Measuring hardware performance counters on ARCHER2 using LIKWID

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LIKWID overview



Command line performance tools:

- Probe hardware topology
 likwid-topology print thread, cache and NUMA topology
- Process & thread placement control
 likwid-pin pin application threads
 likwid-mpirun pin application processes
- Hardware performance monitoring
 likwid-perfctr measure hardware performance counters
 likwid-mpirun likwid-perfctr for MPI applications
- Microbenchmarking
 - likwid-bench evaluate upper performance bounds using assembly microkernels

LIKWID CLI applications User applications Python API Lua API Marker API Lua RT LIKWID core C API and GPU API* LIKWID suid daemon | perf_event | Hwloc | Pinning lib CUDA*

Linux OS Kernel

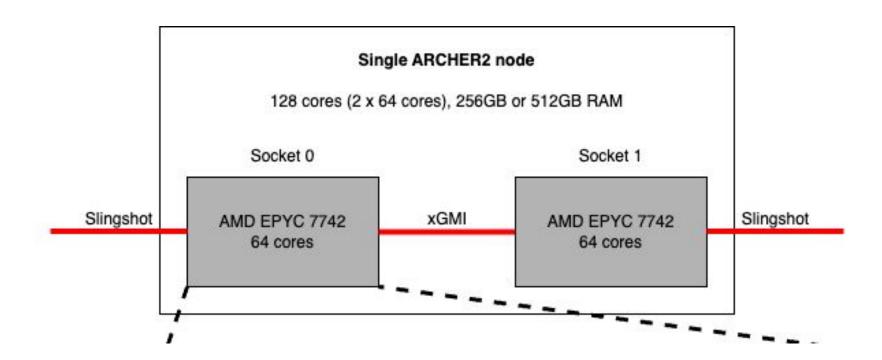
LIKWID 5 Tools Architecture

Nvidia GPUs



ARCHER2 Node





likwid-topology - hardware threads



CPU name: AMD EPYC 7742 64-Core Processor CPU type: AMD K17 (Zen2) architecture

CPU stepping: 0

Hardware Thread Topology

Sockets: 2

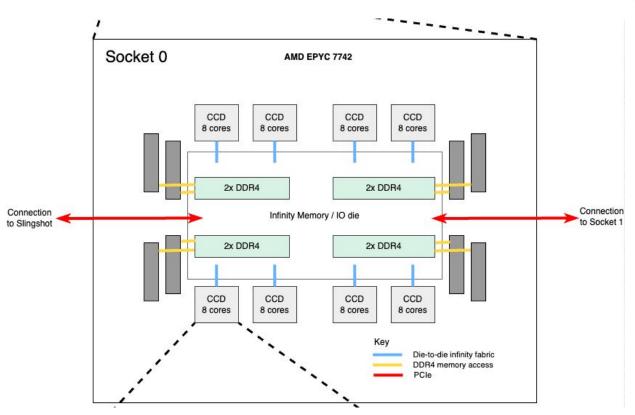
CPU dies: 2

Cores per socket:64 Threads per core:2

HWThread	Thread	Core	Die	Socket	Available
0	0	0	0	0	*
1	0	1	0	0	*
2	0	2	0	0	*
•					
•					
126	0	126	0	1	*
127	0	127	0	1	*
128	1	0	0	0	*
129	1	1	0	0	*
130	1	2	0	0	*
•					
254	1	126	0	1	*
255	1	127	0	1	*

ARCHER2 processor





likwid-topology - NUMA

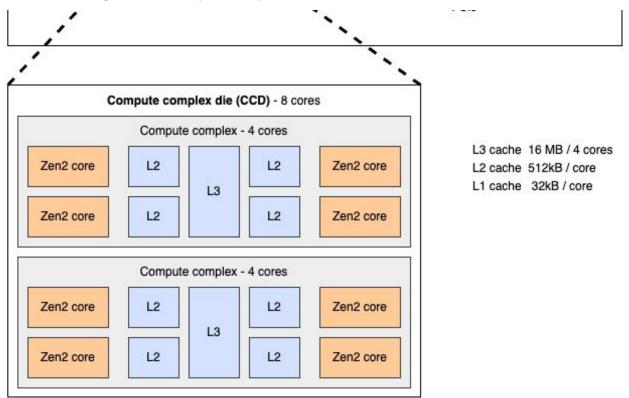


******************************* NUMA Topology **************************** NUMA domains: Domain: (0 128 1 129 2 130 3 131 4 132 5 133 6 134 7 135 8 136 9 137 10 138 11 139 12 140 13 141 14 142 15 143) Processors: Distances: 10 12 12 12 32 32 32 32 Free memory: 29215.3 MB Total memory: 63749.9 MB Domain: (16 144 17 145 18 146 19 147 20 148 21 149 22 150 23 151 24 152 25 153 26 154 27 155 28 156 29 157 30 158 31 159) Processors: Distances: 12 10 12 12 32 32 32 32 Free memory: 27638.5 MB 64504.1 MB Total memory: Domain: (96 224 97 225 98 226 99 227 100 228 101 229 102 230 103 231 104 232 105 233 106 234 107 235 108 236 109 237 110 238 Processors: 111 239) Distances: 32 32 32 32 12 12 10 12 29890.9 MB Free memory: 64504.1 MB Total memory: Domain: Processors: (112 240 113 241 114 242 115 243 116 244 117 245 118 246 119 247 120 248 121 249 122 250 123 251 124 252 125 253 126 254 127 255) Distances: 32 32 32 32 12 12 12 10 Free memory: 32468.4 MB Total memory: 64494.1 MB

