# Prolog

-Recursive Rules

-List Data structure

## **Prolog Features**

- Prolog has the power of designing knowledge based systems at very ease.
- Prolog can be interfaced with almost all high level languages. E.g. with Python, C Java etc.; <a href="https://wiki.python.org/moin/IntegratingPythonWithOtherLanguages#Prolog">https://wiki.python.org/moin/IntegratingPythonWithOtherLanguages#Prolog</a>
- Prolog is very much suited for parsing data or parsing other languages (natural and computer ).
- Prolog is also very popular as a Game Description language with very powerful libraries.

For installation of Prolog editor in Eclipse for SWI prolog, see instructions at:

https://sewiki.iai.uni-bonn.de/research/pdt/docs/download https://sewiki.iai.uni-bonn.de/research/pdt/docs/start

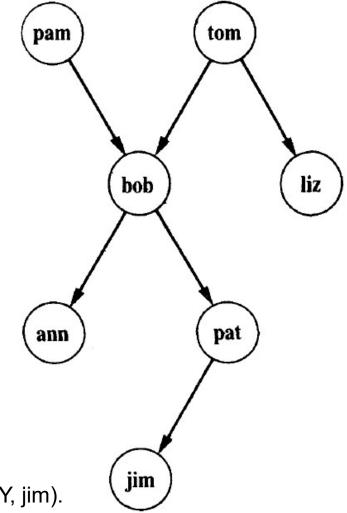
## SWI tutorials

- CLP(FD) Constraint Logic Programming over Finite Domains (<a href="http://www.pathwayslms.com/swipltuts/clpfd/clpfd.html">http://www.pathwayslms.com/swipltuts/clpfd/clpfd.html</a>)
- Using Definite Clause Grammar in Prolog (<a href="http://www.pathwayslms.com/swipltuts/dcg/">http://www.pathwayslms.com/swipltuts/dcg/</a>)
- https://www.swi-prolog.org/features.html

## Example 2: A Family Tree

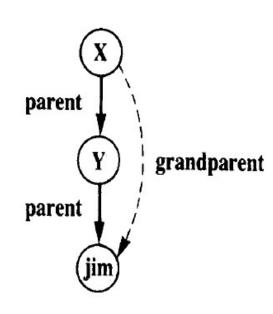
```
parent( pam, bob).
parent( tom, bob).
parent( tom, liz).
parent( bob, ann).
parent( bob, pat).
parent( pat, jim).
female(pam).
male(tom).
male(bob).
female(liz).
female(pat).
female(ann).
male(jim).
```

- ?- parent(Y, jim), parent(X,Y).
- ?- parent(tom, X), parent(X,Y).
- ?- parent(pam, X), parent(X,Y), parent(Y, jim).



## Recursive Rule

Figure 1.2 The grandparent relation expressed as a composition of two parent relations.



#### Recursive Rules

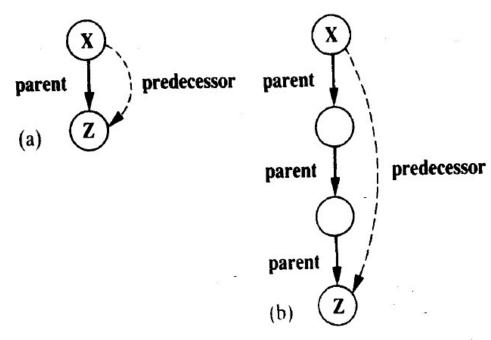
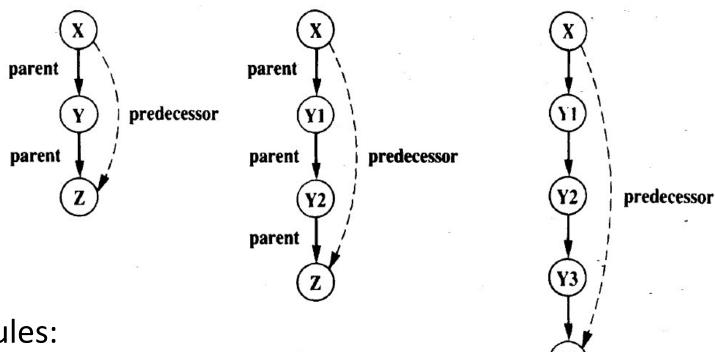


Figure 1.5 Examples of the predecessor relation: (a) X is a *direct* predecessor of Z; (b) X is an indirect predecessor of Z.

