# Deterministic Dynamic Programming IV: Practical and Computational Aspects

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## **Outline**

Algorithm and Implementation Issues

2 Implementable Algorithm(s)

# Algorithm and Implementation I

### Previously ...

- We had taken the description of the optimal solution of the RCK problem as far as possible.
- ullet In this example, we were able to describe the properties of v and  $\pi$  with features that increase with additional regularity assumptions on primitives, U and f.
- To actually solve for the optimal outcomes, we need to resort to numerical approximation and computation.
- First we need to set up a strategy for approximation and computation — i.e. develop an algorithm.

## Algorithm and Implementation II

Let  $\epsilon > 0$ . Algorithm:

- **①** Pick some initial guess for  $v_n: X \to \mathbb{R}$ , where, n = 0.
- Solve and evaluate the Bellman operator:

$$v_{n+1}(k) = \max_{k' \in \Gamma(k)} \{ U[f(k) - k'] + \beta v_n(k') \},$$

for every  $k \in X$ . Also store:

$$\pi_n(k) = \arg\max_{k' \in \Gamma(k)} \{ U[f(k) - k'] + \beta v_n(k') \}.$$

- **3** Calculate distance between consecutive updates:  $d(v_{n+1}, v_n)$ .
- While  $d(v_{n+1}, v_n) \ge \epsilon$ , repeat Steps 2-4.

# Algorithm and Implementation III

#### Implementation issues:

- How to represent state space  $X = [\underline{k}, \overline{k}]$  (infinite set) on computer (finite storage)?
- How to represent  $v_n$ , n = 0, 1, ..., each a function, an element of an infinite dimensional space?
- How to compute the "max" operator in Step 2?
- How to represent and store  $\pi_n$ , also a function?
- What is the appropriate metric *d*?

# Implementable Algorithm(s) I

Let  $\epsilon > 0$ . Algorithm:

- **①** Discretized state space  $\hat{X} = [\underline{k} < \cdots < \overline{k}]$  on computer.
- 2 Pick some initial approximate guess for  $\hat{v}_n: \hat{X} \to \mathbb{R}$ , where, n=0.
- Solve and evaluate the Bellman operator (approximate maximization):

$$\hat{v}_{n+1}(k) = \hat{T}(\hat{v}_n)(k) = \max_{k' \in \Gamma(k)} \{ U[f(k) - k'] + \beta \hat{v}_n(k') \},$$

for every  $k \in \hat{X}$ . Also store:

$$\hat{\pi}_n(k) = \arg\max_{k' \in \Gamma(k)} \{ U[f(k) - k'] + \beta \hat{v}_n(k') \}.$$

- Calculate distance between consecutive updates:  $d(\hat{v}_{n+1}, \hat{v}_n)$ .
- **5** While  $d(\hat{v}_{n+1}, \hat{v}_n) \geq \epsilon$ , repeat Steps 2-4.

# Implementable Algorithm(s) II

v = T(v) is true, but unknown and nonanalytic value function.

Two other sources of errors, on top of  $\epsilon > 0$ :

- ullet Step 2: Given  $\hat{X}$ , error from function approximation,  $d(\hat{v},v)$ ; and
- Step 3: (Aggregate) error arising from approximate maximization scheme:  $d(\hat{T}(v), T(v))$ .

# Implementable Algorithm(s) III

We will consider two ways to tackle Steps 2 and 3:

- Method A:
  - Step 2A: Represent as finite  $\hat{v}_n=[\hat{v}_n(\underline{k})<\cdots<\hat{v}_n(\overline{k})]$ , each  $n=0,1,\ldots$
  - Step 3A: Given Step 2A, and  $\hat{X}$  a finite set,  $\max$  operator in  $\hat{T}$  is just a "Table-lookup" maximization problem.

# Implementable Algorithm(s) IV

#### Method B:

- Step 2B: Represent infinite  $\hat{v}_n$  with known families of analytic functions. But store finite list  $[\hat{v}_n(\underline{k}) < \cdots < \hat{v}_n(\overline{k})]$ , each  $n=0,1,\ldots$
- Step 3B: Given Step 2A, and  $\hat{X}$  a finite set, we can always interpolate for points not contained in  $[\hat{v}_n(\underline{k})<\cdots<\hat{v}_n(\overline{k})]$ . Apply gradient/non-gradient based optimization algorithms to evaluate  $\max$  operator in  $\hat{T}$ .

# Implementable Algorithm(s) V

- Method A is known as the discretization method. Requires a lot of points in  $\hat{X}$  to get an accurate approximate solution. Costly if X is multidimensional state space!
- Method B:
  - often implemented using projection methods on bounded and continuous function spaces.
  - Optimization doable in a continuous way since the (approximate) objective function is now continuous.
  - Requires less points in  $\hat{X}$  to get accurate approximate solutions.
  - In practice, curse of dimensionality still a problem.

# Implementable Algorithm(s) VI

- Now we look at Method A first.
- Then we consider Method B in the next set of lectures.

... Go to attached set of Codes for this week.