Approximate nearest neighbours search

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Refresh

- Vector and Metric space
- How similar and different LSA and word2vec?
- DSSM from MS and Yandex

Agenda

- ANNS (not ANNs)
  - Clustering
  - Proximity graphs
  - Trees (second lecture)

Before we start...

What’s wrong with inverted index in terms of data structure?

Do you know the difference of $O(N)$, $O_A(N)$, $E(N)$?

Approximate Nearest Neighbours Search

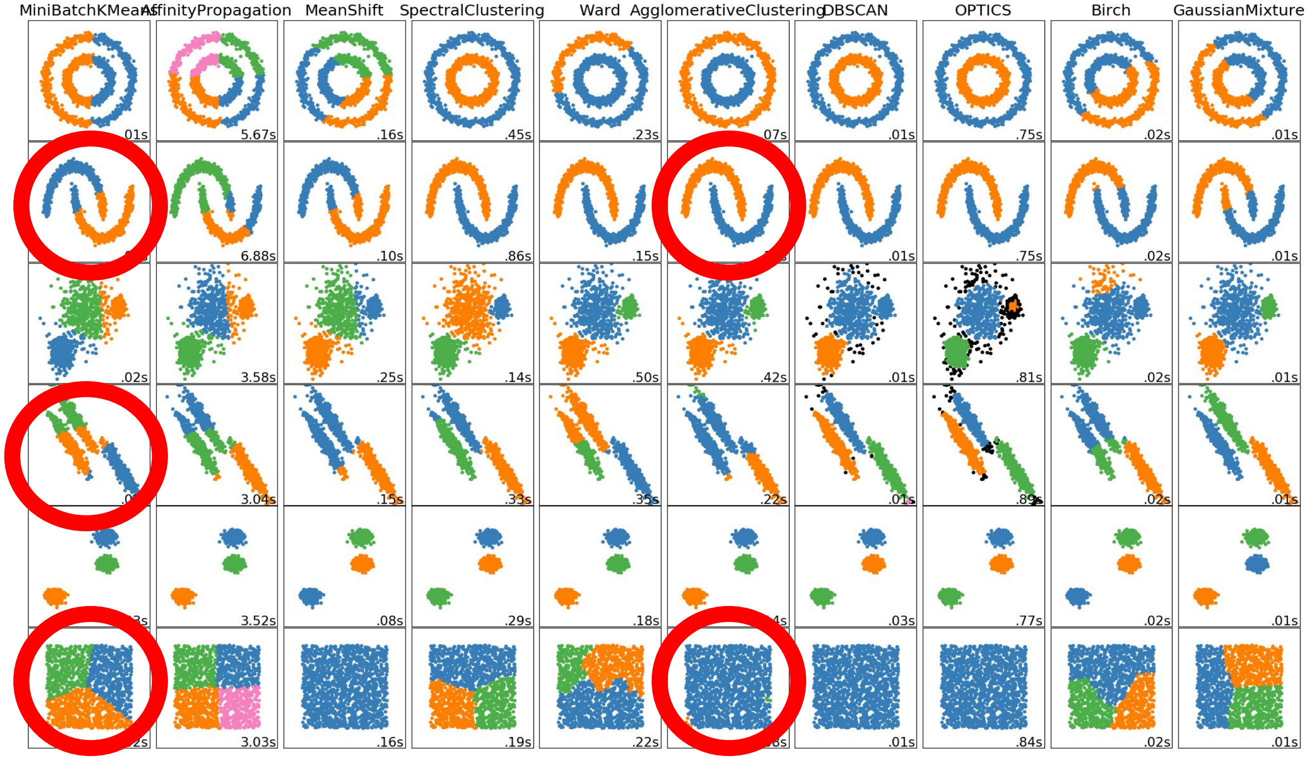
Approximation for k-NN search

1. Pre-select **k*c** elements from approximate neighbourhood (pre-ranking set).
2. Then select and re-rank relevant ones.

