

→ Transport layer

↳ TCP and UDP

Transmission
control protocol

↳ User Datagram Protocol

Q: How TCP set up a connection?

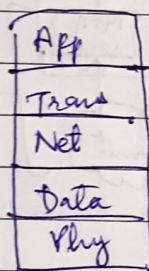
Q: How TCP segments look like?

Q: How two computers reliably transfer data?

↳ Checksum, cyclic redundancy, message authentication codes or protocols

→ TCP/IP protocol suite - application layer, transport layer, network layer, data-link layer, physical layer

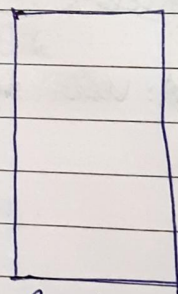
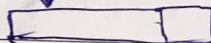
App



Sender

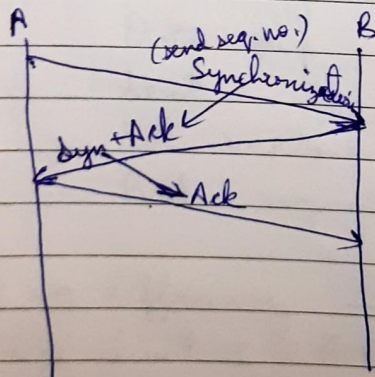
↓ Bytestream

Data | Port TCP segment



Receiver

↳ TCP is reliable (end to end reliability)



3-way handshake
supported by TCP

↳ Initial sequence number (ISN) generated by OS varies from 0 → 4294967295

→ Connection tear down

→ Properties of Transport layer

→ supports stream of byte → reliable byte delivery service

→ reliable delivery - acknowledgement

↳ indicates correct delivery

- checksum detect corrupted data

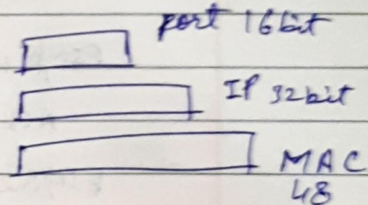
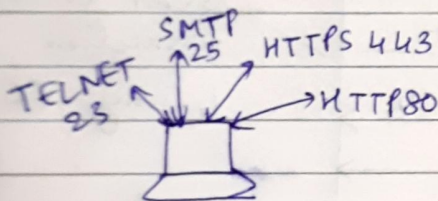
- sequence number to detect missing data

