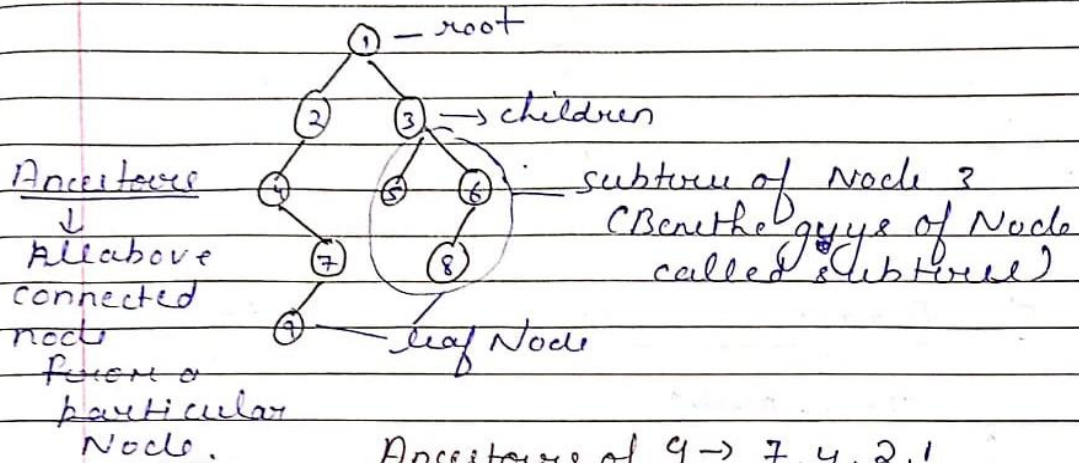


# Tree (BY – Stiver + Love)

Tree → Tree is a Hierarchical data structure.



Ancestors of 9 → 7, 4, 2, 1  
but 3 is not ancestor of 9.

## Types of BT

- 1) Full BT - either has 0 or 2 children.
- 2) Complete BT -
  - i) all levels are completely filled except the last level.
  - The last node has all nodes on left as possible.
- 3) Perfect BT - all leaf nodes are at same level.
- 4) Balanced BT - height of tree at max  $\log(N)$  ∴  $N$  is no. of nodes.

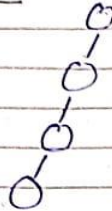
$$\text{If } n=8 \quad \log(8) = \log(2^3) = 3 \log_2 2 = 3$$

Maximum height of tree = 3

used for finding  
complexity

Date: / /  
Page No.

5) Degenerate Tree - (skew Tree like LL)  
if  $n=4$



# Binary Representation of in C++

```

struct Node {
    int data;
    struct Node* left;
    struct Node* Right;
}
  
```

```

Node (int val) {
    data = val;
    left = Right = NULL;
}
  
```

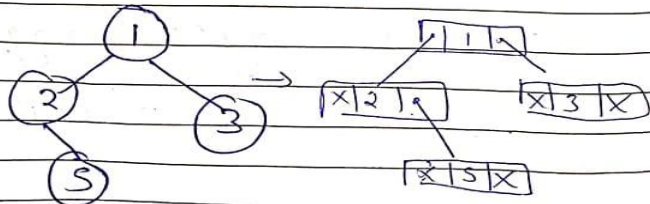
Constructors

Main() {

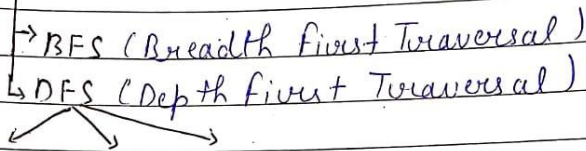
```

    struct Node* root = newnode(1)
    root->left = newnode(2)
    root->Right = newnode(3)
    root->left->right = newnode(5);
}
  
```

in C++  
(NULL)



## Traversal Technique



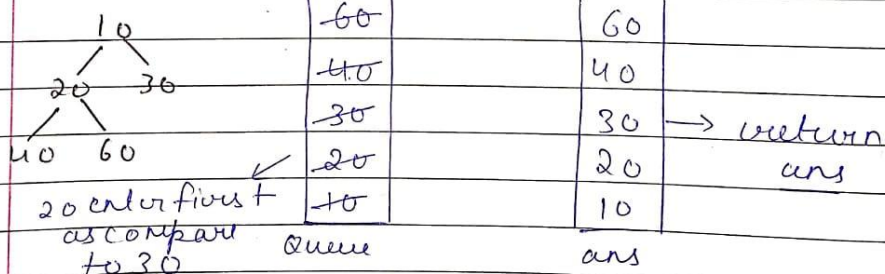
Inorder Preorder Post Order

### # Level Order Traversal / BFS

#### Type 1 - Level Order Traversal

Note - For level order Traversal, we always use queue data structure (FIFO)  
 And to show our level order Traversal in Type 1, we use vector<int> ans.

