ALGORITHMS DATA STRUCTURE

Exercise 1: Inventory Management System

Steps:

1. Understand the Problem:

Why Data Structures and Algorithms are Essential in Handling Large Inventories

Data structures and algorithms are crucial for efficiently managing large inventories because they help:

- Optimize Storage: Efficient data structures use memory effectively, allowing you to store large amounts of data without wasting space.
- Speed Up Retrieval: Appropriate data structures enable fast retrieval of information, which is essential for operations like searching for a product.
- Improve Performance: Algorithms determine the speed of operations like adding, updating, and deleting products, which can significantly impact the system's performance.

Types of Data Structures Suitable for This Problem

- ArrayList: Good for storing a list of products where the order matters and quick access by index is needed. However, adding/removing elements in the middle can be slow.
- HashMap: Excellent for fast access, addition, and deletion based on a key (e.g., productId). Provides average-case constant time complexity for these operations.

2. Setup:

```
package com.example.inventory;

public class Product {
    private String productId;
    private String productName;
    private int quantity;
    private double price;

    public Product(String productId, String productName, int
    quantity, double price) {
        this.productId = productId;
        this.productName = productName;
        this.quantity = quantity;
        this.price = price;
    }

    // Getters and setters
    public String getProductId() { return productId; }
```

```
public void setProductId(String productId) { this.productId =
productId; }
  public String getProductName() { return productName; }
  public void setProductName(String productName) { this.productName
= productName; }
  public int getQuantity() { return quantity; }
  public void setQuantity(int quantity) { this.quantity = quantity;}
}
  public double getPrice() { return price; }
  public void setPrice(double price) { this.price = price; }
}
```

```
import com.example.inventory.Inventory;
import com.example.inventory.Product;

public class Main {
    public static void main(String[] args) {
        Inventory inventory = new Inventory();

        // Add products
        Product product1 = new Product("1", "Laptop", 10, 999.99);
        Product product2 = new Product("2", "Smartphone", 50, 499.99);
        inventory.addProduct(product1);
        inventory.addProduct(product2);

        // Update product
```

```
product1.setPrice(949.99);
    inventory.updateProduct(product1);

// Get product
    Product retrievedProduct = inventory.getProduct("1");
    System.out.println("Retrieved Product: " +
retrievedProduct.getProductName() + ", Price: " +
retrievedProduct.getPrice());

// Delete product
    inventory.deleteProduct("2");

// Attempt to get deleted product
    Product deletedProduct = inventory.getProduct("2");
    if (deletedProduct == null) {
        System.out.println("Product with ID 2 has been deleted");
    }
}
```

Time Complexity of Each Operation

1. Add Product:

Operation: addProduct

Time Complexity: O(1) (Average case)

2. Update Product:

Time Complexity: O(1) (Average case)

- 3. Delete Product: **Time Complexity:** O(1) (Average case)
- Get Product: Time Complexity: O(1) (Average case).
 Optimization Suggestions
- **Load Factor Management:** Ensure the HashMap load factor is well-managed to avoid excessive collisions, which can degrade performance.
- **Concurrency:** For a multi-threaded environment, consider using ConcurrentHashMap to handle concurrent access to the inventory.
- **Batch Operations:** For large-scale operations, consider batch processing to reduce the overhead of multiple single operations.

Exercise 2: E-commerce Platform Search Function

1: Understand Asymptotic Notation

Big O Notation

- Definition: Big O notation is a mathematical representation used to describe the upper bound of an algorithm's running time or space requirements in terms of the input size. It provides a high-level understanding of the algorithm's efficiency.
- Purpose: It helps in analyzing the performance of algorithms, especially for large inputs, by focusing on the most significant factors that affect scalability.

