Security Risks in Zero Knowledge Proof Cryptocurrencies



Zhiniang Peng of 360 Core Security

Who am I

Zhiniang Peng
Ph.D. in cryptography
Security researcher @Qihoo 360
Twitter: @edwardzpeng

Research areas:

Software security

Applied cryptography

Threat hunting

Outlines

Introduction

Security Risks in ZKP Cryptocurrencies

Implementation vulnerability

Trust risk

Info leak in Tx

Crypto fail

Others

ZKP for hackers

Key theft

Selling hacked database

Selling Oday

Conclusion

The Privacy of Bitcoin

Bitcoin: decentralized digital currency

Public verifiable, No anonymity

Privacy issues:

Personal cash flow

Account balance

Money becomes unequal:

Black money

Souvenir Coin

Privacy VS Public Verifiability

From	Alice
То	Bob
Amount	1

From	Scrooge			
То	Donald			
Amount	2			

From	Bob			
То	Eve			
Amount	1			

Privacy → Encryption?

From	Enc(A)		
То	Enc(B)		
Amount	Enc(1)		

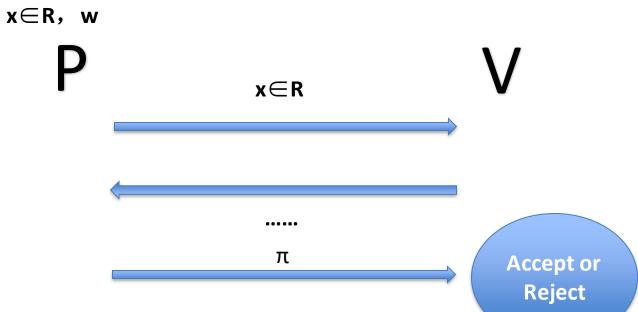
Enc(S)
Enc(D)
Enc(2)

_

From	Enc(B)
То	Enc(E)
Amount	Enc(1)

Encryption Conflict with Public verifiability





Completeness: If the statement is correct, the prover can persuade the verifier

Soundness: If the statement is wrong, the prover cannot persuade the verifier

Zero knowledge: The verifier is unable to obtain any other information except that the statement is correct

Non interactive zero knowledge proof (NIZK)

If factorization is difficult, forwary NP language there is a NIZK

ZKP with Bitcoin

_	_							
	From	Enc(A)	From	Enc(S)	From	Enc(B)	From	C ₁
	То	Enc(B)	То	Enc(D)	То	Enc(E)	То	C ₂
	Amount	Enc(1)	Amount	Enc(2)	Amount	Enc(1)	Amount	C ₃
	Proof	π	Proof	π'	Proof	π''	Proof	π'''

Encryption all the transactions

Proof π :

Non interactive zero knowledge proof

Prove that encrypted transactions are legal

Encryption Conflict with Public verifiability (Solved!)

