

# # Coin Change Problem

if arr = 

1	2	3
0	1	2

 target = 5

you have to find minimum no. of element  $\rightarrow$  to reach the target sum.

$\{1, 1, 1, 1, 1\} \rightarrow 5$   
 $\{1, 2, 1, 1\} \rightarrow 4$

$\{2, 2, 1\} \rightarrow 3$

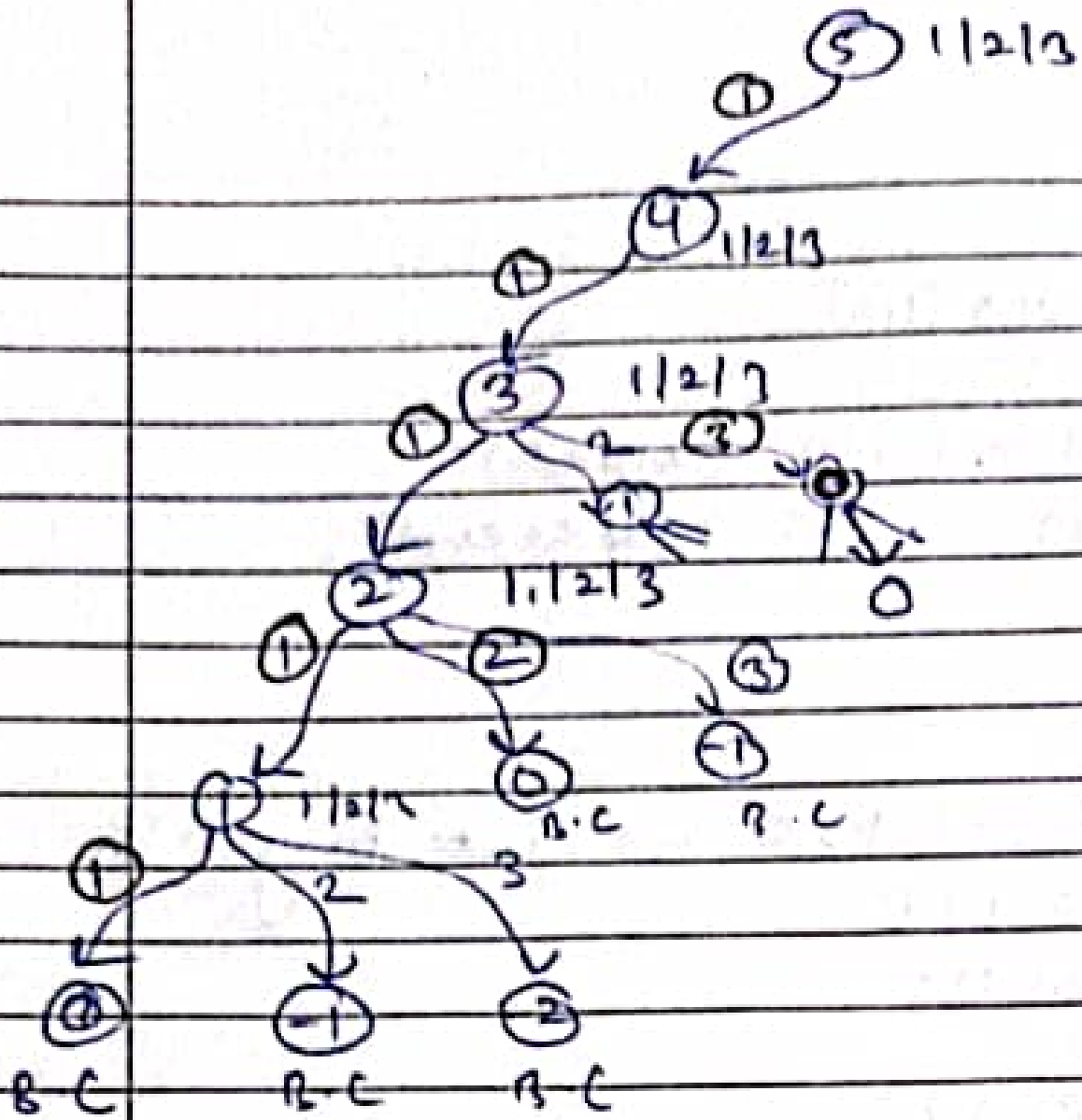
$\{3, 1, 1\} \rightarrow 3$

$\{3, 2\} \rightarrow 2$

$\rightarrow$  minimum no. of element to reach target sum.