## Recursive Rules



#### • Rules:

predecessor(X,Y) :- parent(X,Y).

predecessor(X,Y) :- parent(X, Z), predecessor (Z, Y).

## Recursive Definition

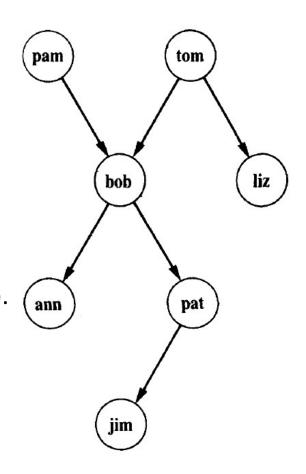
```
parent( pam, bob).
parent( tom, bob).
parent( tom, liz).
parent( bob, ann).
parent( bob, pat).
parent( pat, jim).
```

#### • Rules:

jim.

no.

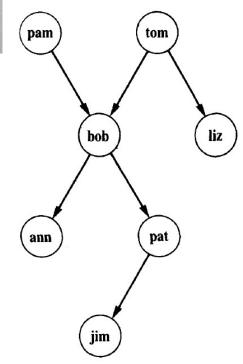
```
predecessor(X,Y) :- parent(X,Y).
predecessor(X,Y) :- parent(X, Z), predecessor (Z, Y).
?-predecessor(pam, Y), write(Y).
X=bob;
ann;
pat;
```



pred (X, Y):- parent (X, Y).

pred (X, Y):- parent (X, Z), pred (Z, Y).

?- pred (pam, jim).



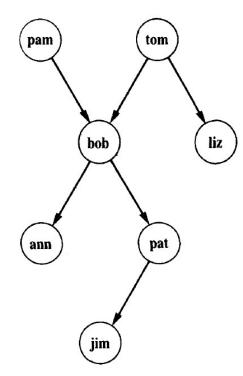
```
pred (X,Y):- parent (X,Y).

pred (X,Y):- parent (X,Z), pred (Z,Y).

?- pred (pam, jim).

(X=pam)
(Y=jim)

perent (pam,Z), pred (Z,jim).
```



pred (X,Y):- parent (X,Y).

pred (X,Y):- parent (X,Z), fred (Z,Y).

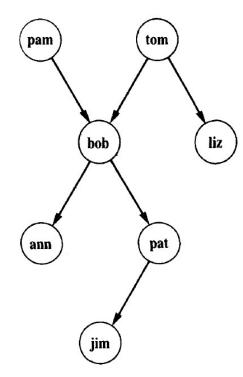
?- pred (pam, jim).

X= pam
Y= jim

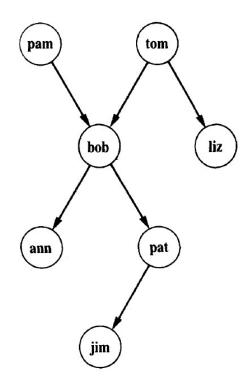
perent (pam, Z), pred (Z, jim).

Z=bob

parend (paim, bob), pred (bob, jim).



```
pred (X,Y):- parent (X,Y).
  pred (X, Y): - parent (X, Z), pred (Z, Y).
```



```
pred (X, Y): - parent (X, Y).
                 pred (x, y):- parent (x, Z), pred (Z, Y).
         ?- pred (pam, jim).

= pam
Y= jim

perent (pam, Z), pred (Z, jim).

Z=bob

parend (paim, bob), pred (bob, jim).

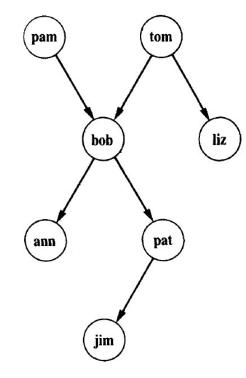
x= bob
Y= jim

parend (bob, ZI), pred (ZI, jim).
zi= ann
zi= pat

zi= pat

parent (bob, pat), pred (pat, jim)

pred (ann, jim)
false
```



```
pred (X, Y):- parent (X, Y).
          pred (X, Y): - parent (X, Z), pred (Z, Y).
          ?- pred (pam, jim).
       parent (pam, Z), pred (Z, jim).
      parent (bob, ZI), pred (ZI, jim)
parent (ann),

pred (ann, jim)

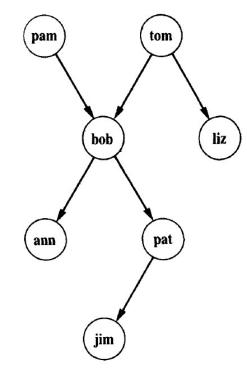
false

parent (bob, pat), pred (pat, jim)

parent (pat, jim)

force

force
```



