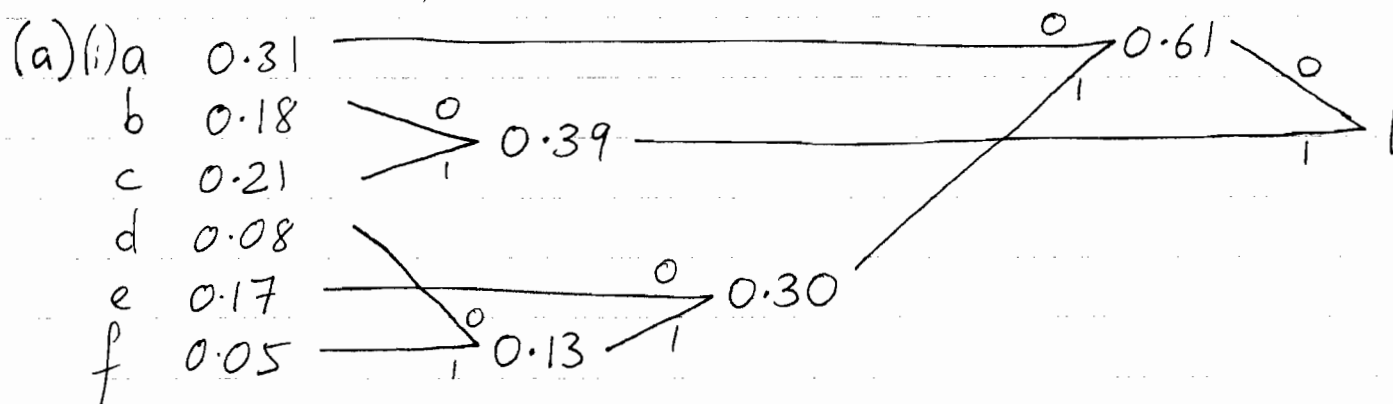


INFO THEORY - paper 8, 2003 q 10

NAD



a	00	2 bits
b	10	2 bits
c	11	2 bits
d	0110	4 bits
e	010	3 bits
f	0111	4 bits

(ii) average bits per symbol = $(0.31 + 0.18 + 0.21) \times 2$
 $+ 0.17 \times 3$
 $+ (0.08 + 0.05) \times 4$
 $= 1.40 + 0.51 + 0.52$
 $= 2.43$

this is greater than the entropy
 (it cannot be less than because this is impossible;
 it cannot be equal because the probabilities are not
 all of the form $\frac{1}{2^n}$, $n \in \mathbb{Z}^+$)

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(b)(i)

$$8+8+8+24+12+4 = 64$$

Y \ X				
	2	3	4	
C	$\frac{1}{8}$	$\frac{3}{8}$	$\frac{1}{16}$	$\frac{9}{16}$
M	$\frac{1}{8}$	$\frac{1}{8}$	$\frac{3}{16}$	$\frac{7}{16}$
	$\frac{1}{4}$	$\frac{1}{2}$	$\frac{1}{4}$	

$$\begin{aligned}\frac{8}{64} &= \frac{1}{8} \\ \frac{12}{64} &= \frac{3}{16} \\ \frac{4}{64} &= \frac{1}{16} \\ \frac{24}{64} &= \frac{3}{8}\end{aligned}$$

$$\begin{aligned}H(X) &= 2 \times \frac{1}{4} \log 4 + \frac{1}{2} \log 2 \\ &= \frac{1}{2} \times 2 + \frac{1}{2} \times 1 \\ &= \frac{3}{2} = 1.500\end{aligned}$$

$$\begin{aligned}H(Y) &= \frac{9}{16} \log \frac{16}{9} + \frac{7}{16} \log \frac{16}{7} \\ &= 0.467 + 0.522 \\ &= 0.989\end{aligned}$$

$$\begin{aligned}H(X, Y) &= 3 \times \frac{1}{8} \log 8 + \frac{3}{8} \log \frac{8}{3} + \frac{3}{16} \log \frac{16}{3} + \frac{1}{16} \log 16 \\ &= 3 \times 0.375 + 0.531 + 0.453 + 0.250 \\ &= 2.359 \quad (\text{or } 2.358 \text{ if you add then round to 3dp})\end{aligned}$$

$$\begin{aligned}H(X|Y) &= H(X, Y) - H(Y) \\ &= 1.370 \quad (\text{or } 1.369)\end{aligned}$$

$$\begin{aligned}H(Y|X) &= H(X, Y) - H(X) \\ &= 0.859 \quad (\text{or } 0.858)\end{aligned}$$

$$\begin{aligned}I(X; Y) &= H(X) + H(Y) - H(X, Y) \\ &= 0.130 \quad (\text{or } 0.129)\end{aligned}$$

(ii) There are several possible answers. The three which require the least re-calculation are:

- (A) both rows identical
- (B) rows multiplicatively related in the ratio $\frac{9}{16} : \frac{7}{16}$
- (C) one row all zeros.

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(A)

8	16	8
8	16	8

$\frac{1}{8}$	$\frac{1}{4}$	$\frac{1}{8}$
$\frac{1}{8}$	$\frac{1}{4}$	$\frac{1}{8}$

$$\begin{aligned} H(X,Y) &= 4 \times \frac{1}{8} \log 8 + 2 \times \frac{1}{4} \log 4 \\ &= \frac{1}{2} \times 3 + \frac{1}{2} \times 2 \\ &= 2.5 \end{aligned}$$

(B)

9	18	9
7	14	7

			$\frac{9}{16}$
			$\frac{7}{16}$
$\frac{1}{4}$	$\frac{1}{2}$	$\frac{1}{4}$	

$$H(X,Y) = H(X) + H(Y) \quad \text{which are the same as in part (i)}$$

$$\begin{aligned} \therefore &= 1.500 + 0.989 \\ &= 2.489 \end{aligned}$$

(C)

16	32	16
0	0	0

			1
			0
$\frac{1}{4}$	$\frac{1}{2}$	$\frac{1}{4}$	

$$\begin{aligned} H(X,Y) &= H(X) + H(Y) \quad \text{where } H(X) \text{ is as in part (i)} \\ &\quad \text{and } H(Y) = 0 \text{ by inspection.} \\ &= 1.500 \end{aligned}$$

Any answer of this form is acceptable

x	2x	x
16-x	32-2x	16-x

$$x \in \{0, 1, 2, \dots, 16\}$$

but only $x \in \{0, 7, 8, 9, 16\}$ allows for easy calculation of $H(X,Y)$

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Anyone who has paid attention in the course can do very well on this question. BUT parts (ii) of both (a) and (b) can be done much more rapidly by a student who understood the material thoroughly. Indeed (b)(ii) will be beyond a student who has not understood but has merely memorised.

- (a) tests parts of Ch 5 & an understanding of Ch 2
(b) tests the core relationships in information theory

MARKING SCHEME

(a)(i) correct method 3
a correct code 2
5

(ii) correct method 1
correct answer (2.43) 1
that it is GREATER THAN (not = nor <) 1
3

(b)(i) $H(x), H(y), H(x,y), H(x|y), H(y|x), I(x;y)$ - 1 mark each 6
correct method for calculating $H(a)$ 1
" " " " " $H(a|b)$ 1
" " " " " $I(a;b)$ 1
9

(ii) a correct table 2
a correct value of $H(x,y)$ for the table 1
3

20