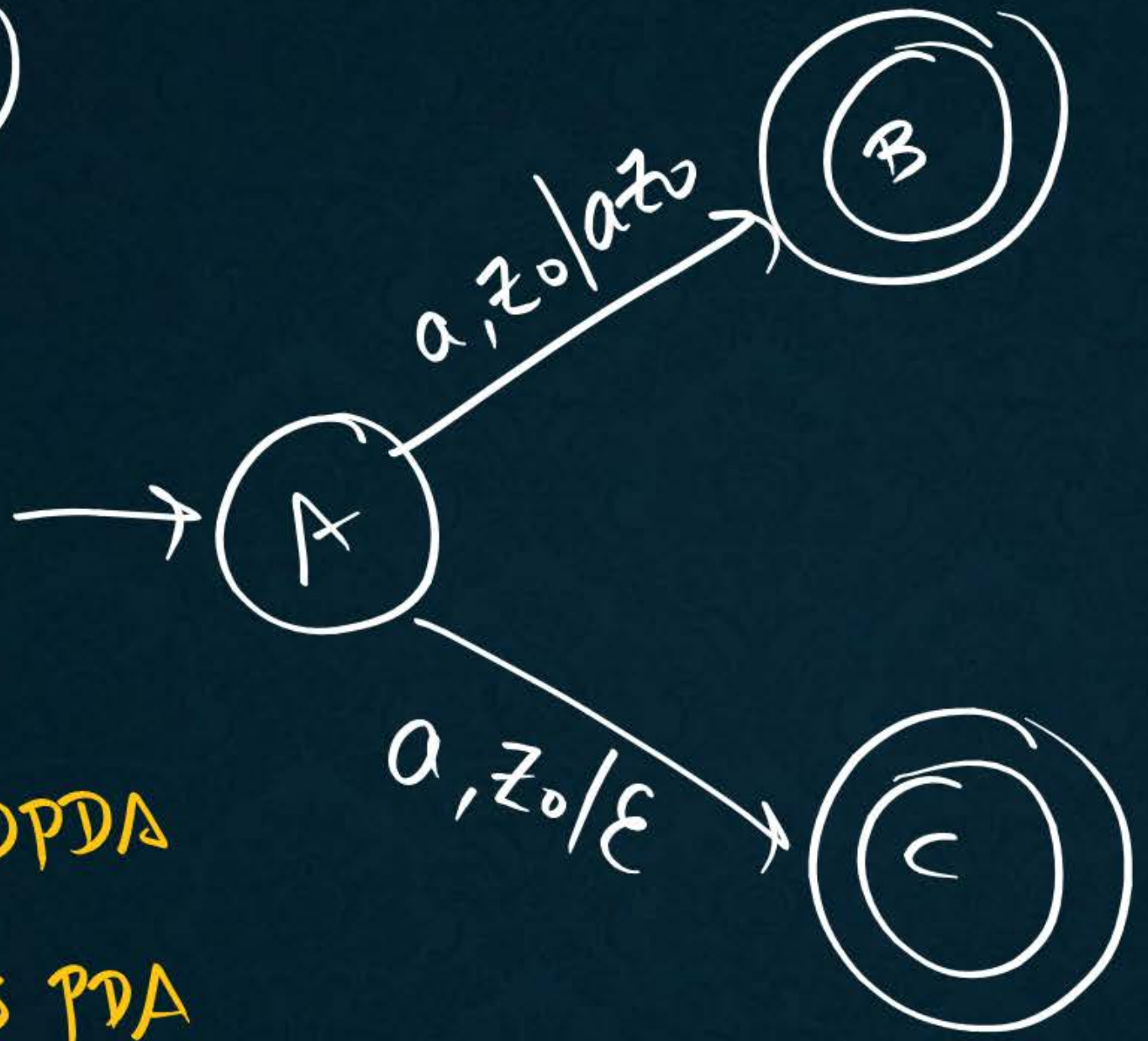
⑤



DPDA ✓
PDA ✓



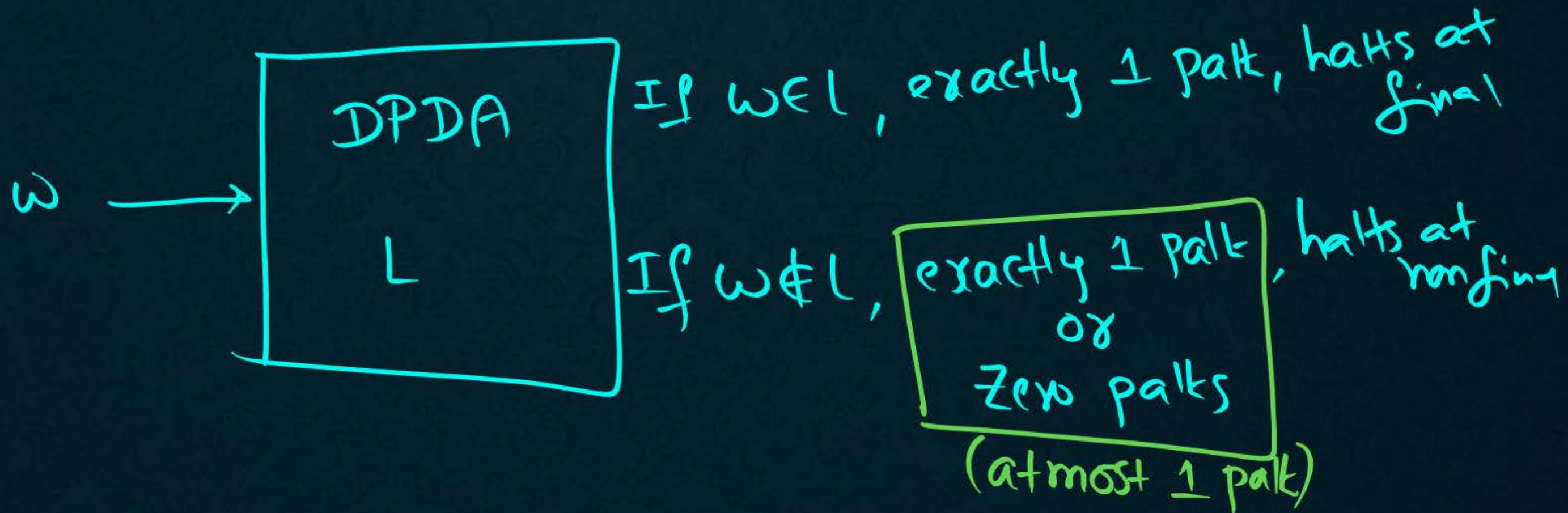
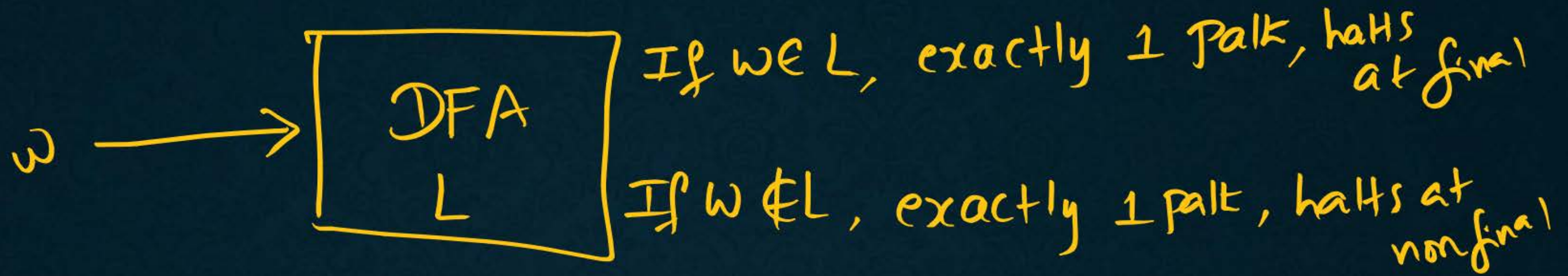
⑥



Not DPDA
It is PDA

$QX\Sigma X\Gamma \rightarrow QX\Gamma^*$
fails

$QX\Sigma_\epsilon X\Gamma^* \rightarrow QX\Gamma^*$
 $\delta(A, a, z_0) = (B, az_0)$
 $(C, \text{OR } \epsilon)$



Zero paths

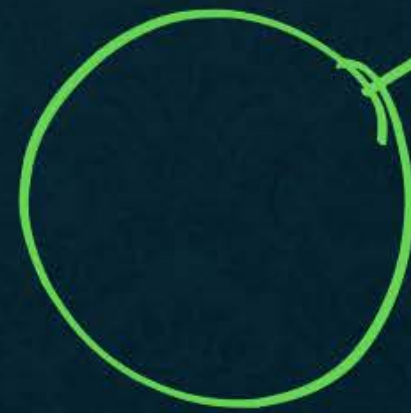


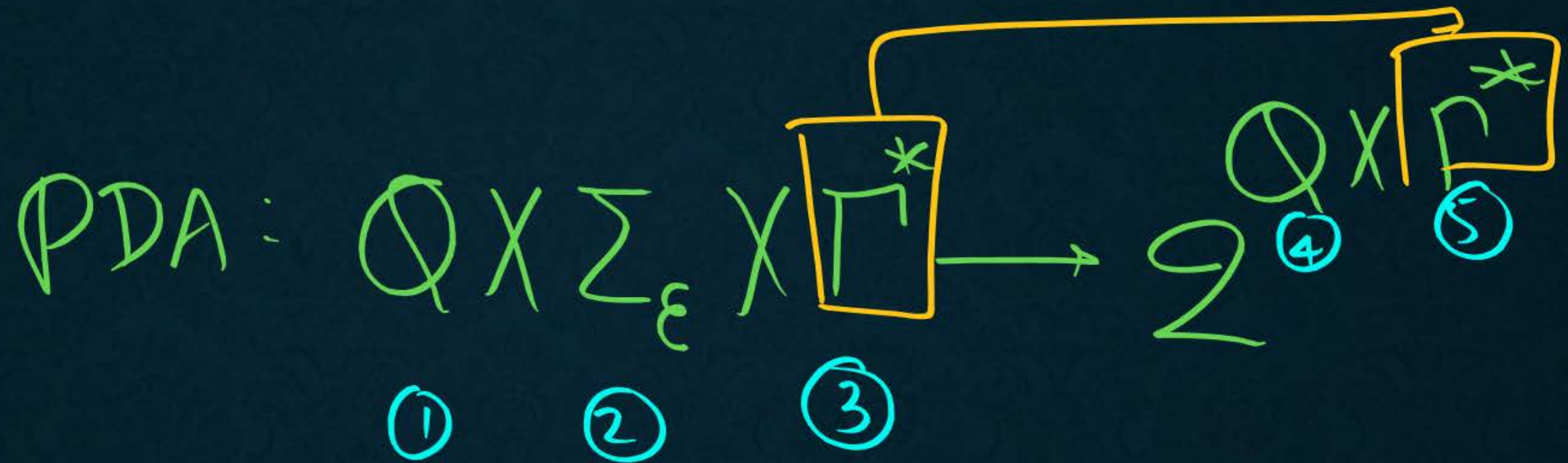
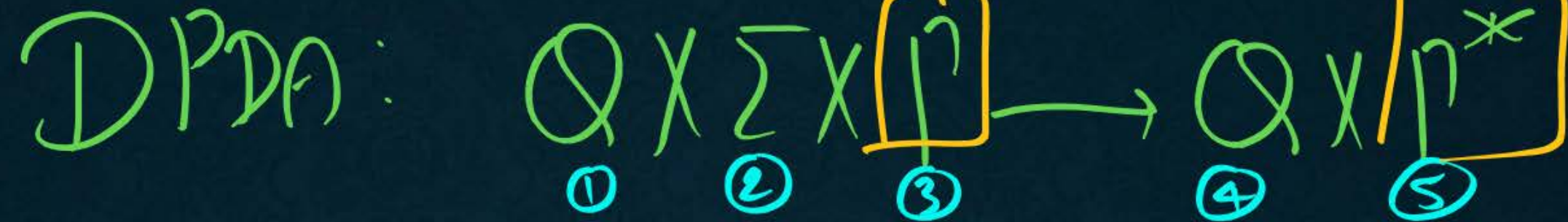
always

