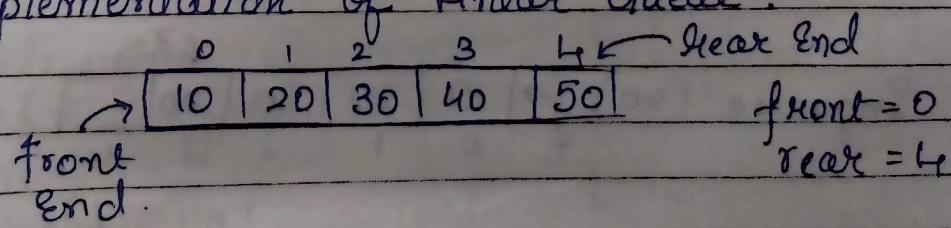


## 2. Queues

The Concept:

1. People standing outside the ticketing window of a cinema hall.  
The first person in the line will get the ticket first & thus will be the first one to move out of it.
  2. People moving on an escalator. The people who got on the escalator first will be the first one to step out of it.
- ✓ Stores data in an ordered manner.
  - ✓ It is a FIFO Data Structure.  
    → First In First Out.
  - ✓ Elements added from rear end.
  - ✓ Elements are removed from front end.

ARRAY Implementation of Linear Queue.



In Stack we use one variable top.  
where as in Queue we use  
front Variable for front end.  
rear Variable for rear end.

## Operations On Queues

1. Insertion / Enqueue .
2. Deletion / Dequeue .
3. traversing / Display .
4. peek .

Note: For Queue .

Array should be used in horizontal Manner .

1. Insertion / Enqueue .

→ Adding new element to the queue .

① Initial Scenario .

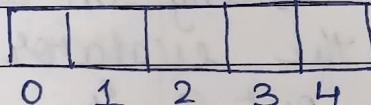
$$\text{MAX} = 5$$

$$\text{front} = \underline{\text{front}} \quad \text{rear} = \underline{\text{rear}} - 1$$

$$(f)$$

$$(r)$$

$$f \searrow$$

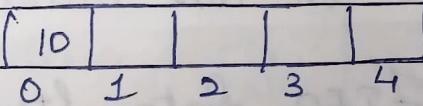


$\gamma \gamma \gamma$  ;

$$q[\gamma] = \text{val};$$

$f = r = -1 \leftarrow \text{Queue is Empty}$  .

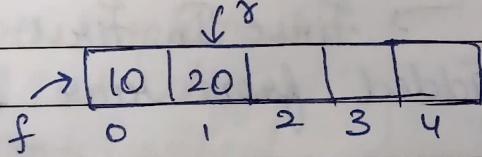
② Enqueue (10)



$\gamma \gamma \gamma$  ;

$$q[\gamma] = 10$$

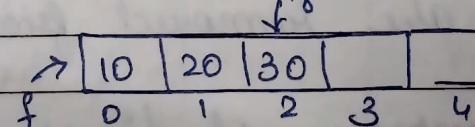
③ Enqueue (20)



$\gamma \gamma \gamma$  ;

$$q[\gamma] = 20$$

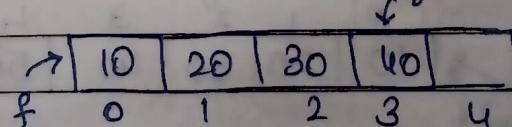
④ Enqueue (30)



$\gamma \gamma \gamma$  ;

$$q[\gamma] = 30$$

⑤ Enqueue (40)

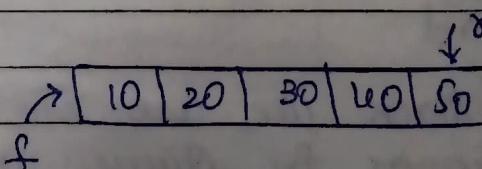


$\gamma \gamma \gamma$  ;

$$q[\gamma] = 40$$

⑥

Enqueue (50)



$\gamma \gamma \gamma$  ;

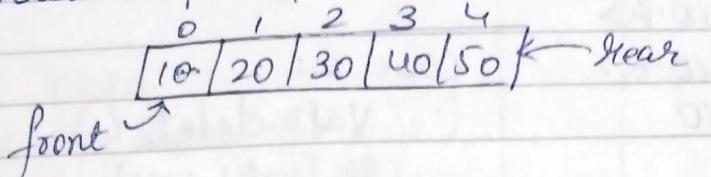
$$q[\gamma] = 50$$

⑦

Enqueue (60) → Overflow Condition

∴ rear == MAX-1 is Overflow Condition .

