

Recurrence Relation

* Problem Solving Approach

Components of a Recurrence Relation:

- i) Base Case: What when we stop or know the answers to
- ii) Recursive Formula: What we continue doing

* Towers of Hanoi

For n disk	m moves
0	0
1	1
2	3
3	7
.	.
.	.
.	.
k	$2^k - 1$

- For $n=2$, \rightarrow When solving T_n , we know solⁿ of T_{n-1}
- S-1: Move ^{top}_{n-1} disks to middle (aux) towers (T_{n-1})
- S-2: Move largest disk to destination (final tower) (1)
- S-3) .. ^{top}_{mid} $n-1$ disk from aux to dest (T_{n-1})

$$\begin{aligned} \therefore T_n &= T_{n-1} + 1 + T_{n-1} \\ \therefore T_n &= 2T_{n-1} + 1 \quad \xrightarrow{\text{Open Form Solution}} \\ &\quad \hookrightarrow \text{Getting } n \text{ doesn't help} \\ &= 2(T_{n-2} + 1) + 1 \\ &= 2T_{n-2} + 2 + 1 \\ &= 2^2(2T_{n-3} + 1) + 2 + 1 \\ &= 2^3(2T_{n-4} + 1) + 4 + 2 + 1 \\ &\quad \vdots \\ &= 2^n(2T_0 + 1) + 2^{n-2} + \dots + 4 + 2 + 1 \\ &= 2^n(0 + 1) + 2^{n-2} + \dots + 4 + 2 + 1 \\ &= 2^{n-1} + 2^{n-2} + \dots + 4 + 2 + 1 \\ &= \frac{2^n - 1}{2 - 1} \\ \therefore T_n &= 2^n - 1 \end{aligned}$$

* Using Mathematical Induction:

Basis: $T_0 = 2^0 - 1 = 0$

Hypothesis: Let, $2^n T_n = 2^n - 1 \quad / \quad T_k = 2^k - 1$
we have to prove,

Induction: $T_{k+1} = 2^{k+1} - 1$

$$\begin{aligned}\therefore T_{k+1} &= 2(T_k + 1) + 1 - 2T_k + 1 \\ &= 2(2^k - 1) + 1 \\ &= 2^{k+1} - 1\end{aligned}$$

m

C-2, w-1

24/06/25

TOH Variation

④ Double Towers of Hanoi:



- S-1 : Move top $n-1$ pairs of disks to aux tower (DT_{n-1})
S-2 : Move largest two pairs of disks to dest tower (2)
S-3 : Move top $n-1$ pairs of disks from aux to dest tower (DT_{n-1})

$$\begin{aligned} DT_n &= DT_{n-1} + 2 + DT_{n-1} \\ \therefore DT_n &= 2DT_{n-1} + 2 = 2(DT_{n-1} + 1) \\ &= 2(2DT_{n-2} + 2) + 2 \\ &= 2^2 DT_{n-2} + 2^2 + 2 \\ &= 2^2 (2DT_{n-3} + 2) + 2^2 + 2 \\ &= 2^3 DT_{n-3} + 2^3 + 2^2 + 2 \end{aligned}$$

$$\begin{aligned}
 &= 2^n DT_{n-n} + 2^n + \dots + 2^3 + 2^2 + 2 \\
 &= 2^n (0) + 2^n + 2 \dots + 2^3 + 2^2 + 2 \\
 &= 2^n + \dots + 2^3 + 2^2 + 2 \\
 &= 2(2^{n-1} + \dots + 2^2 + 2 + 1)
 \end{aligned}$$

$\therefore DT_n = 2(2^n - 1) = 2 \times \text{Single Tower of Hanoi}$

self study : TTOH

* Using Mathematica - Mathematical Induction:

Prove using Mathematical Induction that, for DTOT with n pairs of disks, minimum number of $2(2^n - 1)$ moves are required

Basis: For, $n = 0$, $DT_0 = 2(2^0 - 1) = 0$

Hypothesis: Let, $DT_n = 2(2^n - 1)$ for $n = n$

Induction: $DT_{n+1} = 2(DT_n + 2)$
 $= 2(2^n \times [2(2^n - 1)] + 2)$
 $= 2(2^n - 2) + 2$
 $= 2(2^n - 2 + 1) = 2(2^n - 1)$ (Proved)

Lines in the Plane : Max numbers of regions using n lines

1 Line	2 Region
2 "	4 "
3 "	7 "
4 "	11 "

$$\therefore \text{Ans} T_n = T_{n-1} + n$$

Hence,

$$T_0 = 1$$

$$\begin{aligned}
 \therefore T_n &= T_{n-1} + n = (T_{n-2} + n-1) + n \\
 &= \{T_{n-3} + (n-2)\} + 2n - 1 \\
 &= \{T_0 + n - (n-1)\} + n(n-1) - 1 - 2 - \dots \\
 &= 1 + 1 + (n
 \end{aligned}$$

C-3, W-1

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Lines in Surface:

$$L_n = L_{n-1} + n \text{ (open)}$$

$$= L_{n-2} + (n-1) + n$$

$$= L_{n-3} + (n-2) + (n-1) + n$$

:

:

$$= \cancel{L_{n-n}} + \cancel{\{n}} + (n-1) + n$$

$$= L_0 + 1 + 2 + 3 + \dots + (n-1) + n$$

$$= 1 + \frac{n(n+1)}{2} \text{ (closed)}$$

* Using Mathematical Induction:

$$\text{Basis: } L_0 = 1 + \frac{0(0+1)}{2} = 1$$

$$\text{Hypothesis: Let, } L_{n-1} = 1 + \frac{n(n-1)(n-1+1)}{2} = 1 + \frac{(n-1)n}{2}$$

$$\text{Induction: } L_n = L_{n-1} + n$$

$$= 1 + \frac{(n-1)n}{2} + n$$

$$= \frac{2 + n^2 - n + 2n}{2} = \frac{n^2 + n + 2}{2}$$

$$= \frac{n(n+1)}{2} + \frac{2}{2} = \frac{n(n+1)}{2} + 1 \text{ (Proved)}$$

Variations

If n th line passes through $(n-1)$ th intersection point, $L_n = 2n$

Zig:

No. of Zigs (n)	No. of Lines (2n)	Zigs Region z_n	Line Region (L_{2n})	Lost Region $\binom{L_{2n}}{2n}$
0	2	1	4	2
1	4	2	11	4
2	6	7		

$$\therefore z_n = L_{2n} - 2n = \text{Line Region} - \text{Numbers of Lost Region}$$

$$\begin{aligned}\therefore z_n &= \frac{2n(2n+1)}{2} + 1 - 2n \\ &= 2n^2 + n + 1 - 2n \\ &= 2n^2 - n + 1\end{aligned}$$

#Self: Inductive Proof

* Zigzag:

No. of Zigzags (n)	No. of Lines (3n)	Zigzag Region (22n)	Lines Region (L _{3n})	Regions Less + (22L _{3n} - 22n)
0	0	1	1	0
1	3	2	7	5
2	6	12	22	10

$$\therefore 22n = L_{3n} - 5n$$

$$= \frac{3n(3n+1)}{2} - 5n + 1$$

$$= \frac{9n^2 + 3n - 10n + 2}{2} = \frac{9n^2 - 7n + 2}{2}$$

$$\therefore 22n = \frac{9n^2 - 7n + 2}{2}$$

* W on M Shape:

n	4_n	w_n	L_{4n}	$L_{4n} - w_n$
0	0	1	1	0
1	4	2	11	9
2	8	19	37	18

$$\begin{aligned}w_n &= L_{4n} - 9_n \\&= \frac{4n(n+1)}{2} - 9n + 1 \\&= 8n^2 + 2n - 9n + 1\end{aligned}$$

$$\therefore w_n = 8n^2 - 7n + 1$$

* Using Mathematical Induction:

$$\text{Basis: } w_0 = 8(0)^2 - 7(0) + 1 = 1$$

$$8(n-1)^2 - 7(n-1) + 1$$

$$\text{Hypothesis: } w_{n-1} = 8(n-1)^2 - 7(n-1) + 1$$

$$8n^2 - 7n + 1$$

$$\text{Induction: } w_n = L_{4n} - 9n =$$

\$

C-4, W-2

30/6/25

Open form of lines in planes (variations)

ZZ:

$\frac{n-2}{n-2}$	$\frac{n-1}{n-1}$	$\frac{n}{n}$	$\frac{ZZ_n}{2}$
-1	0	1	2
0	1	2	12
1	2	3	31
2	3	4	59

Regions inclosed

10 10 - (-1)
10 10 - 0
19 20 - 1
28 30 - 1

$$= 10(n-1) - (n-2)$$

→ to find use

$$ZZ_n - ZZ_{n-1}$$

$$\therefore ZZ_n = ZZ_{n-1} + 9n - 8 \text{ (open)}$$

* Mathematical Induction:

$$\text{Basis: } Z_2 = \frac{9(1)^2 - 7(1) + 2}{2} = 2$$

$$\text{Hypothesis: } Z_{n-1} = \frac{9(n-1)^2 - 7(n-1) + 2}{2}$$

$$\text{Induction: } Z_n = Z_{n-1} + 9n - 8$$

$$= \frac{9(n-1)^2 - 7(n-1) + 2}{2} + 9n - 8$$

$$= \frac{9n^2 - 18n + 9 - 7n + 7 + 2 + 18n - 16}{2}$$

$$= \frac{9n^2 - 7n + 2}{2}$$

= (Proved)

* Check for big and small

C-5, w-2

01/07/25

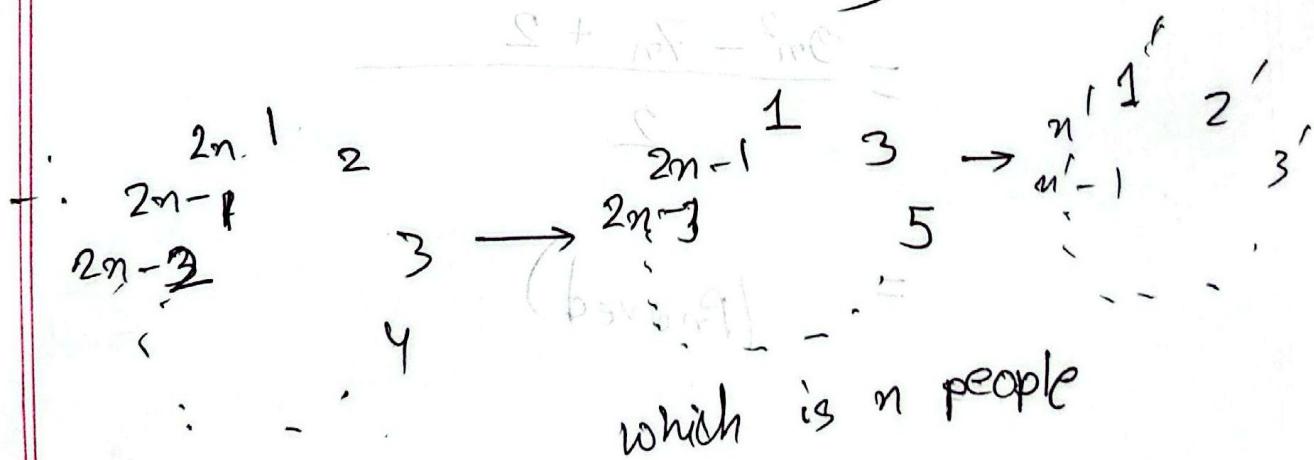
Josephus Problem

* In a cyclic fashion, find the surviving person

$J(n)$ → Survivor for n numbers of people

For $2n$ numbers of People

After 1st round → all even places are missing



∴ $J(n)$ placed person shifts to $2J(n) - 1$

$$\therefore J(2n) = 2J(n) - 1 \quad (\text{even case})$$

$J(1) = 1 \rightarrow 1$ person always survives
(Base case)

\therefore For $2n+1$ people:

Same as ~~area~~ $2n$

Except 1 dies

$$\therefore 2n+1 - n - 1 = n$$

$\therefore J(n) - 2$ is placed at $2(J(2J(n)) + 1)$

$$\therefore J(2n+1) = 2J(n) + 1$$

(odd case)

$$\# n=25$$

$$J(25) = 2J(12) + 1 = 2(2J(6) - 1) + 1$$

$$= 4(2J(3) - 1) - 2 + 1$$

$$= 8J(3) - 4 - 1$$

$$= 8(2J(1) + 1) - 5$$

$$= 8(2 + 1) - 5$$

$$= 19$$

Closed

Open form:

$$J(2^m + i) = 2i + 1 ; [\text{ } 2^m < n]$$

Here,

n	1	2	3	4	5	6	7	8	9	-
$J(n)$	1	1	3	1	3	5	7	1	3	-

2^n
will always
reset to 1

All the outputs
are odd

* Using Mathematical Induction: $J(n) = 2^n + 1$ where $n = 2^m + 2$

Basis: Induction is on m , if $m=0$, then $2=0$ | Hypothesis:

$$\therefore J(2^m + 2) = 2^n + 1$$

$$\text{on, } J(2^0 + 0) = 2(0) + 1$$

$$\therefore J(1) = 1$$

~~Hypothesis and Induction:~~

If $2^m + 2 = 2n$ (Even) and $m > 0$

$$\therefore J(2^m + 2) = 2 \left(2J(2^{m-1} + \frac{2}{2}) \right); \left[\because J(2n) = 2J(n) - 1 \right]$$

$$= 2 \left(2 \frac{2}{2} + 1 \right) - 1; \left[\text{Assuming } C = \frac{2}{2} \right]$$

$$= 2^n + 2 - 1$$

$$= 2^n + 1$$

If $2^m + 2 = 2n + 1$ (odd) and $m > 0$

Now, 2 must be odd $\rightarrow 2^m$ ^{always} will be even; odd + even = odd
 $\therefore 2^m + 2 = \text{even} + \text{odd} = \text{odd}$

