

$$2^1 - 1 = 3 \text{ prim}$$

$$2^3 - 1 = 7 \text{ prim}$$

$$2^5 - 1 = 31 \text{ prim}$$

$$2^7 - 1 = 127 \text{ prim}$$

$$2^{11} - 1 = 2047 = 23 \cdot 89$$

The CrypTool Script

Cryptography, Mathematics, and More

Prof. Bernhard Esslinger
and the CrypTool Development Team

Background reading for CrypTool
the free e-learning program
(with number theory samples for Sage)

The CrypTool Script: Cryptography, Mathematics, and More

Background reading for CrypTool
the free e-learning program
(with number theory code samples for Sage)

(11th edition – distributed with CrypTool version 1.4.31)

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and the CrypTool Team, 1998-2013
Frankfurt am Main, Germany

www.cryptool.org

February 27, 2013

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This also includes the code of the Sage samples in this document.

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Overview about the Content of the CrypTool Script

The rapid spread of the Internet has lead to intensified research in the technologies involved, especially within the area of cryptography where a good deal of new knowledge has arisen.

In this *script accompanying the CrypTool programs* you will find predominantly mathematically oriented information on using cryptographic procedures. Also included are many sample code pieces written in the computer algebra system **Sage** (see appendix A.7). The main chapters have been written by various **authors** (see appendix A.8) and are therefore independent from one another. At the end of most chapters you will find references and web links. The sections have been enriched with many *footnotes*. Within the footnotes you can see where the described functions can be called in the different CrypTool versions.

The first chapter explains the principles of symmetric and asymmetric **encryption** and definitions for their resistancy.

Because of didactic reasons the second chapter gives an exhaustive overview about **paper and pencil encryption methods**.

Big parts of this script are dedicated to the fascinating topic of **prime numbers** (chap. 3). Using numerous examples, **modular arithmetic** and **elementary number theory** (chap. 4) are introduced. The features of the **RSA procedure** are a key aspect.

By reading chapter 5 you'll gain an insight into the mathematical ideas and concepts behind **modern cryptography**.

Chapter 6 gives an overview about the status of attacks against modern **hash algorithms** and is then shortly devoted to **digital signatures**, which are an essential component of e-business applications.

Chapter 7 describes **elliptic curves**: They could be used as an alternative to RSA and in addition are extremely well suited for implementation on smartcards.

Chapter 8 describes **homomorphic crypto functions**: They are a modern research topic which got especial attention in the course of cloud computing.

The last chapter **Crypto2020** discusses threats for currently used cryptographic methods and introduces alternative research approaches to achieve long-term security of cryptographic schemes.

Whereas the CrypTool *e-learning programs* motivate and teach you how to use cryptography in practice, the *script* provides those interested in the subject with a deeper understanding of the mathematical algorithms used – trying to do it in an instructive way.

Within the **appendices** A.1, A.2, A.3, and A.4 you can gain a fast overview about the functions delivered by the different CrypTool variants via:

- the function list and the menu tree of CrypTool 1 (CT1),
- the function list and the templates in CrypTool 2 (CT2),
- the function list of JCrypTool (JCT), and
- the function list of CrypTool-Online (CTO).

The authors would like to take this opportunity to thank their colleagues in the particular companies and at the universities of Bochum, Darmstadt, Frankfurt, Gießen, Karlsruhe and Siegen.

As with the e-learning program CrypTool, the quality of the script is enhanced by your suggestions and ideas for improvement. We look forward to your feedback.

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Preface to the 11th Edition of the CrypTool Script

Starting in the year 2000 this script became part of the CrypTool 1 (CT1) package. The script's goal is to explain some mathematical topics in more detail, but still in a way which is easy to understand.

Topics from mathematics and cryptography have been meaningfully split up and for each topic an extra chapter has been written which can be read on its own. This enables developers/authors to contribute independently of each other. Naturally there are many more interesting topics in cryptography which could be discussed in greater depth – therefore this selection is only one of many possible ways.

The later editorial work in TeX added footnotes and cross linkages between different sections.

This edition completely updated the TeX sources of the document, and of course the content of the script was amended, corrected, and updated with some topics, e.g.:

- the definitions of the strength of security functions (chap. 1.1),
- the list of movies or novels, in which cryptography or number theory played a major role (see appendix A.5),
- the overviews of all functions in for CrypTool 2 (CT2), for JCrypTool (JCT), and for CrypTool-Online (CTO),
- further Sage scripts for cryptography, and the appendix A.7 about using the computer algebra system Sage,
- the section about the Goldbach conjecture (see 3.8.3),
- the section about shared primes in RSA modules used in reality (see ??),xxxxxxxxxxxxxxxxxxxxx■
- the section about RSA fixed points (see 4.19.7), and
- the chapter about homomorphic encryption (see 8).

The script was delivered the first time with the CT1 package in version 1.2.01. Since then it has been expanded and revised in almost every new version of CT1.

In the meantime the CT project gets feedback and testimonials from almost all countries of the planet.

Acknowledgment

At this point I'd like to thank explicitly the following people who particularly contributed to the CrypTool project. They applied their very special talents and showed really great engagement:

- Mr. Henrik Koy
- Mr. Jörg-Cornelius Schneider
- Mr. Florian Marchal
- Dr. Peer Wichmann
- Hr. Dominik Schadow
- Staff of Prof. Johannes Buchmann, Prof. Claudia Eckert, Prof. Alexander May, Prof. Torben Weis and especially Prof. Arno Wacker.

Also I want to thank all the many people not mentioned here for their hard work (mostly carried out in their spare time).

Thanks also to the readers who sent us feedback.

I hope that many readers have fun with this script and that they get out of it more interest and greater understanding of this modern but also very ancient topic.

Bernhard Esslinger

Frankfurt and Siegen (Germany), January 2013

PS:

We'd be glad if further authors would show up to add further chapters, e.g. about the Riemann Zeta function, hash functions, random numbers, or the design and attack of crypto protocols (like SSL).

Introduction – How do the Script and the Program Play together?

This CrypTool script

This document is delivered together with the open-source programs of the CrypTool project.

The articles in this script are largely self-contained and can also be read independently of the CrypTool programs.

Chapters 5 (Modern Cryptography), 7 (Elliptic Curves), and 8 (Homomorphic Encryption) require a deeper knowledge in mathematics, while the other chapters should be understandable with a school leaving certificate.

The authors have attempted to describe cryptography for a broad audience – without being mathematically incorrect. We believe that this didactic pretension is the best way to promote the awareness for IT security and the readiness to use standardized modern cryptography.

The programs **CrypTool 1**, **CrypTool 2** and **JCrypTool**

CrypTool 1 (CT1) is an educational program with a comprehensive online help enabling you to use and analyse cryptographic procedures within a unified graphical user interface. The online help in CrypTool 1 contains both instructions how to use the program and explanations about the methods itself (both not as detailed and in another structure than the CT script).

CrypTool 1 and the successor versions CrypTool 2 (CT2) and JCrypTool (JCT) are used world-wide for training in companies and teaching at schools and universities.

CrypTool-Online

Also part of the CT project is the website CrypTool-Online (CTO) (<http://www.cryptool-online.org>), where you can check and apply cryptographic methods within a browser or on a smartphone. The scope of CTO is far below from the standalone programs CT1, CT2 and JCT.

MTC3

MTC3 is the abbreviation for MysteryTwister C3, an international cryptography contest (<http://www.mysterytwisterc3.org>), which is also based on the CT project. Here you can find cryptographic riddles in four categories, a high-score list and a moderated forum. In the meantime more than 3500 active users participate, and more than 140 challenges are offered (currently 89 of them are solved by at least one user).

The Computer Algebra System Sage

Sage is a comprehensive open-source CAS package which can be used to easily program the mathematical methods explained in this script. A speciality of this CAS is, that the scripting language is Python.

Sage becomes more and more the standard open-source CAS system at universities.

The Pupil's Crypto Courses

Within this initiative, one and two 2 day courses in cryptology are offered for pupils and teachers in order to show how attractive MINT subjects like mathematics, computer science and especially cryptology are. The course agenda is a virtual secret agent story.

In the meantime, these courses took place for several years in Germany in different towns.

All course material is freely available on net (<http://www.cryptoool.org/schuelerkrypto/>).

All software used is free software (using mostly CT1 and CT2).

As all course material is currently only available in German we'd be happy if someone could translate the course material and build an according course in English.

Acknowledgment

I am deeply grateful to all the people helping with their impressive commitment who have made this global project so successful.

Bernhard Esslinger

Frankfurt and Siegen (Germany), January 2013

Chapter 1

Security Definitions and Encryption Procedures

(Bernhard Esslinger, Joerg-Cornelius Schneider, May 1999; Updates Dec. 2001, Feb. 2003, June 2005, July 2007, Jan. 2010, Dec. 2012)

This chapter introduces the topic in a more descriptive way without using too much mathematics.

The purpose of encryption is to change data in such a way that only an authorized recipient is able to reconstruct the plaintext. This allows us to transmit data without worrying about it getting into unauthorized hands. Authorized recipients possess a piece of secret information – called the key – which allows them to decrypt the data while it remains hidden from everyone else.

For explanations in the following we use the notation from Figure 1.1:

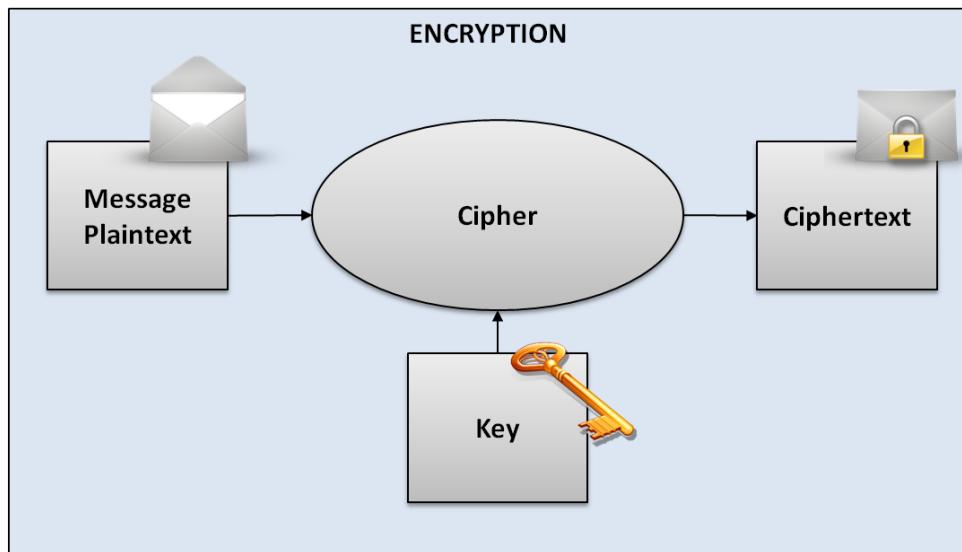


Figure 1.1: Common notations when using ciphers

First we present the ideas how the security of cryptosystems is defined.

1.1 Security definitions

Explain it to me, I will forget it. Show it to me, maybe I will remember it. Let me do it, and I will be good at it.
--

Quote 1: Saying from India

Modern cryptography is heavily based on mathematical theory and computer science practice. Cryptographic algorithms are designed around computational hardness assumptions, making such algorithms hard to break in practice by any adversary.

Depending on the adversary's capabilities there are mainly two basic notations of security distinguished in literature (see e.g. *Contemporary Cryptography* [Oppliger2011]):

- **Computational, conditional or practical security**

A cipher is computationally secure if it is theoretically possible to break such a system but it is infeasible to do so by any known practical means. Theoretical advances (e.g., improvements in integer factorization algorithms) and faster computing technology require these solutions to be continually adapted.

Even using the best known algorithm for breaking it will require so much resources (e.g., 1,000,000 years) that practically the cryptosystem is secure.

So this concept is based on assumptions of the adversary's limited computing power and the current state of science.

- **Information-theoretical or unconditional security**

A cipher is unconditionally secure if it is secure no matter how much resources (time, space) the attacker has. So even in the case where an adversary has unlimited resources for solving a problem e.g. breaking a cipher. An adversary has not sufficient information for doing that because he is unable to gain any meaningful data from a ciphertext.

There exist information-theoretically secure schemes that provably cannot be broken even with unlimited computing power – an example is the *One Time Pad* (OTP).

As the OTP is information-theoretically secure it derives its security solely from information theory and is secure even with unlimited computing power at the adversary's disposal. However, OTP has several practical disadvantages (the key used must be used only once, randomly selected and must be at least as long as the message being protected), which means that it is hardly used except in closed environments such as for the hot wire between Moscow and Washington.

Two more concepts are sometimes used:

- **Provable security** This means that breaking such a cryptographic system is as difficult as solving some supposedly difficult problem e.g. discrete logarithm computation, discrete square root computation, very large integer factorization.

Example: Currently we know that RSA is at most as difficult as factorization, but we cannot prove that its exactly as difficult as factorization. So RSA has no proven minimum security. Or in other words: We cannot prove, that if RSA (the cryptosystem) is broken, that then factorization (the hard mathematical problem) can be solved.

The Rabin cryptosystem was the first cryptosystem which could be proven to be computationally equivalent to a hard problem.

- **Ad-hoc security** A cryptographic system has this security feature if it is not worth to try to break into such a system because of inadequate price of data with comparison to price of work needed to do so. Or an attack can't be done in sufficiently short time (see *Handbook of Applied Cryptography* [Menezes2001]).

Example: This may apply if a message relevant for the stock market will be published tomorrow and you need a year to break it.

For good procedures used today the time taken to break them is so long that it is practically impossible to do so, and these procedures can therefore be considered (practically) secure – from a pure algorithm's point of view.

We basically distinguish between symmetric and asymmetric encryption procedures. The book of Bruce Schneier [Schneier1996] offers a good overview of the different encryption algorithms.

1.2 Symmetric encryption¹

“Transparency. That’s the best one can hope for in a technologically advanced society ... otherwise you will just be manipulated.”

Quote 2: Daniel Suarez²

For *symmetric* encryption sender and recipient must be in possession of a common (secret) key which they have exchanged before actually starting to communicate. The sender uses this key to encrypt the message and the recipient uses it to decrypt it.

All classical methods are of this type. Examples can be found within the CT programs, in chapter 2 (“Paper and Pencil Encryption Methods”) of this script, or in [Nichols1996]. In this chapter however, we want to consider only modern symmetric mechanisms.

The advantages of symmetric algorithms are the high speed with which data can be encrypted and decrypted. One disadvantage is the need for key management. In order to communicate with one another confidentially, sender and recipient must have exchanged a key using a secure channel before actually starting to communicate. Spontaneous communication between individuals who have never met therefore seems virtually impossible. If everyone wants to communicate with everyone else spontaneously at any time in a network of n subscribers, each

¹With CrypTool 1 (**CT1**) you can execute the following modern symmetric encryption algorithms (using the menu path **Crypt \ Symmetric (modern)**):

IDEA, RC2, RC4, DES (ECB), DES (CBC), Triple-DES (ECB), Triple-DES (CBC), MARS (AES candidate), RC6 (AES candidate), Serpent (AES candidate), Twofish (AES candidate), Rijndael (official AES algorithm).

With CrypTool 2 (**CT2**) you can execute the following modern symmetric encryption algorithms (using in the Startcenter **Templates \ Cryptography \ Modern \ Symmetric**):

AES, DES, PRESENT, RC2, RC4, SDES, TEA, Triple-DES, Twofish.

In JCrypTool (**JCT**) you can execute the following modern symmetric encryption algorithms:

AES, Rijndael, Camellia, DES, Dragon, IDEA, LFSR, MARS, Misty1, RC2, RC5, RC6, SAFER++, SAFER++, Serpent, Shacal, Shacal2, Twofish.

²Daniel Suarez, “Freedom”, Dutton Adult, 2010, Chapter 5, “Getting with the Program”, p. 63, Price.

subscriber must have previously exchanged a key with each of the other $n - 1$ subscribers. A total of $n(n - 1)/2$ keys must therefore be exchanged.

The most well-known symmetric encryption procedure is the DES algorithm. The DES algorithm has been developed by IBM in collaboration with the National Security Agency (NSA), and was published as a standard in 1975. Despite the fact that the procedure is relatively old, no effective attack on it has yet been detected. The most effective way of attacking consists of testing (almost) all possible keys until the right one is found (*brute-force-attack*). Due to the relatively short key length of effectively 56 bits (64 bits, which however include 8 parity bits), numerous messages encrypted using DES have in the past been broken. Therefore, the procedure can not be considered secure any longer. Alternatives to the DES procedure include IDEA, Triple-DES and especially AES.

Up-to-the-minute procedure is the symmetric AES standard. The associated Rijndael algorithm was declared winner of the AES award on October 2nd, 2000 and thus succeeds the DES procedure.

More details and further references about the AES algorithms and the AES candidates of the last round can be found i.e. within the online help of CrypTool.³

1.2.1 Results about cryptanalysis of AES

Below you will find some results, which have recently called into question the security of the AES algorithm – from our point of view these doubts practically still remain unfounded . The following information is based particularly on the original papers and the articles [Wobst-iX2002] and [Lucks-DuD2002].

AES with a minimum key length of 128 bit is still in the long run sufficiently secure against brute-force attacks – as long as the quantum computers aren't powerful enough. When announced as new standard AES was immune against all known cryptanalytic attacks, mostly based on statistical considerations and earlier applied to DES: using pairs of clear and cipher texts expressions are constructed, which are not completely at random, so they allow conclusions to the used keys. These attacks required unrealistically large amounts of intercepted data.

Cryptanalysts already label methods as “academic success” or as “cryptanalytic attack” if they are theoretically faster than the complete testing of all keys (brute force analysis). In the case of AES with the maximal key length (256 bit) exhaustive key search on average needs 2^{255} encryption operations. A cryptanalytic attack needs to be better than this. At present between 2^{75} and 2^{90} encryption operations are estimated to be performable only just for organizations, for example a security agency.

In their 2001-paper Ferguson, Schroeppel and Whiting [Ferguson2001] presented a new method of symmetric codes cryptanalysis: They described AES with a closed formula (in the form of a continued fraction) which was possible because of the ”relatively” clear structure of AES. This formula consists of around 1000 trillion terms of a sum - so it does not help concrete practical cryptanalysis. Nevertheless curiosity in the academic community was awakened. It was already known, that the 128-bit AES could be described as an over-determined system of about 8000 quadratic equations (over an algebraic number field) with about 1600 variables (some of them are the bits of the wanted key) – equation systems of that size are in practice not solvable. This special equation system is relatively sparse, so only very few of the quadratic

³CrypTool 1 online help: The index head-word **AES** leads to the 3 help pages: **AES candidates**, **The AES winner Rijndael** and **The Rijndael encryption algorithm**

A comprehensive description of AES including C code can be found in [Haan2008].

terms (there are about 1,280,000 are possible quadratic terms in total) appear in the equation system.

The mathematicians Courtois and Pieprzyk [Courtois2002] published a paper in 2002, which got a great deal of attention amongst the cryptology community: The pair had further developed the XL-method (eXtended Linearization), introduced at Eurocrypt 2000 by Shamir et al., to create the so called XSL-method (eXtended Sparse Linearization). The XL-method is a heuristic technique, which in some cases manages to solve big non-linear equation systems and which was till then used to analyze an asymmetric algorithm (HFE). The innovation of Courtois and Pieprzyk was, to apply the XL-method on symmetric codes: the XSL-method can be applied to very specific equation systems. A 256-bit AES could be attacked in roughly 2^{230} steps. This is still a purely academic attack, but also a direction pointer for a complete class of block ciphers. The major problem with this attack is that until now nobody has worked out, under what conditions it is successful: the authors specify in their paper necessary conditions, but it is not known, which conditions are sufficient. There are two very new aspects of this attack: firstly this attack is not based on statistics but on algebra. So attacks seem to be possible, where only very small amounts of ciphertext are available. Secondly the security of a product algorithm⁴ does not exponentially increase with the number of rounds.

Currently there is a large amount of research in this area: for example Murphy and Robshaw presented a paper at Crypto 2002 [Robshaw2002a], which could dramatically improve cryptanalysis: the burden for a 128-bit key was estimated at about 2^{100} steps by describing AES as a special case of an algorithm called BES (Big Encryption System), which has an especially "round" structure. But even 2^{100} steps are beyond what is achievable in the foreseeable future. Using a 256 bit key the authors estimate that a XSL-attack will require 2^{200} operations.

More details can be found in the Web links section at "AES or Rijndael cryptosystem".

So for AES-256 the attack is much more effective than brute-force but still far more away from any computing power which could be accessible in the short-to-long term.

The discussion is very controversial at the moment: Don Coppersmith (one of the inventors of DES) for example queries the practicability of the attack because XLS would provide no solution for AES [Coppersmith2002]. This implies that then the optimization of Murphy and Robshaw [Robshaw2002b] would not work.

In 2009 Biryukov und Khovratovich [Biryukov2009] published another theoretical attack on AES. This attack uses different methods from the ones described above. They applied methods from hash function cryptanalysis (local collisions and boomerang switching) to construct a related-key attack on AES-256. I. e. the attacker not only needs to be able to encrypt arbitrary data (chosen plain text), in addition he needs to be able to manipulate the unknown key (related-key).

Based on those assumptions, the effort to find a AES-256 key is reduced to 2^{119} time and 2^{77} memory (considering asymmetric complexity). In the case of AES-192 the attack is even less practical, for AES-128 the authors do not provide an attack.

⁴A ciphertext can be used as input for another encryption algorithm. A cascade cipher is build up as a composition of different encryption transformations. The overall cipher is called product algorithm or cascade cipher (sometimes depending whether the used keys are statistically dependent or not).

Cascading does not always improve the security.

This process is also used *within* modern algorithms: They usually combine simple and, considered at its own, cryptologically relatively insecure single steps in several rounds into an efficient overall procedure. Most block ciphers (e.g. DES, IDEA) are cascade ciphers.

Also serial usage of the same cipher with different keys (like with Triple-DES) is called cascade cipher.

1.2.2 Current status of brute-force attacks on symmetric algorithms

The current status of brute-force attacks on symmetric encryption algorithms can be explained with the block cipher RC5.

Brute-force (exhaustive search, trial-and-error) means to completely examine all keys of the key space: so no special analysis methods have to be used. Instead, the ciphertext is decrypted with all possible keys⁵ and for each resulting text it is checked, whether this is a meaningful clear text⁶. A key length of 64 bit means at most $2^{64} = 18,446,744,073,709,551,616$ or about 18 trillion (GB) / 18 quintillion (US) keys to check.

Companies like RSA Security provided so-called cipher challenges in order to quantify the security offered by well-known symmetric ciphers as DES, Triple-DES or RC5.⁷ They offered prizes for those who managed to decipher ciphertexts, encrypted with different algorithms and different key lengths, and to unveil the symmetric key (under controlled conditions). So theoretical estimates can be confirmed.

It is well-known, that the “old” standard algorithm DES with a fixed key length of 56 bit is no more secure: This was demonstrated already in January 1999 by the Electronic Frontier Foundation (EFF). With their specialized computer Deep Crack they cracked a DES encrypted message within less than a day.⁸

The current record for strong symmetric algorithms unveiled a key 64 bit long. The algorithm used was RC5, a block cipher with variable key size.

The RC5-64 challenge has been solved by the distributed.net team after 5 years.⁹ In total 331,252 individuals co-operated over the internet to find the key.¹⁰ More than 15 trillion (GB) / 15 quintillion (US) keys were checked, until they found the right key.

This makes clear, that symmetric algorithms (even if they have no cryptographic weakness) using keys of size 64 bit are no more appropriate to keep sensible data private.

There are also cipher challenges¹¹ for asymmetric algorithms (please see chapter 4.11.4).

⁵With CT1 you can also perform brute-force attacks of modern symmetric algorithms (using the menu path **Analysis \ Symmetric Encryption (modern)**): Here the weakest knowledge of an attacker is assumed, he performs a ciphertext-only attack.

With CT2 you can also perform brute-force attacks (using the templates under **Cryptanalysis \ Modern**). Highly powerful is the KeySearcher component, which can be used to distribute the calculations to many different computers.

⁶If the cleartext is written in a natural language and at least 100 B long, this check also can be performed automatically.

To achieve a result in an appropriate time with a single PC you should mark not more than 24 bit of the key as unknown.

⁷<http://www.rsasecurity.com/rsalabs/challenges/secretkey/index.html>

Unfortunately, in May 2007 RSA Inc announced that they will not confirm the correctness of the not yet solved RC5-72 challenge.

⁸<http://www.rsasecurity.com/rsalabs/challenges/des3/index.html>

⁹<http://distributed.net/pressroom/news-20020926.html>

¹⁰An overview of current distributed computing projects can be found here:

<http://distributedcomputing.info/>

¹¹A wide spectrum of simple and complex, symmetric and asymmetric crypto riddles are included in the “MysteryTwister C3” challenge: <http://www.mysterytwisterc3.org>.

1.3 Asymmetric encryption¹²

In the case of *asymmetric* encryption each subscriber has a personal pair of keys consisting of a *secret* key and a *public* key. The public key, as its name implies, is made public, e.g. in a key directory on the Internet (this kind of “bill-board” is also called just directory or public key ring).

If Alice¹³ wants to communicate with Bob, then she finds Bob’s public key in the directory and uses it to encrypt her message to him. She then sends this ciphertext to Bob, who is then able to decrypt it again using his secret key. As only Bob knows his secret key, only he can decrypt messages addressed to him. Even Alice who sends the message cannot restore plaintext from the (encrypted) message she has sent. Of course, you must first ensure that the public key cannot be used to derive the private key.

Such a procedure can be demonstrated using a series of thief-proof letter boxes. If I have composed a message, I then look for the letter box of the recipient and post the letter through it. After that, I can no longer read or change the message myself, because only the legitimate recipient has the key for the letter box.

The advantage of asymmetric procedures is the easier key management. Let’s look again at a network with n subscribers. In order to ensure that each subscriber can establish an encrypted connection to each other subscriber, each subscriber must possess a pair of keys. We therefore need $2n$ keys or n pairs of keys. Furthermore, no secure channel is needed before messages are transmitted, because all the information required in order to communicate confidentially can be sent openly. In this case, you simply¹⁴ have to pay attention to the accuracy (integrity and authenticity) of the public key. Disadvantage: Pure asymmetric procedures take a lot longer to perform than symmetric ones.

The most well-known asymmetric procedure is the RSA algorithm¹⁵, named after its developers Ronald Rivest, Adi Shamir and Leonard Adleman. The RSA algorithm was published in 1978. The concept of asymmetric encryption was first introduced by Whitfield Diffie and Martin Hellman in 1976. Today, the ElGamal procedures also play a decisive role, particularly the Schnorr variant in the DSA (Digital Signature Algorithm).

¹²With CT1 you can execute RSA encryption and decryption (using the menu path **Crypt \ Asymmetric**). In both cases you must select a RSA key pair. Only in the case of decryption the secret RSA key is necessary – so here you are asked to enter the PIN.

With CT2 you can also perform asymmetric methods (using the templates under **Cryptography \ Modern**). JCT offers asymmetric methods like RSA both within the **Visuals** menu of the Default Perspective as well as within the Algorithm Perspective.

¹³In order to describe cryptographic protocols participants are often named Alice, Bob, ... (see [Schneier1996, p. 23]). Alice and Bob perform all 2-person-protocols. Alice will initiate all protocols and Bob answers. The attackers are named Eve (eavesdropper) and Mallory (malicious active attacker).

¹⁴That this is also not trivial is explained e.g. in xxxx . Besides the requirements for the key generation it has to be considered that nowadays also (public-key) infrastructures are targets of attacks.

¹⁵The RSA algorithm is extensively described in chapter 4.10 and later within this script. The topical research results concerning RSA are described in chapter 4.11.

The RSA cryptosystem can be executed in many variations with CT1 (using the menu path **Individual Procedures \ RSA Cryptosystem \ RSA Demonstration**).

1.4 Hybrid procedures¹⁶

In order to benefit from the advantages of symmetric and asymmetric techniques together, hybrid procedures are usually used (for encryption) in practice.

In this case the bulk data is encrypted using symmetric procedures: The key used for this is a secret session key generated by the sender randomly¹⁷ that is only used for this message.

This session key is then encrypted using the asymmetric procedure, and transmitted to the recipient together with the message.

Recipients can determine the session key using their private keys and then use the session key to encrypt the message.

In this way, we can benefit from the easy key management of asymmetric procedures (using public/private keys) and we benefit from the efficiency of symmetric procedures to encrypt large quantities of data (using secret keys).

¹⁶Within CT1 you can find this technique using the menu path **Crypt \ Hybrid**: There you can follow the single steps and its dependencies with concrete numbers. The variant with RSA as the asymmetric algorithm is graphically visualized; the variant with ECC uses the standard dialogs. In both cases AES is used as the symmetric algorithm.

JCT offers hybrid methods like ECIES within the Algorithm Perspective under **Algorithms \ Hybrid Ciphers**.

¹⁷An important part of cryptographically secure techniques is to generate random numbers. Within CT1 you can check out different random number generators using the menu path **Indiv. Procedures \ Generate Random Numbers**. Using the menu path **Analysis \ Analyze Randomness** you can apply different test methods for random data to binary documents.

CT1 concentrates on cryptographically strong **pseudo** random number generators. “Real” random sources are only included via calling the integrated Secude generator.

JCT offers **pseudo** random number generators both within the Default Perspective in the menu **Algorithms \ Random Number Generator** as well as within the Algorithm Perspective.

There is an old saying inside the US National Security Agency (NSA):
”Attacks always get better; they never get worse.”

Quote 3: IETF¹⁸

1.5 Cryptanalysis and symmetric ciphers for educational purposes¹⁹

Compared to public-key ciphers based on mathematics like RSA, the structure of AES and most other modern symmetric ciphers (like DES, IDEA or Present), is very complex and cannot be explained as easily as RSA.

So simplified variants of modern symmetric ciphers were developed for educational purposes in order to allow beginners to perform encryption and decryption by hand and gain a better understanding of how the algorithms work in detail. These simplified variants also help to understand and apply the according cryptanalysis methods.

One of these variants is e.g. S-AES (Simplified-AES) by Prof. Ed Schaefer and his students [Musa-etal2003].²⁰ Another one is Mini-AES [Phan2002] (see chapter 1.7.1 “Mini-AES”):

- Edward F. Schaefer: *A Simplified Data Encryption Standard Algorithm* [Schaefer1996].
- Raphael Chung-Wei Phan: *Mini Advanced Encryption Standard (Mini-AES): A Testbed for Cryptanalysis Students* [Phan2002].
- Raphael Chung-Wei Phan: *Impossible differential cryptanalysis of Mini-AES* [Phan2003].
- Mohammad A. Musa, Edward F. Schaefer, Stephen Wedig: *A simplified AES algorithm and its linear and differential cryptanalyses* [Musa-etal2003].
- Nick Hoffman: *A SIMPLIFIED IDEA ALGORITHM* [Hoffman2006].
- S. Davod. Mansoori, H. Khaleghei Bizaki: *On the vulnerability of Simplified AES Algorithm Against Linear Cryptanalysis* [Mansoori-etal2007].

1.6 Further information

Beside the information you can find in the following chapters, in many other books and on a good number of websites, the online help of all CrypTool variants also offer very many details about the symmetric and asymmetric encryption methods.

¹⁸<http://tools.ietf.org/html/rfc4270>

¹⁹A good starting point to learn cryptanalysis is the book from Mark Stamp [Stamp2007].

Very high-level and concentrating on analysing symmetric block ciphers only is the article from von Bruce Schneier [Schneier2000].

Several of the cipher challenges at “MysteryTwister C3” (<http://www.mysterytwisterc3.org>) are also well suitable for educational purposes.

²⁰See the two articles with the same title “Devising a Better Way to Teach and Learn the Advanced Encryption Standard” at the university’s news at

- <http://math.scu.edu/~eschaefe/getfile.pdf>
- <http://www.scu.edu/cas/research/cryptography.cfm>

1.7 Appendix: Examples using Sage

Below is Sage source code related to contents of the chapter 1 (“Security Definitions and Encryption Procedures”).

Further details concerning cryptosystems within Sage (e.g. S-DES) can be found e.g. in the thesis of Minh Van Nguyen [Nguyen2009].

1.7.1 Mini-AES

The Sage module `crypto/block_cipher/miniaes.py` supports Mini-AES to allow students to explore the inner working of a modern block cipher.

Mini-AES, originally described at [Phan2002], is a simplified variant of the Advanced Encryption Standard (AES) to be used for cryptography education.

How to use Mini-AES is exhaustively described at the Sage reference page:

http://www.sagemath.org/doc/reference/sage/crypto/block_cipher/miniaes.html.

The following Sage code from the release tour of Sage 4.1²¹ calls the implementation of the Mini-AES.

Further example code for Mini-AES can be found in [Nguyen2009, chap. 6.5 and appendix D].

1.7.2 Further symmetric crypto algorithms in Sage

The Sage Reference v5.4.1²² lists the following further ciphers:

- Linear feedback shift register (LFSR),
- ShrinkingGeneratorCryptosystem,
- Blum-Blum-Shub (BBS) pseudorandom bit generator,
- Lattice related functions.

²¹See <http://mvngu.wordpress.com/2009/07/12/sage-4-1-released/>

²²See <http://www.sagemath.org/doc/reference/sage/crypto/>

Sage sample 1.1 Encryption and decryption with Mini-AES

```
# We can encrypt a plaintext using Mini-AES as follows:
sage: from sage.crypto.block_cipher.miniaes import MiniAES
sage: maes = MiniAES()
sage: K = FiniteField(16, "x")
sage: MS = MatrixSpace(K, 2, 2)
sage: P = MS([K("x^3 + x"), K("x^2 + 1"), K("x^2 + x"), K("x^3 + x^2")]); P
[x^3 + x  x^2 + 1]
[x^2 + x  x^3 + x^2]
sage: key = MS([K("x^3 + x^2"), K("x^3 + x"), K("x^3 + x^2 + x"), K("x^2 + x + 1")]); key
[x^3 + x^2  x^3 + x]
[x^3 + x^2 + x  x^2 + x + 1]
sage: C = maes.encrypt(P, key); C
[x  x^2 + x]
[x^3 + x^2 + x  x^3 + x]

# Here is the decryption process:
sage: plaintext = maes.decrypt(C, key)
sage: plaintext == P
True

# We can also work directly with binary strings:
sage: from sage.crypto.block_cipher.miniaes import MiniAES
sage: maes = MiniAES()
sage: bin = BinaryStrings()
sage: key = bin.encoding("KE"); key
0100101101000101
sage: P = bin.encoding("Encrypt this secret message!")
sage: C = maes(P, key, algorithm="encrypt")
sage: plaintext = maes(C, key, algorithm="decrypt")
sage: plaintext == P
True

# Or work with integers n such that 0 <= n <= 15:
sage: from sage.crypto.block_cipher.miniaes import MiniAES
sage: maes = MiniAES()
sage: P = [n for n in xrange(16)]; P
[0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15]
sage: key = [2, 3, 11, 0]; key
[2, 3, 11, 0]
sage: P = maes.integer_to_binary(P)
sage: key = maes.integer_to_binary(key)
sage: C = maes(P, key, algorithm="encrypt")
sage: plaintext = maes(C, key, algorithm="decrypt")
sage: plaintext == P
True
```

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Web Links

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<http://www.cryptosystem.net/aes>
2. AES discussion groups at NIST
<http://csrc.nist.gov/archive/aes/>
3. distributed.net: “RC5-64 has been solved”
http://www.distributed.net/Pressroom_press-rc5-64
4. RSA: “The RSA Secret Key Challenge”
<http://www.rsa.com/rsalabs/node.asp?id=2100>
5. RSA: “DES Challenge”
<http://www.rsa.com/rsalabs/node.asp?id=2108>
6. Further links can be found at the CrypTool homepage
<http://www.cryptool.org>

All links have been confirmed at December 12, 2012.

Chapter 2

Paper and Pencil Encryption Methods

(Christine Stötzel, April 2004; Updates: B.+C. Esslinger, June 2005; Updates Minh Van Nguyen and B. Esslinger, Nov. 2009, June 2010, Jan. 2013)

Few persons can be made to believe that it is not quite an easy thing to invent a method of secret writing which shall baffle investigation. Yet it may be roundly asserted that human ingenuity cannot concoct a cipher which human ingenuity cannot resolve.

Quote 4: Edgar Allan Poe: A Few Words on Secret Writing, 1841

The following chapter provides a broad overview of paper and pencil methods¹ each with references to deeper information. All techniques that people can apply manually to en- and decipher a message are embraced by this term. These methods were and still are especially popular with secret services, as a writing pad and a pencil – in contrast to electronic aids – are totally unsuspicious.

The first paper and pencil methods already arose about 3000 years ago, but new procedures were developed during the past century, too. All paper and pencil methods are a matter of symmetric methods. Even the earliest encryption algorithms use the basic principles such as transposition, substitution, block construction and their combinations. Hence it is worthwhile to closely consider this “ancient” methods especially under didactic aspects.

Methods to be successful and wide-spread had to fulfill some attributes which are equally required for modern algorithms:

- Exhaustive description, almost standardization (including special cases, padding, etc.).
- Good balance between security and usability (because methods being too complicated were error-prone or unacceptably slow).

¹The footnotes of this chapter describe how these methods can be performed using the offline programs CrypTool 1 (**CT1**), CrypTool 2 (**CT2**) and JCrypTool (**JCT**). See the appendices A.1, A.2, and A.3.

Many of the methods can also be performed within a browser, e.g. on the website CrypTool-Online (**CTO**) (<http://www.cryptool-online.org>). See the appendix A.4.

Additionally the last sub chapter (2.5) contains example code for the computer-algebra system **Sage**.

2.1 Transposition ciphers

Encrypting a message by means of transposition does not change the original characters of this message, only their order is modified (transposition = exchange).²

2.1.1 Introductory samples of different transposition ciphers

- **Rail Fence cipher**³ [Singh1999]: The characters of a message are alternately written in two (or more) lines, creating a zigzag pattern. The resulting ciphertext is read out line by line.

This is more a children's method.

See table 2.1.

Plaintext⁴: an example of transposition

n	x	m	l	o	t	a	s	o	i	i	n
a	e	a	p	e	f	r	n	p	s	t	o

Table 2.1: Rail Fence cipher

Ciphertext⁵: NXMO TASOI INAEA PEFRN PSTO

- **Scytale**⁶ [Singh1999]: This method was probably used since 600 B.C. – a description of how it operated is not known from before Plutarch (50-120 B.C.).
A long strip of paper is wrapped around a wooden cylinder and then the message is written along the length of this strip. The ciphertext is produced by unwinding the strip.
- **Grille cipher** [Goebel2003]: Both parties use identical stencils. Line by line, their holes are filled with plaintext that is read out column by column to produce the ciphertext. If there is plaintext left, the procedure is repeated.⁷

²Sometimes, the name permutation is used to describe how characters, groups of characters or columns of the plaintext are exchanged, e.g. $(1, 2, 3, 4, 5) \Leftrightarrow (3, 4, 2, 1, 5)$.

³This method can directly be found in CT1 at the menu item **Encrypt/Decrypt \ Symmetric (classic) \ Scytale / Rail Fence**. You can simulate this method also under the menu **Encrypt/Decrypt \ Symmetric (classic) \ Permutation**: For a Rail Fence with 2 lines use as key "B,A" and accept the default settings (only one permutation, where your input is done line-by-line and the output is taken column-by-column). Using the key "A,B" would start the zigzag pattern below in the way, that the first letter is written into the first line instead of the second line.

⁴: If the alphabet only uses 26 letters, we write from now onwards the plaintext in small letters and the ciphertext in capital letters.

⁵The letters of the cleartext are – as used historically – grouped within blocks of 5 letters. It does not matter if the (constant) block length is different or no blank is inserted.

⁶This method can directly be found in CT1 at the menu item **Encrypt/Decrypt \ Symmetric (classic) \ Scytale / Rail Fence**. As this method is a special case of a simple columnar transposition, you also can simulate it in CT1 under the menu **Encrypt/Decrypt \ Symmetric (classic) \ Permutation**: For the Scytale within the dialog box only the first permutation is used. If the wood has e.g. 4 angles use as key "1,2,3,4". This is equivalent to write the text horizontally in blocks of 4 letters in a matrix and to read it out vertically . Because the key is in an ascending order, the Scytale is denoted as an identical permutation. And because writing and read-out is done only once it is a simple (and no double) permutation.

In CT2 you can find the Scytale within the templates **Cryptography \ Classical**.

⁷This method cannot be simulated with a pure column transposition.

- **Turning grille** [Savard1999]: The German army used turning grilles during WW1.⁸ A square grille serves as a stencil, a quarter of its fields being holes. The first part of the message is written on a piece of paper through these holes, then the grille is rotated by 90 degrees and the user can write down the second part of the message, etc. But this method does only work, if the holes are chosen carefully: Every field has to be used, and no field may be used twice, either. The ciphertext is read out line by line.

In the example for a turning grille in table 2.2 you can write 4 times 16 characters of the cleartext on a piece of paper:

O	-	-	-	-	O	-	-
-	-	-	-	O	-	-	O
-	-	-	-	O	-	-	O
-	-	-	O	-	-	-	-
-	-	-	-	-	O	-	-
O	-	O	-	-	-	O	-
-	O	-	-	-	-	-	O
-	-	-	O	O	-	-	-

Table 2.2: 8x8 turning grille

2.1.2 Column and row transposition⁹

- **Simple columnar transposition** [Savard1999]: First of all, a keyword is chosen, that is written above the columns of a table. This table is filled with the text to be encrypted line by line. Then the columns are rearranged by sorting the letters of the keyword alphabetically. Afterwards the columns are read out from left to right to build the ciphertext.¹⁰

See table 2.3.

Plaintext: an example of transposition

K	E	Y
a	n	e
x	a	m
p	l	e
o	f	t
r	a	n
s	p	o
s	i	t
i	o	n

Table 2.3: Simple columnar transposition

⁸The turning grille was already invented in 1881 by Eduard Fleissner von Wostrowitz.

A good visualization can be found under www.turning-grille.com.

In JCT you can find it in the default perspective via the menu item **Visuals \ Grille**.

⁹Most of the following methods can be simulated in CT1 under the menu **Encrypt/Decrypt \ Symmetric (classic) \ Permutation**.

¹⁰Using CT1: Choose a key for the 1st permutation, input line by line, permute and output column by column.

Transposition key: K=2; E=1; Y=3.

Ciphertext: NALFA PIOAX PORSS IEMET NOTN

- **AMSCO cipher** [ACA2002]: The characters of the plaintext are written in alternating groups of one respectively two letters into a grille. Then the columns are swapped and the text can be read out.
- **Double column transposition** [Savard1999] : Double columnar transposition was frequently used during WW2 and during the Cold War. Two simple columnar transpositions with different keys are executed successively.¹¹
If the keys are different and long enough (at least each 20 characters), then this is even for today's computer a real challenge.¹²
- **Column transposition, General Luigi Sacco** [Savard1999]: The columns of a table are numbered according to the letters of the keyword. The plaintext is entered line by line, in the first line up to column number one, in the second line up to column number two, etc. Again, the ciphertext is read out in columns.

See table 2.4.

Plaintext: an example of transposition

C 1	O 5	L 2	U 6	M 3	N 4
a					
n	e	x			
a	m	p	l	e	
o	f	t	r	a	n
s	p				
o	s	i	t		
i	o	n			

Table 2.4: Columnar transposition (General Luigi Sacco)

Ciphertext: ANAOS OIEMF PSOXP TINLR TEAN

- **Column transposition, French army in WW1** [Savard1999]: After executing a simple columnar transposition, diagonal rows are read out.
- **Row transposition** [Savard1999]: The plaintext is divided into blocks of equal length and a keyword is chosen. Now the letters of the keyword are numbered and permutation is done only within each block according to this numbering.¹³

¹¹Using CT1: Choose a key for the 1st permutation, input line by line, permute and output column by column. Then choose a (different) key for the 2nd permutation, input line by line, permute and output column by column.

¹²MTC3 offers an according level-X challenge called "Double Columnar Transposition":

<http://www.mysterytwisterc3.org/en/challenges/level-x/double-column-transposition>

¹³Using CT1: Choose a key for 1st permutation, input line by line, permute column by column and output line by line.

2.1.3 Further transposition algorithm ciphers

- **Geometric figures** [Goebel2003]: Write the message into a grille following one pattern and read it out using another.
- **Union Route Cipher** [Goebel2003]: The Union Route Cipher derives from Civil War. This method does not rearrange letters of a given plaintext, but whole words. Particularly sensitive names and terms are substituted by codewords which are recorded in codebooks together with the existing routes. A route determines the size of a grille and the pattern that is used to read out the ciphertext. In addition, a number of filler words is defined.
- **Nihilist Transposition** [ACA2002]: Insert the plaintext into a square grille and write the same keyword above the columns and next to the lines. As this keyword is sorted alphabetically, the contents of the grille are rearranged, too. Read out the ciphertext line by line.

See table 2.5.

Plaintext: an example of transposition

	W	O	R	D	S		D	O	R	S	W	
W	a	n	e	x	a		D	s	p	o	i	s
O	m	p	l	e	o		O	e	p	l	o	m
R	f	t	r	a	n		R	a	t	r	n	f
D	s	p	o	s	i		S	n	i	o	-	t
S	t	i	o	n	-		W	x	n	e	a	a

Table 2.5: Nihilist transposition¹⁴

Ciphertext: SPOIS EPLOM ATRNF NIOTX NEAA

- **Cadenus cipher** [ACA2002]: Cadenus is a form of columnar transposition that uses two keywords.
The 1st keyword is used to swap columns.
The 2nd keyword is used to define the initial letter of each column: this 2nd keyword is a permutation of the used alphabet. This permutation is written on the left of the first column. Afterwards, each column is moved (wrap-around) so that it begins with the letter, which is in the same line as the key letter of the first keyword within the second keyword. Ciphertext is read out line by line.

See table 2.6.

Plaintext: cadenus is a form of columnar transposition using a keyword

Ciphertext:

SAASR PIFIU LONNS KTGWN EDOOA TDNNU IISFA OMYOC ROUCM AERRS

¹⁴After filling the matrix with the cleartext you get the left block. After switching rows and columns you get the right block

¹⁵Within the 2nd block of three chars those chars are printed bold which are at the top of the 3rd block after applying the 2nd key word.

	K	E	Y	E	K	Y	E	K	Y
A	c	a	d	a	c	d	s	a	a
D	e	n	u	n	e	u	s	r	p
X	s	i	s	i	s	s	i	f	i
K	a	f	o	f	a	o	u	l	o
C	r	m	o	m	r	o	n	n	s
W	f	c	o	c	f	o	k	t	g
N	l	u	m	u	l	m	w	n	e
S	n	a	r	a	n	r	d	o	o
Y	t	r	a	r	t	a	a	t	d
E	n	s	p	s	n	p	n	n	u
D	o	s	i	s	o	i	i	i	s
T	t	i	o	i	t	o	f	a	o
U	n	u	s	u	n	s	m	y	o
B	i	n	g	n	i	g	c	r	o
R	a	k	e	k	a	e	u	c	m
G	y	w	o	w	y	o	a	e	r
H	r	d	-	d	r	-	r	s	-

Table 2.6: Cadenus cipher¹⁵

2.2 Substitution ciphers

2.2.1 Monoalphabetic substitution ciphers

Monoalphabetic substitution assigns one character of the ciphertext alphabet to each plaintext character. This mapping remains unchanged during the whole process of encryption.

- **General monoalphabetic substitution / Random letter pairs**¹⁶ [Singh1999]: The substitution occurs by a given assignment of single letters.
- **Atbash cipher**¹⁷ [Singh1999]: Replace the first letter of the alphabet by the last letter of the alphabet, the second one by the last but one, etc.
- **Shift cipher, for example Caesar cipher**¹⁸ [Singh1999]: Plaintext alphabet and ciphertext alphabet are shifted against each other by a determined number of letters. Using the Caesar cipher means shifting letters about three positions.

Plaintext: three positions to the right

Ciphertext: WKUHH SRVLWLRQV WR WKH ULJKW

- **Affine cipher**¹⁹: This is a generalization of the shift cipher. A plaintext character is

¹⁶This cipher can be simulated in CT1 under the menu **Encrypt/Decrypt \ Symmetric (classic) \ Substitution / Atbash**.

In CT2 you can find these methods within the templates **Cryptography \ Classical**.

¹⁷This cipher can be simulated in CT1 under the menu **Encrypt/Decrypt \ Symmetric (classic) \ Substitution / Atbash**.

¹⁸In CT1 this method can be found at three different places in the menu tree:

- **Encrypt/Decrypt \ Symmetric (classic) \ Caesar / ROT13**
- **Analysis \ Symmetric Encryption (classic) \ Ciphertext only \ Caesar**
- **Indiv. Procedures \ Visualization of Algorithms \ Caesar**.

¹⁹Some according Sage samples are implemented at 2.5.2.3.

first substituted by another character and then the result is encrypted using the shift cipher. The name “affine cipher” was chosen because its encryption and decryption can be described as affine or linear function.

- **Substitution with symbols** [Singh1999], for instance the so-called “freemason cipher”: Each letter is replaced with a symbol.
- **Variants:** Fill characters, intentional mistakes [Singh1999].
- **Nihilist substitution**²⁰ [ACA2002]: Insert the alphabet into a 5x5-matrix to assign each letter the number built from row and column number. A keyword is chosen and placed above the columns of a second matrix (grille). The plaintext is written row by row into the grille. The ciphertext results from adding the numbers of the plaintext and the numbers of the keyword. Numbers between 100 and 110 are transformed to numbers between 00 and 10, so that each letter is represented by a two-digit number.

See table 2.7.

Plaintext: an example of substitution

	1	2	3	4	5
1	S	U	B	T	I
2	O	N	A	C	D
3	E	F	G	H	K
4	L	M	P	Q	R
5	V	W	X	Y	Z

	K	E	Y
	(35)	(31)	(54)
a	n	e	
(58)	(53)	(85)	
x	a	m	
(88)	(54)	(96)	
p	l	e	
(78)	(72)	(85)	
o	f	s	
(56)	(63)	(65)	
u	b	s	
(47)	(44)	(65)	
t	i	t	
(49)	(46)	(68)	
u	t	i	
(47)	(55)	(69)	
o	n		
(56)	(53)		

Table 2.7: Nihilist substitution

Ciphertext: 58 53 85 88 54 96 78 72 85 56 63 65 47 44 65 49 46 68 47 55 69 56 53

²⁰An animation of this Nihilist method can be found in CT1 at the menu item **Indiv. Procedures \ Visualization of Algorithms \ Nihilist**.

- **Codes** [Singh1999]: In the course of time, codebooks were used again and again. A codebook assigns a codeword, a symbol or a number to every possible **word** of a message. Only if both parties hold identical codebooks and if the assignment of codewords to plaintext words is not revealed, a successful and secret communication can take place.
- **Nomenclator** [Singh1999]: A nomenclator refers to techniques that combine the use of a cipher algorithm with a codebook. Often the encryption system is based upon a ciphertext alphabet. This alphabet is used to encrypt (via substitution) the bigger part of the message. Particularly frequent or top-secret words are replaced by a limited number of codewords existing besides the ciphertext alphabet.
- **Map cipher** [ThinkQuest1999]: This method constitutes a combination of substitution and steganography²¹. Plaintext characters are replaced by symbols which are arranged in a map following certain rules
- **Straddling Checkerboard** [Goebel2003]: A 3x10 matrix is filled with the letters of the used alphabet and two arbitrary digits or special characters as follows: The different letters of a keyword and the remaining characters are written into the grille. The columns are numbered 0 to 9, the second and the third line are numbered 1 and 2. Each plaintext character is replaced by the corresponding digit, respectively the corresponding pair of digits. As “1” and “2” are the first digits of the possible two-digit-numbers, they are not used as single digits.

See table 2.8.

Plaintext: an example of substitution

	0	1	2	3	4	5	6	7	8	9
	K	-	-	E	Y	W	O	R	D	A
1	B	C	F	G	H	I	J	L	M	N
2	P	Q	S	T	U	V	X	Z	.	/

Table 2.8: Straddling checkerboard with password “Keyword”

Ciphertext: 91932 69182 01736 12222 41022 23152 32423 15619

Besides, “1” and “2” are the most commonly used digits, but this feature is removed by the following technique.

It is ostentatious, how often the numbers 1 and 2 appear, but this will be fixed with the following version.

- **Straddling Checkerboard, variant** [Goebel2003] : This variant of the straddling checkerboard was developed by Soviet spies during WW2. Ernesto (Ché) Guevara and Fidel Castro allegedly used this cipher for their secret communication. A grille is filled with the alphabet (number of columns = length of keyword), and two arbitrary digits are chosen as reserved to indicate the second and third line of a 3x10-matrix (see above). Now the grille is traversed column by column and the single letters are transferred row by row into the matrix: For a faster encryption, the eight most common letters (ENIRSATO) are assigned the digits from 0 to 9, the reserved 2 digits are not assigned. The remaining

²¹Instead of encrypting a message, pure steganography tries to conceal its existence.

letters are provided with combinations of digits one after another and are inserted into the grille.

See table 2.9.

Plaintext: an example of substitution

	K	E	Y	W	O	R	D
Grille	A	B	C	F	G	H	I
	J	L	M	N	P	Q	S
	T	U	V	X	Z	.	/

	0	1	2	3	4	5	6	7	8	9
Matrix	A	T	E	-	N	O	R	-	I	S
3	K	J	B	L	U	Y	C	M	V	W
7	F	X	G	P	Z	H	Q	.	D	/

Table 2.9: Variant of the straddling checkerboard

Ciphertext: 04271 03773 33257 09343 29181 34185 4

- **Ché Guevara cipher:** A special variant is the cipher used by Ché Guevara (with an additional substitution step and a slightly changed checkerboard):
 - * The seven most frequent letters in Spanish are distributed in the first row.
 - * Four instead of three rows are used.
 - * So one could encrypt $10 * 4 - 4 = 36$ different characters.
- **Tri-Digital cipher** [ACA2002]: A keyword with ten letters is used to create a numeric key by numbering its letters corresponding to their alphabetical order. This key is written above the columns of the 3x10-matrix. This matrix is filled line by line with the alphabet as follows: The different letters of a keyword are inserted first, followed by the remaining letters. The last column is left out. Plaintext characters are substituted with numbers, the number of the last column is used to separate words.
- **Baconian cipher** [ACA2002]: Assign a five-digit binary code to every letter and to 6 numbers or special characters (for example 00000 = A, 00001 = B, etc.) and replace the plaintext characters with this binary code. Now use a second, unsuspicious message to hide the ciphertext inside of it. This may happen by upper and lower case or italicized letters: e.g. all letters of the unsuspicious message below a binary “1” are capitalised. Overall this is obtrusive.

See table 2.10.

2.2.2 Homophonic substitution ciphers

Homophonic methods constitute a special form of monoalphabetic substitution. Each character of the plaintext alphabet is assigned several ciphertext characters.

message	F	I	G	H	T
ciphertext	00101	01000	00110	00111	10011
unuspicious message	it is w	ar man	the su	ni ssh	in ing
Baconian Cipher	itIsW	aRman	thESu	niSSH	IniNG

Table 2.10: Baconian cipher

- **Homophonic monoalphabetic substitution**²² [Singh1999]: Each language has a typical frequency distribution of letters. To conceal this distribution, each plaintext letter is assigned several ciphertext characters. The number of ciphertext characters assigned depends on the frequency of the letter to be encrypted.
- **Beale cipher** [Singh1999]: The Beale cipher is a book cipher that numbers the words of a keytext. These numbers replace the cleartext letters by the words' initial letters.
- **Grandpré cipher** [Savard1999]: A square grille with 10 columns (other layouts are possible, too) is filled with ten words. The initial letters should result in an eleventh word. As columns and rows are numbered from 0 to 9, letters can be replaced by two-digit numbers. It is obvious that with the table having a hundred fields, most letters can be represented by more than one number. You should keep in mind that those ten words have to contain all letters of the plaintext alphabet.
- **Book cipher**: The words of a message are substituted by triples “page-line-position”. This method requires a detailed agreement of which book to use, especially regarding the edition (layout, error correction, etc.).

2.2.3 Polygraphic substitution ciphers

Polygraphic techniques do not work by replacing single characters, but by replacing whole groups of characters. In most cases, these groups are diagrams, trigrams or syllables.

- **“Great Chiffre”** [Singh1999]: This cipher was used by Louis XIV. and was not solved until the end of the nineteenth century. Cryptograms consisted of 587 different numbers, every number representing a syllable. The inventors of the “Great Chiffre” (Rossignol, father and son) constructed additional traps to increase security. For example, a number could assign a different meaning to or delete the preceding one.
- **Playfair cipher**²³ [Singh1999]: A 5x5-matrix is filled with the plaintext characters. For example, the different letters of a keyword are inserted first, followed by the remaining letters. The plaintext is divided into pairs, these digraphs are encrypted using the following rules:
 1. If both letters can be found in the same column, they are replaced by the letters underneath.
 2. If both letters can be found in the same row, take the letters to their right.

²²This cipher can be simulated in CT1 under the menu **Encrypt/Decrypt \Symmetric (classic)\ Homophone**.

²³In CT1 you can call this method under the menu **Encrypt/Decrypt \ Symmetric (classic) \ Playfair**.
In JCT you can find it in the default perspective via the menu item **Algorithms \ Classic \ Playfair**.

3. If both letters of the digraph are in different columns and rows, the replacement letters are obtained by scanning along the row of the first letter up to the column where the other letter occurs and vice versa.
4. Double letters are treated by special rules, if they appear in one digraph. They can be separated by a filler, for example.

See table 2.11.

Unformated Plaintext: plaintext letters are encrypted in pairs

Formated Plaintext: pl ai nt ex tl et te rs ar ee ncryptd in pairs

Formated Plaintext: plaintext letters are x encrypted in pairs

Formated Plaintext: pl ai nt ex tl et te rs ar ex en cr yp te di np ai rs

K	E	Y	W	O
R	D	A	B	C
F	G	H	I	L
M	N	P	Q	S
T	U	V	X	Z

Table 2.11: 5x5 Playfair matrix with password “Keyword”

Ciphertext: SHBHM UWUZF KUUKC MBDWU DURDA VUKBG PQBHC M

- **Trigraphic Playfair:** A 5x5-matrix is filled with the alphabet (see above) and the plaintext is divided into trigraphs. Trigraphs are encrypted according to the following rules:
 1. Three equal letters are substituted by three equal letters. It is the letter on the right underneath the original letter.
 2. A trigraph with two different letters is encrypted like a digraph in Playfair.
 3. If a trigraph contains three different characters, very complex rules come into effect.
See [Savard1999]
- **Substituting digraphs by symbols** [Savard1999]: Giovanni Battista della Porta, 15th century. He created a 20x20-matrix that contained one symbol for every possible combination of letters (his alphabet did not comprise more than twenty letters).
- **Four square cipher** [Savard1999]: This method is similar to Playfair, because it is based on a system of coordinates whose four quadrants are each filled with the alphabet. The layout of letters can differ from quadrant to quadrant. To encipher a message, act in the following way: Look up the first plaintext letter in the first quadrant and the second one in the third quadrant. These two letters are opposite corners of a rectangle and the ciphertext letters can be found in quadrant number two and four.

See table 2.12.

Plaintext: plaintext letters are encrypted in pairs

Ciphertext: MWYQW XQINO VNKGC ZWPZF FGZPM DIICC GRVCS

d	w	x	y	m	E	P	T	O	L
r	q	e	k	i	C	V	I	Q	Z
u	v	h	p	s	R	M	A	G	U
a	l	b	z	n	F	W	Y	H	S
g	c	o	f	t	B	N	D	X	K
Q	T	B	L	E	v	q	i	p	g
Z	H	N	D	X	s	t	u	o	h
P	M	I	Y	C	n	r	d	x	y
V	S	K	W	O	b	l	w	m	f
U	A	F	R	G	c	z	k	a	e

Table 2.12: Four square cipher

- **Two square cipher** [Savard1999]: The two square cipher resembles the four square cipher, but the matrix is reduced to two quadrants. Are both letters of the digraph part of the same row, they are just exchanged. Otherwise, the plaintext letters are considered as opposite corners of a rectangle and substituted by the other vertices. Quadrants can be arranged horizontal and vertical.
- **Tri square cipher** [ACA2002]: Three quadrants are filled with the same alphabet. The first plaintext letter is looked up in the first quadrant and can be encrypted with every letter of that column. The second plaintext letter is looked up in the second quadrant (diagonally across) and can be encrypted with every letter of that row. Between these two ciphertext characters, the letter at the intersection point is set.
- **Dockyard cipher** [Savard1999]: Used by the German navy during WW2.

2.2.4 Polyalphabetic substitution ciphers

Concerning polyalphabetic substitution, the assignment of ciphertext characters to plaintext characters is not static, but changes during the process of encryption (depending on the key).

- **Vigenère**²⁴ [Singh1999]: Each plaintext character is encrypted with a different ciphertext alphabet that is determined by the characters of a keyword (the so-called Vigenère tableau serves auxiliary means). If the plaintext is longer than the key, the latter is repeated.

See table 2.13.

- **Interrupted key:** The key is not repeated continuously, but starts again with every new word of the message.
- **Autokey**²⁵ [Savard1999]: After using the agreed key, use the message itself as a key.
See table 2.14.

²⁴In CT1 you can call this method under the menu **Encrypt/Decrypt \ Symmetric (classic) \ Vigenère**.

In CT2 you can find this method within the templates **Cryptography \ Classical**.

In JCT you can find it in the default perspective via the menu item **Algorithms \ Classic \ Vigenère**.

²⁵In CT2 you can find this method within the templates **Cryptography \ Classical**.

In JCT you can find it in the default perspective via the menu item **Algorithms \ Classic \ Autokey Vigenère**.

Plaintext:	the	alphabet	is	changing
Key:	KEY	KEYKEYKE	YK	EYKEYKEY
Ciphertext:	DLC	KPNREZOX	GC	GFKRESRE

-	A B C D E F G H I J K L M N O P Q R S T U V W X Y Z
A	A B C D E F G H I J K L M N O P Q R S T U V W X Y Z
B	B C D E F G H I J K L M N O P Q R S T U V W X Y Z A
C	C D E F G H I J K L M N O P Q R S T U V W X Y Z A B
D	D E F G H I J K L M N O P Q R S T U V W X Y Z A B C
E	E F G H I J K L M N O P Q R S T U V W X Y Z A B C D
F	F G H I J K L M N O P Q R S T U V W X Y Z A B C D E
G	G H I J K L M N O P Q R S T U V W X Y Z A B C D E F
H	H I J K L M N O P Q R S T U V W X Y Z A B C D E F G
I	I J K L M N O P Q R S T U V W X Y Z A B C D E F G H
J	J K L M N O P Q R S T U V W X Y Z A B C D E F G H I
K	K L M N O P Q R S T U V W X Y Z A B C D E F G H I J
...

Table 2.13: Vigenère tableau

Plaintext:	the	alphabet	is	changing
Key:	KEY	THEALPHA	BE	TISCHANG
Ciphertext:	DLC	TSTHLQLT	JW	VPSPNIAM

Table 2.14: Autokey variant of Vigenère

- **Progressive key** [Savard1999]: The key changes during the process of encryption. With every repetition, the characters of the keyword are shifted about one position. “KEY” becomes “LFZ”.
- **Gronsfeld** [Savard1999]: Variant of Vigenère that uses a numeric key.
- **Beaufort** [Savard1999]: Variant of Vigenère, the key is subtracted, not added. The ciphertext alphabets may be written backwards.
- **Porta** [ACA2002]: Variant of Vigenère with only 13 alphabets. As a consequence, two letters of the keyword are assigned the same ciphertext alphabet and the first and the second half of the alphabet are reciprocal.
- **Slidefair** [ACA2002]: This method can be used as a variant of Vigenère, Gronsfeld or Beaufort. Slidefair does encrypt digraphs according to the following rules: Look up the first letter in the plaintext alphabet above the tableau. Then look up the second one in the row belonging to the corresponding keyword letter. These two letters make up opposite corners of an imaginary rectangle. The letters at the two remaining corners substitute the digraph.

• Superposition

- **Running-key cipher**: A keytext (for example out of a book) is added to the plaintext.

- **Superposition with numbers:** A sequence or a number of sufficient length (for example pi) is added.
- **One Time Pad (OTP)**²⁶: A sequence of bytes is XORed byte-by-byte to the plaintext. This is a generalization of Vigenère’s mechanism and it was the first information-theoretically secure scheme (see chapter 1.1 “Security definitions”). To fulfill this claim the pad must be random and it must be used only once (to eliminate any semblance of pattern from the ciphertext).
Reason: Given ciphertext C, plaintext P, pad K and two plaintexts encrypted with the same key: $C_1 = P_1 \oplus K$; $C_2 = P_2 \oplus K$;
Thus, $C_1 \oplus C_2 = (P_1 \oplus K) \oplus (P_2 \oplus K) = P_1 \oplus P_2$;
which effectively could leak the plaintexts.²⁷
- **Phillips cipher** [ACA2002]: The alphabet is filled into a square table with 5 columns. Seven more tables are generated by first shifting the first row one position towards the bottom, then shifting the second row towards the bottom. The plaintext is divided into blocks of five which are encrypted with one matrix each. Letters are substituted by the ones on their right and underneath.
- **Ragbaby cipher** [ACA2002]: Construct an alphabet with 24 characters. Then number the plaintext characters, starting the numeration of the first word with “1”, the numeration of the second one with “2” and so forth. Number 25 corresponds to number 1. Each letter of the message is encrypted by shifting it the corresponding positions to the right.

See table 2.15.

alphabet: KEYWORDABCFGHILMNPSTUVXZ

Plaintext:	t h e	a l p h a b e t	i s	c h a n g i n g
Numbering:	1 2 3	2 3 4 5 6 7 8 9	3 4	4 5 6 7 8 9 10 11
Ciphertext:	U L O	C P V P I M C O	N X	I P I Z T X Y X

Table 2.15: Ragbaby cipher

2.3 Combining substitution and transposition

In the history of cryptography one often comes across combinations of the previous mentioned methods.

- **ADFG(V)X**²⁸ [Singh1999]: ADFG(V)X-encryption was developed in Germany during WW1. The alphabet is filled into a 5x5 or 6x6 matrix, and columns and rows are marked with the letters ADFGX and V, depending on the size of the grille. Each plaintext

²⁶CT1 offers this method under the menu **Encrypt/Decrypt \ Symmetric (classic) \ Vernam / OTP**.

In CT2 you can find this method within the templates **Cryptography \ Classical**.

In JCT you can find it in the default perspective via the menu item **Algorithms \ Classic \ XOR**.

²⁷In JCT you can play with an automatic cryptanalysis of running-key ciphertexts using the Viterbi analysis under the menu **Visuals \ Viterbi**.

²⁸In CT1 you can call this method under the menu **Encrypt/Decrypt \ Symmetric (classic) \ ADFGVX**.
In CT2 you can find this method within the templates **Cryptography \ Classical**.

character is substituted by the corresponding pair of letters. Finally, a (row-) transposition cipher is performed on the resulting text.

- **Fractionation** [Savard1999]: Generic term for all kinds of methods that encrypt one plaintext character by several ciphertext characters and then apply a transposition cipher to this ciphertext so that ciphertext characters originally belonging to each other are separated.
 - **Bifid/Polybius square/checkerboard** [Goebel2003]: Bifid encryption is the basic form of fractionation. A 5x5 matrix is filled with the plaintext alphabet (see Playfair encryption), rows and columns are numbered, so that each cleartext character can be substituted by a pair of digits. Mostly the plaintext is divided into blocks of equal length. The length of blocks (here 5) is another configuration parameter of this cipher. Block-by-block all line numbers are read out first, followed by all numbers naming the columns. To obtain the ciphertext, the digits are pairwise transformed into letters again. The numbers can be any permutation of (1,2,3,4,5), which is one key of configuration parameter of this cipher. Instead of numbering rows and columns, a keyword can be used, too.

See table 2.16.

	2	4	5	1	3
1	K	E	Y	W	O
4	R	D	A	B	C
2	F	G	H	I	L
3	M	N	P	Q	S
5	T	U	V	X	Z

Plaintext:	combi	nings	ubsti	tutio	nandt	ransp	ositi
Rows:	41342	32323	54352	55521	34345	44333	13252
Columns:	33211	41443	41321	24213	45442	25435	33121

Table 2.16: Bifid cipher

41342 32323 54352 55521 34345 44333 13252 33211 41443 41321 24213 45442 25435
33121

Ciphertext: BNLLL UPHVI NNUCS OHLMW BDNOI GINUR HCZQI

- **Trifid** [Savard1999]: 27 characters (alphabet + 1 special character) may be represented by a triple consisting of the digits 1 to 3. The message to be encrypted is divided into blocks of three and the relevant triple is written underneath each plaintext character as a column. The resulting numbers below the plaintext blocks are read out line by line and are substituted with the corresponding characters.
- **Bazeries cipher** [ACA2002]: The plaintext alphabet is filled into a 5x5-matrix column by column, a second matrix is filled line by line with a keyword (a number smaller than a million) followed by the remaining letters of the alphabet. Then the message is divided into blocks of arbitrary length and their characters' order is inverted. Finally, each letter is substituted – according to its position in the original matrix – by its counterpart in the second matrix.

See table 2.17.

Plaintext: combining substitution and transposition

Keyword: 900.004 (nine hundred thousand and four)

a	f	l	q	v	N	I	E	H	U
b	g	m	r	w	D	R	T	O	S
c	h	n	s	x	A	F	B	C	G
d	i	o	t	y	K	L	M	P	Q
e	k	p	u	z	V	W	X	Y	Z

com	bini	ngs	ub	stitu	tiona	ndt	ran	sposi	ti	on
moc	inib	sgn	bu	utits	anoit	tdn	nar	isops	it	no
TMA	LBLD	CRB	DY	YPLPC	NBMLP	PKB	BNO	LCMXC	LP	BM

Table 2.17: Bazeries cipher

- **Digrafid cipher** [ACA2002]: To substitute digraphs, the following table is used (to simplify matters, the alphabet is used in its original form). Look up the first letter of the digraph in the horizontal alphabet and write down the column number. Then look up the second letter in the vertical alphabet and write down the corresponding line number. Between these two numbers, the number at the intersection point is set. Afterwards, the triple are written vertically underneath the digraphs that are arranged in groups of three. The three digit numbers arising horizontally are transformed back into digraphs.

Remark: This cipher only works with complete blocks of 3 pairs of cleartext characters. For a complete description, it is necessary to explain how sender and receiver handle texts which fill in the last block only 1-5 characters. The possibilities range from ignoring a last and incomplete block to padding it with random characters or with characters predefined in advance.

See table 2.18.

1	2	3	4	5	6	7	8	9			
A	B	C	D	E	F	G	H	I	1	2	3
J	K	L	M	N	O	P	Q	R	4	5	6
S	T	U	V	W	X	Y	Z	.	7	8	9
									A	J	S
									B	K	T
									C	L	U
									D	M	V
									E	N	W
									F	O	X
									G	P	Y
									H	Q	Z
									I	R	.
											9

co mb in	in gs ub	st it ut	io na nd	tr an sp	os it io
3 4 9	9 7 3	1 9 3	9 5 5	2 1 1	6 9 9
2 4 2	2 3 7	9 3 9	2 4 4	8 2 8	6 3 2
6 2 5	5 1 2	2 2 2	6 1 4	9 5 7	1 2 6
LI KB FN	.C BY EB	SU I. BK	RN KD FD	BA HQ RP	X. FT AO

Table 2.18: Digrafid cipher

- **Nicodemus cipher** [ACA2002]: First of all, a simple columnar transposition is carried out. Before reading out the columns, the message is encrypted additionally by Vigenère (all letters of a column are enciphered with the corresponding keyword letter). The ciphertext is read out in vertical blocks.

See table 2.19.

Plaintext: combining substitution and transposition

Ciphertext: SMRYX MLSCC KLEZG YSRVW JSKDX RLBYN WMYDG N

2.4 Further methods

- **“Pinprick encryption”** [Singh1999]: For centuries, this simple encryption method has been put into practice for different reasons (actually steganography). During the Victorian Age, for example, small holes underneath letters in newspaper articles marked the characters of a plaintext, as sending a newspaper was much more cheaper than the postage on a letter.
- **Stencil:** Stencils (Cardboard with holes) are also known as “Cardinal-Richelieu-Key”. Sender and receiver have to agree upon a text. Above this text, a stencil is laid and the letters that remain visible make up the ciphertext.
- **Card games** [Savard1999]: The key is created by means of a pack of cards and rules that are agreed upon in advance. All methods mentioned in this paragraph are designed as paper and pencil methods, i.e. they are applicable without electronic aid. A pack of cards is unsuspicious to outsiders, shuffling the deck provides a certain amount of coincidence,

K	E	Y	E	K	Y	E	K	Y
c	o	m	o	c	m	S	M	K
b	i	n	i	b	n	M	L	L
i	n	g	n	i	g	R	S	E
s	u	b	u	s	b	Y	C	Z
s	t	i	t	s	i	X	C	G
t	u	t	u	z	t	Y	J	R
i	o	n	o	i	n	S	S	L
a	n	d	n	a	d	R	K	B
t	r	a	r	t	a	V	D	Y
n	s	p	s	n	p	W	X	N
o	s	i	s	o	i	W	Y	G
t	i	o	i	t	o	M	D	N

Table 2.19: Nicodemus cipher

cards can be transformed into numbers easily and a transposition cipher can be carried out without any further aid.

- **Solitaire cipher (Bruce Schneier)**²⁹ [Schneier1999]: Sender and receiver have to own a deck of cards shuffled in the same manner. A key stream is generated that has to consist of as many characters as the message to be encrypted.

The algorithm to generate the key is based on a shuffled deck of 54 cards (Ace, 2 - 10, jack, queen, king in four suits and two jokers). The pack of cards is held face up:

1. Swap the first joker with the card beneath it.
2. Move the second joker two cards down.
3. Now swap the cards above the first joker with those below the second one.
4. Look at the bottom card and convert it into a number from 1 to 53 (bridge order of suits: clubs, diamonds, hearts, spades; joker = 53). Write down this number and count down as many cards starting with the top card. These cards are swapped with the remaining cards, only the bottom card remains untouched.
5. Look at the top card and convert it into a number, too. Count down as many cards starting with the top card.
6. Write down the number of the following card. This card is converted into your first keystream character. As we need numbers from 1 to 26 to match the letters of our alphabet, clubs and hearts correspond to the numbers 1 to 13, diamonds and spades to 14 to 26. If your output card is a joker, start again.

For each keystream character you like to generate, these six steps have to be carried out. This procedure is – manually – very lengthy (4 h for 300 characters, dependant on your exercise) and requires high concentration.

Encryption takes place by addition modulo 26. Encryption is relatively fast compared to the key stream generation.

This P&P cipher creates a key stream which is so good, that even nowadays it is hard to crack the cipher if you don't know the originally sortierted card deck (ciphertext-only attack).

²⁹In CT1 you can call this method under the menu **Encrypt/Decrypt \ Symmetric (classic) \ Solitaire**.

- **Mirdek cipher (Paul Crowley)** [Crowley2000]: Even though this method is quite complicated, the author provides a very good example to illustrate the procedure.
- **Playing Card cipher (John Savard)** [Savard1999]: This algorithm uses a shuffled deck of 52 cards (no joker). Separate rules describe how to shuffle the deck. A keystream is created via the following steps:
 1. The pack of cards lies in front of the user, top down. Cards are turned up and dealt out in a row until the total of the cards is 8 or more.
 2. If the last card dealt out is a J, Q or K, write down its value, otherwise write down the sum of the cards dealt out (a number between 8 and 17). In a second row, deal out that number of cards.
 3. The remaining cards are dealt out in rows under the second row. The first one ends under the lowest card of the top row, the second one under the next lowest card, and so on. If there are two identical cards, red is lower than black.
 4. The cards dealt out under step 3 are collected column by column, starting with the column under the lowest card. The first card that is picked up becomes the bottom card (face up).
 5. The cards dealt out in step 1 and 2 are picked up, beginning with the last card.
 6. The deck is turned over, the top card is now the bottom card (face down). Afterwards, steps 1 to 6 are repeated twice.

To generate a keystream character, write down the first card not being J, Q or K. Count down that number of cards. The card selected has to be between 1 and 10. Now repeat these steps beginning with the last card. These two numbers are added and the last digit of the sum is your keystream character.

- **VIC cipher** [Savard1999]: This is a highly complicated but relatively secure paper and pencil method. It has been developed and applied by Soviet spies. Amongst other things, the user had to create ten pseudo-random numbers out of a date, the first words of a sentence and any five-digit number. A straddling checkerboard is part of the encryption, too. A detailed description can be found under [Savard1999].

2.5 Appendix: Examples using Sage

In the following section some classic ciphers are implemented using the open source computer algebra system **Sage**³⁰. The code was tested with Sage version 5.3. All ciphers are explained in chapter 2 (“Paper and Pencil Encryption Methods”).

To make the sample code³¹ of the ciphers easier to understand, we used the structure and the naming conventions shown in the Figure 2.1 below:

- Encryption consists of the two steps encoding and enciphering.
 - Encoding adapts the letters in the given plaintext P to the case defined in the given alphabet, and all non-alphabet characters are filtered out.
 - Enciphering creates the ciphertext C .
- Decryption also consists of two steps: deciphering and decoding.
 - Decoding is only necessary if the symbols in the alphabet are not ASCII characters.

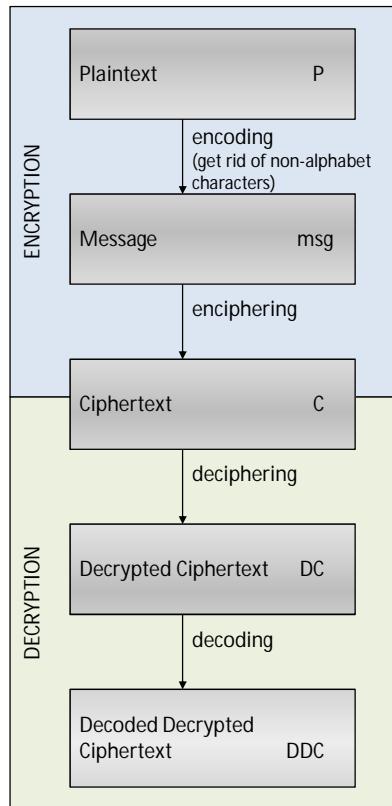


Figure 2.1: Structure and naming convention of the Sage cipher code examples

³⁰A first introduction to the CAS Sage can be found in the appendix A.7.

³¹Further examples with Sage concerning classic crypto methods can be found e.g.:

- at http://www.sagemath.org/doc/constructions/linear_codes.html#classical-ciphers
- in the thesis of Minh Van Nguyen [Nguyen2009]

2.5.1 Transposition ciphers

Transposition ciphers are implemented in the Sage class

```
sage.crypto.classical.TranspositionCryptosystem
```

To construct and work with a transposition cipher, we first need to determine the alphabet that contains the symbols used to build the space of our plaintext and ciphertext. Typically, this alphabet will be the upper-case letters of the English alphabet, which can be accessed via the function

```
sage.monoids.string_monoid.AlphabeticStrings
```

We then need to decide on the block length of a block permutation, which is the length of the row vector to be used in the simple columns transposition. This row vector is our key, and it specifies a permutation of a plaintext.

