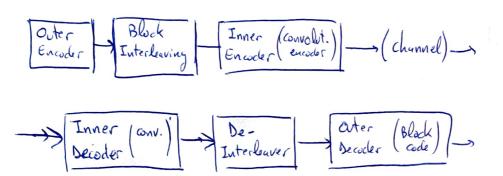
Lecture 8: Turbo Codes

Classic Concatenated Coding It was a series concatenation



Principal Characteristics

- * Parallel, instead of series, concatenated coding
- * Use of convolutional codes
- * Pseudo-random interleaving
- * Iterative (SISO) decoding soft in, soft out decoding

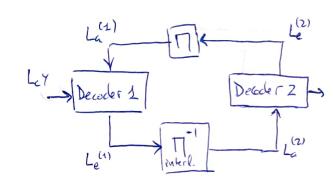
Turbo Principle

The of turbo decoder

*We have a first decoder and it gives a first Ley Decoder 2]

Output and the after the interleaving it goes

to a second decodeto a second decoder



- * Instead of deciding the final result, we send it to the first decoder, completing in this way one iteration.
- * Then, the first decoder receives a "cleaned" information (cleaned by 2nd decoder) We continue doing iterations.
- * Where nothing changes from one iteration to another the process stops and it produces the final result.

Turbo Encoder Architecture

UK Convol. Par Puncturing Interleaver Convol. Px

- @ with what we have described the rate well be very long so there is a puncturing where we discard (with precise procedure) some bits
- 15 The final result is send on the channel.

- : Uk is the original information and it is send to the channel in different ways
- 1 It is send directly
- @ The same UK is send to a first convolutional encoder (Systematic Conv. Encod.) which produces the parity check bets Pk that are send
- 3 The same Ux is send to an interleaver (a mix in time domain between bits) and to another convolutional encoder that produces the second parity check Pk that is also send.
 - # RSC -> Recursive Systematic Convolutional enoclers

SISO decoder

Soft in - soft out -> We are working with the real values taken by the sampler * We work with logaritus because is more convenient

by the prob-of the value assume -1

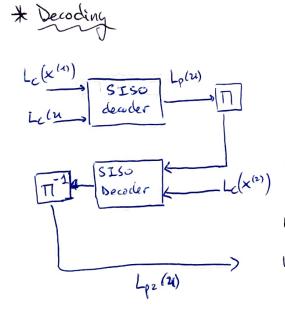
L(u) =
$$\log \left(\frac{P(u=+1)}{P(u=-1)} \right) \Rightarrow \int_{-\infty}^{\infty} \frac{P(u=+1)}{P(u=-1)} > 1$$

L(u) \(\dots \) \(\frac{P(u=+1)}{P(u=-1)} \) \(\dots \) \(\frac{P(u=+1)}{P(u=-1)} \) \(\dots \

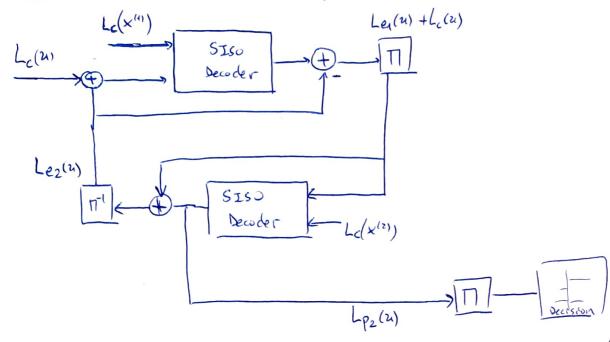
The sign of L(21) decide who is the winner.

La At the input we have
$$L_c(u) \rightarrow \log$$
 likehood satio associated to the original information $L_c(x^{(1)}) \rightarrow u$ " u u u to the parity check control

Lo Output: It produces the posteriori value related to the original information Using BCJR algorithm evaluate the max. a posteriori prob. Lp(u) - This info. is more affordable than Lc(u)



Lo In the first SISO deader we apply channel transmitted information Lacon and the channel information related to the first parity check control. La(x11) Apply the procedure and obtain a posterior info related to the original data Lp(n) Ly Apply the permutation and give the value to the second siso dec. which uses the second parity check control Ly Rearrange the bit in the proper order Is We have the output Lp2(2) so we can decide the final result * Architecture of iterative decoding



* Here we give to the first siso dec. the information given by the channel Lean and the info. of the channel related to the first parity check control Le(x(1)).

