

## CS201A/201: Math for CS I/Discrete Mathematics midsem

Max marks: 100

Time: 120 mins.

20-Sep-2014

1. *Answer all 4 questions. The paper has 2 pages. Hints are part of the question. Use them.*
2. **Please start each answer to a question on a fresh page. And keep answers of parts of a question together.**
3. *Just writing a number/final value will not get you credit. You must calculate/justify your answers.*
4. *You can consult **only your own handwritten notes and my notes that I put up on the course site.** Nothing else is allowed.*

1. (a) Term  $T_n$  is given by the following equation:

$$T_n = \frac{1}{2!} + \frac{2}{3!} + \dots + \frac{n}{(n+1)!}$$

First conjecture a simpler formula for  $T_n$  only in terms of  $n$  and then prove that it calculates  $T_n$  using induction.

### **Solution:**

We first calculate  $T_n$  for a few values of  $n$ .

$n$	$T_n$
1	$\frac{1}{2}$
2	$\frac{5}{6}$
3	$\frac{23}{24}$
4	$\frac{119}{120}$

We see that the numerator is 1 less than the denominator so a reasonable conjecture is:

$$T_n = \frac{(n+1)! - 1}{(n+1)!}.$$

We now prove this by induction.

It is clearly true for  $n = 1$ , (actually from the table also for  $n = 2, 3, 4$  as well).

Assume true for  $n$ . That is:  $T_n = \frac{1}{2!} + \frac{2}{3!} + \dots + \frac{n}{(n+1)!} = \frac{(n+1)! - 1}{(n+1)!}$ .

For  $(n + 1)$  we get:

$$\begin{aligned}
 T_{n+1} &= \frac{1}{2!} + \frac{2}{3!} + \dots + \frac{n}{(n+1)!} + \frac{n+1}{(n+2)!} \\
 &= \frac{(n+1)! - 1}{(n+1)!} + \frac{n+1}{(n+2)!} \\
 &= \frac{(n+2)! - (n+2) + (n+1)}{(n+2)!} \\
 &= \frac{(n+2)! - 1}{(n+2)!}
 \end{aligned}$$

- (b) Consider the sequence defined by  $a_n = a_{n-1} + a_{n-2} + a_{n-3}$  with  $a_1 = 1$ ,  $a_2 = 2$ ,  $a_3 = 3$ . i) Argue using induction that  $a_n < 2^n$  clearly stating the inductive hypothesis. ii) What is the difference between the inductive hypothesis in part (a) and the current hypothesis?

**Solution:**

The hypothesis is clearly true for  $n = 1, 2, 3$  and for  $n = 4$  we get  $a_4 = a_3 + a_2 + a_1 = 3 + 2 + 1 = 6 < 2^4$ .

Inductive hypothesis: Assume hypothesis is true for  $n$ ,  $(n - 1)$ ,  $(n - 2)$ . Then for  $n + 1$  we get:

$$\begin{aligned}
 a_{n+1} &= a_n + a_{n-1} + a_{n-2} \\
 &< 2^n + 2^{n-1} + 2^{n-2} \\
 &< 2^{n-2}(2^2 + 2 + 1) \\
 &< 2^{n-2}7 \\
 &< 2^{n-2}2^3 \\
 &< 2^{n+1}
 \end{aligned}$$

The inductive hypothesis is a version of the strong inductive hypothesis since it is assumed for  $n$ ,  $n - 1$ ,  $n - 2$  unlike the normal case when it is assumed to hold only for  $n$ .

- (c) i. Argue that the sum of a rational number and an irrational number is always irrational.

**Solution:**

Let  $r$  be the irrational number and  $q = \frac{m}{n}$  the rational number. Proof is by contradiction. Assume  $r + q$  is rational. So,  $r + q = \frac{a}{b}$ . Then we get  $r = \frac{a}{b} - \frac{m}{n} = \frac{an - bm}{bn}$ . Since  $a, b, m, n \in \mathbb{Z}$   $r = \frac{c}{d}$  where  $c, d \in \mathbb{Z}$ . This is a contradiction since  $r$  is irrational.

- ii. On the other hand, show that the sum of two irrational numbers can be rational.

**Solution:**

Consider  $x = \frac{1}{4} + \sqrt{2}$  and  $y = \frac{1}{4} - \sqrt{2}$ . Then by (i) both  $x, y$  are irrational but  $x + y = \frac{1}{2}$  a rational. On the other hand  $\sqrt{2} + \sqrt{2} = 2\sqrt{2}$  is irrational since if it was rational then  $\sqrt{2}$  would be rational. So, a sum of irrationals can be either rational or irrational.

