

Question 2 Notes

A development from 1998 P11 q8, with some basic regular algebra to get into the swing, then an example. They may tackle the q. by reasoning directly, in which case play it by ear. Provided they partition as asked it will come out OK.

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SolutionDefinitions

- a)  $E + F = \{w \in S^* \mid w \in E \text{ or } w \in F\}$
- b)  $EF = \{ww' \in S^* \mid w \in E, w' \in F\}$
- c)  $E^* = \{\varepsilon\} + \{w_1 w_2 \dots w_n \mid w_i \in E, n \in \mathbb{N}^+, \textcircled{3} 1 \leq i \leq n\}$
- d)  $= 1 + \{w_1 \mid w_1 \in E\}$   
 $+ \{w_1 w_2 \dots w_n \mid n \geq 2, w_i \in E, 1 \leq i \leq n\}$   
 $= 1 + E + E \cdot \{w_2 \dots w_n \mid n \geq 2, w_i \in E, 2 \leq i \leq n\}$

Question 2 Solution ctd)

$$i) \quad \therefore E^* = 1 + EE^*$$

(easier with a better definition, as in notes)

$$ii) \quad E(FE)^* = \{ w_0 (v_1 w_1 v_2 w_2 \dots v_n w_n) \mid n \geq 0, \\ w_0 \in E, w_i \in E, v_i \in F, 1 \leq i \leq n \}$$

$$= \{ (w_0 v_1 w_1 v_2 \dots w_{n-1} v_n) w_n \mid n \geq 0, \\ w_n \in E, w_{i-1} \in E, v_i \in F, 1 \leq i \leq n \}$$

$$= (EF)^* E \quad (4)$$

Kleene's Theorem

An event accepted by some DFA is regular over input alphabet.

(1)

Given  $M^* = \begin{pmatrix} (A+BD^*C)^* & A^*B(D+CA^*B)^* \\ D^*C(A+BD^*C)^* & (D+CA^*B)^* \end{pmatrix}$

# Maths for Computation Theory

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## Question 2      Solution ctd)

where

$$M = \begin{array}{cc|cc} & \alpha & \beta & \gamma & \delta \\ \alpha & b & 0 & a & 0 \\ \beta & 0 & a & 0 & b \\ \hline \gamma & 0 & a & 0 & b \\ \delta & b & 0 & a & 0 \end{array}$$

Require to compute  $A^*B(D + CA^*B)^*$ .

Here  $A^* = \begin{pmatrix} b^* & 0 \\ 0 & a^* \end{pmatrix}$ ,  $A^*B = \begin{pmatrix} b^*a & 0 \\ 0 & a^*b \end{pmatrix}$

$$CA^*B = \begin{pmatrix} 0 & a \\ b & 0 \end{pmatrix} \begin{pmatrix} b^*a & 0 \\ 0 & a^*b \end{pmatrix}$$

$$= \begin{pmatrix} 0 & aa^*b \\ bb^*a & 0 \end{pmatrix}$$

$$\therefore D + CA^*B = \begin{pmatrix} 0 & b + aa^*b \\ a + bb^*a & 0 \end{pmatrix}$$

$$= \begin{pmatrix} 0 & a^*b \\ b^*a & 0 \end{pmatrix} \quad \text{by i)}$$

# Maths for Computation Theory

(6)

## Question 2 Solution ctd)

Hence

$$(D + CA^*B)^* = \begin{pmatrix} (a^*bba)^* & a^*b(baa^*b)^* \\ ? & ? \end{pmatrix}$$

and the required matrix

$$= \begin{pmatrix} b^*a & 0 \\ 0 & a^*b \end{pmatrix} \begin{pmatrix} (a^*bba)^* & a^*b(baa^*b)^* \\ ? & ? \end{pmatrix}$$

$$= \begin{pmatrix} b^*a(a^*bba)^* & b^*aa^*b(baa^*b)^* \\ ? & ? \end{pmatrix}$$

$$\text{Hence event accepted} = (M^*)_{\alpha\beta} + (M^*)_{\alpha\delta}$$

$$= b^*a \left\{ (a^*bba)^* + a^*b(baa^*b)^* \right\}$$

$$= b^*a(a^*bba)^* \{1 + a^*b\} \quad \text{by ii)}$$