- flow control prevent overrunning receiver

Insequence → Data delivered to application in sequence transmitted

checksum → covers header & data

→ Port



Port 16 bit

$$2^{16} - 1 = 65535$$

We know ports 0 - 1023 by IANA

↳ Internet Assigned Number

~~from~~

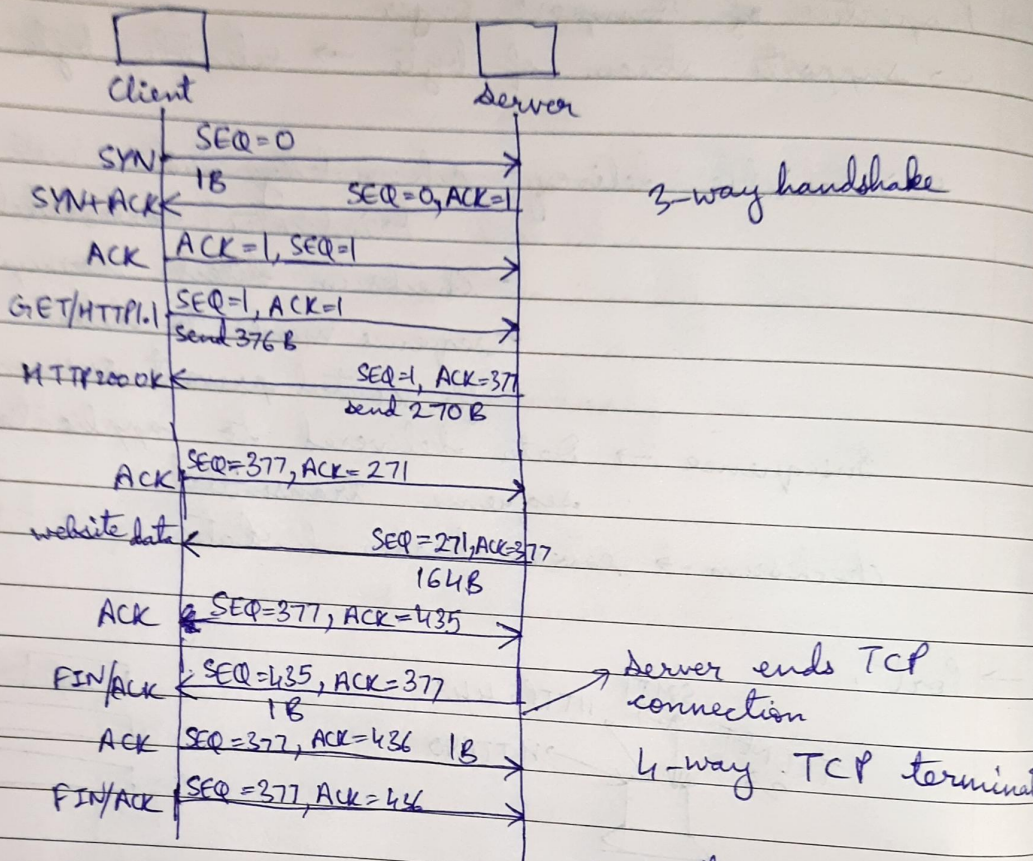
Registered ports 1024 - 49151

Private port 49152 - 65535

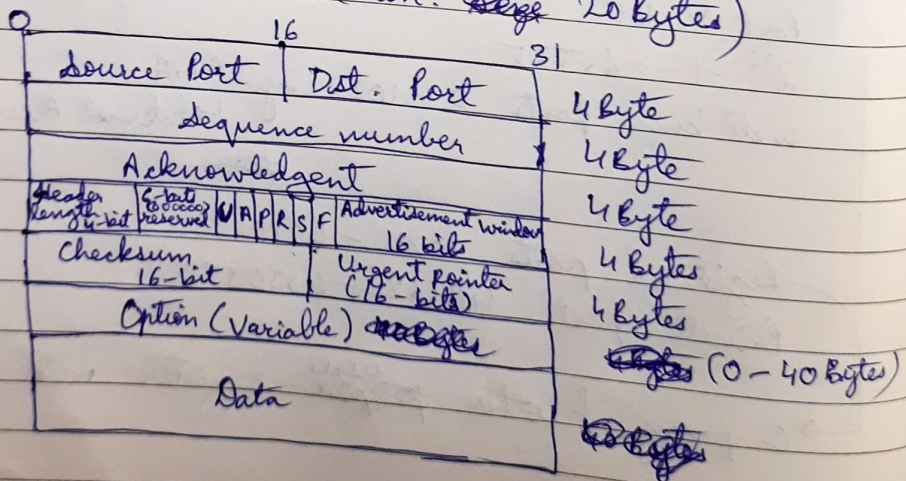
→ Company & other ~~people~~ ^{users} register with ICANN

Connection setup time + Data transfer time + Connection tear down time
 ↳ send "finish"

→ Sequence no., Cumulative Ack



→ TCP Header format (Min. ^{header length} 20 bytes)



9 ↳ U A P R S F flags uses

Urgent Ack PSH Reset Sync Fin

→ TCP checksum

→ TCP header + TCP body + pseudo IP header

Pseudo IP header

Source IP (32 bits)		
Dest. IP (32 bits)		
Fixed 8 bit	Protocol field 8 bits	TCP segment length

- source IP
- Destination IP
- TCP/UDP segment length
- Protocol (Type of protocol)
- Fixed 8-bit (0000 0000)

→ Header length

min - 20 Bytes = 4×5

max - 60 Bytes = 4×15

(multiple of 4)

So use 4-bits to store 0 to 15

If not multiple of 4, add padding bits

→ PSK received

→ ~~PSK~~ Window size used in flow control

↳ contains the size of the receiving window of the sender.