Best, Average, and Worst-Case Scenarios

- Best Case: The scenario where the algorithm performs the minimum number of operations. For search operations, it's typically when the target element is at the beginning of the array.
- Average Case: The scenario that represents the expected performance of the algorithm across all possible inputs. For search operations, it's the average number of operations required to find an element.
- Worst Case: The scenario where the algorithm performs the maximum number of operations. For search operations, it's typically when the target element is at the end of the array or not present.
- 2. Setup:
- 3. Implementation:

```
package com.example.ecommerce;

public class Product {
    private String productId;
    private String productName;
    private String category;

public Product(String productId, String productName, String category) {
        this.productId = productId;
        this.productName = productName;
        this.category = category;
    }

    // Getters
    public String getProductId() { return productId; }
    public String getProductName() { return productName; }
    public String getCategory() { return category; }
}
```

```
package com.example.ecommerce;

public class LinearSearch {
    public Product linearSearch(Product[] products, String productId) {
        for (Product product : products) {
            if (product.getProductId().equals(productId)) {
                return product;
            }
        }
        return null;
    }
}
```

```
package com.example.ecommerce;
```

```
import java.util.Arrays;
import java.util.Comparator;

public class BinarySearch {
    public Product binarySearch(Product[] products, String productId) {
        Arrays.sort(products, Comparator.comparing(Product::getProductId));
        int left = 0;
        int right = products.length - 1;

        while (left <= right) {
            int mid = left + (right - left) / 2;
            int cmp = products[mid].getProductId().compareTo(productId);

        if (cmp == 0) {
            return products[mid];
        } else if (cmp < 0) {
            left = mid + 1;
        } else {
            right = mid - 1;
        }
        return null;
    }
}</pre>
```

```
import com.example.ecommerce.*;

public class Main {
    public static void main(String[] args) {
        Product[] products = {
            new Product("1", "Laptop", "Electronics"),
            new Product("3", "Smartphone", "Electronics"),
            new Product("4", "Monitor", "Electronics"),
            new Product("5", "Keyboard", "Accessories")
      };

      LinearSearch linearSearch = new LinearSearch();
      BinarySearch binarySearch = new BinarySearch();

      // Test Linear Search
      Product foundProduct = linearSearch.linearSearch(products, "3");
      if (foundProduct!= null) {
            System.out.println("Linear Search: Found " +
            foundProduct.getProductName());
      } else {
            System.out.println("Linear Search: Product not found");
      }

      // Test Binary Search
      foundProduct = binarySearch.binarySearch(products, "3");
      if (foundProduct!= null) {
            System.out.println("Binary Search: Found " +
      foundProduct.getProductName());
      } else {
            System.out.println("Binary Search: Found " +
            foundProduct.getProductName());
      } else {
            System.out.println("Binary Search: Product not found");
      }
}
```

4: Analysis

1. Compare Time Complexity:

- Linear Search:
 - Best Case: O(1) (Element found at the beginning)
 - Average Case: O(n)
 - Worst Case: O(n) (Element found at the end or not present)
- Binary Search:
 - Best Case: O(1) (Element found at the middle initially)
 - Average Case: O(log n)
 - Worst Case: O(log n) (Element not present or at the end of the divided intervals)

2. Suitability Analysis:

- Linear Search:
 - Suitable for small arrays or unsorted data.
 - Simpler to implement and requires no additional preprocessing.
- Binary Search:
 - More efficient for large arrays due to logarithmic time complexity.
 - Requires the array to be sorted, which adds an initial overhead but significantly improves search performance for frequent queries.

Exercise 3: Sorting Customer Orders

Steps:

1. Understand Sorting Algorithms:

Bubble Sort

- Description: Bubble Sort repeatedly steps through the list, compares adjacent elements, and swaps them if they are in the wrong order. The process repeats until the list is sorted.
- Time Complexity: O(n^2) for the worst and average case, O(n) for the best case (when the list is already sorted).

Insertion Sort

• Description: Insertion Sort builds the sorted array one item at a time by repeatedly picking the next item and inserting it into its correct position.

• Time Complexity: O(n^2) for the worst and average case, O(n) for the best case (when the list is already sorted).

Quick Sort

- Description: Quick Sort selects a 'pivot' element and partitions the array into elements less than the pivot and elements greater than the pivot. It recursively sorts the partitions.
- Time Complexity: O(n log n) on average, O(n^2) in the worst case (when the pivot selection is poor, such as always selecting the smallest or largest element).

Merge Sort

- Description: Merge Sort divides the list into halves, recursively sorts each half, and then merges the sorted halves back together.
- Time Complexity: O(n log n) for all cases.

```
package com.example.orders;

public class BubbleSort {
    public void bubbleSort(Order[] orders) {
        int n = orders.length;
        for (int i = 0; i < n - 1; i++) {
              for (int j = 0; j < n - i - 1; j++) {
                  if (orders[j].getTotalPrice() > orders[j + 1].getTotalPrice()) {
```

```
import com.example.orders.*;
import java.util.Arrays;

public class Main {
    public static void main(String[] args) {
        Order[] orders = {
            new Order("1", "Arya", 250.75),
            new Order("2", "Mahi", 100.50),
            new Order("3", "Harsh", 320.40),
            new Order("4", "Rahul", 150.30),
            new Order("5", "Anu", 210.20)
        };

        BubbleSort bubbleSort = new BubbleSort();
        QuickSort quickSort = new QuickSort();

        // Test Bubble Sort
        Order[] bubbleSortedOrders = Arrays.copyOf(orders, orders.length);
```

