

# Assignment - I

3

7 Construct CFG for the following languages.

① Balanced parenthesis.

$$L = \{w \in \{(), \{\}^*\}$$

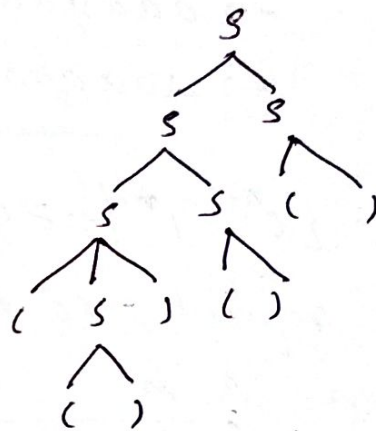
The required CFG in  $\{S\}, \{(), \{\}, R, S\}$

where R is

$$S \rightarrow SS / (S)()$$

$$\text{let } w = (( )) ( ) ( )$$

$$\begin{aligned} S &\rightarrow SS \\ &\rightarrow SSS \\ &\rightarrow (S)SS \\ &\rightarrow (())SS \\ &\rightarrow (())(()) \\ &\quad \underline{\underline{\quad \quad \quad}} \end{aligned}$$



②  $L = \{w \in \{a, b\}^* \mid w \text{ contains substring } ab\}$

$$L = \{ab, aab, abb, bab, \dots\}$$

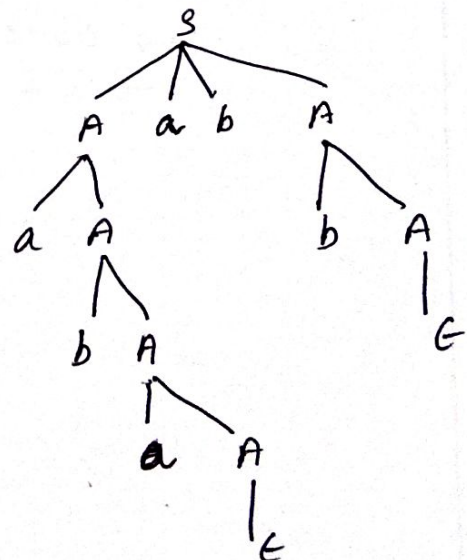
The required CFG in  $\{A, S\}, \{a, b\}, R, S\}$

where R is

$$\begin{aligned} S &\rightarrow AabA \\ A &\rightarrow aA / bA / \epsilon \end{aligned}$$

$$\text{let } w = abadbb$$

$$\begin{aligned} S &\rightarrow AabA \\ &\rightarrow aAabA \\ &\rightarrow abAabA \\ &\rightarrow abaAabA \\ &\rightarrow abaabbbA \\ &\rightarrow \underline{\underline{abaabb}} \end{aligned}$$



iii)  $L = \{a^i b^j \mid 2i = 3j + 1\}$

The required CFG in  $\{\{S\}, \{a, b\}, R, S\}$

where  $R$  is

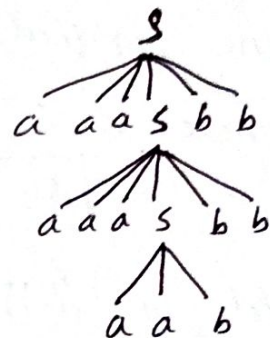
$$S \rightarrow aaasbb / aab$$

Let  $w = aaaaaaabb bbb$

$$S \rightarrow aaasbb$$

$$\rightarrow aaaaaa s bbb$$

$$\rightarrow aaaaaa aab bbb$$



iv)  $L = \{0^{n+2} 1^n \mid n \geq 1\}$

The required CFG in  $\{\{S\}, \{0, 1\}, R, S\}$

where  $R$  is

$$S \rightarrow 0S1 / 0001$$

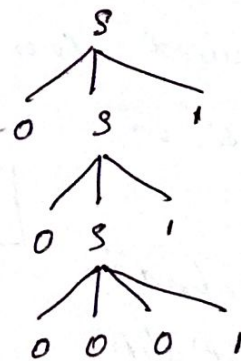
$$L = \{0001, 000011, 00000111, \dots\}$$