- Locality sensitive hashing
- **Search trees and supporting data structures**
- Vector compression, clustering, inverted indexing
- **Proximity graphs**

Hierarchical clustering and Inverted index revised

How clustering differ?



https://scikit-learn.org/stable/modules/clustering.html

Linkage criteria

Single linkage (smallest distance) ~ DBSCAN

Complete linkage (maximum distance)

Minimum energy (variance grows slowly in we merge)

Average distance and centroid-based approaches — kMeans

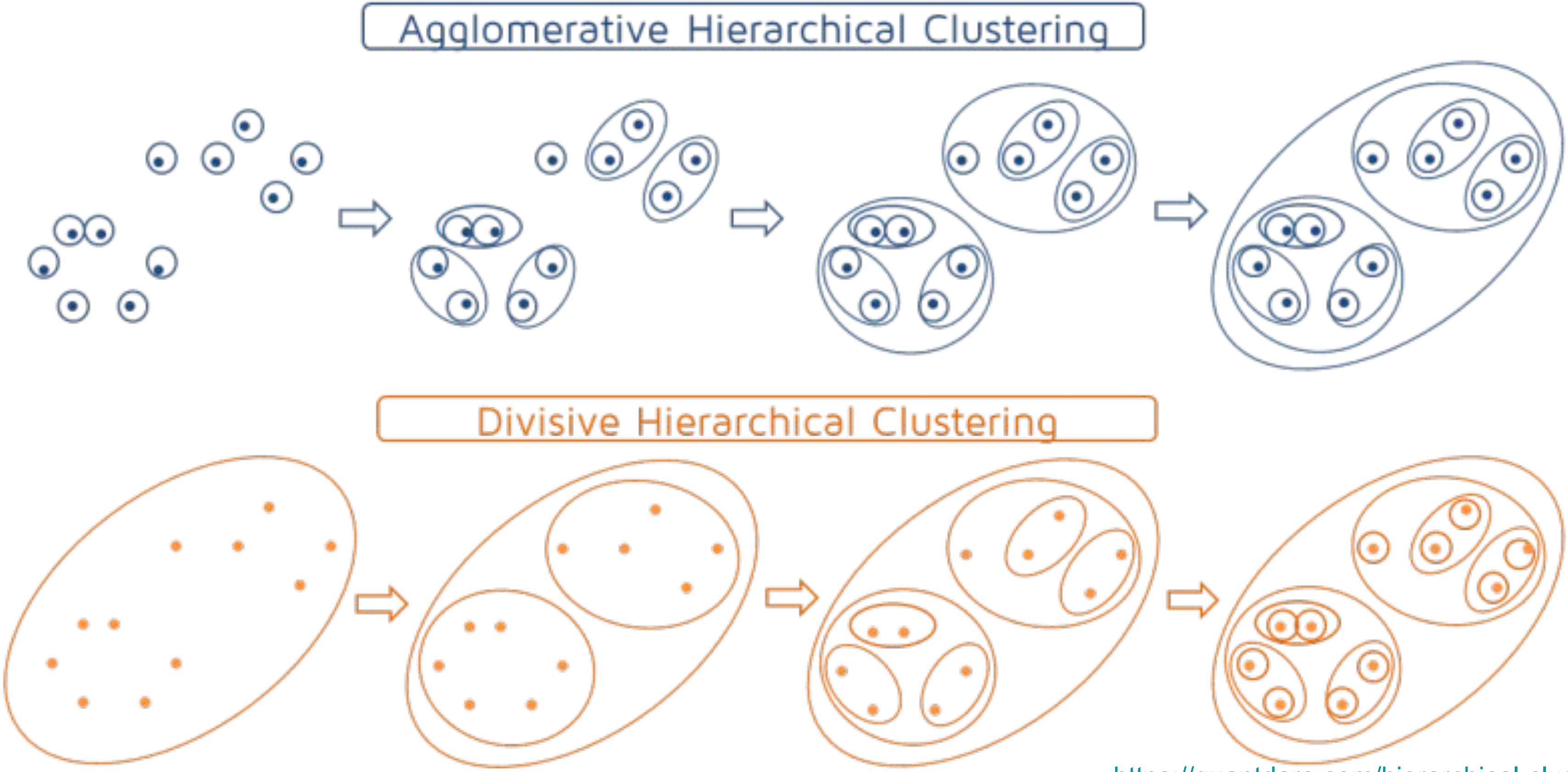
Why do we cluster?