- (d)  $a, b, c, d \in \mathbb{N}$ . Both  $a, b$  are divisible by both  $c, d$  (written,  $c|a, c|b$  and  $d|a, d|b$ ). Also,  $c \nmid d$  ( $c$  does not divide  $d$ ). Someone gives the following proof that  $cd$  divides both  $a$  and  $b$ .

*Proof.* Since  $d|a$  and  $d|b$  there exist  $x, y$  such that:

$$a = xd \text{ and } b = yd \quad (1)$$

. Also,  $c$  divides both  $a, b$  so we have  $c|dx$  and  $c|dy$ . But since  $c \nmid d$  it must be the case that  $c|x$  and  $c|y$  that is  $x = uc$  and  $y = vc$  for some  $u, v \in \mathbb{N}$ . Substituting for  $x$  and  $y$  in (1) we get  $a = ucd$  and  $b = vcd$ . So,  $cd$  divides both  $a, b$ .  $\square$

- i. Is the above proof correct? If not where and what precisely is the error?

**Solution:**

The proof is wrong. The error is in the claim: if  $c|xd$  and  $c \nmid d$  then  $c|x$ . This is easily seen to be false. Consider a case where  $c > d$  and  $c > x$ , for example  $c = 12, d = 3, x = 8$ . While  $c|xd$  and  $c \nmid d$  it is not the case that  $c|x$ .

- ii. Produce a counterexample for the result stated.

**Solution:**

A counter example can be constructed from the above observation: Let  $a = 16, b = 24, c = 8, d = 4$ . Then  $x = 4, y = 6$  and  $c \nmid x, c \nmid y$  though both  $c, d$  divide both  $a, b$  and  $c \nmid d$ .

$$[(2,3),(6,2),(4,3),(3,2)=25]$$

2. (a) Let  $S$  be a set. How many binary relations are possible if the cardinality of  $S$  is 3?

**Solution:**

Since  $|S| = 3, |S \times S| = 9$ . A relation is any subset of  $S \times S$ . So the number of relations is  $|\mathcal{P}(S \times S)| = 2^9 = 512$ .

- (b) Let  $f : S_1 \rightarrow S_2, g : S_2 \rightarrow S_3, h : S_1 \rightarrow S_3$  are functions and  $h = g \circ f$  that is  $h(x) = g(f(x)), x \in S_1$ .
- i. Argue that if  $h$  is surjective and  $f$  is total and injective, then  $g$  must be surjective.

**Solution:**

Since  $h$  is surjective any  $z \in S_3$  has one or more pre-image elements in  $S_1$ . Let these be  $x_1, \dots, x_m$ , i.e.  $h(x_i) = z$ ,  $i = 1..m$ . Since  $h(x) = g(f(x))$  and  $f$  is total and injective  $y_i = f(x_i)$ ,  $i = 1..m$  and  $y_i \neq y_j$  when  $i \neq j$ . Clearly, one of these  $y_i$ s must be such that  $g(y_i) = z$  due to the definition of  $h$  as  $h = g(f(x))$  and therefore  $g$  is a surjection since  $z$  is any element in  $S_3$ .

- ii. Let  $h$  be injective and  $g$  total. Show that  $f$  must be injective. Give a counterexample to show how this claim can fail if  $g$  is not total.

**Solution:**

For  $h$  to be injective distinct elements in  $S_1$  must map to distinct elements in  $S_3$ . This is impossible if distinct elements in  $S_1$  map to the same element in  $S_2$  since  $h(x) = g(f(x))$  so distinct elements in  $S_2$  map to the same element in  $S_3$  and  $h$  fails to be injective. So,  $f$  must be injective.

It is also clear that if  $g$  is not total  $f$  can map just two elements, say  $x_1, x_2 \in S_1$ , to the same element  $y \in S_2$  and all others to distinct elements of  $S_2$  and let  $g$  be partial with no image for  $y$  and mapping the rest to distinct elements of  $S_3$ . Then  $h$  is injective while  $f$  is not. So, for  $f$  to be injective it is necessary that  $g$  is total.