Let  $w = 00000111$

$$S \rightarrow 0S1$$

$$\rightarrow 00S11$$

$$\rightarrow 00000111$$





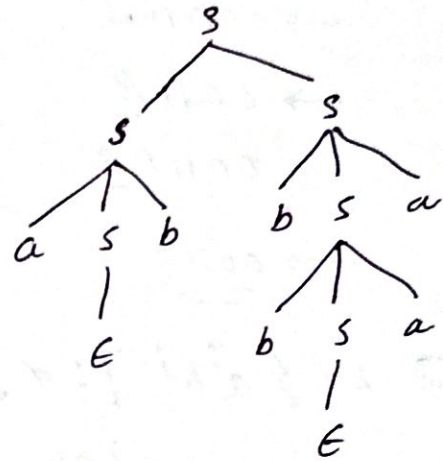
$$\textcircled{v} L = \{ w \in \{a, b\}^* \mid n_a(w) = n_b(w) \}$$

$$L = \{ abab, baba, abbbbaa, \dots \}$$

$$S \rightarrow SS \mid aSb \mid bSa \mid \epsilon$$

eg: abbbbaa

$$\begin{aligned} S &\rightarrow \underline{SS} \\ &\rightarrow a \underline{S} b S \\ &\rightarrow a \underline{b} S \\ &\rightarrow a b b \underline{S} a \\ &\rightarrow a b b b \underline{S} a a \\ &\rightarrow \underline{a b b b a a} \end{aligned}$$



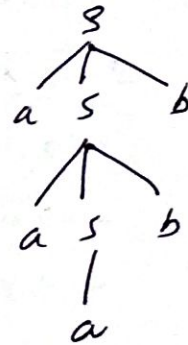
$$\textcircled{vi} L = \{ a^i b^j \mid i = j+1 \}$$

$$L = \{ a, aab, aaab, \dots \}$$

$$S \rightarrow a \mid aSb$$

eg: aaabbb

$$\begin{aligned} S &\rightarrow a \underline{S} b \\ &\rightarrow a a \underline{S} b b \\ &\rightarrow \underline{a a a} b b b \end{aligned}$$

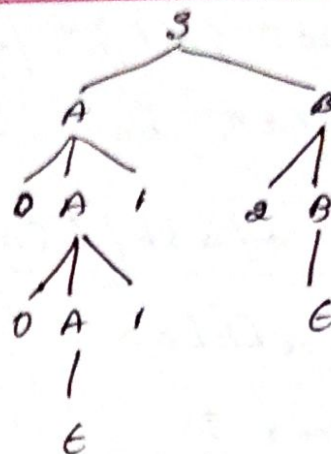


$$\textcircled{vii} L = \{ 0^m 1^n 2^n \mid m \geq 0, n \geq 0 \}$$

$$\begin{aligned} S &\rightarrow AB \\ A &\rightarrow 0A \mid \epsilon \\ B &\rightarrow 2B \mid \epsilon \end{aligned}$$

eg: 00112

$S \rightarrow \underline{A}B$   
 $\rightarrow 0\underline{A}1B$   
 $\rightarrow 00\underline{A}11B$   
 $\rightarrow 0011\underline{B}$   
 $\rightarrow 00112\underline{B}$   
 $\rightarrow \underline{00112}$

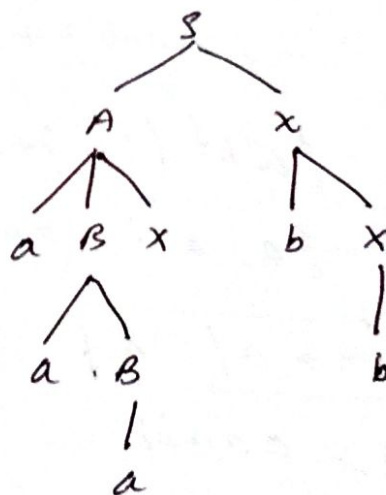


⑦  $L = \{ a^i b^j \mid i \neq j, i \geq 0, j \geq 0 \}$

$S \rightarrow BA / Ax$   
 $A \rightarrow aAb / \epsilon$   
 $B \rightarrow aB / a$   
 $x \rightarrow bx / b$

e.g: aaabbb

$S \rightarrow \underline{A}x$   
 $\rightarrow a\underline{B}x$   
 $\rightarrow aa\underline{B}x$   
 $\rightarrow aaa\underline{x}$   
 $\rightarrow aaab\underline{x}$   
 $\rightarrow \underline{aaabbb}$



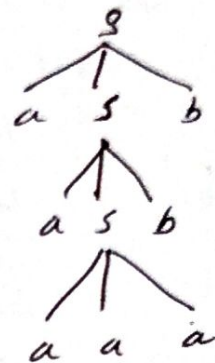
(1x)  $L = \{a^n b^{n-3} \mid n \geq 3\}$