For a flat list we run $O(N)$ comparisons to find kNN

For $\sqrt{N}$ similar* clusters we can pick one closest**
for $O(\sqrt{N})$ and find ***k*** NNs*** in $O(\sqrt{N})$.

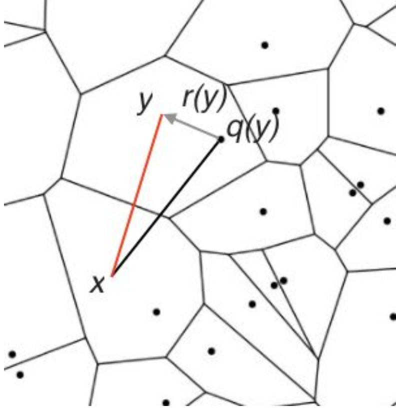
For two layers of $\sqrt[3]{N}$ …

How do we cluster?

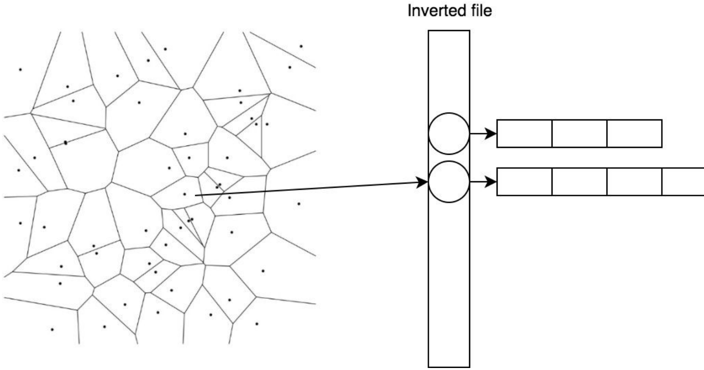


https://quantdare.com/hierarchical-clustering/

Revised IVF. FAISS (Facebook AI similarity search)

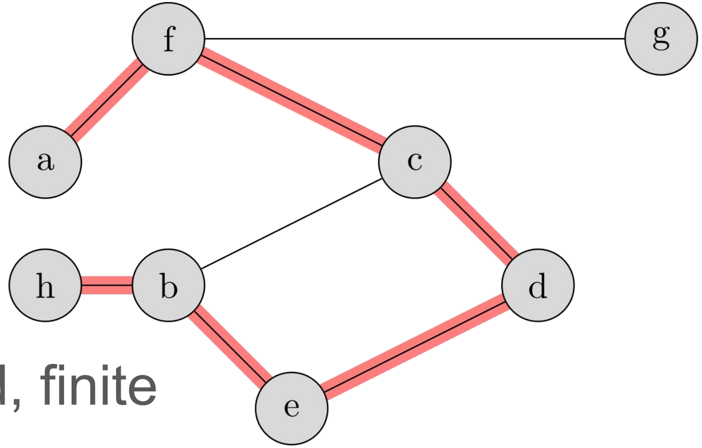


- Uses Voronoi diagram clusters. Vectors are approximated with **centroids** (ADC - asymmetric distance computation)
- Build **inverted index** for points in clusters
- Vector compression: product quantizer
  - Split R^{128} into 8 groups of 16 floats
  - Perform 256-means clustering of these “sub-vectors” and encode with 1 byte each