↳ it advertises how much data in bytes the sender can receive without acknowledgment

↳ changes dynamically → use option if more than 16 bits

→ Option → Timestamp

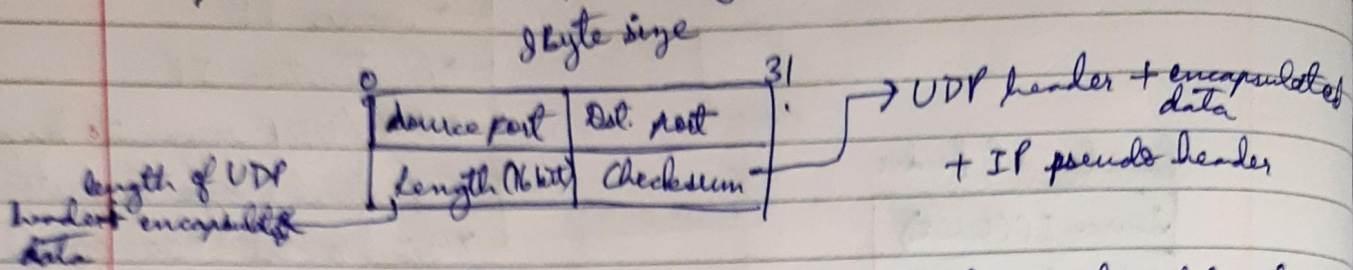
→ Window size extension

→ Parameter negotiation

→ Padding

→ Max. segment size

→ UDP - User Datagram Protocol
↳ Simplest transport layer Protocol



- ↳ Takes data from network layer & adds header.
- ↳ Broadcast, streaming, etc.
- ↳ Connection less protocol
- ↳ Stateless
- ↳ unreliable, fast
- ↳ min. transport service
- ↳ almost null protocol
- ↳ don't guarantee in-order delivery
- ↳ doesn't provide congestion control
- ↳ avoid TCP overhead
- ↳ Applications :-

→ DNS

→ RIP & OSPF

↓
Routing
information
protocol

→ TFTP

→ Streaming, video conferencing

→ SNMP

→ Boot & DHCP

→ Flow-control both in transport and data-link layer
↓
windowing concept

↳ Flow-control is the process of managing the rate of data transmission between 2 nodes (to prevent loss of data) to prevent a fast

sender from overwhelming a slow receiver

↳ Congestion control is different from flow-control

↳ if congestion occurs then control flow of data

↳ stop & wait flow control technique

↳ Windowing flow control technique

↳ stop & wait

↳ simple protocol

↳ sender transmits single frame at a time & then waits for acknowledgement

↳ receiver transmits ack.

Q: In stop & wait protocol, every 4th packet is lost & we need to send total 10 packets. How many transmissions it took to send all the packets?

Ans 13

1 2 3 4 4 5 6 7 7 8 9 10 10
↑ . ↑ . ↑ .

Q: In stop & wait protocol if error prob. is p & no. of packets to send is n , how many packets we have to send?

Ans

$$n + np + np^2 + \dots = \frac{n}{1-p}$$

↳ Demerits of stop & wait

↳ transmission & propagation delay, channel under utilized, no-pipelining

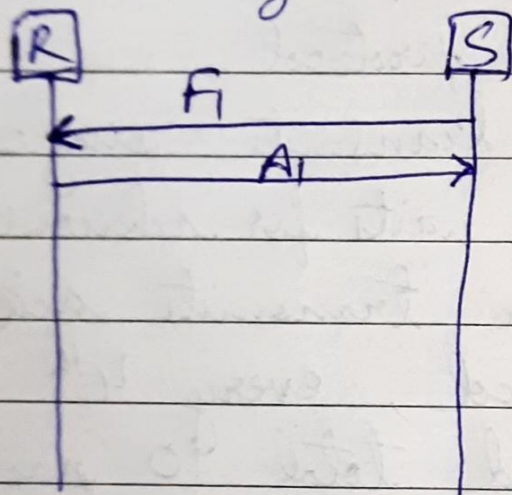
↳ if ack. lost even if data is sent, duplicate data is sent

↳ delayed ack

→ Sliding window

→ Method of flow control in which a receiver gives a transmitter permission to transmit data until a window is full.

→ Windowing :- Controls how much information is transferred from one end to the other.
eg:- Window size = 1



Window size = 3

