

The candidate is being asked to prove the Pumping Lemma (or rather, its contrapositive).

Suppose  $L$  has the given property. We assume  $L$  is regular and derive a contradiction:

If  $L$  is regular, then  $L = L(M)$ , the language accepted by some DFA  $M$ .

Consider  $l \triangleq$  number of states in  $M$ .

If  $w \in L$  with  $\text{length}(w) \geq l$ , say

$w = a_1 a_2 \dots a_n$  ( $n \geq l$ ), since  $w$  is accepted by  $M$ , we have a transition sequence in  $M$  of the form

$$s_M = q_0 \xrightarrow{a_1} q_1 \xrightarrow{a_2} q_2 \dots \xrightarrow{a_l} q_l \dots \xrightarrow{a_n} q_n \in \text{Accept}_M$$

$\underbrace{\hspace{10em}}_{l+1 \text{ states}}$

Since  $l+1$  is one more than the number of states in  $M$ ,  $q_0, \dots, q_l$  can't all be distinct states: so  $q_i = q_j$  for some  $0 \leq i < j \leq l$  and the above transition sequence looks like

$$s_M = q_0 \xrightarrow{u_1} q_i = q_j \xrightarrow{u_2} q_n \in \text{Accept}_M$$

$\underbrace{\hspace{2em}}_v \curvearrowright^*$

where  $u_1 \triangleq a_1 \dots a_i$ ,  $v = a_{i+1} \dots a_j$ ,  $u_2 = a_{j+1} \dots a_n$ . Thus  $\text{length}(u_1) = i \leq l$ ,  $\text{length}(v) = j - i \geq 1$  and  $u_1 v^n u_2 \in L(M) = L$ , for all  $n$ . Contradicting the property of  $L$ .

2 (a) Not regular. Show that  $L_1$  has the property mentioned in the first part of the question.

Given any  $l \geq 1$ , consider

$$w = a^l b a^l b \in L_1$$

$$\text{length}(w) = 2(l+1) \geq l \quad \checkmark$$

For any splitting  $w = u_1 v u_2$  with  $\text{length}(u_1 v) \leq l$ ,  $\text{length}(v) \geq 1$ , can only have

$$u_1 = a^p$$

$$v = a^q$$

$$u_2 = a^{l-p-q} b a^l b$$

} some  $p, q$  with  $p+q \leq l$  &  $q \geq 1$

So  $u_1 v^0 u_2 = a^{l-q} b a^l b \notin L_1$  because if  $a^{l-q} b a^l b = uu$  for some  $u$ , then  $u$  ends in  $b$  (cos  $uu$  does), so can only have  $u = a^{l-q} b$ ,  $u = a^l b$ ; but  $q \geq 1$ , so  $a^{l-q} b \neq a^l b$ . So  $L$  does not have the Pumping Lemma Property, so is not regular.

(5)

1 (b) Is regular. For taking  $w = \epsilon$ , we have that  $L_2 \supseteq \{\epsilon v \epsilon \mid v \in \Sigma^*\} = \Sigma^*$  and hence  $L_2 = \Sigma^*$  - which is regular: e.g.



is a FSA accepting  $\Sigma^*$ .

(3)