Ex -



```
vector<int> levelOrder(Node* root)
{
    // To show our level order Traversal
    vector<int> ans ;
    // if root == Null means return an empty ans
    if(root==NULL)return ans ;
    // Note for level order Traversal , we always use Queue Data Structure (FIFO)
    queue<Node*> q;
    // Push root into queue
    q.push(root);

    // if our queue is not empty then check for left or right child
    while(!q.empty())
    {
        // Make a node of Node* (Node pointer) data type , to stor root
        Node* node=q.front();

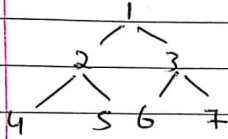
        ans.push_back(node->data);
        if(node->left!=NULL)q.push(node->left);
        if(node->right!=NULL)q.push(node->right);
        q.pop();
    }
    return ans ;
}
```

## Type-2 Level Order Traversal

Here we use queue data structure,

But to show we use vector<vector<int>>ans

Ex



size q = 1  
q = 2  
q = 4

7  
6  
5  
4  
3  
2  
1

Queue  
(q)

1  
[4 5 6 7]  
[2, 3]  
[1]

(ans) vector of vector

\*\*\*

```

vector<vector<int>> levelOrder(TreeNode* root)
{
    // Vector of vector to store level wise value
    vector<vector<int>>ans;
    if(root==NULL) return ans;

    // Create queue because in Level order traversal we use queue
    queue<TreeNode*> q;
    // Queue me root daal denge
    q.push(root); // Queue me root daal denge
    while(!q.empty())
    {
        /* Find queue ka size kyuki hr new ko insert krne ke bdd hum lined up in queue
        lenge , phle wali value queue se bahar ho jaegiii */

        int size=q.size();
        // Level defined is lie because they store lined up element
        vector<int>level;
        for(int i=0;i<size;i++)
        {
            // first in first out , means first value ko node me store krleenge or queue
            // se pop krdenenge
            TreeNode* node=q.front();
            q.pop();

            if(node->left!=NULL)q.push(node->left);
            if(node->right!=NULL)q.push(node->right);

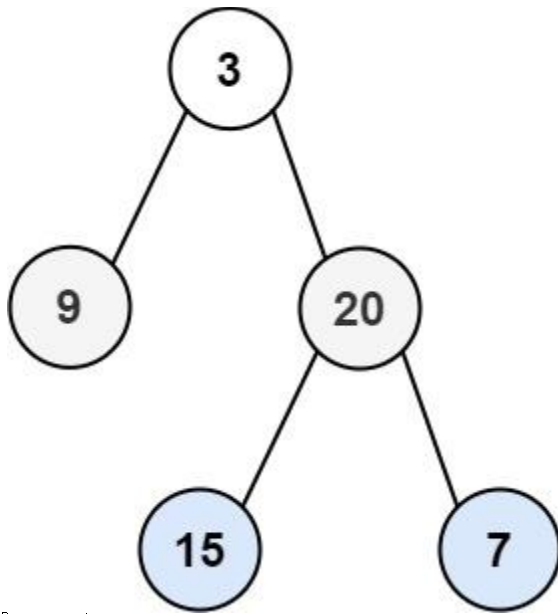
            // level - store lined up value node ki sbhi values ko hum level me daal denge
            level.push_back(node->val);
            // Then pop hone ke bdd hum usko ans me store krdenge level - lined up value ko
        }
        // Phle hum level me pushed up krenge , thn ans vector me level koo
        ans.push_back(level);
    }
    return ans;
}
  
```



## \*\*\* Reverse Level Order Traversal:

Given a binary tree of size N, find its reverse level order traversal. ie- the traversal must begin from the last level.

