

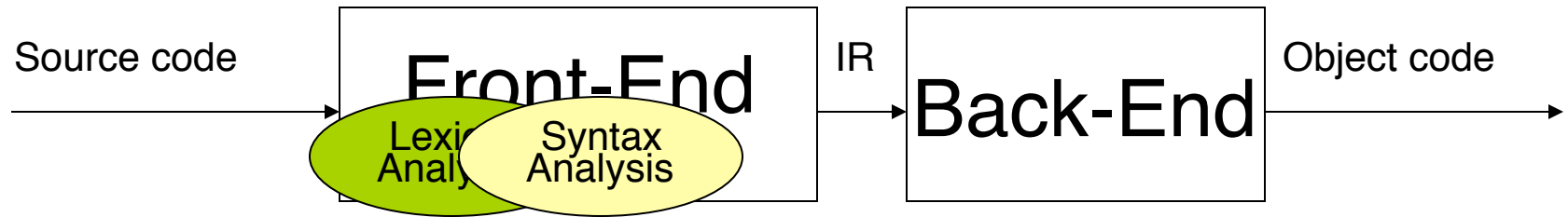
Compiler Design

Lecture 6: Syntax Analysis Bottom-Up Parsing (part IV)

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based on the slides of the course book

Parsing (Syntax Analysis)



■ Syntax Analysis:

- Derivation and parse trees ✓
- Top-down parsing ✓
- **Bottom-up parsing**

Outline

- Introduction
- Shift-reduce Parsing
- **LR Parsing**
 - LR(0)
 - LR(1)
 - **SLR**
 - LALR
- Ambiguity
- Error-handling

SLR(1)

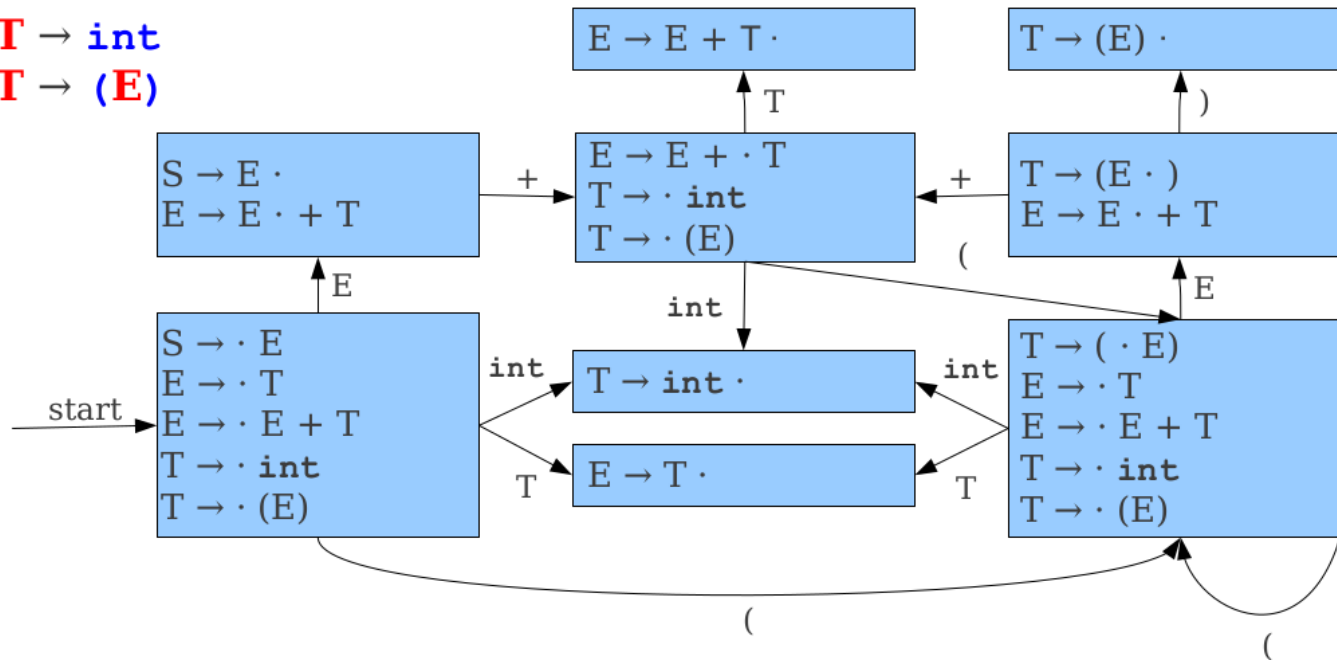
- Simple LR(1)
- Minor modification to LR(0) automaton that uses lookahead to avoid shift/reduce conflicts.
- Idea: Only reduce $A \rightarrow \omega$ if the next token t is in $\text{FOLLOW}(A)$.
- Automaton identical to LR(0) automaton; only change is when we choose to reduce.

SLR(1) Parsing

S → **E**
E → **T**
E → **E** + **T**
T → **int**
T → (**E**)

