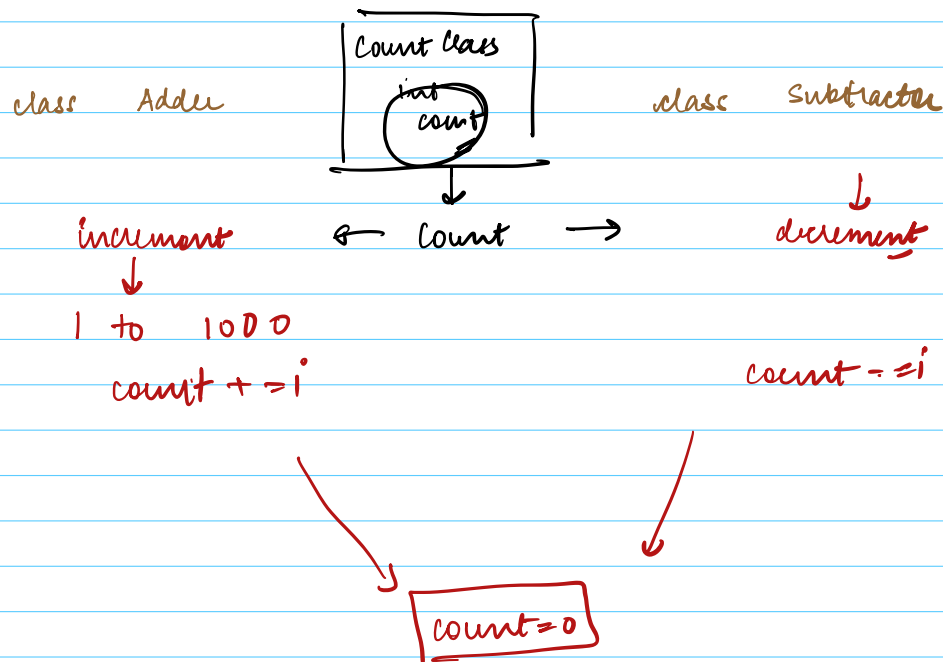


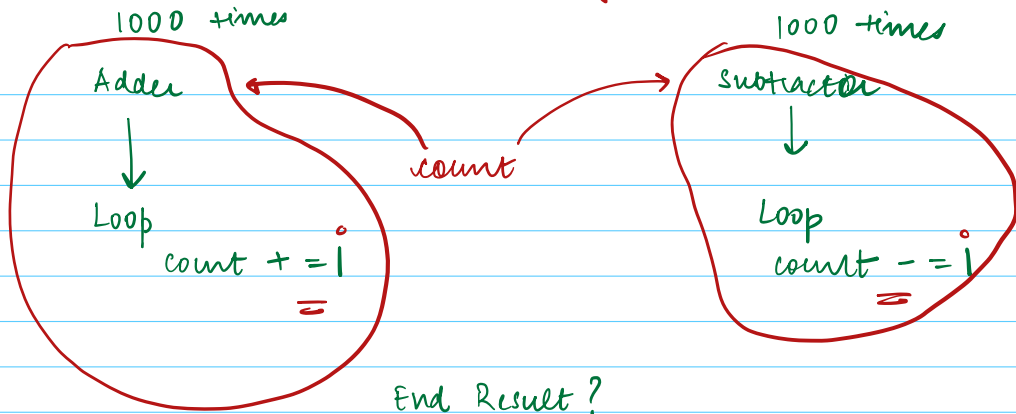
AGENDA

- 1) Problems with Multi-threading programs
 - (i) Adder Subtractor problem
 - (ii) Problem of Synchronization?
 - (iii) conditions leading to problem

- 2) Solutions
 - (i) Theoretical solution
 - (ii) Practical solution
 - a) Mutex
 - b) Synchronized
 - c) Atomic Variables
 - d) Semaphores.