The following first example of transposition ciphers has block length 14, and the key is build in a way, that every letter in the plaintext is shifted to the right by two characters, with wrap around at the end of the block. That is the encryption process. The decryption process is shifting each letter of the ciphertext to the left by $14 - 2 = 12$.

Sage sample 2.1 Simple transposition by shifting (key and inverse key explicitly given)

```
sage: # transposition cipher using a block length of 14
sage: T = TranspositionCryptosystem(AlphabeticStrings(), 14)
sage: # given plaintext
sage: P   = "a b c d e f g h i j k l m n"
sage: # encryption key
sage: key = [3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 1, 2]
sage:
sage: # encode plaintext (get rid of non-alphabet chars, convert lower-case to upper-case)
sage: msg = T.encoding(P)
sage: # encrypt plaintext by shifting to the left by 2 letters (do it in two steps)
sage: E   = T(key)
sage: C   = E(msg); C
CDEFGHIJKLMNOPAB
sage:
sage: # decrypt ciphertext by shifting to the left by 12 letters
sage: keyInv = [13, 14, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12]
sage: D   = T(keyInv)
sage: D(C)
ABCDEFGHIJKLMN
sage:
sage: # Representation of key and inverse key as permutations
sage: E
(1,3,5,7,9,11,13)(2,4,6,8,10,12,14)
sage: D
(1,13,11,9,7,5,3)(2,14,12,10,8,6,4)
```

The second example of transposition ciphers is also a simple shifting column transposition. But now the code is a little bit more automated: The keys are generated from the shift parameter.

Sage sample 2.2 Simple transposition by shifting (key and inverse key constructed with “range”)

```
sage: # transposition cipher using a block length of 14, code more variable
sage: keylen = 14
sage: shift = 2
sage: A = AlphabeticStrings()
sage: T = TranspositionCryptosystem(A, keylen)
sage:
sage: # construct the plaintext string from the first 14 letters of the alphabet plus blanks
sage: # plaintext   = "A B C D E F G H I J K L M N"
sage: A.gens()
(A, B, C, D, E, F, G, H, I, J, K, L, M, N, O, P, Q, R, S, T, U, V, W, X, Y, Z)
sage: P=''
sage: for i in range(keylen): P=P + " " + str(A.gen(i))
....:
sage: P
'A B C D E F G H I J K L M N'
sage:
sage: # encryption key
sage: # key = [3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 1, 2]
sage: key = [(i+shift).mod(keylen) + 1 for i in range(keylen)]; key
[3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 1, 2]
sage:
sage: # encode plaintext (get rid of non-alphabet chars)
sage: msg = T.encoding(P)
sage: # encrypt plaintext by shifting to the left by 2 letters (do it in one step)
sage: C = T.enciphering(key, msg); C
CDEFGHIJKLMNOPAB
sage:
sage: # decrypt ciphertext by shifting to the left by 12 letters
sage: # keyInv = [13, 14, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12]
sage: shiftInv=keylen-shift;
sage: keyInv = [(i+shiftInv).mod(keylen) + 1 for i in range(keylen)]; keyInv
[13, 14, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12]
sage: DC = T.enciphering(keyInv, C); DC
ABCDEFGHIJKLMN
sage:
sage: # decryption using the "deciphering method with key" instead of "enciphering with keyInv"
sage: # using the deciphering method requires to change the type of the variable key
sage: DC = T.deciphering(T(key).key(), C); DC
ABCDEFGHIJKLMN
sage:
sage: # representation of key and inverse key as permutations
sage: T(key)
(1,3,5,7,9,11,13)(2,4,6,8,10,12,14)
sage: T(key).key()
(1,3,5,7,9,11,13)(2,4,6,8,10,12,14)
sage: T(keyInv)
(1,13,11,9,7,5,3)(2,14,12,10,8,6,4)
```

In the third example of transposition ciphers we use an arbitrary permutation as key in the encryption and decryption processes in order to scramble the characters within each block (block length = number of columns in a simple column transposition). If the block length is n , then the key must be a permutation on n symbols. The following example uses the method `random_key()` of the class `TranspositionCryptosystem`. Each call to `random_key()` produces a different key. Note that therefore your results (key and ciphertext) may be different from the following example.

Sage sample 2.3 Simple column transposition with randomly generated (permutation) key

```
sage: # Remark: Enciphering here requires, that the length of msg is a multiple of keylen
sage: keylen = 14    # length of key
sage: A = AlphabeticStrings()
sage: T = TranspositionCryptosystem(A, keylen); T
Transposition cryptosystem on Free alphabetic string monoid on A-Z of block length 14
sage:
sage: P = "a b c d e f g h i j k l m n o p q r s t u v w x y z a b"
sage: key = T.random_key(); key
(1,2,3,13,6,5,4,12,7)(11,14)
sage: msg = T.encoding(P); msg
ABCDEFGHIJKLMNOPQRSTUVWXYZAB
sage: C = T.enciphering(key, msg); C
BCMLDEAHIJNGFKPQAZRSOVWXBUTY
sage: # decryption using the "deciphering method with key" instead of "enciphering with keyInv"
ssage: DC = T.deciphering(key, C); DC
ABCDEFGHIJKLMNOPQRSTUVWXYZAB
sage:
sage: # Just another way of decryption: Using "enciphering" with the inverse key
sage: keyInv = T.inverse_key(key); keyInv
(1,7,12,4,5,6,13,3,2)(11,14)
sage: DC = T.enciphering(keyInv, C); DC
ABCDEFGHIJKLMNOPQRSTUVWXYZAB
sage:
sage: # Test correctness of decryption
sage: msg == DC
True
```

The fourth example of transposition ciphers additionally shows the key space of a simple column transposition.

Sage sample 2.4 Simple column transposition (showing the key_space)

```
sage: keylen = 14 # length of key
sage: A = AlphabeticStrings()
sage: T = TranspositionCryptosystem(A, keylen); T
Transposition cryptosystem on Free alphabetic string monoid on A-Z of block length 14
sage: T.key_space()
Symmetric group of order 14! as a permutation group
sage: # Remark: The key space is not quite correct as also permutations shorter than keylen are counted.
sage:
sage: P = "a b c d e f g h i j k l m n o p q r s t u v w x y z a b"
sage: key = T.random_key(); key
(1,2,7)(3,9)(4,5,10,12,8,13,11)(6,14)
sage: msg = T.encoding(P); msg
ABCDEFGHIJKLMNOPQRSTUVWXYZAB
sage:
sage: # enciphering in one and in two steps
sage: C = T.enciphering(key, msg); C
BGIEJNAMCLDHKFPUWSXBOAQZRVYT
sage:
sage: enc = T(key); enc.key()
(1,2,7)(3,9)(4,5,10,12,8,13,11)(6,14)
sage: C = enc(msg); C
BGIEJNAMCLDHKFPUWSXBOAQZRVYT
sage:
sage: # deciphering
sage: DC = T.deciphering(key, C); DC
ABCDEFGHIJKLMNOPQRSTUVWXYZAB
```

2.5.2 Substitution ciphers

Substitution cryptosystems are implemented in Sage in the class

```
sage.crypto.classical.SubstitutionCryptosystem
```

The following code sample uses Sage to construct a substitution cipher with a random key. A random key can be generated using the method `random_key()` of the class `SubstitutionCryptosystem`. Different keys determine different substitution ciphers. So each call to `random_key()` returns different results.

Sage sample 2.5 Monoalphabetic substitution with randomly generated key

```
sage: # plaintext/ciphertext alphabet
sage: A = AlphabeticStrings()
sage: S = SubstitutionCryptosystem(A)
sage:
sage: P = "Substitute this with something else better."
sage: key = S.random_key(); key
INZDHFUXJPATQOYLKSWGVEMRB
sage:
sage: # method encoding can be called from A or from T
sage: msg = A.encoding(P); msg
SUBSTITUTETHISWITHSOMETHINGELSEBETTER
sage: C = S.enciphering(key, msg); C
WVNWGJGVGHGXJWCJGXWYQHGXJOUHTWHNHGGHS
sage:
sage: # We now decrypt the ciphertext to recover our plaintext.
sage:
sage: DC = S.deciphering(key, C); DC
SUBSTITUTETHISWITHSOMETHINGELSEBETTER
sage: msg == DC
True
```

2.5.2.1 Caesar cipher

The following example uses Sage to construct a Caesar cipher.

Sage sample 2.6 Caesar (substitution by shifting the alphabet; key explicitly given, step-by-step approach)

```
sage: # plaintext/ciphertext alphabet
sage: A = AlphabeticStrings()
sage: P = "Shift the alphabet three positions to the right."
sage:
sage: # construct Caesar cipher
sage: S = SubstitutionCryptosystem(A)
sage: key = A([3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 18, 19, \
....:           20, 21, 22, 23, 24, 25, 0, 1, 2])
sage:
sage: # encrypt message
sage: msg      = A.encoding(P); msg
SHIFTTHEALPHABETTHREEPOSITIONSTOTHERIGHT
sage: encrypt = S(key); encrypt
DEFGHIJKLMNOPQRSTUVWXYZABC
sage: C      = encrypt(msg); C
VKLIWWKHOSKDEHWWKUHHSRVLWLRQVWRWKHULJKW
sage:
sage: # Next, we recover the plaintext.
sage: # decrypt message
sage: keyInv = A([23, 24, 25, 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, \
....:           14, 15, 16, 17, 18, 19, 20, 21, 22])
sage: decrypt = S(keyInv); decrypt
XYZABCDEFGHIJKLMNOPQRSTUVWXYZ
sage: DC      = decrypt(C); DC
SHIFTTHEALPHABETTHREEPOSITIONSTOTHERIGHT
sage: msg == DC
True
```

The second Caesar sample does the same, but the code is more sophisticated/automated/flexible.

Sage sample 2.7 Caesar (substitution by shifting the alphabet; substitution key generated)

```
sage: # plaintext/ciphertext alphabet
sage: A = AlphabeticStrings()
sage: keylen = len(A.gens()); keylen
26
sage: shift = 3
sage: P = "Shift the alphabet three positions to the right."
sage:
sage: # construct Caesar cipher
sage: S = SubstitutionCryptosystem(A)
sage: S
Substitution cryptosystem on Free alphabetic string monoid on A-Z
sage: # key = A([3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 18, 19, \
sage: #           20, 21, 22, 23, 24, 25, 0, 1, 2])
sage: key = [(i+shift).mod(keylen) for i in range(keylen)];
sage: key = A(key); key
DEFGHIJKLMNOPQRSTUVWXYZABC
sage: len(key)
26
sage:
sage: # encrypt message
sage: msg      = A.encoding(P); msg
SHIFTTHEALPHABETTHREEPOSITIONSTOTHERIGHT
sage: C       = S.enciphering(key, msg); C
VKLIWWKHDOSKDEHWWKUHHHSRVLWLRQVWRWKHULJKW
sage:
sage: # Next, we recover the plaintext.
sage: # decrypt message
sage: # keyInv = A([23, 24, 25, 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, \
sage: #           14, 15, 16, 17, 18, 19, 20, 21, 22])
sage: shiftInv=keylen-shift;
sage: keyInv = [(i+shiftInv).mod(keylen) for i in range(keylen)];
sage: keyInv = A(keyInv); keyInv
XYZABCDEFGHIJKLMNOPQRSTUVWXYZ
sage: DC      = S.enciphering(keyInv, C); DC
SHIFTTHEALPHABETTHREEPOSITIONSTOTHERIGHT
sage:
sage: # Just another way of decryption: Using "deciphering" with the key
sage: DC      = S.deciphering(key, C); DC
SHIFTTHEALPHABETTHREEPOSITIONSTOTHERIGHT
sage:
sage: msg == DC
True
```

2.5.2.2 Shift cipher

The shift cipher can also be thought of as a generalization of the Caesar cipher. While the Caesar cipher restricts us to shift exactly three positions along an alphabet, the shift cipher allows us to shift any number of positions along the alphabet.

In the above samples we applied the `SubstitutionCryptosystem` and build Caesar as a special kind of substitution. In contrast here Caesar can be build as a special kind of the shift cipher.

The shift cipher is implemented directly in the Sage class

```
sage.crypto.classical.ShiftCryptosystem
```

In the following example, we construct a shift cipher over the capital letters of the English alphabet. We then encrypt a plaintext P by shifting it 12 positions along the alphabet. Finally, we decrypt the ciphertext C and make sure that the result (DC) is indeed the original plaintext. Shifting is a special way of substitution.

Sage sample 2.8 A shift cipher over the upper-case letters of the English alphabet

```
sage: # construct Shift cipher directly
sage: shiftcipher = ShiftCryptosystem(AlphabeticStrings()); shiftcipher
Shift cryptosystem on Free alphabetic string monoid on A-Z
sage: P = shiftcipher.encoding("Shift me any number of positions."); P
SHIFTMEANYNUMBEROFPOSITIONS
sage: key = 12  # shift can be any integer number
sage:
sage: # shift the plaintext by 12 positions to get the ciphertext
sage: C = shiftcipher.enciphering(key, P); C
ETURFYQMZKZGYNQDARBAEUFUAZE
sage:
sage: # decrypt the ciphertext and ensure that it is the original plaintext
sage: DC = shiftcipher.deciphering(key, C); DC
SHIFTMEANYNUMBEROFPOSITIONS
sage: DC == P
True
```

The Caesar cipher is simply a shift cipher whose shifting key is 3. In the next example, we use the shift cipher to create a Caesar cipher over the capital letters of the English alphabet.

Sage sample 2.9 Constructing the Caesar cipher using the shift cipher

```
sage: # create a Caesar cipher
sage: caesarcipher = ShiftCryptosystem(AlphabeticStrings())
sage: P = caesarcipher.encoding("Shift the alphabet by three positions to the right."); P
SHIFTTHEALPHABETBYTHREEPOSITIONSTOTHERIGHT
sage:
sage: key = 3  # shift the plaintext by exactly 3 positions
sage: C = caesarcipher.enciphering(key, P); C
VKLIWWKHDKOSKDEHWEBWKUHHSRVLWLRQVWRWKHULJKW
sage:
sage: # decrypt the ciphertext and ensure that it is the original plaintext
sage: DC = caesarcipher.deciphering(key, C); DC
SHIFTTHEALPHABETBYTHREEPOSITIONSTOTHERIGHT
sage: DC == P
True
```

2.5.2.3 Affine cipher

The affine cipher is implemented in the Sage class

```
sage.crypto.classical.AffineCryptosystem
```

In the following example, we construct an affine cipher $c_i = b * p_i + a$ with key $(3, 13)$ and use this key to encrypt a given plaintext $P = (p_1, p_2, \dots, p_n)$. The plaintext is then decrypted and the result DC is compared to the original plaintext.

Sage sample 2.10 An affine cipher with key $(3, 13)$

```
sage: # create an affine cipher
sage: affineCipher = AffineCryptosystem(AlphabeticStrings()); affineCipher
Affine cryptosystem on Free alphabetic string monoid on A-Z
sage: P = affineCipher.encoding("The affine cryptosystem.")
sage: P
THEAFFINECRYPTOSYSTEM
sage:
sage: # encrypt the plaintext using the key (3, 13)
sage: a, b = (3, 13)
sage: C = affineCipher.enciphering(a, b, P)
sage: C
SIZNCCLAZTMHGSDPHPSZX
sage:
sage: # decrypt the ciphertext and make sure that it is equivalent to the original plaintext
sage: DC = affineCipher.deciphering(a, b, C)
sage: DC
THEAFFINECRYPTOSYSTEM
sage: DC == P
True
```

We can also construct a shift cipher using the affine cipher. To do so, we need to restrict keys of the affine cipher be of the form $(1, b)$ where b is any non-negative integer. For instance, we can work through Sage example 2.8 on page 43 as follows:

Sage sample 2.11 Constructing a shift cipher using the affine cipher

```
sage: # construct a shift cipher
sage: shiftcipher = AffineCryptosystem(AlphabeticStrings()); shiftcipher
Affine cryptosystem on Free alphabetic string monoid on A-Z
sage: P = shiftcipher.encoding("Shift me any number of positions.")
sage: P
SHIFTMEANYNUMBEROFPPOSITIONS
sage:
sage: # shift the plaintext by 12 positions to get the ciphertext
sage: a, b = (1, 12)
sage: C = shiftcipher.enciphering(a, b, P)
sage: C
ETURFYQMZKGYNQDARBAEUFUAZE
sage:
sage: # decrypt the ciphertext and ensure that it is the original plaintext
sage: DC = shiftcipher.deciphering(a, b, C); P
SHIFTMEANYNUMBEROFPPOSITIONS
sage: DC == P
True
```

We can also use the affine cipher to create the Caesar cipher. To do so, the encryption/decryption key must be $(1, 3)$. In the next example, we work through Sage example 2.9 on page 43 using the affine cipher.

Sage sample 2.12 Constructing the Caesar cipher using the affine cipher

```
sage: # create a Caesar cipher
sage: caesarcipher = AffineCryptosystem(AlphabeticStrings())
sage: P = caesarcipher.encoding("Shift the alphabet by three positions to the right.")
sage: P
SHIFTTHEALPHABETBYTHREEPOSITIONSTOTHERIGHT
sage:
sage: # shift the plaintext by 3 positions
sage: a, b = (1, 3)
sage: C = caesarcipher.enciphering(a, b, P)
sage: C
VKLIWWKHOSKDEHWEBKUHHHSRVLWLRQVWRWKHULJKW
sage:
sage: # decrypt the ciphertext and ensure that it is the original plaintext
sage: DC = caesarcipher.deciphering(a, b, C)
sage: DC
SHIFTTHEALPHABETBYTHREEPOSITIONSTOTHERIGHT
sage: DC == P
True
```

2.5.2.4 Substitution with symbols

In the following Sage example the symbols are from the binary number system. A monoalphabetic substitution cipher with a binary alphabet has very little security: Because the plaintext/ciphertext alphabet has only the two elements 0 and 1, there are only two keys possible: (0 1) and (1 0).

Remark: The key of a general substitution cipher contains all symbols of the alphabet exactly once.

Sage sample 2.13 Monoalphabetic substitution with a binary alphabet

```
sage: # the plaintext/ciphertext alphabet
sage: B = BinaryStrings()
sage: # substitution cipher over the alphabet B; no keylen argument possible
sage: S = SubstitutionCryptosystem(B); S
Substitution cryptosystem on Free binary string monoid
sage: # To get a substitute for each symbol, key has always the length of the alphabet
sage: key = S.random_key(); key
10
sage: len(key)
2
sage: P = "Working with binary numbers."
sage: # encryption
sage: msg = B.encoding(P); msg
010101101101110111001001101011010010110111001100111001000000111011101101\
00101110100011010000010000001100010011010010110111001100001011100100111100100\
1000000110111001110101011011010110001001100101011100100111001100101110
sage: C = S.enciphering(key, msg); C
1010100010010000100011011001010010010110100100011001100011011111000100010010\
110100010111001011110111110011101100101101001000110011110100011011000011011\
01111100100011000101010010010110110101000110110001100110010001
sage: # decryption
sage: DC = S.deciphering(key, C); DC
0101011011011101110010011010110110010110111001100111001000000111011101101\
00101110100011010000010000001100010011010010110111001100001011100100111100100\
1000000110111001110101011011010110001001100101011100100111001100101110
sage: msg == DC
True
```

Remark: Currently `S` has no attribute `key`, and I found no way to transform the binary sequence `DC` back to ASCII.

The second sample of a monoalphabetic substitution with symbols uses a larger alphabet as plaintext/ciphertext space as the first sample. Here the hexadecimal number system is used as substitution alphabet.

Sage sample 2.14 Monoalphabetic substitution with a hexadecimal alphabet (and decoding in Python)

```
sage: A = HexadecimalStrings()
sage: S = SubstitutionCryptosystem(A)
sage: key = S.random_key(); key
2b56a4e701c98df3
sage: len(key)
16
sage: # Number of possible keys
sage: factorial(len(key))
20922789888000
sage: P = "Working with a larger alphabet."
sage:
sage: msg = A.encoding(P); msg
576f726b696e6720776974682061206c617267657220616c7068616265742e
sage: C = S.enciphering(key, msg); C
47e375e9e1efe75277e17ae052eb52e8eb75e7e47552ebe872e0ebe5e47a5f
sage: DC = S.deciphering(key, C); DC
576f726b696e6720776974682061206c617267657220616c7068616265742e
sage: msg == DC
True
sage:
sage: # Conversion hex back to ASCII:
sage: # - AlphabeticStrings() and HexadecimalStrings() don't have according methods.
sage: # - So we used Python directly.
sage: import binascii
sage: DDC = binascii.a2b_hex(repr(DC)); DDC
'Working with a larger alphabet.'
sage:
sage: P == DDC
True
```

2.5.2.5 Vigenère cipher

The Vigenère cipher is implemented in the Sage class

```
sage.crypto.classical.VigenereCryptosystem
```

For our ciphertext/plaintext space, we can work with the upper-case letters of the English alphabet, the binary number system, the octal number system, or the hexadecimal number system. Here is an example using the class `AlphabeticStrings`, which implements the English capital letters.

Sage sample 2.15 Vigenère cipher

```
sage: # construct Vigenere cipher
sage: keylen = 14
sage: A = AlphabeticStrings()
sage: V = VigenereCryptosystem(A, keylen); V
Vigenere cryptosystem on Free alphabetic string monoid on A-Z of period 14
sage:
sage: # alternative could be a given key: key = A('ABCDEFGHIJKLMN'); key
sage: key = V.random_key(); key
WSSSEEGVVAARUD
sage: len(key)
14
sage:
sage: # encoding
sage: P = "The Vigener cipher is polyalphabetic."
sage: len(P)
38
sage: msg = V.encoding(P); msg      # alternative: msg = A.encoding(P); msg
THEVIGENERECIPHERISPOLYALPHABETIC
sage:
sage: # encryption [2 alternative ways (in two steps or in one): both work]
sage: # encrypt = V(key); encrypt
sage: # C = encrypt(msg); C
sage: C = V.enciphering(key, msg); C
PZWMKIZRETCSDWJAWTUGTALGBDXLAG
sage:
sage: # decryption
sage: DC = V.deciphering(key, C); DC
THEVIGENERECIPHERISPOLYALPHABETIC
sage: msg == DC
True
```

2.5.3 Hill cipher

The Hill [Hill1929, Hill1931] or matrix cipher is more mathematically sophisticated than other ciphers mentioned in this chapter. The encryption/decryption key of this cipher is an invertible square matrix, and the plaintext/ciphertext is processed also as a matrix. The encryption and decryption processes use matrix operations modulo 26. The Hill cipher is implemented in the Sage class

```
sage.crypto.classical.HillCryptosystem
```

In the following example our plaintext/ciphertext space is the capital letters of the English alphabet. In the Hill cipher, each letter of this alphabet is assigned a unique integer modulo 26. The size of matrix (also called its dimension) is not restricted by the cipher.

Remark: Comparing the Hill implementation in CrypTool v1.4.31 and in Sage v5.3:

- Sage offers fast commandline operations; CT1 offers its functionality within a GUI.
- Sage offers any dimension; CT1 is restricted to a matrix size between 1 and 10.
- Sage allows negative numbers in the key matrix, and converts them automatically into appropriate non-negative numbers; CT1 doesn't allow negative numbers in the key matrix.
- Sage always sets the first alphabet character to 0;
Sage only allows the 26 capital letters as alphabet;
and it uses only the multiplication variant plaintext row vector * key matrix:
 $C = P^*M$.
- CT1 offers to choose also 1 as value for the first alphabet character; you can combine your alphabet within the text options dialog; and it offers to use a reverse multiplication variant too.

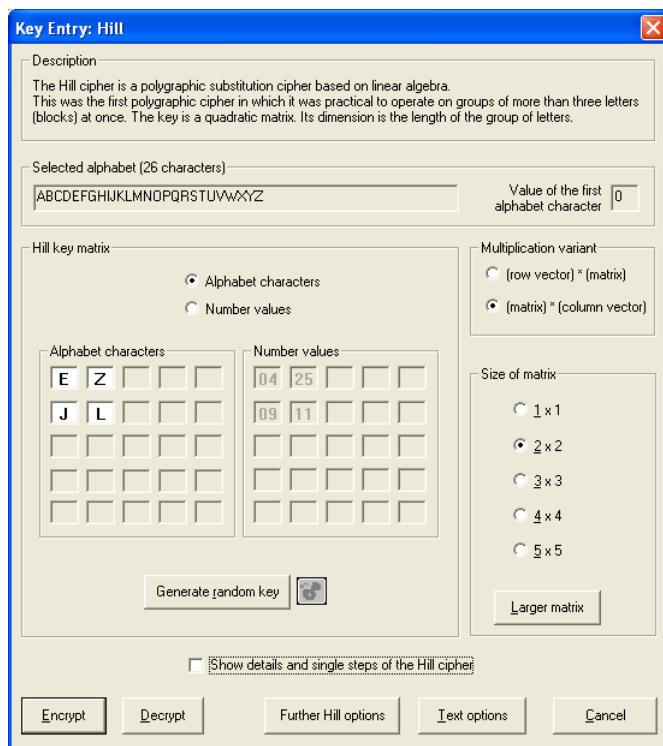


Figure 2.2: Hill dialog in CT1 with the operations and options available

Sage sample 2.16 Hill cipher

```
sage: # construct a Hill cipher
sage: keylen = 19    # keylen = 3 # Alternative key length with non-random small key
sage: A = AlphabeticStrings()
sage: H = HillCryptosystem(A, keylen); H
Hill cryptosystem on Free alphabetic string monoid on A-Z of block length 19
sage:
sage: # To create key non-randomly, HKS is necessary [even H.key_space() is not enough].
sage: # HKS = H.key_space()
sage: # key = HKS([[1,0,1],[0,1,1],[2,2,3]]); key
sage:
sage: # Random key creation
sage: key = H.random_key(); key
[10  7  5  2  0  6 10 23 15  7 17 19 18  2  9 12  0 10 11]
[23  1  1 10  4  9 21  1 25 22 19  8 17 22 15  8 12 25 22]
[ 4 12 16 15  1 12 24  5  9 13  5 15  8 21 23 24 22 20  6]
[ 5 11  6  7  3 12  8  9 21 20  9  4 16 18 10  3  2 23 18]
[ 8 22 14 14 20 13 21 19  3 13  2 11 13 23  9 25 25  6  8]
[24 25  8 24  7 18  3 20  6 11 25  5  6 19  7 24  2  4 10]
[15 25 11  1  4  7 11 24 20  2 18  4  9  8 12 19 24  0 12]
[14  6  2  9 11 20 13  4 10 11  4 23 14 22 14 16  9 12 18]
[12 10 21  5 21 15 16 17 19 20  1  1 15  5  0  2 23  4 14]
[21 15 15 16 15 20  4 10 25  7 15  4  7 12 24  9 19 10  6]
[25 15  2  3 17 23 21 16  8 18 23  4 22 11 15 19  6  0 15]
[14 23  9  3 18 15 10 18  7  5 12 23 11  9 22 21 20  4 14]
[ 3  6  8 13 20 16 11  1 13 10  4 21 25 15 12  3  0 11 18]
[21 25 14  6 11  3 21  0 19 17  5  8  5  4  9  2 23 19 15]
[ 8 11  9 11 20 15  6  1  3 18 18 22 16 17  6  3 15 11  2]
[21 15  5 22  2  9  0  4 22 10  2 10 19 19 17 19  1 21  4]
[ 7 17  9  2 15  5 14  3  6  9 12 12 22 15  8  4 21 14 19]
[19 14 24 19  7  5 22 22 13 14  7 18 17 19 25  2  1 23  6]
[ 2  6 14 22 17  7 23  6 22  7 13 20  0 14 23 17  6  1 12]
sage:
sage: # encoding and encryption
sage: P = "Hill or matrix cipher uses matrix operations."
sage: len(P)
45
sage: # implementation requires: Length of msg is a multiple of matrix dimension (block_length)
sage: msg = H.encoding(P); msg
HILLORMATRIXCIPHERUSESMATRIXOPERATIONS
sage: len(msg)
38
sage:
sage: # encryption
sage: C = H.enciphering(key, msg); C
CRWCKPRVYXBRZTNZCTQWFWSDBCHABGMNEHVP
sage:
sage: # decryption
sage: DC = H.deciphering(key, C); DC
HILLORMATRIXCIPHERUSESMATRIXOPERATIONS
sage: msg == DC
True
sage:
sage: # alternative decryption using inverse matrix
sage: keyInv = H.inverse_key(key); keyInv
[ 6 23  1 23  3 12 17 22  6 16 22 14 18  3  1 10 21 16 20]
[18 23 15 25 24 23  7  4 10  7 21  7  9  0 13 22  5  5 23]
...
[10 11 12  6 11 17 13  9 19 16 14 24  4  8  5 16 18 20  1]
[19 16 16 21  1 19  7 12  3 18  1 17  7 10 24 21  7 16 11]
sage: DC      = H.enciphering(keyInv, C); DC
HILLORMATRIXCIPHERUSESMATRIXOPERATIONS
```

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Chapter 3

Prime Numbers

(Bernhard Esslinger, May 1999; Updates Nov. 2000, Dec. 2001, June 2003, May 2005, March 2006, June 2007, Jan. 2010, Dec. 2012)

Progress requires exchange of knowledge.

Quote 5: Albert Einstein¹

3.1 What are prime numbers?

Prime numbers are whole, positive numbers greater than or equal to 2 that can only be divided by 1 and themselves. All other natural numbers (composite numbers) can be formed by multiplying prime numbers.

The *natural* numbers $\mathbb{N} = \{1, 2, 3, 4, \dots\}$ thus comprise

- the number 1 (the unit value)
- the primes and
- the composite numbers.

Prime numbers are particularly important for three reasons:

- In number theory, they are considered to be the basic components of natural numbers, upon which numerous brilliant mathematical ideas are based.
- They are of extreme practical importance in modern cryptography (public key cryptography). The most common public key procedure, invented at the end of the 1970's, is RSA encryption. Only using (large) prime numbers for particular parameters can you guarantee that an algorithm is secure, both for the RSA procedure and for even more modern procedures (e.g. elliptic curves).
- The search for the largest known prime numbers does not have any practical usage known to date, but requires the best computers, is an excellent benchmark (possibility for determining the performance of computers) and leads to new calculation methods on many

¹Albert Einstein, German physicist and Nobel Prize winner, Mar 14, 1879 – Apr 14, 1955.

computers
 (see also: <http://www.mersenne.org/prime.htm>).

Many people have been fascinated by prime numbers over the past two millennia. Ambition to make new discoveries about prime numbers has often resulted in brilliant ideas and conclusions. The following section provides an easily comprehensible introduction to the basics of prime numbers. We will also explain what is known about the distribution (density, number of prime numbers in particular intervals) of prime numbers and how prime number tests work.

3.2 Prime numbers in mathematics

Every whole number has a factor. The number 1 only has one factor, itself, whereas the number 12 has the six factors 1, 2, 3, 4, 6, 12. Many numbers can only be divided by themselves and by 1. With respect to multiplication, these are the “atoms” in the area of numbers. Such numbers are called prime numbers.

In mathematics, a slightly different (but equivalent) definition is used.

Definition 3.2.1. A whole number $p \in \mathbb{N}$ is called prime if $p > 1$ and p only possesses the trivial factors ± 1 and $\pm p$.

By definition, the number 1 is not a prime number. In the following sections, p will always denote a prime number.

The sequence of prime numbers starts with

2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 41, 43, 47, 53, 59, 61, 67, 71, 73, 79, 83, 89, 97, ⋯.

The first 100 numbers include precisely 25 prime numbers. After this, the percentage of primes constantly decreases. Prime numbers can be factorized in a uniquely *trivial* way:

$$5 = 1 \cdot 5, \quad 17 = 1 \cdot 17, \quad 1013 = 1 \cdot 1013, \quad 1,296,409 = 1 \cdot 1,296,409.$$

All numbers that have 2 or more factors not equal 1 are called *composite* numbers. These include

$$4 = 2 \cdot 2, \quad 6 = 2 \cdot 3$$

as well as numbers that *look like primes*, but are in fact composite:

$$91 = 7 \cdot 13, \quad 161 = 7 \cdot 23, \quad 767 = 13 \cdot 59.$$

Theorem 3.2.1. Each whole number m greater than 1 possesses a lowest factor greater than 1. This is a prime number p . Unless m is a prime number itself, then: p is less than or equal to the square root of m .

All whole numbers greater than 1 can be expressed as a product of prime numbers — in a unique way. This is the claim of the **1st fundamental theorem of number theory** (= fundamental theorem of arithmetic = fundamental building block of all positive integers).

Theorem 3.2.2. Each element n of the natural numbers greater than 1 can be written as the product $n = p_1 \cdot p_2 \dots p_m$ of prime numbers. If two such factorizations

$$n = p_1 \cdot p_2 \cdot \dots \cdot p_m = p'_1 \cdot p'_2 \cdot \dots \cdot p'_{m'}$$

are given, then they can be reordered such that $m = m'$ and for all i : $p_i = p'_i$. (p_1, p_2, \dots, p_m are called the prime factors of n).

In other words: each natural number other than 1 can be written as a product of prime numbers in precisely one way, if we ignore the order of the factors. The factors are therefore unique (the *expression as a product of factors* is unique)! For example,

$$60 = 2 \cdot 2 \cdot 3 \cdot 5 = 2^2 \cdot 3^1 \cdot 5^1.$$

And this — other than changing the order of the factors — is the only way in which the number 60 can be factorized. If you allow numbers other than primes as factors, there are several ways of factorizing integers and the **uniqueness** is lost:

$$60 = 1 \cdot 60 = 2 \cdot 30 = 4 \cdot 15 = 5 \cdot 12 = 6 \cdot 10 = 2 \cdot 3 \cdot 10 = 2 \cdot 5 \cdot 6 = 3 \cdot 4 \cdot 5 = \dots .$$

The following section is aimed more at those familiar with mathematical logic: The 1st fundamental theorem only appears to be obvious . We can construct numerous other sets of numbers (i.e. other than positive whole numbers greater than 1), for which numbers in the set cannot be expressed uniquely as a product of the prime numbers of the set: In the set $M = \{1, 5, 10, 15, 20, \dots\}$ there is no equivalent to the fundamental theorem under multiplication. The first five prime numbers of this sequence are 5, 10, 15, 20, 30 (note: 10 is prime, because 5 is not a factor of 10 in this set — the result is not an element of the given basic set M). Because the following applies in M :

$$100 = 5 \cdot 20 = 10 \cdot 10$$

and 5, 10, 20 are all prime numbers in this set, the expression as a product of prime factors is not unique here.

3.3 How many prime numbers are there?

For the natural numbers, the primes can be compared to elements in chemistry or the elementary particles in physics (see [Blum1999, p. 22]).

Although there are only 92 natural chemical elements, the number of prime numbers is unlimited. Even the Greek, Euclid² knew this in the third century B.C.

Theorem 3.3.1 (Euclid³). *The sequence of prime numbers does not discontinue. Therefore, the quantity of prime numbers is infinite.*

His proof that there is an infinite number of primes is still considered to be a brilliant mathematical consideration and conclusion today (**proof by contradiction**). He assumed that there is only a finite number of primes and therefore a largest prime number. Based on this assumption, he drew logical conclusions until he obtained an obvious contradiction. This meant that something must be wrong. As there were no mistakes in the chain of conclusions, it could

²Euclid, a Greek mathematician of 4th and 3rd century B.C. He worked at the Egyptian academy of Alexandria and wrote “The Elements”, the most well known systematically textbook of the Greek mathematics.

³The common usage of the term does not denote Euclid as the inventor of the theorem rather; the true inventor is merely not as prominent. The theorem has already been distinguished and proven in Euclid's Elements (Book IX, theorem 20). The phraseology is remarkable due to the fact that the word infinite is not used. The text reads as followed

Oí πρῶτοι ἀριθμοὶ πλείους εἰσὶ παντὸς τοῦ προτεύθεντος πλήθος πρώτων ἀριθμῶν,

the English translation of which is: the prime numbers are more than any previously existing amount of prime numbers.

only be the assumption that was wrong. Therefore, there must be an infinite number of primes!

Proof according to Euclid (proof by contradiction)

Assumption: There is a *finite* number of primes.

Conclusion: Then these can be listed $p_1 < p_2 < p_3 < \dots < p_n$, where n is the (finite) number of prime numbers. p_n is therefore the largest prime. Euclid now looks at the number $a = p_1 \cdot p_2 \cdots p_n + 1$. This number cannot be a prime number because it is not included in our list of primes. It must therefore be divisible by a prime, i.e. there is a natural number i between 1 and n , such that p_i divides the number a . Of course, p_i also divides the product $a - 1 = p_1 \cdot p_2 \cdots p_n$, because p_i is a factor of $a - 1$. Since p_i divides the numbers a and $a - 1$, it also divides the difference of these numbers. Thus: p_i divides $a - (a - 1) = 1$. p_i must therefore divide 1, which is impossible.

Contradiction: Our assumption was false.

Thus there is an *infinite* number of primes (Cross-reference: overview under 3.9 of the number of prime numbers in various intervals). \square

Here we should perhaps mention yet another fact which is initially somewhat surprising. Namely, in the prime numbers sequence p_1, p_2, \dots , gaps between prime numbers can have an individually determined length n . It is undeniable that under the n succession of natural numbers

$$(n+1)! + 2, \dots, (n+1)! + (n+1),$$

none of them is a prime number since in order, the numbers $2, 3, \dots, (n+1)$ are comprised respectively as real divisors. ($n!$ means the product of the first n natural numbers therefore $n! = n * (n-1) * \dots * 2 * 1$).

3.4 The search for extremely large primes

The largest prime numbers known today have several millions digits, which is too big for us to imagine. The number of elementary particles in the universe is estimated to be “only” a 80-digit number (See: overview under 3.11 of various orders of magnitude / dimensions).

3.4.1 The 20+ largest known primes (as of July 2012)

The following table contains the biggest, currently known primes and a description of its particular number type⁴:

⁴An up-to-date version can be found in the internet at <http://primes.utm.edu/largest.html> and at <http://primes.utm.edu/mersenne/index.html>.

	Definition	Decimal Digits	When	Description
1	$2^{43,112,609} - 1$	12,978,189	2008	Mersenne, 47th known
2	$2^{42,643,801} - 1$	12,837,064	2009	Mersenne, 46th known
3	$2^{37,156,667} - 1$	11,185,272	2008	Mersenne, 45th known
4	$2^{32,582,657} - 1$	9,808,358	2006	Mersenne, 44th known
5	$2^{30,402,457} - 1$	9,152,052	2005	Mersenne, 43rd known
6	$2^{25,964,951} - 1$	7,816,230	2005	Mersenne, 42nd known
7	$2^{24,036,583} - 1$	7,235,733	2004	Mersenne, M-41
8	$2^{20,996,011} - 1$	6,320,430	2003	Mersenne, M-40
9	$2^{13,466,917} - 1$	4,053,946	2001	Mersenne, M-39
10	$19,249 \cdot 2^{13,018,586} + 1$	3,918,990	2007	Generalized Mersenne ⁵
11	$27,653 \cdot 2^{9,167,433} + 1$	2,759,677	2005	Generalized Mersenne
12	$28,433 \cdot 2^{7,830,457} + 1$	2,357,207	2004	Generalized Mersenne
13	$2^{6,972,593} - 1$	2,098,960	1999	Mersenne, M-38
14	$5,359 \cdot 2^{5,054,502} + 1$	1,521,561	2003	Generalized Mersenne
15	$4,847 \cdot 2^{3,321,063} + 1$	999,744	2005	Generalized Mersenne
16	$3 \cdot 2^{3,136,255} - 1$	944,108	2007	Generalized Mersenne
17	$2^{3,021,377} - 1$	909,526	1998	Mersenne, M-37
18	$2^{2,976,221} - 1$	895,932	1997	Mersenne, M-36
19	$222,361 \cdot 2^{2,854,840} + 1$	859,398	2006	Generalized Mersenne
20	$1,372,930^{131,072} + 1$	804,474	2003	Generalized Fermat ⁶
21	$1,361,244^{131,072} + 1$	803,988	2004	Generalized Fermat
22	$1,176,694^{131,072} + 1$	795,695	2003	Generalized Fermat
23	$342,673 \cdot 2^{2,639,439} - 1$	794,556	2007	Generalized Mersenne

Table 3.1: The 20+ largest known primes and its particular number types (as of July 2012)

The largest known prime is a Mersenne prime, found by the GIMPS project (chapter 3.4.2).

Within the largest known primes there are also numbers of the type generalized Mersenne number (chapter 3.6.2) and generalized Fermat numbers (chapter 3.6.5).

3.4.2 Special number types – Mersenne numbers and Mersenne primes

Almost all known huge prime numbers are special candidates, called *Mersenne numbers*⁷ of the form $2^p - 1$, where p is a prime. Not all Mersenne numbers are prime:

$$\begin{aligned} 2^2 - 1 &= 3 &\Rightarrow \text{prime} \\ 2^3 - 1 &= 7 &\Rightarrow \text{prime} \\ 2^5 - 1 &= 31 &\Rightarrow \text{prime} \\ 2^7 - 1 &= 127 &\Rightarrow \text{prime} \\ 2^{11} - 1 &= 2,047 = 23 \cdot 89 &\Rightarrow \text{NOT prime!} \end{aligned}$$

Even Mersenne knew that not all Mersenne numbers are prime (see exponent $p = 11$). A prime Mersenne number is called Mersenne prime number.

However, he is to be thanked for the interesting conclusion that a number of the form $2^n - 1$ cannot be a prime number if n is a composite number:

Theorem 3.4.1 (Mersenne). *If $2^n - 1$ is a prime number, then n is also a prime number.*

Proof

The theorem of Mersenne can be proved by contradiction. We therefore assume that there exists a composite natural number n (with real factorization) $n = n_1 \cdot n_2$, with the property that $2^n - 1$ is a prime number.

From

$$\begin{aligned} (x^r - 1)((x^r)^{s-1} + (x^r)^{s-2} + \cdots + x^r + 1) &= ((x^r)^s + (x^r)^{s-1} + (x^r)^{s-2} + \cdots + x^r) \\ &\quad - ((x^r)^{s-1} + (x^r)^{s-2} + \cdots + x^r + 1) \\ &= (x^r)^s - 1 = x^{rs} - 1, \end{aligned}$$

we conclude

$$2^{n_1 n_2} - 1 = (2^{n_1} - 1)((2^{n_1})^{n_2-1} + (2^{n_1})^{n_2-2} + \cdots + 2^{n_1} + 1).$$

Because $2^n - 1$ is a prime number, one of the above two factors on the right-hand side must be equal to 1. This is the case if and only if $n_1 = 1$ or $n_2 = 1$. But this contradicts our assumption.

⁵This number was found within the distributed computing project “Seventeen or Bust” (SoB) (<http://www.seventeenorbust.com>) at March 26, 2007. While the well known GIMPS project (chapter 3.4.2) searches for bigger and bigger of the infinitely many primes, there is a chance, that the SoB project could have been completed its task sometime.

The SoB project tries to prove computationally, that the number $k = 78,557$ is the smallest Sierpinski number (John Selfridge proved in 1962, that 78,557 is a Sierpinski number).

The famous Polish mathematician Waclaw Sierpinski (1882 to 1969) proved in 1960, that there exist infinitely many odd integers k , which fulfill the following property: For all Sierpinski numbers k it is true: All numbers $N = k \cdot 2^n + 1$ are composite for all integers $n \geq 1$ (Sierpinski’s Composite Number Theorem, <http://mathworld.wolfram.com/SierpinskisCompositeNumberTheorem.html>).

When the project started in 2002 there have been 17 possible candidates < 78557 (this is the reason for the project’s name “Seventeen or Bust”). It is sufficient to find one single counter-example, to exclude a candidate k , which means to find a single $n \geq 1$, where $N = k \cdot 2^n + 1$ is prime. So it is only a byproduct of this task that this also generates new monster primes.

⁶Generalized Fermat number: $1,372,930^{131,072} + 1 = 1,372,930^{(2^{17})} + 1$

⁷Marin Mersenne, French priest and mathematician, Sep 08, 1588 – Sep 01, 1648.

Therefore the assumption is false. This means that there exists no composite number n , such that $2^n - 1$ is a prime. \square

Unfortunately this theorem only applies in one direction (the inverse statement does not apply, no equivalence): that means that there exist prime exponent for which the Mersenne number is **not** prime (see the above example $2^{11} - 1$, where 11 is prime, but $2^{11} - 1$ not).

Mersenne claimed that $2^{67} - 1$ is a prime number. There is also a mathematical history behind this claim: it first took over 200 years before Edouard Lucas (1842-1891) proved that this number is composite. However, he argued indirectly and did not name any of the factors. In 1903, Cole⁸ showed which factors make up this composite number:

$$2^{67} - 1 = 147,573,952,589,676,412,927 = 193,707,721 \cdot 761,838,257,287.$$

He admitted to having worked 20 years on the factorization (expression as a product of prime factors)⁹ of this 21-digit decimal number!

Due to the fact that the exponents of the Mersenne numbers do not use all natural numbers, but only the primes, the *experimental space* is limited considerably. The currently known Mersenne prime numbers have the exponents¹⁰

2; 3; 5; 7; 13; 17; 19; 31; 61; 89; 107; 127; 521; 607; 1,279; 2,203; 2,281; 3,217;
4,253; 4,423; 9,689; 9,941; 11,213; 19,937; 21,701; 23,207; 44,497; 86,243; 110,503;
132,049; 216,091; 756,839; 859,433; 1,257,787; 1,398,269; 2,976,221; 3,021,377;
6,972,593; 13,466,917; 20,996,011; 24,036,583; 25,964,951; 30,402,457;
32,582,657; 37,156,667; 42,643,801; 43,112,609.

Thus 47 Mersenne prime numbers are currently known.

The 19th number with the exponent 4,253 was the first with at least 1,000 digits in decimal system (the mathematician Samuel Yates coined the expression *titanic* prime for this; it was discovered by Hurwitz in 1961); the 27th number with the exponent 44,497 was the first with at least 10,000 digits in the decimal system (Yates coined the expression *gigantic* prime for this. These names are now long outdated).

For the first 41 Mersenne prime numbers we know that this list is complete. The exponents until the 47th Mersenne prime number have not yet been checked completely.¹¹

Right now (July 2012) all prime exponents smaller than 24,212,801 have been tested and double-checked¹²: So we can be certain, that this is really the 41th Mersenne prime number and that there are no smaller undiscovered Mersenne primes (it is common usage to use the notation M-nn not until it is proven, that the nn-th known Mersenne prime is really the nn-th Mersenne prime). Here are some examples in more detail:

⁸Frank Nelson Cole, American mathematician, Sep. 20, 1861 – May 26, 1926.

⁹Using CT1 you can factorize numbers in the following way: menu **Indiv. Procedures \ RSA Cryptosystem \ Factorization of a Number**.

CT1 can factorize with the quadratic sieve (QS) on a single PC in a reasonable time numbers no longer than 250 bit. Numbers bigger than 1024 bits are currently not accepted by CT1.

With CT2 you can use both the QS and the GNFS (general number field sieve), and CT2 can distribute the calculations on several computers.

The current factorization records are listed in chapter 4.11.4.

¹⁰Landon Curt Noll lists all known Mersenne primes including its date of discovery and its value as number and as word: <http://www.isthe.com/chongo/tech/math/prime/mersenne.html>

Also see: <http://www.utm.edu/>.

¹¹The current status of the check can be found at: <http://www.mersenne.org/status.htm>.

Hints, how the primality of a number can be checked, are in chapter 3.5, prime number tests.

¹²See home page of the GIMPS project: http://www.mersenne.org/report_milestones.

M-37 – January 1998

The 37th Mersenne prime,

$$2^{3,021,377} - 1$$

was found in January 1998 and has 909,526 digits in the decimal system, which corresponds to 33 pages in the newspaper!

M-38 – June 1999

The 38th Mersenne prime, called M-38,

$$2^{6,972,593} - 1$$

was discovered in June 1999 and has 2,098,960 digits in the decimal system (that corresponds to around 77 pages in the newspaper).

M-39 – December 2001

The 39th Mersenne prime, called M-39,

$$2^{13,466,917} - 1,$$

was published at December 6, 2001 – more exactly, the verification of this number, found at November 14, 2001 by the Canadian student Michael Cameron, was successfully completed. This number has about 4 million decimal digits (exactly 4,053,946 digits). Trying only to print this number

$$(924947738006701322247758 \dots 1130073855470256259071)$$

would require around 200 pages in the Financial Times.

GIMPS

The GIMPS project (Great Internet Mersenne Prime Search) was founded in 1996 by George Woltman to search for new largest Mersenne primes (<http://www.mersenne.org>). Further explanations about this number type can be found under Mersenne numbers and Mersenne primes.

Right now the GIMPS project has discovered 13 largest Mersenne primes so far, including the largest known prime number at all.

Table 3.2 contains these Mersenne record primes.^{13,14}

Richard Crandall discovered the advanced transform algorithm used by the GIMPS program. George Woltman implemented Crandall's algorithm in machine language, thereby producing a prime-search program of unprecedented efficiency, and that work led to the successful GIMPS project.

On June 1st, 2003 a possible Mersenne prime was reported to the GIMPS server, which was checked afterwards as usual, before it was to be published. Unfortunately mid June the initiator

¹³An up-to-date version can be found in the internet at <http://www.mersenne.org/history.htm>.

¹⁴Always, when a new record is published in the respective forums the same and often ironic discussions start: Does this kind of research have a deeper sense? Can this result be applied for anything useful? The answer is, that we don't know it yet. In fundamental research one cannot see at once whether and how it brings mankind forward.

Definition	Decimal Digits	When	Who
$2^{43,112,609} - 1$	12,978,189	August 23, 2008	Edson Smith
$2^{42,643,801} - 1$	12,837,064	April 12, 2009	Odd Magnar Strindmo
$2^{37,156,667} - 1$	11,185,272	September 6, 2008	Hans-Michael Elvenich
$2^{32,582,657} - 1$	9,808,358	September 4, 2006	Curtis Cooper/Steven Boone
$2^{30,402,457} - 1$	9,152,052	December 15, 2005	Curtis Cooper/Steven Boone
$2^{25,964,951} - 1$	7,816,230	February 18, 2005	Martin Nowak
$2^{24,036,583} - 1$	7,235,733	May 15, 2004	Josh Findley
$2^{20,996,011} - 1$	6,320,430	November 17, 2003	Michael Shafer
$2^{13,466,917} - 1$	4,053,946	November 14, 2001	Michael Cameron
$2^{6,972,593} - 1$	2,098,960	June 1, 1999	Nayan Hajratwala
$2^{3,021,377} - 1$	909,526	January 27, 1998	Roland Clarkson
$2^{2,976,221} - 1$	895,932	August 24, 1997	Gordon Spence
$2^{1,398,269} - 1$	420,921	November 1996	Joel Armengaud

Table 3.2: The largest primes found by the GIMPS project (as of July 2012)

and GIMPS project leader George Woltman had to tell, that two independent verification runs proved the number was composite. This was the first false positive report of a client in 7 years.

Now more than 130,000 volunteers, amateurs and experts, participate in the GIMPS project. They connect their computers into the so called “primenet”, organized by the company entropia.

3.4.3 Challenge of the Electronic Frontier Foundation (EFF)

This search is also spurred on by a competition started by the non-profit organization EFF (Electronic Frontier Foundation) using the means of an unknown donor. The participants are rewarded with a total of 500,000 USD if they find the longest prime number. In promoting this project, the unknown donor is not looking for the quickest computer, but rather wants to draw people’s attention to the opportunities offered by *cooperative networking*
<http://www.eff.org/awards/coop>

The discoverer of M-38 received 50,000 USD from the EFF for discovering the first prime with more than 1 million decimal digits.

For the next prize of 100,000 USD offered by EFF for a proven prime with more than 10 million decimal digits, Edson Smith qualified, who found the number $2^{43,112,609} - 1$ within the GIMPS project.

According to the EFF rules for their prizes they offer in the next stage 150,000 USD for a proven prime with more than 100 million decimal digits.

Edouard Lucas (1842-1891) held the record for the longest prime number for over 70 years by proving that $2^{127} - 1$ is prime. No new record is likely to last that long.

3.5 Prime number tests¹⁵

In order to implement secure encryption procedures we need extremely large prime numbers (in the region of $2^{2,048}$, i.e. numbers with 600 digits in the decimal system!).

If we look for the prime factors in order to decide whether a number is prime, then the search takes too long, if even the smallest prime factor is enormous. Factorizing numbers using systematic computational division or using the sieve of Eratosthenes is only feasible using current computers for numbers with up to around 20 digits in the decimal system. The biggest number factorized into its 2 almost equal prime factors has 232 digits (see RSA-768 in chapter 4.11.4).

However, if we know something about the *construction* of the number in question, there are extremely highly developed procedures that are much quicker. These procedures can determine the primality attribute of a number, but they cannot determine the prime factors of a number, if it is compound.

In the 17th century, Fermat¹⁶ wrote to Mersenne that he presumed that all numbers of the form

$$f(n) = 2^{2^n} + 1$$

are prime for all whole numbers $n \geq 0$ (see below, chapter 3.6.4).

As early as in the 19th century, it was discovered that the 29-digit number

$$f(7) = 2^{2^7} + 1$$

is not prime. However, it was not until 1970 that Morrison/Billhart managed to factorize it.

$$\begin{aligned} f(7) &= 340,282,366,920,938,463,463,374,607,431,768,211,457 \\ &= 59,649,589,127,497,217 \cdot 5,704,689,200,685,129,054,721 \end{aligned}$$

Despite Fermat was wrong with this supposition, he is the originator of an important theorem in this area: Many rapid prime number tests are based on the (little) Fermat theorem put forward by Fermat in 1640 (see chapter 4.8.3).

Theorem 3.5.1 (“little” Fermat). *Let p be a prime number and a be any whole number, then for all a*

$$a^p \equiv a \pmod{p}.$$

This could also be formulated as follows:

Let p be a prime number and a be any whole number that is not a multiple of p (also $a \not\equiv 0 \pmod{p}$), then $a^{p-1} \equiv 1 \pmod{p}$.

If you are not used to calculate with remainders (modulo), please simply accept the theorem or first read chapter 4 “Introduction to Elementary Number Theory with Examples”. What is important here is that this sentence implies that if this equation is not met for any whole number

¹⁵With the educational tool for number theory **NT** you can apply the tests of Fermat and of Miller-Rabin: See NT learning units 3.2 and 3.3, pages 3-11/11.

NT can be called in CT1 via the menu path **Indiv. Procedures \ Number Theory Interactive \ Learning Tool for Number Theory**. See also appendix A.6.

CT2 contains a visualisation of these methods within the tutorial “**World of Primes**”.

¹⁶Pierre de Fermat, French mathematician, Aug 17, 1601 – Jan 12, 1665.

a , then p is not a prime! The tests (e.g. for the first formulation) can easily be performed using the *test basis* $a = 2$.

This gives us a criterion for non-prime numbers, i.e. a negative test, but no proof that a number a is prime. Unfortunately Fermat's theorem does not apply — otherwise we would have a simple proof of the prime number property (or to put it in other words, we would have a simple prime number criterion).

Pseudo prime numbers

Numbers n that have the property

$$2^n \equiv 2 \pmod{n}$$

but are not prime are called *pseudo prime numbers* (i.e. the exponent is not a prime). The first pseudo prime number is

$$341 = 11 \cdot 31.$$

Carmichael numbers

There are pseudo prime numbers n that pass the Fermat test

$$a^{n-1} \equiv 1 \pmod{n}$$

with all bases a which are relatively prime to n [$\gcd(a, n) = 1$], despite these numbers n are not prime: These numbers are called *Carmichael numbers*. The first of these is

$$561 = 3 \cdot 11 \cdot 17.$$

Sample: The number to be tested is 561. Because $561 = 3 \cdot 11 \cdot 17$ it is:

The test condition $a^{560} \pmod{561} = 1$ is satisfied for $a = 2, 4, 5, 7, \dots$,
but not for $a = 3, 6, 9, 11, 12, 15, 17, 18, 21, 22, \dots$.

This means the test condition must not be satisfied for multiples of the prime factors 3, 11 or 17.

The test applied for $a = 3$ results in: $3^{560} \pmod{561} = 375$.

The test applied for $a = 5$ results in: $5^{560} \pmod{561} = 1$.

Strong pseudo prime numbers

A stronger test is provided by Miller/Rabin¹⁷: It is only passed by so-called *strong pseudo prime numbers*. Again, there are strong pseudo prime numbers that are not primes, but this is much less often the case than for (simple) pseudo prime numbers or for Carmichael numbers. The smallest strong pseudo prime number base 2 is

$$15,841 = 7 \cdot 31 \cdot 73.$$

¹⁷In 1976 an efficient probabilistic primality test was published by Prof. Rabin, based on a number theoretic result of Prof. Miller from the year before.

Prof. Miller worked at the Carnegie-Mellon University, School of Computer Science. Prof. Rabin, born in 1931, worked at the Harvard and Hebrew University.

If you test all 4 bases, 2, 3, 5 and 7, you will find only one strong pseudo prime number up to $25 \cdot 10^9$, i.e. a number that passes the test and yet is not a prime number.

More extensive mathematics behind the Rabin test delivers the probability that the number examined is prime (such probabilities are currently around 10^{-60}).

Detailed descriptions of tests for finding out whether a number is prime can be found on Web sites such as:

<http://www.utm.edu/research/primes/mersenne.shtml>
<http://www.utm.edu/research/primes/prove/index.html>

3.6 Special types of numbers and the search for a formula for primes

There are currently no useful, open (i.e. not recursive) formulae known that only deliver prime numbers (recursive means that in order to calculate the function the same function is used with a smaller variable). Mathematicians would be happy if they could find a formula that leaves gaps (i.e. does not deliver all prime numbers) but does not deliver any composite (non-prime) numbers.

Ideally, we would like, for the number n , to immediately be able to obtain the n -th prime number, i.e. for $f(8) = 19$ or for $f(52) = 239$.

Ideas for this can be found at

http://www.utm.edu/research/primes/notes/faq/p_n.html.

Cross-reference: the table under 3.10 contains the precise values for the n th prime numbers for selected n .

For “prime number formulae” usually very special types of numbers are used. The following enumeration contains the most common ideas for “prime number formulae”, and what our current knowledge is about very big elements of the number series: Is their primality proven? If their are compound numbers could their prime factors be determined?

3.6.1 Mersenne numbers $f(n) = 2^n - 1$ for n prime

As shown above, this formula seems to deliver relatively large prime numbers but - as for $n = 11$ [$f(n) = 2,047$] - it is repeatedly the case that the result even with prime exponents is not prime. Today, all the Mersenne primes having less than around 4,000,000 digits are known (M-39):

<http://perso.wanadoo.fr/yves.gallot/primes/index.html>

3.6.2 Generalized Mersenne numbers $f(k, n) = k \cdot 2^n \pm 1$ for n prime and k small prime / Proth numbers¹⁸

This first generalization of the Mersenne numbers creates the so called Proth numbers. There are (for small k) extremely quick prime number tests (see [Knuth1981]). This can be performed in practice using software such as the Proths software from Yves Gallot:

¹⁸Their names come from the French farmer Franois Proth (1852-1879). More famous as the Proth primes is the related Sierpinski problem: Find all numbers k , so that $k * 2^n + 1$ is composite for all $n \geq 0$. See table 3.1.

<http://www.prothsearch.net/index.html>.

3.6.3 Generalized Mersenne numbers $f(b, n) = b^n \pm 1$ / The Cunningham project

This is another possible generalisation of the Mersenne numbers. The **Cunningham project** determines the factors of all composite numbers that are formed as follows:

$$f(b, n) = b^n \pm 1 \quad \text{for } b = 2, 3, 5, 6, 7, 10, 11, 12$$

(b is not equal to multiples of bases already used, such as 4, 8, 9).

Details of this can be found at:

<http://www.cerias.purdue.edu/homes/ssw/cun>

3.6.4 Fermat numbers¹⁹ $f(n) = 2^{2^n} + 1$

As mentioned above in chapter 3.5, Fermat wrote to Mersenne regarding his assumption, that all numbers of this type are primes. This assumption was disproved by Euler (1732). The prime 641 divides $f(5)$.²⁰

$f(0) = 2^{2^0} + 1 = 2^1 + 1 = 3$	\mapsto prime
$f(1) = 2^{2^1} + 1 = 2^2 + 1 = 5$	\mapsto prime
$f(2) = 2^{2^2} + 1 = 2^4 + 1 = 17$	\mapsto prime
$f(3) = 2^{2^3} + 1 = 2^8 + 1 = 257$	\mapsto prime
$f(4) = 2^{2^4} + 1 = 2^{16} + 1 = 65,537$	\mapsto prime
$f(5) = 2^{2^5} + 1 = 2^{32} + 1 = 4,294,967,297 = 641 \cdot 6,700,417$	\mapsto NOT prime!
$f(6) = 2^{2^6} + 1 = 2^{64} + 1 = 18,446,744,073,709,551,617$ $= 274,177 \cdot 67,280,421,310,721$	\mapsto NOT prime!
$f(7) = 2^{2^7} + 1 = 2^{128} + 1 = (\text{see page 63})$	\mapsto NOT prime!

Within the project “Distributed Search for Fermat Number Dividers” offered by Leonid Durman there is also progress in finding new monster primes:

<http://www.fermatsearch.org/>

This website links to other web pages in Russian, Italian and German.

The discovered factors can be compound integers or primes.

On February 22, 2003 John Cosgrave discovered

- the largest composite Fermat number to date and
- the largest prime non-simple Mersenne number so far with 645,817 decimal digits.

The Fermat number

$$f(2, 145, 351) = 2^{(2^{2,145}, 351)} + 1$$

¹⁹The Fermat prime numbers play a role in circle division. As proven by Gauss a regular p -edge can only be constructed with the use of a pair of compasses and a ruler, when p is a Fermat prime number.

²⁰Surprisingly this number can easily be found by using Fermat’s theorem (see e.g. [Scheid1994, p. 176])

is divisible by the prime

$$p = 3 * 2^{2,145,353} + 1$$

At that time this prime p was the largest known prime generalized Mersenne number and the 5th largest known prime number at all.

This work was done using NewPGen from Paul Jobling's, PRP from George Woltman's, Proth from Yves Gallot's programs and also the Proth-Gallot group at St. Patrick's College, Dublin.

More details are in

http://www.fermatsearch.org/history/cosgrave_record.htm/

3.6.5 Generalized Fermat numbers²¹ $f(b, n) = b^{2^n} + 1$

Generalized Fermat numbers are more numerous than Mersenne numbers of a equal size and many of them are waiting to be discovered to fill the big gaps between the Mersenne primes already found or still undiscovered. Progress in number theory made it possible that numbers, where the representation is not limited to the base 2, can be tested at almost the same speed than a Mersenne number.

Yves Gallot wrote the program Proth.exe to investigate generalized Fermat numbers.

Using this program at February 16, 2003 Michael Angel discovered the largest of them till then with 628,808 digits, which at that time became the 5th largest known prime number:

$$b^{2^{17}} + 1 = 62,722^{131,072} + 1.$$

More details are in

<http://primes.utm.edu/top20/page.php?id=12>

3.6.6 Carmichael numbers

As mentioned above in chapter 3.5 not all Carmichael numbers are prime.

3.6.7 Pseudo prime numbers

See above in chapter 3.5.

3.6.8 Strong pseudo prime numbers

See above in chapter 3.5.

²¹The base of this power is no longer restricted to 2. Even more generic would be: $f(b, c, n) = b^{c^n} \pm 1$

3.6.9 Idea based on Euclid's proof $p_1 \cdot p_2 \cdots p_n + 1$

This idea is based on Euclid's proof that there are infinite many prime numbers.

$2 \cdot 3 + 1$	$= 7$	\mapsto prime
$2 \cdot 3 \cdot 5 + 1$	$= 31$	\mapsto prime
$2 \cdot 3 \cdot 5 \cdot 7 + 1$	$= 211$	\mapsto prime
$2 \cdot 3 \cdots 11 + 1$	$= 2,311$	\mapsto prime
$2 \cdot 3 \cdots 13 + 1$	$= 59 \cdot 509$	\mapsto NOT prime!
$2 \cdot 3 \cdots 17 + 1$	$= 19 \cdot 97 \cdot 277$	\mapsto NOT prime!

3.6.10 As above but -1 except $+1$: $p_1 \cdot p_2 \cdots p_n - 1$

$2 \cdot 3 - 1$	$= 5$	\mapsto prime
$2 \cdot 3 \cdot 5 - 1$	$= 29$	\mapsto prime
$2 \cdot 3 \cdots 7 - 1$	$= 11 \cdot 19$	\mapsto NOT prime!
$2 \cdot 3 \cdots 11 - 1$	$= 2,309$	\mapsto prime
$2 \cdot 3 \cdots 13 - 1$	$= 30,029$	\mapsto prime
$2 \cdot 3 \cdots 17 - 1$	$= 61 \cdot 8,369$	\mapsto NOT prime!

3.6.11 Euclidean numbers $e_n = e_0 \cdot e_1 \cdots e_{n-1} + 1$ with $n \geq 1$ and $e_0 := 1$

e_{n-1} is not the $(n-1)$ th prime number, but the number previously found here. Unfortunately this formula is not open but recursive. The sequence starts with

$e_1 = 1 + 1$	$= 2$	\mapsto prime
$e_2 = e_1 + 1$	$= 3$	\mapsto prime
$e_3 = e_1 \cdot e_2 + 1$	$= 7$	\mapsto prime
$e_4 = e_1 \cdot e_2 \cdot e_3 + 1$	$= 43$	\mapsto prime
$e_5 = e_1 \cdot e_2 \cdots e_4 + 1$	$= 13 \cdot 139$	\mapsto NOT prime!
$e_6 = e_1 \cdot e_2 \cdots e_5 + 1$	$= 3,263,443$	\mapsto prime
$e_7 = e_1 \cdot e_2 \cdots e_6 + 1$	$= 547 \cdot 607 \cdot 1,033 \cdot 31,051$	\mapsto NOT prime!
$e_8 = e_1 \cdot e_2 \cdots e_7 + 1$	$= 29,881 \cdot 67,003 \cdot 9,119,521 \cdot 6,212,157,481$	\mapsto NOT prime!

Also e_9, \dots, e_{17} are composite, which means that this formula is not particularly useful.

Comment:

However, what is special about these numbers is that any pair of them does not have a common factor other than 1²². Therefore they are *relatively prime*.

²²This can easily be shown via the following rule for the *greatest common divisor gcd* with $\gcd(a, b) = \gcd(b - \lfloor b/a \rfloor \cdot a, a)$.

We have for $i < j$:

$$\gcd(e_i, e_j) \leq \gcd(e_1 \cdots e_i \cdots e_{j-1}, e_j) = \gcd(e_j - e_1 \cdots e_i \cdots e_{j-1}, e_1 \cdots e_i \cdots e_{j-1}) = \gcd(1, e_1 \cdots e_i \cdots e_{j-1}) = 1.$$

See page 157.

3.6.12 $f(n) = n^2 + n + 41$

This sequence starts off very *promisingly*, but is far from being a proof.

$f(0) = 41$	\mapsto prime
$f(1) = 43$	\mapsto prime
$f(2) = 47$	\mapsto prime
$f(3) = 53$	\mapsto prime
$f(4) = 61$	\mapsto prime
$f(5) = 71$	\mapsto prime
$f(6) = 83$	\mapsto prime
$f(7) = 97$	\mapsto prime
\vdots	
$f(33) = 1,163$	\mapsto prime
$f(34) = 1,231$	\mapsto prime
$f(35) = 1,301$	\mapsto prime
$f(36) = 1,373$	\mapsto prime
$f(37) = 1,447$	\mapsto prime
$f(38) = 1,523$	\mapsto prime
$f(39) = 1,601$	\mapsto prime
$f(40) = 1681 = 41 \cdot 41$	\mapsto NOT prime!
$f(41) = 1763 = 41 \cdot 43$	\mapsto NOT prime!

The first 40 values are prime numbers (which have the obvious regularity that their difference starts with 2 and increases by 2 each time), but the 41th and 42th values are not prime numbers. It is easy to see that $f(41)$ cannot be a prime number: $f(41) = 41^2 + 41 + 41 = 41(41 + 1 + 1) = 41 \cdot 43$.

3.6.13 $f(n) = n^2 - 79 \cdot n + 1,601$

This function delivers prime numbers for all values from $n = 0$ to $n = 79$.²³ Unfortunately $f(80) = 1,681 = 11 \cdot 151$ is not a prime number. To this date, no function has been found that delivers more prime numbers in a row. On the other hand, each prime occurs twice (first in the decreasing then in the increasing sequence), which means that the algorithm delivers a total of 40 different prime values (these are the same ones as delivered by the function in chapter 3.6.12)²⁴.

²³See chapter 3.14, “Appendix: Examples using Sage” for the source code to compute the table using Sage.

²⁴Another quadratic polynom, which delivers these primes, is: $f(n) = n^2 - 9 \cdot n + 61$.

Among the first 1000 sequence elements more than 50% are prime (See chapter 3.14, “Appendix: Examples using Sage”).

$f(0) = 1.601 \rightarrow \text{prim}$	$f(26) = 223 \rightarrow \text{prim}$
$f(1) = 1.523 \rightarrow \text{prim}$	$f(27) = 197 \rightarrow \text{prim}$
$f(2) = 1.447 \rightarrow \text{prim}$	$f(28) = 173 \rightarrow \text{prim}$
$f(3) = 1.373 \rightarrow \text{prim}$	$f(29) = 151 \rightarrow \text{prim}$
$f(4) = 1.301 \rightarrow \text{prim}$	$f(30) = 131 \rightarrow \text{prim}$
$f(5) = 1.231 \rightarrow \text{prim}$	$f(31) = 113 \rightarrow \text{prim}$
$f(6) = 1.163 \rightarrow \text{prim}$	$f(32) = 97 \rightarrow \text{prim}$
$f(7) = 1.097 \rightarrow \text{prim}$	$f(33) = 83 \rightarrow \text{prim}$
$f(8) = 1.033 \rightarrow \text{prim}$	$f(34) = 71 \rightarrow \text{prim}$
$f(9) = 971 \rightarrow \text{prim}$	$f(35) = 61 \rightarrow \text{prim}$
$f(10) = 911 \rightarrow \text{prim}$	$f(36) = 53 \rightarrow \text{prim}$
$f(11) = 853 \rightarrow \text{prim}$	$f(37) = 47 \rightarrow \text{prim}$
$f(12) = 797 \rightarrow \text{prim}$	$f(38) = 43 \rightarrow \text{prim}$
$f(13) = 743 \rightarrow \text{prim}$	$f(39) = 41 \rightarrow \text{prim}$
$f(14) = 691 \rightarrow \text{prim}$	$f(40) = 41 \rightarrow \text{prim}$
$f(15) = 641 \rightarrow \text{prim}$	$f(41) = 43 \rightarrow \text{prim}$
$f(16) = 593 \rightarrow \text{prim}$	$f(42) = 47 \rightarrow \text{prim}$
$f(17) = 547 \rightarrow \text{prim}$	$f(43) = 53 \rightarrow \text{prim}$
$f(18) = 503 \rightarrow \text{prim}$	\dots
$f(19) = 461 \rightarrow \text{prim}$	$f(77) = 1.447 \rightarrow \text{prim}$
$f(20) = 421 \rightarrow \text{prim}$	$f(78) = 1.523 \rightarrow \text{prim}$
$f(21) = 383 \rightarrow \text{prim}$	$f(79) = 1.601 \rightarrow \text{prim}$
$f(22) = 347 \rightarrow \text{prim}$	$f(80) = 41 \cdot 41 \rightarrow \text{NOT prim!}$
$f(21) = 383 \rightarrow \text{prim}$	$f(81) = 41 \cdot 43 \rightarrow \text{NOT prim!}$
$f(22) = 347 \rightarrow \text{prim}$	$f(82) = 1.847 \rightarrow \text{prim}$
$f(23) = 313 \rightarrow \text{prim}$	$f(83) = 1.933 \rightarrow \text{prim}$
$f(24) = 281 \rightarrow \text{prim}$	$f(84) = 43 \cdot 47 \rightarrow \text{NOT prim!}$
$f(25) = 251 \rightarrow \text{prim}$	

3.6.14 Polynomial functions $f(x) = a_nx^n + a_{n-1}x^{n-1} + \dots + a_1x^1 + a_0$ (a_i in \mathbb{Z} , $n \geq 1$)

There exists no such polynomial that for all x in \mathbb{Z} only delivers prime values. For a proof of this, please refer to [Padberg1996, p. 83 f.], where you will also find further details about prime number formulae.

This means there is no hope in looking for further formulae (functions) similar to that in chap. 3.6.12 or chap. 3.6.13.

3.6.15 Catalan's Mersenne conjecture²⁵

Catalan conjectured²⁶ that C_4 and any further term in this sequence is a prime:

$$\begin{aligned} C_0 &= 2, \\ C_1 &= 2^{C_0} - 1, \\ C_2 &= 2^{C_1} - 1, \\ C_3 &= 2^{C_2} - 1, \\ C_4 &= 2^{C_3} - 1, \dots \end{aligned}$$

This sequence is defined recursively and increases extremely quickly. Does it only consist of primes?

$C(0) = 2$	\mapsto prime
$C(1) = 2^2 - 1 = 3$	\mapsto prime
$C(2) = 2^3 - 1 = 7$	\mapsto prime
$C(3) = 2^7 - 1 = 127$	\mapsto prime
$C(4) = 2^{127} - 1 = 170,141,183,460,469,231,731,687,303,715,884,105,727$	\mapsto prime

It is not (yet) known whether $C_5 = 2^{C_4} - 1$ and all higher elements are prime. In any case, it has not been proved that this formula delivers only primes.

It seems very unlikely that C_5 (or many of the larger terms) would be prime.

So this could be another example of Guy's "law of small numbers"²⁷.

3.6.16 Double Mersenne primes

The above Catalan-Mersenne numbers are from C_2 a subset of the double Mersenne primes.²⁸ A double Mersenne prime is a Mersenne prime of the form

$$M_{M_p} = 2^{2^p - 1} - 1$$

where p is a Mersenne prime exponent and M_p ist a prime Mersenne number.

The first values of p for which M_p is prime are $p = 2, 3, 5, 7, 13, 17, 19, 31, 61, 89, 107, 127, 521, \dots$ (see above).

M_{M_p} is known to be prime for $p = 2, 3, 5, 7$, and there has the according values: 7, 127, 2.147.483.647, 170.141.183.460.469.231.731.687.303.715.884.105.727.²⁹

²⁵Eugene Charles Catalan, Belgian mathematician, May 5, 1814–Feb 14, 1894.

After him, the so-called *Catalan numbers* $A(n) = (1/(n+1)) * (2n)!/(n!)^2$

$= 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012, 742900, 2674440, 9694845, \dots$ are named.

²⁶See <http://www.utm.edu/research/primes/mersenne.shtml> under "Conjectures and Unsolved Problems".

²⁷<http://primes.utm.edu/glossary/page.php?sort=LawOfSmall>

²⁸http://en.wikipedia.org/wiki/Catalan%20%93Mersenne_number

²⁹Sage calculates this with:

```
sage: for p in (2,3,5,7): Mp=(2^p)-1; MMp=(2^Mp)-1; B=is_prime(MMp); print p, Mp, MMp, B;
....:
2 3 7  True
3 7 127  True
5 31 2147483647  True
7 127 170141183460469231731687303715884105727  True
```

For $p = 11, 13, 17, 19$, and 31 , the corresponding double Mersenne numbers are not prime.

The next candidate for the next double Mersenne prime is $M_{M_{61}} = 2^{2305843009213693951} - 1$. Being approximately $1.695 * 10^{694,127,911,065,419,641}$, this number — like C_5 — is far too large for any currently known primality test to be successfully applied.

3.7 Density and distribution of the primes

As Euclid discovered, there is an infinite number of primes. However, some infinite sets are *denser* than others.

Within the set of natural numbers, there is an infinite number of even, uneven and square numbers. How to compare the “density” of two infinite sets is shown with even and square numbers.

The following proves that the even numbers are distributed more densely than square ones.³⁰

- the size of the n th element:

The n th element of the even numbers is $2n$; the n th element of the square numbers is n^2 .

Because for all $n > 2$: $2n < n^2$, the n th even number occurs much earlier than the n th square number.

- the numbers of values that are less than or equal to a certain *maximum value* x in \mathbb{R} are: There are $\lfloor x/2 \rfloor$ such even numbers and $\lfloor \sqrt{x} \rfloor$ square numbers. Because for all $x > 6$ the value $\lfloor x/2 \rfloor$ is greater than the largest integer smaller or equal to the square root of x , the even numbers are distributed more densely.

The value of the n -th prime $P(n)$

Theorem 3.7.1. *For large n : The value of the n -th prime $P(n)$ is asymptotic to $n \cdot \ln(n)$, i.e. the limit of the relation $P(n)/(n \cdot \ln n)$ is equal to 1 if n tends to infinity.*

For $n > 5$, $P(n)$ lies between $2n$ and n^2 . This means that there are fewer prime numbers than even natural numbers, but more prime numbers than square numbers.³¹

The number of prime numbers $PI(x)$

The definition for the number³² $PI(x)$ is similar: It is the number of all primes that do not exceed the maximum value x :

Theorem 3.7.2. *$PI(x)$ is asymptotic to $x/\ln(x)$.*

This is the **prime number theorem**. It was put forward by Legendre³³ and Gauss³⁴ but not proved until over 100 years later.

³⁰Whereas in colloquial language you often can hear, that “there are more” even numbers than square ones, mathematicians say, that from both there are infinitely many, that their sets are equivalent to \mathbb{N} (so both are infinite and countable, i.e. one can assign to each even number and to each square number an integer), but that the set of even numbers is denser than the set of square numbers.

³¹Please refer to the table 3.10

³²Often, instead of $PI(x)$ the convention $\Pi(x)$ is used.

³³Adrien-Marie Legendre, French mathematician, Sep 18, 1752 – Jan 10, 1833.

³⁴Carl Friedrich Gauss, German mathematician and astronomer, Apr 30, 1777–Feb 23, 1855.

The tables under 3.9 show the number of prime numbers in various intervals. The distribution is graphically presented within Figure 3.4 at page 91 within the appendix 3.13.

Formulae for the prime number theorem only apply when n tends to infinity, The formula of Gausee can be replaced by more precise formulae. For $x \geq 67$:

$$\ln(x) - 1,5 < x/PI(x) < \ln(x) - 0,5$$

Given that we know $PI(x) = x/\ln x$ only for very large x (x tending towards infinity), we can create the following overview:

x	$\ln(x)$	$x/\ln(x)$	$PI(x)(counted)$	$PI(x)/(x/\ln(x))$
10^3	6.908	144	168	1.160
10^6	13.816	72,386	78,498	1.085
10^9	20.723	48,254,942	50,847,534	1.054

For a binary number³⁵ x of the length of 250 bits (2^{250} is approximately $= 1.809251 * 10^{75}$) it is:

$$PI(x) = 2^{250}/(250 \cdot \ln 2) \text{ is approximately } = 2^{250}/173.28677 = 1.045810 \cdot 10^{73}.$$

We can therefore expect that the set of numbers with a bit length of less than 250 contains approximately 10^{73} primes (a reassuring result?!).

