CSC006P1M: Design and Analysis of Algorithms Lecture 06 (Divide-and-Conquer Paradigm)

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August 29, 2022

Divide-and-Conquer Paradigm

In Divide-and-Conquer paradigm, we solve a problem recursively, applying three steps at each level of recursion.

- Divide The problem is divided into a number of smaller sub-problems that are smaller instances of the same problem.
- 2 Conquer The sub-problems are conquered by solving them recursively. If the sub-problem sizes are small enough, just solve the sub-problems in straightforward manner.
- Ombine The solutions to the sub-problems are combined into the solutions for the original problem.

Decrease-and-Conquer Paradigm

In Decrease-and-Conquer paradigm, we solve a problem recursively, applying three steps at each level of recursion.

- Decrease The problem instance is reduced to smaller instance of the same problem.
- Conquer The problem is solved by solving smaller instance of the problem.
- Extend The solution of smaller solutions is extended to obtain solution to original problem.

This approach is also known as incremental approach.

Divide-and-Conquer vs Decrease-and-Conquer

- If each problem is divided into two or more sub-problems, then the approach is called divide-and-conquer.
- If each problem is divided into one sub-problem, then the approach is called decrease-and-conquer.

$Decrease-and-Conquer\ Paradigm$

Variations of Decrease-and-Conquer

- Decrease by a constant In each iteration of the algorithm, the size of an instance is reduced by the same constant.
- ② Decrease by a constant factor In each iteration of the algorithm, the size of an instance is reduced by the same constant.
- Variable size decrease The size-reduction pattern varies from one iteration of an algorithm to another.

Some Previous Algorithms

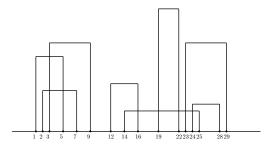
- Evaluating Polynomial (Version 1) Decrease-and-Conquer.
- Evaluating Polynomial (Version 2) Decrease-and-Conquer.
- Evaluating Polynomial (Horner's rule) -Decrease-and-Conquer.
- One-One Mapping Decrease-and-Conquer.
- Gray Code (Version 1) Decrease-and-Conquer.
- MAGNUS problem Decrease-and-Conquer.
- Maximum Consecutive Sum Decrease-and-Conquer.

$\overline{An\ Example}\ of\ a\ Divide-and-Conquer\ Algorithm$

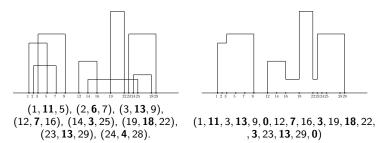
The Skyline Problem

Given the exact locations and shapes of several rectangular buildings in a city, draw the skyline (in two dimensions) of these buildings, eliminating hidden lines.

Example: $(1, \mathbf{11}, 5)$, $(2, \mathbf{6}, 7)$, $(3, \mathbf{13}, 9)$, $(12, \mathbf{7}, 16)$, $(14, \mathbf{3}, 25)$, $(19, \mathbf{18}, 22)$, $(23, \mathbf{13}, 29)$, $(24, \mathbf{4}, 28)$.



12 14 16

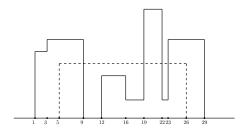




Induction Hypothesis

We know how to solve the problem for n-1 buildings.

And now, we add the n^{th} building. Let B_n be $(5, \mathbf{9}, 26)$.



Previous Configuration:

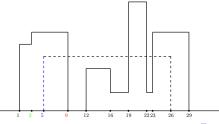
(1, 11, 3, 13, 9, 0, 12, 7, 16, 3, 19, 18, 22, 3, 23, 13, 29, 0)

New Building: $B_n : (5, 9, 26)$.

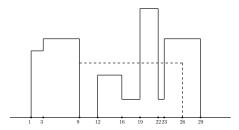
Algorithm:

1. Scan the skyline from left to right and find where the left side of B_n fits. For example, in our case the co-ordinate 5 fits between 3 and 9. ((5, 9, 26);

(1, 11, 3, 13, 9, 0, 12, 7, 16, 3, 19, 18, 22, 3, 23, 13, 29, 0))

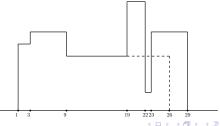


- 2. Scan the horizontal line one after another, and adjust the height of B_n whenever the height of B_n is higher than the existing height.
 - The height of B_n at 5, which is **9**, is covered by the height of the existing skyline from 3 to 9, which is **13**. So, keep it same. (1, 11, 3, 13, 9, 0, 12, 7, 16, 3, 19, 18, 22, 3, 23, 13, 29, 0)

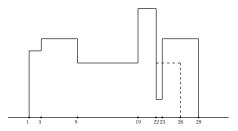


- 3. Scan the horizontal line one after another, and adjust the height of B_n whenever the height of B_n is higher than the existing height.
 - The height of B_n from 9 till 19, which is **9**, is greater than the height of the existing skyline from 9 till 19, which are $\{0,7,3\}$. So, change the configuration.

```
(1, 11, 3, 13, 9, 0, 12, 7, 16, 3, 19, 18, 22, 3, 23, 13, 29, 0)
(1, 11, 3, 13, 9, 9, 19, 18, 22, 3, 23, 13, 29, 0)
```