- (c) Argue that every subset  $S$  of disjoint open intervals in  $\mathbb{R}$  is countable. (*Hint:  $\mathbb{Q}$* )

**Solution:**

Results for infinite sets are generally counter intuitive. It would seem that since  $\mathbb{R}$  is uncountable, all kinds of subsets of it would also be uncountable but that is not true. Since  $S$  is a set of disjoint open intervals so every element of  $S$  is of the form  $(a, b)$ ,  $a < b$ . Every such interval will contain some rationals (actually infinitely many of them). Choose any one of them and associate it with the interval. This allows us to define a bijection  $f : S \rightarrow P \subset \mathbb{Q}$ . Since  $\mathbb{Q}$  is countable any subset of  $\mathbb{Q}$  is also countable. So,  $S$  is countable.

Notice that since the intervals are disjoint the rational points in  $P$  are distinct and can never overlap.

- (d) Show that

$$f(x) = \frac{2x - 1}{2x(1 - x)}$$

is a bijection between  $f : (0, 1) \rightleftarrows \mathbb{R}$  by arguing that  $f$  is an injection and a surjection. (*Note: Cannot use calculus.*) (*Hint: For injection show  $f$  is monotonically increasing and for surjection that every  $y \in \mathbb{R}$  has a pre-image.*)

**Solution:**

There are multiple ways to show a bijection between  $(0, 1)$  and  $\mathbb{R}$ . The common one is to use some form of the *tan* function. Here we are able to do the same with a function that is a ratio of two simple polynomials.

To show that  $f$  is injective it is enough to show that  $f$  is a monotonically increasing function. Then distinct values from  $(0, 1)$  will map to distinct values of  $\mathbb{R}$ . We can break up  $f$  into two fractions (partial fractions method):  $f = \frac{1}{1-x} + \left(\frac{-1}{2x}\right)$ .

First note that  $f(x) \rightarrow -\infty$  as  $x \rightarrow 0$  and  $f(x) \rightarrow \infty$  as  $x \rightarrow 1$ . The fraction  $\frac{1}{1-x}$  has its least value when  $x = 0$  and increases monotonically as  $x \rightarrow 1$  since denominator decreases monotonically. Similarly, the fraction  $\frac{-1}{2x}$  grows from larger negative values to smaller negative values as  $x$  goes from 0 to 1 - that is it is also a monotonically increasing function. The sum of two monotonically increasing functions is monotonically increasing so  $f$  is an injection.

To show that  $f$  is also a surjection let us calculate the pre-image  $x$  for some  $f(x) = y$ . So, we get  $y = \frac{2x-1}{2x(1-x)}$ . We solve for  $x$  in terms of  $y$ . Simplifying we get  $2yx^2 + (2 - 2y)x - 1 = 0$  a quadratic in  $x$ . Solving the quadratic gives:

$$x = \frac{-(2 - 2y) \pm \sqrt{(2 - 2y)^2 + 8y}}{4y}$$

Simplifying the above we get:

$$x = \frac{-1 + y \pm \sqrt{1 + y^2}}{2y}$$

We have to a) determine whether to take the positive or negative sign for the square root and b) since  $y$  is in the denominator we have to see if we can get a better form for the equation. Note that because  $y \in \mathbb{R}$  it can be 0 which is not admissible in the formula above.

For a) let us put  $y = 1$  and we see that choosing the negative sign gives  $\frac{-1}{\sqrt{2}}$  which is outside  $(0, 1)$ . Choosing the positive sign gives a value  $\frac{1}{\sqrt{2}}$  which is within the interval. So, we get:

$$x = \frac{-1 + y + \sqrt{1 + y^2}}{2y}$$

For b) to get rid of  $y$  in the denominator let us multiply both numerator and denom-

inator by the conjugate:

$$\begin{aligned}
 x &= \left( \frac{-1 + y + \sqrt{1 + y^2}}{2y} \right) \left( \frac{-1 + y - \sqrt{1 + y^2}}{-1 + y - \sqrt{1 + y^2}} \right) \\
 &= \frac{1 - 2y + y^2 - 1 - y^2}{2y(-1 + y - \sqrt{1 + y^2})} \\
 &= \frac{-2y}{2y(-1 + y - \sqrt{1 + y^2})} \\
 &= \frac{1}{\sqrt{1 + y^2} - y + 1}
 \end{aligned}$$

For any  $y$  note that  $\sqrt{1 + y^2} - y > 0$  and  $\sqrt{1 + y^2} - y + 1 > 1$  so the value of  $x$  is always between 0 and 1 for any value of  $y$  and this  $x$  is our pre-image. So,  $f$  is a surjection. Since we have already shown it is an injection it is, therefore, a bijection.

