

Satisfiability Checking - WS 2023/2024

Series 3

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Exercise 1

Consider the following formula:

$$\begin{aligned} c_1 : (\neg b \vee c \vee \neg d) \quad \wedge \quad c_2 : (\neg b \vee d) \quad \wedge \quad c_3 : (\neg c \vee \neg d) \quad \wedge \\ c_4 : (\neg a \vee c) \quad \wedge \quad c_5 : (a \vee b) \quad \wedge \quad c_6 : (\neg a \vee \neg c) \quad \wedge \\ c_7 : (a \vee \neg c \vee \neg d) \end{aligned}$$

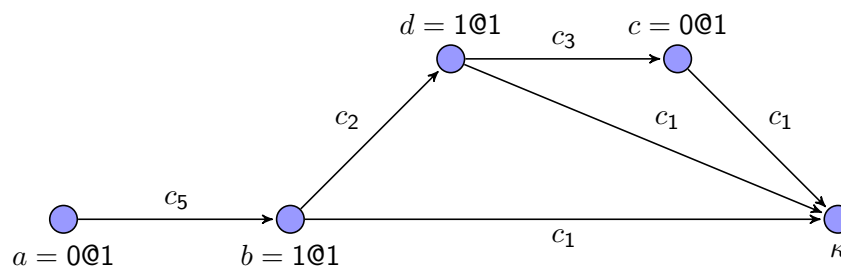
- Give valid initial watch lists for each literal occurring in the formula. Note that there are several possibilities to do so.
- Draw the implication graph for the initial decision using the variable state independent decaying sum (VSIDS) decision heuristic. In the case that the activities of the literals are the same, choose the alphabetically smallest variable first ($a < b < c < d$) and assign it by default to false. Is there a conflict? If so, apply conflict resolution to the conflict which occurred, i.e., use resolution on the conflict clause and the clauses that implied the current assignments, in reverse order back until the first unique implication point is reached. Furthermore, answer the following questions:
 - Which asserting clause is generated?
 - Which assignments would be **undone** using conflict-driven non-chronological backtracking?
 - What does the DPLL algorithm do after backtracking?
- Give the watch lists resulting from the **initial decision**, and also the lists resulting from the **propagation** of the first implication drawn in part b).
BCP阶段
- Is the formula satisfiable? If so, give a satisfying assignment. Otherwise, give an **unsatisfiable core**, i.e., a preferably small subset of the original set of clauses which is unsatisfiable.

Solution:

- Possible initial watch lists for the given formula:

$$\begin{array}{llll} a: & c_5, c_7 & b: & c_5 \\ \neg a: & c_4, c_6 & \neg b: & c_1, c_2 \\ c: & c_1, c_4 & d: & c_2 \\ \neg c: & c_3, c_6, c_7 & \neg d: & c_3 \end{array}$$

- Implication graph for $a = 0@1$:



- Conflicting clause: $cl = c_1 : (\neg b \vee c \vee \neg d)$. Latest assigned literal in cl is c . Variable of c is c . Antecedent of c is c_3 . Resolution of cl and c_3 for c :

$$\frac{cl : (\neg b \vee c \vee \neg d) \quad c_3 : (\neg c \vee \neg d)}{c_8 : (\neg b \vee \neg d)}$$

- $cl = c_8 : (\neg b \vee \neg d)$. Latest assigned literal in cl is $\neg d$. Variable of $\neg d$ is d . Antecedent of d is c_2 . Resolution of cl and c_2 for d :

$$\frac{cl : (\neg b \vee \neg d) \quad c_2 : (\neg b \vee d)}{c_9 : (\neg b)}$$

- The clause c_9 contains exactly one variable that has been assigned in the current decision level and is, hence, asserting (UIP), so we can stop resolution here.

i) c_9 is the result of the conflict resolution and thus, an asserting clause.

ii) All assignments (at decision level 1) get undone.

iii) The solver will assign $b = 0@0$ and propagate this assignment at decision level 0.

c) According to the watch lists of part a) only the clauses for a have to be considered.

- Propagation of $a = 0@1$ in c_5 : implication of assignment $b = 1@1$.
- Propagation of $a = 0@1$ in c_7 : move c_7 from the watch list of a to the list of $\neg d$.

The updated watch lists are:

$$\begin{array}{llll} a: & c_5 & b: & c_5 \\ \neg a: & c_4, c_6 & \neg b: & c_1, c_2 \\ c: & c_1, c_4 & \neg c: & c_3, c_6, c_7 \\ d: & c_2 & \neg d: & c_3, c_7 \end{array}$$

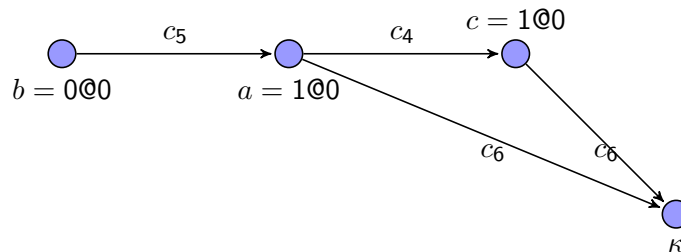
The only implication $b = 1@1$ is propagated in the appropriate clauses as follows.

