

C-1, 2 W-1

22/06/25

CSE - 31

Md. Siam Ansary

Room: 9B02

OH: Sun 12:10 - 1:50

Wed 11:20 - 1:50

Quiz: 3rd, 6th, 9th, 12th week  
Syll: Whatever is done at that time

Front End: Visual representation

Back End: DB

DBMS: Software package to store and manage data

## ACID Property:

- i) Atomic: ~~State~~ Nothing is lost
- ii) Consistency: Value is same
- iii) Isolation: Multiple Read/Write is not possible
- iv) Durability: Durable (idk man)

## Abstraction Level:

- i) Physical: The literal physical record
- ii) Logical: Relationship among data (digital <sup>data type</sup> record)
- iii) View: Front end representation

## Schema:

- i) Physical } Arrangement
- ii) Logical } of levels
- iii) Sub :

Instance: Content at a specific time

- i) Single: ~~Single~~ keeps only one instance
- ii) Oracle: Keeps multiple instances

Data Independence: Data can be changed without affecting other layers

- i) Physical } Exactly what it sounds like
- ii) Logical }

### DB Users:

- i) Application Programmers : Makers
- ii) Sophisticated Users : Analyst
- iii) Specialized Users : Admin
- iv) Naive Users : Everyone Else

Data Definition Language: Language to define the data

- Used At the time of creation
- Primary Key: Unique id data

Data Manipulation Language: Language to manipulate the data

→ Pure: - How a data is manipulated (procedural)

- Which data is manipulated (declarative)

- i) Pure: Algebra

ii) Commercial: SQL

## DB Architecture:

- i) Centralized: Single DB which is connected to every device (small scope)
- ii) Client - Server: DB is stored in a server which is accessed from elsewhere
- iii) Parallel: DB is stored in multiple nodes
- iv) Distributed: Multiple DB in multiple places but are connected

\* check file for queries  
(commands)

variables (n) → string of n characters

select \* from table → select distinct column from table  
↳ all values

↳ table name

→ all

→ display

insert into table values (data-1, data-2, ...)

↳ corresponding

↳ table name

→ self explanatory

\* to enter in specific columns, we need to explicitly specify

select \* from table where 'Condition' and/or 'condition 2'

↳ ~~adds~~ filters only those which fit condition

- - . where Col in (a, b)

↑ same as

- - . where ~~Col in~~ Col = a or Col = b

- - . where Col between (a, b)  $\rightarrow a \leq x \leq b$  is printed

- - . where Col like 'start %. end'

↳ matches like stand and end

[check file for the next]