Let,

$$q = 2n + 1$$

$$\therefore J(2^m + 2) = J(2^m + 2n + 1)$$

$$= J\left(2 \underbrace{(2^{m-1} + n)}_m + 1\right) \rightarrow J(2n+1)$$

$$= 2J(2^{m-1} + n) + 1$$

$$= 2(J(2^{m-1}) + n) + 1; [Hypothesis]$$

$$= 22 + 1; [\because 2 = J(2^{m-1})]$$

$$\begin{aligned} & (2^m + 1)(2^m + 1) \\ & = 2^{2m} + 2^m + 2^m + 1 \\ & = 2^{2m} + 2 \cdot 2^m + 1 \\ & = 2^{2m} + 2^{m+1} + 1 \end{aligned}$$

$$= 2^{2m} + 2^{m+1} + 1$$

$$= 2^{2m} + 2^{m+1} + 1$$

$$\begin{aligned} & (2^m + 1)(2^m + 1) \\ & = 2^{2m} + 2^m + 2^m + 1 \\ & = 2^{2m} + 2 \cdot 2^m + 1 \\ & = 2^{2m} + 2^{m+1} + 1 \end{aligned}$$

$$\begin{aligned} & (2^m + 1)(2^m + 1) \\ & = 2^{2m} + 2^m + 2^m + 1 \\ & = 2^{2m} + 2 \cdot 2^m + 1 \\ & = 2^{2m} + 2^{m+1} + 1 \end{aligned}$$

Binary Properties of Josphus Problem:

When Let

$$n = (b_m \ b_{m-1} \ b_{m-2} \dots \ b_1 \ b_0)_2$$

$$J(n) = (b_{m-1} \ b_{m-2} \dots \ b_1 \ b_0 \ b_m)_2$$

$$\begin{aligned} \therefore 2 &= (0 \ b_{m-1} \ b_{m-2} \dots \ b_1 \ b_0)_2 \\ \therefore 22 &= (b_{m-1} \ b_{m-2} \dots \ b_1 \ b_0 \ 0)_2; [\text{left shift} = \times 2] \\ \therefore 22 + F &= (b_{m-1} \ b_{m-2} \dots \ b_1 \ b_0 \ 1)_2 \\ &= (b_{m-1} \ b_{m-2} \dots \ b_1 \ b_0 \ b_m)_2 \end{aligned}$$

Radic Based Properties of Josephus Problem:

Generalized Version of Josephus:

$$f(1) = \alpha$$

$$f(2n) = 2f(n) + \beta$$

$$f(2n+1) = 2f(n) + \gamma$$

* we solve using Repetition Method:

when $n=1$ for $f(2)$,

$$f(2n) = f(2 \times 1) = 2f(1) + \beta = 2\alpha + \beta$$

for $f(3)$,
 $f(2n+1) = f(2 \times 1 + 1) = 2f(1) + \gamma = 2\alpha + \gamma$

∴ $f(n)$ will always be,

$$\therefore f(n) = A(n)\alpha + B(n)\beta + C(n)\gamma$$

Let,

$$\alpha = 1, \beta = 0, \gamma = 0$$

$$\therefore f(n) = A(n) + 0 + 0 = A(n) \quad (i)$$

$$\therefore A(1) = 1 = \text{def.}$$

$$A(2n) = 2A(n) + \beta = 2A(n) + 0 = 2A(n)$$

$$A(2n+1) = 2A(n) + \gamma = 2A(n)$$

$$* A(8) = 2A(4) = 2(2(A(2))) = 2(2(2A(1))) \\ = 2^3 = 8$$

$$* A(9) = 2(A(2 \times 4 + 1)) = 2A(4) = 2^3 = 8$$

∴ for this recurrence, $A(n) = 2^n$

Let,

$$f(n) = 1$$

$$\therefore f(1) = 1 = \alpha$$

$$\therefore f(2) = 2f(n) + \beta$$

$$\text{or } 1 = 2 + \beta$$

$$\therefore \beta = -1$$

$$\therefore f(2n+1) = 2f(n) + \gamma$$

$$\text{or, } 1 = 2 + \gamma$$

$$\therefore \gamma = -1$$

$$\therefore 1 = A(n) - B(n) - C(n) \quad \dots \text{(ii)}$$

Let,

$$f(n) = n$$

$$\therefore f(1) = 1 = \alpha$$

$$\therefore f(2n) = 2f(n) + \beta$$

$$\text{or, } 2n = 2n + \beta$$

$$\therefore \beta = 0$$

$$\therefore f(2n+1) = 2f(n) + \gamma$$

$$\text{or, } 2n+1 = 2n + \gamma$$

$$\therefore \gamma = 1$$

$$\therefore n = A(n) + C(n) \quad \dots \text{(iii)}$$

$$\text{or, } 2^m + 2 = A(n) + C(n)$$

$$\text{or, } 2^m + 2 = 2^m + C(n)$$

$$\therefore C(n) = 2$$

From eqⁿ (ii)

$$1 = 2^m - B(n) - 2$$

$$\therefore B(n) = \cancel{1+2} = 2^m - 2 - 1$$

$$\therefore f(n) = 2^m \alpha + (2^m - 2 - 1) \beta + 2 \gamma$$

Putting $\alpha = 1, \beta = -1, \gamma = 1$ (Josephus)

$$\begin{aligned}f(n) &= 2^m - 2^m + 2 + 1 + 1 \\&= 22 + 1\end{aligned}$$

Radix based solution of recurrence:

We can only use this iff,

$$\cancel{f(j) = \alpha_j}; [1 \leq j < d]$$

$$f(dn+j) = cf(n) + \beta_j; [0 \leq j < d]$$

$$\cancel{f(1) = \alpha_1} ; \text{ if } d = 3$$

$$f(2) = \alpha_2$$

$$f(3n+0) = 10f(n) + \beta_0$$

$$f(3n+1) = 10f(n) + \beta_1$$

$$f(3n+2) = 10f(n) + \beta_2$$

If $c=d=2$

$$f(1) = \alpha_1$$

$$f(2n+0) = 2f(n) + \beta_0$$

$$f(2n+1) = 2f(n) + \beta_1$$

If we add $\alpha_1 = 1$, $\beta_0 = -1$, $\beta_1 = 1$, we get

josephus

Hence,

$$\text{For } f(n) = (\dots \dots)_d$$

$$= (a_m b_m \dots \dots)_c$$

Sum

A sum of any series

$$a_0 + a_1 + \dots + a_{n-1} + a_n \rightarrow \text{three dot}$$

$$\sum_{i=0, i \leq n} a_k \longrightarrow \text{Sigma Notation}$$

$$\sum_{k=1}^n a_k \longrightarrow \text{Delimited Form}$$

$$\# S_n = \sum_{k=0}^n a_k$$

$$S_0 = 0 \quad a_0$$

$$\therefore S_n = \underbrace{a_0 + a_1 + a_2 + \dots + a_{n-1}}_{S_{n-1}} + a_n$$

$$\therefore S_n = S_{n-1} + a_n$$

For any general series,

Let,

$$R_0 = \alpha$$

$$R_n = R_{n-1} + \beta + \gamma_n \quad , \text{ Case 2: } R_n = n \text{ to find } R(n)$$

