

Introduction to Functional Programming 2004
Paper 12 Question 10 (GMB)

(a) We are given the following code.

```
fun cat (b,f) nil      = b
  | cat (b,f) (x::xs) = f(x,(cat (b,f) xs));
```

The categorically-inclined reader will notice that `cat` defines a list catamorphism. The students have seen this function in the notes in a slightly different form, as the function `foldr`.

(i) $'a * ('b * 'a \rightarrow 'a) \rightarrow 'b \text{ list} \rightarrow 'a$

(ii) `fun filter p = cat ([],(fn (x,xs) => if p(x) then x::xs else xs));`

(iii) `fun cmap f = cat ([],(fn (x,xs) => f(x)::xs));`

(b) We are given the following code.

```
fun ana (p,g) b = if p(b) then
                    []
                  else let val (a,b1)=g(b)
                        in
                          a::(ana (p,g) b1)
                        end;
```

The well-read categorically-inclined reader will recognise that `ana` defines a list *anamorphism*.

(i) $('a \rightarrow \text{bool}) * ('a \rightarrow 'b * 'a) \rightarrow 'a \rightarrow 'b \text{ list}$

(ii)

```
fun zip l = ana ((fn (a,b)          => (a=nil) orelse (b=nil)),
                 (fn (x::xs,y::ys) => ((x,y),(xs,ys)))
                ) l;
```

Note that one may have been tempted to define a value rather than a function, but this falls foul of the value polymorphism restriction of SML'97. No candidates would be penalised for this mistake.

(iii)

```
fun amap f = ana ((fn xs          => xs=nil),
                  (fn (x::xs) => (f(x),xs))
                 );
```