### 3) Deletion / Dequeue

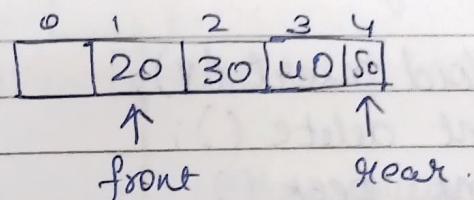


① Queue Empty Condition 1. if ( $\text{front} == -1$ )

② delete one element

$$\text{val} = q[\text{front}] \\ = 10$$

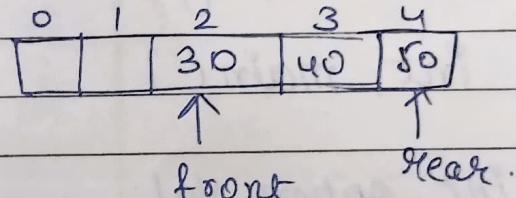
front++  
return val.



③ delete one element

$$\text{val} = q[\text{front}] = 20$$

front++  
return val

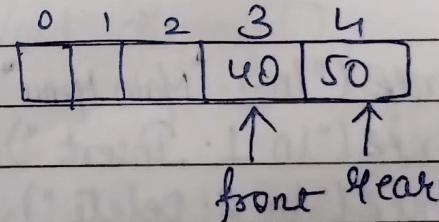


④

delete one element

$$\text{val} = q[\text{front}] = 30$$

front++;  
return val;



⑤

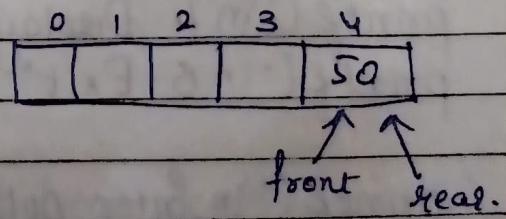
delete one element

$$\text{val} = q[\text{front}] = 40$$

front++;

return val; condition for

if ( $\text{front} == \text{rear}$ ) ← 1 ele.



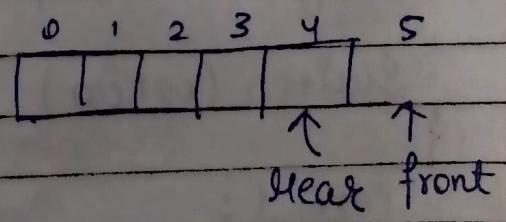
⑥

delete one element

$$\text{val} = q[\text{front}] = 50$$

front++;

return val;



2nd Condition →  
for queue to be empty.

# Codecraft Academy

9359164822

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Q. WAP to implement a Linear Queue.

```
#include <stdio.h>
#include <conio.h>
#define MAX 10
```

```
int queue[MAX];
int front = rear = -1;
```

```
Void insert();
```

```
int delete();
```

```
int peek();
```

```
Void display();
```

```
int main()
```

```
{
```

```
int option, Val;
```

```
do
```

```
{
```

```
printf("In... Main Menu ...");
printf("In 1. Insert");
```

```
printf("In 2. Delete");
```

```
printf("In 3. Peek");
```

```
printf("In 4. Display");
```

```
printf("In 5. Exit");
```

```
printf("In Enter Option:");
```

```
scanf("%d", &option);
```

```
Switch(option)
```

```
{
```

case 1:

```
Insert();
```

```
break;
```

Case 2:

```
Val = delete();
```

```
if (Val != -1)
```

```
pf("In del. ele = %d", Val);
break;
```

Case 3:

```
Val = peek();
```

```
if (Val != -1)
```

```
pf("In 1st ele in queue %d",
Val);
```

```
break;
```

Case 4:

```
display();
```

```
break;
```

```
} // switch ends
```

```
} while (option != 5);
```

```
return 0;
```

```
}
```

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## Void Insert()

```
{  
    int num;  
    printf("In Enter the no:");  
    scanf("%d", &num);  
  
    if (rear == (MAX - 1))  
        printf("In Overflow");  
    else if (front == -1 && rear == -1)  
        {front = rear = 0;}  
    else  
        {rear++;}  
    queue[rear] = num;  
}
```

```
int peek()  
{  
    if (front == -1 ||  
        front > rear)  
    {  
        printf("In Q. is Empty");  
        return -1;  
    }  
    else  
    {  
        return queue[front];  
    }  
}
```

## Int delete()

```
{  
    int val;  
    if (front == -1 || front > rear)  
    {  
        printf("In Underflow");  
        return -1;  
    }  
    else  
    {  
        val = queue[front];  
        front++;  
        if (front > rear)  
            front = rear = -1;  
        return val;  
    }  
}
```

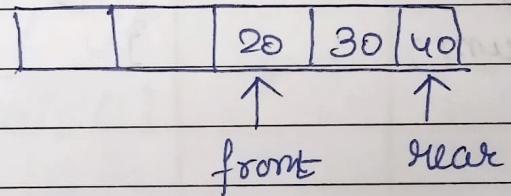
```
Void display()  
{  
    int i;  
    printf("In");  
    if (front == -1 || front > rear)  
        printf("In Q. is Empty");  
    else  
    {  
        for (i = front; i <= rear; i++)  
            printf("It %d", queue[i]);  
    }  
}
```

## Application of Linear Queue

- ✓ Linear queues are used in applications where data elements need to be processed in the order in which they are received.