Case 1:

Let,

$$R_n = 1$$

$$\therefore R_0 = 1$$

$$\therefore R_{n-1} = 1$$

$$\therefore R_0 = 1 = \alpha$$

$$R_n = \alpha + n\beta + \frac{n^2 + n}{2} \gamma$$

\hookrightarrow For math no need to derive,
just use it

[Follow notes]

$$S_n = \sum_{k=0}^n (a + bk)$$

Hence,

$$\begin{aligned} S_0 &= a \\ S_n &= a + (a+b) + (a+2b) + \dots + \underbrace{(a+(n-1)b)}_{(a+nb)} + (a+nb) \\ &= S_{n-1} + (a+nb) \end{aligned}$$

By comparing the eq's with general sum eqⁿ, we

get,

$$\alpha = a$$

$$\beta = a$$

$$\gamma = b$$

$$\begin{aligned} \therefore R_n &= a + an + \frac{n^2 + n}{2} b \\ &= a(n+1) + \frac{n^2 + n}{2} b \end{aligned}$$

+ [More math in lecture notes]

Converting Recurrence to sum:

We assume some things:

- we assume some things:
- we always have to multiply something with recurrence.
We assume we know how to find it.

*TOH:

$$T_0 = 0$$

$$T_n = 2T_{n-1} + 1$$

we multiply both sides with 2^{-n}

$$\therefore \frac{T_0}{2^0} = \frac{0}{2^0} = 0$$

$$\therefore \frac{T_n}{2^n} = \frac{2T_{n-1}}{2^n} + \frac{1}{2^n}$$

$$= \frac{T_{n-1}}{2^{n-1}} + \frac{1}{2^n}$$

$$\text{Let, } S_n = \frac{T_n}{2^n}$$

$$\therefore S_n = \frac{T_{n-1}}{2^{n-1}} + \frac{1}{2^n}$$

$$= S_{n-1} + \frac{1}{2^n}$$

$$= S_{n-2} + 2^{-(n-1)} + 2^{-n}$$

$$= S_{n-3} + 2^{-(n-2)} + 2^{-(n-1)} + 2^{-n} = S_{n-n} + \sum_{k=1}^{n-1} 2^{-k}$$

$$\therefore S_0 + 2^{-1} + 2^{-2} + \dots + 2^{-n}$$

$$= \sum_{k=1}^n 2^{-k}$$

Now,

$$S_n = \sum_{k=1}^n 2^{-k}$$

$$\text{or, } S_n + 2^0 = \sum_{k=1}^n 2^{-k} + 2^0$$

$$\text{or, } S_n = \left(2^0 + 2^{-1} + 2^{-2} + \dots + 2^{-n} \right) - 1$$

$$= \frac{1 \left\{ 1 - \left(\frac{1}{2}\right)^n \right\}}{1 - \frac{1}{2}} - 1 ; \boxed{S_n = \frac{a(1-r^n)}{1-r}}$$

$$= \frac{1 - \frac{1}{2^{n+1}}}{1 - \frac{1}{2}} - 1$$

$$= 2 \left(1 - \frac{1}{2^{n+1}} \right) - 1$$

$$\text{or, } \frac{T_n}{2^n} = 1 - \frac{1}{2^n}$$

$$\therefore T_n = 2^n - 1 \quad \text{upto sum}$$

C-8, W-3

08 /07/25

SUMS

For

$$a_n T_n = b_n T_{n-1} + c_n = \frac{1}{s_n a_n} \left(s_1 b_1 T_0 + \sum_{k=1}^n s_k c_k \right)$$

$$\therefore s_n = \frac{a_{n-1} a_{n-2} a_{n-3} \dots a_1}{b_n b_{n-1} b_{n-2} \dots b_2}$$

↓
summation factors

* Index Change: $\sum_{k=1}^n 2^{-k+1}$ → Let, $j = k-1$

when $k=1, j=0$

" $k=n, j=n-1$

$$\therefore \sum_{k=1}^n 2^{-k+1} = \sum_{j=0}^{n-1} 2^{-j-1}$$

$$= \sum_{k=0}^{n-1} 2^{-k} \left(\frac{1}{1+2^{-1}} - 1 \right) =$$

$$\frac{1}{1+2^{-1}} - 1 = \frac{1}{3}$$

$$\text{Final step} \rightarrow 1 - \sum_{k=0}^{n-1} 2^{-k} = T$$

#Quick Sort: $C_0 = 0$

$$C_n = n+1 + \frac{2}{n} \sum_{k=0}^{n-1} C_k \rightarrow C_n \text{ is number of comparisons}$$

Hence,

$$C_{n-1} = n C_n = n(n+1) + 2 \sum_{k=0}^{n-1} C_k$$

$$\therefore (n-1) C_{n-1} = \cancel{(n-1)^2} + \cancel{n^2+n} - 1 + 2 \sum_{k=0}^{n-1-1} C_k \\ = \cancel{(n-1)} n$$

$$\therefore n C_n - (n-1) C_{n-1} = n^2 + n + 2 \sum_{k=0}^{n-1} C_k - n^2 - \cancel{n} + n$$

$$\therefore n C_n = (n+1) C_{n-1} + 2n$$

$$a_n = n$$

$$b_n = n+1$$

$$d_n = 2n \rightarrow \text{usually } C_n$$

$$\therefore S_n = \frac{a_{n-1} \cdots a_1}{b_n \cdots b_2} = \frac{(n-1)(n-2) \cdots 3, 2, 1}{(n+1)(n)(n-1) \cdots 3} \\ = \frac{2 \cdot 1}{n(n+1)}$$

Harmonic Series: $\sum_{k=1}^n \frac{1}{k}$

Hence,

$$C_n = 2(n+1) \sum_{k=1}^n \frac{1}{k+1} \rightarrow \text{derivation from provided notes}$$

Now,

$$\begin{aligned} \sum_{k=1}^n \frac{1}{k+1} &= \cancel{\left(\frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n+1} \right)} + \frac{1}{n+1} \\ &= \left(1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n} \right) + \frac{1}{n+1} - 1 \\ &= H_n - \frac{n}{n+1} \end{aligned}$$

$$\therefore C_n = 2(n+1)H_n - 2n$$

* Index Replacement: using

Must be in $\sum_{k=0}^n (a + b^k)$ form

- Replace α_k by $n-k$
 - becomes $a(n+1) \frac{n(n+1)}{2} b$

• 10 in the first few years before

$$\# \sum_{k=0}^{\infty} (2k+1) \rightarrow a = 1 \text{ unter best.} \\ b = 2$$

$$\therefore (11) + \frac{10(11)}{2}(2) = 11 + 110 \\ = 121$$

C-9, W-3

10/07/25

Perturbation Method:

Given,

$$S_n = \sum_{0 \leq k \leq n} a_k \rightarrow \text{convert it to closed form}$$

↪ open because we can't put n to find value

$$S_{n+1} = S_n + a_{n+1} + 1 = S_n + \sum_{1 \leq k \leq n+1} a_k ; \text{Step 1: Express in this form}$$

$$= a_0 + \sum_{1 \leq k+1 \leq n} a_{k+1}$$

$$\boxed{S_n + a_{n+1} = a_0 + \sum_{0 \leq k \leq n} a_{k+1} \begin{cases} k+1 = 0 \\ k = 0 \\ k+1 = n+1 \\ k = n \end{cases}}$$

$$= a_0 + S_n$$

↪ Main formula
for Perturbation
Method

Find $S_n = \sum_{0 \leq k \leq n} ax^k$ in closed form

$$S_n = \sum_{0 \leq k \leq n} ax^k$$

$$S_n + ax^{n+1} = ax^0 + \sum_{0 \leq k \leq n} ax^{k+1}$$

$$= a + \cancel{ax^1} \sum_{0 \leq k \leq n} ax^k \cdot x$$

$$\frac{S_n + ax^{n+1}}{x} = a + x \sum_{0 \leq k \leq n} ax^k$$

$$S_n + ax^{n+1} = a + xS_n$$

$$\therefore S_n = \frac{a - ax^{n+1}}{1 - x}$$

$$\therefore S_8 = \frac{3 - 3 \cdot 5^9}{1 - 5} = 1464843$$

$$\# S_n = \sum k \cdot 2^k$$

Hence,

$$\begin{aligned}
 S_n + a & (n+1) \cdot 2^{n+1} = 0 \cdot 2^0 + \sum_{0 \leq k \leq n} (k+1) 2^{k+1} \\
 &= \sum (k 2^{k+1}) + \sum 2^{k+1} \\
 &= \sum k \cdot 2^k \cdot 2 + \sum 2^{k+1} \cdot 2 \\
 &= 2 \sum_{0 \leq k \leq n} k \cdot 2^k + \frac{2 - 2 \cdot 2^{n+1}}{1 - 2}
 \end{aligned}$$

$$S_n + (n+1) 2^{n+1} = 2S_n + \frac{2 - 2 \cdot 2^{n+1}}{1 - 2}$$

$$\therefore S_n = \cancel{(n+1) \cdot 2^{n+1}} + \cancel{2 - 2^{n+1}} \quad \frac{\cancel{2 - 2^{n+1}}}{\cancel{1 - 2}}$$

$$\therefore S_n = 2^{n+1}(n-1) + 2$$

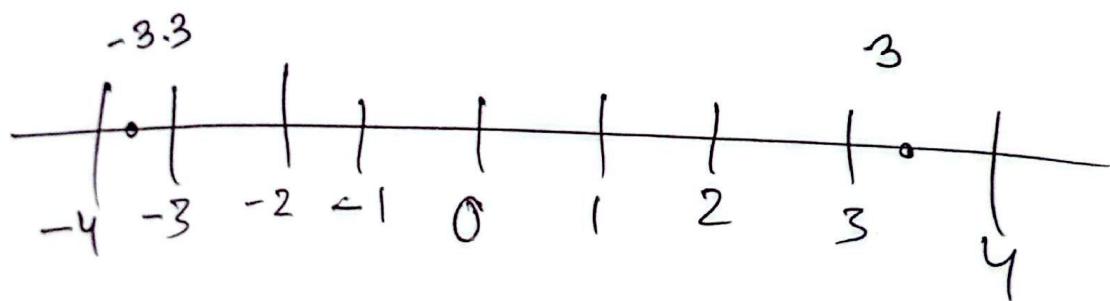
$$\# S_n = \sum k x^k =$$

Chap - 3 : Integer Functions

Floor: $\lfloor 3.3 \rfloor = 3$, $\lfloor -3.3 \rfloor = -4$

Ceiling: $\lceil 3.3 \rceil = 4$, $\lceil -3.3 \rceil = -3$

Mod:



General form of eqⁿ: $\left\lfloor \frac{n}{m} \right\rfloor \times m + n \bmod m = n$

C-10, W-4

14/07/25

Necessary Rules: for $x \in \mathbb{R}$, $n \in \mathbb{Z}$

$$1. \lfloor x \rfloor = n \iff n \leq x < n+1$$

$$2. \lfloor x \rfloor = n \iff x-1 < n \leq x$$

$$3. \lceil x \rceil = n \iff n-1 < x \leq n$$

$$4. \lceil x \rceil = n \iff x \leq n \leq x+1$$

$$5. \lfloor x+n \rfloor = \lfloor x \rfloor + n$$

$$6. n \leq x \iff n \leq \lfloor x \rfloor$$

$$7. x \leq n \iff \lceil x \rceil \leq n$$

#1.

Let,

$$x = 3$$

$$\therefore \lfloor 3 \rfloor = 3$$

$$\therefore n = 3$$

For R.H.S,

$$3 \leq 3 < 3 + 1$$

$$\therefore 3 \leq 3 < 4$$

\therefore when $x \in \mathbb{Z}$, the formula holds

Let,

$$x = 3.5$$

$$\lfloor x \rfloor = \lfloor 3.5 \rfloor = 3 = n$$

For R.H.S,

$$3 \leq 3.5 < 3 + 1$$

$$\therefore 3 \leq 3.5 < 4$$

\therefore when x is floating, the formula holds

$\therefore 1$ is always true

Q. If n is an m -bit integer, then prove

that $m = \lfloor \lg n \rfloor + 1 \rightarrow$ To find how many bits are needed to express a decimal in binary
Hence,

The smallest m bit number is 2^{m-1}

(i) largest m " " " " $2^m - 1$

$$\therefore 2^{m-1} \leq n \leq 2^m - 1$$

or, $2^{m-1} \leq n < 2^m$

or, $\lg(2^{m-1}) \leq \lg n < \lg(2^m)$

or, $m-1 \leq \lg n < m$

or, $m \leq \lg n + 1 < m + 1$

or, $m = \lfloor \lg n + 1 \rfloor ; [\lfloor x \rfloor = n \Leftrightarrow n \leq x < n+1]$

~~or, $\lfloor \lg n + 1 \rfloor \rightarrow$~~

$\therefore m = \lfloor \lg n \rfloor + 1 ; [\lfloor x+n \rfloor = \lfloor x \rfloor + n]$

Q. Prove that, $\lfloor \sqrt{\lfloor x \rfloor} \rfloor = \lfloor \sqrt{x} \rfloor$; $x \in \mathbb{R}, x \geq 0$

Let,

x be an integer

$$\text{L.H.S} = \lfloor \sqrt{\lfloor x \rfloor} \rfloor = \lfloor \sqrt{x} \rfloor = \text{R.H.S}$$