**LeeTcode** - <https://leetcode.com/problems/binary-tree-level-order-traversal-ii/description/>



**Input:** root = [3,9,20,null,null,15,7]

**Output:** [[15,7],[9,20],[3]]

```
class Solution {
public:
    vector<vector<int>> levelOrderBottom(TreeNode* root)
    {
        vector<vector<int>>result;
        if(root==NULL)return result;
        queue<TreeNode*>q;

        q.push(root);

        while(!q.empty())
        {
            int n = q.size();

            vector<int>level;
            for(int i=0;i<n;i++)
            {
                TreeNode* temp = q.front();
                level.push_back(temp->val);
                if(temp->left!=NULL)q.push(temp->left);
                if(temp->right!=NULL)q.push(temp->right);
                q.pop();
            }
            result.push_back(level);
        }
        reverse(result.begin(), result.end());
        return result;
    }
};
```

```

        }
        result.push_back(level);
    }

    reverse(result.begin(), result.end());

    return result;
}
};

```

**Link :** <https://practice.geeksforgeeks.org/problems/reverse-level-order-traversal/1>

## GFG CODE

```

vector<int> reverseLevelOrder(Node *root)
{
    vector<int> ans;
    if(root == NULL) return ans;

    queue<Node *> q;
    q.push(root);

    while(!q.empty())
    {
        // First for root
        Node* t=q.front();
        ans.push_back(t->data);
        q.pop();
        // Then access right first as compare to left , because hume back se
        // traverse krna hai
        if(t->right) q.push(t->right);
        if(t->left) q.push(t->left);
    }
    // Reverse the vector ans
    reverse(ans.begin(), ans.end());
    return ans;
}

```

### \*\*\* Count Maximum No Of Nodes :

Given an integer **i**. Print the **maximum number of nodes** on level **i** of a binary tree.

**Link:**[https://practice.geeksforgeeks.org/problems/introduction-to-trees/1?utm\\_source=youtube&utm\\_medium=collab\\_striver\\_ytdescription&utm\\_campaign=introduction-to-trees](https://practice.geeksforgeeks.org/problems/introduction-to-trees/1?utm_source=youtube&utm_medium=collab_striver_ytdescription&utm_campaign=introduction-to-trees)

#### CODE

```
int countNodes(int i)
{
    if(i==1)
        return 1;
    if(i>1)
    {
        int ans = pow(2,i-1);    bcoz level starting from 2^0 = 1
                                   2^1 = 2
                                   2^2 = 4

        return ans;
    }
}
```

### DFS Traversal :

#### 1 ) Preorder Traversal – (Root,Left,Right)

**Reccursion Code :** [https://practice.geeksforgeeks.org/problems/preorder-traversal/1?utm\\_source=youtube&utm\\_medium=collab\\_striver\\_ytdescription&utm\\_campaign=preorder-traversal](https://practice.geeksforgeeks.org/problems/preorder-traversal/1?utm_source=youtube&utm_medium=collab_striver_ytdescription&utm_campaign=preorder-traversal)

```
vector<int>v;
vector<int> preOrder(Node* root)
{
    if(root!=NULL)
    {
        v.push_back(root->data); // push in v
        preOrder(root->left);
    }
}
```

```

        preOrder(root->right);
    }
    return v;
}
vector<int> preorder(Node* root)
{
    vector<int>res=preOrder(root);
    // we need to clear v , because here we store all elemens from v to res
    v.clear();
    return res;
}

```

**Iterative Code** : <https://leetcode.com/problems/binary-tree-preorder-traversal/submissions/866242110/>

```

vector<int> preorderTraversal(TreeNode* root)
{
    // vector ans lia jisme print krna hai value koo
    vector<int>ans;
    if(root==NULL) return ans;
    // For preorder traversal we use only stack We put root node into stack
    stack<TreeNode*>st;
    st.push(root);

    while(!st.empty())
    {
        // /Humne value ko root me daal kr usko pop kr lia
        root=st.top();
        st.pop();

        ans.push_back(root->val);
        if(root->right!=NULL)st.push(root->right);
        if(root->left!=NULL)st.push(root->left);
    }
    return ans;
}

