Maths for Computation theory 2003 P1098 L4 Paper 10 Solution Section 4, "Relations on a set". Part

of the arsive-requires familiarity

with fart A section 3; "Algebra of
relations".) In order that R be a partial order on A: $(a,b) \in R$, $(b,a) \in R \implies a = b$ 1) In order that R be a total order on A: iii), and "for all a, b e A, $(a,b) \in R$ or $(b,a) \in R$ (or both!) Algebraic formulation of the conditions: LASR; V RoReR; (îi R, R-1 & LA (equivalently, = LA); (iii R J R-1 = (A x A). (11) S= R n R'. In the formalism

5 Mathe for Computation Theory Paper 10 Salution (etd.) Proof: We must show that Sis reflexive, symmetric and transitive. a) $\Delta_A \in R \implies \Delta_A = \Delta_A^{-1} \subseteq R^{-1}$ -) $\triangle_A \subseteq (R_n R^{-1}) = S$. REFLEXIVE $S^{-1} = (R_0 R^{-1})^{-1} = R^{-1}_0 R = S$ SYMMETRIC (SoS) = (RnR-1) o (RnR-1) (RoR) n (R-1 o R-1) = (ROR) n (ROR) C R n R-1 = S TRANSITIVE

Hence S'is indeed an equivalence rel'

6 Mathe for Computation Theory Paper 10 Solution (etd) HIGHLY DESIRABLE to show first that

\[
\text{V} \quad \text{vell-defined}, \quad \text{vel I don't expect}
\] them to do so. Suppose [a] < [b], « e [a], B e [b]. By definition aRb, and by definition of S: «Ra and a Ra; BRb and bRB. BUT «Ra, a Rb, bRB = XRB. Hence the definition of \leq on A/s is udependent of the chrices a e [a], b e [b]. THIS PAGE IS WHAT PRODUCES

THE EXTRA 2 MARKS.

(7 Matter for Computation Theory Hence given [a] e A/s, we may represent it by a. But a Ra. i. [a] < [a], and i) holds. Suppose [a] = [b] and [b] = [a]. By definition, aRb and GRa, hence (a,b) eS. : [a] = [b]. Hence condition (ii) also holds Supose Ca] < Cb] and Cb] < Cc], so that by definition aRb and GRC. Endition ii) holds for R, hence a Re. By definition, [a] < [c]. Hence condition ii) holds for <, and (A/s, <) is a poset.

Maths for Computation theory Paper 10 Solution (etd) is le Z, so 'of xeZ, then by taking q=1 we may show (x,x) ∈ R. is suppose (sing) eR and (y, 2) eR, Say $y=[x, Z=qy], for [,q \in Z]$. Then Z = (pq) nc, so that $(x, z) \in \mathbb{R}$. i. Rissa pre-order on Z. HOWEVER \=-1,-1, -1=1,-1,

so $(1,-1) \in \mathbb{R}$, $(-1,1) \in \mathbb{R}$. $\therefore \mathbb{R}$ nor a partial order on \mathbb{Z} .

(9 Mother for Computation Theory Paper 10 Solution (etd) Let S be the equivalence relation derived from R. Then (x,y) e S 'ff /sc/=/y/. (Z/s, <) behaves in escartly the same way as (M, R), where Pu = R n (M×M) o the restriction of R to the natural numbers. In (Z/s, <) there is a single maximal element (07 = 503, and a single ninimal element []= {1,-1}.