✓

## Limitation of Linear Queue



Though the queue is empty still we cannot insert an element in the queue.

Here it will show Overflow condition  
 $\text{rear} = \text{MAX}-1$  as true.

Hence Circular Queue.

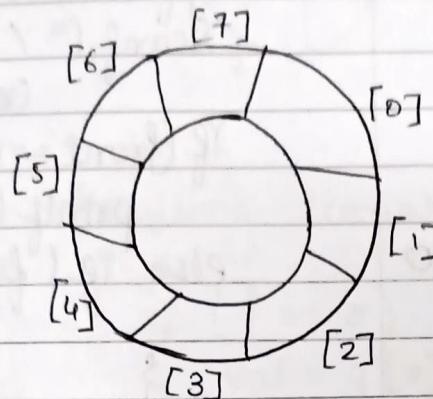
\* The queue data structure can be classified into the following types:

1. Circular Queue
2. Deque
3. Priority Queue
4. Multiple Queue

## \* Circular Queues.

- The first index comes right after the last index.

1) Initially



$$MAX = 8$$

$$front = -1$$

$$rear = -1$$

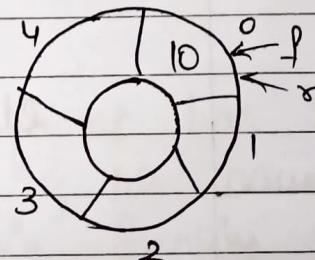
$$q[MAX]$$

2) Insertion

2.1) Insert 10 [1st ele]

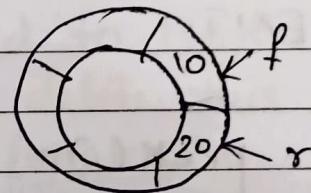
if ( $r == -1 \& \& f == -1$ )  
 $f = r = 0$ .

$$q[r] = Val;$$



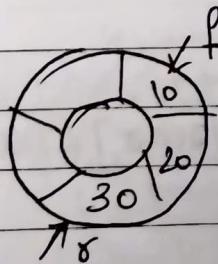
$$r++$$

$$q[r] = 20$$



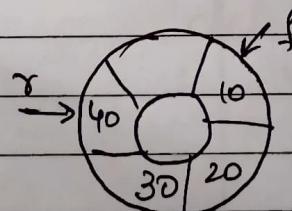
2.3) Insert 30.

$$r++;$$
  
$$q[r] = 30$$



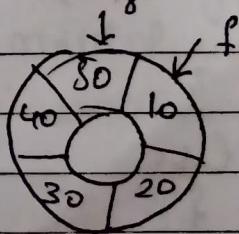
2.4) Insert 40

$$r++;$$
  
$$q[r] = 40;$$



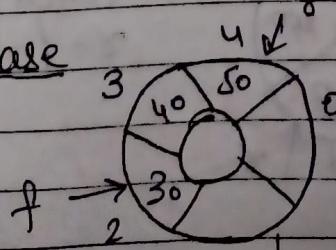
2.5) Insert 50.

$$r++;$$
  
$$q[r] = 50$$



∴ Full condition / overflow condition :-  $rear = MAX - 1$   
and  $front = 0$

Special Case



Here

$$r = MAX - 1$$

but  $f \neq 0$ .

in that case just do  $r = 0$

$$\therefore r = 0$$

$$q[r] = Val.$$

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## Algorithm

1. If  $F == 0 \& R == MAX - 1$

    Write "Overflow"

    Goto Step 4

2. Step If  $F == -1 \& R == -1$

$F = R = 0$

ELSE IF  $R == MAX - 1 \& F != 0$

$R = 0$

ELSE

$R = R + 1$

3.  $Q[R] = VAL$

4. EXIT.

int peek()

if ( $front == -1 \& rear == -1$ )

    Pf("In Queue is Empty");

    return -1;

else

{

    return queue[front];

## Functions

void insert()

{ int num

    Pf("Enter ele to Insert");

    Scanf("%d", &num);

    (rear == front - 1) ;;

    if (front == 0 & rear == MAX - 1)

        printf("In Overflow");

    else if (front == -1 & rear == -1)

        front = rear = 0;

        queue[rear] = num;

}

    else if (rear == MAX - 1 & front != 0)

        rear = 0;

        queue[rear] = num;

}

    else

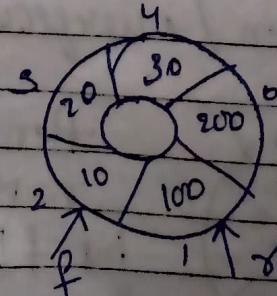
{

        rear++;

        queue[rear] = num;

}

}

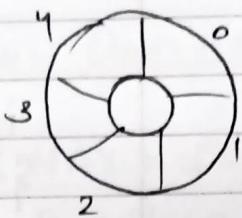


2nd overflow condition .

