

THE ALGEBRAIC GEOMETRY ALGEBRAIC TOPOLOGY DUMP

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CONTENTS

References

10

Part 1. Reading notes on Cox, Little, O’Shea’s *Ideals, Varieties, and Algorithms: An Introduction to Computational Algebraic Geometry and Commutative Algebra*

- 1. Geometry, Algebra, and Algorithms
- 2. Groebner Bases
- 3. Elimination Theory
- 4. The Algebra-Geometry Dictionary
- 5. Polynomial and Rational Functions on a Variety
- 6. Robotics and Automatic Geometric Theorem Proving

Part 2. Reading notes on Cox, Little, O’Shea’s *Using Algebraic Geometry*

- 7. Introduction
- 8. Solving Polynomial Equations
- 9. Resultants
- 10. Computation in Local Rings
- 11.
- 12.
- 13. Polytopes, Resultants, and Equations
- 14. Polyhedral Regions and Polynomials
- 15. Algebraic Coding Theory
- 16. The Berlekamp-Massey-Sakata Decoding Algorithm

Part 3. Algebraic Topology

- 17. Simplicial Complexes

Part 4. Graphs, Finite Graphs

- 18. Graphs, Finite Graphs, Trees

Date: 5 mars 2017.
Key words and phrases. Algebraic Geometry, Algebraic Topology.

ABSTRACT. Everything about Algebraic Geometry, Algebraic Topology

Part 1. Reading notes on Cox, Little, O’Shea’s *Ideals, Varieties, and Algorithms: An Introduction to Computational Algebraic Geometry and Commutative Algebra*

- 1. GEOMETRY, ALGEBRA, AND ALGORITHMS
- 1.1. **Polynomials and Affine Space.** fields are important is that linear algebra works over *any* field
- Definition 1** (2). *set of all polynomials in x_1, \dots, x_n with coefficients in k , denoted $k[x_1, \dots, x_n]$*
polynomial f *divides* polynomial g provided $g = fh$ for some $h \in k[x_1, \dots, x_n]$
- $k[x_1, \dots, x_n]$ satisfies all field axioms except for existence of multiplicative inverses; commutative ring, $k[x_1, \dots, x_n]$ *polynomial ring*
- Exercises for 1.* **Exercise 1.** \mathbb{F}_2 commutative ring since it's an abelian group under addition, commutative in multiplication, and multiplicative identity exists, namely 1. It is a field since for $1 \neq 0$, the multiplicative identity is 1.
- Exercise 2.**
 - (a)
 - (b)
 - (c)
- 1.2. **Affine Varieties.**
- 1.3. **Parametrizations of Affine Varieties.**
- 1.4. **Ideals.**
- 1.5. **Polynomials of One Variable.**

2. GROEBNER BASES

- 2.1. **Introduction.**

2.2. Orderings on the Monomials in $k[x_1, \dots, x_n]$.

2.3. A Division Algorithm in $k[x_1, \dots, x_n]$.

2.4. Monomial Ideals and Dickson’s Lemma.

2.5. The Hilbert Basis Theorem and Groebner Bases.

2.6. Properties of Groebner Bases.

2.7. Buchberger’s Algorithm.

3. ELIMINATION THEORY

3.1. The Elimination and Extension Theorems.

3.2. The Geometry of Elimination.

4. THE ALGEBRA-GEOMETRY DICTIONARY

4.1. Hilbert’s Nullstellensatz.

4.2. Radical Ideals and the Ideal-Variety Correspondence.

5. POLYNOMIAL AND RATIONAL FUNCTIONS ON A VARIETY

5.1. Polynomial Mappings.

6. ROBOTICS AND AUTOMATIC GEOMETRIC THEOREM PROVING

6.1. Geometric Description of Robots.

Part 2. Reading notes on Cox, Little, O’Shea’s *Using Algebraic Geometry*

Using Algebraic Geometry. David A. Cox. John Little. Donal O’Shea. Second Edition. Springer. 2005. ISBN 0-387-20706-6 QA564.C6883 2004

7. INTRODUCTION

7.1. Polynomials and Ideals. *monomial*

(1) (1.1) $x_1^{\alpha_1} \dots x_n^{\alpha_n}$

total degree of x^α is $\alpha_1 + \dots + \alpha_n \equiv |\alpha|$

field k , $k[x_1 \dots x_n]$ collection of all polynomials in $x_1 \dots x_n$ with coefficients k .

polynomials in $k[x_1 \dots x_n]$ can be added and multiplied as usual, so $k[x_1 \dots x_n]$ has structure of commutative ring (with identity)

however, only nonzero constant polynomials have multiplicative inverses in $k[x_1 \dots x_n]$, so $k[x_1 \dots x_n]$ not a field

however set of rational functions $\{f/g|f,g \in k[x_1 \dots x_n], g \neq 0\}$ is a field, denoted $k(x_1 \dots x_n)$

so

$$f = \sum_{\alpha} c_{\alpha} x^{\alpha}$$

where $c_{\alpha} \in k$

so

$$f \in k[x_1 \dots x_n] = \{f|f = \sum_{\alpha} c_{\alpha} x^{\alpha}, x^{\alpha} = x_1^{\alpha_1} \dots x_n^{\alpha_n}, c_{\alpha} \in k\}$$

f homogeneous if all monomials have same total degrees

polynomial f is homogeneous if all monomials have the *same total degree*

Given a collection of polynomials $f_1 \dots f_s \in k[x_1 \dots x_n]$, we can consider all polynomials which can be built up from these by multiplication by arbitrary polynomials and by taking sums

Definition 2 (1.3). *Let $f_1 \dots f_s \in k[x_1 \dots x_n]$
Let $\langle f_1 \dots f_s \rangle = \{p_1 f_1 + \dots + p_s f_s | p_i \in k[x_1 \dots x_n] \text{ for } i = 1 \dots s\}$*

Exercise 1.

(a) $x^2 = x \cdot (x - y^2) + y \cdot (xy)$

(b)

$$p \cdot (x - y^2) = px - py^2$$

and for $pxy = (py)x$

(c)

$$p(y)(x - y^2) = p(y)x - p(y)y^2 \notin \langle x^2, xy \rangle$$

Exercise 2.

$$\sum_{i=1}^s p_i f_i + \sum_{j=1}^s q_j f_j = \sum_{i=1}^s (p_i + q_i) f_i, \quad p_i + q_i \in k[x_1 \dots x_n]$$

$\langle f_1 \dots f_s \rangle$ closed under sums in $k[x_1 \dots x_n]$

If $f \in \langle f_1 \dots f_s \rangle$,

$p \in k[x_1 \dots x_n]$

$$p \cdot f = p \sum_{i=1}^s q_j f_j = \sum_{i=1}^s p q_j f_j, \quad p q_j \in k[x_1 \dots x_n] \text{ so}$$

$$p \cdot f \in \langle f_1 \dots f_s \rangle$$

Done.

