

Transactions

Database Management Systems

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Transaction Concept

- Transactions: Collections of operations that form a single logical unit.
- A transaction is a unit of program execution that accesses and possibly updates various data items.
- Examples
 - Transfer money between bank accounts
 - ATM transactions
 - Book flight/ train via online reservations systems
- Application View (SQL View):

Begin transaction

Sequence of SQL statements (database access operations)

End transaction







Transaction Example

Transaction to transfer ₹500 from account R to account S:

- 1. read(*R*)
- 2. R = R 500
- 3. write(*R*)
- 4. read(*S*)
- 5. S = S + 500
- 6. write(*S*)

- Two main issues to deal with:
 - Failures of various kinds, such as hardware failures and system crashes
 - Concurrent execution of multiple transactions



Example of Fund Transfer

Atomicity requirement

- if the transaction fails after step 3 and before step 6, money will be "lost" leading to an inconsistent database state
 - Failure could be due to software or hardware
- the system should ensure that updates of a partially executed transaction are not reflected in the database
 - **1.** read(*R*)
 - 2. R := R 500
 - **3.** write(*R*)
 - **4. read**(*S*)
 - *5. S*:= *S* + 500
 - **6.** write(*S*)



Example of Fund Transfer (Cont.)

Consistency requirement

- In example, the sum of R and S is unchanged by the execution of the
 - transaction.
- In general, consistency requirements include
 - Explicitly specified integrity constraints
 - primary keys and foreign keys
 - **Implicit** integrity constraints
 - sum of balances of all accounts, minus sum of loan amounts must equal value of cash-inhand
 - A transaction must see a consistent database.
 - During transaction execution, the database may be temporarily inconsistent.
 - When the transaction completes successfully the database must be consistent
 - Erroneous transaction logic can lead to inconsistency

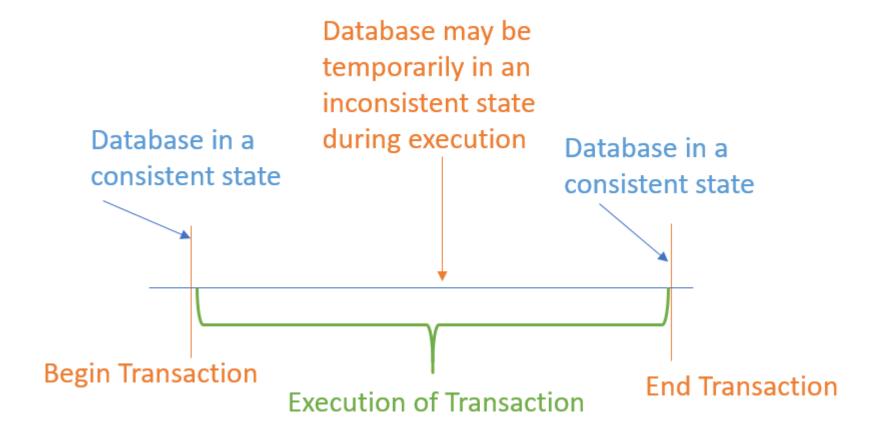


2.
$$R := R - 500$$

- **3.** write(*R*)
- **4. read**(*S*)
- *5. S*:= *S* + 500
- **6.** write(*S*)



Consistency requirement





Example of Fund Transfer (Cont.)

Isolation requirement

if between steps 3 and 6, another transaction T2 is allowed to access the partially updated database, it will see an inconsistent database (the sum R + S will be less than it should be).

```
T1
1. read(R)
2. R := R - 500
3. write(R)
read(R), read(S), print(R+S)
4. read(S)
5. S := S + 500
6. write(S)
```

- Isolation can be ensured trivially by running transactions serially
 - that is, one after the other.
- However, executing multiple transactions concurrently has significant benefits



Example of Fund Transfer (Cont.)

Durability requirement

• once the user has been notified that the transaction has completed (i.e., the transfer of the ₹500 has taken place), the updates to the database by the transaction must persist even if there are software or hardware failures.

- **1.** read(*R*)
- 2. R := R 500
- **3.** write(*R*)
- **4.** read(*S*)
- *5. S*:= *S* + 500
- **6.** write(*S*)



ACID Properties

- To preserve the integrity of data the database system must ensure:
- Atomicity: Either all operations of the transaction are properly reflected in the database or none are.
- **Consistency**: Execution of a transaction in isolation **preserves the consistency** of the database.
- Isolation: Although multiple transactions may execute concurrently, each transaction must be unaware of other concurrently executing transactions.
 - Intermediate transaction results must be hidden from other concurrently executed transactions.
 - That is, for every pair of transactions T_i and T_j , it appears to T_i that either T_j , finished execution before T_i started, or T_j started execution after T_i finished.
- <u>Durability</u>: After a transaction completes successfully, the changes it has made to the database **persist**, even if there are system failures.

