

## SOLUTION NOTES

### Logic and Proof 2003 Paper 5 Question 9 (LCP)

(a) This relates to Lecture 4, on ordered binary decision diagrams. A good ordering is  $P_1 < Q_1 < P_2 < Q_2 < \dots < P_i < Q_i < \dots < R$ . Note that  $\phi(k)$  is true unless  $R$  is false and  $P_i \leftrightarrow Q_i$  is true for all  $i > 0$ . The corresponding OBDD will start with an OBDD for  $P_1 \leftrightarrow Q_1$ , modified such that the final link to false instead leads to true, and the final link for true instead leads to a similarly modified OBDD for  $P_2 \leftrightarrow Q_2$ . The OBDD ends with the trivial one for  $R$ . Ideally, the candidate should supply a diagram such as the one attached (see the file logic-figure.pdf). That one is for the case  $k = 2$ , but the general case is easy to sketch using pen and paper.

(b) One possible ordering is  $P_1 < P_2 < \dots < Q_1 < Q_2 < \dots < R$ . The OBDD will be exponential in  $k$  because the outcome for  $Q_i$  depends upon the value of  $P_i$ . Thus, the OBDD below  $Q_1$  will have to depend upon the values of  $P_1, \dots, P_k$ . There are  $2^k$  possible combinations of those values and therefore at least  $2^k$  decision nodes for  $Q_1$ .

(c) This relates to Lecture 7, resolution. Negating the formula gives

$$(P_1 \leftrightarrow Q_1) \wedge \dots \wedge (P_k \leftrightarrow Q_k) \wedge \neg R.$$

Since  $P \leftrightarrow Q \simeq (P \rightarrow Q) \wedge (Q \rightarrow P)$ , the clauses are

$$\{\neg P_1, Q_1\} \quad \{\neg Q_1, P_1\} \quad \dots \quad \{\neg P_k, Q_k\} \quad \{\neg Q_k, P_k\} \quad \{\neg R\}.$$

(d) The Davis-Putnam procedure would reduce it to the empty set of clauses, indicating that  $\neg\phi(k)$  is satisfiable and that  $\phi(k)$  is not valid. This is the correct outcome, as the formula is obviously falsifiable. The first step of Davis-Putnam would act on the pure literal  $\neg R$  by deleting the clause  $\{\neg R\}$ . Exponentially many case splits then would be necessary. For example, a case split on  $P_1$  would, in the true case, delete the clause  $\{\neg Q_1, P_1\}$  and reduce the clause  $\{\neg P_1, Q_1\}$  to  $\{Q_1\}$ , which would then be deleted as  $Q_1$  would be pure. Further case splits on the remaining variables would eliminate all the clauses. The false case would be similar.