convertible to 1 path

$\sum a_i$ is minimum

assume it is going to lead state

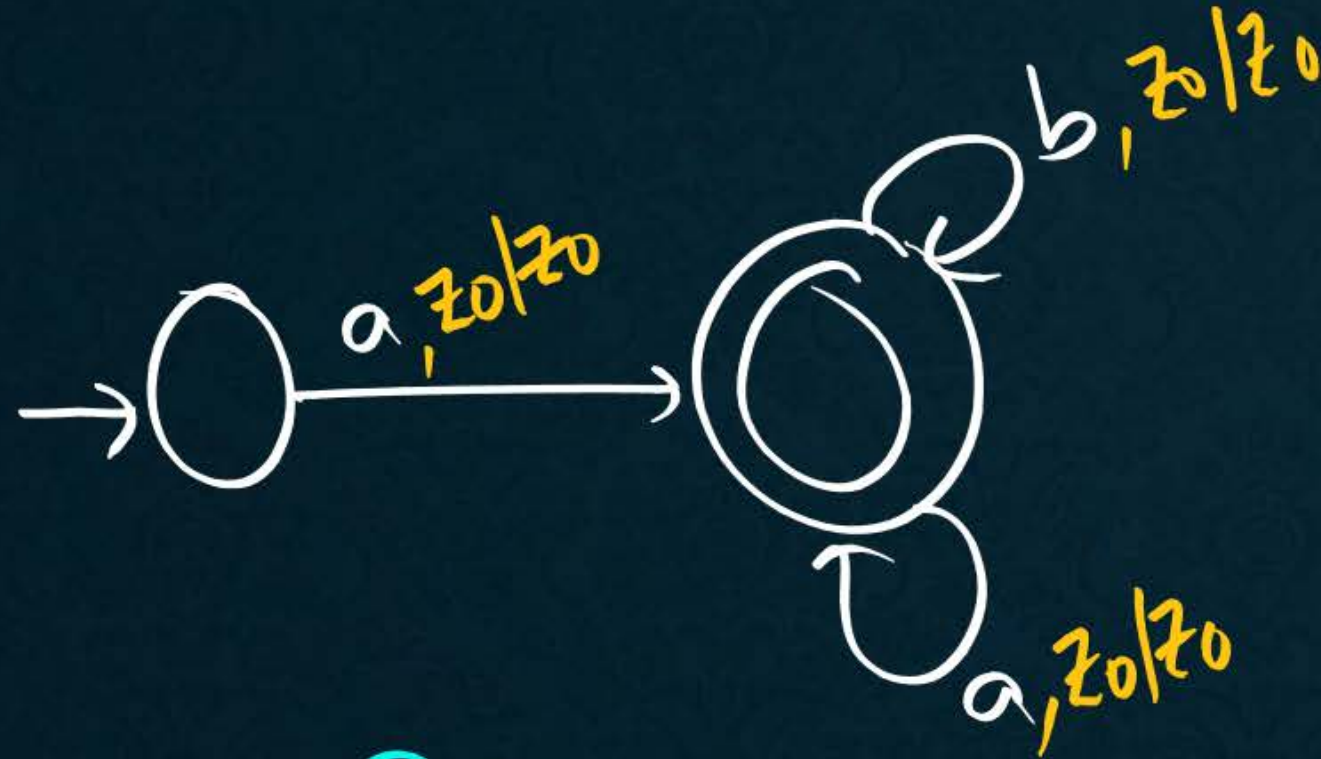
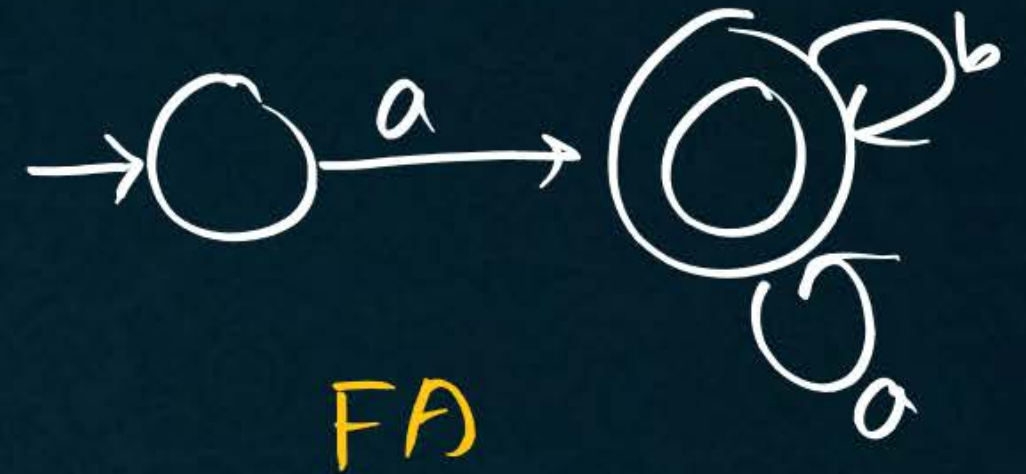