The 2 properties in Ex. 2 are defining properties of ideals in the ring $k[x_1 \dots x_n]$

Definition 3 (1.5). *Let $I \subset k[x_1 \dots x_n]$, $I \neq \emptyset$*

I ideal if

(a) $f + g \in I, \quad \forall f, g \in I$

(b) $pf \in I, \quad \forall f \in I, \text{ arbitrary } p \in k[x_1 \dots x_n]$

Thus $\langle f_1 \dots f_s \rangle$ is an ideal by Ex. 2.

we call it the ideal generated by $f_1 \dots f_s$.

Exercise 3. Suppose \exists ideal J , $f_1 \dots f_s \in J$ s.t. $J \subset \langle f_1 \dots f_s \rangle$

if $f \in \langle f_1 \dots f_s \rangle$, $f = \sum_{i=1}^s p_i f_i$, $p_i \in k[x_1 \dots x_n]$

$\forall i = 1 \dots s$, $p_i f_i \in J$ and so $\sum_{i=1}^s p_i f_i \in J$, by def. of J as an ideal.

$$\langle f_1 \dots f_s \rangle \subseteq J \implies J = \langle f_1 \dots f_s \rangle$$

$\implies \langle f_1 \dots f_s \rangle$ is smallest ideal in $k[x_1 \dots x_n]$ containing $f_1 \dots f_s$

Exercise 4. For $I = \langle f_1 \dots f_s \rangle$

$$J = \langle g_1 \dots g_t \rangle$$

$I = J$ iff $s = t$ and $\forall f \in I, f = \sum_{i=1}^t q_i g_i$ and if $0 = \sum_{i=1}^t q_i g_i, q_i = 0, \forall i = 1 \dots t$, and if $0 = \sum_{i=1}^s p_i f_i, p_i = 0, \forall i = 1 \dots s$

Definition 4 (1.6).

$$\sqrt{I} = \{g \in k[x_1 \dots x_n] | g^m \in I \text{ for some } m \geq 1\}$$

e.g. $x + y \in \sqrt{\langle x^2 + 3xy, 3xy + y^2 \rangle}$
in $\mathbb{Q}[x, y]$ since

$$(x + y)^3 = x(x^2 + 3xy) + y(3xy + y^2) \in \langle x^2 + 3xy, 3xy + y^2 \rangle$$

- (Radical Ideal Property) \forall ideal $I \subset k[x_1 \dots x_n]$, \sqrt{I} ideal, $\sqrt{I} \supset I$
- **(Hilbert basis Thm.)** \forall ideal $I \subset k[x_1 \dots x_n]$
 \exists finite generating set,
i.e. $\exists \{f_1 \dots f_s\} \subset k[x_1 \dots x_n]$ s.t. $I = \langle f_1 \dots f_s \rangle$
- (Division Algorithm in $k[x]$) $\forall f, g \in k[x]$ (EY : in 1 variable)
 $\forall f, g \in k[x]$ (in 1 variable)
 $f = qg + r, \exists!$ quotient q, \exists remainder r

7.2.

7.3. **Gröbner Bases.**

Definition 5 (3.1). *Gröbner basis for $I \equiv G = \{g_1 \dots g_k\} \subset I$ s.t. $\forall f \in I, LT(f)$ divisible by $LT(g_i)$ for some i*

- (Uniqueness of Remainders) let ideal $I \subset k[x_1 \dots x_n]$
division of $f \in k[x_1 \dots x_n]$ by Grö bner basis for I , produces $f = g + r, g \in I$, and no term in r divisible by any element of $LT(I)$

7.4. **Affine Varieties.** affine n -dim. space over k $k^n = \{(a_1 \dots a_n) | a_1 \dots a_n \in k\}$

\forall polynomial $f \in k[x_1 \dots x_n], (a_1 \dots a_n) \in k^n$

$$f : k^n \rightarrow k$$

$$f(a_1 \dots a_n) \text{ s.t. } x_i = a_i \text{ i.e.}$$

if $f = \sum_{\alpha} c_{\alpha} x^{\alpha}$ for $c_{\alpha} \in k$, then

$$f(a_1 \dots a_n) = \sum_{\alpha} c_{\alpha} a^{\alpha} \in k, \text{ where } a^{\alpha} = a_1^{\alpha_1} \dots a_n^{\alpha_n}$$

Definition 6 (4.1). *affine variety $\mathbf{V}(f_1 \dots f_s) = \{(a_1 \dots a_n) | (a_1 \dots a_n) \in k^n, f_1(a_1 \dots a_n) = \dots = f_s(a_1 \dots a_n) = 0\}$
subset $V \subset k^n$ is affine variety if $V = V(f_1 \dots f_s)$ for some $\{f_i\}$, polynomial $f_i \in k[x_1 \dots x_n]$*

- (Equal Ideals Have Equal Varieties) If $\langle f_1 \dots f_s \rangle = \langle g_1 \dots g_t \rangle$ in $k[x_1 \dots x_n]$, then $\mathbf{V}(f_1 \dots f_s) = \mathbf{V}(g_1 \dots g_t)$

so, recap

if $\langle f_1 \dots f_s \rangle = \langle g_1 \dots g_t \rangle$ in $k[x_1 \dots x_n]$,

then $V(f_1 \dots f_s) = V(g_1 \dots g_t)$

Recall Hilbert basis Thm. \forall ideal $I \subset k[x_1 \dots x_n]$

$$I = \langle f_1 \dots f_s \rangle$$

\implies if $I = J$, then $V(I) = V(J)$

think of V defined by I , rather than $f_1 = \dots = f_s = 0$

Exercise 3.

