

# Representations for feasibly approximable functions

Michal Konečný, Eike Neumann

Aston University, Birmingham, United Kingdom

Thirteenth International Conference on Computability and Complexity in Analysis (CCA2016), 15th June 2016

# Introduction

- Given: continuous real function  $f: [-1, 1] \rightarrow \mathbb{R}$
- Problem: compute integral and range
- Generally perceived to be easy in practice
- KO, FRIEDMAN (1982, 1984): integration and range  
# P-hard and NP-hard in general, even for smooth functions:  
There exists a polytime computable smooth  $f: [-1, 1] \rightarrow \mathbb{R}$   
s.t.
  - $i(t) = \int_{-1}^t f(s) ds$  is polytime computable iff  $FP = \# P$ .
  - $m(t) = \max_{s \in [-1, t]} f(s)$  is polytime computable iff  $P = NP$ .

# The elephant in the room

- Where does this discrepancy between “theory” and “practice” come from?
- One explanation: good asymptotic running time in output accuracy does not correspond to practical feasibility.
- Other approach: operators are polytime on well-behaved classes of functions.

# Positive results

- MÜLLER (1987): if  $f$  is polytime real analytic then  $t \mapsto \int_{-1}^t f(s) ds$  and  $t \mapsto \max_{s \in [-1, t]} f(s)$  are polytime computable.
- LABHALLA, LOMBARDI, MOUTAI (2001): Same is true for polytime smooth functions with well-behaved quantitative growth of derivatives (*Gevrey-functions*):

$$\|f^{(n)}\|_{\infty} \leq M \cdot R^n \cdot n^{\alpha n},$$

with  $\alpha > 0$ . Note that  $\alpha = 1$  corresponds to analytic functions.

- KAWAMURA, MÜLLER, RÖSNICK, ZIEGLER (2015): these results translate to uniform results in second-order complexity theory.

# Closure properties of polytime Gevrey functions

The polytime Gevrey functions lack some “obvious” closure properties:

- 1 Not closed under parametric maximisation  
 $t \mapsto \max_{s \in [-1, t]} f(s)$ .
- 2 Not closed under square roots.
- 3 More subtle: the polytime Gevrey functions *are* closed under bounded division (division  $f/g$  where  $g \geq 1$ )...
- 4 but not *uniformly* closed.
- 5 The family of analytic functions

$$\frac{1}{1 + ax^2}$$

is not uniformly polytime Gevrey (in  $\log a$ ).

# Can we do better?

Goal: Find a class of polytime computable functions with better closure properties on which integral and maximum are uniformly polytime computable.

# Feasibly approximable functions

## Theorem (LABHALLA et al., 2001)

*A polytime computable function is  $\alpha$ -Gevrey for some  $\alpha > 0$  if and only if it is feasibly polynomially approximable.*

A function  $f$  is feasibly polynomially approximable iff there is a polytime computable sequence  $(\tilde{f}_n)_n$  of polynomials with

$$\|f - \tilde{f}_n\| < 2^{-n}.$$

⇒ study more general classes of feasibly approximable functions.

- Feasibly approximable  $\simeq$  polytime computable point of a Cauchy representation.

⇒ study more general Cauchy-representations.

- Commonly used representations besides polynomials: rational approximations and splines or piecewise polynomials.

# On the theoretical foundation

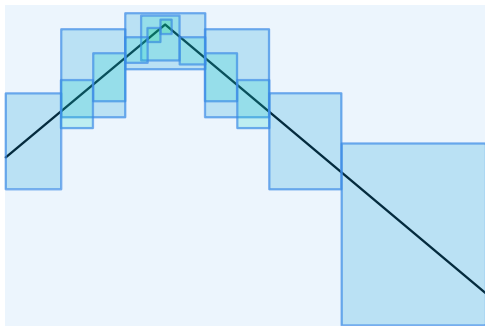
- Officially we work with second-order complexity as introduced by KAWAMURA and COOK (2010).
- We will present very concrete constructions on very concrete representations though...
- So we will be quite informal in this talk.