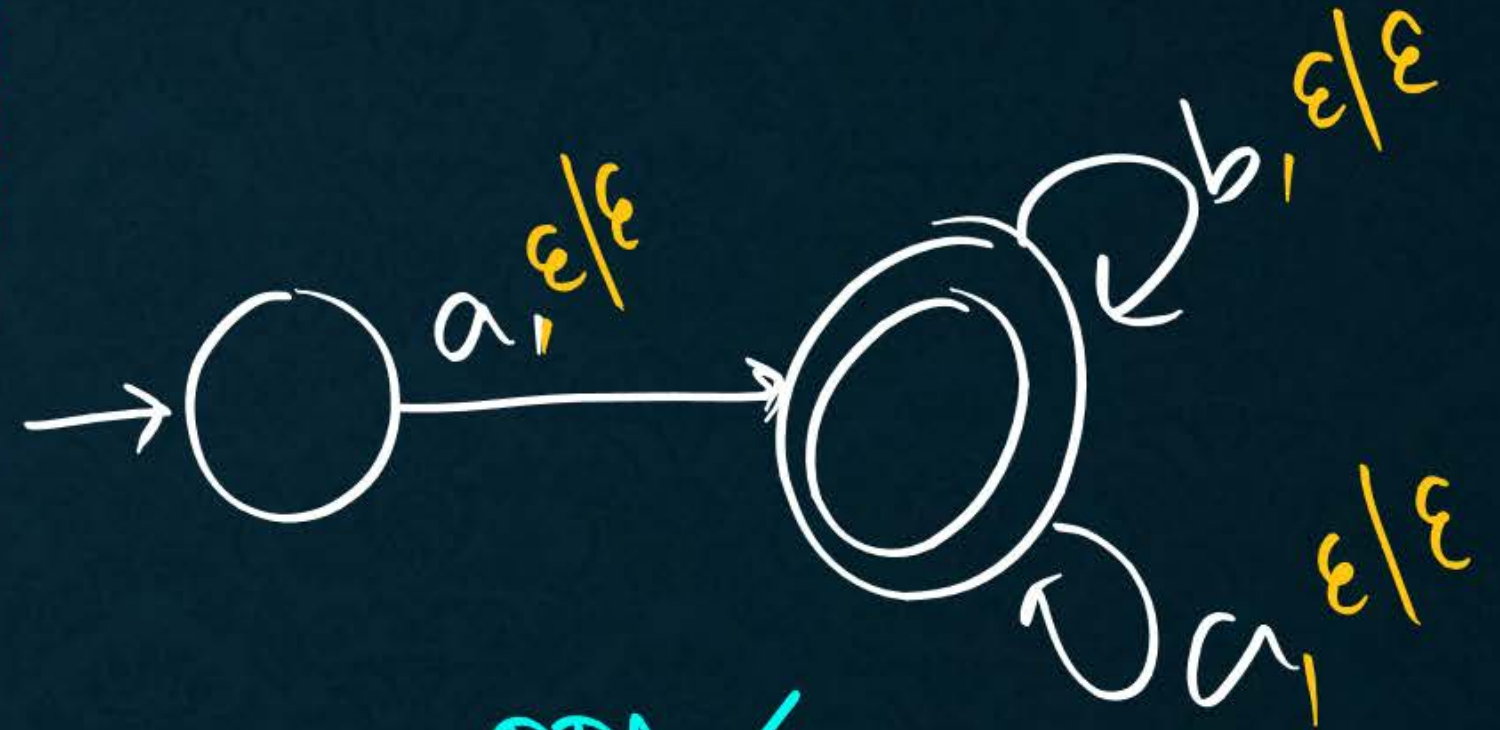
How to find operation?
 ③/⑤

Construction of PDA & DPDA

① $L = a(a+b)^*$



DPDA ✓
PDA ✓

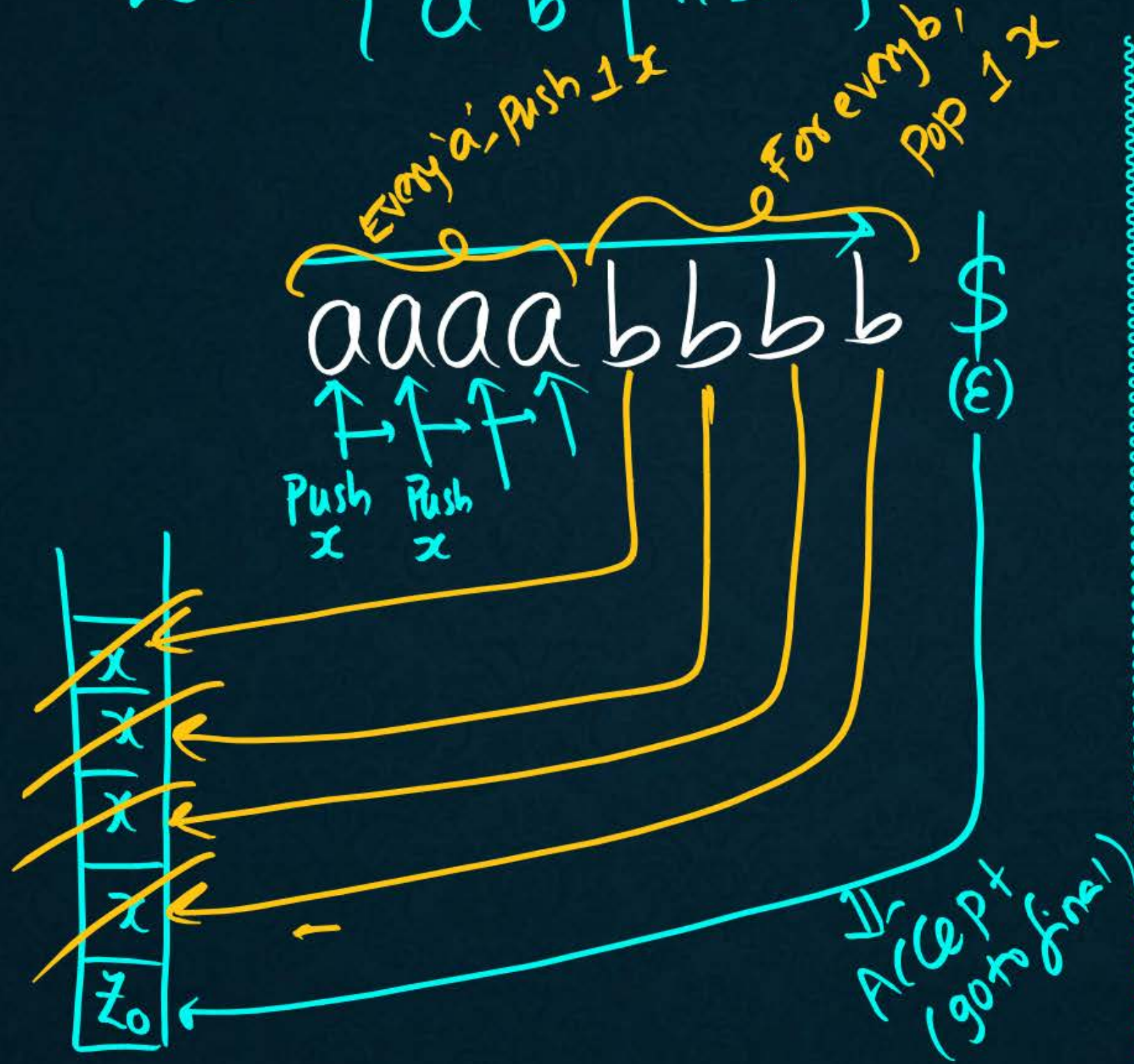


PDA ✓
DPDA ✗

Note : Every FA is convertible to DPDA & PDA

②

$$L = \{ a^n b^n \mid n \geq 1 \}$$



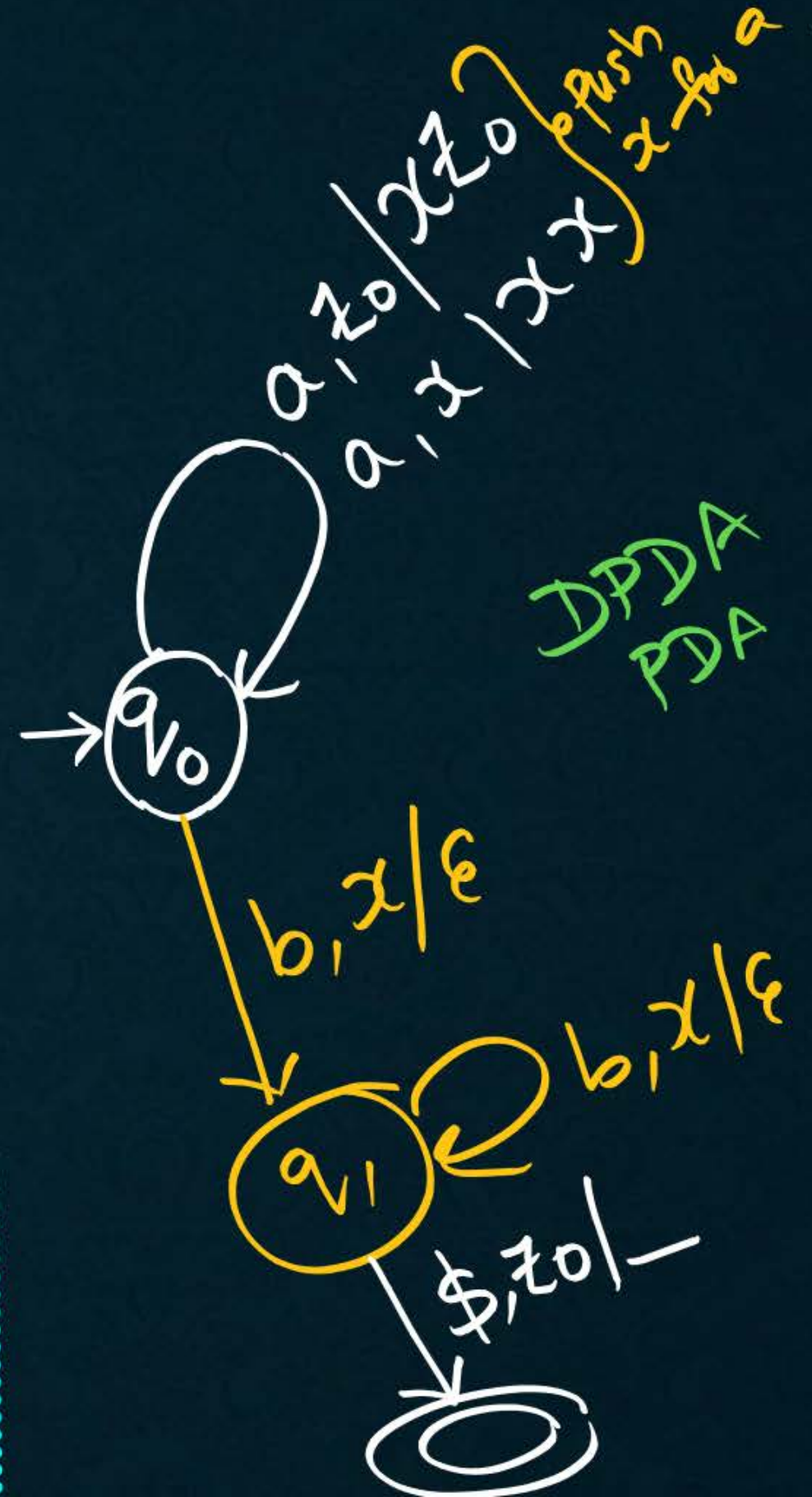
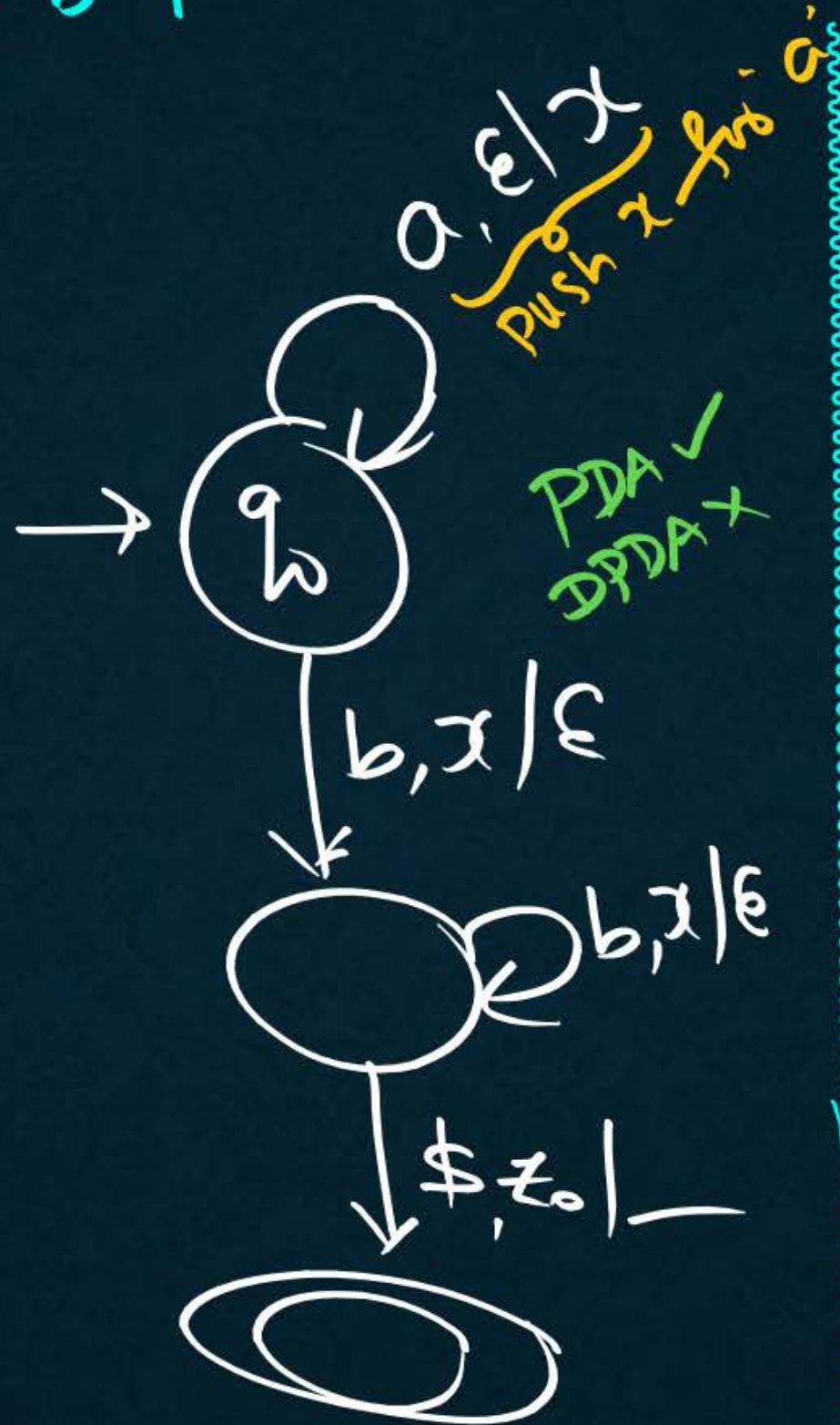
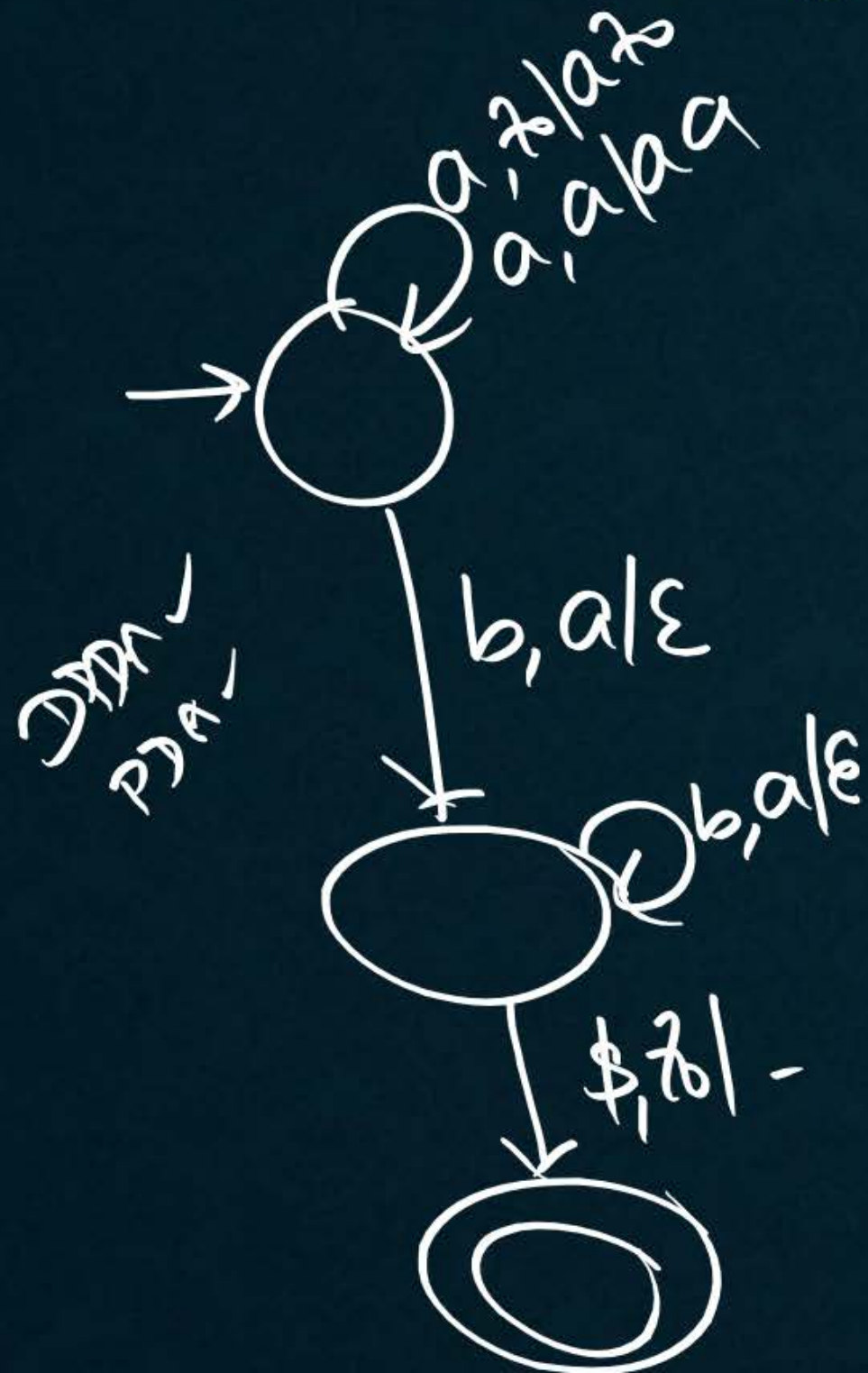
Step 1: Take bigger valid string