$$r = f - 1$$

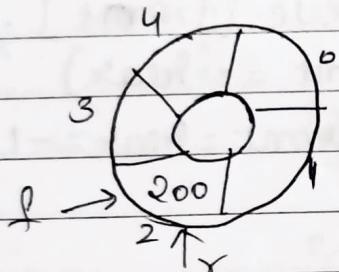
## 2] Deletion Operation On Circular Queue

1) Initial MAX=5  
 $f = -1;$   
 $r = -1;$   
 $q[MAX]$



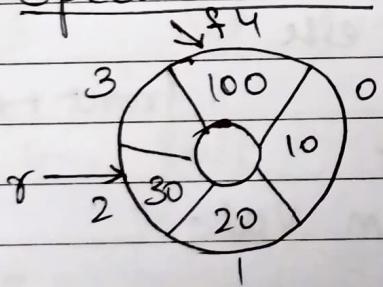
∴ Underflow condition / Queue Empty  
 $f = r = -1$ .

2) Consider the following scenario :- Only 1 ele in Queue

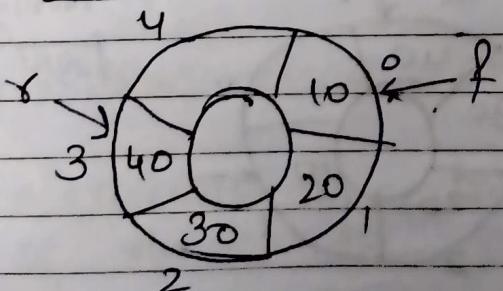


if ( $f == r$ )  
{ val = q[r];  
 f = r = -1.  
} If there is only 1 ele in Queue, store the value in a variable  
Set  $f = r = -1$

3) Special case :- front has reached to MAX-1 & we have to delete this ele.  
∴ Simply Set front = 0.



4) Normal deletion :-



delete element 10  
∴ val = q[f];  
 $f++$

## ALGORITHM

1. If FRONT = -1  
Write "Underflow"  
Goto step 4
2. val = Queue[FRONT]
3. If FRONT = REAR  
FRONT = REAR = -1  
Else  
If FRONT = MAX-1  
SET FRONT = 0  
Else  
SET FRONT++;
4. EXIT

## Void display()

```

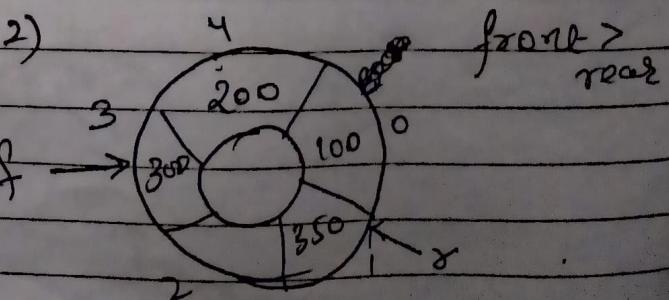
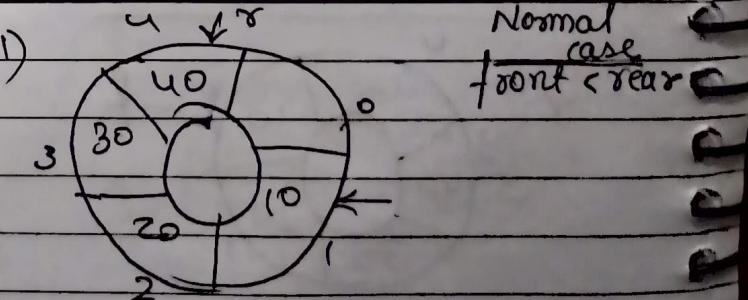
{ int i;
printf("\n");
if (front == -1 && rear == -1)
    printf("In Queue is Empty");
else
{ if (front < rear)
    for (i = front; i <= rear; i++)
        printf("It %d", queue[i]);
    else
    for (i = front; i <= MAX; i++)
        printf("It %d", queue[i]);
    for (i = 0; i <= rear; i++)
        printf("It %d", queue[i]);
}
}

```

```

int delete()
{
    int val;
    if (front == -1 && rear == -1)
        printf("In Underflow");
    return -1;
}
Val = queue[front];
if (front == rear)
    front = rear = -1;
else
{
    if (front == MAX-1)
        front = 0;
    else
        front++;
}
return Val;
}

```



## Circular Queue Updated

### 1) Insertion / Enqueue Operation

1. if  $(\text{rear} + 1) \% \text{MAX} = \text{FRONT}$   
    write "Overflow"

2. if  $\text{front} = -1 \& \text{Rear} = -1$  // 1st ele.  
    Set  $\text{front} = \text{rear} = 0$   
else

    if  $\text{rear} = \text{MAX} - 1 \& \text{front} \neq 0$   
         $\text{rear} = 0;$

    else

$\text{rear} = (\text{rear} + 1) \% \text{MAX}$

3. Set queue [ $\text{rear}$ ] = val

4. Exit.

### DEQUE

✓ pronounced as 'dequeue'

✓ elements are inserted or deleted from either  
of the end.