We can also express this as follows: Let us consider a *random* natural number n . Then the probability that this number is prime is around $1/\ln(n)$. For example, let us take numbers in the region of 10^{16} . Then we must consider $16 \cdot \ln 10 = 36,8$ numbers (on average) until we find a prime. A precise investigation shows: There are 10 prime numbers between $10^{16} - 370$ and $10^{16} - 1$.

Under the heading *How Many Primes Are There* at

<http://www.utm.edu/research/primes/howmany.shtml>

you will find numerous other details.

Using the following Web site:

<http://www.math.Princeton.EDU/~arbooker/nthprime.html>

you can easily determine $PI(x)$.

The **distribution** of primes³⁶ displays several irregularities for which no “system” has yet been found: On the one hand, many occur closely together, like 2 and 3, 11 and 13, 809 and 811, on the other hand large gaps containing no primes also occur. For example, no primes lie between 113 and 127, 293 and 307, 317 and 331, 523 and 541, 773 and 787, 839 and 853 as well as between 887 and 907.

For details, please see:

<http://www.utm.edu/research/primes/notes/gaps.html>

To discover the secrets of these irregularities is precisely part of what motivates mathematicians.

³⁵Number written in the binary system consists only of the digits 0 and 1.

³⁶Some visualizations (plots) of the quantity of primes in different number dimensions can be found in chapter 3.13, “Appendix: Visualization of the quantity of primes in higher ranges”.

Sieve of Eratosthenes

An easy way of calculating all $PI(x)$ primes less than or equal to x is to use the sieve of Eratosthenes. In the 3rd century B.C., he found an extremely easy, automatic way of finding this out. To begin with, you write down all numbers from 2 to x , circle 2, then cross out all multiples of 2. Next, you circle the lowest number that hasn't been circled or crossed out (3) and again cross out all multiples of this number, etc. You only need to continue until you reach the largest number whose square is less than or equal to x .³⁷

Apart from 2, prime numbers are never even. Apart from 2 and 5, prime numbers never end in 2, 5 or 0. So you only need to consider numbers ending in 1, 3, 7, 9 anyway (there are infinite primes ending in these numbers; see [Tietze1973, vol. 1, p. 137]).

Nowadays, you can now find large databases that contain either a large number of primes or the factorization of numerous composite numbers.

Further interesting topics regarding prime numbers

This chapter 3 didn't consider other number theory topics such as divisibility rules, modulus calculation, modular inverses, modular powers, modular roots, Chinese remainder theorem, Euler Phi function or perfect numbers. Some of these topics are considered in the **next chapter** (chapter 4).

³⁷With the educational tool for number theory **NT** you can apply the sieve of Eratosthenes in a computer-aided and guided way: Enter your own number and do the sieving step by step: See the NT learning unit 1.2, pages 6/21 and 7/21.

NT can be called in CT1 via the menu path **Indiv. Procedures \ Number Theory Interactive \ Learning Tool for Number Theory**. See appendix A.6.

CT2 contains a visualisation of this method within the tutorial “**World of Primes**”.

3.8 Notes about primes

The following notes list some interesting theorems, conjectures and open questions about primes, but also some quaint things and overviews.

3.8.1 Proven statements / theorems about primes

- For each number n in \mathbf{N} there are n consecutive natural numbers that are not primes. A proof of this can be found in [Padberg1996, p. 79].
- Paul Erdős³⁸ proved: Between each random number not equal to 1 and its double, there is at least one prime. He was not the first to prove this theorem, but proved it in a much simpler manner than those before him.
- There is a real number a such that the function $f : \mathbf{N} \rightarrow \mathbb{Z}$ where $n \mapsto \lfloor a^{3^n} \rfloor$ only delivers primes³⁹ for all n (see [Padberg1996, p. 82]). Unfortunately, problems arise when we try to determine a (see chapter 3.8.2).⁴⁰
- There are arithmetic prime sequences of arbitrary length.^{41,42}

In 1923 the famous British mathematician Hardy⁴³ compiled the conjecture, that there are arithmetic sequences of arbitrary length, which consist of primes only. This conjecture was proven in 2004 by two young American mathematicians.

At some point every school child learns about arithmetic number series. These are sequences of numbers, for which the difference between any 2 consecutive numbers is equal or constant (an arithmetic sequence must have at least three elements but can also have indefinitely many). In the sample sequence 5, 8, 11, 14, 17, 20 the difference between the series's elements is 3 and the length of the sequence is 6.

Arithmetic series have been known for millennia and one would think they have no more secrets. They get more interesting again, if we impose additional constraints on the series elements - as the prime example shows.

E.g. 5, 17, 29, 41, 53 is an arithmetic prime series which consists of 5 elements and the difference between the elements is always 12.

The sequence is not extendable - the next would be 65, but 65 is not prime (65 is the product of 5 and 13).

How many elements are possible within an arithmetic prime number sequence? Around 1770 the French Joseph-Louis Lagrange and the British Edward Waring investigated this

³⁸Paul Erdős, Hungarian mathematician, Mar 26, 1913–Sep 20, 1996.

³⁹The Gauss bracket $\lfloor x \rfloor$ of a real number x is defined via: $\lfloor x \rfloor$ is the next integer less or equal x .

⁴⁰If someone knows how to prove this, we are very interested to learn about this. Friedhelm Padberg told us, he doesn't have his prove any more.

⁴¹Sources:

- <http://users.cybercity.dk/~ds1522332/math/aprecords.htm> Original source
- <http://primes.utm.edu/glossary/page.php?sort=ArithmetiSequence> Original source
- http://en.wikipedia.org/wiki/Primes_in_arithmetic_progression
- http://en.wikipedia.org/wiki/Problems_involving_arithmetic_progressions
- http://en.wikipedia.org/wiki/Cunningham_chain
- German magazine GEO 10 / 2004: "Experiment mit Folgen"
- <http://www.faz.net> "Hardys Vermutung – Primzahlen ohne Ende" by Heinrich Hemme (July 06, 2004)

⁴²Arithmetic sequences with k primes are called prime arithmetic progressions and therefore their abbreviation is PAP- k or AP- k .

⁴³Godfrey Harold Hardy, British mathematician, Feb 7, 1877–Dec 1, 1947.

question. In 1923 the famous British mathematician Godfrey Harold Hardy and his colleague John Littlewood theorized, that there is no upper limit for the number of elements. But they could not prove this. In 1939 more progress was achieved: The Dutch mathematician Johannes van der Corput was able to prove that there are infinitely many different arithmetic prime number sequences with exactly three elements. Two examples are 3, 5, 7 and 47, 53, 59.

The longest arithmetic prime number sequence known today contains 25 elements. The table 3.3 lists the longest currently known arithmetic prime number sequences with minimal difference⁴⁴.

As a team, the two young⁴⁵ mathematicians Ben Green and Terence Tao, were able in 2004 to prove Hardy's conjecture, which had puzzled mathematicians for over 80 years: It states, that for any arbitrary length there exists an arithmetic prime number series. Additionally they managed to prove, that for any given length there are infinitely many different series.

Green and Tao intended to proof that there are infinitely many arithmetic sequences of length four. For this they considered sets of numbers consisting of primes and so called "near primes". These are numbers with a small set of divisors like numbers which are the product of two primes - these numbers are called "half primes". Thus they managed to considerably simplify their work because about near primes there already existed a lot of useful theorems. Finally they discovered that the results of their theorem were far more reaching than they had assumed and so they were able to prove Hardy's conjecture.

Any one who believes that it is easy to use Green's and Tao's 49 page proof to compute arithmetic prime number series of arbitrary length will soon become disappointed, because the proof is non-constructive. It is a so called proof of existence. This means that these mathematicians have shown "only" that these series exist, but not how to find them in practice.

This means that in the set of the natural numbers there is e.g. a series of one billion primes, which all have the same distance; and there are infinitely many of them. But these sequences lie extremely far beyond the numbers we usually use ("far outside").

⁴⁴In the opposite, http://en.wikipedia.org/wiki/Primes_in_arithmetic_progression lists the "largest known AP-k". Therefore, there the last sequence element is a prime as big as possible.

However, table 3.3 lists the sequences which have the smallest known difference for a given length.

⁴⁵Hardy wrote in his memoirs in 1940, that mathematics - more than all other arts and sciences - is an activity for young people.

At that time, 27-years-old Ben Green from the University of British Columbia in Vancouver and 29-year-old Terence Tao from the University of California in Los Angeles seem to confirm Hardy.

Elements	First element	Distance	When Digits	Discovered by
3	3	2	1	
4	5	6	2	
5	5	6	2	
6	7	30	1909 3	G. Lenaire
7	7	150	1909 3	G. Lenaire
.....				
21	28,112,131,522,731,197,609	9,699,690 = 19#	2008 20	Jaroslaw Wroblewski
22	166,537,312,120,867	96,599,212,710 = 9,95919#	2006 15	Markus Frind
23	403,185,216,600,637	2,124,513,401,010 = 9,52323#	2006 15	Markus Frind,
24	515,486,946,529,943	30,526,020,494,970 = 136,83123#	2008 16	Raanan Chermoni, Jaroslaw Wroblewski
25	6,171,054,912,832,631	81,737,658,082,080 = 366,38423#	2008 16	Raanan Chermoni, Jaroslaw Wroblewski

Table 3.3: Arithmetic prime number sequences with minimal difference (as of Aug. 2012)

If someone wants to discover such sequences he should consider the following thought. The length of a sequence determines the minimal common distance between the single primes of the sequence. Given a sequence with 6 elements the distance between them has to be 30 or a multiple of 30. The number 30 results as the product of all primes smaller than the length of the sequence. So its the product of all primes smaller than 6: $6\# = 5\# = 2 * 3 * 5 = 30$. It is $10\# = 7\# = 2357 = 210$. If you look for a sequence with 15 elements, then the common distance is at least $15\# = 13\# = 2 * 3 * 5 * 7 * 11 * 13 = 30,030$.

This means that the length of an arithmetic prime sequence can be arbitrary big, but the distance between the elements cannot be any arbitrary number. E.g. there is no arithmetic prime sequence with the distance 100, because 100 cannot be divided by 3.

k	k#
2	2
3	6
5	30
7	210
11	2,310
13	30,030
17	510,510
19	9,699,690
23	223,092,870

Table 3.4: Products of the first primes $\leq k$ (called k primorial or $k\#$)

Further restriction:

If you look at arithmetic prime sequences, which fulfill the *additional* requirement, that all primes are *consecutive*⁴⁶, then its getting even more complicated. At the website of Chris Caldwell⁴⁷ you can find further details: The longest known arithmetic prime sequence, consisting only of directly consecutive primes (as of Aug. 2012), has a length of 10, the common distance is

$$10\# = 7\# = 2 * 3 * 5 * 7 = 210$$

and starts with the 93-digit prime

100 9969724697 1424763778 6655587969 8403295093 2468919004 1803603417 7589043417
0334888215 9067229719

⁴⁶They are also called consecutive prime arithmetic progressions and therefore their abbreviation is CPAP-k or CAP-k.

⁴⁷<http://primes.utm.edu/glossary/page.php?sort=ArithmetSequence>

3.8.2 Unproven statements / conjectures about primes / further open questions

- Goldbach⁴⁸ conjectured: Every even natural number greater than 2 can be represented as the sum of two prime numbers.
- Riemann⁴⁹ put forward an important, but still unproved hypothesis⁵⁰ about the location of the nontrivial zeros of the Riemann zeta function. A consequence is an improved estimation within the prime number theorem (distribution of primes) .
- Benford's law^{51,52} does not apply to primes.

According to Benford's law, also called the first-digit law, the single digits in lists of numbers from many (but not all) real-life sources of data, are distributed in a non-uniform way. Especially the leading digit is much more often the digit 1 than any other digit.

Which empirical data applies to this "law" is not completely clear yet. Timo Eckhardt analyzed in his thesis in 2008 extensively attributes of prime numbers. For example, all primes until 7,052,046,499 were described with different bases of the positional notation.

Comparing the bases 3 to 10 the deviation from Benford's law was lowest with base 3. Comparing the first digit for base 10 all digits are almost equally distributed. Analyzing bigger bases showed strong differences.

- The proof (mentioned above in chapter 3.8.1) of the function $f : N \rightarrow Z$ with $n \mapsto \lfloor a^{3^n} \rfloor$ only guarantees the existence of such a number a . How can we determine this number a and will it have a value, making the function also of some practical interest?
- Is there an infinite number of Mersenne prime numbers?
- Is there an infinite number of Fermat prime numbers?
- Does a polynomial time algorithm exist for calculating the prime factors of a number (see [Klee1997, p. 167])? This question can be divided into the three following questions:
 - Does a polynomial time algorithm exist that decides whether a number is prime?
This question has been answered by the AKS algorithm (see chapter 4.11.5.3, "Primes in P": Primality testing is polynomial).
 - Does a polynomial time algorithm exist that calculates for a composite number from how many prime factors it is made up (without calculating these factors)?
 - Does a polynomial time algorithm exist that calculates for a composite number n a non-trivial (i.e. other than 1 and n) factor of n ?⁵³

⁴⁸Christian Goldbach, German mathematician, Mar 18, 1690–Nov 20, 1764.

⁴⁹Bernhard Riemann, German mathematician, Sep 17, 1826–Jul 20, 1866.

⁵⁰http://en.wikipedia.org/wiki/Riemann_hypothesis

⁵¹http://en.wikipedia.org/wiki/Benford%27s_law,
http://arxiv.org/PS_cache/arxiv/pdf/0906/0906.2789v1.pdf.

⁵²Two good didactical articles in German about applications of Benford's law are:

- Rüdiger Baumann: "Ziffernanalyse zwecks Betrugsaufdeckung — Beispiel für kompetenzorientierten und kontextbezogenen Informatikunterricht",
in LOGIN, Informatische Bildung und Computer in der Schule, No. 154/155, 2008, p. 68-72
- Norbert Hungerbühler: "Benfords Gesetz über führende Ziffern",
March 2007, <http://www.educ.ethz.ch/unt/um/mathe/ana/benford>

⁵³Please compare chapters 4.11.5.1 and 4.11.4.

At the end of chapter 4.11.4, section RSA-200 you can see the limits for which numbers the current algorithms for primality testing and for factorization deliver good results.

3.8.3 The Goldbach conjecture

Here we will have a closer look at what is called the **Goldbach Conjecture**⁵⁴.

The weak Goldbach conjecture⁵⁵

Goldbach conjectured in 1742 within a letter to the mathematician Euler:

Every **odd** natural number greater than 5 can be represented as the sum of **exactly three** prime numbers.

Samples: $7 = 3 + 2 + 2$ or $27 = 19 + 5 + 3$ or $27 = 17 + 5 + 5$

This conjecture is still (since more than 250 years) unproven. A preliminary work for a prove could be the prove published recently⁵⁶ by Terence Tao from the University of California. He proved that every **odd** natural number greater than 1 can be represented as the sum of **at most five** prime numbers.

The strong Goldbach conjecture⁵⁷

Goldbach's strong prime number hypothesis was formulated by Euler after a mail exchange with Goldbach. This is now called **the** Goldbach conjecture:

Every **even** natural number greater than 2 can be represented as the sum of **exactly two** prime numbers.

Samples of Goldbach partitions: $8 = 5 + 3$ or $28 = 23 + 5$

Computers have verified⁵⁸ the Goldbach conjecture for all even numbers up to $3 * 10^{17}$ (as at

⁵⁴http://en.wikipedia.org/wiki/Goldbach%27s_conjecture

⁵⁵It is also known as the odd or ternary Goldbach conjecture.

⁵⁶<http://arxiv.org/abs/1201.6656>, submitted on 31 Jan 2012 (v1), last revised 3 Jul 2012 (v4)

⁵⁷It is also known as the even or binary Goldbach conjecture.

⁵⁸It is generally accepted today, that the **Goldbach conjecture** is true, i. e. valid for all even natural numbers greater than 2. In 1999, mathematician Jörg Richstein from the computer sciences institute at the University of Giessen, studied even numbers up to 400 billion ($4 * 10^{14}$) and found no contradictory example ([Richstein1999]). In the meantime further progress was made: See

<http://www.ieeta.pt/~tos/goldbach.html> by Toms Oliveira e Silva,

http://en.wikipedia.org/wiki/Goldbach%27s_conjecture,

<http://primes.utm.edu/glossary/page.php?GoldbachConjecture.html>.

Nevertheless, this does not provide us with general proof.

The fact is that despite all efforts, Goldbach's conjecture has to date not been proven. This leads one to believe that since the pioneer work of the Austrian mathematician Kurt Gödel is well-known, not every true mathematical theorem is provable (see <http://www.mathematik.ch/mathematiker/goedel.html>). Perhaps Goldbach's conjecture was correct, but in any case the proof will never be found. Conversely, that will presumably also remain unproven.

September 2012), but no general proof has yet been found.^{59,60,61}

As bigger an even number is, as more umso such binary Goldbach partitions can be found – in average: For 4 there is only one partition $2 + 2$; for 16 there are two, $3 + 13$ and $5 + 11$. With 100 there are six such partitions: $3 + 97$, $11 + 89$, $17 + 83$, $29 + 71$, $41 + 59$, $47 + 53$.⁶²

Interconnection between the two Goldbach conjectures

If the strong Goldbach conjecture holds, then also the weak is true (so the strong implies the weak conjecture).

The proof for this is relativly simple:

Prerequisite: u is an odd number bigger than 5.

Each such odd number u can be written as sum $u = (u - 3) + 3$. The first summand then is even and ≥ 4 , so it fullfils the prerequisite of the strong Goldbach conjecture, and can be written as sum of two primes p_1 und p_2 (p_1 and p_2 are not necessarily different). So we found a partition of u into the three primes p_1, p_2 and 3. This means, its always possible to find such a sum, where one of the primes is the number 3.

Similarly easy it can be shown, that the weak Goldbach conjecture implies the above mentioned conjecture from Terence Tao (both hold for odd numbers):

- For odd numbers $u > 5$ directly the weak Goldbach conjecture implies that the sum consists at most of five primes.
- For the remaining odd numbers 3 and 5 you can directly check it:
 $3 = 3$ (the “sum” has only one and therefore at most five prime summands);
 $5 = 2 + 3$ (the sum has two and therefore at most five prime summands).

3.8.4 Open questions about twin primes

Twin primes are prime numbers whose difference is 2. Examples include 5 and 7, or 101 and 103, or $1,693,965 \cdot 2^{66,443} \pm 1$.

The biggest known twin pair nowadays is

$$65,516,468,355 \cdot 2^{333,333} \pm 1$$

⁵⁹The English publisher *Faber* and the American publisher *Bloomsbury* issued in 2000 the 1992 published book “Uncle Petros and Goldbach’s Conjecture” by Apostolos Doxiadis. It’s the story of an old maths professor who fails to prove a more than 250 year old puzzle. To boost the sales figures the English and American publishers have offered a prize of 1 million USD, if someone can prove the conjecture – which should be published by 2004 in a well-known mathematical journal.

Surprisingly only British and American citizens are allowed to participate.

⁶⁰The theorem which has come closest so far to Goldbach’s conjecture was proved by Chen Jing-Run in 1966 in a way which is somewhat hard to understand: Each even integer greater than 2 is the sum of one prime and of the product of two primes. E.g.: $20 = 5 + 3 * 5$.

Most of the research about the Goldbach conjecture is collected in the book: “Goldbach Conjecture”, ed. Wang Yuan, 1984, World scientific Series in Pure Maths, Vol. 4.

⁶¹Especially this conjecture makes it clear, that even today we do not have a complete understanding of the deeper connections between addition and multiplication of natural numbers.

⁶²Progress in research about the Goldbach conjectures is often a topic in mathematical press column, like <http://oumathclub.wordpress.com/>.

with 100,355 decimal digits.⁶³

Remark: Triplet primes, however, only occur once: 3, 5, 7. For all other sets of three consecutive uneven numbers, one of them is always divisible by 3 and thus not a prime.

Open questions are:

- The number of twin primes: infinitely many or a limited number?
- Does a formula exist for calculating the number of twin primes per interval?

A big step towards the solution of this problem was made by Dan Goldston, János Pintz and Cem Yıldırım in 2003.⁶⁴ The three mathematicians were investigating the distribution of prime numbers. They could proof, that

$$\liminf_{n \rightarrow \infty} \frac{p_{n+1} - p_n}{\log p_n} = 0,$$

where p_n denotes the n -th prime number.

This means that the smallest limit point (\liminf) of the sequence $\frac{p_{n+1} - p_n}{\log p_n}$ equals zero.

A point is called limit point of a sequence, if there lie in any arbitrarily small neighbourhood of that point infinitely many elements of the sequence.

$\log p_n$ is about the average distance between the prime p_n and the next prime p_{n+1} .

Hence, the term above implies, that there are infinitely many consecutive primes with a gap between them which is arbitrarily small compared to the expected average gap.

Moreover, it was proofed, that

$$p_{n+1} - p_n < (\log p_n)^{8/9}$$

holds true for infinitely many primes⁶⁵.

Those results could be the basis for the proof, that infinitely many twin primes exist.

⁶³<http://primes.utm.edu>, October 2009

⁶⁴D. A. Goldston: "Gaps Between Primes"

<http://www.math.sjsu.edu/~goldston/OberwolfachAbstract.pdf>

See also:

- D. A. Goldstone: "Are There Infinitely Many Twin Primes?"
<http://www.math.sjsu.edu/~goldston/twinprimes.pdf>
- K. Soundararajan: "Small Gaps Between Prime Numbers: The Work Of Goldston-Pintz-Yıldırım"
<http://www.ams.org/bull/2007-44-01/S0273-0979-06-01142-6/S0273-0979-06-01142-6.pdf>

⁶⁵c't magazine 2003, no. 8, page 54

3.8.5 Quaint and interesting things around primes⁶⁶

Primes are not only a very active and serious research area in mathematics. Also a lot of people think about them in their free time and outside the scientific research.

3.8.5.1 Recruitment at Google in 2004

In summer 2004 the company Google used the number e to attract potential employees.^{67,68}

On a prominent billboard in California's Silicon Valley on July 12 there appeared the following mysterious puzzle:

(first 10 digit prime in consecutive digits of e).com

Finding the first 10 digit prime in the decimal expansion of e is not easy, but with various software tools, one can determine that the answer is

7,427,466,391

Then if you visited the website www.7427466391.com, you were presented with an even more difficult puzzle. Figuring this second puzzle out took you to a web page that asks you, to submit your CV to Google. The ad campaign got high attention.

Presumably Google's conceit was that if you're smart enough to solve the puzzles, you're smart enough to work for them. Of course some days after the launch, anyone who really wanted to discover the answers without incurring a headache could merely do a Google search for them, since many solvers immediately posted their solutions online.⁶⁹

3.8.5.2 Contact [movie, 1997] – Primes helping to contact aliens

The movie, directed by Robert Zemeckis, originated from Carl Sagan's book with the same title.

After years of unavailing search the radio astronomer Dr. Ellie Arroway (Jodie Foster) discovers signals from the solar system Vega, 26 light years away. These signals contain the primes in the right order and without a gap. This makes the hero confident, that this message is different from the radio signals which permanently hit earth and which are random and of cosmic origin (radio galaxies, pulsars). In an unmasking scene a politician asks her after that, why these intelligent aliens didn't just speak English ...

⁶⁶Further curious things about primes may be found at:

- <http://primes.utm.edu/curios/home.php>
- <http://www.primzahlen.de/files/theorie/index.htm>.

⁶⁷The base of the natural logarithm e is approximately 2.718 281 828 459. This is one of the most important numbers in all of mathematics like complex analysis, finance, physics and geometry. Now it was used the first time – as far as I know – for marketing or recruitment.

⁶⁸Most of this information is taken from the article “e-number crunching” by John Allen Paulos in TheGuardian, Sept. 30, 2004, and from the web:

- <http://www.mkaz.com/math/google/>
- <http://epramono.blogspot.com/2004/10/7427466391.html>
- <http://mathworld.wolfram.com/news/2004-10-13/google/>
- <http://www.math.temple.edu/~paulos/>.

⁶⁹The second level of the puzzle, which involved finding the 5th term of a given number sequence had nothing to do with primes any more.

Doing communication with absolute strange and unknown beings from deep space is very hard especially because of 2 reasons: First the big distance and therefore the long transfer time makes it impossible to exchange within an average lifetime more than one message in each direction ne after the other. Secondly the first contact must give the receiver of the radio signals a good chance to notice the message and to categorize it as something from intelligent beings. Therefore the aliens send numbers at the beginning of their message, which can be considered as the easiest part of any higher language, and which are not too trivial: so they chose the sequence of primes. These special numbers play such a fundamental role in mathematics that one can assume that they are well known to each species who has the technical know-how to receive radio waves.

The aliens then send a plan to build a mysterious machine ...

3.9 Appendix: Number of prime numbers in various intervals

Ten-sized intervals		Hundred-sized intervals		Thousand-sized intervals	
Interval	Number	Interval	Number	Interval	Number
1-10	4	1-100	25	1-1000	168
11-20	4	101-200	21	1001-2000	135
21-30	2	201-300	16	2001-3000	127
31-40	2	301-400	16	3001-4000	120
41-50	3	401-500	17	4001-5000	119
51-60	2	501-600	14	5001-6000	114
61-70	2	601-700	16	6001-7000	117
71-80	3	701-800	14	7001-8000	107
81-90	2	801-900	15	8001-9000	110
91-100	1	901-1000	14	9001-10000	112

Table 3.5: How many primes exist within the first intervals of tens, of hundreds and of thousands?

Dimension	Interval	Number	Average number per 1000
4	1 - 10,000	1,229	122.900
5	1 - 100,000	9,592	95.920
6	1 - 1,000,000	78,498	78.498
7	1 - 10,000,000	664,579	66.458
8	1 - 100,000,000	5,761,455	57.615
9	1 - 1,000,000,000	50,847,534	50.848
10	1 - 10,000,000,000	455,052,512	45.505

Table 3.6: How many primes exist within the first intervals of dimensions?

A visualization of the number of primes in higher intervals of powers of 10 can be found in chapter 3.13 at page 91.

3.10 Appendix: Indexing prime numbers (n -th prime number)

Index	Precise value	Rounded value	Comment
1	2	2	
2	3	3	
3	5	5	
4	7	7	
5	11	11	
6	13	13	
7	17	17	
8	19	19	
9	23	23	
10	29	29	
100	541	541	
1,000	7,917	7,917	
664,559	9,999,991	9.99999E+06	All prime numbers up to 1E+07 were known at the beginning of the 20th century.
1E+06	15,485,863	1.54859E+07	
6E+06	104,395,301	1.04395E+08	
1E+07	179,424,673	1.79425E+08	
1E+09	22,801,763,489	2.28018E+10	
1E+12	29,996,224,275,833	2.99962E+13	

Table 3.7: List of particular n -th prime numbers $P(n)$

Comment:

With gaps, extremely large prime numbers were discovered at an early stage.

Web links:

<http://www.math.Princeton.EDU/~arbooker/nthprime.html>

http://www.utm.edu/research/primes/notes/by_year.html.

3.11 Appendix: Orders of magnitude / dimensions in reality

In the description of cryptographic protocols and algorithms, numbers occur that are so large or so small that they are inaccessible to our intuitive understanding. It may therefore be useful to provide comparative numbers from the real world around us so that we can develop a feeling for the security of cryptographic algorithms. Some of the numbers listed below originate from [Schwenk1996] and [Schneier1996, p.18].

Probability that you will be hijacked on your next flight	$5.5 \cdot 10^{-6}$
Annual probability of being hit by lightning	10^{-7}
Probability of 6 correct numbers in the lottery	$7.1 \cdot 10^{-8}$
Risk of being hit by a meteorite	$1.6 \cdot 10^{-12}$
Time until the next ice age (in years)	$14,000 = (2^{14})$
Time until the sun dies (in years)	$10^9 = (2^{30})$
Age of the earth (in years)	$10^9 = (2^{30})$
Age of the universe (in years)	$10^{10} = (2^{34})$
Number of molecules within one waterdrop	$10^{20} = (2^{63})$
Number of bacteria living on earth	$10^{30.7} = (2^{102})$
Number of the earth's atoms	$10^{51} = (2^{170})$
Number of the sun's atoms	$10^{57} = (2^{190})$
Number of atoms in the universe (without dark material)	$10^{77} = (2^{265})$
Volume of the universe (in cm^3)	$10^{84} = (2^{280})$

Table 3.8: Likelihoods and dimensions from physics and everyday life

3.12 Appendix: Special values of the binary and decimal system

These values can be used to evaluate from a key length in bit the corresponding number of possible keys and the search effort (if assumed, that e.g. one million keys can be tested within one second).

Binary system	Decimal system
2^{10}	1024
2^{40}	$1.09951 \cdot 10^{12}$
2^{56}	$7.20576 \cdot 10^{16}$
2^{64}	$1.84467 \cdot 10^{19}$
2^{80}	$1.20893 \cdot 10^{24}$
2^{90}	$1.23794 \cdot 10^{27}$
2^{112}	$5.19230 \cdot 10^{33}$
2^{128}	$3.40282 \cdot 10^{38}$
2^{150}	$1.42725 \cdot 10^{45}$
2^{160}	$1.46150 \cdot 10^{48}$
2^{192}	$6.27710 \cdot 10^{57}$
2^{250}	$1.80925 \cdot 10^{75}$
2^{256}	$1.15792 \cdot 10^{77}$
2^{320}	$2.13599 \cdot 10^{96}$
2^{512}	$1.34078 \cdot 10^{154}$
2^{768}	$1.55252 \cdot 10^{231}$
2^{1024}	$1.79769 \cdot 10^{308}$
2^{2048}	$3.23170 \cdot 10^{616}$

Table 3.9: Special values of the binary and decimal systems

Such tables can easily be calculated using computer algebra systems. Here is a code sample for Sage:

Sage sample 3.1 Special values of the binary and decimal systems

```
E = [10, 40, 56, 64, 80, 90, 112, 128, 150, 160, 192, 256, 1024, 2048]
for e in E:
    # print "2^" + str(e), "---", 1.0*(2^e)
    print "2^%4d" % e , " --- ", RR(2^e).n(24)
....:
2^ 10   ---  1024.00
2^ 40   ---  1.09951e12
2^ 56   ---  7.20576e16
2^ 64   ---  1.84467e19
2^ 80   ---  1.20893e24
2^ 90   ---  1.23794e27
2^ 112  ---  5.19230e33
2^ 128  ---  3.40282e38
2^ 150  ---  1.42725e45
2^ 160  ---  1.46150e48
2^ 192  ---  6.27710e57
2^ 256  ---  1.15792e77
2^1024  ---  1.79769e308
2^2048  ---  3.23170e616
```

3.13 Appendix: Visualization of the quantity of primes in higher ranges

The distribution of primes

Between 1 and 10 there are 4 primes. Between 10^3 and 10^4 there are already 1061 primes. In the interval $[10^9, 10^{10}]$ lie $404,204,977 \approx 4 \cdot 10^8$ primes, and in the interval from 10^{19} to 10^{20} there are $1,986,761,935,284,574,233 \approx 1,9 \cdot 10^{18}$ primes.⁷⁰

Why is the difference between the number of primes in the different intervals so big, although the boundaries of the intervals differ only by value 1 of the exponent of the power of 10?

The prime number theorem

The number $\text{PI}(x)$ of primes up to a given number x can approximately be determined by a formula, derived from the so called prime number theorem (see chapter 3.7). $\text{PI}(x)$ denotes the number of primes which are smaller or equal to x . Then the formula is

$$\text{PI}(x) \sim \frac{x}{\ln x}.$$

Note, that this formula only gives an approximation of the number of primes smaller or equal to x . It's getting more exact as the number x increases.

In the following we are using the prime number theorem to examine the distribution of primes.

To understand, why the number of primes is growing so rapidly, although the boundaries of the intervals only differ by the exponent 1, let's have a closer look to both components of the right side of the formula: x and $\ln x$.

The functions x and 10^x

The function x is a straight line. It is shown in figure 3.1a on page 90.

In the next step the function of the boundaries of the intervals are drawn in figure 3.1b on page 90. To get an idea of how the functions look like, the domain of definition was chosen to be from 0 to 10^{10} and from 0 to 10, respectively. You can see, that with increasing exponent x the numbers grow stronger.

The function $\ln x$

In comparison to that we consider the function $\ln x$. The left picture of figure 3.2 on page 90 shows the graph with the domain of definition from 1 to 100. On the right picture the domain of definition was chosen between 1 and 10^{10} .

One can see that the values of the function $\ln x$ grow slowly compared to the growth of the function x . This is visualized by the graph of both functions in one picture shown in figure 3.3 on page 90. In addition to that the graph of the function $\frac{x}{\ln x}$ was drawn in the same figure.

⁷⁰http://en.wikipedia.org/wiki/Prime_number_theorem

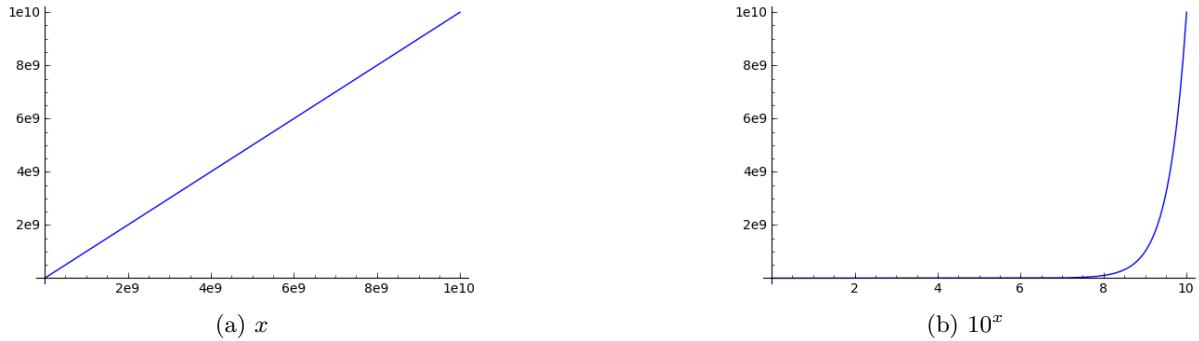


Figure 3.1: Graph of the functions x and 10^x

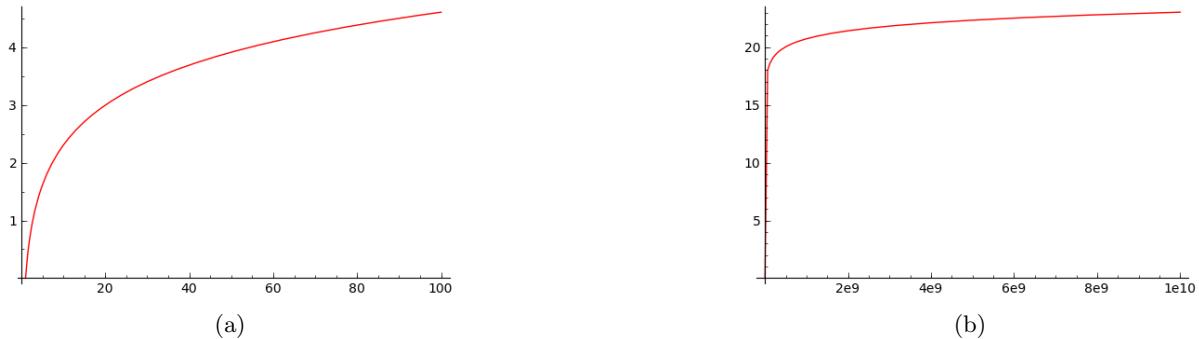


Figure 3.2: Graph of the function $\ln x$ till 100 and till 10^{10}

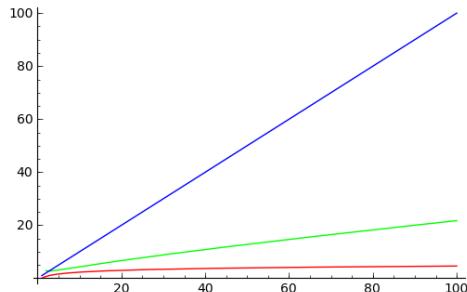


Figure 3.3: The functions x (blue), $\ln x$ (red) and $\frac{x}{\ln x}$ (green)

The function $PI(x) = \frac{x}{\ln x}$

The function $\frac{x}{\ln x}$ consists of the function x as the numerator and the function $\ln x$ in the denominator, which, in comparison to x , increases very slowly. Compared to the number x itself, the number of primes less or equal to x is small. But still, $\frac{x}{\ln x}$ is an increasing function as you can see in figure 3.3 on page 90.

The number of primes in the different intervals

Figure 3.4 visualizes how the number of primes in the intervals $[1, 10^x]$ and $[10^{x-1}, 10^x]$ behave. To calculate it faster, the result of the approximation function is used (not the exact numbers like in the tables in chapter 3.9).

Here for each base 10 exponent two bars are drawn: $\frac{10^x}{\ln 10^x}$ and $\frac{10^x}{\ln 10^x} - \frac{10^{x-1}}{\ln 10^{x-1}}$: The left chart shows the values for the exponents x from 1 to 5, and the right one for x from 1 to 10, where x is the base 10 exponent.

The blue bars represent the overall number of primes up to 10^x . The red bars show how many primes accrue in the interval $[10^{x-1}, 10^x]$, respectively. This makes clear, that the number of primes in intervals of higher exponents keeps growing quite fast.

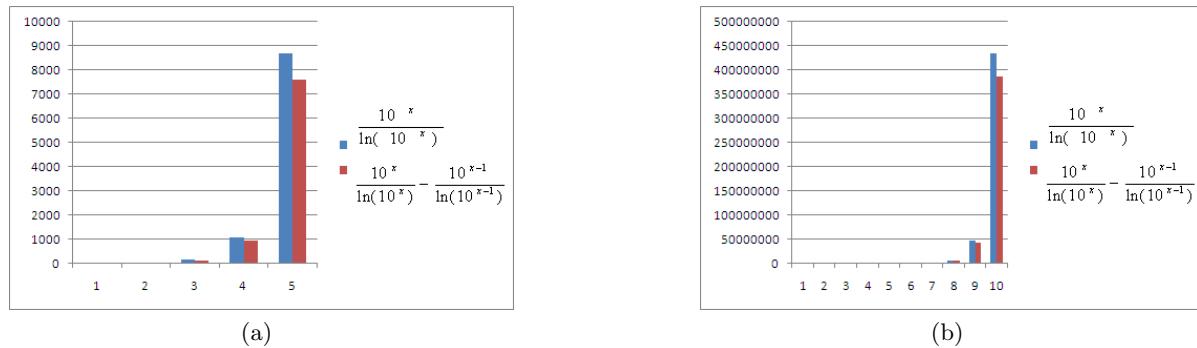


Figure 3.4: Numbers of primes in the interval $[1, 10^x]$ (blue) and in the interval $[10^{x-1}, 10^x]$ (red) for different exponents x

A table containing the number of primes in some dedicated intervals can be found in chapter 3.9 at page 85: For example, within the interval $[1, 10^4]$ there are 1229 primes; thereof are in the interval $[10^3, 10^4]$ $1229 - 168 = 1061$ primes.

Theory about the prime number theorem and the function $\text{PI}(x)$ can be found in chapter 3.7.

Sage sample 3.2 Generation of the graphs of the three functions x , $\log(x)$ and $x/\log(x)$

```
# Definition of function f(x)=x and plots for the domains from 0 to 10^10 and 0 to 100
sage: def f(x):return x
....:
sage: F=plot(f,(0,10^10))
sage: F.plot()

sage: F2=plot(f,(1,100))
sage: F2.plot()

# Definition of function g(x)=10^x and plots for the domain from 0 to 10
sage: def g(x): return 10^x
....:
sage: G=plot(g,(0,10))
sage: G.plot()

# Definition of function h(x)=log(x) and plots for the domains from 1 to 100 and 1 to 10^10
sage: def h(x): return log(x)
....:
sage: H=plot(h,(1,100),color="red")
sage: H.plot()

sage: H2=plot(h,(1,10^10),color="red")
sage: H2.plot()

# Definition of function k(x)=x/log(x) and plots for the domain from 2 to 100
sage: def k(x): return x/log(x)
....:
sage: K=plot(k,(2,100),color="green")
sage: K.plot()

# Plots of the functions f, k and h for the domain of definition up to 100
sage: F2+K+H

# Generation of the data for the bar charts .....
# Determination of the number of primes in the interval [1,10]
sage: pari(10).primepi()-pari(1).primepi()
4

# Determination of the number of primes in the interval [10^3,10^4]
sage: pari(10**4).primepi()-pari(10**3).primepi()
1061

# Determination of the number of primes in the interval [10^8,10^9]
sage: pari(10**9).primepi()-pari(10**8).primepi()
45086079

# (for 10^10: OverflowError: long int too large to convert)
```

3.14 Appendix: Examples using Sage

Below is Sage source code related to contents of the chapter 3 (“Prime Numbers”).

3.14.1 Some basic functions about primes using Sage

This part of the appendix contains Sage code, to perform some simple computations about primes.⁷¹

Sage sample 3.3 Some basic functions about primes

```
# primes (general commands)
# The set of prime numbers
sage: P=Primes(); P
Set of all prime numbers: 2, 3, 5, 7, ...

# Returns the next prime number
sage: next_prime(5)
7

# Returns how many primes <=x are there
sage: pari(10).primepi()
4

# Returns the first x primes
sage: primes_first_n(5)
[2, 3, 5, 7, 11]

# Returns the primes in an interval
sage: list(primes(1,10))
[2, 3, 5, 7]
```

⁷¹See the Sage documentation about Elementary number theory http://www.sagemath.org/doc/constructions/number_theory.html.

3.14.2 Check primality of integers generated by quadratic functions

The following Sage code verifies the primality of integers generated by the function $f(n) = n^2 - 9n + 61$. The code defines a function called `quadratic_prime_formula()` that takes three arguments:

- `start` — An integer which is the lower bound for integers in the sequence `start, start + 1, start + 2, ..., end - 1, end`.
- `end` — An integer which is the upper bound for the integers in the sequence `start, start + 1, start + 2, ..., end - 1, end`.
- `verbose` — (default: `True`) a flag to signify whether to print a message indicating the primality of an integer generated by $f(n)$.

A meaningful modification of this code is to use another function, of which the primality of its function values should be checked.

Sage sample 3.4 Verify the primality of integers generated by a quadratic function

```
def quadratic_prime_formula(start, end, verbose=True):
    print "N -- N^2 - 9*N + 61"
    P = 0 # the number of primes between start and end
    for n in xrange(start, end + 1):
        X = n^2 - 9*n + 61
        if is_prime(X):
            P += 1
            if verbose:
                print str(n) + " -- " + str(X) + " is prime"
        else:
            if verbose:
                print str(n) + " -- " + str(X) + " is NOT prime"
    print "Number of primes: " + str(P)
    print "Percentage of primes: " + str(float((P * 100) / (end - start + 1)))
```

With the following function call we compute the values of $f(n) = n^2 - 9n + 61$ for $n = 0, 1, 2, \dots, 50$ and verify the primality of the generated integers:

```
sage: quadratic_prime_formula(0, 50)
N -- N^2 - 9*N + 61
0 -- 61 is prime
1 -- 53 is prime
2 -- 47 is prime
3 -- 43 is prime
4 -- 41 is prime
5 -- 41 is prime
6 -- 43 is prime
7 -- 47 is prime
8 -- 53 is prime
9 -- 61 is prime
10 -- 71 is prime
11 -- 83 is prime
12 -- 97 is prime
13 -- 113 is prime
14 -- 131 is prime
15 -- 151 is prime
```

```

16 -- 173 is prime
17 -- 197 is prime
18 -- 223 is prime
19 -- 251 is prime
20 -- 281 is prime
21 -- 313 is prime
22 -- 347 is prime
23 -- 383 is prime
24 -- 421 is prime
25 -- 461 is prime
26 -- 503 is prime
27 -- 547 is prime
28 -- 593 is prime
29 -- 641 is prime
30 -- 691 is prime
31 -- 743 is prime
32 -- 797 is prime
33 -- 853 is prime
34 -- 911 is prime
35 -- 971 is prime
36 -- 1033 is prime
37 -- 1097 is prime
38 -- 1163 is prime
39 -- 1231 is prime
40 -- 1301 is prime
41 -- 1373 is prime
42 -- 1447 is prime
43 -- 1523 is prime
44 -- 1601 is prime
45 -- 1681 is NOT prime
46 -- 1763 is NOT prime
47 -- 1847 is prime
48 -- 1933 is prime
49 -- 2021 is NOT prime
50 -- 2111 is prime
Number of primes: 48
Percentage of primes: 94.1176470588

```

The last two lines of the output contain a small statistics. You can see that $f(n)$ generates 48 primes when $0 \leq n \leq 50$, which is approximately 94% of the values generated by $f(n)$.

For larger sequences, it is impractical to print all single messages indicating the primality of integers. In the following Sage session, only the statistics at the end is printed (by setting the verbose parameter to false): the overall number of primes and the percentage of primes, generated by $f(n)$ where $0 \leq n \leq 1000$.

```

sage: quadratic_prime_formula(0, 1000, False)
N -- N^2 - 9*N + 61
Number of primes: 584
Percentage of primes: 58.3416583417

```

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Web links

1. GIMPS (Great Internet Mersenne-Prime Search)
www.mersenne.org is the home page of the GIMPS project,
<http://www.mersenne.org/prime.htm>
2. The Proth Search Page with the Windows program by Yves Gallot
<http://www.utm.edu/research/primes/programs/gallot/index.html>
3. Generalized Fermat Prime Search
<http://primes.utm.edu/top20/page.php?id=12>
4. Distributed Search for Fermat Number Divisors
<http://www.fermatsearch.org/>
5. At the University of Tennessee you will find extensive research results about prime numbers.
<http://www.utm.edu/>
6. The best overview about prime numbers is offered from my point of view by “The Prime Pages” from professor Chris Caldwell.
<http://www.utm.edu/research/primes>
7. Descriptions e.g. about prime number tests
<http://www.utm.edu/research/primes/mersenne.shtml>
<http://www.utm.edu/research/primes/prove/index.html>
8. Showing the n -th prime number $P(n)$
http://www.utm.edu/research/primes/notes/by_year.html
9. The supercomputer manufacturer SGI Cray Research not only employed brilliant mathematicians but also used the prime number tests as benchmarks for its machines.
http://www.isthe.com/chongo/tech/math/prime/prime_press.html
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<http://www.cerias.purdue.edu/homes/ssw/cun/>
11. <http://www.eff.org/awards/coop>
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Chapter 4

Introduction to Elementary Number Theory with Examples

(Bernhard Esslinger, July 2001; Updates: Dec. 2001, June 2002, May 2003, May 2005, March 2006, June 2007, July 2009, Jan. 2010, Dec. 2012)

This “introduction” is for people with a mathematical interest. There is no more pre-knowledge necessary than what you learn in the secondary school.

We intentionally had “beginners” in mind; we did not take the approach of mathematical textbooks, called “introduction”, which cannot be understood at the first reading further than page 5 and which have the real purpose to deliver all information that special monographs can be read.

4.1 Mathematics and cryptography

A large proportion of modern, asymmetric cryptography is based on mathematical knowledge – on the properties (“laws”) of whole numbers, which are investigated in elementary number theory. Here, the word “elementary” means that questions raised in number theory are essentially rooted in the set of natural and whole numbers.

Further mathematical disciplines currently used in cryptography include (see [Bauer1995, p. 2], [Bauer2000, p. 3]) :

- Group theory
- Combination theory
- Complexity theory
- Stochastic (ergodic theory)
- Information theory.

Number theory or arithmetic (the emphasis here is more on the aspect of performing calculations with numbers) was established by Gauss¹ as a special mathematical discipline. Its

¹Carl Friedrich Gauss, German mathematician and astronomer, Apr 30, 1777–Feb 23, 1855.

elementary features include the greatest common divisor² (gcd), congruence (remainder classes), factorization, the Euler-Fermat theorem and primitive roots. However, the most important aspect is prime numbers and their multiplicative operation.

For a long time, number theory was considered to be the epitome of pure research, the ideal example of research in the ivory tower. It delved into “the mysterious laws of the realm of numbers”, giving rise to philosophical considerations as to whether it described elements that exist everywhere in nature or whether it artificially constructed elements (numbers, operators and properties).

We now know that patterns from number theory can be found everywhere in nature. For example, the ratio of rotating counterclockwise and rotating clockwise spirals in a sunflower is equal to two consecutive Fibonacci numbers³, for example 21 : 34.

Also, at the latest when number theory was applied in modern cryptography, it became clear that a discipline that had been regarded as purely theoretical for centuries actually had a practical use. Today, experts in this field are in great demand on the job market.

Applications in (computer) security now use cryptography because this mathematical discipline is simply better and easier to prove than all other “creative” substitution procedures that have been developed over the course of time and better than all sophisticated physical methods such as those used to print bank notes [Beutelspacher1996, p. 4].

This article explains the basics of elementary number theory in a way that you can easily understand. It provides numerous examples and very rarely goes into any proofs (these can be found in mathematical textbooks).

The goal is not to exhaustively explain the number theory findings, but to show the essential procedures. The volume of the content is so oriented that the reader can understand and apply the RSA method.

For this purpose we will use both theory and examples to explain how to perform calculations in finite sets and describe how these techniques are applied in cryptography. Particular attention will be paid to the traditional Diffie-Hellman (DH) and RSA public key procedures.

It was important to me to make verifiable statements about the security of the RSA algorithm, and to add Sage code for as much as possible examples.

²This article deals with the gcd (greatest common divisor) in appendix 4.14.

³The sequence of Fibonacci numbers $(a_i)_{i \in \mathbb{N}}$ is defined by the “recursive” rule $a_1 := a_2 := 1$ and for all numbers $n = 1, 2, 3, \dots$ we define $a_{n+2} := a_{n+1} + a_n$. This historical sequence can be found in many interesting forms in nature (for example, see [Graham1994, p. 290 ff] or the website of Ron Knott, which is devoted to Fibonacci numbers). A lot is known about the Fibonacci sequence and it is used today as an important tool in mathematics.

Mathematics is the queen of sciences and number theory is the queen of mathematics.

Quote 6: Carl Friedrich Gauss

4.2 Introduction to number theory

Number theory arose from interest in positive whole numbers $1, 2, 3, 4, \dots$, also referred to as the set of natural numbers *natural numbers* \mathbb{N} . These are the first mathematical constructs used by human civilization. According to Kronecker⁴, they are a creation of God. In Dedekind's⁵ opinion, they are a creation of the human intellect. Dependent upon one's ideology, this is an unsolvable contradiction or one and the same thing.

In ancient times, no distinction was made between number theory and numerology, which attributed a mystical significance to specific numbers. In the same way as astronomy and chemistry gradually detached themselves from astrology and alchemy during the Renaissance (from the 14th century), number theory also separated itself from numerology.

Number theory has always been a source of fascination – for both amateurs and professional mathematicians. In contrast to other areas of mathematics, many of the problems and theorems in number theory can be understood by non-experts. On the other hand, the solutions to these problems or the prove to the theorems often resisted to the mathematicians for a very long time. It is therefore one thing to pose good questions but quite another matter to find the answer. One example of this is what is known as Fermat's Last (or large) theorem.⁶

Up until the mid 20th century, number theory was considered to be the purest area of mathematics, an area that had no practical use in the real world. This changed with the development of computers and digital communication, as number theory was able to provide several unexpected solutions to real-life tasks. At the same time, advances in information technology allowed specialists in number theory to make huge progress in factorizing large numbers, finding new prime numbers, testing (old) conjectures and solving numerical problems that were previously impossible to solve. Modern number theory is made up of areas such as:

- Elementary number theory
- Algebraic number theory
- Analytic number theory
- Geometric number theory
- Combinatorial number theory
- Numeric number theory
- Probability theory.

⁴Leopold Kronecker, German mathematician, Dec 7, 1823 – Dec 29, 1891

⁵Julius Wilhelm Richard Dedekind, German mathematician, Oct 6, 1831 – Feb 12, 1916.

⁶One of the things we learn in mathematics at school is Pythagoras' theorem, which states the following for a right-angle triangle: $a^2 + b^2 = c^2$, where a and b are the lengths of the sides containing the right angle and c is the length of the hypotenuse. Fermat famously proposed that $a^n + b^n \neq c^n$ for $a, b, c \in \mathbb{N}$ and whole-number exponents $n > 2$. Unfortunately, the letter in which Fermat made the claim did not have enough space for him to prove it. The theorem was not proven until over 300 years later [Wiles1994, p. 433-551].

All of the different areas are concerned with questions regarding whole numbers (both positive and negative whole numbers plus zero). However, they each have different methods of dealing with them.

This article only deals with the area of elementary number theory.

4.2.1 Convention

Unless stated otherwise:

- The letters $a, b, c, d, e, k, n, m, p, q$ are used to present whole numbers.
- The letters i and j represent natural numbers.
- The letters p always represents a prime number.
- The sets $\mathbb{N} = \{1, 2, 3, \dots\}$ and $\mathbb{Z} = \{\dots, -3, -2, -1, 0, 1, 2, 3, \dots\}$ are the *natural* and *whole* numbers respectively.

This isn't magic – it's logic – a puzzle. A lot of the greatest wizards haven't got an ounce of logic.

Quote 7: Joanne K. Rowling⁷

4.3 Prime numbers and the first fundamental theorem of elementary number theory

Many of the problems in elementary number theory are concerned with prime numbers (see chapter 3).

Every whole number has divisors or factors. The number 1 has just one – itself, whereas the number 12 has the six factors 1, 2, 3, 4, 6 and 12.⁸ Many numbers are only divisible by themselves and by 1. When it comes to multiplication, these can be regarded as the “atoms” in the realm of numbers.

Definition 4.3.1. *Prime numbers are natural numbers greater than 1 that can only be divided by 1 and themselves.*

By definition, 1 is not a prime number.

If we write down the prime numbers in ascending order (prime number sequence), then we get:

2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 41, 43, 47, 53, 59, 61, 67, 71, 73, 79, 83, 89, 97, ⋯ .

The first 100 numbers include precisely 25 prime numbers. After this, the percentage of primes decreases, but never reaches zero.

We come across whole numbers that are prime fairly often. In the last decade only, three years were prime: 1993, 1997 and 1999. If they were rare, cryptography would not be able to work with them to the extent it does.

Prime numbers can be factorized in a unique (“*trivial*”) way:

$$\begin{aligned} 5 &= 1 * 5 \\ 17 &= 1 * 17 \\ 1013 &= 1 * 1013 \\ 1,296,409 &= 1 * 1,296,409. \end{aligned}$$

Definition 4.3.2. *Natural numbers greater than 1 that are not prime are called **composite numbers**. These have at least two factors other than 1.*

⁷ Joanne K. Rowling, “Harry Potter and the Philosopher’s Stone”, Bloomsbury, (c) 1997, chapter “Through the trapdoor”, p. 307, by Hermine.

⁸ Due to the fact that 12 has so many factors, this number – and multiples of this number – is often found in everyday life: the 12-hour scale on clocks, the 60 minutes in an hour, the 360-degree scale for measuring angles, etc. If we divide these scales into segments, the segments often turn out to be whole numbers. These are easier to use in mental arithmetic than fractions.

Examples of the decomposition of such numbers into prime factors:

$$\begin{aligned}
 4 &= 2 * 2 \\
 6 &= 2 * 3 \\
 91 &= 7 * 13 \\
 161 &= 7 * 23 \\
 767 &= 13 * 59 \\
 1029 &= 3 * 7^3 \\
 5324 &= 22 * 11^3.
 \end{aligned}$$

Theorem 4.3.1. *Each composite number a has a lowest factor greater than 1. This factor is a prime number p and is less than or equal to the square root of a .*

All whole numbers greater than 1 can be expressed as a product of prime numbers — in a *unique* way.

This is the claim of the 1st *fundamental theorem of number theory* (= fundamental theorem of arithmetic = fundamental building block of all positive integers). This was formulated precisely for the first time by Carl Friedrich Gauss in his *Disquisitiones Arithmeticae* (1801).

Theorem 4.3.2. Gauss 1801 *Every even natural number greater than 1 can be written as the product of prime numbers. Given two such decompositions $a = p_1 * p_2 * \dots * p_n = q_1 * q_2 * \dots * q_m$, these can be resorted such that $n = m$ and, for all i , $p_i = q_i$.*

In other words: Each natural number other than 1 can be written as a product of prime numbers in precisely one way, if we ignore the order of the factors. The factors are therefore unique (the “expression as a product of factors” is unique)!

For example, $60 = 2 * 2 * 3 * 5 = 2^2 * 3 * 5$. And this — other than changing the order of the factors — is the only way in which the number 60 can be factorized.

If you allow numbers other than primes as factors, there are several ways of factorizing integers and the *uniqueness* is lost:

$$60 = 1 * 60 = 2 * 30 = 4 * 15 = 5 * 12 = 6 * 10 = 2 * 3 * 10 = 2 * 5 * 6 = 3 * 4 * 5 = \dots$$

The 1st fundamental theorem only appears to be obvious. We can construct numerous other sets of numbers⁹ for which numbers in the set *cannot* be expressed uniquely as a product of the prime numbers of the set.

In order to make a mathematical statement, therefore, it is important to state not only the operation for which it is defined but also the basic set on which the operation is defined.

For more details on prime numbers (e.g. how “Fermat’s Little Theorem” can be used to test extremely large numbers to determine whether they are prime), please refer to the article on prime numbers, chapter 3 in this script.

⁹These sets are formed especially from the set of natural numbers. An example of this can be found in this script on page 56 at the end of chapter 3.2.

4.4 Divisibility, modulus and remainder classes¹⁰

If whole numbers are added, subtracted or multiplied, the result is always another whole number.

The division of two whole numbers does not always result in a whole number. For example, if we divide 158 by 10 the result is the decimal number 15.8, which is not a whole number!

If, however, we divide 158 by 2 the result 79 is a whole number. In number theory we express this by saying that 158 is *divisible* by 2 but not by 10. In general, we say:

Definition 4.4.1. A whole number n is **divisible** by a whole number d if the quotient n/d is a whole number c such that $n = c * d$.

n is called a *multiple* of d , whereas d is called a *divisor* or *factor* of n .

The mathematical notation for this is $d|n$ (read “ d divides n ”). The notation $d \nmid n$ means that d does not divide the number n .

In our example therefore: $10|158$ but $2 \nmid 158$.

4.4.1 The modulo operation – working with congruence

When we investigate divisibility, it is only the remainder of the division that is important. When dividing a number n by m , we often use the following notation:

$$\frac{n}{m} = c + \frac{r}{m},$$

where c is a whole number and r is a number with the values $0, 1, \dots, m - 1$. This notation is called division with remainder, whereby c is called the whole-number “quotient” and r is the “remainder” of the division.

Example:

$$\frac{19}{7} = 2 + \frac{5}{7} \quad (m = 7, \quad c = 2, \quad r = 5)$$

What do the numbers 5, 12, 19, 26, … have in common for division by 7? The remainder is always $r = 5$. For division by 7, only the following remainders are possible:

$$r = 0, 1, 2, \dots, 6.$$

The numbers that result in the same remainder r when divided by 7 are combined to form the “remainder class r modulo 7”. Two numbers a and b belonging to the same remainder class modulo 7 are said to be “congruent modulo 7”. Or in general:

Definition 4.4.2. The **remainder class r modulo m** is the set of all whole numbers a that have the same remainder r when divided by m .

¹⁰With the educational tool for number theory NT you can have a playful view at the calculation with congruences, discussed in this and the next chapter (see NT learning unit 2.1, pages 2-9/40).

NT can be called in CrypTool via the menu path **Indiv. Procedures \ Number Theory Interactive \ Learning Tool for Number Theory**. See appendix A.6.

CT2 contains a visualisation of these methods within the tutorial “**World of Primes**”.

Example:

Remainder class 0 modulo 4 =

$$\{x|x = 4 * n; n \in \mathbb{Z}\} = \{\dots, -16, -12, -8, -4, 0, 4, 8, 12, 16, \dots\}$$

Remainder class 3 modulo 4 =

$$\{x|x = 4 * n + 3; n \in \mathbb{Z}\} = \{\dots, -13, -9, -5, -1, 3, 7, 11, 15, \dots\}$$

As only the remainders $0, 1, 2, \dots, m-1$ are possible for division modulo m , modular arithmetic works with finite sets. For each modulo m there are precisely m remainder classes.

Definition 4.4.3. Two numbers $a, b \in \mathbb{N}$ are said to be congruent modulo $m \in \mathbb{N}$ if and only if they have the same remainder when divided by m .

We write: $a \equiv b \pmod{m}$ (read a is congruent b modulo m), which means that a and b belong to the same remainder class. The modulo is therefore the divisor. This notation was introduced by Gauss. Although the divisor is usually positive, a and b can also be any whole numbers.

Example:

$19 \equiv 12 \pmod{7}$, because the remainders are equal: $19/7 = 2$ remainder 5 and $12/7 = 1$ remainder 5.

$23103 \equiv 0 \pmod{453}$, because $23103/453 = 51$ remainder 0 and $0/453 = 0$ remainder 0.

Theorem 4.4.1. $a \equiv b \pmod{m}$ if and only if, the difference $(a - b)$ is divisible by m , i.e. if $q \in \mathbb{Z}$ exists with $(a - b) = q * m$.

These two statements are therefore equivalent.¹¹

Therefore: If m divides the difference, there exists a whole number q such that: $a = b + q * m$. As an alternative to the congruence notation, we can also use the divisibility notation: $m|(a-b)$.

Example of equivalent statements:

$35 \equiv 11 \pmod{3} \iff 35 - 11 \equiv 0 \pmod{3}$, where $35 - 11 = 24$ is divisible by 3 without remainder while $35 : 3$ and $11 : 3$ leave the remainder 2.

We can apply the above equivalence in theorem 4.4.1 if we need a quick and easy method of determining whether large numbers are divisible by a certain number.

Example: Is 69,993 divisible by 7?

The number can be written in the form of a difference in which it is clear that each operand is divisible by 7: $69,993 = 70,000 - 7$. Therefore, the difference is also divisible by 7.

Although these considerations and definitions may seem to be rather theoretical, we are so familiar with them in everyday life that we no longer think about the formal procedure. For example, the 24 hours on a clock are represented by the numbers $1, 2, \dots, 12$. We obtain the hours after 12 noon as the remainder of a division by 12 and know immediately that 2 o'clock in the afternoon is the same as 14.00.

This “modular” arithmetic (based on division remainders) forms the basis of asymmetric encryption procedures. Cryptographic calculations are therefore not based on real numbers, as the calculations you performed at school, but rather on character strings with a limited length,

¹¹The above equivalence does apply only to the difference $(a - b)$, not to the sum $(a + b)$!

Example:

$11 \equiv 2 \pmod{3}$, therefore $11 - 2 = 9 \equiv 0 \pmod{3}$; but $11 + 2 = 13$ is not divisible by 3.

The statement in theorem 4.4.1 does not even apply to sums in one direction. It is correct for sums only if the remainder is 0 and only in the following direction: If a divisor divides both summands with no remainder, it also divides the sum with no remainder.

in other words on positive whole numbers that cannot exceed a certain value. This is one of the reasons why we choose a large number m and “calculate modulo m ”. That is, we ignore whole-number multiples of m and, rather than working with a number, we only work with the remainder when this number is divided by m . The result is that all results are in the range 0 to $m - 1$.

4.5 Calculations with finite sets

4.5.1 Laws of modular calculations

From algebra theorems it follows that essential parts of the conventional calculation rules are kept when we proceed to modular calculations over a basic set \mathbb{Z} . For example, addition remains commutative. The same goes for multiplication modulo m . The result of a division¹² is not a fraction but rather a whole number between 0 and $m - 1$.

The known laws apply:

1. Associative law:

$$\begin{aligned} ((a + b) + c) \pmod{m} &\equiv (a + (b + c)) \pmod{m}. \\ ((a * b) * c) \pmod{m} &\equiv (a * (b * c)) \pmod{m}. \end{aligned}$$

2. Commutative law:

$$\begin{aligned} (a + b) \pmod{m} &\equiv (b + a) \pmod{m}. \\ (a * b) \pmod{m} &\equiv (b * a) \pmod{m}. \end{aligned}$$

The associative law and the commutative law apply to both addition and multiplication.

3. Distributive law:

$$(a * (b + c)) \pmod{m} \equiv (a * b + a * c) \pmod{m}.$$

4. Reducibility:

$$\begin{aligned} (a + b) \pmod{m} &\equiv (a \pmod{m} + b \pmod{m}) \pmod{m}. \\ (a * b) \pmod{m} &\equiv (a \pmod{m} * b \pmod{m}) \pmod{m}. \end{aligned}$$

When adding or multiplying the order in which the modulo operation is performed does not matter.

5. Existence of an identity (neutral element):

$$\begin{aligned} (a + 0) \pmod{m} &\equiv (0 + a) \pmod{m} \equiv a \pmod{m}. \\ (a * 1) \pmod{m} &\equiv (1 * a) \pmod{m} \equiv a \pmod{m}. \end{aligned}$$

6. Existence of an inverse element:

For all whole numbers a and m there exists a whole number $-a$ such that:

$$(a + (-a)) \pmod{m} \equiv 0 \pmod{m} \quad (\text{additive inverse}).$$

For each a ($a \not\equiv 0 \pmod{p}$) where p is prime there exists a whole number a^{-1} , such that:

$$(a * a^{-1}) \pmod{p} \equiv 1 \pmod{p} \quad (\text{multiplicative inverse}).$$

7. Closeness¹³:

$$\begin{aligned} a, b \in G \implies (a + b) &\in G. \\ a, b \in G \implies (a * b) &\in G. \end{aligned}$$

¹²The division modulo m is only defined for numbers co-prime to m because other numbers have the same property as zero (this means there is no inverse number). See footnote 16 in chapter 4.6.1 and table 4.3 in chapter 4.6.2.

¹³The property of closeness is always defined in relation to an operation in a set. See chapter 4.15 “Appendix: Forming closed sets”.

8. Transitivity:

$$[a \equiv b \pmod{m}, b \equiv c \pmod{m}] \implies [a \equiv c \pmod{m}].$$

4.5.2 Patterns and structures

In general mathematicians investigate “Structures”. They ask e.g. at $a * x \equiv b \pmod{m}$, which values x can take for given values of a, b, m .

Especially the case is investigated, where the result b of this operation is the neutral element. Then x is the inverse of a regarding this operation.

The way of theory is long — it is short and effective by examples.

Quote 8: Seneca¹⁴

4.6 Examples of modular calculations

As we have already seen:

For two natural numbers a and m , $a \bmod m$ denotes the remainder obtained when we divide a by m . This means that $a \pmod{m}$ is always a number between 0 and $m - 1$.

For example, $1 \equiv 6 \equiv 41 \equiv 1 \pmod{5}$ because the remainder is always 1. Another example is: $2000 \equiv 0 \pmod{4}$ because 4 divides 2000 with no remainder.

Modular arithmetic only contains a limited quantity of non-negative numbers. The number of these is specified by a modulus m . If the modulo is $m = 5$, then only the 5 numbers in the set $\{0, 1, 2, 3, 4\}$ are used.

A calculation result larger than 4 is then reduced “modulo 5”. In other words, it is the remainder when the result is divided by 5. For example, $2 * 4 \equiv 8 \equiv 3 \pmod{5}$ because 3 is the remainder when we divide 8 by 5.

4.6.1 Addition and multiplication

The following shows

- the addition table¹⁵ $\pmod{5}$ and
- the multiplication tables¹⁶ (table 4.1) for mod 5 (table 4.2) and mod 6 (table 4.3).

Example of an addition table

The result when we add 3 and 4 ($\pmod{5}$) is calculated as follows: Calculate $3 + 4 = 7$ and keep subtracting 5 from the result until the result is less than the modulo: $7 - 5 = 2$. Therefore: $3 + 4 \equiv 2 \pmod{5}$.

Example of a multiplication table:

The result of the multiplication $4 * 4 \pmod{5}$ is calculated as follows: $4 * 4 = 16$ and subtract 5 until the result is less than the modulus.

$$16 - 5 = 11; 11 - 5 = 6; 6 - 5 = 1.$$

¹⁴Lucius Annaeus Seneca, philosophical writer and poet, 4 B. C. – 65 A. D.

¹⁵Comment on subtraction modulo 5:

$2 - 4 = -2 \equiv 3 \pmod{5}$.

It is therefore not true modulo 5 that $-2 = 2$ (see chapter 4.16 “Appendix: Comments on modulo subtraction”).

¹⁶Comment on division modulo 6:

Due to the special role of zero as the identity for addition, division by zero is not permitted:

For all a it is $a * 0 = 0$, because $a * 0 = a * (0 + 0) = a * 0 + a * 0$. Obviously 0 has no inverse regarding the multiplication, because if there would be one, it must be $0 = 0 * 0^{-1} = 1$. Also see footnote 12.

$+$	0	1	2	3	4
0	0	1	2	3	4
1	1	2	3	4	0
2	2	3	4	0	1
3	3	4	0	1	2
4	4	0	1	2	3

Table 4.1: Addition table modulo 5

The table directly shows that $4 * 4 \equiv 1 \pmod{5}$ because $16 : 5 = 3$ remainder 1. Multiplication is defined on the set \mathbb{Z} excluding 0.

$*$	1	2	3	4
1	1	2	3	4
2	2	4	1	3
3	3	1	4	2
4	4	3	2	1

Table 4.2: Multiplication table modulo 5

4.6.2 Additive and multiplicative inverses

You can use the tables to read the inverses for each number in relation to addition and multiplication.

The inverse of a number is the number that gives the result 0 when the two numbers are added and 1 when they are multiplied. Thus, the inverse of 4 for addition mod 5 is 1 and the inverse of 4 for multiplication mod 5 is 4 itself, because

$$4 + 1 = 5 \equiv 0 \pmod{5}; \\ 4 * 4 = 16 \equiv 1 \pmod{5}.$$

The inverse of 1 for multiplication mod 5 is 1, while the inverse modulo 5 of 2 is 3 and, since multiplication is commutative, the inverse of 3 is again 2.

If we take a random number and add or multiply another number (here 4) and then add¹⁷ or multiply the corresponding inverse (1 or 4) to the interim result (1 or 3), then the end result is the same as the initial value.

Example:

$$2 + 4 \equiv 6 \equiv 1 \pmod{5}; \quad 1 + 1 \equiv 2 \equiv 2 \pmod{5}, \\ 2 * 4 \equiv 8 \equiv 3 \pmod{5}; \quad 3 * 4 \equiv 12 \equiv 2 \pmod{5}.$$

In the set $\mathbb{Z}_5 = \{0, 1, 2, 3, 4\}$ for the addition and in the set \mathbb{Z}_5^* for the multiplication, all numbers have a **unique** inverse modulo 5.

¹⁷In general $x + y + (-y) \equiv x \pmod{m}$ [$(-y)$ = additive inverse of $y \pmod{m}$].

In the case of modular addition, this is true for every modulo (not just for 5).

This is not the case, however, for modular multiplication.

Theorem 4.6.1. *A natural number a from the set $\{1, \dots, m-1\}$ has one inverse if and only if this number and the modulo m are co-prime¹⁸, in other words if a and m have no common prime factors.*

Since $m = 5$ is prime, the numbers 1 to 4 are relatively prime to 5 and **each** of these numbers has a multiplicative inverse in mod 5.

Table 4.3 shows as a counterexample the multiplication table for mod 6 (since the modulus $m = 6$ is not prime, not all elements from $\mathbb{Z}_6 \setminus \{0\}$ are relatively prime to 6).

*	1	2	3	4	5
1	1	2	3	4	5
2	2	4	0	2	4
3	3	0	3	0	3
4	4	2	0	4	2
5	5	4	3	2	1

Table 4.3: Multiplication table modulo 6

In addition to 0, the numbers 2, 3 and 4 also have no unique inverse (we can also say they have **no** inverse, because the elementary property of an inverse is uniqueness).

The numbers 2, 3 and 4 have the factor 2 or 3 in common with the modulus 6. Only the numbers 1 and 5, which are relatively prime to 6, have multiplicative inverses, namely themselves.

The number of numbers that are relatively prime to the modulus m is the same as the number of numbers that have a multiplicative inverse (see the Euler function $J(m)$ below).

For the two moduli 5 and 6 used in the multiplication tables, this means: the modulus 5 is a prime number itself. In mod 5, therefore, there are exactly $J(5) = 5 - 1 = 4$ numbers that are relatively prime to the modulus, that is all numbers from 1 to 4.

Since 6 is not a prime number, we write it as a product of its factors: $6 = 2 * 3$. In mod 6, therefore, there are exactly $J(6) = (2-1)*(3-1) = 1*2 = 2$ numbers that have a multiplicative inverse, that is 1 and 5.

Although it may seem difficult to calculate the table of multiplicative inverses for large moduli (this only applies to the areas of the table shaded dark grey), we can use Fermat's Little Theorem to create a simple algorithm for this [Pfleeger1997, p. 80]. Quicker algorithms are described, for instance, in [Knuth1998].¹⁹

¹⁸Two whole numbers a and b are co-prime if and only if $\gcd(a, b) = 1$.

If p is prime and a is a random whole number that is not a multiple of p , then p and a are co-prime. Further name to the topic co-prime (with $a_i \in \mathbb{Z}, i = 1, \dots, n$):

1. a_1, a_2, \dots, a_n are *relatively prime*, if $\gcd(a_1, \dots, a_n) = 1$.
2. An even stronger request for more than two numbers is :
 a_1, \dots, a_n are *in pairs relatively prime*, if for all $i = 1, \dots, n$ and $j = 1, \dots, n$ with $i \neq j$: $\gcd(a_i, a_j) = 1$.

Example:

2, 3, 6 are relatively prime, because $\gcd(2, 3, 6) = 1$. They are not in pairs relatively prime, because $\gcd(2, 6) = 2 > 1$.

¹⁹Using Euclid's extended theorem (extended gcd), we can calculate the multiplicative inverse and determine whether numbers have an inverse (see appendix 4.14). Alternatively, we can also use the primitive roots.

Cryptographically not only the unique nature of the inverse is important, but also that the set of possible values has been exhausted.

Theorem 4.6.2. For $a, i \in \{1, \dots, m-1\}$ with $\gcd(a, m) = 1$, then the product $a * i \bmod m$ takes for a certain number a all values from $\{1, \dots, m-1\}$ (exhaustive permutation of the length $m-1$).²⁰

The following three examples²¹ illustrate the properties of multiplicative inverses.

Multiplication table mod 17 (table 4.4) shows for $i = 1, 2, \dots, 18$ the following:

$$(5 * i) / 17 = a \text{ remainder } r \text{ and high-lighted } 5 * i \equiv 1 \pmod{17},$$

$$(6 * i) / 17 = a \text{ remainder } r \text{ and high-lighted } 6 * i \equiv 1 \pmod{17}.$$

We need to **find** the i for which the product remainder $a * i \bmod 17$ with $a = 5$ or $a = 6$ has the value 1.

i	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18
$5 * i$	5	10	15	20	25	30	35	40	45	50	55	60	65	70	75	80	85	90
remainder	5	10	15	3	8	13	1	6	11	16	4	9	14	2	7	12	0	5
$6 * i$	6	12	18	24	30	36	42	48	54	60	66	72	78	84	90	96	102	108
remainder	6	12	1	7	13	2	8	14	3	9	15	4	10	16	5	11	0	6

Table 4.4: Multiplication table modulo 17 (for $a = 5$ and $a = 6$)

Between $i = 1, \dots, m$, all values between $0, \dots, m-1$ occur for the remainders, because both 5 and 6 are also relatively prime to the modulus $m = 17$.

The multiplicative inverse of 5 (mod 17) is 7, while the inverse of 6 (mod 17) is 3.

i	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18
$5 * i$	5	10	15	20	25	30	35	40	45	50	55	60	65	70	75	80	85	90
remainder	5	10	2	7	12	4	9	1	6	11	3	8	0	5	10	2	7	12
$6 * i$	6	12	18	24	30	36	42	48	54	60	66	72	78	84	90	96	102	108
remainder	6	12	5	11	4	10	3	9	2	8	1	7	0	6	12	5	11	4

Table 4.5: Multiplication table modulo 13 (for $a = 5$ and $a = 6$)

Between $i = 1, \dots, m$, all values between $0, \dots, m-1$ occur for the remainders, because both 5 and 6 are relatively prime to the modulus $m = 13$.

The multiplicative inverse of 5 (mod 13) is 8, while the inverse of 6 (mod 13) is 11.

The following table 4.6 contains an example, where the modulus m and the number $a = 6$ are *not* relatively prime.

We have calculated $(5 * i) \pmod{12}$ and $(6 * i) \pmod{12}$. Between $i = 1, \dots, m$, not all values between $0, \dots, m-1$ occur and 6 does not have an inverse mod 12, because 6 and the modulus $m = 12$ are not co-prime.

The multiplicative inverse of 5 (mod 12) is 5. The number 6 has no inverse (mod 12).

²⁰See also theorem 4.9.1 in chapter 4.9, Multiplicative order and primitive roots.

²¹See chapter 4.19.1 “Multiplication table modulo m” for the source code to compute the tables using Sage.

i	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18
$5 * i$	5	10	15	20	25	30	35	40	45	50	55	60	65	70	75	80	85	90
remainder	5	10	3	8	1	6	11	4	9	2	7	0	5	10	3	8	1	6
$6 * i$	6	12	18	24	30	36	42	48	54	60	66	72	78	84	90	96	102	108
remainder	6	0	6	0	6	0	6	0	6	0	6	0	6	0	6	0	6	0

Table 4.6: Multiplication table modulo 12 (for $a = 5$ and $a = 6$)

4.6.3 Raising to the power

In modular arithmetic, raising to the power is defined as repeated multiplication – as usual except that multiplication is now slightly different. We can even apply the usual rules, such as:

$$\begin{aligned} a^{b+c} &= a^b * a^c, \\ (a^b)^c &= a^{b*c} = a^{c*b} = (a^c)^b. \end{aligned}$$

Modular powers work in the same way as modular addition and modular multiplication:

$$3^2 = 9 \equiv 4 \pmod{5}.$$

Even consecutive powers work in the same way:

Example 1:

$$(4^3)^2 = 64^2 \equiv 4096 \equiv 1 \pmod{5}.$$

- (1) We can speed up²² the calculation by reducing the **interim results** modulo 5 but we need to take care because *not* everything will then work in the same way as in standard arithmetic.

$$\begin{aligned} (4^3)^2 &\equiv (4^3 \pmod{5})^2 \pmod{5} \\ &\equiv (64 \pmod{5})^2 \pmod{5} \\ &\equiv 4^2 \pmod{5} \\ &\equiv 16 \equiv 1 \pmod{5}. \end{aligned}$$

- (2) In standard arithmetic, consecutive powers can be reduced to a single power by multiplying the exponents:

$$(4^3)^2 = 4^{3*2} = 4^6 = 4096.$$

This is not quite as simple in modular arithmetic because this would give:

$$(4^3)^2 \equiv 4^{3*2} \pmod{5} \equiv 4^6 \pmod{5} \equiv 4^1 \equiv 4 \pmod{5}.$$

But as we saw above, the correct result is 1 !

- (3) Therefore, the rule is slightly different for consecutive powers in modular arithmetic: We do not multiply the exponents in $(\pmod m)$ but rather in $(\pmod J(m))$.

²²The time required to calculate the multiplication of two numbers normally depends on the length of the numbers. We can observe this if we use the school method to calculate, for instance, $474 * 228$. The time required increases in a quadratic square manner , because we need to multiply $3 * 3$ numbers. The numbers become considerably smaller if we reduce the interim result.

