

Specification & Verification 1: Solution Notes to Question 2

$\vdash \{T\} \text{ SKIP } \{T\}$ SKIP-axiom
 $\vdash \{T \wedge T\} \text{ SKIP } \{T\}$ Precondition strengthening
 $\vdash \{T\} \text{ WHILE } T \text{ DO SKIP } \{T \wedge \neg T\}$ WHILE-rule

WHILE T DO SKIP never terminates so $[T] \text{ WHILE } T \text{ DO SKIP } [T \wedge \neg T]$ is false, hence by soundness of Hoare logic for total correctness it is not the case that $\vdash [T] \text{ WHILE } T \text{ DO SKIP } [T \wedge \neg T]$.

$\{X=x \wedge Y=y\} \text{ TEMP} := X; X := Y; Y := \text{TEMP} \{X=y \wedge Y=x\}$

Is translated to:

$\text{Spec}((\backslash s. s'X'=x \wedge s'Y'=y),$
 $\quad \text{Seq}(\text{Assign}('TEMP', \backslash s. s'X'),$
 $\quad \quad \text{Seq}(\text{Assign}('X', \backslash s. s'Y'), \text{Assign}('Y', \backslash s. s'TEMP'))),$
 $\quad (\backslash s. s'X'=y \wedge s'Y'=x))$

where

$\text{Spec}(p, c, q) = !s1 \ s2. p \ s1 \wedge c(s1, s2) \implies q(s2)$

$\text{Seq}(c1, c2)(s1, s2) = ?x. c1(s1, s) \wedge c2(s, s2)$

$\text{Assign}(v, e)(s1, s2) = (s2 = \backslash x. (x=v \implies e \ s \mid s \ x))$

The following is correctly annotated:

$\{T\} \text{ WHILE } T \text{ DO } \{F\} \text{ SKIP } \{F\}$

As shown above (plus $T \wedge \neg T \implies F$) it is the case that:

$\vdash \{T\} \text{ WHILE } T \text{ DO SKIP } \{F\}$

but the VCs are:

1. $T \implies F$
2. $F \wedge \sim T \implies F$
3. $\sim T \wedge F \implies F$

clearly 1 is false (2 and 3 are true).

The significance is that truth of VCs is a sufficient, but not necessary, condition for the original specification to be proveable.

$[P, Q]$ is the set of commands C such that $\vdash [P]C[Q]$ (i.e. $\{C \mid \vdash [P]C[Q]\}$)

WHILE-law is:

$$\vdash P \wedge S \implies E \geq 0$$

 $[P, P \wedge \sim S] \gg \text{WHILE } S \text{ DO } [P \wedge S \wedge E=n, P \wedge E<n]$

Derivation:

C in $\text{WHILE } S \text{ DO } [P \wedge S \wedge E=n, P \wedge E<n]$
 $\iff C$ in $\{\text{WHILE } S \text{ DO } C' \mid C' \text{ in } [P \wedge S \wedge E=n, P \wedge E<n]\}$
 $\iff C$ in $\{\text{WHILE } S \text{ DO } C' \mid \vdash [P \wedge S \wedge E=n] C' [P \wedge E<n]\}$
 $\iff C$ in $\{\text{WHILE } S \text{ DO } C' \mid \vdash [P] \text{WHILE } S \text{ DO } C' [P \wedge \sim S]\}$
 $\implies \vdash [P] C [P \wedge \sim S]$
 $\iff C$ in $[P, P \wedge \sim S]$