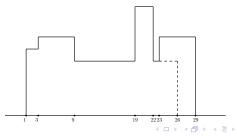


- 4. Scan the horizontal line one after another, and adjust the height of B_n whenever the height of B_n is higher than the existing height.
 - The height of B_n from 19 till 22, which is **9**, is covered by the height of the existing skyline from 19 till 22, which is **18**. So, keep it same. (1, 11, 3, 13, 9, 9, 19, 18, 22, 3, 23, 13, 29, 0)

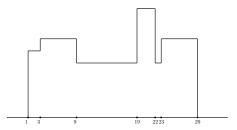


- 5. Scan the horizontal line one after another, and adjust the height of B_n whenever the height of B_n is higher than the existing height.
 - The height of B_n from 22 till 23, which is **9**, is greater than the height of the existing skyline from 22 till 23, which is **3**. So, change the configuration.

```
(1, 11, 3, 13, 9, 9, 19, 18, 22, 3, 23, 13, 29, 0)
(1, 11, 3, 13, 9, 9, 19, 18, 22, 9, 23, 13, 29, 0)
```



- 6. Scan the horizontal line one after another, and adjust the height of B_n whenever the height of B_n is higher than the existing height.
 - The height of B_n from 23 till 26, which is **9**, is covered by the height of the existing skyline from 23 till 26, which is **13**. So, keep it same. (1, 11, 3, 13, 9, 9, 19, 18, 22, 9, 23, 13, 29, 0)



```
How to Combine Pieces? (1, 11, 5), (2, 6, 7), (3, 13, 9), (12, 7, 16), (14, 3, 25), (19, 18, 22), (23, 13, 29), (24, 4, 28). 

• (1, \mathbf{11}, 5), (2, \mathbf{6}, 7) \rightarrow (1, \mathbf{11}, 5, \mathbf{6}, 7)

• (1, \mathbf{11}, 5, \mathbf{6}, 7), (3, \mathbf{13}, 9) \rightarrow (1, \mathbf{11}, 3, \mathbf{13}, 9)

• (1, \mathbf{11}, 3, \mathbf{13}, 9), (12, \mathbf{7}, 16) \rightarrow (1, \mathbf{11}, 3, \mathbf{13}, 9, \mathbf{0}, 12, \mathbf{7}, 16)

:
```

Time Complexity:

$$T(n) = T(n-1) + O(n)$$

$$T(n) = O(n^2)$$

Why?

•
$$T(n) \leq T(n-1) + c \cdot n$$

•
$$T(n-1) \leq T(n-2) + c \cdot (n-1)$$

÷

•
$$T(2) \leq T(1) + c \cdot (2)$$

•
$$T(1) = 0$$

$$T(n) \le c(n+n-1+n-2+\cdots+2) = c(n(n+1)/2-1).$$



Can we do better?

Use Divide-and-Conquer Approach.

Divide them into two sub-problems:

- **1** (1, **11**, 5), (2, **6**, 7), (3, **13**, 9), (12, **7**, 16).
- **2** (14, **3**, 25), (19, **18**, 22), (23, **13**, 29), (24, **4**, 28).

Conquer each sub-problems recursively.

- ② (14, **3**, 25), (19, **18**, 22), (23, **13**, 29), (24, **4**, 28) \rightarrow (14, **3**, 19, **18**, 22, **3**, 23, **13**, 29)

Combine the solutions of each sub-problems to get the final solution.

• (1, **11**, 3, **13**, 9, **0**, 12, **7**, 16, **3**, 19, **18**, 22, **3**, 23, **13**, 29).



Worst-Case Time Complexity:

$$T_w(n) = 2T_w(n/2) + \Theta(n).$$

The Master Theorem

The solution of the recurrence relation $T(n) = aT(n/b) + cn^k$, where a and b are integer constants, $a \ge 1$, $b \ge 2$ and c and k are positive constants, is

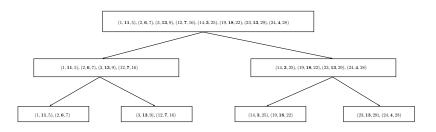
$$T(n) = \begin{cases} \Theta(n^{\log_b a}) & \text{if } a > b^k \\ \Theta(n^k \log_b n) & \text{if } a = b^k \\ \Theta(n^k) & \text{if } a < b^k \end{cases}$$

$$T_w(n) = \Theta(n \lg n).$$

 $T(n) = O(n \lg n).$



Divide



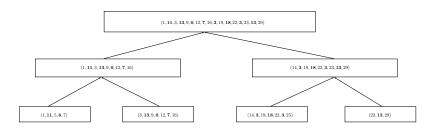
Conquer

(1, 11, 5, 6, 7) (3, 13, 9, 0, 12, 7, 16)

 $(14, \boldsymbol{3}, 19, \boldsymbol{18}, 22, \boldsymbol{3}, 25)$

(23, 13, 29)

Combine



Thank You