Using $J(5) = 4$ gives:

$$(4^3)^2 \equiv 4^{3*2} \pmod{J(5)} \equiv 4^6 \pmod{4} \equiv 4^2 \equiv 16 \equiv 1 \pmod{5}.$$

This delivers the correct result.

Theorem 4.6.3. $(a^b)^c \equiv a^{b*c} \pmod{J(m)}$ (\pmod{m}).

Example 2:

$$3^{28} = 3^{4*7} \equiv 3^{4*7} \pmod{10} \equiv 3^8 \equiv 6561 \equiv 5 \pmod{11}.$$

4.6.4 Fast calculation of high powers

RSA encryption and decryption²³ entails calculating high powers modulo m . For example, the calculation $(100^5) \pmod{3}$ exceeds the 32-bit long integer number range provided we calculate a^n by actually multiplying a with itself n times in line with the definition. In the case of extremely large numbers, even a fast computer chip would take longer than the age of the universe to calculate a single exponential. Luckily, there is an extremely effective shortcut for calculating exponentials (but not for calculating logarithms).

If the expression is divided differently using the rules of modular arithmetic, then the calculation does not even exceed the 16-bit short integer number range:

$$(a^5) \equiv (((a^2) \pmod{m})^2 \pmod{m}) * a \pmod{m}.$$

We can generalize this by representing the exponent as a binary number. For example, the naive method would require 36 multiplications in order to calculate a^n for $n = 37$. However, if we write n in the binary representation as $100101 = 1 * 2^5 + 1 * 2^2 + 1 * 2^0$, then we can rewrite the expression as: $a^{37} = a^{2^5+2^2+2^0} = a^{2^5} * a^{2^2} * a^{2^0}$

Example 3: $87^{43} \pmod{103}$.

Since $43 = 32 + 8 + 2 + 1$, 103 is prime, $43 < J(103)$

and the squares ($\pmod{103}$) can be calculated beforehand

$$\begin{aligned} 87^2 &\equiv 50 \pmod{103}, \\ 87^4 &\equiv 50^2 \equiv 28 \pmod{103}, \\ 87^8 &\equiv 28^2 \equiv 63 \pmod{103}, \\ 87^{16} &\equiv 63^2 \equiv 55 \pmod{103}, \\ 87^{32} &\equiv 55^2 \equiv 38 \pmod{103}. \end{aligned}$$

We have²⁴:

$$\begin{aligned} 87^{43} &\equiv 87^{32+8+2+1} \pmod{103} \\ &\equiv 87^{32} * 87^8 * 87^2 * 87 \pmod{103} \\ &\equiv 38 * 63 * 50 * 87 \equiv 85 \pmod{103}. \end{aligned}$$

²³See chapter 4.10 (“Proof of the RSA procedure with Euler-Fermat”) and chapter 4.13 (“The RSA procedure with actual numbers”).

²⁴See chapter 4.19.2 “Fast exponentiation” for source code implementing the square and multiply method in Sage, which can be used to reproduce the calculations above.

The powers $(a^2)^k$ can be determined easily by means of repeated squaring. As long as a does not change, a computer can calculate them beforehand and – if enough memory is available – save them. In order to then find a^n in each individual case, it now only needs to multiply those $(a^2)^k$ for which there is a one in the k -th position of the binary representation of n . The typical effort is then reduced from 2^{600} to $2 * 600$ multiplications! This frequently used algorithm is called “Square and Multiply”.

4.6.5 Roots and logarithms

The inverses of the powers are also defined. The roots and logarithms are again whole numbers. Yet in contrast to the usual situation, they are not only difficult to calculate but, in the case of large numbers, cannot be calculated at all within a reasonable amount of time.

Let us take the equation $a \equiv b^c \pmod{m}$.

a) Taking the logarithm (determining c) — Discrete logarithm problem²⁵:

If we know a and b of the three numbers a , b and c that meet this equation, then every known method of finding c is approximately just as time-consuming as trying out all m possible values for c one after the other. For a typical m of the order of magnitude of 10^{180} for 600-digit binary numbers, this is a hopeless task. More precisely, for suitably large numbers m , the time required according to current knowledge is proportional to $\exp(C * (\log m[\log \log m]^2)^{1/3})$ with a constant $C > 1$.

b) Calculating the root (determining b):

The situation is similar if b is the unknown variable and we know the values of a and c : If we know the Euler function²⁶ $J(m)$, then we can easily²⁷ calculate d with $c * d \equiv 1 \pmod{J(m)}$ and use theorem 4.6.3 to obtain:

$$a^d \equiv (b^c)^d \equiv b^{c*d} \equiv b^{c*d} \pmod{J(m)} \equiv b^1 \equiv b \pmod{m}$$

the c -th root b of a .

If $J(m)$ cannot be determined²⁸, it is difficult to calculate the c -th root. This forms the basis for the security assumption used by the RSA encryption system (see chapter 4.10 or chapter 5.3.1).

The time required for inverting addition and multiplication, on the other hand, is simply proportional to $\log m$ or $(\log m)^2$. Powers (for a number x calculate x^a with a fixed) and exponents (for a number x calculate a^x with a fixed) are therefore typical one way functions (compare chapters 5.1 and 4.12.1).

4.7 Groups and modular arithmetic in \mathbb{Z}_n and \mathbb{Z}_n^*

Mathematical “groups” play a decisive role in number theory and cryptography. We only talk of groups if, for a defined set and a defined relation (an operation such as addition or

²⁵Further details about the discrete logarithm problem can be found in chapter 5.4.

²⁶See chapter 4.8.2, “The Euler function”.

²⁷See chapter 4.14, “Appendix: gcd and the two algorithms of Euclid”.

²⁸According to the first fundamental theorem of number theory and theorem 4.8.4, we can determine $J(m)$ by reducing m to prime factors.

multiplication), the following properties are fulfilled:

- The set is closed
- A neutral element exists
- An inverse element exists for each element
- The associative law applies.

The abbreviated mathematical notation is $(G, +)$ or $(G, *)$.

Definition 4.7.1. \mathbb{Z}_n :

\mathbb{Z}_n comprises all numbers from 0 to $n - 1$: $\mathbb{Z}_n = \{0, 1, 2, \dots, n - 2, n - 1\}$.

\mathbb{Z}_n is an often used finite group of the natural numbers. It is sometimes also called the *remainder set R modulo n*.

For example, 32-bit computers (standard PCs) only directly work with whole numbers in a finite set, that is the value range $0, 1, 2, \dots, 2^{32} - 1$.

This value range is equivalent to the set $\mathbb{Z}_{2^{32}}$.

4.7.1 Addition in a group

If we define the operation mod+ on such a set where

$$a \text{ mod } + b := (a + b) \pmod{n},$$

then the set \mathbb{Z}_n together with the relation mod+ is a group because the following properties of a group are valid for all elements in \mathbb{Z}_n :

- $a \text{ mod } + b$ is an element of \mathbb{Z}_n (the set is closed),
- $(a \text{ mod } + b) \text{ mod } + c \equiv a \text{ mod } + (b \text{ mod } + c) \quad (\text{mod+ is associative}),$
- the neutral element is 0.
- each element $a \in \mathbb{Z}_n$ has an inverse for this operation, namely $n - a$ (because $a \text{ mod } + (n - a) \equiv a + (n - a) \pmod{n} \equiv n \equiv 0 \pmod{n}$).

Since the operation is commutative, i.e. $(a \text{ mod } + b) = (b \text{ mod } + a)$, this structure is actually a “commutative group”.

4.7.2 Multiplication in a group

If we define the operation mod* on the set \mathbb{Z}_n where

$$a \text{ mod } * b := (a * b) \pmod{n},$$

then \mathbb{Z}_n together with this operation is **usually not a group** because not all properties are fulfilled for each n .

Example:

- a) In \mathbb{Z}_{15} , for example, the element 5 does not have an inverse. That is to say, there is no a with $5 * a \equiv 1 \pmod{15}$. Each modulo product with 5 on this set gives 5, 10 or 0.
- b) In $\mathbb{Z}_{55} \setminus \{0\}$, for example, the elements 5 and 11 do not have multiplicative inverses. That is to say, there is no $a \in \mathbb{Z}_{55}$ such that $5 * a \equiv 1 \pmod{55}$ and no a such that $11 * a \equiv 1 \pmod{55}$. This is because 5 and 11 are not relatively prime to 55. Each modulo product with 5 on this set gives 5, 10, 15, ..., 50 or 0. Each modulo product with 11 on this set gives 11, 22, 33, 44 or 0.

On the other hand, there are subsets of \mathbb{Z}_n that form a group with the operation mod*. If we choose all elements in \mathbb{Z}_n that are relatively prime to n , then this set forms a group with the operation mod*. We call this set \mathbb{Z}_n^* .

Definition 4.7.2. \mathbb{Z}_n^* :

$$\mathbb{Z}_n^* = \{a \in \mathbb{Z}_n \mid \gcd(a, n) = 1\}.$$

\mathbb{Z}_n^* is sometimes also called the reduced remainder set R' modulo n .

Example: For $n = 10 = 2 * 5$ the following applies:

full remainder set $R = \mathbb{Z}_n = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$

reduced remainder set $R' = \mathbb{Z}_n^* = \{1, 3, 7, 9\} \longrightarrow J(n) = 4$.

Comment:

R' or \mathbb{Z}_n^* is always a genuine subset of R or \mathbb{Z}_n because 0 is always an element of R but never an element of R' . Since 1 and $n - 1$ are always relatively prime to n , they are always elements of both sets.

If we select a random element in \mathbb{Z}_n^* and multiply it by every other element in \mathbb{Z}_n^* , then the products²⁹ are all in \mathbb{Z}_n^* , and the results are also a unique permutation of the elements in \mathbb{Z}_n^* . Since 1 is always an element of \mathbb{Z}_n^* , there is a unique “partner” in this set such that the product is 1. In other words:

Theorem 4.7.1. *Each element in \mathbb{Z}_n^* has a multiplicative inverse.*

Example: $a = 3$ modulo 10 with $\mathbb{Z}_n^* = \{1, 3, 7, 9\}$ it holds that $a^{-1} = 7$:

$$\begin{aligned} 3 &\equiv 3 * 1 \pmod{10}, \\ 9 &\equiv 3 * 3 \pmod{10}, \\ 1 &\equiv 3 * 7 \pmod{10}, \\ 7 &\equiv 3 * 9 \pmod{10}. \end{aligned}$$

The unique invertibility is an essential condition for cryptography (see section 4.10).

²⁹This is due to the fact that \mathbb{Z}_n^* is closed with respect to the multiplication and due to the gcd property:
 $[a, b \in \mathbb{Z}_n^*] \Rightarrow [((a * b) \pmod{n}) \in \mathbb{Z}_n^*]$, exactly:
 $[a, b \in \mathbb{Z}_n^*] \Rightarrow [\gcd(a, n) = 1, \gcd(b, n) = 1] \Rightarrow [\gcd(a * b, n) = 1] \Rightarrow [((a * b) \pmod{n}) \in \mathbb{Z}_n^*]$.

Mathematical game theory postulates players who respond rationally. Transactional game theory, on the other hand, deals with games that are not rational, perhaps even **irrational and thereby closer to reality**.

Quote 9: Eric Berne³⁰

4.8 Euler function, Fermat's little theorem and Euler-Fermat

4.8.1 Patterns and structures

As mathematicians investigate the structure $a * x \equiv b \pmod{m}$ (see chapter 4.5.2), so they are interested in the structure $x^a \equiv b \pmod{m}$.

Again here they are interested in the cases, if $b = 1$ (value of the multiplicative inverse) and if $b = x$ (the function $f(x) = x^a \pmod{m}$ has a fixpoint). Concerning RSA fixed points: see 4.19.7.

4.8.2 The Euler function

Given n , the number of numbers from the set $\{1, \dots, n-1\}$ that are relatively prime to n is equal to the value of the Euler³¹ function $J(n)$.

Definition 4.8.1. *The Euler function³² $J(n)$ specifies the number of elements in \mathbb{Z}_n^* .*

$J(n)$ also specifies how many whole numbers have multiplicative inverses in mod n . $J(n)$ can be calculated if we know the prime factors of n .

Theorem 4.8.1. *For a prime number, the following is true: $J(p) = p - 1$.*

Theorem 4.8.2. *If m is the product of two distinct primes, then:*

$$J(p * q) = (p - 1) * (q - 1) \quad \text{or} \quad J(p * q) = J(p) * J(q).$$

This case is important for the RSA procedure.

Theorem 4.8.3. *If $n = p_1 * p_2 * \dots * p_k$ where p_1 to p_k are distinct prime numbers (i.e. no factor occurs more than once), then the following is true (as a generalization of theorem 4.8.2):*

$$J(n) = (p_1 - 1) * (p_2 - 1) * \dots * (p_k - 1).$$

Theorem 4.8.4. *In general, the following is true for every prime number p and every n in \mathbb{N} :*

1. $J(p^n) = p^{n-1} * (p - 1)$.

2. *If $n = p_1^{e_1} * p_2^{e_2} * \dots * p_k^{e_k}$, where p_1 to p_k are distinct prime numbers, then:*

$$J(n) = [(p_1^{e_1-1}) * (p_1 - 1)] * \dots * [(p_k^{e_k-1}) * (p_k - 1)] = n * ((p_1 - 1)/p_1) * \dots * ((p_k - 1)/p_k).$$

³⁰Eric Berne, "Games People Play", rororo, (c) 1964, page 235.

³¹Leonhard Euler, Swiss mathematician, Apr 15, 1707 – Sep 18, 1783

³²Often written as the Euler phi function $\Phi(n)$ or $\phi(n)$.

Example:

- $n = 70 = 2 * 5 * 7 \implies$ using theorem 4.8.3: $J(n) = 1 \cdot 4 \cdot 6 = 24$.
- $n = 9 = 3^2 \implies$ using theorem 4.8.4: $J(n) = 3^1 \cdot 2 = 6$, because $\mathbb{Z}_9^* = \{1, 2, 4, 5, 7, 8\}$.
- $n = 2,701,125 = 3^2 * 5^3 * 7^4 \implies$ using theorem 4.8.4:

$$J(n) = [3^1 * 2] * [5^2 * 4] * [7^3 * 6] = 1,234,800.$$

4.8.3 The theorem of Euler-Fermat

In order to prove the RSA procedure, we need Fermat's theorem and its generalisation (Euler-Fermat theorem) – please see chapter 3.5.

Theorem 4.8.5. Fermat's Little Theorem³³ *Let p be a prime number and a be a random whole number, then:*

$$a^p \equiv a \pmod{p}.$$

An alternative formulation of Fermat's Little Theorem is as follows: Let p be a prime number and a be a random whole number that is relatively prime to p , then:

$$a^{p-1} \equiv 1 \pmod{p}.$$

Theorem 4.8.6. Euler-Fermat theorem (generalization of Fermat's Little Theorem)
For all elements a in the group \mathbb{Z}_n^ (i.e. a and n are natural numbers that are co-prime):*

$$a^{J(n)} \equiv 1 \pmod{n}.$$

This theorem states that if we raise a group element (here a) to the power of the order of the group (here $J(n)$), we always obtain the neutral element for multiplication (the number 1).

The 2nd formulation of Fermat's Little Theorem is derived directly from Euler's theorem if n is a prime number.

If n is the product of two prime numbers, we can - in certain cases - use Euler's theorem to calculate the result of a modular power very quickly. We have: $a^{(p-1)*(q-1)} \equiv 1 \pmod{pq}$.

Examples for calculating a modular power:

- With $2 = 1 * 2$ and $6 = 2 * 3$ where 2 and 3 are both prime; $J(6) = 2$ because only 1 and 5 are relatively prime to 6, we obtain the equation $5^2 \equiv 5^{J(6)} \equiv 1 \pmod{6}$, without having to calculate the power.
- With $792 = 22 * 36$ and $23 * 37 = 851$ where 23 and 37 are both prime, it follows for $31 \in \mathbb{Z}_{851}^*$ that $31^{792} \equiv 31^{J(23*37)} \equiv 31^{J(851)} \equiv 1 \pmod{851}$.

4.8.4 Calculation of the multiplicative inverse

Another interesting application is a special case of determining the multiplicative inverses using the Euler-Fermat theorem (multiplicative inverses are otherwise determined using the extended Euclidean algorithm).

³³Pierre de Fermat, French mathematician, Aug 17, 1601 – Jan 12, 1665.

Example:

Find the multiplicative inverse of 1579 modulo 7351.

According to Euler-Fermat: $a^{J(n)} \equiv 1 \pmod{n}$ for all a in \mathbb{Z}_n^* . If we divide both sides by a , we get: $a^{J(n)-1} \equiv a^{-1} \pmod{n}$. For the special case that the modulo is prime, we have $J(n) = p - 1$. Therefore, the modular inverse is

$$a^{-1} = a^{J(n)-1} \equiv a^{(p-1)-1} \equiv a^{p-2} \pmod{p}.$$

For our example, this means:

Since the modulus 7351 is prime, $p - 2 = 7349$.

$$1579^{-1} \equiv 1579^{7349} \pmod{p}.$$

By cleverly breaking down the exponent, we can calculate this power relatively easily³⁴:

$$7349 = 4096 + 2048 + 1024 + 128 + 32 + 16 + 4 + 1$$

$$1579^{-1} \equiv 4716 \pmod{7351}$$

4.8.5 How many private RSA keys d are there modulo 26

According to theorem 4.6.3, the arithmetic operations of modular expressions are performed in the exponents modulo $J(n)$ rather than modulo n .³⁵

In $a^{e*d} \equiv a^1 \pmod{n}$, if we wish to determine the inverses for the factor e in the exponent, we need to calculate modulo $J(n)$.

Example: (with reference to the RSA algorithm)

If we calculate modulo 26, which set can e and d come from?

Solution: We have $e * d \equiv 1 \pmod{J(26)}$.

The reduced remainder set $R' = \mathbb{Z}_{26}^* = \{1, 3, 5, 7, 9, 11, 15, 17, 19, 21, 23, 25\}$ are the elements in \mathbb{Z}_{26} , which have a multiplicative inverse, that is which are relatively prime to 26 (see 4.7.2).

The reduced remainder set R'' contains only the elements of R' that are relatively prime to $J(n) = 12 : R'' = \{1, 5, 7, 11\}$.

For every e in R'' there exists a d in R'' such that $a \equiv (a^e)^d \pmod{n}$.

For every e in R'' , there exists therefore precisely one element (not necessarily different from e) such that $e * d \equiv 1 \pmod{J(26)}$.

The general case, where n can be any integer (the sample here had n fixed to 26), is considered in chapter 4.19.6. There is a Sage program, calculating the number of all d . For all e that are relatively prime to $J(n)$ we can calculate d as follows using the Euler-Fermat theorem:

For $a^{J(n)} \equiv 1 \pmod{n}$ is the same as saying $a^{J(n)-1} \equiv a^{-1} \pmod{n}$. Therefore

$$d \equiv e^{-1} \pmod{J(n)} \equiv e^{J(J(n))-1} \pmod{J(n)}.$$

³⁴See section 4.6.4, “Fast calculation of high powers”.

³⁵For the following example, we will adopt the usual practice for the RSA procedure of using “ n ” rather than “ m ” to denote the modulus.

The problems of factorizing $n = pq$ with $q \neq p$ and of finding $J(n)$ have a similar degree of difficulty, and if we find a solution for one of the two problems, we also have a solution for the other³⁶ (please compare requisition 3 in section 4.10.1).

4.9 Multiplicative order and primitive roots³⁷

The multiplicative order and the primitive root are two useful constructs (concepts) in elementary number theory.

Mathematicians often ask, in which conditions the repeated application of an operation results in the neutral element (compare Patterns and Structures, chapter 4.8.1).

For the i -times successive modular multiplication of a number a with $i = 1, \dots, m - 1$ the product is the neutral element of the multiplication if and only if a and m are relatively prime.

Definition 4.9.1. *The multiplicative order $\text{ord}_m(a)$ of a whole number $a \pmod{m}$ (where a and m are co-prime) is the smallest whole number i for which $a^i \equiv 1 \pmod{m}$.*

The following table shows that in a multiplicative group (here \mathbb{Z}_{11}^*) not all numbers necessarily have the same order. The orders in this case are 1, 2, 5 and 10 and we notice that:

1. The orders are all factors of 10.
2. The numbers $a = 2, 6, 7$ and 8 have the order 10 - we say that these numbers have the **maximum order** in \mathbb{Z}_{11}^* .

³⁶If we know the factors of $n = p * q$ with $p \neq q$, then $J(n) = (p - 1) * (q - 1) = n - (p + q) + 1$. Additionally the factors p and q are solutions of the quadratic equation: $x^2 - (p + q)x + pq = 0$.

If only n and $J(n)$ are known, then it is: $pq = n$ and $p + q = n - J(n) + 1$. So you get p and q by solving the equation

$$x^2 + (J(n) - n - 1)x + n = 0$$

³⁷With the educational tool for number theory **NT** you can have a playful experience with primitive roots (see learning unit 2.2, pages 10-14/40 and 24-40/40).

NT can be called in CT1 via the menu path **Indiv. Procedures \ Number Theory Interactive \ Learning Tool for Number Theory**. See appendix A.6.

Example 1:

The following table 4.7³⁸ shows the values $a^i \bmod 11$ for the exponents $i = 1, 2, \dots, 10$, and for the bases $a = 1, 2, \dots, 10$ as well as the resulting value $\text{ord}_{11}(a)$ for each a .

	i=1	i=2	i=3	i=4	i=5	i=6	i=7	i=8	i=9	i=10	$\text{ord}_{11}(a)$
a=1	1	1	1	1	1	1	1	1	1	1	1
a=2	2	4	8	5	10	9	7	3	6	1	10
a=3	3	9	5	4	1	3	9	5	4	1	5
a=4	4	5	9	3	1	4	5	9	3	1	5
a=5	5	3	4	9	1	5	3	4	9	1	5
a=6	6	3	7	9	10	5	8	4	2	1	10
a=7	7	5	2	3	10	4	6	9	8	1	10
a=8	8	9	6	4	10	3	2	5	7	1	10
a=9	9	4	3	5	1	9	4	3	5	1	5
a=10	10	1	10	1	10	1	10	1	10	1	2

Table 4.7: Values of $a^i \bmod 11, 1 \leq a, i < 11$ and according order of $a \bmod m$

Table 4.7 shows, for example, that the order of 3 modulo 11 has the value 5.

Definition 4.9.2. If a and m are co-prime and if $\text{ord}_m(a) = J(m)$ (i.e. a has maximum order), then we say that a is a **primitive root** of m .³⁹

Not for every modulo m there is a number a , which is a primitive root. In the table 4.7, only $a = 2, 6, 7$ and 8 are a primitive root with respect to mod 11 ($\text{ord}_{11}(a) = J(11) = 10$).

Using the primitive roots, we can clearly establish the conditions for which powers modulo m have a unique inverse and the calculation in the exponents is manageable.

The following two tables 4.8 and 4.9 show the multiplicative orders and primitive roots modulo 45 and modulo 46.

³⁸The source code to generate table 4.7 with Sage can be found in chapter 4.19.3 “Multiplicative order”.

³⁹In chapter 4.19.4 “Primitive roots” there are Sage programs to calculate primitive roots.

Example 2:

The following table 4.8⁴⁰ shows the values $a^i \bmod 45$ for the exponents $i = 1, 2, \dots, 12$ and for the bases $a = 1, 2, \dots, 12$ as well as the resulting value $\text{ord}_{45}(a)$ for each a .

$a \setminus i$	1	2	3	4	5	6	7	8	9	10	11	12	$\text{ord}_{45}(a)$	$J(45)$
1	1	1	1	1	1	1	1	1	1	1	1	1	1	24
2	2	4	8	16	32	19	38	31	17	34	23	1	12	24
3	3	9	27	36	18	9	27	36	18	9	27	36	—	24
4	4	16	19	31	34	1	4	16	19	31	34	1	6	24
5	5	25	35	40	20	10	5	25	35	40	20	10	—	24
6	6	36	36	36	36	36	36	36	36	36	36	36	—	24
7	7	4	28	16	22	19	43	31	37	34	13	1	12	24
8	8	19	17	1	8	19	17	1	8	19	17	1	4	24
9	9	36	9	36	9	36	9	36	9	36	9	36	—	24
10	10	10	10	10	10	10	10	10	10	10	10	10	—	24
11	11	31	26	16	41	1	11	31	26	16	41	1	6	24
12	12	9	18	36	27	9	18	36	27	9	18	36	—	24

Table 4.8: Values of $a^i \bmod 45$, $1 \leq a, i < 13$

$J(45)$ is calculated using theorem 4.8.4: $J(45) = J(3^2 * 5) = 3^1 * 2 * 4 = 24$.

Since 45 is not a prime, there is no “multiplicative order” for all values of a (for all numbers that are not relatively prime to 45 : 3, 5, 6, 9, 10, 12, …, because $45 = 3^2 * 5$).

Example 3:

Is 7 a primitive root modulo 45?

The necessary, but not sufficient requirement/condition $\text{gcd}(7, 45) = 1$ is fulfilled. Table 4.8 shows that the number $a = 7$ is not a primitive root of 45, because $\text{ord}_{45}(7) = 12 \neq 24 = J(45)$.

⁴⁰The source code to generate table 4.8 with Sage can be found in chapter 4.19.3, “Multiplicative order”.

Example 4:

The following table 4.9⁴¹ answers the question as to whether the number $a = 7$ is a primitive root of 46. The necessary, but not sufficient requirement/condition $\gcd(7, 46) = 1$ is fulfilled.

$J(46)$ is calculated using theorem 4.8.2: $J(46) = J(2 * 23) = 1 * 22 = 22$. The number 7 is a primitive root of 46, because $\text{ord}_{46}(7) = 2 = J(46)$.

$a \setminus i$	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	ord
1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	
2	2	4	8	16	32	18	36	26	6	12	24	2	4	8	16	32	18	36	26	6	12	24	2	
3	3	9	27	35	13	39	25	29	41	31	1	3	9	27	35	13	39	25	29	41	31	1	3	
4	4	16	18	26	12	2	8	32	36	6	24	4	16	18	26	12	2	8	32	36	6	24	4	
5	5	25	33	27	43	31	17	39	11	9	45	41	21	13	19	3	15	29	7	35	37	1	5	
6	6	36	32	8	2	12	26	18	16	4	24	6	36	32	8	2	12	26	18	16	4	24	6	
7	7	3	21	9	17	27	5	35	15	13	45	39	43	25	37	29	19	41	11	31	33	1	7	
8	8	18	6	2	16	36	12	4	32	26	24	8	18	6	2	16	36	12	4	32	26	24	8	
9	9	35	39	29	31	3	27	13	25	41	1	9	35	39	29	31	3	27	13	25	41	1	9	
10	10	8	34	18	42	6	14	2	20	16	22	36	38	12	28	4	40	32	44	26	30	24	10	—
11	11	29	43	13	5	9	7	31	19	25	45	35	17	3	33	41	37	39	15	27	21	1	11	
12	12	6	26	36	18	32	16	8	4	2	24	12	6	26	36	18	32	16	8	4	2	24	12	
13	13	31	35	41	27	29	9	25	3	39	1	13	31	35	41	27	29	9	25	3	39	1	13	
14	14	12	30	6	38	26	42	36	44	18	22	32	34	16	40	8	20	4	10	2	28	24	14	
15	15	41	17	25	7	13	11	27	37	3	45	31	5	29	21	39	33	35	19	9	43	1	15	
16	16	26	2	32	6	4	18	12	8	36	24	16	26	2	32	6	4	18	12	8	36	24	16	
17	17	13	37	31	21	35	43	41	7	27	45	29	33	9	15	25	11	3	5	39	19	1	17	
18	18	2	36	4	26	8	6	16	12	32	24	18	2	36	4	26	8	6	16	12	32	24	18	
19	19	39	5	3	11	25	15	9	33	29	45	27	7	41	43	35	21	31	37	13	17	1	19	
20	20	32	42	12	10	16	44	6	28	8	22	26	14	4	34	36	30	2	40	18	38	24	20	
21	21	27	15	39	37	41	33	3	17	35	45	25	19	31	7	9	5	13	43	29	11	1	21	
22	22	24	22	24	22	24	22	24	22	24	22	24	22	24	22	24	22	24	22	24	22	24	22	
23	23	23	23	23	23	23	23	23	23	23	23	23	23	23	23	23	23	23	23	23	23	23	—	

Table 4.9: Values of $a^i \pmod{46}$, $1 \leq a, i < 24$

Theorem 4.9.1. Given a modulus n and a number a , relative prime to n the following holds: The set $\{a^i \pmod{n} \mid i = 1, \dots, J(n)\}$ equals the multiplicative group Z_n^* if and only if $\text{ord}_n(a) = J(n)$.^{42,43}

The multiplicative group Z_n^* only contains all values from 1 to $n - 1$, if n is prime (see 4.7.2).

⁴¹See chapter 4.19.3, “Multiplicative order” for the source code to generate the table 4.9 using Sage.

⁴²For prime moduli p all a with $0 < a < p$ are of order $J(p) = p - 1$. Compare table 4.8 for an example. In this case $a^i \pmod{n}$ goes through all the values $1, \dots, p - 1$. Exhausting all possible values of the set is an important cryptographic proposition (compare theorem 4.6.2). This determines a permutation $\pi(p - 1)$.

⁴³Table 4.9 demonstrates that for composite moduli n not all a are of maximal order $J(n)$. In this example only 5, 7, 11, 15, 17, 19 and 21 are of order 22.

Example 5: Length of cycles

The following tables 4.10 and 4.11⁴⁴ serve as samples to introduce cycle lengths – this is a topic which goes beyond the multiplicative order.

Cycle here means a sequence of numbers $a^i \bmod n$ with $1 \leq i < n$ for a given a , and a repeating sequence. Within the cycle each number is unique. Generally, cycles don't have to contain the 1 – in the opposite to cycles which have a multiplicative order ≥ 1 (they yet have the 1 always at the end).

With l we low mean the cycle length.

The maximum cycle length l_{max} is $J(n)$.

For the following tables 4.10 and 4.11 $J(n)$ is (according to theorem 4.8.4):

- $J(14) = J(2 * 7) = 1 * 6 = 6$.
- $J(22) = J(2 * 11) = 1 * 10 = 10$.

a) If the multiplicative order exists for a , (indendently whether a is prim) it is: $ord_n(a) = l$.

Samples: The maximum length l_{max} ⁴⁵ is achieved e.g. for:

- $a = 3$ with $l_{max} = ord_{14}(a) = 6$ in table 4.10, or
- $a = 10$ with $l_{max} = ord_{22}(a) = 10$ in table 4.11.

b) Also, if no multiplicative order exists for a , the maximum cycle length can be achieved.⁴⁶

Samples:

- In table 4.10: $l_{max} = J(14) = 6$ for $a = 10, 12$.
- In table 4.11: $l_{max} = J(22) = 10$ for $a = 2, 6, 8, 18$.

$a \setminus i$	1	2	3	4	5	6	7	8	9	10	11	12	13	$ord_{14}(a)$	$J(14)$	l
1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	6	1
2	2	4	8	2	4	8	2	4	8	2	4	8	2	0	6	3
3	3	9	13	11	5	1	3	9	13	11	5	1	3	6	6	6
4	4	2	8	4	2	8	4	2	8	4	2	8	4	0	6	3
5	5	11	13	9	3	1	5	11	13	9	3	1	5	6	6	6
6	6	8	6	8	6	8	6	8	6	8	6	8	6	0	6	2
7	7	7	7	7	7	7	7	7	7	7	7	7	7	0	6	1
8	8	8	8	8	8	8	8	8	8	8	8	8	8	0	6	1
9	9	11	1	9	11	1	9	11	1	9	11	1	9	3	6	3
10	10	2	6	4	12	8	10	2	6	4	12	8	10	0	6	6
11	11	9	1	11	9	1	11	9	1	11	9	1	11	3	6	3
12	12	4	6	2	10	8	12	4	6	2	10	8	12	0	6	6
13	13	1	13	1	13	1	13	1	13	1	13	1	13	2	6	2
14	0	0	0	0	0	0	0	0	0	0	0	0	0	0	6	1
15	1	1	1	1	1	1	1	1	1	1	1	1	1	1	6	1
16	2	4	8	2	4	8	2	4	8	2	4	8	2	0	6	3

Table 4.10: Values of $a^i \bmod 14, 1 \leq a < 17, i < 14$

⁴⁴See chapter 4.19.3, “Multiplicative order” for the source code to generate the tables 4.10 und 4.11 using Sage.

⁴⁵We don't know of a formular telling for which a the length has a maximum.

⁴⁶Here the sequences are built via $a^i \bmod n$ with $1 \leq i < n$. For composite numbers n the sequences never contain all numbers $1, \dots, n - 1$.

This should not be mixed up with RSA, where the “sequence” is built differently, $m^e \bmod n$ with $0 \leq m < n$, and this sequence then takes all numbers $0, \dots, n - 1$ (permutation).

$a \setminus i$	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	$ord_{22}(a)$	l
1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
2	2	4	8	16	10	20	18	14	6	12	2	4	8	16	10	20	18	14	6	12	2	0	10
3	3	9	5	15	1	3	9	5	15	1	3	9	5	15	1	3	9	5	15	1	3	5	5
4	4	16	20	14	12	4	16	20	14	12	4	16	20	14	12	4	16	20	14	12	4	0	5
5	5	3	15	9	1	5	3	15	9	1	5	3	15	9	1	5	3	15	9	1	5	5	5
6	6	14	18	20	10	16	8	4	2	12	6	14	18	20	10	16	8	4	2	12	6	0	10
7	7	5	13	3	21	15	17	9	19	1	7	5	13	3	21	15	17	9	19	1	7	10	10
8	8	20	6	4	10	14	2	16	18	12	8	20	6	4	10	14	2	16	18	12	8	0	10
9	9	15	3	5	1	9	15	3	5	1	9	15	3	5	1	9	15	3	5	1	9	5	5
10	10	12	10	12	10	12	10	12	10	12	10	12	10	12	10	12	10	12	10	12	10	0	2
11	11	11	11	11	11	11	11	11	11	11	11	11	11	11	11	11	11	11	11	11	11	0	1
12	12	12	12	12	12	12	12	12	12	12	12	12	12	12	12	12	12	12	12	12	12	0	1
13	13	15	19	5	21	9	7	3	17	1	13	15	19	5	21	9	7	3	17	1	13	10	10
14	14	20	16	4	12	14	20	16	4	12	14	20	16	4	12	14	20	16	4	12	14	0	5
15	15	5	9	3	1	15	5	9	3	1	15	5	9	3	1	15	5	9	3	1	15	5	5
16	16	14	4	20	12	16	14	4	20	12	16	14	4	20	12	16	14	4	20	12	16	0	5
17	17	3	7	9	21	5	19	15	13	1	17	3	7	9	21	5	19	15	13	1	17	10	10
18	18	16	2	14	10	4	6	20	8	12	18	16	2	14	10	4	6	20	8	12	18	0	10
19	19	9	17	15	21	3	13	5	7	1	19	9	17	15	21	3	13	5	7	1	19	10	10
20	20	4	14	16	12	20	4	14	16	12	20	4	14	16	12	20	4	14	16	12	20	0	5
21	21	1	21	1	21	1	21	1	21	1	21	1	21	1	21	1	21	1	21	1	21	2	2
22	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1
23	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
24	2	4	8	16	10	20	18	14	6	12	2	4	8	16	10	20	18	14	6	12	2	0	10
25	3	9	5	15	1	3	9	5	15	1	3	9	5	15	1	3	9	5	15	1	3	5	5

Table 4.11: Values of $a^i \bmod 22$, $1 \leq a < 26$, $i < 22$

4.10 Proof of the RSA procedure with Euler-Fermat

Using the Euler-Fermat theorem, we can “prove” the RSA⁴⁷ procedure in the group \mathbb{Z}_n^* .

4.10.1 Basic idea of public key cryptography

The basic idea behind public key cryptography is that all participants possess a different pair of keys (P and S) and the public keys for all recipients are published. You can retrieve the public key P for a recipient from a directory just as you would look up some one’s phone number in the phone book. Furthermore, each recipient has a secret key S that is needed in order to decrypt the message and that is not known to anyone else. If the sender wishes to send a message M , he encrypts it using the public key P of the recipient before sending it:

The ciphertext C is determined as $C = E(P; M)$, where E (encryption) is the encryption rule. The recipient uses his private key S to decrypt the message with the decryption rule $D : M = D(S; C)$.

In order to ensure that this system works for every message M , the following four **requirements** must be met:

1. $D(S; E(P; M)) = M$ for every M (invertibility) and M takes “very many” of its possible values.
2. All (S, P) pairs are different for all participants.
3. The time required to derive S from P is at least as high as the time required to decrypt M with no knowledge of S .
4. Both C and M can be calculated relatively easily.

The 1st requirement is a general condition for all cryptographic encryption algorithms.

The prerequisite of the 2nd requirement can easily be met because there is a “very” large number of prime numbers⁴⁸. In addition, that this can be ensured by a central office that issues certificates (see chapter 4.11.5.4, S. 144).

It is this last requirement that makes the procedure actually usable. This is because it is possible to calculate the powers in a linear amount of time (because there is a restriction on the length of the numbers).

Although Whitfield Diffie and Martin Hellman formulated the general method as early as 1976, the actual procedure that met all four requirements was only discovered later by Rivest, Shamir and Adleman.

⁴⁷The RSA procedure is the most common asymmetric cryptography procedure. Developed in 1978 by Ronald Rivest, Adi Shamir and Leonard Adleman, it can be used both for signatures and for encryption. Cryptographers always associate this procedure with the abbreviation “RSA” – the following remark is meant with humor to show that each letter combination can be used with several meanings: In Britain the “Royal Society for the encouragement of Arts, Manufactures & Commerce” is commonly known as the “RSA”.

⁴⁸According to the **prime number theorem** (chapter 3.7.2, p. 72) of Legendre and Gauss there are approximately $n / \ln(n)$ prime numbers up to the number n . This means, for example, that there are $6.5 * 10^{74}$ prime numbers under $n = 2^{256}$ ($= 1.1 * 10^{77}$) and $3.2 * 10^{74}$ prime numbers under $n = 2^{255}$. Between 2^{255} and 2^{256} there are therefore $3.3 * 10^{74}$ prime numbers with precisely 256 bits. This large number is also the reason why we cannot simply save them all.

4.10.2 How the RSA procedure works

The individual steps for implementing the RSA procedure can be described as follows (see [Eckert2003, p. 213 ff] and [Sedgewick1990, p. 338 ff]). Steps 1 to 3 constitute key generation, steps 4 and 5 are the encryption, and steps 6 and 7 are the decryption:

1. Select two distinct random prime numbers^{49,50} p and q and calculate $n = p * q$.⁵¹
The value n is called the RSA modulus.⁵²
 2. Select an arbitrary $e \in \{2, \dots, n - 1\}$ such that⁵³:
 e is relatively prime to $J(n) = (p - 1) * (q - 1)$.
We can then “throw away” p and q .⁵⁴
 3. Select $d \in \{1, \dots, n - 1\}$ with $e * d \equiv 1 \pmod{J(n)}$, i.e. d is the multiplicative inverse of e modulo $J(n)$.^{55,56} We can then “throw away” $J(n)$.
 - (n, e) is the public key P .
 - (n, d) is the private key S (only d must be kept secret).
 4. For encryption, the message represented as a (binary) number is divided into parts such that each part of the number is less than n .
 5. Encryption of the plaintext (or the parts of it) $M \in \{1, \dots, n - 1\}$:
- $$C = E((n, e); M) := M^e \pmod{n}.$$
6. For decryption, the ciphertext represented as a binary number is divided into parts such that each part of the number is less than n .

⁴⁹Compaq introduced the so-called multi-prime method with high marketing effort in 2000. n was the product of two big and one relative small prime: $n = o * p * q$. With theorem 4.8.3 we get: $J(n) = (o - 1) * (p - 1) * (q - 1)$. This method did not assert itself.

One reason probably is, that Compaq claimed a patent on it. Generally there is less understanding in Europe and with the Open Source Initiative, that one can claim patents on algorithms. But there is really no understanding outside the U.S., that one can get a patent for a special case (3 factors) of an algorithm (RSA), although the patent for the general case was almost expired.

JCT contains the multi-prime RSA method both within the **Visuals** menu of the Default Perspective as well as within the Algorithm Perspective.

⁵⁰If the two primes p and q are equal then $(m^e)^d \equiv m \pmod{n}$ is not true for all $m < n$ (although $e * d \equiv 1 \pmod{J(n)}$ is fulfilled). **Example:**

If $n = 5^2$ then according to theorem 4.8.4 it is $J(n) = 5 * 4 = 20$, $e = 3$, $d = 7$, $e * d = 21 \equiv 1 \pmod{J(n)}$. But it is $(5^3)^7 \equiv 0 \pmod{25}$.

⁵¹The GISA (German Information Security Agency) recommends, to choose the prime factors p and q almost the same, but not too close:

$$0.5 < |\log_2(p) - \log_2(q)| < 30.$$

They recommend to generate the primes independently and check that the restriction is fulfilled (see [?]).

⁵²In CT1 the RSA modulo is denoted with a capital “ N ” .

⁵³It is recommended by cryptanalytic reasons, but not necessary to make RSA work, to select e such that:
 $\max(p, q) < e < J(n) - 1$.

⁵⁴The procedure also allows us to select d freely and then calculate e . However, this has practical disadvantages. We usually want to be able to encrypt messages “quickly”, which is why we choose a public exponent e such that it has a short bit length compared to the modulus n and as few binary ones as possible (e.g. $2^{16} + 1$). So a fast exponentiation is possible when encrypting. We want to select the publicly known e to be an advantageous value that allows the exponential calculation to be performed quickly during encryption. The prime numbers 3, 17 and 65537 have proved to be particularly practical for this purpose. The most often used number is $65537 = 2^{16} + 1$, or in binary: $10 \dots 0 \dots 01$ (this number is prime and therefore relatively prime to many other numbers).

⁵⁵For reasons of security, d should not be too small.

⁵⁶We start by determining either d or e depending on the implementation.

7. Decryption of the ciphertext (or the parts of it) $C \in \{1, \dots, n - 1\}$:

$$M = D((n, d); C) := C^d \pmod{n}.$$

The numbers d, e and n are usually extremely large (e. g. d and e 300 bits, n 600 bits).

Comment:

The security of the RSA algorithm depends as with all public key methods on the difficulty to calculate the private key d from the public key (n, e) .

Concretely for the RSA method does this mean:

1. It is hard to calculate $J(n)$ for big compounds n and
2. It is hard to calculate the prime factors of big compounds n (factorization problem).⁵⁷

4.10.3 Proof of requirement 1 (invertibility)

For pairs of keys (n, e) and (n, d) that possess fixed properties in steps 1 to 3 of the RSA procedure, the following must be true for all $M < n$:

$$M \equiv (M^e)^d \pmod{n} \quad \text{with} \quad (M^e)^d = M^{e*d}.$$

This means that the deciphering algorithm above works correctly.

We therefore need to show that:

$$M^{e*d} \equiv M \pmod{n}.$$

We will show this in 3 steps (see [Beutelspacher1996, p. 131ff]).

Step 1:

In the first step we show that: $M^{e*d} \equiv M \pmod{p}$. This results from the requirements and from Euler-Fermat (theorem 4.8.6). Since $n = p * q$ and $J(p * q) = (p - 1) * (q - 1)$ and since e and d are selected in such a way that $e * d \equiv 1 \pmod{J(n)}$, there is a whole number k such that: $e * d = 1 + k * (p - 1) * (q - 1)$.

$$\begin{aligned} M^{e*d} &\equiv M^{1+k*(p-1)*(q-1)} \equiv M * M^{k*(p-1)*(q-1)} \pmod{p} \\ &\equiv M * (M^{p-1})^{k*(q-1)} \pmod{p} \quad \text{based on little Fermat: } M^{p-1} \equiv 1 \pmod{p} \\ &\equiv M * (1)^{k*(q-1)} \pmod{p} \\ &\equiv M \pmod{p} \end{aligned}$$

The requirement for using the simplified Euler-Fermat theorem (theorem 4.8.5) was that M and p are relatively prime.

Since this is not true in general, we need to consider the case when M and p are not relatively prime. Since p is a prime number, this implies that p is a factor of M . But this means:

$$M \equiv 0 \pmod{p}.$$

⁵⁷There is no reason for the concern sometimes mentioned that there are not enough primes: Raising the dimension (exponent) of the modul always offers enough primes to consider – this is visualized in chapter 3.13 “Appendix: Visualization of the quantity of primes in higher ranges”.

If p is a factor of M , then p is also a factor of M^{e*d} . Therefore:

$$M^{e*d} \equiv 0 \pmod{p}.$$

Since p is a factor of both M and M^{e*d} , it is also a factor of their difference:

$$(M^{e*d} - M) \equiv 0 \pmod{p}.$$

And therefore our conjecture is also true in this special case.

Step 2:

In exactly the same way we prove that: $M^{e*d} \equiv M \pmod{q}$.

Step 3:

We now combine the conjectures from step 1 and 2 for $n = p * q$ to show that:

$$M^{e*d} \equiv M \pmod{n} \text{ for all } M < n.$$

From step 1 and 2 we have $(M^{e*d} - M) \equiv 0 \pmod{p}$ and $(M^{e*d} - M) \equiv 0 \pmod{q}$. Therefore, p and q are both factors of the same number $z = (M^{e*d} - M)$. Since p and q are **distinct** prime numbers, their product must also be a factor of this number z . Thus:

$$(M^{e*d} - M) \equiv 0 \pmod{p * q} \quad \text{or} \quad M^{e*d} \equiv M \pmod{p * q} \quad \text{or} \quad M^{e*d} \equiv M \pmod{n}.$$

□

Comment 1:

We can also condense the three steps if we use the theorem 4.8.6 (Euler-Fermat) - i.e. not the simplified theorem where $n = p$ and which corresponds to Fermat's Little Theorem:

$$(M^e)^d \equiv M^{e*d} \equiv M^{(p-1)(q-1)*k+1} \equiv \left(\underbrace{M^{(p-1)(q-1)}}_{\equiv M^{J(n)} \equiv 1 \pmod{n}} \right)^k * M \equiv 1^k * M \equiv M \pmod{n}.$$

Comment 2:

When it comes to signing messages, we perform the same operations but first use the secret key d , followed by the public key e . The RSA procedure can also be used to create digital signatures, because:

$$M \equiv (M^d)^e \pmod{n}.$$

4.11 Regarding the security of the RSA algorithm⁵⁸

There have always been discussions about the suitability of the RSA algorithm for digital signatures and encryption, e. g. after publications of breakthroughs in factorization. Nevertheless the RSA algorithm has become a de-facto standard since it was published more than 20 years ago (compare 7.1).

The security of the RSA algorithm rests — as with all cryptographic methods — on the following 4 central pillars:

- the complexity of the number theoretical problem on which the algorithm is based (here factorization of big numbers),
- the election of fitting parameters (here the length of the module N),
- the adequate usage of the algorithm and key generation and
- the correct implementation of the algorithm.

Usage and key generation are well understood today. Implementation based on long integer arithmetic is very easy.

The following sections examine the RSA algorithm with respect to the first two points.

4.11.1 Complexity

Successful decryption or forgery of a signature — without knowing the private key — requires calculating the e -th root mod n . The private key, this is the multiplicative inverse of e mod $J(n)$, can be easily determined if $J(n)$ is known. $J(n)$ again can be calculated from the prime factors of n . Breaking of RSA therefore cannot be more difficult than factorization of the module n .

The best factorization method known today is a further development of the General Number Field Sieve (GNFS) , which was originally devised to factor only numbers of a special form (like Fermat numbers). The complexity of solving the factorization problem with the GNFS is asymptotically

$$O(l) = e^{c \cdot (l \cdot \ln 2)^{1/3} \cdot (\ln(l \cdot \ln 2))^{2/3} + o(l)}$$

Please refer to: [Lenstra1993] and [Silverman2000]

This formula shows, that the factorization problem belongs to the class of problems with sub-exponential time complexity (i. e. time complexity grows asymptotically not as fast as exponential functions like e^l or 2^l , but strictly slower, e. g. like $e^{\sqrt{l}}$). This classification is all that is currently known; it does not preclude the possibility that the factorization problem can be solved in polynomial time (see 4.11.5.1).

$O(l)$ is the average number of processor steps depending on the bit length l of the number n to be factorized. For the best currently known factorization algorithm the constant $c = (64/9)^{1/173} = 1923$.

The inverse proposition, that the RSA algorithm can be broken only by factorization of n , is still not proven. Most number theorists consider the “RSA problem” and the factorization problem equivalent in terms of time complexity.

Please refer to: *Handbook of Applied Cryptography* [Menezes2001].

⁵⁸Major parts of the first part of chapter 4.11 follow the article “Vorzüge und Grenzen des RSA-Verfahrens” written by F. Bourseau, D. Fox and C. Thiel [Bourseau2002].

4.11.2 Security parameters because of new algorithms

Factorization algorithms⁵⁹

The complexity is basically determined by the length l of the modulus n . Higher values for this major parameter are oriented at the possibilities of the current algorithms for factorization:

- In 1994 a 129-digit RSA modulus (428 bit), published in 1977, was factorized by a distributed implementation of the Quadratic Sieve algorithm (QS), developed 1982 by Pomerance. This effort took 8 months.

Please refer to:

C. Pomerance: *The quadratic sieve factoring algorithm* [Pomerance1984].

- In 1999 a 155-digit modulus (512 bit) was factored with an implementation of the General Number Field Sieve algorithm (GNFS), developed by Buhler, Lenstra and Pomerance. The GNFS is more efficient than QS if n is longer than about 116 decimal digits. This effort took 5 months.

Please refer to:

J.P. Buhler, H.W. Lenstra, C. Pomerance: *Factoring integers with the number field sieve* [Buhler1993].

- Ten years later, end of 2009, a 232-digit modulus (768 bit) was factored by Kleinjung etc. after 2 1/2 years.

Please refer to:

T. Kleinjung, et. al.: *Factorization of a 768-bit RSA modulus* [Kleinjung2010].

This made practically evident that a module length of 768 bit no longer prevents from attackers.

Details about factorization progress since 1999 see chapter 4.11.4.

Lattice base reduction algorithms

The module length l is not the only parameter relevant for security. Beneath requirements from implementation and engineering the sizes and the proportions of the parameters e , d and n are relevant.

According attacks based on lattice reductions are a real threat for (too) simple implementations of RSA. These attacks can be structured into the following four categories:

- Attacks against very small public keys e (e.g. $e = 3$).
- Attacks against relatively small private exponents d (e.g. $d < n^{0.5}$).
- Factorization of the modulus n , if one of the factors p or q is *partly* known.
- Attacks requiring, that a *part* of the private key d is known.

A good overview concerning these attacks can be found in the diploma thesis of Matthias Schneider [SchneiderM2004].

⁵⁹With the educational tool for number theory **NT** you can gather more experience with current factorization algorithms (see learning unit 5.1-5.5, pages 1-15/15).

NT can be called in CT1 via the menu path **Indiv. Procedures \ Number Theory Interactive \ Learning Tool for Number Theory**. See appendix A.6.

The quadratic sieve (QS) can be found in CT1 and CT2; GNFS, the most modern factorisation method for moduli bigger than 130 decimal digits, is only part of CT2 (via the msieve library).

4.11.3 Forecasts about factorization of large integers

Since 1980 a lot of progress has been made. Estimations about the future development of the ability to factor RSA modules vary and depend on some assumptions:

- progression in computing performance (Moore's law: every 18 month the computing power will double) and in grid computing.
- development of new algorithms.

Within the last years the module bit length feasible for factorization increased — even without new algorithms — by 10 bit per year. Larger numbers require not only more time to be factored, but also huge RAM storage for the solutions matrix being used by the best algorithms known today. This need for storage grows like the square root of the computation time, i. e. also sub-exponentially. Because RAM availability increased exponentially in the recent decades, it seems that this should not be the limiting factor.

An estimation of the evolution of secure key lengths was done by Lenstra/Verheul in 1999 [Lenstra1999] (compare figure 7.1 in chapter 7.1).

Within the article [Bourseau2002] Dirk Fox⁶⁰ published his prognosis of an almost linear factorization progression, if all influencing factors are included: Each year the module length feasible for factorization increases by 20 bit on average. So his forecast was below the more optimistic estimations of GISA and NIST.

This forecast by Dirk Fox from the year 2001 seems to prove true by the latest factorization records of RSA-200 and RSA-768 (see chapter 4.11.4). His estimation for the year 2005, to achieve a bit length of 660 bit, was almost a precision landing (compare figure 4.1).

If the forecast withstands in the future then the factorization of an RSA modulus of 1024 bit can be expected in the year 2020.

⁶⁰His company Secorvo Ltd delivered a statement on the recommendation for key length selection published by the GISA (German Information Security Agency). Chapter 2.3.1 of this statement contains a competent and understandable discussion of RSA security (this document exists – to my knowledge – only in German): <http://www.secorvo.de/publikat/stellungnahme-algorithmenempfehlung-020307.pdf>

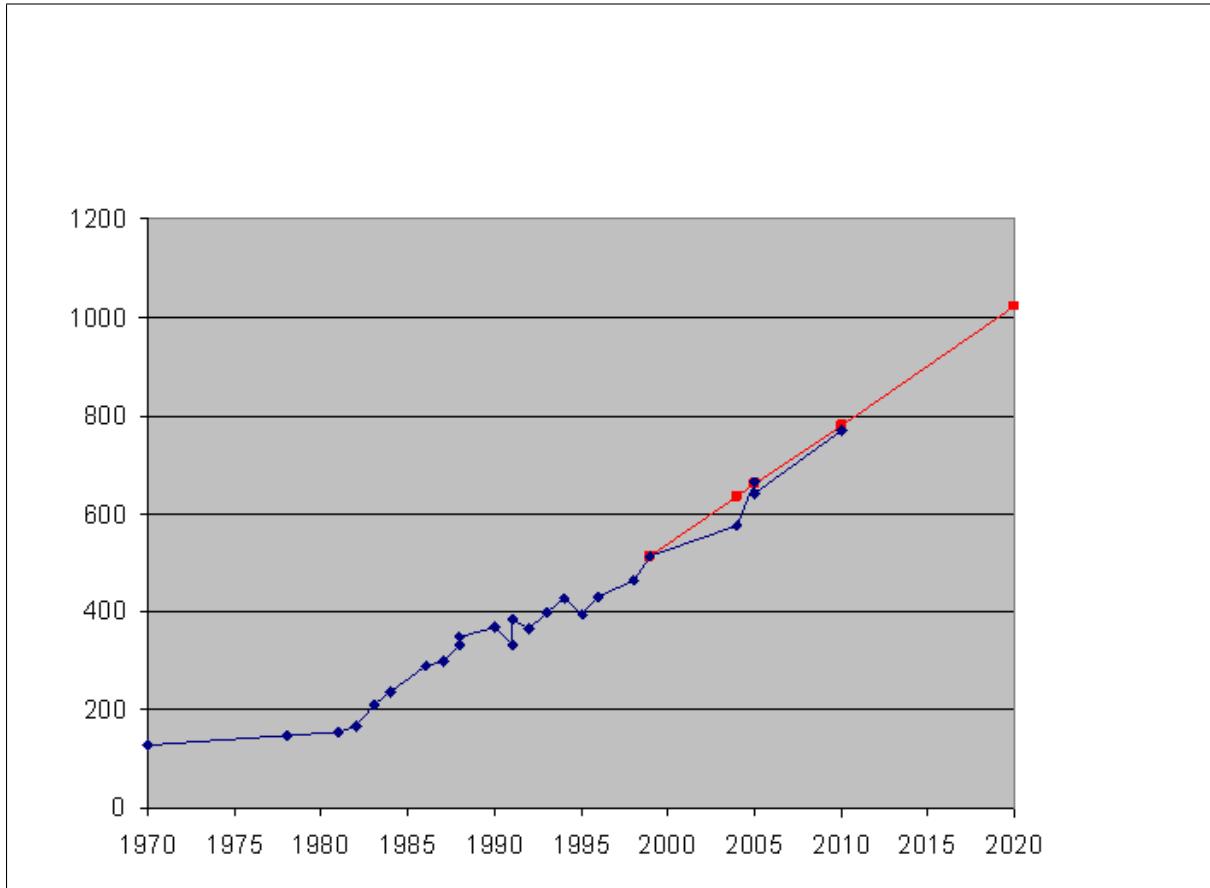


Figure 4.1: Comparison of the published factorization records (blue) and of the predicted development (red) [Source Fox 2001]

To let the possible happen, you again and again have to try the impossible.

Quote 10: Hermann Hesse⁶¹

4.11.4 Status regarding factorization of concrete large numbers

An exhaustive overview about the factoring records of composed integers using different methods can be found on the following web pages:

<http://www.crypto-world.com>

http://www.tutorgig.com/ed/RSA_number The RSA Factoring Challenge

http://en.wikipedia.org/wiki/Integer_factorization_records

http://en.wikipedia.org/wiki/RSA_Factoring_Challenge

The current record (as of Nov. 2012) obtained using the GNFS method (General Number Field Sieve) factorized a general 232 decimal digit into its both prime factors.

The last records⁶² with factorization algorithms for composed numbers are listed in the table 4.12.

	Decimal digits	Binary digits	Factored on	Factored by
RSA-768	232	768	Dec 2010	Thorsten Kleinjung et al.
RSA-200	200	663	May 2005	Jens Franke et al.
RSA-640 ⁶³	193	640	Nov 2005	Jens Franke et al.
RSA-576	174	576	Dec 2003	Jens Franke et al.
RSA-160	160	530	Apr 2003	Jens Franke et al.
RSA-155	155	512	Aug 1999	Herman te Riele et al.
...				
C307	307	1017	May 2007	Jens Franke et al.
C176	176	583	May 2005	Kazumaro Aoki et al.
C158	158	523	Jan 2002	Jens Franke et al.

Table 4.12: The current factoring records (as of Nov. 2012)

⁶¹Hermann Hesse, German/Swiss writer and Nobel Prize winner, July 2, 1877 – August 9, 1962.

⁶²The 'RSA numbers' are certain large semiprime numbers (i.e., numbers with exactly two prime factors). They were generated and published by the company RSA Security: In the RSA Factoring Challenge the prime factors for these numbers are sought.

See <http://www.rsa.com/rsalabs/node.asp?id=2092>.

RSA Labs offers its challenges since the beginning of the 1990th. The first RSA Factoring Challenge labeled the numbers, from RSA-100 to RSA-500, according to their number of decimal digits; the second RSA Factoring Challenge labeled the numbers according to their number of binary digits. Within the second challenge cash prizes were offered for successful factorizations of RSA-576 to RSA-2048 (RSA-576, RSA-640 etc. using 64 bit steps upwards — An exception to this is RSA-617, which was created prior to the change in the numbering scheme). But the RSA challenges ended ahead of time in 2007, RSA Inc. retracted the prize. All till now unsolved RSA challenges of RSA Labs can also be found at the website of the cipher challenge "MysteryTwister C3" (<http://www.mysterytwisterc3.org>).

The 'C numbers' originate from the Cunningham project: <http://www.cerias.purdue.edu/homes/ssw/cun/>. These are factors of Mersenne numbers, which have a very special form. This makes it an order of magnitude easier to factor them as moduli of the same length build for RSA.

⁶³A research group of the GISA solved this challenge which was awarded with 20,000 US dollar using the GNFS method. The researchers needed about five months to divide this number into its both 320 bit long prime factors.

Experiences about the ellipsed time of factorization with the open source software Pari-GP, Sage, CrypTool 1 and CrypTool 2) can be found in “Zeitexperimente zur Faktorisierung”.⁶⁴

Below these last records listed in table 4.12 are explained in more detail⁶⁵:

RSA-155

On August 22, 1999 researchers from the Netherlands found the solution of this RSA challenge. They factorized a 155-digit number into its both 78-digit primes (see chapter 4.11.2).

This 512 bit RSA-155 meant to reach a kind of *magic* border.

C158

On January 18, 2002 researchers at the German University of Bonn⁶⁶ factorized a 158-digit decimal number into its both prime factors (these are build with 73 and 86 decimal digits) using the GNFS method (General Number Field Sieve).

This record got much less attention within the press than the solution of RSA-155.

The task of the researchers from Bonn was not initiated by a challenge, but they wanted to find the last prime factors of the integer $2^{953} - 1$ (see “Wanted List” of the Cunningham Project⁶⁷).

The 6 smaller prime factors, already found before have been:

3, 1907, 425796183929,
1624700279478894385598779655842584377,
3802306738549441324432139091271828121 and
128064886830166671444802576129115872060027.

The first 3 factors can be easily computed⁶⁸. The next three prime factors were found by P. Zimmerman⁶⁹, T. Grandlund⁷⁰ and R. Harley during the years 1999 and 2000 using the elliptic curve factorization method.

The last remaining factor, called “C158”, was known to be composite by then, but its factors

The researchers around Professor Jens Franke (from the University of Bonn, the GISA and the CWI) do not aim on getting cash prizes but in extending the research limits. So statements about the necessary length of a secure RSA modulus are more well-founded.

See <http://www.heise.de/newsticker/meldung/print/65957>.

⁶⁴R.-H. Schulz und H. Witten: “Zeitexperimente zur Faktorisierung”. Ein Beitrag zur Didaktik der Kryptographie”, in LogIn Heft Nr. 166/167 (2010) 113-120, (available currently only in German)
http://bscw.schule.de/pub/bscw.cgi/d864899/Schulz_Witten_Zeit-Experimente.pdf

⁶⁵The two methods, GNFS and SNFS, used to do so are shortly illustrated at the following web pages:

http://en.wikipedia.org/wiki/Special_number_field_sieve
http://en.wikipedia.org/wiki/General_number_field_sieve

⁶⁶http://www.ercim.org/publication/Ercim_News/enw49/franke.html
⁶⁷Cunningham project: <http://www.cerias.purdue.edu/homes/ssw/cun/>

⁶⁸E.g. using CT1 via menu **Indiv. Procedures \ RSA Cryptosystem \ Factorization of a Number**.

CT1 can factorize in a reasonable time numbers no longer than 250 bit (Numbers bigger than 1024 bits are not accepted by CT1). CT2 is able to factorize numbers bigger than 250 bit length.

⁶⁹<http://www.loria.fr/~zimmerma/ecmnet>

⁷⁰<http://www.swox.se/gmp/>

were not known (the following 3 lines contain one number):

```
39505874583265144526419767800614481996020776460304936  
45413937605157935562652945068360972784246821953509354  
4305870490251995655335710209799226484977949442955603
```

The factorization of C158 resulted in the following two 73- and 86-digit prime factors:

```
3388495837466721394368393204672181522815830368604993048084925840555281177
```

and

```
1165882340667125990314837655838327081813101  
2258146392600439520994131344334162924536139.
```

So now all 8 prime factors of $2^{953} - 1$ have been found.

Links:

<http://www.loria.fr/~zimmerma/records/gnfs158>
<http://www.crypto-world.com/FactorRecords.html>
<http://www.crypto-world.com/announcements/c158.txt>

RSA-160

On January 18, 2002 researchers at the German University of Bonn⁷¹ factorized a 160-digit number into its both prime factors (these are build with each 80 decimal digits) using the GNFS method (General Number Field Sieve).

The computations for the factorization of RSA-160 also took place at the German Information Security Agency (GISA) in Bonn.⁷²

The 160-digit decimal number origins from the old challenge list of RSADSI. This number was retracted after RSA-155 (RSA512) had been factorized successfully. The prime factors of RSA-160 were still unknown. So this record of the team of Prof. Franke provides the solution of the old challenge, for which RSADSI didn't award a price anymore.

The composite number called "RSA-160" is (the following 3 lines contain one number):

```
215274110271888970189601520131282542925777358884567598017049  
767677813314521885913567301105977349105960249790711158521430  
2079314665202840140619946994927570407753
```

The factorization of RSA-160 resulted in the following two prime factors:

$$p = 45427892858481394071686190649738831$$
$$656137145778469793250959984709250004157335359$$

⁷¹<http://www.loria.fr/~zimmerma/records/rsa160>
<http://www.loria.fr/~zimmerma/records/factor.html>
<http://www.crypto-world.com/FactorWorld.html>

⁷²Every year the GISA creates a paper to describe which crypto algorithms are feasible to generate digital signatures according to the German signature law – under participation of experts from economy and science. To review signature methods based on the factorization problem the GISA also co-operates with researchers from the University of Bonn. Further information about crypto algorithms can be found on the web page of GISA: <http://www.bsi.bund.de/esig/basics/techbas/krypto/index.htm>

and

$$q = 47388090603832016196633832303788951 \\ 973268922921040957944741354648812028493909367$$

The calculations took place between December 2002 and April 2003.

RSA-200

On May 9, 2005 the research group of Prof. Jens Franke at the German University of Bonn⁷³ announced, that they achieved to factorize a 200-digit number into its both prime factors (these are build with each 100 decimal digits) using the GNFS method (General Number Field Sieve).

The composite number called “RSA-200” is (the following 3 lines contain one number):

$$2799783391122132787082946763872260162107044678695542853756000992932 \\ 6128400107609345671052955360856061822351910951365788637105954482006 \\ 576775098580557613579098734950144178863178946295187237869221823983$$

The factorization of RSA-200 resulted in the following two prime factors:

$$p = 35324619344027701212726049781984643686711974001976 \\ 25023649303468776121253679423200058547956528088349$$

and

$$q = 79258699544783330333470858414800596877379758573642 \\ 19960734330341455767872818152135381409304740185467$$

The calculations took place between December 2003 and May 2005. The factorization done by the group around Bahr, Böhm, Franke, Kleinjung, Montgomery and te Riele lasted almost 17 months. The operating expense of the calculations was about 120,000 MIPS-years⁷⁴.

RSA-768

On December 12, 2009 the research group around Prof. Thorsten Kleinjung⁷⁵ announced, that they achieved to factorize a 232-digit number into its both prime factors (both factors have 116 decimal digits). They used the GNFS method (General Number Field Sieve) in a way where they did “oversieving” on several hundred computers before starting the matrix step.

The composite number called “RSA-768” is (the following 3 lines contain one number):

$$123018668453011775513049495838496272077285356959533479219732245215172640050726 \\ 365751874520219978646938995647494277406384592519255732630345373154826850791702 \\ 6122142913461670429214311602221240479274737794080665351419597459856902143413$$

The factorization of RSA-768 resulted in the following two prime factors (each with 384 bit):

$$p = 3347807169895689878604416984821269081770479498371376856891 \\ 2431388982883793878002287614711652531743087737814467999489$$

⁷³<http://www.loria.fr/~zimmerma/records/rsa200>

⁷⁴A MIPS-year (MY) is the quantity of operations a machine can perform in one year, if the machine constantly achieves one million integer operations per second (MIPS). For illustration: a INTEL Pentium 100 processor achieves about 50 MIPS. To factorize a 2048 bit module it is estimated to need about $8.5 \cdot 10^{40}$ MY.

⁷⁵<http://eprint.iacr.org/2010/006.pdf> [Kleinjung2010]

and

$$q = 3674604366679959042824463379962795263227915816434308764267 \\ 6032283815739666511279233373417143396810270092798736308917$$

The calculations took about 2 1/2 years.⁷⁶

C307 / M1039

In May 2007 Prof. Franke, Prof. Kleinjung (University of Bonn), the Japanese telecommunication company NTT and Prof. Arjen Lenstra (Polytechnical University of Lausanne) announced, that they managed to factorize a 307 digit decimal number into its both prime factors with the SNFS method (Special Number Field Sieve) within 11 months (the two factors have 80 and 227 decimal digits).

The task of the researchers was not initiated by a challenge, but they wanted to find the last prime factors of the Mersenne number $2^{1039} + 1$ (see “Wanted List” of the Cunningham Project⁷⁷).

The number $2^{1039} - 1$ consists of 3 prime factors: The smallest one, $p7 = 5080711$ was already known.⁷⁸

To complete this the second factor (co-divider) “C307” had to be factorized: Till then it was only known, that the last remaining factor was composite, but it was unknown, how many prime factors it had and what are the prime factors. The following 5 lines contain one number:

$$\begin{aligned} C307 = & 1159420574072573064369807148876894640753899791702017724986868353538 \\ & 8224838599667566080006095408005179472053993261230204874402860435302 \\ & 8619141014409345351233471273967988850226307575280937916602855510550 \\ & 0425810771176177610094137970787973806187008437777186828680889844712 \\ & 822002935201806074755451541370711023817 \end{aligned}$$

The factorization of C307 resulted in the following two 80- and 2276-digit prime factors:

$$\begin{aligned} p80 = & 558536666199362912607492046583159449686465270184 \\ & 88637648010052346319853288374753 \end{aligned}$$

⁷⁶This was an “academic effort” – organisations with bigger resources could do it much faster.

⁷⁷Cunningham project: <http://www.cerias.purdue.edu/homes/ssw/cun/>

Cunningham table: <http://homes.cerias.purdue.edu/~ssw/cun/pmain1206>

The numbers in the Cunningham table have the following syntax:

“(2,n)-” means $2^n - 1$; “(2,n)+” means $2^n + 1$.

To describe the magnitude one writes $p < n >$ or $c < n >$: “n” is the number of decimal digits and “p” and “c” tell, whether the number is prime or composite.

$2^{1039} - 1 = p7 * c307 = p7 * p80 * p227$

It is explained more precisely at the page of the Cunningham project:

“2651+ means $2^{651} + 1$ and the size (c209 means 209 decimal digits) of the number which was factored. Then come the new factor(s), the discoverer and the method used. Recently, only the multiple polynomial quadratic sieve (ppmpqs), the elliptic curve method (ecm) and the number field sieve (nfs) have been used. ‘hmpqs’ stands for hypercube multiple polynomial quadratic sieve. Under ‘new factors’, ‘p90’ means a 90-digit prime and ‘c201’ is a 201-digit composite number.”.

⁷⁸This one can also be found using CT1 via menu **Indiv. Procedures \ RSA Cryptosystem \ Factorization of a Number** — with the algorithms of Brent, Williams or Lenstra, which are “relatively” good to separate small factors.

and

$$\begin{aligned} p_{227} = & 207581819464423827645704813703594695162939708007395209881208 \\ & 387037927290903246793823431438841448348825340533447691122230 \\ & 281583276965253760914101891052419938993341097116243589620659 \\ & 72167481161749004803659735573409253205425523689. \end{aligned}$$

So now the number $2^{1039} - 1$ is completely factorized in its 3 prime factors.

Links:

[<http://www.loria.fr/~zimmerma/records/21039->
<http://www.crypto-world.com/announcements/m1039.txt>
<http://www.crypto-world.com/FactorAnnouncements.html>
[http://www1.uni-bonn.de/pressDB/jsp/pressemitteilungsdetails.jsp?
detailjahr=2007&detail=160](http://www1.uni-bonn.de/pressDB/jsp/pressemitteilungsdetails.jsp?detailjahr=2007&detail=160)

Size of factorized numbers compared to primality proven numbers

As you notice the factorized compound numbers built of 2 prime factors are much smaller than the especially structured numbers, for which primality tests are able to decide whether these numbers are prime or not (see chapters 3.4, 3.5 and 3.6).

Length of the current world records in decimal notation:

$$\begin{aligned} [\text{RSA-768 number}] & \longleftrightarrow [\text{47th known Mersenne prime}] \\ 232 & \longleftrightarrow 43,112,609 \\ [\text{see table 4.12}] & \longleftrightarrow [\text{see table 3.1}] \end{aligned}$$

4.11.5 Further current research about primes and factorization

Prime numbers are part of very many topical research areas in number theory and computer science. Progress made with factorization is bigger than was estimated 5 years ago – this is not only due to faster computers but also new mathematical knowledge.

The security of the RSA algorithm is based on the empirical observation that factoring large numbers is a hard problem. A module n (typically, 1024 bit) can be easily constructed as the product of two large primes p, q (typically, 500–600 bit each), by calculating $n = pq$. However, it is a hard problem to extract p, q from n . Without knowing p or q , the private key cannot be calculated.

Thus, any progress in efficiency of factorizing large integers will effect the security of the RSA. As a consequence, the underlying primes p, q and, thus, the module n (1024 bit as of today) have to be increased. In case of a quantum leap in factorization, the RSA algorithm might be compromised.

4.11.5.1 Bernstein's paper and its implication on the security of the RSA algorithm

In his paper “Circuits for integer factorization: a proposal” (<http://cr.yp.to/djb.html>, published November 2001, D. J. Bernstein [Bernstein2001] addresses the problem of factorizing large integers. Therefore, his results are of relevance from a RSA point of view. As a main result Bernstein claims that the implementation of the General Number Field Sieve algorithm (GNFS) can be improved to factor, with the same effort as before, integers with three times more digits.

We note that the definition of *effort* is a crucial point: Bernstein claims that effort is the product of time and costs of the machine (including the memory used). The gist of the paper lies in the fact that he can reduce a big part of factorizing to sorting. Using Schimmler’s scheme, sorting can be optimized by massive parallel computing. At the end of section 3 Bernstein explains this effect: The costs of m^2 parallel computers with a constant amount of memory is a constant times m^2 . The costs of a computer with a single processor and memory of size m^2 is also of the order of m^2 , but with a different constant factor. With m^2 processors in parallel, sorting of m^2 numbers (with Schimmler’s scheme) can be achieved in time m , while a m^2 -memory computer needs time of the order of m^2 . Decreasing memory and increasing the number of processors, the computing time can be reduced by a factor $1/m$ without additional effort in terms of total costs. In section 5 it is said that massive parallel computing can also increase efficiency of factorizing using Lenstra’s elliptic-curve-method (a search algorithm has costs that increase in a quadratic square manner instead of cubically).

We note that all results achieved so far are asymptotic results. This means that they only hold in the limit n to infinity. Unfortunately, there is no upper limit for the residual error (i.e. the difference between the real and the asymptotic value) for finite n — a problem which has already been addressed by the author. As a consequence, one cannot conclude whether the costs (in the sense of Bernstein) for factorizing 1024–2048-bit RSA modules can be significantly reduced.

There is no doubt that Bernstein’s approach is innovative. However, the reduction of computing time under constant costs comes along with a massive use of parallel computing — a scenario which seems not to be realistic yet. For example, formally 1 sec computing time on one machine and $1/1,000,000$ sec time parallel computing time on 1,000,000 machines might have same costs. In reality, it is much harder to realize the second situation, and Bernstein does not take into account the fixed costs, in particular for building a network between all these computers.

Although distributed computing over a large network might help to overcome this problem, realistic costs for data transfer have to be taken into account — a point which was not addressed in Bernstein’s proposal.

As long as there is neither (low cost) hardware nor a distributed computing approach (based on Bernstein’s ideas), there should not be a problem for RSA. It has to be clarified from which magnitude of n on Bernstein’s method could lead to a significant improvement (in the sense of the asymptotic result).

Arjen Lenstra, Adi Shamir et. al. analyzed the paper of Bernstein [Lenstra2002]. In summary they expect a factorization improvement on how much longer the bit length of the keys could be with a factor of 1.17 (instead of factor 3 as proposed by Bernstein).

The abstract of their paper “Analysis of Bernstein’s Factorization Circuit” says:

“... Bernstein proposed a circuit-based implementation of the matrix step of the number

field sieve factorization algorithm. We show that under the non-standard cost function used in [1], these circuits indeed offer an asymptotic improvement over other methods but to a lesser degree than previously claimed: for a given cost, the new method can factor integers that are 1.17 times larger (rather than 3.01). We also propose an improved circuit design based on a new mesh routing algorithm, and show that for factorization of 1024-bit integers the matrix step can, under an optimistic assumption about the matrix size, be completed within a day by a device that costs a few thousand dollars. We conclude that from a practical standpoint, the security of RSA relies exclusively on the hardness of the relation collection step of the number field sieve.”

RSA Security⁷⁹ concludes in its analysis of the Bernstein paper [RSA Security 2002] from April, 8 2002 also – as expected – that RSA is still not compromised.

This is still an ongoing discussion.

When this section was written (June 2002) nothing was publicly known about, how far there exist implementations of his theoretical onsets and how much financing there was for his research project.

Links:

- <http://cr.yp.to/djb.html>
- <http://www.counterpane.com/crypto-gram-0203.html#6>
- <http://www.math.uic.edu>

4.11.5.2 The TWIRL device

In January 2003 Adi Shamir and Eran Tromer from the Weizmann Institute of Science published a preliminary draft called “*Factoring Large Numbers with the TWIRL Device*” raising concerns about the security of key sizes till 1024 bits [Shamir2003].

Their abstract summarizes their results very well: “The security of the RSA cryptosystem depends on the difficulty in factoring large integers. The best current factoring algorithm is the Number Field Sieve (NFS), and its most difficult part is the sieving step. In 1999 a large distributed computation involving thousands of workstations working for many months managed to factor a 512-bit RSA key, but 1024-bit keys were believed to be safe for the next 15-20 years. In this paper we describe a new hardware implementation of the NFS sieving step ... which is 3-4 orders of magnitude more cost effective than the best previously published designs Based on a detailed analysis of all the critical components (but without an actual implementation), we believe that the NFS sieving step for 1024-bit RSA keys can be completed in less than a year with a \$10M device, and that the NFS sieving step for 512-bit RSA keys can be completed in less than ten minutes with a \$10K device. Coupled with recent results about the difficulty of the NFS matrix step ... this raises some concerns about the security of these key sizes.”

A detailed explanation from these two authors also can be found in the RSA Laboratories CryptoBytes [Shamir2003a].

The 3-page article in the DuD issue of June 2003 [Weis2003] contains a very good explanation, how the attack using the Generalized Number Field Sieve (GNFS) works and which progress is made, to factorize numbers. At GNFS we can distinguish 2 general steps: The sieve step (relation collecting) and the matrix reduction. Besides the sieve step is highly parallelizable, it dominates the overall calculation burden. Shamir and Tromer haven’t built a TWIRL device yet, but the estimated costs of 10 till 50 million Euro (in order to factorize a 1024-bit

⁷⁹<http://www.rsasecurity.com/>

number) is not prohibitive for secret agencies or big criminal organizations, because the “costs for a single espionage satellite is estimated e.g. to be several billion USD”. The authors therefore recommend, to get as soon as possible rid of today used sensible RSA, Diffie-Hellman or ElGamal keys up to 1024 bit and to use then keys of at least 2048 bit length. The planned TCPA/Palladium hardware will use 2048-bit RSA keys!

So recommendations like the ones from the GISA (German Information Security Agency) to use higher key lengths are very valid.

4.11.5.3 “Primes in P”: Primality testing is polynomial

In August 2002 the three Indian researchers M. Agrawal, N. Kayal and N. Saxena published the paper “*PRIMES in P*” about a new primality testing algorithm called AKS [Agrawal2002]. They discovered a polynomial time deterministic algorithm for determining if a number is prime or not.

The importance of this discovery is that it provides number theorists with new insights and opportunities for further research. Lots of people over centuries have been looking for a polynomial time test for primality, and this result is a major theoretic breakthrough. It shows that new results can be generated from already known facts.

But even its authors note that other known algorithms may be faster (for example ECPP). The new algorithm works on any integer. For example the GIMPS project uses the Lucas-Lehmer primality test which takes advantage of the special properties of Mersenne numbers. This makes the Lucas-Lehmer test much faster, allowing to test numbers with millions of digits while general purpose algorithms are limited to numbers with a few thousand digits.

