# Introduction to the Theory of Computation Solutions $$\operatorname{\textbf{Ryan}}$$ Dougherty

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#### 4.2

Consider the problem of determining whether a DFA and a regular expression are equivalent. Express this problem as a language and show that it is decidable.

**Solution:** We formulate the problem  $EQ_{DFA,REX} = \{ \langle A, R \rangle \mid A \text{ is a DFA, } R \text{ is a regular expression, and } L(A) = L(R) \}$ . We will design a TM T that decides  $EQ_{DFA,REX}$ :

T = "On input  $\langle A, R \rangle$  where A is a DFA, R is a regular expression:

- 1. Use Theorem 1.54 to convert R into an equivalent DFA B. Therefore, L(B) = L(R).
- 2. Run  $EQ_{DFA}$  on input  $\langle A, B \rangle$ . Output what  $EQ_{DFA}$  outputs."

Since  $EQ_{DFA}$  is decidable, and the conversion from regular expressions to DFAs takes finite time,  $EQ_{DFA,REX}$  is decidable.

#### 4.3

Let  $ALL_{DFA} = \{\langle A \rangle \mid A \text{ is a DFA and } L(A) = \Sigma^* \}$ . Show that  $ALL_{DFA}$  is decidable. **Solution:** We will design a TM T that decides  $ALL_{DFA}$ :  $T = \text{"On input } \langle A \rangle$  where A is a DFA:

- 1. Construct a DFA B such that  $L(A) = \overline{L(B)}$ .
- 2. Run  $E_{DFA}$  on input  $\langle B \rangle$ . Output what  $E_{DFA}$  outputs."

Since  $E_{DFA}$  is decidable,  $ALL_{DFA}$  is decidable.