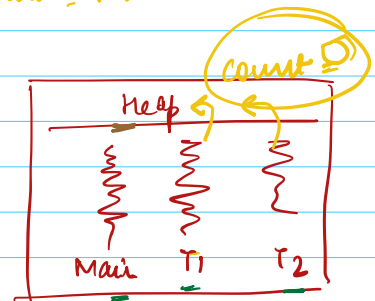
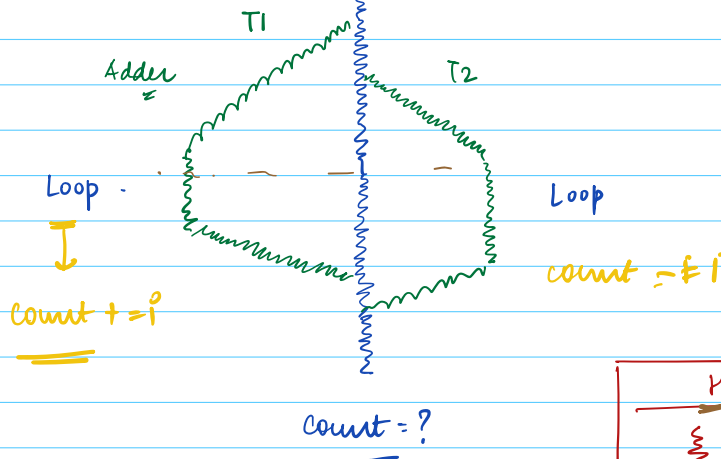


Data sharing



count $c = 0$

main \rightarrow Adder instance
Sub instance

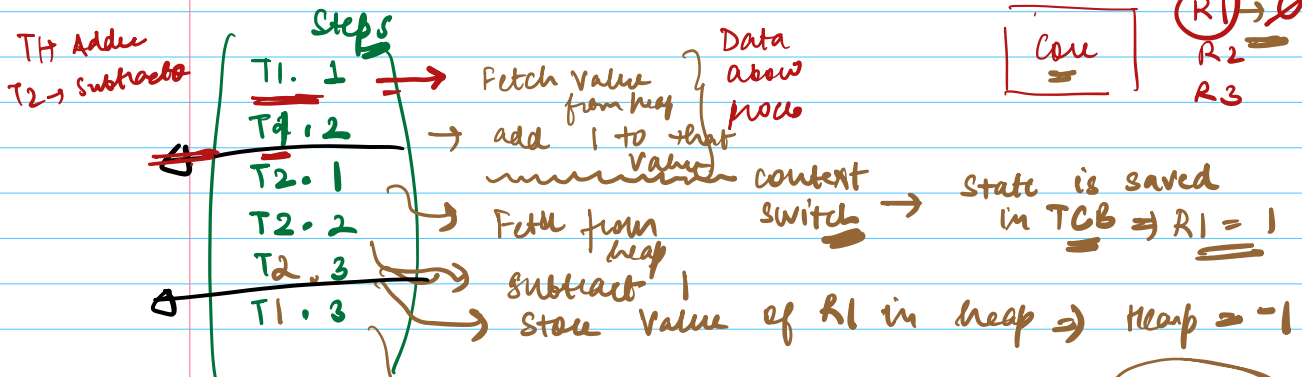


Practical

Memory

PCB

R1 \rightarrow 0
R2
R3



→ store value of R1 in heap ⇒ (heap = 1)

①

Expect = 9

Explanation:

Adder

Subtractor

① From Heap. Load value of count to Register X

① From Heap. Load value of count to Register

② $X += i$ ⇒ Register

② $X -= i$

③ Store X back to PCB

③ Store X back to PCB

PCB count 0

T1. ①

T2. ①

PCB

PCB

$X = 0$

$X = 0$

$X += 1$

$X -= 1$

State will be stored

X

store to heap

Heap

①

store state of PCB to heap

it will override value of heap to -1

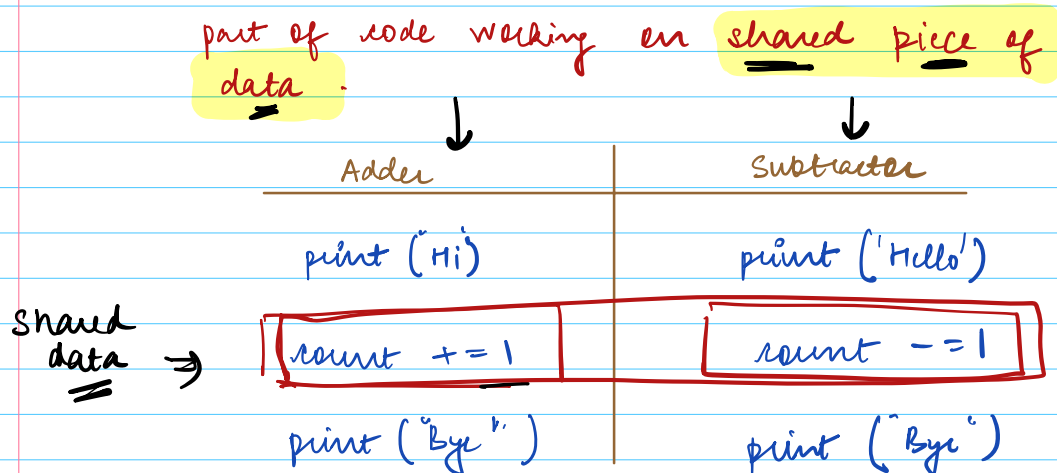
$X = -1$

Problem of Synchronization

If multiple threads work simultaneously on shared data, we might get inconsistent results because of time sharing.

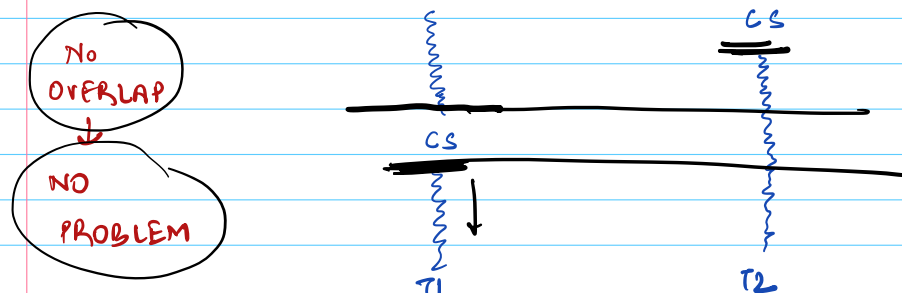
When will synchronisation problems happen?
At the conditions!

① Critical Section



② Race condition

\rightarrow More than one thread trying to enter critical section at same time.



critical section is never paused between

③

Pre-emption:

preemptive scheduling algorithm
↓
OS can context switch before the thread completed its execution.

Threads should be pre-empted while executing critical section

"context-switched"

SOLUTIONS

Theoretically what would you want to happen?

① Mutual Exclusion:

Only 1 thread should execute critical section.

s2

②

Progress -

At least 1 thread should be making progress.
Entire app shouldn't stop.

③

Bounded waiting:

No thread should wait indefinitely to enter its critical area

(starvation)