Current research results on this topic can be found at:

<http://www.mersenne.org/>
<http://fatphil.org/math/aks/> Original paper in English
<http://ls2-www.cs.uni-dortmund.de/lehre/winter200203/kt/material/primes.ps>
Good explanation in German by Thomas Hofmeister.

4.11.5.4 Shared Primesxxxxxxxxxxxxxx

In xxxxxxxxxxxxxxxxxxxxxxxx

It is our choices, that show what we truly are, far more than our abilities.

Quote 11: Joanne K. Rowling⁸⁰

4.12 Applications of asymmetric cryptography using numerical examples

The results of modular arithmetic are used extensively in modern cryptography. Here we will provide a few examples from cryptography using small⁸¹ numbers.

Enciphering a text entails applying a function (mathematical operation) to a character string (number) to generate a different number. Deciphering entails reversing this function, in other words using the distorted image that the function has created from the plaintext in order to restore the original image. For example, the sender could take the plaintext M of a confidential message and add a secret number, the key S , to obtain the ciphertext C :

$$C = M + S.$$

The recipient can reconstruct the plaintext by reversing this operation, in other words by subtracting S :

$$M = C - S.$$

Adding S reliably makes the plaintext impossible to read. However, this encryption is rather weak, because all an interceptor needs to do to calculate the key is obtain a plaintext and the associated ciphertext

$$S = C - M,$$

and can then read any subsequent messages encrypted using S .

The essential reason for this is that subtraction is just as simple an operation as addition.

4.12.1 One way functions

If the key is to be impossible to determine even with knowledge of both the plaintext and the ciphertext, we need a function that is, on the one hand, relatively easy to calculate – we don't want to have problems encrypting messages. On the other hand, the inverse function should exist (otherwise information would be lost during encryption), but should be de facto incalculable.

What are possible candidates for such a **one way function**? We could take multiplication rather than addition, but even primary school children know that the inverse function, division, is only slightly more difficult than multiplication itself. We need to go one step higher in the hierarchy of calculation methods. It is still relatively simple to calculate the power of a number, but the corresponding two reverse functions – *taking roots* (find b in the equation $a = b^c$ when a and c are known) and *calculating logarithms* (find c in the above equation when a and b are known) are so complicated that pupils normally do not learn them at school.

Although a certain structure can still be recognised for addition and multiplication, raising numbers to the power of another or calculating exponentials totally mixes up all the numbers.

⁸⁰ Joanne K. Rowling, “Harry Potter and the Chamber of Secrets”, Bloomsbury, 1998, last chapter “Dobby’s reward”, p. 245, by Dumbledore.

⁸¹ In the RSA procedure, we call numbers “small” if the bit lengths are much less than 1024 bits (i.e. 308 decimal points). In practice, 1024 bits is currently considered the minimum length for a secure RSA modul.

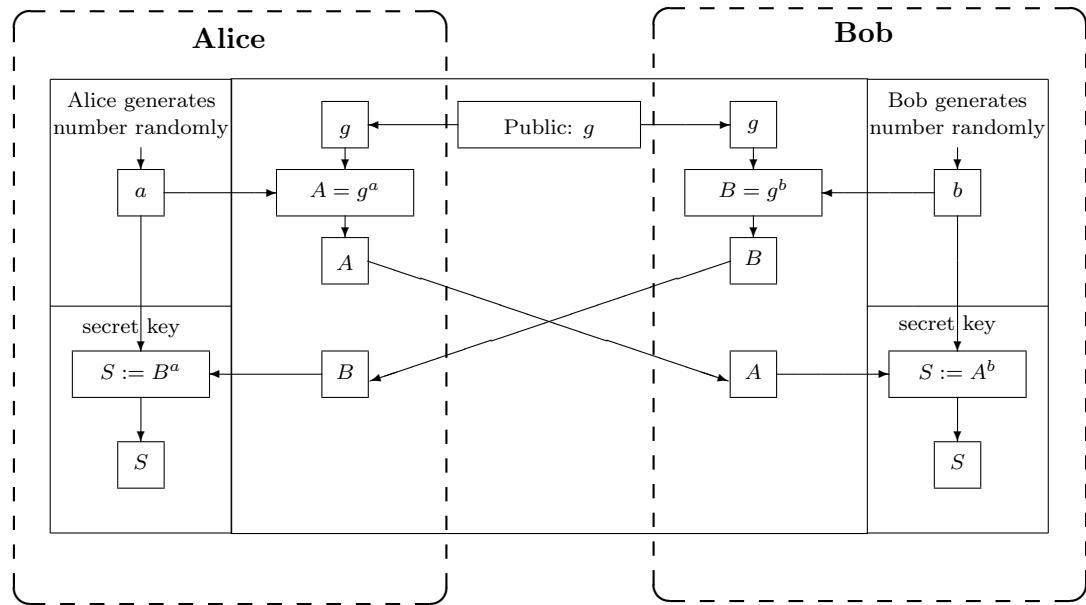
Knowing a few values of the function doesn't tell us much about the function as a whole (in contrast to addition and multiplication).

4.12.2 The Diffie-Hellman key exchange protocol

Whitfield Diffie, Martin E. Hellman and Ralph Merkle developed this DH key exchange protocol in Stanford in 1976.⁸²

Alice and Bob⁸³ use a one way function to obtain a key S , the session key, for subsequent correspondence. This is then a secret that is only known to the two of them. Alice selects a random number a and keeps it secret. She applies a one way function to a to calculate the number $A = g^a$ and sends it to Bob. He does the same, by selecting a secret random number b , calculating $B = g^b$ and sending it to Alice. The number g is random and can be publicly known. Alice applies the one way function together with her secret number a to B , while Bob does the same with his secret number b and the received number A .

The result S is the same in each case because the one way function is commutative: $(g^a)^b = (g^b)^a$. But even Bob cannot reconstruct Alice's secret number a from the data available to him, while Alice cannot determine Bob's secret number b . And a perpetrator who knows g and has intercepted both A and B cannot use this knowledge to determine a, b or S .



Procedure:

Alice and Bob want to negotiate a secret session key S via a channel that may be intercepted.

1. They select a prime number p and a random number g and exchange this information openly.

⁸²With CT1 this exchange protocol has been visualized: You can execute the single steps with concrete numbers using menu **Indiv. Procedures \ Protocols \ Diffie-Hellman Demonstration**.

In JCT you can find it in the default perspective via the menu item **Visuals \ Diffie-Hellman Key Exchange (EC)**.

⁸³Bob and Alice are the default names used for the two authorized participants in a protocol (see [Schneier1996, p. 23]).

2. Alice now selects a , a random number less than p and keeps it secret.
Similarly, Bob selects b , a random number less than p and keeps it secret.
3. Alice now calculates $A \equiv g^a \pmod{p}$.
Bob calculates $B \equiv g^b \pmod{p}$.
4. Alice sends the result A to Bob.
Bob sends the result B to Alice.
5. In order to now determine the session key to be used by both, they both separately raise the respective results they have received to the power of their secret random number modulo p . This means:
 - Alice calculates $S \equiv B^a \pmod{p}$ and
 - Bob calculates $S \equiv A^b \pmod{p}$.

Even if a spy intercepts g, p , and the interim results A and B , he cannot use these in order to determine the used session key used – due to the difficulty of calculating the discrete logarithm⁸⁴.

We will now use an example with (unrealistically) small numbers to illustrate this.

Example using small numbers:

1. Alice and Bob select $g = 11, p = 347$.
2. Alice selects $a = 240$, Bob selects $b = 39$ and they keep a and b secret.
3. Alice calculates $A \equiv g^a \equiv 11^{240} \equiv 49 \pmod{347}$.
Bob calculates $B \equiv g^b \equiv 11^{39} \equiv 285 \pmod{347}$.
4. Alice sends Bob: $A \equiv 49$,
Bob sends Alice: $B \equiv 285$.
5. Alice calculates $B^a \equiv 285^{240} \equiv 268 \pmod{347}$,
Bob calculates $A^b \equiv 49^{39} \equiv 268 \pmod{347}$.

Alice and Bob can now communicate securely using their shared session key. Even if spies were to intercept everything transferred via the connection: $g = 11, p = 347, A = 49$ and $B = 285$, they would not be able to calculate the secret key.

⁸⁴Further details about the discrete logarithm problem can be found in chapter 5.4.

Comment:

In this example using such small numbers, it is easily possible to calculate the discrete logarithms, but with large numbers the discrete logarithm problem^{85,86} is extremely difficult to solve.

To get the discrete logarithms, here we need to calculate:

For Alice: $11^x \equiv 49 \pmod{347}$, that means $\log_{11}(49) \pmod{347}$.

For Bob: $11^y \equiv 285 \pmod{347}$, that means $\log_{11}(285) \pmod{347}$.

⁸⁵You can use Sage to determine the discrete logarithm x that solves the equation $11^x \equiv 49 \pmod{347}$ (here for Alice): `discrete_log(mod(49, 347), mod(11, 347))`. The returned value is 67.

Such number theoretic tasks can also be solved using other tools like PariGP, LiDIA, BC or Mathematica (see the list of web sites in the appendix at the end of this chapter):

- Pari-GP: `znlog(Mod(49,347),Mod(11,347))`.
- LiDIA: `d1(11,49,347)`.
- Mathematica: The general `Solve` function delivers the em tdep message “The equations appear to involve the variables to be solved for in an essentially non-algebraic way”.
- Mathematica: `MultiplicativeOrder[11, 347, 49]`.

All deliver the result 67.

⁸⁶Why have the functions delivered the value 67 for the discrete logarithm of Alice rather than 240 which Alice selected as exponent a ?

The discrete logarithm is the smallest natural exponent that solves the equation $11^x \equiv 49 \pmod{347}$. Both $x = 67$ and $x = 240$ (the number selected in the example) satisfy the equation and can therefore be used to calculate the session key: $285^{240} \equiv 285^{67} \equiv 268 \pmod{347}$. If Alice and Bob had selected a primitive root modulo p as base g , then for every remainder from the set $\{1, 2, \dots, p-1\}$ there is exactly one exponent from the set $\{0, 1, \dots, p-2\}$.

As an aside, there are 172 different primitive roots modulo 347, 32 of which are prime (not necessary). Since the number 11 selected for g in the example is not a primitive root of 347, the remainders do not take all values from the set $\{1, 2, \dots, 346\}$. Thus, for a particular remainder there may be more than one exponent or even no exponent at all in the set $\{0, 1, \dots, 345\}$ that satisfies the equation.

With the relevant Sage commands you find:

`is_prime(347)=True, euler_phi(347)=346, gcd(11,347)=1 and multiplicative_order(mod(11, 347))=173.`

i	$11^i \pmod{347}$	
0	1	
1	11	
2	121	
3	290	
67	49	searched exponent
172	284	
173	1	= multiplicative order of $11^i \pmod{347}$
174	11	
175	121	
176	290	
240	49	searched exponent

Further information can be found in chapter 4.19.4 “Primitive roots”.

4.13 The RSA procedure with actual numbers⁸⁷

“Games are Nature’s way of preparing us to face difficult realities. Are you finally ready to face reality, Sergeant?”

Quote 12: Daniel Suarez⁸⁸

Having described above how the RSA procedure works, we will now work through the steps using actual, but small, numbers.

4.13.1 RSA with small prime numbers and with a number as message

Before applying the RSA procedure to a text, we will first demonstrate it directly using a single number as message.⁸⁹

1. Let the selected prime numbers be $p = 5$ and $q = 11$.
Thus, $n = 55$ and $J(n) = (p - 1) * (q - 1) = 40$.
2. $e = 7$ (e should⁹⁰ lie between 11 and 39, and must be relatively prime to 40).
3. $d = 23$ (since $23 * 7 \equiv 161 \equiv 1 \pmod{40}$),
→ Public key of the recipient: $(55, 7)$,
→ Private key of the recipient: $(55, 23)$.
4. Let the message be the number $M = 2$ (so no division into blocks is required).
5. Encryption: $C \equiv 2^7 \equiv 18 \pmod{55}$.
6. The ciphertext is simply the number $C = 18$ (we therefore do not need to divide it into blocks).
7. Decryption: $M \equiv 18^{23} \equiv 18^{(1+2+4+16)} \equiv 18 * 49 * 36 * 26 \equiv 2 \pmod{55}$.

We will now apply the RSA procedure to a text, first using the upper case alphabet (26 characters), then using the entire ASCII character set as the basis for the messages.

4.13.2 RSA with slightly larger primes and a text of upper case letters

We have the text “ATTACK AT DAWN”, and the characters are coded according to table 4.13.⁹¹

⁸⁷Additional material: Minh Van Nguyen, “Number Theory and the RSA Public Key Cryptosystem”, 2009. An introductory tutorial on using Sage to study elementary number theory and public key cryptography. A didactically very clear article about some basic number theory and Sage usage.

http://nguyenminh2.googlepages.com/sage_numtheory-rsa.pdf.

⁸⁸Daniel Suarez, “Daemon”, Dutton Adult, 2010, Chapter 45, “Respawning”, p. 610, Sobol.

⁸⁹Using CT1 you can solve this with the menu **Indiv. Procedures \ RSA Cryptosystem \ RSA Demonstration**.

⁹⁰See footnote 53 on page 129.

⁹¹Using CT1 you can solve this with the menu **Indiv. Procedures \ RSA Cryptosystem \ RSA Demonstration**. This is also described in the tutorial/scenario in CT1’s online help [Options, specify alphabet, number system, block length 2 and decimal representation].

Character	Numerical value	Character	Numerical value
Blank	0	M	13
A	1	N	14
B	2	O	15
C	3	P	16
D	4	Q	17
E	5	R	18
F	6	S	19
G	7	T	20
H	8	U	21
I	9	V	22
J	10	W	23
K	11	X	24
L	12	Y	25
		Z	26

Table 4.13: Capital letters alphabet

Key generation (steps 1 to 3):

1. $p = 47, q = 79$ ($n = 3713; J(n) = (p - 1) * (q - 1) = 3588$).
2. $e = 37$ (e should⁹² lie between 79 and 3587, and must be relatively prime to 3588).
3. $d = 97$ (since $e * d = 1 \bmod J(n); 37 * 97 \equiv 3589 \equiv 1 \pmod{3588}$).⁹³

4. Encryption:

Text:	A	T	T	A	C	K		A	T		D	A	W	N
Number:	01	20	20	01	03	11	00	01	20	00	04	01	23	14

This 28-digit number is divided into 4-digit parts (because 2626 is still smaller than $n = 3713$):
 0120 2001 0311 0001 2000 0401 2314

All 7 parts are encrypted using: $C \equiv M^{37} \pmod{3713}$:⁹⁴

1404 2932 3536 0001 3284 2280 2235

5. Decryption:

Ciphertext: 1404 2932 3536 0001 3284 2280 2235

This 28-digit number is divided into 4-digit parts.

All 7 parts are decrypted using: $M \equiv C^{97} \pmod{3713}$:

0120 2001 0311 0001 2000 0401 2314

The 2-digit numbers are transformed into capital letters and blanks.

Using the selected values it is easy for a cryptanalyst to derive the secret values from the public parameters $n = 3713$ and $e = 37$ by revealing that $3713 = 47 * 79$.

If n is a 768-bit number, there is, according to present knowledge, little chance of this.

⁹²See footnote 53 on page 129.

⁹³How to compute $d = 97$ using the *extended gcd* algorithm is shown in appendix 4.14

⁹⁴See chapter 4.19.5 “RSA examples with Sage” for source code to do RSA encryption using Sage.

You can also encrypt the message with CT1 via the menu path **Indiv. Procedures \ RSA Cryptosystem \ RSA Demonstration**.

4.13.3 RSA with even larger primes and a text made up of ASCII characters

In real life, the ASCII alphabet is used to code the individual characters of the message as 8-bit numbers.

The idea for this task⁹⁵ is taken from the example in [Eckert2003, p. 271].

Coded in decimal notation, the text “RSA works!” is as follows:

Text:	R	S	A	w	o	r	k	s	!	
Number:	82	83	65	32	119	111	114	107	115	33

We will work through the example in 2 variants. The steps 1 to 3 are common for both.

Key generation (steps 1 to 3):

1. $p = 503$, $q = 509$ ($n = 256,027$; $J(n) = (p - 1)(q - 1) = 255,016 = 2^3 * 127 * 251$).⁹⁶
2. $e = 65,537$
(e should⁹⁷ lie between 509 and 255,015, and must⁹⁸ be relatively prime to 255,016).
3. $d = 231,953$
(since $e \equiv d^{-1} \pmod{J(n)}$: $65,537 * 231,953 \equiv 15,201,503,761 \equiv 1 \pmod{255,016}$).⁹⁹

Variant 1: All ASCII characters are en-/decrypted separately (no blocks are formed).

4. Encryption:

Text:	R	S	A	w	o	r	k	s	!	
Number:	82	83	65	32	119	111	114	107	115	33

The letters are not combined!¹⁰⁰

Each character is encrypted using: $C = M^{65,537} \pmod{256,027}$:¹⁰¹

212984	025546	104529	031692	248407
100412	054196	100184	058179	227433

5. Decryption:

Ciphertext:

212984	025546	104529	031692	248407
100412	054196	100184	058179	227433

Each character is decrypted using: $M \equiv C^{231,953} \pmod{256,027}$:

82	83	65	32	119	111	114	107	115	33
----	----	----	----	-----	-----	-----	-----	-----	----

⁹⁵Using CT1 you can solve this via the menu path **Indiv. Procedures \ RSA Cryptosystem \ RSA Demonstration**.

⁹⁶See chapter 4.19.5 “RSA examples with Sage” for the source code to factorize the number $J(n)$ using Sage. Using CT1 you can solve this with the **Indiv. Procedures \ RSA Cryptosystem \ Factorization of a Number**.

⁹⁷See footnote 53 on page 129.

⁹⁸ e cannot, therefore, be 2, 127 or 251 ($65,537 = 2^{16} + 1$) ($255,016 = 2^3 * 127 * 251$).

In real life, $J(n)$ is not factorized but rather the Euclidean algorithm is used for the selected e to guarantee that $\gcd(e, J(n)) = 1$.

⁹⁹Other possible combinations of (e, d) include: (3, 170, 011), (5, 204, 013), (7, 36, 431).

¹⁰⁰For secure procedures we need large numbers that assume – as far as possible – all values up to $n - 1$. If the possible value set for the numbers in the message is too small, even large prime numbers cannot make the procedure secure. An ASCII character is represented by 8 bits. If we want larger values we must combine several numbers. Two characters need 16 bits, whereby the maximum value that can be represented is 65536. The modulus n must then be greater than $2^{16} = 65536$. This is applied in variant 2. When the numbers are combined, the leading zeros are kept in binary notation (just as if we were to write all numbers with 3 digits in decimal notation above and were then to obtain the sequence 082 083, 065 032, 119 111, 114 107, 115 033).

¹⁰¹See chapter 4.19.5 “RSA examples with Sage” for the source code for RSA exponentiation using Sage.

Variant 2: The ASCII characters are en-/decrypted two at a time as blocks.

In variant 2 the block formation is done in two different sub-variants: (4./5. and 4'./5').).

Text:	R	S	A	w	o	r	k	s	!	
Number:	82	83	65	32	119	111	114	107	115	33

4. Encryption:

Blocks are formed¹⁰² (each ASCII character is encoded into a 8 digit binary number below):

21075 16672 30575 29291 29473¹⁰³

Each block is encrypted using: $C \equiv M^{65,537} \pmod{256,027}$:¹⁰⁴

158721 137346 37358 240130 112898

5. Decryption:

Ciphertext:

158721 137346 37358 240130 112898

Each block is decrypted using: $M \equiv C^{231,953} \pmod{256,027}$:

21075 16672 30575 29291 29473

4'. Encryption:

Blocks are formed: (each ASCII character is encoded into a 3 digit decimal number below):

82083 65032 119111 114107 115033¹⁰⁵

Each block is encrypted using: $C \equiv M^{65,537} \pmod{256,027}$:¹⁰⁶

198967 051405 254571 115318 014251

5'. Decryption:

Ciphertext:

198967 051405 254571 115318 014251

Each block is decrypted using: $M \equiv C^{2473} \pmod{67,519}$:

82083 65032 119111 114107 115033

¹⁰²

single character	binary representation	decimal representation
01010010, 82	01010010 01010011	= 21075
01010011, 83		
01000001, 65	01000001 00100000	= 16672
00100000, 32		
01110111, 119	01110111 01101111	= 30575
01101111, 111		
01110010, 114	01110010 01101011	= 29291
01101011, 107		
01110011, 115	01110011 00100001	= 29473
00100001, 33:		

¹⁰³Using CT1 you can solve this with the menu **Indiv. Procedures \ RSA Cryptosystem \ RSA Demonstration** with the following options: all 256 ASCII characters, b-adic, block length 2 and decimal representation.

¹⁰⁴See chapter 4.19.5 “RSA examples with Sage” for the source code for RSA exponentiation using Sage.

¹⁰⁵The RSA encryption works correctly with the modulus $n = 256,027$ because each ASCII block of two characters will be encoded into a number that is smaller or equal than the number 255, 255.

¹⁰⁶See chapter 4.19.5 “RSA examples with Sage” for the source code for RSA exponentiation using Sage.

4.13.4 A small RSA cipher challenge (1)

The task is taken from [Stinson1995, Exercise 4.6]: The pure solution has been published by Prof. Stinson.¹⁰⁷ However, it is not the result that is important here but rather the individual steps of the solution, that is, the explanation of the cryptanalysis.¹⁰⁸

Two samples of RSA ciphertext are presented in Tables 4.14¹⁰⁹ and 4.15¹¹⁰. Your task is to decrypt them. The public parameters of the system are

$n = 18,923$ and $e = 1261$ (for Table 4.14) and
 $n = 31,313$ and $e = 4913$ (for Table 4.15).

This can be accomplished as follows. First, factor n (which is easy because it is so small). Then compute the exponent d from $J(n)$, and, finally, decrypt the ciphertext. Use the square-and-multiply algorithm to exponentiate modulo n .

In order to translate the plaintext back into ordinary English text, you need to know how alphabetic characters are “encoded” as elements in \mathbb{Z}_n . Each element of \mathbb{Z}_n represents three alphabetic characters as in the following examples:

$$\begin{aligned} \text{DOG} &\mapsto 3 * 26^2 + 14 * 26 + 6 = 2398 \\ \text{CAT} &\mapsto 2 * 26^2 + 0 * 26 + 19 = 1371 \\ \text{ZZZ} &\mapsto 25 * 26^2 + 25 * 26 + 25 = 17,575. \end{aligned}$$

You will have to invert this process as the final step in your program.

The first plaintext was taken from “The Diary of Samuel Marchbanks”, by Robertson Davies, 1947, and the second was taken from “Lake Wobegon Days”, by Garrison Keillor, 1985.

¹⁰⁷ <http://www.cacr.math.uwaterloo.ca/~dstinson/solns.html> or <http://bibd.unl/~stinson/solns.html>

¹⁰⁸ The method of solving the problem is outlined in the scenario of the online help to CT1 and in the presentation on the CT website. If anyone sends us a well prepared exact method of solving the problem, we would be pleased to include it in the documentation.

¹⁰⁹ The numbers of this table can be worked with via Copy and Paste.

¹¹⁰ The numbers of this table are in the online help of CT1 in the chapter “Example illustrating the RSA demonstration”.

12423	11524	7243	7459	14303	6127	10964	16399
9792	13629	14407	18817	18830	13556	3159	16647
5300	13951	81	8986	8007	13167	10022	17213
2264	961	17459	4101	2999	14569	17183	15827
12693	9553	18194	3830	2664	13998	12501	18873
12161	13071	16900	7233	8270	17086	9792	14266
13236	5300	13951	8850	12129	6091	18110	3332
15061	12347	7817	7946	11675	13924	13892	18031
2620	6276	8500	201	8850	11178	16477	10161
3533	13842	7537	12259	18110	44	2364	15570
3460	9886	8687	4481	11231	7547	11383	17910
12867	13203	5102	4742	5053	15407	2976	9330
12192	56	2471	15334	841	13995	17592	13297
2430	9741	11675	424	6686	738	13874	8168
7913	6246	14301	1144	9056	15967	7328	13203
796	195	9872	16979	15404	14130	9105	2001
9792	14251	1498	11296	1105	4502	16979	1105
56	4118	11302	5988	3363	15827	6928	4191
4277	10617	874	13211	11821	3090	18110	44
2364	15570	3460	9886	9988	3798	1158	9872
16979	15404	6127	9872	3652	14838	7437	2540
1367	2512	14407	5053	1521	297	10935	17137
2186	9433	13293	7555	13618	13000	6490	5310
18676	4782	11374	446	4165	11634	3846	14611
2364	6789	11634	4493	4063	4576	17955	7965
11748	14616	11453	17666	925	56	4118	18031
9522	14838	7437	3880	11476	8305	5102	2999
18628	14326	9175	9061	650	18110	8720	15404
2951	722	15334	841	15610	2443	11056	2186

Table 4.14: RSA ciphertext A

6340	8309	14010	8936	27358	25023	16481	25809
23614	7135	24996	30590	27570	26486	30388	9395
27584	14999	4517	12146	29421	26439	1606	17881
25774	7647	23901	7372	25774	18436	12056	13547
7908	8635	2149	1908	22076	7372	8686	1304
4082	11803	5314	107	7359	22470	7372	22827
15698	30317	4685	14696	30388	8671	29956	15705
1417	26905	25809	28347	26277	7897	20240	21519
12437	1108	27106	18743	24144	10685	25234	30155
23005	8267	9917	7994	9694	2149	10042	27705
15930	29748	8635	23645	11738	24591	20240	27212
27486	9741	2149	29329	2149	5501	14015	30155
18154	22319	27705	20321	23254	13624	3249	5443
2149	16975	16087	14600	27705	19386	7325	26277
19554	23614	7553	4734	8091	23973	14015	107
3183	17347	25234	4595	21498	6360	19837	8463
6000	31280	29413	2066	369	23204	8425	7792
25973	4477	30989					

Table 4.15: RSA ciphertext B

4.13.5 A small RSA cipher challenge (2)

The following task is a corrected version from the book written by Prof. Yan [Yan2000, Example 3.3.7, p. 318]. However, it is not the result that is important here but rather the individual steps of the solution, that is, the explanation of the cryptanalysis.¹¹¹

There are three tasks with completely different degrees of difficulty here. In each case we know the ciphertext and the public key (e, n) :

- (a) Known plaintext: find the secret key d using the additionally known original message.
- (b) Ciphertext-only: find d and the plaintext.
- (c) Calculate the RSA modulus, in other words factorization (with no knowledge of the message).

$$n = 63978486879527143858831415041, e = 17579$$

Message¹¹²:

1401202118011200,
1421130205181900,
0118050013010405,
0002250007150400

Cipher:

45411667895024938209259253423,
16597091621432020076311552201,
46468979279750354732637631044,
32870167545903741339819671379

Comment:

The original message consisted of a sentence containing 31 characters (coded with the capital letters alphabet from section 4.13.2). Each group of 16 decimal numbers is then combined to form one number (the last number is filled with zeros). These numbers are raised to the power of e .

When you decrypt the message you must fill the calculated numbers with leading zeros in order to obtain plaintext.

This needs to be stressed because the type of padding is extremely important during implementation and standardization for interoperable algorithms.

¹¹¹The method of solving the problem is outlined in the scenario of the online help to CT1 and in the CrypTool presentation. If anyone sends us a well prepared exact method of solving the problem, we would be pleased to include it in the documentation.

¹¹²The numbers of this table are in the online help of CT1 in the chapter “Example illustrating the RSA demonstration”.

4.14 Appendix: The greatest common divisor (gcd) of whole numbers and the two algorithms of Euclid¹¹³

The greatest common divisor of two natural numbers a and b is an important value that can be calculated very quickly. Here we make use of the fact that if a number c divides the numbers a and b (i.e. there exists an a' and a b' such that $a = a' * c$ and $b = b' * c$), then c also divides the remainder r of a/b . In short notion we can write: If c divides a and b it follows that c divides $r = a - \lfloor a/b \rfloor * b$.¹¹⁴

As the latter statement is valid for each common divisor c of a and b it follows that:

$$\text{gcd}(a, b) = \text{gcd}(a - \lfloor a/b \rfloor * b, b).$$

Using this information, the algorithm for calculating the gcd of two numbers can be written as follows (in pseudo code):

```

INPUT: a,b != 0
1. if ( a < b ) then x = a; a = b; b = x; // Swap a and b (a > b)
2. a = a - int(a/b) * b           // a is smaller than b, the
                                // gcd(a, b) is unchanged
3. if ( a != 0 ) then goto 1.      // a falls after each step and
                                // the algorithm ends when a==0.
OUTPUT "gcd(a,b) = " b      // b is the gcd of the original a and b

```

Also further relationships can be derived from the gcd: For this, we need the set of equations for a and b :

$$\begin{aligned} a &= 1 * a + 0 * b \\ b &= 0 * a + 1 * b, \end{aligned}$$

or, in matrix notation:

$$\begin{pmatrix} a \\ b \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} * \begin{pmatrix} a \\ b \end{pmatrix}.$$

We summarize this information in the extended matrix:

$$\left(\begin{array}{c|cc} a & 1 & 0 \\ b & 0 & 1 \end{array} \right)$$

If we apply the above gcd algorithm to this matrix, we obtain the *extended Euclid algorithm* which can be used to calculate the multiplicative inverse:

¹¹³With the educational tool for number theory **NT** you can see

a) how Euklid's algorithm calculates the gcd (learning unit 1.3, pages 14-19/21) and
 b) how Euklid's enhanced algorithm finds the multiplicative inverse (learning unit 2.2, page 13/40).
NT can be called in CT1 via the menu path **Indiv. Procedures \ Number Theory Interactive \ Learning Tool for Number Theory**. See appendix A.6.

CT2 contains Euklid's extended algorithm within the tutorial “**World of Primes** – Number Theory – Number-theoretic functions: Extended Euklid” bxxxxxxxxxxxxxxxxxxxxxx .

¹¹⁴The Gauss bracket $\lfloor x \rfloor$ of a real number x is defined via: $\lfloor x \rfloor$ is the next integer less or equal x .
 See footnote 117 on page 160.

INPUT: $a, b \neq 0$

0. $x_{1,1} := 1, x_{1,2} := 0, x_{2,1} := 0, x_{2,2} := 1$

$$1. \left(\begin{array}{c|cc} a & x_{1,1} & x_{1,2} \\ b & x_{2,1} & x_{2,2} \end{array} \right) := \left(\begin{array}{cc} 0 & 1 \\ 1 & -\lfloor a/b \rfloor * b \end{array} \right) * \left(\begin{array}{c|cc} a & x_{1,1} & x_{1,2} \\ b & x_{2,1} & x_{2,2} \end{array} \right).$$

2. **if** ($b \neq 0$) **then goto** 1.

OUTPUT: “ $\text{gcd}(a, b) = a * x + b * y$:”, “ $\text{gcd}(a, b) =$ ” b , “ $x =$ ” $x_{2,1}$, “ $y =$ ” $x_{2,2}$

Since this algorithm only performs linear transformations, the same equations always apply

$$\begin{aligned} a &= x_{1,1} * a + x_{1,2} * b \\ b &= x_{2,1} * a + x_{2,2} * b, \end{aligned}$$

We get the extended gcd equation at the end of the algorithm¹¹⁵:

$$\text{gcd}(a, b) = a * x_{2,1} + b * x_{2,2}.$$

Example:

Using the extended gcd we can determine for $e = 37$ the multiplicative inverse number d to modulo 3588 (i.e. $37 * d \equiv 1 \pmod{3588}$):

$$\begin{aligned} 0. \left(\begin{array}{c|cc} 3588 & 1 & 0 \\ 37 & 0 & 1 \end{array} \right) \\ 1. \left(\begin{array}{c|cc} 37 & 1 & 0 \\ 36 & 0 & -96 \end{array} \right) = \left(\begin{array}{cc} 0 & 1 \\ 1 & -(\lfloor 3588/36 \rfloor = 96) * 37 \end{array} \right) * \left(\begin{array}{c|cc} 3588 & 1 & 0 \\ 37 & 0 & 1 \end{array} \right). \\ 2. \left(\begin{array}{c|cc} 36 & 1 & -96 \\ 1 & -1 & 97 \end{array} \right) = \left(\begin{array}{cc} 0 & 1 \\ 1 & -(\lfloor 37/36 \rfloor = 1) * 36 \end{array} \right) * \left(\begin{array}{c|cc} 37 & 1 & 0 \\ 36 & 0 & -96 \end{array} \right). \\ 3. \left(\begin{array}{c|cc} 1 & -1 & 97 \\ 0 & 37 & -3588 \end{array} \right) = \left(\begin{array}{cc} 0 & 1 \\ 1 & -(\lfloor 36/1 \rfloor = 36) * 1 \end{array} \right) * \left(\begin{array}{c|cc} 36 & 1 & -96 \\ 1 & -1 & 97 \end{array} \right). \end{aligned}$$

OUTPUT:

$$\text{gcd}(37, 3588) = a * x + b * y;$$

$$\text{gcd}(37, 3588) = 1, x = -1, y = 97.$$

Thus

1. 37 and 3588 are relatively prime (37 has an inverse modulo 3588).
2. $37 * 97 = (1 * 3588) + 1$ in other words $37 * 97 \equiv 1 \pmod{3588}$.
and therefore the number 97 is the multiplicative inverse to 37 modulo 3588.

¹¹⁵By termination of the gcd algorithm, the program variables a and b contain the values $a = 0$ and $b = \text{gcd}(a, b)$. Please keep in mind, that the program variables are different to the numbers a and b and that they are only relevant for the scope of the algorithm.

4.15 Appendix: Forming closed sets

The property of closeness within a set is always defined in relation to an operation. The following shows how to construct the “closed set” G with respect to the operation $+$ (mod 8) for a given initial set G_0 :

$G_0 = \{2, 3\}$ —— addition of the numbers in G_0 determines further numbers :

$$2 + 3 \equiv 5 \pmod{8} = 5$$

$$2 + 2 \equiv 4 \pmod{8} = 4$$

$$3 + 3 \equiv 6 \pmod{8} = 6$$

$G_1 = \{2, 3, 4, 5, 6\}$ —— addition of the numbers in G_1 determines :

$$3 + 4 \equiv 7 \pmod{8} = 7$$

$$3 + 5 \equiv 8 \pmod{8} = 0$$

$$3 + 6 \equiv 9 \pmod{8} = 1$$

$G_2 = \{0, 1, 2, 3, 4, 5, 6, 7\}$ —— addition of the numbers in G_2 does not extend the set!

$G_3 = G_2$ —— we say : G_2 is closed for addition $\pmod{8}$.

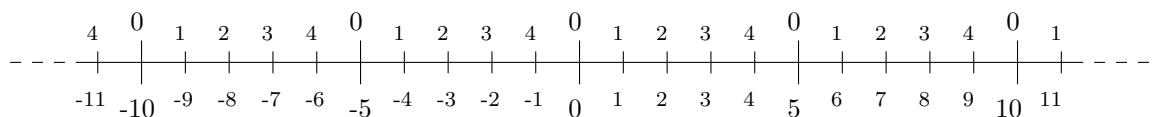
4.16 Appendix: Comments on modulo subtraction

Comment on subtraction modulo 5: $2 - 4 = -2 \equiv 3 \pmod{5}$.

It is therefore not true that $-2 = 2 \pmod{5}$!

People often make the mistake of equating this. You can show this clearly if you place the permutation $(0, 1, 2, 3, 4)$ in \mathbb{Z}_5 , for example from -11 to $+11$, over the range of numbers in \mathbb{Z} .

range of numbers modulo 5



4.17 Appendix: Base representation and base transformation of numbers, estimation of length of digits

For a given number z one may ask how to represent such a number. In general we use representations like $z = 2374$ or $z = \sqrt{2}$. The second number consists of an infinite number of digits and therefore it can never be described precisely by the first representation. You can get around this problem by writing the number symbolically. But if you have to write it in digits, the number must be rounded.

We represent numbers usually in the decimal system (base 10). Computers are working with the binary representation of numbers — only for the display numbers are represented in decimal or sometimes hexadecimal (base 16) form.

This appendix describes how to generate arbitrary base representations of any positive integer and how to determine the number of required digits via the logarithm function.

b -adic sum representation of positive integers

Given base b , each positive integer z can be represented as a b -adic sum

$$z = a_n b^n + a_{n-1} b^{n-1} + \cdots + a_1 b + a_0,$$

where $a_i \in \{0, 1, \dots, b-1\}$, $i = 0, 1, \dots, n$ are called *digits*.

For this sum, it follows that:

- 1) For arbitrary digits a_0, a_1, \dots, a_n it is: $b^{n+1} > a_n b^n + a_{n-1} b^{n-1} + \cdots + a_1 b + a_0$.
- 2) There exist digits a_0, a_1, \dots, a_n (namely $a_i = b-1$ for $i = 0, \dots, n$), following that $b^{n+1} - 1 \leq a_n b^n + a_{n-1} b^{n-1} + \cdots + a_1 b + a_0$.

(Using these inequalities it can be shown that each positive integer can be represented by a b -adic sum).

By writing the digits $a_n a_{n-1} \cdots a_1 a_0$ in a row directly after each other (without the b^i) the usual writing for numbers comes to hand.

Example:

Base $b = 10$: $10278 = 1 \cdot 10^4 + 0 \cdot 10^3 + 2 \cdot 10^2 + 7 \cdot 10^1 + 8$

Base $b = 16$: $FE70A = 15 \cdot 16^4 + 14 \cdot 16^3 + 7 \cdot 16^2 + 0 \cdot 16^1 + 10$.

Number of digits to represent a positive integer

For a positive integer z the length of the b -adic representation can be determined via the following steps. Starting from the inequality $b^{n+1} > z \geq b^n$ we have — after applying the logarithm function on basis b ¹¹⁶: $n+1 > \log_b z \geq n$. Therefore we have $n = \lfloor \log_b z \rfloor$.¹¹⁷ We call $l_b(z)$ the number of required digits to represent the number z on the base b . We have

$$l_b(z) := \lfloor \log_b z \rfloor + 1.$$

¹¹⁶Applying the logarithm formula on base b and b' we have $\log_b z = \log_{b'} z / \log_{b'}(b)$. It is therefore easy using e.g. logarithm tables for the base $b' = 10$ to compute the logarithm of base $b = 2$.

¹¹⁷The function $\lfloor x \rfloor$ determines the next integer smaller than x (in case $x \geq 0$ the digits after the decimal point are truncated).

See footnote 114 on page 157.

Example 1 (decimal→hex):

We compute for the decimal number $z = 234$ (EA in hex) the hexadecimal representation (number base $b = 16$)

$$l_{16}(z) = \lfloor \log_{16}(z) \rfloor + 1 = \lfloor \ln(z)/\ln(16) \rfloor + 1 = \lfloor 1.96... \rfloor + 1 = 1 + 1 = 2.$$

Example 2 (decimal→binary):

We compute for the decimal number $z = 234$ (11101010 in binary) the binary representation (number base $b = 2$)

$$l_2(z) = \lfloor \log_2(z) \rfloor + 1 = \lfloor \ln(z)/\ln(2) \rfloor + 1 = \lfloor 7.87... \rfloor + 1 = 7 + 1 = 8.$$

Example 3 (binary→decimal):

We compute for the decimal number $z = 11101010$ (234 decimal) the decimal representation (number base $b = 10$)

$$l_{10}(z) = \lfloor \log_{10}(z) \rfloor + 1 = \lfloor \ln(z)/\ln(10) \rfloor + 1 = \lfloor 2.36... \rfloor + 1 = 2 + 1 = 3.$$

Algorithm to compute the base representation

Given the number z one can compute the base b representation of z using the following algorithm:

```
input:  $z, b$ 
 $n := 0, z' := z$ 
while  $z' > 0$  do
     $a_n := z' \bmod b,$ 
     $z' := \lfloor z'/b \rfloor$ 
     $n := n + 1$ 
end do
output:  $a_n a_{n-1} \cdots a_1 a_0$  in base  $b$  representation.
```

Example 1 (decimal→hex):

The integer $z = 234$ on the number base 10 will be transformed into the hex representation via $a_0 = 234 \bmod 16 = 10 = A$, $234/16 = 14 = E$,

$$a_1 = 14 \bmod 16 = E$$

and therefore we have EA.

Example 2 (binary→decimal):

The binary number $z = 1000100101110101$ is transformed into the decimal representation via the following steps:

$$1000100101110101 = 1001 \pmod{1010} \implies a_0 = 9, \quad 1000100101110101 / 1010 = 110110111110$$

$$110110111110 = 1000 \pmod{1010} \implies a_1 = 8, \quad 110110111110 / 1010 = 10101111$$

$$10101111 = 1 \pmod{1010} \implies a_2 = 1, \quad 10101111 / 1010 = 100011$$

$$100011 = 101 \pmod{1010} \implies a_3 = 5, \quad 100011 / 1010 = 1$$

$$11 = 11 \pmod{1010} \implies a_4 = 3$$

therefore $z = 35189$.

4.18 Appendix: Interactive presentation about the RSA cipher

The following presentation (last update Nov. 2010) shows the basics of the RSA cipher in an interactive way.

There are three variants:

- Powerpoint 2007 (for download; dynamical, animated)¹¹⁸
- PDF (for download; static, no interaction)¹¹⁹
- Flash (can be started within the browser, requires JavaScript; time-controlled replay)¹²⁰

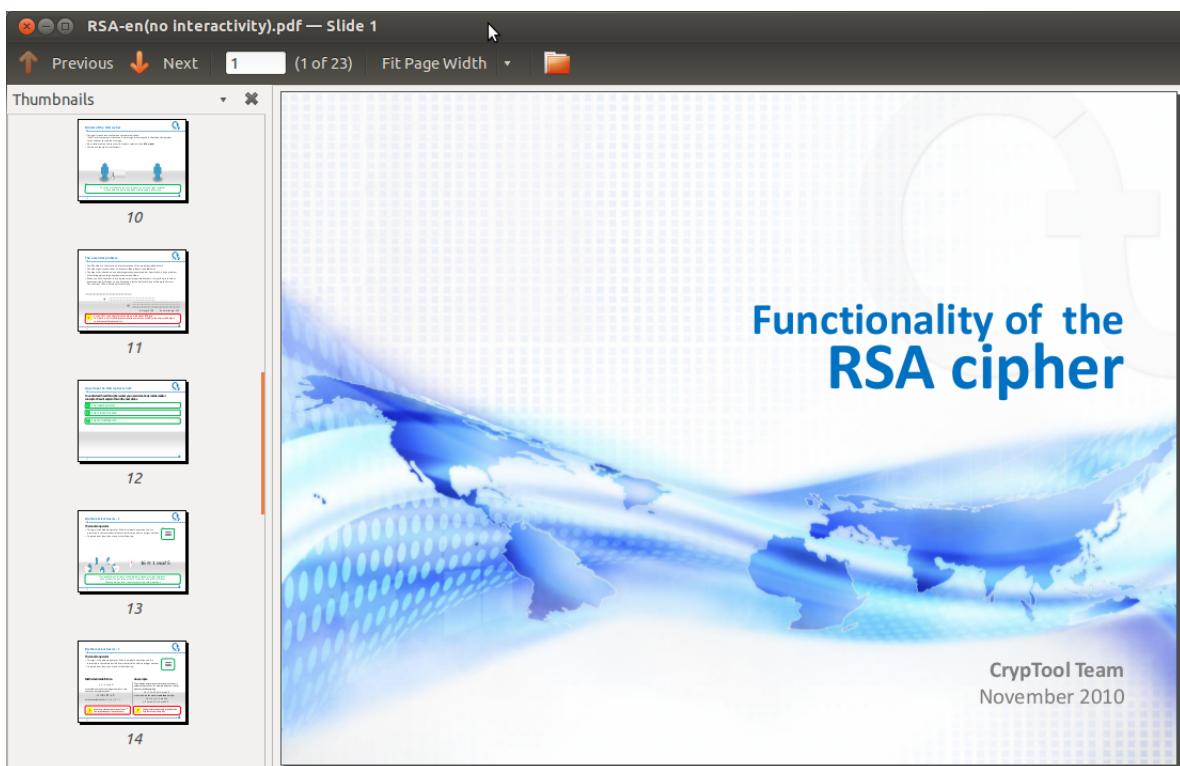


Figure 4.2: Screenshot RSA Presentation (PDF)

¹¹⁸ <http://www.cryptool.org/images/ct1/presentations/RSA/RSA-de.pptx>

¹¹⁹ [http://www.cryptool.org/images/ct1/presentations/RSA/RSA-de\(keine%20Interaktivitaet\).pdf](http://www.cryptool.org/images/ct1/presentations/RSA/RSA-de(keine%20Interaktivitaet).pdf)

¹²⁰ <http://www.cryptool.org/images/ct1/presentations/RSA/RSA-Flash-de/player.html>

4.19 Appendix: Examples using Sage

“She would never be able to tell her parents ... about any of this. She couldn’t tell them about her code-breaking work. About her near death at the hands of the Daemon. About the shadowy entities pulling the strings of her government.”

Quote 13: Daniel Suarez¹²¹

Below you can find Sage source code related to contents of the chapter 4 (“Introduction to Elementary Number Theory with Examples”).

4.19.1 Multiplication table modulo m

The multiplication table 4.4 (from page 113) for $a \times i \pmod{m}$, where $m = 17$, $a = 5$ and $a = 6$, and i ranges over all integers from 0 to 16 can be computed using Sage as follows:

Sage sample 4.1 Multiplication tables for $a \times i \pmod{m}$ with $m = 17$, $a = 5$ and $a = 6$

```
sage: m = 17; a = 5; b = 6
sage: [mod(a * i, m).lift() for i in xrange(m)]
[0, 5, 10, 15, 3, 8, 13, 1, 6, 11, 16, 4, 9, 14, 2, 7, 12]
sage: [mod(b * i, m).lift() for i in xrange(m)]
[0, 6, 12, 1, 7, 13, 2, 8, 14, 3, 9, 15, 4, 10, 16, 5, 11]
```

The function `mod()` returns an object that represents integers modulo m (in our case $m = 17$). From the `Mod` object you can get its single components either with the function `component` or with the function `lift`. We use the method `lift()` to convert that object to an integer representation.

The other multiplication table examples modulo 13 (table 4.5) and modulo 12 (table 4.6) on page 113 can similarly be computed by replacing $m = 17$ with $m = 13$ and $m = 12$ respectively.

4.19.2 Fast exponentiation

The fast exponentiation modulo m can be computed using the Sage function `power_mod()`. The result of this function is an integer. We can compute the exponentiation in the example in chapter “Fast calculation of high powers” on page 115 as follows:

Sage sample 4.2 Fast exponentiation mod $m = 103$

```
sage: a = 87; m = 103
sage: exp = [2, 4, 8, 16, 32, 43]
sage: [power_mod(a, e, m) for e in exp]
[50, 28, 63, 55, 38, 85]
```

¹²¹Daniel Suarez, “Freedom”, Dutton Adult, 2010, Chapter 5, “Crossroad”, p. 229, Philips.

4.19.3 Multiplicative order

The order $\text{ord}_m(a)$ of a number a in the multiplicative group \mathbf{Z}_m^* is the smallest number $i \geq 1$ such that $a^i \equiv 1 \pmod{m}$ holds (see chapter 4.9, “Multiplicative order and primitive roots”). To create table 4.7 on page 123 we can print all exponentiation $a^i \pmod{11}$ as follows:

Sage sample 4.3 Table with all powers $a^i \pmod{m}$ for $m = 11$, $a = 1, \dots, 10$

```
sage: m = 11
sage: for a in xrange(1, m):
....:     print [power_mod(a, i, m) for i in xrange(1, m)]
....:
[1, 1, 1, 1, 1, 1, 1, 1, 1, 1]
[2, 4, 8, 5, 10, 9, 7, 3, 6, 1]
[3, 9, 5, 4, 1, 3, 9, 5, 4, 1]
[4, 5, 9, 3, 1, 4, 5, 9, 3, 1]
[5, 3, 4, 9, 1, 5, 3, 4, 9, 1]
[6, 3, 7, 9, 10, 5, 8, 4, 2, 1]
[7, 5, 2, 3, 10, 4, 6, 9, 8, 1]
[8, 9, 6, 4, 10, 3, 2, 5, 7, 1]
[9, 4, 3, 5, 1, 9, 4, 3, 5, 1]
[10, 1, 10, 1, 10, 1, 10, 1, 10]
```

and including the last column with the order of each $a \pmod{11}$

```
sage: m = 11
sage: for a in xrange(1, m):
....:     lst= [power_mod(a, i, m) for i in xrange(1, m)]
....:     lst.append(multiplicative_order(mod(a,m)))
....:     print lst
....:
[1, 1, 1, 1, 1, 1, 1, 1, 1, 1]
[2, 4, 8, 5, 10, 9, 7, 3, 6, 1, 10]
[3, 9, 5, 4, 1, 3, 9, 5, 4, 1, 5]
[4, 5, 9, 3, 1, 4, 5, 9, 3, 1, 5]
[5, 3, 4, 9, 1, 5, 3, 4, 9, 1, 5]
[6, 3, 7, 9, 10, 5, 8, 4, 2, 1, 10]
[7, 5, 2, 3, 10, 4, 6, 9, 8, 1, 10]
[8, 9, 6, 4, 10, 3, 2, 5, 7, 1, 10]
[9, 4, 3, 5, 1, 9, 4, 3, 5, 1, 5]
[10, 1, 10, 1, 10, 1, 10, 1, 10, 1, 2]
```

Table 4.8 on page 124 gives examples for the order modulo 45 $\text{ord}_{45}(a)$ and the Euler number $J(45)$.

The following Sage code constructs a table similar to that on page 124.

Sage sample 4.4 Table with all powers $a^i \pmod{45}$ for $a = 1, \dots, 12$ plus the order of a

```
sage: m = 45
sage: for a in xrange(1, 13):
....:     lst = [power_mod(a, i, m) for i in xrange(1, 13)]
....:     try:
....:         lst.append(multiplicative_order(mod(a, m)))
....:     except:
....:         lst.append("None")
....:     lst.append(euler_phi(m))
....:     print lst
....:
[1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 24]
[2, 4, 8, 16, 32, 19, 38, 31, 17, 34, 23, 1, 12, 24]
[3, 9, 27, 36, 18, 9, 27, 36, 18, 9, 27, 36, 'None', 24]
[4, 16, 19, 31, 34, 1, 4, 16, 19, 31, 34, 1, 6, 24]
[5, 25, 35, 40, 20, 10, 5, 25, 35, 40, 20, 10, 'None', 24]
[6, 36, 36, 36, 36, 36, 36, 36, 36, 36, 36, 'None', 24]
[7, 4, 28, 16, 22, 19, 43, 31, 37, 34, 13, 1, 12, 24]
[8, 19, 17, 1, 8, 19, 17, 1, 8, 19, 17, 1, 4, 24]
[9, 36, 9, 36, 9, 36, 9, 36, 9, 36, 9, 36, 'None', 24]
[10, 10, 10, 10, 10, 10, 10, 10, 10, 10, 10, 10, 'None', 24]
[11, 31, 26, 16, 41, 1, 11, 31, 26, 16, 41, 1, 6, 24]
[12, 9, 18, 36, 27, 9, 18, 36, 27, 9, 18, 36, 'None', 24]
```

The number $\text{ord}_m(a)$ only exists if a is relatively prime to m , which can be checked with `gcd(a, m)`.

In the above code example, we put the calculation of the multiplicative order within a `try-except` block. This allows Sage to catch any exceptions or errors raised by the function `multiplicative_order()`. If an exception or error is raised in the `try` block, then we know that $\text{ord}_m(a)$ does not exist for that particular value of a , hence in the `except` block we append the string "None" to the row as represented by the object `lst`.

Table 4.9 on page 125 displays exponentiation $a^i \pmod{46}$ as well as the order $\text{ord}_{46}(a)$.

Sage can create that table as follows:

Sage sample 4.5 Table with all powers $a^i \pmod{46}$ for $a = 1, \dots, 23$ plus the order of a

```
sage: m = 46
sage: euler_phi(m)
22
sage: for a in xrange(1, 24):
....:     lst = [power_mod(a, i, m) for i in xrange(1, 24)]
....:     try:
....:         lst.append(multiplicative_order(mod(a, m)))
....:     except:
....:         lst.append("None")
....:     print lst
....:
[1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1]
[2, 4, 8, 16, 32, 18, 36, 26, 6, 12, 24, 2, 4, 8, 16, 32, 18, 36, 26, 6, 12, 24, 2, 'None']
[3, 9, 27, 35, 13, 39, 25, 29, 41, 31, 1, 3, 9, 27, 35, 13, 39, 25, 29, 41, 31, 1, 3, 11]
[4, 16, 18, 26, 12, 2, 8, 32, 36, 6, 24, 4, 16, 18, 26, 12, 2, 8, 32, 36, 6, 24, 4, 'None']
[5, 25, 33, 27, 43, 31, 17, 39, 11, 9, 45, 41, 21, 13, 19, 3, 15, 29, 7, 35, 37, 1, 5, 22]
[6, 36, 32, 8, 2, 12, 26, 18, 16, 4, 24, 6, 36, 32, 8, 2, 12, 26, 18, 16, 4, 24, 6, 'None']
[7, 3, 21, 9, 17, 27, 5, 35, 15, 13, 45, 39, 43, 25, 37, 29, 19, 41, 11, 31, 33, 1, 7, 22]
[8, 18, 6, 2, 16, 36, 12, 4, 32, 26, 24, 8, 18, 6, 2, 16, 36, 12, 4, 32, 26, 24, 8, 'None']
[9, 35, 39, 29, 31, 3, 27, 13, 25, 41, 1, 9, 35, 39, 29, 31, 3, 27, 13, 25, 41, 1, 9, 11]
[10, 8, 34, 18, 42, 6, 14, 2, 20, 16, 22, 36, 38, 12, 28, 4, 40, 32, 44, 26, 30, 24, 10, 'None']
[11, 29, 43, 13, 5, 9, 7, 31, 19, 25, 45, 35, 17, 3, 33, 41, 37, 39, 15, 27, 21, 1, 11, 22]
[12, 6, 26, 36, 18, 32, 16, 8, 4, 2, 24, 12, 6, 26, 36, 18, 32, 16, 8, 4, 2, 24, 12, 'None']
[13, 31, 35, 41, 27, 29, 9, 25, 3, 39, 1, 13, 31, 35, 41, 27, 29, 9, 25, 3, 39, 1, 13, 11]
[14, 12, 30, 6, 38, 26, 42, 36, 44, 18, 22, 32, 34, 16, 40, 8, 20, 4, 10, 2, 28, 24, 14, 'None']
[15, 41, 17, 25, 7, 13, 11, 27, 37, 3, 45, 31, 5, 29, 21, 39, 33, 35, 19, 9, 43, 1, 15, 22]
[16, 26, 2, 32, 6, 4, 18, 12, 8, 36, 24, 16, 26, 2, 32, 6, 4, 18, 12, 8, 36, 24, 16, 'None']
[17, 13, 37, 31, 21, 35, 43, 41, 7, 27, 45, 29, 33, 9, 15, 25, 11, 3, 5, 39, 19, 1, 17, 22]
[18, 2, 36, 4, 26, 8, 6, 16, 12, 32, 24, 18, 2, 36, 4, 26, 8, 6, 16, 12, 32, 24, 18, 'None']
[19, 39, 5, 3, 11, 25, 15, 9, 33, 29, 45, 27, 7, 41, 43, 35, 21, 31, 37, 13, 17, 1, 19, 22]
[20, 32, 42, 12, 10, 16, 44, 6, 28, 8, 22, 26, 14, 4, 34, 36, 30, 2, 40, 18, 38, 24, 20, 'None']
[21, 27, 15, 39, 37, 41, 33, 3, 17, 35, 45, 25, 19, 31, 7, 9, 5, 13, 43, 29, 11, 1, 21, 22]
[22, 24, 22, 24, 22, 24, 22, 24, 22, 24, 22, 24, 22, 24, 22, 24, 22, 24, 22, 24, 22, 'None']
[23, 23, 23, 23, 23, 23, 23, 23, 23, 23, 23, 23, 23, 23, 23, 23, 23, 23, 23, 23, 23, 'None']
```

The following code for generating the tables 4.10 and 4.11 at page 126 f. also delivers the result in a way, that it can be easily processed in LaTeX. The prerequisite is that all content is assigned to one Sage object (here a matrix).¹²²

Sage sample 4.6 Code for tables with all powers $a^i \pmod{m}$ for variable a and i plus order of a and Eulerphi of m

```
def power_mod_order_matrix(m, max_a, max_i):
    r = matrix(ZZ, max_a+1, max_i+3)
    for a in xrange(0, max_a+1):
        r[a, 0] = a
        for i in xrange(1, max_i+1):
            if a==0:
                r[a,i] = i
            else:
                r[a, i] = power_mod(a, i, m)
        try:
            r[a, max_i+1] = multiplicative_order(mod(a, m))
        except:
            r[a, max_i+1] = 0
            r[a, max_i+2] = euler_phi(m)
    return r

print "\n1: m=45;max_i=13;max_a=13";m=45;max_i=13;max_a=13
r = power_mod_order_matrix(m, max_a, max_i);print r;print latex(r)

print "\n2: m=46;max_i=25;max_a=25";m=46;max_i=25;max_a=25
r = power_mod_order_matrix(m, max_a, max_i);print r.str();print latex(r)

print "\n3: m=14;max_i=13;max_a=16";m=14;max_i=13;max_a=16
r = power_mod_order_matrix(m, max_a, max_i);print r;print latex(r)

print "\n4: m=22;max_i=21;max_a=25";m=22;max_i=21;max_a=25
r = power_mod_order_matrix(m, max_a, max_i);print r.str();print latex(r)
...
3: m=14;max_i=13;max_a=16
[ 0  1  2  3  4  5  6  7  8  9 10 11 12 13  0  6]
[ 1  1  1  1  1  1  1  1  1  1  1  1  1  1  1  6]
[ 2  2  4  8  2  4  8  2  4  8  2  4  8  2  0  6]
[ 3  3  9 13 11  5  1  3  9 13 11  5  1  3  6  6]
...
\left(\begin{array}{rrrrrrrrrrrrrr}
0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 & 12 & 13 & 0 & 6 \\
1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 6 \\
2 & 2 & 4 & 8 & 2 & 4 & 8 & 2 & 4 & 8 & 2 & 4 & 8 & 2 & 0 & 6 \\
3 & 3 & 9 & 13 & 11 & 5 & 1 & 3 & 9 & 13 & 11 & 5 & 1 & 3 & 6 & 6
\end{array}\right)
```

¹²²Remark about the Sage program, especially the Sage indices:

- for x in $xrange(2, 5)$ delivers 2,3,4.
- $m = matrix(ZZ, 2, 5)$ has 2 rows and 5 columns.
The cells are named $m(0,0)$ to $m(1,4)$.
- All elements of the matrix have to be numerical, so “0” instead of “None” as in the tables before.
- The output of matrices kann be controlled in Sage with:

```
sage: from sage.matrix.matrix0 import set_max_cols, set_max_rows
sage: set_max_cols(100)
sage: set_max_rows(100)
```

- The length of the cycle in the last column of the tables 4.10 and 4.11 was added manually.

4.19.4 Primitive roots

Computing a primitive root (see chapter 4.9, “Multiplicative order and primitive roots”) in Sage is very straightforward. If n is an integer, the command `primitive_root(n)` computes *one* primitive root of the multiplicative group $(\mathbf{Z}/n\mathbf{Z})^*$, if one exists. Where n is prime, then this is the same as calculating a primitive root of $\mathbf{Z}/n\mathbf{Z}$.

Here, we calculate some primitive roots of a few integers.

Sage sample 4.7 Calculating one primitive root for a given prime

```
sage: primitive_root(4)
3
sage: primitive_root(22)
13
sage: for p in primes(1, 50):
....:     print p, primitive_root(p)
....:
2 1
3 2
5 2
7 3
11 2
13 2
17 3
19 2
23 5
29 2
31 3
37 2
41 6
43 3
47 5
```

If p is prime, then $\mathbf{Z}/p\mathbf{Z}$ has at least one primitive root.

Sometimes we want to compute all the primitive roots of $\mathbf{Z}/p\mathbf{Z}$, not just any primitive root of $\mathbf{Z}/p\mathbf{Z}$. The following function can do this¹²³.

Sage sample 4.8 Function “enum_PrimitiveRoots_of_an_Integer” to calculate all primitive roots for a given number

```
def enum_PrimitiveRoots_of_an_Integer(M):
    """
    Return all the primitive roots of the integer M (if possible).
    """
    try:
        g = primitive_root(M)
    except:
        return None
    targetOrder = euler_phi(M)
    L=[]
    # Stepping through all odd integers from 1 up to M, not including
    # M. So this loop only considers values of i where 1 <= i < M.
    for i in xrange(1,M,2):
        testGen = Mod(g^i,M)
        if testGen.multiplicative_order() == targetOrder:
            L.append(testGen.lift())
    # removing duplicates
    return Set(L)

# AA_Start -- Testcases for enum_PrimitiveRoots_of_an_Integer(M)
print "AA_Start -- Testcases for enum_PrimitiveRoots_of_an_Integer(M)"
M=10; print "1-----Testcase: M = %s" % M
LL = enum_PrimitiveRoots_of_an_Integer(M)
if LL==None:
    print M
else:
    print LL
M=8; print "2-----Testcase: M = %s" % M
# M=8 hat keine primitive root mod m. Checke, ob per try - except abgefangen.
LL = enum_PrimitiveRoots_of_an_Integer(M)
if LL==None:
    print M
else:
    print LL
M=17; print "3-----Testcase: M = %s" % M
LL = enum_PrimitiveRoots_of_an_Integer(M)
if LL==None:
    print M
else:
    print LL
# AA_End -- Testcases

OUTPUT:
AA_Start -- Testcases for enum_PrimitiveRoots_of_an_Integer(M)
1-----Testcase: M = 10
{3, 7}
2-----Testcase: M = 8
8
3-----Testcase: M = 17
{3, 5, 6, 7, 10, 11, 12, 14}
```

¹²³This code was developed in a Sage script file and executed non-interactively. That is why you don't see "sage:" and "...:" at the beginning of the lines like in the Sage samples before.

For example, here is a list of all primitive roots of the prime 541.

Sage sample 4.9 Table with all primitive roots for the given prime 541

```
sage: L=enum_PrimitiveRoots_of_an_Integer(541); L
{2, 517, 10, 523, 13, 14, 527, 528, 18, 531, 24, 539, 30, 37, 40, 51,
54, 55, 59, 62, 65, 67, 68, 72, 73, 77, 83, 86, 87, 91, 94, 96, 98,
99, 107, 113, 114, 116, 117, 126, 127, 128, 131, 132, 138, 150, 152,
153, 156, 158, 163, 176, 181, 183, 184, 195, 197, 199, 206, 208,
210, 213, 218, 220, 223, 224, 244, 248, 250, 257, 258, 259, 260,
261, 263, 267, 269, 270, 271, 272, 274, 278, 280, 281, 282, 283,
284, 291, 293, 297, 317, 318, 321, 323, 328, 331, 333, 335, 342,
344, 346, 357, 358, 360, 365, 378, 383, 385, 388, 389, 391, 403,
409, 410, 413, 414, 415, 424, 425, 427, 428, 434, 442, 443, 445,
447, 450, 454, 455, 458, 464, 468, 469, 473, 474, 476, 479, 482,
486, 487, 490, 501, 504, 511}
sage: len(L)
144
```

With a little bit of programming, we can count how many primitive roots are in a given range of integers. We can check this for all numbers or only for the primes within this range.

Sage sample 4.10 Function “count_PrimitiveRoots_of_an_IntegerRange” to calculate all primitive roots for a given range of integers

```
def count_PrimitiveRoots_of_an_IntegerRange(start, end, bPrimesOnly=True):
    """
    Compute all primitive roots of all numbers between start and end,
    inclusive, and count them.
    If the flag bPrimesOnly is True, it performs primality tests, so it
    allows us to count the number of primes from start to end, inclusive.
    If the flag bPrimesOnly is false, it additionally counts these even
    numbers which have no primitive root.
    """
    nCheckedNumb = 0
    nCheckedNumb_WithoutPrimitivRoots = 0
    nPrimitiveRoots = 0
    for n in xrange(start, end+1):
        if bPrimesOnly:
            if is_prime(n):
                nCheckedNumb += 1
                L = enum_PrimitiveRoots_of_an_Integer(n)
                nPrimitiveRoots += len(L)
            else:
                nCheckedNumb += 1
                L = enum_PrimitiveRoots_of_an_Integer(n)
                if L==None:
                    nCheckedNumb_WithoutPrimitivRoots += 1
                else:
                    nPrimitiveRoots += len(L)

        if bPrimesOnly:
            print "Found all %s" % nPrimitiveRoots + \
                  " primitive roots of %s primes." % nCheckedNumb
        else:
            if nCheckedNumb_WithoutPrimitivRoots == 0:
                print "Found all %s " % nPrimitiveRoots + \
                      "primitive roots of %s numbers." % nCheckedNumb
            else:
                print "Found all %s " % nPrimitiveRoots + \
                      "primitive roots of %s numbers." % \
                      (nCheckedNumb - nCheckedNumb_WithoutPrimitivRoots)
            print "(Total of numbers checked: %s, " % nCheckedNumb + \
                  "Amount of numbers without primitive roots: %s)" % \
                  nCheckedNumb_WithoutPrimitivRoots
```

Using the Sage command `time`, we can also find out how long it takes on our computer.

Sage sample 4.11 Function “`count_PrimitiveRoots_of_an_IntegerRange`”: testcases and output

```
# BB_Start -- Testcases for count_PrimitiveRoots_of_an_IntegerRange(start, end, bPrimesOnly=True)
print "\n\nBB_Start -- Testcases for count_PrimitiveRoots_of_an_IntegerRange(start, end, True)"

print "\n1-----Testcase: (1, 500)"
time count_PrimitiveRoots_of_an_IntegerRange(1, 500)

print "\n2-----Testcase: (5, 6, False)"
time count_PrimitiveRoots_of_an_IntegerRange(5, 6, False)

print "\n3-----Testcase: (1, 500, False)"
time count_PrimitiveRoots_of_an_IntegerRange(1, 500, False)
# BB_End -- Testcases

OUTPUT:
BB_Start -- Testcases for count_PrimitiveRoots_of_an_IntegerRange(start, end, bPrimesOnly=True)

1-----Testcase: (1, 500)
Found all 8070 primitive roots of 95 primes.
Time: CPU 0.94 s, Wall: 0.97 s

2-----Testcase: (5, 6, False)
Found all 3 primitive roots of 2 numbers.
Time: CPU 0.00 s, Wall: 0.00 s

3-----Testcase: (1, 500, False)
Found all 11010 primitive roots of 170 numbers.
(Total of numbers checked: 500, Amount of numbers without primitive roots: 330)
Time: CPU 1.52 s, Wall: 1.59 s
```

Using our custom-defined function `enum_PrimitiveRoots_of_an_Integer`, we can find all primitive roots of one prime integer p .

The following function counts how many primes are in a given range and enumerate all their primitive roots.

From this list of primitive roots, we can determine the smallest and largest primitive root for $\mathbf{Z}/p\mathbf{Z}$, as well as count the number of primitive roots of $\mathbf{Z}/p\mathbf{Z}$.

Sage sample 4.12 Function “`count_PrimitiveRoots_of_a_PrimesRange`” to calculate the number of primitive roots for a given range of primes

```
def count_PrimitiveRoots_of_a_PrimesRange(start, end):
    """
    Compute all primitive roots of all primes between start and end,
    inclusive. This uses a primes iterator.
    """
    nPrimes = 0
    nPrimitiveRoots = 0
    for p in primes(start, end+1):
        L = enum_PrimitiveRoots_of_an_Integer(p)
        print p, len(L)
        nPrimes += 1
        nPrimitiveRoots += len(L)
    print "Found all %s" % nPrimitiveRoots + " primitive roots of %s primes." % nPrimes

# CC_Start -- Testcases for count_PrimitiveRoots_of_a_PrimesRange(start, end)
print "\n\nBB_Start -- Testcases for count_PrimitiveRoots_of_a_PrimesRange(start, end)"
print "-----Testcase: (1, 1500)"
time count_PrimitiveRoots_of_a_PrimesRange(1, 1500)
# CC_End -- Testcases

OUTPUT:
CC_Start -- Testcases for count_PrimitiveRoots_of_a_PrimesRange(start, end)
-----Testcase: (1, 1500)
2 1
3 1
5 2
7 2
11 4
13 4
17 8
19 6
23 10
29 12
31 8
37 12
...
1483 432
1487 742
1489 480
1493 744
1499 636
Found all 62044 primitive roots of 239 primes.
Time: CPU 7.55 s, Wall: 7.85 s
```

A slightly modified version of our function `count_PrimitiveRoots_of_a_PrimesRange`, was used to generate a database of all primitive roots of all primes between 1 and 100,000.

Sage sample 4.13 Code to generate the database with all primitive roots for all primes between 1 and 100,000

```
start = 1
end = 10^5
fileName = "/scratch/mvngu/primroots.dat"
file = open(fileName, "w")
for p in primes(start, end+1):
    L = enum_PrimitiveRoots_of_an_Integer(p)
    print p, len(L)
    # Output to a file. The format is:
    # (1) the prime number p under consideration
    # (2) the number of primitive roots of Z/pZ
    # (3) all the primitive roots of Z/pZ
    file.write(str(p) + " " + str(len(L)) + " " + str(L) + "\n")
    file.flush()
file.close()
```

It took about one day on the machine sage.math to generate the file “primroots.dat” (done in July 2009 by Minh Van Nguyen).

This code and the function `enum_PrimitiveRoots_of_an_Integer` was put in a Sage script file and executed non-interactively.

The file “primroots.dat” is a database of all primitive roots of all primes between 1 and 100,000 inclusive. It is a very large file (about 1 GB uncompressed, and 285 MB compressed with bzip2). You can find the file at <http://sage.math.washington.edu/home/mvngu/doc/primitive-roots/primroots.dat.bz2>.

This database file “primroots.dat” was used then to create three graphics using the following code.

Sage sample 4.14 Code to generate the graphics about the primitive roots

```
sage: # open a database file on primitive roots from 1 to 100,000
sage: file = open("/scratch/mvngu/primroots.dat", "r")
sage: plist = []      # list of all primes from 1 to 100,000
sage: nlist = []      # number of primitive roots modulo prime p
sage: minlist = []    # smallest primitive root modulo prime p
sage: maxlist = []    # largest primitive root modulo prime p
sage: for line in file:
....:     # get a line from the database file and tokenize it for processing
....:     line = line.strip().split(" ", 2)
....:     # extract the prime number p in question
....:     plist.append(Integer(line[0]))
....:     # extract the number of primitive roots modulo p
....:     nlist.append(Integer(line[1]))
....:     # extract the list of all primitive roots modulo p
....:     line = line[-1]
....:     line = line.replace("{", "")
....:     line = line.replace("}", "")
....:     line = line.split(", ")
....:     # sort the list in non-decreasing order
....:     line = [Integer(s) for s in line]
....:     line.sort()
....:     # get the smallest primitive root modulo p
....:     minlist.append(line[0])
....:     # get the largest primitive root modulo p
....:     maxlist.append(line[-1])
....:
sage: file.close()  # close the database file
sage: # plot of number of primitive roots modulo p
sage: nplot = point2d(zip(plist, nlist), pointsize=1)
sage: nplot.axes_labels(["x", "y"])
sage: nplot
sage: # plot of smallest primitive root modulo prime p
sage: minplot = point2d(zip(plist, minlist), pointsize=1)
sage: minplot.axes_labels(["x", "y"])
sage: minplot
sage: # plot of largest primitive root modulo prime p
sage: maxplot = point2d(zip(plist, maxlist), pointsize=1)
sage: maxplot.axes_labels(["x", "y"])
sage: maxplot
```

Figure 4.3 graphs the number of primitive roots for each prime between 1 and 100,000. The x -axis represents primes between 1 and 100,000, while the y -axis counts the number of primitive roots for each prime within that interval.

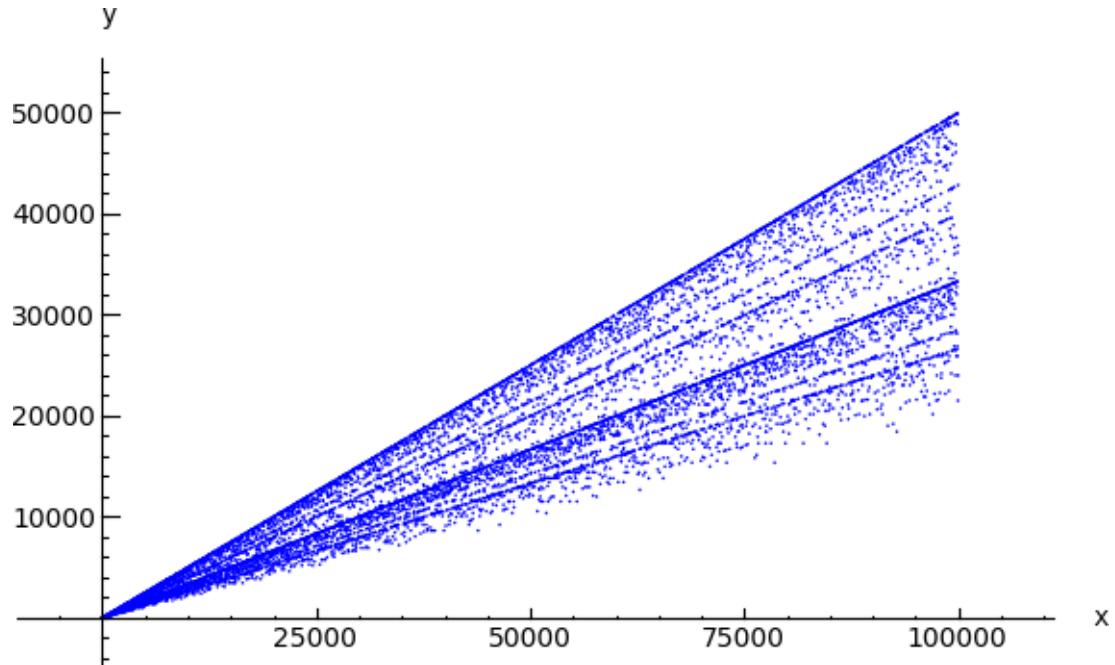


Figure 4.3: The number of primitive roots of all primes between 1 and 100,000.

Figure 4.4 graphs the smallest primitive roots of all primes between 1 and 100,000. The x -axis represents primes between 1 and 100,000. The y -axis represents the smallest primitive root of each prime within that interval.

Figure 4.5 shows a corresponding graph for the largest primitive root of each prime within the above interval.

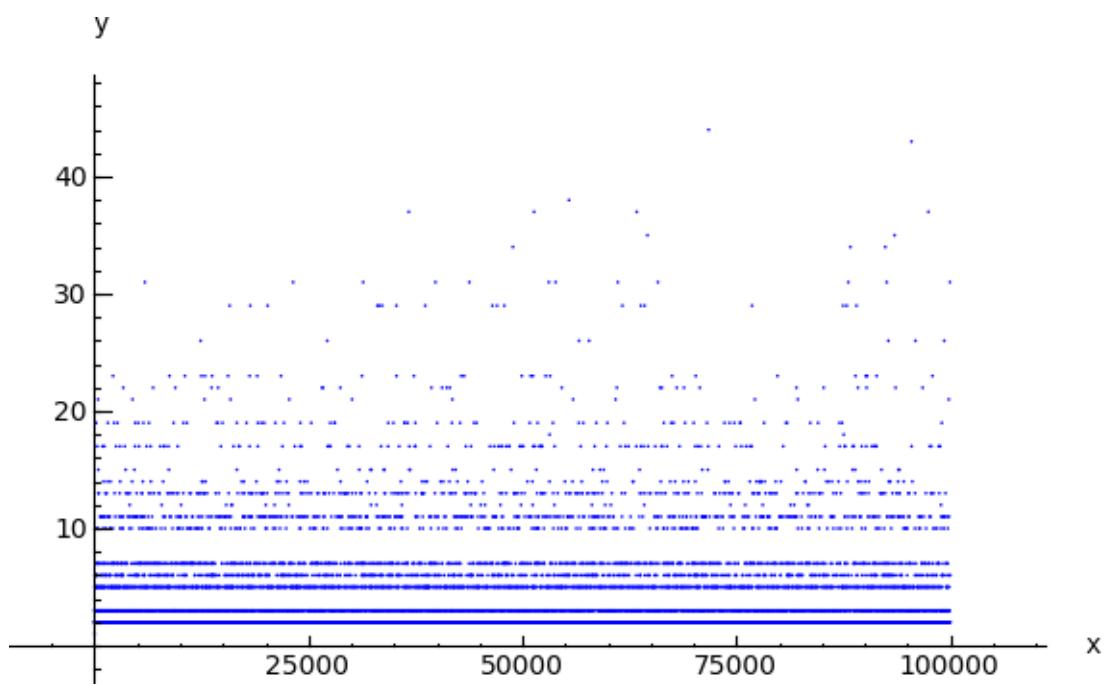


Figure 4.4: The smallest primitive roots of all primes between 1 and 100,000.

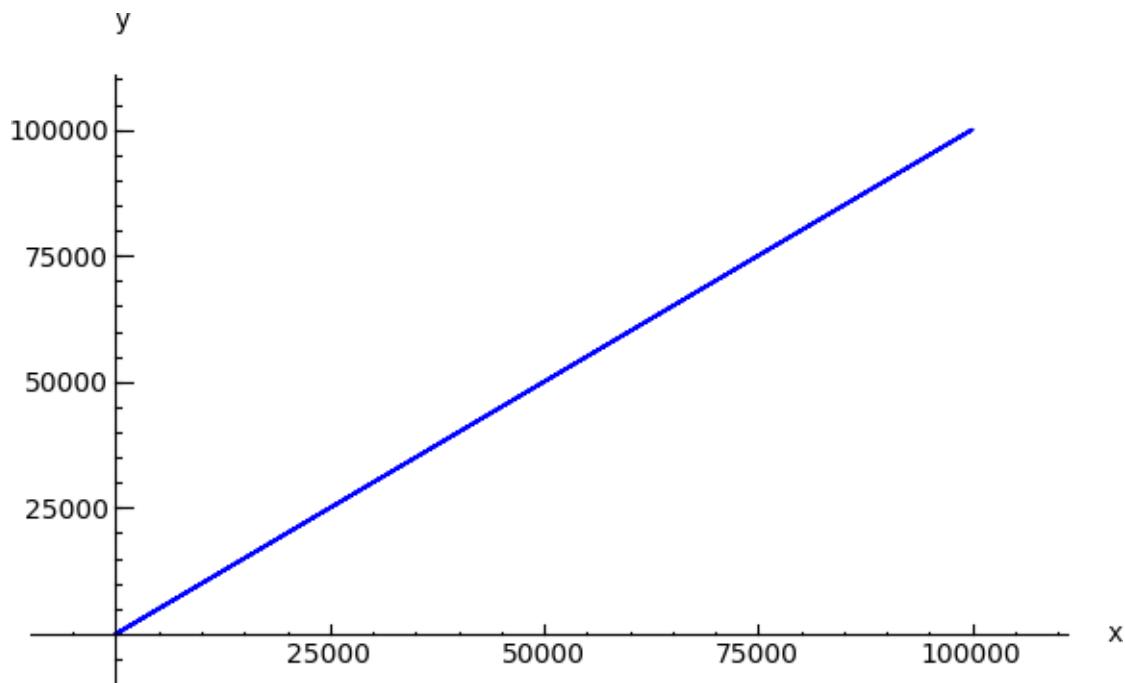


Figure 4.5: The largest primitive roots of all primes between 1 and 100,000.