Note Target Reach krne ke we have 2 way

- 1) target  $\rightarrow 0$  (if target  $= 0$ ) means reach target
- 2)  $0 \rightarrow 5$  (target)  $\rightarrow$  target reach hojaga, but here use Additional Variable, joki 5 ko store krega, jisse hum 0 me update krke 5 ko reach krve.
- original value ko decrease krega 0 se jisme ke lie
- original ko store krke comparison

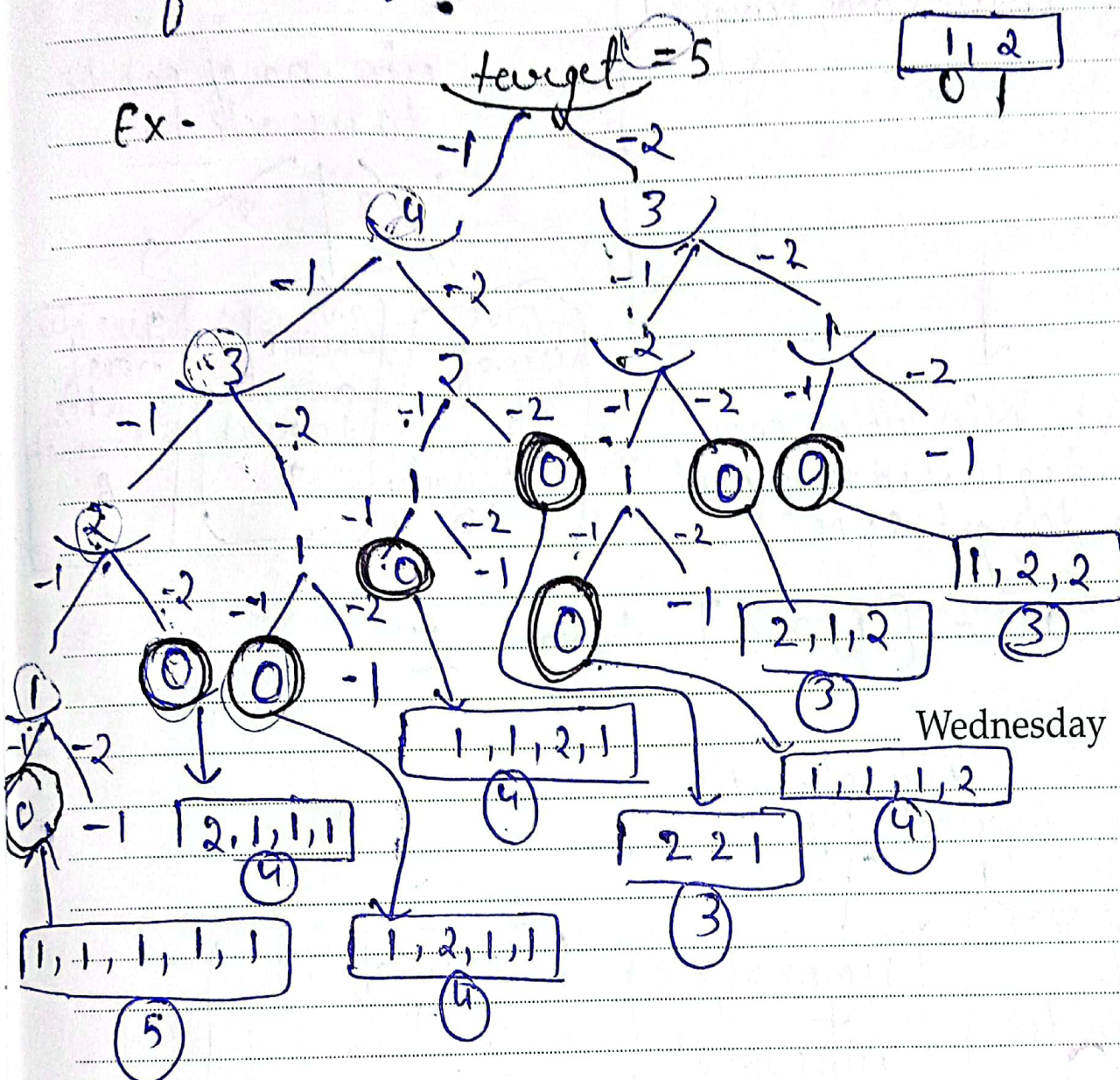




Sun	Mon	Tue	Wed	Thu	Fri	Sat
1	2	3	4	5	6	7
8	9	10	11	12	13	14
15	16	17	18	19	20	21
22	23	24	25	26	27	28
29	30	31				

if both equal means, we reach target

Ex -



## Wednesday

20



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Thursday

Sun	Mon	Tue	Wed	Thu	Fri	Sat
6	7	8	9	10	11	12
13	14	15	16	17	18	19
20	21	22	23	24	25	26
27	28	29	30	31		

