1999 Comm. autom. and Pi Calculus. Questin A, Solution notes.

Rules (bookwak): REACT: $P \xrightarrow{\lambda} P' \qquad Q \xrightarrow{\overline{\lambda}} Q'$ SUM: M+ d.P+N ~P P/Q -> p'/Q' L-PARt: P/Q -> P/Q

R-PARt $\frac{Q \xrightarrow{a} Q'}{P | Q \xrightarrow{a} P | Q'}$ RESt: P -> P'

New a P -> new a P' (d & [a.a]) IDENT: \(\frac{\fir}{\frac{\frac{\frac{\frac{\frac{\frac{\frac{\frac{\frac{\frac{\frac{\frac{\frac{\frac{\frac}\frac{\frac{\frac{\frac{\frac{\frac}{\fir}{\firin}}}}{\firac{\frac{\frac{\frac{\frac{\frac{\frac{\f [5 marks] P/Q => R can be inferred either (1) REACT, from P > P' and Q > Q', with d= I and R=P/Q' or (2) by L-PAR, from P=P, with R=P/10 or (3) by R-PAR+ pm Q=Q', with R=PQ' In the first case, R, is P'/Q'. But by taking $\lambda' = \overline{\lambda}$ we have $Q \xrightarrow{\lambda'} Q'$, $P \xrightarrow{\overline{\lambda}} P'$ and hence by REACT; we infer $Q/P \xrightarrow{\overline{\lambda}} Q'/P'$, which is structurally impress to R, In the second case Ris P/Q. But R-PAR, can be used to infer jum P => P' That BIP = QIP' = R, . The thind case is similar. [5 marks] An example is P = a.P', $Q = \overline{a}.Q'$. If now a $(P|Q) \xrightarrow{\sim} R$, then we must have that $P|Q \stackrel{\sim}{\to} R$ where $R_i = \text{new } a R_i d \not = \int_{Q_i} R_i d r_i$ by RESt, and by R-PARt We get P/nowa Q = P/newa Q = R, by (2) or (3): Then the inference P/newa Q => R = R' 15
tainly easy to (mstruct (and doe not depend on the syntactic (mditing) Note that I cannot be a or R. [5 marks,]