9

ARCHER2 Compute Complex Die (CCD) = 2 Core Complexes (CCX)





likwid-topology - cache



```
*****************************
Cache Topology
****************************
Level:
                      1
Size:
                      32 kB
Type:
                      Data cache
Associativity:
Number of sets:
                      64
Cache line size: 64
                      Non Inclusive
Cache type:
Shared by threads:
Cache groups:
                      ( 0 128 ) ( 1 129 ) ... ( 126 254 ) ( 127 255 )
Level:
Size:
                      512 kB
                      Unified cache
Type:
Associativity:
Number of sets:
                      1924
Cache line size: 64
                      Non Inclusive
Cache type:
Shared by threads:
                      ( 0 128 ) ( 1 129 ) ... ( 126 254 ) ( 127 255 )
Cache groups:
Level:
                      16 MB
Size:
Type:
                      Unified cache
Associativity:
                      16
Number of sets:
                      16384
Cache line size: 64
Cache type:
                      Non Inclusive
Shared by threads:
                      ( 0 128 1 129 2 130 3 131 ) ( 4 132 5 133 6 134 7 135 ) ... ( 120 248 121 249 122 250 123 251 ) ( 124 252 125 253 126
Cache groups:
254 127 255 )
```



likwid-pin - thread pinning



Pin threads to cores based on:

- 1. OS: numbering according to the OS likwid-pin -c 0,2,4,6
- 2. Node: logical numbering over whole node (N:<list>)
 likwid-pin -c N:0-3
- 3. Socket: logical numbering within each socket (S#:<list>)
 likwid-pin -c S0:0-3@S1:0-3
- 4. Die: logical numbering within each die (D#:likwid-pin -c D0:0-63
- 5. Cache group: logical numbering within each last level cache groups (C#:<list>) likwid-pin -c C0:0@C1:0@C2:0@C3:0
- 6. Memory domain: logical numbering within each NUMA domain (M#:<list>)
- 7. Cpuset: logical numbering inside each Linux cpuset (**L prefix**)

Various list syntax options

likwid-pin - thread domains on ARCHER2



```
> likwid-pin -p
Domain N: 0,128,1,129 ... 126,254,127,255

Domain S0: 0,128,1,129 ... 62,190,63,191
Domain S1: 64,192,65,193 ... 126,254,127,255

Domain D0: 0,128,1,129 ... 62,190,63,191
Domain D1: 64,192,65,193 ... 126,254,127,255

Domain C0: 0,128,1,129,2,130,3,131
...
Domain C31: 124,252,125,253,126,254,127,255

Domain M0: 0,128,1,129 ... 14,142,15,143
...
Domain M7: 112,240,113,241 ... 126,254,127,255
```

Note: srun --cpus-per-task=128 --hint=nomultithread likwid-pin -p would exclude 128-255



AMD Zen2 event counters



Core-local counters

3 fixed-purpose:

Counter name	Event name
FIXC0	INST_RETIRED_ANY
FIXC1	ACTUAL_CPU_CLOCK
FIXC2	MAX_CPU_CLOCK

6 general-purpose: PMC0-PMC5, programmable event types

Socket-wide counters

Energy counters:

Counter name	Event name
PWR0	RAPL_CORE_ENERGY
PWR1	RAPL_PKG_ENERGY

Derived metrics and performance groups



Metrics of interest are derived from architecture-specific counters according to LIKWID-provided recipes

Grouped into "performance groups" (event sets & derived metrics):

Event set	Function					
FLOPS_DP	Double Precision MFlops/s					
FLOPS_SP	Single Precision MFlops/s					
L2 L2 cache bandwidth in MBytes						
L3	L3 cache bandwidth in MBytes/s					
MEM	Main memory bandwidth in MBytes/s					
CACHE L1 Data cache miss rate/ratio						
L2CACHE	L2 Data cache miss rate/ratio					
L3CACHE	L3 Data cache miss rate/ratio					
DATA	Load to store ratio					
BRANCH	Branch prediction miss rate/ratio					
TLB	Translation lookaside buffer miss rate/ratio					

AMD Zen2 FLOPS_DP eventset & metrics



SHORT Double Precision MFLOP/s

```
EVENTSET
FIXC1 ACTUAL CPU CLOCK
FIXC2 MAX CPU CLOCK
     RETIRED_INSTRUCTIONS
PMC1
     CPU_CLOCKS_UNHALTED
PMC2
     RETIRED_SSE_AVX_FLOPS_ALL
PMC3 MERGE
METRICS
Runtime (RDTSC) [s] time
Runtime unhalted [s] FIXC1*inverseClock
Clock [MHz] 1.E-06*(FIXC1/FIXC2)/inverseClock
    PMC1/PMC0
CPI
DP [MFLOP/s] 1.0E-06*(PMC2)/time
LONG
Formulas:
CPI = CPU_CLOCKS_UNHALTED/RETIRED_INSTRUCTIONS
DP [MFLOP/s] = 1.0E-06*(RETIRED_SSE_AVX_FLOPS_ALL)/time
```

likwid-perfctr - counter measurement



- Wrapper mode (default):
 - wrap application launch, measure performance counters during execution