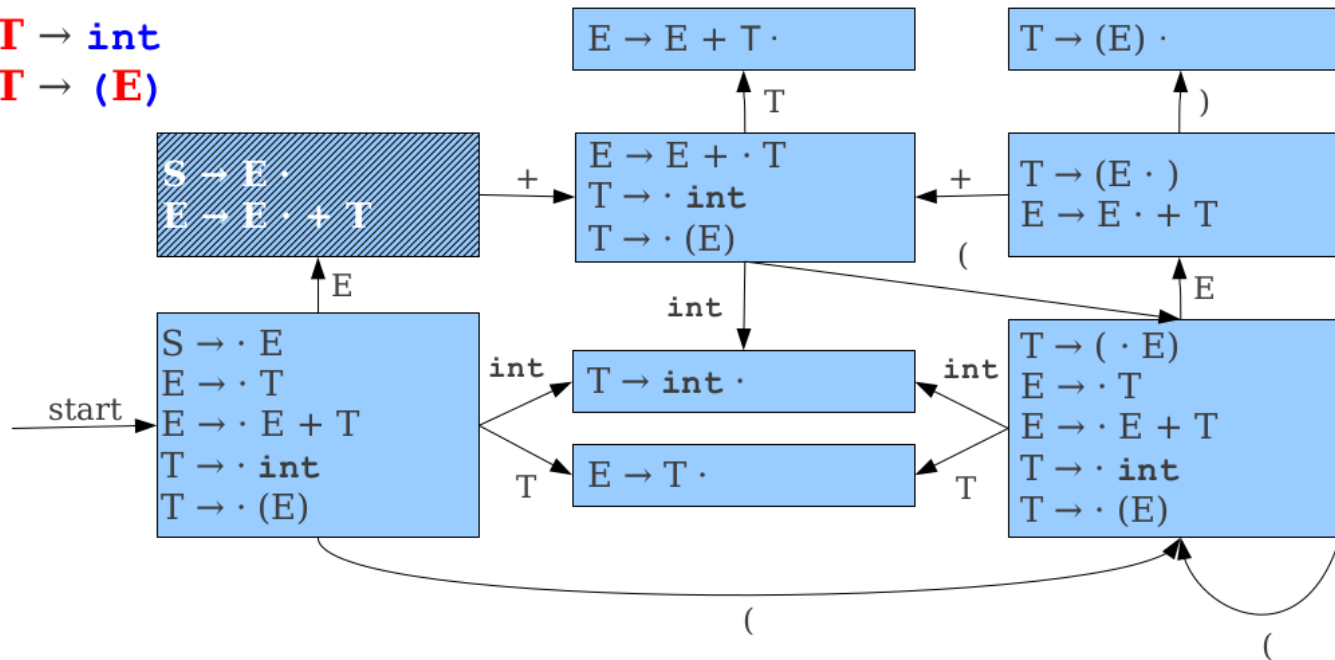
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$S \rightarrow E$
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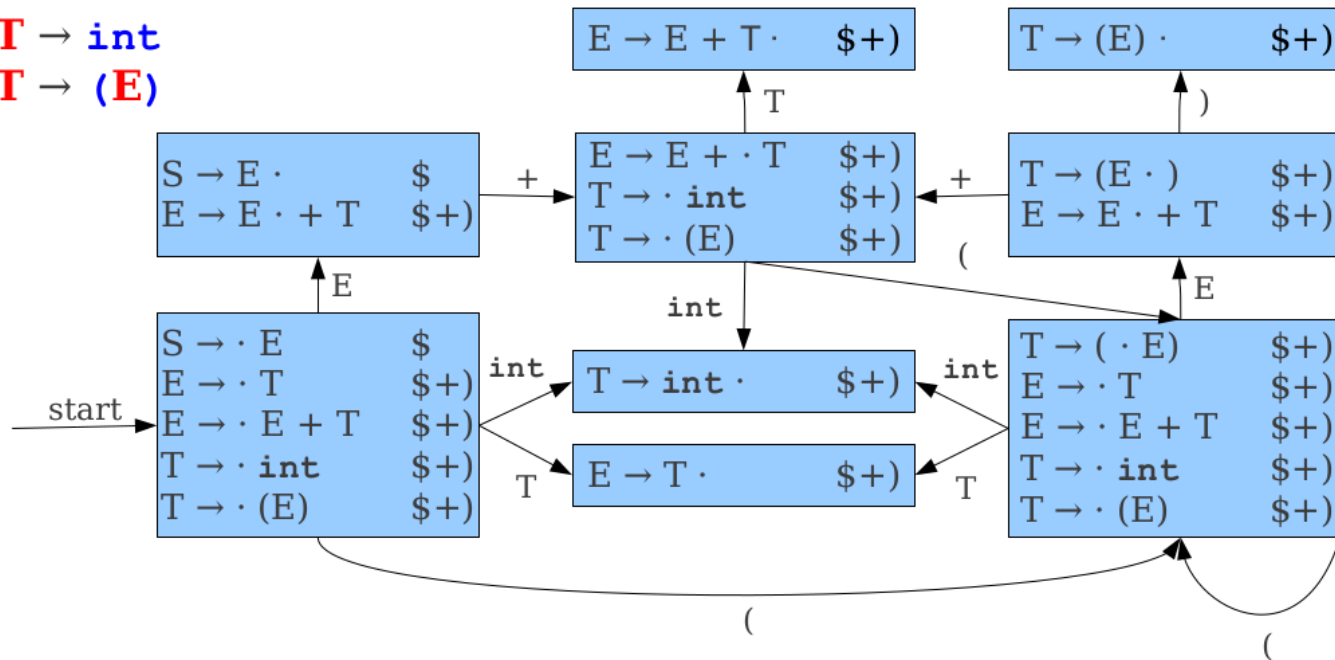
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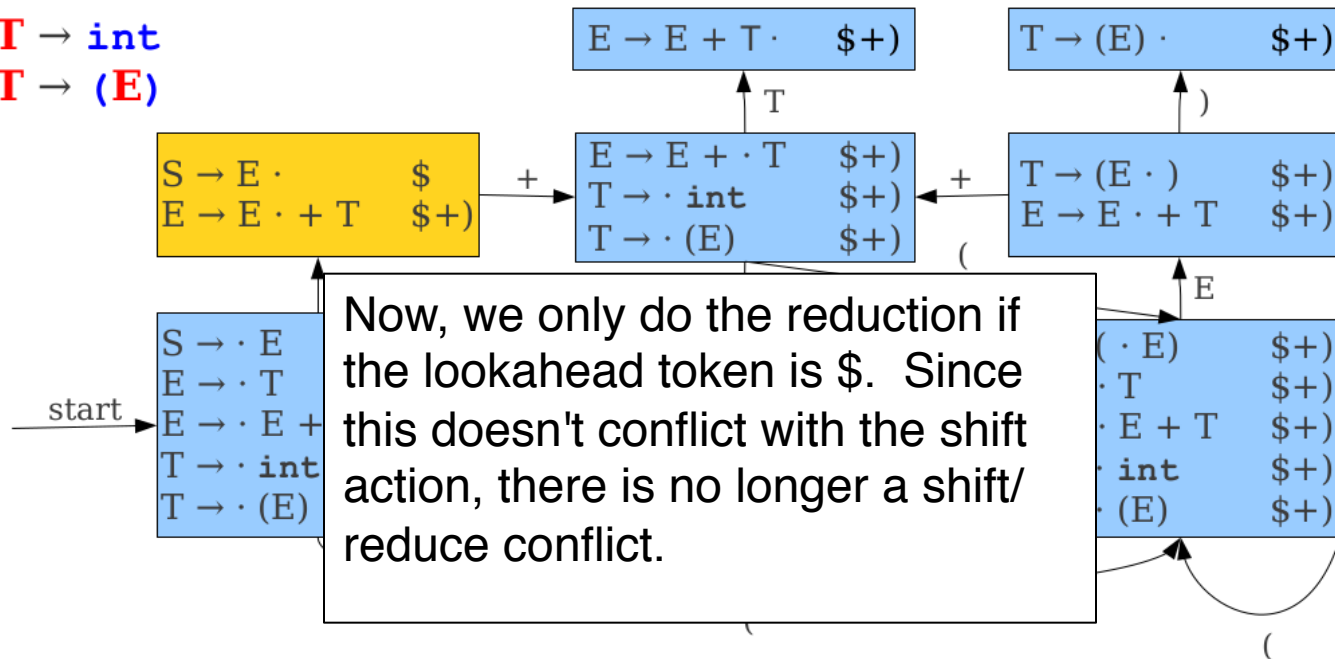
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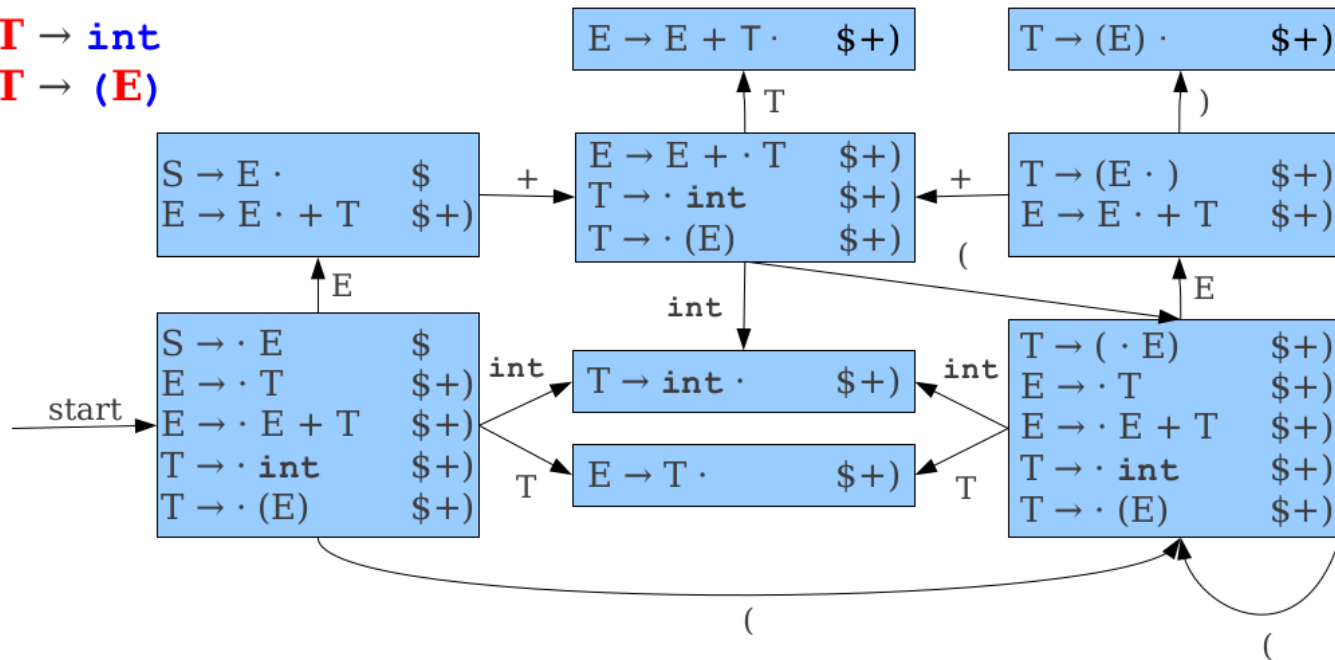
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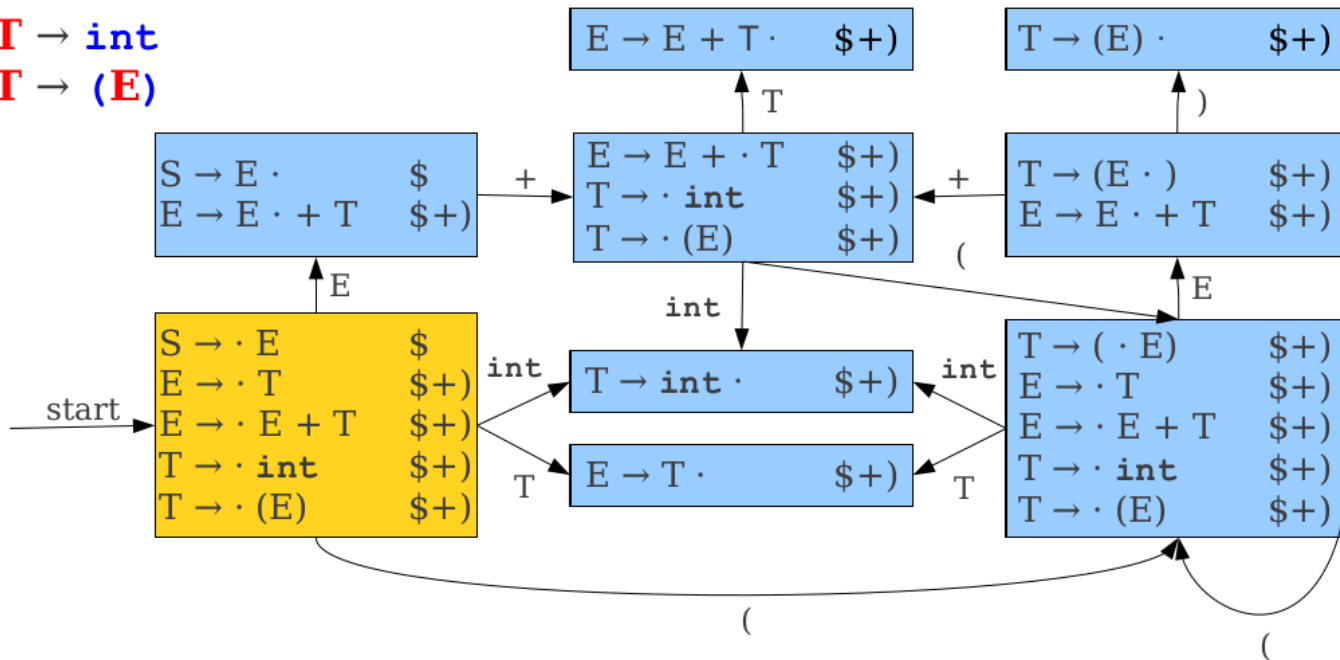