Recall Def. 1.5 Let $I \subset k[x_1 \dots x_n]$

I ideal if $f + g \in I \quad \forall f, g \in I$

$$pf \in I, \quad \forall f \in I \text{ arbitrary } p \in k[x_1 \dots x_n]$$

Let $f, g \in I(V)$

$$(f + g)(a_1 \dots a_n) = f(a_1 \dots a_n) + g(a_1 \dots a_n) = 0 + 0 = 0 \quad f + g \in I(V)$$

$$pf(a_1 \dots a_n) = p(a_1 \dots a_n)f(a_1 \dots a_n) = 0 \quad pf \in I(V)$$

Then $I(V)$ an ideal.

$$V = V(x^2) \text{ in } \mathbb{R}^2$$

$$I = \langle x^2 \rangle \text{ in } \mathbb{R}[x, y], \quad I = \{px^2 | p \in k[x, y]\}$$

$$I \subset I(V), \text{ since } px^2 = 0 \text{ for } x^2 = 0, (0, b), \quad b \in \mathbb{R}$$

But $p(x, y) = x \in I(V)$, as

$$I(V) = \{f \in k[x_1 \dots x_n] | f(a_1 \dots a_n) = 0, \forall (a_1 \dots a_n) \in V\}$$

$$p(0, b) = x = 0$$

But $x \notin I$

Exercise 4. $I \subset \sqrt{I}$

Recall Def. 1.6 $\sqrt{I} = \{g \in k[x_1 \dots x_n] | g^m \in I \text{ for some } m \geq 1\}$

$\forall f \in I, f = f^1, m = 1$, so $f \in \sqrt{I}$, $I \subset \sqrt{I}$

Hilbert basis thm., \forall ideal $I \subset k[x_1 \dots x_n]$ s.t. $I = \langle f_1 \dots f_s \rangle$

$$\left\{ V(I) = \{(a_1 \dots a_n) | (a_1 \dots a_n) \in k^n, f_1(a_1 \dots a_n) = \dots = f_s(a_1 \dots a_n) = 0 \} \right\}$$

$$\mathbf{I}(\mathbf{V}(I)) = \{f \in k[x_1 \dots x_n] | f(a_1 \dots a_n) = 0 \quad \forall (a_1 \dots a_n) \in V(I)\}$$

Let $g \in \sqrt{I}, g^m \in I, g^m = g^{m-1}g$

$$g^m(a_1 \dots a_n) = 0 = g^{m-1}(a_1 \dots a_n)g(a_1 \dots a_n) = 0. \text{ Then } g(a_1 \dots a_n) = 0 \text{ or } g^{m-1}(a_1 \dots a_n) = 0$$

as $g^m \in I$, and $V(I)$ is s.t. $f_1(a_1 \dots a_n) = \dots = f_s(a_1 \dots a_n) = 0$ for $I = \langle f_1 \dots f_s \rangle$

- (Strong Nullstellensatz) if k algebraically closed (e.g. \mathbb{C}), I ideal in $k[x_1 \dots x_n]$, then

$$\mathbf{I}(\mathbf{V}(I)) = \sqrt{I}$$

- (Ideal-variety correspondence) Let k arbitrary field

$$I \subset I(V(I))$$

$$V(I(V)) = V \quad \forall V$$

Additional Exercises for Sec.4. Exercise 6.

8. SOLVING POLYNOMIAL EQUATIONS

8.1.

8.2. **Finite-Dimensional Algebras.** Gröbner basis $G = \{g_1 \dots g_t\}$ of ideal $I \subset k[x_1 \dots x_n]$,

recall def.: Gröbner basis $G = \{g_1 \dots g_t\} \subset I$ of ideal $I, \forall f \in I, LT(f)$ divisible by $LT(g_i)$ for some i

$f \in k[x_1 \dots x_n]$ divide by G produces $f = g + r, g \in I, r$ not divisible by any $LT(I)$ uniqueness of r

$f \in k[x_1 \dots x_n]$ divide by G ,

Recall from Ch. 1, divide $f \in k[x_1 \dots x_n]$ by G , the division algorithm yields

$$(2) \quad (2.1) \quad f = h_1 g_1 + \dots + h_t g_t + \bar{f}^G$$

where remainder \bar{f}^G is a linear combination of monomials $x^{\alpha} \notin \langle LT(I) \rangle$

since Gröbner basis, $f \in I$ iff $\bar{f}^G = 0$

$\forall f \in k[x_1 \dots x_n]$, we have coset $[f] = f + I = \{f + h | h \in I\}$ s.t. $[f] = [g]$ iff $f - g \in I$

We have a 1-to-1 correspondence

remainders \leftrightarrow cosets

$$\overline{f}^G \leftrightarrow [f]$$

algebraic

$$\overline{f}^G + \overline{g}^G \leftrightarrow [f] + [g]$$

$$\overline{\overline{f}^G \cdot \overline{g}^G} \leftrightarrow [f] \cdot [g]$$

$B = \{x^\alpha | x^\alpha \notin \langle \text{LT}(I) \rangle\}$ is a basis of A , basis monomials, standard monomials

20141023 EY's take

$$\forall [f] \in A = k[x_1 \dots x_n]/I, \quad [f] = p_i b_i; \quad b_i \in B = \{x^\alpha | x^\alpha \notin \langle \text{LT}(I) \rangle\}$$

For $I = \langle G \rangle$

$$\text{e.g. } G = \{x^2 + \tfrac{3}{2}xy + \tfrac{1}{2}y^2 - \tfrac{3}{2}x - \tfrac{3}{2}y, xy^2 - x, y^3 - y\}$$

$$\langle \text{LT}(I) \rangle = \langle x^2, xy^2, y^3 \rangle$$

$$\text{e.g. } B = \{1, x, y, xy, y^2\}$$

$$[f] \cdot [g] = [fg]$$

$$\text{e.g. } f = x, g = xy, [fg] = [x^2y]$$

$$\text{now } f = h_1g_1 + \dots + h_tg_t + \overline{f}^G$$

8.3.

8.4. Solving Equations via Eigenvalues and Eigenvectors.

9. RESULTANTS

10. COMPUTATION IN LOCAL RINGS

10.1. Local Rings.

Definition 7 (1.1).

$$k[x_1 \dots x_n]_{\langle x_1 \dots x_n \rangle} \equiv \{ \frac{f}{g} \mid \text{rational functions } \frac{f}{g} \text{ of } x_1 \dots x_n \text{ with } g(p) \neq 0 \text{ at } p \}$$

main properties of $k[x_1 \dots x_n]_{\langle x_1 \dots x_n \rangle}$

Proposition 1 (1.2). *Let $R = k[x_1 \dots x_n]_{\langle x_1 \dots x_n \rangle}$. Then*

(a) *R subring of field of rational functions $k(x_1 \dots x_n) \supset k[x_1 \dots x_n]$*

(b) *Let $M = \langle x_1 \dots x_n \rangle \subset R$ (ideal generated by $x_1 \dots X_n$ in R)*

Then $\forall \frac{f}{g} \in R \setminus M$, $\frac{f}{g}$ unit in R (i.e. multiplicative inverse in R)