#### Lists

- A collection of ordered data.
- Has zero or more elements enclosed by square brackets
   ('[]') and separated by commas (',').

```
    (a) ← a list with one element
    (j) ← an empty list
    (34, tom, [2, 3]] ← a list with 3 elements where the 3<sup>rd</sup> element is a list of 2 elements.
```

Like any object, a list can be unified with a variable

```
|?- [Any, list, 'of elements'] = X.
X = [Any, list, 'of elements']?
yes
```

## Recursive nature of List

e.g. 
$$L = [a, b, c]$$

can be represented as  $\rightarrow$ 
 $L = [a \mid X]$ 

or 
$$(a, X)$$

special  $(a, \cdot (b, \cdot (c, [])))$ .

or 
$$[a,b,c] = [a|[b,c]] = [a,b|[c]]$$
  
=  $[a,b,c|[]]$ 

#### **List Unification**

 Two lists unify if they are the same length and all their elements unify.

```
|?-[a,B,c,D]=[A,b,C,d]. |?-[(a+X),(Y+b)]=[(W+c),(d+b)].
A = a
                          W = a,
B = b_{\prime}
                          X = C
C = C
                          Y = d?
D = d?
                          yes
yes
|?-[[X,a]]=[b,Y].
                        |?-[[a],[B,c],[]]=[X,[b,c],Y].
                          B = b,
no
                          X = [a],
                          Y = [] ?
            Length 2
   Length 1
                          yes
```

#### **Definition of a List**

Lists are recursively defined structures.

"An empty list, [], is a list.

A structure of the form [X, ...] is a list if X is a term and [...] is a list, possibly empty."

- This recursiveness is made explicit by the bar notation
  - [Head|Rem list] ('|' = bottom left PC keyboard character)
- Head must unify with a single term.
- Rem\_list unifies with a list of any length, including an empty list, [].
  - the bar notation turns everything after the Head into a list and unifies it with Tail.

Head and rest of list List (in terms of Head and rest of the list)

```
?-[a,b,c,d]=[X|Y]. ?-[a,b,c,d]=[X|[Y|Z]].
X = a
                   X = a,
Y = [b,c,d]?
                   Y = b
                    Z = [c,d];
yes
                    yes
                         ?-[a,b,c]=[W|[X|[Y|Z]]].
?-[a] = [H|T].
H = a
                    W = a
                         X = b
T = [];
                         Y = C
yes
                           Z = []? yes
?-[] = [H|T]. ?-[a | [b | [c|[]]]] = List.
                    List = [a,b,c]?
no
                    yes
```

## Identifying a list

- **lists:** [a, [], green (bob)]
- lists are recursively defined structures:

"An empty list, [], is a list.

A structure of the form [X, ...] is a list if X is a term and [...] is a list, possibly empty."

 This can be tested using the Head and Rem\_list notation, [H|T], in a recursive rule.

#### Base and Recursive Cases

- A recursive definition, whether in prolog or some other language (including English!) needs two things.
- A definition of when the recursion terminates.
  - Without this the recursion would never stop!
  - This is called the base case: is\_a\_list([]).
  - Almost always comes before recursive clause
- A definition of how we can define the problem in terms of a similar, smaller problem.
  - This is called the *recursive case*:  $is_a_list([_|T]):-is_a_list(T)$ .
- There might be more than one base or recursive case.

#### **Focussed Recursion**

- To ensure that the predicate terminates, the recursive case must move the problem closer to a solution.
  - If it doesn't it will loop infinitely.
- With list processing this means stripping away the Head of a list and recursing on the Tail.
   Head is replaced with an underscore as we

```
is_a_list([\_|T]):-
is_a_list(T).
```

The same focussing has to occur when recursing to find a property or fact.

point ([]).

point ([x : Y]): write (x), point (Y).

?- paint ([1,2,3]).

point ([]).
point ([X | Y]): - write (X), point (Y).

?- paint ([1,2,3]).

$$X=1$$
 $Y=[2,3]$ 

point ([1,2,3])

point ([]).
point ([X | Y]): - write (X), point (Y).