```

## 2 ) Inorder Traversal – (Left,Root,Right)

**Reccursion Code** : <https://leetcode.com/problems/binary-tree-inorder-traversal/submissions/>

```

// Funtion To find Inorder Traversal
void inorder(TreeNode* root, vector<int>& ans){

    if(root == NULL){
        return;
    }
    inorder(root->left, ans); //left subtree
    ans.push_back(root->val); //pushing value of root to ans
}

```



```

        inorder(root->right, ans); //right subtree
    }

    vector<int> inorderTraversal(TreeNode* root) {
        vector<int> ans;
        inorder(root, ans);
        return ans;
    }

```

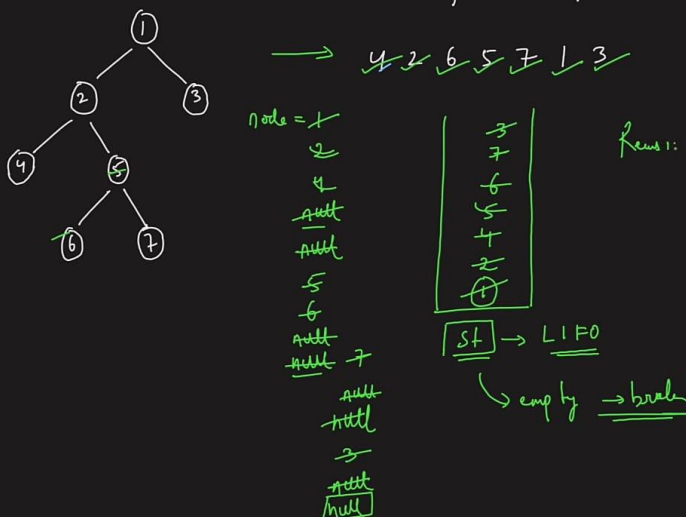
## Iterative Code :

```

vector<int> inorderTraversal(TreeNode* root)
{
    stack<TreeNode*> st;
    TreeNode* node=root;
    vector<int> inorder;

    while (true)
    {
        // If node ke left part not NULL so we goes to left part
        if(node!=NULL)
        {
            st.push(node);
            node = node->left;
        }
        else
        {
            if(st.empty()==true)break; // node ka left or right both null
            node=st.top();
            st.pop();
            inorder.push_back(node->val);
            node = node->right;
        }
    }
    return inorder;
}

```



### 3 ) PostOrder Traversal – (Left, Right, Root)

#### Reccursion Code :

```
// vector<int> &ans - Means i give a address of my ans vector
void recursive(TreeNode* root, vector<int> &ans) {
    if(root) {
        recursive(root->left, ans);
        recursive(root->right, ans);
        ans.push_back(root->val);
    }
}

vector<int> postorderTraversal(TreeNode* root) {
    vector<int> ans;
    recursive(root, ans);
    return ans;
}
```