Super Keys

ID  
Email

ID, Email  
Email, Name

Candidate Key

ID  
Email

PK ~~ID~~ ID  
Alternative Key Email

Can	Car	Engine No	Chassis No.	Reg	Color	Model

Candidate Key: Can Engine No.  
Chassis No.  
Reg

PK: Reg

Multiple attributes  
can be PK

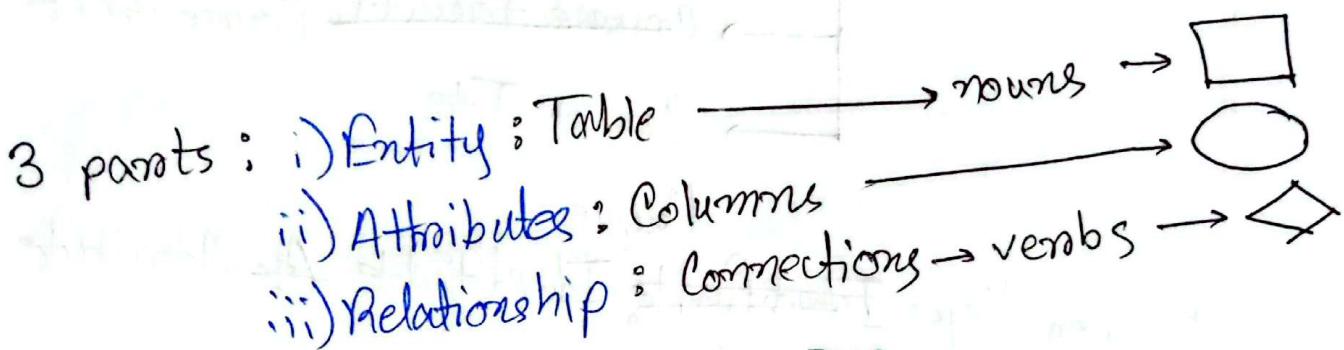
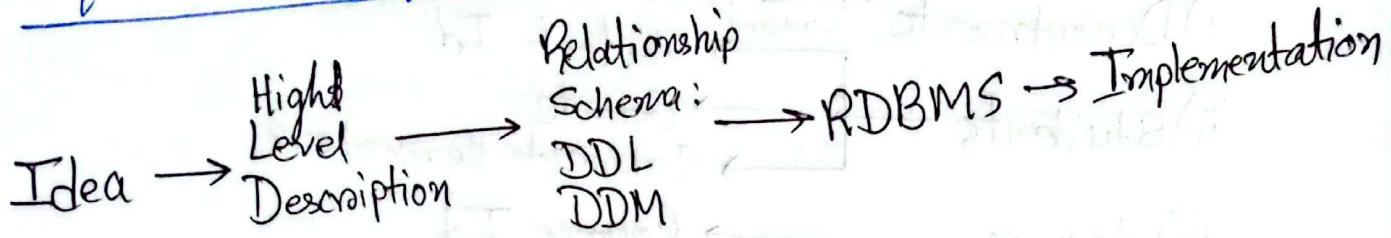
C-2, W-1

# DB theory

25/06/25

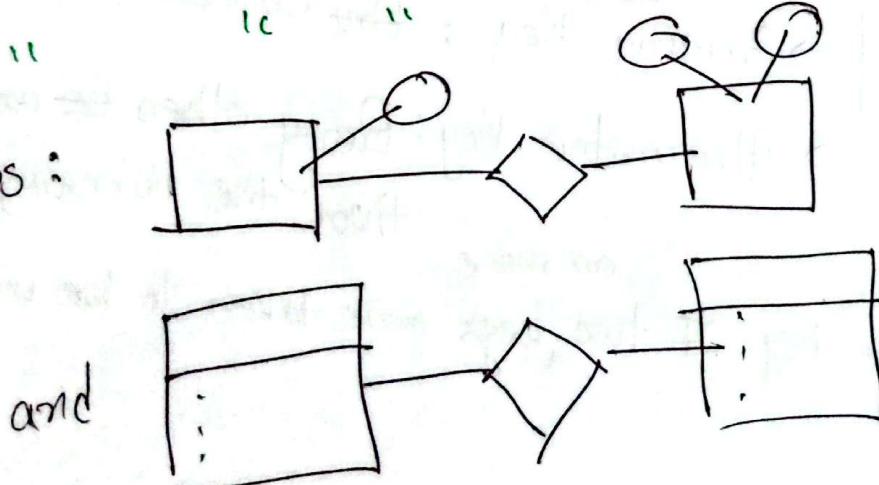
## Entity Relationship Diagram

### \*Requirement Analysis:



\* Some nouns may be the DB  
and " " " attributes

2 meps :



Aust is . - receive grades

classrooms

Entity:

i) Departments

→ labs  
→ No. of students

ii) Students

→ name  
→ student Id  
→ Grade  
→ Current semester

iii) Courses

→ Course Id  
→ Course faculty  
→ Course Credit  
→ Course Title

Unique

\* Super Keys: ~~Identifiable~~, Identifiable At. Attribute

↳ Candidate Key: Shortest Super Keys

↳ Primary Key: That Candidate Key used to identify

↳ Alternative Key: Every other candidate key than the Primary Key

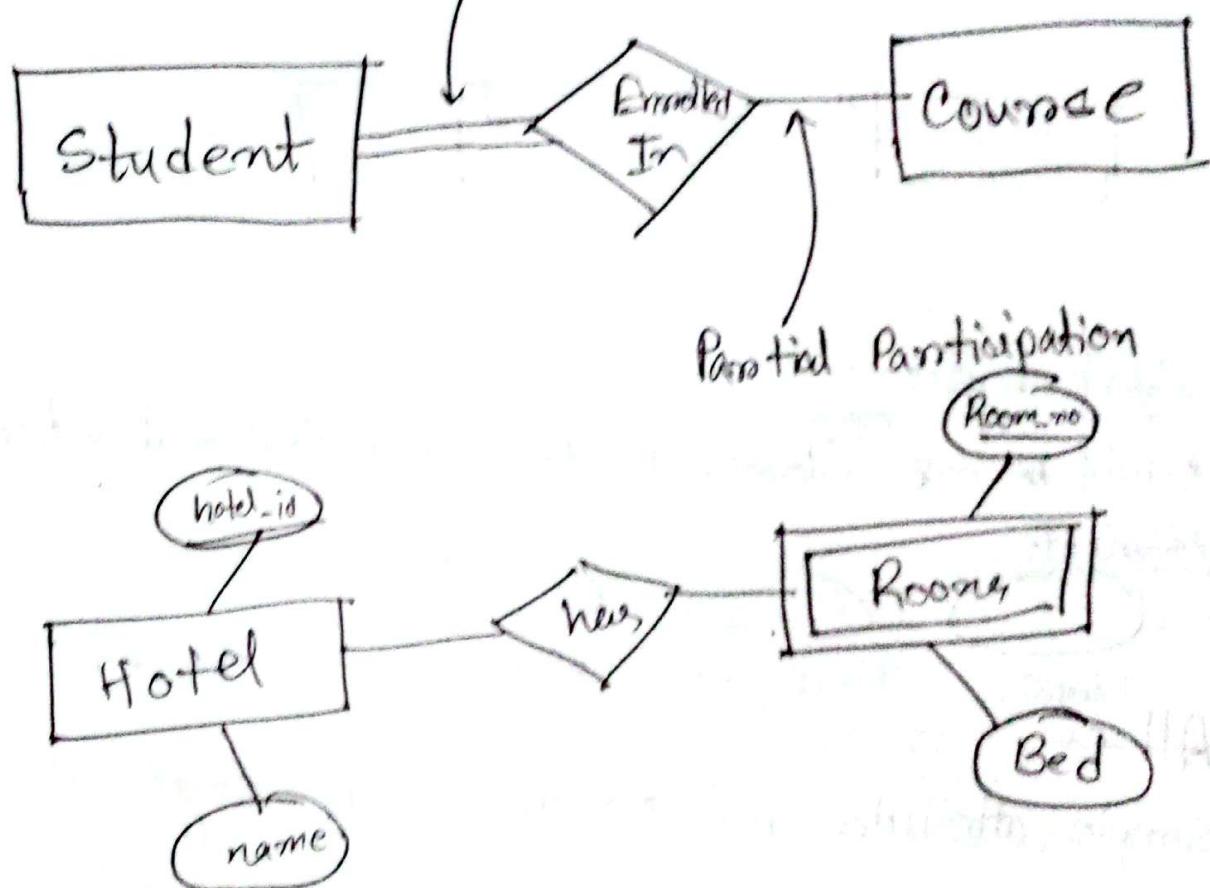
on more  
\* Composite Key: If two keys are taken to be unique

(3, 4-9)

C-5, W-2

DB

02/07/25

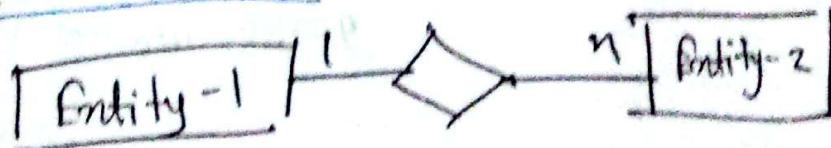


## Cardinality Constraints of Relationship

Many to Many:



One to Many:



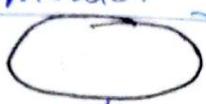
## One to One:



## Weak Entity:

\* will be one (dominant) to many (subordinate)

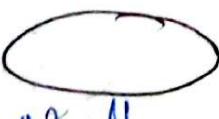
### Attribute:



Single



Multi-value



Null



Derived

### All str

simple attributes will not always be single-valued  
as single-valued attributes can give more info



ID: 20220204003

↓

single  
but not simple

2022 - 02 - 04 - 003

all individually  
gives more info

## Degree of Relationship:

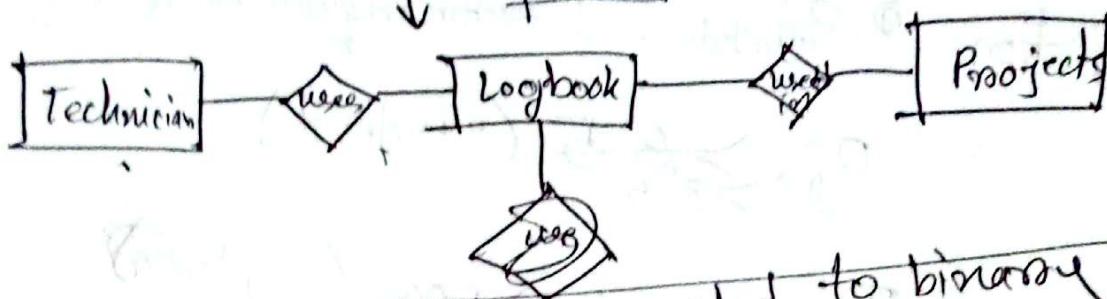
\* Binary :  $A \rightarrow B$

```

classDiagram
    class Technician
    class Logbook
    class Project

    Technician "Temporary" --> Logbook
    Technician --> Project
    Logbook <|-- Logbook
  
```

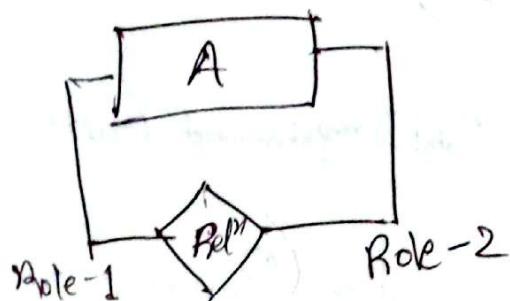
The diagram shows a UML Class Diagram. The 'Technician' class has two associations: one to 'Logbook' and one to 'Project'. A note 'Temporary' is associated with the 'Technician' class. A diamond-shaped multiplicity 'one' is placed between 'Technician' and 'Logbook', and a multiplicity 'zero or one' is placed between 'Logbook' and 'Project'. A note 'Temporary' is also associated with the 'Technician' class.



All ternary can be converted to binary

- \* Better to make binary relationship

A Unary:



N-ary: 3+

## Relational Algebra

select \* [where] query is replaced  
by,

Projection =  $\Pi_{\text{attribute}} [\square \text{ condition requirement (Table)}]$

$\Pi_{\text{gpa} \geq 3.5} (\text{student})$

and  $\Pi_{\text{age} < 25 \wedge \text{salary} > 5000} (\text{employed})$

$\Pi_{\text{name} = "man \%."} (\text{employee})$

\* for specific attributes,

Projection =  $\Pi_{\text{attr}} (\Pi_{\text{attr}} [\square \text{ requirement (table)})$

$\Pi_{\text{name}, \text{age}, \text{salary}} (\Pi_{\text{salary} > 5000} (\text{staff}))$

\* Rename =  $P_{\text{old name}} \rightarrow \text{new name} (\text{Table}) [\text{attr}]$

$P_{\text{new name}} (\text{Table}) [\text{table}]$

Quiz - 1 : 09/07/25

Syllabus : Up to now

## Binary Operation Compatibility

\* ~~UNION~~ UNION or INTERSECTION can only be done with similar data types and numbers of columns