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int	+	(int	+	int	+	int)	\$
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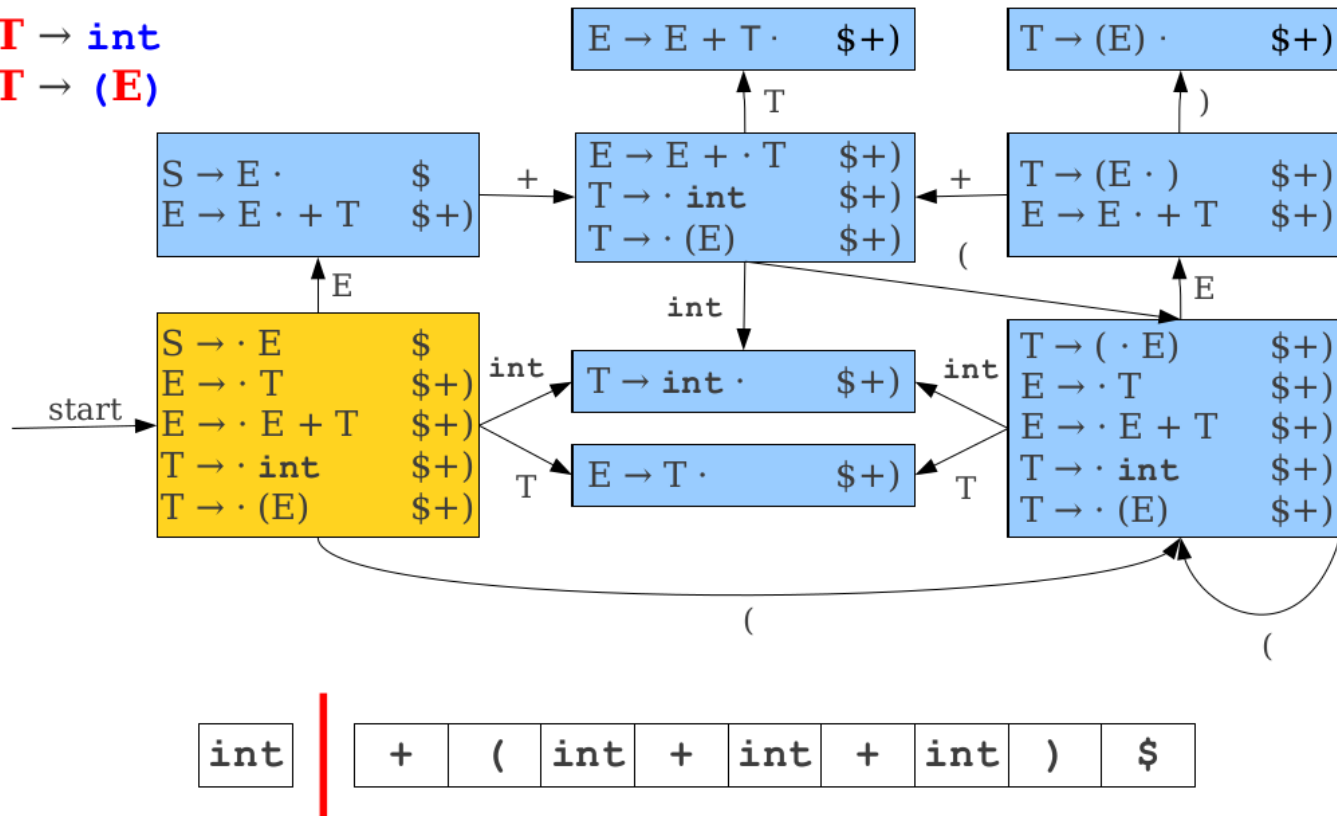
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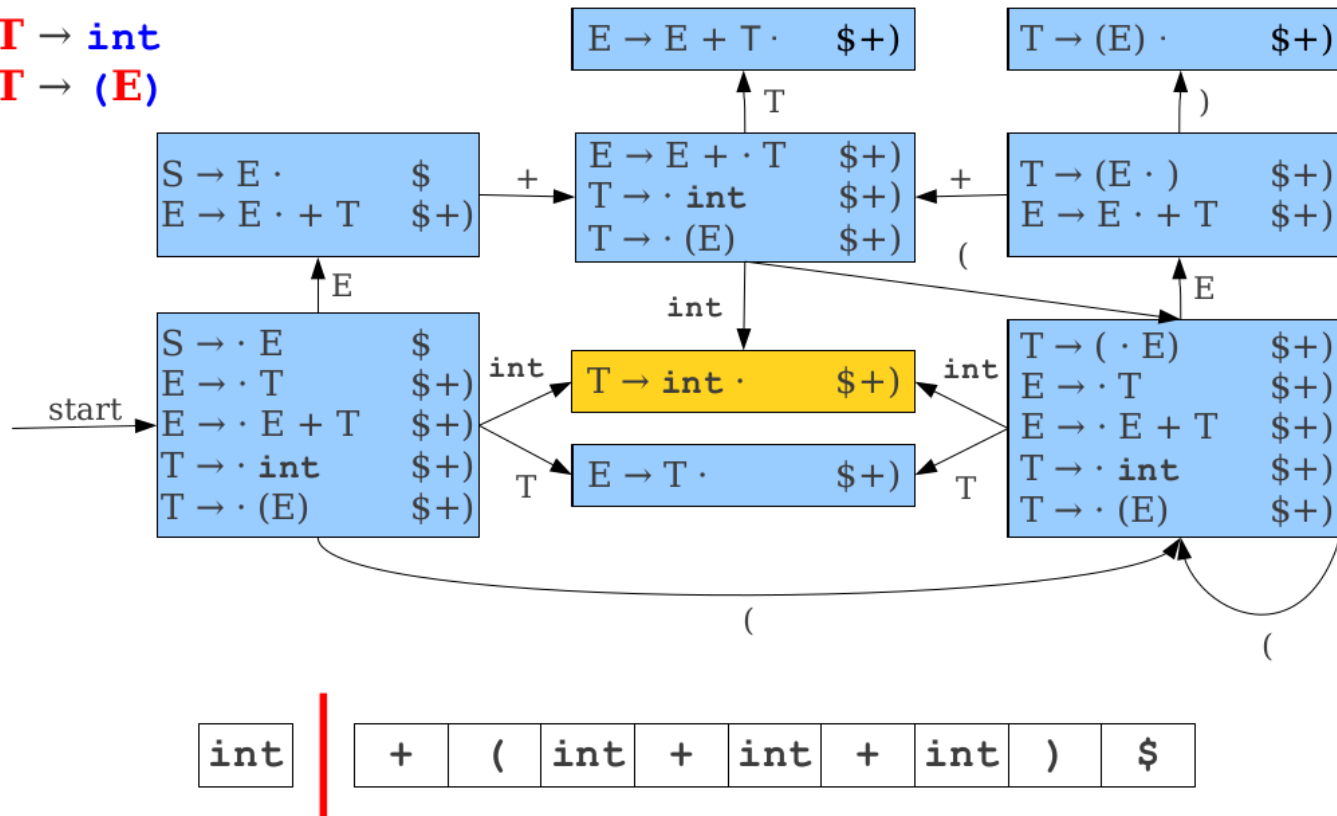
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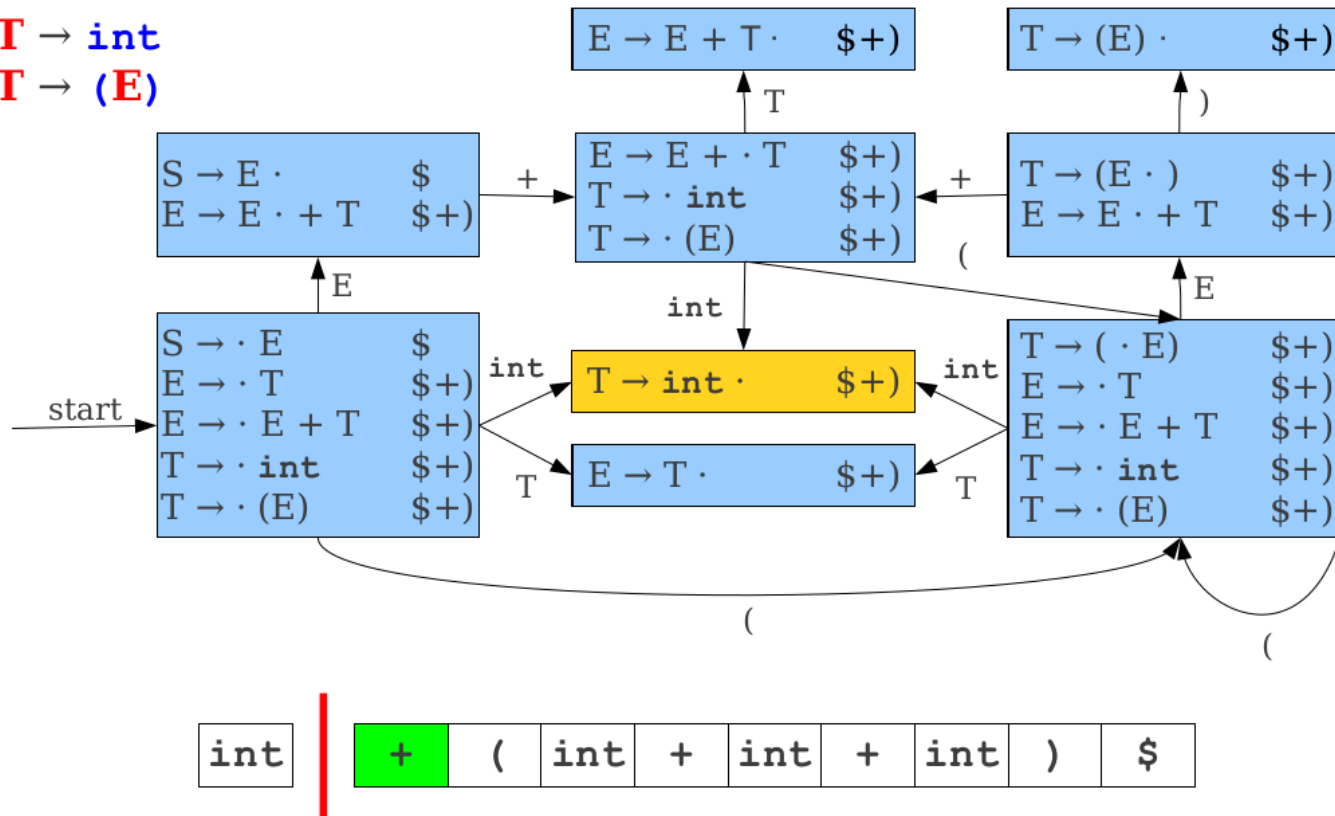
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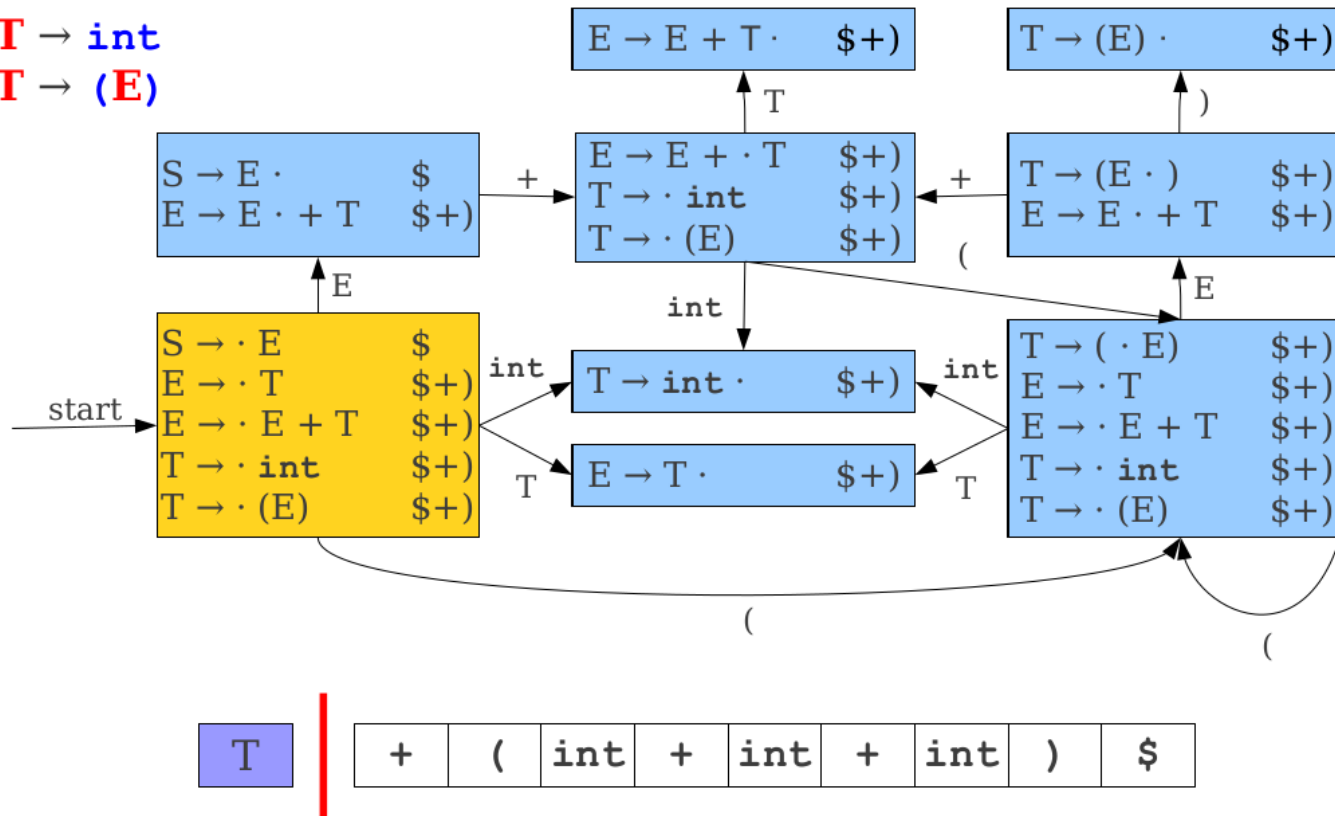
