# Theory of Computation (1) Coding Club

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#### Introduction

- ▶ Theory of Computation? Study of theoretical models of computation.
- ▶ Format? You'll see the following today:
  - ▶ Mathematical definitions with examples
  - ▶ Applications to compilers
- ▶ Why? Modelling behavior and processing strings are major applications of this area.
- ► These Slides? See our GitHub repository. (More topics and applications also available.)
- ▶ If you don't know coding: Check document on GitHub, or find a guide online.

Any questions?



#### Outline

- 1 Basic Definitions
- 2 Deterministic Finite Automata
- Regular Languages
  - Regular Expressions
- 4 Applications to Compilers
  - Flex



# Alphabets and Strings

## Definition (Alphabet)

An alphabet  $\Sigma$  is any nonempty finite set. The members of the alphabet are said to be the symbols of the alphabet.

## Definition (Strings)

A string w over  $\Sigma$  is a finite sequence of symbols from  $\Sigma$ .

 $\epsilon$  is the empty string.

wv is the concatenation of two strings w, v.

 $w^n$  is concatenation of n copies of w. Also,  $w^0 = \epsilon$ .



# Languages

## Definition (Languages)

Language: set of strings over an alphabet.

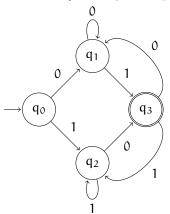
Empty language ∅ contains no strings.

Kleene closure  $\Sigma^*$  of an alphabet  $\Sigma$  is the set of all strings over  $\Sigma$ .



#### Deterministic Finite Automata

Start at  $q_0$  and follow the arrows according to an input string of 0's and 1's.  $q_3$  is a final state. If input ends at  $q_3$ , string is accepted.





#### Deterministic Finite Automata

DFAs: Easy to implement! That one took 30 lines.

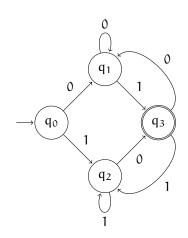
## Definition (Deterministic Finite Automaton)

A DFA has five components:

- 1 Q is a set of states
- $q_0 \in Q$  is the initial state
- $\mathbf{F} \subseteq \mathbf{Q}$  is the set of final states
- $\Sigma$  is the input alphabet
- $\delta$  is the transition function



# Example



- **2** The initial state is  $q_0$ .
- The set of final states is  $\{q_3\}$ .
- **4** The input alphabet  $\Sigma$  is  $\{0, 1\}$ .
- $\delta$  represented by this table:

0	1
q <sub>1</sub>	$q_2$
q <sub>1</sub>	$q_3$
$q_3$	q <sub>2</sub>
q <sub>1</sub>	q <sub>2</sub>
	q <sub>1</sub> q <sub>1</sub> q <sub>3</sub>



# Operations on Languages

#### Definition

Given two languages A and B:

- ▶ Union:  $A \cup B$  contains all strings in A and all strings in B.
- ▶ Concatenation: AB contains all strings ab, where  $a \in A$  and  $b \in B$ .
  - $ightharpoonup A^n$  is concatenation of A with itself n times.
- ► Closure:  $A^*$  is union of  $A^0$ ,  $A^1$ ,  $A^2$ , ...



# Regular Languages

- ▶ Regular languages and DFAs are equivalent!
  - ▶ Any regular language has some DFA that recognizes it.
  - ▶ Any language recognized by a DFA is regular.
- ▶ But what are regular languages?



# Regular Languages

## Definition (Regular Language over $\Sigma$ )

Recursive definition.

- ▶ ∅ is regular
- ▶ For every  $a \in \Sigma$ ,  $\{a\}$  is regular
- ▶ If A is regular, A\* is regular
- ▶ If A, B are regular,  $A \cup B$  and AB are regular
- $\blacktriangleright$  No other languages over  $\Sigma$  are regular



## Regular Expressions

Notation for expressing regular languages.

## Example (Regular Expressions)

 $\alpha b + \alpha^*$  is union of  $\alpha b$  and  $\alpha^*$ 

ab is concatenation of a and b

 $\alpha^*$  is closure of  $\alpha$ 



# Applications to Compilers

- ▶ Traditional compiler design:
  - 1 Lexical Analysis (or Tokenization)
  - 2 Syntax Analysis
  - 3 Semantic Analysis
  - 4 Code Generation
- ▶ Regular languages help with tokenization.
- ▶ Common tool is Flex. Input is regular expression, output is tokenizer.



#### Demo: Calculator

We'll make a basic calculator "language" that has:

- ▶ let keyword to declare variables
- ▶ Identifiers for the variables (letters only)
- ▶ Numbers, floating-point only
- ▶ +, -, \*, / for arithmetic
- ▶ ( and ) for grouping
- ► = for assignment

Paradigm: Input a statement and press enter. If it can be evaluated (i.e. not declaration or assignment), result will be printed.



## Flex Regex Syntax

- ▶ ab matches a followed by b
- ▶ a|b matches "either a or b"
  - ▶ [ab] is equivalent
  - ▶ [a-z] is shorthand for a|b|...|z
  - ► [A-Z] is A|B|...|Z
  - ▶ [0-9] is 0|1|...|9
- ▶ a\* matches "zero or more consecutive occurrences of a"
- ▶ a+ matches "one or more consecutive occurrences of a"
- matches any character



#### That's All!

Most of this was based on the following:

- ▶ Compilers: Principles, Techniques, and Tools (Aho et al.) Chapter 3
- ▶ Introduction to the Theory of Computation (Sipser) Chapters 0, 1