Printing items in a List

?- paint (
$$[1,2,3]$$
).  
 $x=1$   
 $y=[2,3]$   
paint ( $[1,2,3]$ )

Woite(1), point ([2,3]).

point ([]) point ([x | Y]): write (x), point (Y). ?- paint ([1,2,3]).  $\begin{array}{c}
X = 1 \\
Y = [2,3]
\end{array}$ paint ([1 | 2,3]) Woite(1), point ([2,3]). χ= 2 γ= [3]

paint ([2 | 3])

point ([]). point ([x | Y]): write (x), point (Y). ?- paint ([1,2,3]). X=1 Y=[2,3]paint ([1 : 2,3]) Woite(1), point ([2,3]).  $\chi = 2$   $\chi = [3]$ paint ([2 | 3]) waite(2), paint ([3])

```
point ([]).
point ([x : Y]): write (x), point (Y).
 ?- paint ([1,2,3]).
 X=1
Y=[2,3]
point ([1 : 2,3])
  Woite(1), point ([2,3]).
          \chi = 2
\chi = [3]
          paint ([2 | 3])
      waite(2), paint ([3])
            point ([3:[]])
```

```
point ([])
point ([x | Y]): write (x), point (Y).
 ?- paint ([1,2,3]).
  X=1
Y= [2,3]
point ([1:2,3])
  Woite(1), point ([2,3]).
          x= 2
Y= [3]
          paint ([2 | 3])
      waite(2), paint ([3])
        point ([3 | []])
woite (3), point ([])
toue.
```

#### List Processing Predicates: Member/2

- Member/2 is possibly the most used user-defined predicate (i.e. you have to define it every time you want to use it!)
- It checks to see if a term is an element of a list.
  - it returns yes if it is
  - and fails if it isn't.

```
| ?- member(c,[a,b,c,d]).
yes
```

```
member(H,[H|_]).
member(H,[_|T]):-
    member(H,T).
```

- It 1<sup>st</sup> checks if the Head of the list unifies with the first argument.
  - If yes then succeed.
  - · If no then fail first clause.
- The 2<sup>nd</sup> clause ignores the head of the list (which we know doesn't match) and recurses on the Tail.

#### member/2

```
member(X, [X \mid Y]).
member(X, [Z \mid Y]):- member(X, Y).
```

?-member(3,[1, 2, 3,4]).

## List Processing Predicates: Member/2

```
[?- member(ringo,[john,paul,ringo,george]).
Fail(1): member(ringo,[john]_]).
  (2): member(ringo,[_|paul,ringo,george]):-
     Call: member(ringo,[paul,ringo,george]).
          Fail(1): member(ringo,[paull_]).
             (2): member(ringo,[_|ringo,george]):-
           Call: member(ringo,[ringo,george]).
                Succeed(1): member(ringo,[ringo|_]]).
        1) member(H , [ H | _ ]).
        2) member(H , [ _ | T]) :- member(H , T).
```

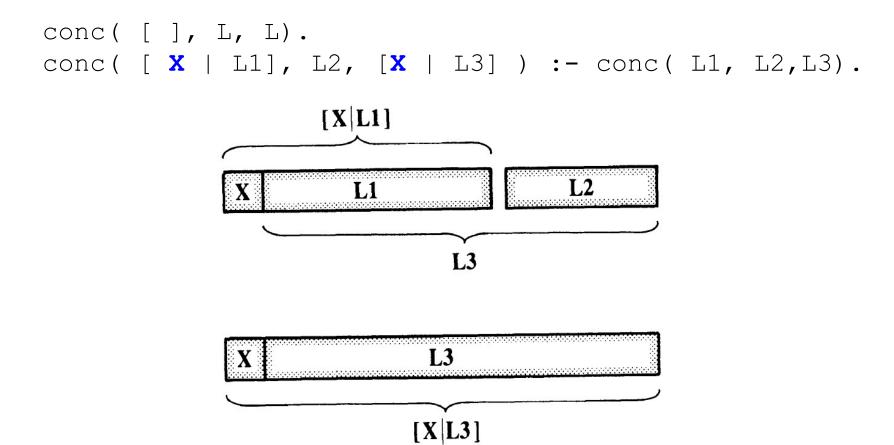


Figure 3.2 Concatenation of lists.

conc ( $\Gamma$ ], L, L). Conc ( $\Gamma$ X; L1], L2,  $\Gamma$ X; L3]):—
Conc ( $\Gamma$ 1,  $\Gamma$ 2,  $\Gamma$ 3). ?- conc ( $\Gamma$ 9,  $\Gamma$ 1,  $\Gamma$ 2,  $\Gamma$ 3],  $\Gamma$ 1.

Conc ([5], L, L).