✓ No element can be added or deleted from  
the middle.

✓ In Computer's memory, a deque is implemented  
using either a circular array or a circular  
doubly linked list.

✓ They are classified into two types.

1. Input Restricted Deque.
2. Output Restricted Deque.

✓ INPUT Restricted Deque.

→ Operations.

- 1) Insertion of an element at the Rear end.
- 2) Deletion of an ele. from both ends

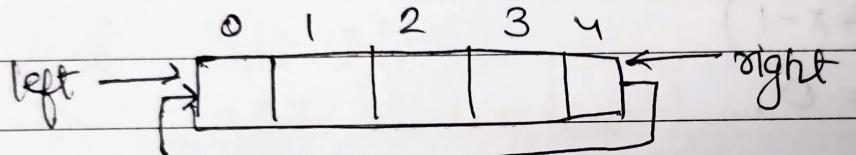
✓ Output Restricted Deque.

→ Operations.

- 1) Insertion of an ~~to~~ element from both ends
- 2) Deletion of an element from front end.

\* Operations Associated with Deque.

- a) empty() :- i.e. Underflow / Empty queue condn.
- b) full() :- i.e. Overflow condition.
- c) enqueueR() :- Add item to the Rear end.
- d) enqueueF() :- Add item to the Front end..
- e) dequeueR() :- Delete item from the Rear End.
- f) dequeueF() :- Delete item from the front End.

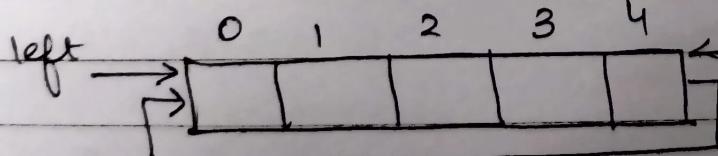
Input Restricted DequeOperations

1) insert\_right

2) delete\_left

3) delete\_right

Imagine the queue as above.

Output Restricted DequeOperations:

1) insert\_right

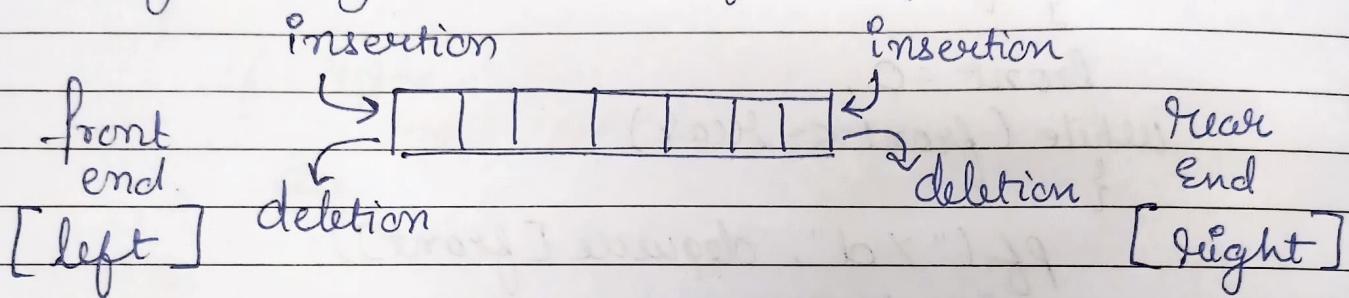
2) insert\_left

3) delete\_left

Image the queue as above

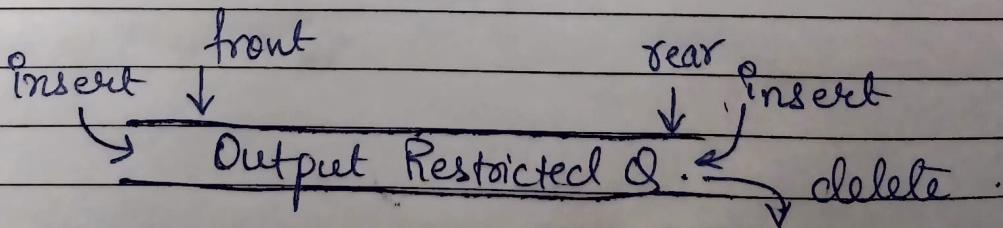
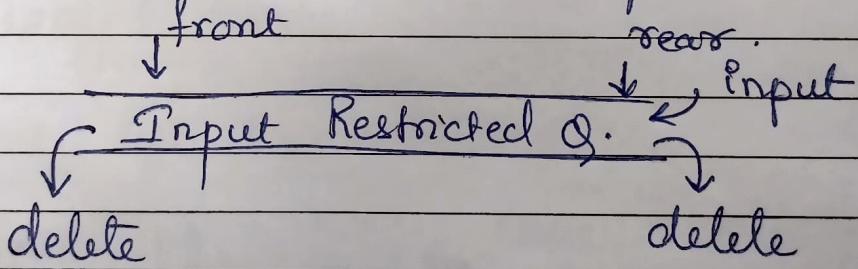
## Double Ended Queue [deque]

- ✓ Linear data structure
- ✓ Insertion & deletion operations are performed from both ends.
- ✓ Generalized version of the queue.