zkSNARKs

High performance requirements for blockchain Performance of Universal ZKP

zkSNARKs

Succinct

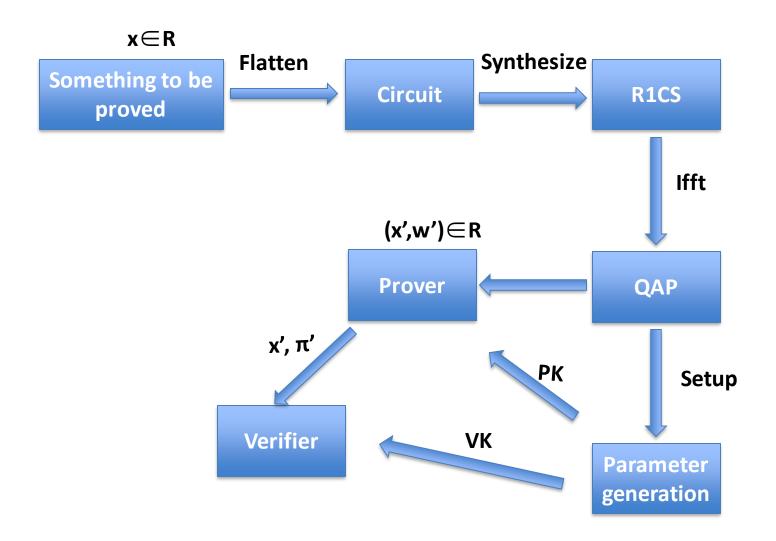
Non-interactive

Argument

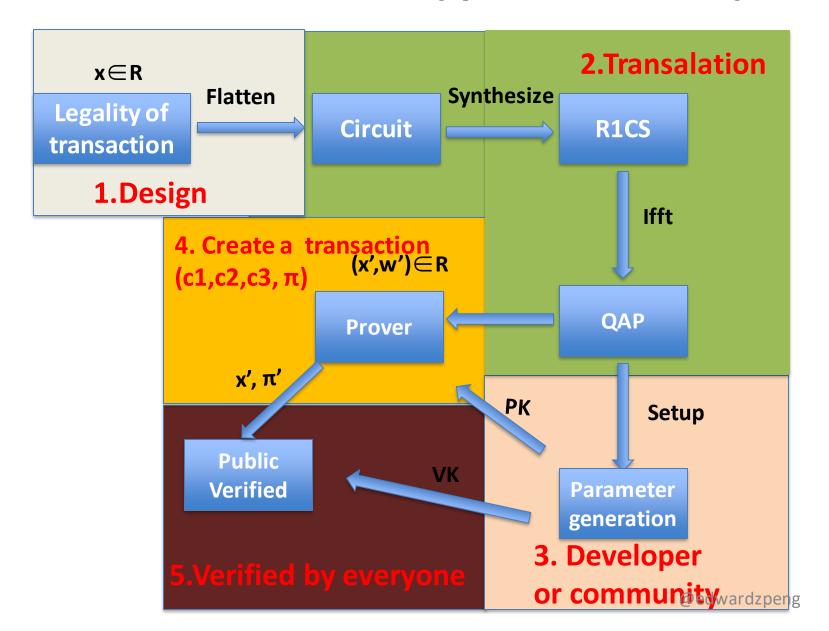
of Knowledge

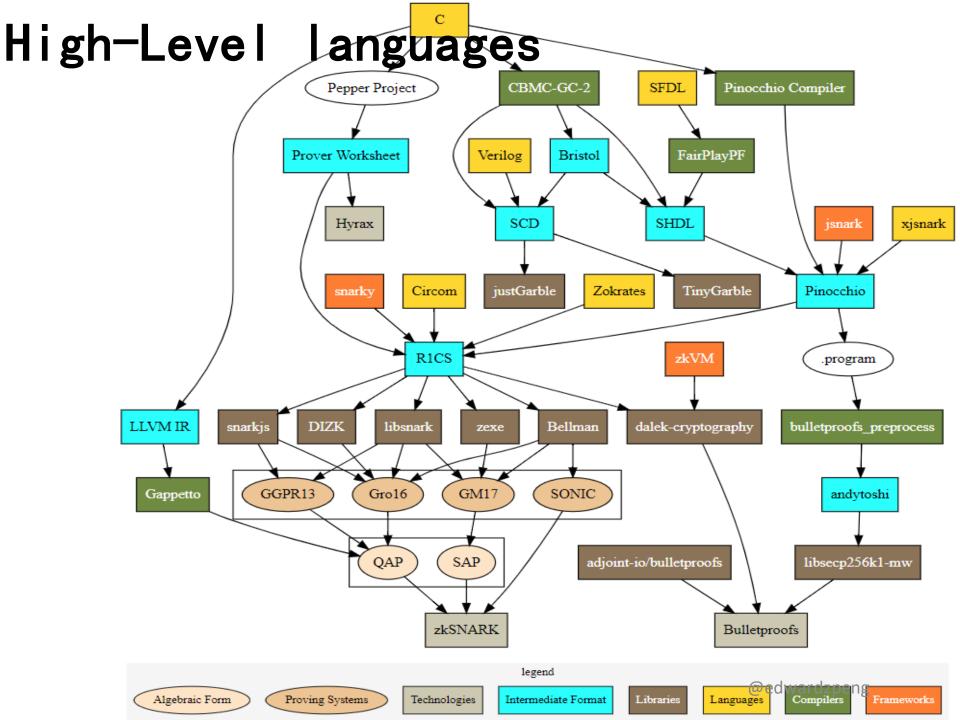
Zero-knowledge

How zkSNARKs works



zkSNARKs in Cryptocurrency





Security Risks in ZKP Cryptocurrencies

Implementation vulnerability

Trust risk

Info leak in Tx

Crypto fail

Others

Implementation vulnerability

Category

Memory corruption ⊗

Most Implemented in : rust \ java \ go (memory safety)

libSNARKs(C++): hard to exploit

Logic bug ©

Protocol Design

Circuit Implementation

Application logic

Crypto Implementation

New crypto -> new bugs!