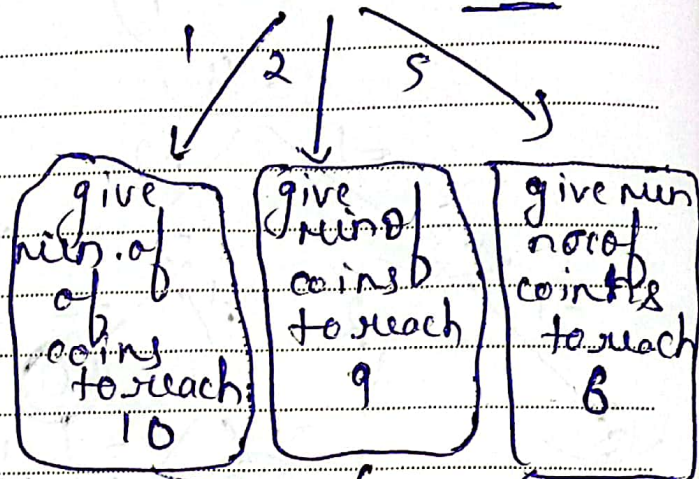
```
int solve(arr, target)
{

```

give min. no. of coins required to reach the target sum

arr = [1, 2, 5] target = 11

min. no. of coins to reach 11

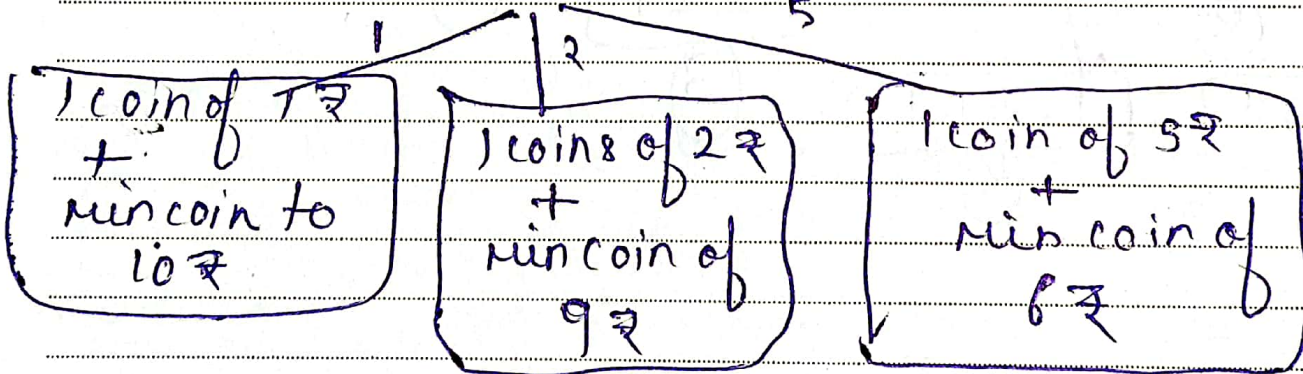


ge sb hure function design hai

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Friday

Min. of coins to reach 11







Cut into segments

ip  $\rightarrow$  N - odd length

Maximum no of segments but only with x, y, z

eg. N = 7  
x = 5  
y = 2  
z = 2

```
int main() {
    int n = 7;
    int x = 5;
    int y = 2;
    int z = 2;
```