- Propagation of $b = 1@1$ in c_1 : move c_1 from the watch list of $\neg b$ to the list of $\neg d$.
- Propagation of $b = 1@1$ in c_2 : assignment $d = 1@1$.

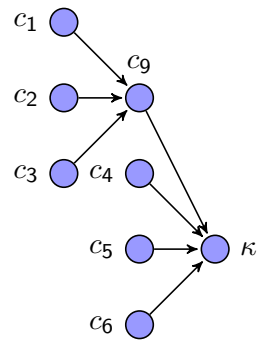
The updated watch lists are:

$$\begin{array}{llll} a: & c_5 & b: & c_5 \\ \neg a: & c_4, c_6 & \neg b: & c_2 \\ c: & c_1, c_4 & \neg c: & c_3, c_6, c_7 \\ d: & c_2 & \neg d: & c_3, c_7, c_1 \end{array}$$

d) The formula is unsatisfiable. When the procedure propagates $b = 0@0$, we get the following implication graph:



Thus the resolution graph for the whole search procedure look as follows:



Unsatisfiable core: $\{c_1, c_2, c_3, c_4, c_5, c_6\}$.

导致empty clause的clause

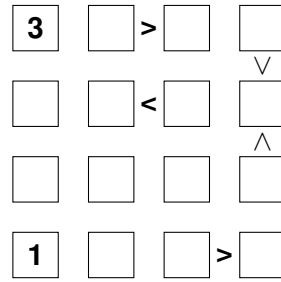
Exercise 2

Consider the game **Unequal** having the following rules:

You have an $n \times n$ square grid; each square may contain a number from 1 to n , and some squares have clue signs between them. Your aim is to fully populate the grid with numbers such that

- each row contains only one occurrence of each number,
- each column contains only one occurrence of each number,
- and all the clue signs are satisfied.

Abbildung 1: An example starting instance of Unequal.



- Formulate the rules of Unequal for an arbitrary square grid in propositional logic. (With no clue signs between the grids and no square fixed to a value as in the example.)
- Formulate the starting instance given by Figure ?? in propositional logic such that the resulting formula is satisfiable iff this instance has a solution according to the game's rules.
- If the formula is not yet in CNF, convert it into CNF.
- Transfer the formula into the standard SAT input format (DIMACS) and let MiniSat solve it. Is the formula satisfiable?

Solution:

a) Variables:

Boolean variables $g_{i,j,k}$ for $i, j, k \in N := \{1, \dots, n\}$. With $g_{i,j,k}$ representing that at the coordinates (i, j) the grid has the number k .

Game rules:

- Each grid has at least one number:

$$\varphi_{grids} := \bigwedge_{i, j \in N} \left(\bigvee_{k \in N} g_{i,j,k} \right)$$

- Each pair of grids within a row has different numbers:

$$\varphi_{rows} := \bigwedge_{\substack{i, j, k, l \in N \\ l > j}} (\neg g_{i,j,k} \vee \neg g_{i,l,k})$$

- Each pair of grids within a column has different numbers:

$$\varphi_{columns} := \bigwedge_{\substack{i, j, k, l \in N \\ l > i}} (\neg g_{i,j,k} \vee \neg g_{l,j,k})$$

We do not need to state that a grid has not more than one number, because it is implied by φ_{rows} and $\varphi_{columns}$, respectively.

b) First, we formulate that the grid at coordinates (r, s) is greater than the grid at coordinates (t, u) .

$$\varphi_{>}(r, s, t, u) := \bigwedge_{\substack{k, l \in N \\ l > k}} (\neg g_{r,s,k} \vee \neg g_{t,u,l})$$

We do not need to cover the case that $l = k$, as φ_{rows} and $\varphi_{columns}$ already exclude this case. Then, the starting instance given by ?? has a solution according to the game's rules, if the following formula is satisfiable.

$$\begin{aligned} \varphi \quad := \quad & \varphi_{grid} \wedge \varphi_{rows} \wedge \varphi_{columns} \\ & \wedge \varphi_{>}(1, 2, 1, 3) \wedge \varphi_{>}(2, 3, 2, 2) \wedge \varphi_{>}(1, 4, 2, 4) \\ & \wedge \varphi_{>}(3, 4, 2, 4) \wedge \varphi_{>}(4, 3, 4, 4) \wedge g_{1,1,3} \wedge g_{4,1,1} \end{aligned}$$

c) φ is already in CNF.

d) The formula is satisfiable.