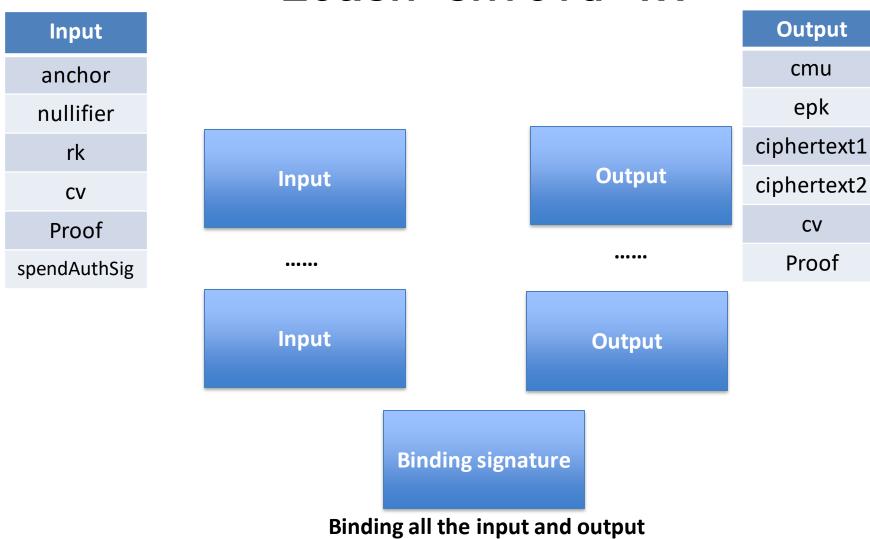
Protocol Design

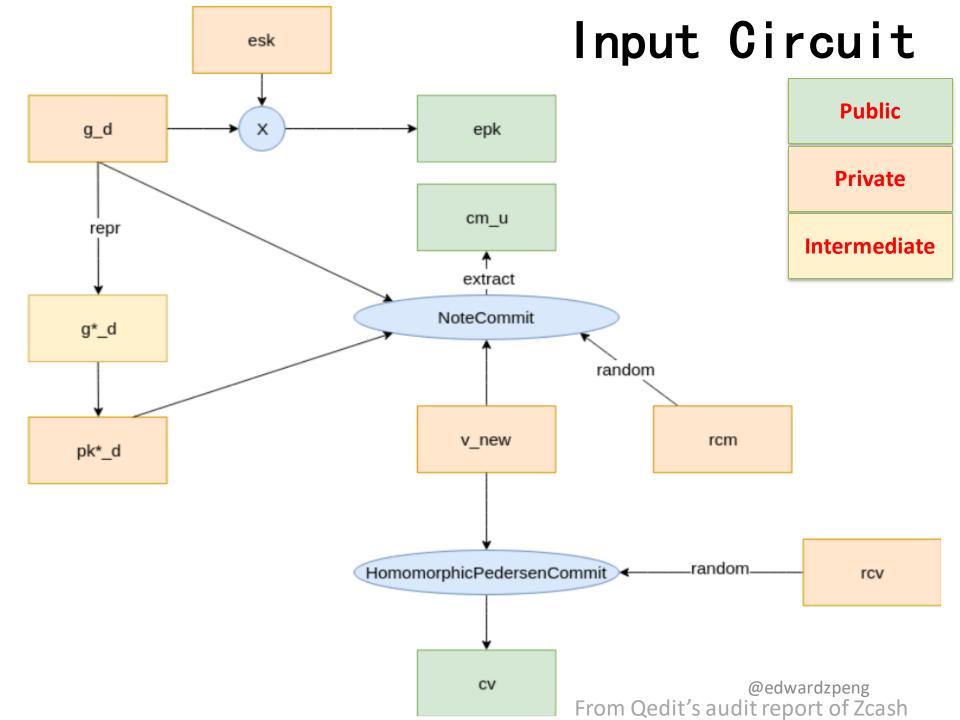
Complex

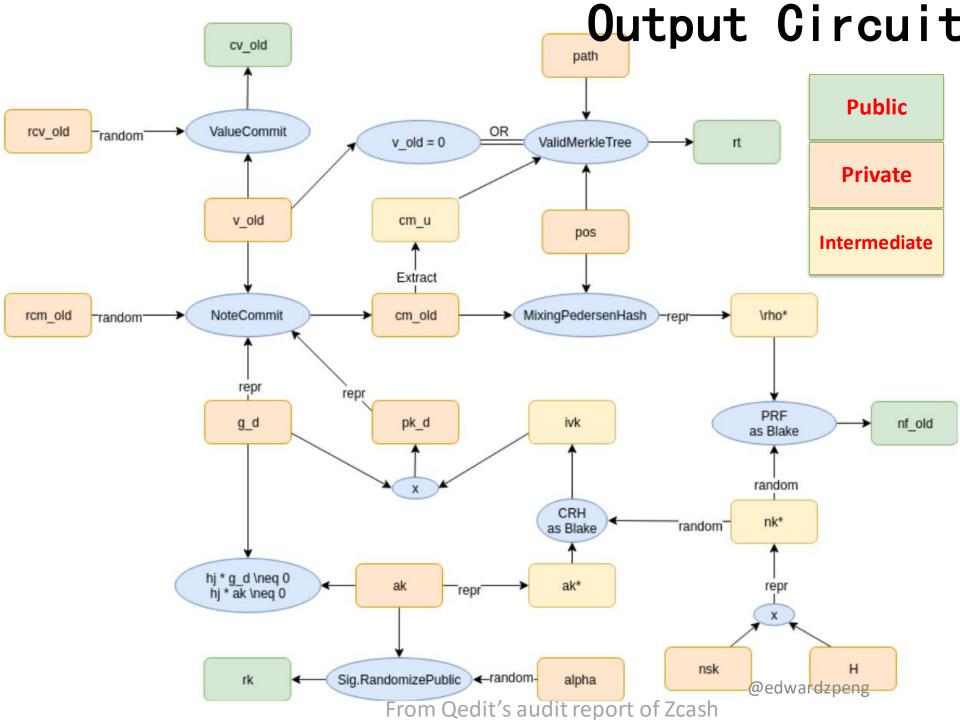
Privacy, performance, ease of use -> Complex Lots of crypto Only expert can review

Zcash for example

Zcash shield TX







Protocol Design bug

Zcash Faerie Gold attack:

Attacker send two note with the same rho

Only one can be used

Fake money

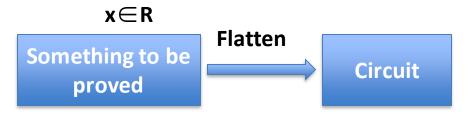
Fix: Force to be different by crypto

URL: https://github.com/zcash/zcash/issues/98

Similar vulnerability also occur in Monero

Circuit Implementation

What is Circuit?



The same logic implemented twice:

In typical programing languages and in Circuit

Inconsistence between two implementation

Result in critical vulnerability

Heavily check in all the popular projects

Bugs in Application

Application developer calls ZKP lib to build application

Due to the lack of sufficient understanding of the ZKP, their code is easy to be vulnerable.

Semaphore double spend

Vulnerability allowing double spend #16



poma opened this issue on 26 Jul · 2 comments



poma commented on 26 Jul • edited ▼



Looks like in Semaphore.sol#L83 we don't check that nullifier length is less than field modulus. So nullifier hash + 21888242871839275222246405745257275088548364400416034343698204186575808495617 will also pass snark proof verification if it fits into uint256, allowing double spend.

Example of 2 transactions:

https://kovan.etherscan.io/tx/0x5e8bf35ad76a086b98698f9d20bd7b6397ccc90aa6f85c1c5debc0262be54 58a

https://kovan.etherscan.io/tx/0x9a47cc8daec9d0a5e9a860ada77730190124f9864a5917dcb8f41773d94cf c1a



To prevent double spend: a unique Nullifier for each Tx Semaphore: If Nullifier n works, Nullifier n+p also works Double spend if you don't check the length @edwardzpeng

Tron double spend

Tron use Librustzcash to build their privacy solution

Directly use functions in the Lib without constrains on variables



Tron nullifier

One transaction has multiple input Does not check the duplicate

```
src/main/java/org/tron/core/actuator/ShieldedTransferActuator.java
@@ -103,9 +103,14 @@ private void executeTransparentTo(byte[] toAddress, long amount) throws Contract
     //record shielded transaction data.
     private void executeShielded(List<SpendDescription> spends, List<ReceiveDescription> receives) {
     private void executeShielded(List<SpendDescription> spends, List<ReceiveDescription> receives)
         throws ContractExeException {
       //handle spends
       for (SpendDescription spend : spends) {
         if (dbManager.getNullfierStore().has(
             new BytesCapsule(spend.getNullifier().toByteArray()).getData())) {
          throw new ContractExeException("double spend");
         dbManager.getNullfierStore().put(new BytesCapsule(spend.getNullifier().toByteArray()));
                                                                                  @edwardzpeng
```

Crypto Implementation

New crypto -> new bugs! LibSNARKs:

R1CS-to-QAP reduction bug

Linear dependent -> soundness error

Fix: Increase redundancy

https://eprint.iacr.org/2015/437.pdf

Zcash side channel attack

Reject attack, Ping attack

Node processing a transaction related with himself:

Side channel in decryption time -> Info leak

Break the anonymity

https://crypto.stanford.edu/timings/pingreject.pdf

Trust Risk

Trust Risk

Basic ideas in zkSNARKs:

Generate the challenge (x) ahead of time keep encrypted (x), discard plaintext (x)

Verifying the proof (A,B,C):

$$A \cdot B = lpha \cdot eta + rac{\sum_{i=0}^{c} a_i (eta u_i(x) + lpha v_i(x) + w_i(x))}{\gamma} \cdot \gamma + C \cdot \delta$$

Undetectable backdoor (x)