```
    int ans = solve(n, x, y, z);
    cout << "Ans is " << ans << endl;
    return 0;
}
```

```
int solve(int n, int x, int y, int z) {
    if (n == 0) {
        return 0;
    }
    ? int ans = 0;
    if (n - x >= 0) {
        ans = solve(n - x, x, y, z) + 1;
    }
    int b = 0;
    if (n - y >= 0) {
        b = solve(n - y, x, y, z) + 1;
    }
}
```

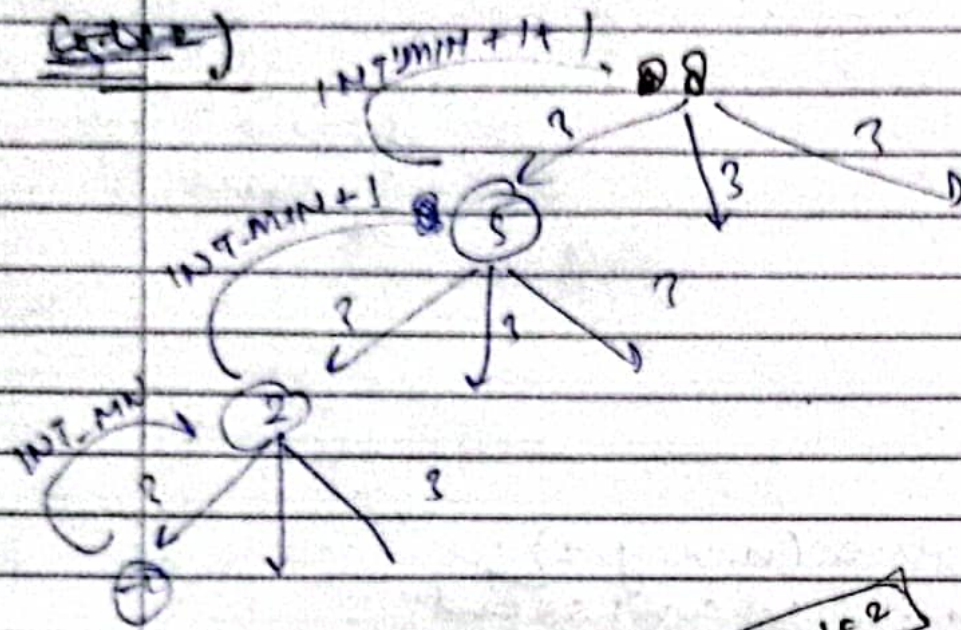
if (ans < 0)  
ans = 0;



```

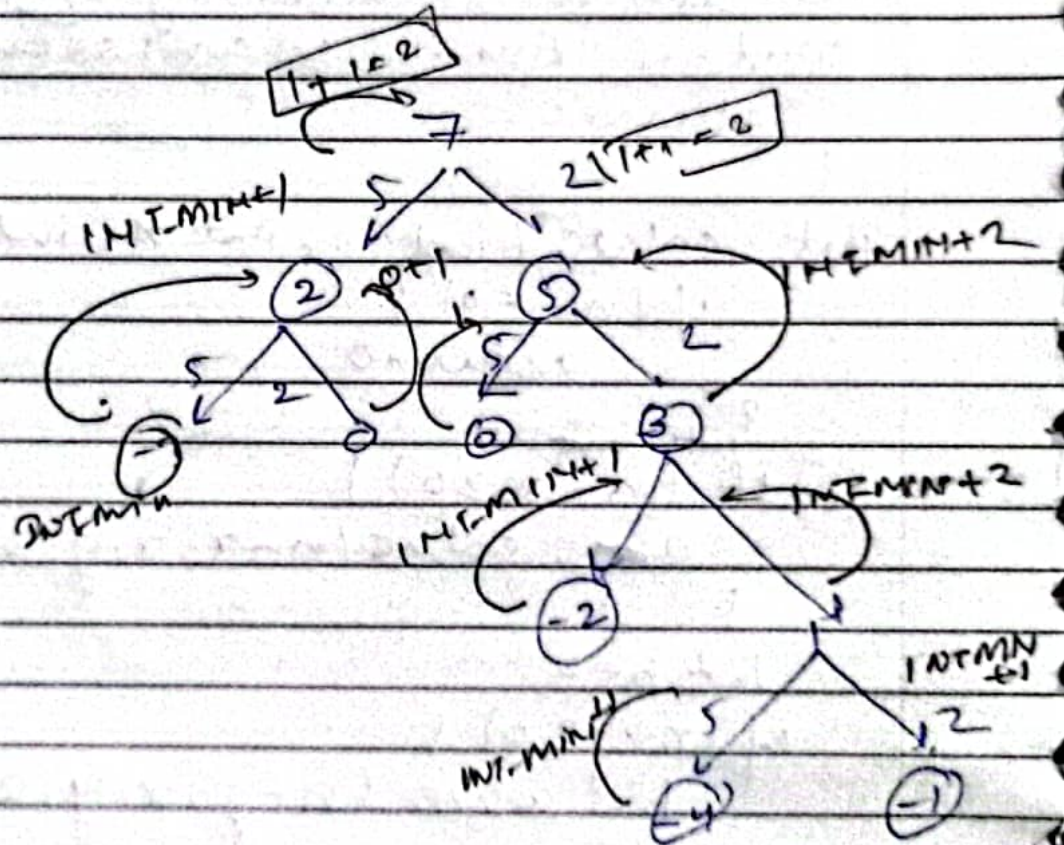
int c = 0;
if (n-2 >= 0) {
    c = solve(n-2, x, y, 2) + 1;
}
int ans = max(a, max(b, c));
return ans;
}