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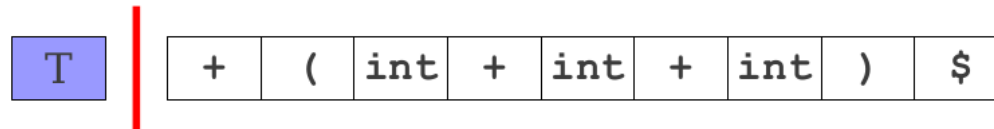
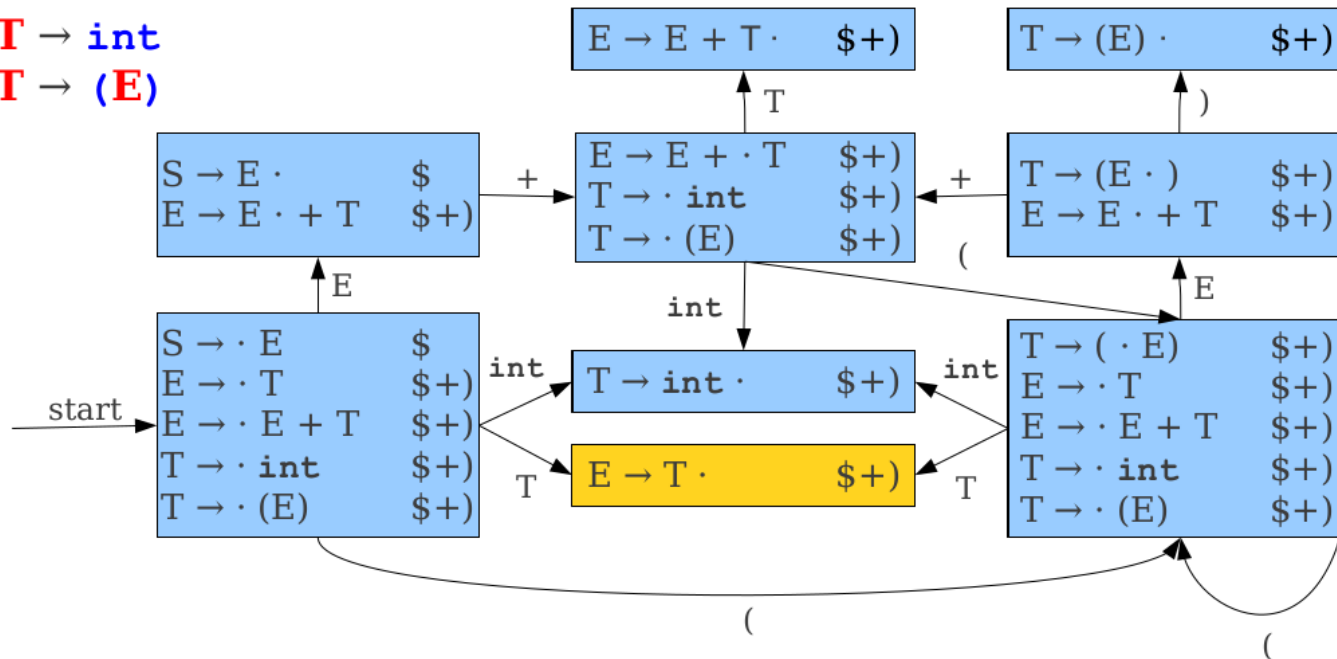
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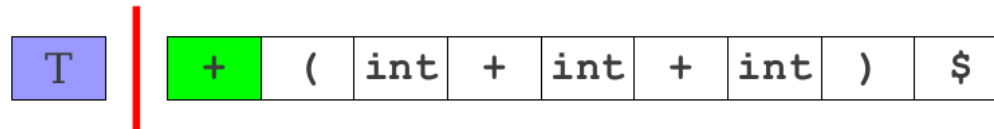
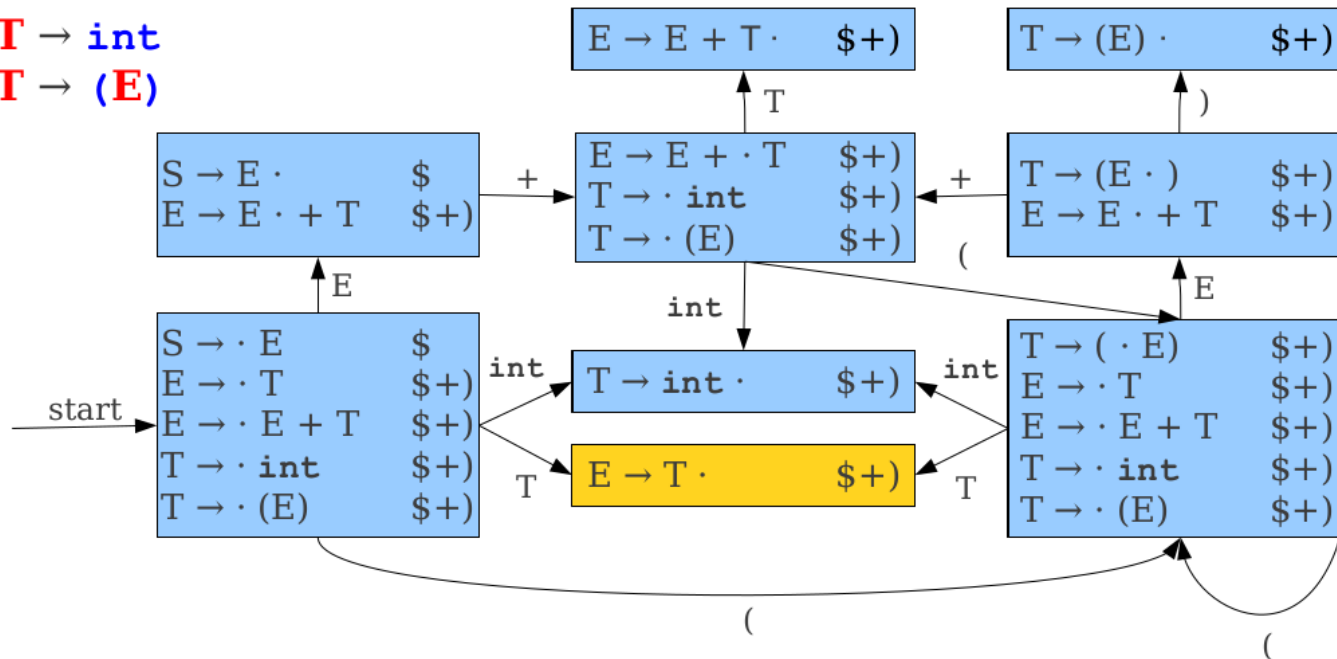
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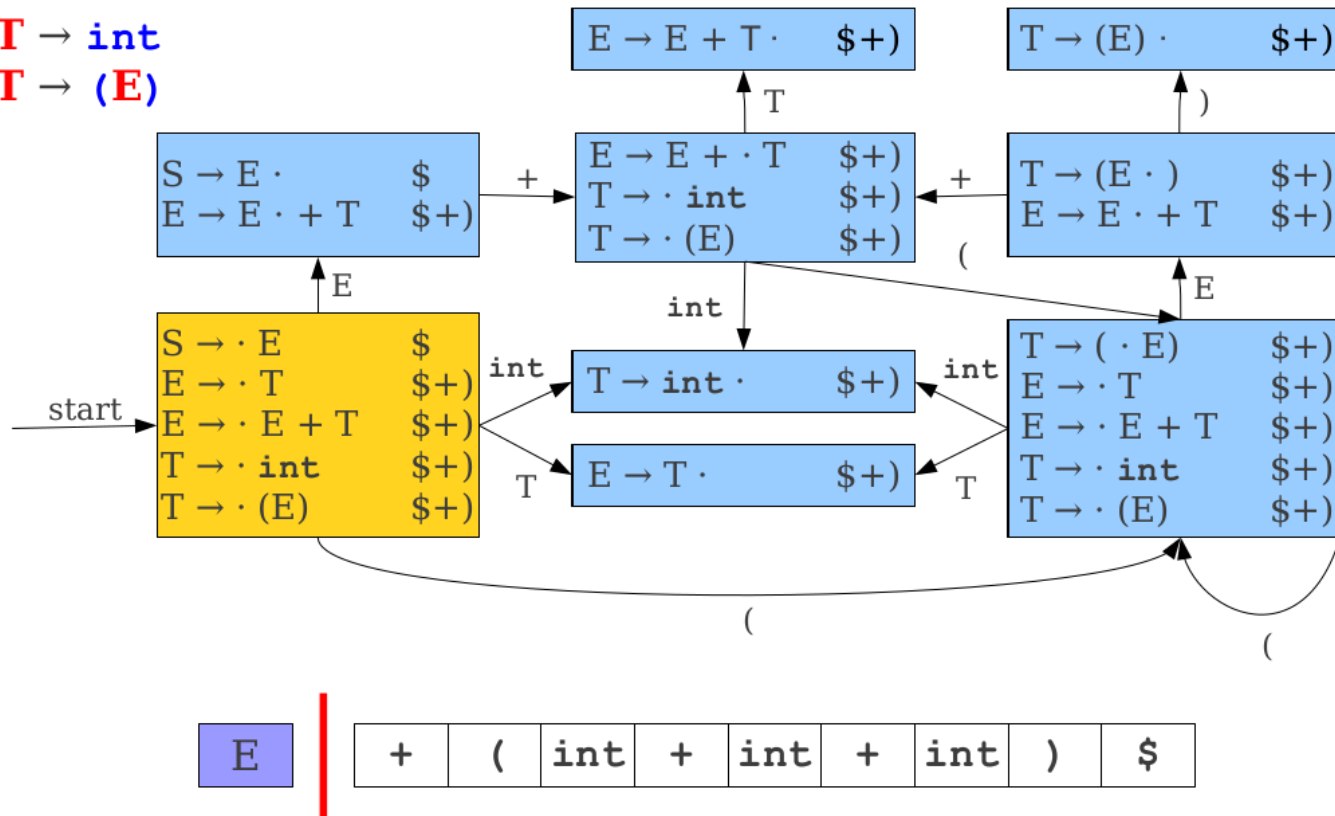
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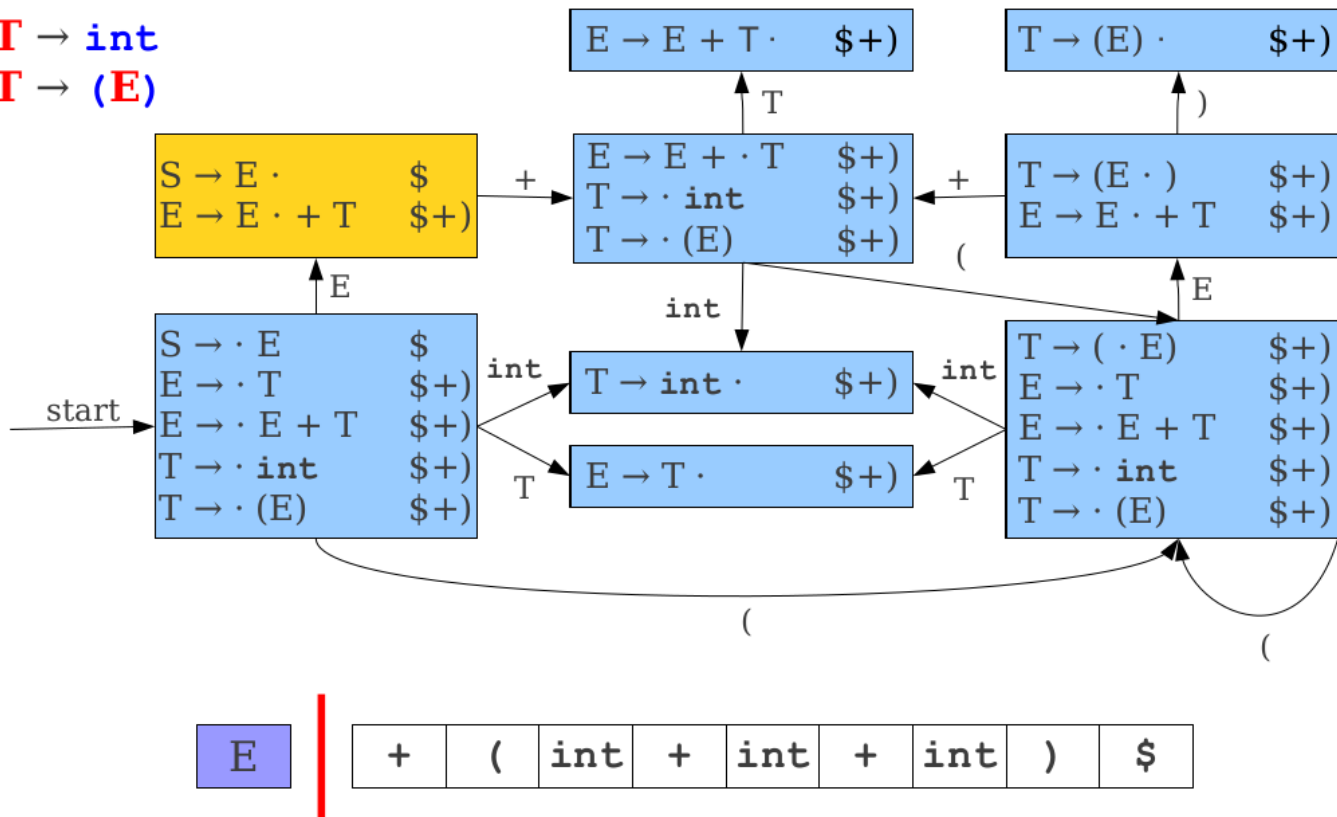
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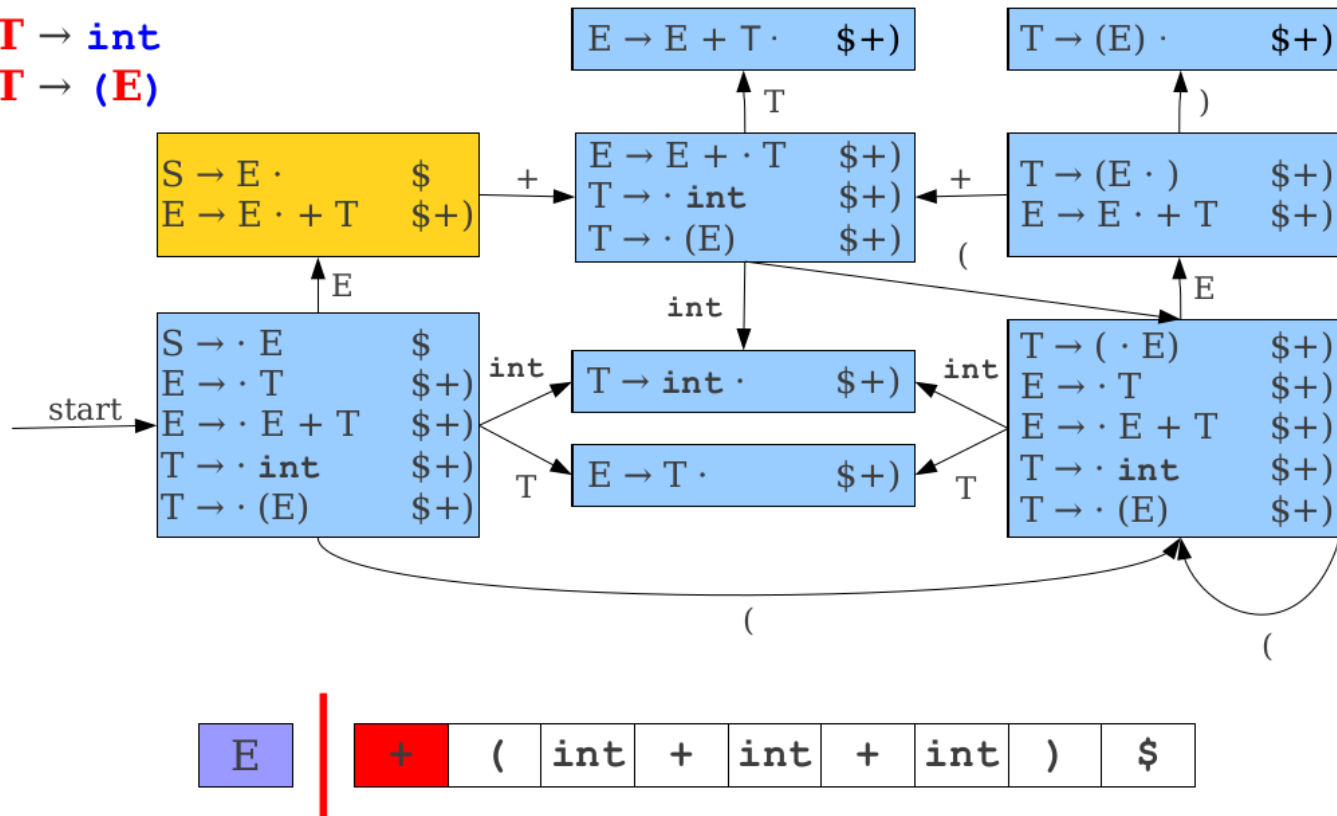