Step 2: Make logic using stack

Step 3: Construct PDA

②

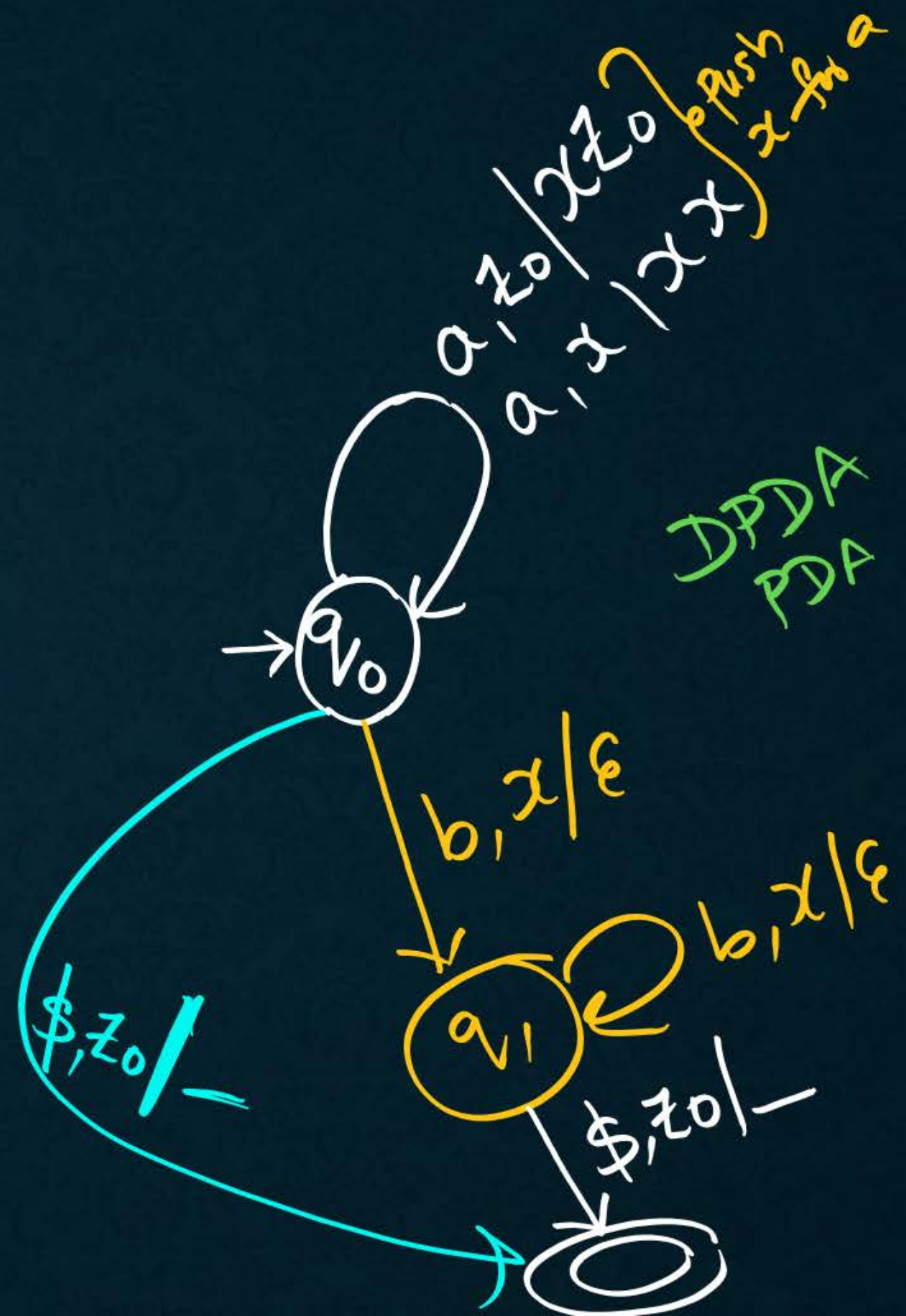
$$L = \{ a^n b^n \mid n \geq 1 \}$$

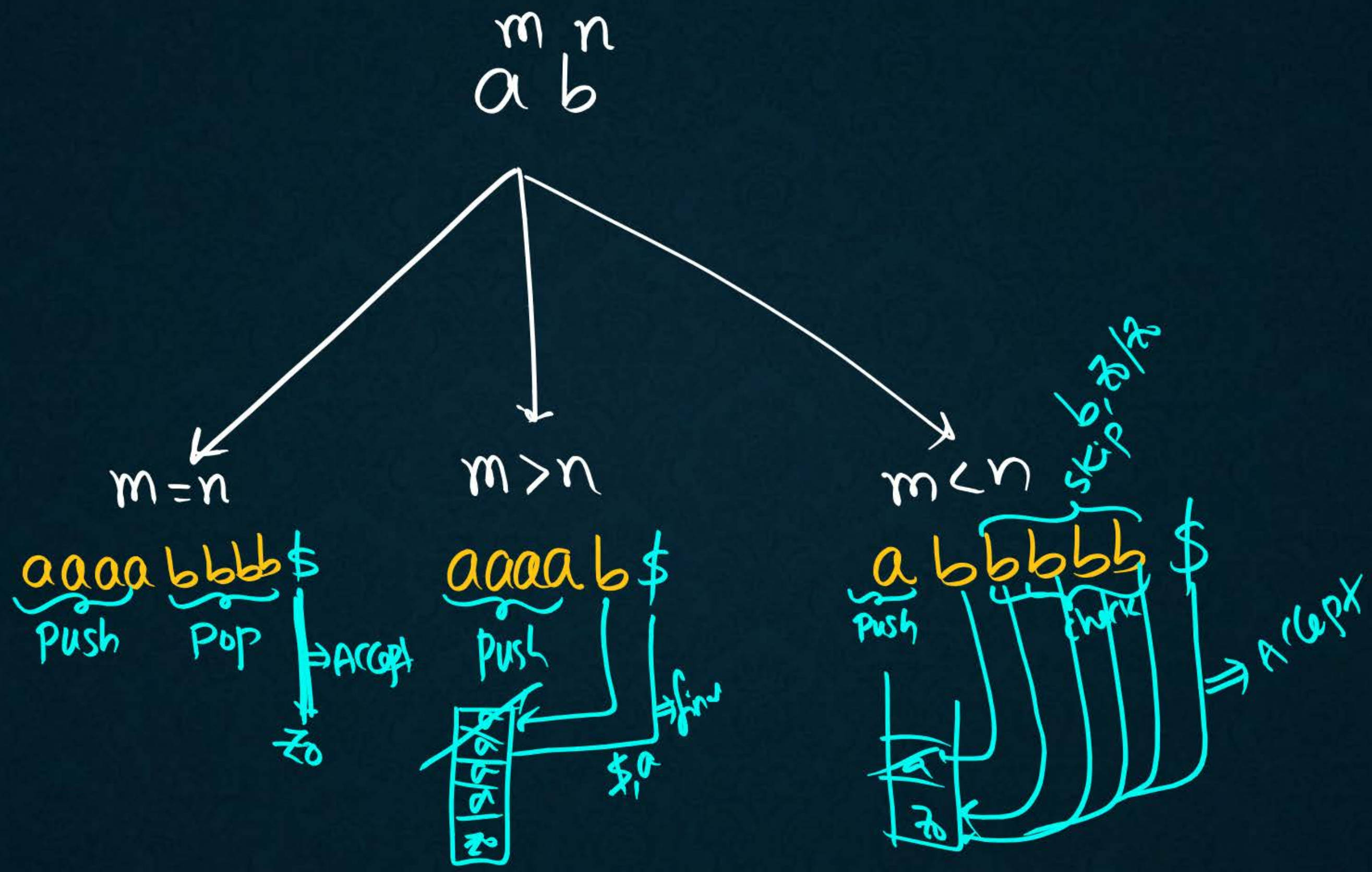


③

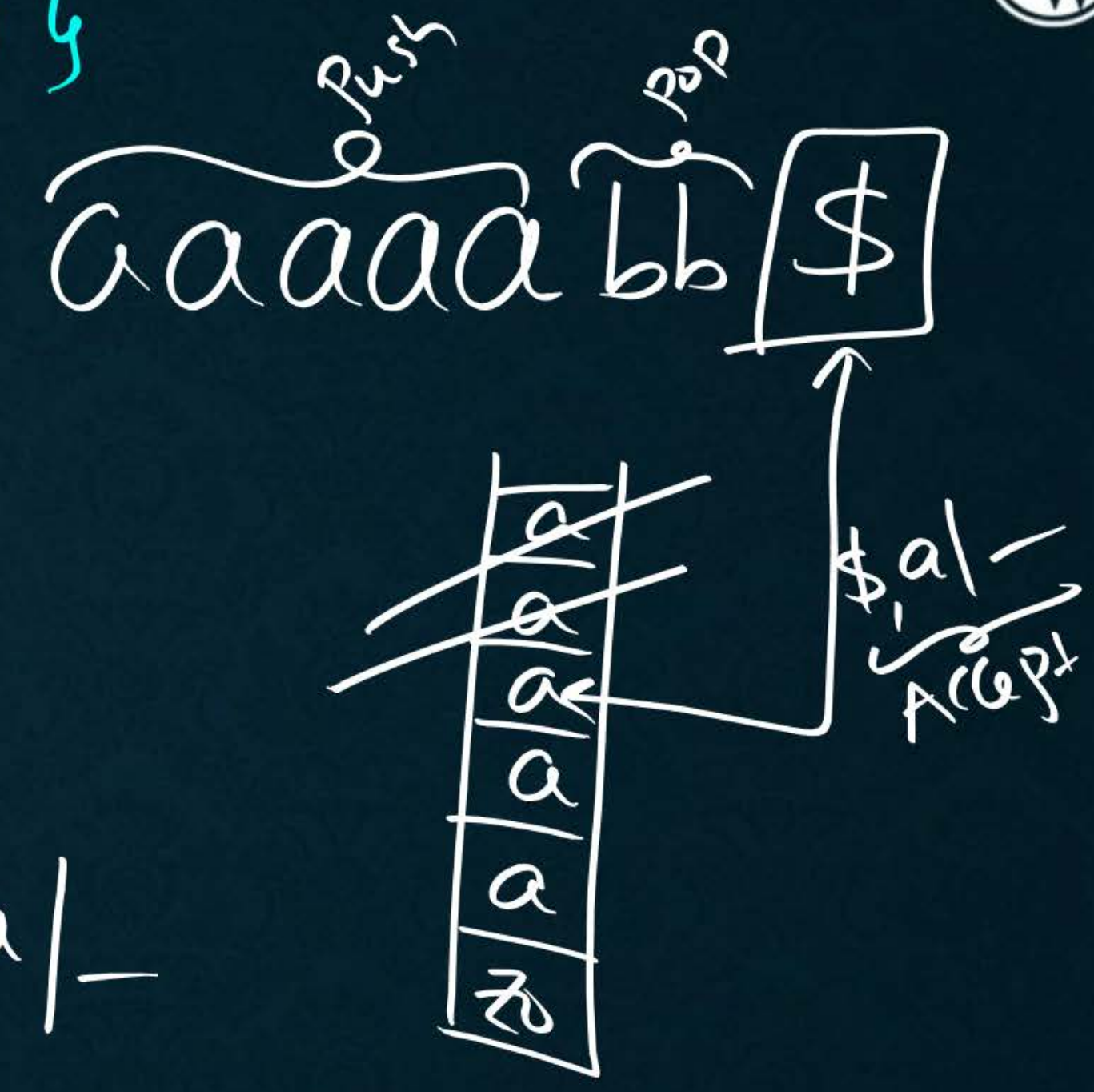
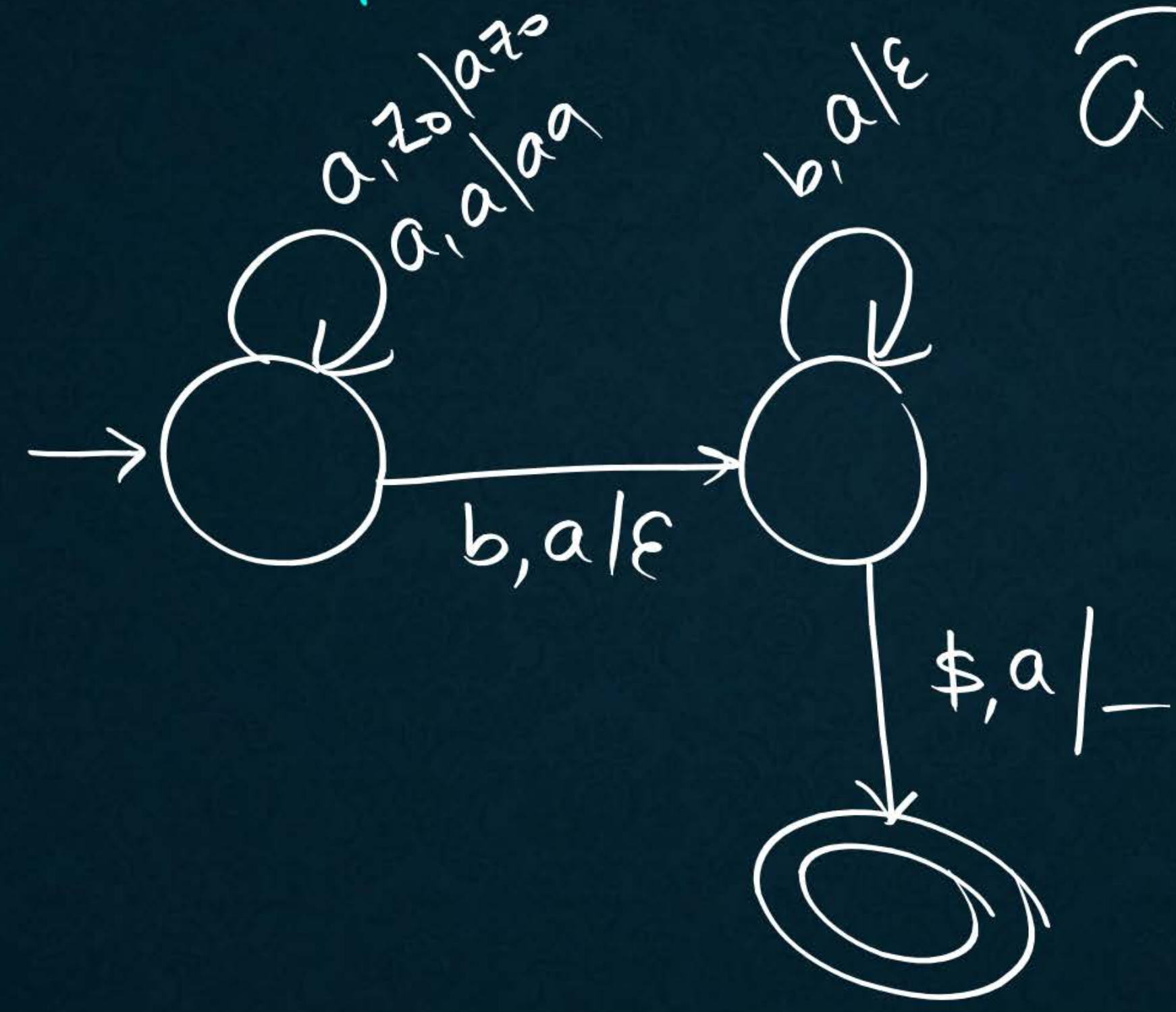
$$L = \{a^n b^n \mid n \geq 0\}$$