④

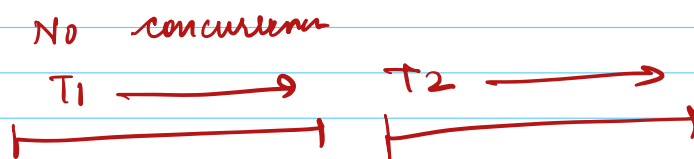
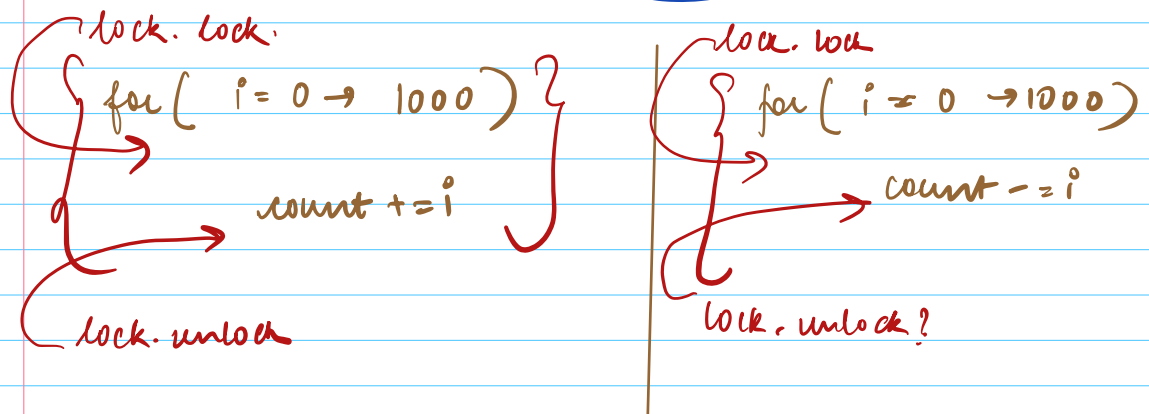
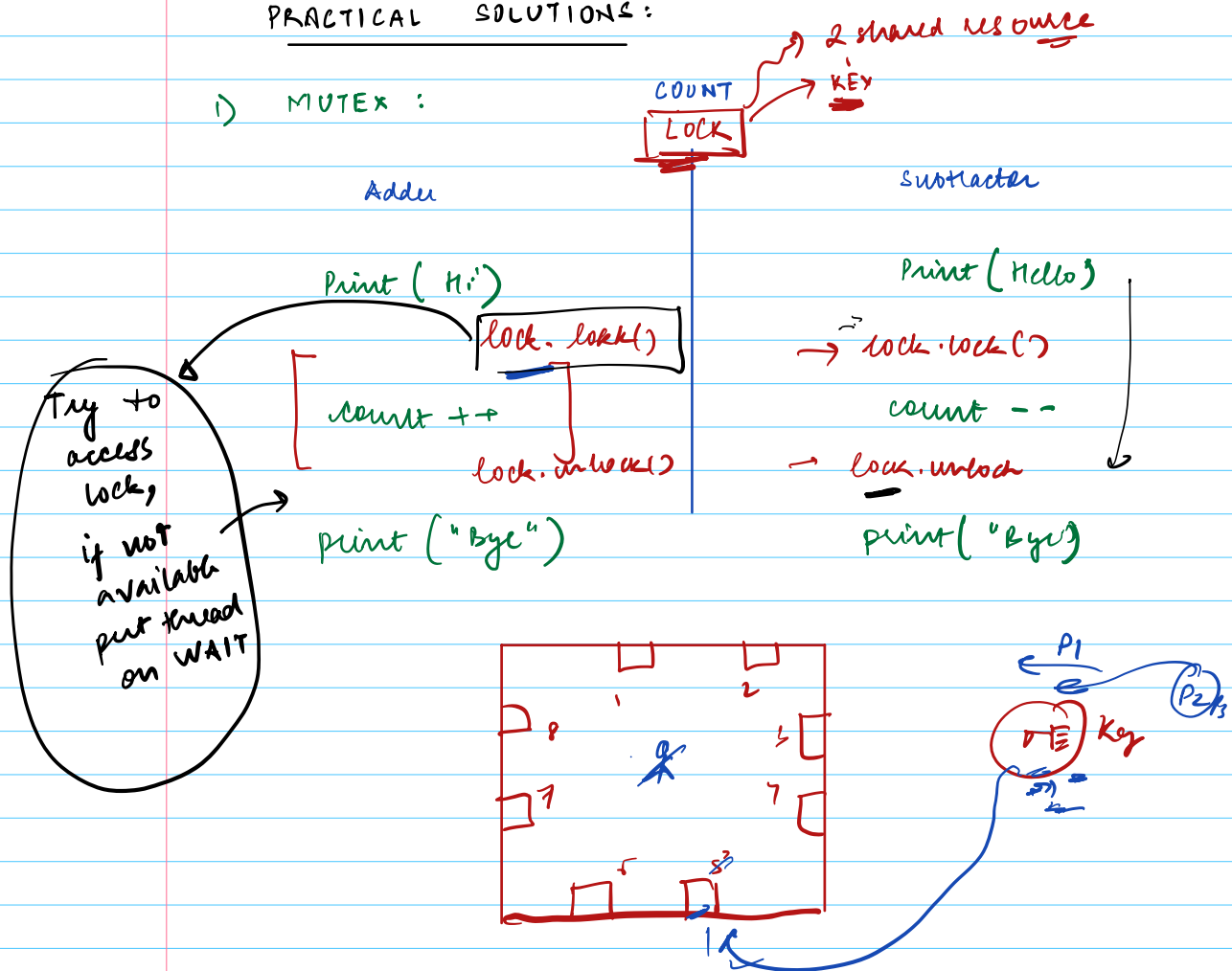
No busy waiting