# Representations: Fun

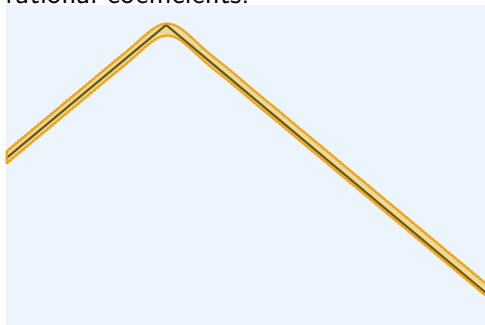
Representation Fun: Real function  $f: [-1, 1] \rightarrow \mathbb{R}$  is represented by  $\varphi: \mathbb{ID} \rightarrow \mathbb{ID}$ , taking dyadic rational intervals to dyadic rational intervals with

- $\varphi(I) \supseteq f(I)$ .
- $(I_n)_n \rightarrow \{x\} \Rightarrow \varphi(I_n) \rightarrow \{f(x)\}$ .



# Representations: Poly

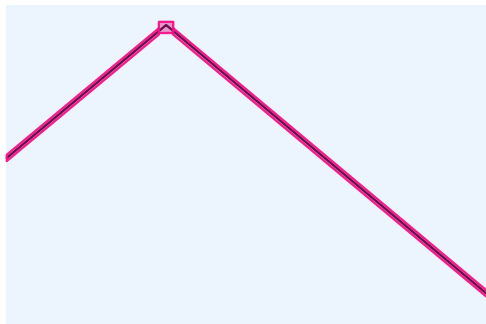
Representation Poly: Real function  $f: [-1, 1] \rightarrow \mathbb{R}$  is represented by fast converging Cauchy sequence of polynomials with dyadic rational coefficients.



# Representations: PPoly

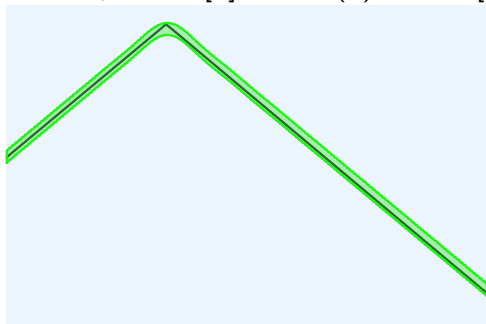
Representation PPoly: Real function  $f: [-1, 1] \rightarrow \mathbb{R}$  is represented by fast converging Cauchy sequence of *piecewise polynomials*:

- $[-1, 1]$  is divided onto rational segments  $[a, b]$ ,  $a, b \in \mathbb{Q}$ .
- On each segment,  $f$  is approximated by a polynomial with dyadic rational coefficients.



# Representations: Frac

Representation Frac: Real function  $f: [-1, 1] \rightarrow \mathbb{R}$  is represented by fast converging Cauchy sequence of rational functions  $P_n/Q_n$  with  $P_n, Q_n \in \mathbb{D}[x]$  and  $Q_n(x) \geq 1$  on  $[-1, 1]$ .



# Relationship between the representations

## Theorem

*We have polytime computable translations*

$$\text{Poly} \rightarrow \text{PPoly} \leftrightarrow \text{Frac} \rightarrow \text{Fun}$$

*none of which reverses unless indicated.*

Most of these proved in LABHALLA et al. (2001) The only “new” result is the translation  $\text{Frac} \rightarrow \text{PPoly}$ .

## Computing $\text{Frac} \rightarrow \text{PPoly}$

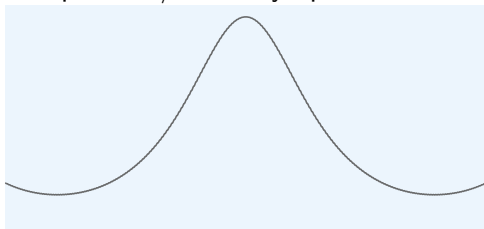
Computing  $1/f$  ( $f \geq 1$ ) w.r.t PPoly:

- Get a piecewise-polynomial approximation  $\tilde{f}$  of  $f$ .

# Computing $\text{Frac} \rightarrow \text{PPoly}$

Computing  $1/f$  ( $f \geq 1$ ) w.r.t PPoly:

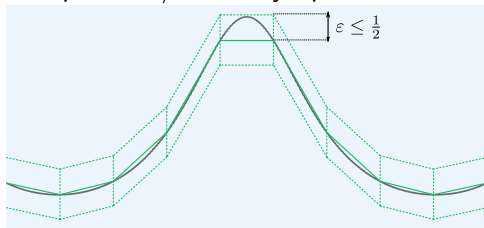
- Get a piecewise-polynomial approximation  $\tilde{f}$  of  $f$ .
- Interpolate  $1/\tilde{f}$  linearly up to error  $\varepsilon < 1/2$ .