You can create proof for any statement if you know x

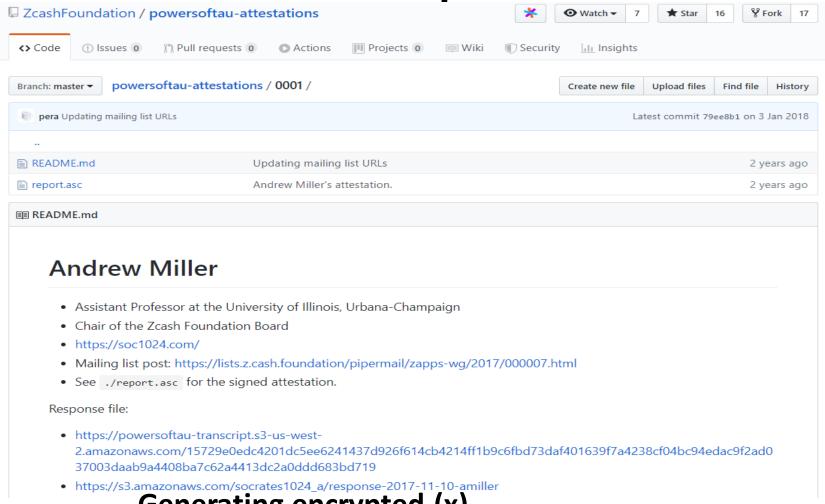
Create fake proof -> create fake money ©

Zero knowledge proof -> undetectable!

To solve Trust risk:

Generate encrypted (x) by secure multiparty computation (MPC) ----> nobody know the plaintext of (x)

Zcash MPC phase 1



Generating encrypted (x)

Public verifiable

More than 100 participants

Secure as long as at least one person is honest

Zcash MPC phase 2

Participants:

- Sean Bowe
 - 4caed5dfc4fd6959b5181c4adc31fe013c58b438c360d30af5a40fe02cc19a9ad56344679524e
 07fbad6166cc44ee35e1f6e70e5267b46ce93d059cad3ab7ed5
- Ian Munoz
 - dd8ca189564354d5c9405b16b32f8dd0ed4644e4184e9fbc9feaccebd732ad8276033318b0b
 70fa5b1f59f49ce3ed42f3f575b9e6219c11d3e28a82ed24ca990
- Jason Davies
 - 851b81478d1ce92c1f7116c8f355c999447e36a151b97f36ff8d5cbf9b98a5b00431b2a087da2
 038d201669dd3a8eba8bbd2d92ea8e348b92e1b51b09fe5bb5f
- Larry Ruane
 - 82641343d6840bcb218a7cef8f46d6f6678fff3692e8ec6b9bfc44245737d0c047dd441fd3608e a1463302a8c9521dcd716df3b114741bcdf4da5add68bf35d3
- Gabor Losonci
 - be9fe35787982dda14b81a796dbab4a5ecfe1bf2d56c6b0186e4552b5a942d61a5378c4c903a a3087226b7bdf0fcf7808eaa6a2b1a1ae0d42793bb9836718a23
- Michael Perklin
 - ec782c5b9b33d98750012e3d1d62f34b490acd2d392a30629e98596198143e414536f30bd94
 14ee01dc26fc7d132ca245e400c738073a68cad9e439a5acb83a6
- Alexey Ermishkin
 - ce07ea6e537f106c0b9b70a5aff9dcc8105dfc1f5f69e32d272f7749e2b41ba728e3ef860d57ab
 933bdfb458c675bd5cd46c08429a1229dc8b21dc8e95f5815b

Generate parameter for specific circuit Some project lack this phase:

-> backdoor again

Personal experience in MPC

Lots of project are doing MPC

Ethereum, Tron

Some personal experience

Zcash: no reply ⊗

Tron: Waiting for 3 month now 🕾

Ethereum:

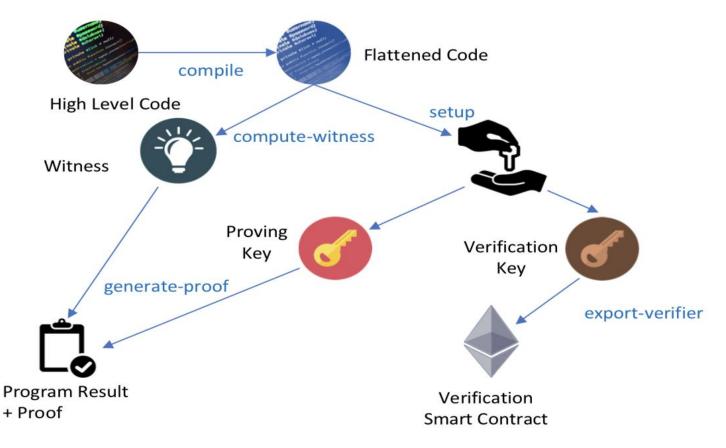
<u>https://github.com/weijiekoh/perpetualpowersoftau/blob/master/0011_zhiniang_response/zhiniang_attestation.txt</u>

MPC process is completely controlled by organizer

Black Box style ⊗

If you're not involved, you can't be 100% sure

demo: ZoKrates



ZoKrates: a toolbox for zkSNARKs on Ethereum

Compile your high level code into verification smart contract

Black box setup: No MPC

Backdoor 😊

Other ZKP systems

	libSNARK	Ligero	Bulletproofs	Hyrax	libSTARK	Aurora	Libra		
	[14]	[6]	[17]	[50]	[9]	[12]			
	O(C)						O(n)		
\mathcal{G}	per-statement		no trusted setup						
	trusted setup						trusted setup		
\mathcal{P}	$O(C \log C)$	$O(C \log C)$	O(C)	$O(C \log C)$	$O(C \log^2 C)$	$O(C \log C)$	O(C)		
ν	O(1)	O(C)	O(C)	$O(\sqrt{n} + d \log C)$	$O(\log^2 C)$	O(C)	$O(d \log C)$		
$ \pi $	O(1)	$O(\sqrt{C})$	$O(\log C)$	$O(\sqrt{n} + d \log C)$	$O(\log^2 C)$	$O(\log^2 C)$	$O(d \log C)$		
\mathcal{G}	1027s	s		NA			210s		
\mathcal{P}	360s	400s	13,000s	1,041s	2,022s	3199s	201s		
\mathcal{V}	0.002s	4s	900s	9.9s	0.044s	15.2s	0.71s		
$ \pi $	0.13KB	1,500KB	5.5KB	185KB	395KB	174.3KB	51KB		