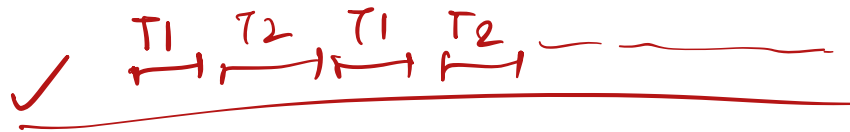
some kind of notification system.

5 min → Mutex, synchronizer

PRACTICAL SOLUTIONS:

1) MUTEX :





2) Synchronized Keyword (Java)

locks are available on all objects

Computer

Adder

```

Print ("Hi")
↓
synchronized (count) {
    [count ++]
}
print ("bye")
  
```

Subtractor

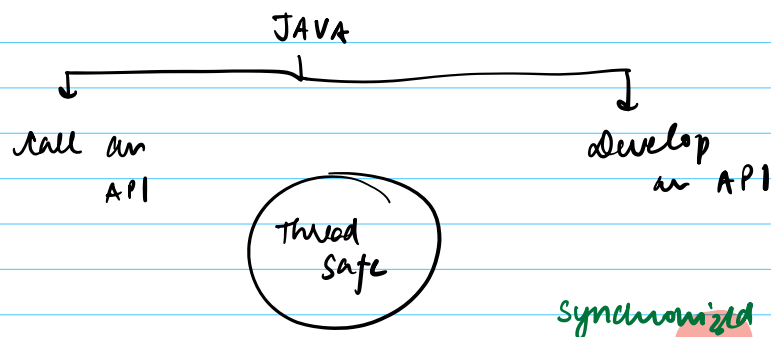
```

Print ("Hello")
↓
synchronized (count) {
    [count --]
}
print ("bye")
  
```

Mutex is an OS concept

Synchronized → Java specific

In java, locks are available on all objects



Counter
 ↓
not thread safe

```

ThreadSafe
  ↑
class Counter {

```

```

    private int count = 0

```

```

    public getCount () {

```

```

        return count
    }

```

```

}

```

instance

```

    synchronized public void increment () {

```

```

        count += 1
    }

```

← only 1 thread will be able to access this method at a particular time

```

    synchronized public void decrement () {

```

```

        count -= 1
    }

```

```

}

```


- ① Hashtable v/s HashMap.
- ② StringBuilder v/s StringBuffer.

<p>Counter</p> <p>static count = 0</p> <p>static increment count() {</p> <p style="padding-left: 40px;">count + 1</p> <p>}</p>	<p>Adder()</p> <p>increment count</p> <p>↓</p> <p>class level</p>
<p>sign decrement()</p> <p style="padding-left: 40px;">count - 1</p>	<p>Subtractor</p> <p>↓</p> <p>instance level</p>