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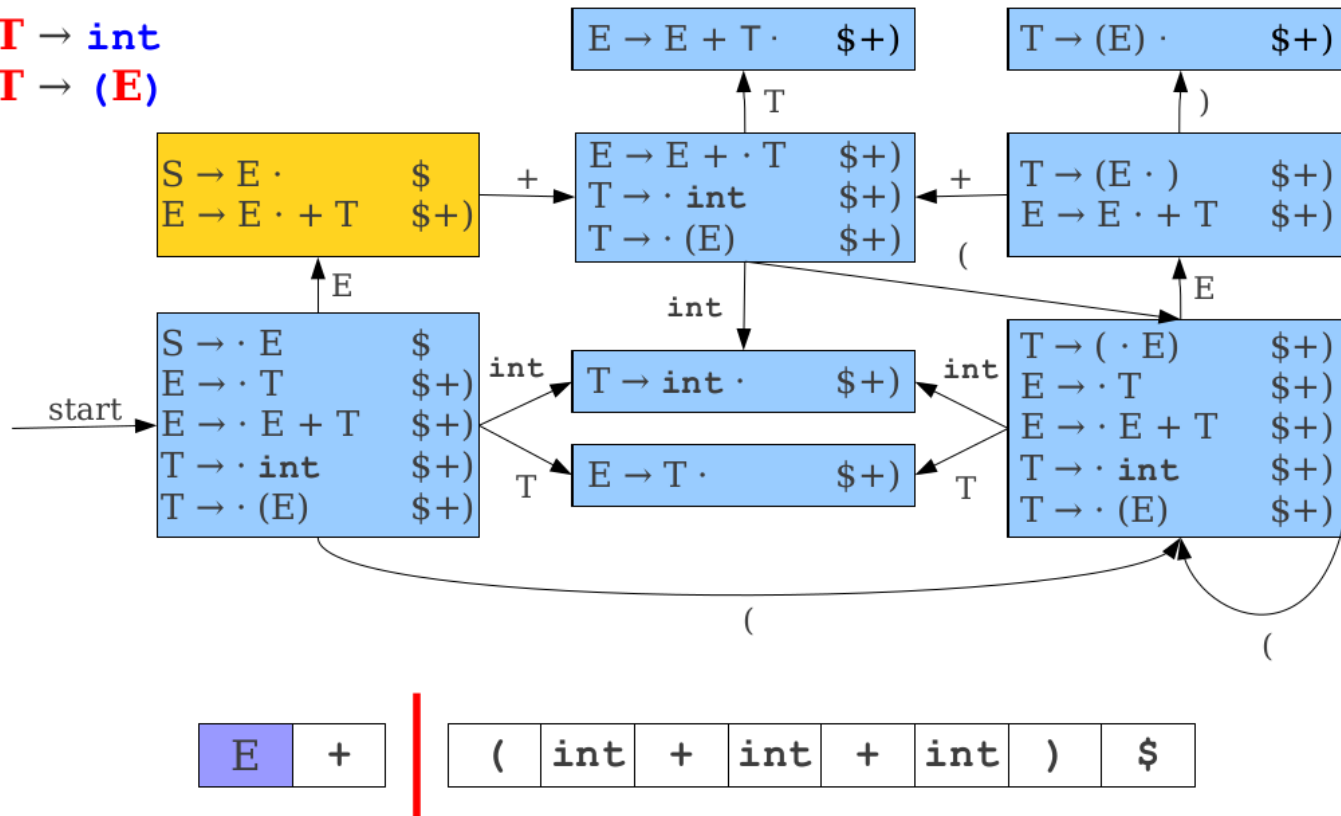
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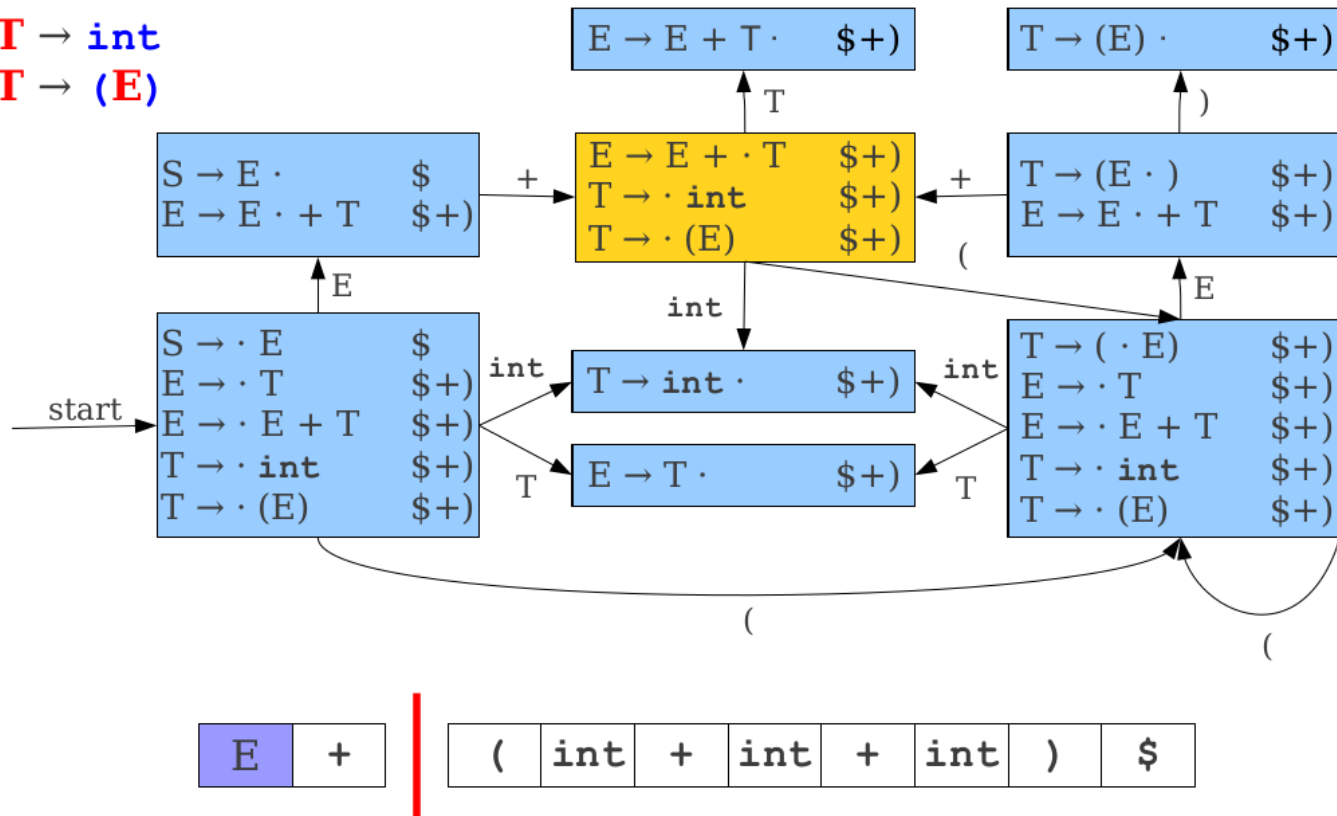
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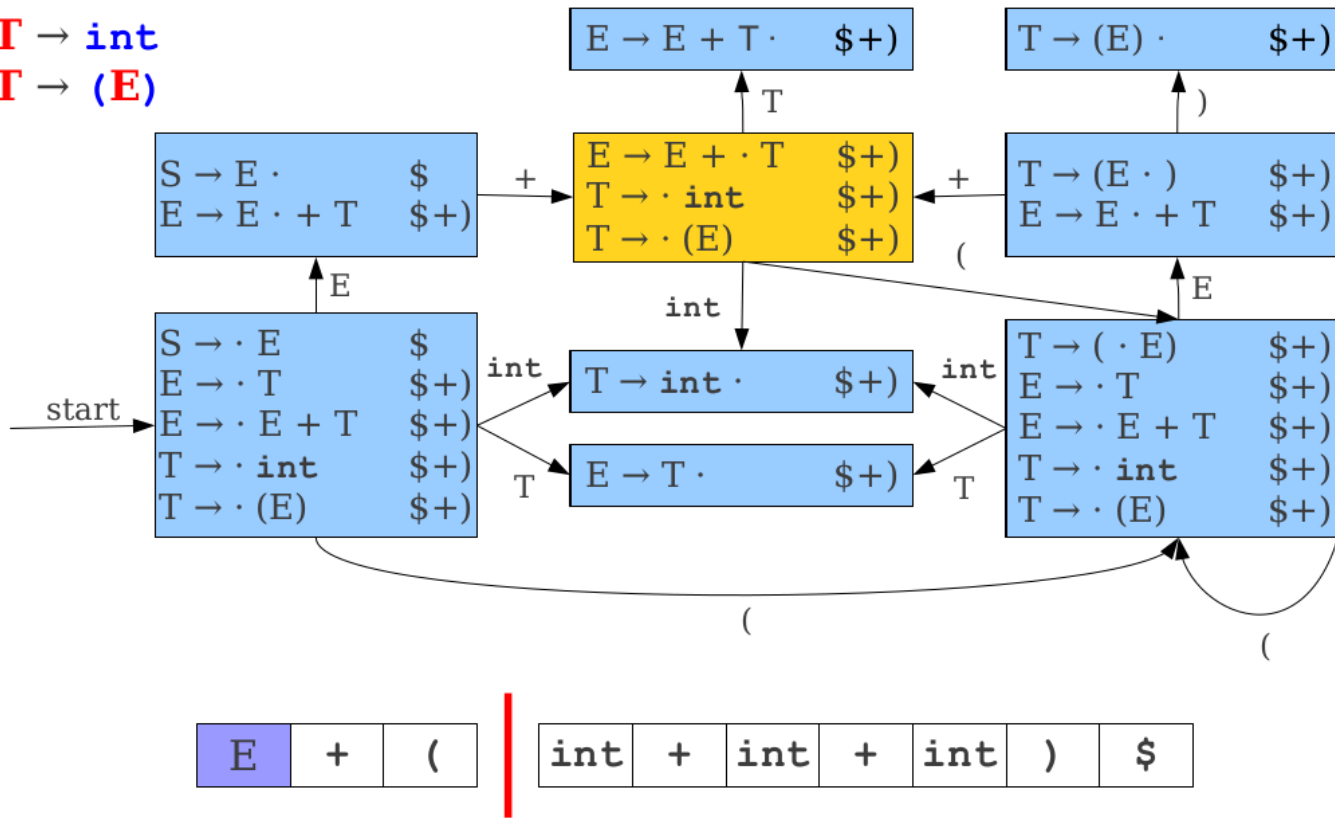
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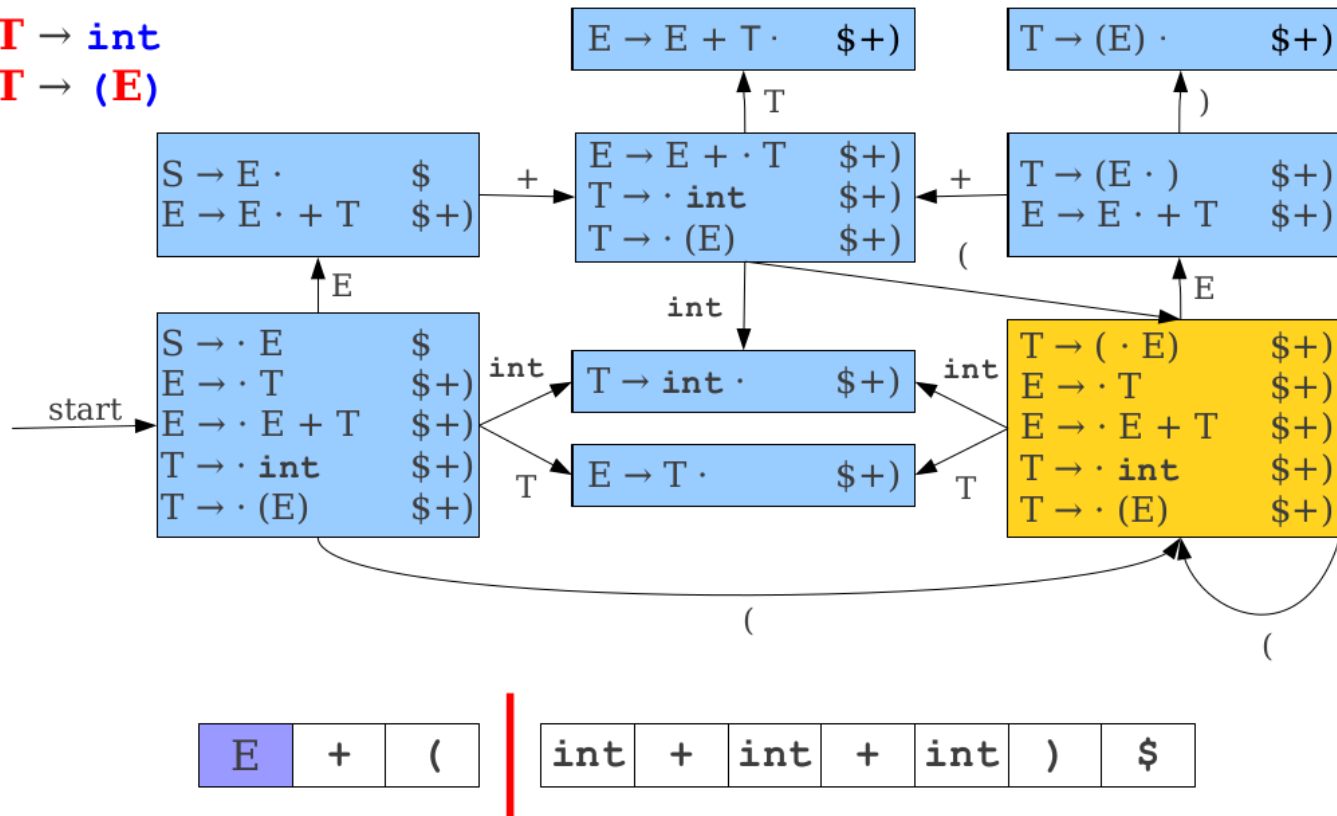
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SLR(1) Parsing