Lots of new schemes have emerged recently,

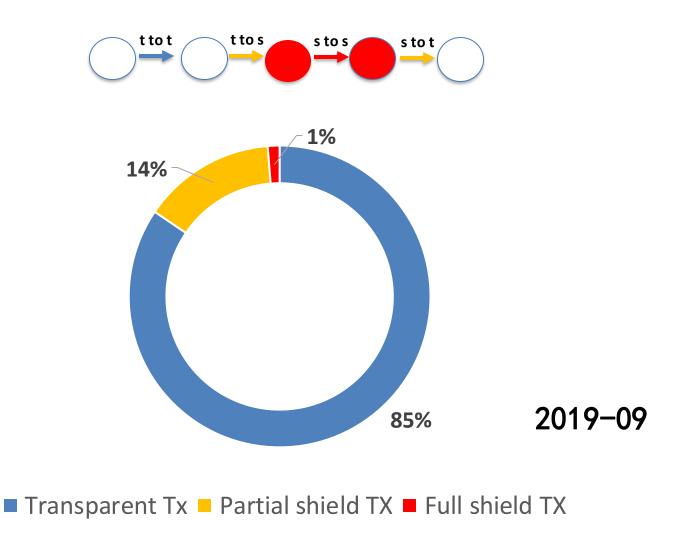
Trade-off between security, performance, trust model.

More options for the future project

@edwardzper

Info leak in Tx

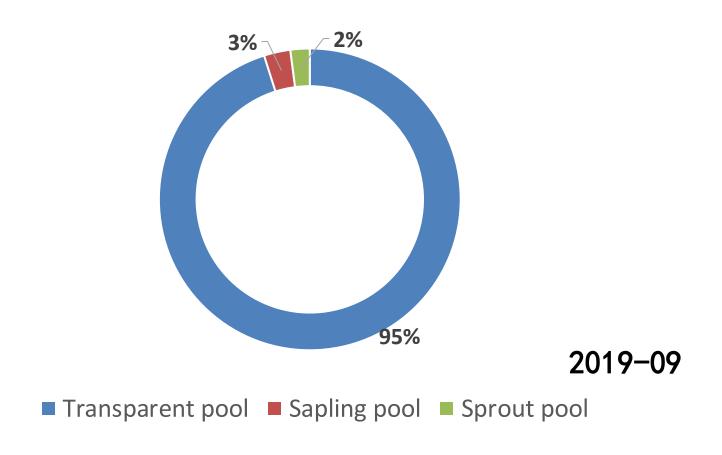
Zcash Transaction Statistics



Most transactions are traceable

@edwardzpeng

Zcash Coin Statistics



An Empirical Analysis of Anonymity in Zcash

George Kappos, Haaroon Yousaf, Mary Maller, and Sarah Meiklejohn
University College London
{georgios.kappos.16,h.yousaf,mary.maller.15,s.meiklejohn}@ucl.ac.uk

On the linkability of Zeash transactions

Jeffrey Quesnelle University of Michigan-Dearborn

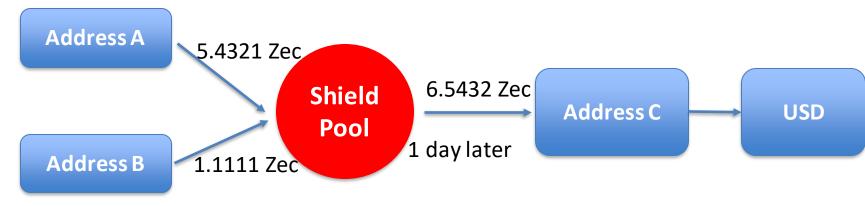
An Empirical Analysis of Traceability in the Monero Blockchain

Malte Möser, Kyle Soska, Ethan Heilman, Kevin Lee, Henry Heffan, Shashvat Srivastava Kyle Hogan, Jason Hennessey, Andrew Miller, Arvind Narayanan, Nicolas Christin

PETS 2018: The 18th Privacy Enhancing Technologies Symposium

@edwardzpeng

Empirical Analysis



Info leak in Tx

Similar user habits

Relevant address

Equal amount

Similar time

Most of the shield transactions are traceable!

Privacy of Lightweight Node

Most individual are using lightweight node

Mobile phone, Embedded Wallet

Do not have all block data

ZKP computation cost heavily

Not so friendly to lightweight node

Improving ©

Needs to hand over the decryption key to a full node

ZKP transaction need to be decrypted by full node

Then relay to the lightweight node

Almost no privacy for lightweight node 🕾

"Privacy" versus "Ease of use"

There is no good solution in the market ardzpeng

Tron scanNote RPC

```
@Override
public void scanNoteByIvk(GrpcAPI.IvkDecryptParameters request,
    StreamObserver<GrpcAPI.DecryptNotes> responseObserver) {
  long startNum = request.getStartBlockIndex();
  long endNum = request.getEndBlockIndex();
  try {
    DecryptNotes decryptNotes = wallet
        .scanNoteByIvk(startNum, endNum, request.getIvk().toByteArray());
    responseObserver.onNext(decryptNotes);
  } catch (BadItemException | ZksnarkException e) {
    responseObserver.onError(getRunTimeException(e));
  responseObserver.onCompleted();
```

Lightweight node Needs to hand over the decryption key to a full node

Info leak in Tx

Reason:

Almost no privacy for lightweight node

Most Exchange only supports transparent transaction

Countermeasures:

Use shield transaction

Use new address each time

Split the amount

Wait for enough time

Best practice:

Only use shield transaction

Crypto Fail

Risk in Crypto

ZKP technology is relatively new

Paper publish in 2016, large scale use in 2017

Some kind of radical

Parameter selection and optimization are also radical

Time will tell the truth

Provable security

Rely on too many hard problems

Some problems are not standard

The proof itself need more auditing

Zcash Counterfeiting Vulnerability

CVE-2019-7167

Discover in 2018/03 (Ariel Gabizon), publish in 2019

Vulnerability in ZKP system [BCTV14]

Anyone can create fake proof -> create fake Zcash

Affecting multiple projects

It take 8 month to fix

Change the whole proof system and upgrade the whole network It's Zero knowledge -> No one knows whether it has been exploited

Zcash official "believe" that it has not been exploited:

Few people have the skill to find this vulnerability

Total amount of Zcash seem remain unchanged

Zcash Counterfeiting Vulnerability

CVE-2019-7167

[BCTV14] doesn't have provable security

[BCTV14] redundant elements in parameter result in fake proof

The basic idea of this attack is quite simple

[BCTV14] upgrade to [Groth16]

Provable security

Another crypto bug found by Bryan Parno before

Happens to be unexploitable

Will it be the last time?