### Types of deque

- 1) Input Restricted Queue    2) Output Restricted Queue



### Operations to be performed on deque

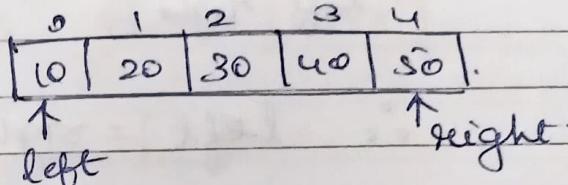
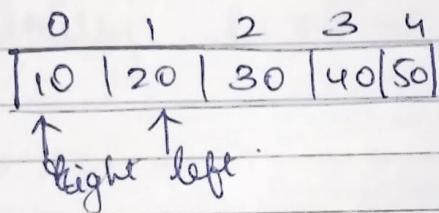
1. Insertion at the front
2. Insertion at the rear. (Same as linear queue)
3. Deletion at front (" " " ")
4. Deletion at rear { " " " " }

\* Insertion at the front end.

Case 1: Full Queue

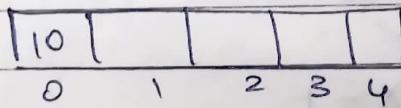
$$\text{left} = \text{right} + 1$$

$$\text{left} = 0 \text{ & right} = \text{MAX} - 1.$$



Case 2: Queue is Empty:

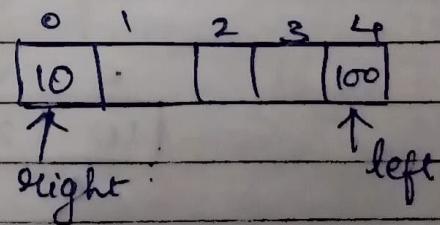
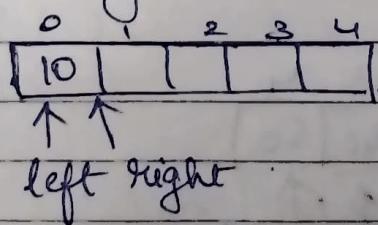
$$\text{left} = \text{right} = -1 \quad \boxed{\quad \quad \quad \quad \quad} \quad \leftarrow \text{Empty Queue}$$



$$\therefore \text{left} = \text{right} = 0$$

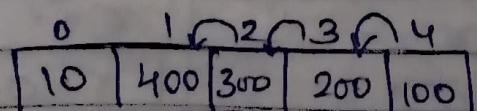
$$q[\text{left}] = \text{val} \text{ (say } 10\text{)}$$

Case 3: Only 1 element



$$\therefore \text{left} = \text{MAX} - 1 \therefore \text{left} = 4.$$

Case 4: Normal Insertion



$$\text{left} --$$

$$q[\text{left}] = 200$$

$$\text{left} --$$

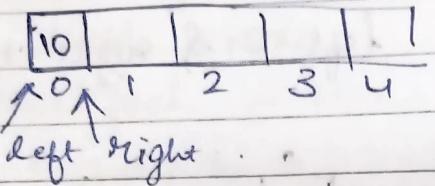
$$q[\text{left}] = 300$$

$$\text{left} --$$

$$q[\text{left}] = 400$$

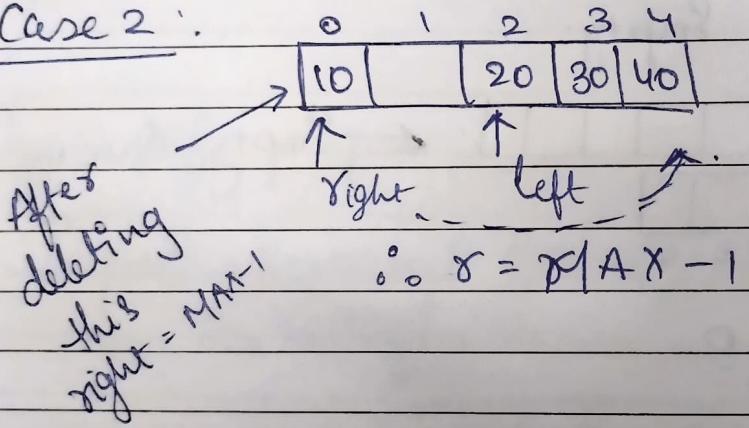
## Deletion from Right End

Case 1: only 1 element



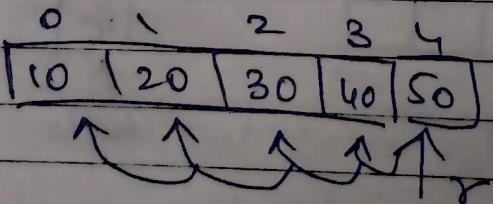
$$\therefore \text{left} = \text{right} = -1;$$