Conc ([5], L, L).

Conc ([5], L, L):

Conc ([5], L, L):

Conc ([5], L, L).

Conc ([5], L, L).

Conc ([I], L, L).

Conc ( $[X|L|], L^2, [X|L|]$ ):—

Conc ( $[I], L^2, L^3$ ).

?- conc ( $[I], L^2, L^3$ ).

conc  $(\Gamma J, L, L)$ . Conc ([X; L1], L2, [X; L3]):Conc (L1, L2, L3) ?- conc ([a,b], [1,2,3], L).
conc ([a|b], [1,2,3], [a|L3]) Conc ([b], [1,2,3], L3) conc ([b|[]], [1,2,3], [b|L31]) Conc ([], [1,2,3], L31)

conc  $(\Gamma J, L, L)$ . Conc ([X; L1], L2, [X; L3]):Conc (L1, L2, L3) ?- conc ([a,b], [1,2,3], L).
conc ([a|b], [1,2,3], [a|L3]) conc ([b], [1,2,3], L3) conc ([b|[]], [1,2,3], [b|L31]) Conc ([], [1,2,3], L31) L31 = [1,2,3]conc([], [1,2,3], [1,2,3]). Soul

## conc (EJ, L, L). Conc ([X; L1], L2, [X; L3]):Conc (L1, L2, L3) ?- conc ([a, b], [1,2,3], L). conc ([a|b], [1,2,3], [a|L3]) Conc ([b], [1,2,3], L3) conc ([b|[]], [1,2,3], [b|L31]) conc ([], [1,2,3], L31) $\int L31 = [1,2,3]$ conc([], [1,2,3], [1,2,3]). b L31 [1,2,37 L= [9, b, 1, 2, 3]

```
conc( [ ], L, L).
conc( [ X | L1], L2, [X | L3] )
:- conc( L1, L2, L3).
```

- Although the conc program looks rather simple it can be used flexibly in many other ways.
- For example, we can use conc in the inverse direction for decomposing a given list into two lists, as follows:

## ?- conc( L1, L2, [a,b,c] ).

no

```
conc( [ ], L, L). conc( [ \mathbf{X} | L1], L2, [\mathbf{X} | L3] ) :- conc( L1, L2,L3).
```

- We can also use this program to look for a certain pattern in a list.
- example, we can find the months that precede and the months that follow a given month say may, as in the following goal:
  - ?- conc( Before, [may | After], [jan,feb,mar,apr,may,jun,jul,aug,sep,oct,nov,dec] ).

```
Before = [jan,feb,mar,apr]
After = [jun,jul,aug,sep,oct,nov,dec].
```

```
conc([], L, L). conc([\mathbf{X} | L1], L2, [\mathbf{X} | L3]) :- conc(L1, L2,L3).
```

 Further we can find the immediate predecessor and the immediate successor of may by asking:

```
?- conc(_, [Month1,may,Month2 | _],
        [jan,feb,mar,apr,may,jun,jul,aug,sep,oct,nov,dec] ).
```

```
Month1 = apr
Month2 = jun
```

```
conc( [ ], L, L). conc( [ \mathbf{X} | L1], L2, [\mathbf{X} | L3] ) :- conc( L1, L2,L3).
```

 we can, for example, delete from some list, L1, everything that follows three successive occurrences of z in L1 together with the three z's. For example:

## List processing in Prolog

```
english_spanish("One", "Uno").
english_spanish("Two", "Dos").
english_spanish("Three", "Tres").
english_spanish("Four", "Cuatro").
english_spanish("Five", "Cinco").
english_spanish("Six", "Seis").
english_spanish("Seven", "Siete").
english_spanish("Eight", "Ocho").
english_spanish("Nine", "Nueve").
english_spanish("Ten", "Diez").
translate_td([], []).
```

```
%Traditional top-down iteration
%% translate_td(?EnglishList,
?SpanishList) is det
```

- One of the advantages of this is its bidirectionality.
- We can either put in a list of English numbers and get Spanish numbers, or vice versa.

```
translate_td([English|EnglishList], [Spanish|SpanishList]) :-
   english_spanish(English, Spanish), translate_td(EnglishList, SpanishList).

?- translate_td(["One", "Two", "Three", "Four", "Five", "Six", "Seven", "Eight",
"Nine", "Ten"], S).
?- translate_td(E, ["Uno", "Dos", "Tres", "Cuatro", "Cinco", "Seis", "Siete", "Ocho",
"Nueve", "Diez"])
```