

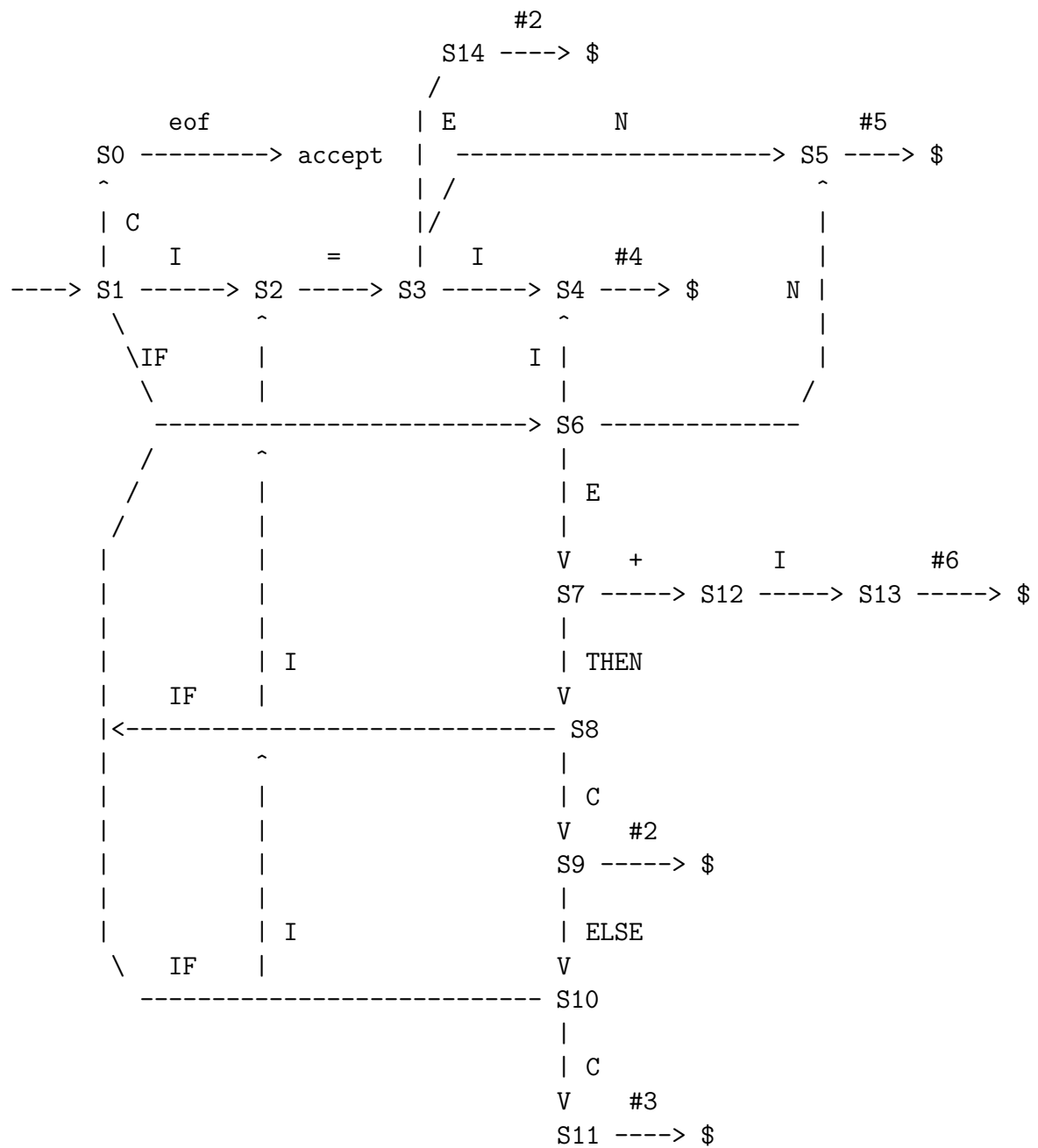
SOLUTION NOTES**Compiler Construction 2003 Paper 4 Question 1 (AM)**

[This is a question on parsing.]

- Firstly the grammar is ambiguous, we will assume the normal resolution to “dangling else”, i.e. that an **else** matches with the innermost of a choice of **ifs**.
- For a recursive descent parser we must adapt E to remove left-recursion, and also left-factor the C cases.

```
• void rdS()
  { lex(); // assume lex() sets global 'token'
    rdC(); if (token != EOF) error();
  }
void rdE()
  { switch (token)
    {
default: error();
case I:
case N: lex();
        if (token == '+') { lex(); rdE(); }
    }
  }
void rdC()
  { switch (token)
default: error();
case I: lex(); if (token != '=') error();
        lex(); rdE(); return;
case IF: lex(); rdE(); if (token != THEN) error();
         lex(); rdC(); if (token != ELSE) return;
         lex(); rdC(); return;
  }
```

- The CSFM for SLR (number productions #0...#6)



Note the shift-reduce conflict for state S9 which, if resolved in favour of shift (i.e. SLR(1)), gives the normal ambiguity resolution above.