Now,
If x is not an integer,

Let,

$$m = \lfloor \sqrt{\lfloor x \rfloor} \rfloor$$

$$\text{or, } m \leq \sqrt{\lfloor x \rfloor} < m+1$$

$$\text{or, } m^2 \leq \lfloor x \rfloor < (m+1)^2$$

$$\text{or, } m^2 \leq x < (m+1)^2 ; [\text{Rule 6}]$$

$$\text{or, } m \leq \sqrt{x} < (m+1)$$

$$\therefore \lfloor \sqrt{x} \rfloor = m$$

$$\therefore \lfloor \sqrt{\lfloor x \rfloor} \rfloor = \lfloor \sqrt{x} \rfloor$$

$$\therefore \text{L.H.S} = \text{R.H.S} \quad (\text{Proved})$$

$$\star n = \left\lfloor \frac{n}{m} \right\rfloor * m + n \bmod m$$

$$\therefore n \bmod m = n - \left\lfloor \frac{n}{m} \right\rfloor * m$$

$$\begin{aligned}\text{Case 1: } 5 \bmod 3 &= 5 - \left\lfloor \frac{5}{3} \right\rfloor * 3 \\ &= 5 - 1 * 3 \\ &= 2\end{aligned}$$

$$\begin{aligned}\text{Case 2: } -5 \bmod 3 &= -5 - \left(\left\lfloor \frac{-5}{3} \right\rfloor * 3 \right) \\ &= -5 - \left(-\left\lceil \frac{5}{3} \right\rceil * 3 \right) \\ &= -5 - (-2 * 3) \\ &= -5 + 6 \\ &= 1\end{aligned}$$

* Special case: $n \bmod 0 = n$

C-II, W-4

15/07/25

Prove that $c(x \bmod y) = cx \bmod cy$

$$\begin{aligned} L.H.S &= c(x \bmod y) \\ &= c\left(x - y \cdot \left\lfloor \frac{x}{y} \right\rfloor\right); \quad \left[n \bmod m = n - m \cdot \left\lfloor \frac{n}{m} \right\rfloor \right] \\ &= cx - cy \cdot \left\lfloor \frac{x}{y} \right\rfloor \\ &= cy - cy \cdot \left\lfloor \frac{cx}{cy} \right\rfloor \\ &= cx \bmod cy \quad ; \quad \boxed{\quad} \end{aligned}$$

Prove that, $(x \bmod ny) \bmod y = x \bmod y$

Case 1: Let, $y = 0$

$$\begin{aligned} L.H.S &= (x \bmod n \cdot 0) \bmod 0 \\ &= (x \bmod 0) \bmod 0 \\ &= x \bmod y^* \end{aligned}$$

B.

Case 2: Let $y \neq 0$,

$$\begin{aligned} L.H.S &= (x \bmod ny) \bmod y \\ &= (x - ny \cdot \left\lfloor \frac{x}{ny} \right\rfloor) \bmod y \end{aligned}$$

$$= \left(x - ny + \left\lfloor \frac{x}{ny} \right\rfloor \right) - y \left\lfloor \frac{x - ny + \left\lfloor \frac{x}{ny} \right\rfloor}{y} \right\rfloor$$

$$= \left(x - ny + \left\lfloor \frac{x}{ny} \right\rfloor \right) - y \left\lfloor \frac{x}{y} - \frac{ny}{y} \left\lfloor \frac{x}{ny} \right\rfloor \right\rfloor \quad \begin{matrix} \text{integers} \\ \text{ny} \\ \text{y} \end{matrix}$$

$$= x - ny \left\lfloor \frac{x}{ny} \right\rfloor - y \left\lfloor \frac{x}{y} \right\rfloor + ny \left\lfloor \frac{x}{ny} \right\rfloor$$

$$= x - y \left\lfloor \frac{x}{y} \right\rfloor$$

$$= x \bmod y$$

= R.H.S

C-12, W-5

21/07/23

Number Theory

$K \nmid m$

$K \mid m \rightarrow m = \text{dividend} \rightarrow \frac{m}{K} \approx \text{some integer}$
 $K = \text{divisor}$

$\text{GCD}(m, n) = K \mid m \text{ and } K \mid n$

$K^{\max m}$

$\text{LCM}(m, n) = K \mid m \mid K \text{ and } K \mid n$

$K^{\min m}$

GCD: Stein's Formula:

$$1. \text{GCD}(0, v) = v$$

2. If u and v are both even then,

$$g(u, v) = 2g(u/2, v)$$

3. If u is even and v is odd then,

$$g(u, v) = g(u/2, v)$$

4. If u is odd but v is even,

$$g(u, v) = g(u, \frac{v}{2})$$

5. If u and v are both odd,

i) if $u \geq v$

$$g(u, v) = g\left(\frac{u-v}{2}, v\right)$$

ii) if $u < v$,

$$g(u, v) = g\left(\frac{v-u}{2}, v\right)$$

$$\# \text{GCD}(12, 18)$$

$$= 2 \text{GCD}(6, 9)$$

$$= 2 \text{GCD}(3, 9)$$

$$= 2 \text{GCD}\left(\frac{9-3}{2}, 9\right)$$

$$= 2 \text{GCD}(3, 3)$$

$$= 2 \text{GCD}\left(\frac{3-3}{2}, 3\right)$$

$$= 2 \text{GCD}(0, 3)$$

$$= 2 \times 3$$

$$= 6$$

Euclid's Formula:

GCD

$$\text{GCD}(u, v) = \begin{cases} \text{GCD}(0, v), & \text{if } u=0 \\ \text{GCD}(v \bmod u, u), & \text{if } u>0 \end{cases}$$

$$\# \text{GCD}(12, 18)$$

$$= \text{GCD}(18 \bmod 12, 12)$$

$$= \text{GCD}(6, 12)$$

$$= \text{GCD}(12 \bmod 6, 6)$$

$$= \text{GCD}(0, 6)$$

$$= 6$$

LCM:

$$m \times n = \text{GCD}(m, n) \times \text{LCM}(m, n)$$

$$\therefore \text{LCM} = \frac{m \times n}{\text{GCD}(m, n)}$$

→ Use this if LCM using
Euclid

Prime Numbers: If divisor is 1 and only itself

Composite Numbers: If more than 2 divisors

Fundamental Theorem of Arithmetic:

There is only one way to write n as a product of primes in non decreasing order (proof by contradiction)

Let, n can be expressed as two ways in product of primes

$$\therefore n = p_1 * p_2 * p_3 * \dots * p_m$$

and, $n = q_1 * q_2 * q_3 * \dots * q_m$ and $p_1 < q_1$

and $p_1 * (p_2 * p_3 * \dots * p_m) = q_1 * q_2 * q_3 * \dots * q_m$

or, $p_2 * p_3 * \dots * p_m = \frac{q_1 * q_2 * q_3 * \dots * q_m}{p_1}$

Hence, L.H.S is an integer

∴ R.H.S must also be an integer

but this can not be the case as q_1 is a prime

and cannot be divided by p_1

$$\therefore p_1 \nmid q_1, p_1 \nmid q_2$$

Similarly, we can prove it for p_2, p_3, \dots, p_m

$$\therefore p_1 \nmid q_1 \rightarrow \text{There is only one way to } \dots$$

Every positive integer can be expressed as

$$n = \prod_p p^{n_p} \quad (\text{Product of primes something})$$

(Self)

$$12 = (2^2 \times 3^1 \times 5^0 \times 7^0 \times \dots)$$

$$18 = (2^1 \times 3^2 \times 5^0 \times 7^0 \times \dots)$$

$$19 = (2^0 \times 3^0 \times 5^0 \times 7^0 \dots 19^1)$$

C-13, W-5

22/02/23

Any positive integer > 1 is a prime or product of prime

OR, Every composite number has a prime divisor.