# Computing $\text{Frac} \rightarrow \text{PPoly}$

Computing  $1/f$  ( $f \geq 1$ ) w.r.t PPoly:

- Get a piecewise-polynomial approximation  $\tilde{f}$  of  $f$ .
- Interpolate  $1/\tilde{f}$  linearly up to error  $\varepsilon < 1/2$ .



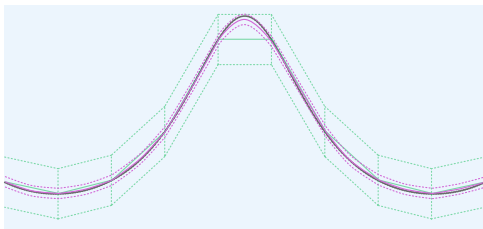


# Computing Frac $\rightarrow$ PPoly

Computing  $1/f$  ( $f \geq 1$ ) w.r.t PPoly:

- Get a piecewise-polynomial approximation  $\tilde{f}$  of  $f$ .
- Interpolate  $1/\tilde{f}$  linearly up to error  $\varepsilon < 1/2$ .
- Apply Newton-Raphson division to improve approximation:

$$f_{n+1} = 2f_n - f_n^2 \tilde{f}.$$



- Degree after  $n$  iterations:  $(2^n - 1)d_{\tilde{f}} + 2^n$ .
- Error after  $n$  iterations:  $\varepsilon^{2^n}$ .

# Computing range and integral

Theorem (LABHALLA et al., 2001)

*Maximisation and Integration (viewed as functionals) are uniformly polytime computable w.r.t. PPoly.*

Corollary

*Maximisation and Integration (viewed as functionals) are uniformly polytime computable w.r.t. Frac.*

# Computing Range

Computing  $\max_{[-1,1]} f$  with respect to Poly:

- Get a polynomial approximation  $\tilde{f}$  with accuracy  $\varepsilon$ .
- Compute the critical points of  $\tilde{f}$ :
  - Compute the derivative  $\tilde{f}'$ .
  - Compute the *separable part*  $\tilde{f}'_s = \tilde{f}' / \gcd(\tilde{f}', \tilde{f}'')$ .
  - Isolate the real roots of  $\tilde{f}'_s$ .
  - Approximate the roots up to sufficient accuracy.
- Take the maximum over the critical points and the boundary points.

For PPoly, do this piecewise.

# Closure properties of PPoly

## Theorem

*The class of feasibly piecewise-polynomially approximable functions is uniformly closed under*

- *Addition, subtraction and multiplication*
- *Bounded division*
- *Composition*
- *Pairwise and parametric maximisation*
- *Taking antiderivatives*
- *Taking square roots*
- *Taking absolute values*

# Ko-Friedman revisited

## Corollary

*If  $f: [-1, 1] \rightarrow \mathbb{R}$  is expressible as a term whose leaves can be written as polytime computable analytic functions and whose nodes use only the operations  $+$ ,  $-$ ,  $\times$ ,  $\div$ ,  $\circ$ ,  $\max$ ,  $\int$ ,  $\sqrt{\cdot}$ ,  $|\cdot|$ , then*

- *$f$  is polytime computable.*
- *$m(t) = \max_{s \in [-1, t]} f(s)$  is polytime computable.*
- *$i(t) = \int_{-1}^t f(s) ds$  is polytime computable.*

# Validation

- How well do these theoretical results translate into practice?
- To validate our results we have implemented the representations and algorithms in Haskell (<http://tinyurl.com/aern2-fnreps>):
- Fun: Haskell type

$\text{Interval MPFloat} \rightarrow \text{Interval MPFloat}.$

- Poly: Haskell type

$\text{Int} \rightarrow (\text{Map Int MPFloat}, \text{Double}).$

- PPoly: Haskell type

$\text{Int} \rightarrow ([(\text{Interval Rational}, \text{Map Int MPFloat})], \text{Double}).$

# Validation

- Polynomial approximations are constructed from Taylor series or using Chebyshev-interpolation via Discrete Cosine Transform.
- Division for PPoly uses Newton iteration, division for Poly uses Chebyshev-interpolation.
- Integration for Fun is just Riemann integration.
- To have a slightly less naive integration algorithm available for Fun we introduced a new representation DFun which encodes a Fun name of a function and its derivative.