 - typically measure on all cores where your application executes
 can instrument using Marker API to measure region(s) kernel(s) only
- Stethoscope mode
 - "listen to" whatever is running on any specified cores for specified duration
 output: aggregate statistics over specified cores and duration

 - difficult to relate to what application code is being executed
- Timeline mode
 - periodically output aggregate statistics over specified cores & time window
 - can gain insight into application phases but still difficult to correlate with exactly what application code is being executed

No code changes / recompilation needed except for Marker API

likwid-mpirun - likwid-perfctr with MPI

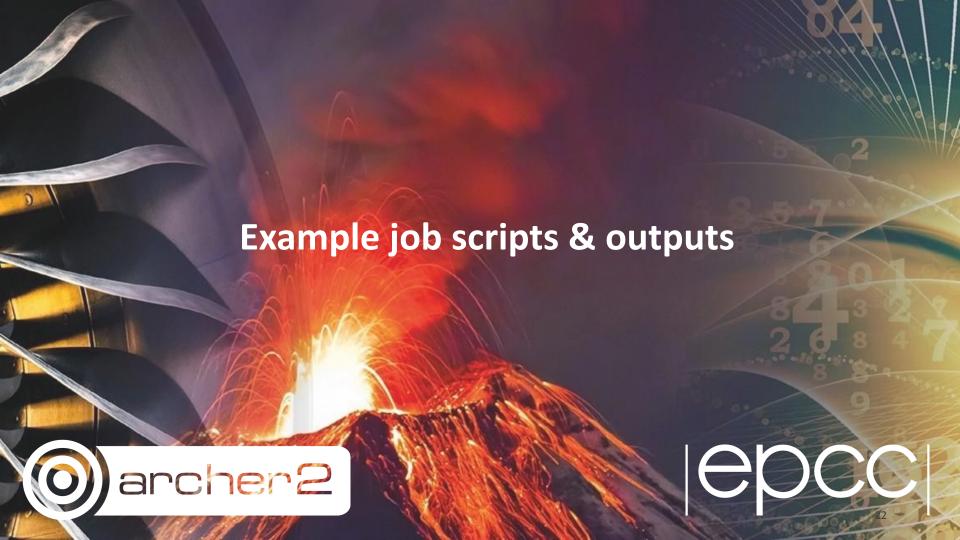


- ·likwid-perfctr supports threaded applications but not MPI
 - could launch likwid-perfctr using srun (see LIKWID MPI Tutorial)
 - statistics not aggregated over processes (could be useful for load imbalance)
- ·likwid-mpirun supports MPI-parallel (incl. MPI + threaded) apps:
 - wrapper mode only (incl. marker API)
 - generates srun commands to launch likwid-perfctr
 - aggregates statistics over processes and threads

likwid-mpirun options



- **-n <count>**: total number of processes
- **-t <count>**: number of threads per process
- -pin <list>: pinning of processes (and their threads if relevant)
 - (mostly) follows likwid-pin syntax
- -g/--group <perf>: performance group (e.g. FLOPS DP, MEM, etc.)
- **--nocpubind**: disable process binding through (ARCHER2 custom option, central LIKWID install only)
 - recommended for convenience control process binding as usual on ARCHÉR2
- -s/--skip <hex>: skip shepherd threads
 - always use -s 0x0 on ARCHER2
- -d/--debug: more information, e.g. on generated srun command
 very useful to ensure desired application placement and measurement
- **--mpiopts**: pass any desired arguments to generated srun commands
- **-nperdomain**: only useful for pure MPI, can accomplish similar using -pin
- -m: activate marker API mode



Example job scriptPure MPI, 2 fully populated nodes (2 x 128 processes)



```
#!/bin/bash
#SBATCH --account=z19
#SBATCH --partition=standard
#SBATCH --qos=short
#SBATCH --time=00:05:00
#SBATCH --nodes=2
#SBATCH --tasks-per-node=128
#SBATCH --tasks-per-task=1
#SBATCH --hint=nomultithread
#SBATCH --distribution=block:block
module load likwid
module load xthi

export OMP_NUM_THREADS=1
export SRUN_CPUS_PER_TASK=$SLURM_CPUS_PER_TASK
likwid-mpirun -n $SLURM_NTASKS --nocpubind -s 0x0 -g FLOPS_DP xthi_mpi &> xthi_mpi.out
```

