

Mission 2: Shhh... It's a Secret (in the Cache)

Aakash Tarang

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1 Introduction

This will be my solution report for the second week assignment of the learner's space Hardware Security course.

The following describe my implementation details

2 Creating a shared memory address

I achieved this using the same way as the video tutorial discussed

```
1  const int SIZE = 4096; //size of shared memory segment
2  const char *name = "OS"; //name of shared memory segment
3  int fd; //file descriptor id, similar to an open file
4  char *shared_mem;
5  fd = shm_open (name, O_CREAT|O_RDWR, 0666) ;
6  ftruncate(fd, SIZE);
7  shared_mem = (char*)mmap (0, SIZE, PROT_READ|PROT_WRITE,
    MAP_SHARED, fd, 0);
```

3 Cache-Based Covert Channel Communication

The following functions implement a covert communication channel by exploiting CPU cache timing side effects. The sender modulates cache states to transmit data, while the receiver infers bits by measuring memory access latency.

3.1 Function send_bit(char* addr, int bit)

This function transmits a single bit by manipulating the cache state of a shared memory location (addr).

- If bit == 1:
 - Calls `maccess(addr)` to bring the memory line into the cache.
 - The receiver observes a **fast access** (cache hit), inferring a logical '1'.
- If bit == 0:

- Skips `maccess(addr)`, leaving the cache line evicted.
- The receiver observes a **slow access** (cache miss), inferring a logical ‘0’.
- `flush(addr)` is done when the bit is zero, so that the latency measurement by the receiver program indicates longer time interval. The receiver program also contains a flush which prepares the cache line for the next bit.
- `usleep(BIT_TIME_US)` enforces synchronization between sender and receiver by introducing a fixed delay.

3.2 Function `send_byte(char* addr, char byte)`

This function transmits a full byte by iteratively calling `send_bit()` for each bit in `byte`.

- Iterates over each of the 8 bits (LSB to MSB) in `byte`.
- Extracts the current bit using `(byte >> i) & 1`.
- Calls `send_bit(addr, bit)` to transmit the bit via cache state modulation.

3.3 Receiver Operation

A receiver reconstructs the transmitted data by:

- Monitoring `addr` and measuring access latency.
- Interpreting **fast accesses** (cached) as ‘1’ and **slow accesses** (uncached) as ‘0’.
- Reassembling received bits into bytes.

3.4 Security Implications

This method is a classic example of a **timing-based side-channel attack**, with applications in:

- Exfiltrating data across privilege boundaries (e.g., Spectre attacks).
- Covert inter-process communication (e.g., malware communication).

Mitigations include cache partitioning, constant-time programming, and aggressive cache flushing.

4 Synchronization Pattern

A sync pattern is described at the beginning of the `sender.c` file

```
1 #define SYNC_PATTERN 0b10101010
```

This sync pattern is transmitted first as, it’s used as a synchronization signal so the receiver knows when the actual message starts. Without this, the receiver might misinterpret noise or timing mismatches.

5 Transmit the Message

```
1
2 for (int i = 0; i < msg_size; i++) {
3     send_byte(&shared_mem[OFFSET], msg[i]);
4 }
```

Here, the sender loops through each character in the message stored in `msg[]`. Each character is a byte, and it gets sent using the same `send_byte()` function. That function breaks the byte into 8 bits and sends each bit using Flush+Reload.

So if your message was "Hi", this loop would send the ASCII value of 'H', then 'i', each broken into bits. and we finally end the message with the following flag.

6 Reciever Code

Bit Detection using Flush+Reload

The receiver detects transmitted bits by leveraging cache access timing differences via the Flush+Reload side-channel. The core of this detection logic is implemented in the `detect_bit` and `receive_byte` functions:

```
1 int detect_bit(char *addr) {
2     usleep(BIT_TIME_US / 2);
3     unsigned long long start = rdtsc();
4     maccess(addr);
5     unsigned long long end = rdtsc();
6     flush(addr); // flush for next bit
7     return (end - start) < THRESHOLD ? 1 : 0;
8 }
```

Explanation:

- `usleep(BIT_TIME_US/2)`: Waits for a fixed bit interval to align with sender's transmission.
- `rdtsc()` before and after `maccess(addr)` measures access latency.
- If the memory is in cache (access is fast), it indicates the sender accessed it — representing a bit 1.
- If not in cache (access is slow), it indicates the sender did not access it — representing a bit 0.
- `flush(addr)`: Evicts the memory address from all cache levels for the next bit.

The `receive_byte` function assembles 8 such bits into a byte:

```
1 char receive_byte(char *addr) {
2     char byte = 0;
3     for (int i = 0; i < 8; i++) {
4         int bit = detect_bit(addr);
5         byte |= (bit << i);
6     }
```

```
6     }  
7     return byte;  
8 }
```

It sequentially calls `detect_bit()` for each bit of the byte and shifts the result into position. This forms a complete ASCII character, enabling reconstruction of the original message sent by the sender.