\* U shows only once distinctly

\* CROSS JOIN = Cartesian Product (X)

↳ Cartesian Product repeats column

$$P_B \rightarrow \partial P_{R,B}(R)$$

↳ This creates new column if similar column

$$P_B \rightarrow s.B(s)$$

#  $\Pi_{\text{name}} (\sigma_{\text{Branch} = \text{"Gulshan"} \wedge \text{Borrower}.\text{Loan\_No} = \text{Loan}.\text{Loan\_No}} \text{Branch} = \text{"Gulshan"} \text{ (result\_table)})$

SELECT name FROM Borrower CROSS JOIN  
Loan WHERE Branch = "Gulshan" AND ~~Branch~~  
AND Borrower.Loan-No = Loan.Loan-No.

\* Other name of inner join is natural join ( $\bowtie$ )

\*  $\bowtie$  Inner Join

~~$\bowtie$~~  Left "

~~$\bowtie$~~  Right "

~~$\bowtie$~~  Full "

\* Another way of writing:  $\theta$  join

$\Pi_{\text{attr}} (\text{Table1} \bowtie \text{Table2})$  condition

\*  $\text{G}_{\text{F(A)}}(R)$

```

    graph TD
      GF[G_{F(A)}(R)] -- "Group By" --> Table[Table]
      GF -- "Function" --> Table
      GF -- "Aggregate Symbol" --> Table
  
```

\* Many Big queries can be summed up using,

\* Let,

$$\begin{array}{ccc} \text{Condition} = T_1 & \longleftarrow & \text{Condition 1} \\ \dashv & & \dashv \\ \dashv & \longleftarrow & \text{Condition 2} \\ & \vdots & \end{array}$$

$$\pi_{\cdot} = (\sigma_{T_1 \wedge T_2}(T_3))$$

## Database Security (theoretical)

### \* Database in different capacities:

- i) Application Support: Version control etc
- ii) Secured Storage of sensitive Information
- iii) Online Transaction Processing (OLTP): Transactions that happen online
- iv) Data warehousing: Large data stored for future analysis

### \* Security Layers:

- i) Server Level Security: Physical security of servers
- ii) Network Level Security: Encryptions
- iii) Operating System Security: OS of the db servers

## \*Types of Backup:

- i) Full Backups : Complete copy of the data
- ii) Differential Backup: Chunks of change
- iii) Transaction Log Backups: Logs are kept

Quiz - 2 : ~~Wet~~ 30/07/25

Syllabus : Upto today

## Classification of Physical Storage Media

### \* Priorities:

- i) Speed
- ii) Cost
- iii) Reliability

Storage: i) Volatile: loses data when power is off  
ii) Non-volatile: does not lose data when power is off

### \* Storage Hierarchy:

- i) Primary : Volatile → High Speed
- ii) Secondary } Non-volatile → Slow speed
- iii) Tertiary }

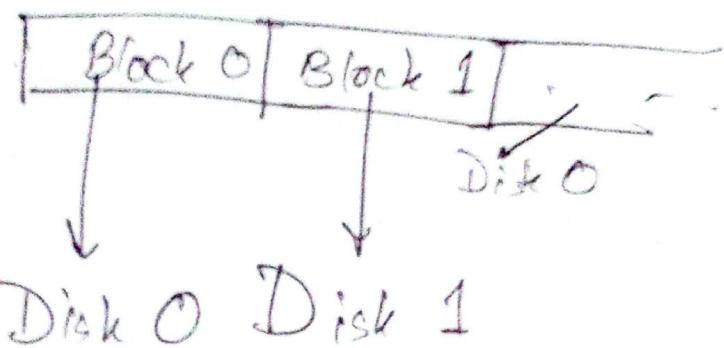
### \* Mass storage

RAID: Redundant Array of Independent Disks

→ Multiple smaller disks together to make bigger

\* RAID level 0: Striping:

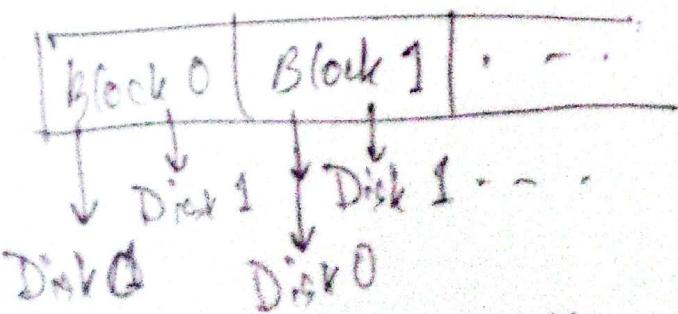
→ Data is divided and stored in different disks



→ No redundancy so no backup

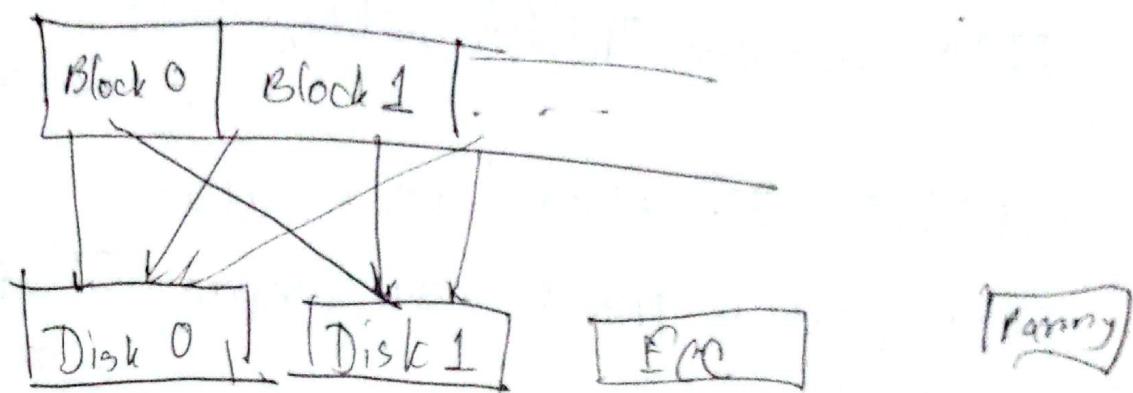
↳ ∴ No full recovery

\* RAID level 1: Data is stored in one disk but duplicated



→ same read but less write

RAID Level 2: Striping but with common connections



Implementations:  
i) Software  
ii) Hardware

Mean time between failures,  $MBTF = \frac{\sum (\text{downtime} - \text{uptime})}{\text{No. of failures}}$

## Transaction Concept

A unit of program, that possibly updates various data items

- \* **ACID**
  - has to restore/prime everything when it crashes
  - single transaction at an account at once
  - values must be equal
  - all one notification
    - ↳ may be temporarily inconsistent

### Transaction State:

- i) Active
- ii) Partially Committed
- iii) Failed
- iv) Aborted
- v) Committed

\* Concurrent Executions: Multiple transactions at once

\* Schedules: Sequence of instructions

$$T_2: A = 100$$

$$\text{TEMP} = 100 \times 0.1 = 10$$

$$A = 100 - \text{TEMP} = 90$$

$$B = 200$$

$$B = 200 + 10$$

$$B = 210$$

$$T_1: A = 90$$

$$A = 90 - 50$$

$$A = 40$$

$$B = 210$$

$$B = 210 + 50$$

$$B = 260$$

To solve this issue, we finish A's work of  $T_1$  and  $T_2$  and then B's work of  $T_1$  and  $T_2$



## Query Processing

- \* Parsing and translation
  - User queries in high level language
  - Translated expression for physical level
  - Check syntax and create parse tree
- \* Evaluation → Create Evaluation plan
- \* Optimization → Choosing the ~~best one~~ one with the least cost

Select name from Emp where salary > 10000

$\sigma_{\text{salary} > 10000}(\pi_{\text{name}}(\text{Emp}))$

$\pi_{\text{name}}(\sigma_{\text{salary} > 10000}(\text{Emp}))$