Example job script Pure MPI, 2 nodes, 2 ranks per node, 1 per socket (processor)



```
#!/bin/bash
#SBATCH --account=z19
#SBATCH --partition=standard
#SBATCH --gos=short
#SBATCH --time=00:05:00
#SBATCH --nodes=2
#SBATCH --tasks-per-node=2
#SBATCH --cpus-per-task=64
#SBATCH --hint=nomultithread
#SBATCH --distribution=block:block
module load likwid
module load xthi
export OMP NUM THREADS=1
export SRUN CPUS PER TASK=SSLURM CPUS PER TASK
likwid-mpirun -n $SLURM_NTASKS -pin N:0_N:64 --nocpubind -s 0x0 -q FLOPS_DP --debug xthi_mpi &> xthi_mpi.out
# Alternatively:
likwid-mpirun -n $SLURM_NTASKS -pin S0:0_S1:0 --nocpubind -s 0x0 -g FLOPS_DP --debug xthi_mpi &> xthi_mpi.out
```

Application (xthi) output Pure MPI, 2 nodes, 2 ranks per node, 1 per socket (processor)



```
Node summary for
                  2 nodes:
Node
       0, hostname nid001068, mpi 2, omp
                                           1, executable xthi_mpi
                                           1, executable xthi_mpi
Node
       1, hostname nid001079, mpi
                                  2, omp
MPI summary: 4 ranks
       0, rank
                 0, thread 0, (affinity =
                                              0)
Node
       0, rank 1, thread 0, (affinity =
Node
                                             64)
Node
       1, rank 2, thread 0, (affinity =
                                              0)
Node
       1, rank
                 3, thread
                             0, (affinity =
                                             64)
```

likwid-mpirun results for application (cp2k H2O-32) Pure MPI, 2 nodes, 2 ranks per node, 1 per socket (processor)



Event	Counter	nid001068:0:0	nid0	01068:1:64	nid	1001079:2:0	nid	001079:3:64
ACTUAL_CPU_CLOCK	FIXC1	272843810277	2	73383588225	2	73224152138	 	273645884841
MAX_CPU_CLOCK	FIXC2	307238142692	3	07671324996	3	07777646157		308110758689
RETIRED_INSTRUCTIONS	PMC0	639095695379	6	12418860306	6	18138505835	1	624705982405
CPU_CLOCKS_UNHALTED	PMC1	4325299091000	43	34074578504	43	29865216924	4	331611334344
RETIRED_SSE_AVX_FLOPS_ALL	. PMC2	379776213171	3	45443299984	3	62489794174		378163333211
MERGE	PMC3	4644337115725824	46443 -+	37115725824	+	174604081152 	+	337115725824
 Event	+ Count	+	 	 Min		 Max		+
	+	+	ا ++		ı +			+
ACTUAL_CPU_CLOCK STA	T FIXO	:1 10930974	435481 I	2728438	10277 I	2736458	84841	2.732744e+
MAX_CPU_CLOCK STAT	FIXC	•						
RETIRED_INSTRUCTIONS S	TAT PMC	0 24943590	943925	6124188	60306	6390956	95379	6.235898e+
CPU_CLOCKS_UNHALTED ST	AT PMC	1 173208502	220772	43252990	91000 j	43340745	78504	43302125551
RETIRED_SSE_AVX_FLOPS_ALL	STAT PMC	2 14658726	640540	3454432	99984	3797762	13171	3664681601
MERGE STAT	PMC	3 187180859512	258624	46443371157	25824	47850746046	81152	46795214878146 +
+	+		+	+		+		
Metric ni	d001068:0:0 +	nid001068:1:64	nia001 +	0/9:2:0 ni		1:3:64		
Runtime (RDTSC) [s]	148.7140	148.7171	1	47.6779	147	.6825		
Runtime unhalted [s]	121.4966	121.7369	1	21.6676	121	.8564		
Clock [MHz]	1994.2933	1995.4260	19	93.5457	1994	.4471		
CPI	6.7678	7.0770	•	7.0047		.9338		
DP [MFLOP/s] +	2553.7354	2322.8217	24 +	54.5974 +	2560	0.6509 +		
Metric	-+ Sum	+	Max	+ Avg	+ %ile	25 %ile	+ 50	%ile 75
Runtime (RDTSC) [s] STAT	-+ 592.7915	147.6779 14	 48.7171	+ 148.1979	+	6779 147.	6825	148.7140
Runtime unhalted [s] STAT	486.7575	121.4966 12	21.8564	121.6894	121.	4966 121.	6676	121.7369
Clock [MHz] STAT	7977.7121	1993.5457 199	95.4260	1994.4280	1993.	5457 1994.	2933	1994.4471
CPI STAT	27.7833	6.7678	7.0770	6.9458	6.	7678 6.	9338	7.0047
DP [MFLOP/s] STAT		2322.8217 256						

