1999 Introduction to Functional Programming

Model Answer

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1. datatype 'a tree = Lf
                      | Br of 'a*'a tree*'a tree;
2. fun post Lf
     | post (Br(v,t1,t2)) = post t1 @ post t2 @ [v];
3. fun post2 t = let fun postord (Lf, vs) = vs
                         | postord (Br(v,t1,t2), vs) =
                                    postord (t1, postord (t2, v::vs))
                  in
                     postord(t,[])
                  end;
4. We prove that \forall t, l.post(t) @ l = postord(t, l), by induction on t.
   (a) Base case. t = Lf.
                              post(Lf) @ l = [] @ l
                                           = postord(Lf, l)
   (b) Inductive case. t = Br(v, t1, t2).
                  post(Br(v,t1,t2)) @ l = (post(t1) @ (post(t2) @ [v])) @ l
                                       = post(t1) @ (post(t2) @ ([v] @ l))
                                        = post(t1) @ (post(t2) @ (v :: l))
                                        = post(t1) @ (postord(t2, v :: l))
                                        = post(t1, postord(t2, v :: l))
                                        = postord(Br(v, t1, t2), l)
5. datatype 'a mtree = MTree of 'a * ('a mtree) list;
6. fun map(f,[])=[]
     | map(f,x::xs)=f(x)::map(f,xs);
  fun foldr(f,[],e)
     | foldr(f,x::xs,e) = f(x,foldr(f,xs,e));
  fun post3 (MTree(n,[])) = [n]
     | post3 (MTree(n,ts)) = foldr(op@,map(post3,ts),[])@[n];
```