4.19.5 RSA examples with Sage

Below is Sage source code for the simple RSA examples in section 4.13 (“The RSA procedure with actual numbers”).

Example on page 150:

The RSA exponentiation $M^{37} \pmod{3713}$ on message $M = 120$ can be calculated in Sage as follows:

```
sage: power_mod(120, 37, 3713)
1404
```

Example on page 151:

The factorization of $J(256027) = 255016 = 2^3 * 127 * 251$ can be calculated using Sage as follows:

Sage sample 4.15 Factoring a number

```
sage: factor(255016)
2^3 * 127 * 251
```

Example on page 151:

Sage can do RSA encryption as follows:

Sage sample 4.16 RSA encryption by modular exponentiation of a number (used as message)

```
sage: A = [82, 83, 65, 32, 119, 111, 114, 107, 115, 33]
sage: e = 65537; m = 256027
sage: [power_mod(a, e, m) for a in A]
[212984, 25546, 104529, 31692, 248407, 100412, 54196, 100184, 58179, 227433]
```

Example on page 152:

RSA encryption using Sage:

```
sage: A = [21075, 16672, 30575, 29291, 29473]
sage: e = 65537; m = 256027
sage: [power_mod(a, e, m) for a in A]
[158721, 137346, 37358, 240130, 112898]
```

Example on page 152:

RSA encryption using Sage:

```
sage: A = [82083, 65032, 119111, 114107, 115033]
sage: e = 65537; m = 256027
sage: [power_mod(a, e, m) for a in A]
[198967, 51405, 254571, 115318, 14251]
```

4.19.6 How many private RSA keys d exist within a given modulo range?

The RSA encryption procedure was described in section 4.10.2 (“How the RSA procedure works”). Steps 1 to 3 constitute key generation, steps 4 and 5 are the encryption:

1. Select two distinct random prime numbers p and q and calculate $n = p * q$.
The value n is called the RSA modulus.
2. Select an arbitrary $e \in \{2, \dots, n - 1\}$ such that:
 e is relatively prime to $J(n) = (p - 1) * (q - 1)$.
We can then “throw away” p and q .
3. Select $d \in \{1, \dots, n - 1\}$ with $e * d \equiv 1 \pmod{J(n)}$,
i.e. d is the multiplicative inverse of e modulo $J(n)$. We can then “throw away” $J(n)$.
 $\rightarrow (n, e)$ is the public key P .
 $\rightarrow (n, d)$ is the private key S (only d must be kept secret).
4. For encryption, the message represented as a (binary) number is divided into parts such that each part of the number represents a number less than n .
5. Encryption of the plaintext (or the parts of it) $M \in \{1, \dots, n - 1\}$:

$$C = E((n, e); M) := M^e \pmod{n}.$$

The default way to crack a given RSA ciphertext C would be to use the public key of the recipient and to try to factorize n . Then you can go through the steps 2 and 3 and generate the private key e , which is normally used to decrypt a ciphertext.

According to the “prime number theorem”¹²⁴ the number of prime numbers $PI(x)$ is asymptotic to $x/\ln(x)$. Between 1 and a given n there are about $n/\ln(n)$ different primes.

If you don't want to use factorization but ask the question like in classic encryption, you may want to find out: How many possible private keys (n, d) are there for a given key size range $n \in [a, b]$?¹²⁵

Sage source code 4.17 below defining the function `count_Number_of_RSA_Keys` can answer this question concretely (if the modulus is not too big).¹²⁶

As there are many more private keys (n, d) within a bigger range of values for n , even brute-force factoring is much more efficient as brute-force trying all the keys.

¹²⁴See section 3.7.2 (“Density and distribution of the primes”).

¹²⁵Chapter 4.8.5 (“How many private RSA keys d are there modulo 26”), p. 121 deals with the special case $n = 26$.

¹²⁶a) Calling `sage: count_Number_of_RSA_Keys(100, 1000)` means to consider the interval $[100, 1000]$ for n .

n is defined by the two primes $p, q : n = p * q$.

So here one prime can have the maximal value 500 because $2 * 500 = 1000$ (while then the other prime will have the smallest possible prime value 2).

The number of possible combinations of primes is $comb = 258$.

The number of primes in the given range is 143.

The number of private keys is 34,816.

b) Calling `sage: count_Number_of_RSA_Keys(100, 100, True)` has the following output:

- Number of private keys for modulus in a given range: 0

- Number of primes in a given range: 0

The reason for that is, that with this call only $n = 100$ is considered, and the function investigates only semiprime n : 100 is not semi prime, this means 100 is not the product of only two primes.

Sage sample 4.17 How many private RSA keys d are there if you know a range for the public key n?

```
def count_Number_of_RSA_Keys(start, end, Verbose=False):
    """
    How many private RSA keys (n,d) exist, if only modulus N is given, and start <= N <= end?
    (prime_range(u,o) delivers all primes >=u und < o).
    """
    a = start
    b = end
    s = 0
    comb = 0
    for p in prime_range(1, b/2+1):
        for q in prime_range(p + 1, b/2+1):
            if a <= p * q and p * q <= b:
                comb = comb +1
                s = s + (euler_phi(euler_phi(p * q))-1)
            if Verbose:
                print "p=%s, " % p + "q=%s, " % q + "s=%s" % s
    print "Number of private keys d for modulus in a given range: %s" % s + " (comb=%s), " % comb

    # Just for comparison: How many primes are in this range?
    s = 0
    for p in prime_range(a, b+1):
        if Verbose:
            print "a=%s, " % a + "b=%s, " % b + "p=%s" % p
        s = s + 1
    print "Number of primes in a given range: %s" % s

print "\n\nDD_Start -- Testcases for count_Number_of_RSA_Keys(start, end)"
print "\n-----Testcase: (100, 1000) [Should deliver 34.816]"
time count_Number_of_RSA_Keys(100, 1000)
print "\n-----Testcase: (100, 107, True) [Should deliver 23]"
time count_Number_of_RSA_Keys(100, 107, True)
u = 10^3; o = 10^4;
print "\n-----Testcase: (%s, " % u + "%s) [Should deliver 3.260.044]" % o
time count_Number_of_RSA_Keys(u, o)

OUTPUT:
DD_Start -- Testcases for count_Number_of_RSA_Keys(start, end)

-----Testcase: (100, 1000) [Should deliver 34.816]
Number of private keys d for modulus in a given range: 34816 (comb=258),
Number of primes in a given range: 143
Time: CPU 0.03 s, Wall: 0.04 s

-----Testcase: (100, 107, True) [Should deliver 23]
p=2, q=53, s=23
Number of private keys d for modulus in a given range: 23 (comb=1),
a=100, b=107, p=101
a=100, b=107, p=103
a=100, b=107, p=107
Number of primes in a given range: 3
Time: CPU 0.00 s, Wall: 0.00 s

-----Testcase: (1000, 10000) [Should deliver 3,260,044]
Number of private keys d for modulus in a given range: 3260044 (comb=2312),
Number of primes in a given range: 1061
Time: CPU 0.63 s, Wall: 0.66 s
```

4.19.7 RSA fixed points $m^e = m \bmod n$ mit $m \in \{1, \dots, n-1\}$

Also encryption methods can have fixed – cleartext messages where the according ciphertext matches the original. In mathematics, variables mapped by the algorithm (function) onto themselves are called fixed points. In cryptography the according messages are called “unconcealed messages”.

Generally speaking: The more fixed points an encryption algorithm contains, the easier it is to break it.

With the RSA procedure: $n = pq$ is the product of two different prime numbers, and there exists e where $\gcd(e, (p-1)(q-1)) = 1$. The encryption is then $c = m^e \bmod n$. A fixed point in the RSA procedure is the message m , where: $m = m^e \bmod n$. The result of the encryption is the given message.

When the size of n is sufficiently big, the probability of the occurrence of fixed points in RSA is very small – as illustrated in Figure 4.6: In average, we found not more than 40 fixed points.

Students often presume the occurrence of fixed points high, because they counter a “relatively” large number of examples when experimenting with **small** prime numbers, as $m = 0, 1$ and $n-1$ are also always fixed points.

In practice, where large prime numbers are chosen, fixed points have no significance for the security of RSA. Therefore, this paragraph refers more to the mathematical questions.¹²⁷

4.19.7.1 The number of RSA fixed points

In this section we show how many RSA fixed points there are for $m \in \{1, \dots, n-1\}$.

Theorem 4.19.1. *The number of the fixed points $m^e = m \bmod n$ with $m \in \{1, \dots, n-1\}$ is $\gcd(p-1, e-1) \cdot \gcd(q-1, e-1)$.*

Proof. Given $m^e = m \bmod n$. According to the CRT, the following statements are equivalent:

$$[m^e = m \bmod n] \Leftrightarrow [m^e = m \bmod p \text{ and } m^e = m \bmod q]$$

This decomposition is equivalent to:

$$m^{e-1} = 1 \bmod p \text{ and } m^{e-1} = 1 \bmod q.$$

We consider $m^{e-1} = 1 \bmod p$ and search all $(e-1)$ roots of unity¹²⁸ in \mathbb{Z}_p^* .

It holds: \mathbb{Z}_p^* for p prime is cyclic. \Rightarrow A generator g exists which produces \mathbb{Z}_p^* : $\mathbb{Z}_p^* = \langle g \rangle$.

¹²⁷Thanks to Taras Shevchenko for gathering parts of the content of this chapter and to Volker Simon for writing the Sage program 4.18 “Getfixpoints”.

¹²⁸- In mathematics, a **root of unity** is a number x that equals 1 when raised to some integer power n .

- An n -th root of unity x is **primitive** if it is not a k -th root of unity for all integers k smaller than n :

$$x^n = 1 \text{ and } x^k \neq 1 \quad (k = 1, 2, 3, \dots, n-1)$$

- If F is a finite field and n is a positive integer, then a n th-root of unity in F is a solution of the equation

$$x^n - 1 = 0 \text{ in } F$$

The following theorem from [Katzenbeisser2001, Pg. 69] characterizes all $(e - 1)$ -th roots of unity in \mathbb{Z}_p^* :

Theorem 4.19.2. g^α is exactly then $(e - 1)$ -th root of unity in \mathbb{Z}_p^* , when $(e - 1)\alpha = 0 \pmod{p-1}$. There are $\gcd(p-1, e-1)$ of these.

Proof. The first thesis results directly from the small theorem from Fermat:

$$g^{\alpha(e-1)} = 1 \pmod{p} \Rightarrow \alpha(e-1) = 0 \pmod{p-1}$$

Let $\delta = \gcd(p-1, e-1)$. $\alpha(e-1) = 0 \pmod{p-1}$ implies $\frac{\alpha(e-1)}{\delta} = 0 \pmod{\frac{p-1}{\delta}}$. Since $\frac{e-1}{\delta}$ and $\frac{p-1}{\delta}$ are coprime (each was reduced by the gcd of their corresponding numerator), α must be a multiple of $\frac{p-1}{\delta}$.

$$\alpha \frac{p-1}{\delta} \text{ with } \alpha = 1, \dots, \delta$$

These δ different powers then correspond to the $(e - 1)$ -th roots of unity $g^{\alpha \frac{p-1}{\delta}} \pmod{p}$ in \mathbb{Z}_p^* . \square

Analog for q : For $m^{e-1} = 1 \pmod{q}$ we then have $\gcd(q-1, e-1)$ many of $(e - 1)$ -th roots of unity.

The number of combinations of the $(e - 1)$ -th root of unity in \mathbb{Z}_p^* and \mathbb{Z}_q^* gives the total quantity of RSA fixed points: $m^e = m \pmod{n}$ with $m \in \{1, \dots, n-1\}$:

$$\gcd(p-1, e-1) \cdot \gcd(q-1, e-1)$$

Adding $m = 0$ to the above, results in the theorem 4.19.3:

Theorem 4.19.3. If $m \in \{0, \dots, n-1\}$, then the quantity of the RSA fixed points is:

$$(\gcd(p-1, e-1) + 1) \cdot (\gcd(q-1, e-1) + 1)$$

\square

4.19.7.2 Lower bound for the quantity of RSA fixed points

In the following section, we show that there is a lower bound for the quantity of RSA fixed points. This lower bound 6 exists when the two different RSA prime numbers are the smallest possible values (2 and 3).

Thesis 1: $p = 2, q = 3$

The quantity of RSA fixed points for $p = 2$ and $q = 3$ is

$$\underbrace{(\gcd(p-1, e-1) + 1)}_{=1} \cdot \underbrace{(\gcd(q-1, e-1) + 1)}_{=2} = 2 \cdot 3 = 6$$

Thesis 2: $p \neq q; p > 2, q > 2$

The quantity of RSA fixed points for $p \neq q; p, q > 2$ is ≥ 9 .

Proof. Since p and q are prime, $(p - 1)$ and $(q - 1)$ for $p, q > 2$ are even.

The RSA algorithm requires to choose e so that $1 < e < \phi(n) = (p - 1)(q - 1)$ and $\gcd(e, (p - 1)(q - 1)) = 1$

Since $(p - 1)$ and $(q - 1)$ are even, e is odd $\Rightarrow e - 1$ is even.

Since $(p - 1)$ and $(e - 1)$ are even, then:

$$\gcd(p - 1, e - 1) \geq 2$$

$$\Rightarrow (\gcd(p - 1, e - 1) + 1) \geq 3 \text{ and } (\gcd(q - 1, e - 1) + 1) \geq 3$$

$$\Rightarrow (\gcd(p - 1, e - 1) + 1) \cdot (\gcd(q - 1, e - 1) + 1) \geq 9$$

□

Samples:

For $(e, n) = (17, 6)$, all six possible messages $\{0, 1, 2, 3, 4, 5\}$ are fixed points (for $n = 6$, it is independent from the value of e).

For $(e, n) = (17, 10)$, all 10 possible messages are fixed points.

For $(e, n) = (19, 10)$, only 6 of the 10 possible messages are fixed points.

4.19.7.3 Unfortunate choice of e

In this section, we show that with $e = 1 + \text{lcm}(p - 1, q - 1)$ each encryption results in a fixed point (independently of the size of p, q , or n); and then we broaden this to all unfortunate choices of e .

If $e = 1$, then for all m : $c = m^e = m$. This is the trivial case.

Thesis 1: $p, q > 2$

If $e = 1 + \text{lcm}(p - 1, q - 1)$, then for all $m \in \{1, \dots, n - 1\}$: $m^e = m \pmod{n}$.

Proof.

Given:

- $e \cdot d = 1 \pmod{\phi(n)}$ or $e \cdot d = 1 \pmod{\text{lcm}(p - 1, q - 1)}$
- $m^x \pmod{n} = m^{x \pmod{\phi(n)}} \pmod{n}$

Encryption of messages:

$c = m^e \pmod{n}$, where c is the ciphertext and m is the plaintext.

Decryption of messages:

$m' = c^d \pmod{n}$, where d is the multiplicative inverse of e .

We will show: $c = m \pmod{n}$ for the chosen e .

$$\begin{aligned} c &= m^e \pmod{n} \\ c &= m^{1+\text{lcm}(p-1,q-1)} \pmod{n} \\ c &= m^1 \cdot m^{k \cdot (p-1) \cdot (q-1)} \pmod{n} \\ c &= m^1 \cdot m^{[k \cdot \phi(n)]} \pmod{\phi(n)} \pmod{n} \\ c &= m^1 \cdot m^0 = m \pmod{n} \end{aligned}$$

□

Example: Fixed point property for all m :

Given $n = p \cdot q = 13 \cdot 37 = 481$

$$\Rightarrow \phi(n) = (p-1)(q-1) = 12 \cdot 36 = 432$$

$$\Rightarrow e = \text{lcm}(p-1, q-1) + 1 = \text{lcm}(12, 36) + 1 = 36 + 1 = 37.$$

With $m \in \{4, 6, 7, 480\}$ we get in $m^e \bmod n$ as:

$$4^{37} \bmod 481 = 4$$

$$6^{37} \bmod 481 = 6$$

$$7^{37} \bmod 481 = 7$$

$$480^{37} \bmod 481 = 480$$

There is not just the one single e (see above), where all $m \in \{1, \dots, n-1\}$ have the fixed point property $m^e = m \bmod n$.¹²⁹

Theorem 4.19.4. *The complete fixed point property of all m is valid for every $e = j \cdot \text{lcm}(p-1, q-1) + 1$, where $j = 0, 1, 2, 3, 4, \dots$ to $e \leq \phi(n)$.*

Example: Further values for e with fixed point properties:

Given $n = p \cdot q = 13 \cdot 37 = 481$ with $\text{lcm}(p-1, q-1) = \text{lcm}(12, 36) = 36$.

Then, e can have the following values: $e = j \cdot \text{lcm}(p-1, q-1) + 1$ for $j = 0, 1, 2, \dots, 11$:

$$\Rightarrow e \in \{1, 37, 73, 109, 145, 181, 217, 253, 289, 325, 361, 397\}.$$

Starting $j = 12$, the following is valid: $e = 12 \cdot \text{lcm}(12, 36) + 1 = 432 + 1 = 433 > 432 = \phi(n)$.

Checking the four values above for m with $e = 217$, the results are:

$$4^{217} \bmod 481 = 4$$

$$6^{217} \bmod 481 = 6$$

$$7^{217} \bmod 481 = 7$$

$$480^{217} \bmod 481 = 480$$

Theorem 4.19.5. *The number of possible values for e with $m^e = m \bmod n$ may be computed with the following:*

$$[\text{Quantity } e] = \left\lfloor \frac{\phi(n)}{\text{lcm}(p-1, q-1) + 1} \right\rfloor + 1 = \frac{\phi(n)}{\text{lcm}(p-1, q-1)}$$

In our example, this results in $\frac{432}{\text{lcm}(12, 36)} = 12$ different values for e , where $m^e = m \bmod n$ for all m in \mathbb{Z}_{481} .

4.19.7.4 An empirical estimate of the quantity of fixed points for growing moduli

In this section, we make an empirical estimate of the quantity of fixed points for growing moduli (and e not weak).

¹²⁹Sometimes these e , which make any message to a fixed point, are called a “weak key” (e, n) of the RSA algorithm. This notation is different to the “weak keys” k in DES, where **every** message m relates to itself if the encryption is done twice. To my knowledge, for larger n the RSA procedure does not have weaks in this meaning: $(m^e)^e = m$.

For this, we randomly choose p and q from the six following ranges each characterized by its lower and upper bound: $(2^2, 2^{10})$, $(2^{10}, 2^{20})$, $(2^{20}, 2^{40})$, $(2^{40}, 2^{80})$, $(2^{80}, 2^{160})$, $(2^{160}, 2^{320})$.

10 attempts were made for each range. For the exponent e , the standard value $e = 2^{16} + 1$ was always chosen. The quantity of fixed points for all 60 attempts was computed with the program 4.18 “Getfixpoints.sage”.

The following five sets contain the randomly chosen value pairs (p,q) of the first five ranges.

$$From(2^2, 2^{10}) : (p, q) \in \{(127, 947), (349, 809), (47, 461), (587, 151), (19, 23), (709, 509), (653, 11), (859, 523), (823, 811), (83, 331)\}$$

$$From(2^{10}, 2^{20}) : (p, q) \in \{(447401, 526283), (474223, 973757), (100829, 126757), (35803, 116933), (577751, 598783), (558121, 607337), (950233, 248167), (451103, 73009), (235787, 164429), (433267, 287939)\}$$

$$From(2^{20}, 2^{40}) : (p, q) \in \{(58569604997, 321367332149), (286573447351, 636576727223), (134703821971, 134220414529), (161234614601, 711682765579), (19367840881, 804790726361), (932891507377, 521129503333), (337186437739, 426034644493), (986529569219, 604515928397), (276825557171, 654134442649), (639276602353, 1069979301731)\}$$

$$From(2^{40}, 2^{80}) : (p, q) \in \{(667530919106151273090539, 287940270633610590682889), (437090557112369481760661, 590040807609821698387141), (1131921188937480863054851, 813935599673320990215139) (874130181777177966406673, 632270193935624953596331), (599303355925474677078809, 717005631177936134003029), (752829320004631398659063, 714134510643836818718761), (1046313315092743492917349, 835721729660755006973833), (877161707568112212806617, 42831503328261105793649), (575464819450637793425803, 5425832051159043433027), (321404337099945148592363, 992663778486687980443879)\}$$

$From(2^{80}, 2^{160}) : (p, q) \in \{(838952969674957834783403492645269831354775774659, 694309130163549038783972189350416942879771871411), (981985107290629501374187748859961786804311564643, 178616495258601001174141825667078950281544628693), (614446632627716919862227545890890553330513965359, 761232454374959264696945191327265643178491649141), (1421756952722008095585945863962560425554707936337, 986781711714138924140285492105143175328486228197), (862346475785474165539441761205023498091366178341, 438589995804600940885415547506719456975478582911), (1034081318899669345416602574034081247538053001533, 1207032778571434704618111297072774884748706223447), (308083812465705343620096534684980088954958466893, 350597371862294596793629011464584694618569736021), (830376326124356299120963861338027196931951857769, 924874232653136669722297184352059466357375363191), (85600581120154590810189237569820706006659829231, 297064381842806596646150718828138629443319259829), (1358984492013516052055790129324581847590275909129, 609402294805414245544586792657989060761523960427)\}$

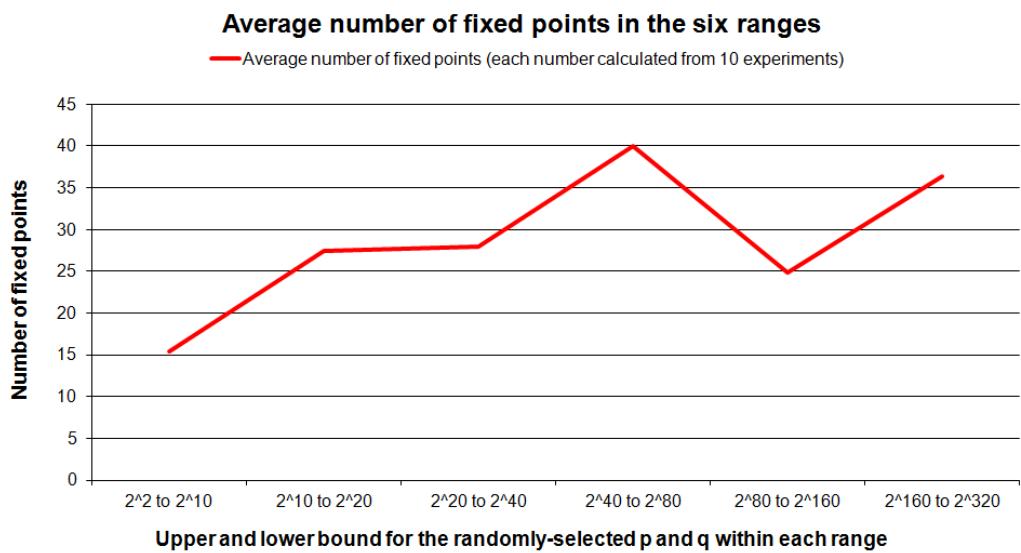


Figure 4.6: An empirical estimate of the quantity of fixed points for growing moduli

Figure 4.6 shows that within the six ranges of size, the average number of fixed points was not higher than 40.

4.19.7.5 Example: Determining all fixed points for a specific public RSA key

The exercise is to determine all fixed points for $(n, e) = (866959, 17)$.

Solution:

We start by factoring n : $866959 = 811 \cdot 1069$.

The quantity of RSA fixed points results from the theorem 4.19.3:

$$(gcd(p-1, e-1)+1) \cdot (gcd(q-1, e-1)+1) = (gcd(811-1, 17-1)+1) \cdot (gcd(1069-1, 17-1)+1) = (2+1) \cdot (4+1) = 15$$

Sage program 4.18 “Getfixpoints” returns the following 15 fixed points for $(n, e) = (866959, 17)$:

0	1	23518	23519	47037
188964	212482	236000	654477	843440
843441	630959	677995	819922	866958

Example:

Confirm that $m = 843441$ is actually a fixed point: $843441^{17} \bmod 866959 = 843441$

Meaning of the Variables in the Sage Code 4.18:

- `gen_f_p = r.multiplicative_generator()`
`r` is a residue class ring modulo p and `multiplicative_generator()` returns a generator element that was created by the ring modulo p .
- `power_mod(gen_f_p, Integer(i*(p-1)/gcd_p), p)`
The `power_mod` function raises a number m to the power of e and returns the results modulo n .
E.g.: `power_mod(m, e, n) := m^e \bmod n`
- `numpy.append(fp, power_mod(gen_f_p, Integer(i*(p-1)/gcd_p), p))`
The `append` function extends an array (`fp`) by an additional element.
- `crt(Integer(r), Integer(s), Integer(p), Integer(q))`
`CRT` is the acronym for Chinese Remainder Theorem. `crt(r, s, p, q)` solves the congruences $x = r \bmod p$ and $x = s \bmod q$ with the help of the Chinese Remainder Theorem.

Sage sample 4.18 Determining all Fixed Points for a Specific Public RSA Key

```
import numpy

print "--- Search for fixpoints in Textbook-RSA given p, q, e ---";
fp=numpy.array([0])
fq=numpy.array([0])

#Edit e,p,q here
###EDIT BEGIN###
e=17;
p=811;
q=1069;
###EDIT END###

n=p*q;
print "Prime p: ",p;
print "Prime q: ",q;
print "Modul n: ",n;
print "Public exponent e: ", e;

r=Integer(p)
gen_f_p = r.multiplicative_generator(); print "\nGenerator of f_p: ",gen_f_p;
s=Integer(q)
gen_f_q = s.multiplicative_generator(); print "Generator of f_q: ",gen_f_q;

gcd_p = gcd(e-1,p-1)
gcd_q = gcd(e-1,q-1)
print "\ngcd(e-1,p-1): ", gcd_p;
print "gcd(e-1,q-1): ", gcd_q;

print "\nNumber of fixpoints: ",(gcd_p+1)*(gcd_q+1);
#Calculating fixpoints modulo F_p
#run i from 0 until gcd(e-1,p-1):
#g^( i*(p-1) / (gcd(e-1,p-1)) ) mod p

print "\nFixpoints modulo p";
print "0 (trivial fixpoint added manually)";
i=0;
for i in range(gcd_p):
    fix_p = power_mod(gen_f_p,Integer(i*(p-1)/gcd_p),p); print fix_p;
    fp = numpy.append(fp,fix_p)

print "\nFixpoints modulo q";
print "0 (trivial fixpoint added manually)";
j=0;
for j in range(gcd_q):
    fix_q = power_mod(gen_f_q,Integer(j*(q-1)/gcd_q),q); print fix_q;
    fq = numpy.append(fq,fix_q);

print "\nFixpoints for the public RSA key (n,e) = (", n, ", ", e, ")"
for r in fp:
    for s in fq:
        print crt(Integer(r),Integer(s),Integer(p),Integer(q))

print "\nRemark: You can verify each fixpoint with power_mod(m,e,n).";
```

4.20 Appendix: List of the definitions and theorems formulated in this chapter

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Here, everything revolves around Fibonacci numbers.
<http://www.mcs.surrey.ac.uk/personal/R.Knott/Fibonacci/fib.html>
2. CrypTool,
Open source e-learning software to illustrate cryptography and cryptanalysis
<http://www.cryptool.de>,
<http://www.cryptool.org>,
<http://www.cryptool.com>
3. Mathematica,
Commercial mathematics package
<http://www.wolfram.com>
4. LiDIA,
Extensive library containing number-theory functions and the LC interpreter. Maintenance stopped.
<http://www.informatik.tu-darmstadt.de/TI/LiDIA>
5. BC,
Interpreter with number-theory functions
<http://www.gnu.org/software/bc/bc.html>
6. Pari-GP,
Excellent, fast, free interpreter with number theoretical functions.
<http://pari.math.u-bordeaux.fr/>
<http://en.wikipedia.org/wiki/PARI/GP>
Resources for PARI/GP at Karim Belabas's website:
<http://www.math.u-bordeaux.fr/~belabas/pari/>
7. Only after I had completed this article, did I come across the website of Mr. Münchenbach, which interactively and didactically uses elementary number theory to provide a sophisticated description of the fundamental mathematical thought processes. It was created for a teaching project in the 11th grade of the technical grammar school (unfortunately only available in German):
<http://www.hydrargyrum.de/kryptographie>
8. Web site of Mr. Wagner, who is responsible for the development of the curriculum of computer science in one of the German federal states (Länder). Here you can get hold of a collection of texts and (Java-)programs (available only in German):
<http://www.saar.de/~awa/kryptolo.htm>
9. GISA,
German Information Security Agency
<http://www.bsi.bund.de>
10. Factorization records and challenges,
<http://www.crypto-world.com/>
<http://www.crypto-world.com/FactorWorld.html>, Webseite von Scott Contini
<http://www.loria.fr/~zimmerma/records/factor.html>

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<http://www.rsa.com/rsalabs/node.asp?id=2092>

11. The Cunningham Project,

<http://www.cerias.purdue.edu/homes/ssw/cun/>

12. Sage,

Excellent, open source computer algebra system with Python as script language, used to build the code samples in this chapter. See the introduction in chapter A.7.

<http://www.sagemath.org/>

http://en.wikipedia.org/wiki/Sage_%28mathematics_software%29

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Chapter 5

The Mathematical Ideas behind Modern Cryptography¹

(Oyono R. / Esslinger B./ Schneider J., Sep. 2000; Updates Nov. 2000, Feb. 2003, Apr. 2007, Mar. 2010, Jan. 2013)

I don't know if its getting better, if we change it,
but I know, that we have to change it, if it should become better.

Quote 14: Georg Christoph Lichtenberg²

5.1 One way functions with trapdoor and complexity classes

A **one way function** is a function that can be calculated efficiently, but whose inverse is extremely complicated and practically impossible to calculate.

To put it more precisely: A one way function is a mapping f from a set X to a set Y , such that $f(x)$ can be calculated easily for each element x of X , whereas for (almost) every y from Y it is practically impossible to find an inverse image x (i.e. an x where $f(x) = y$).

An everyday example of a one way function is a telephone book: the function to be performed is to assign a name to the corresponding telephone number. This can be done easily due to the fact that the names are sorted alphabetically. However, the inverse function - assigning a name to a given number - is obviously difficult if you only have a telephone book available.

One way functions play a decisive role in cryptography. Almost all cryptographic terms can be rephrased using the term one way function. Let's take for example public key encryption (asymmetric cryptography):

Each subscriber T to the system is assigned a private key d_T and what is known as a public key e_T . These keys must have the following property (public key property):

¹With the educational tool for number theory NT you can apply some of the methods introduced here (RSA, Rabin, DH, ElGamal) (see learning unit 4.2 and 4.3, pages 9-17/17).

NT can be called in CT1 via the menu path **Indiv. Procedures \ Number Theory Interactive \ Learning Tool for Number Theory**. See appendix A.6.

You can find according functions within the programs CT1, CT2 and JCT: see the list of the included functions within appendix A.1, A.2 and A.3.

²Georg Christoph Lichtenberg, German writer and physicist (1742-1799),
(also see: http://en.wikipedia.org/wiki/Georg_Christoph_Lichtenberg)

For an opponent who knows the public key e_T , it is practically impossible to determine the private key d_T .

In order to construct useful public key procedures, therefore, we look for a one way function that is “easy” to calculate in one direction, but is “difficult” (practically impossible) to calculate in the other direction, provided that a particular piece of additional information (trapdoor) is not available. This additional piece of information allows the inverse to be found efficiently. Such functions are called **trapdoor one way functions**. In the above case, d_T is the trapdoor information.

In this process, we describe a problem as “easy” if it can be solved in polynomial time as a function of the length of the input. If the length of the input is n bits, then the time for calculating the function is proportional to n^a , where a is a constant. We say that the complexity of such problems is $O(n^a)$ [Landau- or Big-O notation].

If you compare two functions 2^n and n^a , where a is a constant, then there always exists a value for n , from which for all further n applies: $n^a < 2^n$. The function n^a has a lower complexity. Sample: for $a = 5$ the following applies: from the length $n = 23$, 2^n is greater than n^5 ; for further n 2^n clearly increases more quickly [$(2^{22} = 4,194,304, 22^5 = 5,153,632)$, $(2^{23} = 8,388,608, 23^5 = 6,436,343)$, $(2^{24} = 16,777,216, 24^5 = 7,962,624)$].

The term “practically impossible” is slightly less precise. In general, we can say that a problem cannot be solved efficiently, if the time required to solve it increases more quickly than the polynomial time as a function of the size of the input. If, for example, the length of the input is n bits and the time required for calculating the function is proportional to 2^n , then the following currently applies: the function practically cannot be calculated for $n > 80$.

In order to develop a public key procedure that can be implemented in practice, it is therefore necessary to discover a suitable trapdoor one way function.

In order to tidy things up among this confusing multitude of possible problems and their complexities, we group problems with similar complexities into classes.

The most important complexity classes are the classes **P** and **NP**:

- The class **P**: This class contains those problems that can be solved in a polynomial amount of time.
- The class **NP**: The definition of this class doesn't look at the time required to solve a problem, but rather at the time required to verify a given solution. The class **NP** consists of those problems for which a given solution can be verified in a polynomial amount of time. Hereby, the term **NP** “non-deterministic” means polynomial and is based on a calculation model, i.e. on a computer that only exists in theory and can “guess” correct solutions non-deterministically then verify them in polynomial time.

The class **P** is contained in the class **NP**. A well-known unsolved problem is the question whether or not $\mathbf{P} \neq \mathbf{NP}$ is true, i.e. whether or not **P** is a true subset. An important property of the class **NP** is that it also contains what are known as “**NP**-complete” problems. These are problems that represent the class **NP** as follows: If a “good” algorithm for such a problem exists, then “good” algorithms exist for all problems from **NP**. In particular: if **P** only contained one complete problem, i.e. if a polynomial solution algorithm existed for this problem, then **P** would be equal to **NP**. In this sense, the **NP**-complete problems are the most difficult problems in **NP**.

Many cryptographic protocols are formed in such a way that the “good” subscribers only have to solve problems from **P**, whereas a perpetrator is faced with problems from **NP**.

Unfortunately, we do not yet know whether one way functions actually exist. However, we can prove that one way functions exist if and only if $\mathbf{P} \neq \mathbf{NP}$ [Balcazar1988, S.63].

Some mathematicians have again and again claimed to have proven this equivalence, but so far the claims have always turned out to be false [Hesselink2001].

A number of algorithms have been suggested for public key procedures. In many cases - although they at first appeared promising - it was discovered that they could be solved polynomially. The most famous failed applicant is the knapsack with trapdoor, suggested by Ralph Merkle [Merkle1978].

5.2 Knapsack problem as a basis for public key procedures

5.2.1 Knapsack problem

You are given n objects G_1, \dots, G_n with the weights g_1, \dots, g_n and the values w_1, \dots, w_n . The aim is to carry away as much as possible in terms of value while restricted to an upper weight limit g . You therefore need to find a subset of $\{G_1, \dots, G_n\}$, i.e. $\{G_{i_1}, \dots, G_{i_k}\}$, so that $w_{i_1} + \dots + w_{i_k}$ is maximised under the condition $g_{i_1} + \dots + g_{i_k} \leq g$.

Such questions are called **NP**-complete problems (not deterministically polynomial) that are difficult to calculate.

A special case of the knapsack problem is:

Given the natural numbers a_1, \dots, a_n and g , find $x_1, \dots, x_n \in \{0, 1\}$ where $g = \sum_{i=1}^n x_i a_i$ (i.e. where $g_i = a_i = w_i$ is selected). This problem is also called a **0-1 knapsack problem** and is identified with $K(a_1, \dots, a_n; g)$.

Two 0-1 knapsack problems $K(a_1, \dots, a_n; g)$ and $K(a'_1, \dots, a'_n; g')$ are called congruent if two co-prime numbers w and m exist in such a way that

1. $m > \max\{\sum_{i=1}^n a_i, \sum_{i=1}^n a'_i\}$,
2. $g \equiv wg' \pmod{m}$,
3. $a_i \equiv wa'_i \pmod{m}$ for all $i = 1, \dots, n$.

Comment:

Congruent 0-1 knapsack problems have the same solutions. No quick algorithm is known for clarifying the question as to whether two 0-1 knapsack problems are congruent.

A 0-1 knapsack problem can be solved by testing the 2^n possibilities for x_1, \dots, x_n . The best method requires $O(2^{n/2})$ operations, which for $n = 100$ with $2^{100} \approx 1.27 \cdot 10^{30}$ and $2^{n/2} \approx 1.13 \cdot 10^{15}$ represents an insurmountable hurdle for computers. However, for special a_1, \dots, a_n the solution is quite easy to find, e.g. for $a_i = 2^{i-1}$. The binary representation of g immediately delivers x_1, \dots, x_n . In general, the a 0-1 knapsack problem can be solved easily if a permutation³ π of $1, \dots, n$ exists with $a_{\pi(j)} > \sum_{i=1}^{j-1} a_{\pi(i)}$. If, in addition, π is the identity, i.e. $\pi(i) = i$ for $i = 1, 2, \dots, n$, then the sequence a_1, \dots, a_n is said to be super-increasing. Crypto procedure 5.1 solves the knapsack problem with a super-increasing sequence in the time of $O(n)$.

³A permutation π of the numbers $1, \dots, n$ is a change in the order in which these numbers are listed. For example, a permutation π of $(1, 2, 3)$ is $(3, 1, 2)$, i.e. $\pi(1) = 3$, $\pi(2) = 1$ and $\pi(3) = 2$.

Crypto Procedure 5.1 Solving knapsack problems with super-increasing weights

```
for i = n to 1 do
    if  $T \geq a_i$  then
         $T := T - s_i$ 
         $x_i := 1$ 
    else
         $x_i := 0$ 
    if  $T = 0$  then
         $X := (x_1, \dots, x_n)$  is the solution.
    else
        No solution exists.
```

5.2.2 Merkle-Hellman knapsack encryption

In 1978, Merkle and Hellman [Merkle1978] specified a public key encryption procedure that is based on “defamiliarizing” the easy 0-1 knapsack problem with a super-increasing sequence into a congruent one with a super-increasing sequence. It is a block ciphering that ciphers an n -bit plaintext each time it runs, see crypto procedure 5.2 for the details.

Crypto Procedure 5.2 Merkle-Hellman (based on knapsack problems)

Let (a_1, \dots, a_n) be super-increasing. Let m and w be two co-prime numbers with $m > \sum_{i=1}^n a_i$ and $1 \leq w \leq m - 1$. Select \bar{w} with $w\bar{w} \equiv 1 \pmod{m}$ the modular inverse of w and set $b_i := wa_i \pmod{m}$, $0 \leq b_i < m$ for $i = 1, \dots, n$, and verify whether the sequence b_1, \dots, b_n is not super-increasing. A permutation $b_{\pi(1)}, \dots, b_{\pi(n)}$ of b_1, \dots, b_n is then published and the inverse permutation μ to π is defined secretly. A sender writes his/her message in blocks $(x_1^{(j)}, \dots, x_n^{(j)})$ of binary numbers n in length, calculates

$$g^{(j)} := \sum_{i=1}^n x_i^{(j)} b_{\pi(i)}$$

and sends $g^{(j)}$, ($j = 1, 2, \dots$).

The owner of the key calculates

$$G^{(j)} := \bar{w}g^{(j)} \pmod{m}, \quad 0 \leq G^{(j)} < m$$

and obtains the $x_{\mu(i)}^{(j)} \in \{0, 1\}$ (and thus also the $x_i^{(j)}$) from

$$\begin{aligned} G^{(j)} &\equiv \bar{w}g^{(j)} = \sum_{i=1}^n x_i^{(j)} b_{\pi(i)} \bar{w} \equiv \sum_{i=1}^n x_i^{(j)} a_{\pi(i)} \pmod{m} \\ &= \sum_{i=1}^n x_{\mu(i)}^{(j)} a_{\pi(\mu(i))} = \sum_{i=1}^n x_{\mu(i)}^{(j)} a_i \pmod{m}, \end{aligned}$$

by solving the easier 0-1 knapsack problems $K(a_1, \dots, a_n; G^{(j)})$ with super-increasing sequence a_1, \dots, a_n .

In 1982, Shamir [Shamir1982] specified an algorithm for breaking the system in polynomial time without solving the general knapsack problem. Len Adleman [Adleman1982] and Jeff

Lagarias [Lagarias1983] specified an algorithm for breaking the twice iterated Merkle-Hellman knapsack encryption procedure in polynomial time. Ernst Brickell [Brickell1985] then specified an algorithm for breaking multiply iterated Merkle-Hellman knapsack encryption procedures in polynomial time. This made this procedure unsuitable as an encryption procedure. It therefore delivers a one way function whose trapdoor information (defamiliarization of the 0-1 knapsack problem) could be discovered by an eavesdropper.

5.3 Decomposition into prime factors as a basis for public key procedures

5.3.1 The RSA procedure^{4,5}

As early as 1978, R. Rivest, A. Shamir, L. Adleman [RSA1978] introduced the most important asymmetric cryptography procedure to date.

Crypto Procedure 5.3 RSA (based on the factorization problem)

Key generation:

Let p and q be two different prime numbers and $N = pq$. Let e be any prime number relative to $\phi(N)$, i.e. $\gcd(e, \phi(N)) = 1$. Using the Euclidean algorithm, we calculate the natural number $d < \phi(N)$, such that

$$ed \equiv 1 \pmod{\phi(N)}.$$

whereby ϕ is the **Euler phi Function**.

The output text is divided into blocks and encrypted, whereby each block has a binary value $x^{(j)} \leq N$.

Public key:

$$N, e.$$

Private key:

$$d.$$

Encryption:

$$y = e_T(x) = x^e \pmod{N}.$$

Decryption:

$$d_T(y) = y^d \pmod{N}.$$

Comment:

The Euler phi function is defined as: $\phi(N)$ is the number of natural numbers that do not have a common factor with N $x \leq N$. Two natural numbers a and b are co-prime if $\gcd(a, b) = 1$.

For the Euler phi function it holds that:

$$\phi(1) = 1, \phi(2) = 1, \phi(3) = 2, \phi(4) = 2, \phi(6) = 2, \phi(10) = 4, \phi(15) = 8.$$

For example, $\phi(24) = 8$, because

$$|\{x < 24 : \gcd(x, 24) = 1\}| = |\{1, 5, 7, 11, 13, 17, 19, 23\}|.$$

⁴Please compare chapters 4.10, ff.

⁵Using CT1 you can gain practical experience with the RSA procedure via the menu **Indiv.Procedures \ RSA Cryptosystem \ RSA Demonstration**.

You can find RSA also in CT2 and JCT.

If p is a prime number, then $\phi(p) = p - 1$.

If we know the various prime factors p_1, \dots, p_k of N , then

$$\phi(N) = N \cdot \left(1 - \frac{1}{p_1}\right) \cdots \left(1 - \frac{1}{p_k}\right).^6$$

Table 5.1 gives the values up to 15. In the case of $N = pq$,

$$\phi(N) = pq(1 - 1/p)(1 - 1/q) = p(1 - 1/p)q(1 - 1/q) = (p - 1)(q - 1).$$

n	$\phi(n)$	The natural numbers that are co-prime to n and less than n .
1	1	1
2	1	1
3	2	1, 2
4	2	1, 3
5	4	1, 2, 3, 4
6	2	1, 5
7	6	1, 2, 3, 4, 5, 6
8	4	1, 3, 5, 7
9	6	1, 2, 4, 5, 7, 8
10	4	1, 3, 7, 9
15	8	1, 2, 4, 7, 8, 11, 13, 14

Table 5.1: Euler phi function

The function e_T is a one way function whose trapdoor information is the decomposition into primes of N .

At the moment, no algorithm is known that can factorize two prime numbers sufficiently quickly for extremely large values (e.g. for several hundred decimal places). The quickest algorithms known today [Stinson1995] factorize a compound whole number N in a time period proportional to $L(N) = e^{\sqrt{\ln(N) \ln(\ln(N))}}$. Some example values can be found in table 5.2.

N	10^{50}	10^{100}	10^{150}	10^{200}	10^{250}	10^{300}
$L(N)$	$1.42 \cdot 10^{10}$	$2.34 \cdot 10^{15}$	$3.26 \cdot 10^{19}$	$1.20 \cdot 10^{23}$	$1.86 \cdot 10^{26}$	$1.53 \cdot 10^{29}$

Table 5.2: $L(N)$ value table [Factorization effort related to the modul length]

To this date, it has not been proved that the problem of breaking RSA is equivalent to the factorization problem. Nevertheless, it is clear that the RSA procedure will no longer be safe if the factorization problem is “solved”.⁷

⁶Further formulas for the Euler phi function are in chapter 4.8.2 “The Euler function”.

⁷In 2000 the authors assumed that values of the order magnitude 100 to 200 decimal places are currently safe. They estimates that the current computer technology indicates that a number with 100 decimal places could be factorized in approximately two weeks at justifiable costs, and using an expensive configuration (e.g. of around 10 million US dollars), a number with 150 decimal places could be factorized in about a year, and a 200-digit number should remain impossible to factorize for a long time to come, unless there is a mathematical breakthrough. However, you can never be sure that there won’t be a mathematical breakthrough tomorrow. How easy it is to guess the future wrong is shown by the factorization of RSA-200 (see chapter 4.11.4) – completely without a “mathematical breakthrough”.

5.3.2 Rabin public key procedure (1979)

The Rabin public key procedure (crypto procedure 5.4) has been shown to be equivalent to breaking the factorization problem. Unfortunately, this procedure is susceptible to chosen-ciphertext attacks.

Crypto Procedure 5.4 Rabin (based on the factorization problem)

Let p and q be two different prime numbers with $p, q \equiv 3 \pmod{4}$ and $n = pq$. Let $0 \leq B \leq n-1$.

Public key:

$$e = (n, B).$$

Private key:

$$d = (p, q).$$

Encryption:

$$y = e_T(x) = x(x + B) \pmod{n}.$$

Decryption:

$$d_T(y) = \sqrt{y + B^2/4} - B/2 \pmod{n}.$$

Caution: Because $p, q \equiv 3 \pmod{4}$ the encryption is easy to calculate (if the key is known). This is not the case for $p \equiv 1 \pmod{4}$. In addition, the encryption function is not injective: There are precisely four different source codes that have $e_T(x)$ as inverse image: $x, -x-B, \omega(x+B/2)-B/2, -\omega(x+B/2)-B/2$, where ω is one of the four roots of unity. The source codes therefore must be redundant for the encryption to remain unique!

Backdoor information is the decomposition into prime numbers of $n = pq$.

5.4 The discrete logarithm as basis for public key procedures⁸

Discrete logarithms form the basis for a large number of algorithms for public-key procedures.

5.4.1 The discrete logarithm in \mathbb{Z}_p

Let p be a prime number and let $g \in \mathbb{Z}_p^* = \{0, 1, \dots, p-1\}$. Then the discrete exponential function base g is defined as

$$e_g : k \longrightarrow y := g^k \pmod{p}, \quad 1 \leq k \leq p-1.$$

The inverse function is called a discrete logarithm function \log_g ; the following holds:

$$\log_g(g^k) = k.$$

The problem of the discrete logarithm (in \mathbb{Z}_p^*) is understood to be as follows:

Given p, g and y , determine k such that $y = g^k \pmod{p}$.

⁸With the educational tool for number theory **NT** you can play with the distribution of the discrete logarithm values and apply Shank's baby-step-giant-step method: See learning units 6.1-6.3, pages 1-6/6.
NT can be called in CT1 via the menu path **Indiv. Procedures \ Number Theory Interactive \ Learning Tool for Number Theory**. See appendix A.6.

It is much more difficult to calculate the discrete logarithm than to evaluate the discrete exponential function (see chapter 4.9). Table 5.3 lists several procedures for calculating the discrete logarithm and their complexity [Stinson1995].

Name	Complexity
Baby-step-giant-step	$O(\sqrt{p})$
Silver-Pohlig-Hellman	polynomial in q , the greatest prime factor of $p - 1$.
Index-Calculus	$O(e^{(1+o(1))\sqrt{\ln(p)\ln(\ln(p))}})$

Table 5.3: Procedures for calculating the discrete logarithm over \mathbb{Z}_p^*

The current record (as of April 2007) for calculating discrete logarithms was established in February 2007 by the group Kleinjung, Franke and Bahr at University of Bonn.⁹ Kleinjung calculated the discrete logarithm modulo a 160 digit prime number p and generator g :

$$\begin{aligned} p &= \lfloor 10^{159}\pi \rfloor + 119849 \\ &= 314159265358979323846264338327950288419716939937510582097494 \\ &\quad 459230781640628620899862803482534211706798214808651328230664 \\ &\quad 7093844609550582231725359408128481237299 \\ g &= 2 \end{aligned}$$

The discrete logarithms k of the following integer y was determined:¹⁰

$$\begin{aligned} y &= \lfloor 10^{159}e \rfloor \\ &= 271828182845904523536028747135266249775724709369995957496696 \\ &\quad 762772407663035354759457138217852516642742746639193200305992 \\ &\quad 1817413596629043572900334295260595630738 \\ k &= \log_g(y) \mod p \\ &= 829897164650348970518646802640757844024961469323126472198531 \\ &\quad 845186895984026448342666252850466126881437617381653942624307 \\ &\quad 537679319636711561053526082423513665596 \end{aligned}$$

The search was performed with GNFS method (General Number Field Sieve, Index-Calculus) and took about 17 CPU years on 3.2 GHz Xeon machines.

5.4.2 Diffie-Hellman key agreement¹¹

The mechanisms and algorithms of classical cryptography only take effect when the subscribers have already exchanged the secret key. In classical cryptography you cannot avoid exchanging secrets without encrypting them. Transmission safety here must be achieved using non-cryptographic methods. We say that we need a secret channel for exchanging secrets. This channel can be realised either physically or organisationally.

What is revolutionary about modern cryptography is, amongst other things, that you no longer

⁹[http://www.nabble.com/Discrete-logarithms-in-GF\(p\)----160-digits-t3175622.html](http://www.nabble.com/Discrete-logarithms-in-GF(p)----160-digits-t3175622.html)

¹⁰The integer y was chosen as the first 159 digits of the Euler number e .

¹¹With CT1 this exchange protocol has been visualized: you can execute the single steps with concrete numbers using menu **Indiv. Procedures \ Protocols \ Diffie-Hellman Demonstration**.

In JCT you can find it in the default perspective via the menu item **Visuals \ Diffie-Hellman Key Exchange (EC)**.

need secret channels: You can agree secret keys using non-secret, i.e. public channels. One protocol that solves this problem is that of Diffie and Hellman (crypto procedure 5.5).

Crypto Procedure 5.5 Diffie-Hellman key agreement

Two subscribers A and B want to agree on a joint secret key.

Let p be a prime number and g a natural number. These two numbers do not need to be secret. The two subscribers then select a secret number a and b from which they calculate the values $\alpha = g^a \pmod p$ and $\beta = g^b \pmod p$. They then exchange the numbers α and β . To end with, the two subscribers calculate the received value to the power of their secret value to get $\beta^a \pmod p$ and $\alpha^b \pmod p$.

Thus

$$\beta^a \equiv (g^b)^a \equiv g^{ba} \equiv g^{ab} \equiv (g^a)^b \equiv \alpha^b \pmod p$$

The safety of the **Diffie-Hellman protocol** is closely connected to calculating the discrete logarithm mod p . It is even thought that these problems are equivalent.

5.4.3 ElGamal public key encryption procedure in \mathbb{Z}_p^*

By varying the Diffie-Hellman key agreement protocol slightly, you can obtain an asymmetric encryption algorithm, crypto procedure 5.6. This observation was made by Taher ElGamal.

Crypto Procedure 5.6 ElGamal (based on the discrete logarithm problem)

Let p be a prime number such that the discrete logarithm in \mathbb{Z}_p^* is difficult to compute. Let $\alpha \in \mathbb{Z}_p^*$ be a primitive element. Let $a \in \mathbb{N}$ and $\beta = \alpha^a \pmod p$.

Public key:

$$p, \alpha, \beta.$$

Private key:

$$a.$$

Let $k \in \mathbb{Z}_{p-1}$ be a random number and $x \in \mathbb{Z}_p^*$ the plaintext.

Encryption:

$$e_T(x, k) = (y_1, y_2),$$

where

$$y_1 = \alpha^k \pmod p$$

and

$$y_2 = x\beta^k \pmod p.$$

Decryption:

$$d_T(y_1, y_2) = y_2(y_1^a)^{-1} \pmod p$$

5.4.4 Generalised ElGamal public key encryption procedure

The discrete logarithm can be generalised in any number of finite groups (G, \circ) . The following provides several properties of G , that make the discrete logarithm problem difficult.

Calculating the discrete exponential function Let G be a group with the operation \circ and $g \in G$. The (discrete) exponential function base g is defined as

$$e_g : k \mapsto g^k, \quad \text{for all } k \in \mathbb{N}.$$

where

$$g^k := \underbrace{g \circ \dots \circ g}_{k \text{ times}}.$$

The exponential function is easy to calculate:

Lemma. *The power g^k can be calculated in at most $2\log_2 k$ group operations.*

Proof

Let $k = 2^n + k_{n-1}2^{n-1} + \dots + k_12 + k_0$ be the binary representation of k . Then $n \leq \log_2(k)$, because $2^n \leq k < 2^{n+1}$. k can be written in the form $k = 2k' + k_0$ with $k' = 2^{n-1} + k_{n-1}2^{n-2} + \dots + k_1$. Thus

$$g^k = g^{2k'+k_0} = (g^{k'})^2 g^{k_0}.$$

We therefore obtain g^k from $g^{k'}$ by squaring and then multiplying by g . The claim is thus proved by induction to n . \square

Problem of the discrete logarithm

Let G by a finite group with the operation \circ . Let $\alpha \in G$ and $\beta \in H = \{\alpha^i : i \geq 0\}$.

We need to find a unique $a \in \mathbb{N}$ with $0 \leq a \leq |H| - 1$ and $\beta = \alpha^a$.

We define a as $\log_\alpha(\beta)$.

Calculating the discrete logarithm A simple procedure for calculating the discrete logarithm of a group element, that is considerably more efficient than simply trying all possible values for k , is the baby-step-giant-step algorithm.

Theorem 5.4.1. [baby-step-giant-step algorithm] *Let G be a group and $g \in G$. Let n be the smallest natural number with $|G| \leq n^2$. Then the discrete logarithm of an element $h \in G$ can be calculated base g by generating two lists each containing n elements and comparing these lists. In order to calculate these lists, we need $2n$ group operations.*

Proof

First create the two lists

Giant-step list: $\{1, g^n, g^{2n}, \dots, g^{n \cdot n}\}$,

Baby-step list: $\{hg^{-1}, hg^{-2}, \dots, hg^{-n}\}$.

If $g^{jn} = hg^{-i}$, i.e. $h = g^{i+jn}$, then the problem is solved. If the lists are disjoint, then h cannot be represented as g^{i+jn} , $i, j \leq n$. As all powers of g are thus recorded, the logarithm problem does not have a solution. \square

You can use the baby-step-giant-step algorithm to demonstrate that it is much more difficult to calculate the discrete logarithm than to calculate the discrete exponential function. If the numbers that occur have approximately 1000 bits in length, then you only need around 2000 multiplications to calculate g^k but around $2^{500} \approx 10^{150}$ operations to calculate the discrete logarithm using the baby-step-giant-step algorithm.

In addition to the baby-step-giant-step algorithm, there are also numerous other procedures for calculating the discrete logarithm [Stinson1995].

The theorem from Silver-Pohlig-Hellman In finite Abelian groups, the discrete logarithm problem can be reduced to groups of a lower order.

Theorem 5.4.2. [Silver-Pohlig-Hellman] Let G be a finite Abelian group with $|G| = p_1^{a_1} p_2^{a_2} \cdot \dots \cdot p_s^{a_s}$. The discrete logarithm in G can then be reduced to solving logarithm problems in groups of the order p_1, \dots, p_s .

If $|G|$ contains a “dominant” prime factor p , then the complexity of the logarithm problem is approximately

$$O(\sqrt{p}).$$

Therefore, if the logarithm problem is to be made difficult, the order of the group used G should have a large prime factor. In particular, if the discrete exponential function in the group \mathbb{Z}_p^* is to be a one way function, then $p - 1$ must be a large prime factor. In this case a generalized ElGamal procedure can be defined (crypto procedure 5.7).

Crypto Procedure 5.7 Generalized ElGamal (based on the factorization problem)

Let G be a finite group with operation \circ , and let $\alpha \in G$, so that the discrete logarithm in $H = \{\alpha^i : i \geq 0\}$ is difficult. Let a with $0 \leq a \leq |H| - 1$ and let $\beta = \alpha^a$.

Public key:

$$\alpha, \beta.$$

Private key:

$$a.$$

Let $k \in \mathbb{Z}_{|H|}$ be a random number and $x \in G$ be a plaintext.

Encryption:

$$e_T(x, k) = (y_1, y_2),$$

where

$$y_1 = \alpha^k$$

and

$$y_2 = x \circ \beta^k.$$

Decryption:

$$d_T(y_1, y_2) = y_2 \circ (y_1^a)^{-1}$$

Elliptic curves provide useful groups for public key encryption procedures.

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Chapter 6

Hash Functions and Digital Signatures

(Schneider J. / Esslinger B. / Koy H., June 2002; Updates: Feb. 2003, June 2005, July 2009, Nov. 2012)

We can make everything out of this world, but we cannot create a world, where humans in some ten thousand years can think: 'Ok, now it is enough. Everything should stay like it is. Let's do no changes any more, don't do inventions any more, because it cannot become better, and if, then we don't want this.'

Quote 15: Stanislaw Lem¹

The aim of digital signatures is to guarantee the following two points:

- User authenticity:
It can be checked whether a message really does come from a particular person.
- Message integrity:
It can be checked whether the message has been changed (on route).

An asymmetric technique is used again (see encryption procedures). Participants who wish to generate a digital signature for a document must possess a pair of keys. They use their secret key to generate signatures and the recipient uses the sender's public key to verify whether the signature is correct. As before, it must be impossible to use the public key to derive the secret key.

In detail, a *Signature procedure*² looks like this:

Senders use their message and secret key to calculate the digital signature for the message. Compared to hand-written signatures, digital signatures therefore have the advantage that they also depend on the document to be signed. Signatures from one and the same participant are different unless the signed documents are completely identical. Even inserting a blank in the

¹This was the answer of Stanislaw Lem to heavy critics at his philosophical main book "Summa Technologiae", 1964, where he thought about the possibility of an evolution creating artificial intelligence.

²With CT1 you can also generate and check digital signatures: Using the submenus of the main menu **Digital Signatures / PKI** or using menu **Indiv. Procedures \ RSA Cryptosystem \ Signature Demonstration (Signature Generation)**. Also with JCT (in the default and the algorithm perspektive) its possible to create different kinds of electronic signatures.

text would lead to a different signature. The recipient of the message would therefore detect any injury to the message integrity as this would mean that the signature no longer matches the document and is shown to be incorrect when verified.

The document is sent to the recipient together with the signature. The recipient can then use the sender's public key, the document and the signature to establish whether or not the signature is correct. The procedure we just described has in practice, however, a decisive disadvantage. The signature would be approximately as long as the document itself. To prevent an unnecessary increase in data traffic, and also for reasons of performance, we apply a cryptographic hash function³ to the document – before signing. The output of the hash function will then be signed.

6.1 Hash functions

A *hash function*⁴ maps a message of any length to a string of characters with a constant size, the hash value.

6.1.1 Requirements for hash functions

Cryptographically secure hash functions fulfill the following three requirements (the order is in a way that the requirements increase):

- Resistance against 1st pre-image attacks:

It should be practically impossible, for a given number, to find a message that has precisely this number as hash value.

Given (fix): hash value H' ,

Searched: message m , so that: $H(m) = H'$.

- Resistance against 2nd pre-image attacks:

It should be practically impossible, for a given message, to find another message, which has precisely the same hash value.

Given (fix): message m_1 [and so the hash value $H_1 = H(m_1)$],

Searched: message m_2 , so that: $H(m_2) = H_1$.

- Collision resistance:

It should be practically impossible to find any two messages with the same hash value (it

³Hash functions are implemented within CT1 at several places.

Using menus **Individual Procedures \ Hash** and **Analysis \ Hash** you have the possibilities

- to apply one of 6 hash functions to the content of the current window,
- to calculate the hash value of a file,
- to test, how changes to a text change the according hash value,
- to calculate a key from a password according to the PKCS#5 standard,
- to calculate HMACs from a text and a secret key, and
- to perform a simulation, how digital signatures could be attacked by a targeted search for hash value collisions.

CT2 and JCT also contain different hash methods: See the functions' lists within the appendix A.2 and A.3.

⁴Hash algorithms compute a condensed representation of electronic data (message). When a message is input to a hash algorithm, the result is an output called a message digest. The message digests typically range in length from 128 to 512 bits, depending on the algorithm. Secure hash algorithms are typically used with other cryptographic algorithms, such as digital signature algorithms and keyed-hash message authentication codes, or in the generation of random numbers (bits).

doesn't matter what hash value).

Searched: 2 messages m1 and m2, so that: $H(m1) = H(m2)$.

6.1.2 Current attacks against hash functions like SHA-1

So far, no formal proof has been found that perfectly secure cryptographic hash functions exist.

During the past several years no new attacks against hash algorithms came up, and so the candidates that had not yet shown any weaknesses in their structure in practice (e.g. SHA-1⁵ or RIPEMD-160⁶) were trusted.

At Crypto 2004 (August 2004)⁷ this safety-feeling was disputed: Chinese researchers published collision attacks against MD4, SHA-0 and parts of SHA-1. This globally caused new motivation to engage in new hash attack methods.

The initially published result reduced the expected complexity for one SHA-1 collision search from 2^{80} (brute-force) to 2^{69} [Wang2005]. More recent announcements claim to further reduce the required effort to 2^{63} [Wang2005b] and 2^{52} [McDonald2009]. This would bring collision attacks into the practical realm, as similar efforts have been mastered in the past (s. 1.2.2).

According to our current knowledge there is no need to run scared. But in the future digital signatures should use longer hash values and/or other hash algorithms.

Already before Crypto 2004 the U.S. National Institute of Standards and Technology (NIST) announced, to discontinue SHA-1 in the next few years. So it is recommended not to use SHA-1 for new products generating digital signatures. The SHA-2 family [FIPS180-3] provides stronger algorithms. To address new findings in cryptanalysis, NIST has opened a competition to develop a new cryptographic hash algorithm in 2008: In October 2012 Keccak was announced as "SHA-3"⁸.

Further information about this topic can be found in the article "Hash cracked – The consequences of the successful attacks on SHA-1" by Reinhard Wobst and Jürgen Schmidt⁹ by Heise Security.

⁵SHA-1 is a 160 bit hash function specified in FIPS 180-1 (by NIST), ANSI X9.30 Part 2 and [FIPS186].

SHA means Secure Hash Algorithm, and is widely used, e.g. with DSA, RSA or ECDSA.

The current standard [FIPS180-3] defines four secure hash algorithms – SHA-1, SHA-256, SHA-384, and SHA-512. For these hash algorithms there are also validation tests defined in the test suite FIPS 140-2.

The output length of the SHA algorithms was enhanced because of the possibility of birthday attacks: these make n-bit AES and a 2n-bit hash roughly equivalent:

- 128-bit AES – SHA-256
- 192-bit AES – SHA-384
- 256-bit AES – SHA-512.

With CT1 you can comprehend the birthday attack on digital signatures:

using the menu **Analysis \ Hash \ Attack on the Hash Value of the Digital Signature**.

CT2 contains an MD5 collider.

⁶RIPEMD-160, RIPEMD-128 and the optional extension RIPEMD-256 have object identifiers defined by the ISO-identified organization TeleTrusT, both as hash algorithm and in combination with RSA. RIPEMD-160 is also part of the ISO/IEC international standard ISO/IEC 10118-3:1998 on dedicated hash functions, together with RIPEMD-128 and SHA-1. Further details:

- <http://www.esat.kuleuven.ac.be/~bosselaer/ripemd160.html>
- <http://www.ietf.org/rfc/rfc2857.txt> ("The Use of HMAC-RIPEMD-160-96 within ESP and AH").

⁷<http://www.iacr.org/conferences/crypto2004/>

⁸<http://csrc.nist.gov/groups/ST/hash/sha-3/>

⁹<http://www.heise.de/security/artikel/56634>.

Further references are e.g.:

<http://www.bsi.bund.de/esig/basics/techbas/krypto/index.htm>

<http://csrc.nist.gov/CryptoToolkit/tkhash.html>.

6.1.3 Signing with hash functions

“Manipulation was Sobol’s speciality ... the main investigation should branch off and try to discover Sobol’s master plan.”

Quote 16: Daniel Suarez¹⁰

The signature procedure with hash functions¹¹ is as follows: Rather than signing the actual document, the sender now first calculates the hash value of the message and signs this. The recipient also calculates the hash value of the message (the algorithm used must be known), then verifies whether the signature sent with the message is a correct signature of the hash value. If this is the case, the signature is verified to be correct. This means that the message is authentic, because we have assumed that knowledge of the public key does not enable you to derive the secret key. However, you would need this secret key to sign messages in another name.

Some digital signature schemes are based on asymmetric *encryption* procedures, the most prominent example being the RSA system, which can be used for signing by performing the private key operation on the hash value of the document to be signed.

Other digital signature schemes were developed exclusively for this purpose, as the DSA (Digital Signature Algorithm), and are not directly connected with a corresponding encryption scheme.

Both, RSA and DSA signature are discussed in more detail in the following two sections. After that we go one step further and show how digital signatures can be used to create the digital equivalent of ID cards. This is called Public Key Certification.

6.2 RSA signatures

As mentioned in the comment at the end of section 4.10.3 it is possible to perform the RSA private and public key operation in reverse order, i. e. raising M to the power of d and then to the power of e ($\bmod N$) yields M again. Based on this simple fact, RSA can be used as a signature scheme.

The RSA signature S for a message M is created by performing the private key operation:

$$S \equiv M^d \pmod{N}$$

In order to verify, the corresponding public key operation is performed on the signature S and the result is compared with message M :

$$S^e \equiv (M^d)^e \equiv (M^e)^d \equiv M \pmod{N}$$

If the result matches the message M , then the signature is accepted by the verifier, otherwise the message has been tampered with, or was never signed by the holder of d .

As explained above, signatures are not performed on the message itself, but on a cryptographic hash value of the message. To prevent certain attacks on the signature procedure (alone or in combination with encryption) it is necessary to format the hash value before doing the exponentiation, as described in the PKCS#1 (Public Key Cryptography Standard #1 [PKCS1]).

¹⁰Daniel Suarez, “Daemon”, Dutton Adult, 2010, Chapter 14, “Meme Payload”, p. 142, Ross.

¹¹Compare: http://en.wikipedia.org/wiki/Digital_signature.

The fact that this standard had to be revised recently, after being in use for several years, can serve as an example of how difficult it is to get the details of cryptography right.

6.3 DSA signatures

In August of 1991, the U.S. National Institute of Standards and Technology (NIST) proposed a digital signature algorithm (DSA), which was subsequently adopted as a U.S. Federal Information Processing Standard (FIPS 186 [FIPS186]).

The algorithm is a variant of the ElGamal scheme. Its security is based on the Discrete Logarithm Problem. The DSA public and private key and its procedures for signature and verification are summarised in crypto procedure 6.1.

Crypto Procedure 6.1 DSA signature

Public Key

p prime

q 160-bit prime factor of $p - 1$

$g = h^{(p-1)/q} \pmod{p}$, where $h < p - 1$ and $h^{(p-1)/q} > 1 \pmod{p}$

$y \equiv g^x \pmod{p}$

Remark: Parameters p, q and g can be shared among a group of users.

Private Key

$x < q$ (a 160-bit number)

Signing

m the message to be signed

k choose at random, less than q

$r = (g^k \pmod{p}) \pmod{q}$

$s = (k^{-1}(\text{SHA-1}(m) + xr)) \pmod{q}$

Remark:

- (s, r) is the signature.
- The security of the signature depends not only on the mathematical properties, but also on using a good random source for k .
- SHA-1 is a 160-bit hash function.

Verifying

$w = s^{-1} \pmod{q}$

$u_1 = (\text{SHA-1}(m)w) \pmod{q}$

$u_2 = (rw) \pmod{q}$

$v = (g^{u_1}y^{u_2}) \pmod{p}$

Remark: If $v = r$, then the signature is verified.

While DSA was specifically designed, so that it can be exported from countries regulating export of encryption soft and hardware (like the U.S. at the time when it was specified), it has been noted [Schneier1996, p. 490], that the operations involved in DSA can be used to emulate RSA and ElGamal encryption.

6.4 Public key certification

The aim of public key certification is to guarantee the connection between a public key and a user and to make it traceable for external parties. In cases in which it is impossible to ensure that a public key really belongs to a particular person, many protocols are no longer secure, even if the individual cryptographic modules cannot be broken.

6.4.1 Impersonation attacks

Assume Charlie has two pairs of keys (PK_1, SK_1) and (PK_2, SK_2), where SK denotes the secret key and PK the public key. Further assume that he manages to palm off PK_1 on Alice as Bob's public key and PK_2 on Bob as Alice's public key (by falsifying a public key directory).

Then he can attack as follows:

- Alice wants to send a message to Bob. She encrypts it using PK_1 because she thinks that this is Bob's public key. She then signs the message using her secret key and sends it.
- Charlie intercepts the message, removes the signature and decrypts the message using SK_1 . If he wants to, he can then change the message in any way he likes. He then encrypts the message again, but this time using Bob's genuine public key, which he has taken from a public key directory, signs the message using SK_2 and forwards it to Bob.
- Bob verifies the signature using PK_2 and will reach the conclusion that the signature is correct. He then decrypts the message using his secret key.

In this way Charlie can listen in on communication between Alice and Bob and change the exchanged messages without them noticing. The attack will also work if Charlie only has one pair of keys.

Another name for this type of attack is “man-in-the-middle attack”. Users are promised protection against this type of attack by public-key certification, which is intended to guarantee the authenticity of public keys. The most common certification method is the X.509 standard.

6.4.2 X.509 certificate

Each participant who wants to have an X.509 certificate ([X.509]) verifying that his public key belongs to a real person consults what is known as a certification authority (CA)¹². He proves his identity to this CA (for example by showing his ID). The CA then issues him an electronic document (certificate) which essentially contains the name of the certificate-holder and the name of the CA, the certificate-holder's public key and the validity period of the certificate. The CA then signs the certificate using its secret key.

Anyone can now use the CA's public key to verify whether a certificate is falsified. The CA therefore guarantees that a public key belongs to a particular user.

This procedure is only secure as long as it can be guaranteed that the CA's public key is correct. For this reason, each CA has its public key certified by another CA that is superior in the hierarchy. In the upper hierarchy level there is usually only one CA, which can of course then have its key certified by another CA. It must therefore transfer its key securely in another way. In the case of many software products that work with certificates (such as the Microsoft

¹²Often called trust center, if the certificates are not only offered to a closed user group.

and Netscape Web browsers), the certificates of these root CAs are permanently embedded in the program right from the start and cannot be changed by users at a later stage. However, (public) CA keys, in particularly those of the root entity, can also be secured by means of making them available publicly.

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Chapter 7

Elliptic Curves

(Filipovics B. / Büger M. / Esslinger B. / Oyono R., April 2000, Updates: Dec. 2001, June 2002, Mar. 2003, November 2009)

7.1 Elliptic curve cryptography – a high-performance substitute for RSA?

In many business sectors secure and efficient data transfer is essential. In particular, the RSA algorithm is used in many applications. Although the security of RSA is beyond doubt, the evolution in computing power has caused a growth in the necessary key length. Today, 1024-bit RSA keys are standard, but the GISA (German Information Security Agency) recommends the usage of 2048-bit keys from 2006 on (compare section 4.11). The fact that most chips on smart cards cannot process keys extending 1024 bit shows that there is a need for alternatives. Elliptic curve cryptography (ECC) can be such an alternative in the area of asymmetric cryptography.

The efficiency of a cryptographic algorithm depends on the key length and the calculation effort that is necessary to provide a prescribed level of security. The major advantage of ECC compared to RSA is that it requires much shorter key lengths. If we assume that the computing power increases by Moore's law (i. e. it doubles every 18 months)¹, then the evolution of the key lengths for secure communication will be as figure 7.1 [Lenstra1999] (source: Arjen Lenstra and Eric Verheul: <http://cryptosavvy.com/table.htm>).

In addition, a digital signature can be processed 10-times faster with ECC than with RSA. However, verification of a given signature is still more efficient with RSA than with ECC. Refer to figure 7.2 (source: Dr. J. Merkle, Elliptic Curve Cryptography Workshop, 2001) for a comparison. The reason is that RSA public keys can be chosen relatively small as long as the secret key is long enough.

Nevertheless, thin clients like smart cards usually have to store the (long) secret key and have to process a digital signature rather than verify one. Therefore, there is a clear advantage in using ECC in terms of efficiency.

Nowadays, the major problem with ECC implementations is the lack of standardization. There is only one way to implement RSA, but there are many ways for ECC: One can work with different sets of numbers, different (elliptic) curves — described by parameters² — , and

¹empirical knowledge by Gordon Moore, co-founder of Intel, 1965

²see chapter 7.4

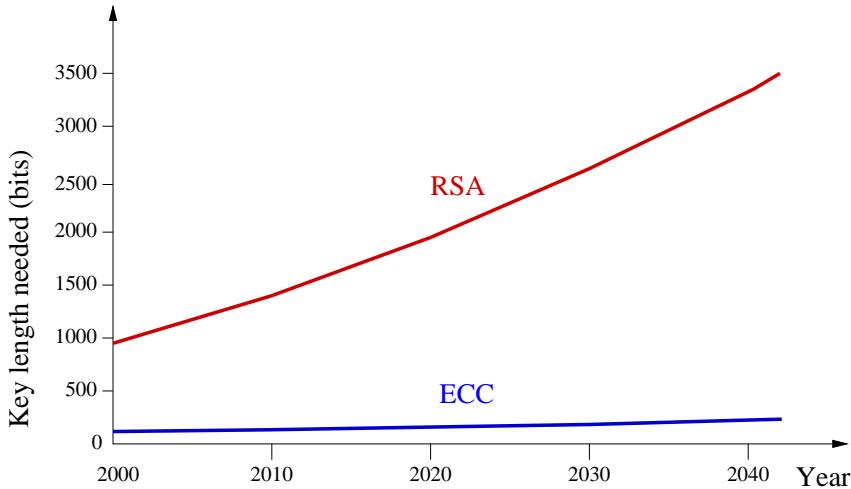


Figure 7.1: Prognosis of the key lengths to be regarded safe for RSA and Elliptic Curves

a variety of representations of the elements on the curve. Each choice has its advantages and disadvantages, and one can certainly construct the most efficient for each application. However, this causes problems in interoperability. But if all ECC-tools should be able to communicate with each other, they will have to support all different algorithms, which might put the advantage of efficient computation and the need of less storage capacity to the contrary.

Therefore, international standardization organizations like IEEE (P1363), ASC (ANSI X9.62, X9.63), ISO/IEC as well as major players like RSA labs or Certicom have recently started standardization initiatives. While the IEEE only describes the different implementations, the ASC has explicitly stated 10 elliptic curves and recommends their usage. The advantage of the ASC approach is that one needs only a single byte to indicate which curve is meant. However, it is not yet clear whether the ASC curves will become a de facto standard.

Although we see no need to replace RSA in any application today³, one should take the usage of ECC-based tools into consideration whenever a new system is set up — in particular, when the tool should be available beyond 2005⁴.

7.2 Elliptic curves – history

Mathematicians have been researching elliptic curves for over 100 years. Over the course of time, many lengthy and mathematically complex results have been found and published which are connected to elliptic curves. A mathematician would say that elliptic curves (or the mathematics behind them) are widely understood. This research was originally purely mathematical. That is to say, elliptic curves were investigated, for example, in the mathematical areas of number theory and algebraic geometry, which are generally highly abstract. Even in the recent past, elliptic curves played an important role in pure mathematics. In 1993 and 1994, Andrew Wiles published mathematical works that triggered enthusiasm far beyond the specialist audience. In these works, he proved a conjecture put forward in the 1960's. To put it short, this conjecture was concerned with the connection between elliptic curves and what are called module forms.

³Current information about the security of the RSA algorithm can be found in chapter 4.11.

⁴Compare the recommendation of GISA: “Fitting Crypto Algorithms” from October 24th, 2002.

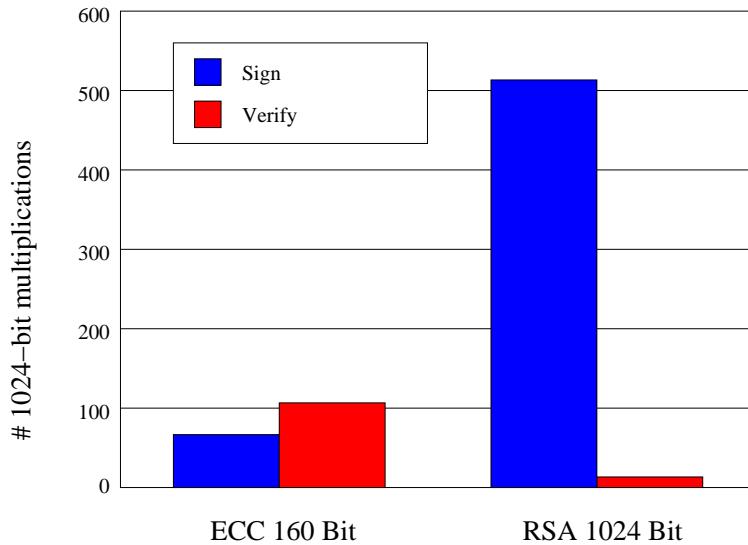


Figure 7.2: Comparison of signing and verification time for RSA and Elliptic Curves

What is particularly interesting for most people is that the works of Wiles also proved the famous second theorem of Fermat. Mathematicians had spent centuries (Fermat lived from 1601 to 1665) trying to find a strict proof of this theorem. Understandably, therefore, Wiles' proof got a good response. Fermat formulated his theorem as follows (written in the border of a book):

Cubum autem in duos cubos, aut quadratoquadratum in duos quadratoquadratos, et generaliter nullam in infinitum ultra quadratum potestatem in duos ejusdem nominis fas est dividere: cuius rei demonstrationem mirabilem sane detexi. Hanc marginis exiguitas non caperet.

With a free translation, using the denotation of modern mathematics, this means:
No positive whole numbers x, y and z greater than zero exist such that $x^n + y^n = z^n$ for $n > 2$. I have found an amazing proof of this fact, but there is too little space within the confines of this book to include it.

This is truly amazing: A statement that is relatively simple to understand (we are referring to Fermat's second theorem here) could only be proved after such a long period of time, although Fermat himself claimed to have found a proof. What's more, the proof found by Wiles is extremely extensive (all of Wiles publications connected with the proof made up a book in themselves). This should therefore make it obvious that elliptic curves are generally based on highly complex mathematics.