Debug outputPure MPI, 2 nodes, 2 ranks per node, 1 per socket (processor)



```
EXEC (Rank 0): likwid-perfctr -s 0x0 -C L:N:0-0 -g
ACTUAL_CPU_CLOCK:FIXC1,MAX_CPU_CLOCK:FIXC2,RETIRED_INSTRUCTIONS:PMC0,CPU_CLOCKS_UNHALTED:PMC1,RETIRED_SSE_AVX_FLO
PS_ALL:PMC2,MERGE:PMC3 -o temporary1.csv cp2k

EXEC (Rank 1): likwid-perfctr -s 0x0 -C L:N:0-0 -g
ACTUAL_CPU_CLOCK:FIXC1,MAX_CPU_CLOCK:FIXC2,RETIRED_INSTRUCTIONS:PMC0,CPU_CLOCKS_UNHALTED:PMC1,RETIRED_SSE_AVX_FLO
PS_ALL:PMC2,MERGE:PMC3 -o temporary2.csv cp2k

EXEC: srun --ntasks-per-node=2 --ntasks=4 --mpi=cray_shasta --nodes=2 .likwidscript_122134.txt
```

.likwidscript_122134.txt (bash script):

myNodeLocalRank = 0 or 1 on each node

if myNodeLocalRank == 0 then EXEC (Rank 0) if myNodeLocalRank == 1 then EXEC (Rank 1)

Example job scriptPure OpenMP, fully populated node (128 threads)



```
#!/bin/bash
#SBATCH --account=z19
#SBATCH --partition=standard
#SBATCH --gos=short
#SBATCH --time=00:05:00
#SBATCH --nodes=1
#SBATCH --tasks-per-node=1
#SBATCH --cpus-per-task=128
#SBATCH --hint=nomultithread
#SBATCH --distribution=block:block
module load likwid
module load xthi
export OMP_NUM_THREADS=128
export OMP PLACES=cores
export SRUN_CPUS_PER_TASK=$SLURM_CPUS_PER_TASK
likwid-mpirun -n 1 -t 128 --nocpubind -s 0x0 -q FLOPS_DP --debug xthi &> xthi.out
```

Debug outputPure OpenMP, fully populated node (128 threads)



```
EXEC (Rank 0): likwid-perfctr -s 0x0 -C L:N:0-127 -g

ACTUAL_CPU_CLOCK:FIXC1, MAX_CPU_CLOCK:FIXC2, RETIRED_INSTRUCTIONS:PMC0, CPU_CLOCKS_UNHALTED:PMC1, RETIRED_SSE_AVX_FL0

PS_ALL:PMC2, MERGE:PMC3 -o temporary1.csv xthi

EXEC: srun --ntasks=1 --mpi=cray_shasta --ntasks-per-node=1 --nodes=1 .likwidscript_231548.txt
```

Don't need likwid-mpirun!

Example job script using likwid-perfctr Pure OpenMP, fully populated node (128 threads)



```
#!/bin/bash

#SBATCH --account=z19
#SBATCH --partition=standard
#SBATCH --qos=short
#SBATCH --time=00:05:00
#SBATCH --nodes=1

module load likwid
module load xthi

export OMP_NUM_THREADS=128
export OMP_PLACES=cores

likwid-perfctr -C N:0-127 -s 0x0 -g FLOPS_DP xthi &> xthi.out

# Alternatively:
likwid-perfctr -C E:N:128:1:2 -s 0x0 -g FLOPS_DP xthi &> xthi.out
```

Example job script Pure OpenMP, 4 threads, 1 per CCX



```
#!/bin/bash
#SBATCH --account=z19
#SBATCH --partition=standard
#SBATCH --gos=short
#SBATCH --time=00:05:00
#SBATCH --nodes=1
module load likwid
module load xthi
export OMP_NUM_THREADS=4
export OMP_PLACES=cores
likwid-perfctr -C 0,4,8,12 -s 0x0 -g FLOPS_DP xthi &> xthi.out
# Alternatively:
likwid-perfctr -C C0:0@C1:0@C2:0@C3:0 -s 0x0 -q FLOPS_DP xthi &> xthi.out
# Alternatively:
likwid-perfctr -C E:N:4:1:8 -s 0x0 -g FLOPS_DP xthi &> xthi.out
```

Application (xthi) output Pure OpenMP, 4 threads, 1 per CCX

```
epcc
```

```
Node summary for
                  1 nodes:
       0, hostname nid001380, mpi 1, omp 4, executable xthi
Node
MPI summary: 1 ranks
       0, rank 0, thread 0, (affinity =
Node
                                             0)
Node
       0, rank
                 0, thread 1, (affinity =
                                             4)
       0, rank 0, thread 2, (affinity =
                                             8)
Node
                 0, thread 3, (affinity =
Node
       0, rank
                                            12)
```

likwid-perfctr results for application (cp2k H2O-32) Pure OpenMP, 4 threads, 1 per CCX



	Event	Counter	HWThread 0	HWThread 4	HWThread 8	HWThread 12
i	ACTUAL_CPU_CLOCK	FIXC1	288225142494	277342582674	276504761957	275708683034
	MAX_CPU_CLOCK	FIXC2	324520870098	312204162837	311334613420	310553053741
	RETIRED_INSTRUCTIONS	PMC0	642746850133	597585919597	590879312097	591524771907
	CPU_CLOCKS_UNHALTED	PMC1	4579460146688	4429904890720	4409955410336	4404089955280
	RETIRED_SSE_AVX_FLOPS_ALL	PMC2	380268064831	368436346965	360781276795	359314236216
1	MERGE	PMC3	5348024557502464	5348024557502464	5207287069147136	5207287069147136