#### Iterative Code : Using 1 Stack

```
vector<int> postorderTraversal(TreeNode* root)
{
    vector<int> ans;
    TreeNode* curr = root; // Curr me root
    stack<TreeNode*> s;
    if(root==NULL)
    {
        return ans;
    }
    //
    while(curr!=NULL || !s.empty())
    {
        if(curr!=NULL) // if curr null nhi hai to we go on left left
        {
            s.push(curr);
            curr = curr->left;
        }

        else
        {
            // When curr left goel null , so now we goona move on right

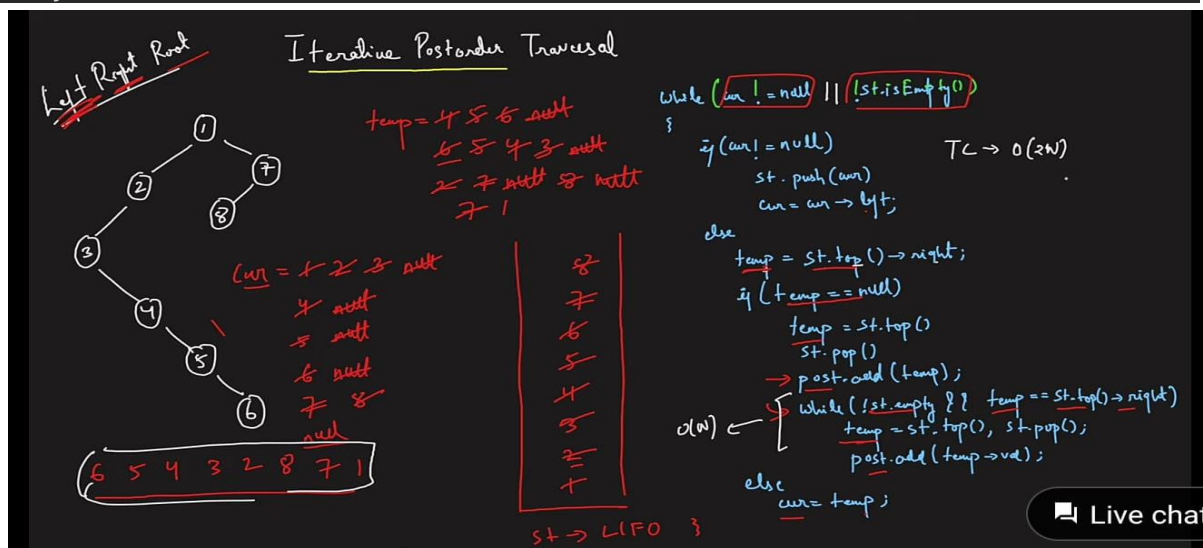
```

```

        // point to second last on right element
        TreeNode* temp = s.top()->right;
        // But if right also null
        if(temp == NULL){
            temp = s.top();
            s.pop();
            ans.push_back(temp->val);
        }
        /* Actually element add hone ke bdd stack se pop ho jaate hai then , hum temp
        or stack me top elemt ka compare kreng if both equal , then we temp ko pop
        and push into vector */

        while(!s.empty() && temp == s.top()->right) {
            temp = s.top();
            s.pop();
            ans.push_back(temp->val);
        }
    }
    else{
        curr = temp;
    }
}
return ans;
}

```



## Iterative Code : Using 2 Stack

```

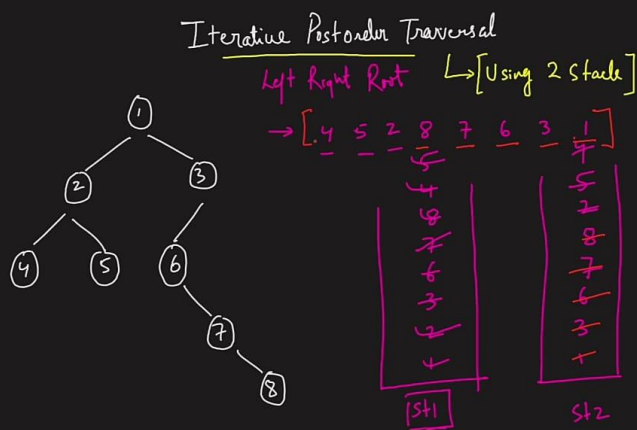
vector<int> postorderTraversal(TreeNode* root)
{
    vector<int> ans;
    if(root == NULL) return ans;
    stack<TreeNode*> st1, st2;
}

