CO322: DS & A Linear Data Structures

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What is to come

- Look at some simple data structures
- Most of it you know and have used before :)
- Mostly a recap of what was discussed

Linked list

Limitation of arrays:

- Need to know the size at the start! (at least some good estimate)
- Can not go beyond that!

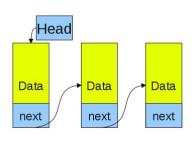
Advantages of arrays:

- Contiguous in memory (good cache locality if accessed sequentially)
- Very limited bookkeeping is required (only need to know the start of the array and its size)

Alternative \Rightarrow Linked list Used when we do not know the size in advance and size varies a lot! (need to think before selecting)



Linked list ...



- Head points to the first element (ideal for sequential)
- Add/remove items as required
- Each element contains a pointer to the next
- No limit on how many elements we can have
- Additional storage for pointers

In your program you need to store, on an average N integers. In the worst case this can be 2N. Should you use an array or linked list? Explain. (assume 32bit addresses).

Linked list in Java

```
import java.util.*;
class Linked {
   public static void main(String [] args) {
LinkedList<String> list = new LinkedList<String>();
for(int i=0; i<10; i++)</pre>
   list.add(Integer.toString(i));
list.remove();
System.out.println("List content: "+ list);
```

```
$ java Linked
List content: [1, 2, 3, 4, 5, 6, 7, 8, 9]
```

Linked list in Java interface

from:

http://docs.oracle.com/javase/7/docs/api/java/util/LinkedList.html

- ▶ add(E e)
- add(int index, E e)
- addFirst(E e)
- contains(E e)
- get(int index)
- **...**

Note: Somewhat advanced than the simple structure we discussed. Stemming from programmer needs! You might need different representations to support implementing the above interface (example: two pointers to find head and tail.)

Linked list variations

Some issues of *singly-linked list* and variations to resolve issues:

- Difficult to remove from the middle (you know the element, want ro remove)
 - Doubly-linked list: Each node has two pointers: next and previous.
- You need to worry about corner cases
 - Circular-linked list: Last node points to the first node
- Combination of both variations: Circular-doubly-linked list

Queue data structure

Basic idea:

- First-in First-out (FIFO) data structure
- Linked list with tail insert and head removal

should we have a doubly-linked list for this? should we have a circular-linked list for this?

Or can be implemented using an array

Queue operations:

- ▶ add(E e)
- remove(E e)
- peek()

Used when every you need FIFO processing.

Think of some Examples?



Example application: background code

```
class Student {
 public String eNumber, fieldOfChoice;
 public double gpa;
 Student(String eNumber, String fieldOfChoice, double
     gpa) {
   this.eNumber = eNumber;
   this.fieldOfChoice = fieldOfChoice;
   this.gpa
                    = gpa;
 public void display() {
   System.out.println(this.eNumber + ": GPA = " +
       this.gpa + " wants to do " + this.fieldOfChoice);
 public boolean preferComp() {
   return new
       String("Computer").equals(this.fieldOfChoice);
```

Example application: background code

```
import java.util.*;
class QueueEx {
 static String [] fields = {"Chemical", "Civil",
     "Computer",
   "Electrical", "Mechanical", "Production"};
 static int comp = 5;
 static String getFieldOfChoice() {
   return fields[(int)(Math.random() * 10 % 6)]:
 static double getGPA() {
  return Math.random() * 4 ;
 /* more to come */
```

Example application: main

```
public static void main(String [] args) {
 Queue < Student > queue = new LinkedList < Student > ();
 Student s:
 for(int i=0; i<400; i++) {</pre>
   s = new Student("E/11/0" + Integer.toString(i),
   getFieldOfChoice(),
   getGPA());
   queue.add(s);
 while(!queue.isEmpty()) { /* note isEmpty */
   s=queue.remove();
   if((s.preferComp()) && (comp-- > 0)) {
     s.display();
   } /* while */ } /* main ends */
} /* end of class */
```

Broken code! We do not need FIFO; need priority queue

Priority Queues

Oder or removal depends on priority

- ▶ Internally the list is kept sorted. When inserting (or adding) a new element add to it's correct place. (recall insertion sort)(insert O(N), removal O(1))
- ▶ Else we can insert to the tail and locate the highest element when removing? (insert O(1), removal O(N))
- Which method is better? Why?