+- +-	Event	İ	Counter	Ī	Sum	+ Min	+- +-	Max	+- +-	Avg	+
+-	ACTUAL_CPU_CLOCK STAT MAX_CPU_CLOCK STAT RETIRED_INSTRUCTIONS STAT CPU_CLOCKS_UNHALTED STAT RETIRED_SSE_AVX_FLOPS_ALL STAT	 	FIXC1 FIXC2 PMC0 PMC1 PMC2	 	1117781170159 1258612700096 2422736853734 17823410403024 1468799924807	275708683034 275708683034 310553053741 590879312097 4404089955280 359314236216	+ · · · · · · · · · · · · · · · · · · ·	288225142494 324520870098 642746850133 4579460146688 380268064831	+- 	2.794453e+11 314653175024 6.056842e+11 4455852600756 3.672000e+11	i
į	MERGE STAT	į	PMC3	į	21110623253299200	5207287069147136	į	5348024557502464	į	5277655813324800	į

	Metric	HWThre	ad 0 HW	Γhread 4 H	WThread 8	HWThread 12
1111+	Runtime (RDTSC) [s] Runtime unhalted [s] Clock [MHz] CPI DP [MFLOP/s]	158. 128. 1994. 1994. 2396.	7059 3473 5 5010 19 1248 0556 23		158.7059 123.1282 1994.4366 7.4634 2273.2701	158.7059 122.7737 1993.6993 7.4453 2264.0263
+	Metric	1	Sum	Min	Max	Avg
i	Runtime (RDTSC) [s] Runtime unhalted [s] Clock [MHz] STAT CPI STAT DP [MFLOP/s] STAT	STAT 6 STAT 6 7	634.8236 497.7505 977.5450 29.4465	158.7059 122.7737 1993.6993 7.1248 2264.0263	158.7059 128.3473 1994.9081 7.4634 2396.0556	158.7059 124.4376 1994.3863 7.3616

Example job script

Hybrid MPI + OpenMP, 1 node, 2 ranks per node, 64 threads per rank



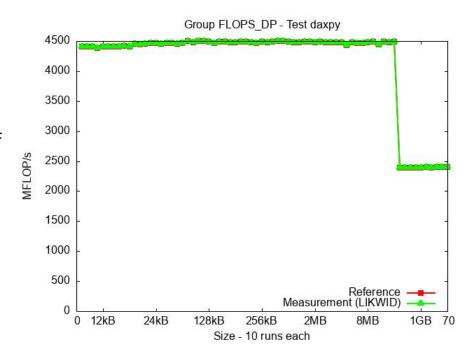
```
#!/bin/bash
#SBATCH --job-name=likwid
#SBATCH --account=z19
#SBATCH --partition=standard
#SBATCH -- gos=short
#SBATCH --time=00:05:00
#SBATCH --nodes=1
#SBATCH --tasks-per-node=2
#SBATCH --cpus-per-task=64
#SBATCH --hint=nomultithread
#SBATCH --distribution=block:block
module load likwid
module load xthi
export SRUN_CPUS_PER_TASK=$SLURM_CPUS_PER_TASK
export OMP NUM THREADS=SSLURM CPUS PER TASK
export OMP PLACES=cores
likwid-mpirun -n $SLURM_NTASKS -t 64 --nocpubind -s 0x0 -q FLOPS_DP --debug xthi &> xthi.out
# Alternative:
likwid-mpirun -n $SLURM_NTASKS -pin N:0-63_N:64-127 --nocpubind -s 0x0 -q FLOPS_DP --debug xthi &> xthi.out
```



Validation through microbenchmarking



- likwid-bench provides a microbenchmark suite
- Tests for different performance groups
- Assembly microkernels with known number of instructions for validation of values measured with likwid-perfctr
- Provide estimates of performance upper bounds based on concrete hardware performance rather than theoretical, input to roofline analysis
- Can be useful for validating during prototyping of own assembly kernels





Using LIKWID MARKER API



- Main use case: monitor regions of interest at fine granularity
- Need to recompile with -DLIKWID PERFMON
- and with -I \$LIKWID DIR/include
- Link with -L \$LIKWID_DIR/lib -llikwid and adjust LD LIBRARY PATH
- Uses the LD_PRELOAD mechanism, so the application must be dynamically linked
- Run likwid-perfctr or likwid-mpirun with the -m flag
- Avoid calling very frequently: calls add overhead
- Can define multiple regions