```

```

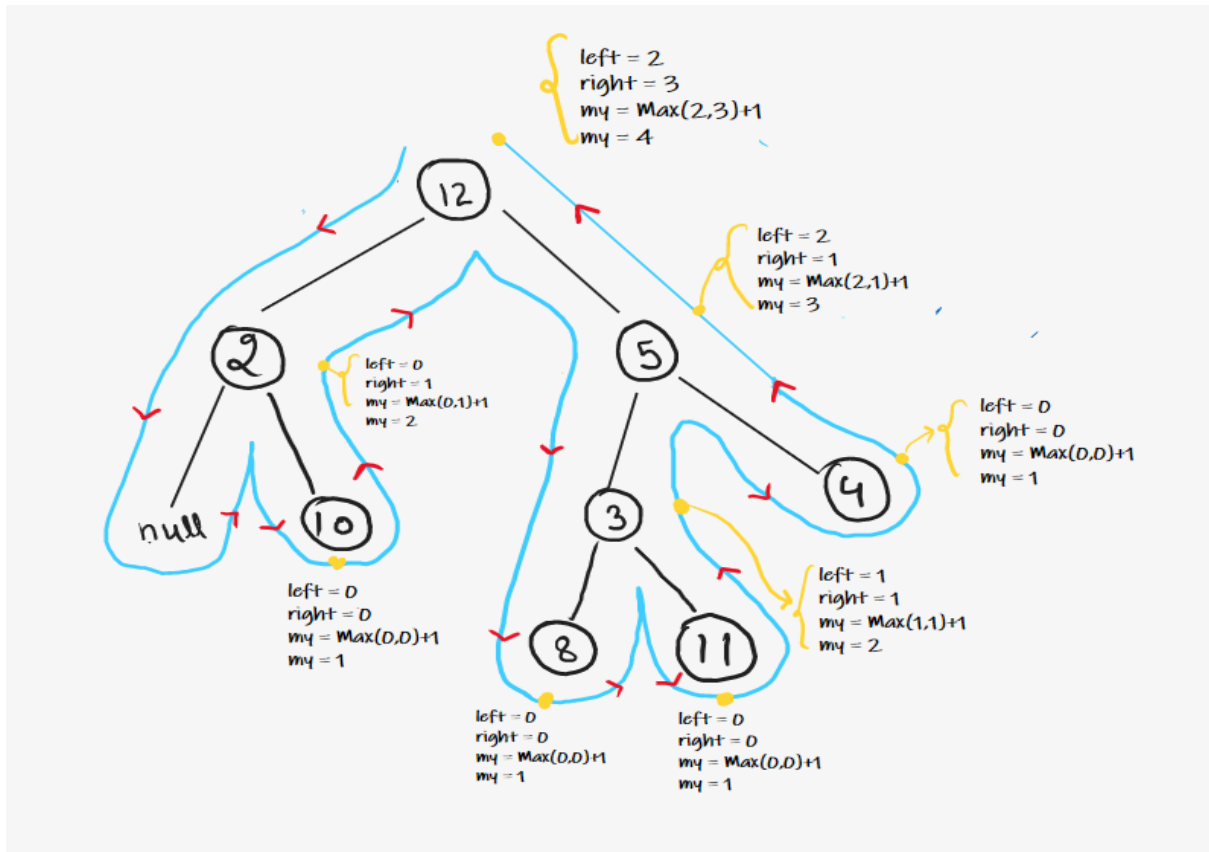
st1.push(root);
while(!st1.empty())
{
    root=st1.top();
    st1.pop();
    st2.push(root);
    // First insert into from left
    if(root->left!=NULL)
    {
        st1.push(root->left);
    }
    // Second insert into from right
    if(root->right!=NULL)
    {
        st1.push(root->right);
    }
}
// st2 stack ke elements ko ek ek krke ans me daalte rhenge
while(!st2.empty())
{
    ans.push_back(st2.top()->val);
    st2.pop();
}
return ans;
}

```



# Problems

**Q1- Maximum Depth Of Binary Tree** - A binary tree's **maximum depth** is the number of nodes along the longest path from the root node down to the farthest leaf node.



## Code - Using Recursion

```
int maxDepth(TreeNode* root)
{
    if(root==NULL) return 0;

    int leftheight = maxDepth(root->left);
    int rightheight = maxDepth(root->right);
    int totalheight = max(leftheight,rightheight)+1;
    return totalheight;
}
```

## Intuition + Approach: Using LEVEL ORDER TRAVERSAL

If we observe carefully, the depth of the Binary Tree is the number of levels in the binary tree. So, if we simply do a level order traversal on the binary tree and keep a count of the number of levels, it will be our answer.

In this example, if we start traversing the tree level-wise, then we can reach at max Level 4, so our answer is 4. Because the maximum depth we can achieve is indicated by the last level at which we can travel.

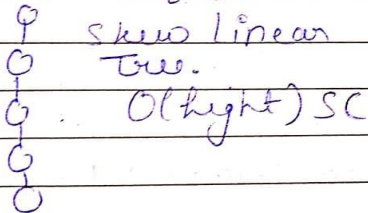
## Maximum Depth in Binary Tree

This problem is solved by using Recursion as well as Level order Traversal.

It does not use any space complexity, but it uses Auxiliary stack space. which after execution is eventually removed from stack space.

Auxiliary Space (compl - exity)  $O(\text{height})$

Worst Case if given



which does not use Recursion, but it uses queue data structure

It takes  $SC \rightarrow O(N)$

in worst case.

And worst case if given