(c) *M maximal ideal in R*

Exercise 1. if $p = (a_1 \dots a_n) \in k^n$, $R = \{ \frac{f}{g} \mid f, g \in k[x_1 \dots x_n], g(p) \neq 0 \}$

(a) R subring of field of rational functions $k(x_1 \dots x_n)$

(b) Let M ideal generated by $x_1 - a_1 \dots x_n - a_n$ in R

Then $\forall \frac{f}{g} \in R \setminus M$, $\frac{f}{g}$ unit in R (i.e. multiplicative inverse in R)

(c) M maximal ideal in R

Proof. let $p = (a_1 \dots a_n) \in k^n$

let $g_1(p) \neq 0$, $g_2(p) \neq 0$

$$\begin{aligned} \frac{f_1}{g_1} + \frac{f_2}{g_2} &= \frac{f_1g_2 + f_2g_1}{g_1g_2} & g_1(p)g_2(p) \neq 0 \text{ so } \frac{f_1}{g_1} + \frac{f_2}{g_2} &\in R \\ \frac{f_1}{g_1} \cdot \frac{f_2}{g_2} &= \frac{f_1f_2}{g_1g_2} & g_1(p)g_2(p) \neq 0 \text{ so } \frac{f_1}{g_1} \frac{f_2}{g_2} &\in R \end{aligned}$$

$$f = \frac{f}{1} \in R, \quad \forall f \in k[x_1 \dots x_n], \text{ so } k[x_1 \dots x_n] \subset R$$

□

EY : 20141027, to recap,

Let $V = k^n$

Let $p = (a_1 \dots a_n)$

single pt. $\{p\}$ is (an example of) a variety

$$I(\{p\}) = \{x_1 - a_1 \dots x_n - a_n\} \subset k[x_1 \dots x_n]$$

$$R \equiv k[x_1 \dots x_n]_{\langle x_1 - a_1 \dots x_n - a_n \rangle}$$

$$R = \{ \frac{f}{g} \mid \text{rational function } \frac{f}{g} \text{ of } x_1 \dots x_n, g(p) \neq 0, p = (a_1 \dots a_n) \}$$

Prop. 1.2. properties

(a) R subring of field of rational functions $k(x_1 \dots x_n) \quad k(x_1 \dots x_n) \subset R$

(b) $M = \langle x_1 \dots a_1 \dots x_n - a_n \rangle \subset R$. ideal generated by $x_1 - a_1 \dots x_n - a_n$

Then $\forall \frac{f}{g} \in R \setminus M$, $\frac{f}{g}$ unit in R (\exists multiplicative inverse in R)

(c) M maximal ideal in R .

in R we allow denominators that are not elements of this ideal $I(\{p\})$

Definition 8 (1.3). *local ring is a ring that has exactly 1 maximal ideal*

Proposition 2 (1.4). *ring R with proper ideal $M \subset R$ is local ring if $\forall \frac{f}{g} \in R \setminus M$ is unit in R*

localization Ex. 8, Ex. 9

parametrization

Exercise 2.

$$\begin{aligned} x &= x(t) = \frac{-2t^2}{1+t^2} \\ y &= y(t) = \frac{2t}{1+t^2} \end{aligned}$$

$$\begin{aligned} k[t]_{\langle t \rangle} &= \frac{-2t^2}{1+t^2} \text{ rational function of } t. \quad 1+t^2 \neq 0 \\ \text{if } k &= \mathbb{C} \text{ or } \mathbb{R} \end{aligned}$$

Consider set of convergent power series in n variables

$$(3) \qquad (1.5) \qquad k\{x_1 \dots x_n\} = \{ \sum_{\alpha \in \mathbb{Z}_{\geq 0}^n} c_\alpha x^\alpha \mid c_\alpha \in k, \text{ series converges in some open } U \ni 0 \in k^n \}$$

Consider set $k[[x_1 \dots x_n]]$ of formal power series

$$(4) \qquad (1.6) \qquad k[[x_1 \dots x_n]] = \{ \sum_{\alpha \in \mathbb{Z}_{\geq 0}^n} c_\alpha x^\alpha \mid c_\alpha \in k \} \text{ series need not converge}$$

variety V

$$k[x_1 \dots x_n]/\mathbf{I}(V) \qquad \text{variety } V$$

10.2. Multiplicities and Milnor Numbers. if I ideal in $k[x_1 \dots x_n]$, then denote $Ik[x_1 \dots x_n]_{\langle x_1 \dots x_n \rangle}$ ideal generated by I in larger ring $k[x_1 \dots x_n]_{\langle x_1 \dots x_n \rangle}$

Definition 9 (2.1). *Let I 0-dim. ideal in $k[x_1 \dots x_n]$, so $V(I)$ consists of finitely many pts. in k^n . Assume $(0 \dots 0) \in V(I)$ multiplicity of $(0 \dots 0) \in V(I)$ is*

$$\dim_k k[x_1 \dots x_n]_{\langle x_1 \dots x_n \rangle} / Ik[x_1 \dots x_n]_{\langle x_1 \dots x_n \rangle}$$

generally, if $p = (a_1 \dots a_n) \in V(I)$
multiplicity of p , $m(p) = \dim k[x_1 \dots x_n]_M / Ik[x_1 \dots x_n]_M$

$$\dim k[x_1 \dots x_n]_M / Ik[x_1 \dots x_n]_M$$

localizing $k[x_1 \dots x_n]$ at maximal ideal $M = I(\{p\}) = \langle x_1 - a_1 \dots x_n - a_n \rangle$

11.

12.

13. POLYTOPES, RESULTANTS, AND EQUATIONS

14. POLYHEDRAL REGIONS AND POLYNOMIALS

14.1. Integer Programming. Prop. 1.12.

Suppose 2 customers A, B ship to same location
A: ship 400 kg pallet taking up $2\,m^3$ volume
B: ship 500 kg pallet taking up $3\,m^3$ volume

shipping firm trucks carry up to 3700 kg, up to $20\,m^3$

B’s product more perishable, paying \$ 15 per pallet

A pays \$ 11 per pallet
How many pallets from A, B each in truck to maximize revenues?

$$\begin{aligned} (5) \qquad \qquad \qquad (1.1) \qquad \qquad \qquad & 4A + 5B \leq 37 \\ & 2A + 3B \leq 20 \\ & A, B \in \mathbb{Z}_{\geq 0}^* \end{aligned}$$

maximize $11A + 15B$

integer programming.
max. or min. value of some linear function

$$l(A_1 \dots A_n) = \sum_{i=1}^n c_i A_i$$

on set $(A_1 \dots A_n) \in \mathbb{Z}_{\geq 0}^n$ s.t.