Case 2:



$$\therefore \gamma = \gamma / \Delta x - 1$$

Case 3: Normal deletion from rear end.  
right



$$\gamma = 4.$$

delete(50)	delete(40)	delete(30)	delete(20)	delete(10)
$\gamma --$				
$\boxed{\gamma = 3}$	$\boxed{\gamma = 2}$	$\boxed{\gamma = 1}$	$\boxed{\gamma = 0}$	$\boxed{\gamma = -1}$

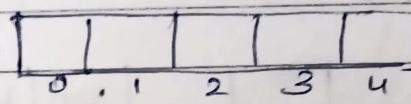
## Input Restricted Queue

Operations which can be performed.

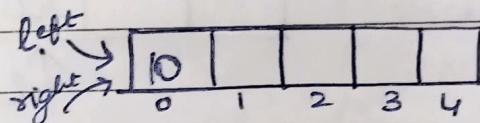
- 1) insertR()
- 2) deleteL()
- 3) deleteR()

Operations are called in the following manner.

Initially  $l = r = -1$  Max = 5 deque [MAX].



1) insertR(10)



insertR(20)

insertR(30)

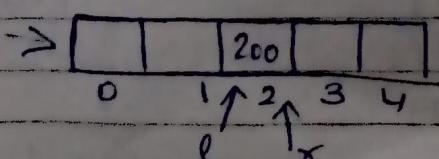
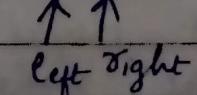
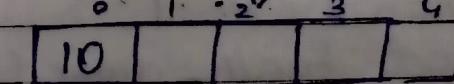
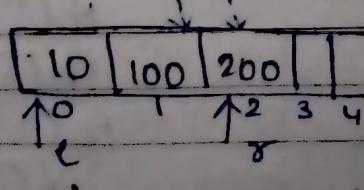
insertR(40)      insertR(50)

deleteR()

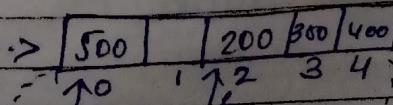
[10 20 30 40 50]

left

right

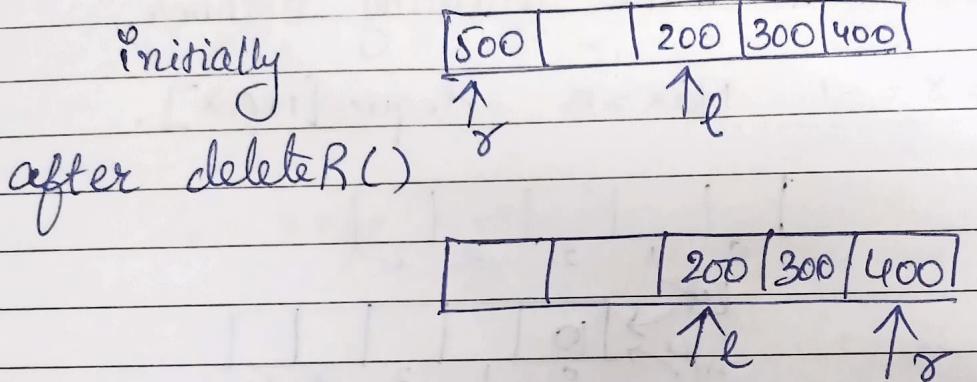


Now Perform 3 times InsertR --> [500 200 300 400]



Now, if we perform deleteR()

& which is pointing to 0th index  
will be pointing to MAX-1 next.



Similarly Students can perform operations on  
Output Restricted Queue.

WAP to implement Input & Output Restricted deque.

#include <stdio.h>

#include <conio.h>

#define MAX 10

int deque[MAX];

int l = r = -1;

Void Input\_deque();

Void Output\_deque();

Void Insert\_right();

Void Insert\_left();

Void delete\_right();

Void delete\_left();

Void display();

```
int main()
{
    int option;
    clrscr();
}
```

Pf("... MAIN MENU ...");

Pf(" -- 1. Input Restricted");

Pf(" -- 2. Output Restricted");

Pf(" Enter option");

Sf(" %d ", &option);

switch(option)

{

case 1 : Input\_deque();  
break;

case 2 : Output\_deque();

3 "main" break;

3 "main"

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Void Input-deque()

{  
int option;  
do

{ pf("... Menu ...");  
(" 1. Insert Right");  
(" 2. Delete Left");  
(" 3. Delete Right");  
(" 4. Display");  
(" 5. Quit");  
(" Enter option");  
sf("%d", &option);  
switch (option)

{ Case 1:

insert-right();  
break;

Case 2:

delete-left();  
break;

Case 3:

delete-right();  
break;

Case 4:

display();  
break;

g

} while (option != 5);

Void output-deque()

