Solution Notes

[Syllabus: "lexical and syntax analysis"]

- (a) A context-free grammar is a 4-tuple (N, T, S, R) where N and T are disjoint sets of respectively non-terminal and terminal symbols, $S \in N$ is the start symbol, and R is a set of rules of the form $\alpha \to u_1 \dots u_k$ where $\alpha \in N$ and $u_1, \dots, u_n \in (N \cup T)$.
- (b) The grammar is as above, the language it generates that subset of T^* (strings of tokens from T) which can be generated from S using a finite number of rules.
- (c) The grammar is ambiguous because the string 1+1+1 is may be generated in two separate ways:

$$E \rightarrow E + E \rightarrow 1 + E \rightarrow \cdots \rightarrow 1 + 1 + 1$$

and

$$E \rightarrow E + E \rightarrow E + 1 \rightarrow \cdots \rightarrow 1 + 1 + 1$$
.

(d)
$$(N,T,S,R)$$
 where $N=\{E\},T=\{1,2,X,+,*,-\},S=E$ and
$$R=\{E\to 1,E\to 2,E\to X,E\to E+E,E\to E*E,E\to -E\}$$

(e) (i)

```
E ::= F | E+F | E*F
F ::= -P | P
P ::= 1 | 2 | X
```

(e) (ii)

```
E ::= F | F+E | F*E
F ::= -P | P
P ::= 1 | 2 | X
```

(e) (iii)

```
E ::= F | E+F
F ::= -G | G
G ::= P * G | P
P ::= 1 | 2 | X
```

(f) I

```
ParseTree rdE()
{    ParseTree x = rdF();
    while (token == "+") { nexttoken(); x = mkplus(x, rdF(); }
    return x;
}
```

```
ParseTree rdF()
      if (token == "-") { nexttoken(); return mkneg(x, rdG(); }
        else return rdG();
    }
ParseTree rdG()
    { ParseTree x = rdP();
        if (token == "*") { nexttoken(); x = mktimes(x, rdG(); }
        return x;
    }
ParseTree rdP()
      ParseTree x;
        switch (token)
        { case '1':
            case '2':
            case 'X': x = mkprimitivenode(token);
                     nexttoken();
                     return x;
           default: error("expecting 1/2/X");
       }
    }
ParseTree ReadSentence()
    { ParseTree x = rdE();
        if (token != EOF) error("junk after program");
    }
```