Discrete Mathematics - Question 2000

The following fragment of ML implements Stein's algorithm for evaluating the Greatest Common Divisor, (a,b), of two natural numbers, a and b:

```
fun stein a b c =
  if a = b then a * c
  else
     if (a \mod 2) = 0 then
       if (b mod 2) = 0 then stein (a div 2) (b div 2) (c * 2)
       else stein (a div 2) b c
     else
       if (b \mod 2) = 0 then stein a (b \operatorname{div} 2) c
       else
         if a > b then stein (a - b) b c
         else stein (b - a) a c;
fun gcd a b = stein a b 1;
The following fragment implements the same algorithm in Java:
static int stein (int a, int b, int c) {
       while (a != b)
              switch (((a & 1) << 1) + (b & 1)) {
              case 0:
                     a >>= 1; b >>= 1; c <<= 1;
                    break;
              case 1:
                    a >>= 1;
                    break;
              case 2:
                    b >>= 1;
                    break;
              case 3:
                    if (a > b) a -= b;
                    else \{a = b-a; b -= a; \};
             };
       return a * c;
}
static int gcd (int a, int b) {
      return stein (a, b, 1);
}
Prove that, at each iteration within the Stein algorithm, the product (a,b).c remains invariant.
```

Prove that, at each iteration within the Stein algorithm, the product (a,b).c remains invariant.

Observing that the procedure starts with c=1 and concludes by returning a.c when a=b, deduce that the algorithm correctly calculates the Greatest Common Divisor.

Show also that after two iterations the product a.b is reduced by at least a factor of 2.

Deduce that Stein's algorithm is at least as efficient as Euclid's algorithm.

[8 marks]

[6 marks]

Answer

```
 (2u,2v) = 2.(u,v), \ (2u,2v+1) = (u, 2v+1), \ (2u+1,2v) = (2u+1,v), \ (u,v) = (u-v,u) = (u-v,v).  Invariant starts as (a,b).1 and ends as (a,a).c = a.c which is the final value returned.  u.v \le 2u.2v/2, \ u.(2v+1) \le 2u.(2v+1)/2, \ (2u+1).v \le (2u+1).2v/2, \ (u-v)(2v+1) = (2u-2v)(2v+1)/2 \le (2u+1)(2v+1)/2.  If a < 2^n and b < 2^n then a.b < 2^{2n} and the algorithm concludes in at most 4n steps. Hence O(\log a).
```

- 6 cases: a=b, alberer, aeverabodd, adddbeve, albodd asbjets
afbodd bka.