[3,(4,4),5,(4,5)=25]

3. (a) Five distinct boys and four distinct girls have to line up in a row such that no two girls are next to each other. How many such arrangements are possible?

**Solution:**

Imagine we arrange the 5 boys as below where E stands for empty.

E B1 E B2 E B3 E B4 E B5 E

The girls can occupy any of 6 E positions so we have  ${}^6P_4$  arrangements of girls for each arrangement of boys. So, the total number is:  ${}^5P_5 \cdot {}^6P_4 = 120 \times 360 = 43200$ .

- (b) Given a standard deck of cards (52 cards, 4 suits - ♡, ♣, ◇, ♠) what is the minimum number of cards that must be drawn to ensure:
- Two cards each from any two suits in the selection.
  - Three ♠ cards in the selection.

**Solution:**

i. Simple pigeonhole application. Worst case: one card of each suit (i.e. 4). Fifth card will give us 2 cards of one suit. Now worst case is we get 11 more cards of the same suit. Giving us a total of 16 cards. The 17th card will make at least 2 cards from 2 suits. So we must select 17 cards.

Another way to count: 13 cards of one suit followed by one card from each of the other suits gives us 16 cards. The 17th card will ensure at least 2 cards from 2 suits.

ii. You may not draw a spade for the first 39 cards after which the next 3 cards will be ♠. So a selection of 42 cards is enough to ensure you have at least 3 ♠. Note that

it does not matter if we draw 1 or 2 ♠ cards in between. The minimum cards needed to ensure 3 ♠ cards remains the same.

- (c) Consider the ordered sequence  $(x_1, x_2, \dots, x_m)$  of positive integers (i.e.  $x_i \in \mathbb{N}$ ,  $i = 1..m$ ) that satisfy the equation  $x_1 + x_2 + \dots + x_m = n$ . How many such sequences are possible? (*Hint: Bijection principle.*)

**Solution:**

Let us write  $n$  as a sequence of  $n$  1s with an empty space between them shown by  $e$  and one 0 each at the two ends. So, this is how it will look:

01e1e1e1e...1e10

Now imagine distributing  $m - 1$  0s in the  $n - 1$  empty places marked by  $e$  and adding the 1s between any two 0s. This will give us  $m$  positive integers that add up to  $n$ . So, all possible ways of distributing  $m - 1$  0s in the  $n - 1$  slots marked by  $e$  will give the total number of sequences that are possible. This number is  ${}^{n-1}C_{m-1}$ . Note that there is a bijection between the binary string and the sequence. Since 0 cannot occur consecutively all  $x_i$  are positive.

- (d) Using the *principle of inclusion and exclusion* calculate the number of integers from 1 to 100 that are not divisible by 2, 3, or 5.

**Solution:**

Let  $S = \{1, 2, \dots, 100\}$ ,  $S_i$  be the set of numbers in  $S$  divisible by  $i$ ,  $S_{i,j}$  be the set of numbers for  $S$  that are divisible by  $i$  and by  $j$ , similarly,  $S_{i,j,k}$ . Then by PIE

$$S_{none} = |S| - (|S_2| + |S_3| + |S_5|) + (|S_{2,3}| + |S_{2,5}| + |S_{3,5}|) - |S_{2,3,5}|$$

Then by simple counting  $|S_2| = 50$ ,  $|S_3| = 33$ ,  $|S_5| = 20$ . For  $|S_{i,j}|$  we have to count how many are divisible by  $i \times j$ . So,  $|S_{2,3}| = 16$ ,  $|S_{2,5}| = 10$ ,  $|S_{3,5}| = 6$ ,  $|S_{2,3,5}| = 3$ . So,

$$S_{none} = 100 - (50 + 33 + 20) + (16 + 10 + 6) - 3 = 26.$$

[3,(3,3),10,6=25]

4. (a) Solve the recurrence relation:  $a_n = 9a_{n-1} - 26a_{n-2} + 24a_{n-3}$  for  $n > 2$  with  $a_0 = 0$ ,  $a_1 = 1$ ,  $a_2 = 10$  to obtain an expression for  $a_n$ .