Using LIKWID MARKER API



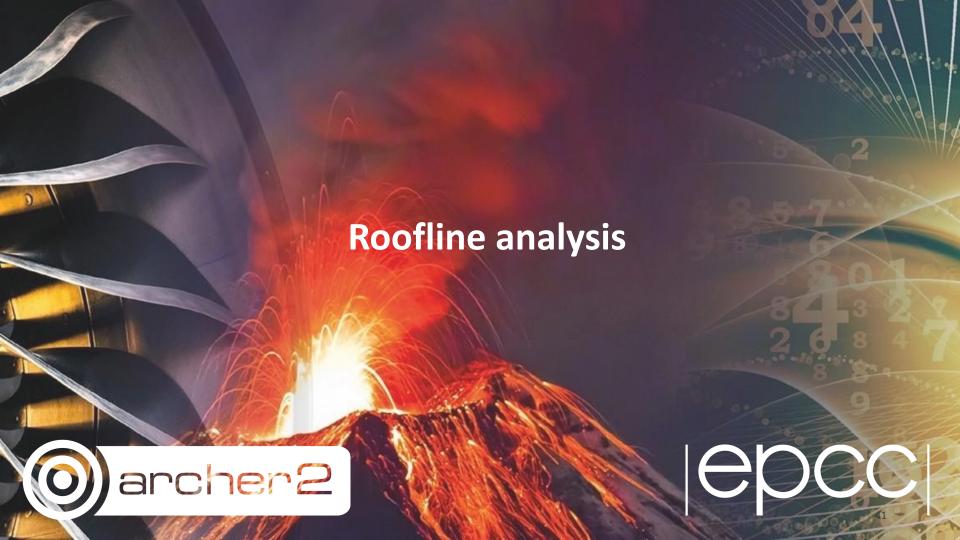
• Core C API (API also define for Fortran and a few other language)

```
... // other includes
#include <likwid-marker.h>
LIKWID MARKER INIT; // macro call to setup measurement system
LIKWID MARKER REGISTER ("myregion"); // to reduce overhead
// if OpenMP is used: called the above and start/stop
// in a parallel region
LIKWID MARKER START ("myregion");
... // code region of interest
LIKWID MARKER STOP ("myregion");
LIKWID MARKER CLOSE;
```

Serial hello world example output (w -g MEM) CC



CPU name: AMD EPYC 7742 64-Core Processor CPU type: AMD K17 (Zen2) architecture CPU clock: 2.25 GHz								
Hello world								
Region hello, Group 1: MEM								
Region Info		HW	HWThread 0		Ī			
RDTSC Runtime [s] call count		0.000013						
· · · · · · · · · · · · · · · · · · ·								
1	<u>Event</u>		Counter		HWThread 0			
ACTUAL_CPU_CLOCK MAX_CPU_CLOCK RETIRED_INSTRUCTIONS CPU_CLOCKS_UNHALTED DRAM_CHANNEL_0 DRAM_CHANNEL_1		FIXC1 FIXC2 PMC0 PMC1 DFC0 DFC1		114358 75082 7011 14666 447 524				
į	Metric			ŀ	HWThread 0			
	Runtime (RDTSC) [s] Runtime unhalted [s] Clock [MHz] CPI Memory bandwidth [MBytes/s] Memory data volume [GBytes]		1	.315996e-05 0.0001 3426.9943 2.0919 4722.2028 0.0001				



Roofline (Williams et al., 2009, CACM 52(4):65-76)



Roofline: An Insightful Visual Performance Model for Floating-Point Programs and Multicore Architectures*

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ABSTRACT

We propose an easy-to-understand, visual performance model that offers insights to programmers and architects on improving parallel software and hardware for floating point computations.

1. INTRODUCTION

Conventional wisdom in computer architecture led to homogeneous designs. Nearly every desktop and server computer uses caches, pipelining, superscalar instruction issue, and out-of-order execution. Although the instruction sets varied, the microprocessors were all from the same school of design.

The switch to multicore means that microprocessors will become more diverse, since there is no conventional wisdom yet for them. For example, some offer many simple processors versus fewer complex processors, some depend on multithreading, and some even replace caches with explicitly addressed local stores. Manufacturers will likely offer multiple products with differing number of cores to cover multiple price-performance points, since the cores per chip will likely double every two years [4].

While diversity may be understandable in this time of uncertainty, it exacerbates the already difficult job of programmers, compiler writers, and even architects. Hence, an easy-to-understand model that offers performance guidelines could be especially valuable.

A model need not be perfect, just insightful. For example, the 3Cs model for caches is an analogy [19]. It is not a perfect model,

limited by the serial portion of a parallel program. It has been recently applied to heterogeneous multicore computers [4][18].

3. THE ROOFLINE MODEL

We believe that for the recent past and foreseeable future, off-chip memory bandwidth will often be the constraining resource[23]. Hence, we want a model that relates processor performance to offchip memory traffic.

Towards that goal, we use the term operational intensity to mean operations per byte of DRAM traffic. We define total bytes accessed as those that go to the main memory after they have been filtered by the cache hierarchy. That is, we measure traffic between the caches and memory rather than between the processor and the caches. Thus, operational intensity suggests the DRAM bandwidth needed by a kernel on a particular computer.

We use operational intensity instead of the terms arithmetic intensity [16] or machine balance [8][11] for two reasons. First, arithmetic intensity and machine balance measure traffic between the processor and cache, whereas we want to measure traffic between the caches and DRAM. This subtle change allows us to include memory optimizations of a computer into our bound and bottleneck model. Second, we think the model will work with kernels where the operations are not arithmetic (see Section 7), so we needed a more general term than arithmetic.