Let,

smallest numbers which is not a prime or product of prime
 $\rightarrow 2 \text{ to } (n-1)$ is a prime or product of prime

i.e. We can say

$n = a \times b$, where a is a composite number
 b " "

or

Now, if a, b are factors of n , then $2 \leq a, b < n$

Hence, if a, b are both primes,

then n becomes products of primes

Again, if a, b are neither prime,
then a and b can be expressed as product of prime

∴ the no becomes products of prime

- Our assumption contradicts that n is a product of prime
- There is no such n .

(Proved)

Euclid's Theorem: There are infinitely many prime numbers:

Let,

there are \star a finite number of primes.

$$\therefore S = \{2, 3, 5, 7, \dots, p_k\}$$

Again, let,

$$A = (2 \times 3 \times 5 \times 7 \times \dots \times p_k) + 1$$

A is not divisible by any of the prime numbers $2, 3, 5, 7, \dots, p_k$
because A-1 is " " the " "

$\therefore A$ cannot be expressed as product of primes

$\therefore A$ is prime & greater than p_k
(proved)

Euler Number:

$$e_n = e_0 + e_1 * e_2 + \dots + e_{n-1} + 1$$

↳ open form

$$e_0 = 1$$

$$\therefore e_1 = e_0 + 1 \\ = 1 + 1 = 2$$

Usually they are primes

But not always

$$\therefore e_5 = 1807 = 13 * \cancel{13} 139$$

→ For closed form:

$$e_n = e_0 * e_1 * \dots * e_{n-1} + 1$$

$$\text{or, } e_n = e_1 * e_2 * \dots * e_{n-1} + 1$$

Let,

$$e_1 * e_2 * \dots * e_{n-1} = p_{n-1}$$

$$\therefore e_n = p_{n-1} + 1$$

$$\therefore p_n = e_1 * e_2 * \dots * e_n$$

$$\text{Or, } p_{n+2} = p_{n+1} + e_n$$

$$\text{Or, } p_n = (p_{n-1}) (p_{n-1} + 1) \dots \text{ (i)}$$

Now,

$$e_n = p_{n-1} + 1$$

$$\hookrightarrow e_2 = e_{n-1} + 1$$

$$\therefore e_{n-1} = p_{n-2} + 1$$

$$\hookrightarrow = p_{n-2} \cdot e_{n-1} + 1$$

$$= (e_{n-1} + 1)(e_{n-1}) + 1$$

$$= e_{n-1}^2 - e_{n-1} + 1$$

* Sieve of Eratosthenes: Finding primes in
any specific range

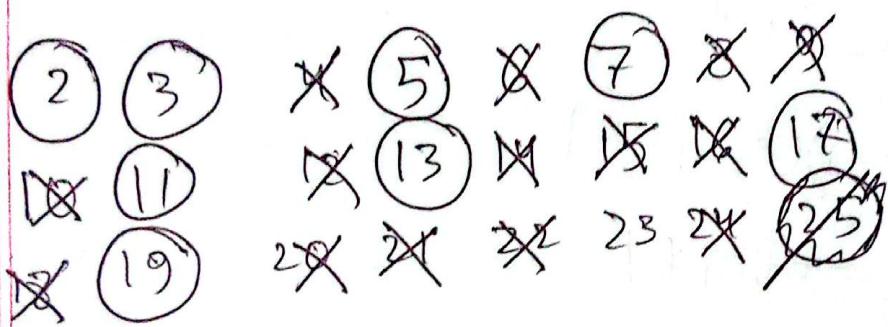
$$\# 2 - 25$$

Let,

$$25 = n$$

$$\therefore \sqrt{n} = 5$$

Here,



Step 1: Remove all
Multiples of 2

Step 2: Remove all
Multiples of the
next uncrossed
numbers

Step 3: Repeat until
(all) uncrossed
 $= \sqrt{n}$

* Mersenne Number: $2^p - 1$ (where p is a prime)

Not all Mersenne Numbers are prime

But the ones that are, are called Mersenne Prime



C-14, W-5

24/07/23

Factorial Factors:

$E_p(n!)$ = max^m power of prime p is $n!$

$$E_2(4!) = 4! = 1 \times 2 \times 3 \times 4 \\ = 1 \times 2 \times 3 \times 2 \times 2$$

$$= 2^3 \times 3 \\ = 8 \quad \text{→ how many of } p \text{ we can find}$$

Generalized Version:

$$\# E_2(10!) \rightarrow p=2 \\ n=10$$

Divisible by	1 to n										Buckets of 2
	1	2	3	4	5	6	7	8	9	10	
2	X		X	X	X	X					$\lfloor \frac{10}{2} \rfloor = 5$
4				X				X			$\lfloor \frac{10}{4} \rfloor = 2$
8						X					$\lfloor \frac{10}{8} \rfloor = 1$
Power of 2	0	1	0	2	0	10	30	1		8	

$2^4 = 16$
so we take
 $2^3 = 8$

$$\therefore E_p(n!) = \left\lfloor \frac{n}{p} \right\rfloor + \left\lfloor \frac{n}{p^2} \right\rfloor + \left\lfloor \frac{n}{p^3} \right\rfloor + \dots + \left\lfloor \frac{n}{p^k} \right\rfloor ; p^k \leq n$$

* For $E_2(n)$,

$$E_2(n) = n - v_2(n) \quad \text{where } v_2 \text{ is the function of finding number of 1 in binary}$$

* For Upper bound:

$$\begin{aligned} E_p(n!) &= \left\lfloor \frac{n}{p} \right\rfloor + \left\lfloor \frac{n}{p^2} \right\rfloor + \dots + \left\lfloor \frac{n}{p^k} \right\rfloor \\ &\leq \frac{n}{p} + \frac{n}{p^2} + \frac{n}{p^3} + \dots = \frac{n}{p} \left(1 + \frac{1}{p} + \frac{1}{p^2} + \dots \right) \\ &= \frac{n}{p} \left(\frac{1}{1 - \frac{1}{p}} \right) \end{aligned}$$

$$= \frac{n}{p} \left(\frac{p}{p-1} \right)$$

$$= \frac{n}{p-1}$$

→ Actual ans
is smaller

$$E_p(n!) \leq$$

Relative Primality: $m \perp n$

then $\rightarrow \gcd(m, n) = 1$

*Stern-Brocot Tree: Binary tree where,

if left ancestor $= \frac{m}{n}$

if right " $= \frac{m'}{n'}$

then, child (median) $= \frac{m+m'}{n+n'}$

we find relative primes & using this

* Always start with $\frac{0}{1}$ and $\frac{1}{0}$

* Stern Brocot tree always follows $m'n - mn' = 1$

Basis: Initially, $\frac{m}{n} = \frac{0}{1}$, $\frac{m'}{n'} = \frac{1}{0}$

$$\therefore m'n - mn' = 1 \times 1 - 0 \times 0 = 1$$

Hypothesis: Let, $m'n - mn' = 1$ is true for two consecutive fractions $\frac{m}{n}$ and $\frac{m'}{n'}$