3. Finally, by introducing additional variables; rewrite linear constraint inequalities as equalities. The new variables are called “slack variables”

$$(6) \qquad \qquad \qquad (1.4) \qquad \qquad \qquad a_{ij}A_j = b_i, \quad A_j \in \mathbb{Z}_{\geq 0}$$

introduce indeterminate $z_i, \quad \forall$ equation in (1.4)

$$z_i^{a_{ij}A_j} = z_i^{b_i}$$

m constraints

$$\prod_{i=1}^m z_i^{a_{ij}A_j} = \prod_{i=1}^m z_i^{b_i} = \left(\prod_{i=1}^m z_i^{a_{ij}} \right)^{A_j}$$

Proposition 3 (1.6). *Let k field, define $\varphi : k[w_1 \dots w_n] \rightarrow k[z_1 \dots z_m]$ by*

$$\varphi(w_j) = \prod_{i=1}^m z_i^{a_{ij}} \qquad \forall j = 1 \dots n$$

and

$$\varphi(g(w_1 \dots w_n)) = g(\varphi(w_1) \dots \varphi(w_n))$$

\forall general polynomial $g \in k[w_1 \dots w_n]$

Then $(A_1 \dots A_n)$ integer pt. in feasible region iff $\varphi : w_1^{A_1} \dots w_n^{A_n} \mapsto z_1^{b_1} \dots z_m^{b_m}$

Exercise 3.
Now

$$\begin{aligned} \varphi(w_j) &= \prod_{i=1}^m z_i^{a_{ij}} \\ z_i^{a_{ij}A_j} &= z_i^{b_i} \end{aligned}$$

If $(A_1 \dots A_n)$ an integer pt. in feasible region, $a_{ij}A_j = b_i$

$$z_i^{a_{ij}A_j} = z_i^{b_i} = \prod_{j=1}^n z_i^{a_{ij}A_j} \implies \prod_{j=1}^n \prod_{i=1}^m (z_i^{a_{ij}})^{A_j} = \prod_{i=1}^m z_i^{b_i} = \prod_{j=1}^n \varphi(w_j)^{A_j} = \prod_{j=1}^n \varphi(w_j)^{A_j} = \varphi \left(\prod_{j=1}^n w_j^{A_j} \right) = \prod_{i=1}^m z_i^{b_i}$$

since $\varphi(g(w_1 \dots w_n)) = g(\varphi(w_1) \dots \varphi(w_n))$

If $\varphi : \prod_{j=1}^n w_j^{A_j} \mapsto \prod_{i=1}^m z_i^{b_i}$

$$\varphi \left(\prod_{j=1}^n w_j^{A_j} \right) = \prod_{j=1}^n (\varphi(w_j))^{A_j} = \prod_{i=1}^m z_i^{b_i} = \prod_{j=1}^n \left(\prod_{i=1}^m z_i^{a_{ij}} \right)^{A_j} \implies \prod_{j=1}^n z_i^{a_{ij}A_j} = z_i^{b_i}$$

or $a_{ij}A_j = b_i$. So $(A_1 \dots A_n)$ integer pt.

Exercise 4.

$$\prod_{i=1}^m z_i^{b_i} = \prod_{i=1}^m \prod_{j=1}^n z_i^{a_{ij}A_j} = \prod_{j=1}^n \left(\prod_{i=1}^m z_i^{a_{ij}} \right)^{A_j} = \prod_{j=1}^n \varphi(w_j)^{A_j} = \varphi \left(\prod_{j=1}^n w_j^{A_j} \right)$$

So if given $(b_1 \dots b_m) \in \mathbb{Z}^m$, and for a given a_{ij} , $a_{ij}A_j = b_i$

For $m \leq n$, then a_{ij} is surjective, so $\exists A_j$ s.t. $\prod_{i=1}^m z_i^{b_i} = \varphi \left(\prod_{j=1}^n w_j^{A_j} \right)$

Proposition 4 (1.8). *Suppose $f_1 \dots f_n \in k[z_1 \dots z_m]$ given
Fix monomial order in $k[z_1 \dots z_n, w_1 \dots w_n]$ with elimination property:
 \forall monomial containing 1 of z_i greater than any monomial containing only w_j*

Let \mathcal{G} Gröbner basis for ideal
$$I = \langle f_1 - w_1 \dots f_n - w_n \rangle \subset k[z_1 \dots z_m, w_1 \dots w_n]$$

$\forall f \in k[z_1 \dots z_m]$, let $\overline{f}^{\mathcal{G}}$ be remainder on division of f by \mathcal{G}
Then

- (a) polynomial f s.t. $f \in k[f_1 \dots f_n]$ iff $g = \overline{f}^{\mathcal{G}} \in k[w_1 \dots w_n]$
- (b) if $f \in k[f_1 \dots f_n]$ as in part (a),
 $g = \overline{f}^{\mathcal{G}} \in k[w_1 \dots w_n]$
then $f = g(f_1 \dots f_n)$, giving an expression for f as polynomial in f_j
- (c) if $\forall f_i, f$ monomials, $f \in k[f_1 \dots f_n]$,
then g also a monomial.

14.2. Integer Programming and Combinatorics.

15. ALGEBRAIC CODING THEORY

16. THE BERLEKAMP-MASSEY-SAKATA DECODING ALGORITHM

Gröbner Bases, Martin R. Albrecht of the DTU Crypto Group

Part 3. Algebraic Topology

cf. Bredon (1997) [3]

17. SIMPLICIAL COMPLEXES

cf. pp. 245, from Sec. 21 Simplicial Complexes of Ch. 4 Homology Theory in Bredon (1997) [3]
 $\mathbf{v}_0, \dots \mathbf{v}_n \in \mathbb{R}^\infty$, "affinely independent" if they span an affine n -plane, i.e.

$$\text{if } \left(\sum_{i=0}^n \lambda_i \mathbf{v}_i = 0, \sum_{i=0}^n \lambda_i = 0 \right), \text{ then } \implies \forall \lambda_i = 0$$

If not, then, e.g. $\lambda_0 \neq 0$, assume $\lambda_0 = -1$, and solve the equations to get

$$\mathbf{v}_0 = \sum_{i=1}^n \lambda_i \mathbf{v}_i$$
$$\sum_{i=1}^n \lambda_i = 1$$

i.e. \mathbf{v}_0 is in affine space spanned by $\mathbf{v}_1 \dots \mathbf{v}_n$.