# Maximisation revisited

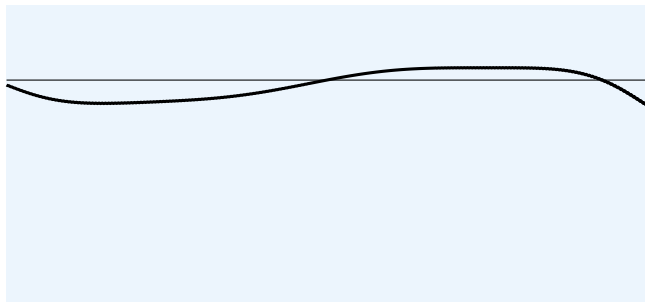
- How to compute the maximum in practice?
- Initial attempt: use root isolation. Compute separable part using signed subresultant sequences.
- Quadratic blowup in coefficient size in the signed subresultant sequence turns out to be a bit much.

⇒ avoid computing separable part.



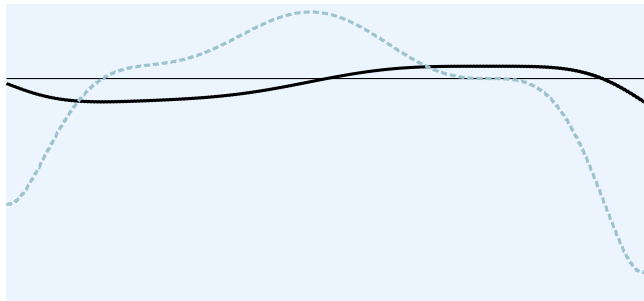
# Maximisation algorithm

- Get a polynomial approximation.



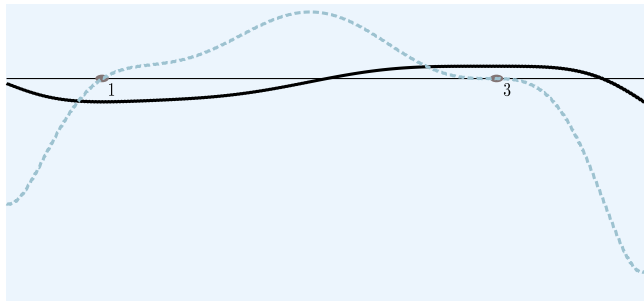
# Maximisation algorithm

- Get a polynomial approximation.
- Compute the derivative.



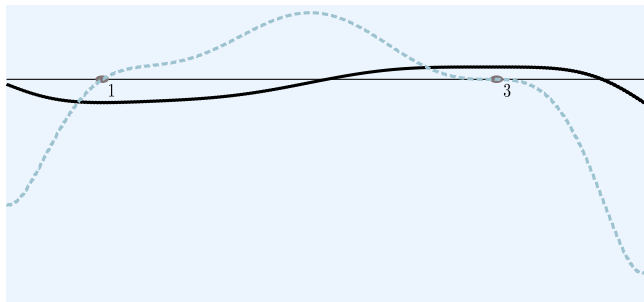
# Maximisation algorithm

- Get a polynomial approximation.
- Compute the derivative.



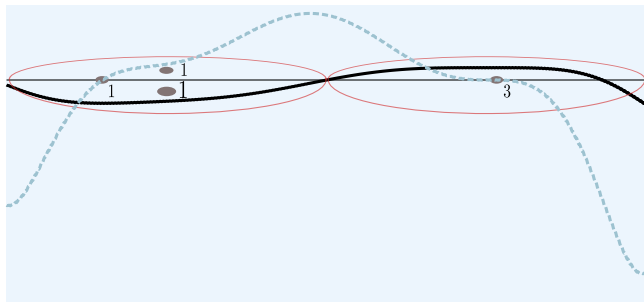
# Maximisation algorithm

- Get a polynomial approximation.
- Compute the derivative.
- Estimate the maximum on the interval using a Lipschitz constant.



