SOLUTION NOTES

Concurrent Systems and Applications 2002 Paper 5 Question 4 (TLH)

- (a) The suspend() method causes the target thread to pause execution immediately. It may do when while it is holding a lock either one acquired explicitly by the application or one used internally in the implementation of the JVM. Either case can result in deadlock. 'Lost wake up' problems can also arise if a thread invokes suspend on itself after determining that it should not proceed into a critical section at that time.
- (b) See example file Barrier. java. Points to note:
 - (i) At all times the simple idiom of a while loop calling wait() has been used along with notifyAll() to wake waiting threads.
 - (ii) At the cost of some extra programming notify() could be used instead note the correspondence between the calls to notifyAll() and the threads that those calls wake.
 - (iii) InterruptedException is propagated since there is no clear way to deal with errors here.
 - (iv) The two methods are entirely symmetric, as you would expect. Each has two sections: during the first the thread competes with those of the same kind, essentially picking which will be paired up next, and then during the second it waits for a partner.
 - (v) This structure means that there is no need for an explicit shared structure to hold the id values they are passed between the two threads concerned in the a and b fields.
- (c) See example files Shop.java, Customer.java and Barber.java. This is a 'classic' concurrency problem, usually attributed to Dijkstra. In a computing context the barbers represent devices performing operations on behalf of clients (the customers). During service there's a 1-1 association between clients and devices. After service the device must wait for the client to retrieve results (leave the chair) before moving to another client. Points to note:
 - (i) The Barrier class is used to pair up customers with barbers.
 - (ii) The shop, as defined here, serves only to identify the barrier to use it does not need to identify the barbers or customers using it.
 - (iii) The haircutFinished and customerLeft fields are protected by the mutual-exclusion lock on the instance of barber they essentially denote when the barber and the customer have respectively finished using the chair.

(iv) It is important to be clear on how the notify methods are used in conjunction with mutual-exclusion locks — that is, that the synchronized regions here in Customer acquire the lock on the associated Barber object

```
Barrier.java:
public class Barrier {
    boolean waitingA = false;
    boolean waitingB = false;
    Object a, b;
    public synchronized Object enterA (Object id)
throws InterruptedException
    {
while (waitingA) {
    wait();
}
waitingA = true;
a = id;
notifyAll();
while (!waitingB) {
    wait();
}
waitingA = false;
notifyAll();
return b;
    }
    public synchronized Object enterB (Object id)
throws InterruptedException
while (waitingB) {
    wait ();
waitingB = true;
b = id;
notifyAll();
while (!waitingA) {
    wait ();
waitingB = false;
notifyAll();
return a;
}
```

```
Shop.java
public class Shop {
    Barrier w = new Barrier ();
}
Customer.java
public class Customer {
    public Barber getHaircut (Shop s)
throws InterruptedException
Barber b = (Barber) (s.w.enterA (this));
synchronized (b) {
    while (!b.haircutFinished) {
b.wait();
    }
}
return b;
    public void leaveChair (Barber b)
throws InterruptedException
    {
synchronized (b) {
    b.customerLeft = true;
    b.notifyAll ();
}
    }
}
Barber.java
public class Barber {
    boolean haircutFinished;
    boolean customerLeft;
    public Customer getCustomer (Shop s)
throws InterruptedException
Customer c;
haircutFinished = false;
customerLeft = false;
c = (Customer) (s.w.enterB (this));
```

```
return c;
    }
    public synchronized void finishedCustomer (Customer c)
throws InterruptedException
    {
    haircutFinished = true;
    notifyAll();
    while (!customerLeft) {
        wait();
}
    }
}
```