Approach #2. Proximity graphs

Graphs cheat sheet



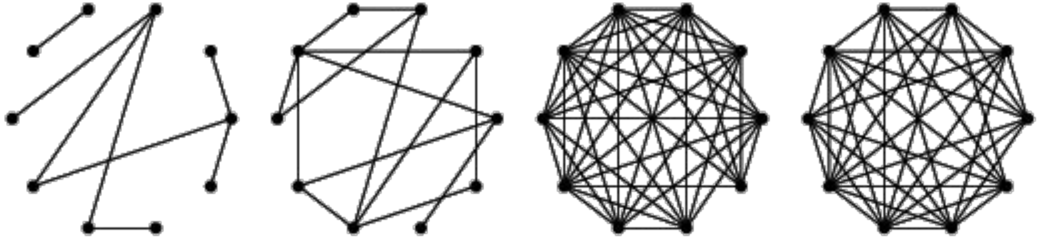
Graph - G = (V, E), can be weighted, directed, finite

[Simple] **path** - sequence of vertices and edges

Degree of vertex - number of incident edges

Graph diameter - longest shortest path between a pair of vertices

Random graph



Some random process (uniform, Gaussian, …) generates edges.

Almost every graph in the world. *Previously* considered as a model for social networks.

Small average shortest path - which is good for **search.**

Small clustering coefficient (defines how close are neighborhoods to cliques) - which is **bad** for **NN search.**

$$C(v) = \frac{e(v)}{deg(v)\,(deg(v) - 1)\,/2}$$

$$\widetilde{C} = \frac{1}{N}\sum_{i=1}^{N} C(i)$$

Regular graphs

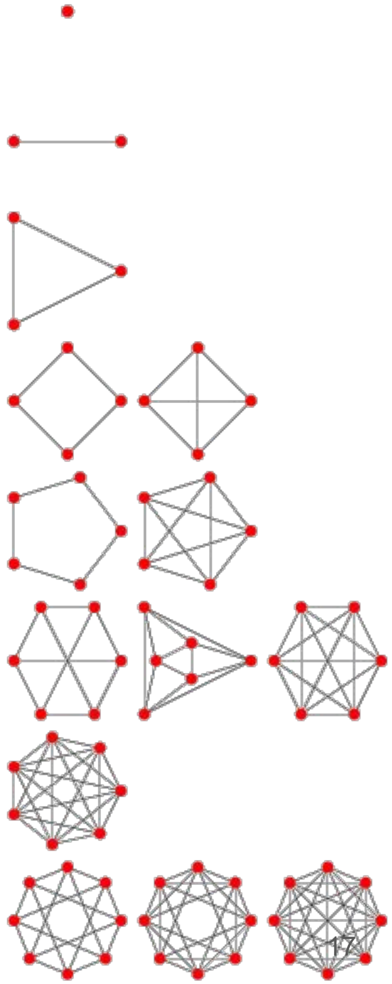
K-regular graph is a graph with **`deg(v) = K`** for any **v**.

Used to model big homogeneous networks.

Can also be random (as there are multiple K-regular graphs on the same size)

Big diameter - which is **bad** for **search**

Big clustering coefficient - which is **good** for **NN search**



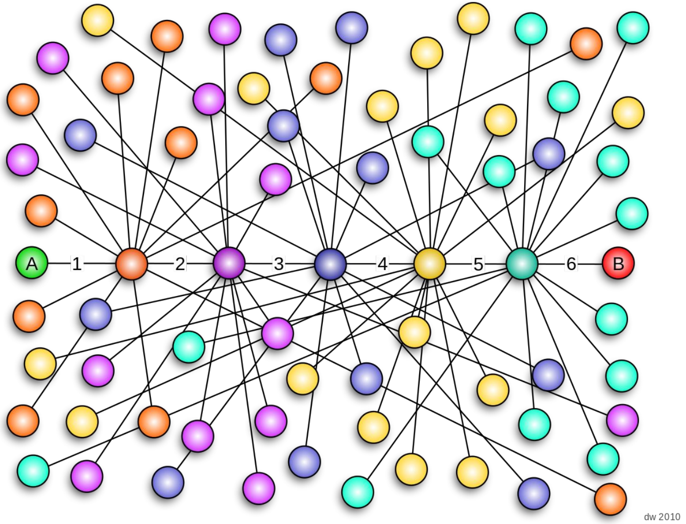
Small World experiment by Stanley Milgram, 1967

Initially it was considered, that social graph is kind of regular.