If $\mathbf{v}_0, \dots \mathbf{v}_n$ affinely independent, then

(7)
$$\sigma = (\mathbf{v}_0, \dots \mathbf{v}_n) = \left\{ \sum_{i=0}^n \lambda_i \mathbf{v}_i \mid \sum_{i=0}^n \lambda_i = 1, \lambda_i \geq 0 \right\}$$

is "affine simplex" spanned by \mathbf{v}_i ; also convex hull of \mathbf{v}_i .

$\forall k \leq n$, k -face of σ is any affine simplex of form $(\mathbf{v}_{i_1}, \dots \mathbf{v}_{i_k})$, where vertices all distinct, so are affinely independent.

Definition 10. (geometric) simplicial complex $K :=$ collection of affine simplices s.t.

- (1) $\sigma \in K \implies$ any face of $\sigma \in K$; and
- (2) $\sigma, \tau \in K \implies \sigma \bigcap \tau$ is a face of both σ and τ , or $\sigma \bigcap \tau = \emptyset$

If K simplicial complex, $|K| = \bigcup \{ \sigma \mid \sigma \in K \} \equiv$ "polyhedron" of K

Definition 11 (Def. 21.2 of Bredon (1997) [3]). *polyhedron* $:=$ space X if \exists homeomorphism $h : |K| \xrightarrow{\sim} X$ for some simplicial complex K . h, K is triangulation of X ; (map h , complex K)

Let K finite simplicial complex.
Choose ordering of vertices $\mathbf{v}_0, \mathbf{v}_1 \dots$ of K .
If $\sigma = (\mathbf{v}_{\sigma_0}, \dots \mathbf{v}_{\sigma_n})$ is simplex of K , where $\sigma_0 < \dots < \sigma_n$, then
let $f_\sigma : \Delta_n \rightarrow |K|$ be

$$f_\sigma = [\mathbf{v}_{\sigma_b}, \dots \mathbf{v}_{\sigma_n}]$$

in notation of Def. 1.2. Bredon (1997) [3].
Then this gives CW-complex structure on $|K|$ with f_σ as characteristic maps.

Part 4. Graphs, Finite Graphs

18. GRAPHS, FINITE GRAPHS, TREES

Serre (1980) [4]
cf. Chapter I. Trees and Amalgams, Section 1 Amalgams, Subsection 1.1 Direct limits of Serre (1980) [4]
Let $(G_i)_{i \in I}$, family of groups.
 \forall pair (i, j) , let F_{ij} = set of homomorphisms of G_i into G_j
Want: group $G = \varinjlim G_i$ and

$$\{f_i \mid f_i : G_i \rightarrow G\} \text{ s.t. } f_j \circ f = f_i \quad \forall f \in F_{ij}$$

group G and family $\{f_i\}$ universal in that
(*) if H group, if $\{h_i \mid h_i : G_i \rightarrow H; h_j \circ f = h_i \quad \forall f \in F_{ij}\}$,
then $\exists ! h : G \rightarrow H$ s.t. $h_i = h \circ f_i$
i.e. $\text{Hom}(G, H) \simeq \varprojlim \text{Hom}(G_i, H)$, the inverse limit being taken relative to F_{ij} .
i.e. G direct limit of G_i relative to the F_{ij} .

Proposition 5. $\exists !$ pair G , family $(f_i)_{i \in I}$, i.e. (pair consisting of $G, (f_i)_{i \in I}$, unique up to unique isomorphism.

Proof. Define G by generators and relations.
Take generating family to be disjoint union of those for G_i .
relations - xyz^{-1} where $x, y, z \in G_i, z = xy \in G_i$
 xy^{-1} where $x \in G_i, y \in G_j, y = f(x)$ for at least $f \in F_{ij}$.
Thus, existence of $G, \{f_i\}$.
 G represents functor $H \mapsto \varprojlim \text{Hom}(G_i, H)$.
Thus, uniqueness (also from universal property). □

e.g. groups A, G_1, G_2 , homomorphisms $f_1 : A \rightarrow G_1$.
 $f_2 : A \rightarrow G_2$
 G obtained by amalgamating A in G_1, G_2 by $f_1, f_2 \equiv G_1 *_A G_2$.

1 can have $G = \{1\}$, even though f_1, f_2 non-trivial.

Application: (Van Kampen Thm.)
Let topological space X be covered by open U_1, U_2 .

Suppose $U_1, U_2, U_{12} = U_1 \bigcap U_2$ arcwise connected.

Let basept. $x \in U_{12}$.
Then $\pi_1(X; x)$ obtained by taking 3 groups

$$\pi_1(U_1; x), \pi_1(U_2; x), \pi_1(U_{12}; x)$$

and amalagamating them according to homomorphism

$$\pi_1(U_{12}; x) \rightarrow \pi_1(U_1; x)$$

$$\pi_1(U_{12}; x) \rightarrow \pi_1(U_2; x)$$

Exercise 1. Let homomorphisms $f_1 : A \rightarrow G_1$ amalgam $G = G_1 *_A G_2$.

$$f_2 : A \rightarrow G_2$$

Define subgroups A^n, G_1^n, G_2^n , of A, G_1, G_2 recursively by

$$A^1 = \{1\}$$

$$G_1^1 = \{1\}$$

$$G_2^1 = \{1\}$$

A^n = subgroup of A generated by $f_1^{-1}(G_1^{n-1})$ and $f_2^{-1}(G_2^{n-1})$

G_1^n = subgroup of G_1 generated by $f_1(A^n)$

Let A^∞, G_i^∞ be unions of A^n, G_i^n resp.

Show that f_i defines injection $A/A^\infty \rightarrow G_i/G_i^\infty$.

So the amalgamation is $G \simeq G_1/G_1^\infty *_A/A^\infty G_2/G_2^\infty$.

Take the first induction case (for intuition about the solution).

$$A^2 = \langle f_1^{-1}(G_1^1), f_2^{-1}(G_2^1) \rangle = \langle f_1^{-1}(\{1\}), f_2^{-1}(\{1\}) \rangle$$

$$G_i^2 = f_i(A^2)$$

Let $f_i(a) = f_i(b) \in G_i/G_i^\infty$; $a, b \in A/A^\infty$.

Then since $f_i(a), f_i(b) \in G_i/G_i^\infty$, $f_i(a), f_i(b) \in \{gG_i^\infty | g \in G_i\}$ (quotient is defined to be the set of all left cosets of G_i^∞ , which has to be a normal subgroup for G_i/G_i^∞ to be a quotient group).

Since $a, b \in A/A^\infty$, suppose we take $a, b \in A$.