* We do the permutation and give to the second siso dec.

- * Instead of deciding inmediatly we come back to the feedback loop.
- => Important: To avoid positive feedback we have to remove the information that was already given. => Extrinsic information
- => Extrinsic information. The extrinsic into is found by substructing the corresponding input from the output
 - AIt is necessary to subtract the info. that is already available at the decoder to prevent positive feedback
 - The extriusic info is the amount of new info gained by the current decoder step

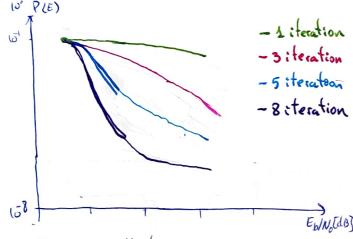
Performance

* According to Shannon the requirements to have a good performance

La Random codes and large blocks => Many bits

Lo Soft Decoding > Turbo uses convolutional codes (soft decoding can be implemented)
using Vitrebi adjointhm

* the behaviour respect the iterations

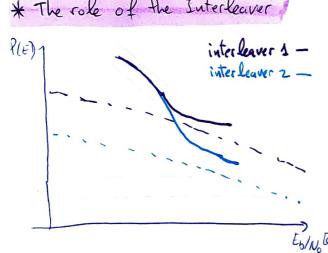


We can see that the behaviour of P(E) is Composed by two parts: Waterfall and Error Floor

- The first part is the waterfall and we can see that increasing the SNR we decrease the P(E) a lot
- The second part is the Error Floor and we can see that even if we increase SNR the improvement of the performance is not very important.

Finally, we see that increasing the iterations leads to a better performances, besides we don't need many iterations to reach to a very good performance.

* The role of the Interleaver



We can see here that if we change the. interleaver, the floor will be different During the the waterfall the interleaver is not important, it only affects strongly the performance of the floor

The interleaver is related to the permutation => but is not clear what is the best permutation to be applied

* Random cocling argument:

- e) Truly random codes approaches capacity => but not feasible
- o) Turbo codes appear random -> but it allows practical decoding

* Distance spectrum argument

- e) Traditionally the code design is focused on => Maximizing Minimum Distance
- o) With Turbo Edes we have a minimum distance that is not so big the goal is: => Reduce the multiplicity of low weight code words

Ly when I transmit a code word there are many other codewords that are close to the transmitted one = We try to reduce it

Ly & with smal duin a remarkable performance can be achieve at low SNR

EXIT Chart (Extrinsic Information Transfer)

* It is a way to analyze the performance and the behaviour of a Turbo Codes with respect the iterations

* In axis. Shows the mutual information at the imput of the SISO Decoder The mutual information between the original data and the information given in the input of SISO dec.

* I e axis. Shows the mutual information at the output of SISO Decocler. Mutual info between output and the original information.

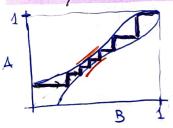
If we increase . In (mutual info at input), the Ie (mutual info at output) will be also increased

* Behaviour with SWR - SNR = 1.5 dB

- SNR = 0,6 dB - SNR = 0,3 dB Increasing the SNR I will obtain better performances

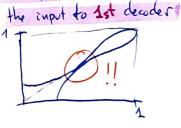
The mutual info. at output is increased when SNR is increased because the noise will be less important

* Analysis of Turbo Codes



A routput of so decoder becomes the input to decoder

B - Output of and decoder becomes



We are working with the same code for the two decoders. The ortput of the first idea becomes the input of the second dec. and then the output of the second code becomes the input of the first dec. This ⇒ The path shows us what happens with the iterations between one code repeats iteratively.

> We use the EXIT chart in this way to analize the convergence

=> When the part named the tunnel is reduced we need more iterations to pass through it

If the tunnel is closed the convergence is not graranteed.

The goal of the iterations is to reach to the nutral info equal to 1.

The EXIT chart is the best way to analyze the behaviour of the code with respect iterations and final good