Anyway that's enough about the role of elliptic curves in pure mathematics. In 1985 Neal Koblitz and Victor Miller independently suggested using elliptic curves in cryptography. Elliptic curves have thus also found a concrete practical application. Another interesting area of application for elliptic curves is for factorizing whole numbers (the RSA cryptographic system is based on the difficulty/complexity of finding prime factors of an extremely large number;

compare section 4.11.). In this area, procedures based on elliptic curves have been investigated and used since 1987 (compare section 7.8).

There are also prime number tests based on elliptic curves.

Elliptic curves are used differently in the various areas. Encryption procedures based on elliptic curves are based on the difficulty of a problem known as elliptic curve discrete logarithm. The factorization of whole numbers uses the fact that a large number of elliptic curves can be generated for a natural composite number n with several prime factors; however, these curves are not then groups for composite n . More information about this can be found under the chapter 7.8.

7.3 Elliptic curves – mathematical basics

This section provides information about *groups* and *fields*.

7.3.1 Groups

Because the term *group* is used differently in everyday language than in mathematics, we will, for reasons of completeness, begin by introducing the essential statement of the formal definition of a group:

- A group is a non-empty set G on which an operation “.”. The set G is closed under this operation, which means that for any two elements a, b taken from G , performing the operation on them gives an element in G , i.e. $ab = a \cdot b$ lies in G .
- For all elements a, b and c in G : $(ab)c = a(bc)$ (associative law).
- There exists an element e in G that behaves neutrally with respect to the operation \cdot . That means that for all a in the set G : $ae = ea = a$.
- For each element a in G there exists a so-called inverse⁵ element a^{-1} in G such that: $aa^{-1} = a^{-1}a = e$.

If also $ab = ba$ (commutative law) for all a, b in G , then we call the group an *Abelian* group.

Since we may define different operations on the same set, we distinguish them by giving them different names (e.g. + addition or \cdot multiplication).

The simplest example of an (Abelian) group is the group of whole numbers under the standard operation of addition. The set of whole numbers is denoted as \mathbb{Z} . \mathbb{Z} has an infinite number of elements, because $\mathbb{Z} = \{\dots, -4, -3, -2, -1, 0, 1, 2, 3, 4, \dots\}$. For example, the operation of $1 + 2$ lies in \mathbb{Z} , for $1 + 2 = 3$ and 3 lies in \mathbb{Z} . The neutral element in the group \mathbb{Z} is 0 . The inverse element of 3 is -3 , for $3 + (-3) = 0$.

For our purpose, so-called *finite* groups play an important role. This means that these exists a set \mathcal{M} with a fixed number of elements and an operation $+$ such that the above conditions are fulfilled. One example of this is any set \mathbb{Z}_n where $\mathbb{Z}_n = \{0, 1, 2, 3, \dots, n-1\}$, n is a positive whole number and the operation is addition mod n , i.e. a and b in \mathbb{Z}_n are subject to the operation $a + b$ mod n .

⁵The inverse is uniquely determined because if $x, y \in G$ are each inverse to a , i.e. $ax = xa = e$ and $ay = ya = e$, then $x = xe = x(ay) = (xa)y = ey = y$.

Cyclic groups Cyclic groups⁶ are those groups G' that possess an element g from which the group operation can be used to generate all other elements in the group. This means that for each element a in G' there exists a positive whole number i such that if g is subject to the operation i times (i.e. “ $g \cdot i$ ”), $g + g + \cdots + g = a$ (additive group) or $g^i = g \cdot g \cdots g = a$ (multiplicative group). The element g is the *generator* of the cyclic group — each element in G' can be generated using g and the operation.

Group order Now to the order of an element of the group: Let a be in G . The smallest positive whole number r for which a subject to the operation with itself r times is the neutral element of the group G' (i.e.: $r \cdot a = a + a + \cdots + a = e$ respectively $a^r = e$), is called the *order* of a .

The order of the group is the number of elements in the set G .

7.3.2 Fields

In mathematics, one is often interested in sets on which at least two (group) operations are defined — frequently called addition and multiplication. Most prominent are so called fields.

A field is understood to be a set K with two operations (denoted as $+$ and \cdot) which fulfills the following conditions:

- The set K forms an Abelian group together with the operation $+$ (addition), where 0 is the neutral element of the operation $+$.
- The set $K \setminus \{0\}$ also forms an Abelian group together with the operation \cdot (multiplication).
- For all elements a, b and c in K , we have $c \cdot (a + b) = c \cdot a + c \cdot b$ and $(a + b) \cdot c = a \cdot c + b \cdot c$ (distributive law).

Fields may contain an infinite number of elements (e.g. the field of real numbers). They are called *infinite* fields. In contrast we call a field *finite*, if it contains only a finite number of elements (e.g. $\mathbb{Z}_p = \{0, 1, 2, 3, \dots, p - 1\}$, where p is a prime. \mathbb{Z}_p with addition mod p and multiplication mod p).

Characteristic of a field Let K be a field and 1 be the neutral element of K with respect to the multiplicative operation “ \cdot ”. Then the characteristic of K is said to be the order of 1 with respect to the additive operation. This means that the characteristic of K is the smallest positive integer n such that

$$\underbrace{1 + 1 + \cdots + 1}_{n \text{ times}} = 0.$$

If there is no such n , i.e. if $1 + 1 + \cdots + 1 \neq 0$ no matter how many 1s we add, then we call K a field with characteristic 0.

Thus, fields with characteristic 0 are infinite since they contain the (pairwise distinct) elements $1, 1 + 1, 1 + 1 + 1, \dots$. On the other hand, fields with finite characteristic may be finite or infinite.

If the characteristic is finite, it has to be prime. This fact can easily be proved: Assume $n = pq$, $p, q < n$, is the characteristic of a field K . By definition of n , the elements $\bar{p} =$

⁶Cyclic groups can be in general also endless like the additive group of the integer numbers. We consider here only finite cyclic groups.

$\underbrace{1+1+\cdots+1}_{p \text{ times}}, \bar{q} = \underbrace{1+1+\cdots+1}_{q \text{ times}}$ of K are not equal to 0. Thus, there exist inverse elements

$\bar{p}^{-1}, \bar{q}^{-1}$ with respect to multiplication. It follows that $(\bar{p}\bar{q})(\bar{p}^{-1}\bar{q}^{-1}) = 1$, which contradicts the fact that $\bar{p}\bar{q} = \bar{n} = \underbrace{1+1+\cdots+1}_{n \text{ times}} = 0$ and, hence, $(\underbrace{\bar{p}\bar{q}}_{=0})(\bar{p}^{-1}\bar{q}^{-1}) = 0$.

Comment:

The field of real numbers has the characteristic 0; the field \mathbb{Z}_p has the characteristic p . If p is not prime, \mathbb{Z}_p is not a field at all.

The most simple field is $\mathbb{Z}_2 = \{0, 1\}$. It contains only two elements, the neutral elements with respect to addition and multiplication. In particular, we have $0+0=0$, $0+1=1+0=1$, $1+1=0$, $1 \cdot 1=1$, $0 \cdot 0=0 \cdot 1=1 \cdot 0=0$.

Finite Fields As mentioned above, each finite field has a characteristic $p \neq 0$, where p is a prime. On the other hand, given a prime p there is a field which has exactly p elements, that is \mathbb{Z}_p .

However, the number of elements of a field need not be prime in general. For example, it is not hard to construct a field with 4 elements⁷.

One can show that the order of any field is a prime power (i.e. the power of a prime number). On the other hand, we can construct a field with p^n elements for any given prime p and positive integer n . Since two fields that have the same number of elements can not be distinguished⁸, we say that there is **the field with p^n elements** and denote it by $GF(p^n)$. Here GF stands for *Galois Field* to commemorate the French Mathematician Galois.

The fields $GF(p)$ of prime order play a prominent role. They are called prime fields and often denoted by \mathbb{Z}_p ⁹.

7.4 Elliptic curves in cryptography

In cryptography elliptic curve are a useful tool. Such curves are described by some equation. A detailed analysis has shown that curves of the form¹⁰

$$F(x_1, x_2, x_3) = -x_1^3 + x_2^2 x_3 + a_1 x_1 x_2 x_3 - a_2 x_1^2 x_3 + a_3 x_2 x_3^2 - a_4 x_1 x_3^2 - a_6 x_3^3 = 0, \quad (7.1)$$

⁷The set $K = \{0, 1, a, b\}$ fitted with the operation defined in the tabular below is a field:

$+$	0	1	a	b
0	0	1	a	b
1	1	0	b	a
a	a	b	0	1
b	b	a	1	0

und

\cdot	0	1	a	b
0	0	0	0	0
1	0	1	a	b
a	0	a	b	1
b	0	b	1	a

⁸If K, K' are fields with $k = p^n$ elements, then there is a one-to-one map $\varphi : K \rightarrow K'$, that respects the arithmetic of the field. Such a map is called an isomorphy. Isomorphic fields mathematically behave in the same way so that it makes no sense to distinguish between them. For example, \mathbb{Z}_2 und $K' = \{\text{ZERO}, \text{ONE}\}$ with zero-element *ZERO* and one-element *ONE* are isomorphic. We note that mathematical objects are only defined by their mathematical properties.

⁹For prime fields additive as well as multiplicative group are cyclic. Furthermore, each field $GF(p^n)$ contains a subfield that is isomorphic to the prime field \mathbb{Z}_p .

¹⁰This curve is given by the zeros of a polynomial F of degree three in three variables. In general, expressions of the form $P = \sum_{i_1, \dots, i_n \in \mathbb{N}_0} a_{i_1 \dots i_n} x_1^{i_1} \dots x_n^{i_n}$ with coefficients $a_{i_1 \dots i_n} \in K$ are called polynomials in n variables x_1, \dots, x_n with underlying field K , if $\deg P := \max\{i_1 + \dots + i_n : a_{i_1 \dots i_n} \neq 0\}$ is finite, i.e. the sum has only finitely many non-zero terms (monomials). The sum of the exponents of the variables of each term of the sum is at most 3, at least one term of the sum has a single variable with 3 as value of the according exponent.

are especially useful. The variables x_1, x_2, x_3 and parameters a_1, \dots, a_4, a_6 are elements of a given field K , which has certain properties that make it useful from the cryptographic point of view. The underlying field K might be the well known field of real numbers or some finite field (see last section). In order to obtain a curve that is useful for cryptography, the parameters have to be chosen in a way that the following conditions hold

$$\frac{\partial F}{\partial x_1} \neq 0, \quad \frac{\partial F}{\partial x_2} \neq 0, \quad \frac{\partial F}{\partial x_3} \neq 0.$$

We identify points on the curve that can be derived from each other by multiplying each component with some scalar. This makes sense since (x_1, x_2, x_3) solves (7.1) if and only if $\alpha(x_1, x_2, x_3)$ ($\alpha \neq 0$) does. Formally, this means that we consider classes of equivalent points instead of single points, where points are called equivalent if one is the scalar multiple of the other one. If we put $x_3 = 0$ in the basic equation (7.1), then this equation collapses to $-x_1^3 = 0$, leading to $x_1 = 0$. Thus, the equivalence class which includes the element $(0, 1, 0)$ is the only one that contains a point with $x_3 = 0$. For all points on the curve that are not equivalent to $(0, 1, 0)$, we may apply the following transformation

$$K \times K \times (K \setminus \{0\}) \ni (x_1, x_2, x_3) \mapsto (x, y) := \left(\frac{x_1}{x_3}, \frac{x_2}{x_3} \right) \in K \times K,$$

which reduces the number of variables to two instead of three. We note that the basic equation (7.1) $F(x_1, x_2, x_3) = 0$ was chosen in a way that this transformation leads to the famous so-called Weierstrass-Equation¹¹ holds

$$y^2 + a_1xy + a_3y = x^3 + a_2x^2 + a_4x + a_6. \quad (7.2)$$

Since all but one point (i.e. equivalence class) of the elliptic curve can be described using equation (7.2), this equation is often called the elliptic equation, and its solutions written as

$$E = \{(x, y) \in K \times K \mid y^2 + a_1xy + a_3y = x^3 + a_2x^2 + a_4x + a_6\} \cup \{\mathcal{O}\}.$$

Here, \mathcal{O} represents the point $(0, 1, 0)$ that is loosely speaking mapped to infinity by the transformation (division by x_3) that reduces the three variables to two.

In contrast to figure 7.3 only finite fields $K = GF(p^n)$ are used in elliptic curve cryptography. The reason is loosely speaking that in modern communication engineering data processing is always based on discrete data (simply because computers accept only discrete data).

For practical reasons, it turned out to be useful to take either $GF(p)$ with a large prime p or $GF(2^n)$ with a (large) positive integer n . Using $GF(p)$ has the advantage of providing a relatively simple arithmetic; on the other hand $GF(2^n)$ allows a binary representation of each element that supports the way computers work. Other fields like, for example, $GF(7^n)$ do not have any of these advantages and are, thus, not considered, although there is no mathematical reason why they should not.

A coordinate transformation can result in a simpler version¹² of the Weierstrass equation. Depending whether $p > 3$, different transformations are used, and we obtain

- in case of $GF(p)$, $p > 3$, the elliptic curve equation of the form

$$y^2 = x^3 + ax + b \quad (7.3)$$

with $4a^3 + 27b^2 \neq 0$

¹¹Karl Weierstrass, 31.10.1815–19.12.1897, German mathematician, famous for his rigorous formal approach to mathematics.

¹²Such a coordinate transformation is combination of a rotation and a dilatation of the coordinate system without changing the elliptic curve itself.

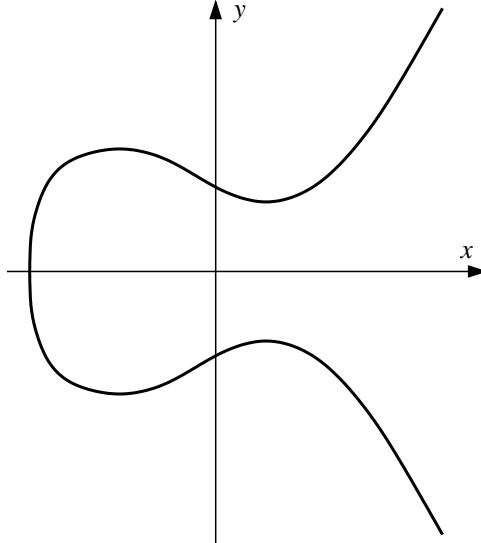


Figure 7.3: Example of an elliptic curve with the real numbers as underlying field.

- in case of $GF(2^n)$ the elliptic curve equation of the form

$$y^2 + xy = x^3 + ax^2 + b \quad (7.4)$$

with $b \neq 0^{13}$.

This conditions on the parameters a, b ensure that the elliptic equation can be used in the context of cryptography¹⁴.

Let $|E|$ denote the number of elements of an elliptic curve E given an underlying field $GF(k)$ (for practical reasons either $k = p$ with p prim or $k = 2^n$). Then Hasse's theorem[Silverman1986] yields $||E| - k - 1| \leq 2 \cdot \sqrt{k}$. This Inequality is equivalent to $k + 1 - 2\sqrt{k} < |E| < k + 1 + 2\sqrt{k}$. In particular, this means that the number of elements of an elliptic curve is approximately k (for large k).

7.5 Operating on the elliptic curve

In order to work with elliptic curves in practice, we define an operation (often written in an additive way $+$) on the set of points on the curve. If we have a curve over the field $GF(p)$, we define the commutative operation $+$ by

1. $P + \mathcal{O} = \mathcal{O} + P = P$ for all $P \in E$,
2. for $P = (x, y)$ and $Q = (x, -y)$ we set $P + Q = \mathcal{O}$,
3. for $P_1 = (x_1, x_2), P_2 = (x_2, y_2) \in E$ with $P_1, P_2 \neq \mathcal{O}$ and $(x_2, y_2) \neq (x_1, -y_1)$ we set $P_3 := P_1 + P_2$, $P_3 = (x_3, y_3)$ defined by

$$x_3 := -x_1 - x_2 + \lambda^2, \quad y_3 := -y_1 + \lambda(x_1 - x_3)$$

¹³The form (7.3) is called the standard form of the Weierstrass-equation. If the characteristic of the field is 2 or 3, we obtain $4 = 0$ respectively $27 = 0$, which means that the condition on parameters a, b collapse. Loosely speaking, this is the reason why the transformation to the standard form does not work in these cases.

¹⁴Formally we call such curves non singular.

with the auxiliary quotient

$$\lambda := \begin{cases} \frac{y_1 - y_2}{x_1 - x_2} & \text{if } P_1 \neq P_2, \\ \frac{3x_1^2 + a}{2y_1} & \text{if } P_1 = P_2. \end{cases}$$

In particular, we obtain $-P = (x, -y)$ for $P = (x, y) \in E$.

If we deal with a curve over the field $GF(2^n)$, we define the operation $+$ in an analogous way by

1. $P + \mathcal{O} = \mathcal{O} + P = P$ for all $P \in E$,
2. for $P = (x, y)$ and $Q = (x, x + y)$ we set $P + Q = \mathcal{O}$,
3. for $P_1 = (x_1, x_2), P_2 = (x_2, y_2) \in E$ with $P_1, P_2 \neq \mathcal{O}$ and $(x_2, y_2) \neq (x_1, x_1 + y_1)$ we set $P_3 := P_1 + P_2$, $P_3 = (x_3, y_3)$ defined by

$$x_3 := -x_1 + x_2 + \lambda + \lambda^2 + a, \quad y_3 := y_1 + x_3 + \lambda(x_1 + x_3)$$

with auxiliary quotient

$$\lambda := \begin{cases} \frac{y_1 + y_2}{x_1 + x_2} & \text{if } P_1 \neq P_2, \\ x_1 + \frac{y_1}{x_1} & \text{if } P_1 = P_2. \end{cases}$$

In particular, we obtain $-P = (x, -y)$ for $P = (x, y) \in E$.

(Note that $-(-P) = (x, x + (x + y)) = (x, 2x + y) = (x, y)$, since the underlying field has characteristic 2.)¹⁵

One can verify that $+$ defines a group operation on the set $E \cap \{\mathcal{O}\}$. In particular this means that the sum of two points is again a point on the elliptic curve. How his operation works is geometrically visualized in the following section.

¹⁵An animation of the addition of points on elliptic curves can be found on the Certicom homepage
<http://www.certicom.com/index.php/ecc-tutorial>
 See also the web link about the Java-Tutorial at the end of this chapter.

How to add points on an elliptic curve

The following figures show how points on an elliptic curve over the field of real numbers are summed up using affine coordinates. We note that the point infinity \mathcal{O} cannot be shown in the affine plane.

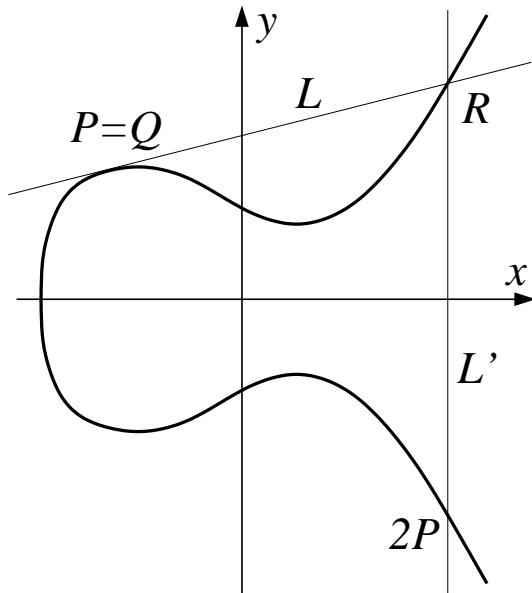


Figure 7.4: Doubling of a point

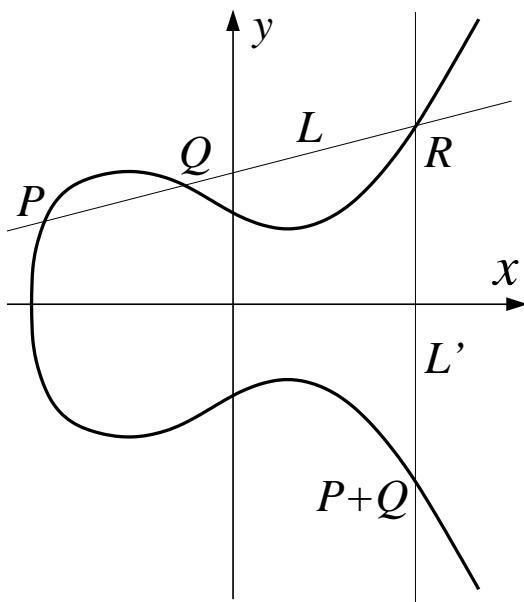


Figure 7.5: Summing up two different points over the real number field

7.6 Security of elliptic-curve-cryptography: The ECDLP

As mentioned above in section 7.4, we only consider elliptic curves over the finite¹⁶ fields $GF(2^n)$ or $GF(p)$ (for a large prime p). This means that all parameters that describe the curve are taken from this underlying field. If E is an elliptic curve over such a field and P is a point on the curve E , then we can derive for all positive integers m

$$mP := \underbrace{P + P + \cdots + P}_{m \text{ times}}.$$

Looking on this operation from the cryptographic point of view, it turns out to be very interesting by the following reason: On the one hand one needs only $\log m$ operations to calculate mP — one simply has to calculate $P, 2P, 2^2P, 2^3P, \dots$, write m in a binary form and finally add all these multiples 2^kP of P with respect to the binary representation of m — on the other hand it seems to be very hard to find m given P and $Q = mP$ on E . Of course, we may simply calculate $P, 2P, 3P, 4P, 5P, \dots$ and compare each of them with Q . But this will take as much as m operations.

Yet there is no algorithm known that efficiently derives m given P and Q . The best algorithms known so far need about \sqrt{q} operations where q is the (largest) prime factor of $p - 1$, in case the underlying field is $GF(p)$; here m should be between 1 and q liegen so that one needs at most $\log q$ operations to calculate mP . However, the quotient $\frac{\sqrt{q}}{\log q}$ tends to $+\infty$ very fast for large q .

If we choose the parameters sufficiently large (for example, let p be prime and at least 160 bits long), an computer will easily be able to calculate mP (in less than a second). The *inverse problem* however, to derive m from mP and P , can (still) not be solved in reasonable time.

This problem is known as the “Elliptic Curve Discrete Logarithm Problem” (for short ECDLP).

In elliptic curve cryptography we formally look at points on the elliptic curve as elements of a group with point addition $+$ as operation. Furthermore, we use only elliptic curves that have a sufficiently large number of points. However, in special cases curves may be weak and not useful due to other reasons. For such special cases the ECDLP can be much easier to solve than in the general case. This means that one has to look carefully at the parameters when choosing an elliptic curve for cryptographic applications.

Not useful for cryptography are *a-normal* (that are curves over \mathbb{Z}_p , for which the set **E** consists of exactly p elements) and *supersingular* curves (that are curves, for which the ECDLP can be reduced to the “normal” discrete logarithms in another, smaller finite field). This means that there are cryptographically useful and non-useful elliptic curves. Given the parameters a and b , it is possible to determine whether a curve is useful or not. In many publications one can find parameters that turned out to be useful for cryptography. The open (scientific) discussion guarantees that these results take into account latest research.

Given a secure curve, the time that is needed to solve the ECDLP is strongly correlated with parameter p in case $GF(p)$ respectively n in case of $GF(2^n)$. The larger these parameters become, the more time an attacker needs to solve the ECDLP — at least with the best algorithms known so far. Experts recommend bit-lengths of 200 for p for secure curves. A comparison with RSA modulus length shows why elliptic curves are so interesting for applications. We note that the computation effort for signing and encryption is closely related to the bit-length of the

¹⁶Discrete in contrast to continuous.

parameters. In addition the initiation process, i.e. the generation of the private-public-key-pair, becomes more complicated the larger p is. Thus, one looks for the smallest parameters that still come along with the security required. It is remarkable that a length of 200 bits for p is sufficient to construct a *good* elliptic curve that is as secure as RSA with a 1024 bit RSA modulus (as far as we know today). For short, the reason for this advantage of ECC lies in the fact that the best algorithms known for solving the ECDLP need exponential time while the best algorithms for factorizing are sub-exponential (number field sieve, quadratic sieve or factorizing with elliptic curves). Hence, the parameters for a cryptosystem that is based on the problem of *factorizing large integers* have to be larger than the parameters for a system based on ECDLP.

7.7 Encryption and signing with elliptic curves

The *elliptic curve discrete logarithm problem* (ECDLP) is the basis for elliptic curve cryptography. Based on this problem, there are different signature schemes. In order to apply one of these, we need:

- An elliptic curve \mathbf{E} with an underlying field $GF(p^n)$.
- A prime $q \neq p$ and a point G on the elliptic curve \mathbf{E} with order q . This means that $qG = \mathcal{O}$ and $rG \neq \mathcal{O}$ for all $r \in \{1, 2, \dots, q-1\}$. Thus q is a factor of the group order (i.e. the number of elements) $\#\mathbf{E}$ of E . Since q is prime, G generates a cyclic sub-group of \mathbf{E} of order q .

The parameters mentioned are often called *Domain* parameter. They describe the elliptic curve \mathbf{E} and the cyclic sub-group of \mathbf{E} on which the signature scheme is based.

7.7.1 Encryption

Using elliptic curves one can construct a key exchange protocol based on the Diffie-Hellman protocol (see chapter 5.4.2). The key exchanged can be used for a subsequent symmetric encryption. We note that in contrast to RSA there is no pair of private and public key that can be used for encryption and decryption!

In the notation of elliptic curves, the Diffie-Hellman protocol reads as follows: First both partners (A und B) agree on a group G and an integer q . Then they choose $r_A, r_B \in \{1, 2, \dots, q-1\}$ at random, derive the points $R_A = r_A G$, $R_B = r_B G$ on the elliptic curve and exchange them (using an insecure channel). After that A easily obtains $R = r_A R_B$; B gets the same point ($R = r_B R_A$) by calculating $r_B R_A = r_B r_A G = r_A r_B G = R$. We note that R_A, R_B are easy to derive as long as r_A respectively r_B are known G . However, the inverse operation, to get R_A respectively R_B from r_A respectively r_B is hard.

Using the best algorithms known so far, it is impossible for any attacker to obtain R without knowing either r_A or r_B — otherwise he would have to solve the ECDLP.

In order to prohibit a “Man-in-the-middle” attack, one may sign the values G, q, R_A, R_B as described in chapter 6.4.1.

7.7.2 Signing

Using the DSA signature scheme, one can proceed as follows: The signing party chooses a (non-trivial) number $s \in \mathbb{Z}_q$, which will be the private key, and publishes q , G and $R = sG$. We note that s cannot be obtained from G and R are not sufficient — a fact on which the security of the signature scheme is based.

Given the message m , which should be signed, one first constructs a digital finger print using a hash-algorithm h such that $h(m)$ has its values in $\{0, 1, 2, \dots, q - 1\}$. Thus, $h(m)$ can be considered as an Element of \mathbb{Z}_q . Then the signing party chooses a random number $r \in \mathbb{Z}_q$ and derives $R = (r_1, r_2) = rG$. We note that the first component r_1 of R is an element of $GF(p^n)$. This component will then be projected onto \mathbb{Z}_q , i.e. in case of $n = 1$ it is interpreted as the remainder of an element of $\{0, 1, \dots, p - 1\}$ divided by q . This projection of r_1 onto \mathbb{Z}_q is denoted by \bar{r}_1 . Then one determines $x \in \mathbb{Z}_q$ such that

$$rx - s\bar{r}_1 - h(m) = 0.$$

The triple (m, r_1, x) is then published as the digital signature of message m .

7.7.3 Signature verification

In order to verify a signature, one has to build $u_1 = h(m)/x$, $u_2 = \bar{r}_1/x$ (in \mathbb{Z}_q and derive

$$V = u_1G + u_2Q.$$

Since we have $Q = sG$, the point $V = (v_1, v_2)$ satisfies $v_1 = u_1 + u_2s$. We note that this operations take place in the field $GF(p^n)$. The projection of $GF(p^n)$ on \mathbb{Z}_q mentioned above should be chosen in such a way that $\bar{v}_1 = u_1 + u_2s$ is an element of \mathbb{Z}_q . Then it follows that

$$\bar{v}_1 = u_1 + u_2s = h(m)/x + \bar{r}_1s/x = (h(m) + \bar{r}_1s)/x = rx/x = r.$$

Since $R = rG$, we obtain $\bar{v}_1 = \bar{r}_1$, i.e. R and V coincide modulo the projection onto \mathbb{Z}_q .

7.8 Factorization using elliptic curves

There are factorization¹⁷ algorithms based on elliptic curves¹⁸. More precisely, these procedures exploit the fact that elliptic curves can be defined over \mathbb{Z}_n (n composite number). Elliptic curves over \mathbb{Z}_n do not form a group, because not every point on such an elliptic curve has an inverse point. This is connected with the fact that - if n is a composite number - there exist elements in \mathbb{Z}_n that do not have an inverse with respect to multiplication mod n . In order to add two points on an elliptic curve over \mathbb{Z}_n , we can calculate in the same way as on elliptic curves over \mathbb{Z}_p . Addition of two points (on an elliptic curve over \mathbb{Z}_n), however, fails if and only if a factor of n has been found. The reason for this is that the procedure for adding points on elliptic curves

¹⁷Especially John M. Pollard was involved in the development of many different factorization algorithms; also at factorization with ECC he was one of the leading heads. As an employee of British Telekom he never published much. At the RSA data Security Conference in 1999 he was awarded for his “outstanding contributions in mathematics”: http://www.eff.org/Privacy/Crypto_misc/DESCracker/HTML/19990118_rsa_awards.html. In 1987 H.W. Lenstra published an often used factorization algorithm, based on elliptic curves (see [Lenstra1987]).

¹⁸The biggest compound numbers currently factorized with elliptic curves have about 80 decimal digits:

<http://www.loria.fr/~zimmerma/records/top50.html>.

See also the web link about the ECMNET project at the end of this chapter.

gives elements in \mathbb{Z}_n and calculates the inverse elements for these (with respect to multiplication mod n) in \mathbb{Z}_n . The extended Euclidean algorithm is used here. If the addition of two points (that lie of an elliptic curve over \mathbb{Z}_n) gives an element in \mathbb{Z}_n that does not have an inverse element in \mathbb{Z}_n , then the extended Euclidean algorithm delivers a genuine factor of n .

Factorization using elliptic curves thus principally works as follows: Random curves over \mathbb{Z}_n are selected, as well as random points (that lie on this curve) and add them; you thus obtain points that also lie on the curve or find a factor of n . Factorization algorithms based on elliptic curves therefore work probabilistically. The opportunity of defining large number of elliptic curves over \mathbb{Z}_n allows you to increase the probability of finding two points which you can add to obtain a factor of n . These procedures are therefore highly suitable for parallelization.

7.9 Implementing elliptic curves for educational purposes

There are not many free programms offering ECC under a graphical user interface. The following subsections explain which according functionality is available in CrypTool and in Sage.

7.9.1 CrypTool

CT1 offers elliptic curves for the digital signature function¹⁹ and for ECC-AES hybrid encryption²⁰.

It implements the basic algorithms for group operations, for generating elliptic curves, for importing and exporting parameters for elliptic curves over finite fields with p elements (p prime). The algorithms have been implemented in ANSI C and comply with draft no. 8 of the IEEE P1363 work group *Standard Specifications for Public Key Cryptography*

<http://grouper.ieee.org/groups/1363>.

The procedure implements the cryptographic primitives for generating and verifying signatures for the variations of Nyberg-Rueppel signatures and DSA signatures based on elliptic curves.

Step-by-step point addition on elliptic curves is visualized in CT1 and JCT.²¹

7.9.2 Sage

In Sage elliptic curves are described at

http://www.sagemath.org/doc/constructions/elliptic_curves.html²².

Additionally there is an exhaustive, interactive ECC tutorial by Maike Massierer. This interactive introduction to Elliptic Curve Cryptography is built up as a Sage notebook.

Sage notebooks are running after a logon within a browser^{23,24}.

¹⁹The dialog box, which appears in CT1 after clicking the menu **Digital Signatures/PKI \ Sign Message**, offers the EC methods ECSP-DSA and ECSP-NR.

²⁰Within CT1 you can find this technique using the menu path **Crypt \ Hybrid**.

²¹CT1: menu **Digital Signatures/PKI \ Signature Demonstration (Signature Generation)**, JCT (default perspective): menu **Visuals \ Elliptic Curve Calculations**.

²²According Sage samples can be found at the "Published Worksheets" at <http://www.sagenb.org/pub/>
- about Elliptic Curve: <http://www.sagenb.org/home/pub/606/>
- about Elliptic Curve El Gamal: <http://www.sagenb.org/home/pub/104/>, or at
- the "Elliptic Curve Cryptography (ECC) Tutorial"
<http://www.williamstein.org/simuw06/notes/notes/node12.html>

²³If you installed Sage on your own (Unix) server, you first have to enter at the command line the command **notebook()**.

²⁴The ECC notebook of Maike Massierer needs the KASH3 library: Therefore e.g. with Sage 4.2.1 the package "kash3-2008-07-31.spkg" has to be installed (command **sage -i**).

The ECC notebook of Massierer^{25,26} consists of 8 parts (title page plus 7 chapters) and aims to let even beginners understand what elliptic curves are:

0. ECC Notebook (title page and contents)
1. Introduction and Overview
2. Motivation for the use of Elliptic Curves in Cryptography
3. Elliptic Curves in Cryptography
4. Cryptographic Protocols in ECC
5. Domain Parameter Generation for ECC Systems
6. Conclusion and Further Topics
7. References

²⁵Instructions to use an interactive Sage notebook:

- Public Sage servers like <http://sage.mathematik.uni-siegen.de:8000> or <http://www.sagenb.org/> often offer running samples as "Published Worksheets", which you can run and download without login on. These worksheets are listed if you click on "Published" in the above right corner.
- Worksheets using the `interact` command currently need some additional todos for a user to work correctly: sign-in, make a copy and execute all commands again.
This works as follows (described for the sagenb server and for the ECC tutorial):
 - Sign-up for a Sage notebook account at <http://sagenb.org/register> and sign in at <http://sagenb.org/>.
 - Open the worksheet <http://sagenb.org/home/pub/1126/>. This contains the table of contents of the interactive ECC notebook. From here you can navigate via a click to the different chapters of the document.
 - In the very top left corner, click `Edit a copy` in order to create your own copy of the worksheet.
 - Sometimes it's necessary at the beginning to re-evaluate the worksheet. Click in the left upper corner on `Action -> Evaluate all`.
 - Some of the applications still do not always work after opening a worksheet. Instead of nice output, they show lots of (blue) error messages. This normally can be solved quickly by clicking the gray "%hide" string: Then you get the code behind the graphics. Simply generate the graphics again with Shift-Enter.
Even after doing this, the graphics code does not always disappear. Instead, it sometimes turns gray. Should this happen, click on the gray text, then click somewhere outside of the text box. The code will then disappear and leave you with a nice layout of the worksheet.
 - Some of the ECC tutorial's content uses a special math fonts that are not installed by default with most browsers. When you notice that formulas are not displayed correctly or even get an error message about missing fonts from your browser, you need to install the jsMath fonts for a better layout.
See <http://www.math.union.edu/~dpvc/jsMath/> and <http://pubpages.unh.edu/~jsh3/jsMath/>. After installing these fonts you can see the jsMath symbol at the lower border of your browser. If you click this symbol you can find the download page for the TIFF fonts. This fonts installation has to be done at every PC.
 - According to the Sage-support newsgroup there is work underway to create a system for using `@interact` completely outside of the Sage notebook (JS code within a static html pages).

²⁶Since 2008 this ECC notebook can be found at <http://sage.mathematik.uni-siegen.de:8000/home/pub/45/> (#45 to #52). To logon in Siegen you have to allow port 8000 and Cookies.
Since 2009 an updated version of this ECC notebook can be found at <http://sagenb.org/home/pub/1126/> (#1126 to #1133).

7.10 Patent aspects

If the field $GF(2^n)$ is used instead of the prime field $GF(p)$, one has to make substantial changes in the implementation. The advantage of $GF(2^n)$ lies in the fact that calculations in $GF(2^n)$ can be implemented very efficiently using the binary representation. In particular, divisions are much easier to process compared to $GF(p)$ (this is particularly important in the signature scheme mentioned above where a division is needed for processing a signature as well as for the verification).

In order to achieve maximal gain in efficiency, one may choose a field that allows special basis like polynomial basis (useful for software implementations) or normal basis (best for hardware implementations). For special n (like, for example, $n = 163, 179, 181$) one may even combine both advantages. However, they are still non-standard.

Sometimes only the first component and one additional bit is used as representation of a point on the elliptic curve instead of the full two components. Since the first component together with the additional bit is sufficient to derive the full point, this representation minimizes the memory capacity needed. In particular, for normal basis this point compression can be implemented efficiently. In addition, the cryptographic protocols themselves become more effective. A disadvantage is, however, that *point compression* can be used for about half of all elliptic curves only and is protected under US patent (US Patent 6141420, Certicon), causing additional costs. In the general case $GF(p^n)$ (and also in case $n = 1$) often so called affine or projective co-ordinates are used. Depending on the application, these co-ordinates may result in a gain in efficiency as well.

A comprehensive description of all implementations and their advantages and disadvantages would go far beyond the scope of this paper. We only want to state that there is a variety of possible implementations for elliptic curve cryptography, much more than for RSA. Therefore, there are serious efforts to reduce this large to a small number of standard implementations. Some standardization committees even try to reduce the complexity by focusing on a small number of (prescribed) curves (ASC-approach).

Today it is still not clear whether these standardization initiatives will be successful or not. However, without agreed standards, ECC is not likely to become a real alternative for RSA. The committees might be forced to act fast if there was a break-through in factorization.

7.11 Elliptic curves in use

Today elliptic curve cryptography is already in use. A prominent example is the information network Bonn-Berlin²⁷, used for the exchange of strictly confidential documents between different German federal governmental institutions in Berlin and Bonn. With the help of ECC a high security solution could be realized. Interoperability, however, played only a minor role.

In Austria ECC has been massively launched: A bank card with digital signature function.

Both examples show the typical range of application for elliptic curve cryptography: For high security solutions and for implementations on smartcards in which the key length is crucial (because of physical memory available).

²⁷The Informationsverbund Bonn-Berlin (IVBB) connects governmental institutions in the old and new German capital.

http://www.cio.bund.de/cln_094/sid_92C19118CBA5A021AFD1ABAEC15D2B77/DE/IT-Angebot/IT-Infrastrukturen/IVBB/ivbb_inhalt.html

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Chapter 8

Homomorphic Ciphers

(Martin Franz, January 2013)

8.1 Introduction

Homomorphic ciphers are public-key cryptosystems with special properties. They allow performing certain arithmetic operations on encrypted ciphertexts, without knowing the corresponding plaintexts and without having to decrypt the ciphertexts first. These special properties have led to a huge amount of applications for homomorphic ciphers, e.g. in the domain of cloud computing. A very famous and relatively new cryptosystem with homomorphic properties is the Paillier cryptosystem. But also some of the older and well established cryptosystems, such as ElGamal or RSA, have homomorphic properties.

8.2 Origin of the term “homomorphic”

We first clarify the meaning and the origin of the term “homomorphic”. This term in cryptography is derived from its counterpart in mathematics: in mathematics, a homomorphism is a structure-preserving map between two algebraic structures. In the common sense this means, that a homomorphism $f : X \rightarrow Y$ maps the structure of X to the structure of Y . Using an example, this can be easily illustrated: Let $(X, +)$ and $(Y, *)$ two algebraic groups with group operations $+$ and $*$, respectively. A homomorphism $f : X \rightarrow Y$ maps any given $x \in X$ to a value $y \in Y$, in a way that it holds:

$$f(x_1 + x_2) = f(x_1) * f(x_2)$$

for any two x_1, x_2 in X . This means, that for any two values x_1, x_2 it does not matter whether we first compute their sum (group operation of X) and then apply f (this is the left side of the above given equation); or, whether we first apply f to the values x_1, x_2 , and then compute their product in Y , thus apply the group operation of Y . Please note that the operations $+$ and $*$ were chosen here only as an example, they always depend on the algebraic group they belong to. Naturally, the same relation holds for homomorphisms between groups with the same group operation.

Example: Let $X = \mathbb{Z}$ be the set of integer values. The set \mathbb{Z} together with the addition operation forms an algebraic group $G_1 = (\mathbb{Z}, +)$. Similarly, the real values \mathbb{R} without the value zero together with the multiplication operation form a group $G_2 = (\mathbb{R} \setminus \{0\}, *)$. The function

$f : \mathbb{Z} \rightarrow \mathbb{R} \setminus \{0\}$, $z \mapsto e^z$ is a homomorphism, since for all $z_1, z_2 \in \mathbb{Z}$ it holds: $f(z_1 + z_2) = e^{(z_1 + z_2)} = f(z_1) * f(z_2)$. On the contrary, $f : \mathbb{Z} \rightarrow \mathbb{R} \setminus \{0\}$, $z \mapsto z^2$ is an example for a function which is not a homomorphism.

8.3 Decryption function is a homomorphism

In the remainder of this chapter we will consider public-key cryptosystems with a special property, namely that its decryption function is a homomorphism. A public-key cryptosystem with this property will be called homomorphic.

Let us for now assume, the above described homomorphism f is the decryption function of a known cryptosystem. This means that we can perform certain algebraic operations in the ciphertext space, knowing which effects this will have on their plaintexts. Following the above given example:

Y corresponds to the set of cipher texts, X is the set of plaintexts. For two plaintexts x_1, x_2 with corresponding ciphertexts y_1, y_2 it holds:

$$f(y_1 * y_2) = f(y_1) + f(y_2) = x_1 + x_2$$

This equation can be interpreted as follows: If we multiply two ciphertexts y_1, y_2 with each other and subsequently decrypt their product, then we will obtain the sum of the originally encrypted values x_1 and x_2 . Everybody can – without knowledge of the plaintexts, without having to decrypt and even without knowing the private decryption key – compute a product of the two ciphertexts and knows that upon decryption the owner of the private key will obtain the sum of the two originally encrypted plaintexts.

8.4 Examples of homomorphic ciphers

8.4.1 Paillier cryptosystem

The most famous cryptosystem with homomorphic properties is the one by Paillier [Paillier1999]. First we will see how the Paillier key generation process works. After that, we will show that the Paillier cryptosystem indeed has homomorphic properties.

8.4.1.1 Key generation

First, we generate two random prime numbers p, q in a way that their product $n = pq$ forms a valid RSA modulus. As for common RSA, the value n should have a bit length of at least 1024 bits.

Using the prime values p and q , we can compute the value $\lambda = lcm(p - 1, q - 1)$. lcm here denotes the least common multiple. The RSA modulus n will now be the public key, while the private key is the value λ .

8.4.1.2 Encryption

Let m be the message which will be encrypted, where m is taken from the plaintext space \mathbb{Z}_n . For each encryption, we first choose a random element r from the plaintext space \mathbb{Z}_n .

Subsequently, using the public key, we compute the ciphertext n as:

$$c = E(m, r) = (n + 1)^m * r^n \bmod n^2$$

8.4.1.3 Decryption

Given the private key λ and a ciphertext $c \in \mathbb{Z}_{n^2}^*$, we first compute $S = c^\lambda \bmod n^2$ and subsequently $T = \phi(n)^{(-1)} \bmod n^2$, where ϕ denotes the Euler function.

Finally, we compute the plaintext $m = D(c) = (S - 1)/n * T \bmod n$.

8.4.1.4 Homomorphic property

We will now show that the Paillier cryptosystem has the homomorphic property as described above. For this, we will use E to denote the encryption and D to denote the decryption function of the Paillier cryptosystem. For simplicity, we set $g := n + 1$. For any two plaintexts m_1, m_2 and random values r_1, r_2 we obtain ciphertexts c_1, c_2 as

$$c_1 = g^{m_1} * r_1^n \bmod n^2 \text{ and } c_2 = g^{m_2} * r_2^n \bmod n^2,$$

respectively. Now it is easy to see that for the product $c_3 = c_1 * c_2$ it holds

$$c_3 = (g^{m_1} * r_1^n \bmod n^2) * (g^{m_2} * r_2^n \bmod n^2) = g^{m_1+m_2} * (r_1 * r_2)^n \bmod n^2 = E(m_1 + m_2, r_1 * r_2)$$

Thus, the product of two given ciphertexts is in fact a valid ciphertext, namely the encryption of the sum of the originally encrypted messages. Now it is straightforward to see that the decryption function is a homomorphism. Given two plaintexts m_1, m_2 it holds

$$D(E(m_1, r_1) * E(m_2, r_2)) = D(E(m_1 + m_2, r_1 r_2)) = m_1 + m_2 = D(E(m_1, r_1)) + D(E(m_2, r_2))$$

8.4.2 Other cryptosystems

Also older public-key cryptosystems can have homomorphic properties. Both the ElGamal cryptosystem and RSA constitute famous examples. We will show their homomorphic properties by means of some easy examples.

8.4.2.1 RSA

Let (e, n) be the public RSA key (e the public encryption exponent, n the RSA modulus). For any two messages m_1, m_2 we obtain encryptions $c_1 = m_1^e \bmod n$ and $c_2 = m_2^e \bmod n$. Now for the product of these two encryptions it holds: $c_1 * c_2 = m_1^e * m_2^e \bmod n = (m_1 * m_2)^e \bmod n$. Thus, we obtain an encryption of the product of the two messages m_1 and m_2 . As it is straightforward to see, this property holds for any two plaintexts m_1, m_2 and similar as for Paillier, the decryption function is a homomorphism. As we have seen here, RSA is an example for a homomorphism, where both groups have the same group operation.

8.4.2.2 ElGamal

Similar to RSA we can also show the homomorphic properties of the ElGamal cryptosystem. Let (p, g, K) the public key while the private key is k (thus, it holds $g^k \bmod p = K$). For any

two messages m_1, m_2 and random values r, s we obtain encryptions $(R, c_1) = (K^r \bmod p, m_1 * g^r \bmod p)$ and $(S, c_2) = (K^s \bmod p, m_2 * g^s \bmod p)$. As for RSA, we verify that their product $(R * S, c_1 * c_2)$ is an encryption of $m_1 * m_2$. Again it is straightforward to see that the decryption function is a homomorphism.

8.5 Applications

The homomorphic property of the Paillier cryptosystem can be used to add two encrypted values or to multiply any value under encryption with a known constant (note that the multiplication corresponds to the repeated application of the addition operation). This makes homomorphic ciphers to important and easy to use base primitives in cryptographic applications.

1. One of these applications is the so called “Electronic Voting”. Electronic voting allows a large number of voters to submit their ballots in an encrypted form. This is important in situations, where the voters cannot come together to the same location. This happens, for example, if the voters can only communicate over the internet via email. If the voting behavior of the single parties should remain secret, then the use of homomorphic ciphers is a good solution to this problem. The main principle of electronic voting using homomorphic ciphers is as follows.

All voters encrypt their ballots, using homomorphic encryption (see at the left site of the screenshot). The screenshot depicts the next steps (1 to 3):

- All voters (on the left in the figure 8.1) encrypt the value 1 if they opt positive and the value 0, if opposed to the decision.
 - Using the homomorphic property, one can compute the sum of all encrypted ballots. Since this happens on encrypted values, the voting behavior of all participants remains secret.
 - At the end, the result of the election is determined and published, this happens by decrypting the sum which was computed using the homomorphic property.
2. A second application of homomorphic ciphers is “Secure Multiparty Computation”. Here, two or more parties can compute any commonly known function. Each of the parties provides one or more of the inputs for the function to be computed. The goal of the secure computations is to keep all private inputs secret, while only the result of the function is revealed. The use of homomorphic encryption helps to perform these computations on encrypted data. However, since the Paillier encryption only allows to compute additions of encrypted values (and, e.g. no multiplications can be performed), a number of additional methods and techniques have to be applied. The Wikipedia page [SMC] offers a great start for reading more about this topic and more advanced techniques for Secure Multiparty Computation.
 3. Furthermore it is expected, that homomorphic encryption will provide great advantages in the areas of “Cloud Computing”. Using so called “fully-homomorphic encryption” [HomEnc] it will be possible to run large applications on external servers only on encrypted data. For this, necessarily one needs to be able to perform both arithmetic operations, the addition and the multiplication, on encrypted data (in contrast to Paillier encryption, which only allows performing additions). Such a crypto system was first presented in 2009 [Gentry2009].

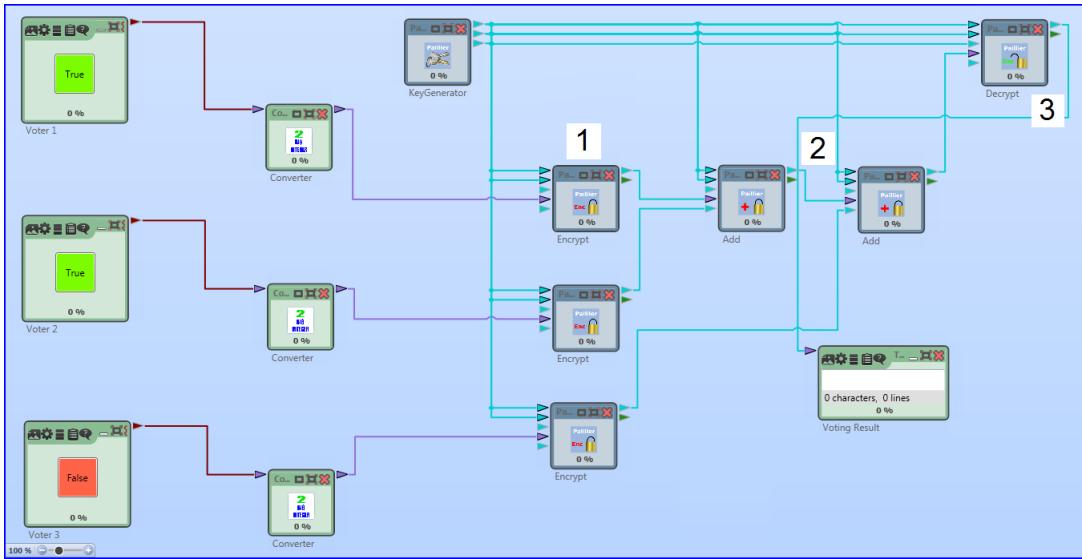


Figure 8.1: Voting example for Paillier

8.6 Homomorphic ciphers in CrypTool

8.6.1 CrypTool 2

In CrypTool 2 you will find an implementation of the Paillier cryptosystem. Among the available components, there are components for key generation (Paillier Key Generator), an example for encryption and decryption with Paillier (called Paillier Text), as well as examples which illustrate the homomorphic properties of the cryptosystem (Paillier Addition, Paillier Blinding and Paillier Voting).

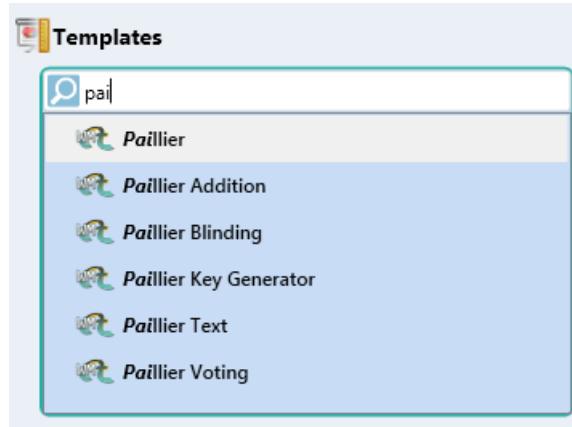


Figure 8.2: Paillier cryptosystem in CrypTool 2

8.6.2 JCrypTool

In JCrypTool there is an implementation (see Figure 8.3), which visualizes the homomorphic properties of various cryptosystem. For RSA and Paillier it shows, that multiplications, for RSA, and additions for Paillier, respectively, can be performed on encrypted values. For the fully-homomorphic cryptosystem by Gentry it is possible to perform both multiplications, as well as additions on encrypted values.

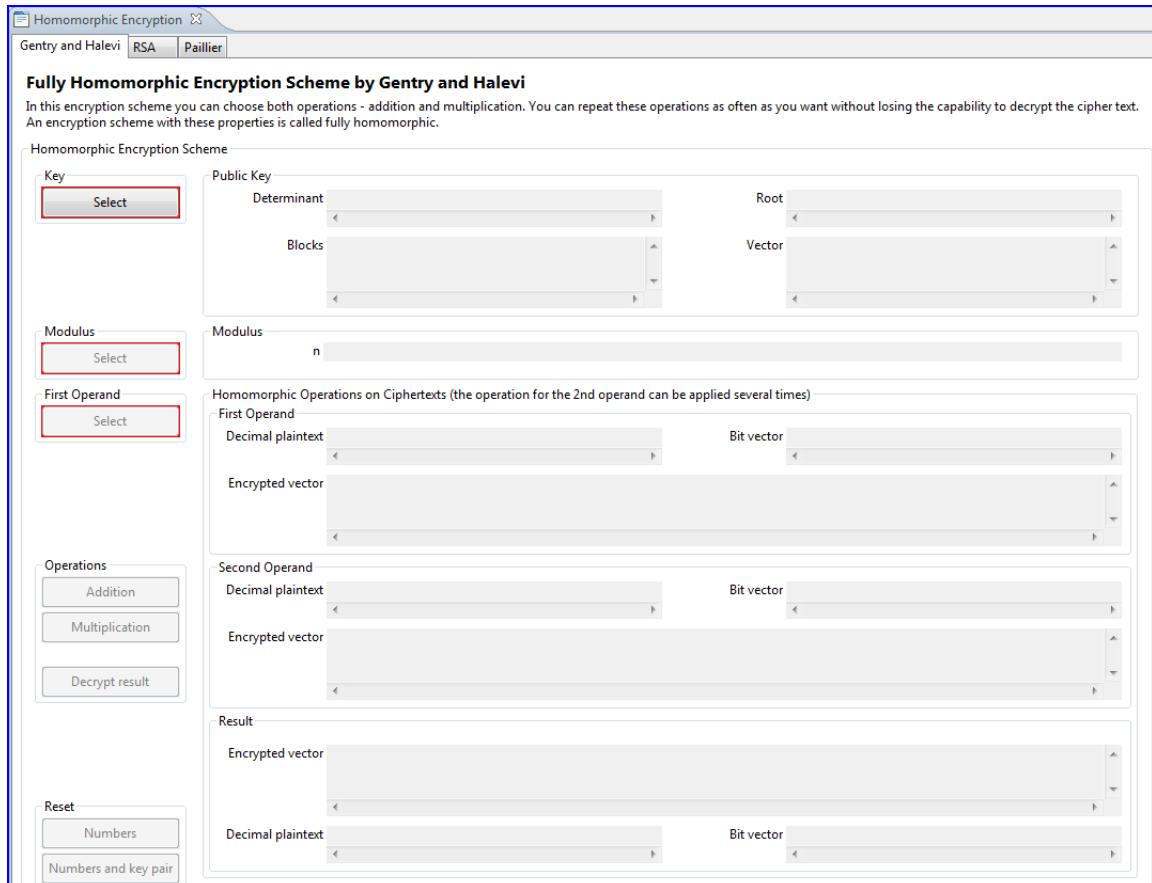


Figure 8.3: Visualization of homomorphic properties in JCrypTool

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Chapter 9

Crypto 2020 — Perspectives for Long-Term Cryptographic Security

(Johannes Buchmann, Erik Dahmen, Alexander May and Ulrich Vollmer, TU Darmstadt, May 2007)

Cryptography is a basic building block of all IT security solutions. Yet, for how long are the cryptographic tools we use today going to remain secure? Is this time long enough to ensure the confidentiality of medical data, to name just one example? Even in the short-term, the potential for havoc is great if certain keys are broken. Just think of the digital signatures that protect the authenticity of automatic updates for the Windows operating system.

9.1 Widely used schemes

In 1978, Rivest, Shamir and Adleman suggested the RSA public key encryption and signature schemes [?]. RSA is still the most widely used public key scheme. The security of RSA depends on the difficulty of factoring so-called RSA moduli which are products of two large prime numbers. In their 1978 paper, the inventors of RSA suggested the use of RSA moduli with 200 decimal digits for long-term security. Later, the company RSA Security published a list of RSA moduli of increasing size, the RSA challenge numbers. RSA Security offered prizes totaling \$ 635,000 for the factorization of these numbers, cf. www.rsasecurity.com/rsalabs/.

In 2005, that is 27 years after the invention of RSA, Bahr, Boehm, Franke, and Kleinjung from Bochum University managed to factor a 200 digit RSA challenge number (www.mat.uniroma2.it/~eal/rsa640.txt). A key with size originally thought to be secure for a very long time was broken with a computation that took them just five months. This illustrates the tremendous progress factoring technology has made within the last 30 years. This progress is based on break-through mathematical ideas — e.g. the number field sieve proposed by John Pollard — as well as significant developments in computer hardware and software implementation technology.¹

In 2000, Lenstra and Verheul developed an extrapolation formula that is supposed to help us forecast the security one can achieve with RSA and other important cryptographic schemes in the long term (www.keylength.com). The formula suggests the use of 850 digit RSA moduli

¹Please compare chapter 4.11 Considerations regarding the security of the RSA algorithm, and especially chapters 4.11.4 and 4.11.5.

if one wishes to protect data for the next 30 years. This corresponds to a 3072 bit RSA key.

Yet, even a well thought out extrapolation formula is no security guarantee! At any time, a brilliant mathematical idea can allow us to factor large numbers easily, and destroy the security of RSA. In 1996, Peter Shor showed that a quantum computer — a new type of computer that leverages the laws of quantum mechanics to speed up certain types of computations — can in principle be used for the fast factorization of large numbers [?]. Despite intensive research in the area, it is still too early to judge whether we are ever going to be able to build quantum computers of sufficient capacity to apply Shor's algorithm to numbers of relevant size.² Recent announcements of significant progress in this area made by the start-up company D-Wave (www.dwavesys.com) have been greeted with a lot of scepticism, even ridicule.

The development of attacks on another frequently used scheme called DSA (Digital Signature Algorithm) and the Elliptic Curve Cryptography (ECC) class of schemes moves in analogy to those on RSA. The security of these schemes depends on the difficulty of computing discrete logarithms. Even today, there is significant algorithmic progress. Quantum computers would render these schemes insecure.

What's the state of affairs with the so-called secret-key encryption schemes? In 1977, DES was introduced as the Data Encryption Standard [?]. Twenty-one years later, the Electronic Frontier Foundation (EFF) built the special purpose machine Deep Crack which needed just 56 hours to break a DES key. The problem with DES was that it used keys which were too short. It seems that the inventors of DES did not foresee the speed of hardware development. The Advanced Encryption Standard AES [?], successor to DES, is deemed secure at the moment even though there are interesting, if still inefficient, methods to attack AES with algebraic methods.

9.2 Preparation for tomorrow

Is the security of today's cryptography measuring up to its increasing importance? The experience shows: Carefully designed and implemented cryptographic schemes have a life time of five to twenty years. Whoever uses RSA, ECC or AES for short-term protection of data may feel safe. Moreover, it is also possible to achieve long-term authenticity, integrity and non-reputability of data, e.g., using the multiple signature scheme suggested by Sönke Maseberg [?].

However, current schemes cannot guarantee long-term confidentiality. And what is to be done in twenty years from now? What should we do if, quasi over-night, unexpected mathematical progress renders an important cryptographic scheme insecure? Three things are necessary to prepare us for this event:

- a pool of secure alternative cryptographic schemes,

²Required qubits for attacks on RSA, DSA and ECDSA using key with a bit length n:

RSA	$2n + 3$
DSA	$2n + 3$
ECDSA 2^n	$\sim 2n + 8 \log n$
ECDSA p	$\sim 4n$

Please compare chapter 5.3 in "SicAri – Eine Sicherheitsplattform und deren Werkzeuge für die ubiquitäre Internetnutzung, KB2.1 – Abschlussbericht, Übersicht über Angriffe auf relevante kryptographische Verfahren", version 1.0, Mai 17, 2005, Prof. Dr. Johannes Buchmann et al., TUD-KryptC and cv cryptovision GmbH (http://www.cdc.informatik.tu-darmstadt.de/~schepers/kb_21_angriffe.pdf) and the dissertation of Axel Schmidt at the same faculty.

- infrastructures that enable us to exchange one cryptographic scheme for another, easily and quickly, and
- methods that ensure long-term confidentiality.

For many years, the cryptography group at the Technische Universität Darmstadt and its spin-off, the company FlexSecure (www.flexsecure.de), have worked to provide these tools. The trust center software FlexiTrust which is employed by the German National Root Certification Authority and the German Country Signing Authority offers an infrastructure within which cryptographic schemes can be easily exchanged. The open source library FlexiProvider implements a multitude of cryptographic schemes. Lately, we have intensified our research into “Post Quantum Cryptography” seeking cryptographic schemes which remain secure even in the event that powerful quantum computers are built.

The security of public key cryptography traditionally rests on the difficulty of the solution of certain mathematical problems. Today, the following alternatives to the factorization and discrete logarithm problems are discussed in depth: the decoding problem, the shortest and closest vector problem in lattices, and the problem of solving large systems of multivariate quadratic equations. It is conjectured that quantum computers offer little advantage if we try to solve these problems efficiently.

9.3 New mathematical problems

Let us look at these alternatives a little more closely. The first encryption scheme based on the decoding problem was proposed by McEliece [?]. The background: Error-correcting codes are used to transmit or store electronic data in such a way that they remain undistorted even if a small number of bits are changed in transit or on the storage media. This property is used in, e.g., compact discs (CDs). The data on a CD can be reconstructed even if the disc has been slightly scratched.

In a code-based encryption scheme a message is encrypted by adding a fixed number of errors to (i.e. flipping a fixed numbers of bits of) the encoded message. Decoding requires knowledge of a suitable decoding procedure which eliminates these errors efficiently. This method is the secret key. Code-based encryption is in general very efficient. At the moment, research focus on the question which codes lead to secure encryption schemes with keys which are as small as possible.

Encryption on the basis of lattice problems is very similar to that on the basis of error-correcting codes. Lattices are regular structures of points in space. For instance, the points where the lines on squared paper cross form a two-dimensional lattice. For cryptographic usage, the dimension of the lattices is chosen to be much larger. Encryption works as follows: The plain-text is used to construct a lattice point which is then slightly distorted in such a way that it is no longer a lattice point, but close to one. Whoever knows a secret about the lattice is able to find this lattice point in the vicinity of the given point in space. The lattice point in turn yields the plain text. A particularly efficient lattice based encryption scheme is NTRU Encrypt (www.ntru.com). However, because NTRU was introduced fairly recently (in 1998), and its specification underwent several changes due to a variety of attacks, more cryptanalytic scrutiny is required to achieve confidence in its security.

9.4 New signatures

In 1979, Ralph Merkle proposed a remarkable framework for new signature schemes in his PhD thesis [?]. Contrary to all other signature schemes, its security does not rest on the difficulty of a number-theoretic, algebraic or geometric problem. The only thing it requires is something which other signature schemes need anyway: a cryptographically secure hash function and a secure pseudo-random number generator. Each new hash function leads to a new signature algorithm. In consequence, the Merkle scheme has the potential to solve the problem of long-term availability of digital signature schemes.

Merkle uses in his construction so-called One-Time Signatures: Each new signature requires a new signing key and a new verification key. The idea Merkle had was to reduce the validity of many verification keys using a hash tree to the validity of a unique public hash value. When generating keys for the Merkle scheme one has to determine the number of signatures one can make with it in advance. For a long time this seemed a significant disadvantage. In [?], however, a variant of Merkle's scheme was proposed which allows to compute 2^{40} signatures with a single key pair.

9.5 Quantum cryptography – a way out of the impasse?

From the point of view of today's state of the art of cryptography, the problem of long-term confidentiality remains unsolved: There is *no* practical method to protect the confidentiality of an encrypted message over a very long period of time.

One way out of that dilemma may be to employ quantum cryptography: it allows for key agreement schemes (of very long keys for one-time pads) whose security is guaranteed by the laws of quantum mechanics, cf., e.g., [?]. At the moment, however, quantum cryptography is still rather inefficient, and it is unclear which cryptographic functionalities can be implemented on top of it.

9.6 Conclusion

What's on the balance sheet of today's cryptography? We have good tools to ensure short and medium term security. Software developers can employ these tools in their applications with good conscience as long as they make sure that components can quickly be exchanged when they become insecure.

In order to guarantee IT security for the future, too, we need to prepare a portfolio of secure cryptographic schemes. This portfolio needs to contain schemes which are suitable for the world of ubiquitous computing with many less powerful computers. It also needs to contain schemes which remain secure in the event that powerful quantum computers are built. Several promising candidates have been discussed in this article. They need to be studied carefully and prepared for use in everyday scenarios. The question how to ensure long-term confidentiality remains an important open research problem upon which cryptographic research should focus.

Appendix A

Appendix

- 1 CrypTool 1 Menu Tree
- 2 CrypTool 2 Templates
- 3 JCrypTool Functions
- 4 CrypTool-Online Functions
- 5 Bibliography of Movies and Fictional Literature with Relation to Cryptography
- 6 Learning Tool for Elementary Number Theory
- 7 Using Sage with this Script
- 8 Authors of the CrypTool Script

A.1 CrypTool 1 Menus

This appendix contains at the following page the complete menu tree of CrypTool version 1.4.31¹.

The main menu of CT1 contains both, generic service functions in the six main menu items

- File
- Edit
- View
- Options
- Window
- Help,

and the actual crypto functions in the following four main menus

- Encrypt / Decrypt
- Digital Signature / PKI
- Individual Procedures
- Analysis.

Within **Individual Procedures** you find visualizations of single algorithms and of protocols. Some procedures are implemented both for a fast performance (mostly under the main menu **Encrypt/Decrypt**) and for a step-by-step visualization.

Which of the menu items in CrypTool 1 are active (that means not greyed), depends on the type of the currently active document window: The brute-force analysis for DES e. g. is only available, if the active window is opened in the hexadecimal view. On the other hand the menu item “Generate Random Numbers...” is always available (even if no document is opened).

¹Since 2010, changes for the stable CrypTool 1 (**CT1**) focus mainly on bugfixes. Many new developments go into the two successor versions CrypTool 2 (**CT2**) and JCrypTool (**JCT**):

- Web site CT2: <http://www.cryptool.org/en/ct2-documentation-en>
- Web site JCT: <http://www.cryptool.org/en/jct-volunteer-en>

These successor versions are currently (Nov. 2012) still beta versions; but they are close to their first release and they are stable enough since quite a while to be used by end-users.

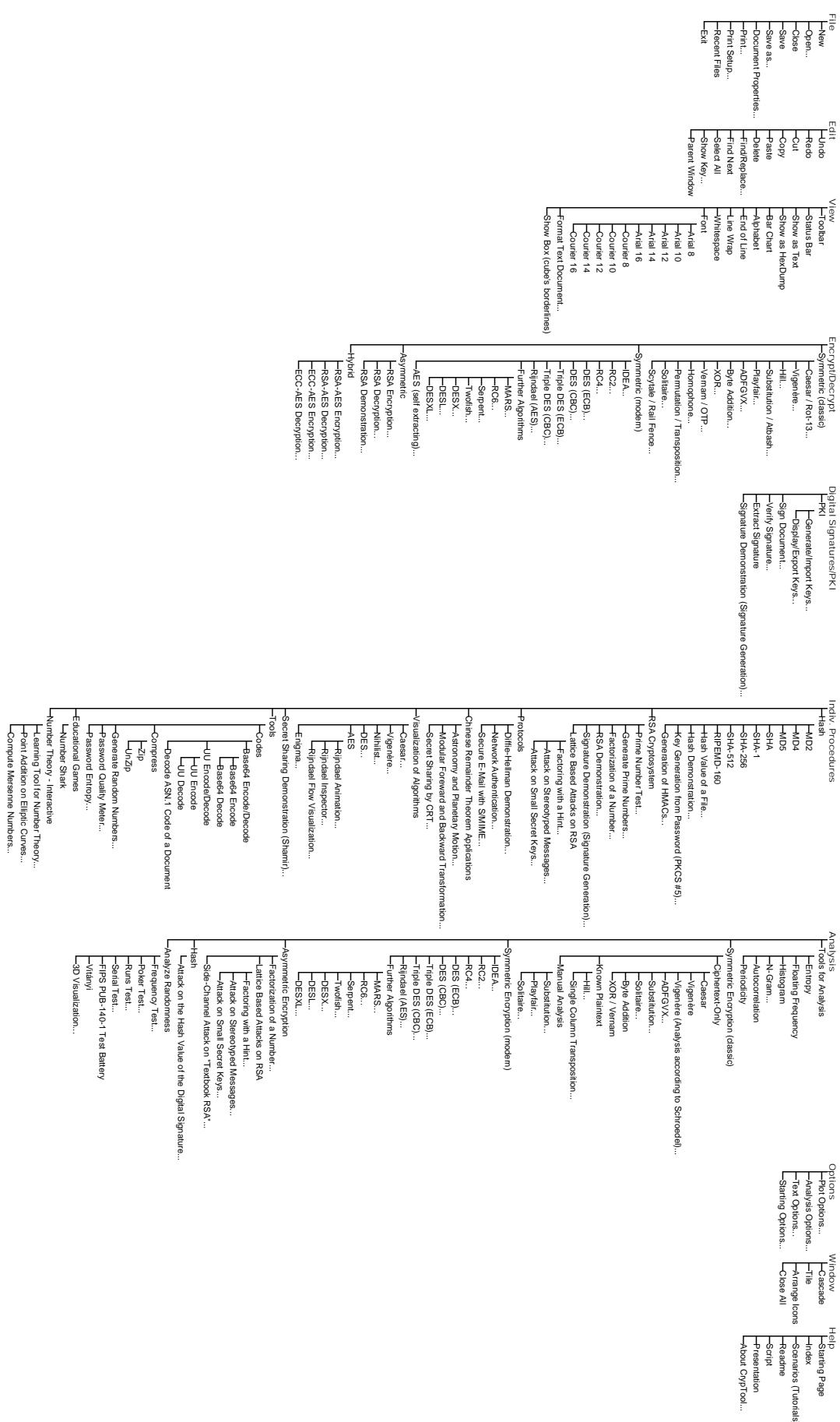


Figure A.1: Complete overview of the menu tree of CT1 (CrypTool 1.4.31)

A.2 CrypTool 2 Templates

This appendix contains on the following pages the complete tree with all templates in CT2.² When you start CT2 it first shows the Startcenter.

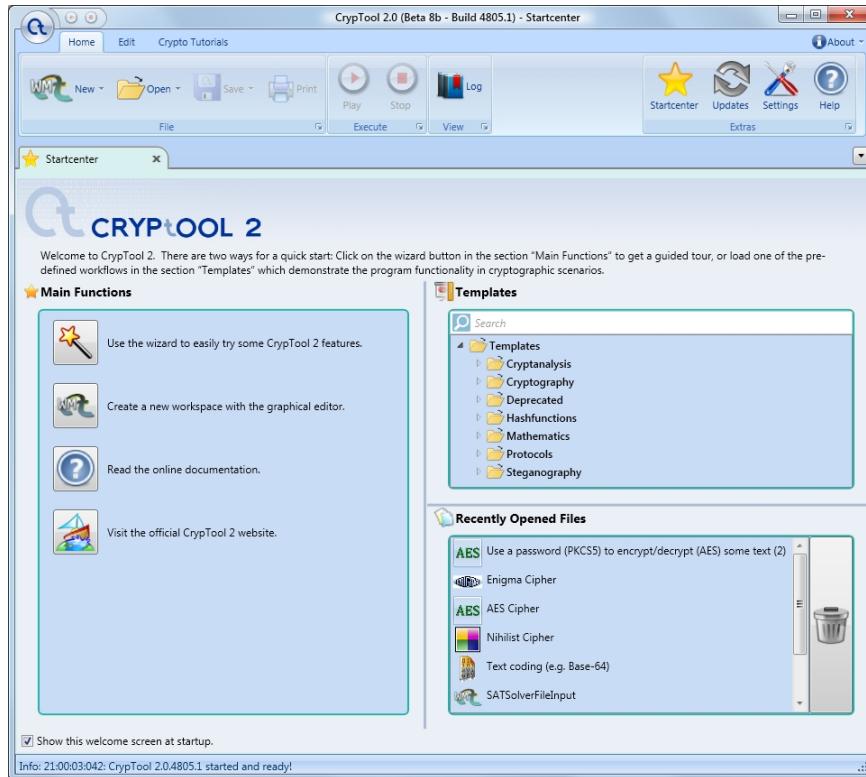


Figure A.2: Startcenter in CT2 (Beta 8b, May 2012)

Within the Startcenter you have three main ways to use the implemented functionality:

- via the Wizard, which leads you to the functions.
- via the Workspace, where you can compose the components (e.g. an encryption method, a text input function, ...) by yourself according to the visual programming concept.
- via the template tree, which offers ready-to-run workflows.

The Wizard asks questions so you can get to the desired scenarios (e.g. base64 coding) and then runs the according functions. You can afterwards save the chosen scenario as a normal template including the entries you used.

The empty workspace allows to drag on it any component from the navigation bar on the left and connect these components in the way you like. Most of the crypto functionality implemented in CT2 is offered using these components (e.g. Enigma, AES).

The template tree contains at least one template for each implemented component. The offered templates contain read-to-run workflows. If you change e.g. within the AES template your input, you can see at once, how the output is modified dynamically (e.g. adding another block via padding, influence of the chosen chaining, ...).

²You can find further information about CT2 at: <http://www.cryptool.org/en/ct2-documentation-en>

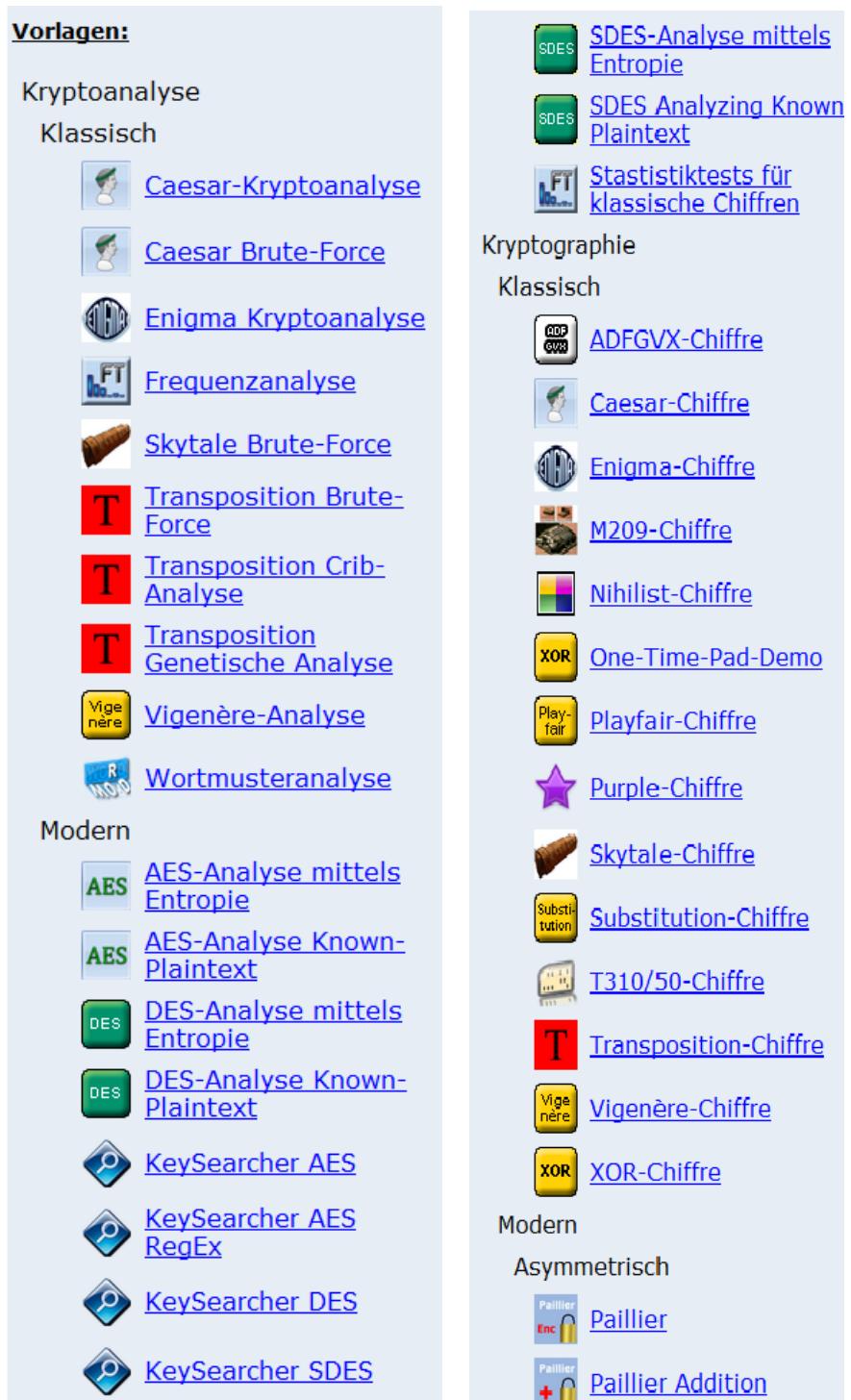


Figure A.3: Screenshot of the template tree of CT2 (NB4882.1, July 2012), Part 1

A.3 JCrypTool Functions

On the following pages this appendix contains a list of all functions in JCrypTool.³

When you start JCT the first time it shows the Welcome window.

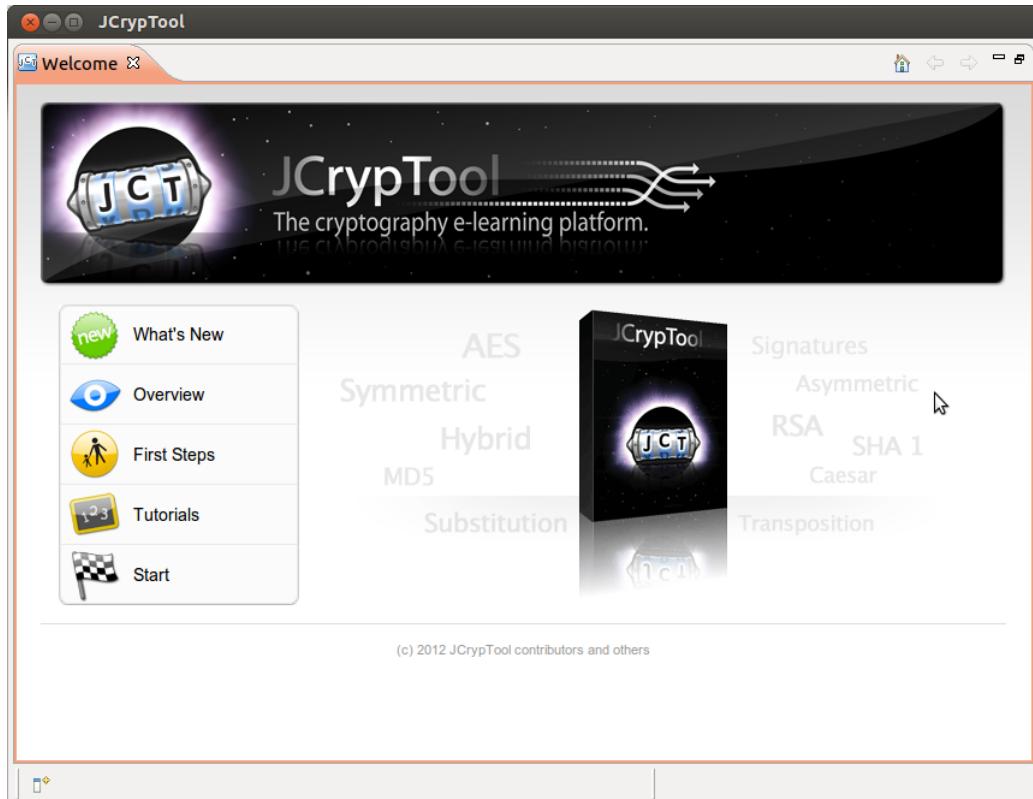


Figure A.4: Welcome screenshot in JCT (RC6, July 2012)

After pressing “Start” you can directly use the different functions. The functions implemented in JCT are presented in two different perspectives:

- Default Perspective
- Functional Perspective

All functions of the **Default Perspective** can be found both in the menus and in the navigation bar called “Crypto Explorer” (at the right side). The Default Perspective contains all important methods like classic transposition or modern AES, and many visualizations (e.g. Diffie-Hellman key exchange or calculations on elliptic curves).

