Last dass:
$$\{x_n y_{n\eta_0} iid$$

$$S_{m} = x_1 + \cdots + x_n$$

Central limit Theorem

$$\frac{S_{n}-n\mu}{\sqrt{n}} \stackrel{d}{\longrightarrow} N(0,\sigma^{2})$$
, as $n\to\infty$

Weak Law Of Carge Numbers

$$\frac{S_{n}}{m} \stackrel{p.}{\longrightarrow} \mu , \text{ as } n \rightarrow \infty$$

$$Sample mean \qquad True mean$$

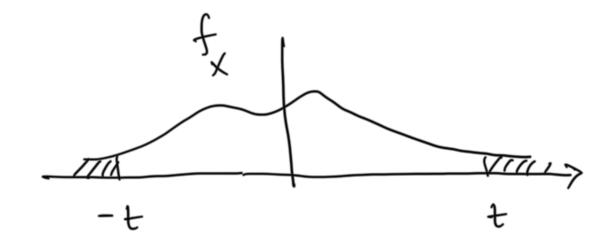
$$P(|S_{n} - \mu| \gg E) \leq \sqrt{\frac{2\sigma^{2}}{\pi n} S^{2}} \stackrel{e}{\longleftarrow} -D$$

approximation

$$P(|\frac{s_m}{m} - \mu| > \epsilon) \leq \frac{\sigma^2}{m\epsilon^2}$$

$$\bigcirc$$

Tail Bounds: IP (|x 17 t)



R. H. S of (1) is better than R. H. S of (2) (swaver for large n)

* asymptotic result (as n > 0)

* no decision to make; we are

simply looking at a sequence Sn

* no algoritanic interaction.

 \bigvee

at this stage we really cannot appreciate why R.H.S (2) being better than R.H.S (1) is great.

aval. Connect Regret in multi-armed bandits
to tail bounds.

Muet armed bandit problem

· arms $A = \{1, \dots, k3\}$

- · Distributions are P,,..., Pk
- · Play an arm At and receive rewat

- · Means are H3 = Ep. [X]
- Best reward $\mu_{\star} = \max_{i} \mu_{i}$
- · Best arm i* = argmax fix

Regret =
$$R_m = E \left[\sum_{t=1}^{m} (\mu_k - \chi_t) \right]$$

Lemma: Let
$$\Delta_a = \mu_{\chi} - \mu_a$$
 be the

sub-optimality gap
$$m = \sum_{n=1}^{\infty} T_{n}$$
 random Variable

Indicator function $T_{\text{gcond } y} = 1$, cond is the

= 0 , cond is passe

number of times arm a was chosen in n rounds.

$$R_m = \sum_{a \in A} \# [T_a(n)] \Delta_a$$

Proof:

$$R_n = \mathbb{E} \left[\sum_{t=1}^{m} C_{t} + X_{t} \right]$$

Since at each time t, only one arm is chosen one $\mathbb{F}_{\{A_s=1\}}$ is

$$T_{\xi_{A_{5}=23}} + T_{\xi_{A_{5}=23}} + \cdots + T_{\xi_{A_{5}=k3}} = 1$$

$$R_n = \mathbb{E} \left[\sum_{t=1}^n (k_t - x_t) \cdot 1 \right]$$

$$= \mathbb{E} \left[\sum_{t=1}^{\infty} (\mu_{*} - x_{t}) (\mathbb{I}_{A_{t}=13}^{+} + \mathbb{I}_{A_{t}=23}^{+} + \cdots + \mathbb{I}_{A_{t}=k7}^{+}) \right]$$

$$= \mathbb{E} \left[\sum_{\alpha \in A} \sum_{t=1}^{n} (h_{x} - X_{t}) \mathbb{I}_{\xi_{A_{t}} = \alpha 3} \right]$$

using Tower property of expectation

[[X 1 Y] or #[Y 1 X]

Tower property: condition on At

#[# [\(\frac{\frac{\gamma}{\gamma}}{\alpha\epsilon \frac{\gamma}{\tau}} \]

#[# [\(\frac{\frac{\gamma}{\gamma}}{\alpha\epsilon \frac{\gamma}{\tau}} \]

#[# [\(\frac{\frac{\gamma}{\gamma}}{\gamma} \frac{\gamma}{\gamma} \]

#[\(\frac{\gamma}{\gamma} \frac{\gamma}{\gamma} \frac{\gamma}{\gamma} \]

#[\(\frac{\gamma}{\gamma} \frac{\gamma}{\gamma} \frac{\gamma}{\gamma} \frac{\gamma}{\gamma} \frac{\gamma}{\gamma} \frac{\gamma}{\gamma} \]

Moving & [- 1 At] inside summation

$$= \mathbb{E} \left[\sum_{a \in A} \mathbb{E} \left[C_{H_{x}} - x_{t} \right] \mathbb{I}_{A_{t}=a_{3}} |_{A_{t}} \right]$$

E [g(x) f(y) ly] = f(y) E[g(x) ly]

$$= \mathbb{E} \left[\sum_{\text{at } A \in \mathbb{N}} \mathbb{E} \left[\sum_{\text{t} = 1} \mathbb{E} \left[\sum_{\text{t} = 1$$

$$= \mathbb{E} \left[\sum_{a \in A} \frac{\pi}{2} I_{A_{t} = a} C_{H_{t}} - H_{A_{t}} \right]$$

=
$$\mathbb{E} \begin{bmatrix} 2 & \tilde{\Sigma} \\ q \in A \end{bmatrix} = \begin{bmatrix} (\mu_x - \mu_a) \end{bmatrix}$$

$$= \mathbb{E} \left[\begin{array}{ccc} \mathbf{Z} & \mathbf{\widetilde{Z}} & \mathbf{\mathbb{I}} \\ \mathbf{A} \in \mathbf{A} & \mathbf{t} = 1 \end{array} \right] \times \mathbb{A}_{\mathbf{t}} = \mathbf{A}_{\mathbf{J}} \cdot \Delta_{\mathbf{A}}$$

moving # [] inside summation

Arm 1
$$\gamma_{11}(\omega) - \gamma_{12}(\omega) - \cdots$$

Moral: Number of times we play the non-optimal arm needs to be kept down.

Arm 1

Arm 2

with the source of the second of the second

* : True Mean

· : Sample Mean

Tail Bound: P(|Sample Mean - True Mean 1 > E) 2 tail mass