$L = \{a^3, a^4b, a^5b^2, a^6b^3, \dots\}$

$S \rightarrow aSb / aaa$

eg:  $w = a^5b^2$

$$\begin{aligned} S &\rightarrow aSb \\ &\rightarrow aaSbb \\ &\rightarrow aaaaaabb \\ &= \end{aligned}$$



(x)  $(011+1)^* (01)^*$

$L = \{\epsilon, 010101, 01111, 101, 01101, \dots\}$

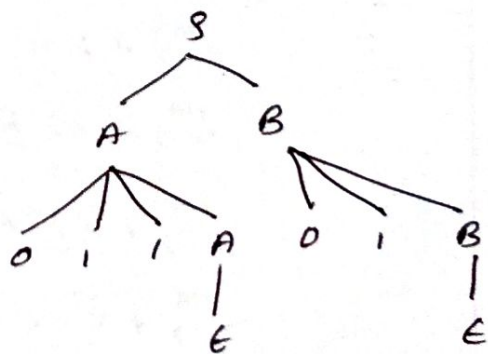
$S \rightarrow AB$

$A \rightarrow 011A / 1A / \epsilon$

$B \rightarrow 01B / \epsilon$

eg:  $01101$

$$\begin{aligned} S &\rightarrow \underline{AB} \\ &\rightarrow 011\underline{AB} \\ &\rightarrow 011\underline{B} \\ &\rightarrow 01101\underline{B} \\ &\rightarrow 01101 \\ &= \end{aligned}$$





2) Consider the grammar.

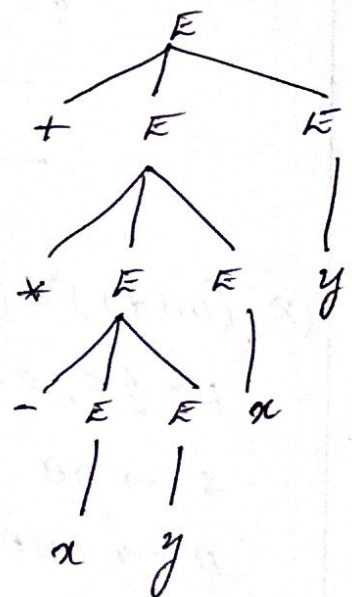
$$E \rightarrow +EE \mid *EE \mid -EE \mid x \mid y$$

Find the LMD, RMD & parse tree for the string  
 $+* - xyxy$

LMD :

$$\begin{aligned}
 S &\xRightarrow{lm} + \underline{E} E \\
 &\xRightarrow{lm} + * \underline{E} E E \\
 &\xRightarrow{lm} + * - \underline{E} E E E \\
 &\xRightarrow{lm} + * - x \underline{E} E E \\
 &\xRightarrow{lm} + * - xy \underline{E} E \\
 &\xRightarrow{lm} + * - xy x \underline{E} \\
 &\xRightarrow{lm} + * - xy xy \\
 &=
 \end{aligned}$$

Parse Tree :



RMD :

$$\begin{aligned}
 S &\xRightarrow{rm} + E \underline{E} \\
 &\xRightarrow{rm} + E y \\
 &\xRightarrow{rm} + * E \underline{E} y \\
 &\xRightarrow{rm} + * E x y \\
 &\xRightarrow{rm} + * - E \underline{E} x y \\
 &\xRightarrow{rm} + * - E y x y \\
 &\xRightarrow{rm} + * - xy xy \\
 &=
 \end{aligned}$$

3) Consider the following grammar with production

$$S \rightarrow AbB$$

$$A \rightarrow aA/\epsilon$$

$$B \rightarrow aB/bB/\epsilon$$

Find the LMD, RMD & parse tree for the string aaabab

LMD :

$$S \xRightarrow{lm} \underline{AbB}$$

$$\xRightarrow{lm} a\underline{AbB}$$

$$\xRightarrow{lm} aa\underline{AbB}$$

$$\xRightarrow{lm} aaa\underline{AbB}$$

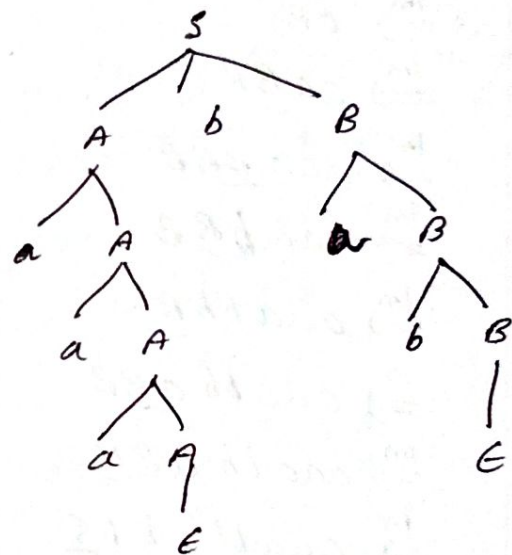
$$\xRightarrow{lm} aaab\underline{B}$$

$$\xRightarrow{lm} aaaba\underline{B}$$

$$\xRightarrow{lm} aaabab\underline{B}$$

$$\Rightarrow \underline{aaabab}$$

Parse Tree :