Experiment discovered (even with some questions to method) that even graph is **highly clustered**, **average path length is small.**

Was a basis for 6 handshakes rule.

New type of graphs was suggests: small world networks.



Small world network

Most vertices are not neighbours (small degree means *sparse* graph).

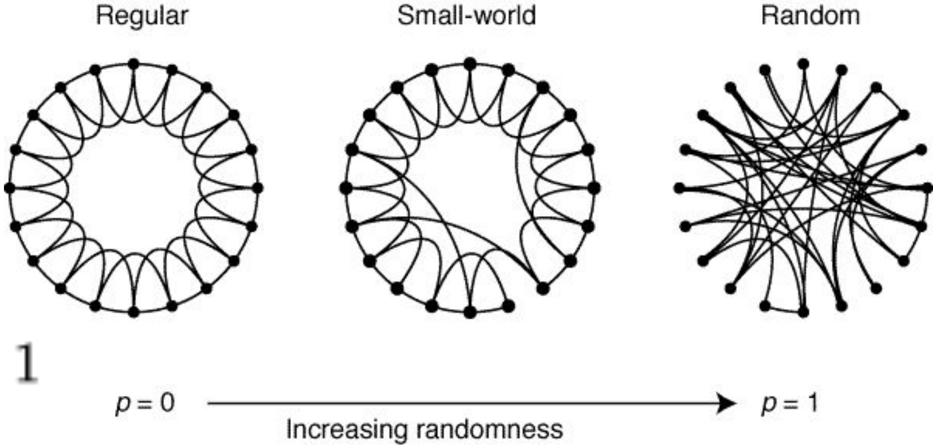
Nevertheless, small number of hops needed to reach any other node.

Typical path length ***L*** between 2 random nodes (of ***N***): $L \propto \log N$

Many real world networks are like this: internet, wiki, social graphs, power grids, brain cells. Although not all real networks like SW: many-generation networks, classmate graphs.

Watts–Strogatz model and **Kleinberg model** are how we describe and build SW networks

Watts–Strogatz model



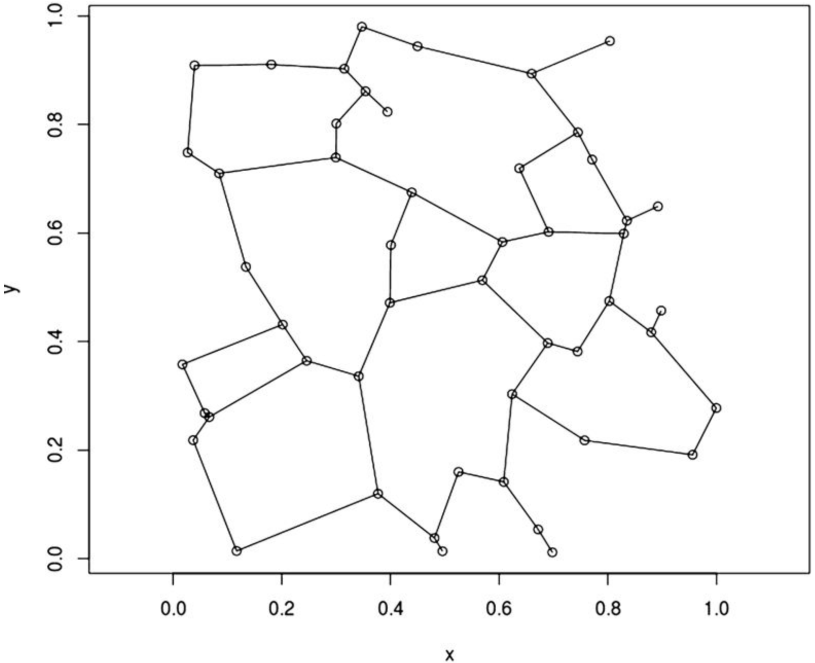
Given ***N*** nodes and ***K***-”regularity” (average degree ***K***) $N \gg K \gg \ln N \gg 1$

Given parameter ***p*** from [0, 1].