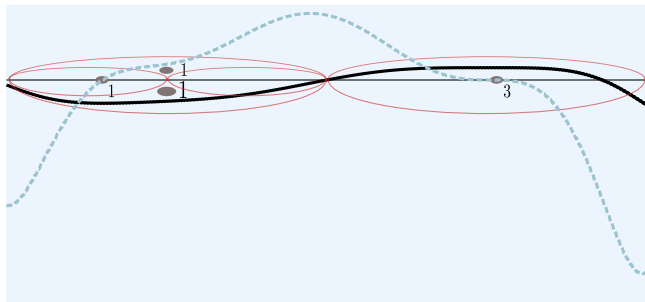
# Maximisation algorithm

- Get a polynomial approximation.
- Compute the derivative.
- Estimate the maximum on the interval using a Lipschitz constant.
- If the estimate is not accurate enough, split the interval in two, and count the complex roots around the sub-intervals.



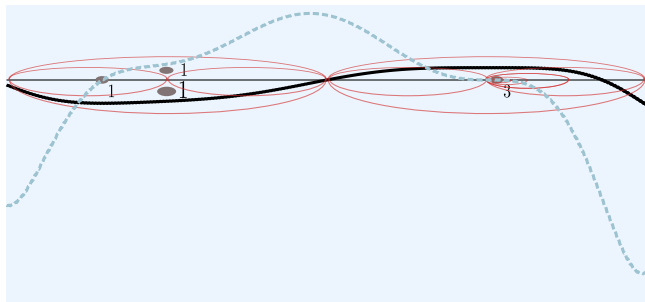
# Maximisation algorithm

- Get a polynomial approximation.
- Compute the derivative.
- Estimate the maximum on the interval using a Lipschitz constant.
- If the estimate is not accurate enough, split the interval in two, and count the complex roots around the sub-intervals.

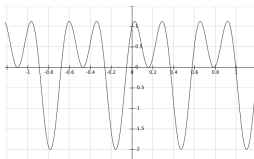


# Maximisation algorithm

- Get a polynomial approximation.
- Compute the derivative.
- Estimate the maximum on the interval using a Lipschitz constant.
- If the estimate is not accurate enough, split the interval in two, and count the complex roots around the sub-intervals.

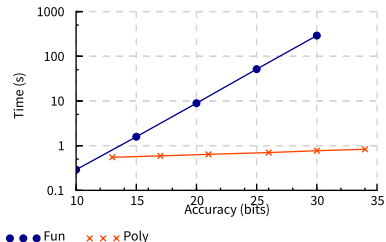


# Benchmarks: A simple analytic function

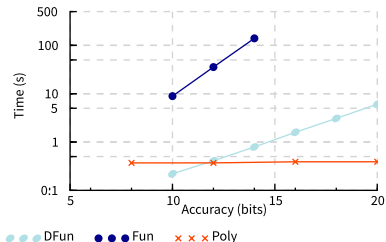


$$f(x) = \sin(10x) + \cos(20x)$$

Maximum:

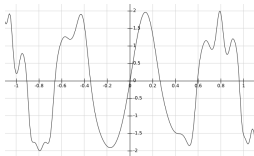


Integral:



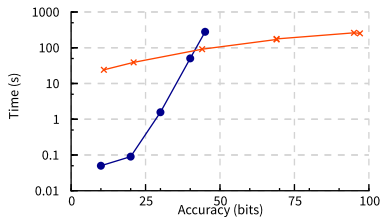


# Benchmarks: A more complicated analytic function involving composition



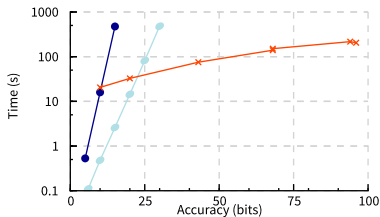
$$f(x) = \sin(10x + \sin(20x^2)) + \sin(10x)$$

Maximum:



● ● ● Fun    × × × Poly

Integral:



● ● ● DFun    ● ● ● Fun    × × × Poly

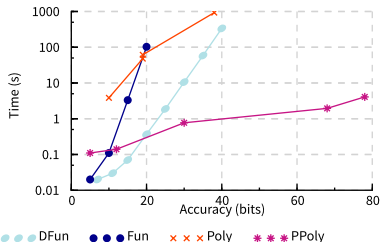
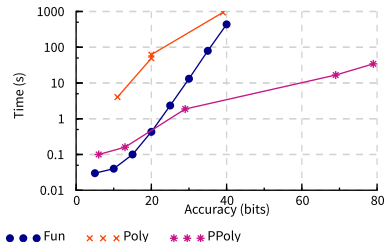
# Benchmarks: An simple analytic function with singularities near the origin