**Solution:**

This requires a systematic application of the rules to solve such recurrences. Let

$A(s) = \sum_{n=0}^{\infty} a_n x^n$ . Multiply recurrence by  $x^n$  and sum for  $n > 2$ :

$$\sum_{n=3}^{\infty} a_n x^n - 9 \sum_{n=3}^{\infty} a_{n-1} x^n + 26 \sum_{n=3}^{\infty} a_{n-2} x^n - 24 \sum_{n=3}^{\infty} a_{n-3} x^n = 0.$$

Write in terms of  $A(x)$ :

$$(A(x) - a_0 - a_1 x - a_2 x^2) - 9x(A(x) - a_0 - a_1 x) + 26x^2(A(x) - a_0) - 24x^3 A(x).$$

Simplify:

$$\begin{aligned} A(x)(1 - 9x + 26x^2 + 24x^3) &= a_0 + a_1 x + a_2 x^2 - 9a_0 x - 9a_1 x^2 + 26a_0 x^2 \\ A(x) &= \frac{a_0 + (a_1 - 9a_0)x + (a_2 - 9a_1 + 26a_0)x^2}{1 - 9x + 26x^2 + 24x^3} \end{aligned} \quad (1)$$

Break into partial fractions:

$$\begin{aligned} A(x) &= \frac{c_1}{1 - 2x} + \frac{c_2}{1 - 3x} + \frac{c_3}{1 - 4x} \\ &= \sum_{n=0}^{\infty} c_1 2^n x^n + \sum_{n=0}^{\infty} c_2 3^n x^n + \sum_{n=0}^{\infty} c_3 4^n x^n \\ &= \sum_{n=0}^{\infty} (c_1 2^n + c_2 3^n + c_3 4^n) x^n \end{aligned}$$

So,  $a_n = c_1 2^n + c_2 3^n + c_3 4^n$ . To calculate  $c_1, c_2, c_3$  use the given initial conditions. Substituting  $a_0 = 0$ ,  $a_1 = 1$  and  $a_2 = 10$  in (1) above:

$$A(x) = \frac{x + x^2}{(1 - 2x)(1 - 3x)(1 - 4x)} = \frac{c_1}{1 - 2x} + \frac{c_2}{1 - 3x} + \frac{c_3}{1 - 4x}$$

Comparing coefficients the simultaneous equations we get are:

$$\begin{aligned} c_1 + c_2 + c_3 &= 0 \\ 7c_1 + 6c_2 + 5c_3 &= -1 \\ 12c_1 + 8c_2 + 6c_3 &= 1 \end{aligned}$$

Solving the above:  $c_1 = \frac{3}{2}$ ,  $c_2 = -4$ ,  $c_3 = \frac{5}{2}$ .

- (b) The Bubble sort algorithm sorts a sequence  $S$  of length  $n$  by making  $(n - 1)$  bubble passes over the sequence. In a single bubble pass the algorithm compares  $S(i), S(i + 1)$  and switches the elements if  $S(i) > S(i + 1)$  for  $n = 1..(n - 1)$ . Note  $S(i)$  stands for the  $i^{th}$  element in the sequence.

Derive a recurrence for the number of comparisons in the Bubble sort algorithm and solve the recurrence.



**Solution:**

Let  $T(n)$  be the number of comparisons required to Bubble sort a sequence of size  $n$ . The difference between the number of comparisons for  $T(n-1)$  and  $T(n)$  is one bubble pass which can take  $(n-1)$  comparisons in the worst case. So, the recurrence is:  $T(n) = T(n-1) + n - 1$ , with  $T(1) = 0$ . This is most easily solved by unfolding and substitution:

$$\begin{aligned}
 T(n) &= T(n-1) + n - 1 \\
 &= T(n-2) + n - 2 + n - 1 = T(n-2) + 2n - (1 + 2) \\
 &\dots \\
 &= T(1) + (n-1)n - (1 + 2 + \dots + (n-1)) \\
 &= 0 + n(n-1) - \frac{n(n-1)}{2} \\
 &= \frac{n(n-1)}{2}
 \end{aligned}$$

So, solution is  $\frac{n(n-1)}{2}$

- (c) Many algorithms have the following recurrence structure:  $a(n) = d.a(n/d) + e$  where  $d, e \in \mathbb{N}$ ,  $d > 1$  and  $e > 0$ . Assuming  $n = d^k$ ,  $k \in \mathbb{N}$  what is the solution to the recurrence.

**Solution:**

We can solve this by unfolding and substitution.

$$\begin{aligned}
 a(n) &= d.a\left(\frac{n}{d}\right) + e \\
 &= d^2.a\left(\frac{n}{d^2}\right) + 2e \\
 &= d^3.a\left(\frac{n}{d^3}\right) + 3e \\
 &\dots \\
 &= d^{\log_d(n)}a(1) + \log_d(n).e \\
 &= a(1)n + \log_d(n)e
 \end{aligned}$$

[10,(4,4),7=25]