17 Explain the following with example.

[Relationship types.

There are 3 types of relationships that can exist b/w

two entities. * One to one relationship: Buch a relationship exists when each record of one table in related to only one of

the other table. eg: If there are two entities person (Id. name, age, address)

p 'parrfort' (parrport-id, parrport-no) so each person can have only one payfort & each payfort belongs to only one person

* One to many or Many to one: Such a relationship exists when each record one table can be related to one or more than one record of the other table.

e.g. of there are two entity type 'curtomer' of 'account' Then

each 'eustomer' can have more man one 'account' but each account is held by only one 'enstoner'.

* May to may relationship! Such a relationship with when each record of the first table can be related to one or more than one record of the second table of a single record of the second table can be related to one or more than one record of the first table. eg: If there are two onlity type 'customer' of 'Product' then each customer coin buy more Thou one product & a product can be bought by many diff. emboner.

B Role, names of Rewriter Relationship: Each outily type that farticipates in a relationship type plays a particular role in the relationship.

The role name signifies the role that a participating entity from the entity type plays in each relationship means instance. It helps to explain what the relationship means e.g. in the works-FOR relationship type, Employee plays the role of employee or worker of DEPAREMENT plays the role of department or employer. Is some cases the same ety) entity type participate more than one in a relationship type in diff. roles in such cases the sole name becomes essential for distinguishing The meaning of each participation. Buch relationship types one called recursive relationship. advantural constraint: The constraints me determined from the miniworld situation that the relationship eg: Suffort each emfloyer must work for only one 27 Explain the characteristicis of relations along with relational model notations. Also briefly describe relational model constraints. characteristis of Relations: Dodoving of the relationship in highe.

A relation is set of typics. There is no farticular order for elements in a set. Hence typics in a relationship do not have for any farticular order.

I Ordering of values within a tuple: Means ordering of values associated with wondraids in the relation R. The order of attributes of their values is not that infortant as long as the correspondance b/w attribute of values maintain alternative b/m attributes of values maintain alternatives.

₹ Values & Hulls in the typle: Each value in the type his an automin value ie not divisible into comfonents. Therefore composit of amultagle value attributes not allowed because of this froperty model in sometimes called flat relational model.

Relational model constraints;

I following are the notations used for relational model. R-Relational schema R (A1, A2, A3--. An) of degree N.

If The upper cause letters Q. R. s denotes relation name.

I The lower care letters 9, 8, 8 represents relation states

of t. u. v represents ty tuples.

Relationed model constraints:

Combrant con be divaded into 3 mais categories.

O model based: Combails that are inherent in the data model are called as model based constraints.

eg: 40 duplicate tuples combraints.

3 Schema barcel: Construct that com be directly expressed in the schema of data model. which are typically specified by using DDL. eg: SSN- Primary ky.

(3) Application based: Lowbraints that cannot be expressed of Perspected through applied programs.

frograms.

e.g. Salay: employee of manager.

3] Explain nexted of wo-related Hested queries with example. In nested quevies a query in written inside a query. The Terult of inner query is used in execution of outer query. There merted queries are complete blocks within I another set. The outer part quevies called outer query 4 inner part quevies called inner query. The nested query can offer along with where dans or From dan or SELECT class. Wormally in nested over the day of the second of t query are the comparison operator 'IN' or 'EQUAL TO' eg. To find the names of employee who are dept. 20 103. SELECT E. ename.

SELECT E. ename.

FROM Employee E

WHERE E. id EN (SELECTED. id from dept. D. where D. ID = 103). Co-veloted norted quevous are nested quevous which ove independent or depend only on the same row of our outer query being an embedded query.

I am outer query being an employee who are dept. id 103 e.g! To find the name of employee who are dept. SELECT E. Name WHERE EXISTS (SELECT * FROM dept D. where D. id = 103

AND Drid = E.id)

Explain with example Exists of HOT EXISTS functions The EXISTS for in son in used to check whether the result of a w-related guery in empty or not. The runt of EXISTS in a Bodian value TREE of The nexted query result contacints arleast one high or FALSE of the nexted query result combains no typles. eg Retrieve the name of each employee who has a depudent with the same first name of its the same sex as the employee. SELECT L. F. Name, L. LNAME FROM EMPLOYEE AS E WHERE EXISTS (select * FROM DEPENDENS AS D WHERE E.SSN= D.ESSN AND B.SEX = D.SEX AND E. FName = D. Dogadent Name); EXISTS of MOTERISTS are sypteally used in conjunction

EXISTS of NOTEXISTS are supreally wild in conjunction with a wo-related nested guery.

eg, Retrieve the names of employee who have no defeated series frame, kname

FROM EMPLOYEE

WHERE HOT EXISTS (SELECT *

FROM DEPENDENT

WHERE SON = ESSN)

what are assertions of triggers in sqt?

while sqt program to create an assertion to

specify the constraint that the salary of employee must not be greater than the salary of the depotedent manager. The employer for in the Company database When a countraints involves 2 (or) more tables The table courtraint mechanism in sometimes hard of results many not some as rejected. To wer such intuation in sat supports the creation of arrentions which are combraints not associated with only one table. And an assertion statement should ensure a certain condition with always out in the databare. DBMs always checks the assertion whenever modifications are done in the corresponding table.