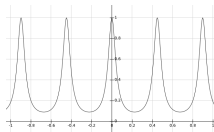
$$f(x) = \frac{x}{1 + 100x^2}$$

Maximum:

Integral:

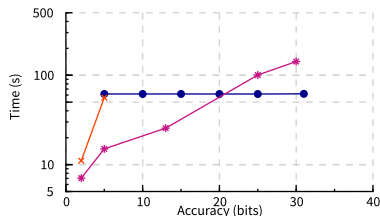


# Benchmarks: A more complicated analytic function with singularities near the origin



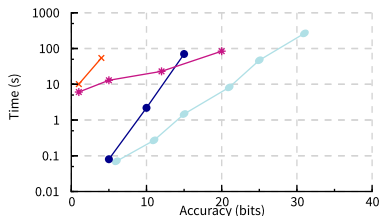
$$f(x) = \frac{1}{1 + 10 \sin^2(7x)}$$

Maximum:



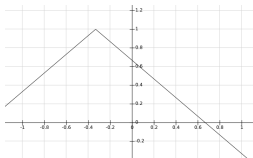
● ● ● Fun    × × × Poly    \* \* \* PPoly

Integral:



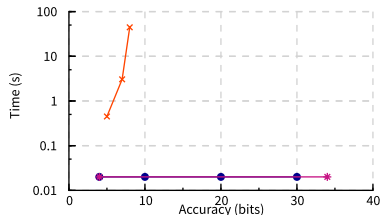
● ● ● DFun    ● ● ● Fun    × × × Poly    \* \* \* PPoly

# Benchmarks: A very simple non-smooth function



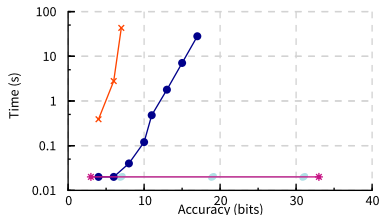
$$f(x) = 1 - |x + 1/3|$$

Maximum:



●●● Fun    ××× Poly    \* \* \* PPoly

Integral:



●●● DFun    ●●● Fun    ××× Poly    \* \* \* PPoly

# Conclusion

- In terms of polytime reducibility, *piecewise polynomials* are better than polynomials and equivalent to rational functions.
- In contrast to the standard function space representation, they render integration and maximisation polytime computable.
- Thus, we've found a class of polytime computable functions which
  - is uniformly closed under integration and maximisation.
  - enjoys nice further closure properties.
- For the functions we tested, polytime computability and practical feasibility seem to match quite well.
  - with some exceptions...

# Future work

- Improve division algorithm.
- Improve PPoly-integration.
- Introduce representations for multivariate functions and study differential equations.
- Is there a nice characterisation of class of feasibly piecewise-polynomially approximable functions?

# References



Akitoshi Kawamura and Stephen A. Cook.  
Complexity theory for operators in analysis.  
*ACM Transactions on Computation Theory*, 4(2):5, 2012.



Ker-I Ko and Harvey Friedman.  
Computational complexity of real functions.  
*Theoretical Computer Science*, 20:323–352, 1982.



Akitoshi Kawamura, Norbert Müller, Carsten Rösnick, and Martin Ziegler.  
Computational benefit of smoothness: Parameterized bit-complexity of numerical operators on analytic functions and Gevrey's hierarchy.  
*Journal of Complexity*, 31(5):689 – 714, 2015.



Ker-I Ko.  
*Complexity Theory of Real Functions*.  
Birkhäuser, 1991.



S. Labhalla, H.Lombardi, and E.Moutai.  
Espaces métriques rationnellement présentés et complexité, le cas de l'espace des fonctions réelles uniformément continues sur un intervalle compact.  
*Theoretical Computer Science*, 250:265–332, 2001.



Norbert Th. Müller.  
Uniform computational complexity of Taylor series.  
In *Automata, Languages and Programming*, volume 267 of *Lecture Notes in Computer Science*, pages 435–444. Springer, 1987.