Induction: The new median $= \frac{m+m'}{n+n'}$

First case: $m = m+m'$, $n = n+n'$

$$\therefore (m+m')n - m(n+n') = mn + m'n - mn - m'n = m'n - mn' = 1$$

Again, $m = (m+m')$, $n = n+n'$

$$\therefore (m+m')(n+n') - m(n+n') - m'(n+n') = 1$$

$\frac{m}{n} < \frac{m+m'}{n+n'} < \frac{m'}{n'},$; $\frac{m}{n} < \frac{m'}{n'}$ and all values are non-negative

Hence,

$$\frac{m+m'}{n+n'} - \frac{m}{n} = \frac{mn + m'n' - mn - mn'}{n(n+n')} = \frac{1}{n(n+n')} > 0$$

$$\therefore \frac{m+m'}{n+n'} < \frac{m}{n}$$

Again,

$$\frac{m'}{n'} - \frac{m+m'}{n+n'} = \frac{m'n + m'n' - mn' - m'n'}{n(n+n')} = \frac{1}{n(n+n')} > 0$$

$$\therefore \frac{m+m'}{n+n'} < \frac{m'}{n'}$$

(Proved)

— —

C-15, W-6

28/07/23

Binomial Coefficients

→ use this form $\binom{n}{k}$

$$\binom{n}{k} = {}^n C_k = \frac{n!}{k!(n-k)!} = \frac{n(n-1)(n-2)\dots(n-k+1)}{k!}; k > 0$$

from n elements
we choose k in each subset

$$\frac{n!}{k!(n-k)!} = \frac{n(n-1)(n-2)\dots(n-k+1)}{k!}; k > 0$$

$$L.H.S = \frac{n!}{k!(n-k)!}$$

$$= \frac{n \times (n-1)(n-2)\dots(n-k+1)(n-k)(n-k-1)\dots 2 \cdot 1}{k! \cdot (n-k)(n-k-1)\dots 2 \cdot 1}$$

$$= \frac{n(n-1)(n-2)\dots(n-k+1)}{k!}$$

$$\binom{n}{0} = 1 \quad \binom{n}{1} = n$$

$$\binom{n}{k} = 0 \text{ if } k < 0 \text{ or } k > n$$

* Pascal's Triangle :

$$\binom{n}{k} = \binom{n}{n-k} \rightarrow \text{Symmetry identity}$$

$$\underline{\text{R.H.S}} = \frac{n!}{(n-k)! (n-(n-k))!}$$

$$= \frac{n!}{(n-k)! k!}$$

$$= \underline{\text{L.H.S}}$$

$$\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1} \rightarrow \text{Absorption Identity}$$

$$\underline{\text{L.H.S}} = \frac{n!}{k! (n-k)!} = \frac{n(n-1)!}{k(k-1)! (n-1) \cancel{(n-k+1)}!}$$

$$= \frac{n}{k} \cdot \frac{(n-1)!}{\cancel{(k-1)!} \cancel{(n-1)-(k-1)}!}$$

$$= \frac{n}{k} \cdot \frac{\cancel{(n-1)!}}{\cancel{(k-1)!} \cancel{(n-1)}} \left(\frac{n-1}{k-1} \right)$$

$$\cancel{k(n-k)} \left(\frac{n}{k} \right) = \cancel{(n-k)} \left(\frac{n}{n-k} \right)$$

$$k(n-k) \left(\frac{n}{n-k} \right) = n \left(\frac{n-1}{k} \right)$$

$$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1} \leftarrow \text{Addition formula}$$

C-16, W-6

29/07/25

$$\# \binom{5}{3} = \left(\cancel{\binom{5}{2}} + \cancel{\binom{4}{2}} \right) \binom{4}{3} + \left(\cancel{\binom{4}{2}} \right); [\text{Addition}]$$

to find
this

②

Use this

$$\sum_{0 \leq k \leq n} \binom{n+k}{k} = \binom{n+n+1}{n}$$

$$\sum_{0 \leq k \leq n} \binom{k}{m} = \binom{n+1}{m+1}$$

$$\# \sum_{k \leq n} \binom{m+k}{k} = \binom{m+n+1}{n}; [k \leq n]$$

$$L.H.S = \sum_{0 \leq k \leq n} \binom{m+n}{m+k-k} = \sum_{-m \leq k \leq n} \binom{m+k}{m} = \sum_{0 \leq k \leq m+n} \binom{k}{m}$$

$$= \binom{m+n+1}{m+1} = \binom{m+n+1}{m+n+1-m-1}$$

$$= \binom{m+n+1}{n} = R.H.S$$

(Proved)

C-17, W-2

04/07/25

Probability Models

For $S = \{H, T\}$

$$\therefore P(H) = \frac{1}{2}$$

$$P(\text{favorable event}) = \frac{|\text{favorable event}|}{|\text{total event}|}$$

* For eve sample spaces E, F

$$P(E \cap F) =$$

Also written
as $P(EF)$

$$P(E \cup F) = P(E) + P(F) - P(E \cap F)$$

↳ OR
 $P(E \cap F) = P(E) + P(F) - P(E \cup F)$

$$P(E \cup F)^c = 1 - P(E \cup F)$$

→ For mutually exclusive,

$$P(E \cup F) = P(E) + P(F)$$

General Formula :

$$P(E_1 \cup E_2 \cup \dots \cup E_n)$$

$$= \sum P(E_i) - \sum P(E_i E_j) + \sum P(E_i E_j E_k) - \dots + (-1)^n \sum P(E_i E_j \dots E_n)$$

Conditional Probability :

Probability of E given F

$$P(E|F) = \frac{P(EF)}{P(F)}$$

Independent :

$$P(EF) = P(E)P(F)$$

Bayes Rule / Formula:

$$P(E|F) = \frac{P(F|E) \cdot P(E)}{P(F)}$$

↑ Likelihood ↓ Evidence

↓ Posterior Probability ↓ Prior Probability

Hence,

$$P(F) = P(EF) + P(E^c F)$$

$$= P(F|E) \cdot P(E) + P(F|E^c) \cdot P(E^c)$$

only when two cases

General Form:

$$P(F_j|E) = \frac{P(E|F_j) P(F_j)}{\sum_{i=1}^n P(E|F_i) P(F_i)} = \frac{P(EF_j)}{P(E)}$$

for j-th event

* $\frac{(n-1)!}{n!}$ for finding out wrong