CREATE ASSERTIONS. SALARY LONSFRAINTS

enter (NOT Exists (selut * Grom employer Et, employer on, dept. D. where E salary > on. salary AND MGR.SSN- E.SSN));

A tryger in a database object that is associated with the table, it will be activated with the table it will be activated when a defined action is executed for he table. The original can be executed when eve run the following statements
ENSERT, DEDATE DELETE I it can be invoked before or after the event.

B) Drile trigger in sqt to call a stored procedure.

SNFORM-ESUPERVISOR() whenever a new record is created or updated their whether the employee's salary in greater than the salary of his or her direct supervisor in the company database. CREATE TRIGGER Salay - Violation BEFORE INSERS DR UPDATE of salary, supervisor-SSN ON Employee For each row when (NEW. salary > (SELECT salary FROM employee WHERE SSN= NEW, Supervisor-SSN)) INFORM-SUPERVISOR (NEW_SUPENVBOY - SEN, NEW_SCN)); Final the following wir to stored procedure. * (reating a simple stored procedure To Calling a stored procedure. Stored procedures are created to porform one ore more Done operations on database. It is nothing but the group of SQL statements that accepts some i/p int se form of parameters of performs Some task of may or may not returns a value. The stored procedure must have a nam, This Movel procedure has the name ShowHumber Of Orders' otherwise it just contains on SQL statement shal-

in precompiled & stored at The server.

CREATE PROCEDORE Show Humber Of Orders SELECT c. cid, (. cnami, cours (x) FROM Cartomers C. orders o WHERE C.O.d = O.C.Pd GROUP BY Cied & (name. Calling a stored procedure can be called in interactive SQL with the CARE statement. 8) Describe the architecture of JOBL. Java Application JDBC API JOBL Diver manager JOBC DONEY JOBC DINER JOBC Outa source

SQL Server

erive cranager! This dass manages a list of database drivers. Matches connection requests from the java applia with the proper database driver using communication gub protocol. The 1st driver that recognizes a certain subprotocol under JOBC will be med to atablish a Driver; This Enterface handles the communication with the dutabane connection. database server, Abstracts the details with working with driver objects.

Connection: This interface with all methods for contacting a database. The connection object repress communication context is all communication with database is through connection object only. Statement: Some derived interface accept parameters in addition to executing stored procedures. Result set: There object, hold data retrieved from a databal after of execute an son guery wring statement objects. It acts as an interestor to allow to move Through its dalar SOL Exception: This class handles any errors that occur. In a databan application a) Briefly aplain on following individuals steps that we required to submit a database query to a data source of to ratrieve the results. . All drivers we managed by the Driver manager dars. JOBL Oriver Management: · Loading a JARC driver! In the jova code class. forHame ("Oracle / idbe. drner. Oracle drner"); When starting the jour applin D-jdéc. drivers = oracle /sdbc. enver.

Connections:

Interacts with a data source through session. Each connection identifies a logical session.

TOBC URL:

jdbc; < Subprotocol >: < other parameters >

Exocating SAL Statements: Three lift ways of executing sal statements:

- · Statement
- · Prepared Statement.
- · Callable Statement.

Statements are both static & dynamic 3ax statements. Prepared statements are semi-static 80% statements. callable statements are stored procedures.

101 Describe Three-Tier architecture:

Printation Tier

Application Tier Database Fier

Dotaban Tier: At Min Her, the dalaban revides along with its growy processing longuages. Also have the relations that define the data of their constraints at This level. Application Fier: At this tier revides the apptor server & The programs that accen the database for a user this applin tier represents an abstracted view of the database End was one unware of any accidance of the database beyond The applin. At the other end the database tier in not aware of any other mer beyond the applin tier. Hence the applin layer tites sits in the middle of

datesbare.
User Fier: End wers operate on this teer of they know nothing about any existence of the database begins this layer. At this layer multiple views of the database for database can be provided by the application.
All views are generated by applies that yeride in the application that