Zcash Hash Collusion Attack

Zerocash hash collusion attack

Commitment: COMMrcm and COMMs

Need to be computationally bind

Truncated to 128bits, result in 64bits security

Result in double spend

Potentially creating arbitrary amount of currency

Other issues

Perfect Zero Knowledge?

Perfect zero knowledge scheme in Mathematics

[Groth16]

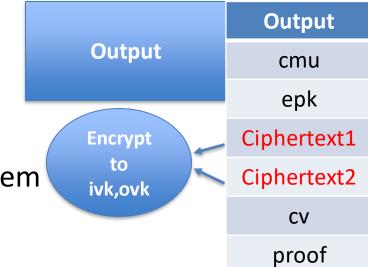
Info leak in real usage:

Zcash:

Proof is perfect zero knowledge

Ciphertext security rely on hard problem

Transaction is not zero knowledge



Perfect Zero Knowledge ≠ Perfect Privacy

Maybe 20 or 30 years latter, the encryption is broken

Your privacy will gone, remember every thing is on the chain ⊗

Unlinkability of diversified addresses

Rely on the DDH problem in JubJub curve
Not a standard ECC curve, optimize for ZKP performance
Maybe broken someday in the future

All the transaction is on the chain, Unlinkability will not remain forever.

Side Channel Attack

Always ignored in traditional software security

hard to exploit

Not directly result in compromise

Very important in a privacy system

Privacy is directly compromised

Groth16 side channel attack

Timing attack:

To compute a proof, you need to calculate (A,B,C)

$$A = lpha + rac{\sum_{i=0}^m a_i u_i(x) + r \delta$$

$$B = eta + rac{\sum_{i=0}^m a_i v_i(x) + s \delta}{}$$

$$C = rac{\sum_{i=\ell+1}^m a_i (eta u_i(x) + lpha v_i(x) + w_i(x)) + h(x)t(x)}{\delta} + As + rB - rs\delta$$

Computing time of A,B,C strongly related with private input a

Other side channel: Cache

ZKP for Hackers

Key Theft

Bob need a signing key of some authority, He found Alice on the Dark net.

But they don't trust each other, and there is no an trusted third party.

How can they fairly make a deal without trust?

Smart Contract: Key Theft



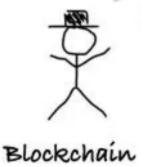




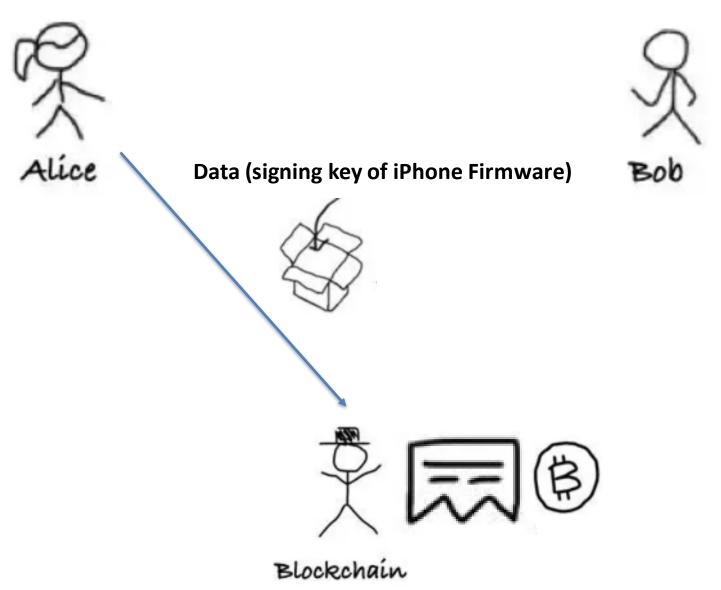
Setup an Contract

If data x satisfied f(x)=0,

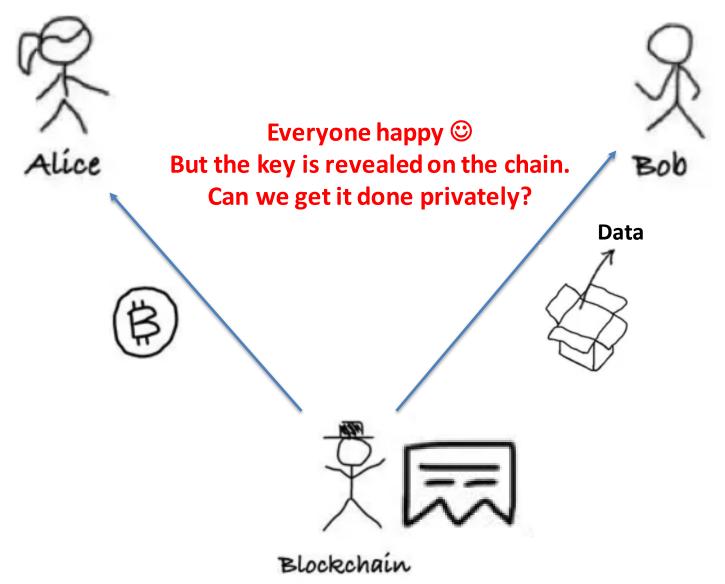
then send money to Alice



Smart Contract: Key Theft



Smart Contract: Key Theft



Zero Knowledge contingent Payment (ZKCP)



Data (signing key of iPhone Firmware)