Compare Time Complexity:

• Bubble Sort:

Best Case: O(n) (already sorted)

Average Case: O(n^2)Worst Case: O(n^2)

Quick Sort:

Best Case: O(n log n)Average Case: O(n log n)

Worst Case: O(n^2) (poor pivot selection)

Why Quick Sort is Generally Preferred:

- **Efficiency:** Quick Sort is generally faster than Bubble Sort for large datasets due to its O(n log n) average time complexity.
- **Practical Performance:** Despite its O(n^2) worst-case complexity, Quick Sort's performance can be improved with good pivot selection strategies (like choosing the median or using randomized pivots).
- **Memory Usage:** Quick Sort typically has better memory usage compared to other O(n log n) algorithms like Merge Sort, as it is an in-place sorting algorithm.

Exercise 4: Employee Management System

Steps:

1. Understand Array Representation:

Array Representation in Memory

- Contiguous Memory Allocation: Arrays are stored in contiguous memory locations. This means that elements are stored sequentially in adjacent memory addresses.
- Fixed Size: The size of an array is fixed at the time of creation. This means that the number of elements it can hold is determined when the array is instantiated.
- Indexing: Arrays allow direct access to any element using its index, making retrieval operations very efficient (O(1) time complexity).

Advantages of Arrays

- Fast Access: Constant time access to elements using indices.
- Memory Efficiency: Minimal overhead compared to some other data structures.
- Predictability: Memory allocation is straightforward and predictable.

2. Setup:

```
package com.example.employees;
import java.util.Arrays;

public class EmployeeManager {
    private Employee[] employees;
    private int size;

    public EmployeeManager(int capacity) {
        employees = new Employee[capacity];
        size = 0;
    }

    // Add an employee
    public void addEmployee(Employee employee) {
        if (size >= employees.length) {
    }
}
```

```
public void traverseEmployees() {
public void deleteEmployee(int employeeId) {
       if (employees[i].getEmployeeId() == employeeId) {
```

```
import com.example.employees.*;

public class Main {
   public static void main(String[] args) {
        EmployeeManager manager = new EmployeeManager(10);

        Employee emp1 = new Employee(1, "Alice", "Manager", 75000);
        Employee emp2 = new Employee(2, "Bob", "Developer", 50000);
        Employee emp3 = new Employee(3, "Charlie", "Analyst", 60000);

        manager.addEmployee(emp1);
```

```
manager.addEmployee(emp2);
manager.addEmployee(emp3);

System.out.println("All Employees:");
manager.traverseEmployees();

System.out.println("\nSearching for Employee with ID 2:");
Employee searchedEmployee = manager.searchEmployee(2);
System.out.println(searchedEmployee);

System.out.println("\nDeleting Employee with ID 2:");
manager.deleteEmployee(2);
manager.traverseEmployees();
}
```

Analyze Time Complexity:

- Add Operation:
 - o Time Complexity: O(1) (inserting at the end of the array, assuming there is space).
- Search Operation:
 - o Time Complexity: O(n) (linear search through the array).
- Traverse Operation:
 - o Time Complexity: O(n) (printing all elements).
- Delete Operation:
 - o Time Complexity: O(n) (linear search to find the element, then shifting elements).

Limitations of Arrays:

- Fixed Size: Arrays have a fixed size, which can be a limitation if the number of elements changes frequently.
- Inefficient Deletions/Insertions: Deleting or inserting elements (other than at the end) requires shifting elements, which can be inefficient (O(n) time complexity).
- Memory Allocation: Allocating large arrays may lead to memory wastage if not fully utilized

Exercise 5: Task Management System.

Steps:

1. Understand Linked Lists:

Types of Linked Lists

- 1. Singly Linked List:
 - o Each node contains data and a reference to the next node.

- The last node points to null, indicating the end of the list.
- o Operations: Add, delete, search, and traverse.

2. Doubly Linked List:

- Each node contains data, a reference to the next node, and a reference to the previous node.
- Allows traversal in both directions.
- o More complex and requires more memory compared to singly linked lists.