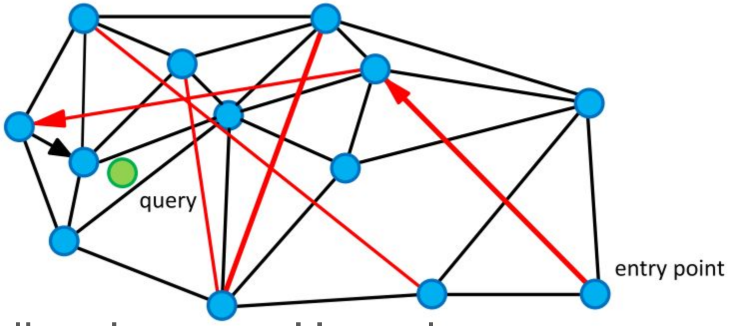
1) Construct a regular ring lattice.
2) take every edge connecting **vertex** to its **K/2 rightmost neighbors**, and rewire it with probability ***p***. Rewiring is done by **replacing destination** with vertex ***k*** (chosen **uniformly** at random from all possible nodes while avoiding self-loops and duplication).

Proximity graphs

A proximity graph is a simply a graph in which two vertices are connected by an edge if and only if the vertices satisfy particular geometric requirements.



[Navigable small world networks]()



Idea is similar to [skip-lists]().

We can also measure **distance** (e.g. dot product, Euclidian, L_k-norm, Humming, Levenstein, …) between *query* and *current vertex*. Originally *Delaunay graph* needed to converge for exact search, but ANNS allows other small-world graphs.

Building:

1. One-by-one insertion via kNN search. Distant edges are created in the beginning.

Search:

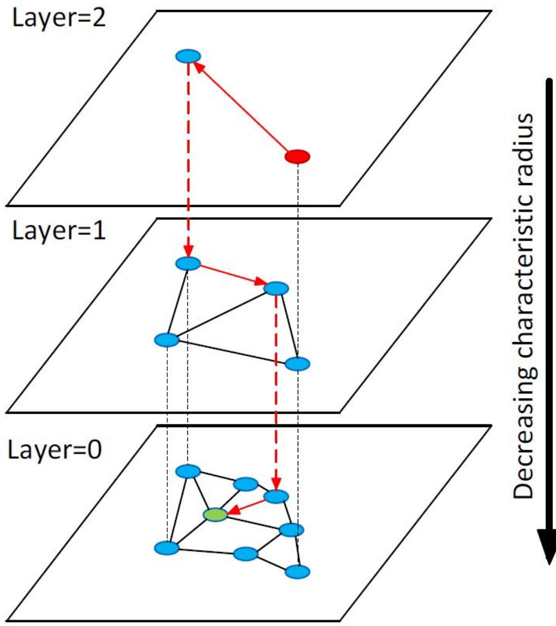
1. Perform greedy search. Move to the neighbour vertex **closest to query**
2. Update NN set on each step until it converges

Hierarchical navigable small world (github)

Layer 0 holds complete NSW network

Ideas:

- Better **start search** from a node with **high degree**
- Higher layer has longer links (skip-list!)
- Decrease layer size exponentially



Highlight:

1) **search procedure requires <u>only</u> `dist(u, v)` function**
2) **<u>No embedding</u>, hyperplanes, centroids of whatever needed**

To read

[An Introduction to Proximity Graphs]

[Efficient and robust approximate nearest neighbor search using Hierarchical Navigable Small World graphs]