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Analysis of SLR(1)

- Exploits lookahead in a small space.
 - Small automaton – same number of states as in as LR(0).
 - Works on many more grammars than LR(0)
- Too weak for most grammars: lose context from not having extra states.

The Limits of SLR(1)

$S \rightarrow E$
 $E \rightarrow L = R$
 $E \rightarrow R$
 $L \rightarrow id$
 $L \rightarrow *R$
 $R \rightarrow L$

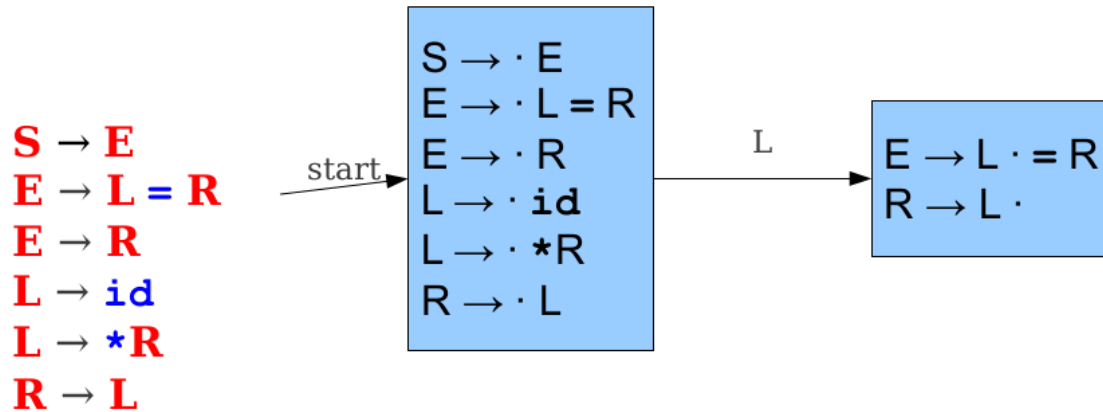
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L → *******R**
R → **L**

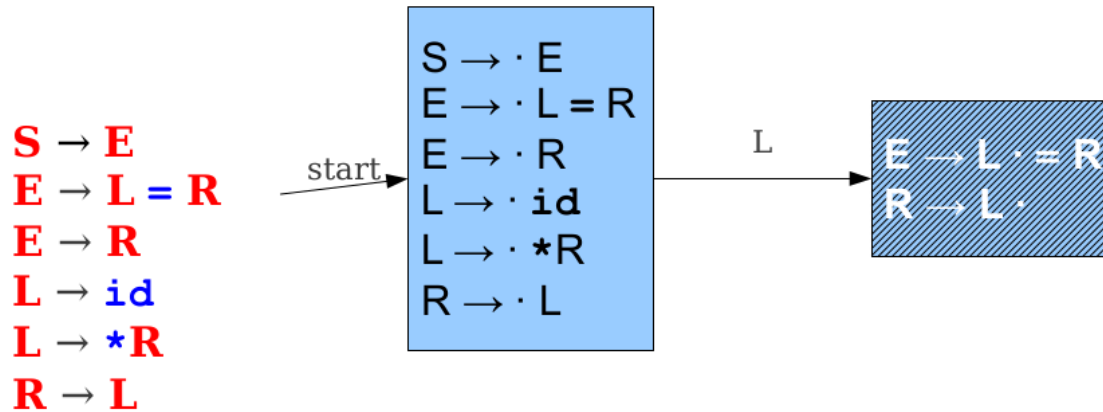
start →

S → · E
E → · L = R
E → · R
L → · id
L → · * R
R → · L

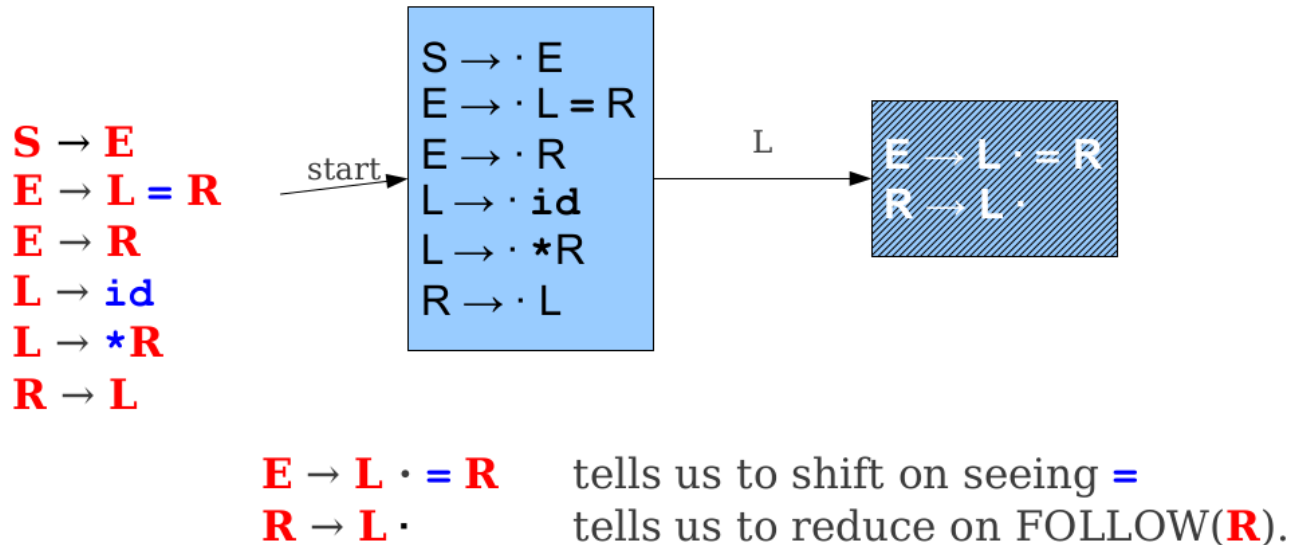
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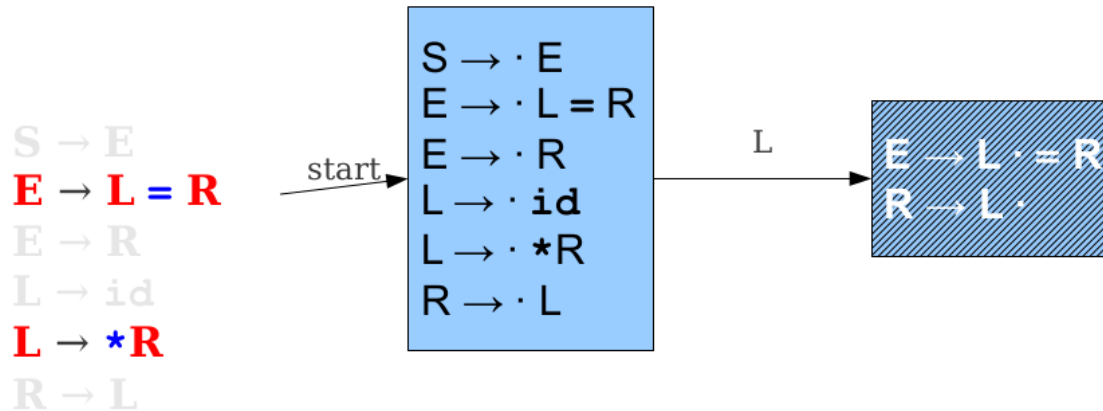
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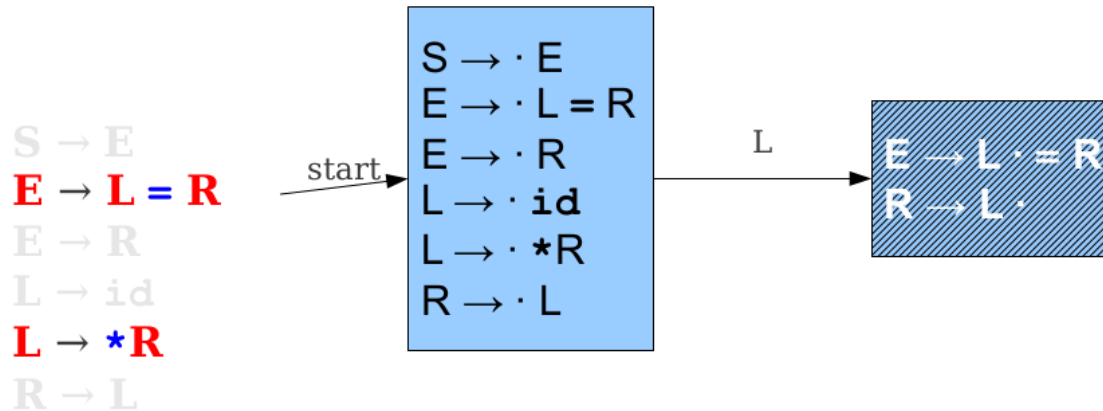


The Limits of SLR(1)



$E \rightarrow L \cdot = R$ tells us to shift on seeing =
 $R \rightarrow L \cdot$ tells us to reduce on FOLLOW(**R**).

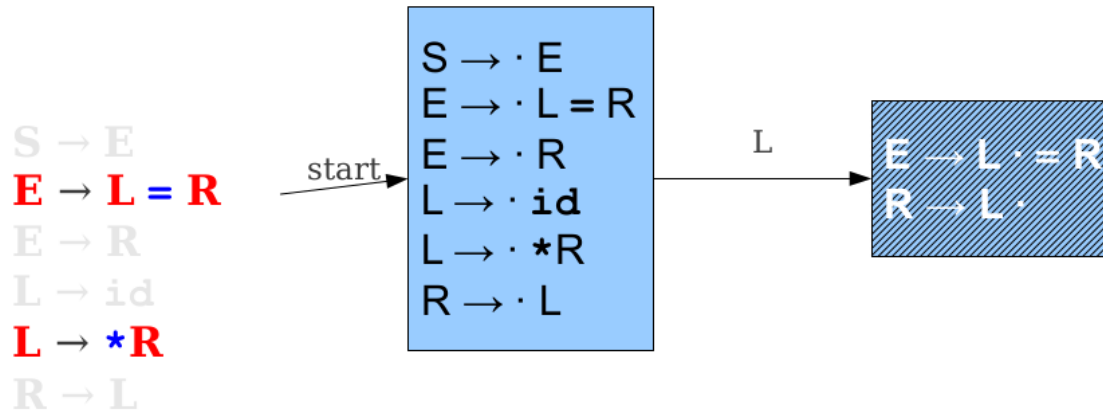
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= ∈ FOLLOW(**R**).

The Limits of SLR(1)



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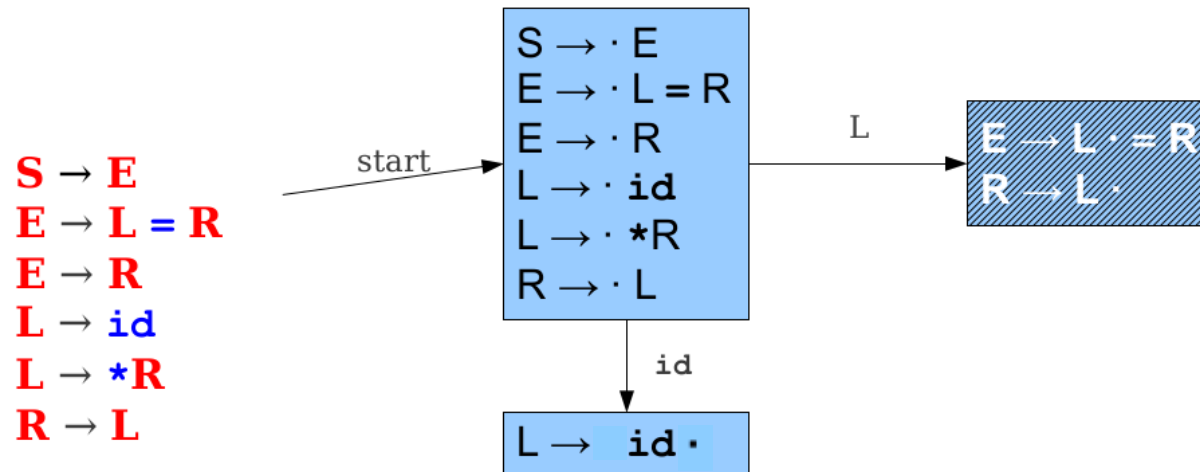
$= \in \text{FOLLOW}(R)$.

We have a conflict!

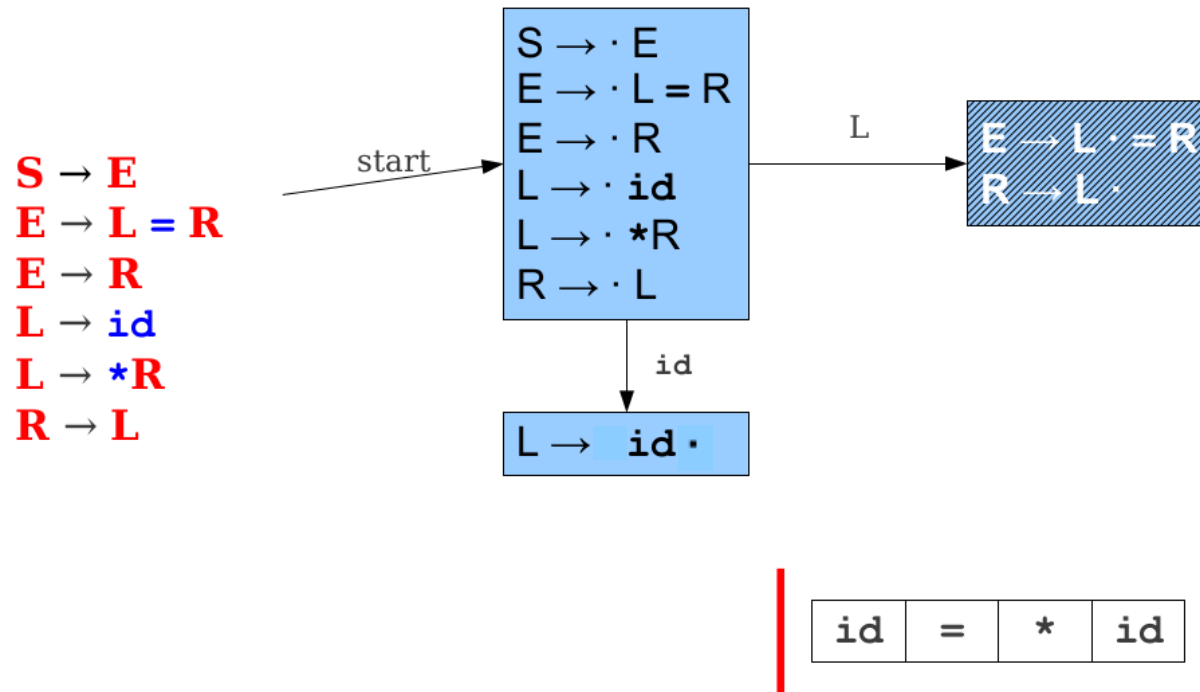
Why is SLR(1) Weak?

