

Computer Systems Modelling

RJG

[Relates to Poisson Process]

①. The necessary conditions are

(i)  $N(0) = 0$

(ii) The numbers of events in disjoint time intervals are independent

(iii) The distribution of the number of events in a given interval depends only on the interval's length (and not its location).

(iv)  $P(N(h) = 0) = 1 - \lambda h + o(h)$  as  $h \rightarrow 0$

(v)  $P(N(h) = 1) = \lambda h + o(h)$  as  $h \rightarrow 0$

(vi)  $P(N(h) > 1) = o(h)$  as  $h \rightarrow 0$

② Divide  $[0, t]$  into  $n$  equal length sub-intervals - each of length  $t/n$

Now for large  $n$

$$P(\text{a sub-interval contains 1 event}) \approx \lambda \frac{t}{n} \approx 1 - P(\text{a sub-interval contains no events})$$

numbers of events in sub-intervals are independent

Total number of events occurring is given by a Binomial

random variable with parameters  $n$  and  $p = \lambda t/n$

So, for large  $n$ , the number of events in  $(0, t)$ ,  $N(t)$ , follows a

Poisson distribution with mean  $np = n \times t/n = \lambda t$ . So  $P(N(t) = j) = \frac{e^{-\lambda t} (\lambda t)^j}{j!}$   
 $j = 0, 1, 2, \dots$

③ We have that

$$P(X_1 > t) = P(N(t) = 0) = e^{-\lambda t}$$

$$\text{So } X_1 \sim \text{Exp}(\lambda)$$

and

$$\begin{aligned} P(X_2 > t \mid X_1 = s) &= P(0 \text{ events in } (s, t+s) \mid X_1 = s) \\ &= P(0 \text{ events in } (s, t+s)) \\ &= e^{-\lambda t} \end{aligned}$$

So,  $X_1$  and  $X_2$  are independent  
 with  $X_2 \sim \text{Exp}(\lambda)$ .

(4) Set  $S_n = \sum_{i=1}^n X_i$

Then,

$$P(S_n \leq t) = P(N(t) \geq n)$$

$$= \sum_{j=n}^{\infty} e^{-\lambda t} \frac{(\lambda t)^j}{j!}$$

Differentiate wrt  $t$  to find density,

$f_n(t)$ , of  $S_n(t)$  as

$$f_n(t) = \sum_{j=n}^{\infty} (\lambda) e^{-\lambda t} \frac{(\lambda t)^{j-1}}{(j-1)!} - \sum_{j=n}^{\infty} \lambda e^{-\lambda t} \frac{(\lambda t)^j}{j!}$$

$$= \sum_{j=n}^{\infty} \lambda e^{-\lambda t} \frac{(\lambda t)^{j-1}}{(j-1)!} - \sum_{j=n}^{\infty} \lambda e^{-\lambda t} \frac{(\lambda t)^j}{j!}$$

$$= \lambda e^{-\lambda t} \frac{(\lambda t)^{n-1}}{(n-1)!}$$

(5) Algorithm

① Let  $t = 0$ ,  $I = 0$

② Generate (pseudo) random number,  $U$

③  $t = t - \frac{1}{\lambda} \log U$ , If  $t > T$ , stop

④  $I = I + 1$ ,  $S(I) = t$

⑤ Go to step ②

$I$  is number of event  
 $S(I)$  is time of  $I^{\text{th}}$  event.