And suppose we take

$$f_i(a) = f_i(a)G_i^\infty = f_i(a)f_i(A^{n_a}) = f_i(aA^{n_a})$$

$$f_i(b) = f_i(b)G_i^\infty = f_i(b)f_i(A^{n_b}) = f_i(bA^{n_b})$$

Taking f_i^{-1} (recall for group homomorphisms, they map inverse of element of 1st. group to inverse of image of this element).

$aA^{n_a} = bA^{n_b} \in A/A^\infty$ (This is okay as we've "quotiented out A^∞ ; so indeed, they're equal)

cf. Subsection 1.2 Structure of amalgams of Serre (1980) [4]

Suppose given group A , family of groups $(G_i)_{i \in I}$, and, $\forall i \in I$, injective homomorphism $A \rightarrow G_i$.

$*_A G_i \equiv$ direct limit (cf. no. 1.1) of family (A, G_i) with respect to these homomorphisms, call it *sum* (in category theory sense, i.e. product) of G_i with A amalgamated.

e.g. $A = \{1\}$,

$*G_i \equiv$ free product of G_i .

18.0.1. *reduced word*. $\forall i \in I$, choose set S_i of right coset representations of G_i modulo A ,

assume $1 \in S_i$,

$(a, s) \mapsto as$ is bijection of $A \times S_i$ onto G_i ,

$A \times (S_i - \{1\}) \rightarrow G_i - A$ (onto)

Let $\mathbf{i} = (i_1 \dots i_n)$, $n \geq 0$, $i_j \in I$, s.t.

$$(8) \quad i_m \neq i_{m+1} \text{ for } 1 \leq m \leq n-1$$

cf. (T) of Serre (1980) [4].

So *reduced word* m is defined as

$$m = (a; s_1 \dots s_n)$$

where $a \in A, s_1 \in S_{i_1} \dots s_n \in S_{i_n}$, and $s - j \neq 1 \forall j$.

$f \equiv$ canonical homomorphism of A into group $G = *_A G_i$

$f_i \equiv$ canonical homomorphism of G_i into group $G = *_A G_i$

EY : 20170611 (Further explanations, basic examples, from me):

Given $A, \{G_i\}_{i \in I}$, injective (group) homomorphisms $\{f_i : A \rightarrow G_i\}_i$.

$G_i \setminus f_i(A) = \{f_i(A)g | g \in G_i\}$.

Right coset representation of $f_i(A)g \mapsto g$.

e.g. $A, G_1, G_2, f_1 : A \rightarrow G_1$.

$$f_2 : A \rightarrow G_2$$

$$G_1 \setminus f_1(A) = \{f_1(A)g | g \in G_1\}$$

$$G_2 \setminus f_2(A) = \{f_2(A)g | g \in G_2\}$$

$\mathbf{i} = (i_1 \dots i_n)$, $i_j \in I$, $i_m \neq i_{m+1}$ for $1 \leq m \leq n-1$.

Consider (1212...12)

$m = (a; f_1 g_2 f_3 g_4 \dots f_{2n-1} g_{2n})$ where f 's $\in S_1 \subset G_1$, g 's $\in S_2 \subset G_2$.

and so

Definition 12 (reduced word). *reduced word* of type \mathbf{i} , m ,

$$(9) \quad m = (a; s_1 \dots s_n)$$

where $a \in A, s_1 \in S_{i_1}, \dots s_n \in S_{i_n}$, $s_j \neq 1 \quad \forall j$,

$\mathbf{i} = (i_1 \dots i_n)$, $i_j \in I$, s.t. $i_m \neq i_{m+1}$ for $1 \leq m \leq n-1$,

with $S_i = \{g | g \in f_i(A)g \in f_i(A)G_i\}$

Theorem 1 (1 of Serre (1980) [4]). $\forall g \in G, \exists$ sequence \mathbf{i} s.t. $i_m \neq i_{m+1}$ for $1 \leq m \leq n-1$ and *reduced word*

$$m = (a; s_1 \dots s_n)$$

of type \mathbf{i} s.t.

$$g = f(a)f_{i_1}(s_1) \dots f_{i_n}(s_n)$$

Furthermore, \mathbf{i} and m unique.

Remark. Thm. 1 implies $f; f_i$ injective.

Then identify A and G_i with images $f(A), f_i(G_i)$ in G , and reduced decomposition (*) of $g \in G$

$$g = as_1 \dots s_n, \quad a \in A, s_1 \in S_{i_1} - \{1\} \dots s_n \in S_{i_n} - \{1\}$$

Likewise, $G_i \cap G_j = A$ if $i \neq j$.

In particular, $S_i - \{1\}$ pairwise disjoint in G .

Proof. Let $X_i \equiv$ set of reduced words of type \mathbf{i} , $X = \coprod X_i$.

Make G act on X .

In view of universal property of G , sufficient to make $\forall i, G_i$ act,

check action induced on A doesn't depend on i

Suppose then that $i \in I$, and let $Y_i =$ set of reduced words of form $(1; s_1 \dots s_n)$, with $i_1 \neq i$.

EY : 20170611

Recall that

$$S_i = \{g | g \in f_i(A)g \in f_i(A)G_i\}$$

$$A \times S_i \rightarrow G_i \text{ onto}$$

$$A \times (S_i - \{1\}) \rightarrow G_i - A \text{ onto}$$

$$(a, s) \mapsto as \text{ bijection}$$

Let $Y_i =$ set of reduced words of form $(1; s_1 \dots s_n) = \{(1; s_1 \dots s_n) | 1 \in A; s_1 \in S_{i_1} \dots s_n \in S_{i_n}; \mathbf{i} = (i_1 \dots i_n), i_j \in I \text{ s.t. } i_m \neq i_{m+1} \text{ for } 1 \leq m \leq n-1\}$.

$$A \times Y_i \rightarrow X = \coprod_i X_i$$

$$(a, (1; s_1 \dots s_n)) \mapsto (a; s_1 \dots s_n)$$

$$A \times \{S_i - \{1\}\} \times Y_i \rightarrow X$$

$$((a, s), (1; s_1 \dots s_n)) \mapsto (a; s, s_1 \dots s_n)$$

and remember that $X_i =$ set of reduced words of type \mathbf{i} .

It's clear that this yields a bijection $A \times Y_i \bigcup A \times (S_i - \{1\}) \times Y_i \rightarrow X$.