2. Setup:.

```
package com.example.tasks;

public class TaskManager {
   private Node head;

   private class Node {
       Task task;
       Node next;

      Node(Task task) {
       this.task = task;
       this.next = null;
      }
   }
}
```

```
import com.example.tasks.*;

public class Main {
    public static void main(String[] args) {
        TaskManager manager = new TaskManager();

        Task task1 = new Task(1, "Design Database", "Pending");
        Task task2 = new Task(2, "Develop API", "In Progress");
        Task task3 = new Task(3, "Test Application", "Pending");

        manager.addTask(task1);
        manager.addTask(task2);
        manager.addTask(task3);

        System.out.println("All Tasks:");
        manager.traverseTasks();

        System.out.println("\nSearching for Task with ID 2:");
        Task searchedTask = manager.searchTask(2);
        System.out.println(searchedTask);

        System.out.println("\nDeleting Task with ID 2:");
        manager.deleteTask(2);
        manager.traverseTasks();
    }
}
```

Analyze Time Complexity:

- Add Operation:
 - o Time Complexity: O(n) (inserting at the end of the list requires traversing the list).
- Search Operation:
 - o Time Complexity: O(n) (linear search through the list).
- Traverse Operation:
 - o Time Complexity: O(n) (printing all elements).
- Delete Operation:
 - o Time Complexity: O(n) (linear search to find the element, then adjusting pointers).

Advantages of Linked Lists Over Arrays:

- Dynamic Size: Linked lists can grow and shrink dynamically without needing to allocate or deallocate memory for the entire list.
- Efficient Insertions/Deletions: Insertions and deletions at the beginning or middle of the list are more efficient (O(1) for beginning insertions/deletions) compared to arrays, which require shifting elements (O(n) time complexity).

Exercise 6: Library Management System

Scenario:

You are developing a library management system where users can search for books by title or author.

Steps:

- 1. Understand Search Algorithms:
- 2. Setup:.
- 3. Implementation:

```
package com.example.library;
import java.util.List;

public class Library {
    private List<Book> books;

    public Library(List<Book> books) {
        this.books = books;
    }

    public Book linearSearchByTitle(String title) {
        for (Book book : books) {
            if (book.getTitle().equalsIgnoreCase(title)) {
                return book;
            }
        }
}
```

```
return null;
}

public Book binarySearchByTitle(String title) {
   int left = 0;
   int right = books.size() - 1;

   while (left <= right) {
      int mid = left + (right - left) / 2;
      Book midBook = books.get(mid);

   int cmp = midBook.getTitle().compareToIgnoreCase(title);

   if (cmp == 0) {
      return midBook;
   } else if (cmp < 0) {
      left = mid + 1;
   } else {
      right = mid - 1;
   }
   return null;
}</pre>
```

Compare Time Complexity:

- Linear Search:
 - Time Complexity: O(n)
 - Suitable for: Unsorted or small datasets.
- Binary Search:
 - o Time Complexity: O(log n)
 - Suitable for: Large, sorted datasets.

When to Use Each Algorithm:

- Linear Search: Use when the dataset is unsorted or the list size is small, as it does not require sorting and is straightforward to implement.
- Binary Search: Use when the dataset is large and sorted, as it significantly reduces the number of comparisons needed to find the desired element.

Exercise 7: Financial Forecasting

Steps:

1. Understand Recursive Algorithms:

Recursion is a process where a function calls itself directly or indirectly to solve a problem. It can simplify certain problems by breaking them down into smaller sub-problems of the same type.

- Base Case: The condition under which the recursion stops.
- Recursive Case: The part of the function that calls itself with a modified argument.

2. Setup:

```
package com.example.financialforecasting;

public class FinancialForecasting {
    // Recursive method to predict future value
    public static double predictFutureValue(double initialValue,
    double growthRate, int years) {
        // Base case: if years is 0, return the initial value
        if (years == 0) {
            return initialValue;
        }

        // Recursive case: calculate future value for the next year
        double previousYearValue = predictFutureValue(initialValue,
        growthRate, years - 1);
        return previousYearValue * (1 + growthRate);
    }
}
```

```
public static void main(String[] args) {
         double initialValue = 1000.0;
         double growthRate = 0.05; // 5% annual growth rate
         int years = 10;

         double futureValue = predictFutureValue(initialValue,
growthRate, years);
         System.out.println("Future value after " + years + " years: "
+ futureValue);
    }
}
```

```
public class FinancialForecastingOptimized {
       if (years == 0) {
       if (memo.containsKey(years)) {
           return memo.get(years);
       memo.put(years, currentValue);
   public static void main(String[] args) {
       double initialValue = 1000.0;
       Map<Integer, Double> memo = new HashMap<>();
```

Time Complexity:

• The time complexity of the recursive algorithm is O(n), where n is the number of years. This is because the method is called recursively n times.

Optimizing the Recursive Solution:

Memoization: Store the results of previous calculations to avoid redundant computations.
 This technique can transform a simple recursive solution into a more efficient one by reducing the time complexity.

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