

(a) $M_1 =_{\beta} M_2$ holds if there is a finite chain of reductions between M_1 & M_2 :

$$M_1 - \dots - M_2$$

where each $-$ is either \rightarrow or \leftarrow and:

- $M_1 \rightarrow M_2$ holds if M_2 is obtained from M_1 by replacing a subterm of the form

$$\begin{cases} (\lambda x: \tau(M)) N \\ \text{or } (\Lambda \alpha(M)) \tau \end{cases} \quad \text{by} \quad \begin{cases} M[N/x] \\ N[\tau/\alpha] \end{cases}$$

- $M_1 \leftarrow M_2$ holds if $M_1 \rightarrow M_2$ does.

To decide $M_1 =_{\beta} M_2$:

- reduce M_1 & M_2 to β -normal forms N_1 & N_2 (because M_1 & M_2 are typeable, such n.f.'s always exist (by Strong Normalisation result) and are unique (by Church-Rosser result)).
- compare N_1 & N_2 up to α -equivalence (which is decidable).

This doesn't work for untyped terms, because they may have no β -normal form (e.g. $(\lambda x: \tau(xx)) \lambda x: \tau(xx)$ reduces to itself).

1b) We can take

$$I \triangleq \Lambda \alpha (\lambda x: \alpha (\Lambda \alpha_1 (\lambda y: \forall \alpha_2 (\alpha_2 \rightarrow \alpha_1) (y \alpha x))))$$

(i) Putting $\Gamma = [x:\alpha, y:\forall\alpha_2(\alpha_2 \rightarrow \alpha_1)]$, we have

$$\begin{array}{c}
 \frac{}{\Gamma \vdash y: \forall\alpha_2(\alpha_2 \rightarrow \alpha_1)} \text{(var)} \\
 \frac{}{\Gamma \vdash y\alpha: \alpha \rightarrow \alpha_1} \text{(spec)} \quad \frac{}{\Gamma \vdash x: \alpha} \text{(var)} \\
 \frac{}{\Gamma \vdash y\alpha x: \alpha_1} \text{(app)} \\
 \frac{}{x: \alpha \vdash \lambda y: \forall\alpha_2(\alpha_2 \rightarrow \alpha_1)(y\alpha x): (\forall\alpha_2(\alpha_2 \rightarrow \alpha_1)) \rightarrow \alpha_1} \text{(fn)} \\
 \frac{}{x: \alpha \vdash \lambda\alpha_1(\dots) : \omega} \text{(gen)} \\
 \frac{}{\Phi \vdash \lambda x: \alpha(\dots) : \alpha \rightarrow \omega} \text{(fn)} \\
 \frac{}{\Phi \vdash I: \forall\alpha(\alpha \rightarrow \omega)} \text{(gen)}
 \end{array}$$

(ii) Suppose $\Phi \vdash M_i : \tau$ ($i=1,2$).

As mentioned in part (a), we therefore know that $M_i \rightarrow^* N_i$ for (unique) β -normal forms N_i ($i=1,2$). Then from the definition of I we get

$$I\tau M_i \alpha x \rightarrow^* x \tau M_i \rightarrow^* x \tau N_i$$

and $x \tau N_i$ is in normal form, because N_i is.

So if $M_1 =_\beta M_2$, then $I\tau M_1 \alpha x =_\beta I\tau M_2 \alpha x$ so the latter terms reduce to the same normal form, so from above $x \tau N_1 = x \tau N_2$ and hence $N_1 = N_2$. Therefore $M_1 =_\beta N_1 = N_2 =_\beta M_2$, i.e. $M_1 =_\beta M_2$, as required.

Commentary

This question concerns material covered in lectures 6 & 7 of the course.

Part (a) is bookwork.

Part (b) is a new problem. Given the definition of ω , the candidate has to discover a term I satisfying (i): there is an obvious candidate once one uses techniques illustrated by example in the lectures. Part (ii) then follows if one takes the hint; perhaps the hardest part is to realise that N_1 and N_2 are α -equivalent if $x \in N_1$ and $x \in N_2$ are.