- With LR(1), incredible contextual information.
 - Lookaheads at each state only possible after applying the productions that could get us there.
- With SLR(1), minimal context.
 - FOLLOW(A) means “what could follow A somewhere in the grammar?,” even if in a particular state A couldn't possibly have that symbol after it.

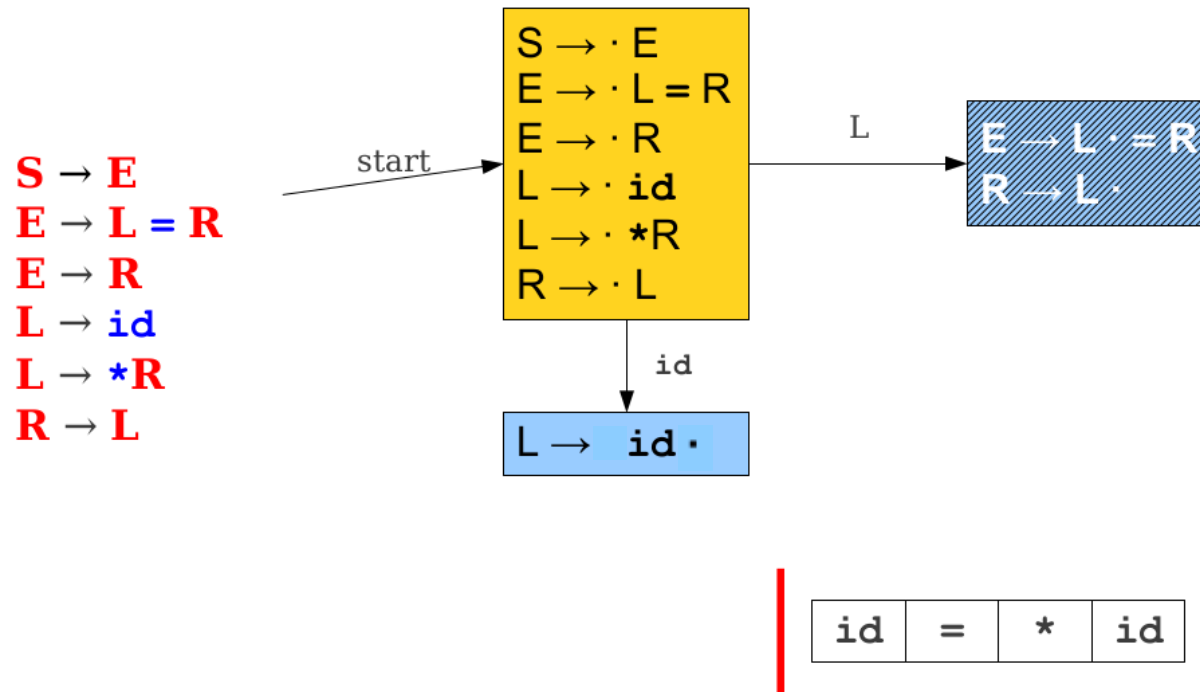
A Lack of Context



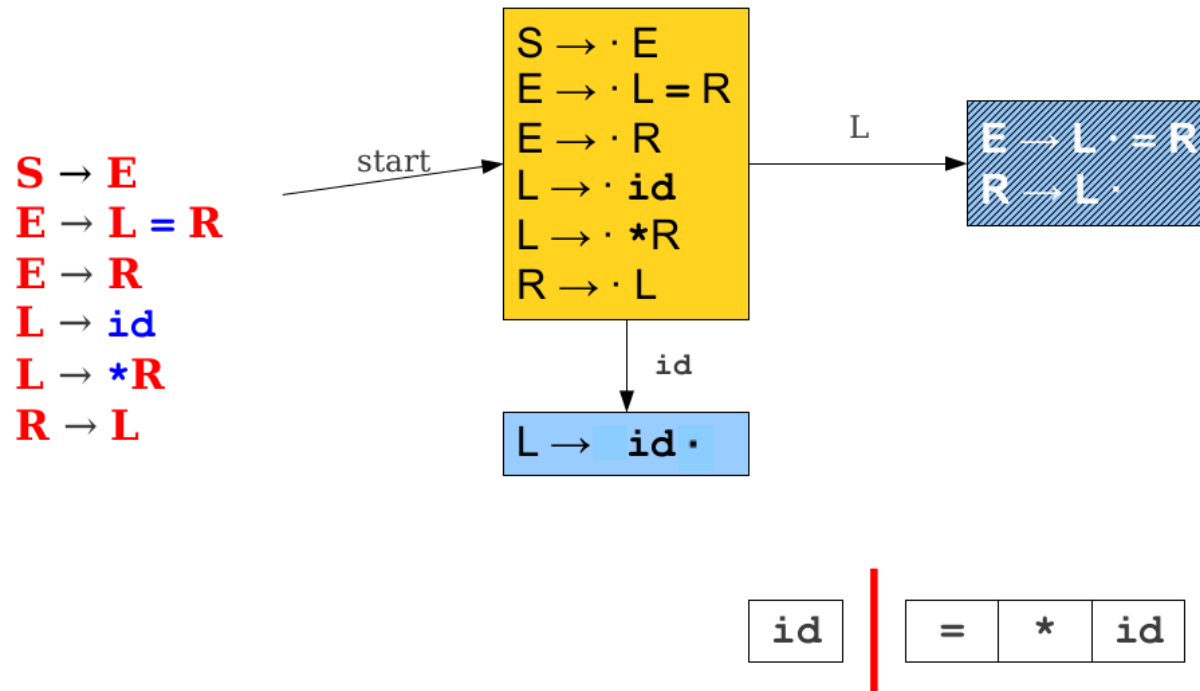
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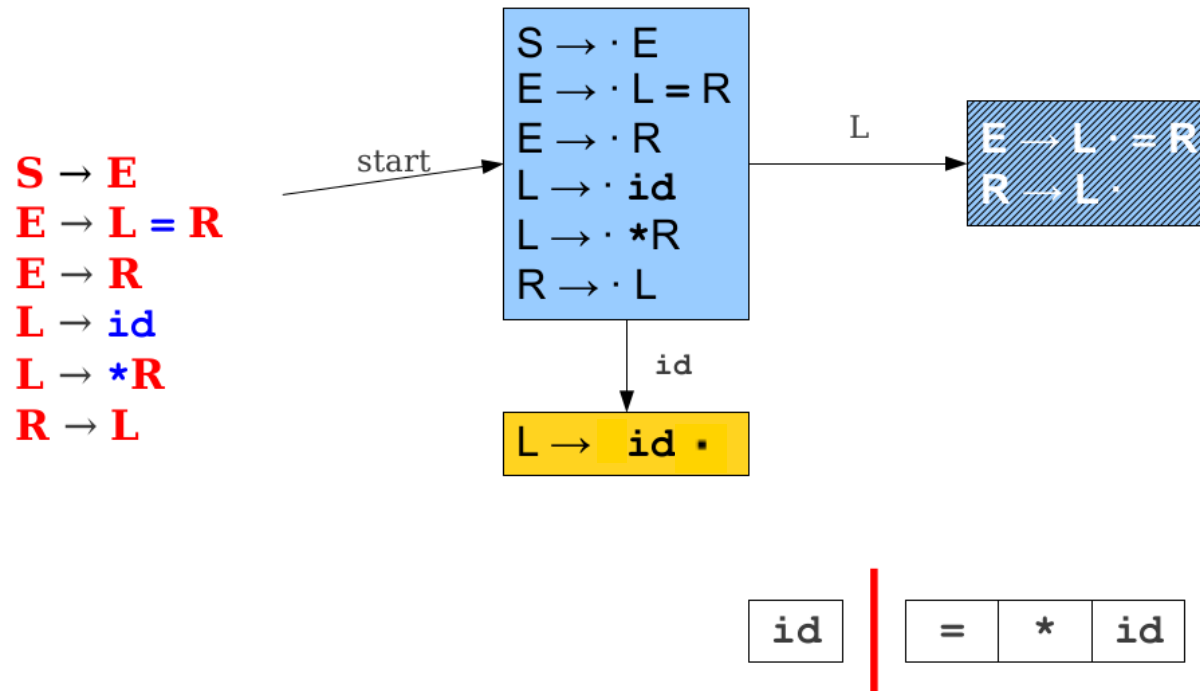
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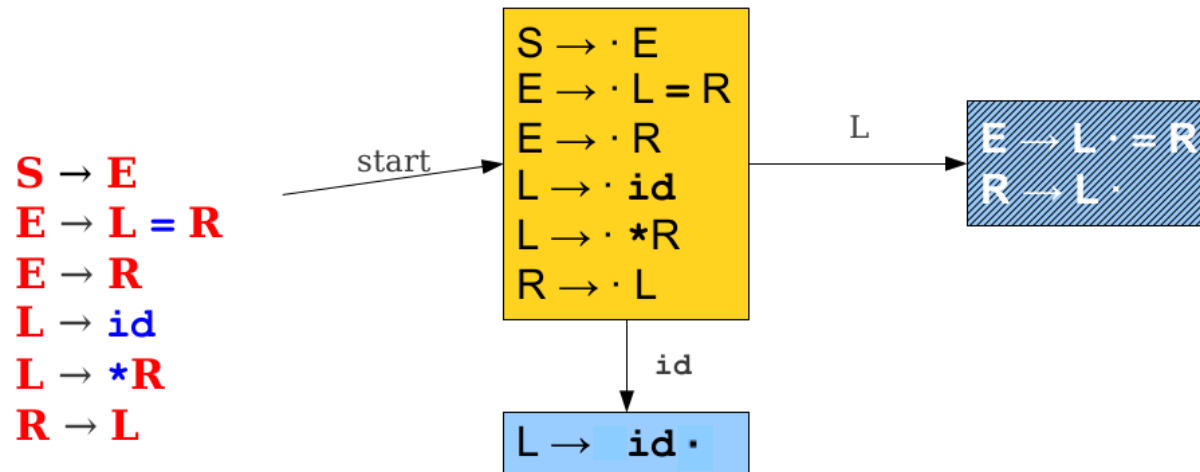
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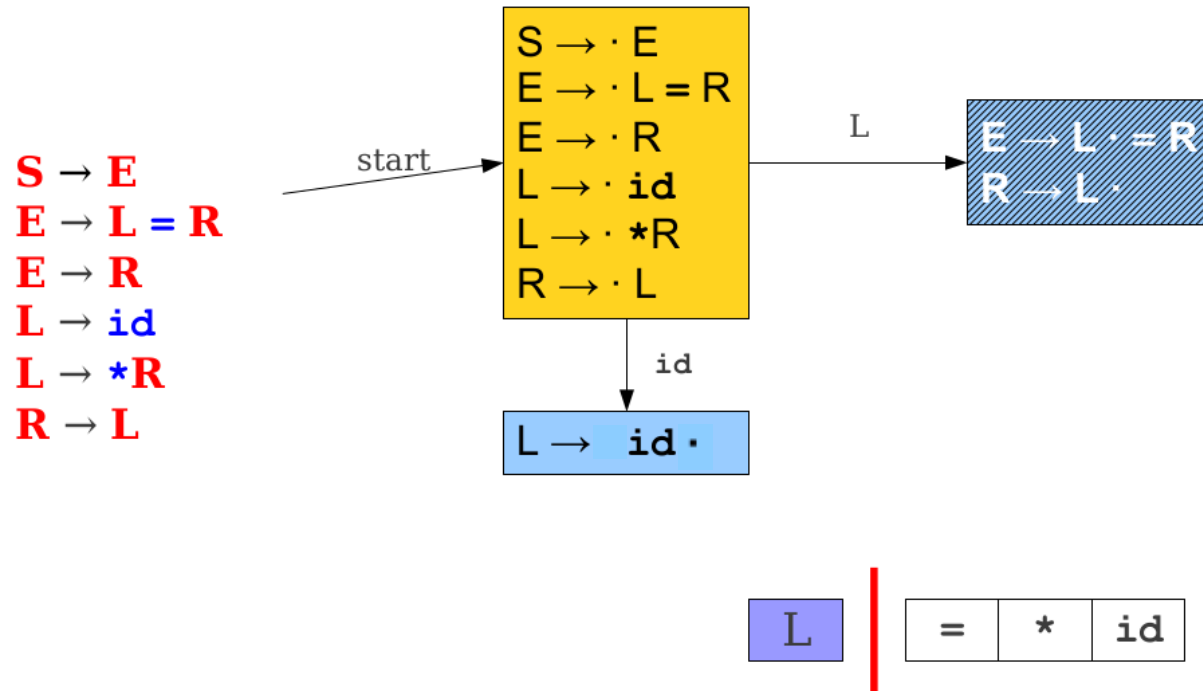