$$= \textcircled{2} \cup \{\epsilon\}$$

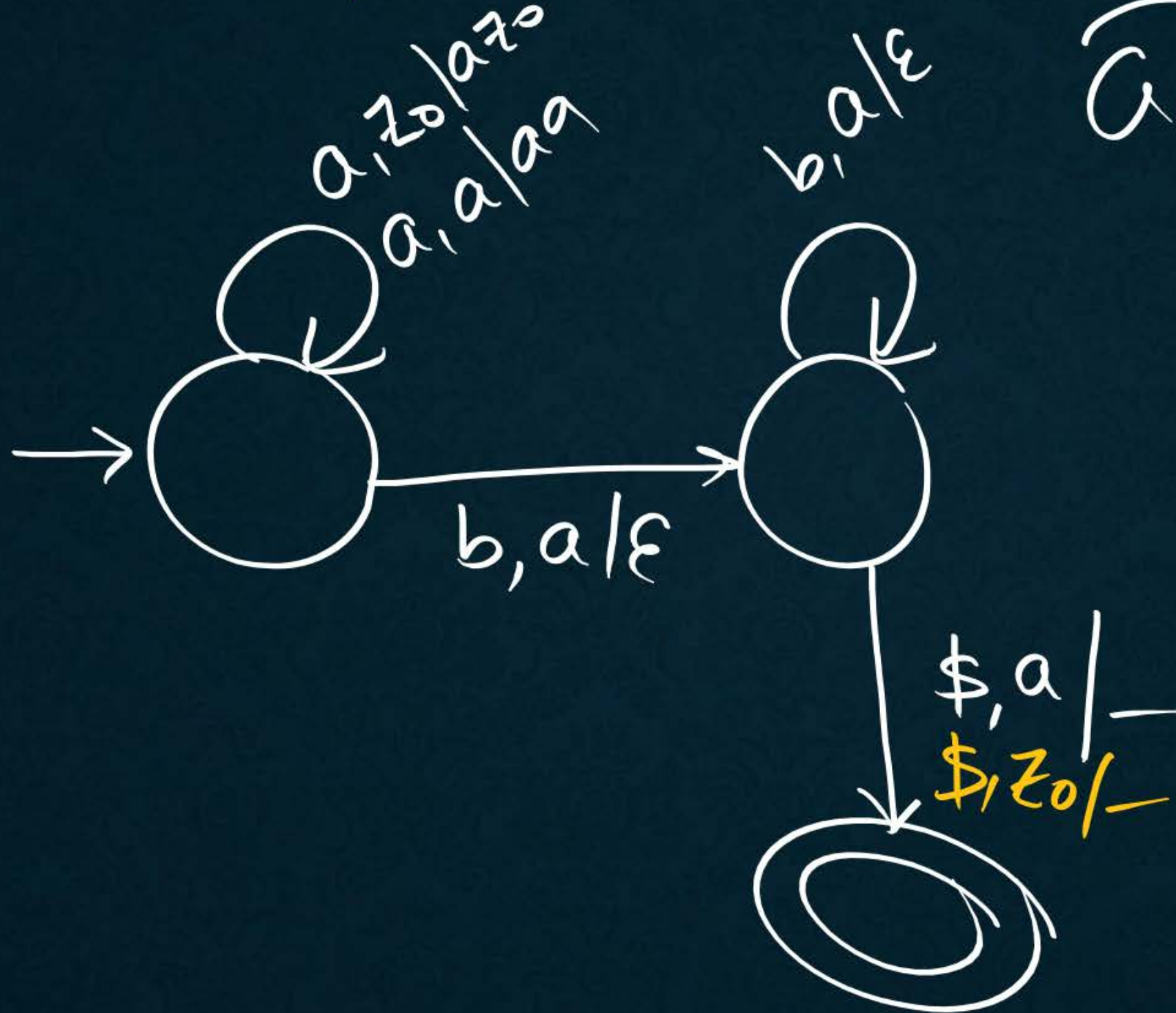




④ $\{a^m b^n \mid m > n, m, n \geq 1\}$

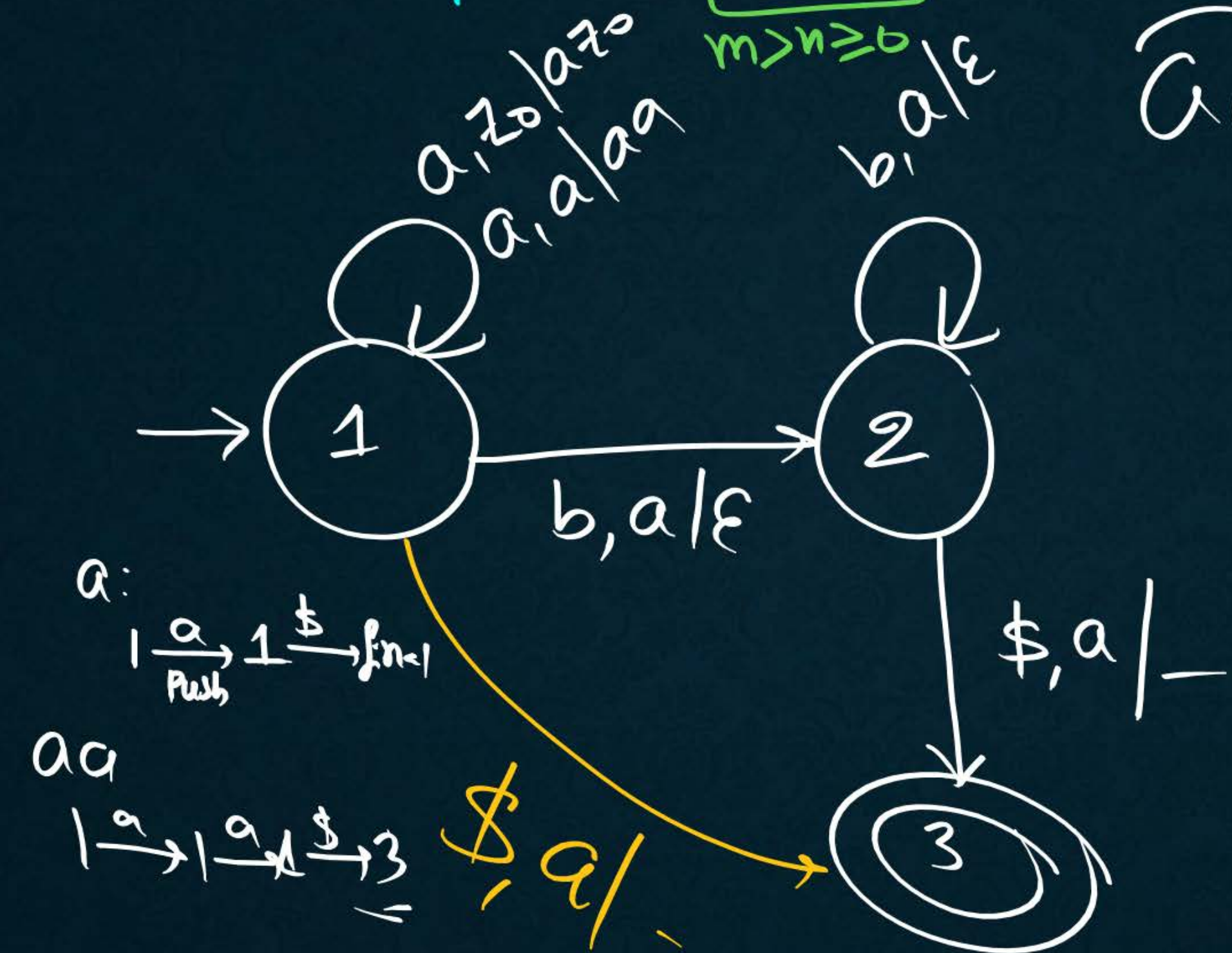


5) $\{a^m b^n \mid m \geq n, m, n \geq 1\}$



⑥ $\{a^m b^n \mid \underbrace{\#a's > \#b's}_{m > n}, \underbrace{m, n \geq 0}_{m > n \geq 0}\} = \textcircled{4} \cup \textcircled{4}^?$

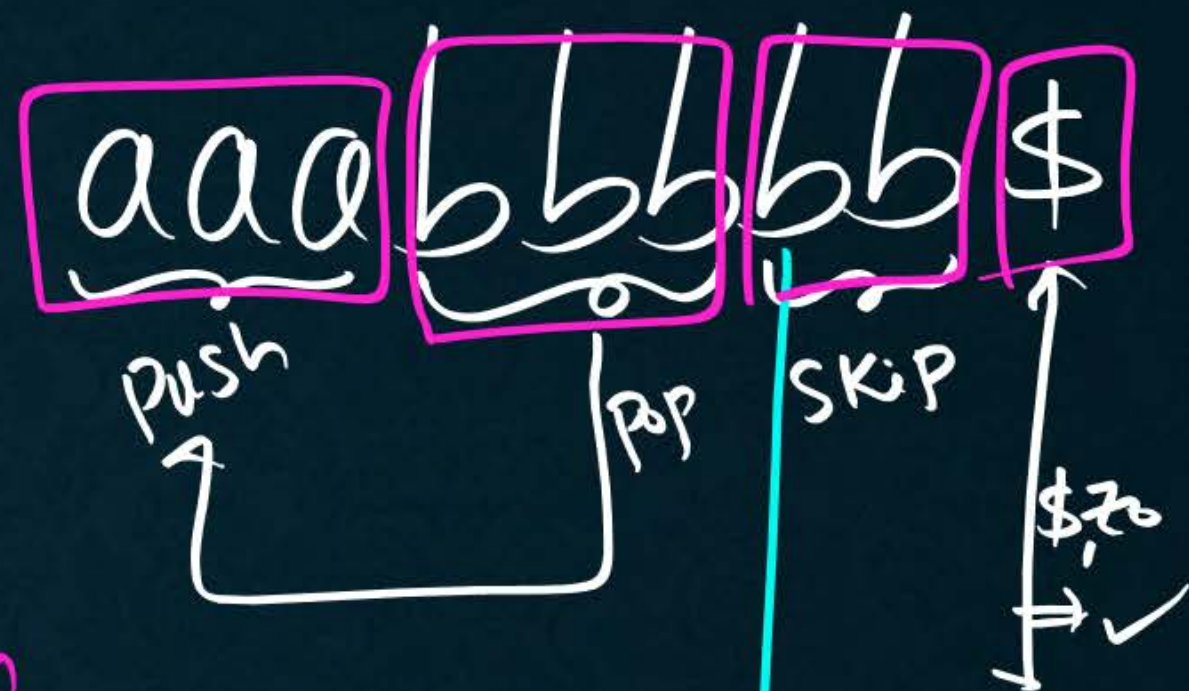
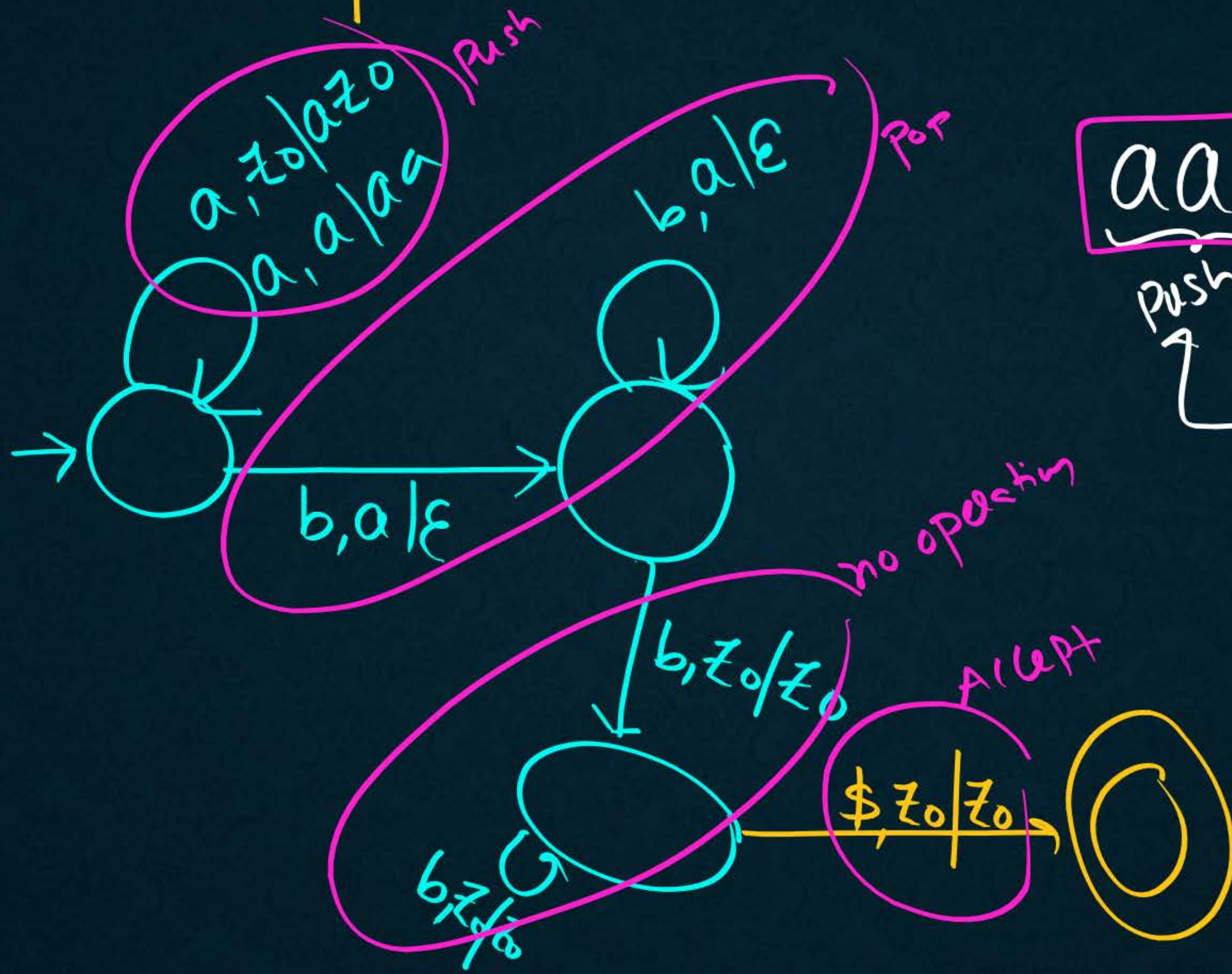
$\underbrace{aaaaa}_{\text{push}} \underbrace{bb}_{\text{pop}} \boxed{\$}$



