## Optimising Compilers 2002

## Solution Notes y2002p8.tex

$$(VAR) \frac{\Gamma[x:t] \vdash x:t,\emptyset}{\Gamma[x:t] \vdash e:t',F}$$

$$(LAM) \frac{\Gamma[x:t] \vdash e:t',F}{\Gamma \vdash \lambda x.e:t \xrightarrow{F} t',\emptyset}$$

$$(APP) \frac{\Gamma \vdash e_1:t \xrightarrow{F''} t',F \qquad \Gamma \vdash e_2:t,F'}{\Gamma \vdash e_1e_2:t',F \cup F' \cup F''}$$

$$(REF) \frac{\Gamma \vdash e:int,F}{\Gamma \vdash ref\ e:intref,F \cup \{A\}}$$

$$(DEREF) \frac{\Gamma \vdash e:int,F \cup \{A\}}{\Gamma \vdash !e:int,F \cup \{R\}}$$

$$(ASS) \frac{\Gamma \vdash e_1:intref,F \qquad \Gamma \vdash e_2:int,F'}{\Gamma \vdash e_1:=e_2:int,F \cup F' \cup \{W\}}$$

$$(COND) \frac{\Gamma \vdash e_1:int,F \qquad \Gamma \vdash e_2:t,F' \qquad \Gamma \vdash e_3:t,F''}{\Gamma \vdash if\ e_1\ then\ e_2\ else\ e_3:t,F \cup F' \cup F''}$$

Types and effects, also issues and problems arising:

- $e_1 \equiv if !g then ref 1 else g$ This typechecks simply and enables one to deduce  $\vdash e_1 : intref, \{A, R\}$ . No problems, note the union in if-then-else.
- $e_2 \equiv \lambda x.if !g then ref x else g$ This also typechecks simply and enables one to deduce  $\vdash e_2 : int \xrightarrow{\{A,R\}} intref, \emptyset$ . No problems, note the latent side effects, and no immediate ones.
- $e_3 \equiv if !g then \lambda x.ref x else \lambda x.g$ This is problematic, unless one adds a clause

(SUB) 
$$\frac{\Gamma \vdash e : t \stackrel{F'}{\rightarrow} t', F}{\Gamma \vdash e : t \stackrel{F''}{\rightarrow} t', F} \text{ (provided } F' \subseteq F'')$$

in which case we can deduce  $\vdash e_3 : int \xrightarrow{\{A\}} intref$ ,  $\{R\}$ . The issue is that the system given above requires the *same* t for the if-then-else consequents and the t here are functions which differ in latent effect (int  $\xrightarrow{\{A\}} intref$  versus int  $\xrightarrow{\{\}} intref$ ).