The proposed model ties together floating-point performance,

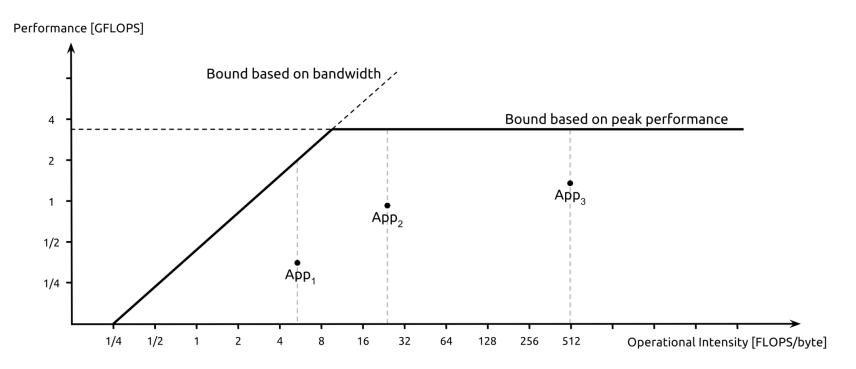
Roofline Analysis with LIKWID



- Roofline model is an intuitive visual performance model
- Performance is limited by peak attainable memory performance (e.g. obtained by running the STREAM benchmark) and by peak FLOPS which define the limiting Rooflines
- Y-axis: performance (FLOPS)
- X-axis operational intensity of the application (FLOPS/byte)
- Can plot applications, functions or loops as points
- Distance from the Roofline represents optimisation potential
- •There can be multiple Rooflines (e.g. for single vs double precision)

Roofline Analysis with LIKWID





[image from: https://en.wikipedia.org/wiki/Roofline model]

Using likwid-bench for Roofline analysis



- Roofline does only tell that a loop is memory or compute bound and shows roughly optimisation potential, but does not tell what those optimisations are
- Get memory Roofline by running a version of STREAM or estimate from number of DRAM channels and frequency (or use likwid-bench)
- •Get compute Roofline: can roughly estimate by using number of cores times number of operations (from CPU frequency)
- Measure your application (as further above) or region of interest (with Marker API)
- Select optimisations based on the resulting position of the point

Conclusions



- LIKWID is a powerful yet lightweight and portable tool suite with focus on performance engineering
- Allows to measure hardware performance counters on ARCHER2
- We encourage you to give it a go!
- To get started use the ARCHER2 documentation: <u>https://docs.archer2.ac.uk/data-tools/likwid/</u> and The official LIKWID Wiki for more in depth information: <u>https://github.com/RRZE-HPC/likwid/wiki</u>

Q: overheads?

Counters implemented in hardware so virtually no overheads, but possible overheads in (post-execution?) aggregation - not sure how scales to very large process counts

https://ieeexplore.ieee.org/document/7103452 (also includes comparison with PAPI)

Overhead from marker API might be significant if very frequent calls (entry/exit of instrumented regions)

Q: LIKWID vs PAPI?

https://arxiv.org/pdf/1004.4431

	LIKWID	PAPI
Dependencies	Needs system headers of Linux 2.6 kernel. No other external dependencies.	Needs kernel patches depending on platform and architecture. No patches necessary on Linux kernels > 2.6.31.
Installation	Build system based on make only. Install documentation 10 lines. Build configuration in a single text file (21 lines).	Install documentation is 582 lines (3.7.2) and 397 lines (4.0.0). The installation of PAPI for this comparison was not without problems.
Command line tools	Core is a collection of command line tools which are intended to be used standalone.	Collection of small utilities. These utilities are not supposed to be used as standalone tools. There are many PAPI-based tools available from other sources.
User API support	Simple API for configuring named code regions. API only turns counters on and off. Configuration of events and output of results is still based on the command line tool.	Comparatively high-level API. Events must be configured in the code.
Library support	While it can be used as library this was not initially intended.	Mature and well tested library API for building own tooling.
Topology information	Listing of thread and cache topology. Results are extracted from cpuid and presented in and accessible way as text and ASCII art. Nondata caches are omitted. No output of TLB information.	Information also based on cpuid. Utility outputs all caches (including TLBs). No output of shared cache information. Thread topology only as accumulated counts of HW threads and Cores. No mapping from processor Ids to thread topology.
Thread and process pinning	There is a dedicated tool for pinning processes and threads in a portable and simple manner. This tool is intended to be used together with likwid-perfCtr	No support for pinning.
Multicore support	Multiple cores can be measured simultaneously. Binding of threads or processes to correct cores is the responsibility of the user.	No explicit support for multicore measurements.
Uncore support	Uncore events are handled by applying socket locks, which prevent multiple measurements in threaded mode.	No explicit support for measuring shared resources.
Event abstraction	Preconfigured event sets (so-called event groups) with derived metrics.	Abstraction through papi_events, which map to native events.
Platform support	Supports only x86-based processors on Linux with 2.6 kernel.	Supports a wide range of architectures on various platforms (dedicated support for HPC systems like BlueGene or Cray XT3/4/5) with various operating systems (Linux, FreeBSD, and Windows).
Correlated measurements	LIKWID can measure performance counters only	PAPI-C can be extended to measure and correlate various data like, e.g., fan speeds or temperatures.

Table I: Comparison between LIKWID and PAPI