

1999

p6a6  
AM  
p13q7

# Compiler Construction

Compiling Techniques cmptech4.tex

(a) It is desired to obtain an unambiguous context-free grammar  $G'$  which generates the same strings as the grammar  $G$  with start symbol  $S$ .

$$\begin{aligned} S &\rightarrow E \\ E &\rightarrow (E) \mid [E] \mid E * E \mid a \mid b \mid c \\ (E) &\rightarrow (+E) \\ [E] &\rightarrow [-E] \end{aligned}$$

Define a suitable  $G'$  or explain why  $G$  already satisfies the criterion. [5 marks]

(b) Write a context-free (Type 2) grammar which describes floating point numbers of the form  $[\pm]dd^*[.d^*][e[\pm]dd^*]$  where  $d$  stands for decimal digit and  $d^*$  stands for zero or more decimal digits.  $[\dots]$  means that the enclosed item is optionally present in the floating point number. [7 marks]

(c) Explain, giving brief reasons, whether it would be possible to construct a regular (Type 3) grammar which recognised the same syntax of numbers as your answer to part (b). [1 marks]

(d) Sketch a recursive descent parser for the grammar  $H$  with start symbol  $S$ :

$$\begin{aligned} P &\rightarrow 1 \mid 2 \mid (E) \\ E &\rightarrow P \mid E - P \\ S &\rightarrow E \text{ eof} \end{aligned}$$

assuming a routine `lex()` which sets variable `token` to one of '1', '2', '(', ')', '-', or eof. [7 marks]

~~Is this too long?~~

## Model Answer

(a)

$$\begin{aligned} S &\rightarrow E \\ E &\rightarrow F \mid E * F \\ F &\rightarrow (+E) \mid [-E] \mid a \mid b \mid c \end{aligned}$$

(b)

$$\begin{aligned} S &\rightarrow +F \mid -F \mid F \\ F &\rightarrow DS \mid DS . \mid DS . DS \mid \\ &\quad DS E \mid DS . E \mid DS . DS E \\ E &\rightarrow e + DS \mid e - DS \mid e DS \\ DS &\rightarrow D \mid D DS \\ D &\rightarrow 0 \mid 1 \mid \dots \mid 9 \end{aligned}$$

(c) Yes, the grammar is regular, it is only the expression of it above in (b) which is Type 2, hence there exists a Type 3 form

A → a  
A → a B

(d)

```
void RdP()
{ switch (token)
  { case '(': lex(); RdE();
    if (token != ')') error("expected ')");
    lex(); return;
    case '1': lex(); return;
    case '2': lex(); return;
    default: error("unexpected token");
  }
}

void RdE()
{ RdP();
  for (;;) switch (token)
  { case '-': lex(); RdP(); continue;
    default: return;
  }
}

void RdS()
{ RdE();
  switch (token)
  { case 'eof': return;
    default: error("just (not eof) after E");
  }
}
```

Lose lots of marks for calling RdE() as first call within RdE(). Gain a mark for dealing with eof correctly.