{  
int option;  
do

{ pf("... MENU ...");  
(" 1. Insert Right");  
(" 2. Insert Left");  
(" 3. Delete Left");  
(" 4. Display");  
(" 5. Quit");  
(" Enter option");  
sf("%d", &option);  
switch (option)

{

Case 1:

insert-right();  
break;

Case 2:

insert-left();  
break;

Case 3:

delete-left();  
break;

Case 4:

display();  
break;

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} while (option != 5);

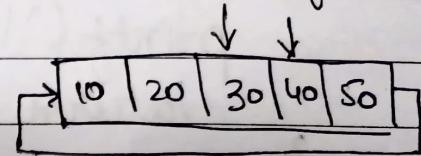
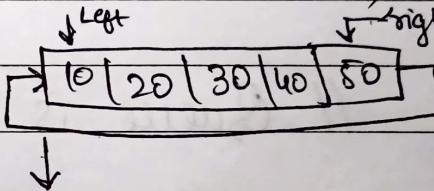
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Void insert - right()

right (rear)  
left (front)

{  
int val;

Pf("Enter value to be added");  
Sf("y-d", &val);



if (left == 0 && right == MAX-1) || (left == right + 1))

{

Pf("In overflow");

return;

}

if (left == -1) // Queue is initially empty

{  
left = right = 0;

}

else

{

    if (right == MAX-1)

        right = 0;

    else

        right ++;

}

    dequeue [right] = val;

4.

Void insert-left()

{

    int val;

    printf("In Enter the Value to be added");

    scanf("%d", &val);

    if (left == 0 & right == MAX-1) || (left == right + 1))

{

        printf("In overflow");

        return;

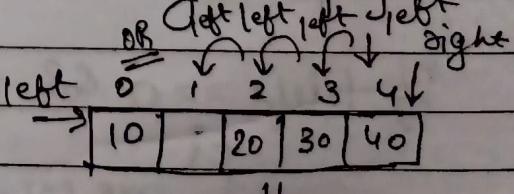
}

    if (left == -1) // Queue is initially empty

{

        left = right = 0;

}



else

{

    if (left == 0) // only 1 element

        left = MAX-1;

    else

        left = left - 1;

3

    dequeue [left] = val;

Void delete - left()

{  
if (left == -1)  
{ pf("Underflow");  
return;  
}

printf("The deleted 'ele is : %d", deque[left]);

if (left == right) // Queue has only one element

left = -1;

right = -1;

}

else

{

if (left == MAX - 1)

left = 0;

else

left = left + 1;

}

}

Void delete - right()

{

if (left == -1)

Pf("In Underflow");  
return;

}

printf("In The ele deleted is: %d", deque[right]);

{

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if (left == right) || deque has only one ele

    left = -1;

    right = -1;

else

{

    if (right == 0)

        right = MAX - 1;

    else

        right = right - 1;

}

}

## Display

{ void display ()

    int front = left, rear = right;

    if (front == -1)

{

        printf("In Queue is Empty!");

        return;

}

        printf("In The element of the queue are:");

    if (front <= rear)

        while (front <= rear)

            printf("%d", deque[front]);

            front++;

}

}

else

{

while (front <= MAX - 1)

{

printf("x.d", deque[front]);

front++;

}

front = 0;

while (front <= rear)

{

pf("x.d", dequeue[front]);

front++;

}

}

printf("\n");

}

## Priority Queue

- ✓ Arranges elements based on their priority values.
- ✓ Elements with higher priority values are typically removed or retrieved before elements with lower priority values.
- ✓ Each element has a priority value associated with it. When we add an item, it is inserted in a position based on its priority value.
- ✓ General rules of processing.
  - An element with higher priority is processed before an element with a lower priority.
  - Two elements with the same priority are processed on a First-Come-First-Served (FCFS) basis.
- ✓ Types of Priority Queue:- Ascending & Descending.
- ✓ Application of Priority Queue
  - CPU Scheduling
  - Data compression in Huffman code.
  - Event driven simulation such as Customer Waiting queue.
  - Finding kth largest / smallest ele.
  - All queue applications where priority is involved.
  - Graph algo. like Prim's Minimum Spanning tree, Dijkstra's shortest path algo.

Advantages: 1) Helps access the elements faster way.  
2) Dynamic in Nature.  
3) Efficient.  
4) Included in real-time systems.

Disadvantages:-  
1) High complexity  
2) High consumption of memory.

### Questions

1. Define Linear Queue ? Give its applications.
2. WAP to perform basic operations on Linear Queue using Arrays. OR. Array Implementation of Linear Queue.
3. Give Limitation of Linear Queue.
4. WAP to calculate no. of items in a queue.
5. What is a priority queue ? Give its applications
6. Explain the concept of Circular Queue ?. How is it better than a linear Queue ?
7. WAP to Implement Circular Queue using Arrays .
8. Give applications of Circular Queue .
9. WAP to Implement Input & Output Restricted deque.