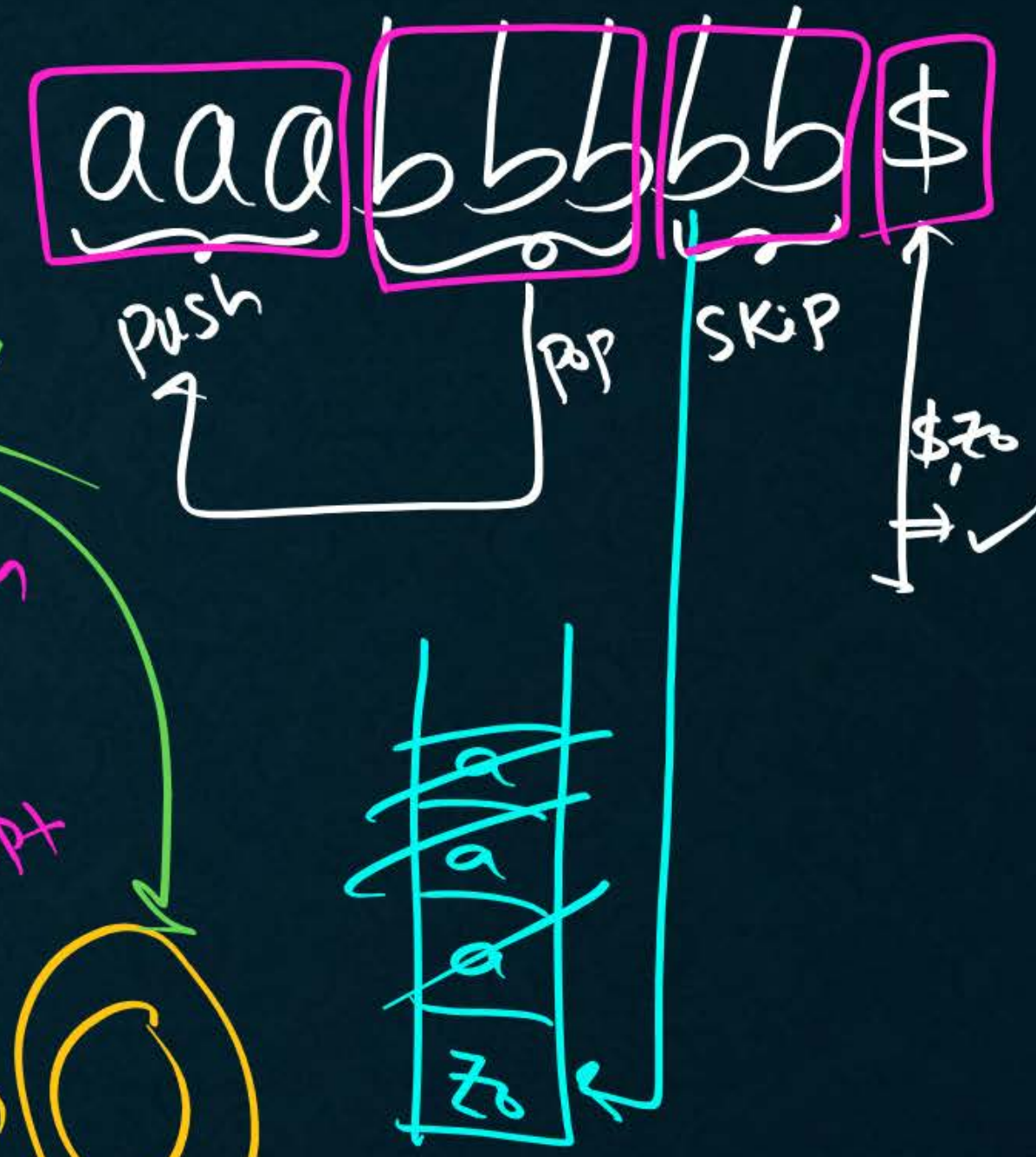
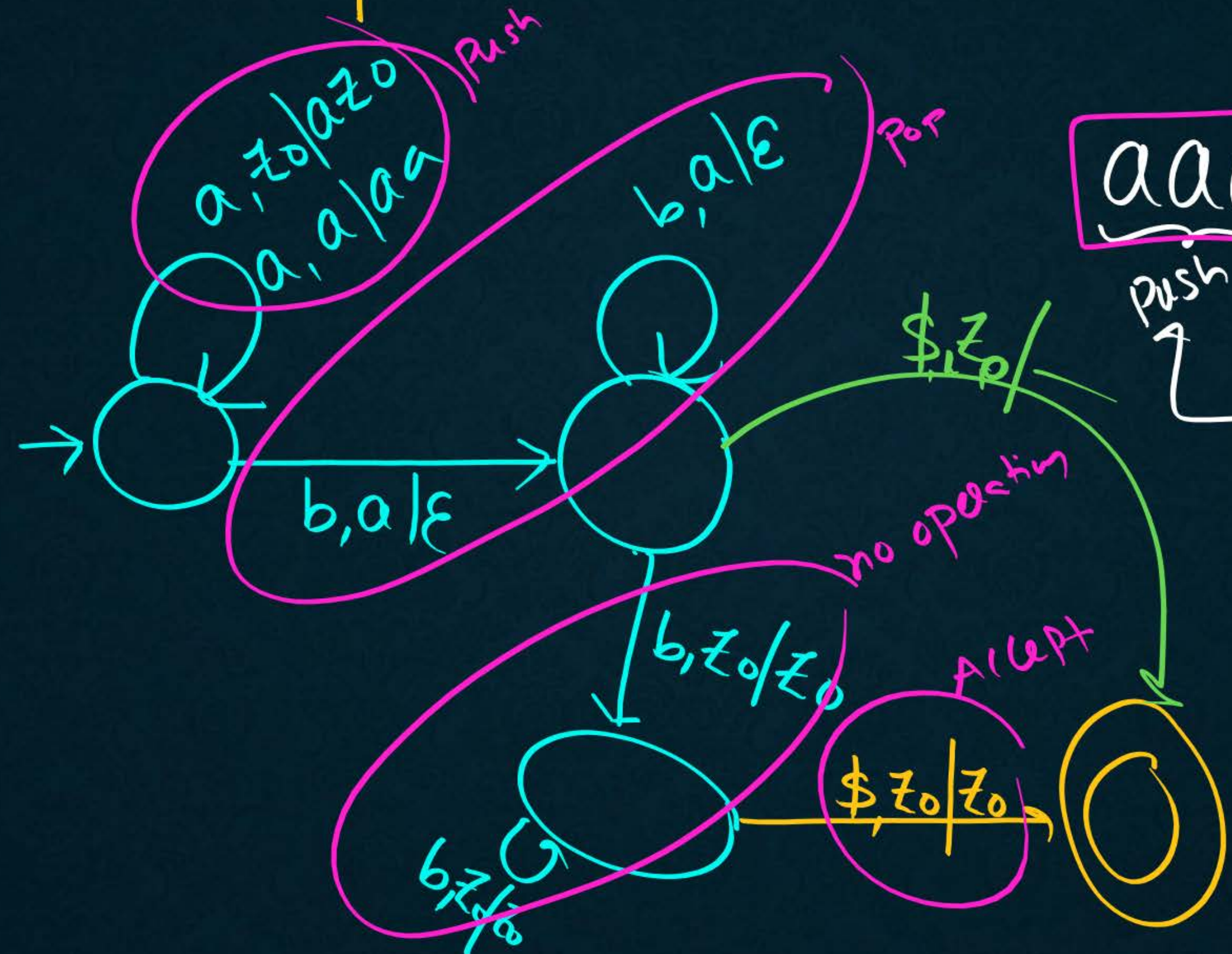
$$\underbrace{m > n}, \underbrace{m, n \geq 0}$$