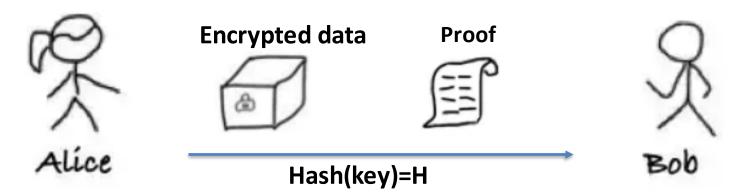
Encrypted data





Proof: The data satisfied f(x)=0

Hash(Key)=H









Bob

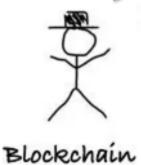


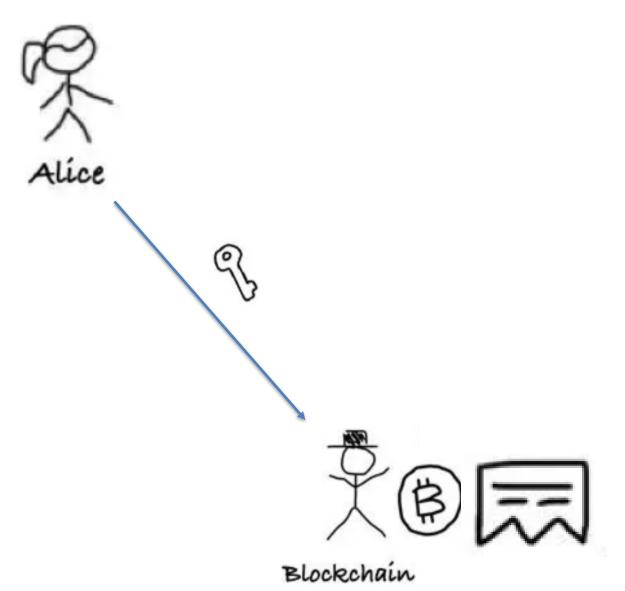


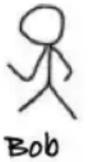
Setup an Contract

If K satisfied Hash(K)=h,

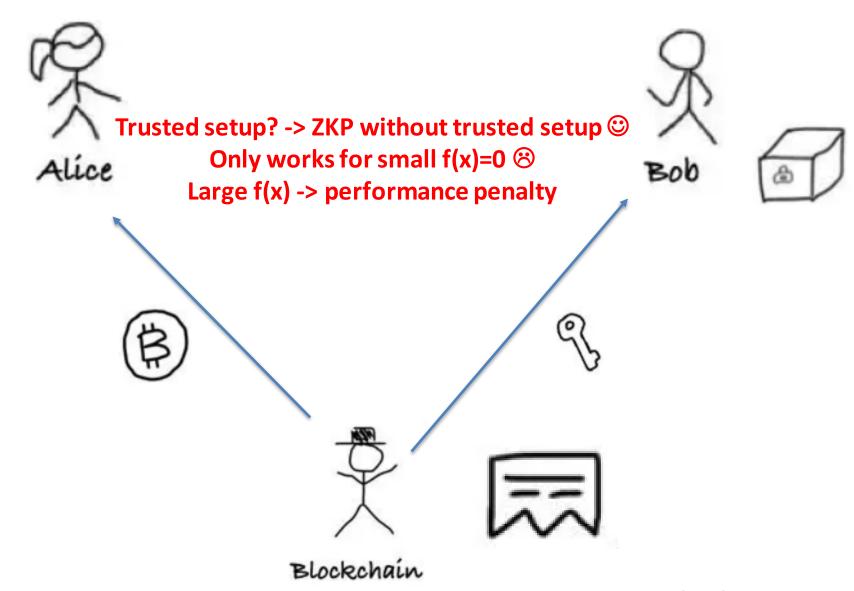
then send money to Alice



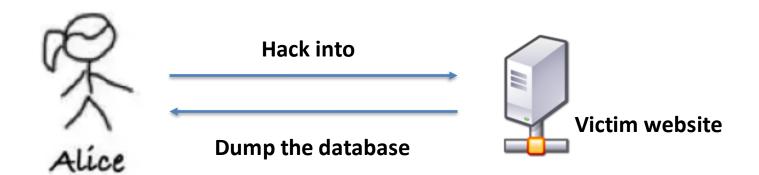








Selling Hacked Database



	Username	Telephone	Email	Password
X1 X2	Edward	77777	<u>e@360.cn</u>	edward
	Zhiniang	88888	<u>z@360.cn</u>	123456
Xn	Peng	99999	a@360.cn	201911

Fake database?

What if the database is a fake one?

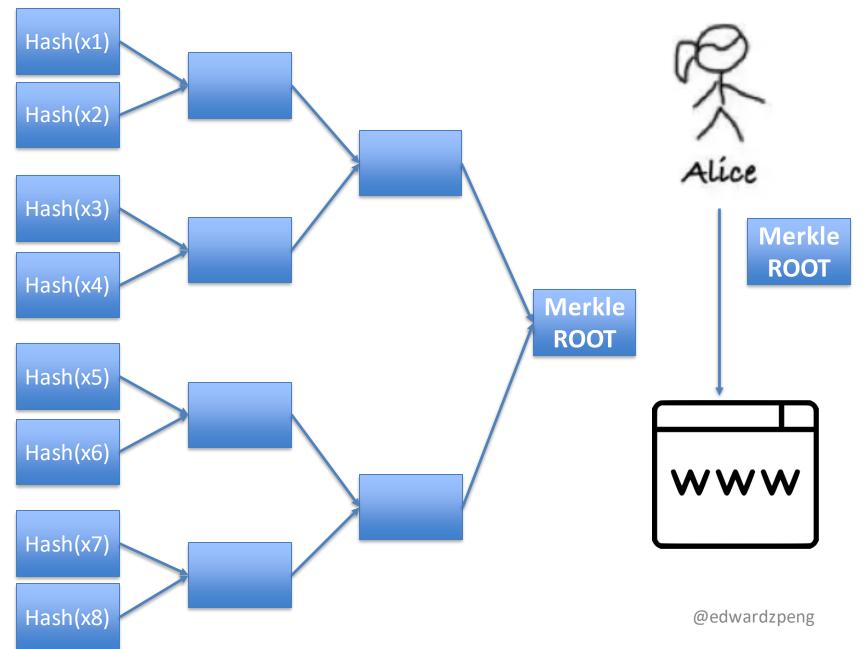
There is no function f(x) to measure the correctness

Solution:

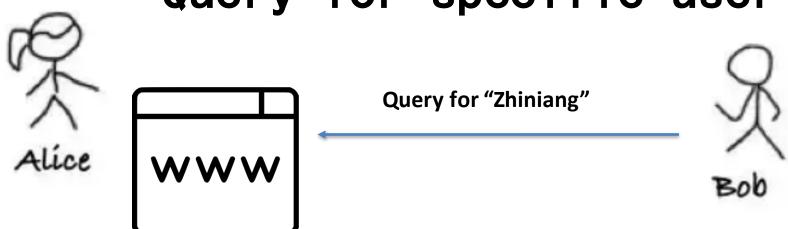
Alice can commit to the database

Bob buy some data to check

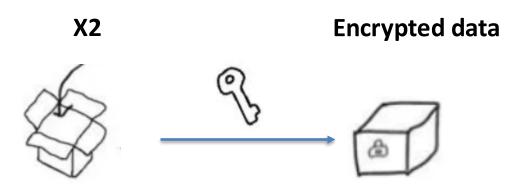
Publish a commitment



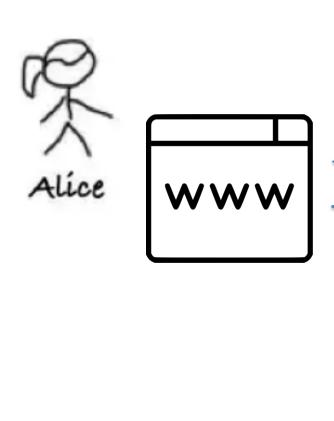
Query for specific user



Search the database X2 contain "Zhiniang"



Reply



Query for "Zhiniang"



Bob



Encrypted X2



Proof:
X2 contain "Zhiniang"
X2 is in the Merkle Tree
Hash(Key)=H





Bob

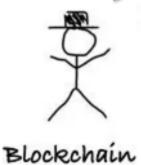


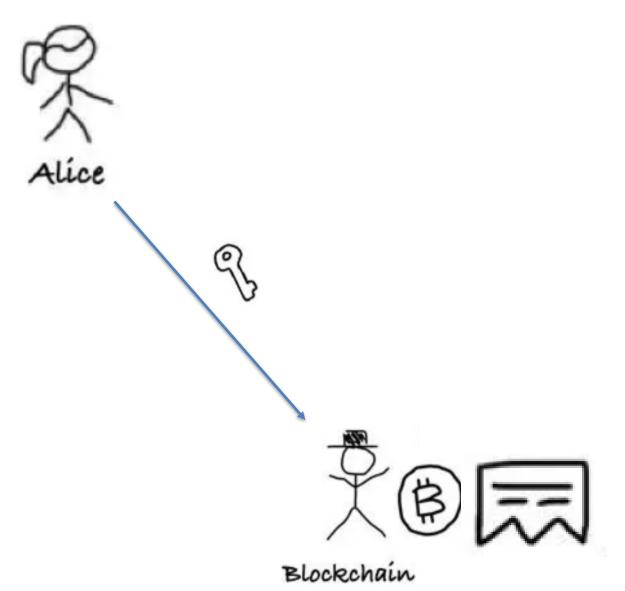


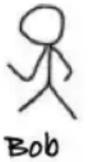
Setup an Contract

If K satisfied Hash(K)=h,

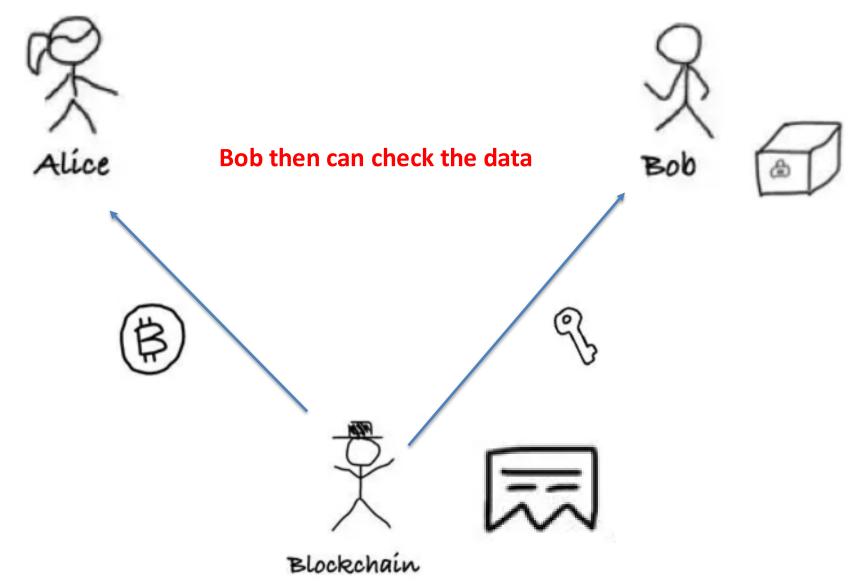
then send money to Alice











Performance Problem

The performance?

The database may be more than 100G

Directly use ZKP on them will be a extremely slow

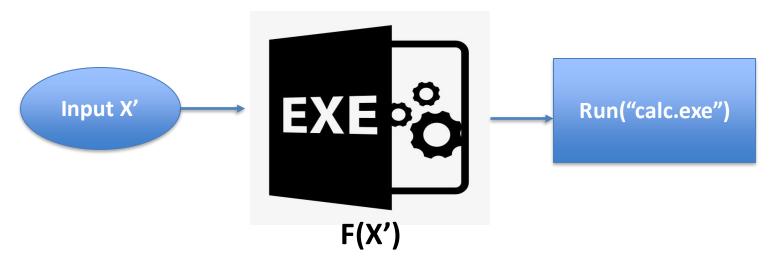
Efficient fair exchange protocol:

Fairswap: https://eprint.iacr.org/2018/740.pdf

zk-pod: https://github.com/sec-bit/zkPoD-node

Performance is reasonable

Can Alice fairly sell a Oday exploit?



Theoretically possible

Prove F(X')=Run("calc.exe")

Then sell X'

Difficulty

Simulate F(X) correctly

Performance

zkSNARKs for Von Neumann Architecture

https://eprint.iacr.org/2013/879.pdf

@edwardzpeng

Conclusions

ZKP technology is relatively new and develops rapidlyThere are risks, but it is improving

ZKP cryptocurrency can provide a strong anonymity

If you use it really carefully

ZKP application for hackers

Try to build your own protocol

Thanks