RMD :

$$S \xRightarrow{rm} Ab\underline{B}$$

$$\xRightarrow{rm} Aba\underline{B}$$

$$\xRightarrow{rm} Abab\underline{B}$$

$$\xRightarrow{rm} \underline{Abab}$$

$$\xRightarrow{rm} a\underline{Abab}$$

$$\xRightarrow{rm} aa\underline{Abab}$$

$$\xRightarrow{rm} aaa\underline{Abab}$$

$$\xRightarrow{rm} \underline{aaabab}$$



41 Let  $G$  be the following grammar.

$$S \rightarrow aB \mid bA$$

$$A \rightarrow a \mid aS \mid bAA$$

$$B \rightarrow b \mid bS \mid aBB \quad ; \text{ string: } aaabbaabba$$

Find LMD, RMD & parse Tree.

LMD:

$$S \xRightarrow{lm} aB$$

$$\xRightarrow{lm} aaB$$

$$\xRightarrow{lm} aaaBB$$

$$\xRightarrow{lm} aaaBB$$

$$\xRightarrow{lm} aaabbB$$

$$\xRightarrow{lm} aaabbAB$$

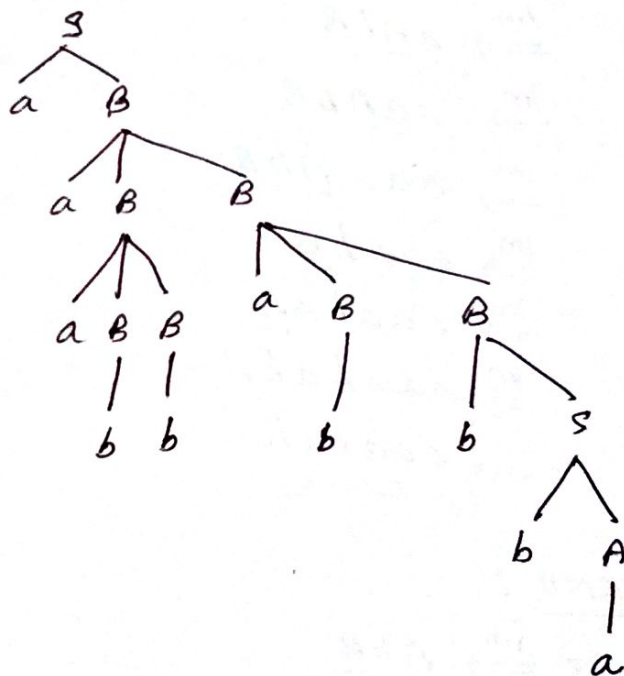
$$\xRightarrow{lm} aaabbabB$$

$$\xRightarrow{lm} aaabbabbs$$

$$\xRightarrow{lm} aaabbabbbA$$

$$\xRightarrow{lm} \underline{aaabbabbbba}$$

Parse Tree:





$$\xRightarrow{rm} aaa \underline{B} a b b b a$$

$$\xRightarrow{rm} aaa \underline{B} b a b b b a$$

$$\xRightarrow{rm} aa a b b a b b b a$$

5) Show that the following grammar is ambiguous.