Let $x \in X$. Then $x \in X_{\mathbf{i}}$ for some \mathbf{i} . So x is a reduced word of type \mathbf{i} : $x = (a; s_1 \dots s_n)$. Then clearly $x = (a; s_1 \dots s_n) \mapsto (a, (1; s_1 \dots s_n)) \in A \times Y_i$.

cf. pp. 13, Sec. 2. Trees, 2.1 Graphs of Serre (1980) [4]

Definition 13 (1. of Serre (1980) [4]). ***graph*** $\Gamma = (X, Y, Y \rightarrow X \times X, Y \rightarrow Y)$, where $\text{set } X = \text{vert } \Gamma$
 $\text{set } Y = \text{edge } \Gamma$

$$Y \rightarrow X \times X$$

$$y \mapsto (o(y), t(y))$$

$$Y \rightarrow Y$$

$$y \mapsto \bar{y}$$

s.t. $\forall y \in Y, \bar{\bar{y}} = y, \bar{y} \neq y, o(y) = t(\bar{y})$.
vertex $P \in X$ of Γ .

(oriented) edge $y \in Y, \bar{y} \equiv$ inverse edge.

origin of $y :=$ vertex $o(y) = t(\bar{y})$.

terminus of $y :=$ vertex $t(y) = o(\bar{y})$

extremities of $y := \{o(y), t(y)\}$

If 2 vertices **adjacent**, they're extremities of some edge.

orientation of graph $\Gamma = Y_+ \subset Y = \text{edge } \Gamma$ s.t. $Y = Y_+ \coprod \bar{Y}_+$. It always exists.

oriented graph defined, up to isomorphism, by giving 2 sets X, Y_+ and $Y_+ \rightarrow X \times X$.

corresponding set of edges is $Y = Y_+ \coprod \bar{Y}_+$ where $\bar{Y}_+ \equiv$ copy of Y_+

18.0.2. *Realization of a Graph.* cf. Realization of a Graph in Serre (1980) [4].

Let graph $\Gamma, X = \text{vert } \Gamma, Y = \text{edge } \Gamma$.

topological space $T = X \coprod Y \times [0, 1]$, where X, Y provided with discrete topology.

Let R be finest equivalence relation on T for which

$$(10) \quad \begin{aligned} (y, t) &\equiv (\bar{y}, 1 - t) \\ (y, 0) &\equiv o(y) & \forall y \in Y, \forall t \in [0, 1] \\ (y, 1) &\equiv t(y) \end{aligned}$$

quotient space $\text{real}(\Gamma) = T/R$ is *realization* of graph Γ . (realization is a functor which commutes with direct limits).

Let $n \in \mathbb{Z}^+$. Consider oriented graph of $n+1$ vertices $0, 1, \dots, n$,

orientation given by n edges $[i, i+1], 0 \leq i < n, o([i, i+1]) = i$

$$t([i, i+1]) = i+1$$

For $n \geq 1$,

$(y_1 \dots y_n)$ sequence of edges $y_i = c([i-1, i])$ s.t.

$$t(y_i) = o(y_{i+1}), \quad 1 \leq i < n \text{ determine } c$$

If $P_i = c(i)$,

c is a path from P_0 to P_n , and P_0 and P_n are *extremities of the path* c .

pair of form $(y_i, y_{i+1}) = (y_i, \bar{y}_i)$ in path is **backtracking**.

path (of length $n-2$), from P_0 to P_n given (for $n > 2$) by $(y_1 \dots y_{i-1}, y_{i+2} \dots y_n)$

If \exists path from P to Q in Γ, \exists one without backtracking (by induction)

direct limit $\text{Path}_\infty = \varinjlim \text{Path}_n$ provides notion of infinite path.

□ $\text{Path}_\infty \ni$ infinite sequence (y_1, y_2, \dots) of edges s.t. $t(y_i) = o(y_{i+1}) \quad \forall i \geq 1$.

Definition 15 (connected graph; Def. 3 of Serre (1980) [4]). *graph connected* if \forall 2 vertices, 2 vertices are extremities of at least 1 path.

maximal connected subgraphs (under relation of inclusion) are connected components of graph.

18.0.3. *Circuits.* Let $n \in \mathbb{Z}^+, n \geq 1$.

Consider

set of vertices $\mathbb{Z}/n\mathbb{Z}$, orientation given by n edges $[i, i+1], (i \in \mathbb{Z}/n\mathbb{Z})$ with $o([i, i+1]) = i$

$$t([i, i+1]) = i+1$$

Definition 16 (circuit; Def. 4 of Serre (1980) [4]). *circuit (length n) in graph is subgraph isomorphic to Circ_n .*

i.e. subgraph = path $(y_1 \dots y_n)$, without backtracking, s.t. $P_i = t(y_i), (1 \leq i \leq n)$ distinct, s.t. $P_n = o(y_1)$

$n=1$ case: $\text{Circ}_1, \mathbb{Z}/\mathbb{Z} = \{0\}, 1 \text{ edge}, [0, 1], 0 \in \mathbb{Z}/1\mathbb{Z}, o([0, 1]) = 0$

$$t([0, 1]) = 1$$

Note Circ_1 has automorphism of order 2, which changes its orientation, i.e.

\exists automorphism $\sigma \in \text{Aut}(\text{Circ}_1)$ s.t. $|\sigma| = 2$, i.e. $\sigma^2 = 1$.

loop := circuit of length 1; so loop $\in \text{Circ}_1$.

path $(y_1), P_1 = t(y_1) = o(y_1)$.

$n=2$ case: $\text{Circ}_2, \mathbb{Z}/2\mathbb{Z} = \{0, 1\}, 2 \text{ edges } [0, 1], [1, 2],$

path $(y_1, y_2), (1 \leq i \leq 2), P_1 = t(y_1)$

$$P_2 = t(y_2) = o(y_1)$$

18.1. **Combinatorial graphs.** Let $(X, S) \equiv$ simplicial complex of dim. ≤ 1 , with

$X \equiv$ set

$S \equiv$ set of subsets of X with 1 or 2 elements, containing all the 1-element subsets.

associates with it a graph $\Gamma = (X, \{(P, Q)\})$.

X is its set of vertices.

edges = $\{(P, Q) \in X \times X \text{ s.t. } P \neq Q, \{P, Q\} \in S, \text{ with } \overline{(P, Q)} = (Q, P)$

$$o(P, Q) = P$$

$$t(P, Q) = Q$$

In this graph, 2 edges with same origin and same terminus are equal. This is equivalent to (see following Def.)

Definition 17 (combinatorial; Def. 5 of Serre (1980) [4]). *graph is combinatorial if it has no circuit of length ≤ 2*

Conversely, it's easy to see that
every combinatorial graph Γ derived (up to isomorphism) by construction above from simplicial complex (X, S) , where
 $X = \text{vert}\Gamma$
 $S = \text{set of subset } \{P, Q\} \text{ of } X \text{ s.t. } P \text{ and } Q \text{ either adjacent or equal.}$

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