```

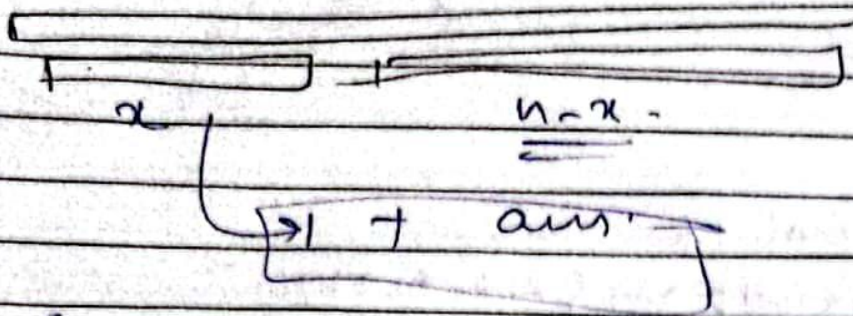


eg.

$n=2$   
 $x=5$   
 $y=2$   
 $z=2$







Reason of  $1 + ans$

3) Maximum sum of non adjacent elements

i/p. 

2	1	4	9
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Returns the max sum of subsequence in which no two elements are adjacent

$$\left. \begin{array}{l} 2 + 4 = 6 \\ 1 + 9 = 10 \\ 2 + 9 = 11 \end{array} \right\}$$

$i = 0$   

2	4	4	9
---	---	---	---

Question me Sub sequence ki bt kia hai mtlb sub sequence me element ka preset or absent hota hai

inc  $i = i + 2$   

2	1	4	9
---	---	---	---

exc  $i$   

2	1	4	9
---	---	---	---

means Use Include / Exclude Technique

inc  

2	1	4	9	$i$
---	---	---	---	-----

exc  $i$   

2	1	4	9
---	---	---	---

B.C.