$$S \rightarrow aB/bA$$

$$A \rightarrow aS/bAA/a$$

$$B \rightarrow bS/aBB/b$$

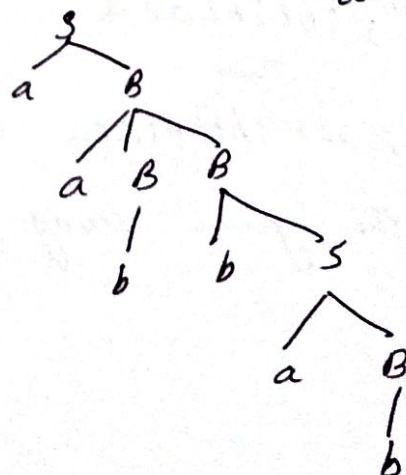
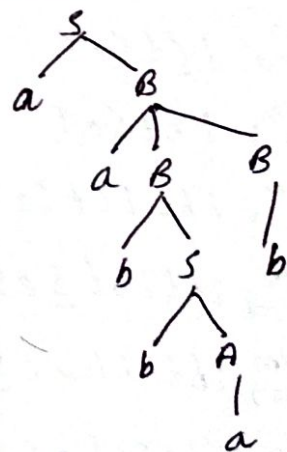
A grammar  $G = (V, \Sigma, R, S)$  is said to be ambiguous if for any string  $w \in L(G)$  there exists more than one left most derivation or more than one right most derivation or more than one distinct parse tree.

$$w = aabbab$$

$$\begin{aligned} \text{LMD's: } S &\xRightarrow{lm} a\underline{B} \\ &\xRightarrow{lm} aa\underline{BB} \\ &\xRightarrow{lm} aab\underline{S}B \\ &\xRightarrow{lm} aabb\underline{AB} \\ &\xRightarrow{lm} aabb\underline{aB} \\ &\xRightarrow{lm} aabb\underline{ab} \end{aligned}$$

$$\begin{aligned} S &\xRightarrow{lm} a\underline{B} \\ &\xRightarrow{lm} aa\underline{BB} \\ &\xRightarrow{lm} aab\underline{B} \\ &\xRightarrow{lm} aabb\underline{S} \\ &\xRightarrow{lm} aabb\underline{aB} \\ &\xRightarrow{lm} aabb\underline{ab} \end{aligned}$$

parse tree:



6) Is the following grammar ambiguous.

$$s \rightarrow ict\bar{s}/ilt\bar{s}es/a$$

$$e \rightarrow b$$

$$\phi \quad w = ibtibtaea$$

$\Rightarrow$  LMD's:

$$s \xRightarrow{lm} i\bar{c}t\bar{s}$$

$$\xRightarrow{lm} i\bar{b}t\bar{s}$$

$$\xRightarrow{lm} i\bar{b}ti\bar{c}t\bar{s}es$$

$$\xRightarrow{lm} i\bar{b}ti\bar{b}t\bar{s}es$$

$$\xRightarrow{lm} i\bar{b}ti\bar{b}tae\bar{s}$$

$$\xRightarrow{lm} i\bar{b}ti\bar{b}taea$$

$$s \xRightarrow{lm} i\bar{c}t\bar{s}es$$

$$\xRightarrow{lm} i\bar{b}t\bar{s}es$$

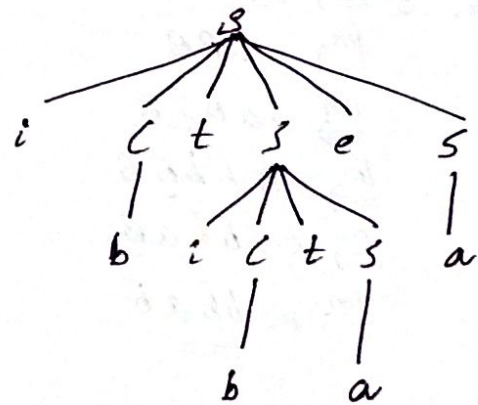
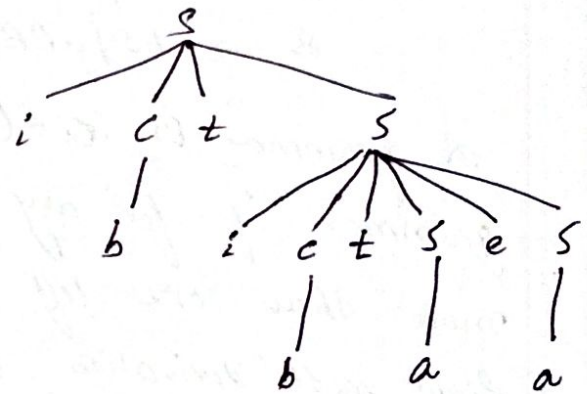
$$\xRightarrow{lm} i\bar{b}ti\bar{c}t\bar{s}es$$

$$\xRightarrow{lm} i\bar{b}ti\bar{b}t\bar{s}es$$

$$\xRightarrow{lm} i\bar{b}ti\bar{b}tae\bar{s}$$

$$\xRightarrow{lm} i\bar{b}ti\bar{b}taea$$

Parse Tree's



Since,  $w = ibtibtaea$  has 2 LMD's & 2 parse tree's in the given string, hence it is ambiguous.