Priority Queues in Java

```
class StudentComp implements Comparator<Student> {
 public int compare(Student a, Student b) {
   return Double.compare(b.gpa, a.gpa); /* since we want
       descending */
/* part of the main method. Whole main is in the next
   slide
*/
 PriorityQueue<Student> queue = new
   PriorityQueue<Student>(2, new StudentComp());
```

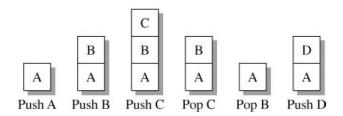
Priority Queues Example

```
public static void main(String [] args) {
 PriorityQueue<Student> queue = new
     PriorityQueue<Student>(2, new StudentComp());
 Student s;
 for(int i=0; i<400; i++) {</pre>
   s = new Student("E/11/0" + Integer.toString(i),
       getFieldOfChoice(), getGPA());
   queue.add(s);
 while(!queue.isEmpty()) {
   s=queue.remove();
   if((s.preferComp()) && (comp-- > 0)) {
     s.display();
}/* main */
```

The stack data structure

Basic idea:

- "list of items that are accessible from only one end"
- Last-In First-Out (LIFO) data structure
- ► Two main operations: push and pop
- Other operations: isEmpty, peek, isFull ...
- Can be implemented using arrays or linked lists



Stack application

```
int factorial(int x) {
  if(x == 1) return 1;
  return x * factorial(x-1);
}
```

- A stack is used to postpone obligations, which must later be fulfilled in reverse order.
- Recall implementing tower of Hanoi using non-recursive implementation.

Stack: Java implementation

```
public interface Stack{
  void push(Object item);
  Object pop();
  boolean isEmpty();
  /* the interface can be implemented in many ways */
}
```

Stack: Implementation of interface

```
public class StackArray implements Stack {
 private Object [] array;
 private int size = 0,capacity = 5;
 public StackArray() {
   array = new Object [capacity];
 public void push(Object item) {
   if(size < capacity) array[size++] = item;</pre>
   else throw new IllegalStateException("Full array");
 public Object pop() {
   if(size != 0) return array[--size];
   return null;
 public boolean isEmpty() { return (size == 0); }
```

Stack: Implementation of interface

Stack: Implementation of interface cnt..

```
private Node top; /* stack top */
public StackList() { top = null; }
public void push(Object item) {
 top = new Node(item, top);
} /* end push */
public Object pop() {
 if(!isEmpty()) {
   Object t = top.item;
   top = top.next;
   return t;
 return null;
} /* end pop */
public boolean isEmpty() { return top == null; }
```

Use of a stack: Example 1: Converting to binary

Convert the first argument into binary:

```
class Binary {
 public static void main(String [] args) {
   int rem, a = Integer.parseInt(args[0]);
   Stack s = new StackList():
   while(a >= 1) {
     rem = a % 2:
     s.push(rem);
     a = a / 2;
   while(!s.isEmpty())
   System.out.print(s.pop());
```

Use of a stack: Example 2: Check brackets

Algorithm:

- Scan the string from left to right
- When you find an opening bracket push it
- When you find closing bracket pop and see if it matches
- At the end of the string, stack should be empty.

```
public static boolean checkBalance(String str) {
  Stack s = new StackList();
  char [] array = str.toCharArray();
  /* implement the rest */
```

Use of a stack: Example 2: Check brackets code

```
public static boolean checkBalance(String str) {
 Stack s = new StackList();
 char [] array = str.toCharArray();
 for(int i=0; i < array.length; i++) {</pre>
   if(array[i] == '[' || array[i] == '(')
       s.push(array[i]);
     if(array[i] == ']') if(s.pop() != '[') return false;
     if(array[i] == ')') if(s.pop() != '(') return false;
   } /* for( */
   return s.isEmpty();
```

What is the run-time complexity of this code? (N char and M bracket types)

Different types of expressions

Expression type	Example
Postfix	a b +
Infix	a+b
Prefix	+ a b

- ▶ In a postfix evaluation format for an arithmetic expression, an operator comes after its operands.
- In a infix evaluation format for an arithmetic expression, an operator comes between its operands.
- ▶ In a prefix evaluation format for an arithmetic expression, an operator comes before its operands.

Infix to postfix

Algorithm:

- Read the expression from left to right
- If you find a operator push to stack until the operand has less precedence than the one on top of the stack.
- When the operator has less precedence pop till stack is empty or operator has higher precedence
- Operands are put as is.
- If the stack is not empty, pop them.

Expression	Output	Stack
A + B * C - D A + B * C - D		
A + B * C - D	A	
A + B * C - D	АВ	+
A + B * C - D	АВС	+ *
A + B * C - D	A B C * +	-
A + B * C - D	A B C * + D -	

Evaluation of a postfix expression

Algorithm:

- Read the expression from left to right
- push the operands to stack
- when you find an operator, pop two operands from the stack, do the operating and push the result.
- ▶ At the end of the expression, stack should have only one value which is the evaluation of the expression.

Original expression: A + B * C - D

Expression	Stack	Current Operation
A B C * + D - A B C * + D - A B C * + D - A B C * + D - A B C * + D - A B C * + D - A B C * + D - A B C * + D -		
A B C * $+$ D $-$	АВС	
A B C * $+$ D $-$	A	B * C
A B C * $+$ D $-$	A (B * C)	
A B C $*$ + D -	A + (B * C)	
A B C * + D -	A + (B * C) D	
A B C * + D -		A + (B * C) - D

Exercises

- Implement a Java function int evaluate(String postfixExpression) which will evaluate the given postfix expression
- 2. Implement a Java function *String evaluate(String postfixExpression)* which will convert the infix expression into postfix.