$$\underbrace{m > n \geq 0}$$

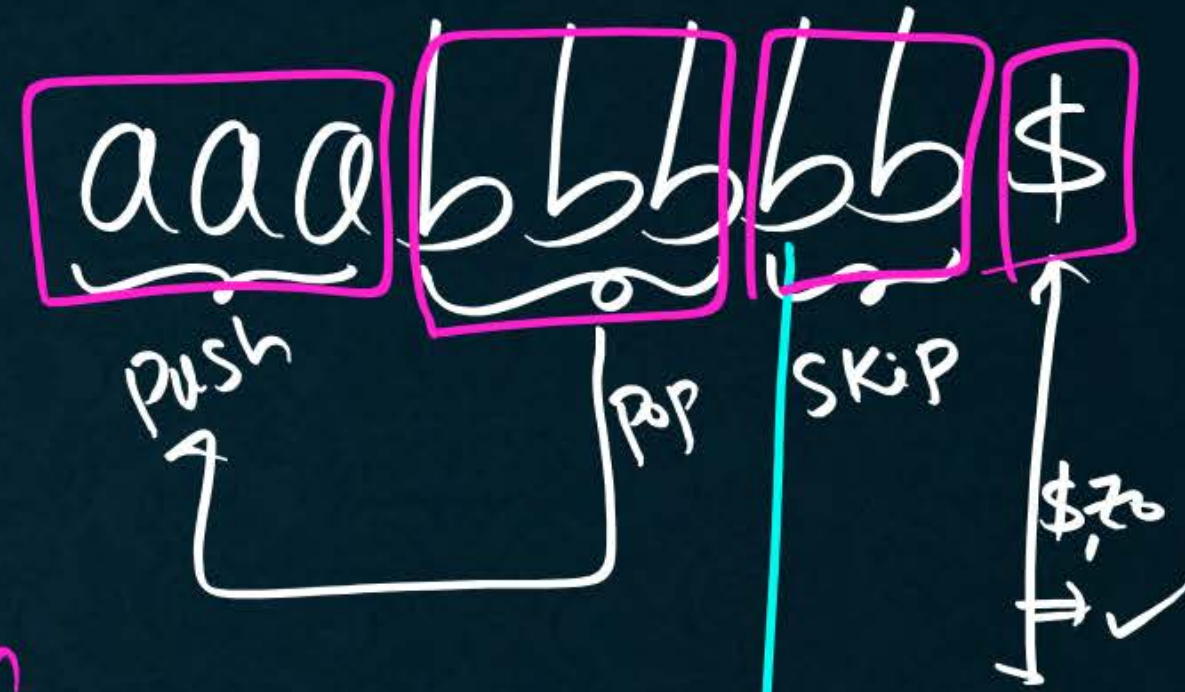
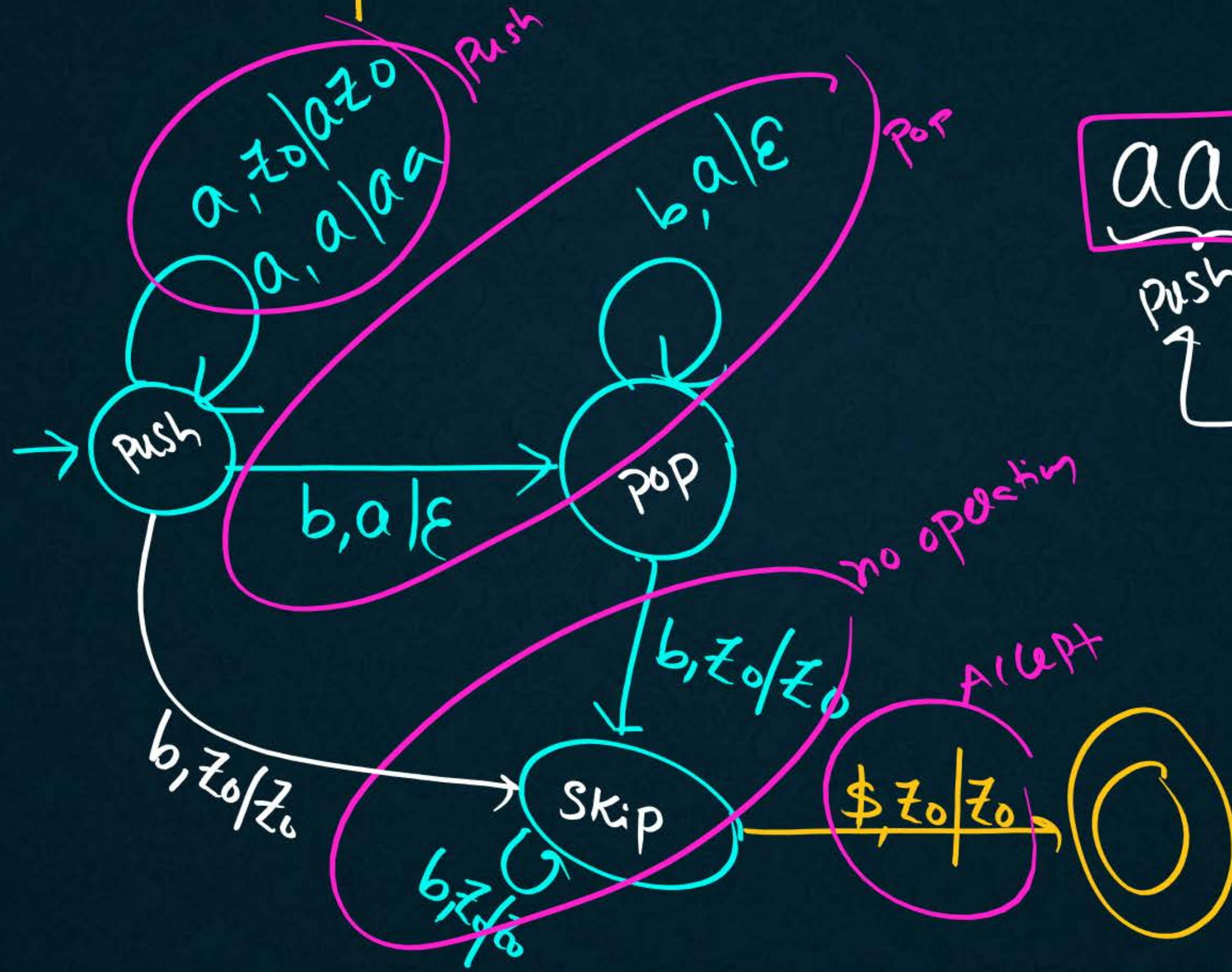
$$\{a^m b^n \mid m < n, \quad m, n \geq 1\}$$


⑧ $\{a^m b^n \mid m \leq n, m, n \geq 1\}$



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$\{a^m b^n \mid m < n, m, n \geq 0\} = \emptyset \cup \{b^+\}$



Design
DPQ

⑩

$$\{a^m b^n \mid m, n \geq 1, \underbrace{m \neq n}_{m < n \text{ or } m > n}\}$$