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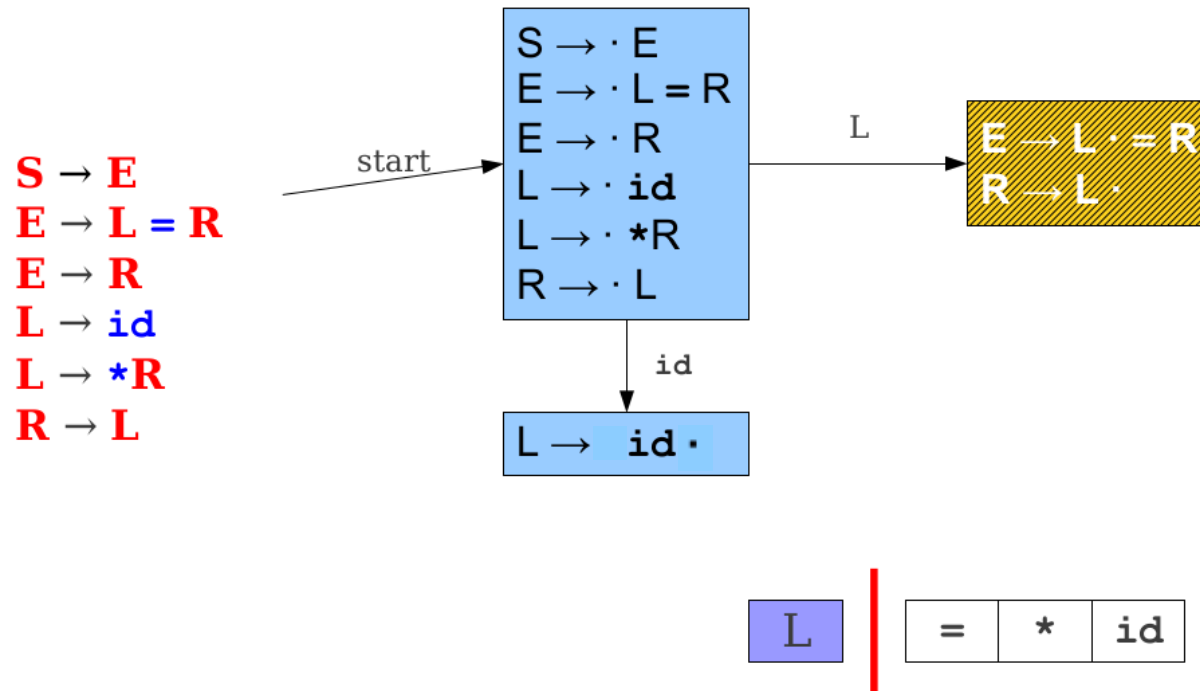


| = * id

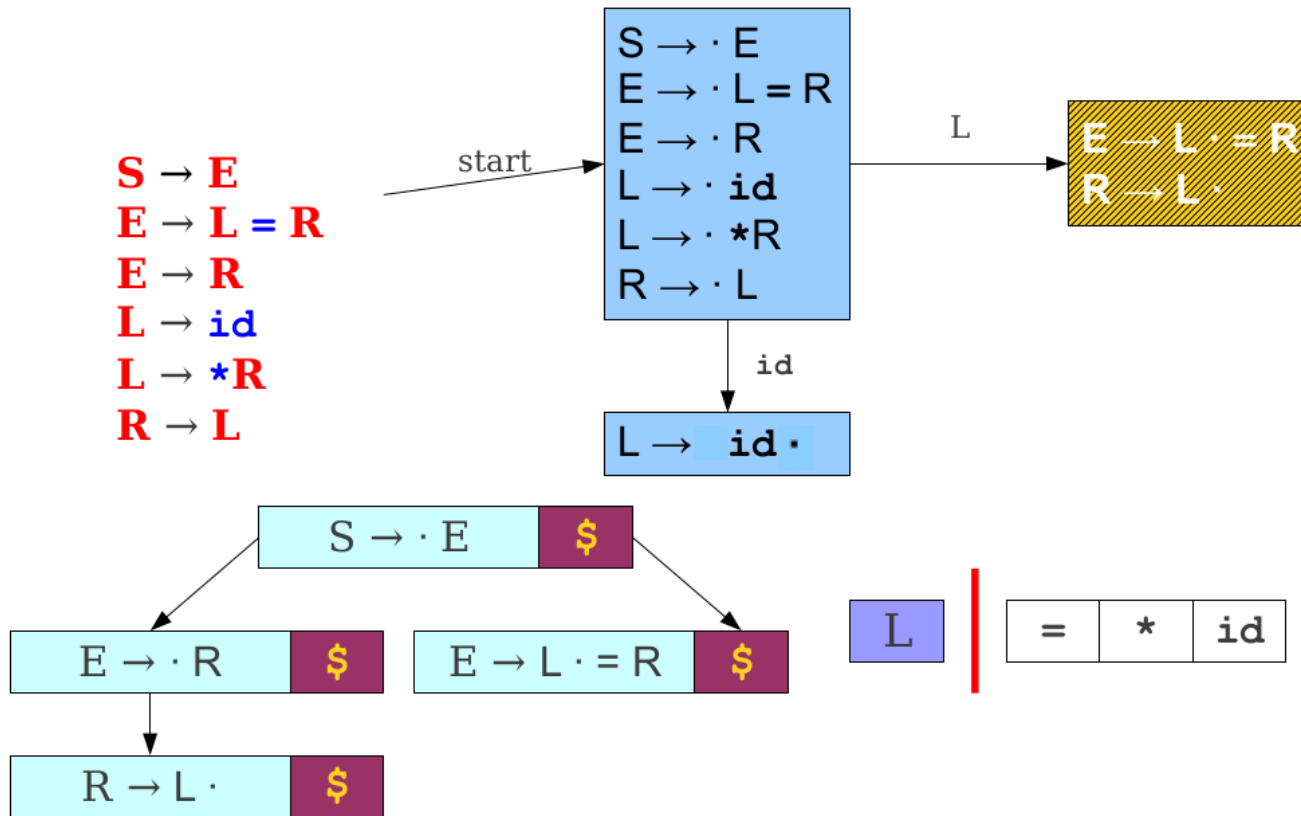
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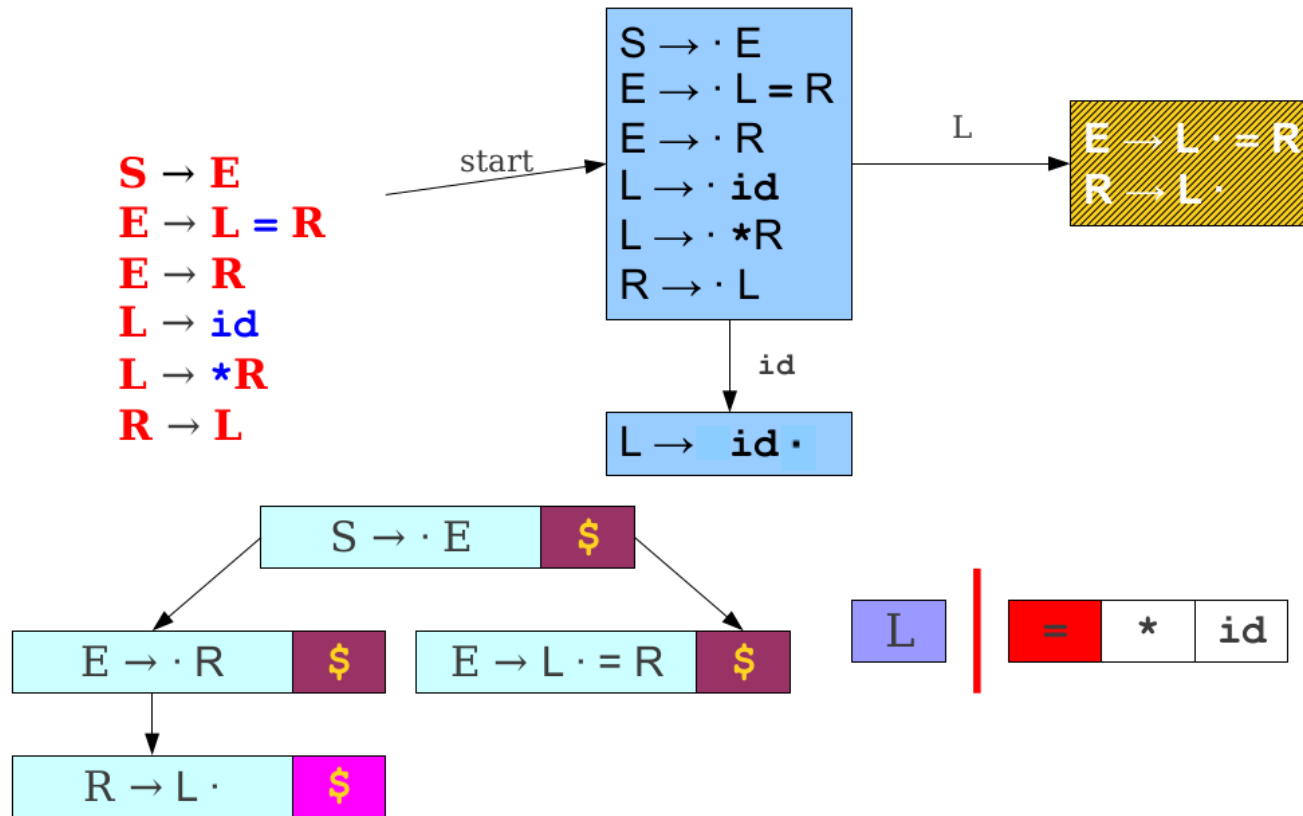
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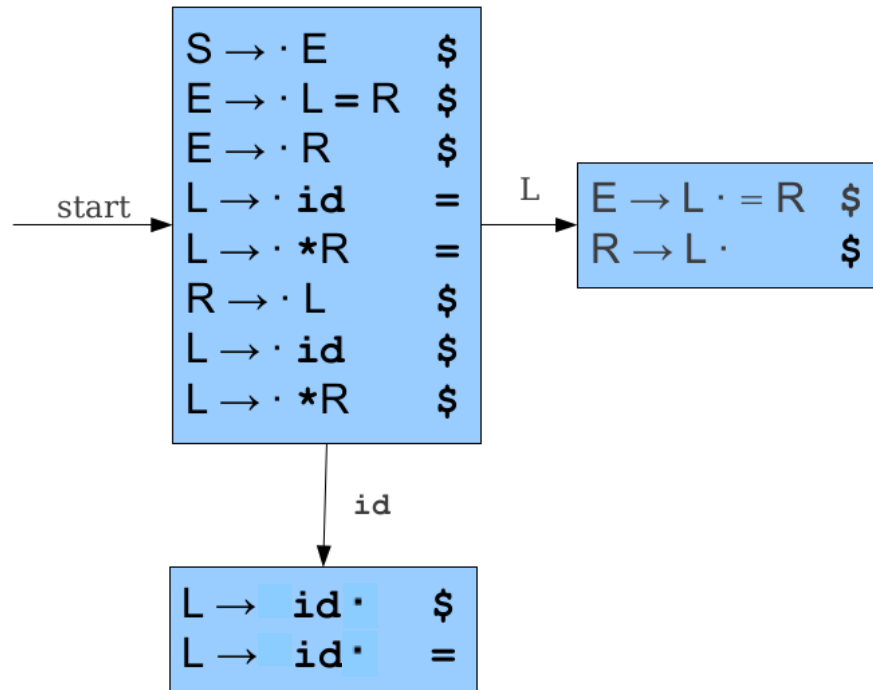


For Reference: LR(1) States

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LR(1) and SLR(1)

- SLR(1) is weak because it has no contextual information.
- LR(1) is impractical because its contextual information makes the automaton too big.
- Can we retain the LR(1) automaton's contextual information without all its states?

Review of LR(1)

- Each state in an LR(1) automaton is a combination of an LR(0) state and lookahead information.
- Two LR(1) items have the same core if they are identical except for lookahead.

$T \rightarrow (\cdot E)$	\$
$E \rightarrow \cdot E + T$)
$E \rightarrow \cdot T$)
$T \rightarrow \cdot \text{int}$)
$T \rightarrow \cdot (E)$)

$T \rightarrow (\cdot E)$)
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$T \rightarrow \cdot (E)$)

A Surprisingly Powerful Idea

- In an LR(1) automaton, we have multiple states with the same core but different lookahead.
- What if we merge all these states together?
- This is called LALR(1)
 - Lookahead(1) LR(0)

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Question?