

Reg. Lang. & FA 2005 - Paper 2 Question 9 (AMP)

(a) Specification of M :

- alphabet of input symbols: $\Sigma_1 \cap \Sigma_2$

- set of states: $\text{States}_{M_1} \times \text{States}_{M_2}$

- start state: (s_{M_1}, s_{M_2})

- accepting states: $\{(q, q') \mid q \in \text{Accept}_{M_1} \text{ \& } q' \in \text{Accept}_{M_2}\}$

- transitions: $\Delta_M((q, q'), a) = \{(q_1, q'_1) \mid q_1 \in \Delta_{M_1}(q, a) \text{ \& } q'_1 \in \Delta_{M_2}(q', a)\}$

Note that if M_1, M_2 are deterministic
 $(\forall a, q \in \text{States}_{M_i}, \exists! q' \Delta_{M_i}(q, a) = \{q'\})$
 then so is M .

Thus $a_0 \dots a_{n-1} \in L(M)$

iff $(s_{M_1}, s_{M_2}) \xrightarrow{a_0} (q_1, q'_1) \rightarrow \dots \xrightarrow{a_{n-1}} (q_n, q'_n) \in \text{Accept}_M$
 for some $q_1, q'_1, \dots, q_n, q'_n$

iff $(s_{M_1} \xrightarrow{a_0} q_1 \rightarrow \dots \xrightarrow{a_{n-1}} q_n \in \text{Accept}_{M_1} \text{ for some } q_1, \dots, q_n)$
 $\& (s_{M_2} \xrightarrow{a_0} q'_1 \rightarrow \dots \xrightarrow{a_{n-1}} q'_n \in \text{Accept}_{M_2} \text{ for some } q'_1, \dots, q'_n)$

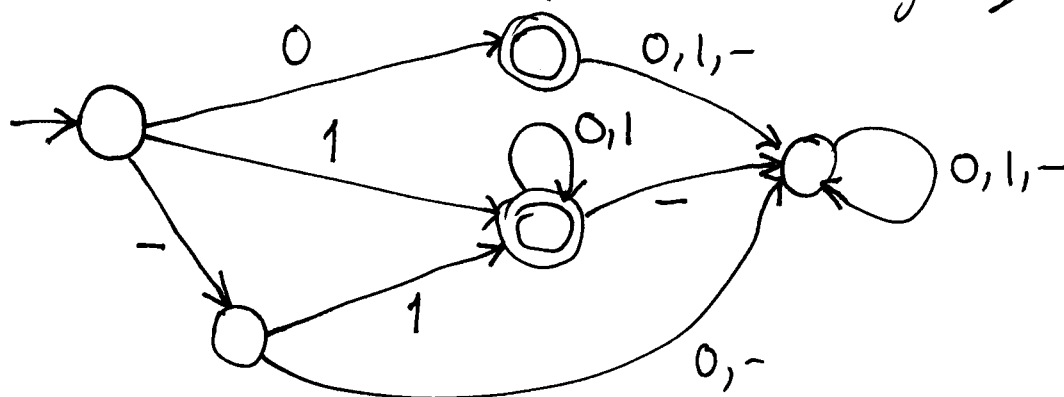
iff $a_0 \dots a_{n-1} \in L(M_1) \text{ \& } a_0 \dots a_{n-1} \in L(M_2)$

Thus $L(M) = L(M_1) \cap L(M_2) = L_1 \cap L_2$.

Note that this holds whether or not M_1, M_2 are deterministic.

(b)

- (i) P generates all strings in $1\{0,1\}^*$
 so N generates all strings in $\{0\} \cup 1\{0,1\}^* \cup -1\{0,1\}^*$
 A DFA that accepts this language:



- (ii) Regular expression(s) determining $\{0\} \cup 1\{0,1\}^* \cup -1\{0,1\}^*$:

$$0 \mid 1(011)^* \mid -1(011)^*$$

- (also $0 \mid (\epsilon \mid -)1(011)^*$, etc.)

- (iii) A CFG is regular if it is either left or right linear. Right linearity means every production is either of the form
- $$x \rightarrow uy \quad (x, y \text{ non-terminals, } u \text{ a string of terminals})$$
- or
- $$x \rightarrow u \quad (x \text{ non-terminal, } u \text{ a string of terminals})$$
- (Left linearity is defined symmetrically.)

1 The languages generated by regular CFGs is equal to the collection of regular languages (i.e. those accepted by some deterministic finite automaton).

2 The CFG in this question is evidently neither left linear (because of the production $N \rightarrow -P$) nor right linear (because of the production $P \rightarrow PO$), so is not regular.

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