All functions of the **Functional Perspective** can be found in the navigation bar called “Algorithms” (in this perspective also at the right side). The Functional Perspective contains all detail settings of the various algorithms, it especially offers post-quantum computing algorithms.

³You can find further information about JCT at: <http://www.cryptoool.org/en/jct-volunteer-en>
This list was generated using the CT Portal website.

Function	JCryptTool*	Path in JCT
ADFGVX	D	Algorithms\ Classic\ ADFGVX
Autokey Vigenère	D	Algorithms\ Classic\ Autokey-Vigenère
Bifid	D	Algorithms\ Classic\ Bifid
Caesar	D	Algorithms\ Classic\ Caesar
Double Box	D	Algorithms\ Classic\ Double Box
Playfair	D	Algorithms\ Classic\ Playfair
Substitution	D	Algorithms\ Classic\ Substitution
Transposition	D	Algorithms\ Classic\ Transposition
Vigenère	D	Algorithms\ Classic\ Vigenère
XOR	D	Algorithms\ Classic\ XOR
3DES	F	Algorithms\ Block Ciphers\ DESede
AES	D\F	Algorithms\ Symmetric\ AES Algorithms\ Block Ciphers\ Rijndael
Camellia	F	Algorithms\ Block Ciphers\ Camellia
Dragon	D	Algorithms\ Symmetric\ Dragon
ECIES	F	Algorithms\ Hybrid Ciphers\ ECIES
ElGamal	D\F	Algorithms\ Asymmetric\ ElGamal Algorithms\ Asymmetric Block Ciphers\ ElGamal
IDEA	D\F	Algorithms\ Symmetric\ IDEA Algorithms\ Block Ciphers\ IDEA
LFSR	D	Algorithms\ Symmetric\ LFSR
MARS	F	Algorithms\ Block Ciphers\ MARS
McEliece Fujisaki	F	Algorithms\ Hybrid Ciphers\ McElieceFujisakiCipher
McEliece Kobara Imai	F	Algorithms\ Hybrid Ciphers\ McElieceKobaraImaiCipher
McEliece PKCS	F	Algorithms\ Asymmetric Block Ciphers\ McEliecePKCS
McEliece Pointcheval	F	Algorithms\ Hybrid Ciphers\ McEliecePointchevalCipher
MeRSA	F	Algorithms\ Asymmetric Block Ciphers\ MeRSA
Misty-1 (Kasumi)	F	Algorithms\ Block Ciphers\ Misty1
MpRSA	F	Algorithms\ Asymmetric Block Ciphers\ MpRSA
Niederreiter	F	Algorithms\ Asymmetric Block Ciphers\ Niederreiter
RC2	F	Algorithms\ Block Ciphers\ RC2
RC5	F	Algorithms\ Block Ciphers\ RC5
RC6	D\F	Algorithms\ Symmetric\ RC6 Algorithms\ Block Ciphers\ RC6
RSA	D\F	Algorithms\ Asymmetric\ RSA Algorithms\ Asymmetric Block Ciphers\ RSA_PKCS1_v1_5 Algorithms\ Asymmetric Block Ciphers\ RSA_PKCS1_v2_1
SAFER+	F	Algorithms\ Block Ciphers\ SAFER+
SAFER++	F	Algorithms\ Block Ciphers\ SAFER++
Serpent	F	Algorithms\ Block Ciphers\ Serpent
Shacal	F	Algorithms\ Block Ciphers\ Shacal
Shacal2	F	Algorithms\ Block Ciphers\ Shacal2
Twofish	F	Algorithms\ Block Ciphers\ Twofish
XML-Security	D	Algorithms\ XML Security\
CBC MACs	F	Algorithms\ Message Authentication Codes\ CBCMac
Cipher-based MACs (CMAC)	F	Algorithms\ Message Authentication Codes\ CMac
DHA-256	F	Algorithms\ Message Digests\ DHA256
FORK-256	F	Algorithms\ Message Digests\ FORK256
HMACs	D\F	Algorithms\ MAC\ HMacMD5 Algorithms\ Message Authentication Codes\ Hmac
MD4	F	Algorithms\ Message Digests\ MD4
MD5	D\F	Algorithms\ Hash\ MD5 Algorithms\ Message Digests\ MD5
PKCS#5	F	Algorithms\ Password-based ciphers\
RIPEMD	F	Algorithms\ Message Digests\ RIPEMD
SHA3 candidates	D	Algorithms\ Hash\ SHA3-Candidates

Figure A.5: Screenshot of the functions of JCT (RC6, July 2012), Part 1

Function	JCrypTool*	Path in JCT
SHA family	D\F	Algorithms\ Hash\ SHA Algorithms\ Message Digests\ SHA
Tiger	F	Algorithms\ Message Digests\ Tiger
Two-Track-MAC	F	Algorithms\ Message Authentication Codes\ TwoTrackMac\
VSH	F	Algorithms\ Message Digests\ VSH
CMSS	F	Algorithms\ Signatures\ CMSSSignature\
DSA	D\F	Algorithms\ Signature\ DSA Algorithms\ Signatures\ DSASignature\
ECDSA	F	Algorithms\ Signatures\ ECDSASignature\
GMSS	F	Algorithms\ Signatures\ GMSSSignature\
IQDSA	F	Algorithms\ Signatures\ IQDSASignature\
IQGQ	F	Algorithms\ Signatures\ IQGQSignature\
IQRDSA	F	Algorithms\ Signatures\ IQRDSSignature\
Merkle OTS	F	Algorithms\ Signatures\ MerkleOTSSignature\
MeRSA	F	Algorithms\ Signatures\ MeRSA
MpRSA	F	Algorithms\ Signatures\ MpRSA
Niederreiter CFS	F	Algorithms\ Signatures\ NiederreiterCFS
PKCS#1 (RSA Signatures)	F	Algorithms\ Signatures\ RSASignaturePKCS1v15\
RSASSA-PSS	F	Algorithms\ Signatures\ RSASSA-PSS
SHA1 ECNR	F	Algorithms\ Signatures\ SHA1\ECNR
SSL (MD5 and SHA1 with RSA)	F	Algorithms\ Signatures\ SHA1\SSL_MD5andSHA1withRSA
Blum Blum Shub (B.B.S.)	F	Algorithms\ Pseudo Random Number Generators\ BBSRandom
Elliptic Curve PRNG	F	Algorithms\ Pseudo Random Number Generators\ ECPRNG
SHA1	D	Algorithms\ Random Number Generator\ SHA1
SHA1 PRNG	F	Algorithms\ Pseudo Random Number Generators\ SHA1PRNG
Entropy	D	Analysis\ Entropy Analysis
Frequency Test	D	Analysis\ Frequency Analysis
Friedman Test	D	Analysis\ Friedman Test
Transposition Analysis	D	Analysis\ Transposition Analysis
Vigenère Analysis	D	Analysis\ Vigenère Breaker
Ant Colony Optimization	D	Visuals\ Ant Colony Optimization
Chinese Remainder Theorem	D	Visuals\ Chinese Remainder Theorem
Differential Power Analysis	D	Visuals\ Differential Power Analysis \ Double and Add
Diffie-Hellman Key Exchange	D	Visuals\ Diffie-Hellman Key Exchange (EC)
ElGamal	D	Visuals\ ElGamal Cryptosystem
Elliptic Curve Cryptography	D	Visuals\ ECC Demonstration
Extended Euclidian	D	Visuals\ Extended Euclidian
Feige Fiat Shamir	D	Visuals\ Feige Fiat Shamir
Fiat Shamir	D	Visuals\ Fiat Shamir
Graph Isomorphism	D	Visuals\ Graph Isomorphism
Grille	D	Visuals\ Grille
Homomorphic Encryption	D	Visuals\ Homomorphic Encryption
Kleptography	D	Visuals\ Kleptography
Magic Door	D	Visuals\ Magic Door
Multipartite Key Exchange	D	Visuals\ Multipartite Key Exchange
RSA	D	Visuals\ RSA Cryptosystem
Shamir's Secret Sharing	D	Visuals\ Shamir's Secret Sharing
Verifiable Secret Sharing (VSS)	D	Visuals\ Verifiable Secret Sharing
Simple Power Analysis	D	Visuals\ Simple Power Analysis \ Square and Multiply
Viterbi	D	Visuals\ Viterbi
XML-Security	D	Algorithms\ XML Security\
Number Shark	D	Games\ Number Shark
Sudoku Solver	D	Games\Sudoku
*	D = Default Perspective F = Functional Perspective	

Figure A.6: Screenshot of the functions of JCT (RC6, July 2012), Part 2

A.4 CrypTool-Online Functions

This appendix contains a list of all functions in CrypTool-Online (CTO).⁴

The following screenshot shows the crypto functions implemented on CTO:

⁴You can find further information about CTO at: www.cryptool-online.org
This list was generated using the functions list at the CT Portal website:
<http://www.cryptool.org/ctp-documentation-en/ctp-functions-en>

Function	JCryptTool*	Path in JCT
ADFGVX	D	Algorithms\ Classic\ ADFGVX
Autokey Vigenère	D	Algorithms\ Classic\ Autokey-Vigenère
Bifid	D	Algorithms\ Classic\ Bifid
Caesar	D	Algorithms\ Classic\ Caesar
Double Box	D	Algorithms\ Classic\ Double Box
Playfair	D	Algorithms\ Classic\ Playfair
Substitution	D	Algorithms\ Classic\ Substitution
Transposition	D	Algorithms\ Classic\ Transposition
Vigenère	D	Algorithms\ Classic\ Vigenère
XOR	D	Algorithms\ Classic\ XOR
3DES	F	Algorithms\ Block Ciphers\ DESede
AES	D\F	Algorithms\ Symmetric\ AES Algorithms\ Block Ciphers\ Rijndael
Camellia	F	Algorithms\ Block Ciphers\ Camellia
Dragon	D	Algorithms\ Symmetric\ Dragon
ECIES	F	Algorithms\ Hybrid Ciphers\ ECIES
ElGamal	D\F	Algorithms\ Asymmetric\ ElGamal Algorithms\ Asymmetric Block Ciphers\ ElGamal
IDEA	D\F	Algorithms\ Symmetric\ IDEA Algorithms\ Block Ciphers\ IDEA
LFSR	D	Algorithms\ Symmetric\ LFSR
MARS	F	Algorithms\ Block Ciphers\ MARS
McEliece Fujisaki	F	Algorithms\ Hybrid Ciphers\ McElieceFujisakiCipher
McEliece Kobara Imai	F	Algorithms\ Hybrid Ciphers\ McElieceKobaraImaiCipher
McEliece PKCS	F	Algorithms\ Asymmetric Block Ciphers\ McEliecePKCS
McEliece Pointcheval	F	Algorithms\ Hybrid Ciphers\ McEliecePointchevalCipher
MeRSA	F	Algorithms\ Asymmetric Block Ciphers\ MeRSA
Misty-1 (Kasumi)	F	Algorithms\ Block Ciphers\ Misty1
MpRSA	F	Algorithms\ Asymmetric Block Ciphers\ MpRSA
Niederreiter	F	Algorithms\ Asymmetric Block Ciphers\ Niederreiter
RC2	F	Algorithms\ Block Ciphers\ RC2
RC5	F	Algorithms\ Block Ciphers\ RC5
RC6	D\F	Algorithms\ Symmetric\ RC6 Algorithms\ Block Ciphers\ RC6
RSA	D\F	Algorithms\ Asymmetric\ RSA Algorithms\ Asymmetric Block Ciphers\ RSA_PKCS1_v1_5 Algorithms\ Asymmetric Block Ciphers\ RSA_PKCS1_v2_1
SAFER+	F	Algorithms\ Block Ciphers\ SAFER+
SAFER++	F	Algorithms\ Block Ciphers\ SAFER++
Serpent	F	Algorithms\ Block Ciphers\ Serpent
Shacal	F	Algorithms\ Block Ciphers\ Shacal
Shacal2	F	Algorithms\ Block Ciphers\ Shacal2
Twofish	F	Algorithms\ Block Ciphers\ Twofish
XML-Security	D	Algorithms\ XML Security\
CBC MACs	F	Algorithms\ Message Authentication Codes\ CBCMac
Cipher-based MACs (CMAC)	F	Algorithms\ Message Authentication Codes\ CMac
DHA-256	F	Algorithms\ Message Digests\ DHA256
FORK-256	F	Algorithms\ Message Digests\ FORK256
HMACs	D\F	Algorithms\ MAC\ HMacMD5 Algorithms\ Message Authentication Codes\ Hmac
MD4	F	Algorithms\ Message Digests\ MD4
MD5	D\F	Algorithms\ Hash\ MD5 Algorithms\ Message Digests\ MD5
PKCS#5	F	Algorithms\ Password-based ciphers\
RIPEMD	F	Algorithms\ Message Digests\ RIPEMD
SHA3 candidates	D	Algorithms\ Hash\ SHA3-Candidates

Figure A.7: Screenshot of the functions of CTO (November 2012)

A.5 Movies and Fictional Literature with Relation to Cryptography

Cryptographic applications – classical as well as modern ones – have been used in literature and movies. In some media they are only mentioned and are a pure admixture; in others they play a primary role and are explained in detail; and sometimes the purpose of the story, which forms the framework, is primarily to transport this knowledge and achieve better motivation. Here is the beginning of an overview.

A.5.1 For Grownups and Teenagers

[Poe1843] Edgar Allan Poe,

The Gold Bug, 1843.

In this short story Poe tells as first-person narrator about his acquaintanceship with the curious Mr. Legrand. They detect a fabulous treasure via a gold bug and a vellum found at the coast of New England.

The cipher consists of 203 cryptic symbols and it proves to be a general monoalphabetic substitution cipher (see chapter 2.2.1). The story tells how they solve the riddle step by step using a combination of semantic and syntax analysis (frequency analysis of single letters in English texts).

In this novel the code breaker Legrand says the famous statement: “Yet it may be roundly asserted that human ingenuity cannot concoct a cipher which human ingenuity cannot resolve – given the according dedication.”

[Verne1885] Jules Verne,

Mathias Sandorf, 1885.

This is one of the most famous novels of the French author Jules Verne (1828-1905), who was called “Father of Science fiction”.

In “Mathias Sandorf” he tells the story of the freedom fighter Earl Sandorf, who is betrayed to the police, but finally he can escape.

The whistle-blowing worked, because his enemies captured and decrypted a secret message sent to him. For decryption they needed a special grille, which they stole from him. This turning grille was a quadratic piece of jig with 6x6 squares, of which 1/4 (nine) were holes (see the turning grille in chapter 2.1.1).

[Kipling1901] Rudyard Kipling,

Kim, 1901.

Rob Slade’s review⁵ of this novel says: “Kipling packed a great deal of information and concept into his stories, and in “Kim” we find The Great Game: espionage and spying. Within the first twenty pages we have authentication by something you have, denial of service, impersonation, stealth, masquerade, role-based authorization (with ad hoc authentication by something you know), eavesdropping, and trust based on data integrity. Later on we get contingency planning against theft and cryptography with key changes.” The book is out of copyright⁶.

⁵See <http://catless.ncl.ac.uk/Risks/24.49.html#subj12>.

⁶You can read it at:

<http://whitewolf.newcastle.edu.au/words/authors/K/KiplingRudyard/prose/Kim/index.html>,

[Doyle1905] Arthur Conan Doyle,

The Adventure of the Dancing Men, 1905.

In this Sherlock Holmes short story (first published in 1903 in the “Strand Magazine”, and then in 1905 in the collection “The Return of Sherlock Holmes” the first time in book-form) Sherlock Holmes has to solve a cipher which at first glance looks like a harmless kid’s picture.

But it proves to be the monoalphabetic substitution cipher (see chapter 2.2.1) of the criminal Abe Slaney. Sherlock Holmes solves the riddle using frequency analysis.

[Sayers1932] Dorothy L. Sayers,

Have his carcase, Harper/Victor Gollancz Ltd., 1932.

In this novel the writer Harriet Vane finds a dead body at the beach. The police believe the death is suicide. Harriet Vane and the elegant amateur sleuth Lord Peter Wimsey together clear of the disgusting murder in this second of Sayers’s famous Harriet Vane mystery series.

This requires to solve a cryptogram. Surprisingly the novel not only describes the Playfair cipher in detail, but also the cryptanalysis of this cipher (see Playfair in chapter 2.2.3).

[Simmel1970] Johannes Mario Simmel,

And Jimmy went to the Rainbow (original title: *Und Jimmy ging zum Regenbogen*), Knaur Verlag, 1970.

The novel plays between 19938 and 1967 in Vienna. The main character Manual Aranda uncovers step by step the past of his murdered father. Important for the plot is an encrypted manuscript, which is decrypted in chapter 33. In the novel the cipher is called “25-fold Caesar cipher”. It is actually a Vigenère cipher with a 25 character key.

A movie of the novel appeared in 1971.

[Crichton1987] Michael Crichton,

Sphere, Pan Books, 1987.

A team of different scientists is send to the ground of the ocean in order to investigate a highly developed 900 m long space ship. The human peculiarities and psychological problems of the researchers surface more and more, because of life threatening events and isolation. There are many mysteries: While the space ship lies on the ground for 300 years, it has English markings and a life of its own, and materializing of the researcher’s imaginations appear. On a computer screen a cipher text appears, which is completely printed in the book. The genius mathematician Harry deciphers the simple helical substitution code.

[Seed1990] Directed by Paul Seed,

House of Cards, 1990.

In this movie Ruth tries to solve the secret, which made her daughter fall silent. Here two young people suffering from autism communicate via 5- and 6-digit primes (see chapter 3). After more than 1 hour the movie contains the following encrypted two series of primes:

21, 383; 176, 081; 18, 199; 113, 933; 150, 377; 304, 523; 113, 933
193, 877; 737, 683; 117, 881; 193, 877

[Robinson1992] Directed by Phil Alden Robinson,
Sneakers, Universal Pictures Film, 1992.

In this movie the “sneakers”, computer experts under their boss Martin Bishop, try to get back the deciphering box SETEC from the “bad guys”. SETEC, invented by a genius mathematician before he was killed, allows to decrypt all codes from any nation.

In the movie the code is not described in any way⁷.

[Baldacci1997] David Baldacci,
Total Control, Mass Market Paperback, 1997.

Jason Archer, executive with a technology company suddenly disappears. Sidney Archer tries to find out about her husband’s surprising death. She gets a clue how the global financial system is abused and that the real control belongs to those with the most money. Here even good passwords don’t help ...

[Natali1997] Directed by Vincenzo Natali,
Cube, Mehra Meh Film, 1997.

In this Canadian low-budget-movie 7 complete strangers of widely varying personality characteristics are involuntarily placed in an kafkaesque maze of cubical rooms containing deadly traps.

To get out the persons have to move through these rooms. To find out which rooms are dangerous, mathematics is crucial: Each cubic room has at its entrance a numerical marking consisting of three sets of three digits. First they deduce that all rooms marked at their entrance with at least one prime number are trapped. Later it comes out that a trapped room can also be marked by a number which is a power of a prime (so traps are p^n , e.g. $128 = 2^7$ or $101 = 101^1 = \text{prime}$, but not $517 = 11 * 47$).

[Becker1998] Directed by Harold Becker,
Mercury Rising, Universal Pictures Film, 1998.

The NSA developed a new cipher, which is pretended to be uncrackable by humans and computers. To test its reliability some programmers hide a message encrypted with this cipher in a puzzle magazine.

Simon, a nine year old autistic boy, cracks the code. Instead of fixing the code, a government agent sends a killer. FBI agent Art Jeffries (Bruce Willis) protects the boy and sets a snare for the killers.

The code is not described in any way.

[Brown1998] Dan Brown,
Digital Fortress, E-Book, 1998.

Dan Brown’s first novel was published in 1998 as e-book, but it was largely unsuccessful then.

The National Security Agency (NSA) uses a huge computer, which enables it to decrypt all messages (needless to say only of criminals and terrorists) within minutes even if they use the most modern encryption methods.

⁷Leonard Adleman (the ”A” within RSA) worked as mathematical consultant for “Sneakers”. He describes the funny story about his contribution at his homepage <http://www.usc.edu/dept/molecular-science/fm-sneakers.htm>. It is assumed that the cipher used everywhere is RSA. According to that within the chip a fast, unknown factorization method is implemented.

An apostate employee invents an unbreakable code and his computer program Diabolus forces the super computer to do self destructing operations. The plot, where also the beautiful computer expert Susan Fletcher has a role, is rather predictable.

The idea, that the NSA or another secret service is able to decrypt any code, is currently popular on several authors: In “Digital Fortress” the super computer has 3 million processors – nevertheless from today’s sight this is by no means sufficient to hack modern ciphers.

[Elsner1999] Dr. C. Elsner,

The Dialogue of the Sisters, c’t, Heise, 1999.

In this short story, which is included in the CrypTool package as PDF file, the sisters confidentially communicate using a variant of RSA (see chapter 4.10 and the following). They are residents of a madhouse being under permanent surveillance.

[Stephenson1999] Neal Stephenson,

Cryptonomicon, Harper, 1999.

This very thick novel deals with cryptography both in WW2 and today. The two heroes from the 40ies are the excellent mathematician and cryptanalyst Lawrence Waterhouse, and the overeager and morphine addicted US marine Bobby Shaftoe. They both are members of the special allied unit 2702, which tries to hack the enemy’s communication codes and at the same time to hide the own existence.

This secretiveness also happens in the present plot, where the grandchildren of the war heroes – the dedicated programmer Randy Waterhouse and the beautiful Amy Shaftoe – team up.

Cryptonomicon is notably heavy for non-technical readers in parts. Several pages are spent explaining in detail some of the concepts behind cryptography. Stephenson added a detailed description of the Solitaire cipher (see chapter 2.4), a paper and pencil encryption algorithm developed by Bruce Schneier which is called “Pontifex” in the book. Another, modern algorithm called “Arethusa” is not explained in detail.

[Elsner2001] Dr. C. Elsner,

The Chinese Labyrinth, c’t, Heise, 2001.

In this short story, which is included in the CrypTool package as PDF file, Marco Polo has to solve problems from number theory within a competition to become a major consultant of the Great Khan. All solutions are included and explained.

[Colfer2001] Eoin Colfer,

Artemis Fowl, Viking, 2001.

In this book for young people the 12 year old Artemis, a genius thief, gets a copy of the top secret “Book of the Elfs”. After he decrypted it with his computer, he finds out things, men never should have known.

The used code is not described in detail or revealed.

[Howard2001] Directed Ron Howard,

A Beautiful Mind, 2001.

This is the film version of Sylvia Nasar’s biography of the game theorist John Nash. After the brilliant but asocial mathematician accepts secret work in cryptography, his life takes

a turn to the nightmarish. His irresistible urge to solve problems becomes a danger for himself and his family. Nash is – within his belief – a most important hacker working for the government.

Details of his way analysing code are not described in any way.

[Apted2001] Directed by Michael Apted,

Enigma, 2001.

This is the film version of Robert Harris’ “historical fiction” *Enigma* (Hutchinson, London, 1995) about the World War II code-breaking work at Bletchley Park in early 1943, when the actual inventor of the analysis Alan Turing (after Polish pre-work) already was in the US. So the fictional mathematician Tom Jericho is the lead character in this spy-thriller. Details of his way analysing the code are not described.

[Isau2003] Ralf Isau,

The Museum of the stolen memories (original title: *Das Museum der gestohlenen Erinnerungen*), Thienemann-Verlag, 1997/2003.

In this exciting novel the last part of the oracle can only be solved with the joined help of the computer community.

The book got several awards and exists in 8 different languages, but not in English yet.

[Brown2003] Dan Brown,

The Da Vinci Code, Doubleday, 2003.

The director of the Louvre is found murdered in his museum in front of a picture of Leonardo da Vinci. And the symbol researcher Robert Langdon is involved in a conspiracy. The plot mentions different classic codes (substitution like Caesar or Vigenère, as well as transposition and number codes). Also there are hints about Schneier and the sunflower. The second part of the book contains a lot of theological considerations.

This book has become one of the most widely read books of all time.

[Hill2003] Tobias Hill,

The Cryptographer, Faber & Faber, 2003.

London 2021: The company SoftMark developed and establish an electronic currency, which guarantees highest security standards by an unbreakable code. The inventor and company founder, called the cryptographer because of his mathematical talent, has become the richest man in the world. But the code was hacked, and in a worldwide economic crisis his company goes bankrupt. Additionally the tax investigator Anna Moore is set on him.

[McBain2004] Scott McBain,

Final Solution, manuscript not published by Harper Collins, 2004 (German version has been published in 2005).

In a near future politicians, chiefs of military and secret services of many different countries take over all the power. With a giant computer network called “Mother” and complete surveillance they want to cement their power and commercialisation of life forever. Humans are only assessed according to their credit rating and globally acting companies elude of any democratic control. Within the thriller the obvious injustice, but also the

realistic likelihood of this development are considered again and again.

With the help of a cryptographer a code to destroy was built into the super computer “Mother”: In a race several people try to start the deactivation (Lars Pedersen, Oswald Plevy, the female American president, the British prime minister and an unknown Finish named Pia, who wants to take revenge for the death of her brother). On the opposite side a killing group acts under the special guidance of the British foreign minister and the boss of the CIA.

[Preston2005] Douglas Preston,

Tyrannosaurus Canyon, Forge Books, 2005.

A very exciting thriller which also struggles with the question why the dinosaurs died off.

Archeologist Stem Weathers is shot in a canyon. Before his murderer appears he gives his notebook to Tom Broadbent, a local animal doctor, coming by accidentally.

The notebook contains on 60 pages only digits. Therefore Tom takes it to Wyman Ford an ex-CIA cryptanalyst, who now lives in a nearby abbey, after his wife was killed in action. Wyman first declines and says that self-invented code are “idiot ciphers” – devised by an idiot and easily crackable by each idiot. The notebook then proves to be not that easy. After intensive analysis he finds out that the digits are no code but the output of an earth radar device showing the picture of a well-preserved T.rex.

After around 250 pages of endless chases a surprising turn comes up: Masago, head of a so-called black-detachment unit of the CIA. He explains: New weapons invented once always have been used. Mankind will kill herself, but its his task to postpone that as far as possible. As head of the LS480 department he will prevent by any means possible that terrorists get any new dangerous biological weapon.

When scanning the dead body of Weathers the murder only found some rock cuttings he took. These rocks are investigated by a young researcher named Melody Crookshank although she doesn't know where the rock cuttings come from. She finds within them a very special kind of virus apparently coming from outer-space.

[Burger2006] Wolfgang Burger,

Heidelberg Lies (original title: Heidelberger Lügen), Piper, 2006.

This detective story playing in the Rhein-Neckar area in Germany has several independent strands and local stories, but mainly it is about Kriminalrat Gerlach from Heidelberg. On page 207 f. the cryptographic reference for one strand is shortly explained: The soldier Hörrle had copied circuit diagrams of a new digital NATO decryption device and the murdered man had tried to sell his perceptions to China.

[Twinig2006] James Twinig,

The Black Sun, HarperCollins, 2006.

A history-based thriller with some artificially constructed elements, dealing also with a treasure hunt to get the hidden uranium of the nazis, and naturally the future of the world depends on today's bad guys being stopped in time ...

Heros are Tom Kirk, a London-based ex-CIA agent and former professional art thief, and Dominique de Lecourt, who loves challenges including riddles and codes.

The only cryptographic parts are a “Sprungcode” (the criminals use this method to communicate via newsletter adverts), steganography (used to hide the Enigma key), and an

Enigma message (containing the encrypted coordinates of the treasure).

At the beginning of the plot an Enigma device is stolen with high efforts which is necessary to let the story play in the constructed way. But in the reality today such a theft is completely needless, as there are great software emulators for the Enigma ...

[**Vidal2006**] Agustin Sanchez Vidal,
Kryptum, Dtv, 2006.

The first novel of the Spanish professor of art history has some similarities with Dan Brown's "The Da Vinci Code" from 2003, but allegedly Vidal started his writing of the novel already in 1996. Vidal's novel is a mixture between historic adventure and mystery thriller. It was a huge success in Spain and Germany. There is currently no English version available.

In the year 1582 Raimundo Randa is waiting to be condemned to death – he was all life long trying to solve a mystery. This mystery is about a parchment with cryptic characters, where a unique power is behind. Around 400 years later the American scientist Sara Toledano is fascinated by this power until she vanishes in Antigua. Her colleague, the cryptographer David Calderon, and her daughter Rachel are searching for her and simultaneously they try to solve the code. But also secret organizations like the NSA chase after the secret of the "last key". They don't hesitate to kill for it.

[**Larsson2006**] Stieg Larsson,
Perdition (original title: *Flickan som lekte med elden*), 2006.

The author was posthumously awarded in 2006 with the Scandinavian thriller award. The super hero Lisbeth Salander uses PGP and occupies herself with mathematical riddles like the Fermat theorem.

[**Schroeder2008**] Rainer M. Schröder,
The Judas Documents (original title: *Die Judas-Papiere*), Arena, 2008.

In the year 1899 Lord Pembroke has three men and one woman in his grip. So they have to follow his order to try to decipher the encrypted messages in the notebook of his dead brother Mortimer and to find the missing gospel according to Judas, which could shock the whole of Christendom. The four people therefor have to solve riddles at many places in the world. The story explains some classic ciphers like Polybius and Freemason.

[**Eschbach2009**] Andreas Eschbach,
A King for Germany (original title: *Ein König für Deutschland*), Lübbe, 2009.

The novel deals with manipulations of electronic voting machines.

Vincent Merrit, a young US-American programmer, is blackmailed, to write such a programme. Beside commercially oriented blackmailers also Massively Multiplayer Online Role-Playing Games (MMORPGs) and Live Action Role Playing (LARP) have a role. Because Merrit assumed that his programme will be misused, he installed a trapdoor: If a party with the name VWM participates at the election, it automatically gets 95 % of the votes ...

The fictional story line is based on many verifiable and well researched facts, which are referenced in footnotes.

While the cryptographic protocols itself could be made secure, their implementation and their organisational management stays susceptible against misuse.

Currently there is no English translation of the book.

[**Juels2009**] Ari Juels,
Tetrakty, Emerald Bay Books, 2009.

The plot exposes the vulnerability of modern computer based identity, authenticity, and security interweaving modern cryptography with classical art and literature. Cryptographer and classicist Ambrose Jerusalem is a UC Berkeley graduate with a beautiful girlfriend and a comfortable future, until the NSA recruits him to track a strange pattern of computer break-ins. Many small pieces provide disturbing evidence that someone has broken RSA encryption. Even more bizarre, a secret cult of latter-day followers of Pythagoras, the great Greek mathematician and philosopher who believed reality could be understood only through a mystical system of numbers, appears to be behind the attacks.

[**Suarez2009**] Daniel Suarez,
Daemon, Penguin Books, 2009, 632 pages.

This is considered as one of the most exciting books during the last few years – it's a near-science-fiction thriller combining developments in the real world and possibilities coming from current research like from the Google X Lab (augmented reality head-mounted displays (HMD) like Google glass, self-driving cars, 3-D printers, ...) to a plausible story.

After the computer genius and game developer Matthew Sobol died a daemon starts acting in the internet, which seemingly ruthless manipulates and trains more and more humans and companies.

By ruling the data everybody seems to be a helpless victim. All the communication of his mercenary soldiers is affected by high-tech and encryption – also the communication between the distributed instances of his incarnation. Core is an MMORPG game (massive multiplayer online role-playing game) which reminds many of WoW. Here also encryption is used, e.g. to advertise the best players: m0wFG3PRCoJVTs7JcgBwsOXb3U7yPxBB

The plot is without redundancy, complex, manifold, very fascinating and with it critics about the plutocrats it also contains concrete social elements. The end is open. And the ideas seem to be realizable in the very next future ...

[**Suarez2010**] Daniel Suarez,
Freedom (TM), Penguin Books, 2010, 486 pages.

“The propulsive, shockingly plausible sequel to the bestseller *Daemon*” (see above). *Freedom (TM)* (*Daemon* #2) patches a number of holes the writer left in the first book. The prose is tighter, the descriptions more direct, the characters are fleshed out, especially Loki. Having laid the groundwork in “*Daemon*”, Suarez uses this foundation in order to explore a new concept of social organization based on empowering information technology and the reasoning why and how the battle runs between the old potentates and the *Daemon* society, which also evolves further already during the story. Cryptography is a natural part of modern technology and modern warfare as described in this book. The new society emerging in ”*Freedom (TM)*” is based on the darknet, an alternative to the internet using fast wireless meshes in order to increase the durability and availability of the network. Despite the story is shocking in some parts, it appears to be realistic and not far away from the parallel usage of modern technology integrated into our modern lives as a virtual world overlaying our real world.

[Burger2011] Wolfgang Burger,

The fifth murderer (original title: *Der fünfte Mörder*), Piper, 2011.

Location & time of the story: Germany / Heidelberg, 1990 - 2009. Episode 7 of the Alexander-Gerlach series. Inspector Alexander Gerlach became almost a victim of a bomb blast when the sport utility vehicle (SUV) of a Bulgarian panderer exploded. Gerlach starts investigating because he wants to prevent a gang warfare, but then his bosses call him off. When the journalist Machatschek supports Gerlach, he communicates with him only via Skype using an add-on encryption program which he believes is the most secure in the world.

[Elsberg2012] Marc Elsberg,

Blackout – Tomorrow its too late (original title: *Blackout – Morgen ist es zu spät*), Blanvalet, 2012, 800 pages.

At a could day in winter all power supply networks in Europe break down. Agencies, energy suppliers and security companies are in the dark and unable to solve the problem. The Italian computer scientist Piero Manzano believes that this caused by terrorists using hackers: All customers use since some years smart meters, electricity meters controlled by software which was manipulated. Despite the integrated security and encryption components they have been hacked, and are out of order by wrong control sequences. The terrifying consequences at various locations are described realistically and excitingly. And in the same way the reactions of the human beings ...

Remark 1: Further samples of cryptology in fictional literature can be found on the following German web page:

http://www.staff.uni-mainz.de/pommeren/Kryptologie99/Klassisch/1_Monoalph/Literat.html

For some older authors (e.g. Jules Verne, Karl May, Arthur Conan Doyle, Edgar Allan Poe) there are links to the original and relevant text pieces.

Remark 2: You can find title pages of some of these books on the web site of Tobias Schrödel, who collects classic books about cryptography:

http://tobiasschroedel.com/crypto_books.php

Remark 3: If you know of further books and movies, where cryptography has a major role then we would be very glad if you could send us the exact title and a short explanation about the movie/book's content. Thanks a lot.

A.5.2 For Kids and Teenagers

The following list contains movies and “kid books”. The kid books contain both stories, and collections of simple ciphers, prepared in a didactic and exciting manner (please send us similar English kid books and kid movies, because at the moment our list contains mostly German kid books):

[Mosesxxxx] [no named author],

Top secret – The Book for Detectives and Spies (original title: Streng geheim – Das Buch für Detektive und Agenten), Edition moses, [no year named].

This is a thin book for small kids with Inspector Fox and Dr. Chicken.

[Arthur196x] Robert Arthur,

The Three Investigators: The Secret Key (German version: Die 3 ????: Der geheime Schlüssel nach Alfred Hitchcock (volume 119)), Kosmos-Verlag (from 1960)

The three detectives Justus, Peter and Bob have to decrypt covered and encrypted messages within this story to find out what is behind the toys of the Copperfield company.

[Para1988] Para,

Ciphers (original title: Geheimschriften), Ravensburger Taschenbuch Verlag, 1988 (1st edition 1977).

On 125 pages filled with a small font this mini format book explains many methods which young children can apply directly to encrypt or hide their messages. A little glossary and a short overview about the usage of encryption methods in history complete this little book.

Right at page 6 it summarizes for beginners in an old fashion style “The Important Things First” about paper&pencil encryption (compare chapter 2):

- “It must be possible to encrypt your messages at any place and at any location with the easiest measures and a small effort in a short time.
- Your cipher must be easy to remember and easy to read for your partners. But strangers should not be able to decrypt them.
Remember: Fastness before finesse, security before carelessness.
- Your message must always be as short and precise as a telegram. Shortness outranks grammar and spelling. Get rid of all needless like salutations or punctuation marks. Preferably use only small or only capital letters.”

[Müller-Michaelis2002] Matthias Müller-Michaelis,

The manual for detectives. Everything you need to know about ciphers, codes, reading tracks and the biggest detectives of the world (original title: Das Handbuch für Detektive. Alles über Geheimsprachen, Codes, Spurenlesen und die großen Detektive dieser Welt), Südwest, 2002.

[Kippenhahn2002] Rudolf Kippenhahn,

Top secret! – How to encrypt messages and to hack codes (original title: Streng geheim! – Wie man Botschaften verschlüsselt und Zahlencodes knackt), rororo, 2002.

In this novel a grandpa, an expert for secret writings teaches his four grandchildren and their friends, how to encrypt messages which nobody should read. Because there is someone who hacks their secrets, the grandpa has to teach them more and more complicated methods.

Within this story, which forms the framework, the most important classic encryption methods and its analysis are explained in a manner exciting and appropriate for children.

[**Harder2003**] Corinna Harder and Jens Schumacher,

Top secret. The big book for detectives (original title: *Streng geheim. Das große Buch der Detektive*), Moses, 2003.

[**Talke-Baisch2003**] Helga Talke and Milena Baisch,

Your mission in the weird villa. Riddle thriller (original title: *Dein Auftrag in der unheimlichen Villa. Kennwort Rätselkrimi*), Loewe, 2003.

From 4th form, <http://www.antolin.de>

Young detectives are faced simple ciphers and codes during their missions.

[**Flessner2004**] Bernd Flessner,

The Three Investigators: Manual for Secret Messages (original title: *Die 3 ????: Handbuch Geheimbotschaften*), Kosmos, 2004.

On 127 pages you learn in an easy and exciting manner, structured by the method types, which secret languages (like the one of the Navajo Indians or dialects) and which secret writings (real encryption or hiding via technical or linguistic steganography) existed and how simple methods can be decrypted.

The author tells where in history the methods were used and in which novel authors used encryption methods [like in Edgar Allan Poe's "The Gold Bug", like with Jules Verne's hero Mathias Sandorf or like with Astrid Lindgren's master detective Blomquist who used the ROR language (similar inserting ciphers are the spoon or the B language)].

This is a didactically excellent introduction for younger teens.

[**Zübert2005**] Directed by Christian Zübert,

The Treasure of the White Hawks (original title: *Der Schatz der weißen Falken*) , 2005.

This exciting adventure movie for kids ties in with the tradition of classics like "Tom Sawyer and Huckleberry Finn" or Enid Blytons "Five Friendse". The plot happens in summer 1981. In an old half tumbledown villa three young kids find the treasure map of the "White Hawks", which they decrypt with the help of a computer. Traced by another gang they aim to go to an old castle.

[**Dittert2011**] Christoph Dittert,

The Three Investigators: Secret Messages (German version: *Die 3 ????: Geheimnisvolle Botschaften*) (volume 160), Kosmos, 2011

In the house of Professor Mathewson an old hand-made book was stolen. The three detectives Justus, Peter and Bob are getting attacked by a ruthless opponent, who seems to be always a step ahead. A major part in this story is played by a palimpsest, an ancient manuscript page, written upon newly. Using X-rays they can make visible again the old text below. Its not only the story which is exciting, but also the way, how the instruction

for a treasure hunt is encrypted. Despite using the simple railfence cipher it's not easy to solve it, as the message is distributed onto two slips and as the printed symbols don't mean single letters.

Remark 1: You can find title pages of many of these kid books on the web site of Tobias Schrödel, who collects classic books about cryptography:

http://tobiasschroedel.com/crypto_books.php

Remark 2: If you know of further books, which address cryptography in a didactic and for children adequate way, then we would be very glad if you could send us the exact book title and a short explanation about the book's content. Thanks a lot.

A.6 Learning Tool for Elementary Number Theory

CT1 contains an interactive educational tool for elementary *number theory*, called “NT”.⁸

The educational tool “NT” (number theory) by Martin Ramberger introduces number theory and visualizes many of the methods and concepts. Where necessary it shows the according mathematical formulas. Often you can apply the mathematical methods dynamically with your own small numerical examples.

The content of this educational tool is mainly based on the books by J. Buchmann and H. Scheid [Buchmann2004, Scheid2003].

This visualized educational tool was built with Authorware 4.⁹

Request for enhancement/upgrade: It would be desirable to update it to a new version Authorware or to use another development platform. If there are developers interested to do this, I'd be more than happy (please send an email at the author of this CrypTool script).

Figures: The figures A.8 till A.15 give you an impression of the educational tool “NT”:

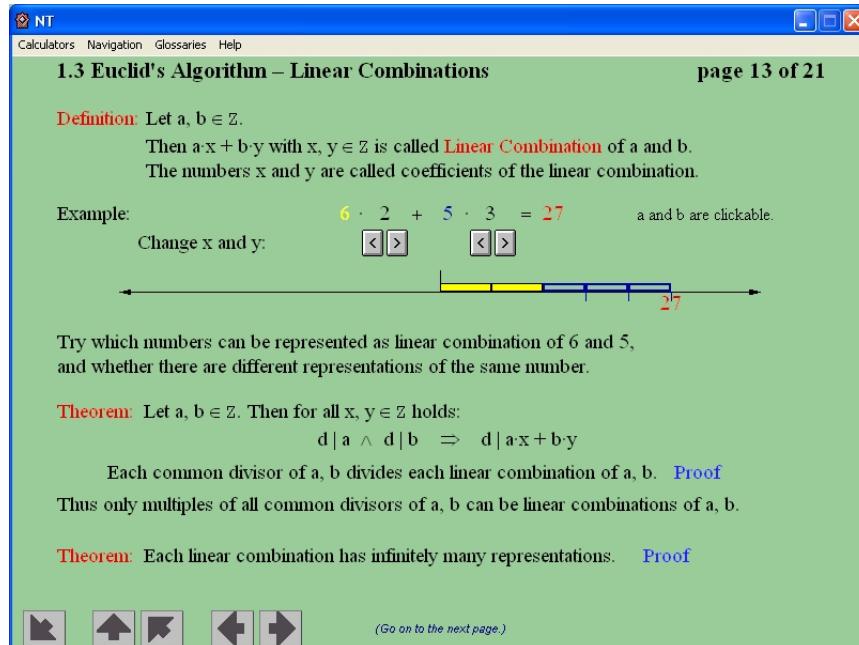


Figure A.8: Each common divisor of two integers also divides all its linear combinations

⁸NT can be called in CT1 via the menu path **Indiv. Procedures \ Number Theory Interactive \ Learning Tool for Number Theory**.

⁹As Authorware is outdated and as the vendor didn't make tools available to easily port it to his successor products, the program ZT will not be further developed.

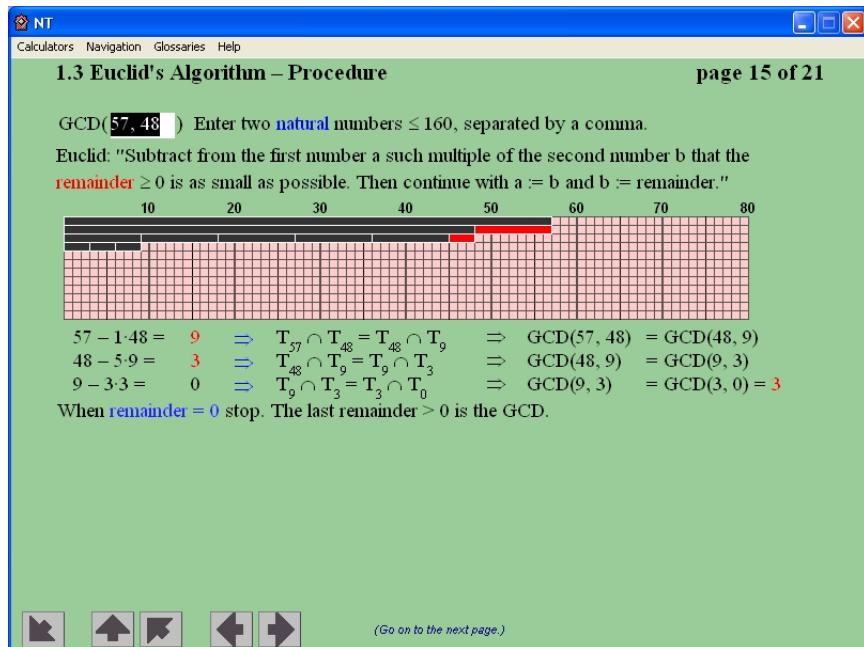


Figure A.9: Euklid's algorithm to determine gcd

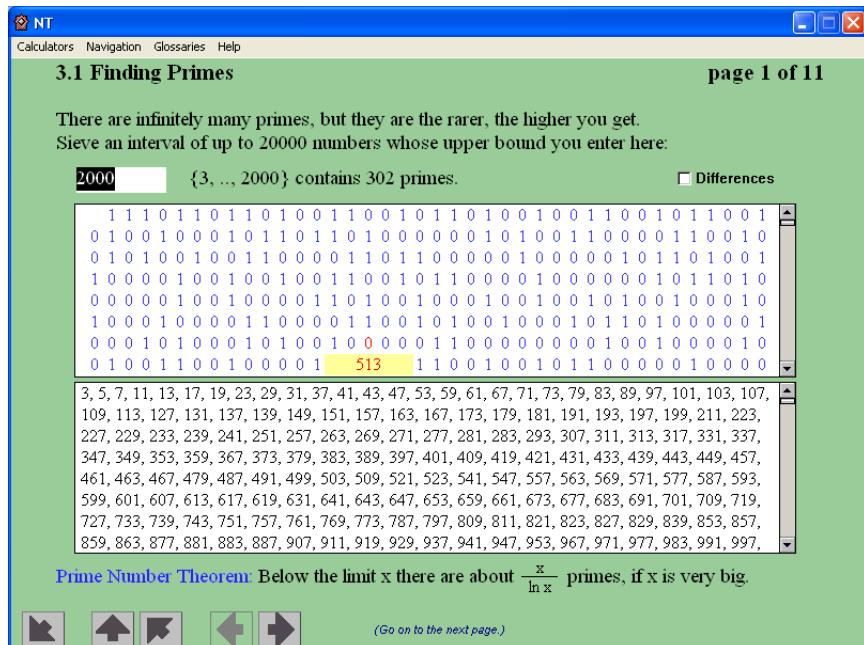


Figure A.10: Distribution of primes and its differences

NT
Calculators Navigation Glossaries Help
3.2 Fermat Test page 4 of 11

Each prime p passes a test that results from Fermat's Little Theorem:
Try for a $b \in \{2, \dots, p-1\}$, if $b^{p-1} \equiv 1 \pmod{p}$.
This test is called **Fermat Test**. Unfortunately some composite numbers pass it as well.

Example: $341 = 11 \cdot 31$, even so is $2^{340} \equiv 1 \pmod{341}$.

A passed test gives no information, one repeats it with a different base b :

$n =$	341	$2^{n-1} \equiv 1 \pmod{n}$	Test passed
GCD(b, n) =	1	<input type="button" value="b"/> < >	

Definition: Let n be a composite number, b coprime to n .
If $b^{n-1} \equiv 1 \pmod{n}$, then one calls

- **Pseudo Prime to Base b**,
- **Liar for** (the primality of) n ,

otherwise one calls b **Witness against** (the primality of) n .

Theorem: If there are any witnesses against n ,
then they make up at least 50% of all $b \in \{1, \dots, n\}$ coprime to n . **Proof**

(Go on to the next page.)

Figure A.11: Finding primes with the prime number test of Fermat

NT
Calculators Navigation Glossaries Help
4.1 Cryptography – Encryption ought to be Reversible page 4 of 17

This method applies addition and multiplication. Its keys are pairs (a, b) .
Try out which pairs make the encryption reversible.

abcdefghijklmnopqrstuvwxyz Nicht alle a funktionieren.

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	
1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	0
N	I	C	H	T	A	L	L	E	A	F	U	N	K	T	I	O	N	I	E	R	E	N				
14	9	3	8	20	0	1	12	12	5	0	1	0	6	21	14	11	20	9	15	14	9	5	18	5	14	

Encryption by the function: $E(x) = ax + b \pmod{27}$ $a = 2$
number of keys: $\varphi(27) \cdot 27$ $b = 0$

B	D	F	H	J	L	N	P	R	T	V	X	Z	A	C	E	G	I	K	M	O	Q	S	U	W	Y	
2	4	6	8	10	12	14	16	18	20	22	24	26	1	3	5	7	9	11	13	15	17	19	21	23	25	0
A	R	F	P	M	B	X	X	J	B	L	0	A	V	M	R	C	A	R	J	I	J	A				
1	18	6	16	13	0	2	24	24	10	0	2	0	12	15	1	22	13	18	3	1	18	10	9	10	1	

Decryption by the inverse function: $D(x) = a' \cdot (x - b') \pmod{27}$ $a' = 14$
It must hold: $a \cdot a' \equiv 1 \pmod{27}$ $b' = 0$

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	
1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	0
N	I	C	H	T	A	L	L	E	A	F	U	N	K	T	I	O	N	I	E	R	E	N				
14	9	3	8	20	0	1	12	12	5	0	1	0	6	21	14	11	20	9	15	14	9	5	18	5	14	

Figure A.12: Reversibility of encryption mechanisms exemplified with additive ciphers

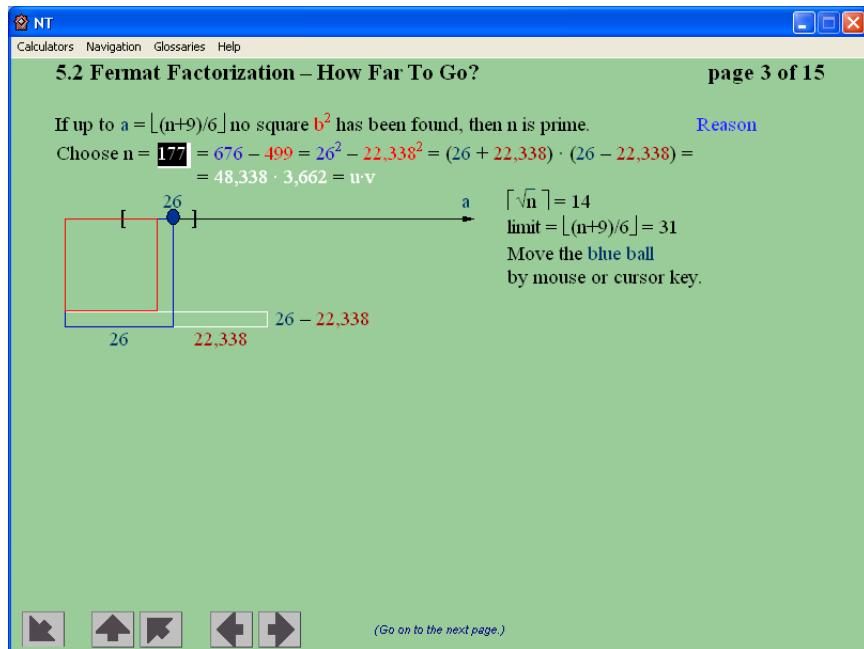


Figure A.13: Fermat factorization using the third binomial formula

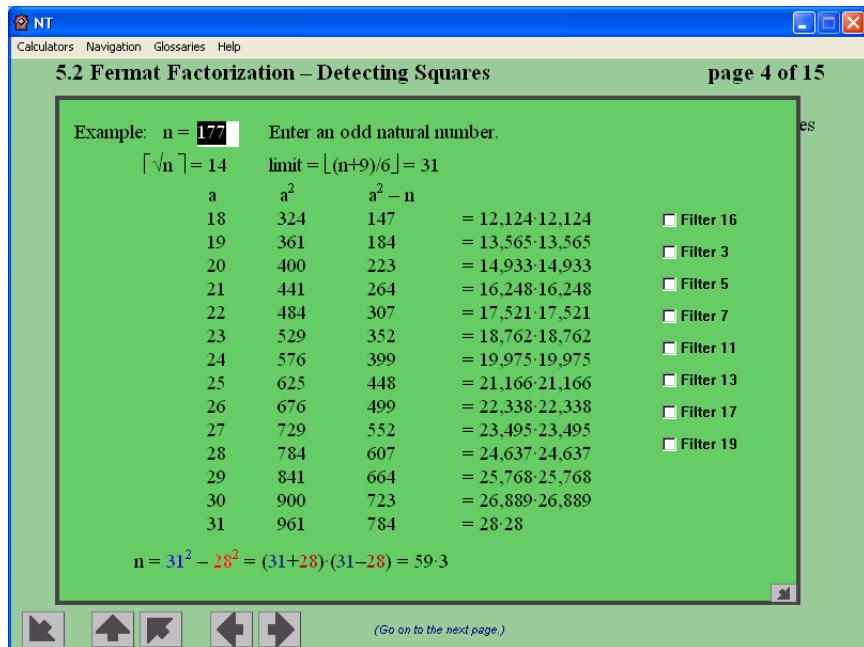


Figure A.14: Fermat factorization: Detecting squares

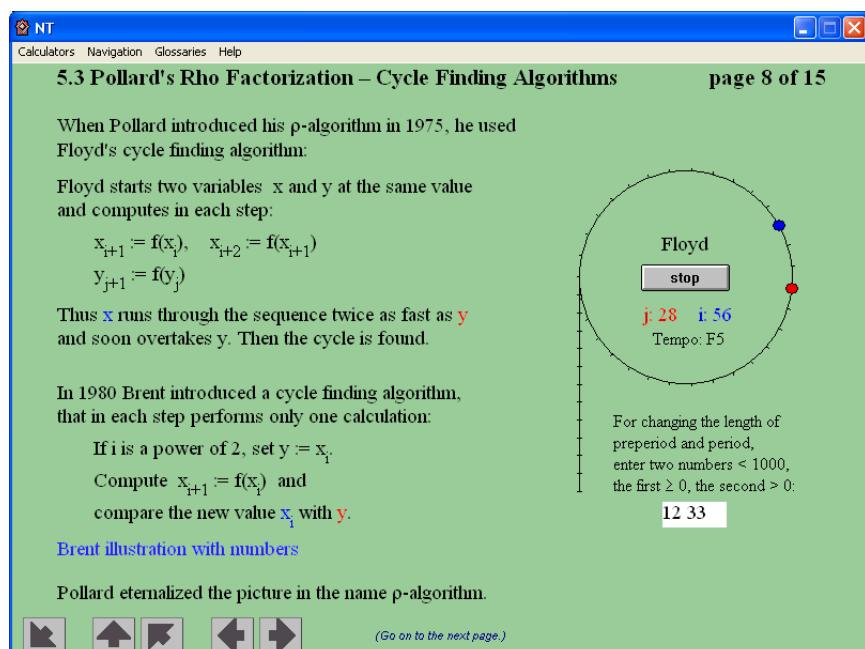


Figure A.15: Pollard's rho factorization: Floyd's cycle finding algorithm

Bibliography

[Buchmann2004] Johannes Buchmann,
Introduction to Cryptography, Springer, 2nd edition, 2004.

[Scheid2003] Harald Scheid,
Zahlentheorie, Spektrum Akademischer Verlag, 3rd edition, 2003.

A.7 Using Sage with this Script

This script includes numerous code samples using Sage. Sage is an open source computer algebra system (CAS) that supports teaching, study and research in mathematics. It combines many high-quality open source packages¹⁰ and provides access to their functionalities via a common interface, namely, a Python¹¹ based programming language.

Sage can be used as a powerful desktop calculator, as a tool to help (undergraduate) students study mathematics, or as a programming environment for prototyping algorithms and research in algorithmic aspects of mathematics.

You can get a quick impression of Sage e.g. with the references in this footnote¹².

The official Sage online documentation¹³ is available at: <http://www.sagemath.org>.

In the meantime there are lots of PDF and HTML documents about using Sage, so we name only a few of them as a good starting point¹⁴.

With respect to studying cryptography, Sage modules can be used to complement a first course in cryptography¹⁵.

Comprehensive introductions into cryptography are in this footnote¹⁶.

Sage user interfaces

Sage is available free of charge and can be downloaded from the following website:

<http://www.sagemath.org>

¹⁰To get an impression of how big Sage is: After downloading the source of Sage 4.1, it took around 5 hours on an average Linux PC to compile the whole system including all libraries. The compiled version occupied 1.8 GB disk space.

¹¹There is also an easy interface to the C language, called Cython, which can be used to substantially speed up functions in Sage.

See http://openwetware.org/wiki/Open_writing_projects/Sage_and_cython_a_brief_introduction.

¹²- “Invitation to Sage” by David Joyner, last update 2009

<http://sage.math.washington.edu/home/wdj/teaching/calc1-sage/an-invitation-to-sage.pdf>

- “The SDSU Sage Tutorial”,

<http://www-rohan.sdsu.edu/~mosulliv/sagetutorial/>

<http://www-rohan.sdsu.edu/~mosulliv/sagetutorial/sagecalc.html>

- “SAGE For Newbies” by Ted Kosan, 2007,

http://sage.math.washington.edu/home/tkasan/newbies_book/sage_for_newbies_v1.23.pdf

¹³The corresponding official PDF documents can be downloaded at

<http://www.sagemath.org/help.html>, <http://www.sagemath.org/doc> and <http://planet.sagemath.org>.

¹⁴- “Library”: <http://www.sagemath.org/library/index.html>,

- “Documentation Project”: <http://wiki.sagemath.org/DocumentationProject>,

- “Teaching”: http://wiki.sagemath.org/Teaching_with_SAGE.

¹⁵- Module sources in the directory SAGE_ROOT/devel/sage-main/sage/crypto.

- Overview, what crypto currently is in Sage:

<http://www.sagemath.org/doc/reference/sage/crypto/>

- Discussions about teaching related aspects of development crypto in Sage:

http://groups.google.com/group/sage-devel/browse_thread/thread/c5572c4d8d42d081

¹⁶- David Kohel’s notes from 2008 are a ready course

<http://www.sagemath.org/library/crypto.pdf> or the same eventually newer at

<http://sage.math.washington.edu/home/wdj/teaching/kohel-crypto.pdf> .

- “Introduction to Cryptography with Open-Source Software”, a very good book from Alasdair McAndrew, CRC, 2011

The default interface to Sage is **command line** based, as shown in figure A.16. However, there is a graphical user interface to the software as well in the form of the Sage notebook (see figure A.17). We can even use Sage **notebooks**¹⁷ online at different servers, without having to install Sage locally, e.g.:

```
http://www.sagenb.org or  
http://sage.mathematik.uni-siegen.de:8000
```

Sage runs under many Linux distributions, Mac OS X, and Windows. For the Windows platform, a complete distribution of Sage currently only runs as a VMware image.

```
Terminal
-----
| Sage Version 4.0.2, Release Date: 2009-06-18
| Type notebook() for the GUI, and license() for information.
-----
sage: alph = AlphabeticStrings(); subcipher = SubstitutionCryptosystem(alph)
sage: key = alph([25 - i for i in xrange(26)]); key
ZYXWVUTSROPONMLKJIHGFEDECBA
sage: encrypt = subcipher(key)
sage: msg = alph.encoding("Substitute this with something else nice."); msg
SUBSTITUTE THIS WITH SOMETHING ELSE NICE
sage: encrypt(msg)
HFYHGRGFGVGSRHDRGSHLNVGSRMTVOHYMRXY
sage:
sage: trancipher = TranspositionCryptosystem(alph, 7)
sage: key = trancipher.random_key(); key
(1,3)(2,7,4,6,5)
sage: trancipher.enciphering(key, msg); msg
BTSIUTSESUITHTTWOISHHEEGTNIEELCSIN
SUBSTITUTE THIS WITH SOMETHING ELSE NICE
sage: 
```

Figure A.16: Sage command line interface

¹⁷Further details about Sage notebooks can be found at chapter 7.9.2 (“Implementing elliptic curves for educational purposes” ⇒ “Sage”).

¹⁸To start the Sage gui locally: Enter notebook() at the Sage prompt, and then your favorite browser (Iceweasel, Firefox, IE, ...) is started e.g. with the URL <http://localhost:8000>.

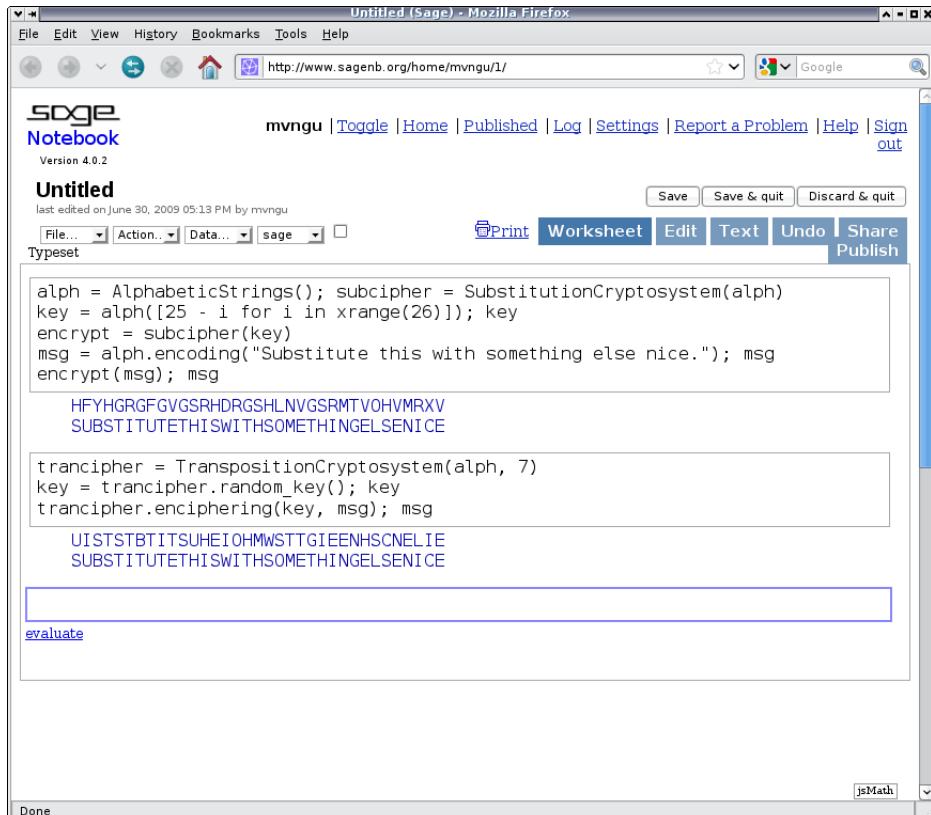


Figure A.17: Sage notebook interface¹⁸

Getting help with using Sage

Upon loading Sage from the command line, we are presented with something similar to the following:

```
mnemonic:~$ sage
-----
| Sage Version 4.1, Release Date: 2009-07-09 |
| Type notebook() for the GUI, and license() for information. |
-----
```

```
sage: help
Type help() for interactive help, or help(object) for help about object.
sage:
sage:
sage: help()
```

Welcome to Python 2.6! This is the online help utility.

If this is your first time using Python, you should definitely check out the tutorial on the Internet at <http://docs.python.org/tutorial/>.

Enter the name of any module, keyword, or topic to get help on writing Python programs and using Python modules. To quit this help utility and return to the interpreter, just type "quit".

To get a list of available modules, keywords, or topics, type "modules", "keywords", or "topics". Each module also comes with a one-line summary

of what it does; to list the modules whose summaries contain a given word such as "spam", type "modules spam".

Plenty of help is provided in the form of the official Sage documentation that is distributed with every release of Sage (see Figure A.18). The official Sage standard documentation includes the

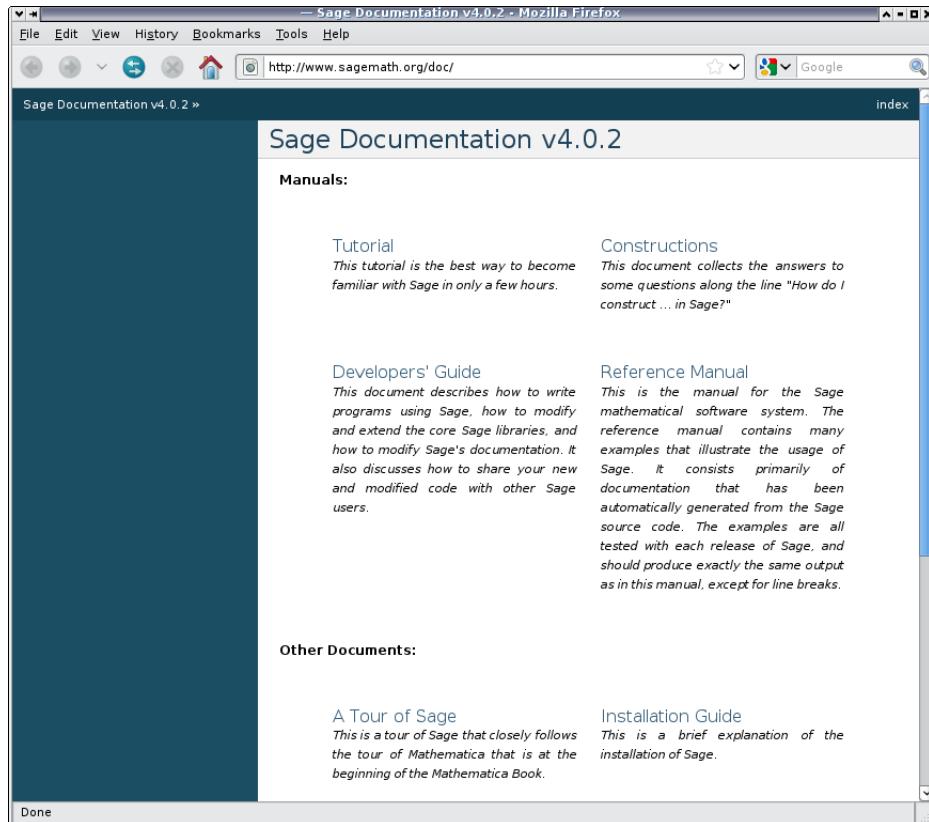


Figure A.18: The Sage standard documentation

following documents:

- Tutorial — This tutorial is designed to help Sage beginners become familiar with Sage. It covers many features that beginners should be familiar with, and takes one to three hours to go through.
- Constructions — This document is in the style of a Sage “cookbook”. It is a collection of answers to questions about constructing various objects in Sage.
- Developers’ Guide — This guide is for developers who want to contribute to the development of Sage. Among other issues, it covers coding style and conventions, modifying the core Sage libraries, modifying the Sage standard documentation, and code review and distribution.
- Reference Manual — This manual provides complete documentation on the major features of Sage. The description of a class is accompanied by numerous code samples. All code samples in the reference manual are tested before each Sage release.
- Installation Guide — This guide explains how to install Sage under various platforms.

- A Tour of Sage — This is a tour of Sage that showcases various features of Sage that are useful for beginners.
- Numerical Sage — This document introduces tools available under Sage that are useful for numerical computation.
- Three Lectures about Explicit Methods in Number Theory Using Sage — This document is about using Sage to perform computations in advanced number theory.

From within a Sage session, we can obtain a list of commands matching some pattern. To do so, we type the first few characters and then press the “Tab” key:

```
sage: Su[TAB]
Subsets           Subwords           SuzukiGroup
SubstitutionCryptosystem  SupersingularModule
```

If we know the exact name of a command, we can use the `help` function to obtain further information on that command, or append the question mark “?” to the command name. For example, the command `help(SubstitutionCryptosystem)` provides documentation on the built-in class `SubstitutionCryptosystem`. We can get documentation on this class with the question mark as follows:

```
sage: SubstitutionCryptosystem?
Type:type
Base Class:<type 'type'>
String Form:<class 'sage.crypto.classical.SubstitutionCryptosystem'>
Namespace:Interactive
File:/home/mvngu/usr/bin/sage-3.4.1/local/lib/python2.5/site-packages/sage/crypto/classical.py
Docstring:

    Create a substitution cryptosystem.

    INPUT:

    - ``S`` - a string monoid over some alphabet

    OUTPUT:

    - A substitution cryptosystem over the alphabet ``S``.

    EXAMPLES::

        sage: M = AlphabeticStrings()
        sage: E = SubstitutionCryptosystem(M)
        sage: E
        Substitution cryptosystem on Free alphabetic string monoid
        on A-Z
        sage: K = M([ 25-i for i in range(26) ])
        sage: K
        ZYXWVUTSRQPONMLKJIHGfedcba
        sage: e = E(K)
        sage: m = M('THECATINTHEHAT')
        sage: e(m)
        GSVXZGRMGSVSZG

    TESTS::

        sage: M = AlphabeticStrings()
```

```
sage: E = SubstitutionCryptosystem(M)
sage: E == loads(dumps(E))
True
```

For further assistance on specific problems, we can also search the archive of the **sage-support** mailing list at

<http://groups.google.com/group/sage-support>

Some examples using the built-in mathematical functions in Sage

Here are a few little examples¹⁹ (all in console mode, for ease) to see what you can do with Sage:

Sage sample A.1 Some small general samples in Sage from different areas in mathematics

```
# * Calculus:  
sage: x=var('x')  
sage: p=diff(exp(x^2),x,10)*exp(-x^2)  
sage: p.simplify_exp()  
1024 x^10 + 23040 x^8 + 161280 x^6 + 403200 x^4 + 302400 x^2 + 30240  
  
# * Linear Algebra:  
sage: M=matrix([[1,2,3],[4,5,6],[7,8,10]])  
sage: c=random_matrix(ZZ,3,1);c  
[ 7 ]  
[-2 ]  
[-2 ]  
sage: b=M*c  
sage: M^-1*b  
[ 7 ]  
[-2 ]  
[-2 ]  
  
# * Number theory:  
sage: p=next_prime(randint(2^49,2^50));p  
1022095718672689  
sage: r=primitive_root(p);r  
7  
sage: pl=log(mod(10^15,p),r);pl  
1004868498084144  
sage: mod(r,p)^pl  
10000000000000000  
  
# * Finite Fields (\url{http://en.wikipedia.org/wiki/Finite_field}):  
sage: F.<x>=GF(2)[]  
sage: G.<a>=GF(2^4,name='a',modulus=x^4+x+1)  
sage: a^2/(a^2+1)  
a^3 + a  
sage: a^100  
a^2 + a + 1  
sage: log(a^2,a^3+1)  
13  
sage: (a^3+1)^13  
a^2
```

¹⁹The examples are from the blog of Dr. Alasdair McAndrew, Victoria University,
<http://amca01.wordpress.com/2008/12/19/sage-an-open-source-mathematics-software-system>

Writing code samples with Sage

When you start using a CAS (computer algebra system) you normally type in the single commands on the command line as in the above example²⁰.

But if you develop your own functions, modify them and call them, then it is much easier to do the development in your own editor, save it to a script file and execute the functions non-interactively on the command line manually. Both ways to develop code were applied in chapter 1.7 (“Appendix: Examples using Sage”), chapter 2.5 (“Appendix: Examples using Sage”), chapter 3.14 (“Appendix: Examples using Sage”) and in chapter 4.19 (“Appendix: Examples using Sage”).

To program and test Sage code using an editor there are two useful commands: `load()` and `attach()`²¹.

Suppose you have a function definition like this:

```
def function(var1):
    """
    DocText.
    """
    ...
    return (L)
```

which has been saved to the file `primroots.sage`.

To load this function into Sage (and do a syntax check at once), use `load()` as follows:

```
sage: load primroots.sage
```

and you can then proceed to use on the command line any variable or function defined in that Sage script²².

Normally we also want to edit our own Sage script and reload the content of the changed script into Sage again. In that case, you can use the command `attach()` (you also can apply `attach()` directly after loading the script, even before having changed the script; and you can even omit `load()`, as this is contained in `attach()`):

²⁰The standard way for presenting Sage code starts the lines with “sage:” and “...”.

```
sage: m = 11
sage: for a in xrange(1, m):
....:     print [power_mod(a, i, m) for i in xrange(1, m)]
....:
```

This script usually uses the above convention for presenting Sage code, if the code doesn’t come from a Sage script. When people copy and paste the Sage code from this script, in order to enter it at the command line, they should leave out “sage:” and “...” from the script (nevertheless in most cases the command prompt can deal with these prefixes correctly).

²¹See Sage tutorial about Programming, chapter “Loading and Attaching Sage files”,
<http://www.sagemath.org/doc/tutorial/programming.html#loading-and-attaching-sage-files>.

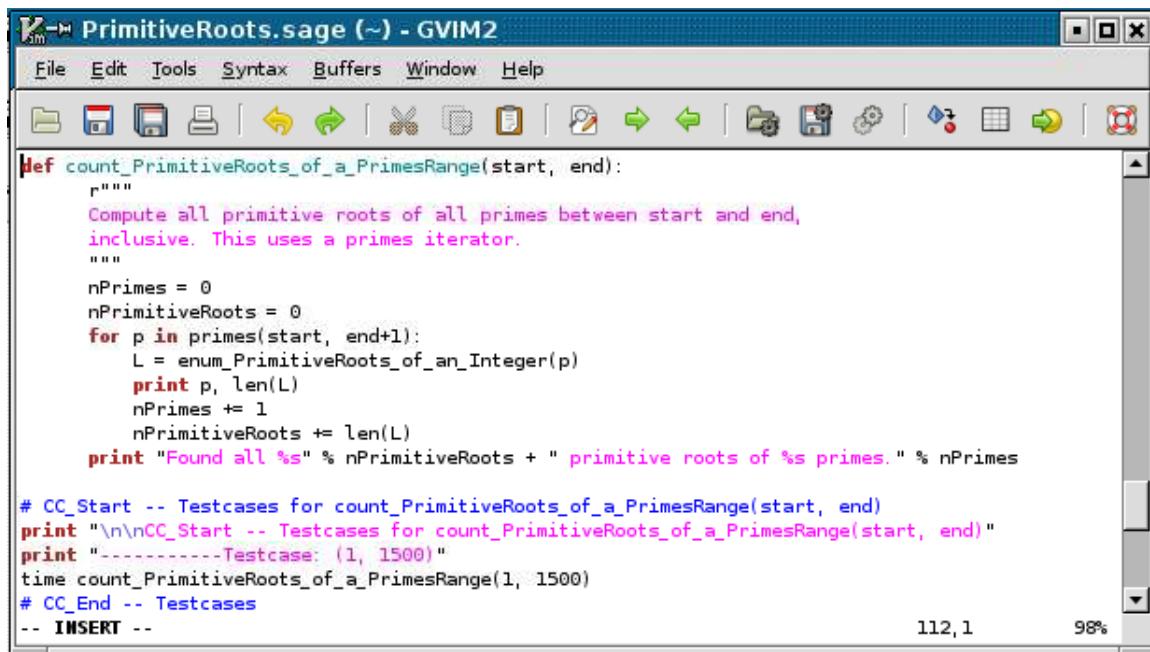
²²Notes:

- Don’t use white spaces in your file name.
- Its recommended that your Sage script has the file extension “.sage”, instead of “.py”. With a Sage script whose file name ends in “.sage”, when you load it into Sage then the default Sage environment is also loaded to make sure that it works as if you have defined your function from the Sage command line. This also applies if you run the script from a bash shell using `$ sage primroots.sage`.
- If you run your script as above, then Sage first parses your script, writes it to another file called “primroots.py” (note the “.py” extension), adds all necessary variables to “primroots.py” as well as writing any import statements to that file. That way, your Sage script is executed as if you had typed the definitions in your script to the Sage command line. An important difference is that all output needs a `print` statement.

```
sage: attach primroots.sage
```

Now edit the Sage script in a text editor, but don't exit Sage. After you saved it within your text editor, the changed function definition is reloaded into the running Sage session after the next typing of Enter (and a syntax check is done at once). This reloading is done automatically for you, provided that all changes to your script have been saved. You can think of the command `attach()` as a way of telling Sage to watch for all changes to a file, and reloading the file again once Sage notices that there have been changes. With this command, you don't have to copy and paste between your text editor and the Sage command line interface.

Here is a picture of Sage code in the editor GVIM with activated code highlighting (see figure A.19).



The screenshot shows a GVIM window titled "PrimitiveRoots.sage (~) - GVIM2". The code in the buffer is:

```
Def count_PrimitiveRoots_of_a_PrimesRange(start, end):
    """
    Compute all primitive roots of all primes between start and end,
    inclusive. This uses a primes iterator.
    """
    nPrimes = 0
    nPrimitiveRoots = 0
    for p in primes(start, end+1):
        L = enum_PrimitiveRoots_of_an_Integer(p)
        print p, len(L)
        nPrimes += 1
        nPrimitiveRoots += len(L)
    print "Found all %s" % nPrimitiveRoots + " primitive roots of %s primes." % nPrimes

# CC_Start -- Testcases for count_PrimitiveRoots_of_a_PrimesRange(start, end)
print "\n\nCC_Start -- Testcases for count_PrimitiveRoots_of_a_PrimesRange(start, end)"
print "-----Testcase: (1, 1500)"
time count_PrimitiveRoots_of_a_PrimesRange(1, 1500)
# CC_End -- Testcases
-- INSERT --
```

The status bar at the bottom right shows "112, 1" and "98%".

Figure A.19: Sage sample shown in an editor with code highlighting

If you prefer to see the output of an attached file as if you would have typed in the commands on the commandline directly (not only what is shown via `print`) then you could use the command `iload()`: Each line is loaded one at a time. To load the next line, you have to press the `Enter` key. You have to repeatedly press the `Enter` key until all lines of the Sage script are loaded into the Sage session.

```
sage: iload primroots.sage
```

Some more hints:

- To get the version of your Sage environment: `version()`
- To move quickly to the Sage code examples in this script,
 - either look in the index at `Sage -> Code examples`,
 - or have a look at the appendix “List of Sage Code Examples”.
- The Sage samples in this script are delivered with CrypTool.
For further details see the end of the overview “List of Sage Code Examples”.

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- The source code of the Sage samples in this script is delivered as Sage program files within the CrypTool-1 setup program. All samples within one chapter are collected in one file. After installing CrypTool 1 you find the Sage samples within the subdirectory `sage` within the CrypTool-1 directory:
 - SAGE-Samples-in-Chap01.sage
 - SAGE-Samples-in-Chap02.sage
 - SAGE-Samples-in-Chap03.sage
 - SAGE-Samples-in-Chap04.sage
- All samples have been tested with Sage Version 5.3 (Release Date 2012-09-08).

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