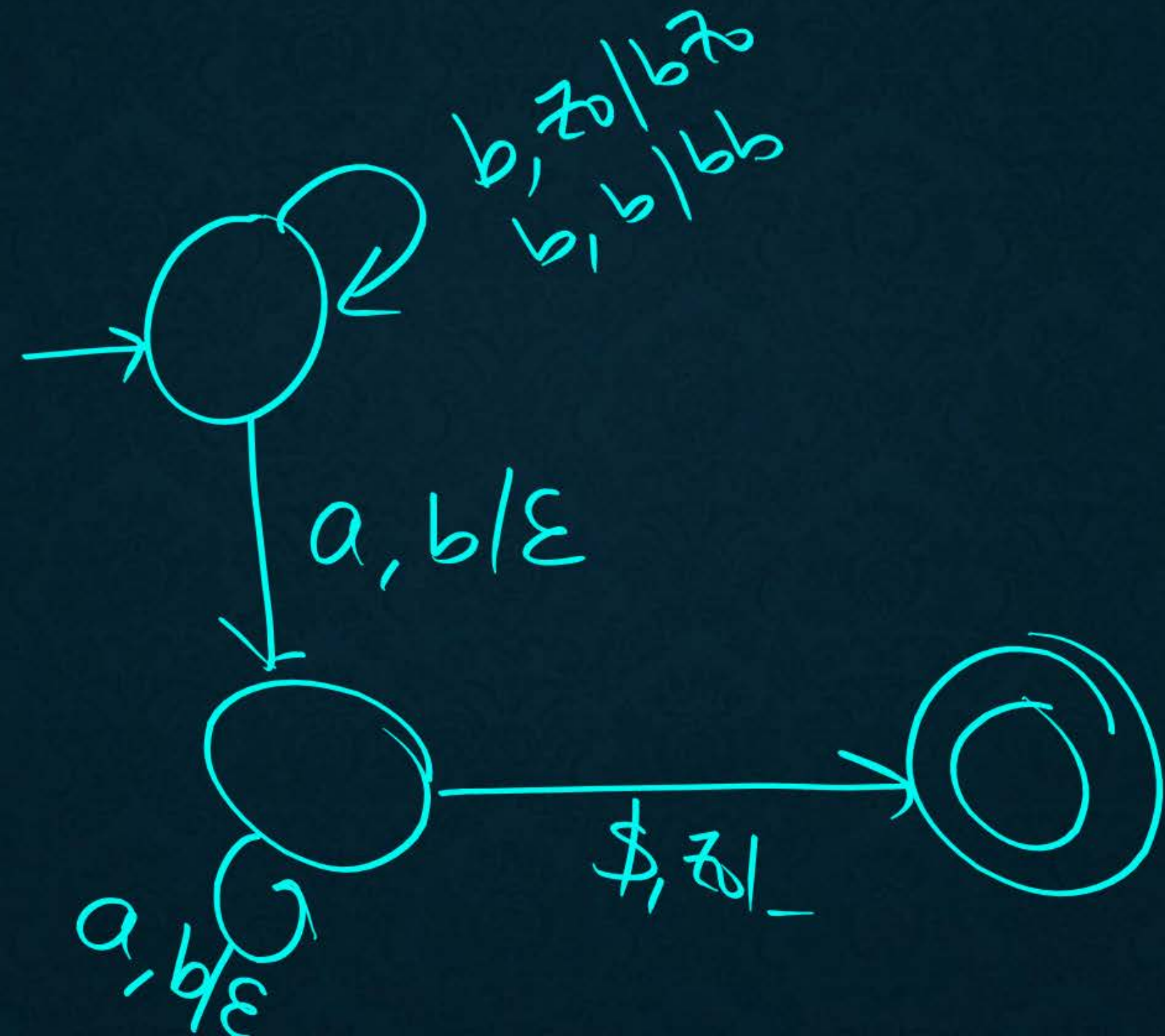
⑪

$$\{a^m b^n \mid m, n \geq 0, m \neq n\}$$

~~Xi~~

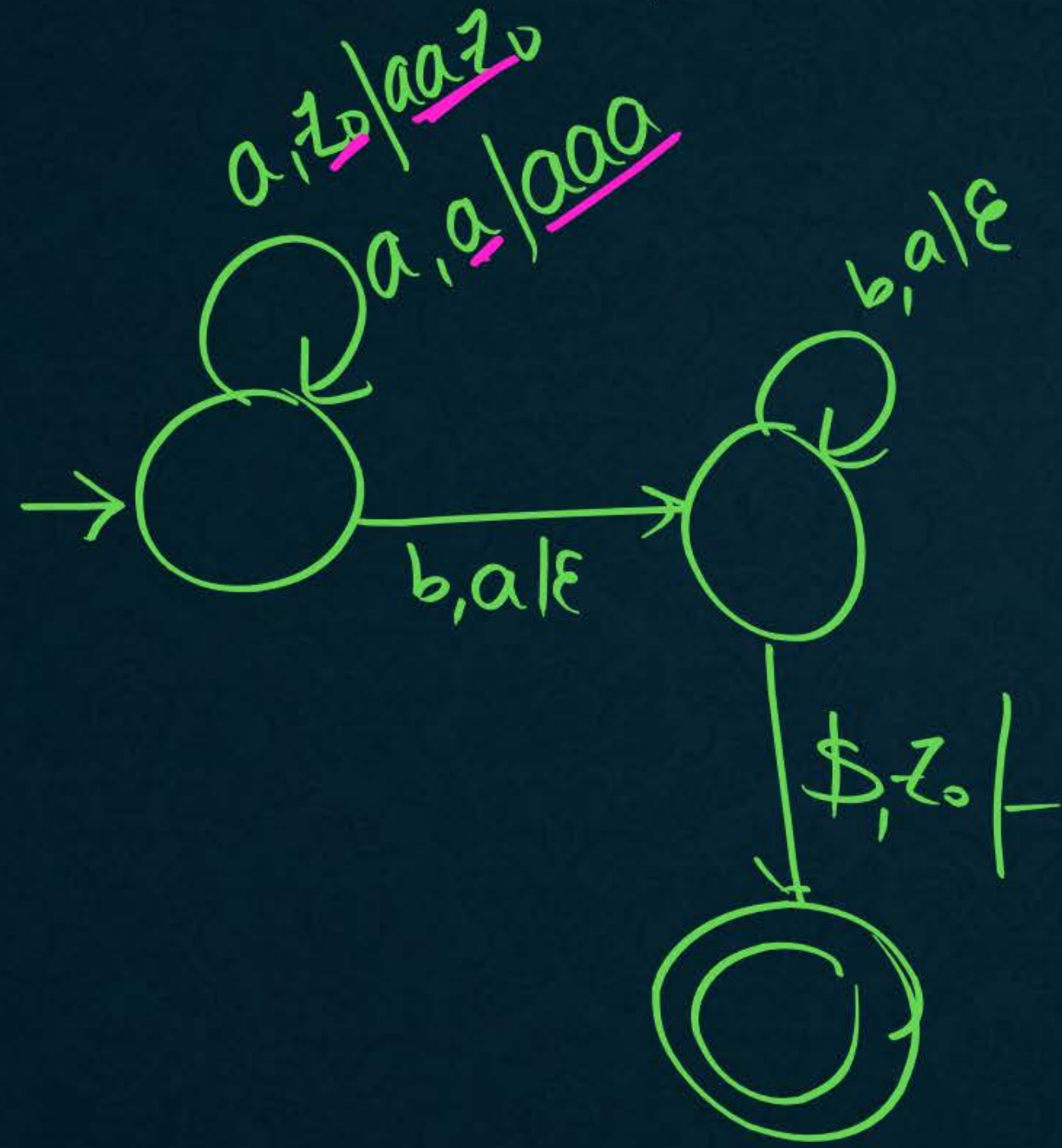
(12) $\{b^n a^n \mid n \geq 1\}$

bbbaaa
push pop

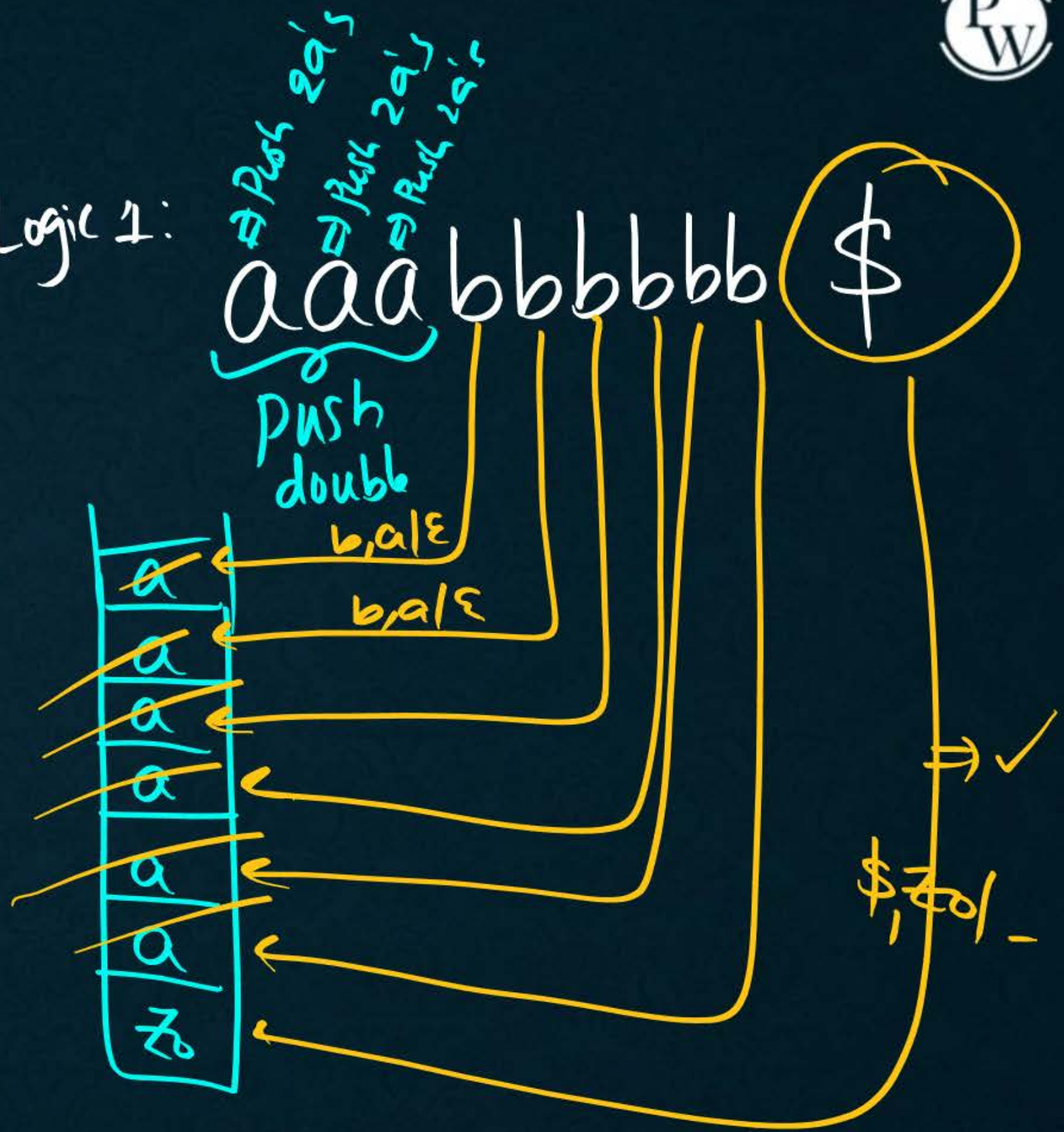


13

$$\{a^n b^n \mid n \geq 1\}$$

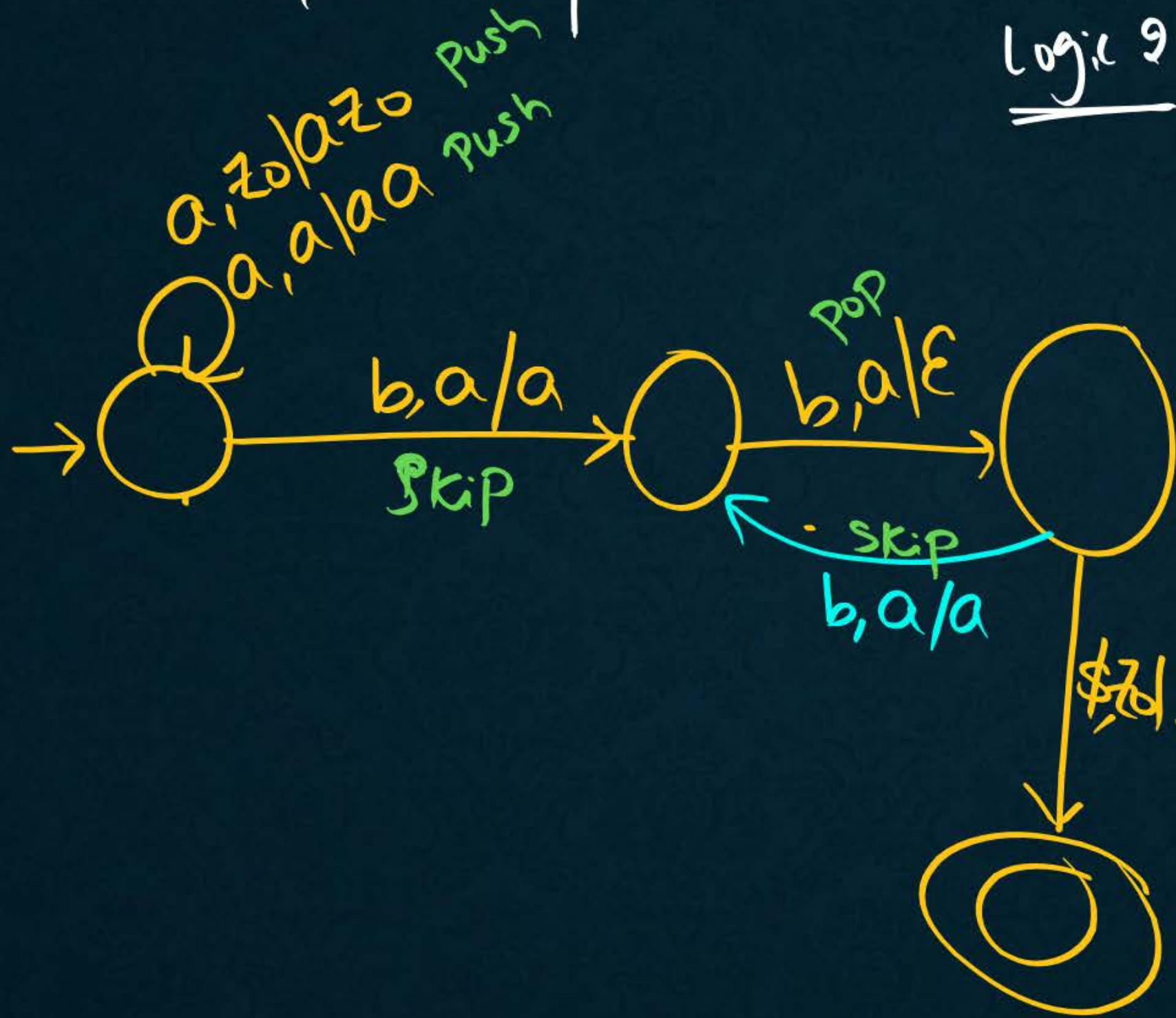


Logic 1:

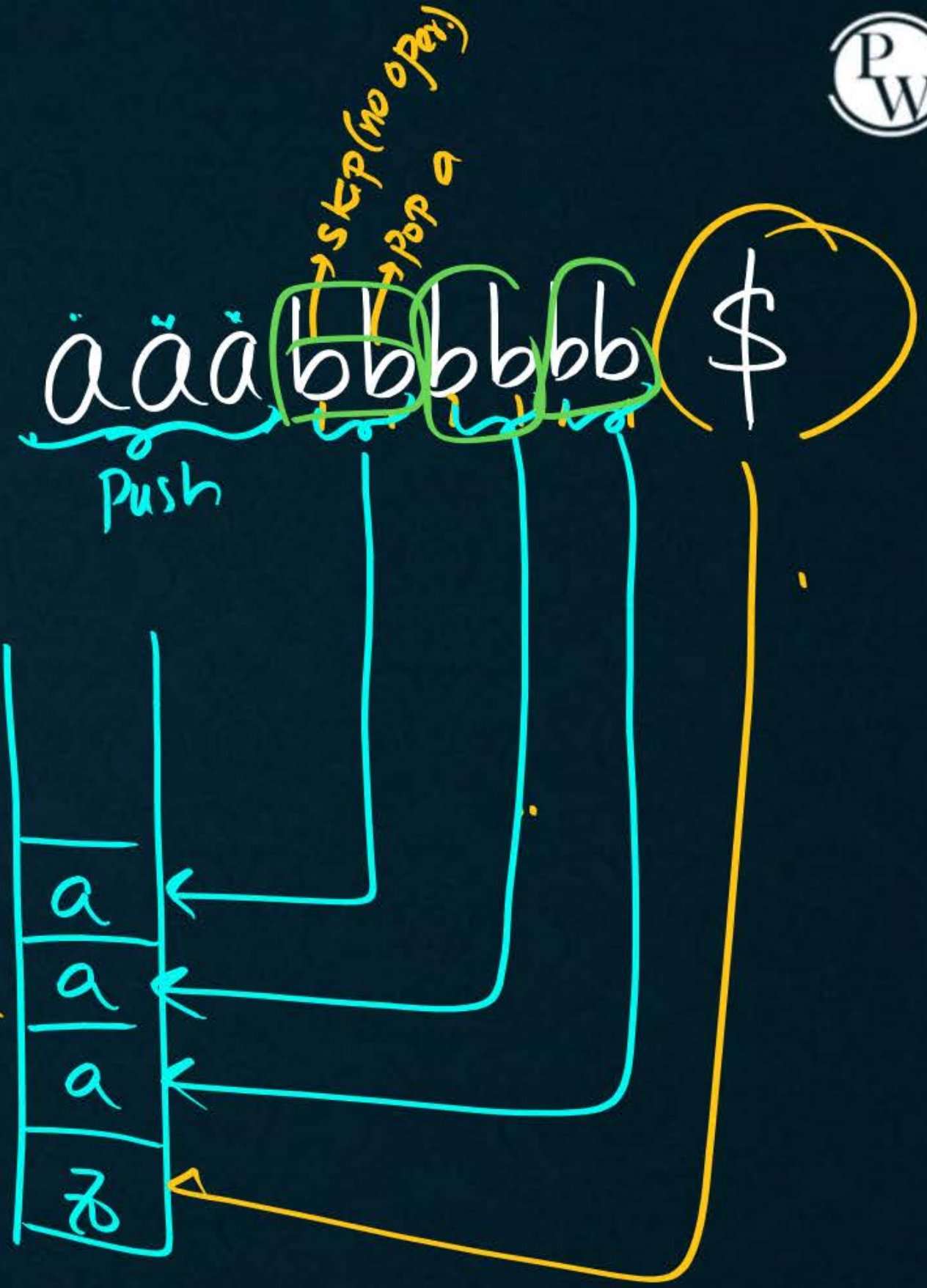


13

$$\{ a^n b^n \mid n \geq 1 \}$$



Logic 2:



H.W.

(14) $\{a^{2n} b^n \mid n \geq 1\}$

(15) $\{a^n b^{3n} \mid n \geq 1\}$

(16) $\{a^{2n} b^{2n} \mid n \geq 1\}$

↑
↑
equal
&
even

(17) $\{w \mid w \in \{a, b\}^*, n_a(w) = n_b(w)\}$

(18) $\{w \mid \text{"}, n_a(w) > n_b(w)\}$

(19) $\{w \mid \text{"}, n_a(w) < n_b(w)\}$

(20) $\{w \mid \text{"}, n_a(w) \leq n_b(w)\}$

(21) $\{w \mid \text{"}, n_a(w) \geq n_b(w)\}$

(22) $\{w \mid \text{"}, n_a(w) \neq n_b(w)\}$

→ DPDA & PDA

